Analysis and Design of Algorithms Lab Assignment -7

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Dijkstra's shortest path algorithm

Code

```
// Keep Changing....@Vi
// Dijkstra's single source shortest path algorithm.
#include <bits/stdc++.h>
using namespace std;
#define V 9
int minDistance(int dist[], bool sptSet[])
    int min = INT_MAX, min_index;
    for (int v = 0; v < V; v++)
        if (sptSet[v] == false && dist[v] <= min)</pre>
            min = dist[v], min_index = v;
    return min_index;
void printSolution(int dist[], int parent[])
    cout<<"Vertex \t Dist. from Source \t Parent\n";</pre>
    for (int i = 0; i < V; i++)
        cout << i << "\t\t" << dist[i] << "\t\t" << parent[i] << endl;</pre>
void dijkstra(int graph[V][V], int src)
    int parent[V];
    // dist[i] will hold the shortest distance from src to i
    int dist[V];
    // sptSet[i] will be true if vertex i is included in shortest path tree
    bool sptSet[V];
    // Initialize all distances as INT_MAX (infinite) and stpSet[] as false
    for (int i = 0; i < V; i++)
        dist[i] = INT_MAX, sptSet[i] = false;
```

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```
dist[src] = 0;
    parent[src] = -1;
    for (int count = 0; count < V - 1; count++)</pre>
        // Pick the minimum distance vertex
        int u = minDistance(dist, sptSet);
        sptSet[u] = true;
        // Update dist value of the adjacent vertices
        for (int v = 0; v < V; v++)
            // Update dist[v] only if is not in sptSet
            if (!sptSet[v] && graph[u][v] && dist[u] != INT_MAX
                && dist[u] + graph[u][v] < dist[v])
                dist[v] = dist[u] + graph[u][v];
                parent[v] = u;
    printSolution(dist, parent);
int main()
    int graph[V][V] = {{0, 4, 0, 0, 0, 0, 0, 8, 0},
                        {4, 0, 8, 0, 0, 0, 0, 11, 0},
                        \{0, 8, 0, 7, 0, 4, 0, 0, 2\},\
                        \{0, 0, 7, 0, 9, 14, 0, 0, 0\},\
                        \{0, 0, 0, 9, 0, 10, 0, 0, 0\},\
                        \{0, 0, 4, 14, 10, 0, 2, 0, 0\},\
                        \{0, 0, 0, 0, 0, 2, 0, 1, 6\},\
                        {8, 11, 0, 0, 0, 0, 1, 0, 7},
                        \{0, 0, 2, 0, 0, 0, 6, 7, 0\}\};
    dijkstra(graph, 0);
    return 0;
```

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Output

Analysis

- Time Complexity of the implementation is $O(V^2)$.
- If the input graph is represented using adjacency list, it can be reduced to $O(E \log V)$ with the help of binary heap.
- Dijkstra's algorithm doesn't work for graphs with negative weight cycles, it may give correct results for a graph with negative edges. For graphs with negative weight edges and cycles, *Bellman–Ford algorithm* can be used.