

STEI - Institut Teknologi Bandung

## Modul 2

# Bits Representation

Representation and Operations

EL3011 Arsitektur Sistem Komputer



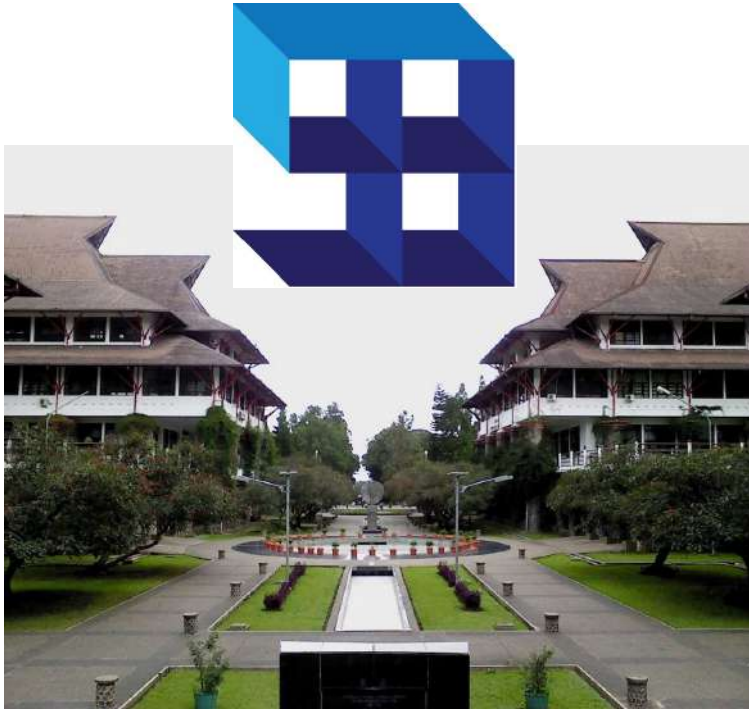
# Contents:

1. Representing information as a bits: <https://www.youtube.com/watch?v=3OuiZ2cAAvo>
2. Bit-level Manipulations: <https://www.youtube.com/watch?v=-aWy5NrYkrl>
3. Representations in memory, pointers, and strings: <https://youtu.be/DHhMuKpZHWY>

**Reading:** Randal E. Bryant, David R. H, Computer Systems A Programmer's Perspective 3<sup>rd</sup> ed [CSAPP], Chapter 2 Representing and Manipulating Information, 21. Information Storage

This module adopted from 15-213 Introduction to Computer Systems Lecture, Carnegie Mellon University, 2020





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## Modul 2. Bits Representation

### 2.1. Representing Information as a Bits

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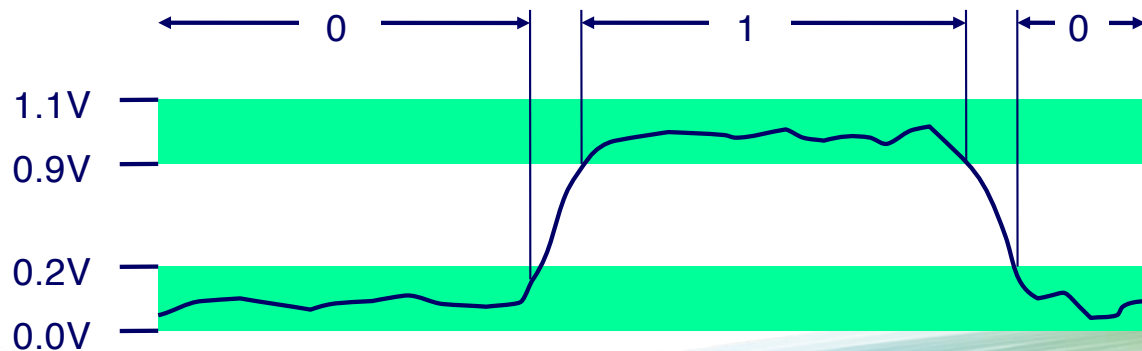
# Everything is bits

- Each bit is 0 or 1
- By encoding/interpreting sets of bits in various ways
  - Computers determine what to do (instructions)
  - ... and represent and manipulate numbers, sets, strings, etc...
- Why bits? Electronic Implementation



# Everything is bits

- Each bit is 0 or 1
- By encoding/interpreting sets of bits in various ways
  - Computers determine what to do (instructions)
  - ... and represent and manipulate numbers, sets, strings, etc...
- Why bits? Electronic Implementation
  - Easy to store with bistable elements
  - Reliably transmitted on noisy and inaccurate wires



For example, can count in binary

- Base 2 Number Representation

- Represent  $13579_{10}$  as  $11010100001011_2$
- Represent  $1.20_{10}$  as  $1.0011001100110011[0011]..._2$
- Represent  $1.3579 \times 10^4$  as  $1.1010100001011_2 \times 2^{13}$



# Encoding Byte Values

- Byte = 8 bits
  - Binary 00000000<sub>2</sub> to 11111111<sub>2</sub>
  - Decimal: 0<sub>10</sub> to 255<sub>10</sub>
  - Hexadecimal 00<sub>16</sub> to FF<sub>16</sub>
    - Base 16 number representation
    - Use characters '0' to '9' and 'A' to 'F'
    - Write FA1D37B<sub>16</sub> in C as
      - 0xFA1D37B
      - 0xfa1d37b

Hex	Decimal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
B	11	1011
C	12	1100
D	13	1101
E	14	1110
F	15	1111

40132: 1001 1100 1100 0100  
                   9      C      C      4



## Example Data Representations

C Data Type	Typical 32-bit	Typical 64-bit	x86-64
<b>char</b>	1	1	1
<b>short</b>	2	2	2
<b>int</b>	4	4	4
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<b>float</b>	4	4	4
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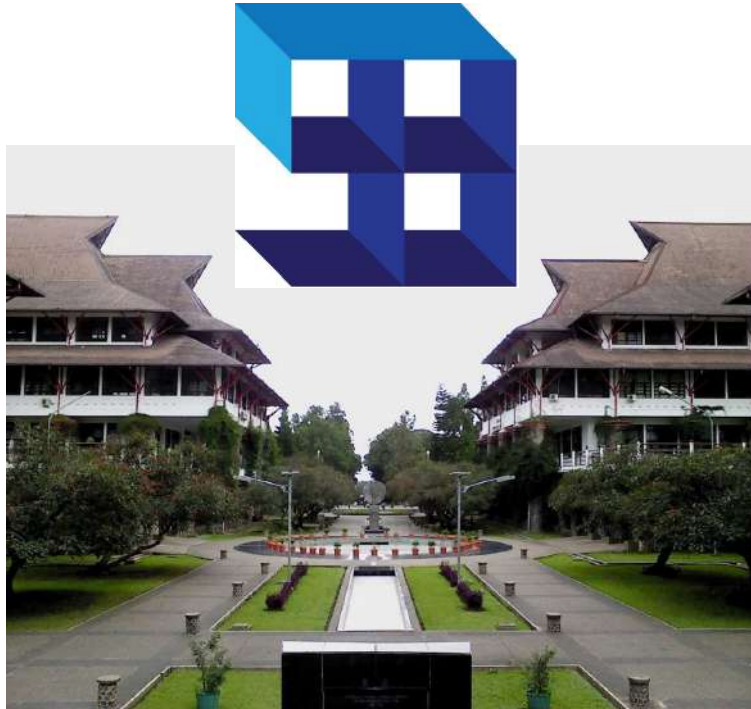




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## Modul 2. Bits Representation

### 2.2. Bit-level Manipulations

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# Boolean Algebra

- Developed by George Boole in 19th Century
  - Algebraic representation of logic
    - Encode “True” as 1 and “False” as 0

And

- $A \& B = 1$  when both  $A=1$  and  $B=1$

$\&$	0	1
0	0	0
1	0	1

Or

- $A | B = 1$  when either  $A=1$  or  $B=1$

$ $	0	1
0	0	1
1	1	1

Not

- $\sim A = 1$  when  $A=0$

$\sim$	
0	1
1	0

Exclusive-Or (Xor)

- $A \wedge B = 1$  when either  $A=1$  or  $B=1$ , but not both

$\wedge$	0	1
0	0	1
1	1	0



# General Boolean Algebras

- Operate on Bit Vectors
  - Operations applied bitwise

01101001	01101001	01101001	01101001
<u>&amp; 01010101</u>	<u>  01010101</u>	<u>^ 01010101</u>	<u>~ 01010101</u>
01000001	01111101	00111100	10101010

- All of the Properties of Boolean Algebra Apply



## Example: Representing & Manipulating Sets

- Representation

- Width  $w$  bit vector represents subsets of  $\{0, \dots, w-1\}$
- $a_j = 1$  if  $j \in A$

- 01101001       $\{0, 3, 5, 6\}$

- 76543210

- 01010101       $\{0, 2, 4, 6\}$

- 76543210

- Operations

- |     |                      |          |                        |
|-----|----------------------|----------|------------------------|
| • & | Intersection         | 01000001 | $\{0, 6\}$             |
| •   | Union                | 01111101 | $\{0, 2, 3, 4, 5, 6\}$ |
| • ^ | Symmetric difference | 00111100 | $\{2, 3, 4, 5\}$       |
| • ~ | Complement           | 10101010 | $\{1, 3, 5, 7\}$       |



# Bit-Level Operations in C

- Operations  $\&$ ,  $|$ ,  $\sim$ ,  $\wedge$  Available in C
  - Apply to any “integral” data type
    - long, int, short, char, unsigned
  - View arguments as bit vectors
  - Arguments applied bit-wise
- Examples (Char data type)
  - $\sim 0x41 \rightarrow 0xBE$ 
    - $\sim 0100\ 0001_2 \rightarrow 1011\ 1110_2$
  - $\sim 0x00 \rightarrow 0xFF$ 
    - $\sim 0000\ 0000_2 \rightarrow 1111\ 1111_2$
  - $0x69 \& 0x55 \rightarrow 0x41$ 
    - $0110\ 1001_2 \& 0101\ 0101_2 \rightarrow 0100\ 0001_2$
  - $0x69 | 0x55 \rightarrow 0x7D$ 
    - $0110\ 1001_2 | 0101\ 0101_2 \rightarrow 0111\ 1101_2$

Hex	Decimal	Binary
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6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
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C	12	1100
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E	14	1110
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# Contrast: Logic Operations in C

- Contrast to Bit-Level Operators
  - **Logic Operations: &&, ||, !**
    - View 0 as “False”
    - Anything nonzero as “True”
    - Always return 0 or 1
    - Early termination
- Examples (char data type)
  - `!0x41` → `0x00`
  - `!0x00` → `0x01`
  - `!!0x41` → `0x01`
  - `0x69 && 0x55` → `0x01`
  - `0x69 || 0x55` → `0x01`
  - `p && *p` (avoids null pointer access)

Watch out for && vs. & (and || vs. |)...  
Super common C programming pitfall!



# Shift Operations

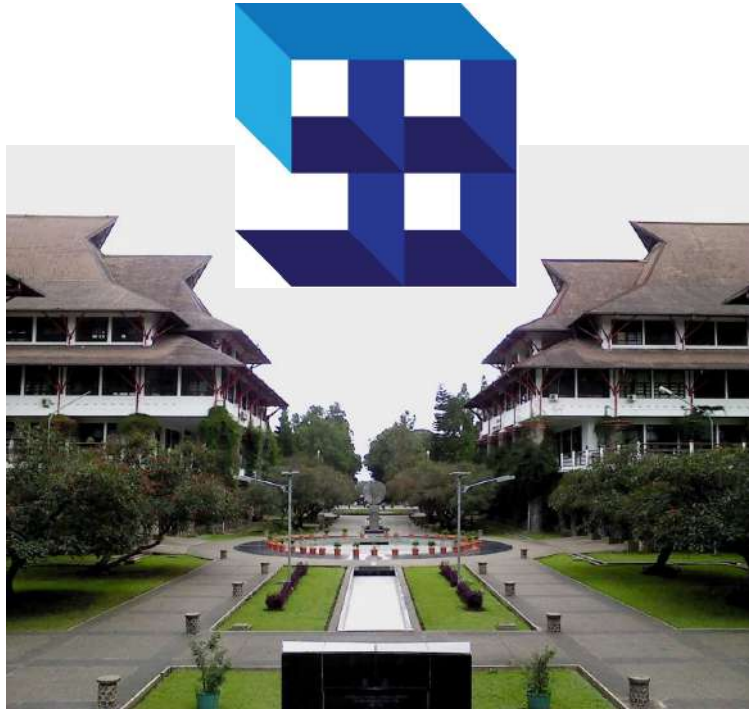
- Left Shift:  $x \ll y$ 
  - Shift bit-vector  $x$  left  $y$  positions
    - Throw away extra bits on left
    - Fill with 0's on right
- Right Shift:  $x \gg y$ 
  - Shift bit-vector  $x$  right  $y$  positions
    - Throw away extra bits on right
  - Logical shift
    - Fill with 0's on left
  - Arithmetic shift
    - Replicate most significant bit on left
- Undefined Behavior
  - Shift amount  $< 0$  or  $\geq$  word size

Argument $x$	01100010
$\ll 3$	00010000
Log. $\gg 2$	00011000
Arith. $\gg 2$	00011000

Argument $x$	10100010
$\ll 3$	00010000
Log. $\gg 2$	00101000
Arith. $\gg 2$	11101000







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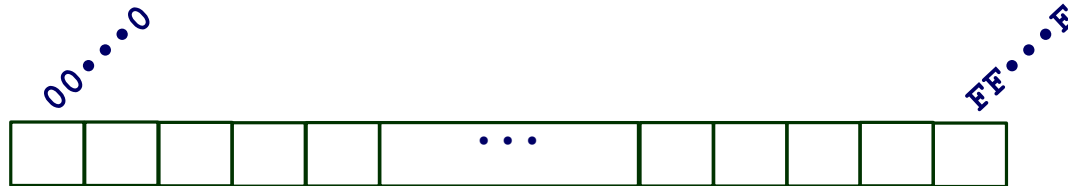
## Modul 2. Bits Representation

### 2.3. Representations in Memory, Pointers, Strings

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# Byte-Oriented Memory Organization



- Programs refer to data by address
  - Conceptually, envision it as a very large array of bytes
    - In reality, it's not, but can think of it that way
  - An address is like an index into that array
    - and, a pointer variable stores an address
- Note: system provides private address spaces to each “process”
  - Think of a process as a program being executed
  - So, a program can clobber its own data, but not that of others



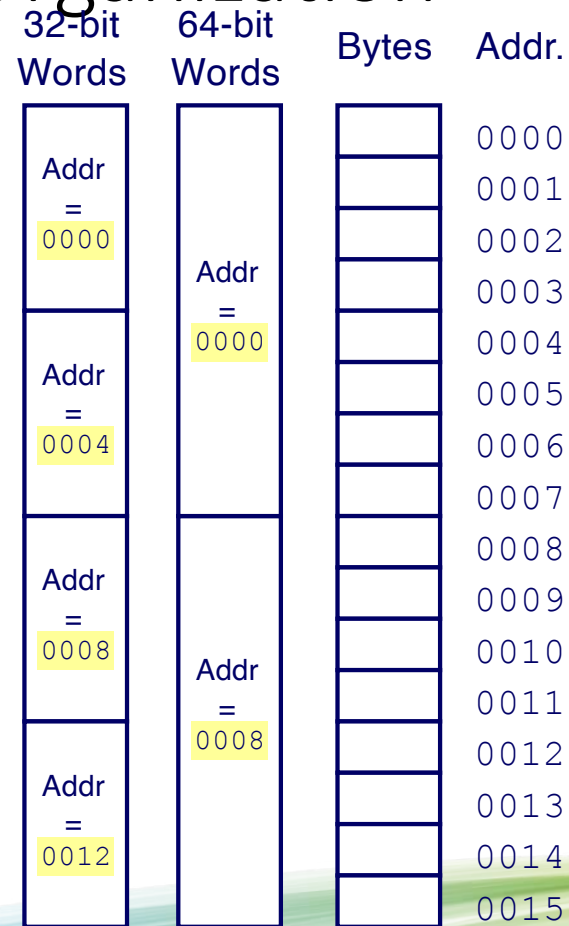
# Machine Words

- Any given computer has a “Word Size”
  - Nominal size of integer-valued data
    - and of addresses
  - Until recently, most machines used 32 bits (4 bytes) as word size
    - Limits addresses to 4GB ( $2^{32}$  bytes)
  - Increasingly, machines have 64-bit word size
    - Potentially, could have 18 EB (exabytes) of addressable memory
    - That's  $18.4 \times 10^{18}$
  - Machines still support multiple data formats
    - Fractions or multiples of word size
    - Always integral number of bytes



# Word-Oriented Memory Organization

- Addresses Specify Byte Locations
  - Address of first byte in word
  - Addresses of successive words differ by 4 (32-bit) or 8 (64-bit)



## Example Data Representations

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# Byte Ordering

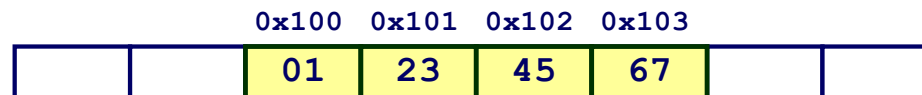
- So, how are the bytes within a multi-byte word ordered in memory?
- Conventions
  - Big Endian: Sun (Oracle SPARC), PPC Mac, *Internet*
    - Least significant byte has highest address
  - Little Endian: *x86*, ARM processors running Android, iOS, and Linux
    - Least significant byte has lowest address



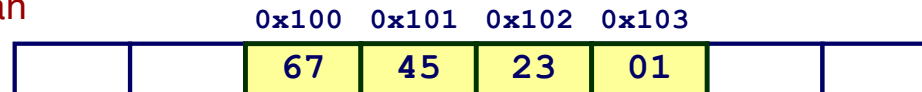
# Byte Ordering Example

- Example
  - Variable x has 4-byte value of 0x01234567
  - Address given by &x is 0x100

Big Endian



Little Endian



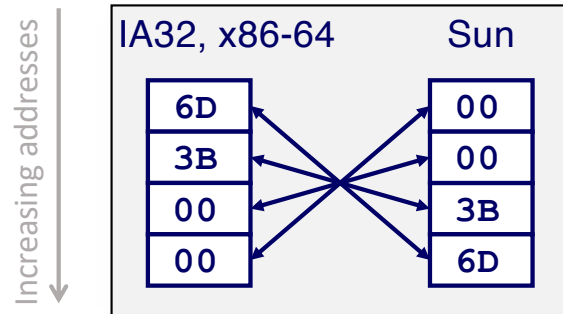
# Representing Integers

Decimal: 15213

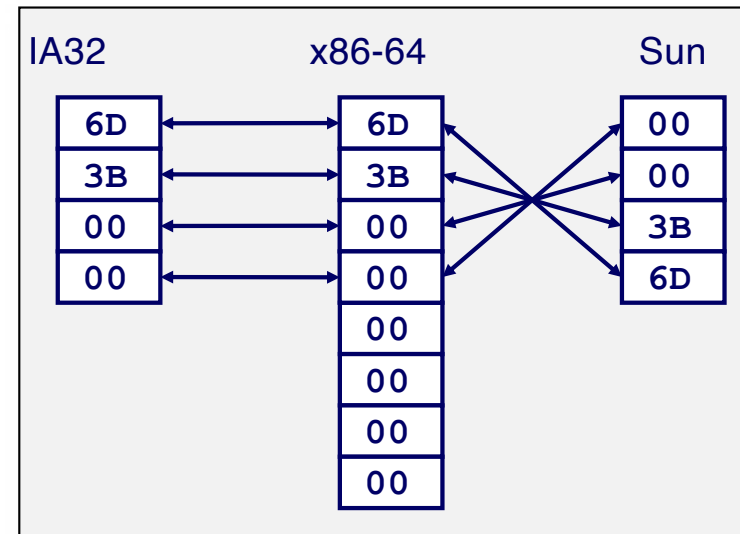
Binary: 0011 1011 0110 1101

Hex: 3 B 6 D

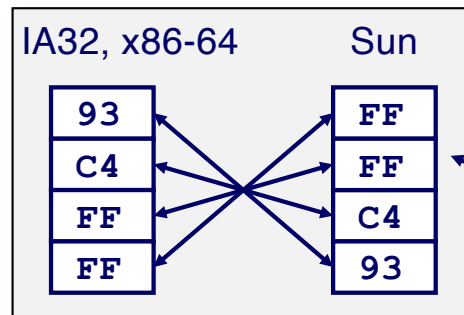
`int A = 15213;`



`long int C = 15213;`



`int B = -15213;`



Two's complement representation





# Examining Data Representations

- Code to Print Byte Representation of Data
  - Casting pointer to unsigned char \* allows treatment as a byte array

```
typedef unsigned char *pointer;  
  
void show_bytes(pointer start, size_t len){  
    size_t i;  
    for (i = 0; i < len; i++)  
        printf("%p\t0x%.2x\n", start+i, start[i]);  
    printf("\n");  
}
```

Printf directives:

%p: Print pointer

%x: Print Hexadecimal



## show\_bytes Execution Example

```
int a = 15213;  
printf("int a = 15213;\n");  
show_bytes((pointer) &a, sizeof(int));
```

Result (Linux x86-64):

```
int a = 15213;  
0x7ffffb7f71dbc 6d  
0x7ffffb7f71dbd 3b  
0x7ffffb7f71dbe 00  
0x7ffffb7f71dbf 00
```



# Representing Pointers

```
int B = -15213;  
int *P = &B;
```

Sun

EF
FF
FB
2C

IA32

AC
28
F5
FF

x86-64

3C
1B
FE
82
FD
7F
00
00

Different compilers & machines assign different locations to objects

Even get different results each time run program



# Representing Strings

```
char S[6] = "18213";
```

- Strings in C
  - Represented by array of characters
  - Each character encoded in ASCII format
    - Standard 7-bit encoding of character set
    - Character "0" has code 0x30
      - Digit  $i$  has code  $0x30+i$
    - *man ascii* for code table
  - String should be null-terminated
    - Final character = 0
- Compatibility
  - Byte ordering not an issue

IA32		Sun
31	↔	31
38	↔	38
32	↔	32
31	↔	31
33	↔	33
00	↔	00



# Reading Byte-Reversed Listings

- Disassembly
  - Text representation of binary machine code
  - Generated by program that reads the machine code
- Example Fragment

Address	Instruction Code	Assembly Rendition
8048365:	5b	pop %ebx
8048366:	81 c3 <u>ab 12 00 00</u>	add \$0x12ab,%ebx
804836c:	83 bb 28 00 00 00 00	cmpl \$0x0,0x28(%ebx)

- Deciphering Numbers

- Value: 0x12ab
- Pad to 32 bits: 0x000012ab
- Split into bytes: 00 00 12 ab
- Reverse: ab 12 00 00



## End of Module #2

1. Representing information as a bits
2. Bit-level Manipulations
3. Representations in memory, pointers, and strings

