

Outline # 3

Algorithm

Algorithm is the step by step procedure to solve a problem

Pseudocode

Pseudocode is the simple English like description of the steps of a procedure and a specification of this procedure using an actual programming language.

Example of Algorithm & Pseudocode

Q: Describe an algorithm for finding maximum value in a finite sequence of integers

Sol:

Algorithm

- (i) Temporarily set the max equal to first integer of sequence
- (ii) Compare the max with the next value of sequence and if the next integer is greater than max value, then set max equal to 2nd integer of sequence.
- (iii) Repeat the (2nd step) if there are more integers in the sequence.
- (iv) Stop when there is no integer left in the sequence. The temp max at this point is the largest integer in the sequence.

Pseudocode

```
( $a_1, a_2, a_3, \dots, a_n$ ) integers sequence  
max =  $a_1$   
for i=2 to n  
    if max <  $a_i$  then max  
        max =  $a_i$ 
```

return max (max is the largest value)

Searching Algorithms

The procedure of locating a specific value from a list of values is called searching algorithm.

(i) Linear Search

Linear Search is also called sequential search.

Linear searching is done by comparing the value, which we want to search, with all values in a sequence or list. The search begins with comparing the value with the first integer. If a_1 (first integer) = x (searched value)

the solution is the location of a_1 , namely 1. If $a_1 \neq x$, then a_2 is compared with x . If $a_2 = x$ then solution is a_2 's location. If $a_2 \neq x$, continue the process by comparing x with each term in list. If the entire list has been searched without locating x , the solution is 0.

~~Algo~~

Procedure : Linear Search)

x from (a₁, a₂, a₃, ..., a_n)

for (i=1, i≤n, i++) → pos(i; i≤n; i++)

{ if ($a_i == x$)

{ loc = i;

break;

else loc = 0;

return loc;

loc = -1

if ($x == a_i$)

{ loc = i;

break;

}

if (loc == -1)

{ S.O.P ("x not found");

else if (loc != -1)

{ S.O.P ("x found at " + index); }

Binary Search

Binary Search algo can be used when the list has values occurring in the order of increasing size (ascending order)

It proceeds by comparing the element to be located in the middle term of list.

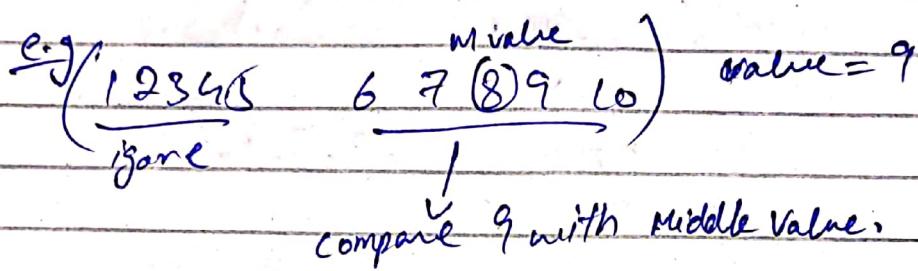
First the list is split into 2 parts.

by comparing the value with the middle value of list.

If value ≠ middle value then we check
if value is greater than middle value or
if value is smaller than middle value.

e.g

if value is greater than middle value, then we ignore the value on left of middle value of list and compare the value with the middle value of remaining list.

e.g. 
compare 9 with middle value.

This process continues until we found the value at the middle value.

Algo

```
int l=left, r=right; size-1  
mid = (l+r)/2;
```

```
while (l <= r) {
```

```
    if (n == arr[mid])
```

```
        S.O.P ("Value found");
```

```
        break;
```

```
    else if (n > arr[mid])
```

```
        l = mid + 1;
```

```
    else if (n < arr[mid])
```

```
        r = mid - 1;
```

```
    if (l > r)
```

```
        S.O.P ("Value not found");
```

Sorting

Sorting is the process of putting the elements into a list in increasing order.

e.g.

unsorted array: 7, 2, 5, 3, 1

sorted array: 1, 2, 3, 5, 7

Bubble Sort

Bubble sort is the simplest sorting algorithm.

It puts a list into increasing order by successively comparing adjacent elements, interchanging them if they are in wrong order.

Algo:

```
for (i=5; i>0; i--)  
{ for (j=0; j<i; i++, j++)  
    if (a[i]<a[j])  
    { temp = a[i];  
     a[i] = a[j];  
     a[j] = temp  
    }  
}
```

e.g. { 8, 2, 3, 4, 5 }
1, 2, 3, 4, 5 ①
2, 1, 3, 4, 5 ②
2, 3, 1, 4, 5 ③
2, 3, 4, 1, 5 ④
2, 3, 4, 5, 1

2, 3, 4, 5, 1 ①
3, 2, 4, 5, 1 ②
3, 4, 2, 5, 1 ③
3, 4, 5, 2, 1

Inserlion Sort

To sort the list with n element, the insertion sort begins with the 2nd element. The insertion sort compares this second element with the first element and inserts it before the first element, if the 2nd element is smaller than first element. And it inserts after the ~~first~~ element if 2nd is greater than 1st element.

Example $\{3, 2, 4, 1, 5\}$

int i, key, j;

for($j=1; j \leq 5; j++$)
 {

 key = arr[j];
 i = j - 1;

 while ($i \geq 0 \& arr[i] > key$)

 arr[i + 1] = arr[i];

 i = i - 1;

 }

 arr[i + 1] = key;

}

Quick Sort

Rules

⇒ First you have to choose a pivot element out of all an array.

pivot element can be any element from array.

e.g. (35) 50 15 25 80 20 90 45

↓
pivot element

⇒ Next you have to chase Two pointers as pointers p & q respectively

(35) 50 15 25 80 20 90 45
↓ ↓
p q

⇒ 'p' element will move to its right side and stops when it will find an element greater than "pivot"

⇒ 'q' element will move to its L.H.S and stops when it will find the lesser element than pivot value.

Why we do it?

just to bring all the greater elements to the right of pivot value and lesser value to its left side

through this we can sort an array through quick sort method.

35 5 4 3 2 1 $\frac{1}{\infty}$ → will stop there
if there is no small element

$\Rightarrow p$ will check \geq
 $\Rightarrow q$ will check \leq

→ This is to stop the P
traversing in case of it
doesn't find greater
element.

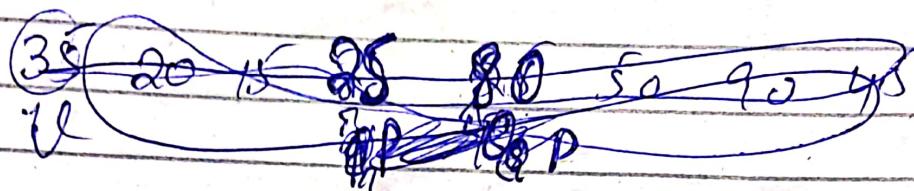
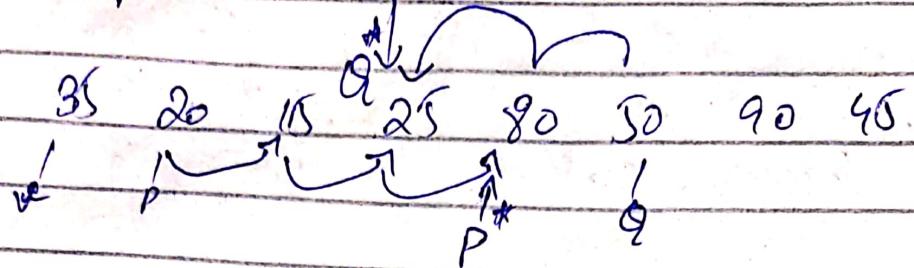
1st pass

(35) 50 15 25 80 20 90 45
↓↑ ↓↑ ↑↑ ↑↑ ↓↑ ↓↑

27 Now p is at 50 and Q is at 20,
 Check that they both crossed each other or
 if not then swap both the value

35 20 15 25 80 50 90 46
✓ ✓ ✓ ✓ ✓ ✓ ✓

Repeat

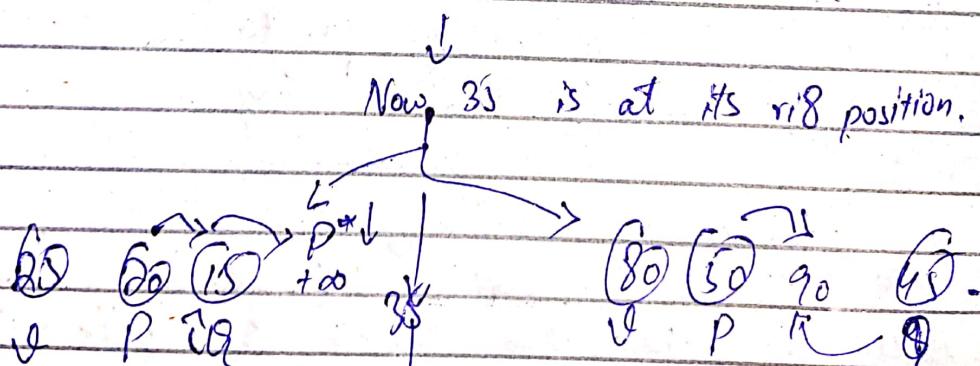


P & Q crossed each other

\Rightarrow Now check whether P & Q have crossed each other or not, if crossed then replace pivot value with Q.

\Rightarrow This condition also holds when both P and Q will be at same position

25 20 15 35 80 50 90 45



Q & P crossed

replace pivot with Q

15 20 25

35

80 50 45 90
V Q P

Now both P & Q crossed

so replace V. & Q

10 20 25 35 45 50 80 90



Integers:

Integers are the set of numbers from 0 to ∞ , or $-\infty$ to 0.

Division

When one integer is divided by a second non zero integer, the quotient may or may not be an integer.

For Example:

$12/3=4$ is an integer

$11/4=2.75$ is not an integer

Division Algorithm

Let a be an integer and d a positive integer. Then there are unique integers q and r such that

$$a = dq + r$$

$\therefore a$ = dividend

d = divisor

q = quotient

r = remainder

$$\Rightarrow \text{'q' is 'a' div 'd'} \Rightarrow \text{'r' is 'a' mod 'd'}$$

Example

Q: What are the quotient & remainder when 101 is divided by 11?

$$q = \text{a div } d$$

$$= 101 / 11$$

$$q = 9$$

$$r = \text{a mod } d$$

$$= 101 \% 11$$

$$= 2$$

Matrices:

$$A = \begin{bmatrix} 2 & 3 & 4 \\ 1 & 5 & 0 \end{bmatrix}$$

Order of matrix

$$= m \times n$$

$$A = 2 \times 3$$

\Rightarrow Scalar matrices = which have same number of rows & columns in both

\Rightarrow Rectangular matrices = which have not same number of rows & columns.

$$A + B = \begin{bmatrix} 2 & 3 \\ 4 & 5 \end{bmatrix} + \begin{bmatrix} 5 & 6 \\ 7 & 8 \end{bmatrix}$$

$$= \begin{bmatrix} 7 & 9 \\ 11 & 13 \end{bmatrix}$$

$$A \cdot B = \begin{bmatrix} 1 & 2 & 3 \\ 2 & 1 & 4 \end{bmatrix} \begin{bmatrix} 1 & 2 \\ 4 & 3 \\ 2 & 0 \end{bmatrix}$$

$$= \begin{bmatrix} 1+8+6 & 2+6+0 \\ 2+4+8 & 4+3+0 \end{bmatrix}$$

$$= \begin{bmatrix} 15 & 8 \\ 14 & 7 \end{bmatrix}$$

\Rightarrow Transpose of a matrix

$$A = \begin{bmatrix} 1 & 2 & 3 \\ a & b & c \end{bmatrix}$$

$$A^T = \begin{bmatrix} 1 & a \\ 2 & b \\ 3 & c \end{bmatrix}$$

: Convert rows into columns
and column into rows
It is transpose.

\Rightarrow Zero-one matrix

$$A = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 1 \end{bmatrix}, B = \begin{bmatrix} 1 & 0 \\ 0 & 0 \\ 0 & 1 \end{bmatrix}$$

\Rightarrow Join of zero-one matrices

$$A = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 1 \end{bmatrix}, \quad B = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 0 \end{bmatrix}$$

$A \vee B$

↑(OR)

~~replacement~~

$$\begin{bmatrix} 1 \vee 1 & 0 \vee 1 & 0 \vee 1 \\ 0 \vee 0 & 1 \vee 1 & 1 \vee 0 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

\Rightarrow Meet of zero-one matrices

$$A = \begin{bmatrix} 1 & 0 & 1 \\ 0 & 0 & 1 \end{bmatrix}, \quad B = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 1 & 1 \end{bmatrix}$$

$$A \wedge B = \begin{bmatrix} 1 \wedge 1 & 0 \wedge 1 & 1 \wedge 0 \\ 0 \wedge 1 & 0 \wedge 1 & 1 \wedge 1 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

⇒ Boolean product zero-one matrices:

$$A = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}, \quad B = \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix}$$

$$A \odot B = \begin{bmatrix} (0 \wedge 1) \vee (1 \wedge 1) & (0 \wedge 1) \vee (1 \wedge 1) \\ (1 \wedge 0) \vee (0 \wedge 1) & (1 \wedge 0) \vee (0 \wedge 1) \end{bmatrix}$$

$$= \begin{bmatrix} (0 \vee 1) & (0 \vee 1) \\ (1 \vee 0) & (1 \vee 0) \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix}$$

⇒ Boolean power of a square zero-one matrix

Find B^2 $\begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix}$

$$B \odot B = \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix}$$

$$= \begin{bmatrix} (1 \wedge 1) \vee (1 \wedge 1) & (1 \wedge 1) \vee (1 \wedge 1) \\ (1 \wedge 1) \vee (1 \wedge 1) & (1 \wedge 1) \vee (1 \wedge 1) \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix}$$