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Formal methods. Using TLC Model Checker

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Introduction to TLC

TLC handles specifications in the standard form

$$Init \wedge \Box[Next]_{vars} \wedge Temporal$$

where

Init is the initial predicate

Next is the next-state action

vars is the tuple of all model variables

Temporal is a temporal formula that specifies a *liveness* condition.



Introduction to TLC

TLC works by generating behaviors that satisfy the specification.

To do this, should be formed a **model** of the specification.

To define a model, first we must assign values to the specification's constant parameters. E.g. for AsyncInterface

CONSTANT Data = {d1, d2, d3}



Modes of TLC

There are two ways to use TLC. The default method is model checking, in which it tries to find all reachable states (that is, all states that can occur in behaviors satisfying the formula).

You can also run TLC in simulation mode, in which it randomly generates behaviors, without trying to check all reachable states.

Second way is useful then we have huge amount of states.

TLC input

The input to TLC consists of a TLA module and configuration. The configuration tells TLC the names of the specification and of the properties to be checked.

For example for HourClock we need specify behavior as:

SPECIFICATION **HC**

OR Initial and next state predicates:

HCini and Hcnext

SPECIFICATION **HC** statement tells TLC that HC is the specification that it should check.

Properties to be checked

For HourClock.tla TLC check the **invariant** property **HCini**
Here invariant **HCini** specifies a state predicate.

In other words, TLC checks that formula **HCini** is an invariant of the specification **HC**, or, that the specification implies $\Box \text{HCini}$.

THEOREM $\text{HC} \Rightarrow \Box \text{HCini}$

$\text{TypeInvariance} \equiv \Box \text{HCini}$

This temporal formula asserting that HCini is always true

THEOREM $\text{HC} \Rightarrow \text{TypeInvariance}$

For **LiveHourClock** we will introduce liveness properties

PROPERTIES AlwaysTick AllTimes TypeInvariance

THEOREM $\text{LSpec} \Rightarrow \text{AlwaysTick} \wedge \text{AllTimes} \wedge \text{TypeInvariance}$

Types of errors to be checked

To find errors in a specification is to verify that it satisfies properties that it should.

You can also run TLC without having it check any property, in which case it will just look for two kinds of errors:

- “Silliness” errors. A silly expression is one like “ $3 + <<1$ ”, “ $2>$ ”, whose meaning is not determined by the semantics of TLA
- Deadlock. it is expressed by the invariance property $\Box(\text{ENABLED } \textit{Next})$



TLC Values


TLC can compute only a restricted class of values, called TLC values. Those values are built from the following four types of primitive values:

Booleans	Values true and false.
Integers	Values like 123.
Strings	Values like "abc".
Model Values	Values introduced in the CONSTANT statement, e.g. {d1, d2, d3}

How TLC Evaluates Expressions

TLC evaluates expressions in a straightforward way, generally evaluating subexpressions “from left to right”. In particular:

- It evaluates $p \wedge q$ by first evaluating p and, if it equals TRUE, then evaluating q .
- It evaluates $p \vee q$ by first evaluating p and, if it equals FALSE, then evaluating q . It evaluates $p \Rightarrow q$ as $\neg p \vee q$.
- It evaluates IF p THEN e_1 ELSE e_2 by first evaluating p , then evaluating either e_1 or e_2 .



Thank you for your attention!
Please ask questions