



Lecture No. 7

Data structures

Linked List

Definition - List

- A list is a collection of items that has a particular order
 - It can have an arbitrary length
 - Objects / elements can be inserted or removed at arbitrary locations in the list
 - A list can be traversed in order one item at a time

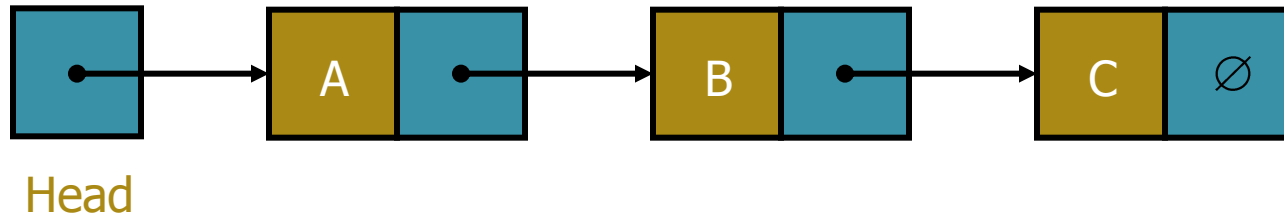
List Overview

- Linked lists
 - Abstract data type (ADT)
- Basic operations of linked lists
 - Insert, find, delete, print, etc.
- Variations of linked lists
 - Singly linked lists
 - Circular linked lists
 - Doubly linked lists
 - Circular doubly linked list

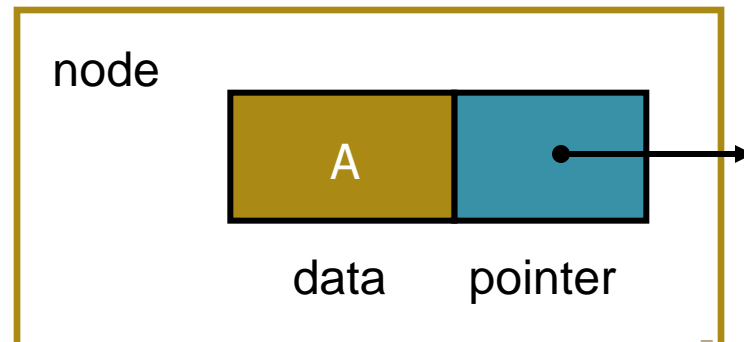
Linked List Terminologies

- **Traversal of List**
 - Means to visit every element or node in the list beginning from first to last.
- **Predecessor and Successor**
 - In the list of elements, for any location n , $(n-1)$ is predecessor and $(n+1)$ is successor.
 - In other words, for any location n in the list, the left element is predecessor and the right element is successor.
 - Also, the first element does not have predecessor and the last element does not have successor.

Linked Lists



- A *linked list* is a series of connected *nodes*
- Each node contains at least
 - A piece of data (any type)
 - Pointer to the next node in the list
- *Head*: pointer to the first node
- The last node points to `NULL`



Lists – Another perspective

□ A list is a **linear** collection of **varying** length of **homogeneous** components.

□ **Homogeneous**: All components are of the same type.

□ **Linear**: Components are ordered in a line (hence called Linear linked lists).



Arrays are lists..

Arrays Vs Lists

- Arrays are lists that have a fixed size in memory.
- The programmer must keep track of the length of the array
- No matter how many elements of the array are used in a program, the array has the same amount of allocated space.
- Array elements are stored in successive memory locations. Also, order of elements stored in array is same logically and physically.

Arrays Vs Lists

- A linked list takes up only as much space in memory as is needed for the length of the list.
- The list expands or contracts as you add or delete elements.
- In linked list the elements are not stored in successive memory location
- Elements can be added to (or deleted from) either end, or added to (or deleted from) the middle of the list.

Array versus Linked Lists

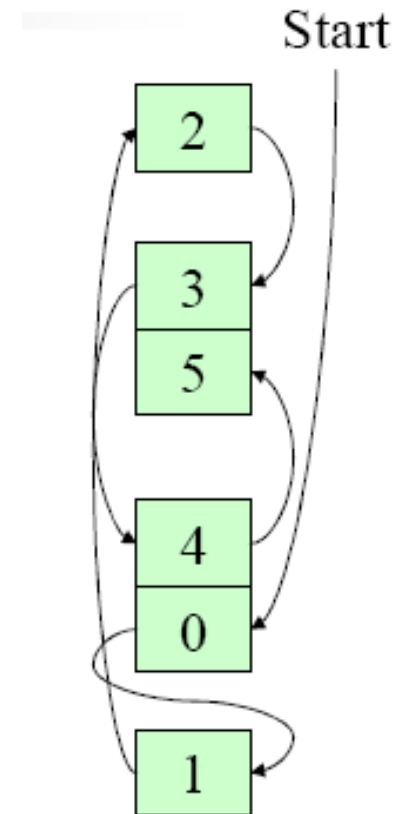
- Linked lists are more complex to code and manage than arrays, but they have some distinct advantages.
 - **Dynamic**: a linked list can easily grow and shrink in size.
 - We don't need to know how many nodes will be in the list. They are created in memory as needed.
 - In contrast, the size of a C++ array is fixed at compilation time.
 - **Easy and fast insertions and deletions**
 - To insert or delete an element in an array, we need to copy to temporary variables to make room for new elements or close the gap caused by deleted elements.
 - With a linked list, no need to move other nodes. Only need to reset some pointers.

Array versus Linked Lists

An Array



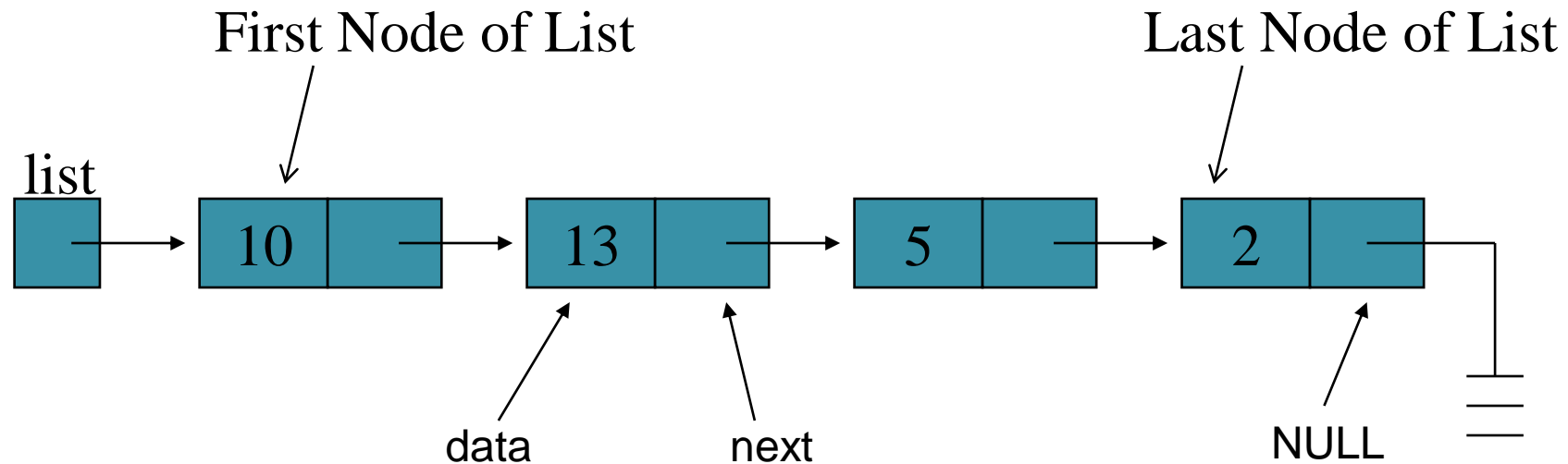
A Linked List



Basic Operations of Linked List

- Operations of Linked List
 - IsEmpty: determine whether or not the list is empty
 - InsertNode: insert a new node at a particular position
 - FindNode: find a node with a given value
 - DeleteNode: delete a node with a given value
 - DisplayList: print all the nodes in the list

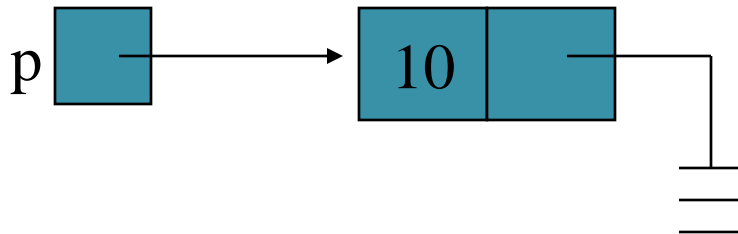
An integer linked list



Creating a List node

```
struct Node {  
    int data;           // data in node  
    Node *next;        // Pointer to next node  
};
```

```
Node *p;  
p = new Node;  
p -> data = 10;  
p -> next = NULL;
```



The NULL pointer

NULL is a special pointer value that does not reference any memory cell.

If a pointer is not currently in use, it should be set to NULL so that one can determine that it is not pointing to a valid address:

```
int *p;  
p = NULL;
```

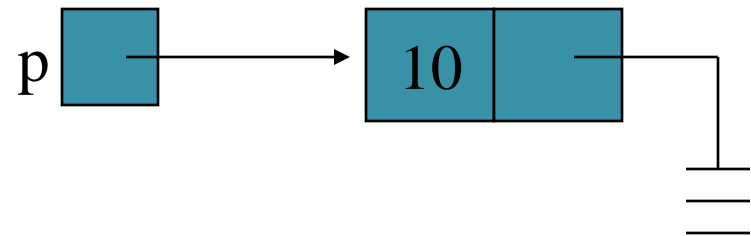
Adding a node to a list

Node *p, *q;

p = new Node;

p -> data = 10;

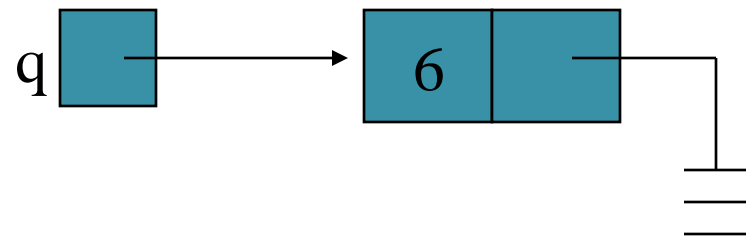
p -> next = NULL;



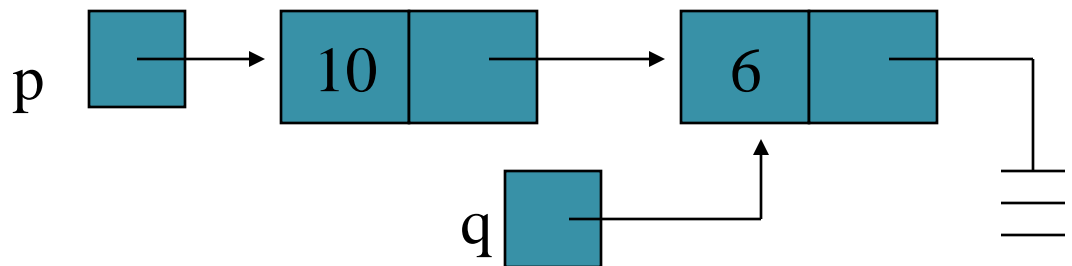
q = new Node;

q -> data = 6;

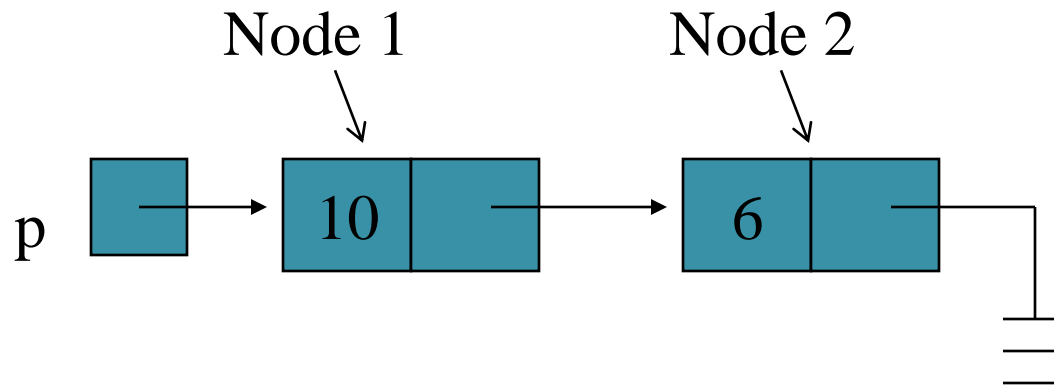
q -> next = NULL;



p -> next = q;



Accessing List Data



Expression

p

p - > data

p - > next

p - > next - > data

p - > next - > next

Value

Pointer to first node (head)

10

Pointer to next node

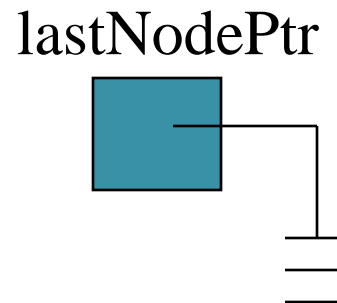
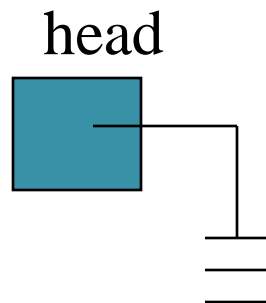
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NULL pointer

Building a list from 1 to n

```
struct Node {  
    int data;  
    Node *next;  
};
```

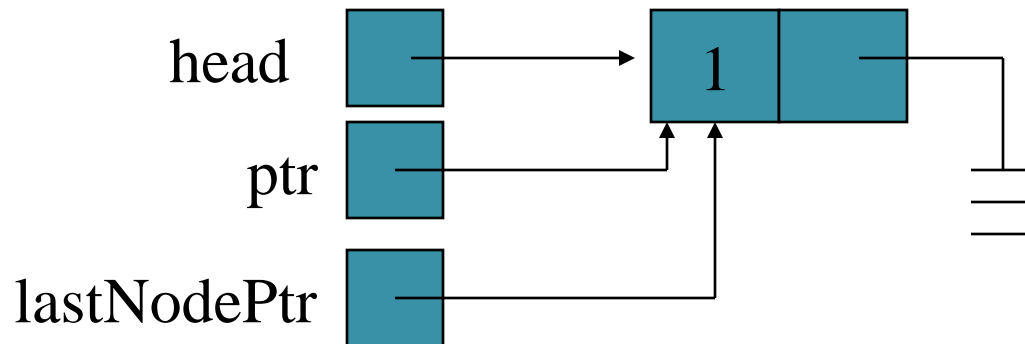
```
Node *head = NULL;           // pointer to the list  
head  
Node *lastNodePtr = NULL;    // pointer to last node  
in list
```



Creating the first node

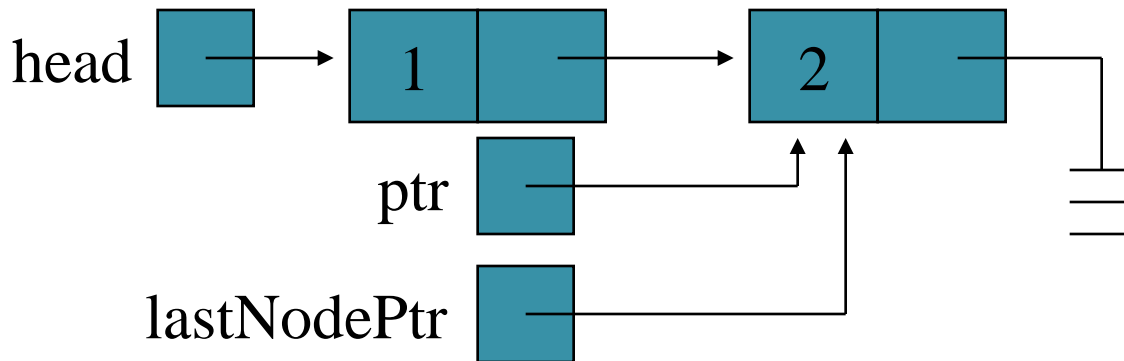
```
Node *ptr;           // declare a pointer to Node
ptr = new Node;       // create a new Node
ptr -> data = 1;
ptr -> next = NULL;
```

```
head = ptr;          // new node is first
lastNodePtr = ptr;   // and last node in list
```

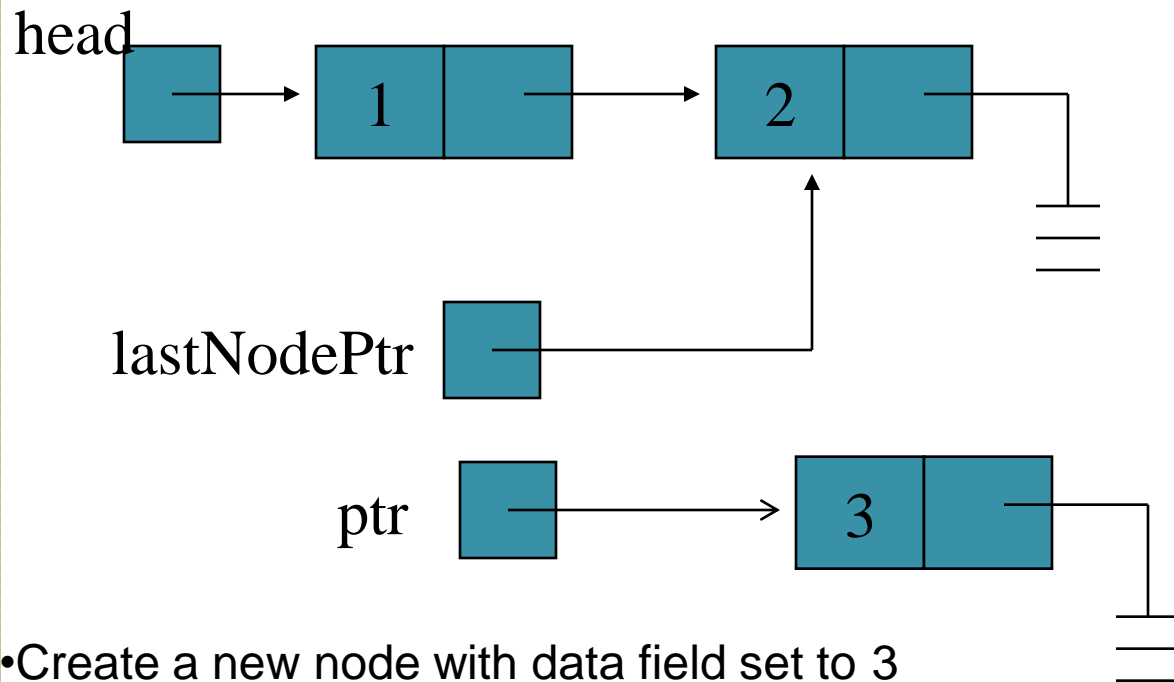
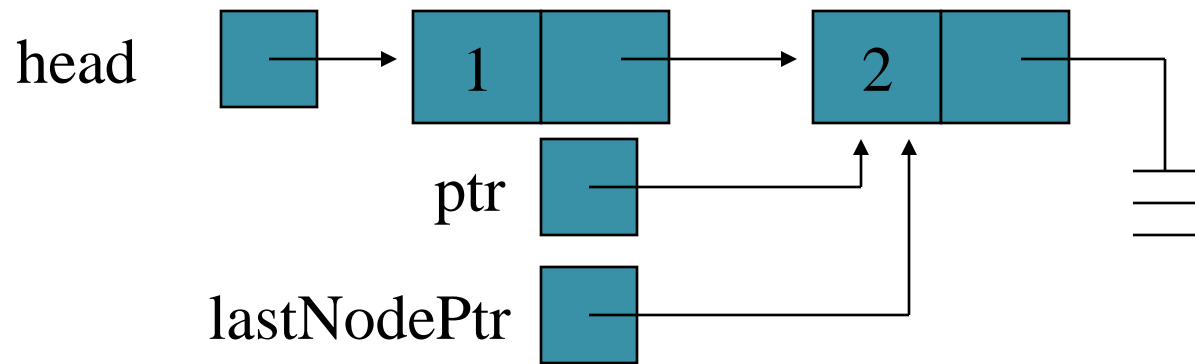


Adding more nodes

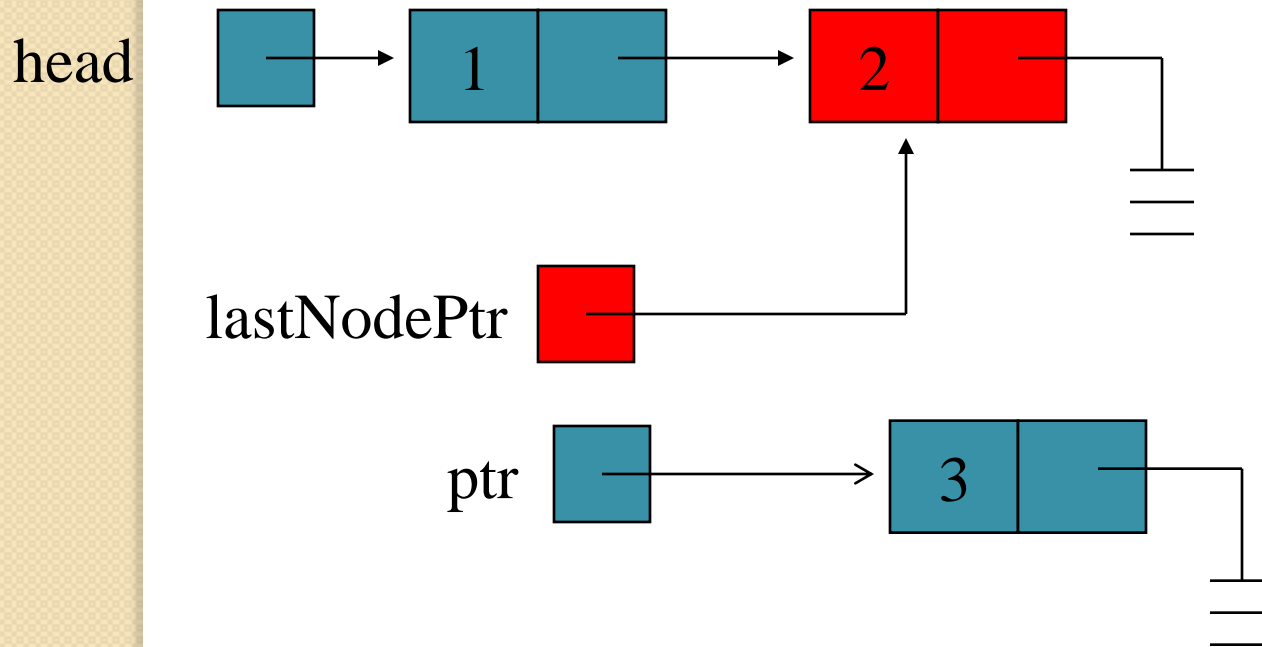
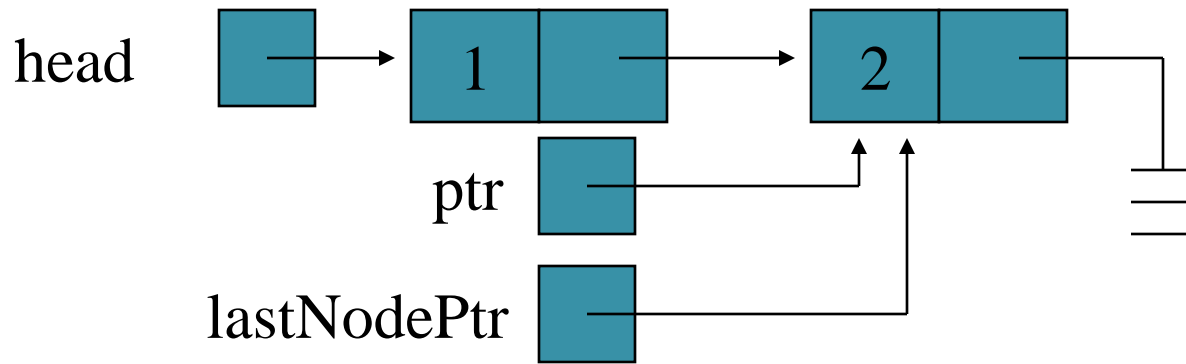
```
for (int i = 2; i <= n; i ++ ) {  
    ptr = new Node;                //create new node  
    ptr -> data = i;  
    ptr -> next = NULL;  
    lastNodePtr -> next = ptr;    // order is  
    lastNodePtr = ptr;            // important  
}
```



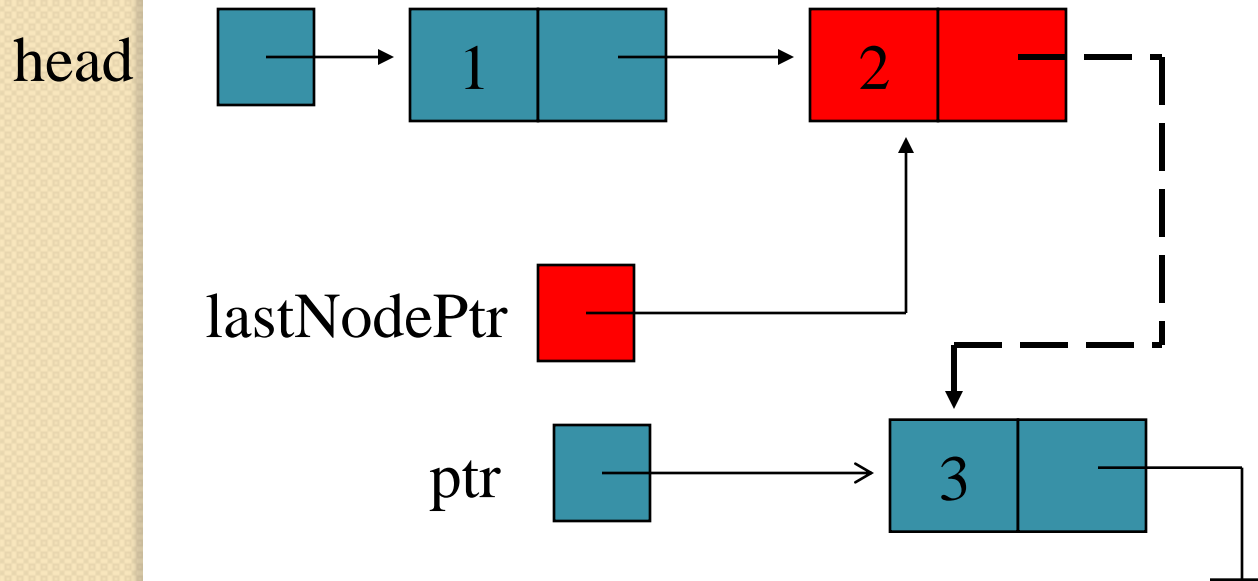
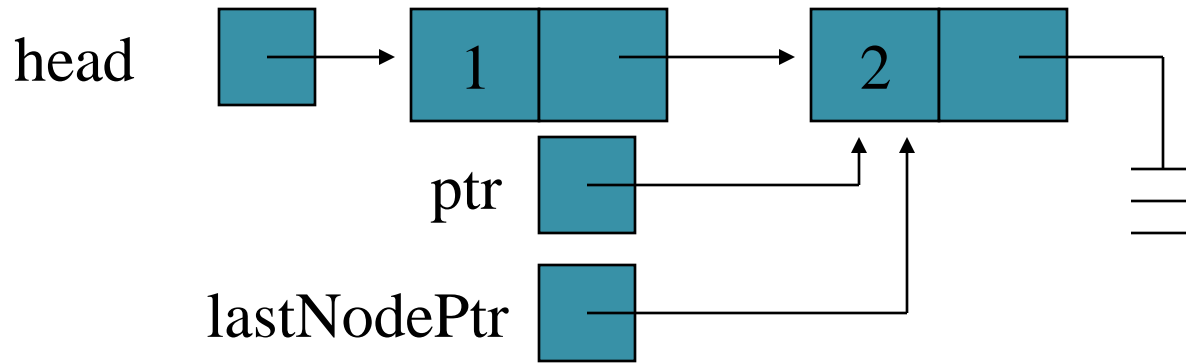
Initially



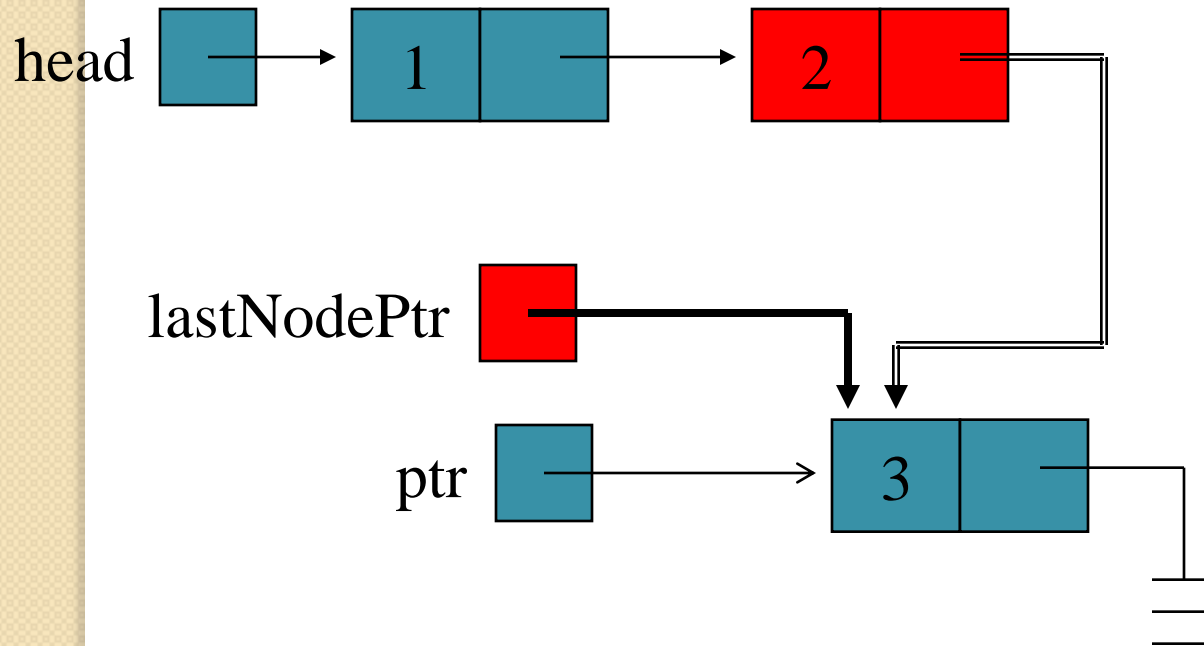
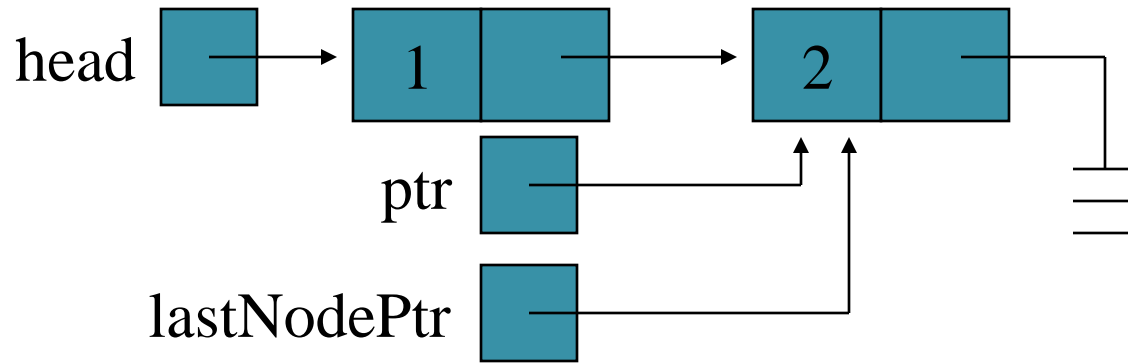
- Create a new node with data field set to 3
- Its next pointer should point to NULL



- The next pointer of the node which was previously last should now point to newly created node
"lastNodePtr->next=ptr"

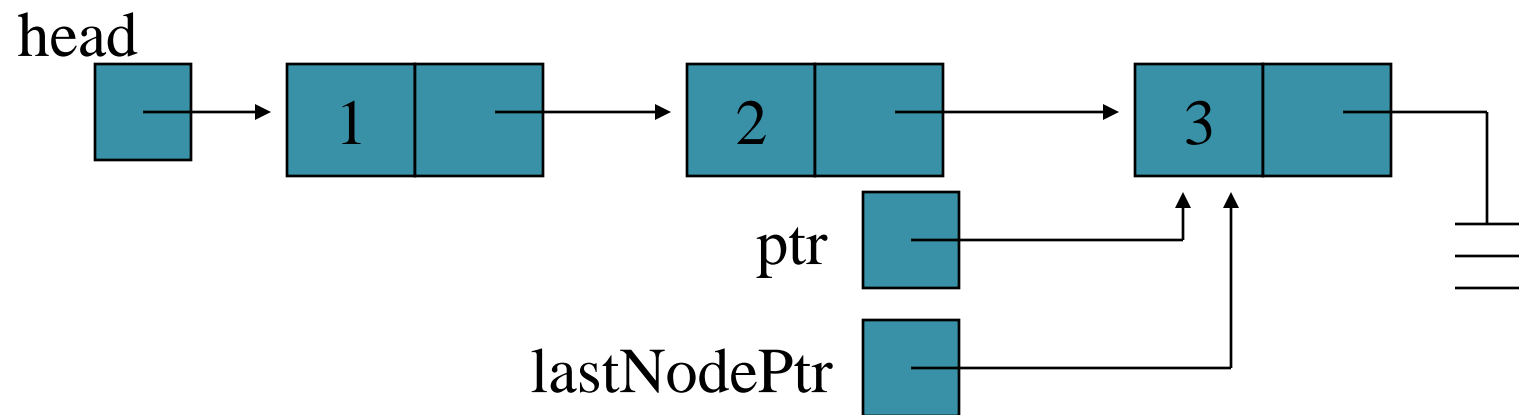


- The next pointer of the node which was previously last should now point to newly created node
"lastNodePtr->next=ptr"
- LastNodePtr should now point to the newly created Node "lastNodePtr = ptr;"

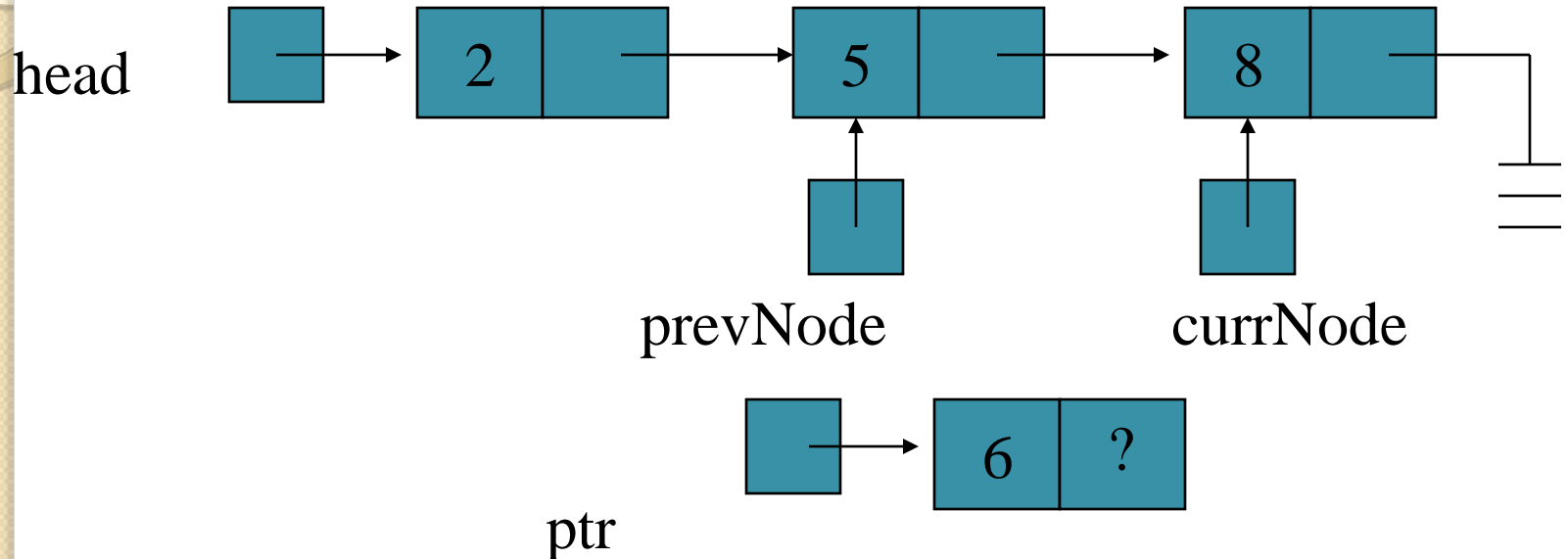


- LastNodePtr should now point to the newly created Node “lastNodePtr = ptr;”

Re-arranging the view



Inserting a node in a list



Step 1: Determine where you want to insert a node.

Step 2: Create a new node:

```
Node *ptr;
```

```
ptr = new Node;
```

```
ptr -> data = 6;
```

```
Node *ptr, *currNode, *prevNode ;
prevNode = head;
ptr = new Node;
ptr->data = 6;
ptr->next = NULL;
currNode = head->next;
While (currNode->data < ptr->data)
{
    prevNode = currNode;
    currNode = currNode->next;
}
```

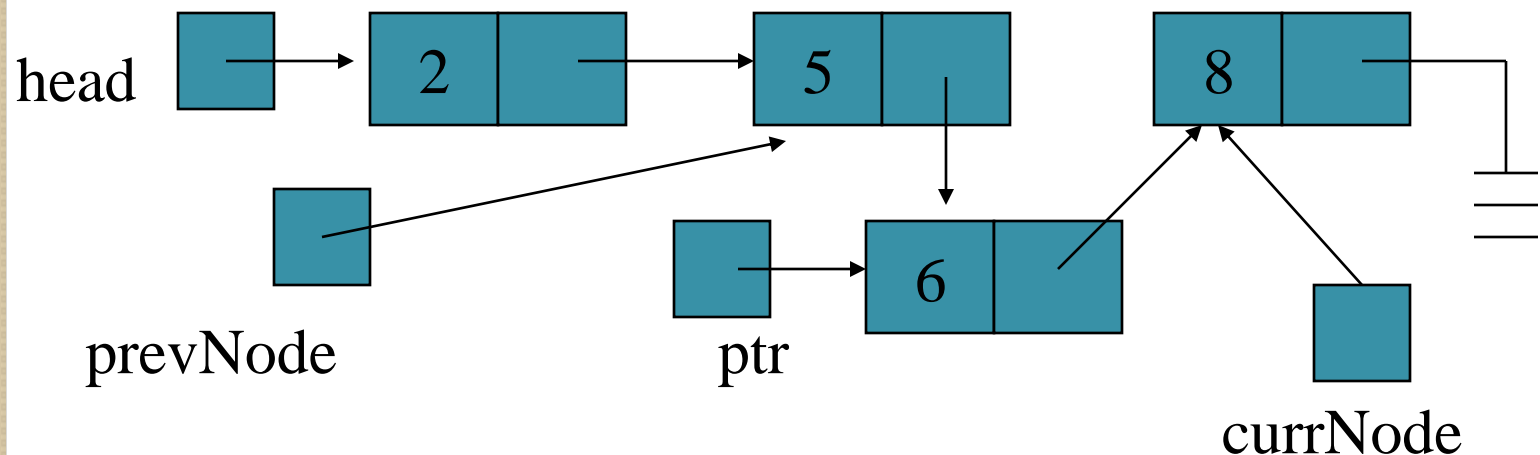
Note:

when this loop terminates **prevNode** and **currNode** are at a place where insertion will take place. Only the “LINKS” or pointers of the list remain to be adjusted

Continuing the insert

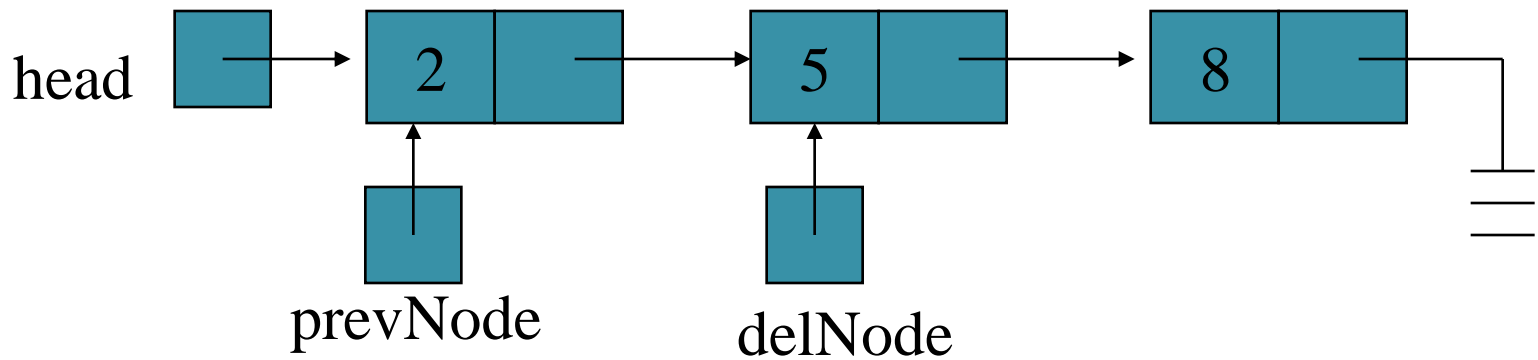
Step 3: Make the new node point to the current Node pointer.
`ptr -> next = currNode;`

Step 4: Make previous node point to the new node:
`prevNode -> next = ptr;`



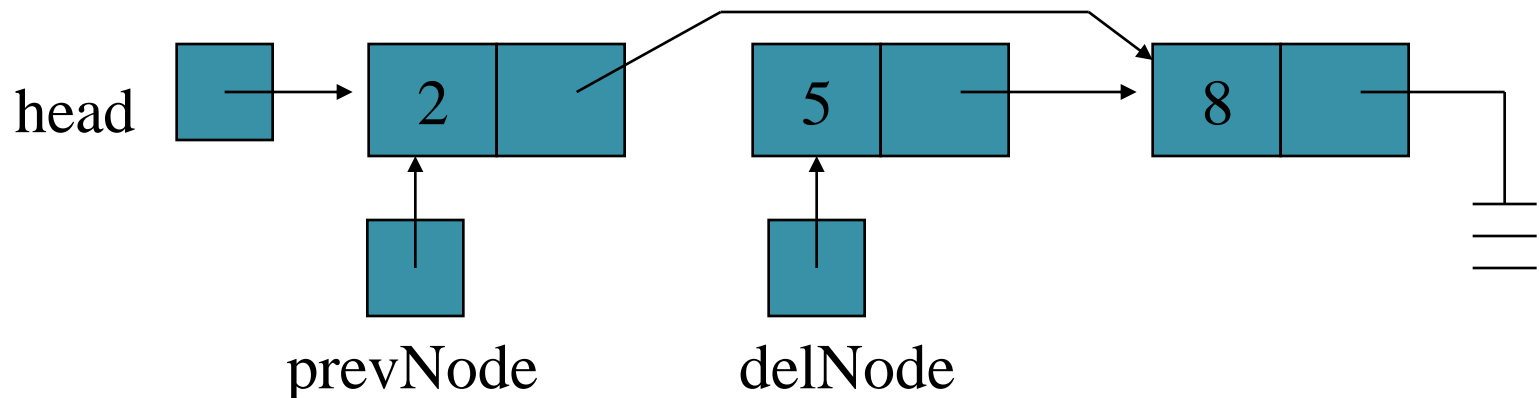
Now The new link has been added in the linked list

Deleting a node from a list



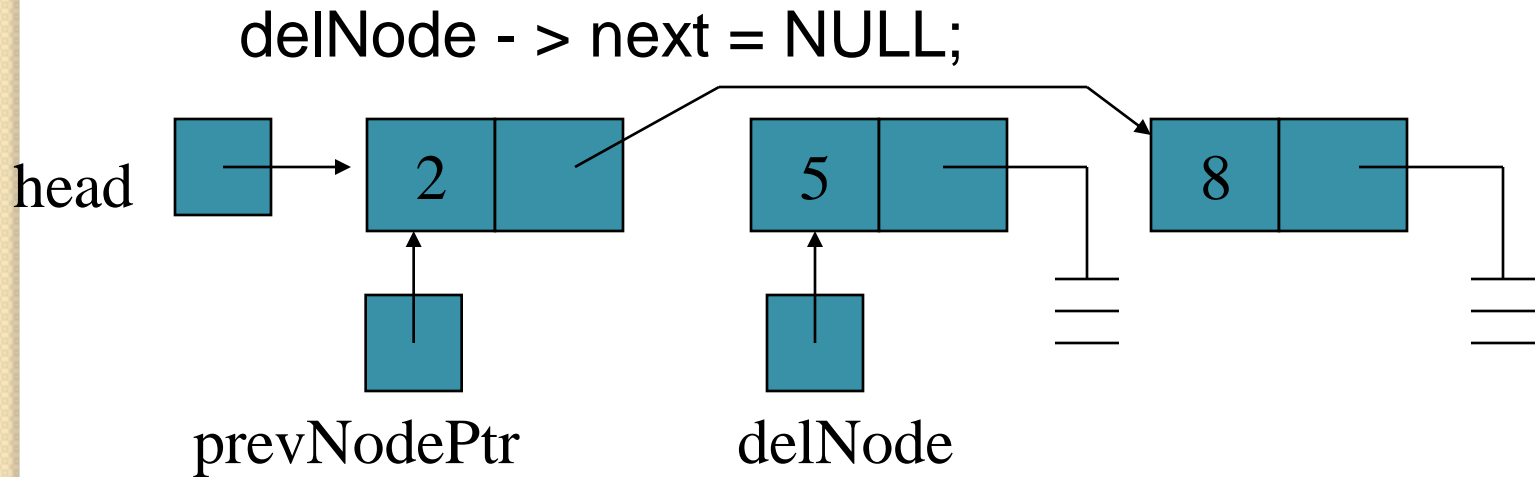
Step 1: Redirect pointer from the Node before the one to be deleted to point to the Node after the one to be deleted.

`prevNode -> next = delNode -> next;`



Finishing the deletion

Step 2: Remove the pointer from the deleted link.



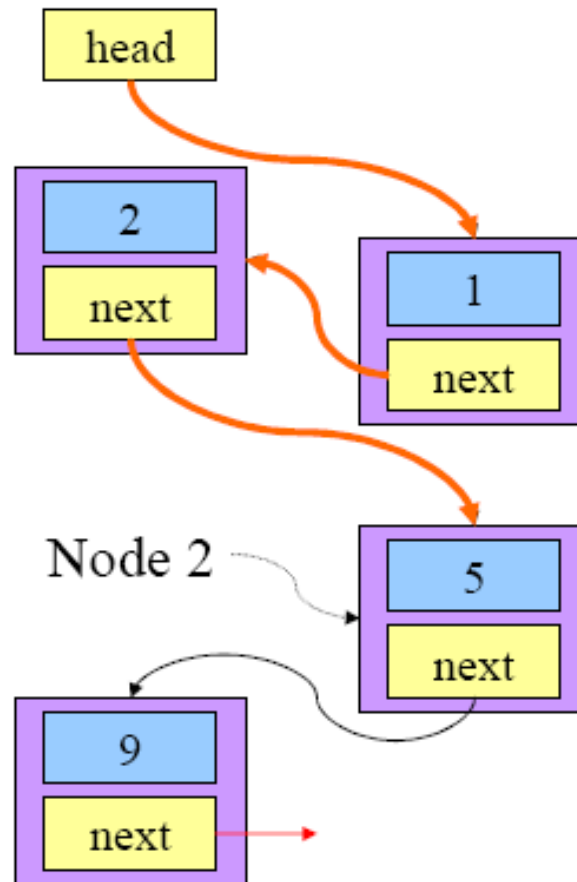
Step 3: Free up the memory used for the deleted node:

`delete delNode;`

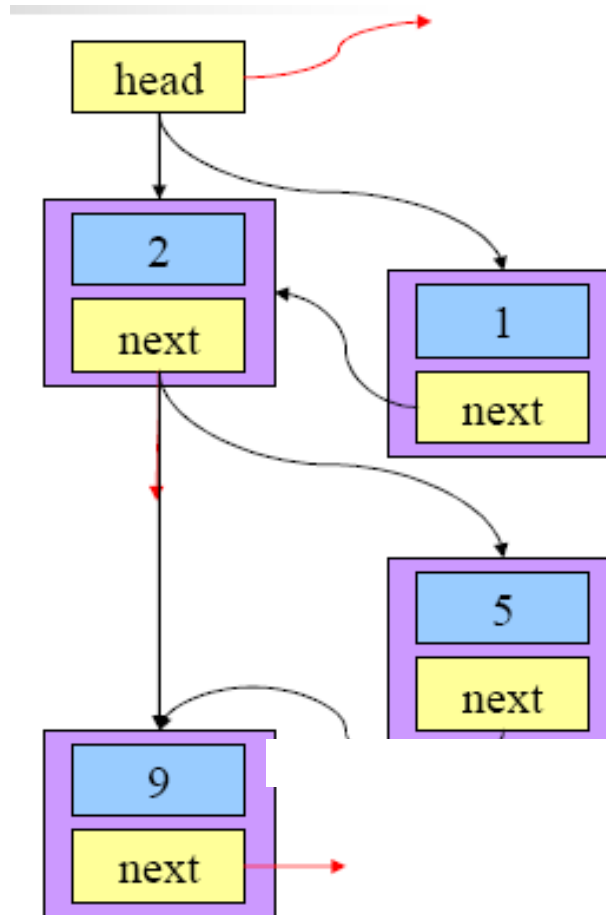


LIST OPERATIONS - SUMMARIZED

Traversing a Linked List



Insertion in a Linked List



Deletion from a Linked List

