# Review Sheet EE 226A

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# 1 Tools and Tricks

need to add bayes theorem graph associated with a markov chain is a tree, then the markov chain is reversible add tools for min max of exp rv

# 2 Elements of Probability Theory

## 2.1 Probability Spaces and Events

## Definition 2.1: Kolmogorov's axioms

For any **probability space**  $(\Omega, \mathcal{F}, P)$ , the function P is called a **probability measure**. It is assumed to satisfy Kolmogorov's axioms:

- *i.*)  $P(A) \ge 0$  for all  $A \in \mathcal{F}$ ;
- ii.)  $P(\Omega) = 1$ ;
- *iii.*) if  $A_1, A_2, \dots \in \mathcal{F}$  are disjoint events, then  $P(\bigcup_{i \geqslant 1} A_i) = \sum_{i \geqslant 1} P(A_i)$ .

The probability space we are working in encodes the model of our experiment, with the **measurable space**  $(\Omega, \mathcal{F})$  being the most fine-grained representation of outcomes we can hope to observe.

#### Theorem 2.2

For a probability space  $(\Omega, \mathcal{F}, P)$ , the probability measure P enjoys the following properties:

- *i.*) Monotonicity: If  $A \subset B$  are events, then  $P(A) \leq P(B)$ .
- *ii.*) Subadditivity (Union bound): If  $(A_i)_{i\geqslant 1}$  is a sequence of events in  $\mathcal F$  and  $A=\bigcup_{i\geqslant 1}A_i$ , then  $P(A)\leqslant \sum_{i\geqslant 1}P(A_i)$ .
- *iii.*) Continuity from below: If  $A_1 \subset A_2 \subset ...$  are events in  $\mathcal{F}$  and  $A = \bigcup_{i \geqslant 1} A_i$ , then  $P(A_i) \to P(A)$ .
- iv.) Continuity from above: If  $A_1 \supset A_2 \supset \dots$  are events in  $\mathfrak F$  and  $A = \bigcap_{i\geqslant 1} A_i$ , then  $P(A_i) \to P(A)$ .

#### Theorem 2.3: Law of total probability

If events  $A_1, A_2, \ldots$  partition  $\Omega$ , then

$$P(B) = \sum_{i\geqslant 1} P(A_i\cap B), \quad B\in \mathfrak{F}.$$

## **Definition 2.4: Infinitely often**

$$\{A_n \text{ infinitely often}\} = \bigcap_{n\geqslant 1} \bigcup_{i\geqslant n} A_i.$$

We should understand  $\{A_n \text{ i.o.}\}$  to be the set of samples  $\omega \in \Omega$  such that  $\omega \in A_i$  for infinitely many  $i \geqslant 1$ .

#### Lemma 2.5: Borel-Cantelli

Let  $A_1, A_2, ...$  be a sequence of events. If

$$\sum_{i \ge 1} P(A_i) < \infty$$

then  $P({A_i infinitely often}) = 0$ .

#### Definition 2.6: Independent events

A collection of events  $A_1, A_2, ...$  are **independent** if

$$P\left(\bigcap_{i\in S}A_i\right)=\prod_{i\in S}P(A_i)$$

for every finite subset  $S \subset \{1, 2, 3, ...\}$ . If  $A_1, A_2, ...$  are independent, then  $A_1^C, A_2, ...$  are independent. By induction, the complements  $A_1^C, A_2^C, ...$  are also independent.

#### Lemma 2.7: Converse to Borel-Cantelli

Let  $A_1, A_2, ...$  be independent events. If

$$\sum_{i\geqslant 1}P(A_i)=\infty,$$

then  $P({A_i infinitely often}) = 1$ .

## Theorem 2.8: Carathéodory's extension theorem

Suppose  $\mathcal{G}$  is a family of subsets of  $\Omega$  that satisfies the following (relatively modest) properties:

- *i.*)  $\emptyset$ ,  $\Omega \in \mathcal{G}$ ;
- *ii.*) if  $A, B \in \mathcal{G}$ , then  $A \cap B \in \mathcal{G}$ ;
- *iii.*) if  $A, B \in \mathcal{G}$ , then there is a *finite* number of *disjoint* sets  $C_1, \ldots, C_n \in \mathcal{G}$  such that  $A \setminus B = \bigcup_{i=1}^n C_i$ . (Note: (*iii*) is weaker than imposing the assumption  $A \in \mathcal{G} \implies A^C \in \mathcal{G}$ .)

The extension theorem says that if we assign numbers (i.e., probabilities) p(A) to the sets  $A \in \mathcal{G}$  so that

- A.  $p(A) \ge 0$  for  $A \in \mathcal{G}$ ;
- B. p(Ω) = 1;
- C. if  $B \in \mathcal{G}$  and  $A_1, A_2, \dots \in \mathcal{G}$  are disjoint with  $B = \bigcup_{i \geqslant 1} A_i$ , then  $p(B) = \sum_{i \geqslant 1} p(A_i)$ ,

then there exists a unique probability measure P on  $\sigma(\mathfrak{G})$  that satisfies A-C and has the property that  $P(A) = \mathfrak{p}(A)$  for all  $A \in \mathfrak{G}$ .

## 2.2 Random Variables and Expectation

## 2.2.1 Random variables and algebraic properties

#### Definition 2.9: Random Variable

We define a random variable to be a function  $X : \Omega \to \overline{\mathbb{R}}$  that satisfies

$$\{\omega \in \Omega : X(\omega) \leqslant \alpha\} \in \mathcal{F} \text{ for each } \alpha \in \overline{\mathbb{R}}.$$

Note that  $\overline{\mathbb{R}} = \mathbb{R} \cup \{-\infty, \infty\}$ .

A function  $X : \Omega \to \overline{\mathbb{R}}$  satisfying the definition above is said to be  $\mathcal{F}$ —**measurable**. If X does not take values  $\pm \infty$  (say, with probability one), then we say it is a **real-valued random variable**.

#### **Proposition 2.10**

If X is a random variable, then pX and  $|X|^p$  are random variables for  $p \in \mathbb{R}$ . Moreover, if X,Y are real-valued random variables, then X + Y, and XY are also random variables.

## **Proposition 2.11**

If  $(X_n)_{n\geqslant 1}$  is a sequence of random variables defined on a common probability space  $(\Omega, \mathcal{F}, P)$ , then

- $\sup_{n\geq 1} X_n$  and  $\inf_{n\geq 1} X_n$  are random variables; and
- $\limsup_{n\to\infty} X_n$  and  $\liminf_{n\to\infty} X_n$  are random variables; and
- if  $\lim_{n\to\infty} X_n$  exists point wise, it is also a random variable.

#### Definition 2.12: Almost sure equivalence of random variables

If X, Y are random variables and  $P(\{\omega : X(\omega) \neq Y(\omega)\}) = 0$ , then we say X = Y almost surely (abbreviated a.s.), or X = Y with probability one.

#### 2.2.2 Distribution functions and distributions

#### **Definition 2.13: Distribution function**

A random variable X on a probability space  $(\Omega, \mathcal{F}, P)$  is described in part by its **distribution function**  $F_X : \mathbb{R} \to [0, 1]$ , defined as

$$F_X(x) := P\{X \leq x\}, x \in \mathbb{R}.$$

#### Theorem 2.14: Properties of the distribution function

A function  $F: \mathbb{R} \to [0,1]$  is the distribution function of a random variable if and only if *i*.) F is nondecreasing

*ii.*) F is right-continuous, that is  $\lim_{y\downarrow x} F(y) = F(x)$ , for all  $x \in \mathbb{R}$ .

Moreover, F is the distribution function of a real-valued random variable if and only if it further holds that

$$\lim_{x\to -\infty} F(x) = 0 \text{ and } \lim_{x\to \infty} F(x) = 1.$$

#### Remark 2.15

Let X be a random variable with distribution function  $F_X$ . It follows by continuity from above and below, respectively, that

$$P\{X=-\infty\}=\lim_{x\to -\infty}F_X(x) \text{ and } P\{X=+\infty\}=1-\lim_{x\to +\infty}F_X(x).$$

These limits are always well-defined by monotonicity of  $F_X$ .

#### Definition 2.16: Law of a random variable

The function

$$L_X(B) \coloneqq P\{X \in B\}, \ B \in \mathfrak{B}_{\overline{\mathbb{R}}}$$

defines a probability measure on  $\overline{\mathbb{R}}$  equipped with the Borel  $\sigma$ -algebra. This function is called the **law** of X and is synonymous with the distribution of X.

## 3 Discrete-Time Markov Chains

## 3.1 Definition of a Markov Chain

#### Definition 3.1: Markov chain

A Markov chain is a process  $(X_n)_{n\geqslant 0}$  satisfying

$$Pr\{X_{n+1} = j \mid X_n = i, X_{n-1} = i_{n-1}, \dots, X_0 = i_0\} = Pr\{X_{n+1} = j \mid X_n = i\}$$

for all  $n \ge 1$  and  $j, i, i_{n-1}, i_0 \in S$ . A Markov chain is said to be **temporally homogeneous** if there are numbers  $(p_{ij})_{i,j \in S}$  such that

$$\Pr\{X_{n+1} = j \mid X_n = i\} = p_{ij}$$

for all  $n \ge 0$  and all states  $i, j \in S$ . The numbers  $(p_{ij})_{i,j \in S}$  are generically referred to as the **transition probabilities** of the Markov chain.

#### **Definition 3.2: Transition matrix**

$$P = \begin{bmatrix} p_{00} & p_{01} & \dots \\ p_{10} & p_{11} & \dots \\ \vdots & \vdots & \ddots \end{bmatrix}$$

The matrix P is called the transition matrix. It is a stochastic matrix, which means it is square with non-negative entries, whose rows sum to one. A Markov chain and transition matrix are equivalent representations of each other.

#### 3.1.1 The Chapman-Kolmogorov equations

## Proposition 3.3: The Chapman-Kolmogorov Equations

We define multi-step transition probabilities

$$P_{ij}^n := \Pr\{X_{n+m} = j \mid X_m = i\}, \quad n, m \geqslant 0.$$

The Chapman-Kolmogorov equations give a recursive formula for computing the n-step transition probabilities.

For all  $m, n \ge 0$  and states i, j,

$$P_{ij}^{\mathfrak{m}+\mathfrak{n}} = \sum_{k} P_{ik}^{\mathfrak{m}} P_{kj}^{\mathfrak{n}}.$$

In particular, we have

$$P^{n} = \begin{bmatrix} P^{n}_{00} & P^{n}_{01} & \dots \\ P^{n}_{10} & P^{n}_{11} & \dots \\ \vdots & \vdots & \ddots \end{bmatrix}$$

where  $P^n$  is the transition matrix P raised to the  $n^{th}$  power. Also note  $P^{m+n} = P^m P^n$ .

#### **Definition 3.4: Irreducible**

A **class** of states is a nonempty set of states such that every state in the set can communicate with one another. These classes form an equivalence relation and partition the state space. We say a Markov chain is irreducible if there is only one class.

## **Definition 3.5: Periodicity**

For a state i, define its period

$$d(i) := gcd\{n \geqslant 1 : P_{i,i}^n > 0\}.$$

States with period 1 are called aperiodic. Periodicity is a class property.

## **Proposition 3.6**

Define the first return time for state  $j \in S$  as

$$T_i := \inf\{n \geqslant 1 : X_n = j\}.$$

Let  $(X_n)_{n\geqslant 0}$  be a Markov chain with transition matrix P. Conditioned on the event  $\{T_j<\infty\}$ , the process  $(X_{n+T_j})_{n\geqslant 0}$  is a Markov chain with transition matrix P and starting state j and is independent of  $X_0,\ldots,X_{T_j}$ .

#### **Proposition 3.7**

A state j is **recurrent** if  $\Pr\{T_j < \infty \mid X_0 = j\} = 1$ , and it is called **transient** if  $\Pr\{T_j < \infty \mid X_0 = j\} < 1$ . Recurrence and transience are class properties.

#### Lemma 3.8

State i is recurrent if and only if  $\sum_{n=1}^{\infty} P_{ii}^n = \infty$ .

## Corollary 3.9

If  $i \leftrightarrow j$  and j is recurrent, then  $Pr\{T_j < \infty | X_0 = i\} = 1$ .

## 3.2 Markov Limit Theorems

## Theorem 3.10: Strong Law of Large Numbers for Markov Chains

Define  $N_j(n)$ ,  $n \ge 1$ , to be the number of transitions into state j, up to and including time n. More precisely,

$$N_i(n) := \#\{1 \leqslant k \leqslant n : X_k = j\}.$$

Also define the expected first return time to be

$$\mu_{jj}\coloneqq \mathbb{E}[T_j|X_0=j].$$

Let  $(X_n)_{n\geqslant 0}$  be a Markov chain starting in state  $X_0=\mathfrak{i}.$  If  $\mathfrak{i}\leftrightarrow\mathfrak{j},$  then

$$\frac{N_j(n)}{n} \to \frac{1}{\mu_{jj}} \ a.s.$$

## Corollary 3.11

For an irreducible Markov chain  $(X_n)_{n\geqslant 0}$ , we have

$$\lim_{n\to\infty}\frac{1}{n}\sum_{k=1}^nP_{ij}^k=\frac{1}{\mu_{jj}}.$$

#### 3.2.1 Stationary distributions and the issue of convergence

#### **Definition 3.12: Stationary Distribution**

A probability distribution  $(\pi_j)_{j \in S}$  is a stationary distribution for a Markov chain with transition probability matrix P if  $\pi_j = \sum_i \pi_i p_{ij}$  for each  $j \in S$ . Equivalently in matrix notation,  $\pi = \pi P$ , when  $\pi$  is considered as a row vector.

#### Definition 3.13: Positive and Null recurrence

A recurrent state j is positive recurrent if  $\mu_{jj} < \infty$ , or null recurrent if  $\mu_{jj} = \infty$ . Positive and null recurrence are class properties.

It is important to note stationary distributions aren't necessarily unique. It is important to note that stationary distribution does not always exist.

#### Theorem 3.14

An irreducible Markov chain satisfies exactly one of the following:

- 1. All states are transient, or all states are null recurrent. In this case,  $\frac{1}{n}\sum_{k=1}^{n}P_{ij}^{k}\to 0$  as  $n\to\infty$  for all states i, j, and no stationary distribution exists.
- 2. All states are positive recurrent. In this case, a unique stationary distribution exists and is given by  $\pi_j = \frac{1}{\mu_{jj}} = \lim_{n \to \infty} \frac{1}{n} \sum_k^n P_{ij}^k.$

#### Theorem 3.15

Let  $(X_n)_{n\geqslant 0}$  be an irreducible, aperiodic, and positive recurrent Markov chain with stationary distribution

 $\pi$ . Then

$$\lim_{n\to\infty}\sum_{j}\left|P^n_{ij}-\pi_j\right|=0 \text{ for all } i\in \mathbb{S}.$$

In particular,  $P_{ij}^n = \pi_j$  for all i, j.

## 3.3 Reversibility and Spectral Gap

#### Definition 3.16: Reversible Markov chain

A Markov chain with transition matrix P and stationary distribution  $\pi$  is reversible if the transition probabilities satisfy

$$\pi_i p_{ij} = \pi_j p_{ji}$$
 for all states i, j.

In this case, we say  $(P, \pi)$  is reversible for convenience.

The equations above are called the detailed balance equations. These equations are also sufficient for reversibility. Hence, if there is a probability distribution  $\pi$  such that the equations are satisfied, then the Markov chain is reversible with stationary distribution  $\pi$ .

## 3.3.1 Spectral gap and trend to equilibrium

#### Definition 3.17: Total variation distance

For probability measures  $\mu, \nu$  on a measurable space  $(\Omega.\mathfrak{F})$ , we define their total variation distance

$$\|\mu - \nu\|_{TV} \coloneqq \sup_{A \in \mathcal{F}} |\mu(A) - \nu(A)|.$$

Also note that when  $\Omega$  is countable, we have

$$\|\mu - \nu\|_{TV} = \frac{1}{2} \sum_{\omega} |\mu(\omega) - \nu(\omega)|.$$

#### Definition 3.18: Spectral gap

Define  $Var_{\pi}(f) = Var(f(X))$  for  $X \sim \pi$  and  $f : S \to \mathbb{R}$ . Also define the function  $Pf : S \to \mathbb{R}$  via the matrix-vector multiplication

$$(Pf)(\mathfrak{i})=\sum_{\mathfrak{j}}\mathfrak{p}_{\mathfrak{i}\mathfrak{j}}f(\mathfrak{j})\text{, }\ \mathfrak{i}\in \mathbb{S}.$$

For a reversible Markov Chain  $(P, \pi)$ , define the spectral gap  $\gamma := 1 - \lambda_2$ , where  $\lambda_2$  is the smallest number satisfying

$$\text{Var}_{\pi}(\text{Pf}) \leqslant \lambda_2 \text{Var}_{\pi}(f) \ \text{ for all } f: \mathbb{S} \to \mathbb{R} \text{ with } \text{Var}_{\pi}(f) < \infty.$$

## Theorem 3.19

If  $(P, \pi)$  is a reversible Markov chain with spectral gap  $\gamma$ , then

$$\|P^n_{i\bullet} - \pi\|^2_{TV} \leqslant \frac{(1-\gamma)^n}{\pi_i} \text{ for each } n \geqslant 0 \text{ and each state i.}$$

Moreover, if the Markov chain is irreducible, aperiodic and has a finite state space, then  $\gamma > 0$ .

## **Definition 3.20**

Let  $\mu$ ,  $\pi$  be probability distributions on state space  $\delta$ . We say that  $\mu$  has density h with respect to  $\pi$  (written  $d\mu = h d\pi$ ) if

$$\mu_j = h(j) \pi_j \ \forall j.$$

## Lemma 3.21

Let  $(P, \pi)$  be a reversible Markov chain. If  $d\mu = hd\pi$ , then  $P^nh$  is the density of  $\mu P^n$  with respect to  $\pi$ .

## 4 Martingales

## 4.1 Definitions and Examples

#### **Definition 4.1**

Let  $(X_n)_{n\geqslant 0}$  be a stochastic process. A process  $(M_n)_{n\geqslant 0}$  is said to be a **martingale with respect to**  $(X_n)_{n\geqslant 0}$  if  $(M_n)_{n\geqslant 0}$  is adapted to  $(X_n)_{n\geqslant 0}$  and, for each  $n\geqslant 0$ ,

(i) 
$$\mathbb{E}|M_n| < \infty$$
;

(ii) 
$$\mathbb{E}[M_{n+1} | X_0, ..., X_n] = M_n$$
.

If the equality in (ii) is replaced by  $\geqslant$  or  $\leqslant$ , then the process is said to be a **submartingale** or **supermartingale**, respectively. The phrase "adapted to" means that  $M_n$  is a measurable function  $(X_0, \ldots, X_n)$  for each  $n \geqslant 0$ .

## **Proposition 4.2**

If  $(M_n)_{n\geqslant 0}$  is a martingale with respect to  $(X_n)_{n\geqslant 0}$ , then for all m>n,

$$\mathbb{E}[M_m \mid X_0, \dots, X_n] = M_n$$
.

If  $(M_n)_{n\geqslant 0}$  is a submartingale or supermartingale, then the equality above is  $\geqslant$  or  $\leqslant$ , respectively.

## 4.2 Stopping Times

#### **Definition 4.3**

A nonnegative integer-valued random variable T is a **stopping time** with respect to  $(X_n)_{n\geqslant 0}$  if, for each  $n\geqslant 0$ , the occurrence of the event  $\{T\leqslant n\}$  is determined entirely by  $(X_0,\ldots,X_n)$ . In other words, the indicator  $1_{\{T\leqslant n\}}$  is a measurable function of  $(X_0,\ldots,X_n)$ .

## **Definition 4.4: Stopped process**

Let  $(X_n)_{n\geqslant 0}$  be a process, and T be a stopping time. If  $(Y_n)_{n\geqslant 0}$  is adapted to  $(X_n)_{n\geqslant 0}$ , then the process  $(Y_{T\wedge n})_{n\geqslant 0}$  is called the **stopped process**. Note that the stopped process satisfies  $Y_{T\wedge n}=Y_n$  for  $n\leqslant T$ , and  $Y_{T\wedge n}=Y_T$  for n>T.

#### 4.2.1 Stopping times and martingales

## **Proposition 4.5**

If  $(M_n)_{n\geqslant 0}$  is a martingale and T is a stopping time, both with respect to  $(X_n)_{n\geqslant 0}$ , then the stopped process  $(M_{T\wedge n})_{n\geqslant 0}$  is also a martingale with respect to  $(X_n)_{n\geqslant 0}$ .

## **Proposition 4.6**

If  $(M_n)_{n\geqslant 0}$  is a submartingale and T is a stopping time, both with respect to  $(X_n)_{n\geqslant 0}$ , then

$$\mathbb{E}[M_0] \leqslant \mathbb{E}[M_{T \wedge n}] \leqslant \mathbb{E}[M_n] \ n \geqslant 0.$$

## **Proposition 4.7: Optimal Stopping Theorem**

Let  $(M_n)_{n\geqslant 0}$  be a submartingale and T be a stopping time, both with respect to  $(X_n)_{n\geqslant 0}$ . If there is a constant  $k<\infty$  such that any one of the following hold

- *i.*)  $T \leq k$  a.s.; or
- ii.)  $|M_n| \le k$  a.s. for each n, and  $Pr\{T < \infty\} = 1$ ; or
- $\mbox{\it iii.}) \ \mathbb{E}[T] < \infty \mbox{ and } |M_n M_{n-1}| \leqslant k \mbox{ a.s. for each } n \geqslant 1,$  then

$$\mathbb{E}[M_0] \leqslant \mathbb{E}[M_T].$$

The inequality above is an equality when  $(M_n)_{n\geqslant 0}$  is a martingale.

## Theorem 4.8: Wald's Identity

Let  $(Y_n)_{n\geqslant 1}$  be adapted to  $(X_n)_{n\geqslant 1}$ . Assume  $Y_{n+1}$  is independent of  $(X_1,\ldots,X_n)$  for each  $n\geqslant 1$ ,  $\sup_{n\geqslant 1}\mathbb{E}|Y_n|<\infty$ , and  $\mathbb{E}[Y_n]=\mu$  for all  $n\geqslant 1$ . If  $T\geqslant 1$  is a stopping time with respect to  $(X_n)_{n\geqslant 1}$  satisfying  $\mathbb{E}[T]<\infty$ , then

$$\mathbb{E}\left[\sum_{n=1}^T Y_n\right] = \mu \mathbb{E}[T].$$

## 5 Poisson Processes

## 5.1 The Exponential Distribution

#### Definition 5.1: Exponential distribution

The **exponential distribution with rate**  $\lambda > 0$ , denote  $Exp(\lambda)$ , has density

$$f(t) = \begin{cases} \lambda e^{-\lambda t} & t \geqslant 0 \\ 0 & t < 0. \end{cases}$$

For  $T \sim Exp(\lambda)$ , the distribution function is given by

$$\Pr\{\mathsf{T}\leqslant\mathsf{t}\} = \begin{cases} 1-e^{-\lambda\mathsf{t}} & \mathsf{t}\geqslant 0\\ 0 & \mathsf{t}<0. \end{cases}$$

We note that  $\mathbb{E}[T] = 1/\lambda$  and  $Var(T) = 1/\lambda^2$ . An important property of exponential random variables is the **memoryless property**. In particular, if  $T \sim Exp(\lambda)$ , then

$$Pr\{T>t+s|T>t\} = \frac{Pr\{T>t+s\}}{Pr\{T>t\}} = \frac{e^{-\lambda(t+s)}}{e^{-\lambda t}} = e^{-\lambda s} = Pr\{T>s\}.$$

## Definition 5.2: Erlang distribution

If,  $T_1, ..., T_k$  are i.i.d. exponential random variables with rate  $\lambda$ , then their sum  $T = T_1 + \cdots + T_k$  has an **Erlang** distribution, with density

$$f_T(t) = \begin{cases} \lambda e^{-\lambda t} \frac{(\lambda t)^{k-1}}{(k-1)!} & t \geqslant 0 \\ 0 & t < 0. \end{cases}$$

#### **5.2** Poisson Processes

## **Definition 5.3: Poisson Process**

Let  $\tau_1, \tau_2, \ldots$  be i.i.d. exponential random variables with rate  $\lambda > 0$  and, for  $n \geqslant 1$ , define  $T_n = \tau_1 + \tau_2 + \cdots + \tau_n$ , with the convention that  $T_0 = 0$ . For each  $t \geqslant 0$ , define the random variable  $N_t = \sup\{n \geqslant 0 : T_n \leqslant t\}$ . The process  $(N_t)_{t\geqslant 0}$  is called a **Poisson process** with rate  $\lambda$ .

This is best thought of as an example of a counting process. A **counting process** is a random process  $(N_t)_{t\geqslant 0}$ , such that (i)  $N_t$  is a non-negative integer for each time  $t\geqslant 0$ ; (ii) the sample paths  $t\mapsto N_t(\omega)$  are non-decreasing in t; and (iii) the sample paths  $t\mapsto N_t(\omega)$  are right-continuous.

#### **Definition 5.4: Poisson distribution**

A random variable X is said to be Poisson distributed with mean  $\lambda \ge 0$  ( $X \sim \text{Poisson}(\lambda)$ ) if X has probability

mass function

$$Pr\{X = k\} = e^{-\lambda} \frac{\lambda^k}{k!}, \ k = 0, 1, 2, ...$$

#### **Proposition 5.5**

If  $(N_t)_{t\geqslant 0}$  is a Poisson process with rate  $\lambda\geqslant 0$ , then for each  $t\geqslant 0$ , we have  $N_t\sim Poisson(\lambda t)$ .

#### Theorem 5.6

Let  $(N_t)_{t\geqslant 0}$  be a Poisson process with rate  $\lambda$ . For any finite collection of distinct time instants  $0=t_0< t_1< \cdots < t_k$ , the increments  $(N_{t_1}-N_{t_0}), \ldots, (N_{t_k}-N_{t_{k-1}})$  are independent with  $(N_{t_i}-N_{t_{i-1}})\sim Poisson(\lambda(t_i-t_{i-1}))$  for each  $1\leqslant i\leqslant k$ .

#### Theorem 5.7: Characterization of Poisson Processes

If  $(N_t)_{t\geqslant 0}$  is a Poisson process, then the following hold:

- 1.  $N_0 = 0$ ;
- 2.  $N_t \sim Poisson(\lambda t) \ \forall t \geqslant 0$ ;
- 3.  $(N_t)_{t\geq 0}$  has independent increments.

Conversely, if these properties hold for a counting process  $(N_t)_{t\geqslant 0}$ , then it is a Poisson process.

## 5.3 Conditioning on Arrivals

#### **Definition 5.8**

Let  $X_1, X_2, ..., X_k$  be a collection of random variables. The **order statistics**  $X_{(1)}, ..., X_{(k)}$  are the random variables defined by sorting the realizations of  $X_1, X_2, ..., X_k$  into increasing order.

#### Theorem 5.9

Let  $(N_t)_{t\geqslant 0}$  be a Poisson process with arrivals  $(T_i)_{i\geqslant 1}$ . Conditioned on the event  $\{N_t=n\}$ , the vector of arrival times  $(T_1,\ldots,T_n)$  has the same distribution as that of order statistics  $(U_{(1)},\ldots,U_{(n)})$ , where  $U_i\sim \text{Unif}(0,t),\ 1\leqslant i\leqslant n$  are independent.

## Theorem 5.10

Let  $(N_t)_{t\geqslant 0}$  be a Poisson process with rate  $\lambda$  and corresponding arrivals  $(T_n)_{n\geqslant 1}$ . For a Borel set  $B\subset [0,\infty)$ , let |B| denotes its Lebesgue volume, and let N(B) denote the number of arrivals in B; i.e.,

$$N(B) = \#\{n \ge 1 : T_n \in B\}.$$

If  $B_1, B_2, \dots \subset [0, \infty)$  are disjoint, bounded Borel sets, then  $N(B_1), N(B_2), \dots$  are independent, with  $N(B_i) \sim$ 

Poisson( $\lambda |B_i|$ ).

## Theorem 5.11: Slivnyak's Theorem

Let  $(N_t)_{t\geqslant 0}$  be a Poisson process with rate  $\lambda$  and let  $x\in (0,\infty)$ . Conditioned on one arrival at time x, the other arrivals form an (unconditional) rate- $\lambda$  Poisson process.

## 6 Continuous-Time Markov Chains

## 6.1 Definitions and Constructions

#### **Definition 6.1**

A process  $(X_t)_{t\geqslant 0}$  taking values in S is a temporally homogeneous **continuous-time Markov chain** if:

- (i) given any initial state  $X_0 = i \in S$ , the sample paths  $t \mapsto X_t$  are a.s. right-continuous (with respect to the discrete topology on S); and
- (ii) for any choice of discrete time instants  $0 \le t_1 < \cdots < t_k < t \le s$  and states  $i_1, i_2, \ldots, i_k, i, j \in S$ , we have the Markov property

$$Pr\{X_s = j \mid X_t = i, X_{t_k} = i_k, ..., X_{t_1} = i_1\} = Pr\{X_{s-t} = j \mid X_0 = i\}.$$

#### Theorem 6.2

Let  $(X_t)_{t\geqslant 0}$  be a continuous-time Markov chain. The transition probabilities satisfy

$$P^{s+t} = P^s P^t$$
 for all  $s, t \ge 0$ ,

and  $\lim_{t\downarrow 0} P^t = I$ . In other words, the transition probabilities  $(P^t)_{t\geqslant 0}$  form a Markov semigroup.

#### Theorem 6.3

Let  $(X_t)_{t\geqslant 0}$  be a continuous-time Markov chain with initial non-absorbing state  $X_0=i$ . The holding time  $T=\inf\{t\geqslant 0: X_t\neq i\}$  has distribution  $T\sim Exp(\lambda_i)$  for  $\lambda_i\geqslant 0$  satisfying

$$P_{ii}^{h} = 1 - h\lambda_i + o(h).$$

Moreover, the next state  $X_T$  is independent of T and has distribution

$$p_{ij} \coloneqq \Pr\{X_T = j \mid X_0 = i\} = \lim_{h \downarrow 0} \frac{P_{ij}^h}{1 - P_{ii}^h}, \quad j \neq i.$$

#### Theorem 6.4

Let  $(X_t)_{t\geqslant 0}$  be a continuous-time Markov chain with transition probabilities  $(P^t)_{t\geqslant 0}$ , starting in non-absorbing state  $X_0=\mathfrak{i}$ , and let  $T=\inf\{t\geqslant 0: X_t\neq \mathfrak{i}\}$  denote the time of the first transition. Conditioned on T and  $X_T=\mathfrak{j}$ , the process  $(X_{T+t})_{t\geqslant 0}$  is a continuous-time Markov chain with transition probabilities  $(P^t)_{t\geqslant 0}$  and starting state  $\mathfrak{j}$ .

#### **Definition 6.5**

The transition probabilities  $(p_{ij})_{i,j\in\mathbb{S}}$  (with  $p_{ii}=0$ ) define a discrete-time Markov chain, known as the **embedded chain**. The parameters  $(\lambda_i)_{i\in\mathbb{S}}$  are called the **transition rates** for the Markov chain,  $\lambda_i$  is

precisely the rate at which the process transitions out of state i.

#### Lemma 6.6

Let  $(p_{ij})_{i,j\in\mathbb{S}}$  be transition probabilities for a discrete-time Markov chain  $(X_n)_{n\geqslant 0}$  starting in non-absorbing state  $X_0=i$ .

(i) The random variable  $N := \inf\{n \ge 0 : X_n \ne i\}$  is geometric with distribution

$$\Pr\{N = k \mid X_0 = i\} = p_{ii}^{k-1}(1 - p_{ii}), k \ge 1$$

(ii) The random variable X<sub>N</sub> is independent of N, and has distribution

$$\Pr\{X_N = j \mid X_0 = i\} = \frac{p_{ii}}{(1 - p_{ii})}, \quad j \neq i.$$

## 6.2 The Infinitesimal Generator

#### **Definition 6.7**

The **infinitesimal generator** for a continuous-time Markov chain  $(X_t)_{t\geqslant 0}$  with transition rates  $(\lambda_i)_{i\in \mathbb{S}}$  is a matrix Q with entries

$$q_{ij} \coloneqq [Q]_{ij} = \begin{cases} \lambda_i p_{ij} & \text{for } j \neq i \\ -\lambda_i & \text{for } j = i, \end{cases}$$

where  $(p_{ij})_{i,j\in\mathbb{S}}$  are the transition probabilities for the embedded chain. In particular,  $\lambda_i = \sum_{j\neq i} q_{ij}$ .

The numbers  $(q_{ij})_{i,j \in S}$  are called the **jump rates** for the Markov chain. Essentially,  $q_{ij}$  describes the rate at which the Markov chain with infinitesimal generator Q transitions from state i to state j  $(j \neq i)$ .

#### 6.2.1 The Kolmogorov differential equations

#### Corollary 6.8

Theorem 6.3 implies the following.

For a continuous-time Markov chain with infinitesimal generator Q, the transition probabilities satisfy

$$P^h_{\text{ii}} = 1 - h \lambda_{\text{i}} + o(h)$$

and

$$P^h_{ij} = hq_{ij} + o(h), \quad j \neq i.$$

## Theorem 6.9: Kolmogorov Differential Equations

Let  $(P^t)_{t\geqslant 0}$  be the transition semigroup for a minimal continuous-time Markov chain with infinitesimal generator Q. The map  $t\mapsto P^t$  is continuously differentiable on  $[0,\infty)$ , and is the (unique) minimal non-negative solution to the differential equations

$$\frac{d}{dt}P^t=QP^t;\;P^0=I \qquad \qquad \text{(Kolmogorov Backward Equation)}$$
 and 
$$\frac{d}{dt}P^t=P^tQ;\;P^0=I. \qquad \qquad \text{(Kolmogorov Forward Equation)}$$

#### Remark 6.10

A Markov semigroup  $(P^t)_{t\geqslant 0}$  with generator Q always satisfies Kolmogorov's backwards equation. In contrast, the forward equation is not satisfied in general, but is satisfied by  $(P_t)_{t\geqslant 0}$  corresponding to the minimal construction of a Markov chain with generator Q.

In the case of finite state space, or more generally bounded transition rates, the Kolmogorov differential equations have a unique solution.

## Corollary 6.11

Let  $(P^t)_{t\geqslant 0}$  be the transition semigroup for a continuous-time Markov chain with infinitesimal general Q. If  $\sup_{i\in S}\lambda_i<\infty$ , then  $P^t=e^{tQ}\coloneqq \sum_{k\geqslant 0}t^k\frac{Q^k}{k!}$  for all  $t\geqslant 0$ .

## 6.2.2 Criteria for non-explosiveness

## **Definition 6.12: Explosiveness**

Define the time of explosion

$$T_{\infty} := \sup_{n \geqslant 1} T_n,$$

$$T_n = \sum_{j=1}^n \tau_j,$$

where  $T_n$  denotes the time of the nth transition. The time of explosion is essentially the time at which an infinite number of transitions have taken place. When a process makes an infinite number of transitions in a finite time, it is known as an **explosion**. We say that a Markov chain is **non-explosive** if  $T_{\infty} = +\infty$  a.s.; otherwise, the chain is said to be **explosive**. Necessary and sufficient conditions for a Markov chain to be non-explosive can be stated in terms of its infinitesimal generator.

#### Theorem 6.13: Reuter's Condition

A Markov chain with infinitesimal generator Q is non-explosive if and only if the only non-negative bounded solution  $\nu = (\nu_i)_{i \in S}$  to  $\nu = Q\nu$  is  $\nu = 0$ .

## 6.3 Continuous-time Markov Limit Theorems

#### **Definition 6.14**

A continuous-time Markov chain is **irreducible** if the embedded chain is irreducible and has at least two states.

## Definition 6.15: Stationary distribution

A stationary distribution for a continuous-time Markov chain with transition probabilities  $(P^t)_{t\geqslant 0}$  is a probability (row) vector p satisfying  $p=pP^t$  for all  $t\geqslant 0$ .

#### Theorem 6.16

Consider a continuous-time Markov chain with infinitesimal generator Q. A probability vector p satisfying  $\sum_i p_i \lambda_i < \infty$  is a stationary distribution if and only if pQ = 0. Moreover, if the chain is irreducible then p is the unique stationary distribution.

#### 6.3.1 Stationary distributions and embedded chains

## Corollary 6.17

Consider and irreducible continuous-time Markov chain with generator Q. The following are equivalent.

- 1. The continuous-time chain has stationary distribution p satisfying  $\sum_i p_i \lambda_i < \infty$ .
- 2. The embedded chain has stationary distribution  $\pi$  satisfying  $\sum_i \pi_i / \lambda_i < \infty$ .

Moreover, if either (and therefore both) are true, then the stationary distributions are unique, and  $\pi_k = C^{-1}p_k\lambda_k$  for all k, where  $C = \sum_i p_i\lambda_i$ .

# 7 Hypothesis Testing

## 7.1 Binary Hypothesis Testing

#### **Definition 7.1: The setup**

On the basis of observing a sample  $\omega \in \Omega$ , we would like to decide whether  $(\Omega, \mathcal{F}, P_0)$  or  $(\Omega, \mathcal{F}, P_1)$  is the better model. The former is called the **null hypothesis**  $H_0$ , and the latter is called the **alternate hypothesis**  $H_1$ .

A **test** is a function  $\hat{H}: \Omega \to \{H_0, H_1\}$  that is measurable in the sense that  $\hat{H}^{-1}(H_0) \in \mathcal{F}$ . Associated with any test  $\hat{H}$  are two fundamental error probabilities: the **Type I error rate** (or, false positive probability), and the **Type II error rate** (or, false negative probability). More precisely,

$$\begin{split} P_0\big\{\hat{H} = H_1\big\} =: \text{Type I error rate (or, false positive probability)} \\ P_1\big\{\hat{H} = H_0\big\} =: \text{Type II error rate (or, false negative probability)}. \end{split}$$

The **power** of a test  $\hat{H}$  is the probability of avoiding a Type II error, and is therefore equal to  $P_1\{\hat{H}=H_1\}$ .

#### 7.1.1 The likelihood ratio

We assume henceforth that  $P_1 \ll P_0$ .

## **Definition 7.2**

The Radon-Nikodym theorem ensures there is a  $\mathcal{F}$ -measurable,  $P_0$ -a.s. unique function  $\Lambda: \Omega \to [0, \infty)$  satisfying the "**change of measure**" identity

$$\mathbb{E}_{P_1}[1_A] = \mathbb{E}_{P_0}[\Lambda 1_A], \text{ for all } A \in \mathcal{F}.$$

The function  $\Lambda$  is called a Radon-Nikodym derivative (usually denoted by  $\frac{dP_1}{dP_0}$ ), but in the context of hypothesis testing it is generally referred to as the **likelihood ratio** because it can be thought of as the relative likelihood of observing a sample  $\omega$  under the different hypotheses  $H_1$  and  $H_0$ .  $\Lambda$  is simply the ratio of densities, or if  $\Omega$  is discrete, then  $\Lambda$  is the ratio of the probability mass functions.

#### 7.1.2 Threshold tests and the error curve

#### **Definition 7.3**

Assume  $P_1 \ll P_0$ , and let  $\eta \geqslant 0$ . The threshold test with threshold  $\eta$ , denote  $\hat{H}_{\eta}$ , is defined according to

$$\hat{H}_{\eta} = \begin{cases} H_1 & \text{ if } \Lambda(\omega) \geqslant \eta \\ H_0 & \text{ if } \Lambda(\omega) < \eta. \end{cases}$$

## Example 7.4: MAP and ML tests

The maximum a posteriori test (MAP) is a threshold test where we have a prior belief that  $H_0$  is true with probability  $\pi_0 < 1$  and  $H_1$  is true with probability  $\pi_1 = 1 - \pi_0$ . The MAP test is the threshold test with threshold  $\eta = \pi_0/\pi_1$ , and has the property that it minimizes the total error rate among all tests. Under prior  $\pi$ , the total error probability for any test  $\hat{H}$  satisfies

$$Pr\left\{\hat{H} \text{ errors}\right\} = \pi_0 P_0 \left\{\hat{H} = H_1\right\} + \pi_1 P_1 \left\{\hat{H} = H_0\right\} \geqslant Pr\left\{\hat{H}_{MAP} \text{ errors}\right\}.$$

The maximum likelihood (ML) test is defined to be the threshold test with  $\eta = 1$ . It minimizes the sum of Type I and Type II error rates, which follows from the previous example when  $\pi_0 = \pi_1$ .

#### **Definition 7.5**

The threshold tests  $(\hat{H}_{\eta})_{\eta\geqslant 0}$  define a function called the **error curve**, which plays a fundamental role in characterizing the best tradeoff between Type I and Type II error rates.

Assume  $P_1 \ll P_0$  and let  $\Lambda$  denote the likelihood ration. The error curve  $\mathfrak{u}: [0,1] \to \mathbb{R}$  is defined via

$$u(\theta)\coloneqq \underset{\eta\geqslant 0}{sup}\left\{P_{1}\left\{\hat{H}_{\eta}=H_{0}\right\}+\eta\left(P_{0}\left\{\hat{H}_{\eta}=H_{1}\right\}-\theta\right)\right\}\text{, }0\leqslant \theta\leqslant 1.$$

Note that, as the pointwise supremum of affine functions, u is a convex function on [0,1].

## 7.1.3 The Neyman-Pearson lemma

We say that a test Ĥ lies above the error curve if

$$P_1 \{ \hat{H} = H_0 \} \ge u(P_0 \{ \hat{H} = H_1 \}),$$

and we say that Ĥ lies on the error curve if this is met with equality.

## Theorem 7.6: Neyman-Pearson Lemma

Assume  $P_1 \ll P_0$ . All tests  $\hat{H}$  lie above the error curve. Moreover, every threshold test  $\hat{H}_{\eta}$  lies on the error curve. The Neyman–Pearson lemma ensures that threshold tests are optimal in the sense that they lie on the error curve, and any other test lies above.

#### Definition 7.7: Randomized threshold test

Fix parameters  $\eta_0, \eta_1 \geqslant 0$  and  $p \in [0,1]$ . The corresponding (randomized) threshold test  $\hat{H}$  is defined by taking  $R \sim \text{Bernoulli}(p)$ , and putting

$$\hat{H}(\omega) = \mathbf{1}_{R=0} \hat{H}_{\eta_0}(\omega) + \mathbf{1}_{R=1} \hat{H}_{\eta_1}(\omega).$$

As a result, by varying the parameters  $\eta_0, \eta_1 \geqslant 0$  and  $p \in [0, 1]$ , any point on the error curve is achievable by a (possibly randomized) threshold test. For a fixed  $\theta \in [0, 1]$ , the **Neyman-Pearson rule** is defined

to be the (possibly randomized) threshold test that achieves Type II error probability  $\mathfrak{u}(\theta)$  subject to the constant that Type I error probability does not exceed  $\theta$ . In other words, subject to a Type I error constraint, the Neyman–Pearson rule is the most powerful test.

#### Definition 7.8: Sufficient statistic

A mapping  $T:\omega\mapsto T(\omega)$  is said to be a **sufficient statistic** if there exists a function  $\nu$  such that  $\Lambda(\omega)=\nu\circ T(\omega)$  for all  $\omega\in\Omega$ .

## 7.2 Sequential Analysis

#### **Definition 7.9**

Let  $P_0$ ,  $P_1$  be probability distributions on, say,  $\mathbb{R}$ , and consider i.i.d. random variables  $(X_n)_{n\geqslant 1}$ , with common distribution Q. Let the null hypothesis be  $H_0: Q=P_0$ , and the alternate hypothesis be  $H_1: Q=P_1$ . Let  $(\hat{H}_n)_{n\geqslant 1}$  be a sequence of tests adapted to  $(X_n)_{n\geqslant 1}$ , and T be a stopping time. This induces Type I and Type II error rates for the **sequential test**  $\hat{H}_T$  equal to

$$\alpha := P_0 \{ \hat{H}_T = H_1 \}$$
, and  $\beta := P_1 \{ \hat{H}_T = H_0 \}$ ,

respectively, where we abuse notation and abbreviate

$$P_{\mathfrak{i}}\left\{\cdot\right\}=Pr\left\{\cdot\mid H_{\mathfrak{i}} \text{ true}\right\}, \quad \mathfrak{i}=0,1.$$

## 7.2.1 Average sample requirements

#### **Definition 7.10**

For probability measures  $u \ll v$  with corresponding likelihood ratio, the relative entropy between u and v is defined as

$$D(\mathfrak{u} \parallel \nu) \coloneqq \mathbb{E}_{\mathfrak{u}}[\log \Lambda] = \mathbb{E}_{\nu}[\Lambda \log \Lambda].$$

The entropy is greater than or equal to 0 due to convexity, with equality if and only if u = v. For two reals  $p, q \in [0, 1]$ , we abuse notation slightly and write

$$D(p \parallel q) \coloneqq p \log \left(\frac{p}{q}\right) + (1-p) \log \left(\frac{1-p}{1-q}\right).$$

#### Theorem 7.11

Let the above notation prevail, and let  $\mathbb{E}_{\mathfrak{i}}[\cdot]$  denote the expectaion under hypothesis  $H_{\mathfrak{i}}$  for  $\mathfrak{i}=0,1$ . If  $\alpha+\beta\leqslant 1$ , it holds that

$$\mathbb{E}_0[\mathsf{T}] \geqslant \frac{\mathsf{D}(\alpha \parallel 1 - \beta)}{\mathsf{D}(\mathsf{P}_0 \parallel \mathsf{P}_1)}, \text{ and } \mathbb{E}_1[\mathsf{T}] \geqslant \frac{\mathsf{D}(1 - \beta \parallel \alpha)}{\mathsf{D}(\mathsf{P}_1 \parallel \mathsf{P}_0)}.$$

Note that T is the number of samples.

## 7.2.2 The sequential probability ratio test

#### **Definition 7.12**

Consider the setting where  $P_0 \neq P_1$ , and assume  $P_1 \ll P_0$ , with likelihood ratio  $\Lambda = \frac{dP_1}{dP_0}$ . Fix two thresholds  $\eta_0 < \eta_1$ . For i.i.d. observations  $(X_n)_{n\geqslant 1}$  as before, define the sequence of likelihoods

$$L_n = \prod_{i=1}^n \Lambda(X_i), \quad n \geqslant 1.$$

Define the stopping time  $T=\inf\{n\geqslant 1: L_n\notin (\eta_0,\eta_1)\}$ , and the corresponding **sequential probability ratio** test

$$\hat{H}_T = \begin{cases} H_0 & \text{if } L_T \leqslant \eta_0 \\ H_1 & \text{if } L_T \geqslant \eta_0. \end{cases}$$

This induces Type I and Type II error rates for the sequential test  $\hat{H}_T$  equal to

$$\alpha \coloneqq P_0 \left\{ \hat{H}_T = H_1 \right\} \text{, and } \beta \coloneqq P_1 \left\{ \hat{H}_T = H_0 \right\}.$$

## **Proposition 7.13**

For thresholds  $0 \leqslant \eta_0 < \eta_1$ , the Type I/II error rates for the corresponding sequential probability ratio test  $\hat{H}_T$  satisfy

$$\frac{\alpha}{1-\beta}\leqslant \frac{1}{\eta_1} \text{ and } \frac{\beta}{1-\alpha}\leqslant \eta_0.$$

## Theorem 7.14

Fix thresholds  $\eta_0 < \eta_1$ , and let  $\alpha, \beta$  denote the Type I/II error rates realized by the corresponding sequential probability ratio test  $\hat{H}_T$ . If  $D(P_1 \parallel P_0) < \infty$  and  $D(P_0 \parallel P_1) < \infty$ , then

$$\mathbb{E}_0[T] \simeq \frac{D(\alpha \parallel 1 - \beta)}{D(P_0 \parallel P_1)} \text{, and } \mathbb{E}_1[T] = \frac{D(1 - \beta \parallel \alpha)}{D(P_1 \mid P_0)}.$$

In this case, the approximations above are fairly accurate, demonstrating (approximate) optimality of the sequential probability ratio test. This assertion of optimality can be made more precise: among all tests with the same power, the sequential probability ratio test requires the fewest samples on average.

## Lemma 7.15

Let  $(Z_n)_{n\geqslant 1}$  be i.i.d. random variables. For given a < b, define the stopping time

$$K = \inf \left\{ n \geqslant 1 : \sum_{i=1}^{n} Z_i \notin (a, b) \right\}.$$

If  $Pr\{|Z_1| > 0\} > 0$ , then  $\mathbb{E}[K] < \infty$ .

This lemma guarantees that any sequential probability ratio test will require finitely many samples on average.