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# RAJSHAHI UNIVERSITY OF ENGINEERING & TECHNOLOGY

Department of Computer Science & Engineering

ALGORITHMS ANALYSIS & DESIGN SESSIONAL

## Matrix Multiplication

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# 1 Matrix Multiplication

## 1.1 Problem Statement

Implement Iterative Matrix Multiplication, Divide-and-Conquer Matrix Multiplication and Strassen's Matrix Multiplication algorithm to multiply two matrix of valid dimension. Compare the 3 algorithms based on various input size on randomly generated matrices. The comparison metric should be the execution time of each matrix multiplication algorithm.

## 1.2 Code

Listing 1: *matrix\_multiplication.cpp*

```
1  #include <chrono>
2  #include <iostream>
3  #include <random>
4  #include <vector>
5  using namespace std;
6  using namespace std::chrono;
7  typedef vector<vector<int>> Matrix;
8  Matrix generate_random_matrix(int size, int min_val = 1, int max_val = 100)
   ↪ {
9      random_device rd;
10     mt19937 gen(rd());
11     uniform_int_distribution<> dis(min_val, max_val);
12
13     Matrix matrix(size, vector<int>(size));
14     for (int i = 0; i < size; ++i) {
15         for (int j = 0; j < size; ++j) {
16             matrix[i][j] = dis(gen);
17         }
18     }
19     return matrix;
20 }
21
22 void iterative_matmul(const Matrix &A, const Matrix &B, Matrix &C, int
   ↪ size) {
23     for (int i = 0; i < size; ++i) {
24         for (int j = 0; j < size; ++j) {
25             C[i][j] = 0;
26             for (int k = 0; k < size; ++k) {
27                 C[i][j] += A[i][k] * B[k][j];
28             }
29         }
30     }
31 }
32
33 void recursive_matmul(const Matrix &A, const Matrix &B, Matrix &C, int
   ↪ row_A,
34                        int col_A, int row_B, int col_B, int row_C, int col_C,
35                        int size) {
```

```

36     if (size == 1) {
37         C[row_C][col_C] += A[row_A][col_A] * B[row_B][col_B];
38         return;
39     }
40
41     int new_size = size / 2;
42
43     recursive_matmul(A, B, C, row_A, col_A, row_B, col_B, row_C, col_C,
44         ↪ new_size);
45     recursive_matmul(A, B, C, row_A, col_A + new_size, row_B + new_size,
46         ↪ col_B,
47         row_C, col_C, new_size);
48     recursive_matmul(A, B, C, row_A, col_A, row_B, col_B + new_size, row_C,
49         ↪ col_C + new_size, new_size);
50     recursive_matmul(A, B, C, row_A, col_A + new_size, row_B + new_size,
51         ↪ col_B + new_size, row_C, col_C + new_size, new_size);
52     recursive_matmul(A, B, C, row_A + new_size, col_A, row_B, col_B,
53         ↪ row_C + new_size, col_C, new_size);
54     recursive_matmul(A, B, C, row_A + new_size, col_A + new_size,
55         ↪ row_B + new_size, col_B, row_C + new_size, col_C,
56         ↪ new_size);
57     recursive_matmul(A, B, C, row_A + new_size, col_A, row_B, col_B +
58         ↪ new_size,
59         row_C + new_size, col_C + new_size, new_size);
60     recursive_matmul(A, B, C, row_A + new_size, col_A + new_size,
61         ↪ row_B + new_size, col_B + new_size, row_C + new_size,
62         ↪ col_C + new_size, new_size);
63 }
64
65 Matrix add(const Matrix &A, const Matrix &B) {
66     int size = A.size();
67     Matrix result(size, vector<int>(size, 0));
68     for (int i = 0; i < size; ++i)
69         for (int j = 0; j < size; ++j)
70             result[i][j] = A[i][j] + B[i][j];
71     return result;
72 }
73
74 Matrix subtract(const Matrix &A, const Matrix &B) {
75     int size = A.size();
76     Matrix result(size, vector<int>(size, 0));
77     for (int i = 0; i < size; ++i)
78         for (int j = 0; j < size; ++j)
79             result[i][j] = A[i][j] - B[i][j];
80     return result;
81 }
82
83 Matrix strassen(const Matrix &A, const Matrix &B) {

```

```

83     int size = A.size();
84     Matrix C(size, vector<int>(size, 0));
85     if (size == 1) {
86         C[0][0] = A[0][0] * B[0][0];
87         return C;
88     }
89
90     int new_size = size / 2;
91
92     Matrix A11(new_size, vector<int>(new_size));
93     Matrix A12(new_size, vector<int>(new_size));
94     Matrix A21(new_size, vector<int>(new_size));
95     Matrix A22(new_size, vector<int>(new_size));
96
97     Matrix B11(new_size, vector<int>(new_size));
98     Matrix B12(new_size, vector<int>(new_size));
99     Matrix B21(new_size, vector<int>(new_size));
100    Matrix B22(new_size, vector<int>(new_size));
101
102    Matrix C11(new_size, vector<int>(new_size));
103    Matrix C12(new_size, vector<int>(new_size));
104    Matrix C21(new_size, vector<int>(new_size));
105    Matrix C22(new_size, vector<int>(new_size));
106
107    Matrix P5(new_size, vector<int>(new_size));
108    Matrix P3(new_size, vector<int>(new_size));
109    Matrix P1(new_size, vector<int>(new_size));
110    Matrix P4(new_size, vector<int>(new_size));
111    Matrix P2(new_size, vector<int>(new_size));
112    Matrix P7(new_size, vector<int>(new_size));
113    Matrix P6(new_size, vector<int>(new_size));
114
115    Matrix S1(new_size, vector<int>(new_size));
116    Matrix S2(new_size, vector<int>(new_size));
117    Matrix S3(new_size, vector<int>(new_size));
118    Matrix S4(new_size, vector<int>(new_size));
119    Matrix S5(new_size, vector<int>(new_size));
120    Matrix S6(new_size, vector<int>(new_size));
121    Matrix S7(new_size, vector<int>(new_size));
122    Matrix S8(new_size, vector<int>(new_size));
123    Matrix S9(new_size, vector<int>(new_size));
124    Matrix S10(new_size, vector<int>(new_size));
125
126    for (int i = 0; i < new_size; ++i) {
127        for (int j = 0; j < new_size; ++j) {
128            A11[i][j] = A[i][j];
129            A12[i][j] = A[i][j + new_size];
130            A21[i][j] = A[i + new_size][j];
131            A22[i][j] = A[i + new_size][j + new_size];
132
133            B11[i][j] = B[i][j];

```

```

134     B12[i][j] = B[i][j + new_size];
135     B21[i][j] = B[i + new_size][j];
136     B22[i][j] = B[i + new_size][j + new_size];
137 }
138 }
139 // P1 = A11 * (B12 - B22)
140 S1 = subtract(B12, B22);
141 P1 = strassen(A11, S1);
142
143 // P2 = (A11 + A12) * B22
144 S2 = add(A11, A12);
145 P2 = strassen(S2, B22);
146
147 // P3 = (A21 + A22) * B11
148 S3 = add(A21, A22);
149 P3 = strassen(S3, B11);
150
151 // P4 = A22 * (B21 - B11)
152 S4 = subtract(B21, B11);
153 P4 = strassen(A22, S4);
154
155 // P5 = (A11 + A22) * (B11 + B22)
156 S5 = add(A11, A22);
157 S6 = add(B11, B22);
158 P5 = strassen(S5, S6);
159
160 // P6 = (A12 - A22) * (B21 + B22)
161 S7 = subtract(A12, A22);
162 S8 = add(B21, B22);
163 P6 = strassen(S7, S8);
164
165 // P7 = (A21 - A11) * (B11 + B12)
166 S9 = subtract(A21, A11);
167 S10 = add(B11, B12);
168 P7 = strassen(S9, S10);
169
170 C11 = subtract(add(add(P5, P4), P6), P4);
171 C12 = add(P1, P2);
172 C21 = add(P3, P4);
173 C22 = subtract(add(add(P5, P1), P7), P3);
174
175 // combine C11, C12, C21, C22 into C
176 for (int i = 0; i < new_size; ++i) {
177     for (int j = 0; j < new_size; ++j) {
178         C[i][j] = C11[i][j];
179         C[i][j + new_size] = C12[i][j];
180         C[i + new_size][j] = C21[i][j];
181         C[i + new_size][j + new_size] = C22[i][j];
182     }
183 }
184 return C;

```

```

185 }
186
187 void print_matrix(const Matrix &matrix) {
188     for (const auto &row : matrix) {
189         for (const auto &val : row) {
190             cout << val << " ";
191         }
192         cout << endl;
193     }
194 }
195
196 int main() {
197     for (int n = 2; n <= 256; n *= 2) {
198         Matrix A = generate_random_matrix(n);
199         Matrix B = generate_random_matrix(n);
200
201         Matrix C_iterative(n, vector<int>(n, 0));
202         Matrix C_recursive(n, vector<int>(n, 0));
203         Matrix C_strassen(n, vector<int>(n, 0));
204         cout << "n = " << n << endl;
205         auto start_time = high_resolution_clock::now();
206         iterative_matmul(A, B, C_iterative, n);
207         auto end_time = high_resolution_clock::now();
208         double time_taken =
209             duration_cast<nanoseconds>(end_time - start_time).count() * 1e-6;
210         cout << "Iterative: " << time_taken << " ms" << endl;
211
212         start_time = high_resolution_clock::now();
213         recursive_matmul(A, B, C_recursive, 0, 0, 0, 0, 0, 0, n);
214         end_time = high_resolution_clock::now();
215         time_taken =
216             duration_cast<nanoseconds>(end_time - start_time).count() * 1e-6;
217         cout << "Divide & Conquer: " << time_taken << " ms" << endl;
218
219         start_time = high_resolution_clock::now();
220         C_strassen = strassen(A, B);
221         end_time = high_resolution_clock::now();
222         time_taken =
223             duration_cast<nanoseconds>(end_time - start_time).count() * 1e-6;
224         cout << "Strassen's: " << time_taken << " ms" << endl;
225     }
226     return 0;
227 }

```



### 1.3 Output

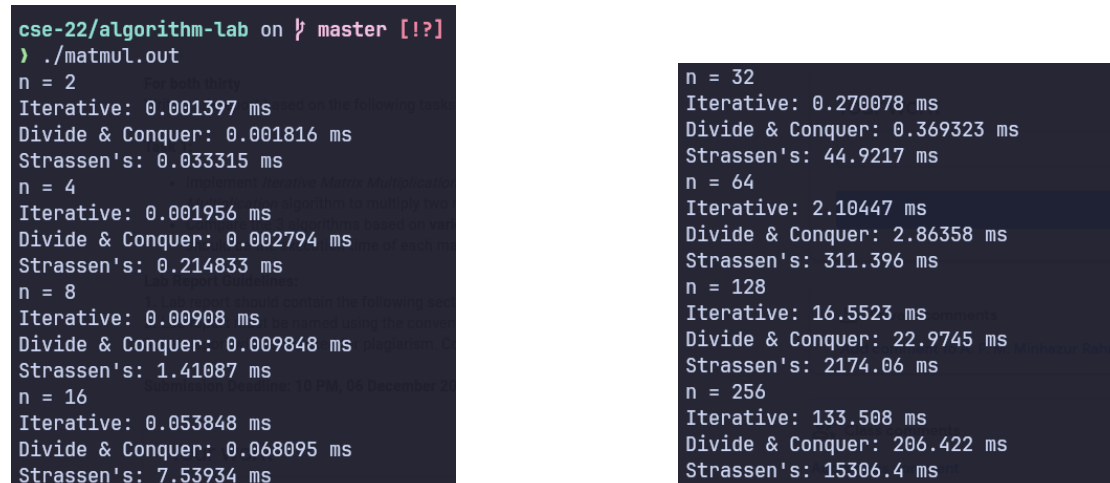


Figure 1: Execution time for matrix multiplication

### 1.4 Result

Table 1: Comparison table for Iterative & Divide & conquer and Strassen's Matrix multiplication algorithm execution time

Input Size	Iterative (ms)	Divide & Conquer (ms)	Strassen's (ms)
$2 \times 2$	0.001397	0.001816	0.033315
$4 \times 4$	0.001956	0.002794	0.214833
$8 \times 8$	0.00908	0.009848	1.41087
$16 \times 16$	0.053848	0.068095	7.53934
$32 \times 32$	0.270078	0.369323	44.9217
$64 \times 64$	2.10447	2.86358	311.396
$128 \times 128$	16.5523	22.9745	2174.06
$256 \times 256$	133.508	206.422	15306.4

### 1.5 Analysis & Discussion

The iterative algorithm demonstrates consistent performance and scales predictably with increasing input size. Its time complexity is  $O(n^3)$ , which aligns with theoretical expectations. As seen in the results, the runtime increases steadily, from 0.001397 ms for  $2 \times 2$  matrices to 133.508 ms for  $256 \times 256$  matrices. Despite its cubic complexity, the iterative approach has minimal overhead, making it an efficient choice for small to medium-sized matrices. Its straightforward implementation also contributes to its practicality in many applications.

The divide-and-conquer (recursive) algorithm shares the same  $O(n^3)$  time complexity but incurs additional overhead due to recursive function calls. This overhead is noticeable for smaller matrices, where it performs slightly worse than the iterative algorithm. For example, at  $2 \times 2$ , the recursive approach takes 0.001816 ms, which is slower than the iterative method's 0.001397 ms. As the input size grows,

the performance gap widens, with the recursive algorithm taking 206.422 ms for  $256 \times 256$  matrices, compared to the iterative method's 133.508 ms. The recursion overhead thus limits its practicality for general use.

Strassen's algorithm theoretically offers better time complexity, approximately  $O(n^{2.81})$ , by reducing the number of multiplications required. However, in practice, it incurs significant overhead due to matrix partitioning, additional additions and subtractions, and increased recursion depth. This overhead dominates for smaller and moderately sized matrices, resulting in much slower performance compared to both iterative and recursive methods. For instance, at  $64 \times 64$ , Strassen's method takes 311.396 ms, compared to 2.10447 ms and 2.86358 ms for the iterative and recursive methods, respectively. For larger sizes, such as  $256 \times 256$ , the overhead becomes even more pronounced, with Strassen's algorithm taking 15306.4 ms, which is substantially slower than the other methods.

The results highlight the impact of algorithmic overhead. While Strassen's algorithm provides theoretical improvements, its practical utility is hindered by its computational and memory overhead. The additional memory requirements for submatrices and intermediate computations can also lead to cache inefficiencies, particularly for larger matrices. These factors explain why Strassen's algorithm underperforms compared to simpler methods in the tested range.

In conclusion, the iterative algorithm remains the most practical and efficient choice for small to medium-sized matrices due to its simplicity and low overhead. While Strassen's algorithm is theoretically advantageous, its overhead makes it unsuitable for smaller matrices, and it only becomes competitive for significantly larger input sizes. Future optimizations, such as hybrid approaches or parallelized implementations, may help mitigate Strassen's overhead and unlock its potential for practical applications involving very large matrices.