

# Data Structures and Algorithms

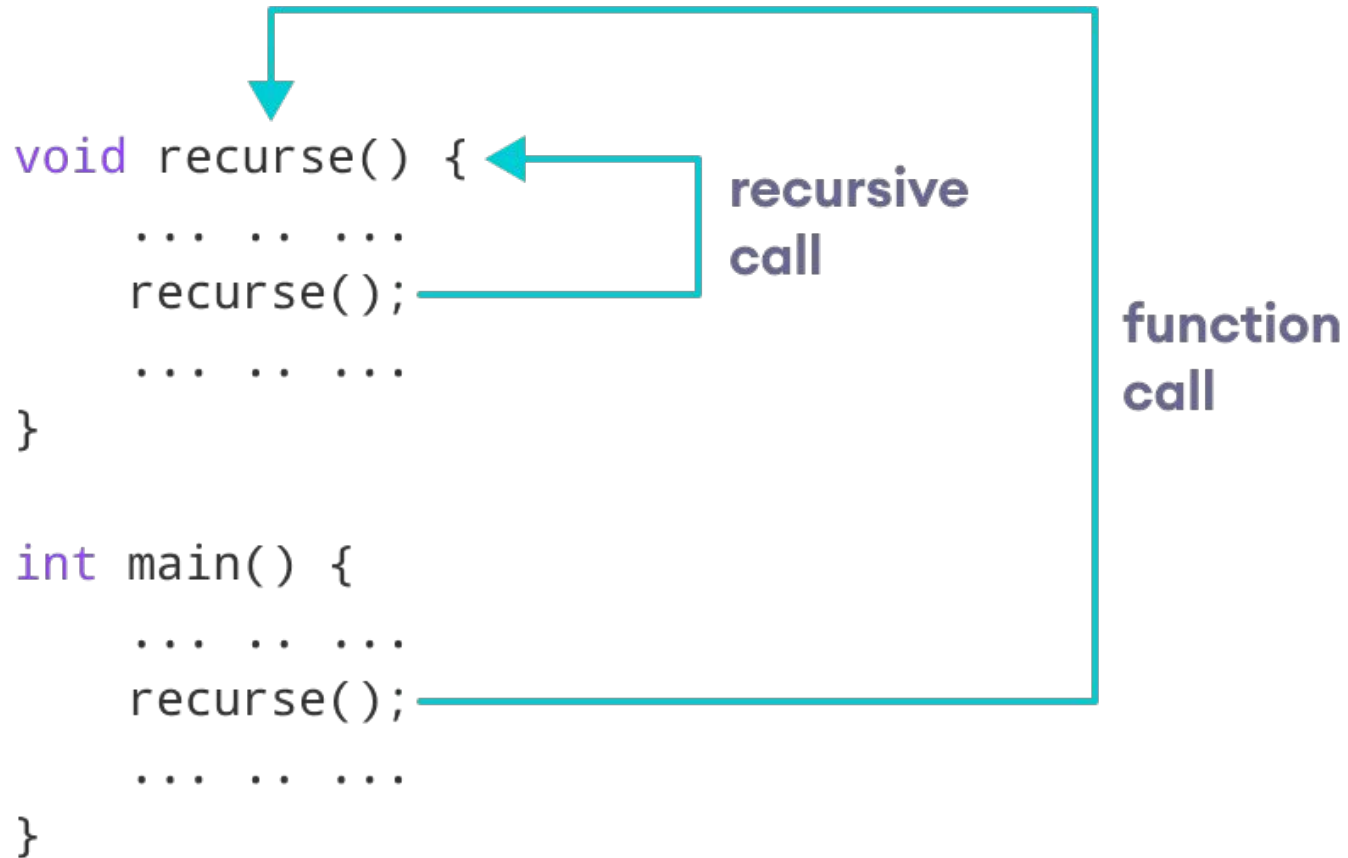
Episode 8

**RECURSION AND SEARCH ALGORITHMS**

# *Recursive Functions*

- A recursive function is a **function that calls itself**
  - With each recursive call, the problem size is reduced
  - The function terminates when **a terminal condition is reached**
- Recursion must terminate; otherwise, we have infinite recursion
- A recursive definition consists of two parts.
  - The first part is called the anchor or the base case.
  - In the second part, rules are given that allow for the construction of new elements out of basic elements.

# *Recursive Functions*



- recursiveFunction:
  - if (**base case**){ // there could be more than one base case  
compute the solution without recursion  
}
  - else {  
break the problem into sub-problem(s) of the same form and call the recursive function on each sub-problem  
  
Assemble the results of the sub-problems  
}

# *Three Musts for Recursion*

- 1. Your code must have a valid case for each input
- 2. Your code must have at least one base case
- 3. When making a recursive call, it must be a simpler instance that progresses towards the base case

# The factorial function !

- The factorial function ! , can be defined in the following manner:

$$n! = \begin{cases} 1 & \text{if } n = 1 \text{ (anchor)} \\ n \cdot (n - 1)! & \text{if } n > 1 \text{ (inductive step)} \end{cases}$$

- $6! = 6 \cdot 5 \cdot 4 \cdot 3 \cdot 2 \cdot 1 = 720$

# The factorial function !

```
int factorial(int n) {
    if (n > 1) {
        return n * factorial(n - 1);
    } else {
        return 1;
    }
}

int main() {
    int n, result;

    cout << "Enter a non-negative number: ";
    cin >> n;

    result = factorial(n);
    cout << "Factorial of " << n << " = " << result;
    return 0;
}
```

factorial(6)

**Call stack**



factorial(5)



**return** 6 x factorial(5)

**Call stack**

factorial(4)



**return** 5 x factorial(4)

**return** 6 x factorial(5)

**Call stack**

factorial(3)



**return** 4 x factorial(3)

**return** 5 x factorial(4)

**return** 6 x factorial(5)

**Call stack**

factorial(2)



**return** 3 x factorial(2)

**return** 4 x factorial(3)

**return** 5 x factorial(4)

**return** 6 x factorial(5)

**Call stack**

factorial(1)



**return** 2 x factorial(1)

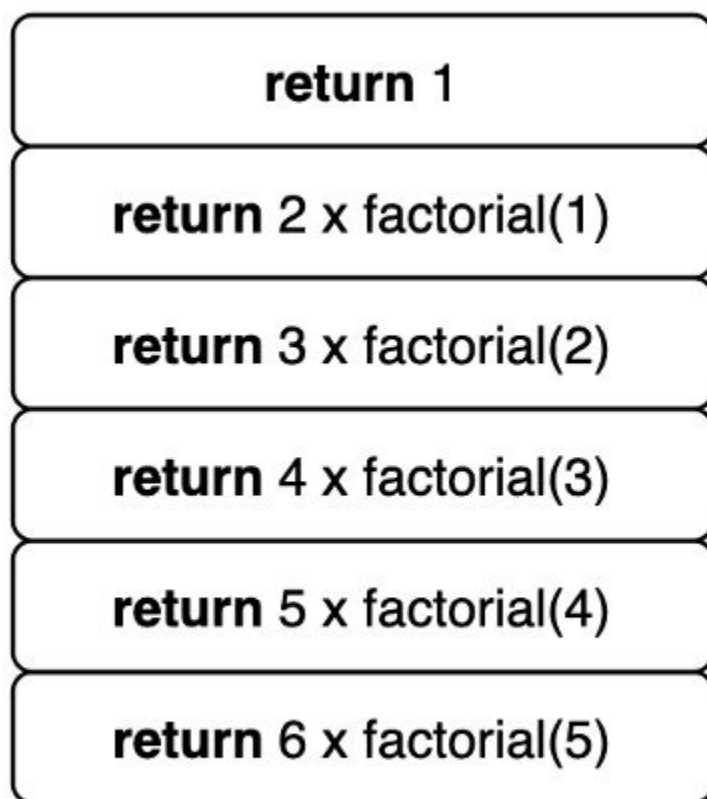
**return** 3 x factorial(2)

**return** 4 x factorial(3)

**return** 5 x factorial(4)

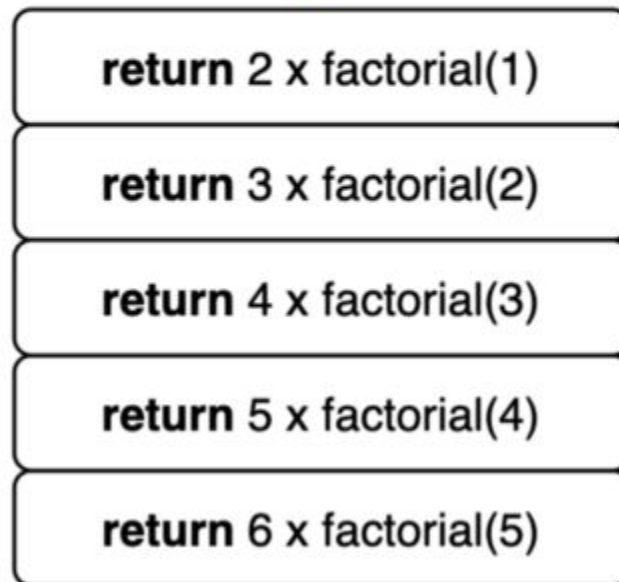
**return** 6 x factorial(5)

**Call stack**

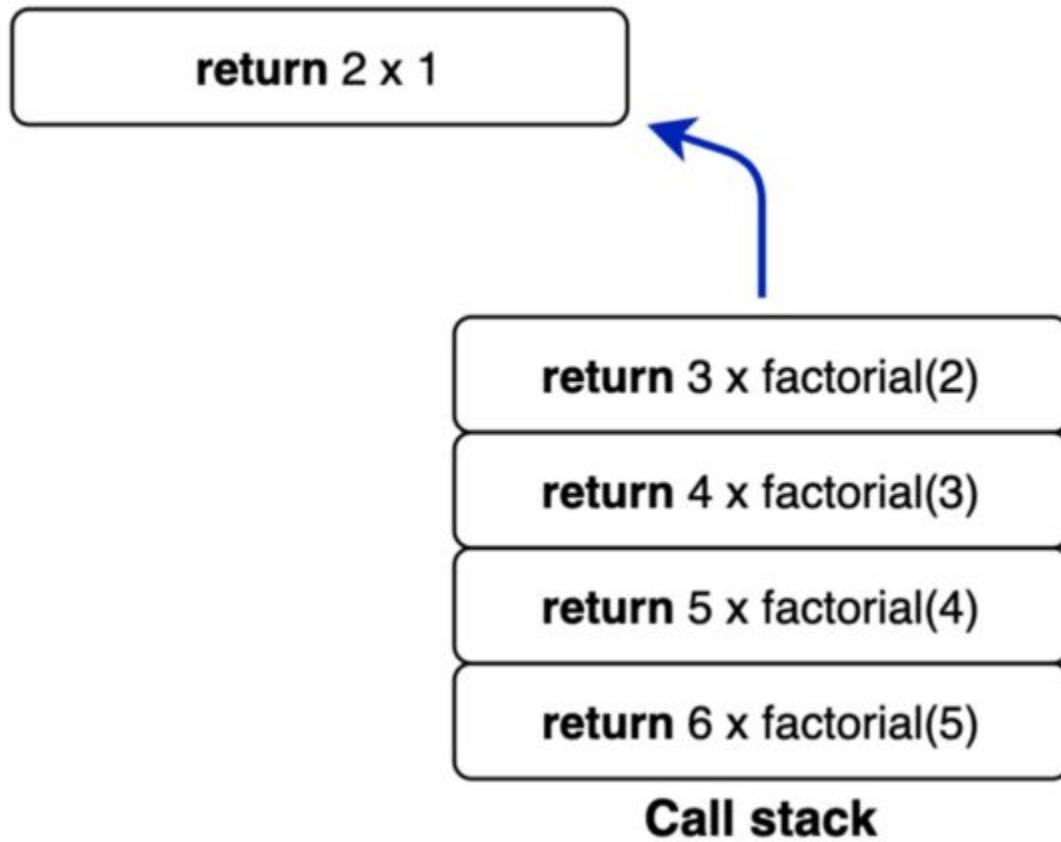


**Call stack**

**return 1**

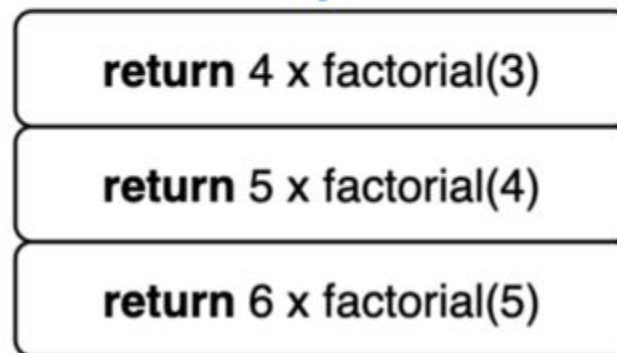


**Call stack**





**return 3 x 2**



**Call stack**

**return 4 x 6**



**return 5 x factorial(4)**  
**return 6 x factorial(5)**

**Call stack**

**return 5 x 24**

**return 6 x factorial(5)**

**Call stack**

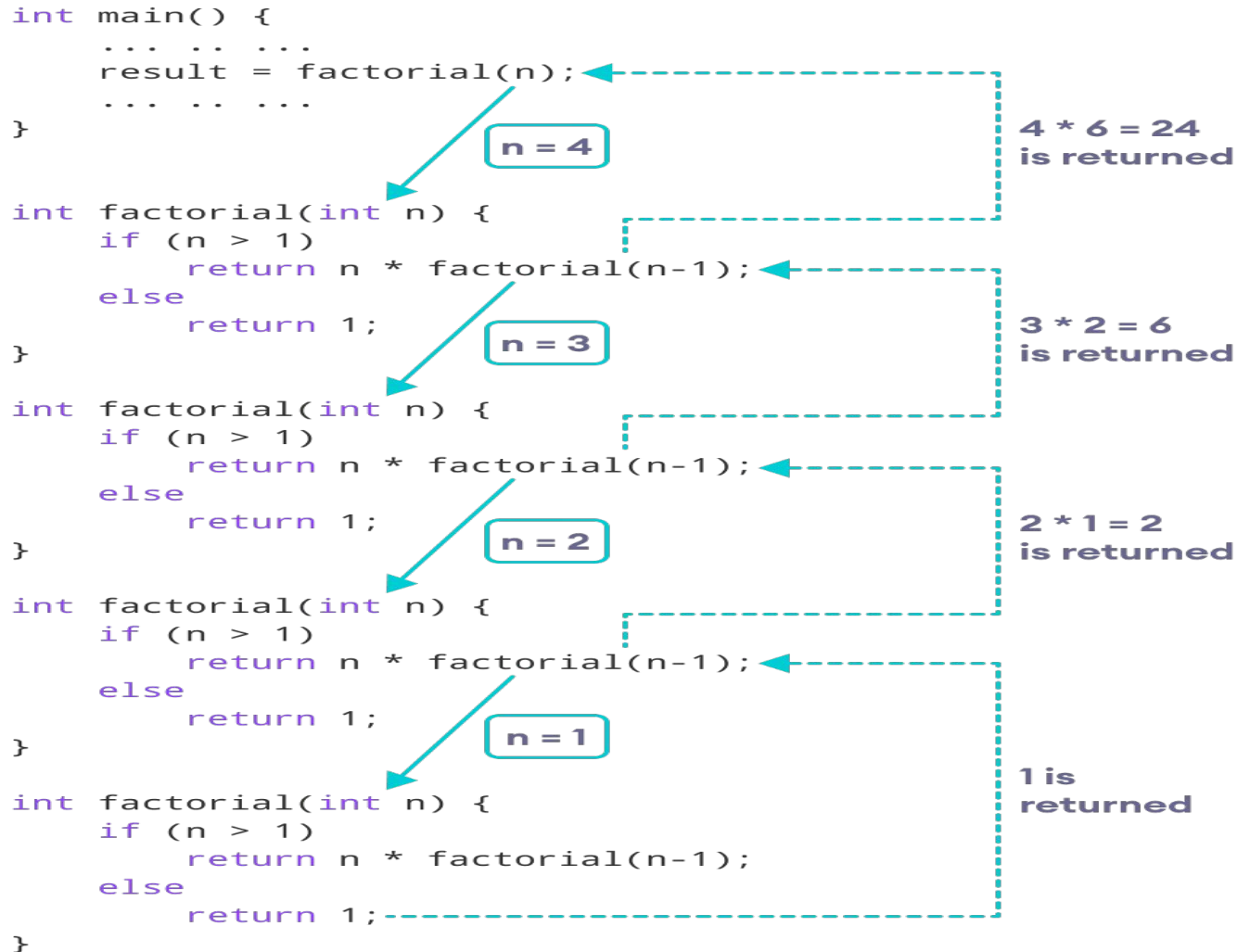


**return** 6 x 120



**Call stack**

# The factorial function !



# Search Algorithms

- **Linear** (*Sequential*) Searching Algorithm
  - Can deal with unsorted arrays
- **Binary** Searching Algorithm
  - Must deal with sorted arrays

# Linear Search Algorithm

- Compare target value with all elements of the array until you find a match or reach the last element in the array
- EX: arr = [25,30,6,17,27,11], n = 6, target = 17

17

25

30

6

17

27

11



Found

# Linear Search Algorithm

```
int linSearch(int list [ ], int listLength, int  
key)  
{  
    int loc;  
  
    for(loc = 0; loc < listLength; loc++)  
    {  
        if(list[loc] == key)  
            return loc;  
    }  
    return -1;  
}
```



# Linear Search Algorithm

- The best case time complexity of linear search takes place when the element to be searched for is on the first index □ Time complexity is 1
- The worst case will take place if:
  - The element to be search is in the last index
  - The element to be search is not present in the list
  - □ Time Complexity is N

# Recursive Linear Search

$$f(a, n, target) \begin{cases} -1 & \text{if } n \leq 0 \\ n-1 & \text{if } target = a[n-1] \\ f(a, n-1, target) & \text{Otherwise} \end{cases}$$

# Recursive Linear Search

```
int linearSearch(int a[], int n, int target) {  
    // Recursive version of linear search  
    if (n <= 0) // an empty list is specified  
        return -1;  
    else {  
        if (a[n-1] == target) //test final  
position  
            return n-1;  
        else // search the rest of the list  
recursively  
            return linearSearch(a, n-1, target);  
    }  
}
```

# Binary Search Algorithm

- Can only be performed on **a sorted list** !!!
- Uses *divide and conquer* technique to search list

# Binary Search Algorithm

- Search item is compared with middle element of list
- If search item  $<$  middle element of list, search is restricted to **first half** of the list
- If search item  $>$  middle element of list, search **second half** of the list
- If search item = middle element, search is complete

# Binary Search Algorithm

- Search for (7)

0	1	2	3	4	5	6	7	8	9	10	11	12
2	4	5	7	8	9	11	14	17	20	27	31	34

0	1	2	3	4	5
2	4	5	7	8	9

3	4	5
7	8	9

3
7

Found

# Binary Search Algorithm

```
int binarySearch(int [ ] list, int listLength, int key)
{
    int first = 0, last = listLength - 1;
    int mid;
    boolean found = false;

    while (first <= last && !found) {
        mid = (first + last) / 2;
        if (list[mid] == key)
            found = true;
        else
            if(list[mid] > key)
                last = mid - 1;
            else
                first = mid + 1;
    }
    if (found)
        return mid;
    else
        return -1;
} //end binarySearch
```

# Binary Search Algorithm

- **The Best Case Complexity** - occurs when the element to search is found in first comparison □ Time complexity is 1
- **The Worst Case Complexity** - occurs, when we have to keep reducing the search space till it has only one element □ Time Complexity is  $\log N$



# Recursive Binary Search

$$f(a, f, l, target) \begin{cases} -1 & \text{if } f > l \\ (f+l)/2 & \text{if } target = (f+l)/2 \\ f(a, f, ((f+l)/2) - 1, target) & \text{if } target < (f+l)/2 \\ f(a, ((f+l)/2) + 1, l, target) & \text{if } target > (f+l)/2 \end{cases}$$

# Recursive Binary Search

```
int binarySearch(int a[], int first, int last, int target)
{
    // Preconditions: a is an array sorted in ascending order, first is the index of
    //the first element to search, last is the index of the last element to //search,
    target is the item to search for.

    if (first > last)
        return -1; // -1 indicates failure of search

    int mid = (first+last)/2;

    if (a[mid] == target)
        return mid;

    else if (target < a[mid]) // left/lower sub-array
        return binarySearch(a, first, mid-1, target);

    else // target must be > a[mid] // right/upper sub-array
        return binarySearch(a, mid+1, last, target);
}
```

# Task

- Implement function that check if an array has two sequential numbers that their sum is K using concept of recursion and binary search where the array is sorted

$F(arr, st, end, K)$  {

false	if (mid = st and sum < K) OR if (mid = end and sum > K)
True	if sum = K
$F(arr, st, mid)$	if sum > K
$F(arr, mid, end)$	if sum < K

- Ex: arr = [6,8,9,10,15] , K = 17 -> True
- Ex: arr = [5,12,20,25,30] , K = 38 -> False

# Bonus

- The same of previous task but use linked list instead of array

Thank You