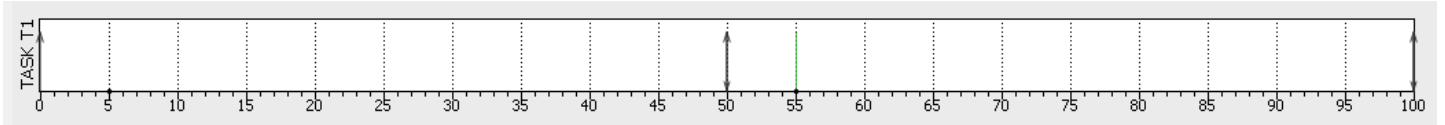


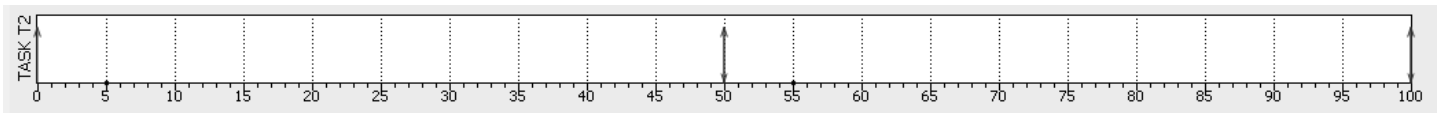
Verifying the system implementation

Tasks :

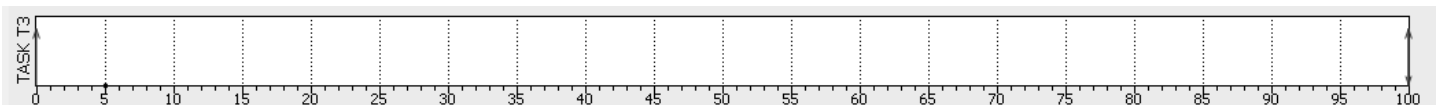
- Task 1: ""Button_1_Monitor"",
{ Periodicity: 50, Deadline: 50, execution time : $2\mu s$ }



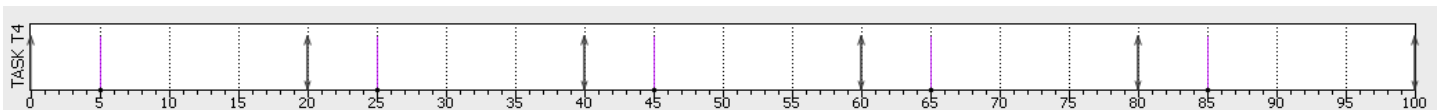
- Task 2: ""Button_2_Monitor"",
{ Periodicity: 50, Deadline: 50, execution time : $2\mu s$ }



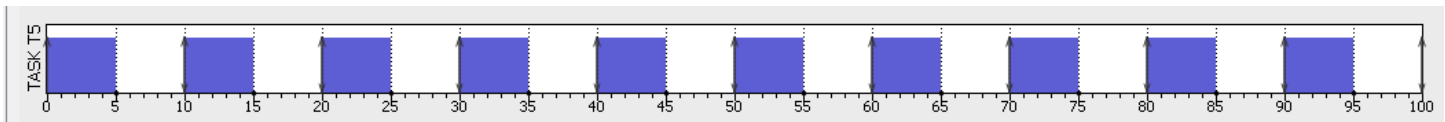
- Task 3: ""Periodic_Transmitter"",
{ Periodicity: 100, Deadline: 100, execution time : $10\mu s$ }



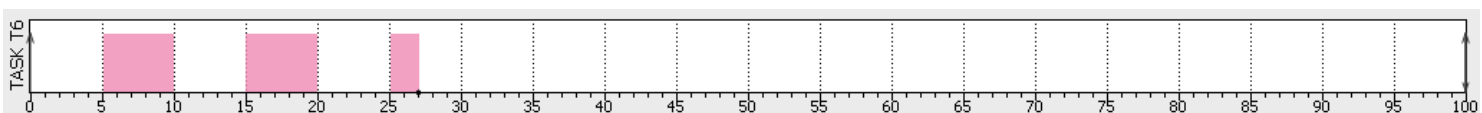
- Task 4: ""Uart_Receiver"",
{ Periodicity: 20, Deadline: 20, execution time : $4\mu s$ }



- Task 5: ""Load_1_Simulation"",
{ Periodicity: 10, Deadline: 10, execution time : $5ms$ }



- Task 6: ""Load_2_Simulation"",
{ Periodicity: 100, Deadline: 100, execution time : $12ms$ }



1. Hyper period :

$$H = LCM(P_i) = LCM(50,100,20,10) = 100$$

That:

- H is the hyper period.
- P_i is all tasks periodicities.

2. CPU load :

$$U = \frac{\sum_{i=1}^n (E_i)}{H}$$

That:

- U is CPU utilization
- H is the hyper period.
- E_i is the i th task total execution time.
- n is the number of tasks.

So,

$$U = \frac{((2*0.002)+(2*0.002)+(1*0.01)+(5*0.004)+(10*5)+(1*12))}{100} = 0.622$$

$$\therefore U\% = 62.2 \%$$

3. Schedulability :

A. Using Rate-Monotonic utilization bound method (URM) :

$$U = \sum_{i=1}^n \frac{E_i}{P_i} \leq n * (2^{\frac{1}{n}} - 1)$$

That:

- U is total utilization.
- H is the hyper period.
- E_i is the i th task execution time.
- P_i is the i th tasks periodicity.
- n is the number of tasks.

$$U = \frac{0.002}{50} + \frac{0.002}{50} + \frac{0.01}{100} + \frac{0.004}{20} + \frac{5}{10} + \frac{12}{100} = 0.62$$

$$URM = 6 * \left(\left(2^{\frac{1}{6}} \right) - 1 \right) = 0.73$$

$$U < URM$$

✓ System is guaranteed schedulable.

B. Using Time-Demand analysis :

For task i :

$$w_i(t) = e_i + \sum_{k=1}^{i-1} \left[\frac{t}{P_k} \right] e_k \quad \text{for } 0 < t \leq P_i$$

That:

- w is the worst response time.
- t is the time instance.
- E_i is the i th task execution time.
- P_i is the i th tasks periodicity.

1. For task 1:

$$w_1(10) = 0.002 + 0 = 0.002$$

$$w_1(20) = 0.002 + 0 = 0.002$$

$$w_1(30) = 0.002 + 0 = 0.002$$

$$w_1(40) = 0.002 + 0 = 0.002$$

$$w_1(50) = 0.002 + 0 = 0.002$$

✓ $w_1(50) < D_1 \rightarrow$ task 1 is schedulable.

2. For task 2:

$$w_2(10) = 0.002 + \frac{10}{50} * 0.002 = 2.4 * 10^{-3}$$

$$w_2(20) = 0.002 + \frac{20}{50} * 0.002 = 2.8 * 10^{-3}$$

$$w_2(30) = 0.002 + \frac{30}{50} * 0.002 = 3.2 * 10^{-3}$$

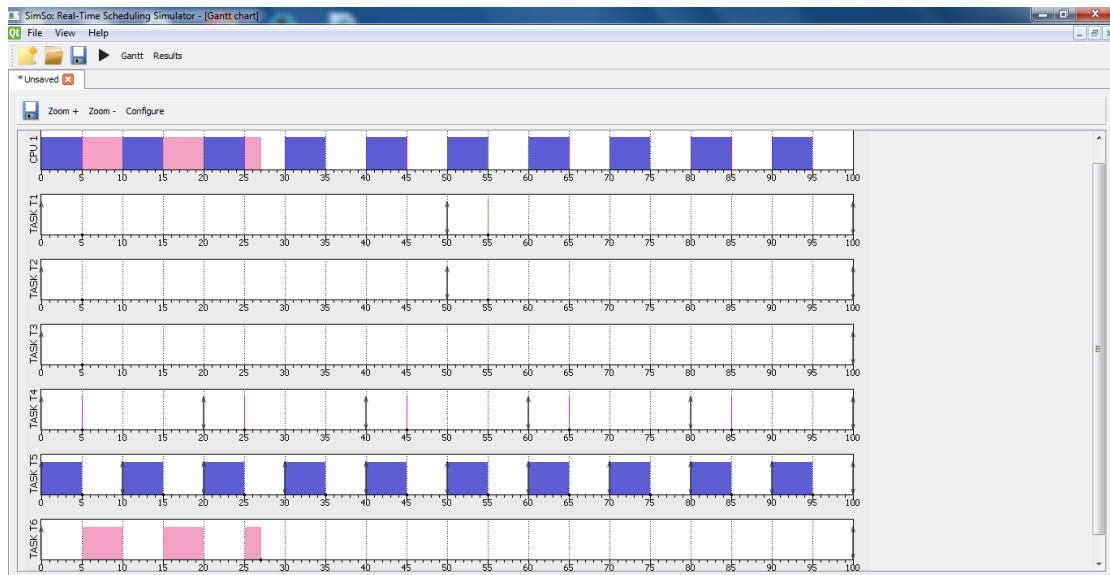
$$w_2(40) = 0.002 + \frac{40}{50} * 0.002 = 3.6 * 10^{-3}$$

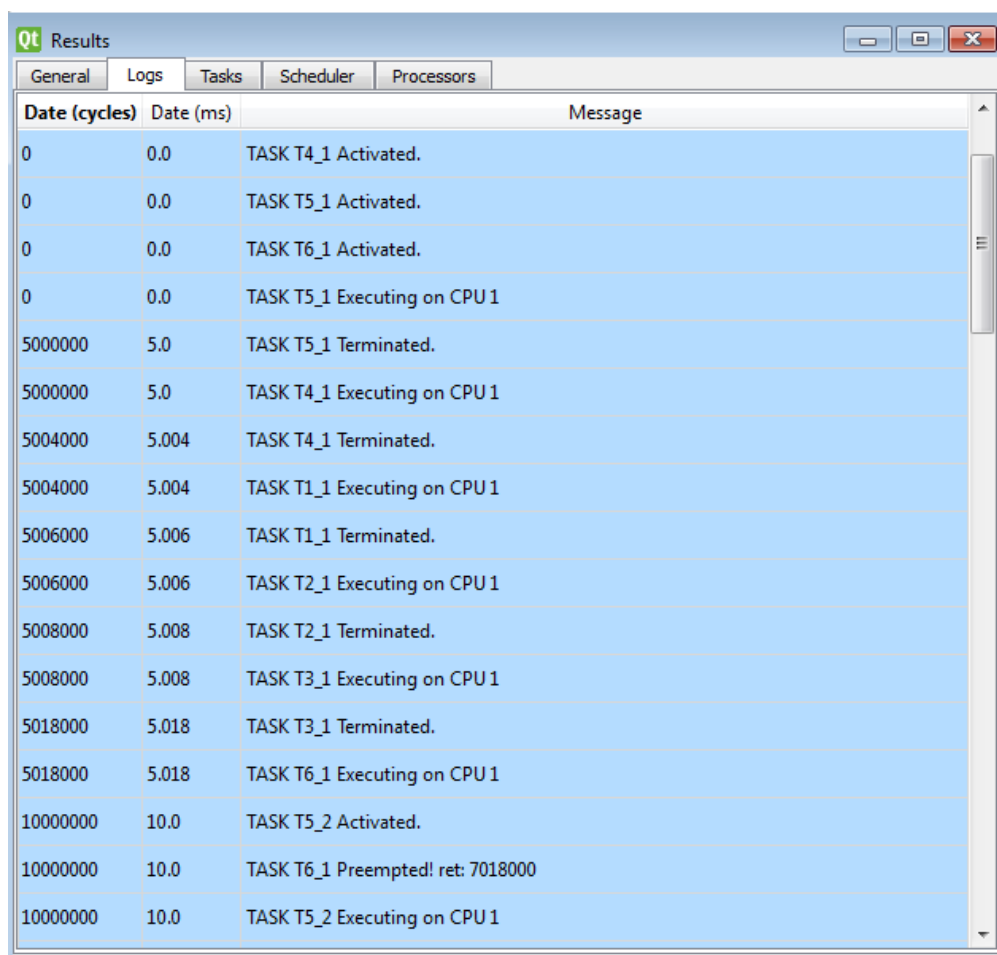
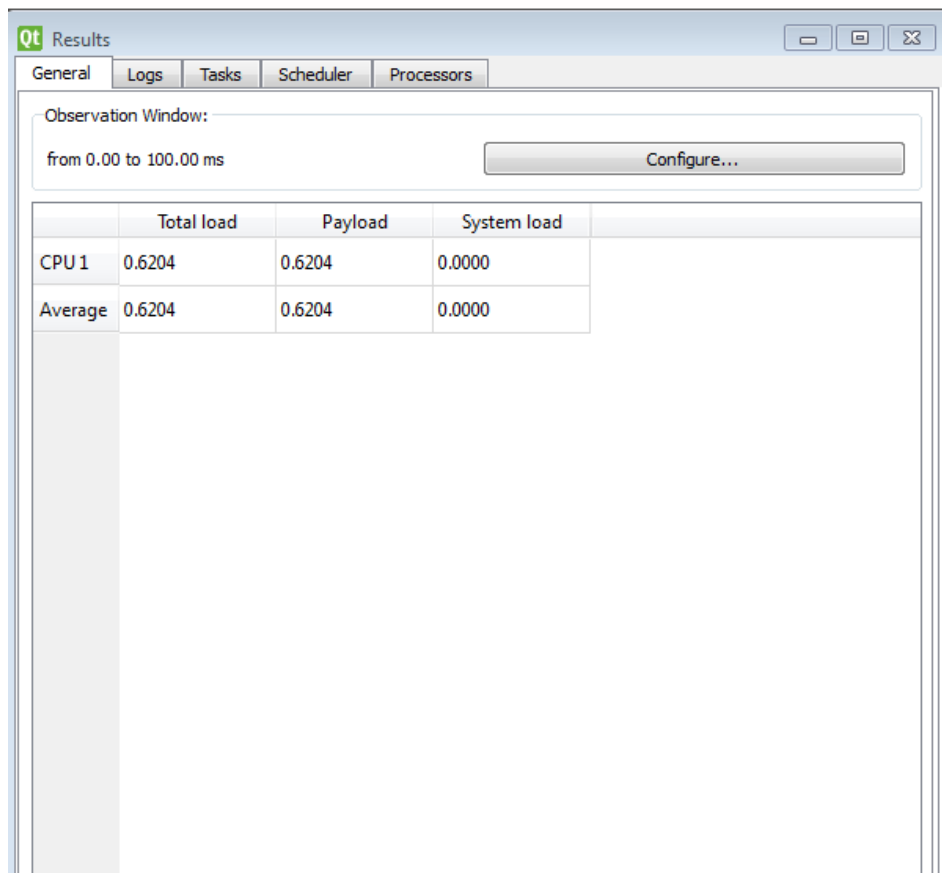
$$w_2(50) = 0.002 + \frac{50}{50} * 0.002 = 4 * 10^{-3}$$

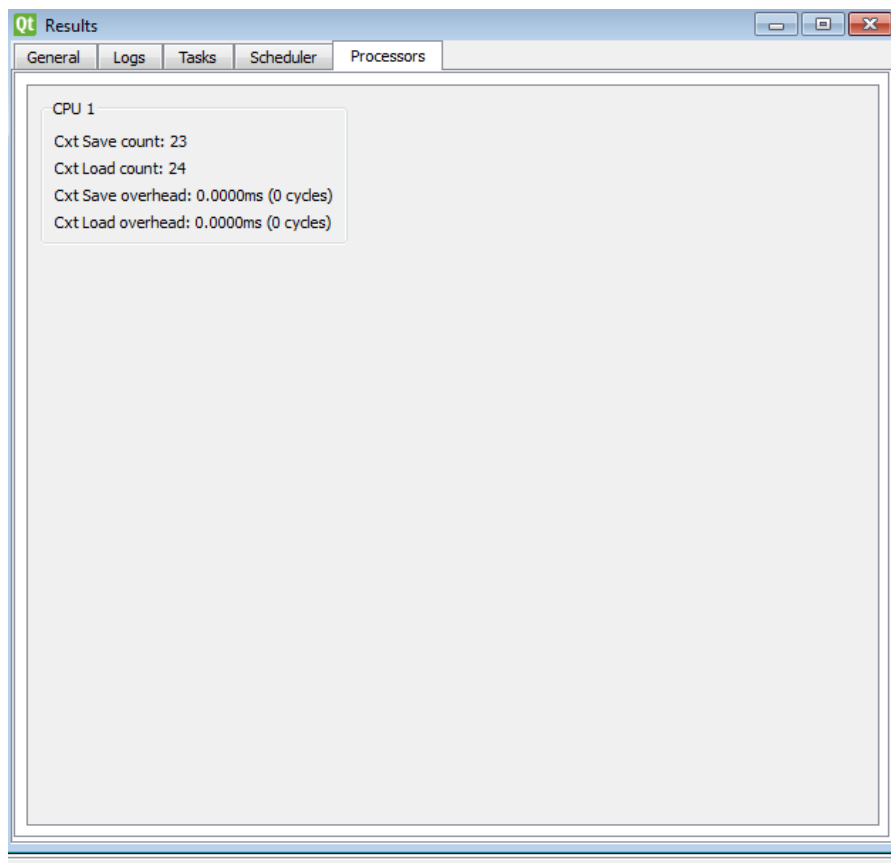
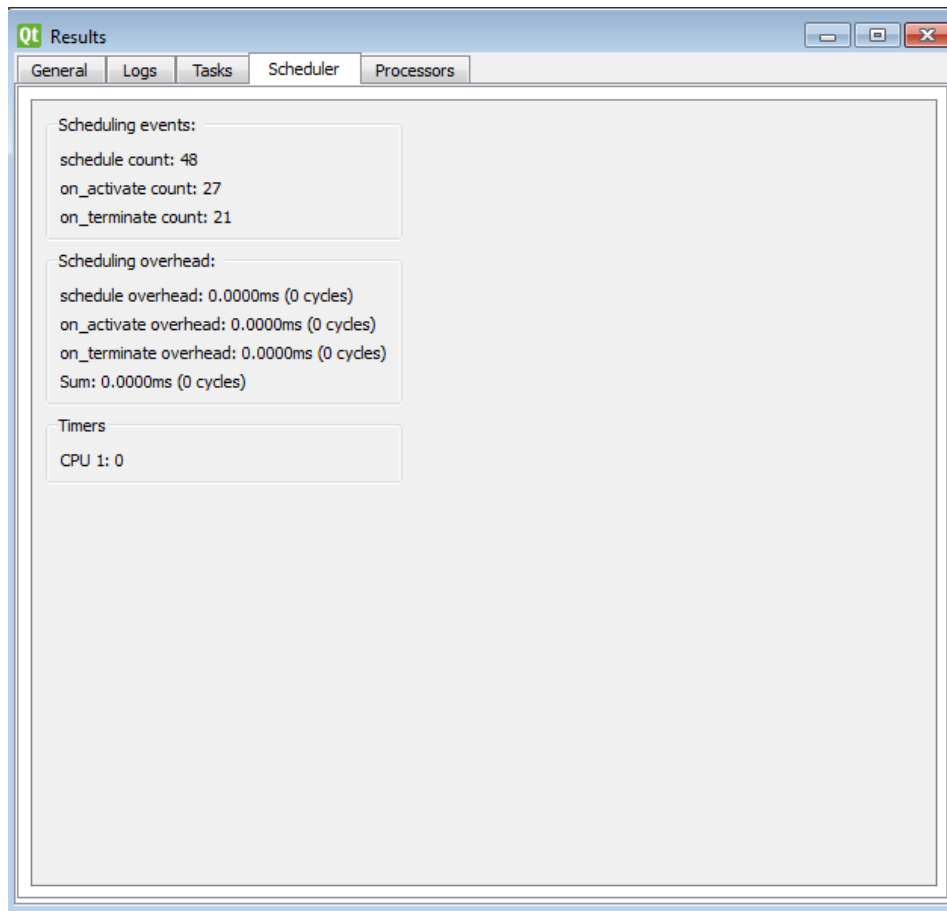
✓ $w_2(50) < D_2 \rightarrow$ task 2 is schedulable.

✓ And so on for rest of tasks...

4. Simso simulation :







5. logic analyzer :

