Agenda

- Hyperledger
- Hyperledger Fabric
- Permissionless Vs Permissioned blockchains
- Fabric component architecture
- Fabric ordering service RAFT
- Distributed System Communicate
- BFT
- PBFT
- DLS
- BDLS
- Hyperledger Fabric Goal



Hyperledger Project

Is an open-source collaborative effort created to advance cross-industry Blockchain technologies.

It is a project under the LINUX Foundation, Hyperledger is not a single project but a collection of projects under the Hyperledger umbrella







Distributed Ledgers



Java-based Ethereum client



Permissionable smart contract machine (EVM)



Enterprise-grade DLT with privacy support



Decentralized identity



Mobile application focus



Permissioned & permissionless support; EVM transaction family

Libraries





















Domain-Specific











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Tools

















Benefits of Hyperledger Fabric



Permissioned network

Establish decentralized trust in a network of known participants rather than an open network of anonymous participants.



Confidential transactions

Expose only the data you want to share to the parties you want to share it with.



Pluggable architecture

Tailor the blockchain to industry needs with a pluggable architecture rather than a one-size-fits-all approach.



Easy to get started

Program smart contracts in the languages your team works in today, instead of learning custom languages and architectures.





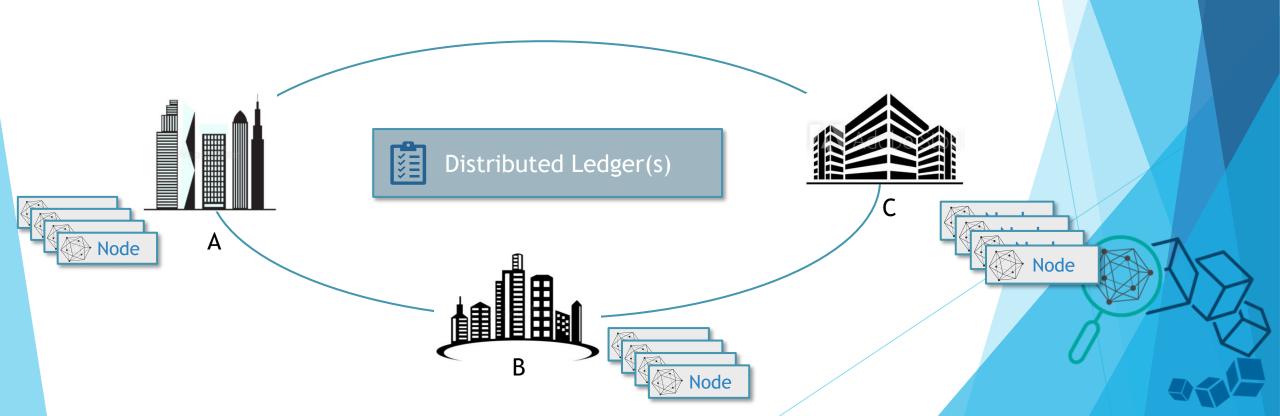
A <u>Distributed Ledger Technology</u> (DLT) framework for building <u>Business</u> Blockchain Applications.



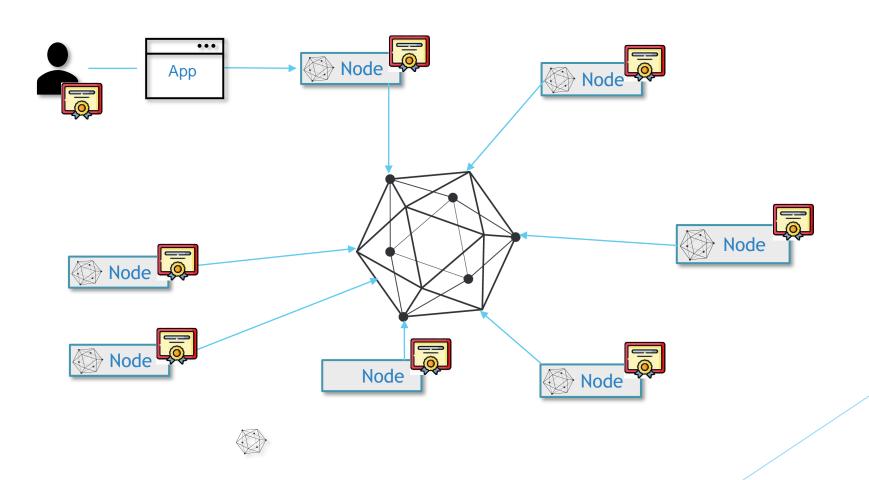




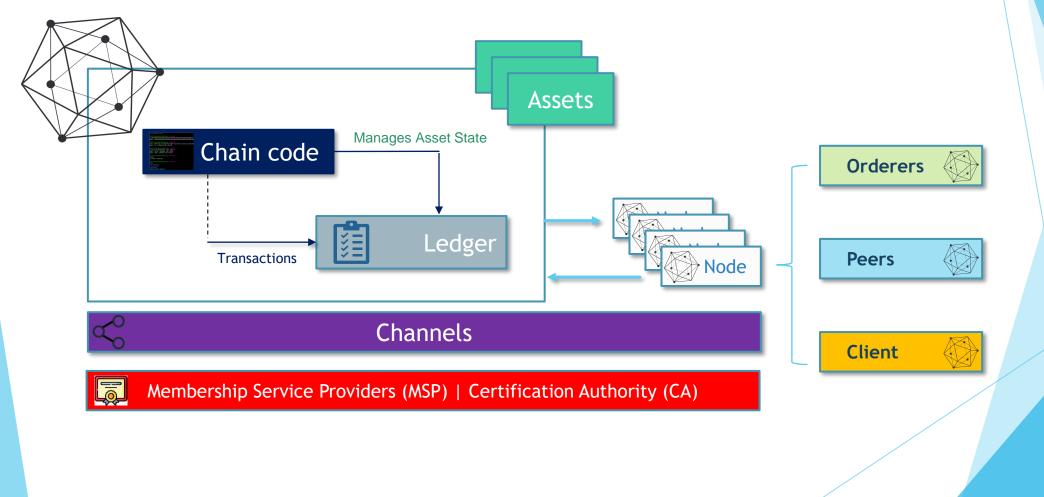
Members = Legally separate entities



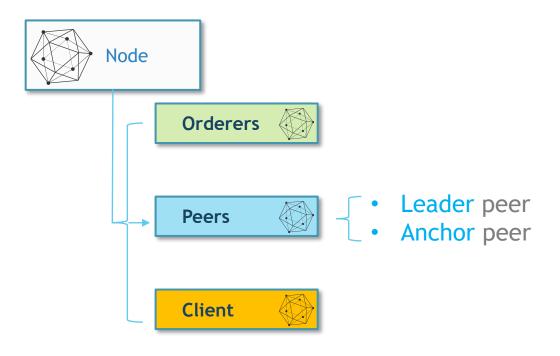












Communication channel of Fabric



Only nodes known outside the organization

Use SDK for the user's endpoints





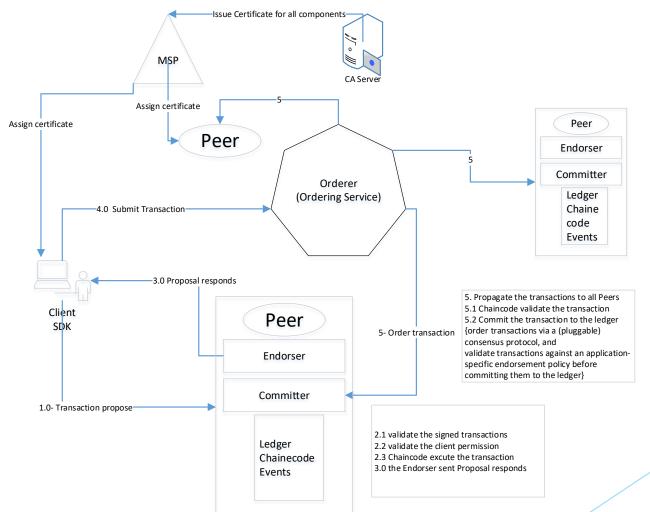


Communication channel of Fabric

- Ordering Service
- Responsible for consistent ledger state across the network
 - Consensus mechanism
 - Ensures order of transactions
- Creates the blocks & guarantees atomic delivery











Implemented with Message Oriented Middleware



The Raft Consensus Algorithm

- Diego Ongaro and John Ousterhout, Stanford University (2014)
 "In Search of an Understandable Consensus Algorithm"
- Managing a replicated log.
- Crash fault-tolerance (CFT)
- Quorum = 1/2N
- It produces a result equivalent to (multi-)Paxos
- Raft is easier for students to learn than Paxos.
 - Leslie Lamport "The part-time parliament" (1989)
 - "Paxos made simple" (2001)
- RAFT & Paxos are non-Byzantine model







Implemented with Message Oriented Middleware



The Raft Consensus Algorithm

- Diego Ongaro and John Ousterhout, Stanford University (2014)
 "In Search of an Understandable Consensus Algorithm"
- Demo: http://thesecretlivesofdata.com/raft/
- Demo the node rest on Raft website:

https://raft.github.io/



"The Byzantine generals problem" Leslie Lamport, Robert Shostak, and Marshall Pease.(1982)

- System with x Byzantine nodes
 3x + 1 total nodes in order to reach consensus
- Potential traitor generals = Byzantine Nodes
- There is no solution in the present of 1/3 or grater percentage potential traitor generals.

The algorithms demonstrated in this paper are only designed to work in a synchronous environment.

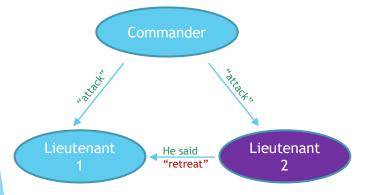


Fig.1 Lieutenant 2 a traitor

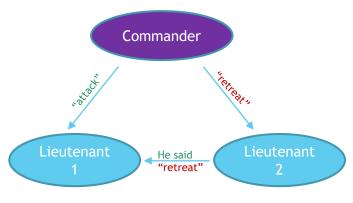


Fig.2 The Commander a traitor





- By "message passing" between one or more other nodes
- Messaging protocol, HTTP, RPC, or a custom protocol.

Synchronous _____ messages will be delivered within some fixed time

Asynchronous network may delay messages infinitely

(**DLS** and **PBFT**) That brought us closer than ever before to breaking the Byzantine + asynchronous barrier.



"Practical byzantine fault tolerance." By Miguel Castro, and Barbara Liskov. (1999)

- Handle f Byzantine faults in a system with 3f + 1 nodes
- Quorum over 2/3 voters.
- Main PBFT algorithm consists of three phases: pre-prepare, prepare, and commit.

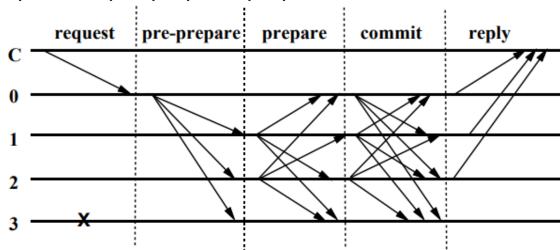


Figure 1: Normal Case Operation





"Consensus in the Presence of Partial Synchrony" by Dwork, Lynch, and Stockmeyer (1988)

- The first known asynchronous Byzantine Consensus solution
- Partial synchrony lies somewhere between Synchronous & Asynchronous.

Two versions of the partial synchrony assumption

- 1. Assume that fixed bounds exist for how long messages take to get delivered. But they are not known a priori.
- 2. Assume the upper bounds for message delivery are known, but they're only guaranteed to hold starting at some unknown time (also called "Global Stabilization Time," GST).

A series of rounds are divided into "trying" and "lock-release" phases.

Liveness is the property of the system continuing to work in case of failures. **Safety** is the agreement of the network on a single state.





"Byzantine Fault Tolerance in Partial Synchronous Networks." Wang, Yongge (2020).

- BDLS consensus based on DLS protocol algorithm.
- Able to achieve consensus with both reduced round complexity and reduced communication complexity.

PBFT
Mesh communication network

BDLS Star networks

HotStuff using threshold cryptography. ———— Facebook's LibraBFT protocol

 Best existing linear communication/ authenticator complexity protocols require at least 7 steps to achieve agreement.

VS

 BDLS participants could reach agreement in 4 steps with linear communication/authenticator complexity to achieve agreement.





Type I $\Delta < \infty$ is unknown.



Type II $\Delta < \infty$ holds eventually, participant knows the value of Δ But this only holds after an unknown time slot Global Stabilization Time (GST).



BDLS is proved to be secure in Type II partial synchronous networks

Attacks against several widely deployed BFT protocols: SUCH AS:

- PBFT
- Tendermint BFT
- Casper FFG

Participants would reach a deadlock before GST and the deadlock could not be removed after GST.





BFT protocols in partial synchronous networks

Steps	PBFT	Tendermint BFT	HotStuff BFT	BDLS	
1	(((a))	ത്ര	(%)	ক্তি))	
2	(C)	(C)	ক্ত ্য)	(M)	
3	(C)	(C)	(%)	ව්))	
4			ිවා)	(((v))	
5			(%)		
6			ক্তি))		
7			(%)		
message complexity	$2n^2 + n$	$2n^2 + n$	7n	4n	
authenticator complexity	O(n2)	O(n)	O(n)	O(n)	

(ஒ) : Leader broadcasts

(9)) : All participants send messages to the leader

: All participants broadcast





Hyperledger Goals



Create enterprise grade, open source, distributed ledger frameworks and code bases to support business transactions



Provide neutral, open, and community-driven infrastructure supported by technical and business governance



Build technical communities to develop blockchain and shared ledger POCs, use cases, field trails and deployments



Educate the public about the market opportunity for blockchain technology

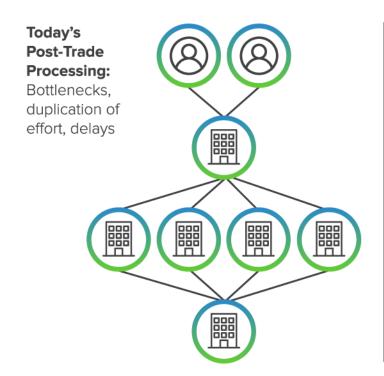


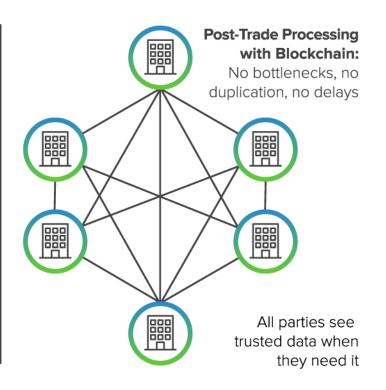
Promote our community of communities taking a toolkit approach with many platforms and frameworks





Financial Services













1. Medical school issues tamperproof credential on blockchain.

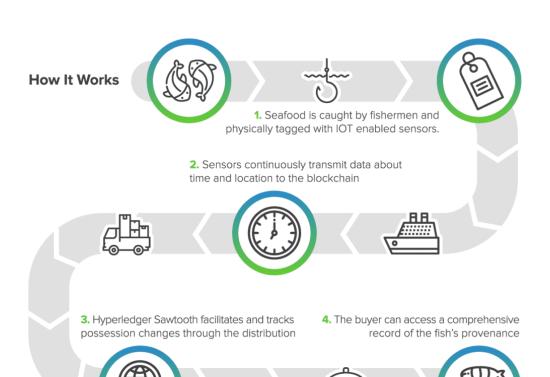
2. Doctor places further credentials on blockchain.

3. Hospital validates all credentials on blockchain, saving many steps and delays.













Thanks

Q/A?



