

# SYMBOLIC CONTROL

## Of a non linear mobile robot



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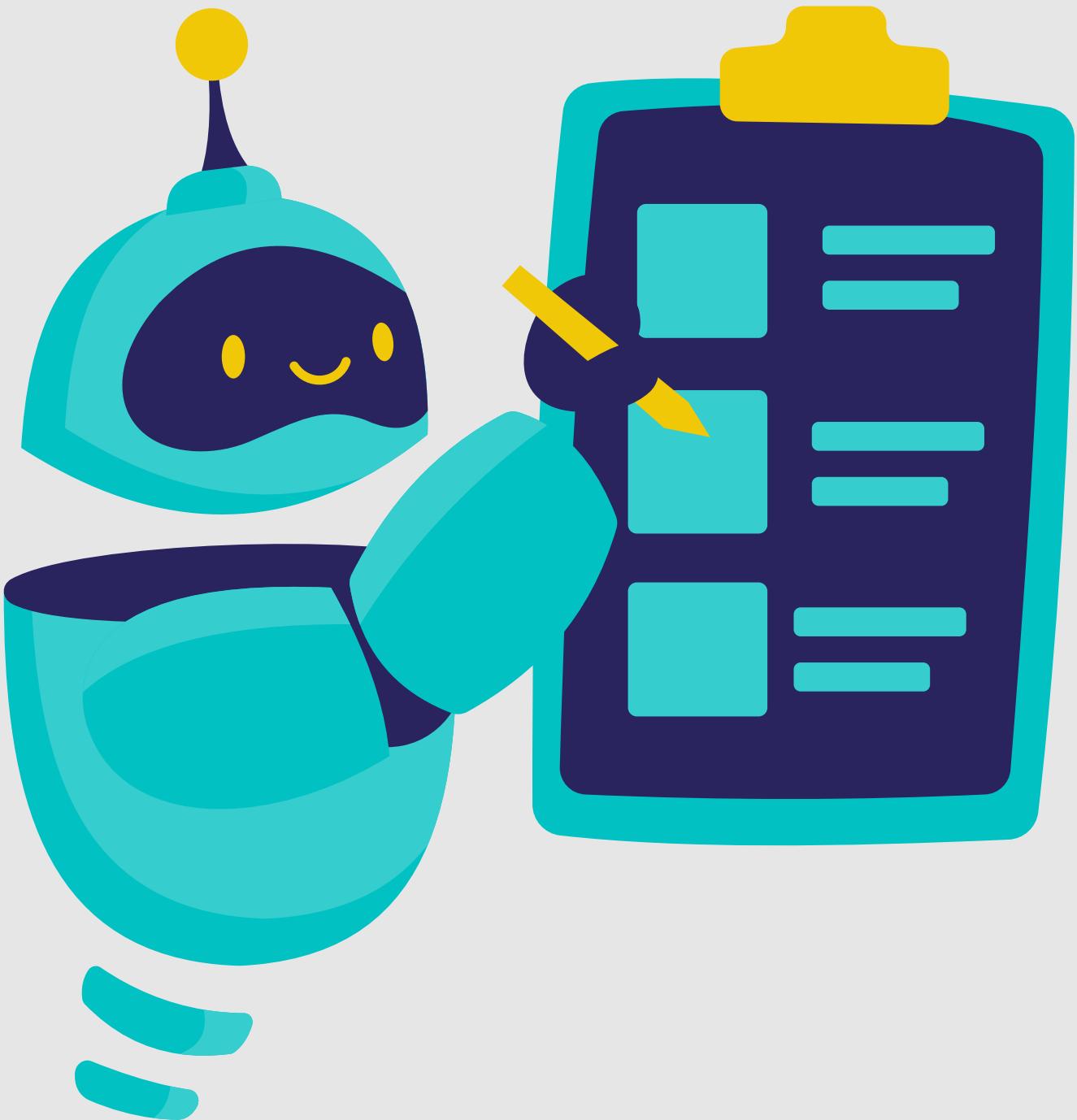


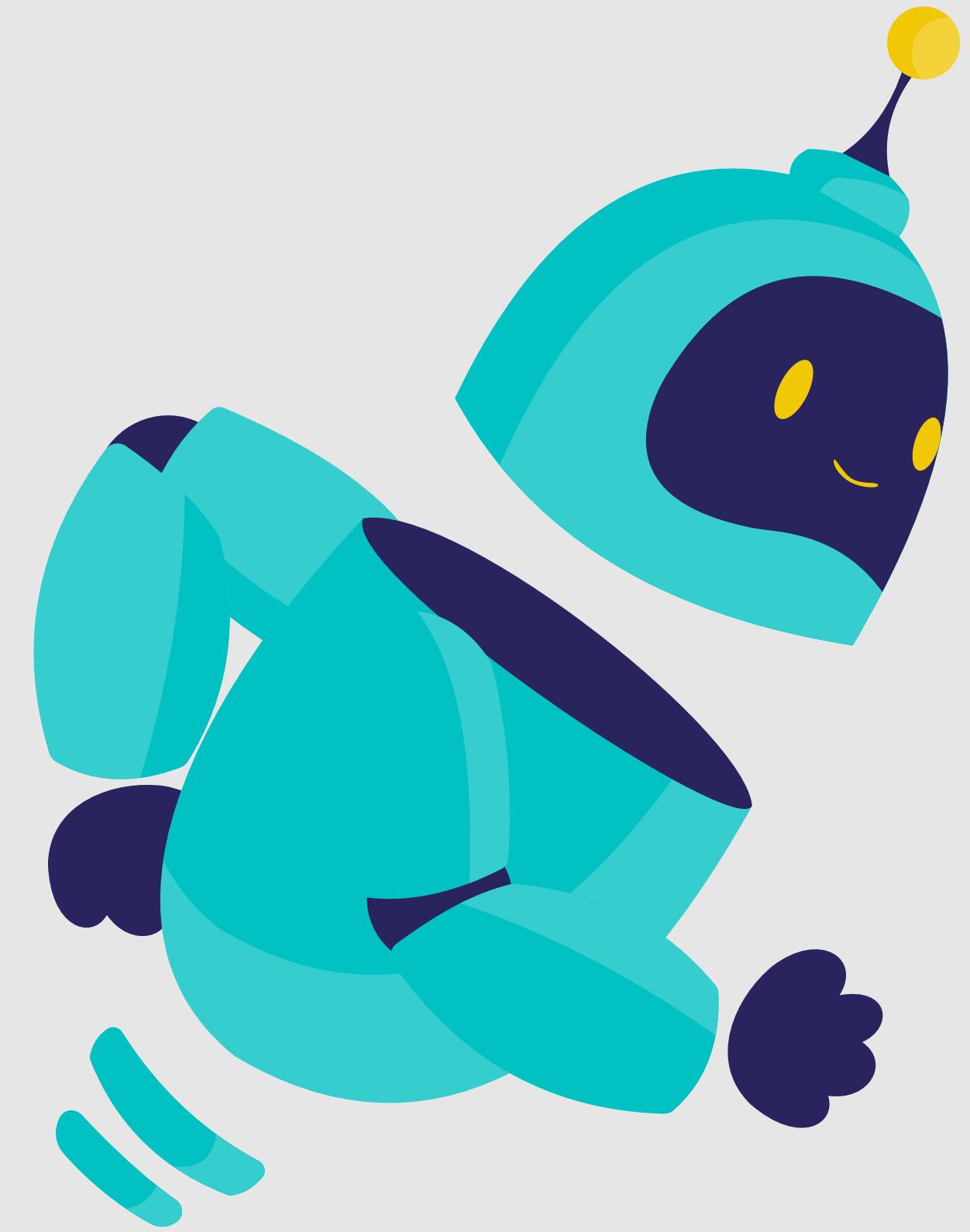
SALIM QADDA



# INTRODUCTION

**FINITE ABSTRACTION CONSTRUCTION  
CONTROLLER SYNTHESIS FOR MULTIPLE  
OBJECTIVES  
VALIDATION THROUGH EXTENSIVE  
NUMERICAL SIMULATIONS**

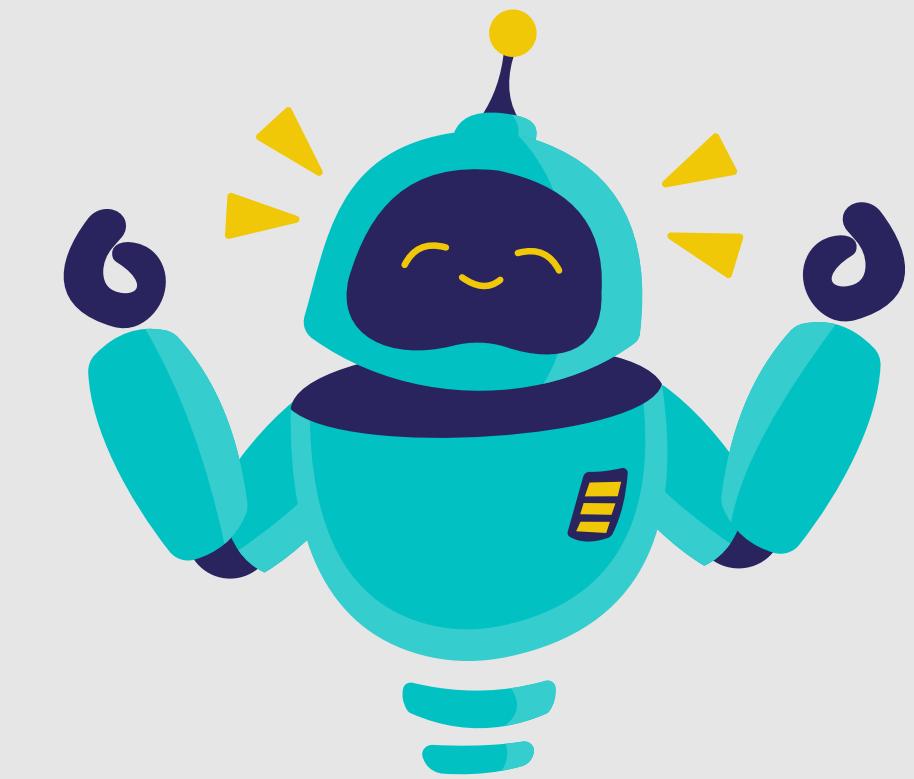




# 2D

## WHY DO WE USE 2D INSTEAD OF 3D?

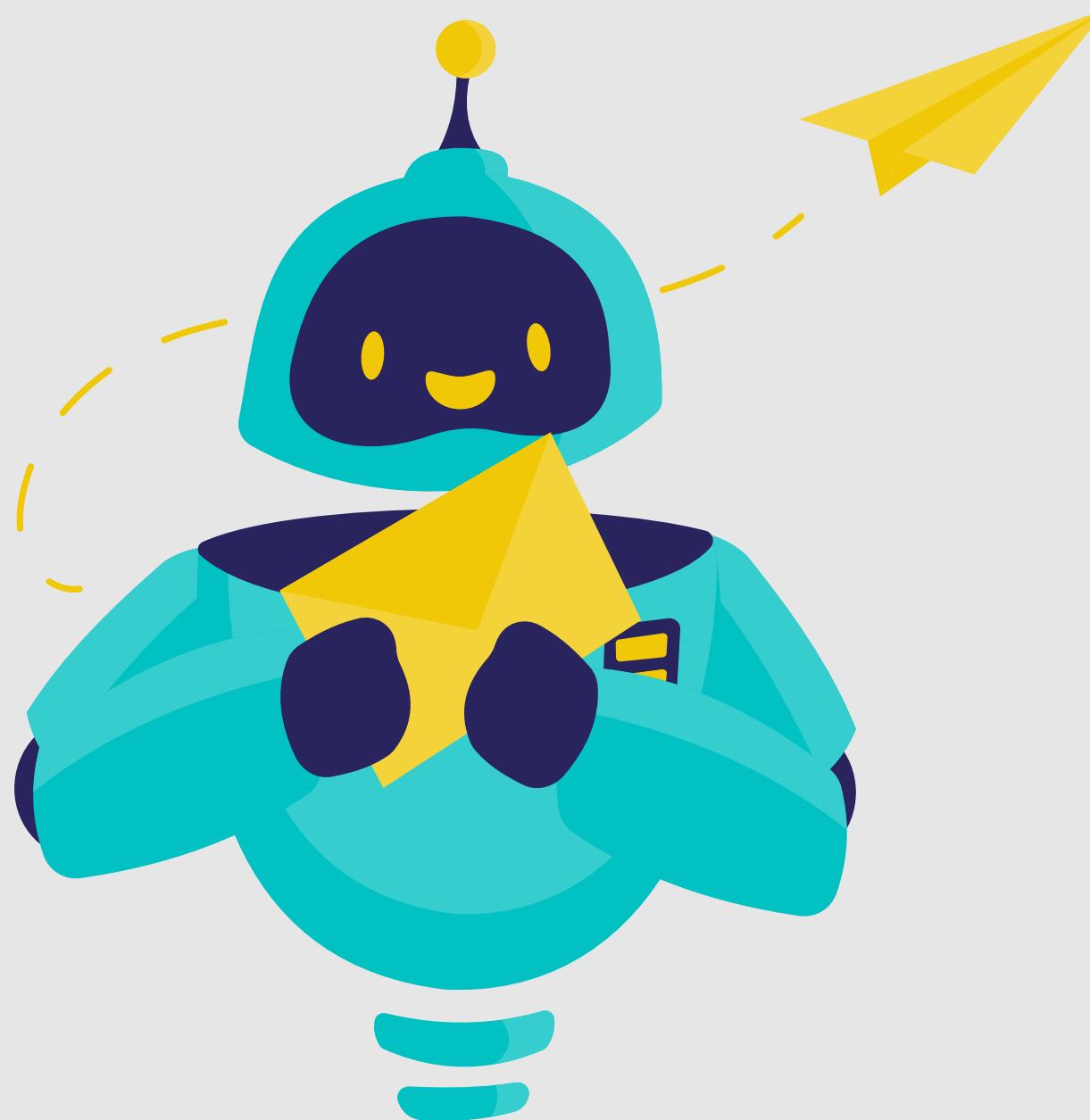
$$\begin{cases} x_1(t+1) = x_1(t) + \tau v_x(t) + \tau w_1(t) \\ x_2(t+1) = x_2(t) + \tau v_y(t) + \tau w_2(t) \end{cases}$$



## COMPLEXITY AND DEBUGGING

# 2D

## OUR MAIN FUNCTION



$$g(\xi, \sigma) = \{\xi^+ \in \Xi \mid Y_{\xi, \sigma} \cap \text{cl}(X_{\xi^+}) \neq \emptyset\}$$

THE ISSUE? COMPUTING ALL STATES TO  
CHECK IF THE CONDITION IS VALID TAKES  
TOO MUCH TIME

# 2D

## INTERVAL CONCRETIZATION

$$(i_{\min}, j_{\min}) = q(Y^{\min}),$$

$$(i_{\max}, j_{\max}) = q(Y^{\max}),$$

$$g(\xi, \sigma) = \{(i, j) \mid i_{\min} \leq i \leq i_{\max}, j_{\min} \leq j \leq j_{\max}\}.$$

# 2D

## PRE STRUGGLE: TAKES TOO LONG

```
def Pre_naive(R):
    result = set()
    for element in all_states:
        for k in range(1, Nvx * Nvy + 1):
            if g(element, k).issubset(R):
                result.add(element)
                break
    return result
```

# 2D SOLUTION: USE G-1

$$Pre(R) = \{ \xi \mid \exists \sigma \in \Sigma, g(\xi, \sigma) \subseteq R \}.$$

Let  $\xi \in Pre(R)$ . Then  $\exists \sigma \in \Sigma$  such that  $g(\xi, \sigma) \subseteq R$ . Take  $u = p(\sigma)$ .

Let  $\xi' \in g(\xi, \sigma) \subseteq R$  and let  $y \in q^{-1}(\xi')$ .

So  $\exists x \in X, \exists w \in W$  such that

$$y = f(x, u, w) = x + Tu + Tw,$$

hence

$$x = y - Tu - Tw = f(y, -u, -w).$$

We have

$$W \text{ is symmetric around } 0 \implies f(y, -u, -w) \in [f(y, -u, \underline{w}), f(y, -u, \overline{w})]$$

Take

$$g^{-1}(\xi', \sigma) = \{ \xi'' \mid Y_{\xi', \sigma}^{-1} \cap \text{cl}(X_{\xi''}) \neq \emptyset \},$$

with

$$Y_{\xi', \sigma}^{-1} = [f(y, -u, \underline{w}), f(y, -u, \overline{w})].$$

# 2D SOLUTION: USE G-1

Since  $x \in Y_{\xi',\sigma}^{-1}$ , we obtain

$$\xi \in g^{-1}(\xi', \sigma).$$

Thus,

$$\xi \in Pre(R) \implies \exists \xi' \in R, \exists \sigma \in \Sigma : \xi \in g^{-1}(\xi', \sigma).$$

Now we simply iterate through  $R$ : take  $\xi \in R$  and

$$\sigma \in \Sigma / \xi' \in g^{-1}(\xi, \sigma)$$

If

$$g(\xi', \sigma) \subseteq R,$$

then

$$\xi' \in Pre(R).$$

**Note.** The mapping  $g$  is not assumed to be bijective. The notation  $g^{-1}$  used above does *not* denote an inverse function;

*Proof written by*

**Salim Qadda**

# 2D REACHABILITY

COMPUTES PRE FOR  
THE SAME STATE  
DURING DIFFERENT  
ITERATIONS:  
REDUNDANT

**Entrées :**  $Q_a \subseteq \Sigma$

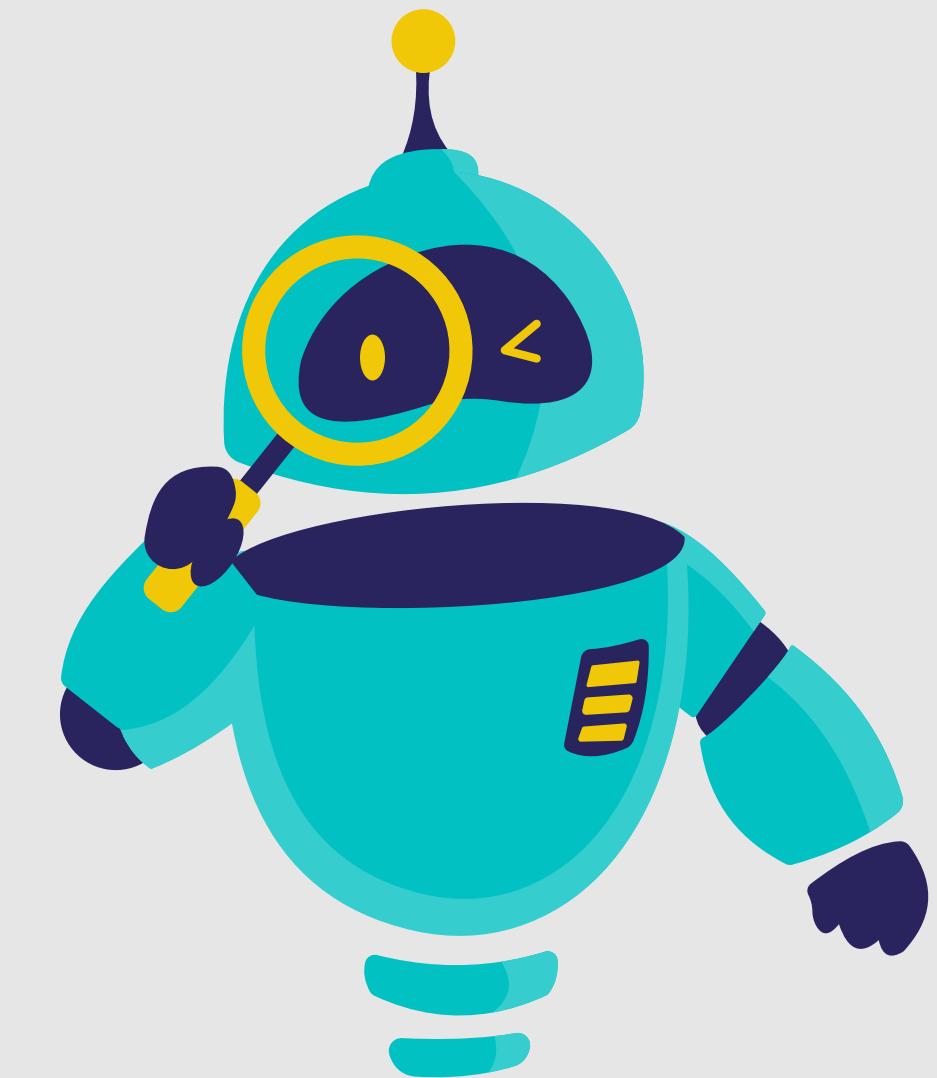
$$R_0 = Q_a$$

**répéter**

$$\quad | \quad R_{k+1} = Q_a \cup \text{Pre}(R_k)$$

**jusqu'à**  $R_{k+1} = R_k$  ;

**Sorties :**  $R^* = R_k$



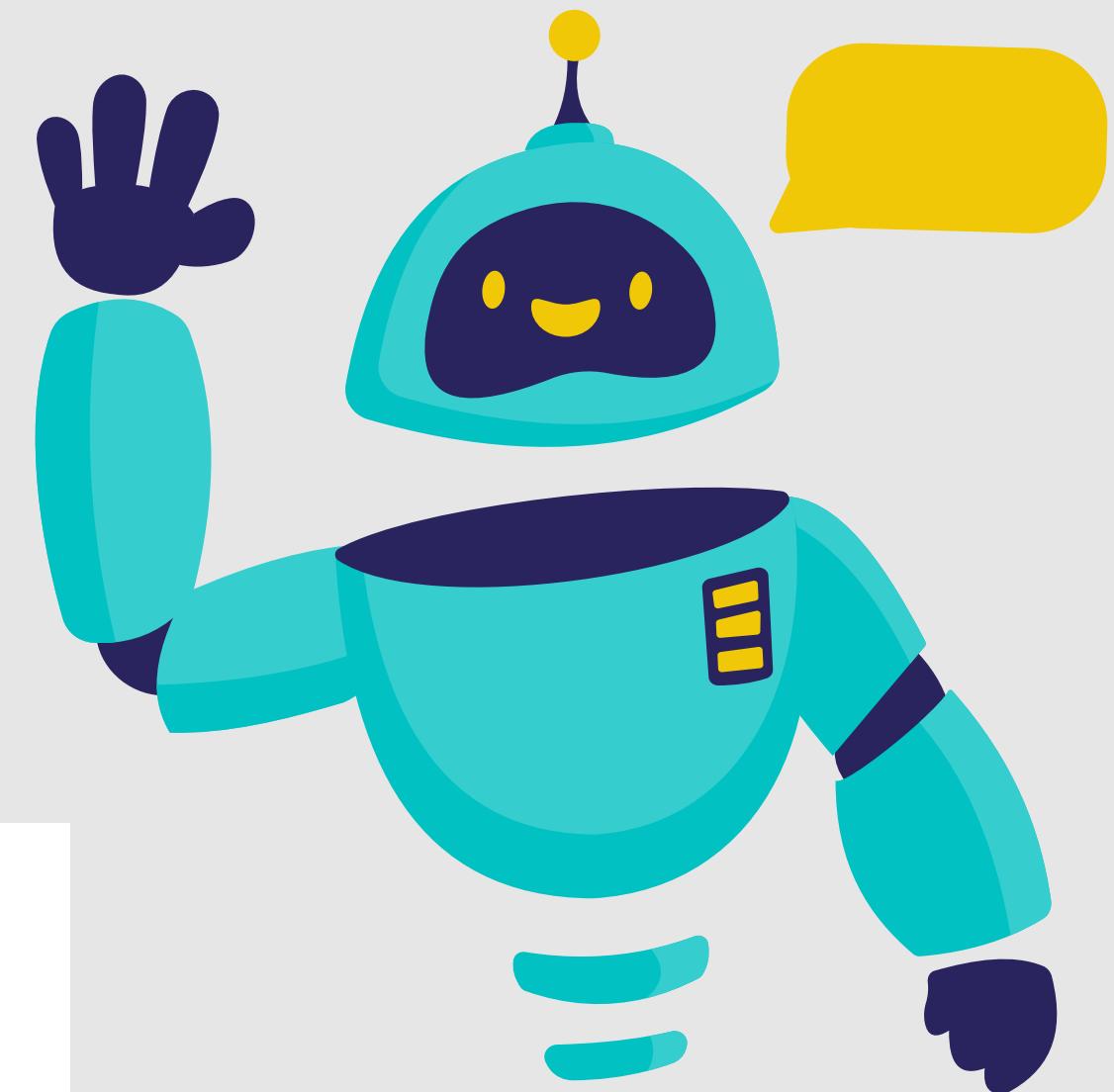
# 2D REACHABILITY

WE CAN USE A QUEUE

# 2D PRODUCT SYSTEM

SAME LOGIC, BUT WE  
DEFINE OUR STATE:

$$(\xi, q) \in E \times Q_{\text{spec}}$$



# 2D

## REACHABILITY – PRODUCT SYSTEM

KEEP THE SAME LOGIC OF REACHABILITY WE USED BEFORE, BUT  
THIS TIME WE USE:

$$Q_r = \{(\xi, q) \in E \times Q_{acc}\}$$



# 2D CONTROLLER

## WE DEFINE OUR CONTROLLER:

$$H(\xi) = \{ \sigma \in \{1, \dots, N_u\} \mid g(\xi, \sigma) \subseteq R \},$$

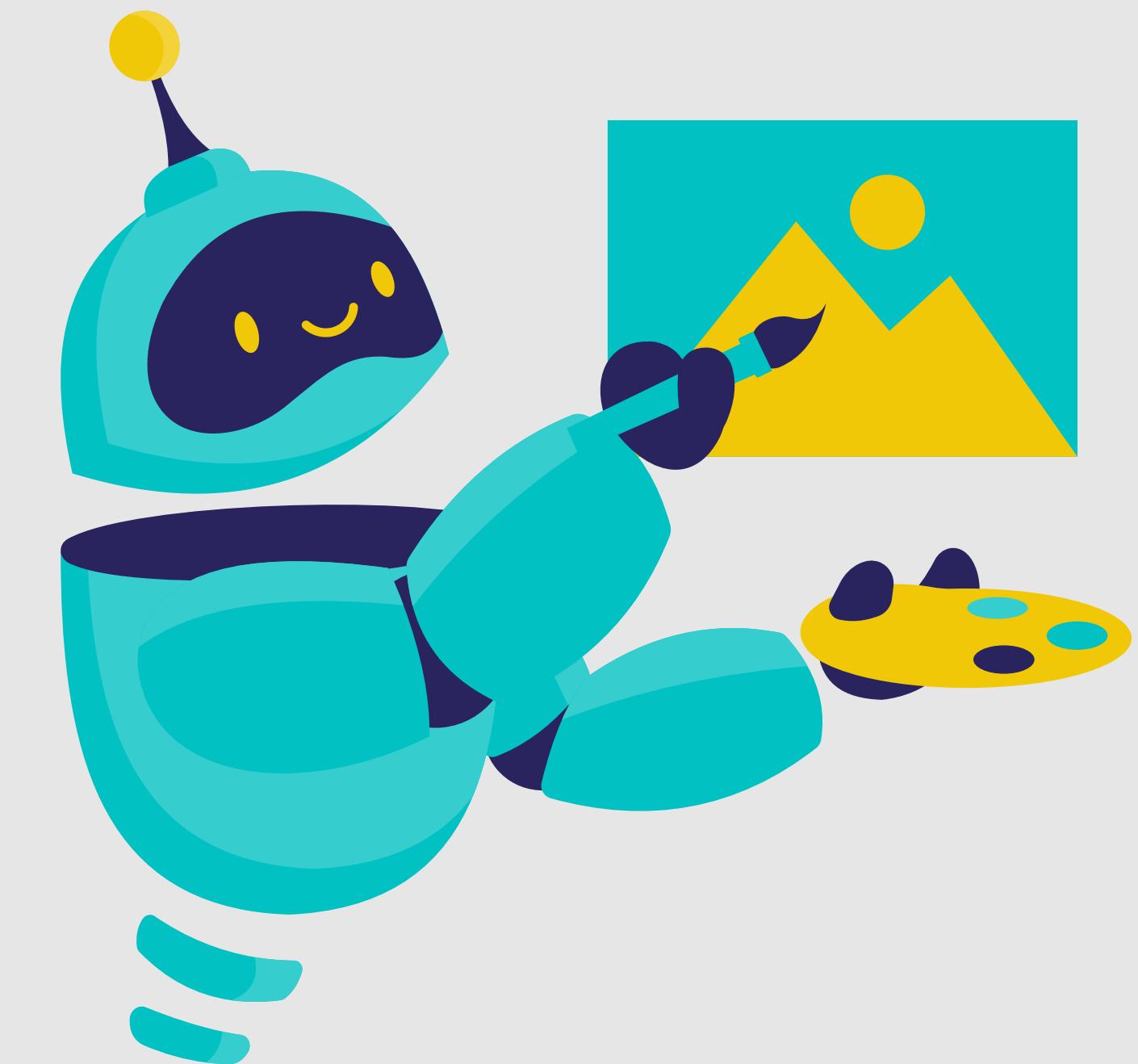
with  $R = \text{Reachability}(Q_r)$



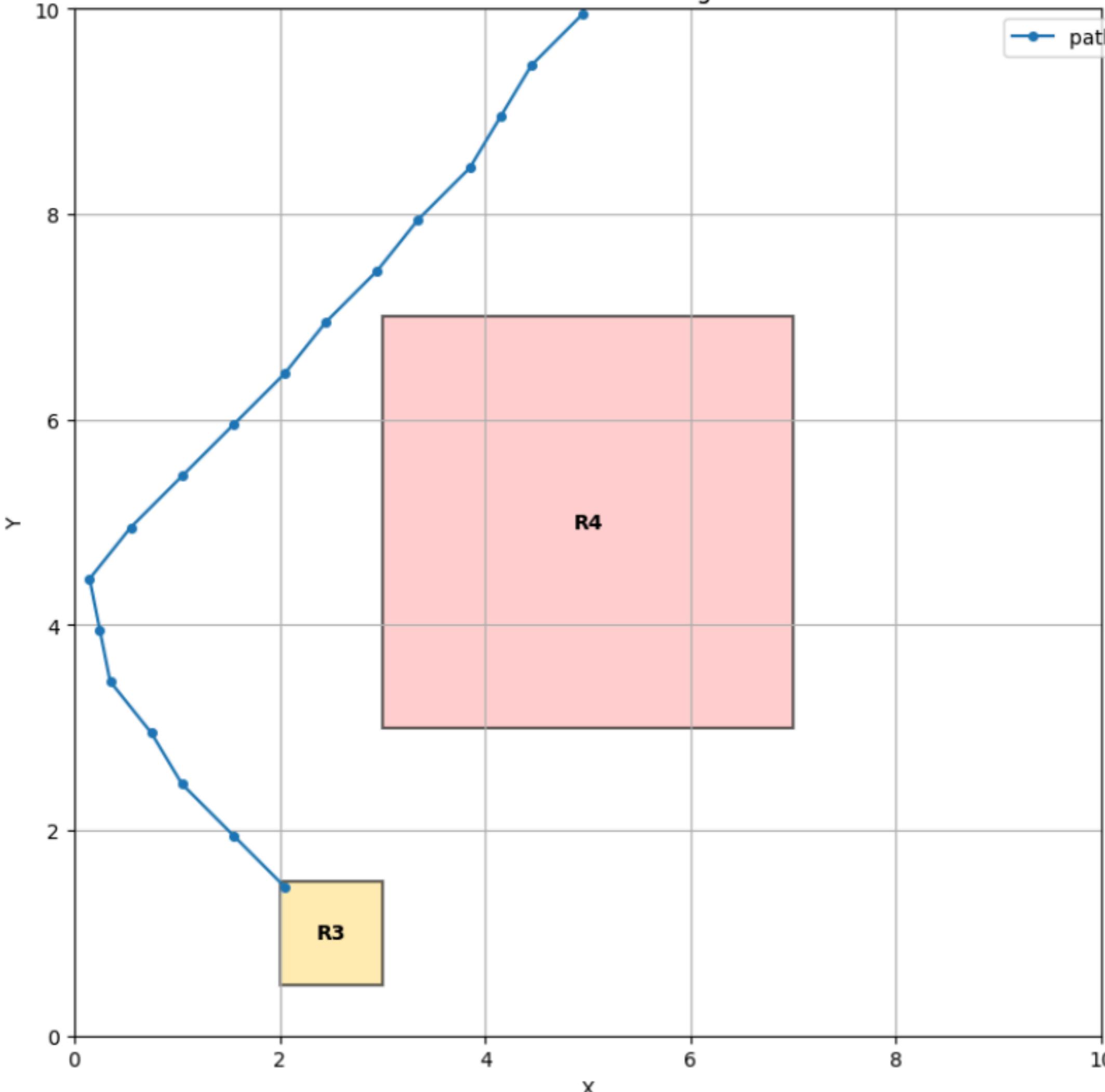
2D

# TIME TO VISUALIZE

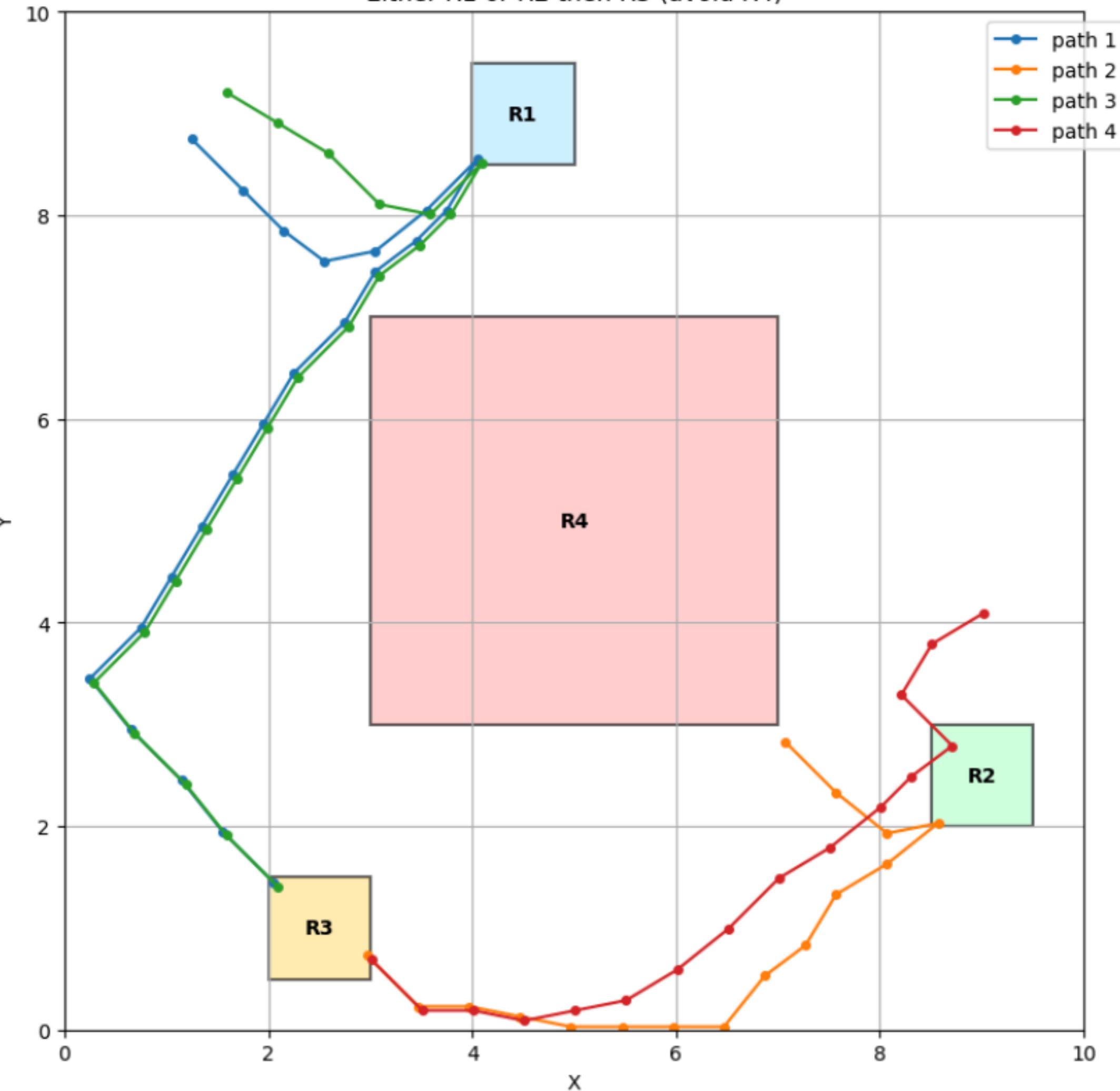
WE APPLY BFS ON  
OUR CONTROLLER  
FOR CLEAN  
VISUALIZATION



Go to R3 while avoiding R4

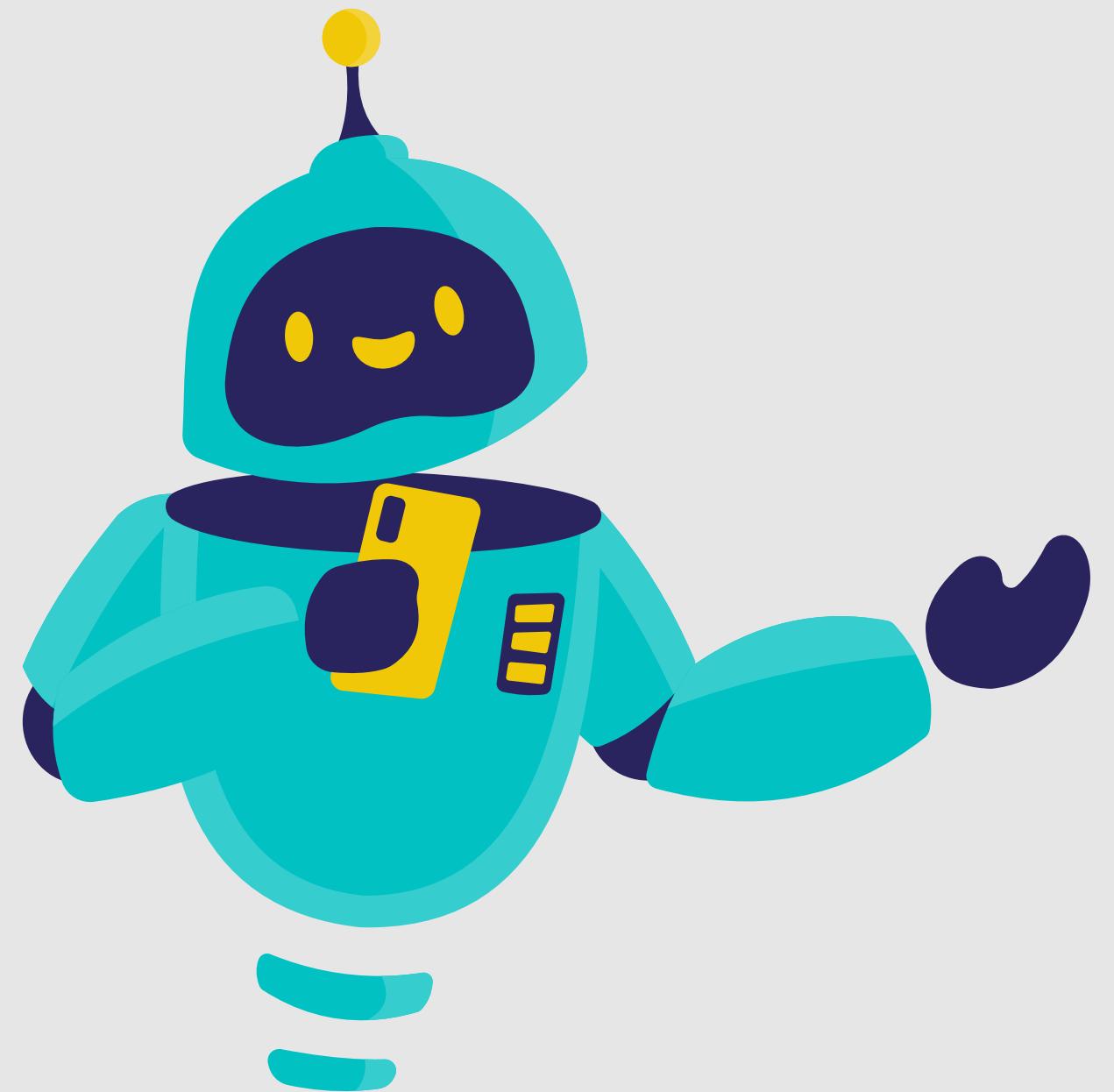


Either R1 or R2 then R3 (avoid R4)

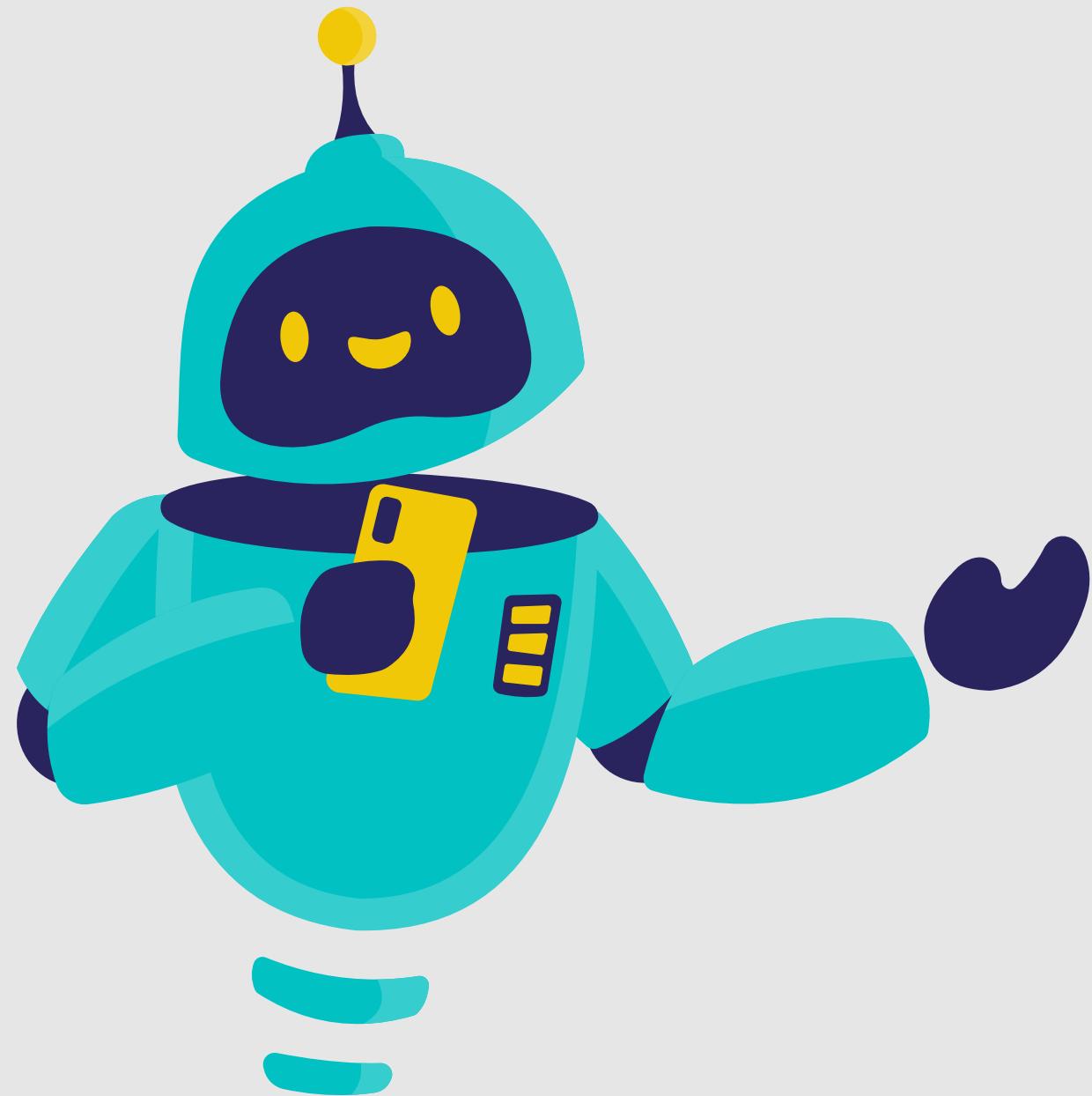
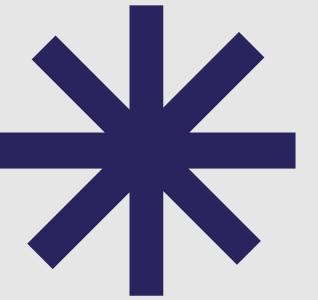




# 3D CONTROLLER \*



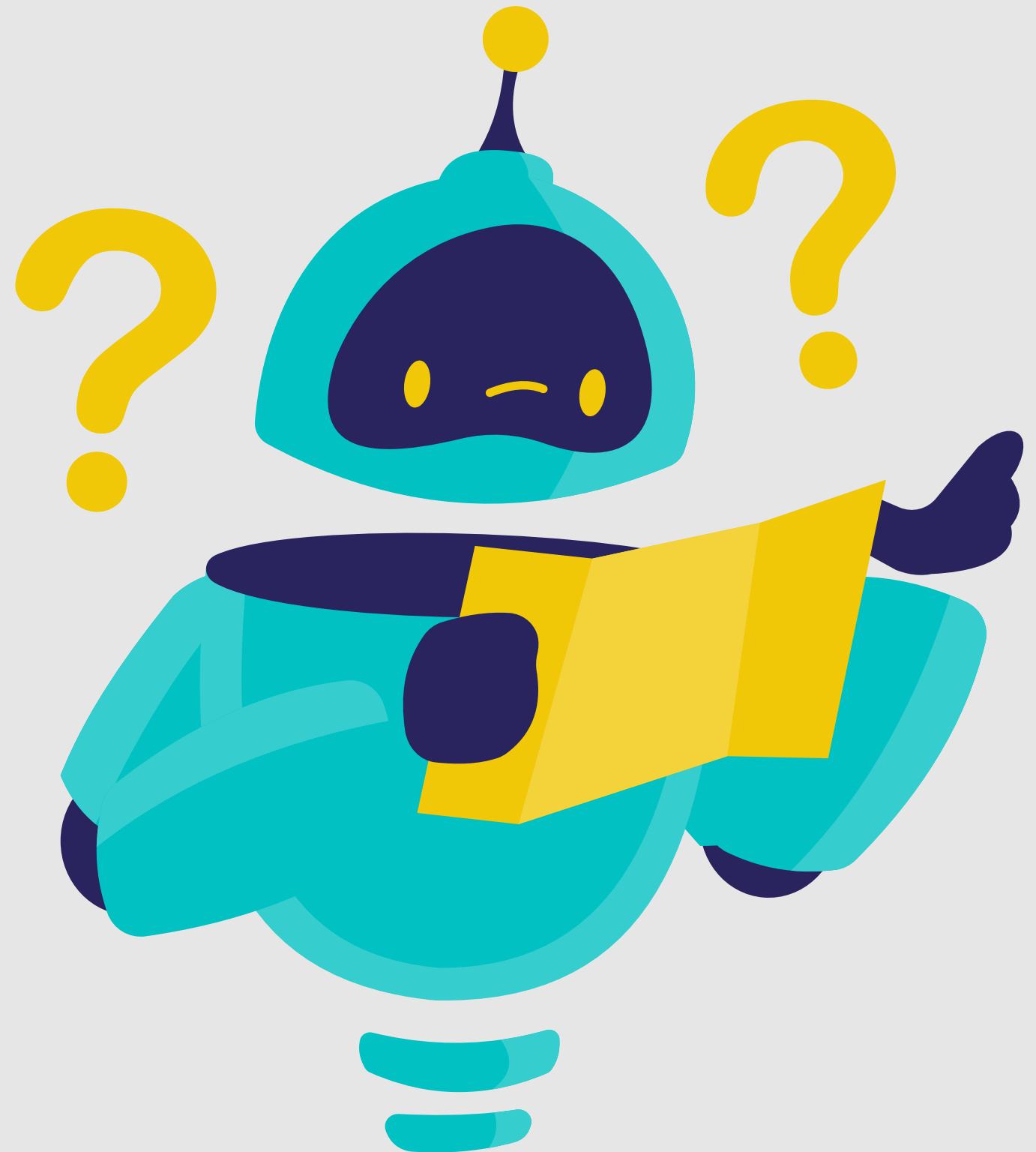
# 3D EQUATION



$$\begin{cases} x_1(t+1) &= x_1(t) + \tau(u_1(t) \cos(x_3(t)) + w_1(t)) \\ x_2(t+1) &= x_2(t) + \tau(u_1(t) \sin(x_3(t)) + w_2(t)) \\ x_3(t+1) &= x_3(t) + \tau(u_2(t) + w_3(t)) \pmod{2\pi} \end{cases}$$

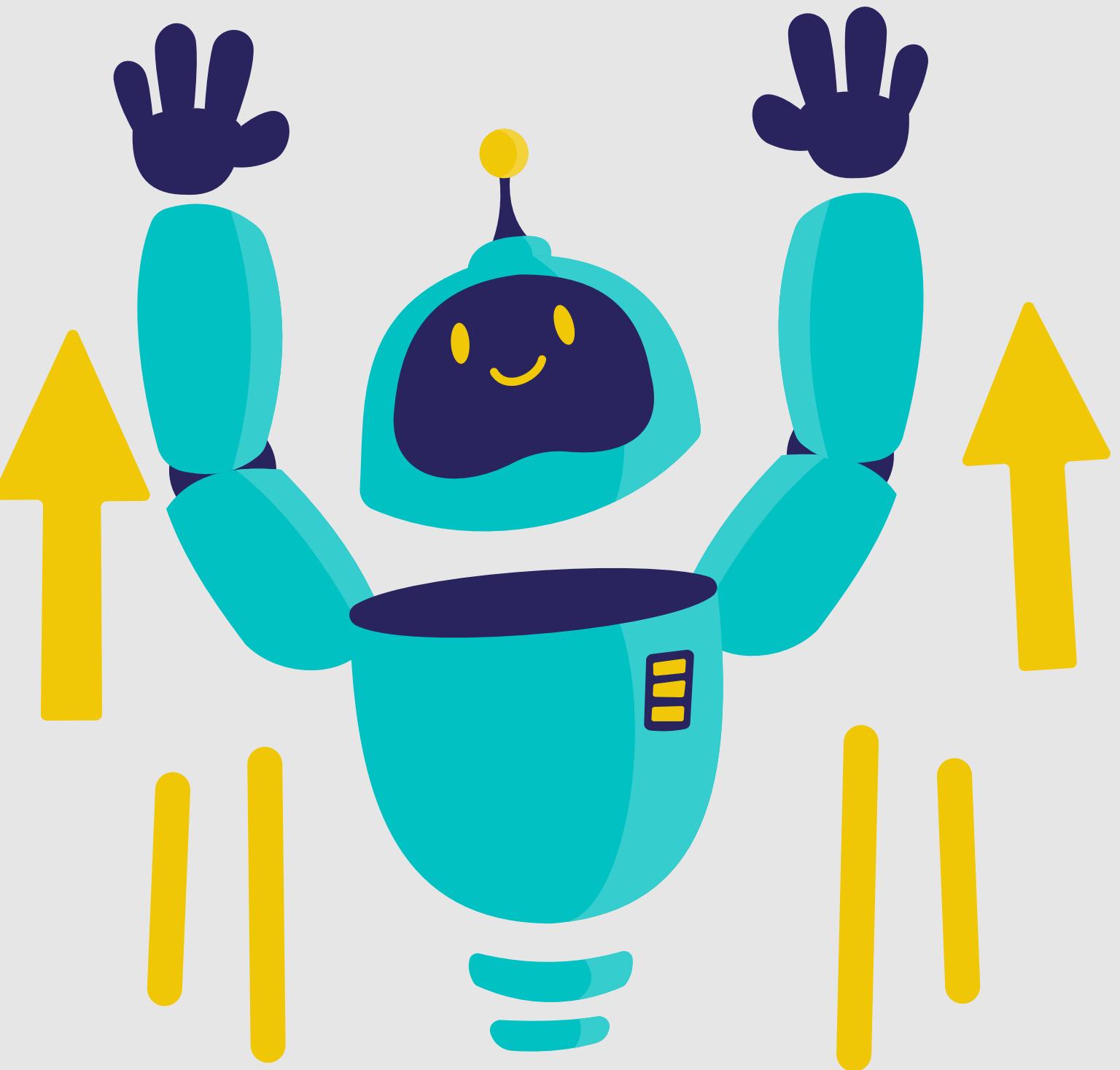
# STRUGGLES

- **TOO SLOW**
- **IT NEVER CONVERGES**
- **IT'S CLOSE TO R3 BUT DOESN'T REACH**
- **THE FILES ARE TOO BIG**
- **IT USES A LOT OF RAM**



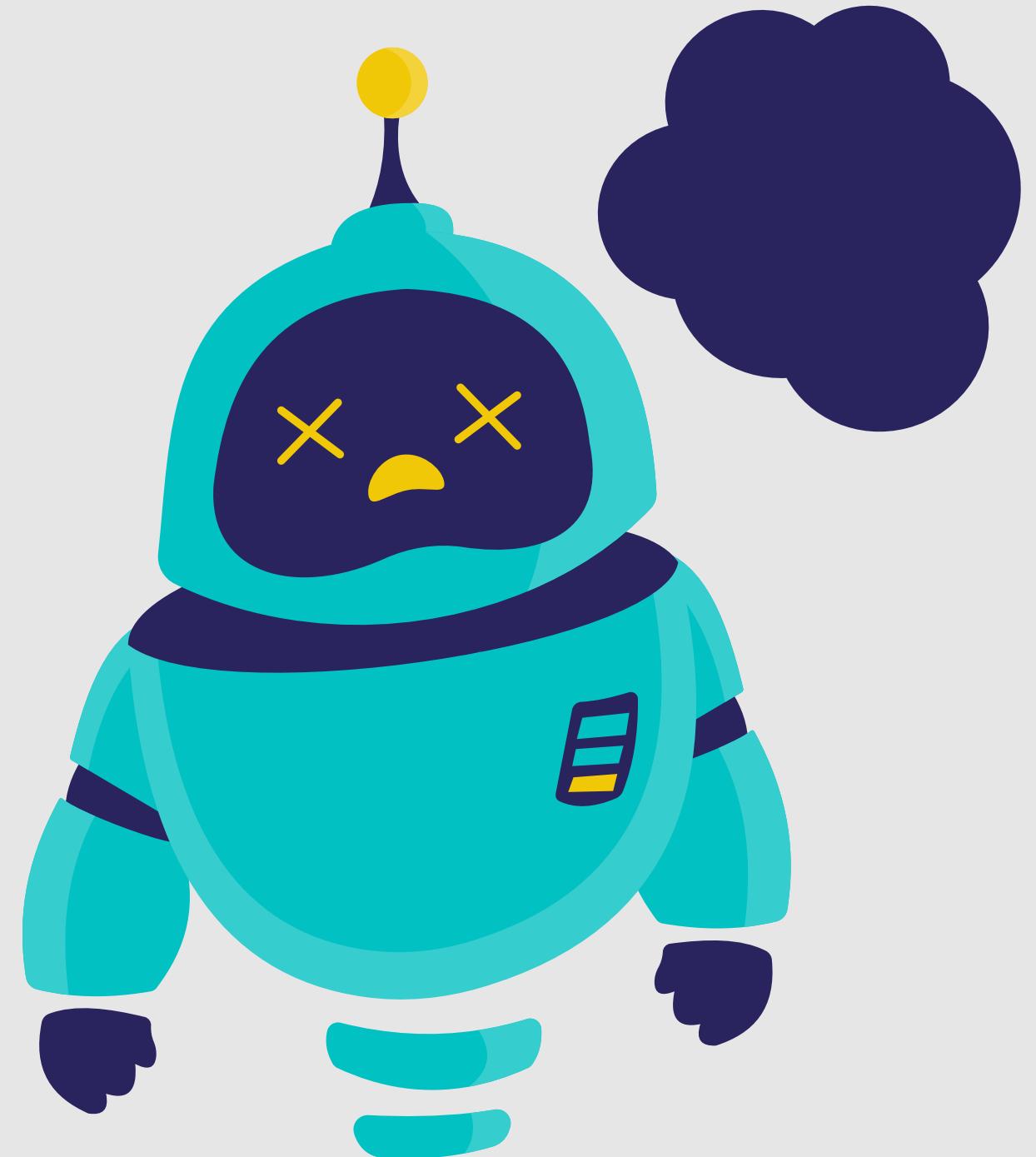
# KAGGLE

**WE HAD SWITCHED TO  
KAGGLE TO GET MORE RAM  
AND PROCESSOR POWER**



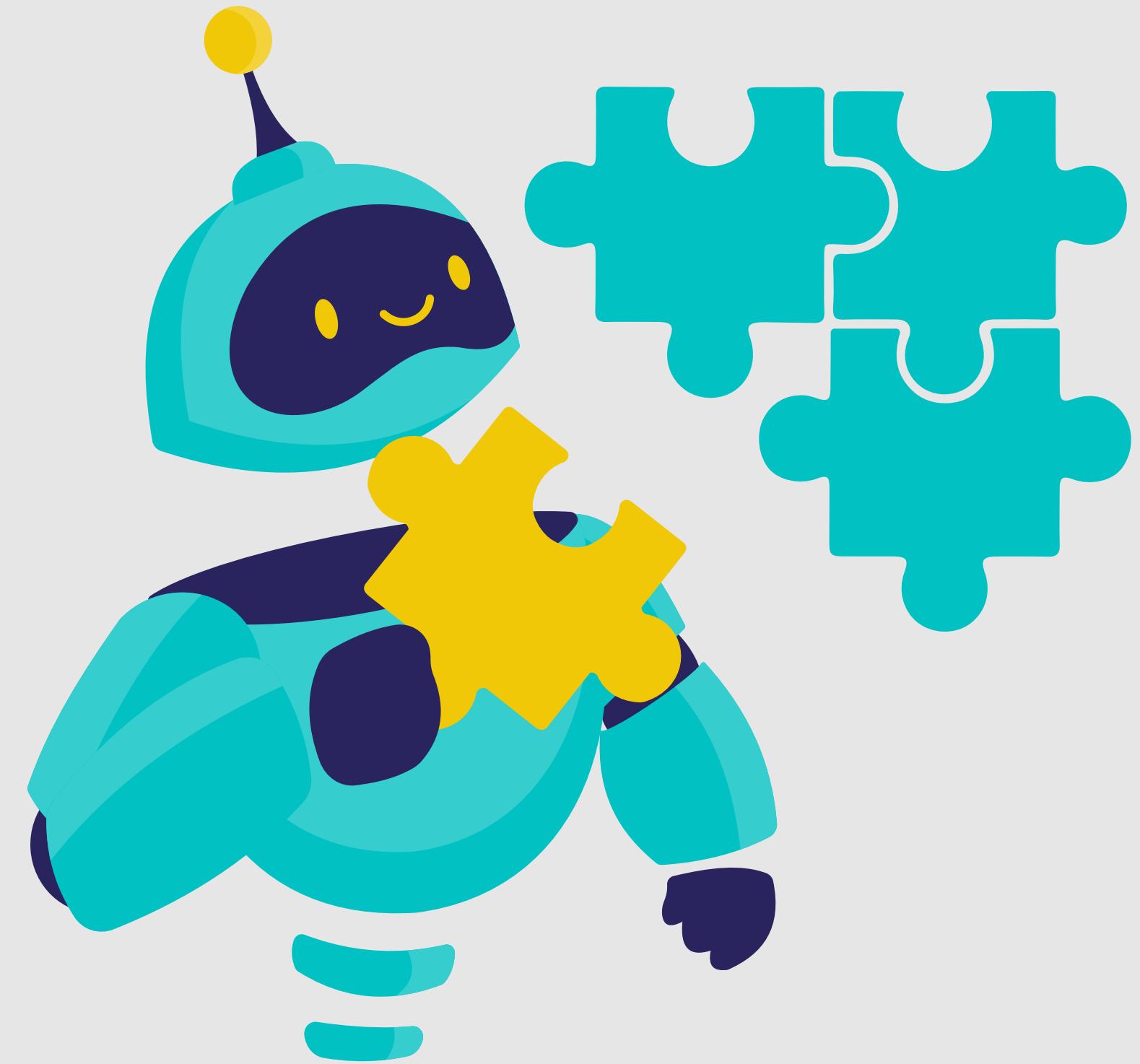
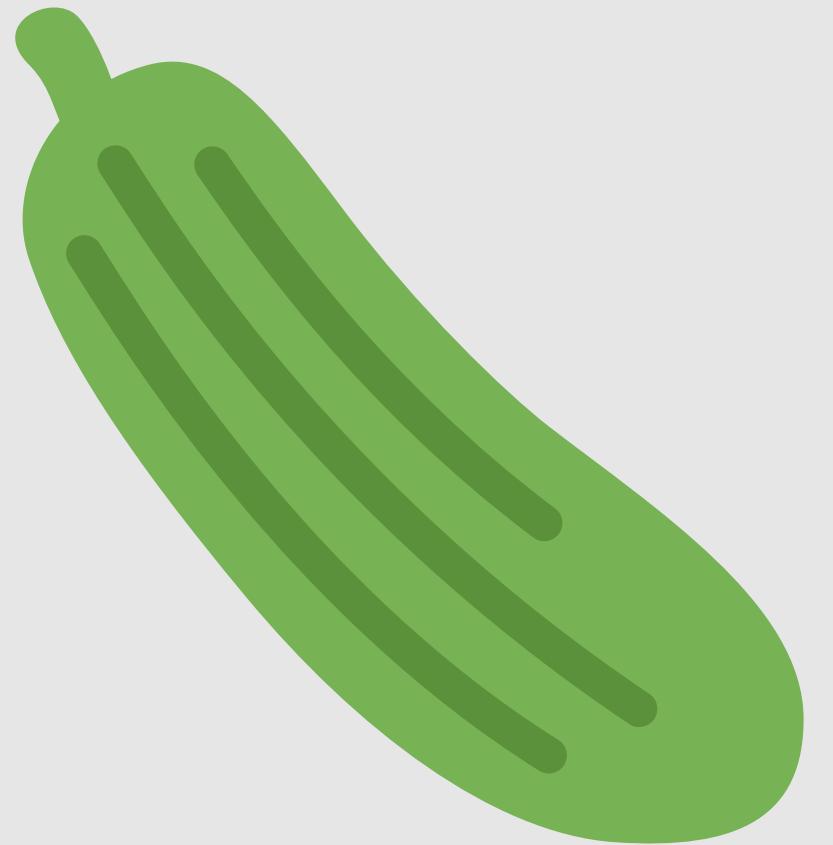
# KAGGLE

**BUT OUR CODE  
CRASHED AFTER 13  
HOURS OF RUNTIME**



# OPTIMISE FILE STORAGE

**PICKLES**



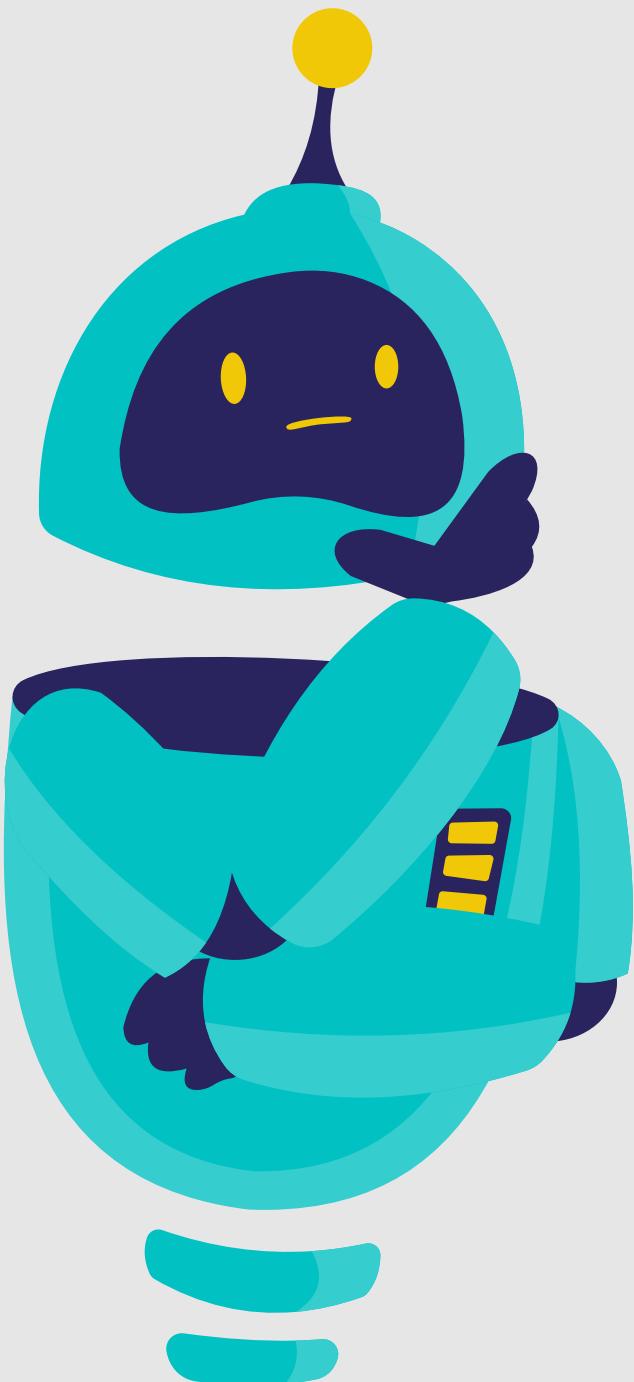
**BIT MASKING** ✗  
**TUPLE → INT** ✓

## PROS:

- **FAST EQUALITY CHECK**
- **SAVES RAM AND MEMORY**

## CONS:

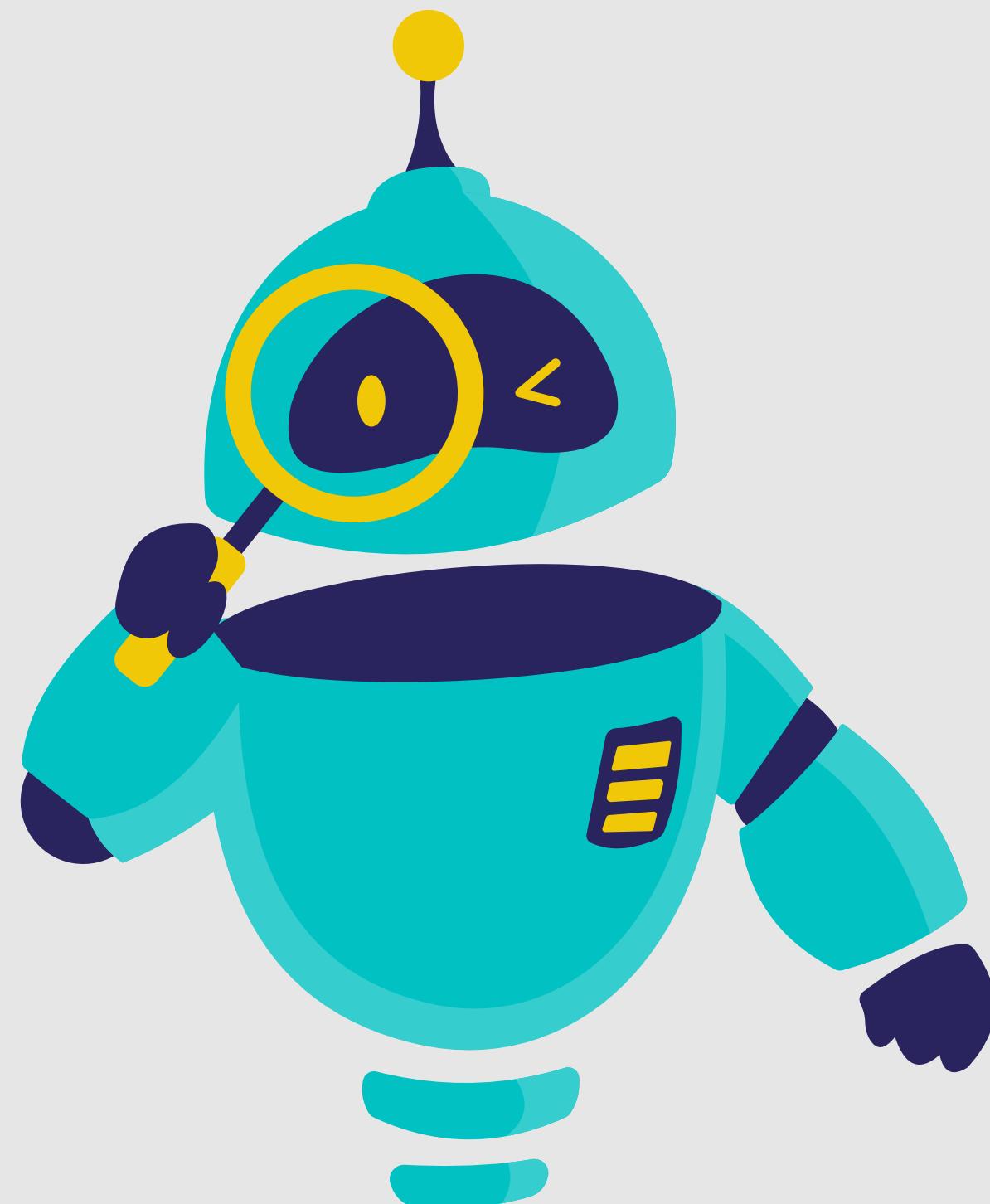
- **DEBUGGING IS HARD**



# COMPUTE LABEL

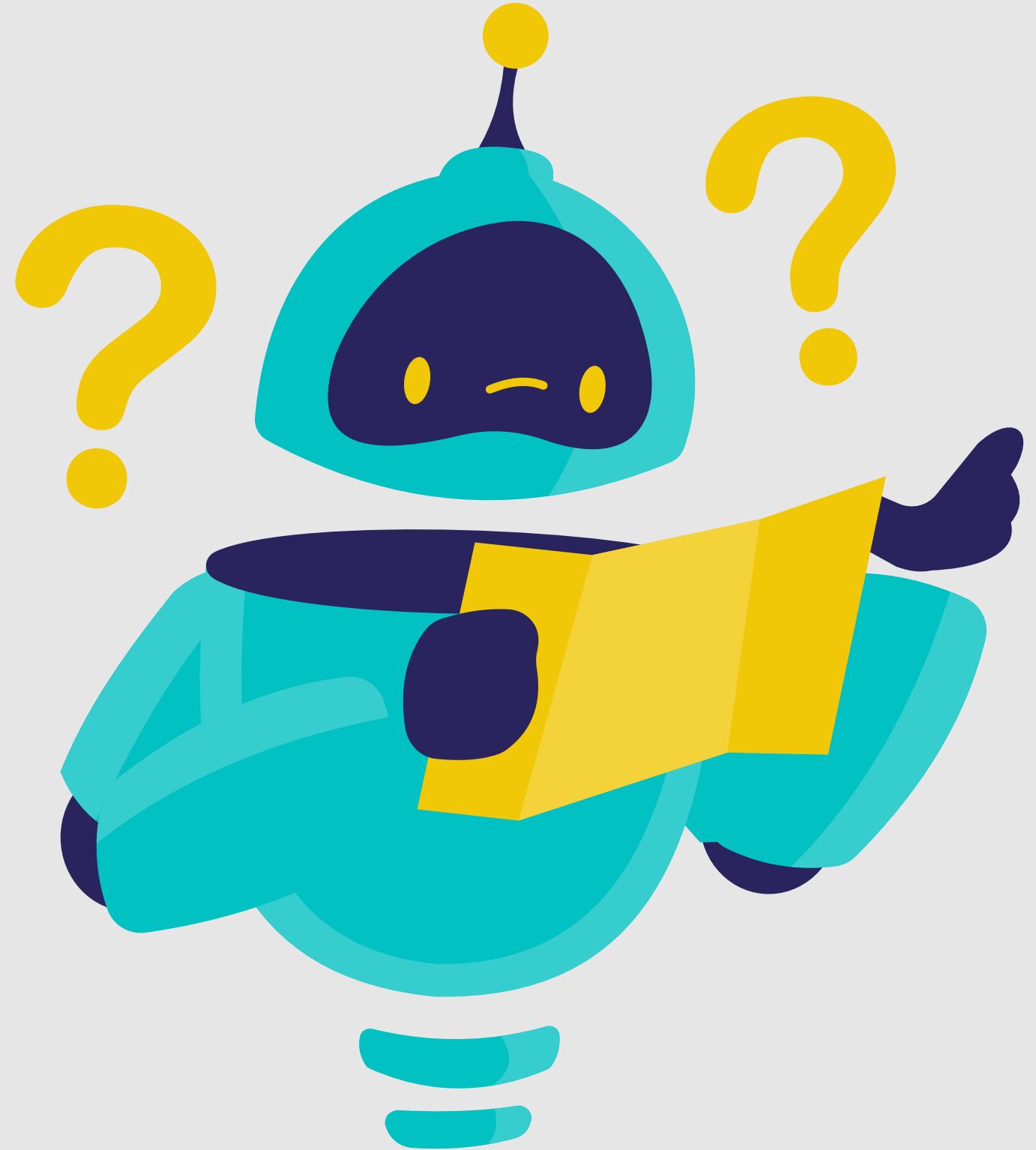
COMPUTE LABEL EACH TIME ✗

USE A LABEL GRID ✓



# LOW DISCRETISATION

**UNFORTUNATELY HERE WE  
DON'T HAVE A CHOICE, WE  
DECREASE THE  
DISCRETISATION**



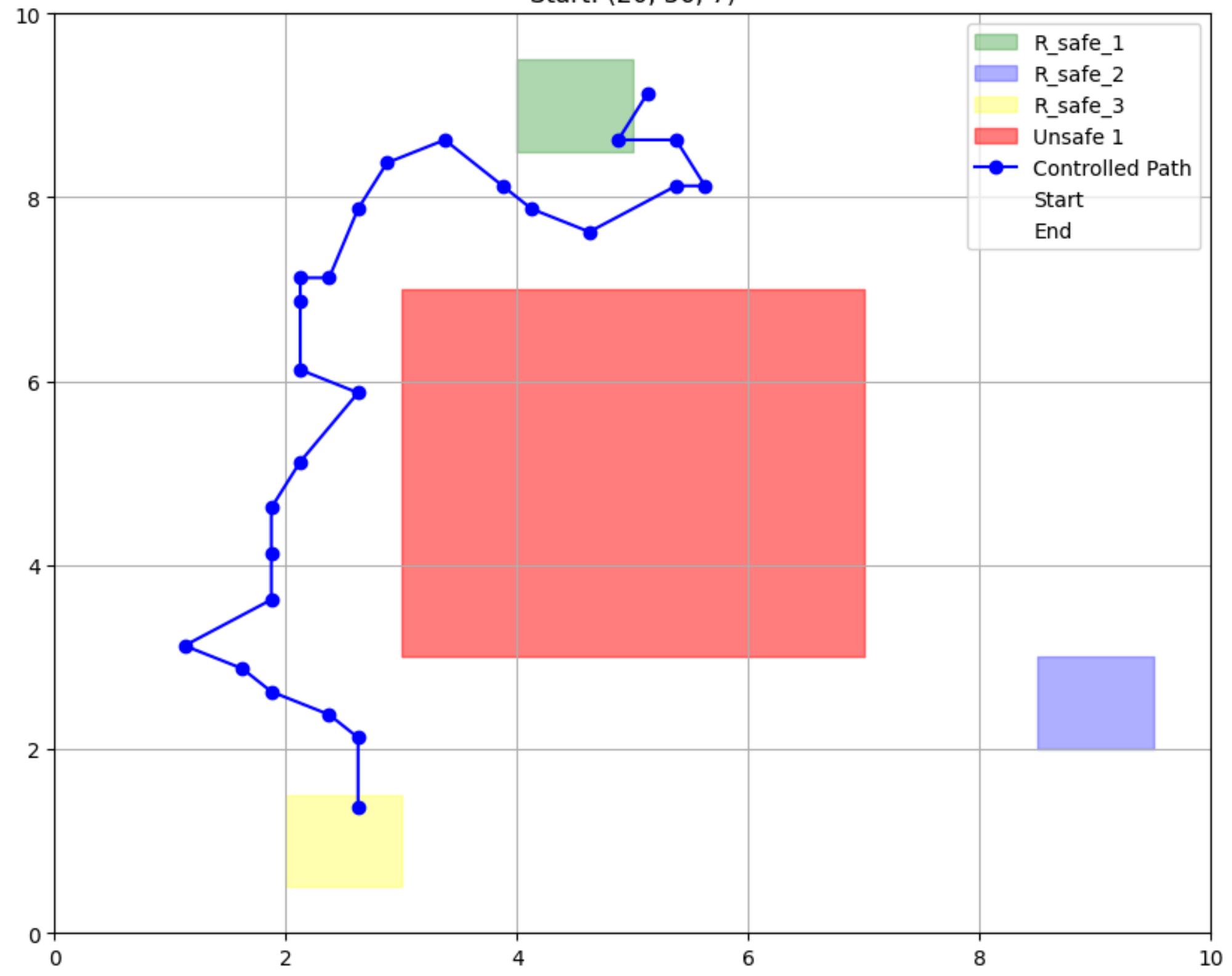
# FINAL BENCHMARK

**NOT COMPILING AFTER 13 HOURS** 

**FINISHING WITHIN 90 MINUTES**



Controller Simulation (original)  
Start: (20, 36, 7)

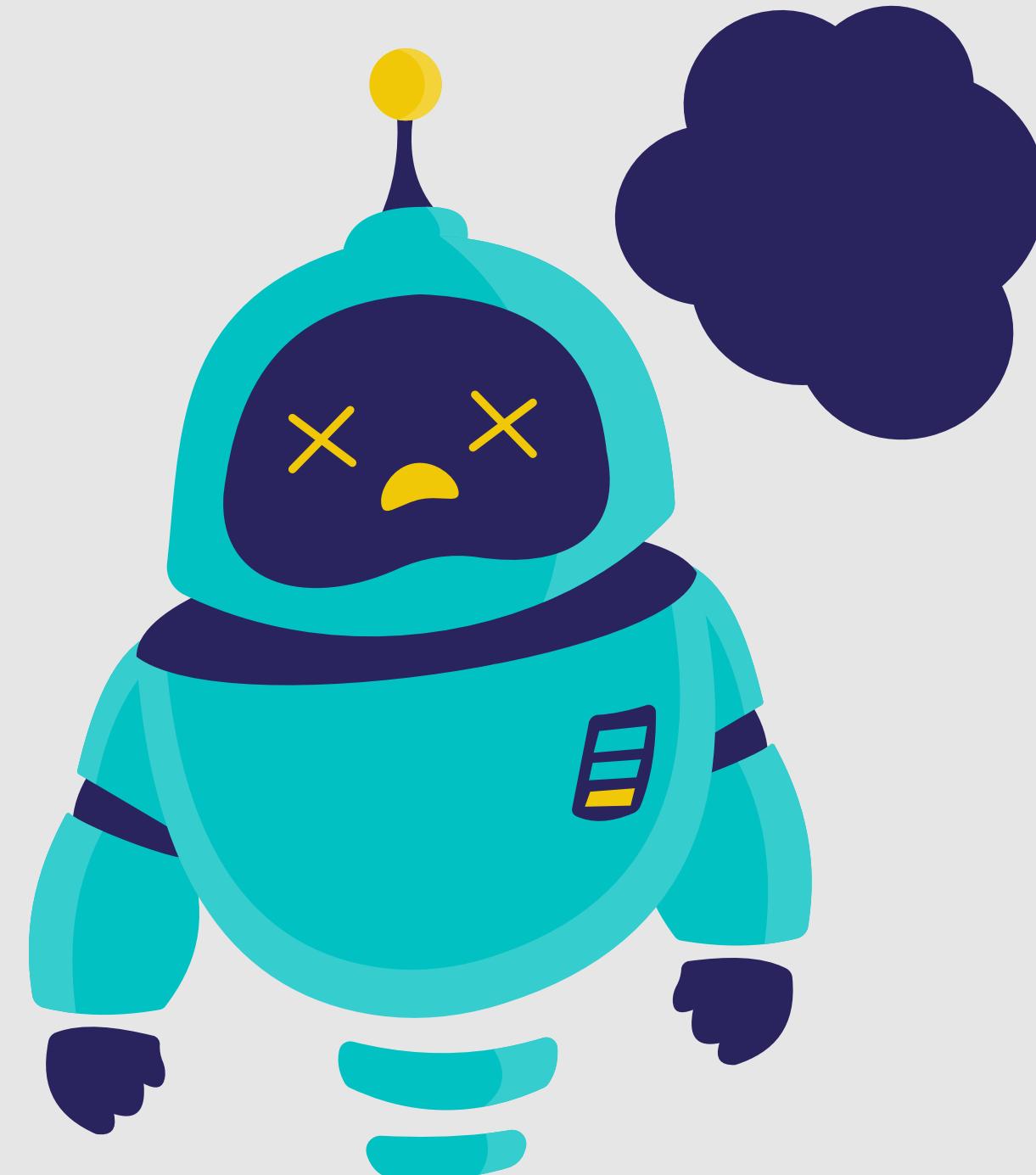




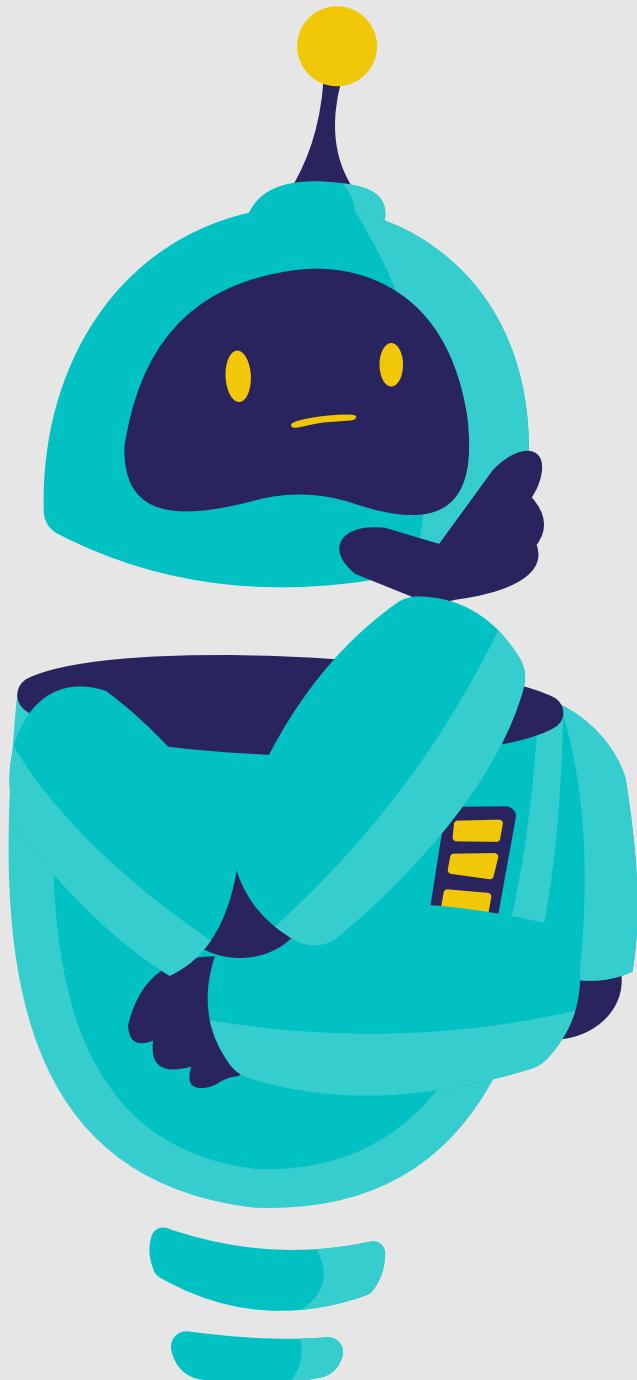
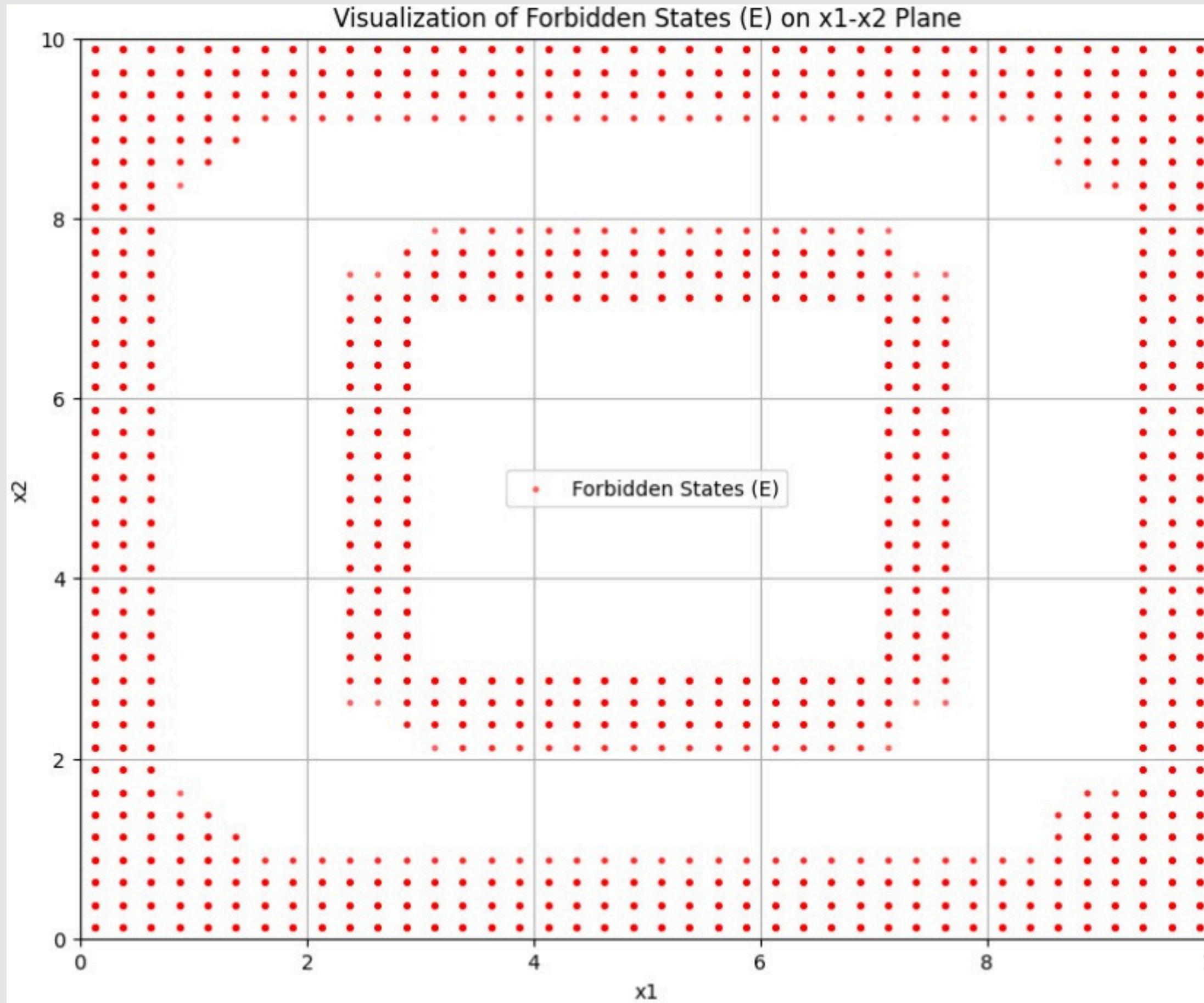
# CONTINUING WITH STRUGGLES

**ANY VS ALL IN**

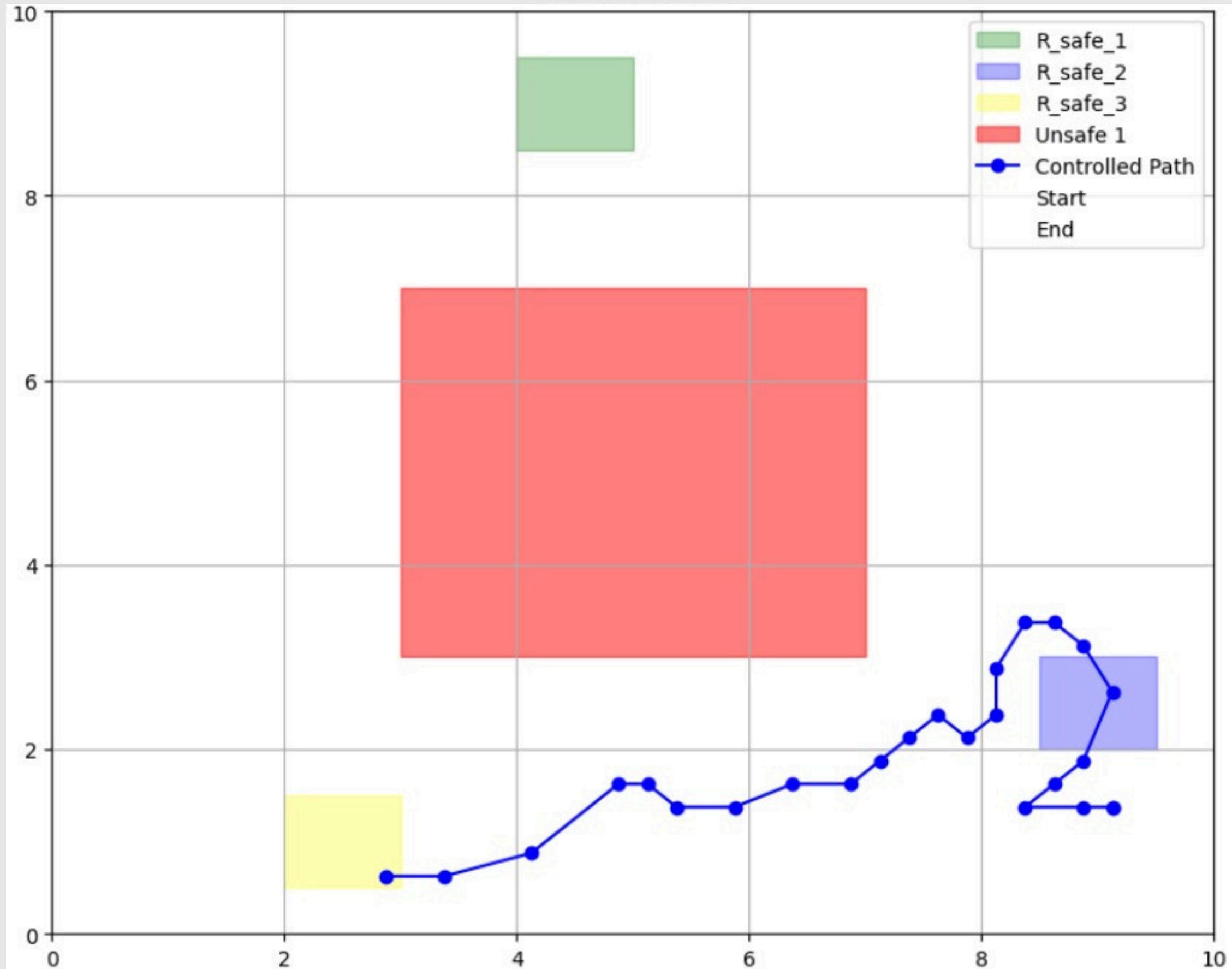
**REACHABILITY REQUIRES ALL, WHY?**



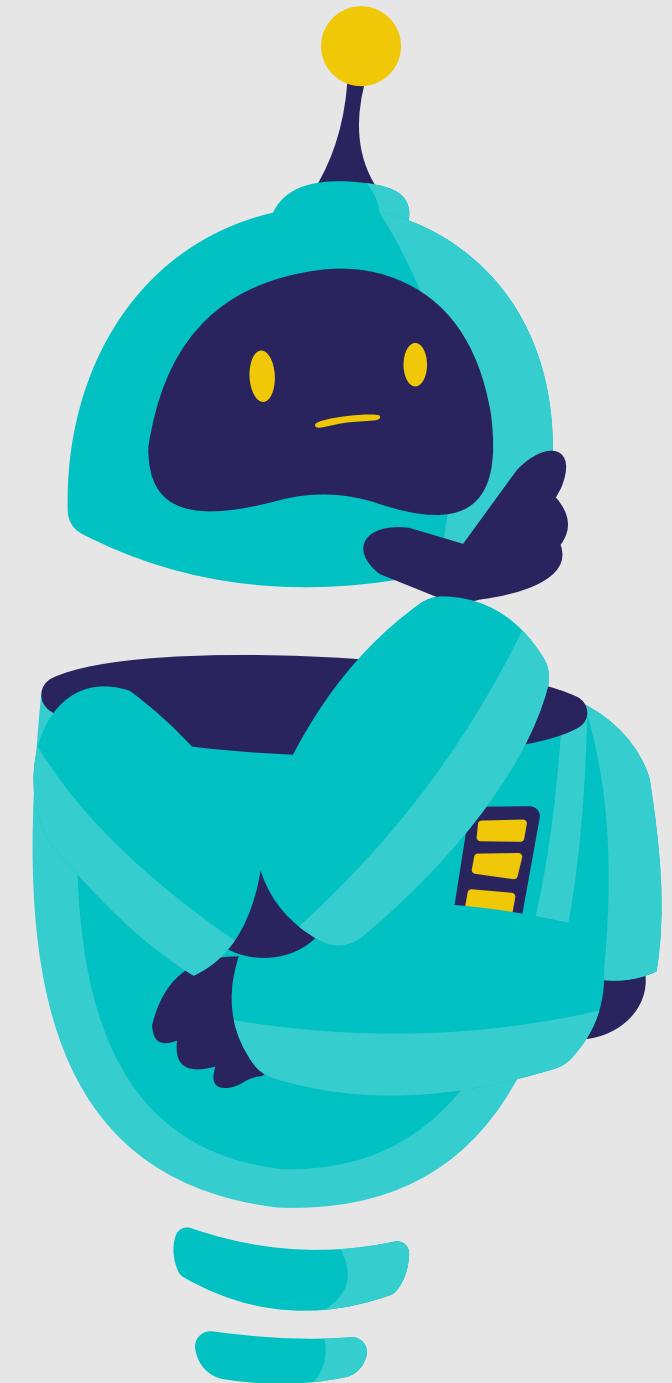
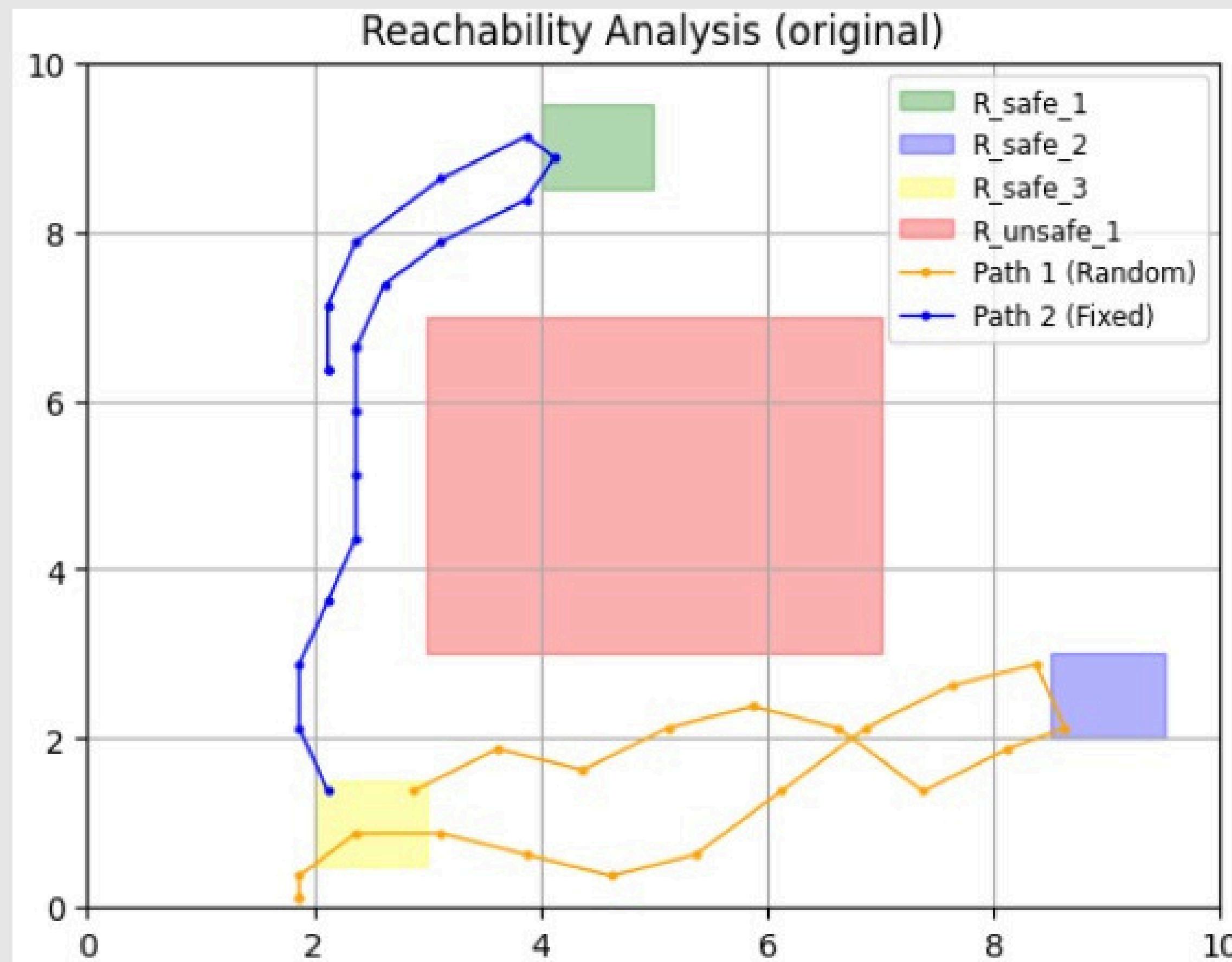
# SAFETY ISSUE: PROBLEM OF CLAMPING



# HOW TO MAKE SIMULATION LOOK BETTER?



# MAYBE BFS??

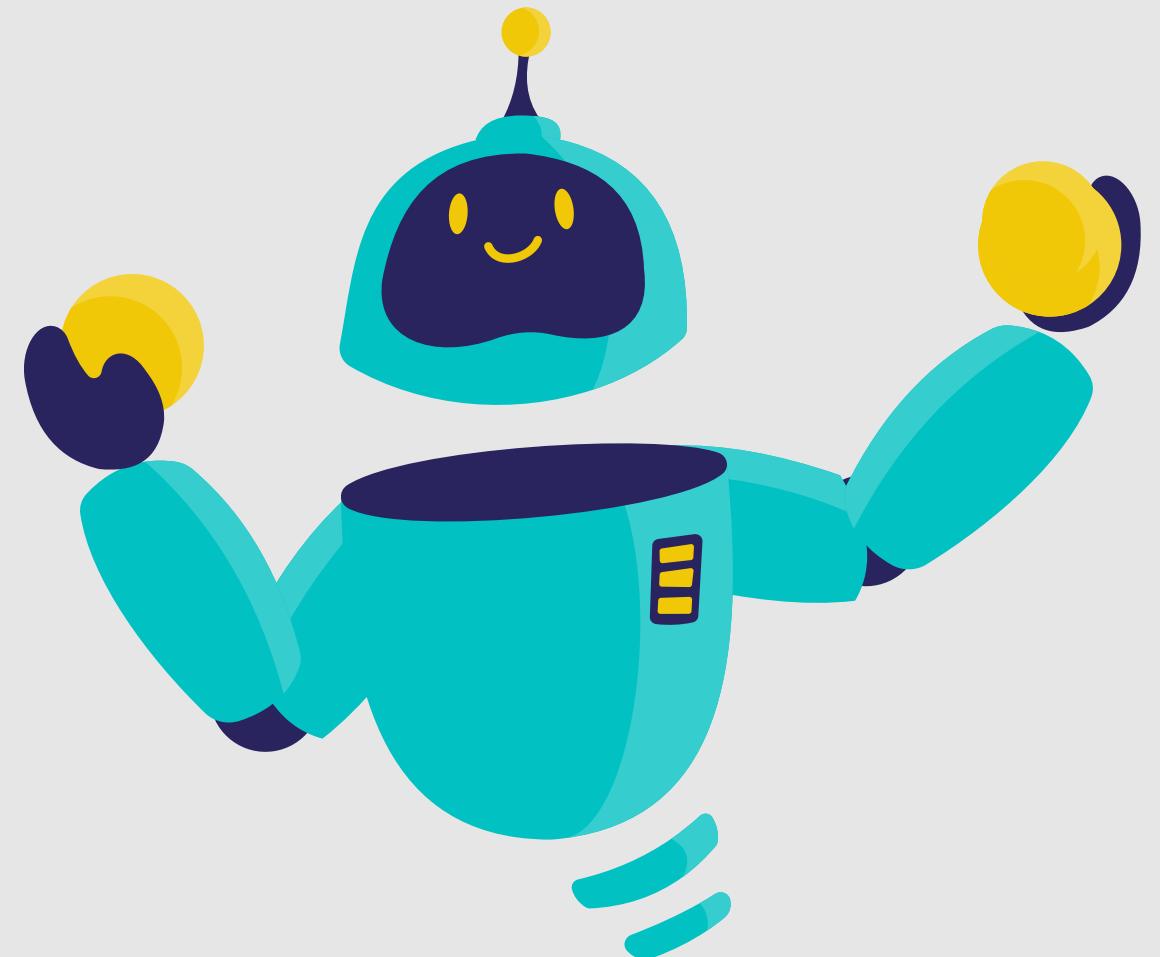




# MULTIPLE SPECIFICATION AUTOMATA

# Why multiple specs?

To evaluate how the controller behaves under different logical constraints,  
we worked with several specification modes rather than just one.



# Specification Families

**Specification modes we implemented:**

- Original spec
- Fairness spec
- Visit-count spec
- Custom JSON-defined spec



# API Integration for Custom Automata

## How it works:

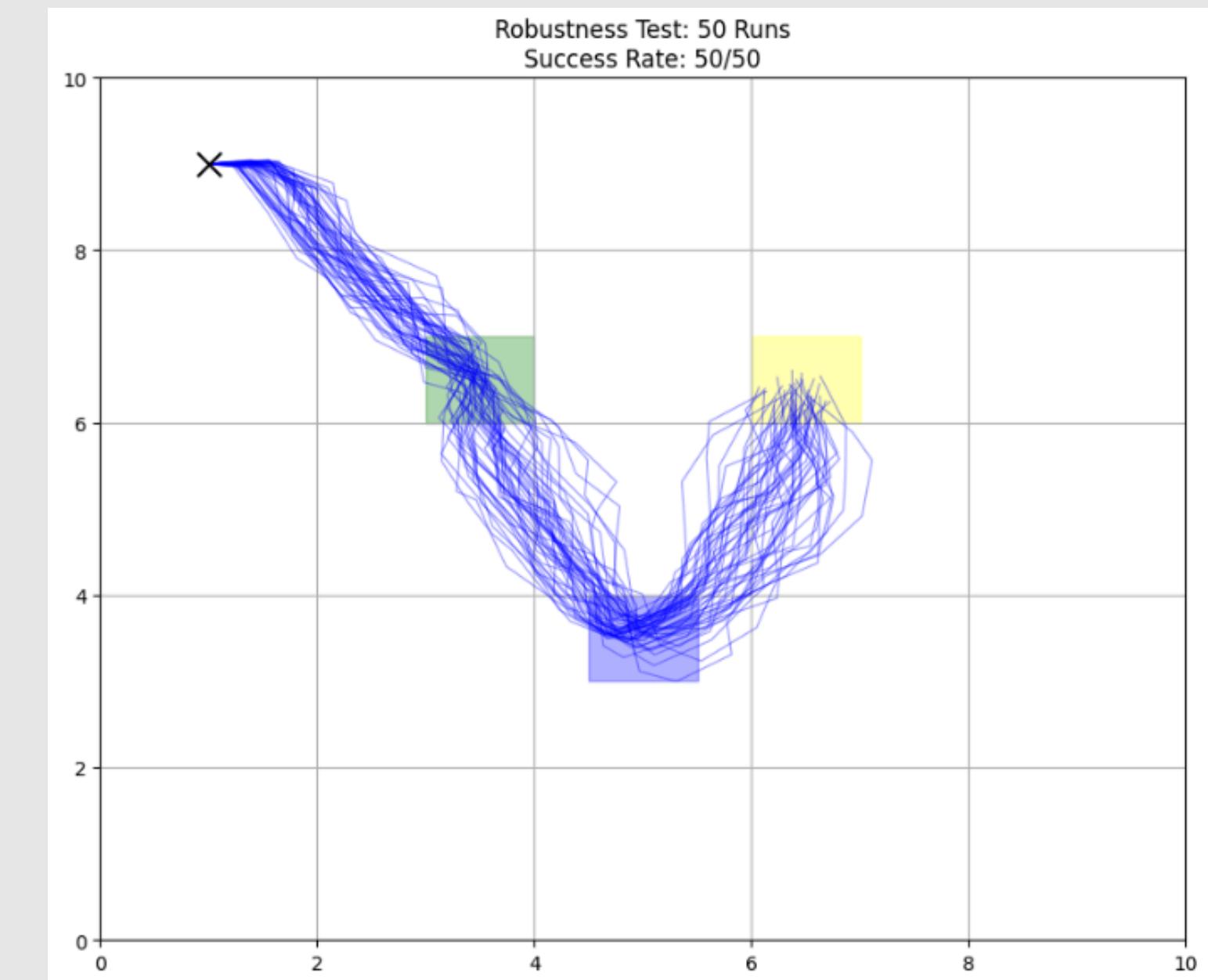
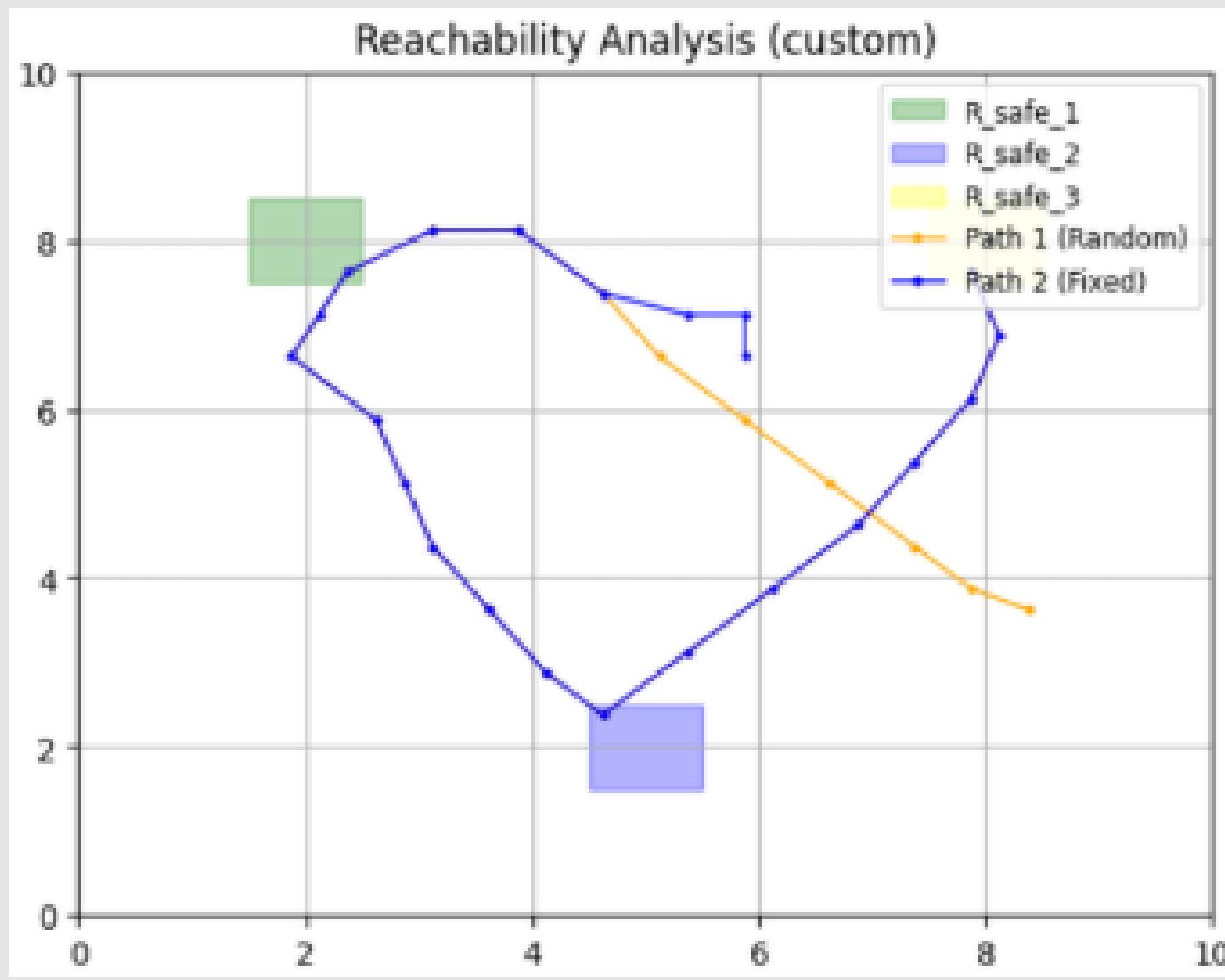
- The user can optionally define regions or let the AI infer them.
- We built a superprompt for the model.
- The model outputs a JSON automaton matching the request.
- Our code processes that JSON and converts it into a controller.

# JSON format

the AI  
outputs

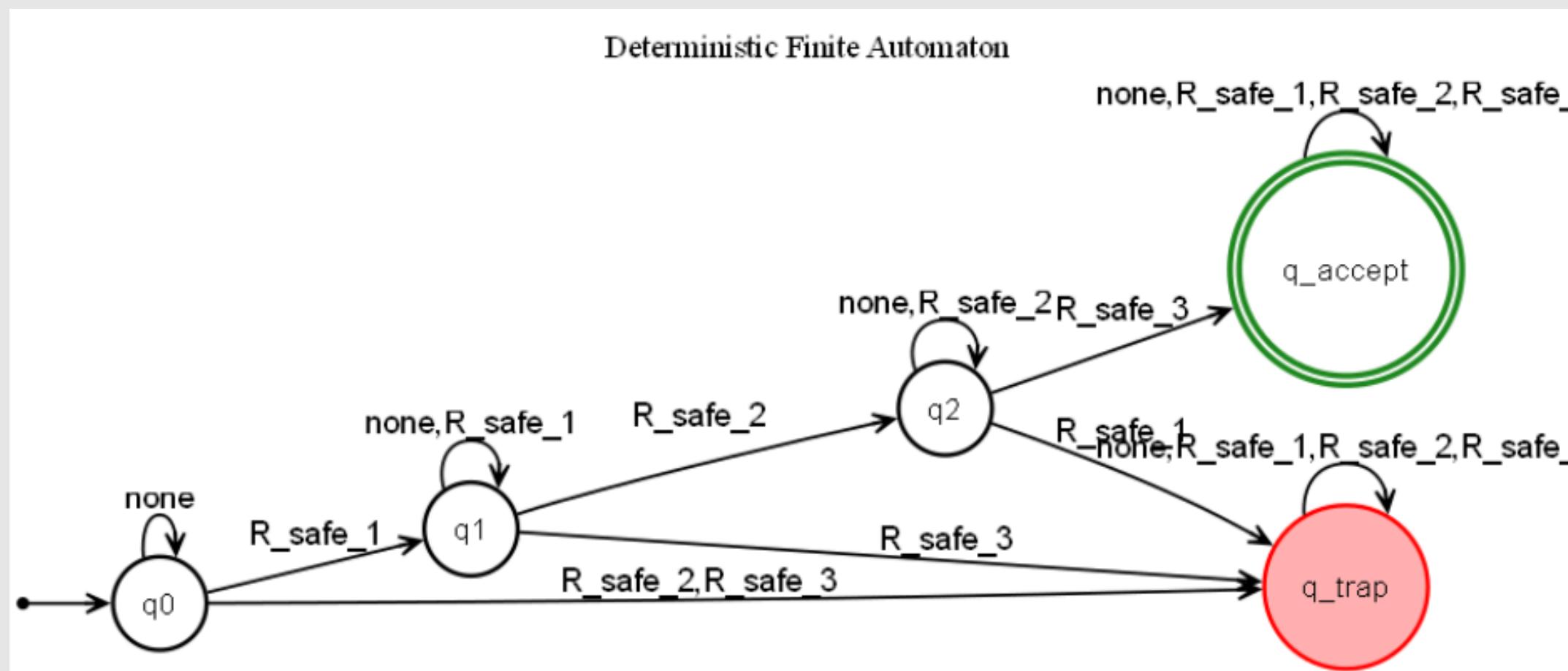
```
{  
  "regions": {  
    "R_safe_1": [[xmin,xmax],[ymin,ymax]],  
    ...  
  },  
  "danger_zones": {  
    "R_unsafe_1": [[xmin,xmax],[ymin,ymax]],  
    ...  
  },  
  "alphabet": ["none", ...],  
  "automaton": {  
    "states": [...],  
    "initial": "q0",  
    "accepting": [...],  
    "transitions": [  
      {"from": "...", "label": "...", "to": "..."},  
      ...  
    ]  
  }  
}
```

# Example: AI-generated automaton for “draw a V shape in the center of the grid”.



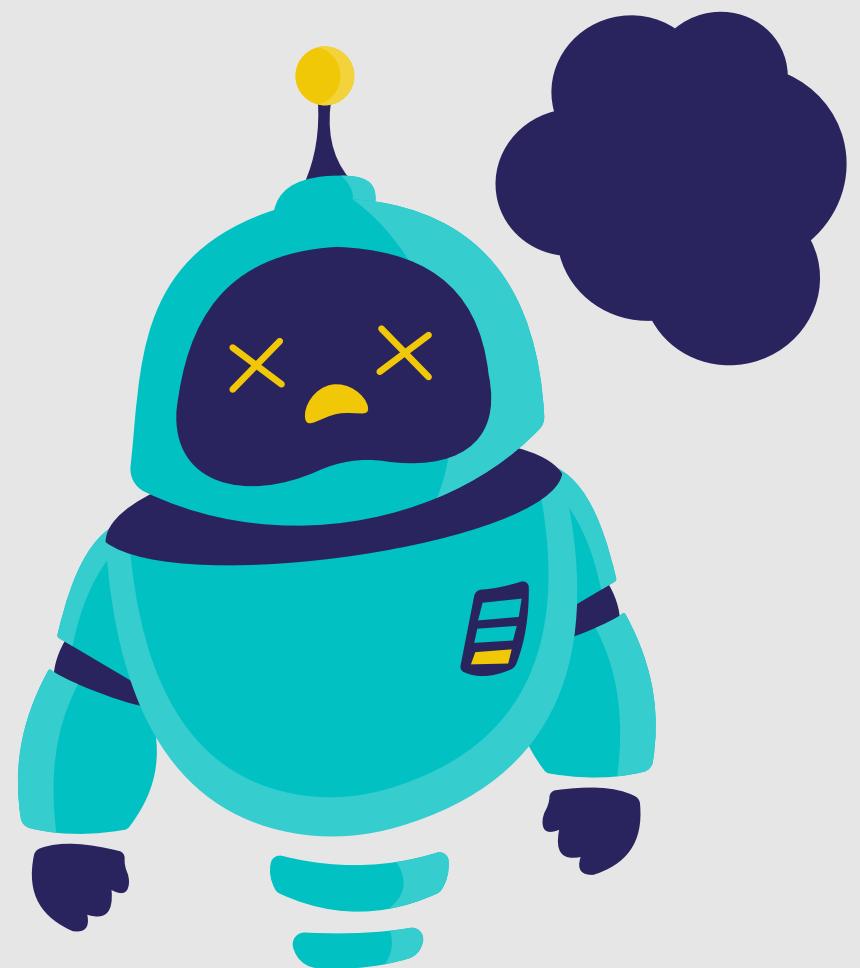
# Code to visualize automaton transition table and diagram

STATE	none	R_safe_1	R_safe_2	R_safe_3
q0	q0	q1	q_trap	q_trap
q1	q1	q1	q2	q_trap
q2	q2	q_trap	q2	q_accept
q_accept	q_accept	q_accept	q_accept	q_accept
q_trap	q_trap	q_trap	q_trap	q_trap



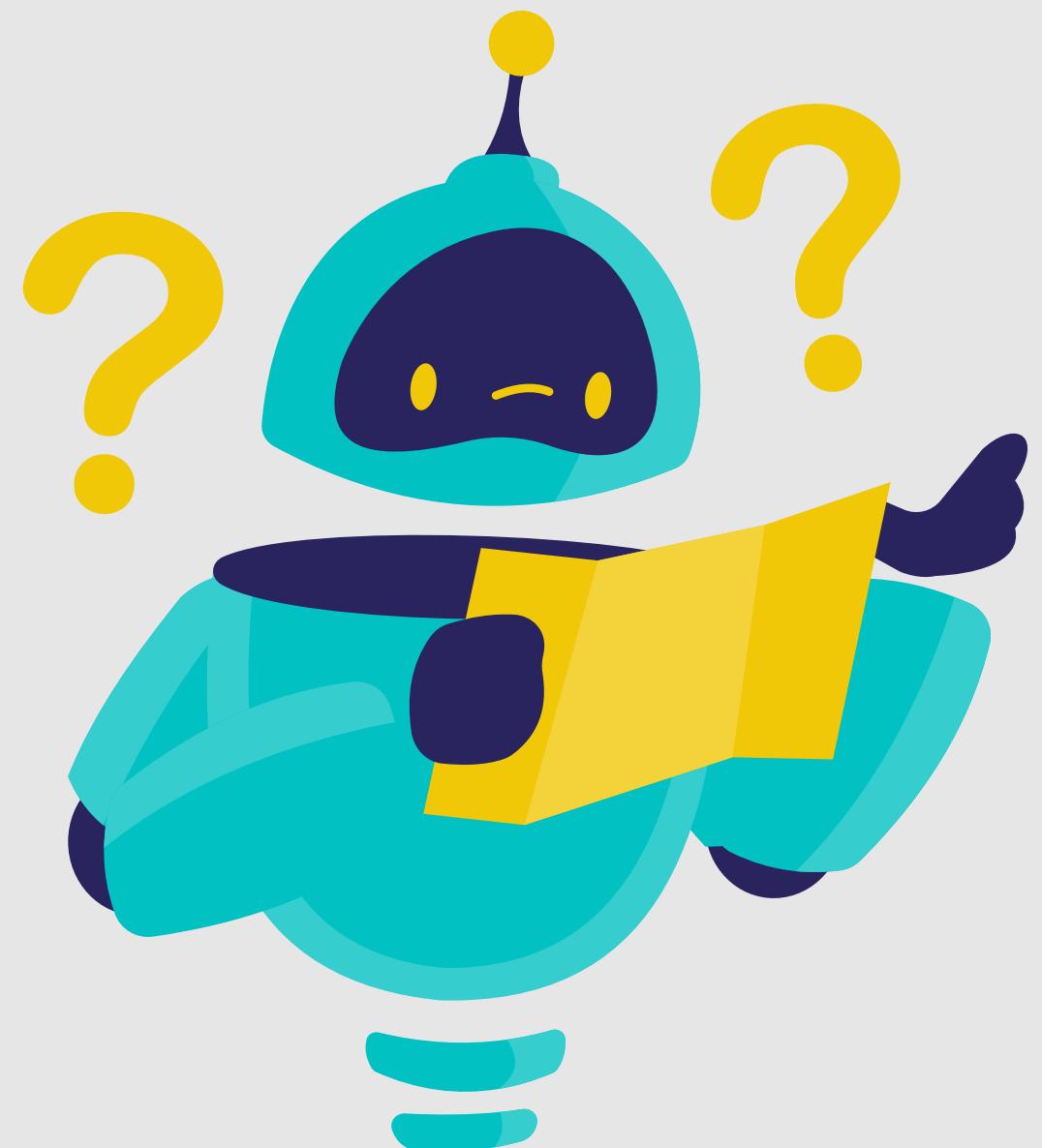
# Problems faced:

- Some automata are too strict and require extremely fine discretization, which makes them slow or even non-computable.
- Example of strict ordering that dramatically increase computation time in 3D but compiles in 2D.

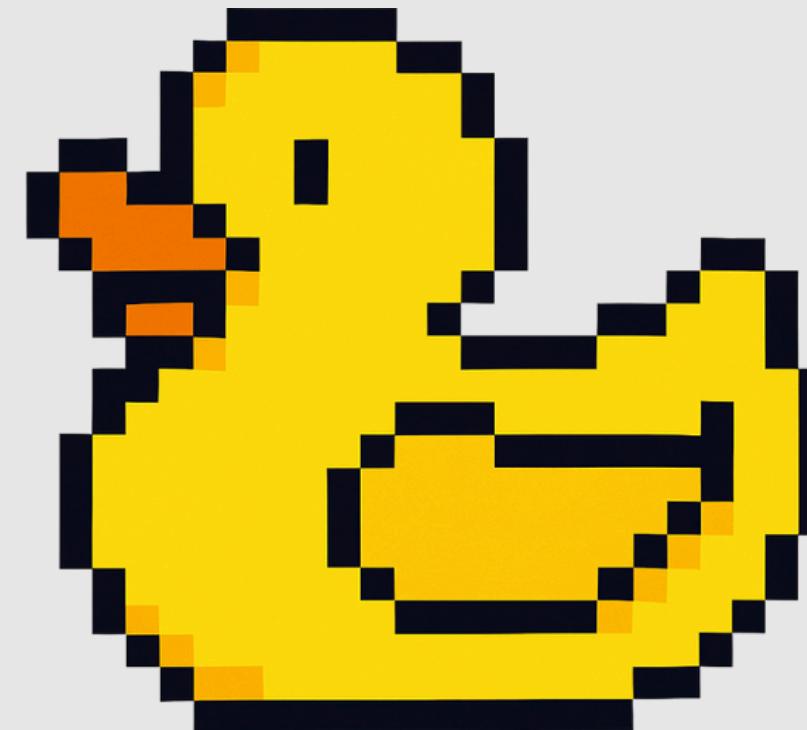
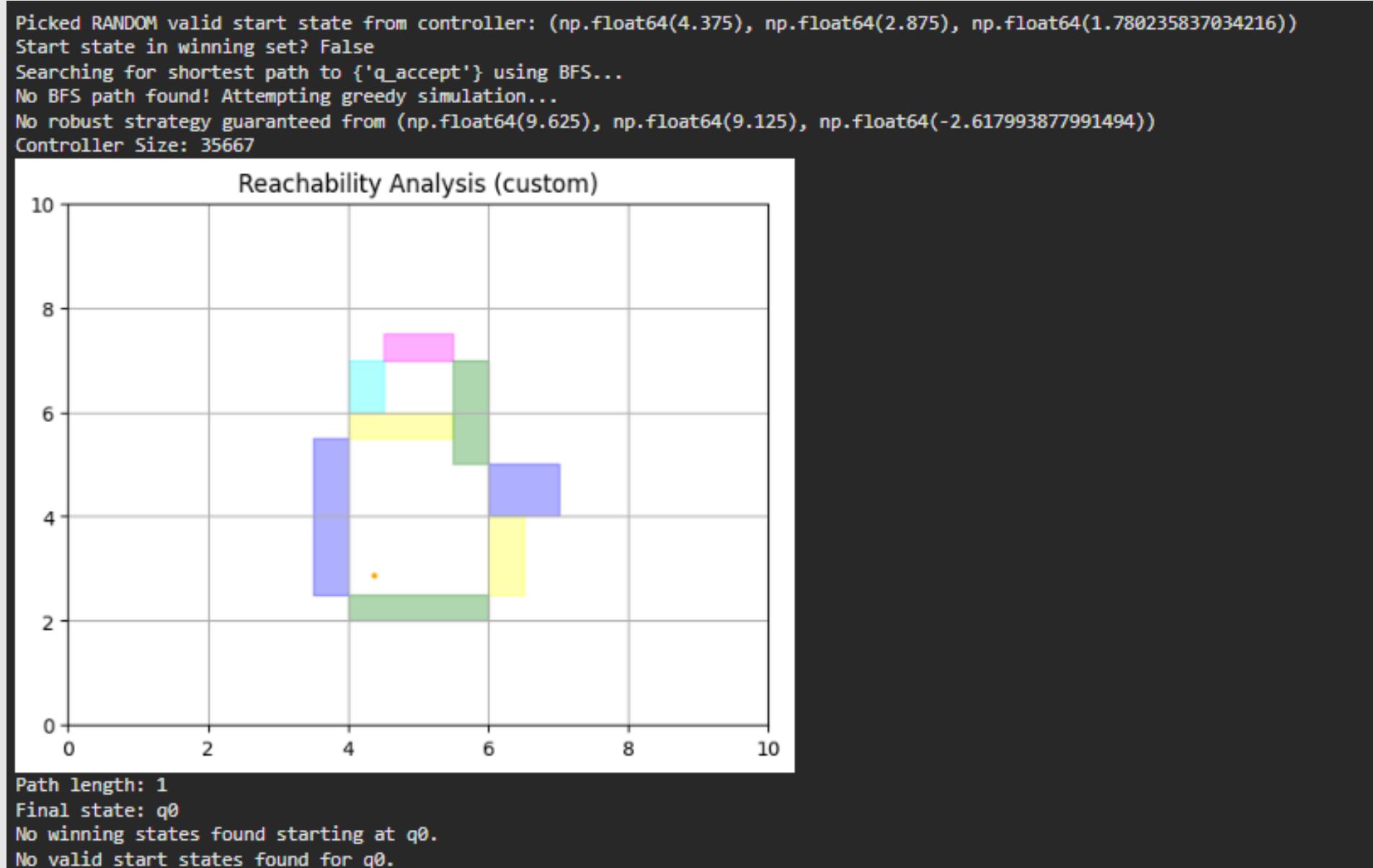


# Problems faced:

- The AI can behave unpredictably if the prompt isn't precise enough.
- We built a super-prompt to limit errors, but the AI can still make mistakes.



# Some other spec automata just don't work because they need too much movement

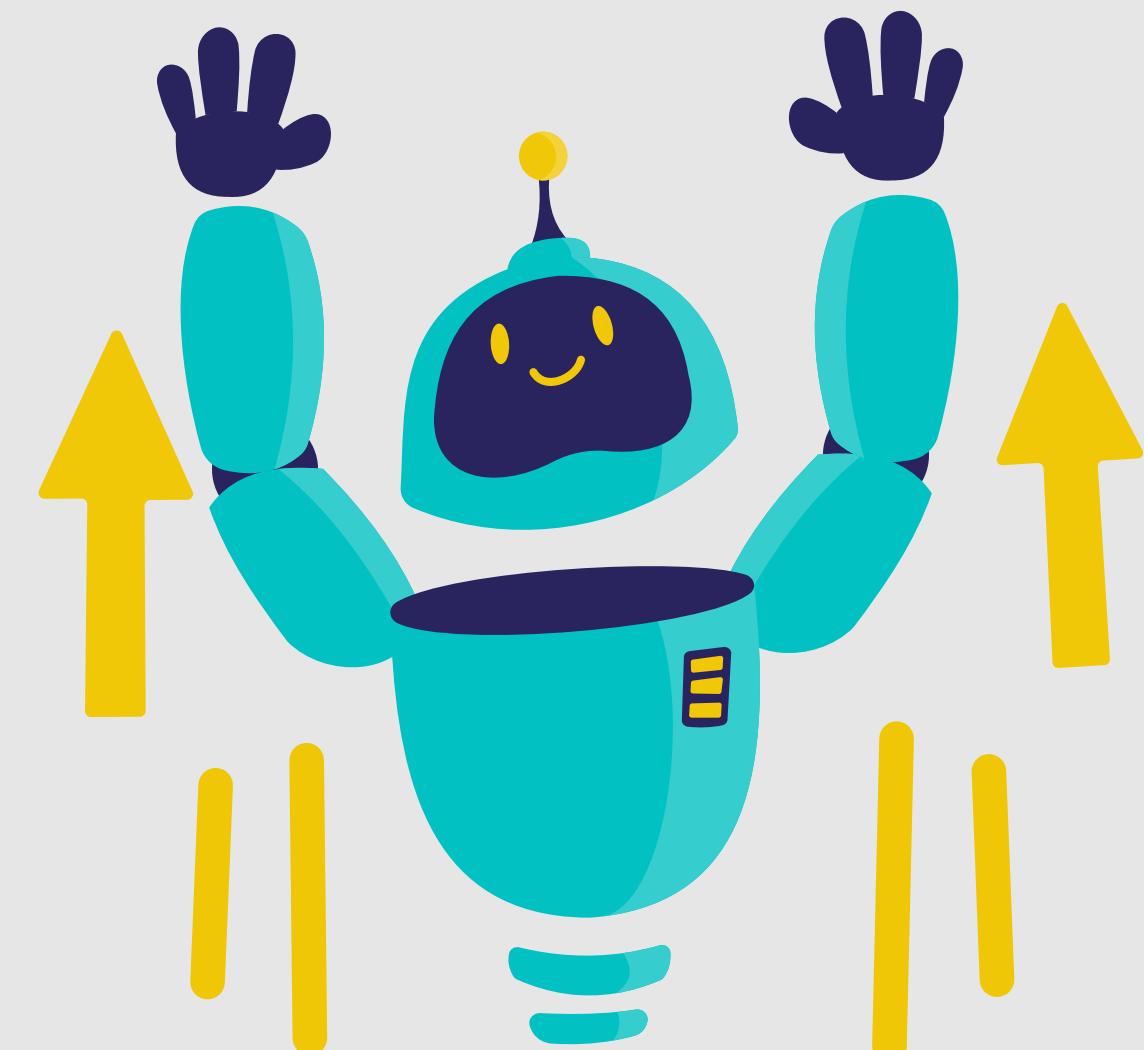


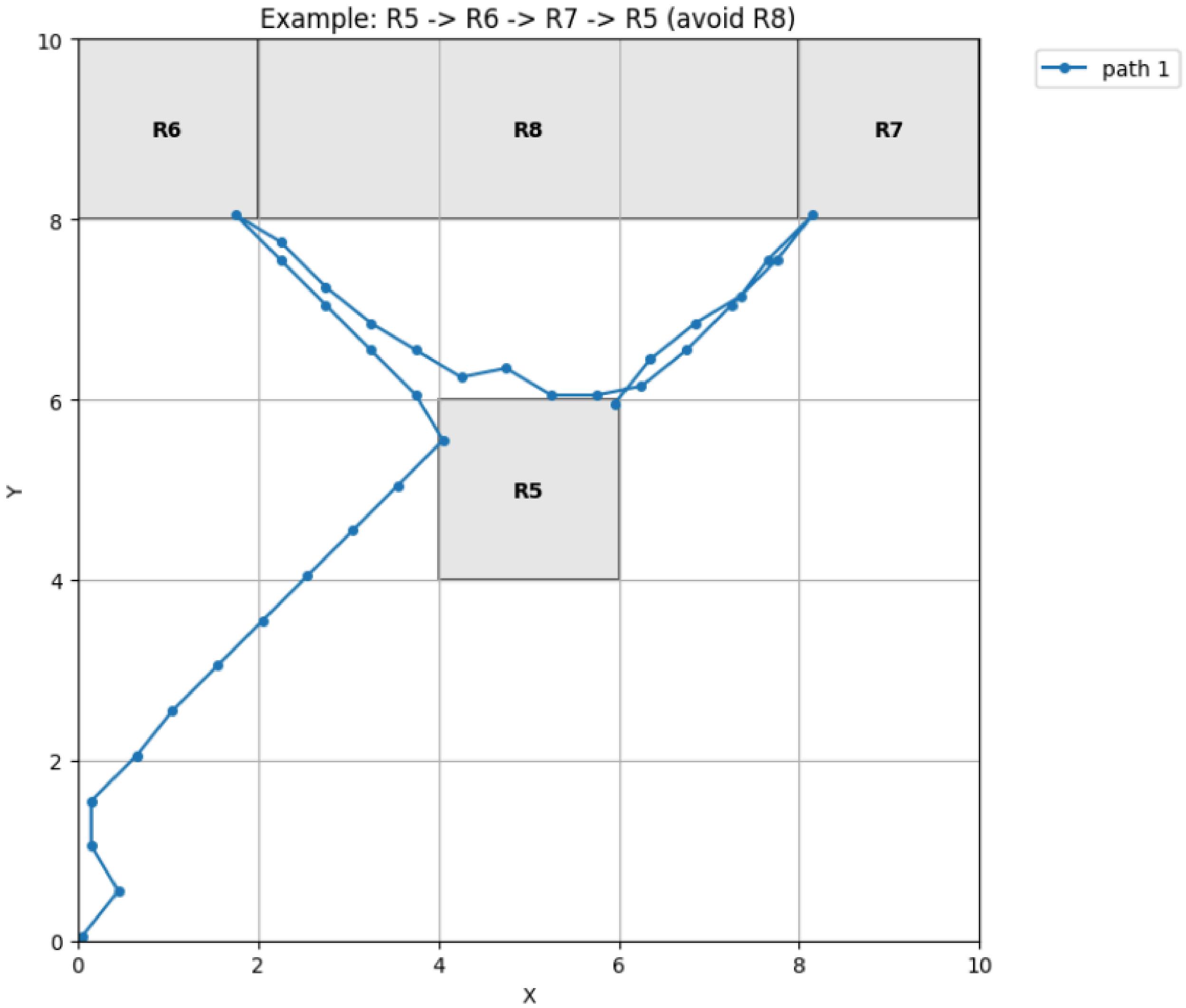
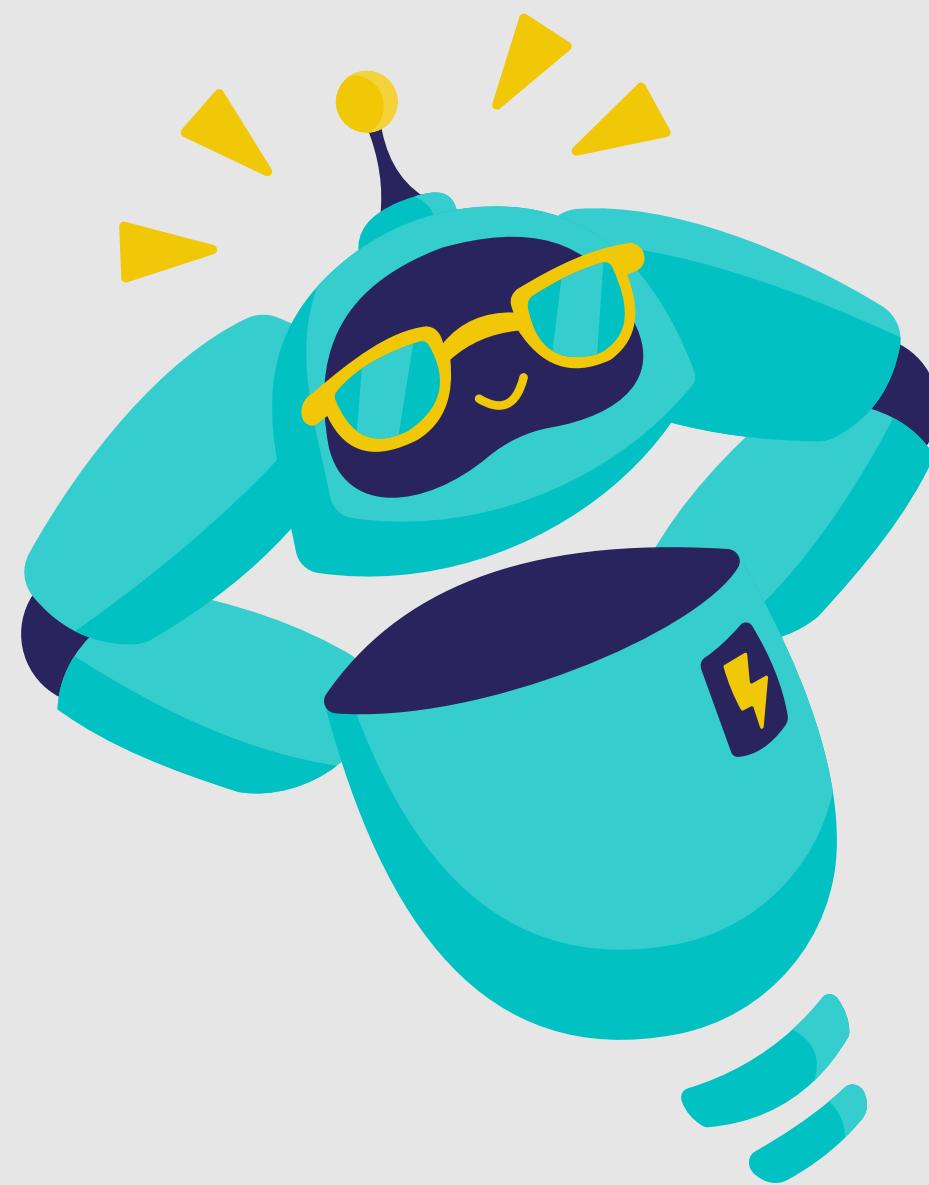


# 2D

## DYNAMIC REGIONS

INSTEAD OF WORKING WITH  
R1, R2, R3 AND R4 AS THEY ARE  
DEFINED IN THE PAPER, LET'S  
GIVE THE USER THE FREEDOM  
TO CHOOSE ANY REGION HE  
WANTS TO CONTROL

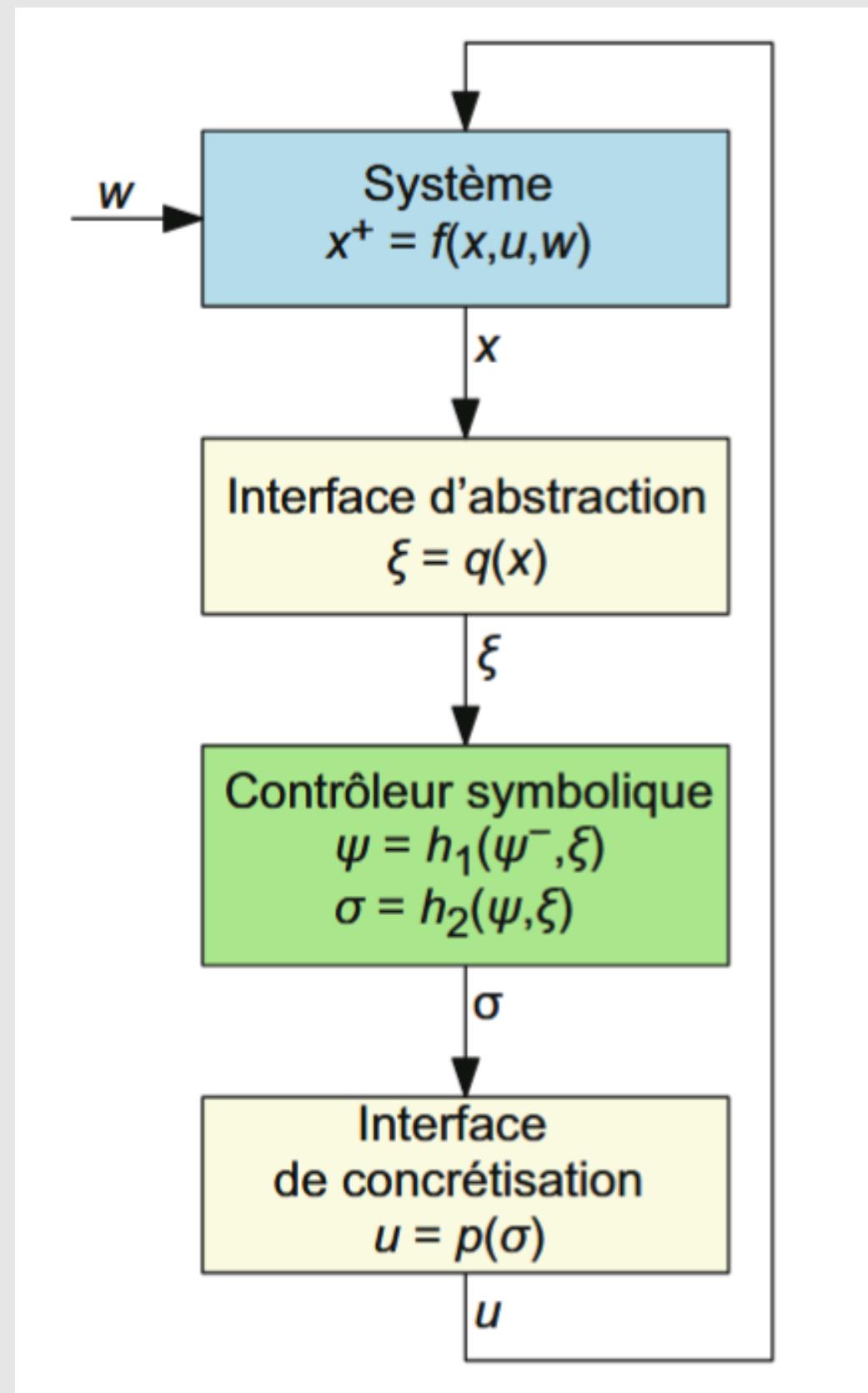




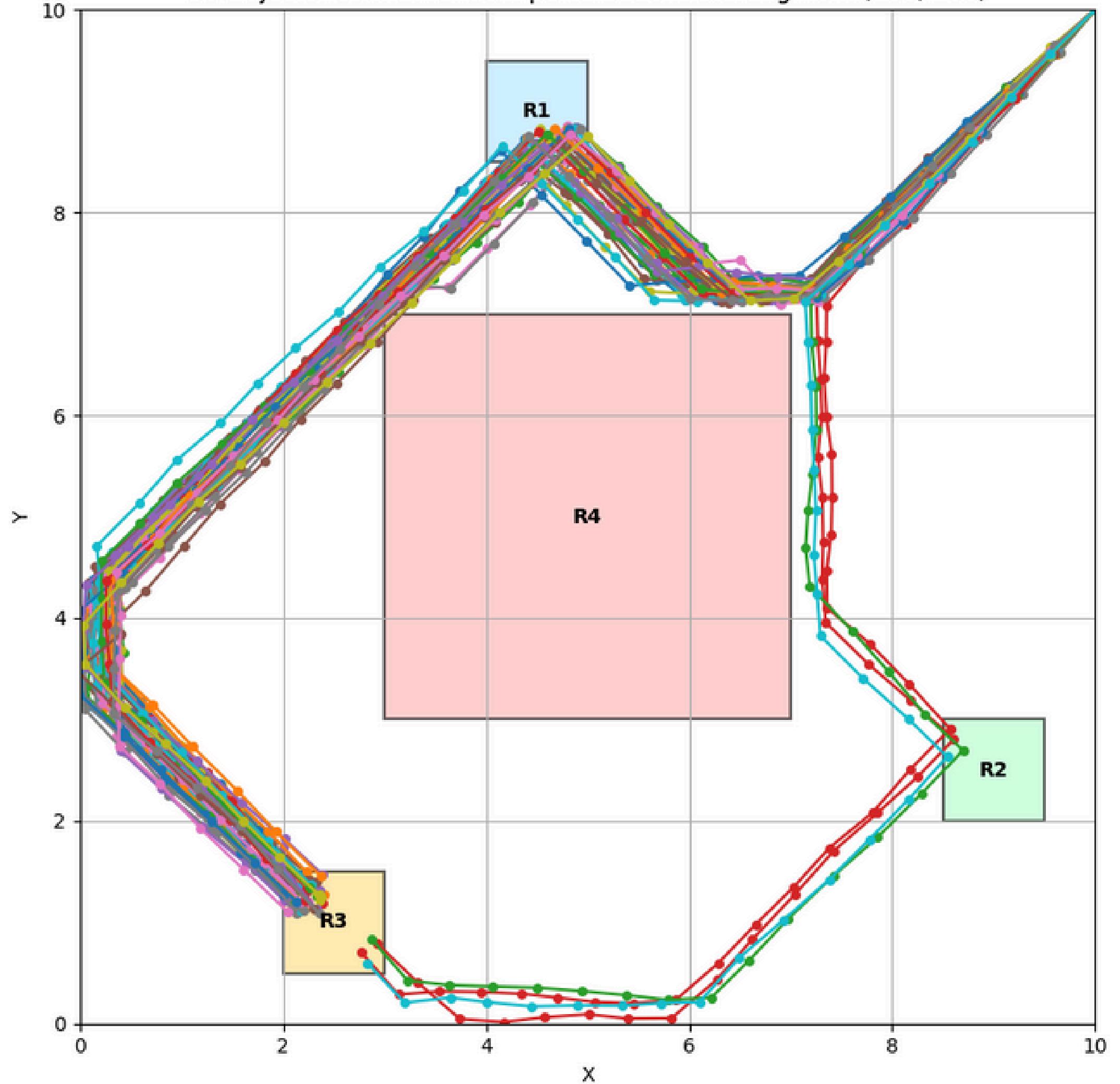
# SIMULATIONS

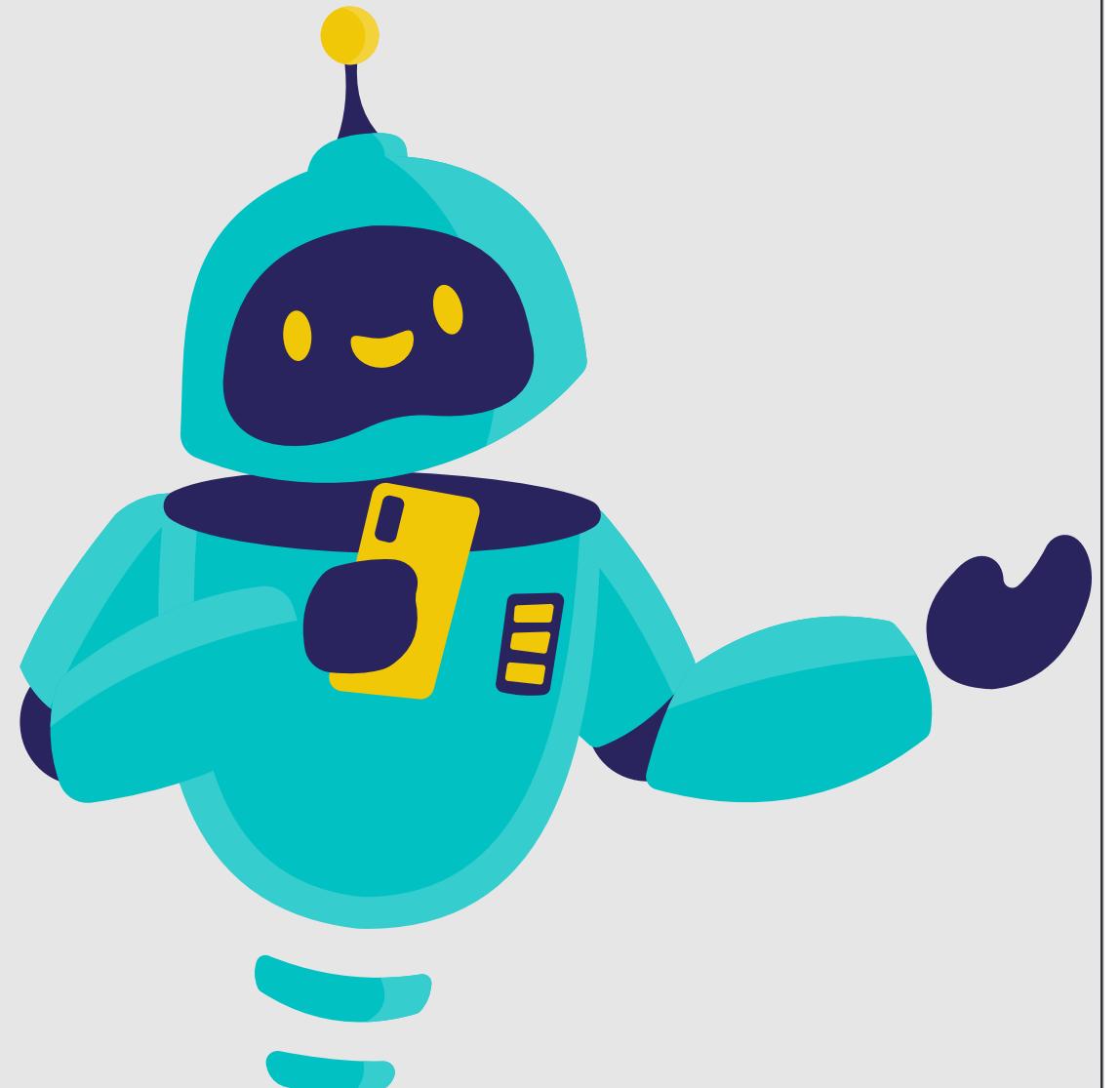
# 2D SIMULATIONS

LET'S SEE HOW IT  
LOOKS LIKE IN 2D,  
USING RANDOM  
PERTURBATION

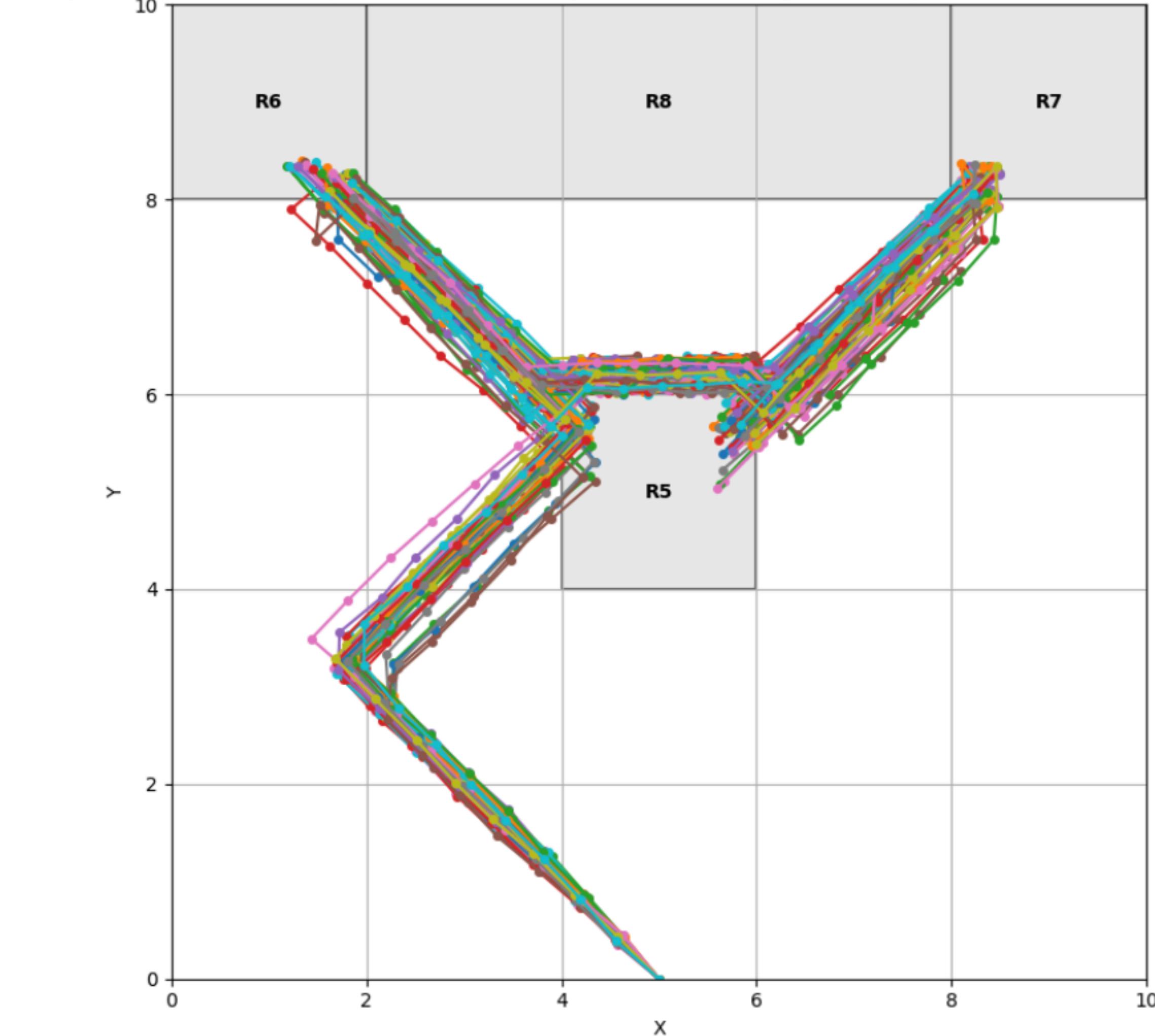


50 trajectories with random perturbations starting from (100, 100)



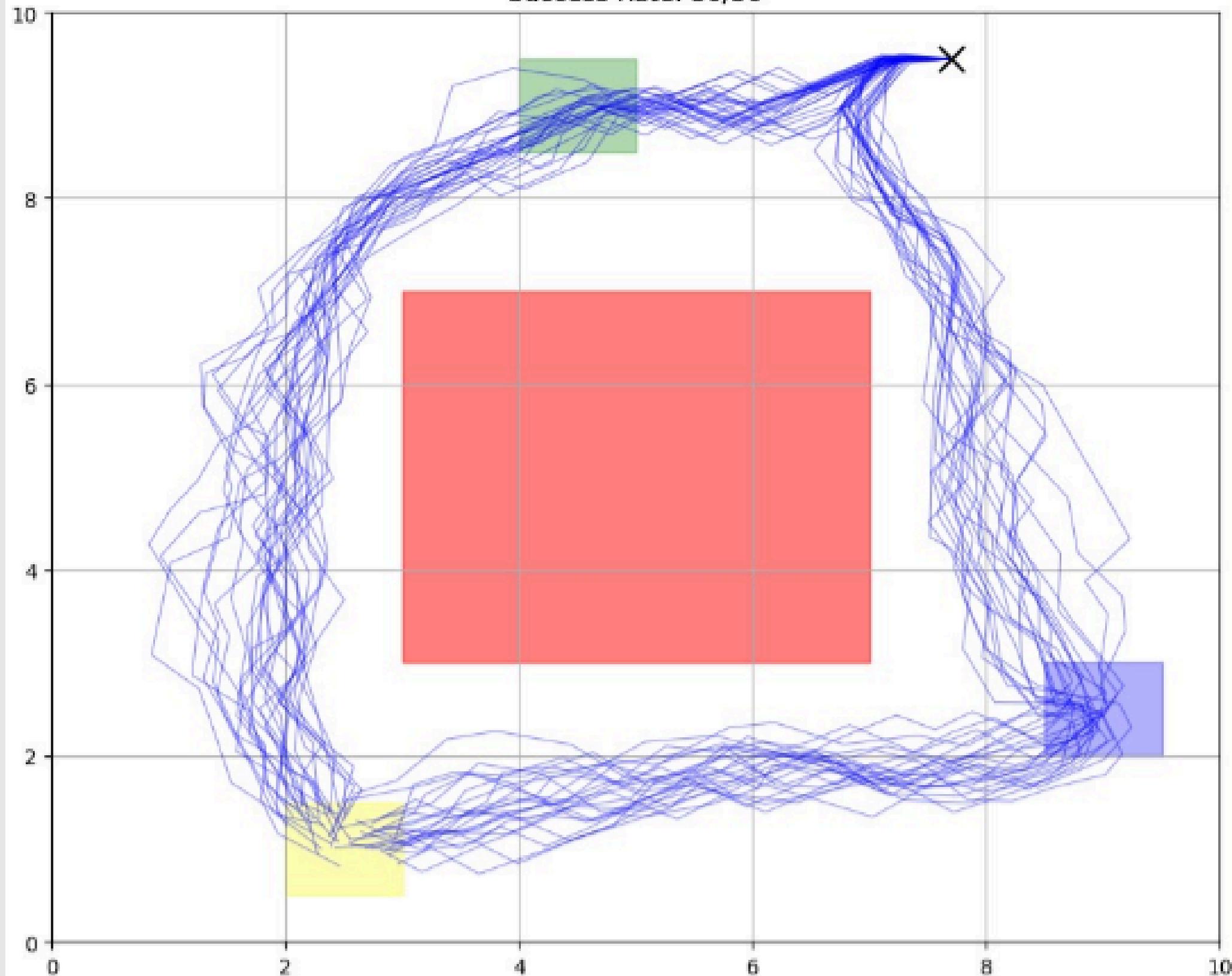


50 trajectories with random perturbations starting from (50, 0) using the 2nd specification (R5 -> R6 -> R7 -> R5)





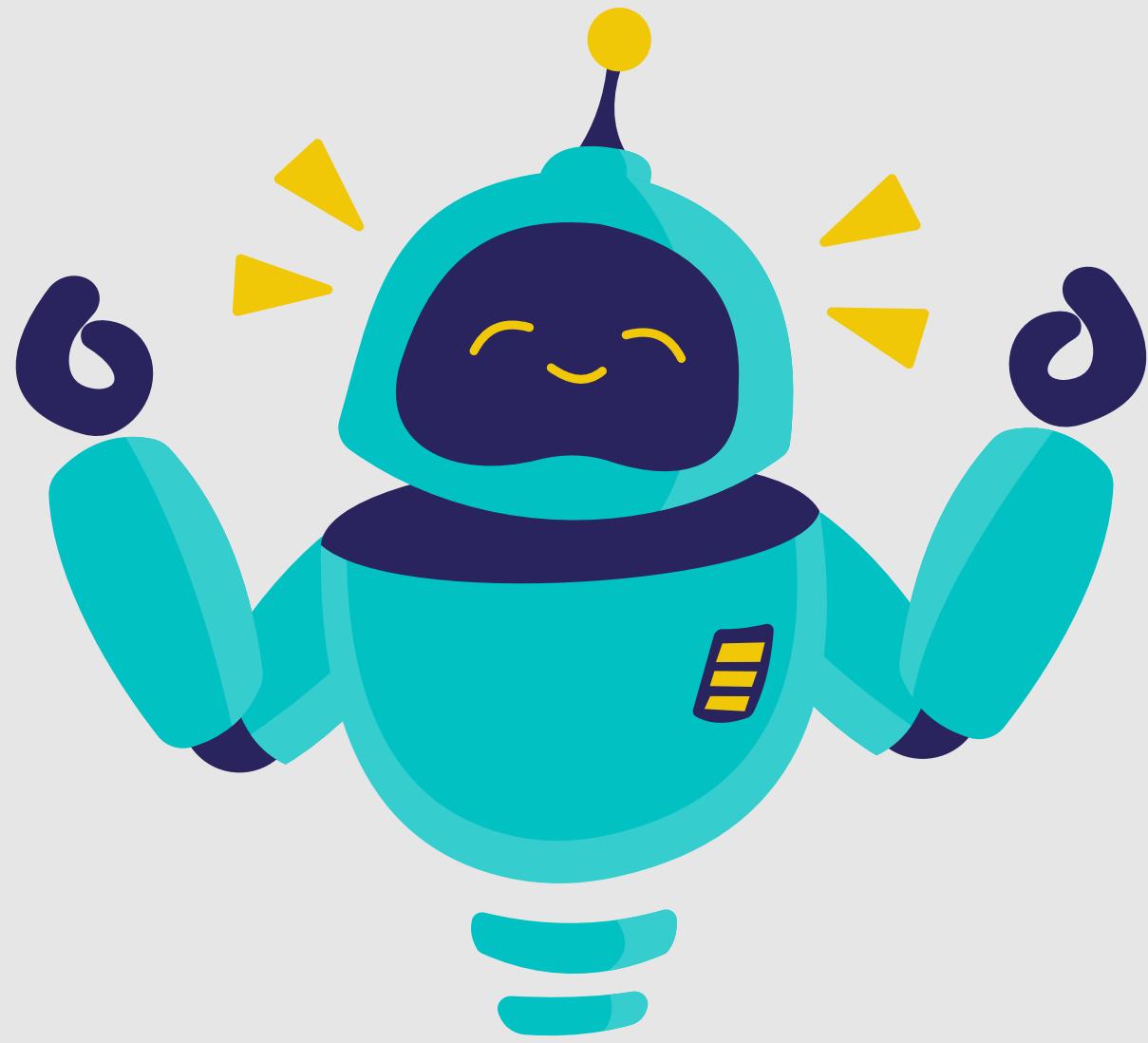
Robustness Test: 50 Runs  
Success Rate: 50/50



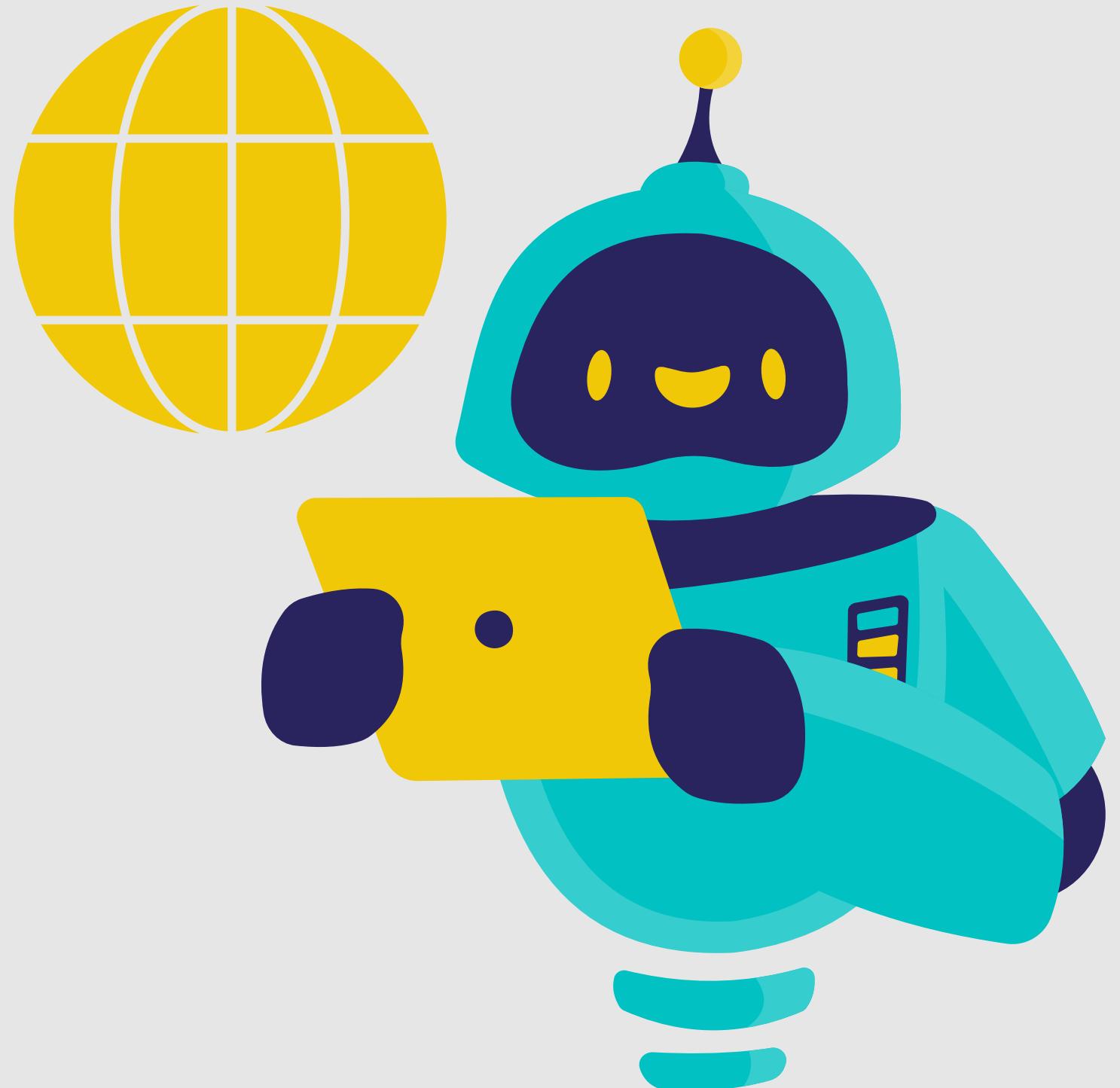
# PYBULLET SIMULATION



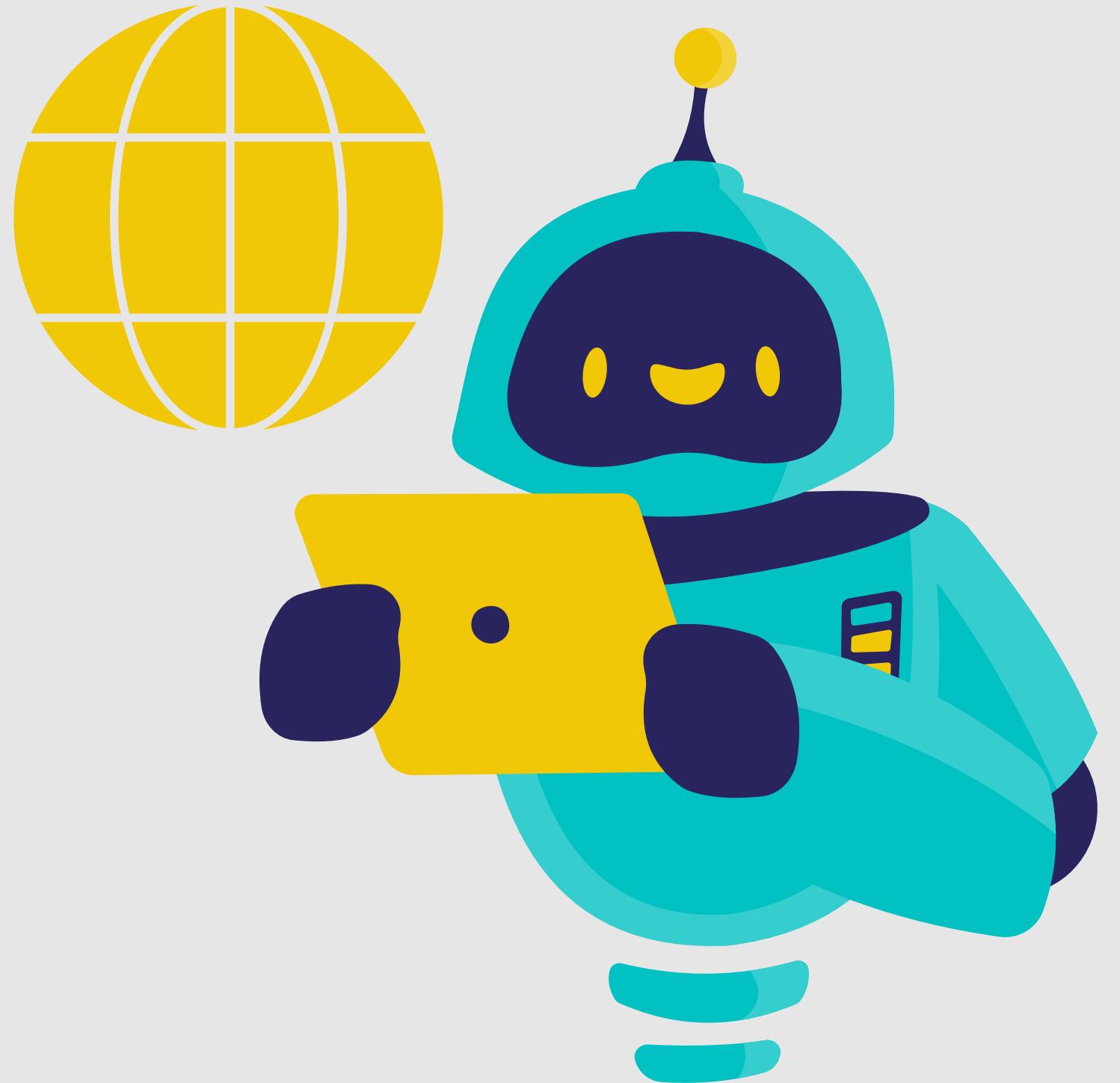
# 3D SIMULATION



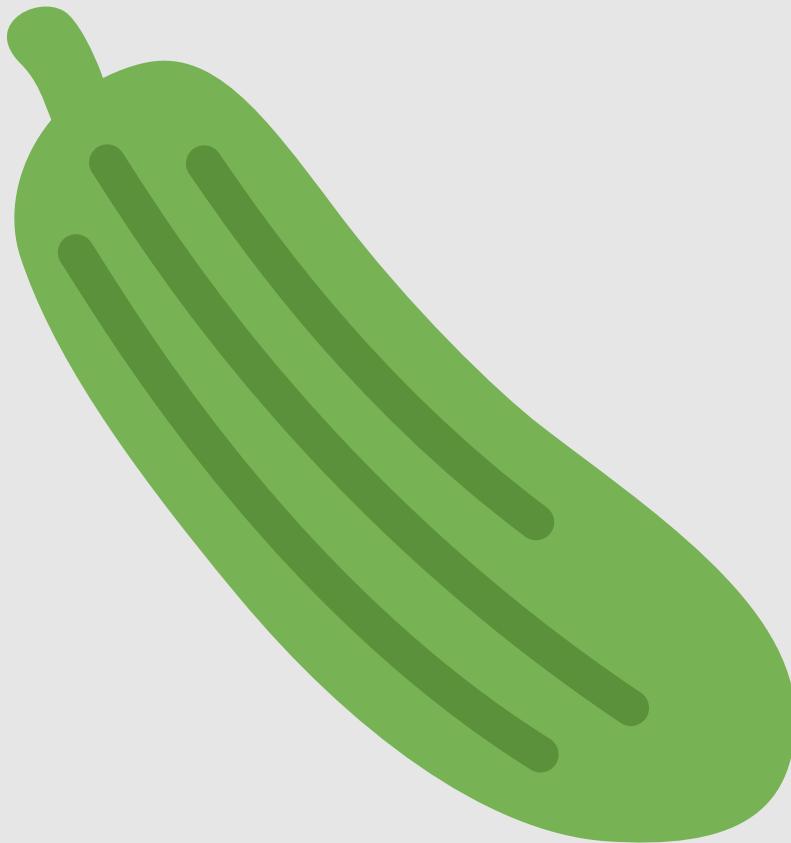
# COMMUNICATION



# COMMUNICATION



**PICKLES**



ANY  
QUESTIONS?

