



## LECTURE-5

# Linear Data Structures

Linked Lists

**CS202: Data Structures (Fall 2025)**

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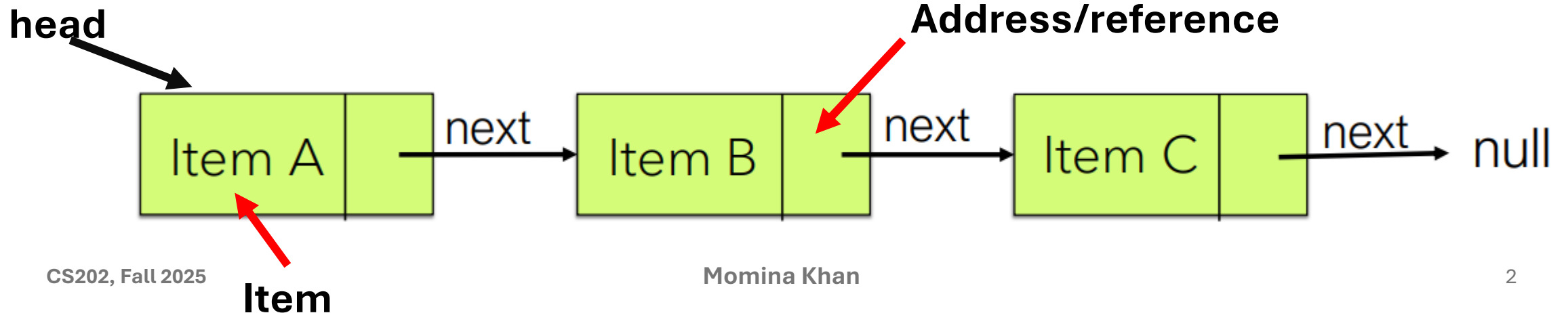
# Quick recap

Can we stretch an array? **NO**

- Dynamic Array (A **hack** around rigid sized array)
- Growing an array means
  - making a bigger
  - Copying contents
  - Discarding old smaller one

**Costly!**

- Linked List



# Agenda

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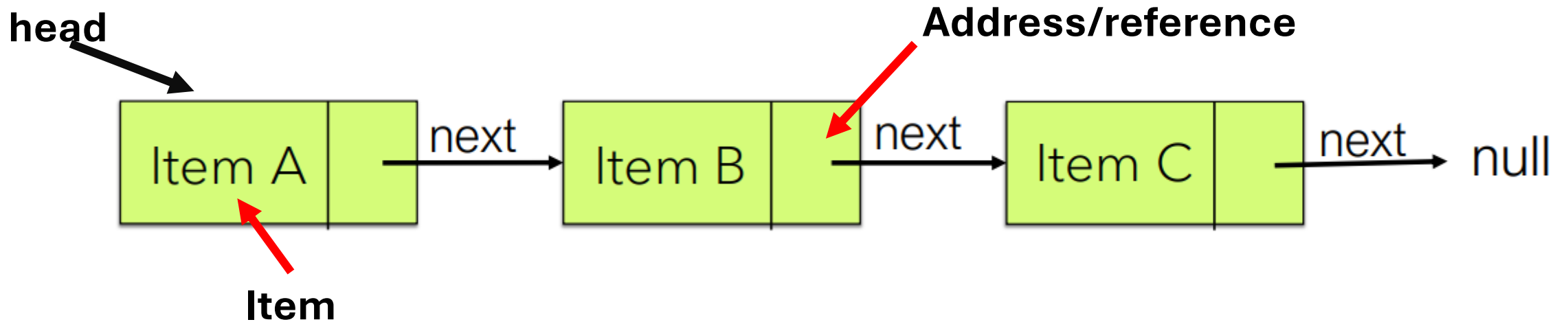
- Linked Lists
- How useful are Singly Linked Lists?
- How do Singly LLs compare with Doubly LLs?
- Circular Lists

# List (Linked List)

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Terms to remember:

- Node,
- link(next pointer),
- <key, value> pair



# Array vs Linked List

Array	Linked List
Contiguous Storage	Non-Contiguous
Fixed Size	Grows with insertions (shrinks too!)
Indexed access	Sequential access (follow the links)
Memory allocation can be compile time and <b>runtime</b>	Memory allocation is <b>runtime</b>

# List (Linked List)

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- Pros:
  - Size flexible; grows on demand
  - No need to allocate extra memory
  - Easier to add elements to the middle
- Cons:
  - Sequential Access (Has to traverse to the end)
  - Cache unfriendliness
  - Memory overhead (extra pointers!) **Caveat!**

# List ADT

```
template <typename T>
```

```
class List {
```

```
Public:
```

```
    virtual void insert( const T& element) = 0;
```

```
    virtual void delete( const T& element) = 0;
```

```
    virtual T find(const T& element ) = 0;
```

```
    virtual T get( int index) = 0;
```

```
    virtual bool isEmpty() = 0;
```

```
    virtual bool size() = 0;
```

**You can enhance this interface with more functions like:**

- **insertAtHead()**
- **deleteAtTail()**

# List (Linked List)

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**//You have access to the head pointer.**

```
void List<T>::insertAtHead(T object) {
```

```
// code comes here!
```

```
}
```

**Always draw a diagram!**

**Corner cases:** List empty, deleting a last node etc.

# List (Unsorted Linked List)

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find(k)	O(.)
insertAtStart(int k)	O(.)
insertAtLoc(int k, int loc)	O(.)
insertAtEnd(int k)	O(.)
deleteFromStart()	O(.)
deleteFromLoc(int loc)	O(.)

**Find the worst-case time complexity in terms of O(.)**

# List (Unsorted Linked List)

find(k)	$O(N)$
insertAtStart(int k)	$O(1)$
insertAtLoc(int k, int loc)	$O(N)$
insertAtEnd(int k)	$O(N)$
deleteFromStart()	$O(1)$
deleteFromLoc(int loc)	$O(N)$

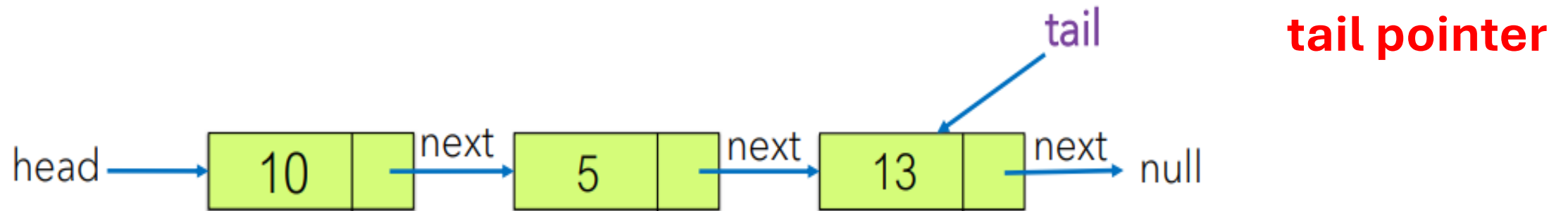
**Most operations are linear since they warrant a list traversal!**

# List (Linked List)

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Can we improve our insertAtEnd() time to something better than linear??

Introduce a second pointer to the end of the list called a ...



What would insertAtEnd(data) function look like with our new design?

# List (Linked List)

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**//You have access to the head pointer and tail pointer.**

```
void List<T>::insertAtEnd(T object) {
```

```
// code comes here!
```

```
}
```

**Always draw a diagram!**

**Corner cases:** List empty, deleting a last node etc.

# List ( **sorted** Linked List)

Which of the following describes the time complexity of the insert(k), delete(k), deleteEnd() and find(k) in a **sorted** singly linked list?

<b>find(k)</b>	<b>O(.)</b>
insert(k)	O(.)
delete(k)	O(.)
deleteEnd()	O(.)

Use the best possible algorithm (e.g., insertion at the beginning vs. end) to find the time complexity **in terms of O**.

# List ( **sorted** Linked List)

Which of the following describes the time complexity of the insert(k), delete(k), deleteEnd() and find(k) in a **sorted** singly linked list?

<b>find(k)</b>	<b>O(N)</b>
insert(k)	O(N)
delete(k)	O(N)
deleteEnd()	O(N)

Use the best possible algorithm (e.g., insertion at the beginning vs. end) to find the time complexity **in terms of O**.

# Singly Linked List – Some Observations

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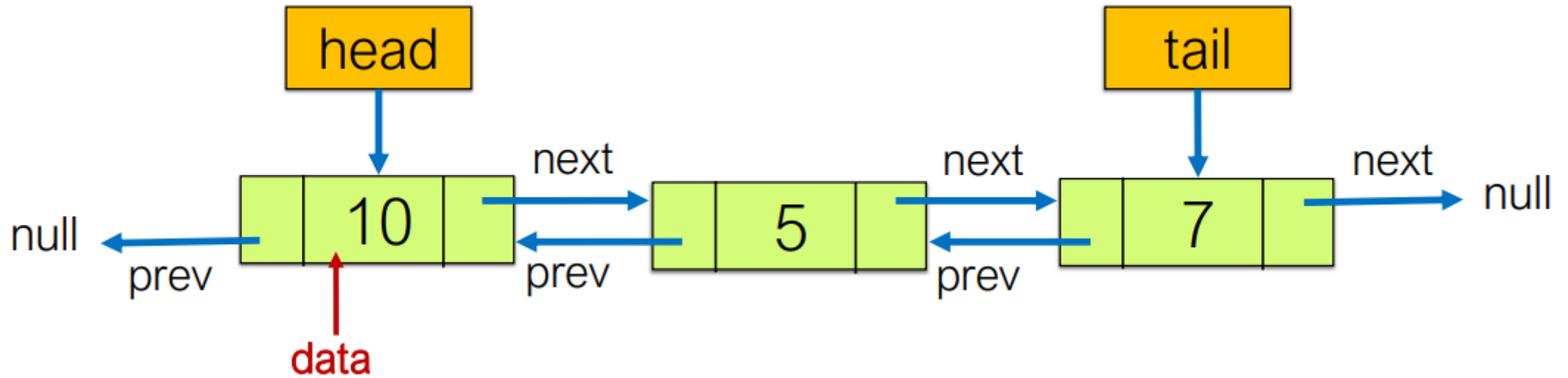
What is the time complexity (best case scenario)

- for inserting an element to the end?  $O(1)$
- for inserting an element at the second-last position?  $O(N)$
- for deleting an element from the end?  $O(1)$
- for deleting an element from the second-last position?  $O(N)$

Running Time is dominated by linear **uni-directional** traversal!

Can we do any better?

# Doubly Linked List



**Can you spot the differences between this and a singly linked list?**

# Doubly Linked List

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- How would the node structure change?
- What benefits are we deriving from this change?
- What remains the same? 😊

# List (Unsorted Doubly Linked List)

find(k)	$O(.)$
insertAtStart(int k)	$O(.)$
insertAtLoc(int k, int loc)	$O(.)$
insertAtEnd(int k)	$O(.)$
deleteFromStart()	$O(.)$
deleteFromLoc(int loc)	$O(.)$

**Find the worst-case time complexity in terms of  $O(.)$**

# List (Unsorted Doubly Linked List)

find(k)	$O(N)$
insertAtStart(int k)	$O(1)$
insertAtLoc(int k, int loc)	$O(N)$
insertAtEnd(int k)	$O(1)$
deleteFromStart()	$O(1)$
deleteFromLoc(int loc)	$O(N)$

**Some operations are still **linear** since they warrant a list traversal!**

# List (Sorted Doubly Linked List)

Which of the following describes the time complexity of the insert(k), delete(k), deleteEnd() and find(k) in a **sorted** doubly linked list?

find(k)	O(.)
insert(k)	O(.)
delete(k)	O(.)
deleteEnd()	O(.)

?

Use the best possible algorithm (e.g., insertion at the beginning vs. end) to find the time complexity **in terms of O**.

# List (Sorted Doubly Linked List)

Which of the following describes the time complexity of the insert(k), delete(k), deleteEnd() and find(k) in a **sorted** doubly linked list?

find(k)	O(N)
insert(k)	O(N)
delete(k)	O(N)
deleteEnd()	O(1)

**Some operations are still **linear** since they warrant a list traversal!**

# Doubly Linked List – Some Observations

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- What is the time complexity ( Big O) for
  - inserting an element to the end?  $O(1)$
  - inserting an element at the second-last position?  $O(1)$
  - deleting an element form the end?  $O(1)$
  - deleting an element from the second-last position?  $O(1)$
  - Searching an element?  $O(N)$
  - Inserting/deleting an element from kth location?  $O(N)$

**Worth noting that the last is  $O(N)$  since it involves searching first!**