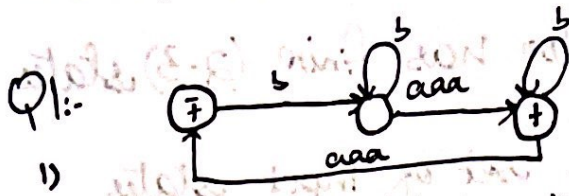


Theory of Automata

Name: Arneel Karter

Roll Number: 20F-0336

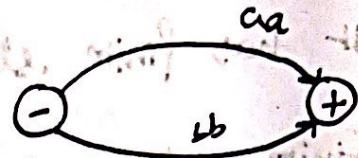
Section: 6C



It can read null as initial & final states are same and self b loop fulfill the many b's condition.

GTG $\emptyset \xrightarrow{b^*aaa b^*}$

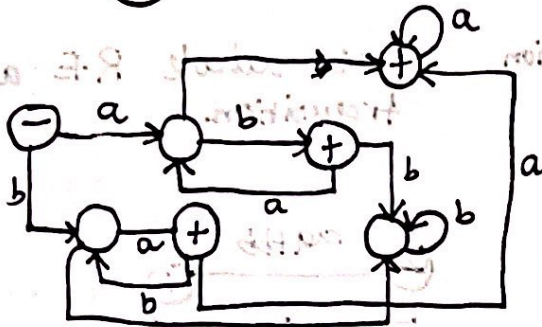
ii



output in clumps and not in single digit

GTG $\emptyset \xrightarrow{aa+bb}$

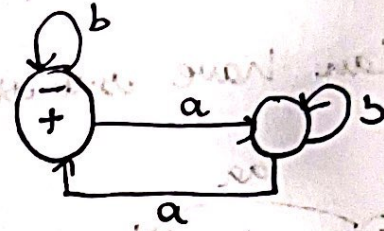
iii



It does not allow multiple reading of same input or does not generate clumps, they only occur at end.

GTG $\emptyset \xrightarrow{(ab)^* + (ba)^*(a^* + b^*)}$

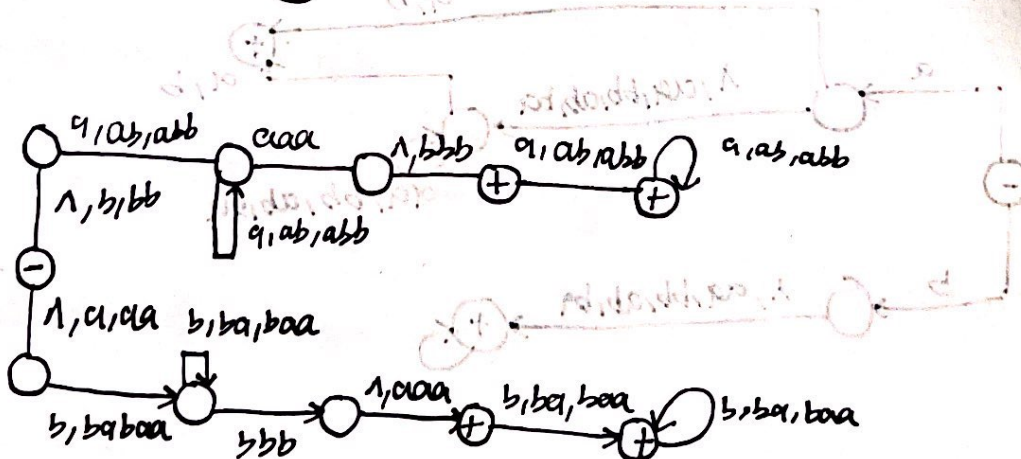
iv



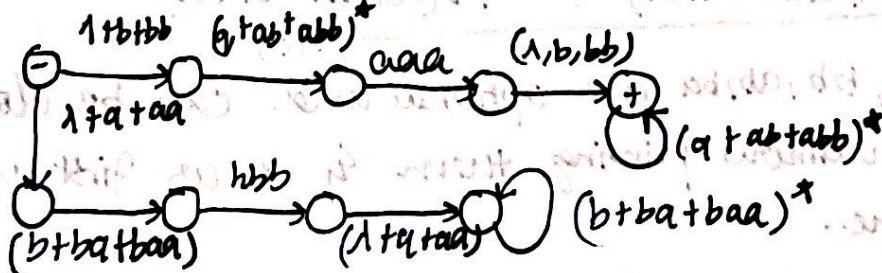
This will always generate even numbers of a as they are looped by a state and can't occur alone.

GTG $\emptyset \xrightarrow{(ab^*ab^* + b)^*}$

v



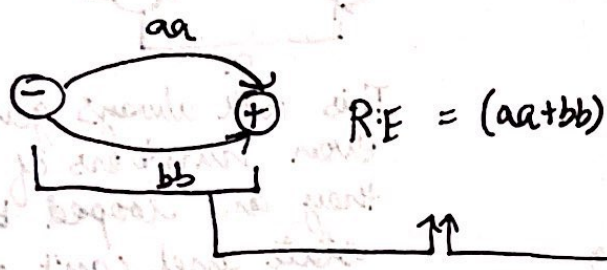
GTG.



Q No 2:

Transition Graph

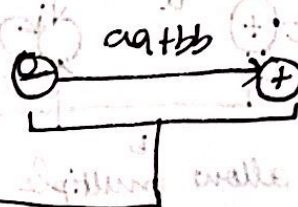
- > It has finite states
- > Finite set of inputs
- > States can be null
- > Null is an input
- > Can have substring transition



General Transition Graph

- It also has finite (2-3) states
- Finite set of input states
- States can't be null.
- Null is an input

has whole R.E as transition.

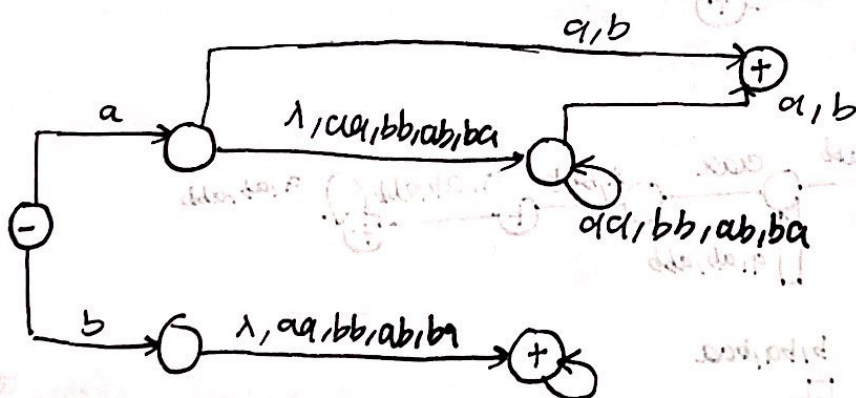


- have symbols like ' , '

uses only $\pi^+ \pi^-$ or $\pi^+ \pi^-$

Q: no: 3

①

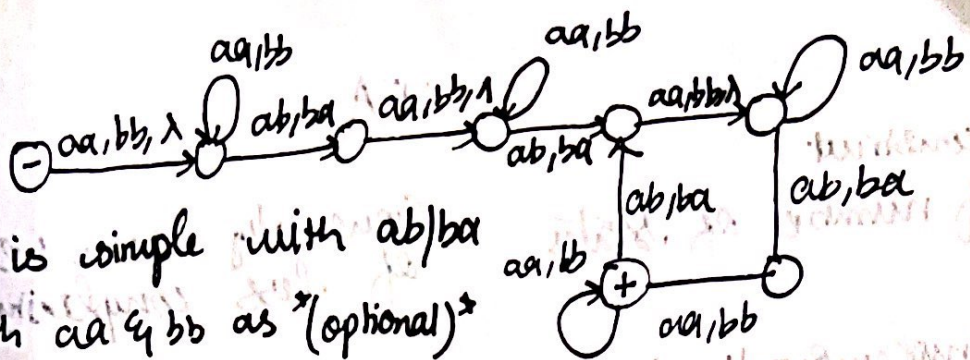


GTC

$$\textcircled{-} a(a^2+bb+ba+ab)^2(ab) + b(a^2+bb+ab+ba)^2 \rightarrow \textcircled{+}$$

aa, bb, ab, ba are optional and can be done with or without printing them & a as first will never be alone.

① T.C

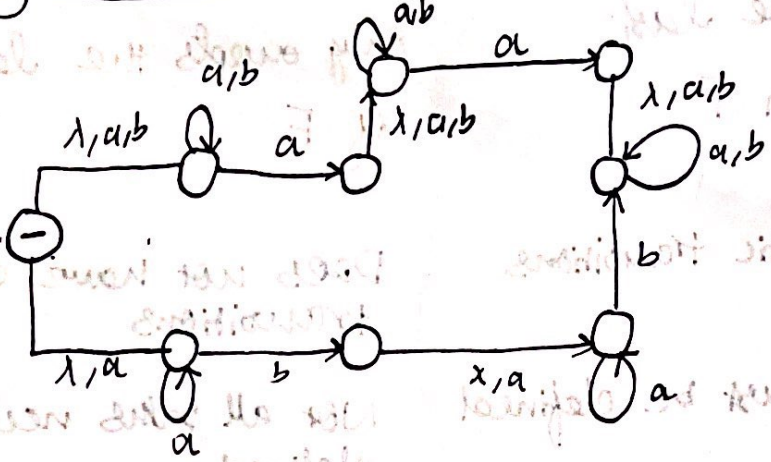


path is simple with ab/ba with aa & bb as $^*(optional)^*$

QTC

$$-(aa+bb)^*(ab+ba)^*(aa+bb)^*(ab+ba)^*(aa+bb)^*(ab+ba)^*(aa+bb)^* \rightarrow +$$

③ T.C



QTC $(a+b)^*a(a+b)^*a(a+b)^* + a^*ba^*b$

$$-(a+b)^*a(a+b)^*a(a+b)^* + a^*ba^*b \rightarrow +$$

This will always produce either two a's or two b's

④ Already done on paper

QTC $-(a+b)^+ \rightarrow +$

⑤ Already drawn on paper

QTC $a(a+b)^* + a(a+b)^*b + \lambda$
 $1(1+0)^* + 1(1+0)^*0 + \epsilon$

Q4:-

DFA

- Harder to construct because of number of states
- Practical implementation is possible
- Accepts input if the last visited state is in F
- Have deterministic transitions
- Every way must be defined

NFA

Generally easier because of less complexity.

limited but possible practical implementation.

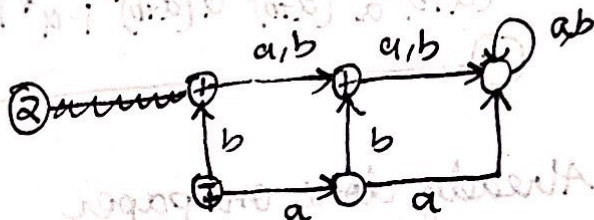
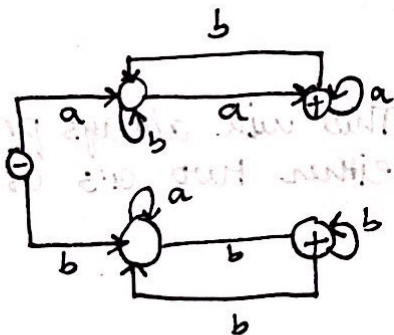
Only accepts the last state of F

Does not have deterministic transitions

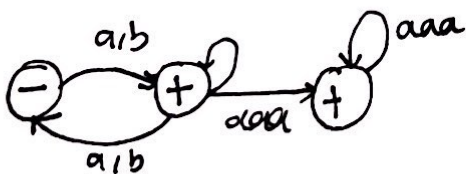
Not all paths need to be defined.

Q5:-

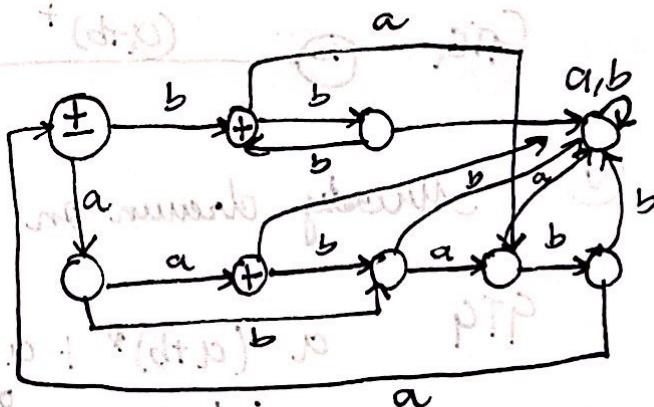
①



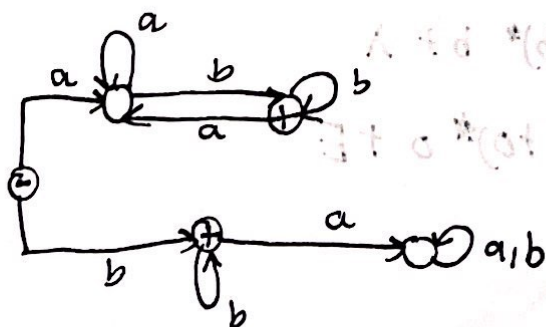
③



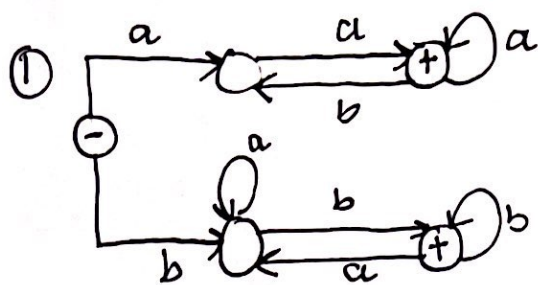
④



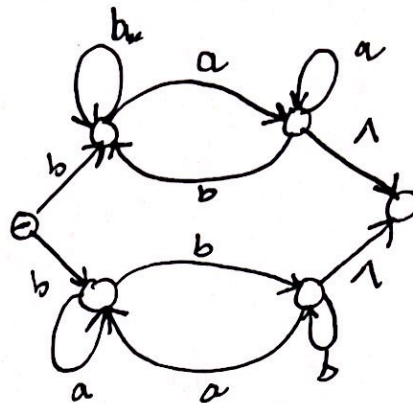
⑤



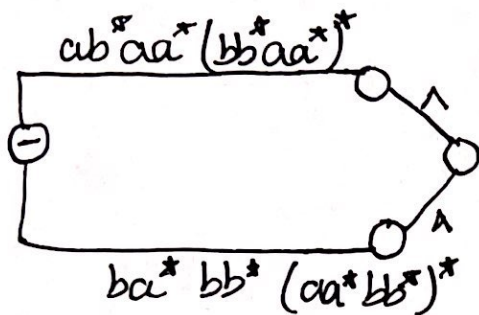
Question 6



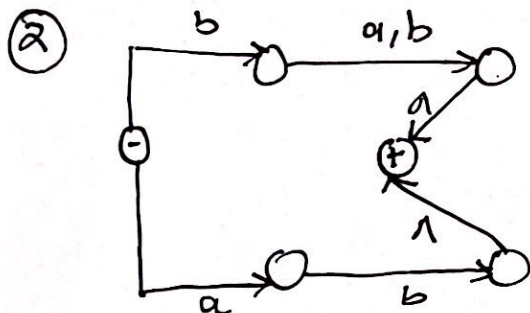
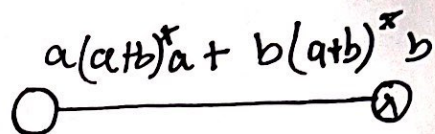
remove final
→



removing first 2



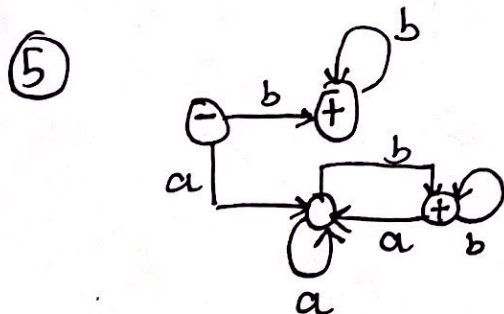
removing
remaining



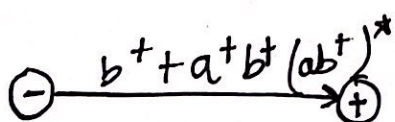
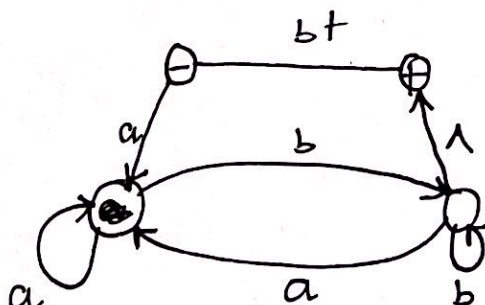
→



$$= (a+tb)((a+tb)(a+tb))^*$$



→



$$= b + a+tb*(ab)^*$$