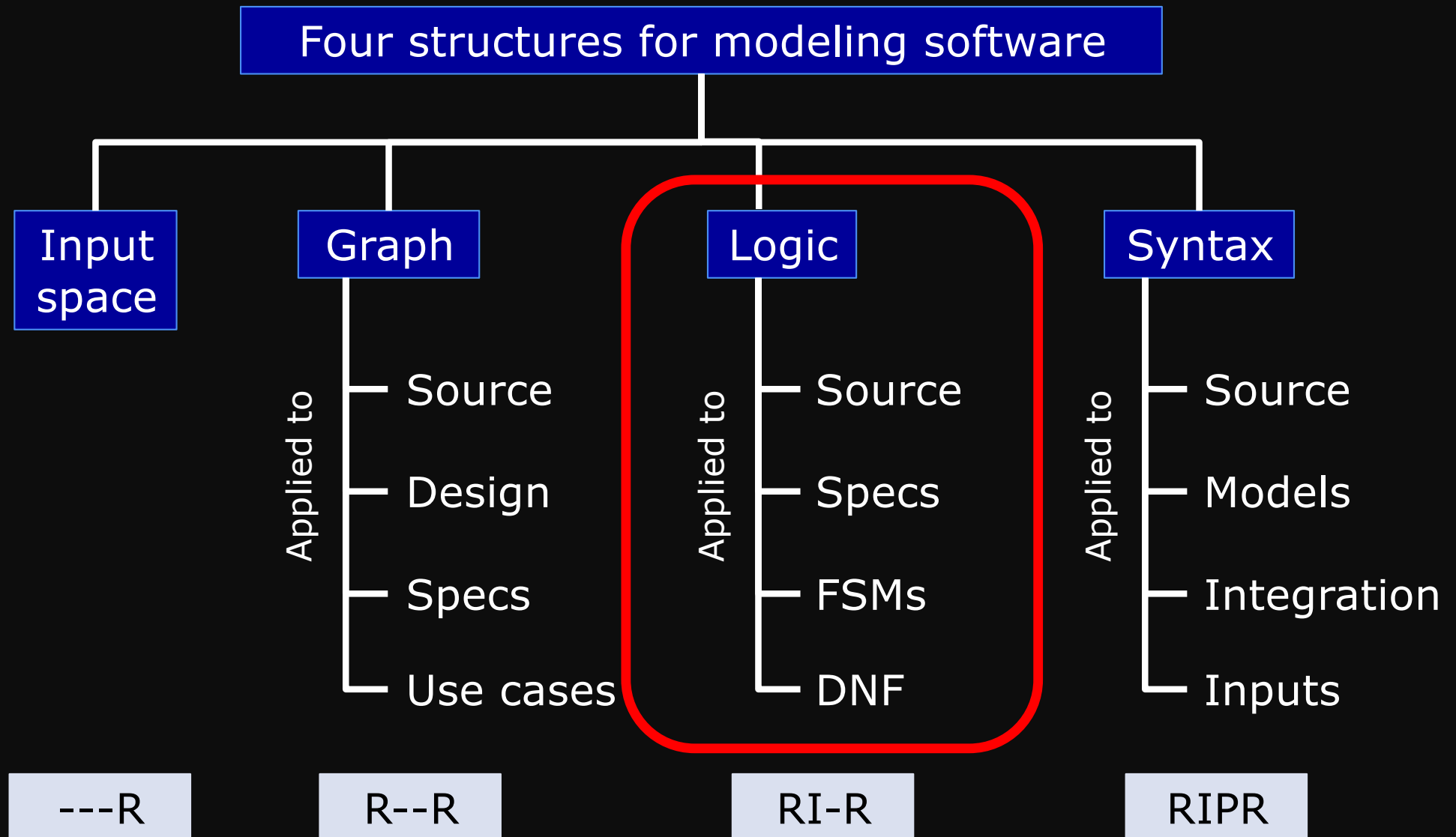


Logic Coverage

CS 3250 Software Testing

[Ammann and Offutt, “Introduction to Software Testing,” Ch. 8]

Structures for Criteria-Based Testing



Overview

- Logic coverage ensures that tests not only **reach** certain locations, but the internal state is **infected** by trying multiple combinations of truth assignments to the expressions
- Covering logic expressions is required by the US Federal Aviation Administration for safety critical avionics software
- Logical expressions can come from many sources
 - Decisions in programs
 - FSMs and statecharts
 - Requirements
 - SQL queries
- Tests are intended to choose some subset of the total number of truth assignments to the expressions

Logic Predicates and Clauses

- **Predicate**: An expression that evaluates to a Boolean value
 - May contain
 - Boolean variable
 - Non-Boolean variables that contain $>$, $<$, $==$, $>=$, $<=$, $!=$
 - Boolean function calls
 - Created by the logical operators

\neg	<i>negation</i> operator
\wedge	<i>and</i> operator
\vee	<i>or</i> operator
\rightarrow	<i>implication</i> operator
\oplus	<i>exclusive or</i> operator
\leftrightarrow	<i>equivalence</i> operator

- **Clause**: A predicate with no logical operators

Let's Refresh Our Memory

Row #	a	b	$\neg a$	$a \wedge b$	$a \vee b$	$a \rightarrow b$	$a \oplus b$	$a \leftrightarrow b$
1	T	T						
2	T							
3		T						
4								

Blank indicates F

Example

Logically equivalent

$$\underline{(a = b)} \vee (\underline{C} \wedge \underline{f(x)})$$

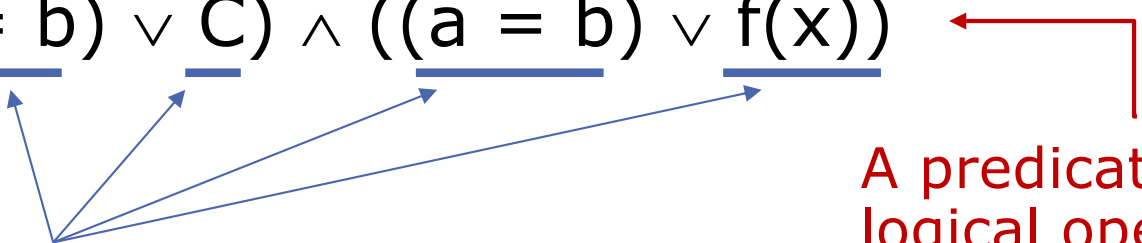

Three clauses

A relational expression ($a = b$)

A boolean variable C

A boolean-valued function $p(x)$

A predicate with
logical operators

$$(\underline{(a = b)} \vee \underline{C}) \wedge (\underline{(a = b)} \vee \underline{f(x)})$$


Three clauses

A relational expression ($a = b$)

A boolean variable C

A boolean-valued function $p(x)$

A predicate with
logical operators

Note on Predicates

- Most predicates have few clauses
- Sources of predicates
 - **Decisions** in program source code

```
public boolean isSatisfactory()
{
    if ((good && fast) || (good && cheap) || (fast && cheap))
        return true;
    else
        return false;
}
```

$(\text{good} \wedge \text{fast}) \vee (\text{good} \wedge \text{cheap}) \vee (\text{fast} \wedge \text{cheap})$

- **Guards** in finite state machines

button2 = true (when gear = park)

$(\text{gear} = \text{park}) \wedge (\text{button2} = \text{true})$

- **Precondition** in specifications

pre: stack not full AND object reference parameter not null

$\neg \text{stackFull()} \wedge (\text{newObj} \neq \text{null})$

Note on Predicates

- Be careful when translating from English

“I am interested in CS6501 **and** CS4501”



Which one ?

(course = CS6501) **AND** (course = CS4501)
(course = CS6501) **OR** (course = CS4501)



From a study of 63 open source programs (>400,000 predicates), most predicates have few clauses [Ammann and Offutt]

- 88.5% have 1 clauses
- 9.5% have 2 clauses
- 1.35% have 3 clauses
- Only .65% have 4 or more

Try to keep the predicate simple and short.
How? Refactor it.

Short Circuit Evaluation

- Impacted by the order of operation
- Evaluate an expression or predicate **until** an outcome is known

$$\underline{((a = b) \vee C)} \boxed{\wedge} \underline{f(x)}$$

If $f(x)$ is evaluated to T, we evaluate $(a = b) \vee C$ which can be T or F

If $f(x)$ is evaluated to F, we stop.
The outcome of the predicate is F

Short Circuit Evaluation

Boolean variable

Method call – something
can change

```
if (isHungry && (time == f(time)))  
{  
  
}
```

If isHungry is evaluated to T, we evaluate
(time == f(time)) which can be T or F

If isHungry is evaluated to F, we stop. The
outcome of the predicate is F

Stop evaluating the predicate when we know the outcome

Logic Coverage Criteria

- We use predicates in testing as follows:
 - Developing a model of the software as one or more predicates
 - Requiring tests to satisfy some combination of clauses
- Abbreviations:
 - P is the set of predicates
 - p is a single predicate in P
 - C is the set of clauses in P
 - C_p is the set of clauses in predicate p
 - c is a single clause in C

Predicate Coverage (PC)

- For each p in P , TR contains two requirements:

- p evaluates to true
- p evaluates to false

“Decision coverage”

$$p = ((a = b) \vee C) \wedge f(x)$$

Need 2 test cases to satisfy PC

a	b	C	f(x)	p
3	T	3	T	T
4	F	3	T	F

- PC does **not** evaluate all the clauses, especially in the presence of short circuit evaluation

Activity #1 (PC)

$$p = a \vee (b \wedge c)$$

Row#	a	b	c	p
1	T	T	T	
2	T	T		
3	T		T	
4	T			
5		T	T	
6		T		
7			T	
8				

List test requirements that satisfy predicate coverage (PC)

Blank indicates F

Activity #2 (PC)

$$p = a \wedge (b \vee c)$$

Row#	a	b	c	p
1	T	T	T	
2	T	T		
3	T		T	
4	T			
5		T	T	
6		T		
7			T	
8				

List test requirements that satisfy predicate coverage (PC)

Blank indicates F

Clause Coverage (CC)

- For each c in C , TR contains two requirements:

- c evaluates to true
- c evaluates to false

“Condition coverage”

$$p = ((a = b) \vee C) \wedge f(x)$$

$(a = b)$ evaluates to **T, F**

C evaluates to **T, F**

$f(x)$ evaluates to **T, F**

a	b	C	f(x)	p
3	T	3	F	F
4	F	3	T	F

Need 2 test cases to satisfy CC

- CC does **not** always ensure PC
- The simplest solution is to test **all combinations**

Activity #1 (CC)

$$p = a \vee (b \wedge c)$$

Row#	a	b	c	p
1	T	T	T	
2	T	T		
3	T		T	
4	T			
5		T	T	
6		T		
7			T	
8				

List test requirements that satisfy clause coverage (CC)

Blank indicates F

Activity #2 (CC)

$$p = a \wedge (b \vee c)$$

Row#	a	b	c	p
1	T	T	T	
2	T	T		
3	T		T	
4	T			
5		T	T	
6		T		
7			T	
8				

List test requirements that satisfy clause coverage (CC)

Blank indicates F

Combinatorial Coverage (CoC)

- Evaluate all possible combination of truth values

“Multiple Condition coverage”

$$p = ((a = b) \vee C) \wedge f(x)$$

a	b	C	f(x)	p
3	T	3	T	T
3	T	3	F	F
3	T	3	F	T
3	T	3	F	F
4	F	3	T	T
4	F	3	F	F
4	F	3	T	F
4	F	3	F	F

Need 2^N test cases to satisfy CoC, where N = number of clauses

Note on CoC

- Coc is simple and comprehensive
- But quite expensive
- 2^N tests, where N is the number of clauses
 - Impractical for predicates with more than 3 or 4 clauses
- The literature has lots of suggestions – some confusing
- The general idea is simple:

Test each clause that makes a big difference ...
"active clause"

Revisit CoC Example

Which clause makes a big difference

$$p = ((a = b) \vee C) \wedge f(x)$$

a	b	C	f(x)	p
3	T	3	T	T
3	T	3	F	F
3	T	F	T	T
3	T	F	F	F
4	F	3	T	T
4	F	3	F	F
4	F	F	T	F
4	F	F	F	F

Active Clauses

- To really test the results of a clause, the clause should be the **determining factor** in the value of the predicate
- **Determination**
 - A clause c_i in predicate p , called the major clause, **determines** p if and only if the values of the remaining minor clauses c_j are such that changing c_i changes the value p
 - That is:
 - **Major clause** – the clause (being considered) that determine the predicate
 - **Minor clause** – all other clauses in the predicate
- This is considered to **make the clause active**

Determination

- **Goal**: Find tests for each clause when the clause determines the value of the predicate
- **Determination**: the conditions under which a clause solely determines the outcome of a predicate
 - Given a **major clause** c_i in a predicate p , c_i determines p if the **minor clauses** $c_j \neq c_i (j \neq i)$
 - Major clause – “**active clause**” – controls the behavior
- Consider $p = a \vee b$
 - If $b = \text{true}$, the value of a does not matter
 - If $b = \text{false}$, the value of a is the determining factor in the value of the predicate

Revisit Coc Example (again)

Which clause determines the predicate

$$p = ((a = b) \vee C) \wedge f(x)$$

a	b	C	f(x)	p
3	T	3	T	T
3	T	3	F	F
3	T	F	T	T
3	T	F	F	F
4	F	3	T	T
4	F	3	F	F
4	F	F	T	F
4	F	F	F	F

f(x) determines
the predicate –
but when ??

Deriving Determination Predicates, Using Mathematical Approach

Deriving (Mathematical Approach) Determination Predicates

$$p = a \wedge (b \vee c)$$

$$\begin{aligned} p_a &= p_{a=\text{true}} \oplus p_{a=\text{false}} \\ &= (\text{true} \wedge (b \vee c)) \oplus (\text{false} \wedge (b \vee c)) \\ &= (b \vee c) \oplus \text{false} \\ &= b \vee c \end{aligned}$$

$$\begin{aligned} p_b &= p_{b=\text{true}} \oplus p_{b=\text{false}} \\ &= (a \wedge (\text{true} \vee c)) \oplus (a \wedge (\text{false} \vee c)) \\ &= (a \wedge \text{true}) \oplus (a \wedge c) \\ &= a \oplus (a \wedge c) \\ &= a \wedge \neg c \end{aligned}$$

$$\begin{aligned} p_c &= p_{c=\text{true}} \oplus p_{c=\text{false}} \\ &= (a \wedge (b \vee \text{true})) \oplus (a \wedge (b \vee \text{false})) \\ &= (a \wedge \text{true}) \oplus (a \wedge b) \\ &= a \oplus (a \wedge b) \\ &= a \wedge \neg b \end{aligned}$$

Deriving (Mathematical Approach) Determination Predicates

$$p = a \wedge (b \vee c)$$

Major clause: **a**

$$\begin{aligned} p_a &= p_{a=\text{true}} \oplus p_{a=\text{false}} \\ &= (\text{true} \wedge (b \vee c)) \oplus (\text{false} \wedge (b \vee c)) \\ &= (b \vee c) \oplus \text{false} \\ &= b \vee c \end{aligned}$$

row	a	b	c	p	p _a	p _b	p _c
1	T	T	T	T	T		
2	T	T		T	T		
3	T		T	T	T		
4	T						
5		T	T		T		
6		T			T		
7			T		T		
8							

(Fill in a table to make it easy to read)

Blank indicates F

Deriving (Mathematical Approach) Determination Predicates

$$p = a \wedge (b \vee c)$$

Major clause: **b**

$$\begin{aligned} p_b &= p_{b=\text{true}} \oplus p_{b=\text{false}} \\ &= (a \wedge (\text{true} \vee c)) \oplus (a \wedge (\text{false} \vee c)) \\ &= (a \wedge \text{true}) \oplus (a \wedge c) \\ &= a \oplus (a \wedge c) \\ &= a \wedge \neg c \end{aligned}$$

row	a	b	c	p	p _a	p _b	p _c
1	T	T	T	T	T		
2	T	T		T	T	T	
3	T		T	T	T		
4	T					T	
5		T	T		T		
6		T			T		
7			T		T		
8							

(Fill in a table to make it easy to read)

Blank indicates F

Deriving (Mathematical Approach) Determination Predicates

$$p = a \wedge (b \vee c)$$

Major clause: **c**

$$\begin{aligned} p_c &= p_{c=\text{true}} \oplus p_{c=\text{false}} \\ &= (a \wedge (b \vee \text{true})) \oplus (a \wedge (b \vee \text{false})) \\ &= (a \wedge \text{true}) \oplus (a \wedge b) \\ &= a \oplus (a \wedge b) \\ &= a \wedge \neg b \end{aligned}$$

row	a	b	c	p	p _a	p _b	p _c
1	T	T	T	T	T		
2	T	T		T	T	T	
3	T		T	T	T		T
4	T					T	T
5		T	T		T		
6		T			T		
7			T		T		
8							

(Fill in a table to make it easy to read)

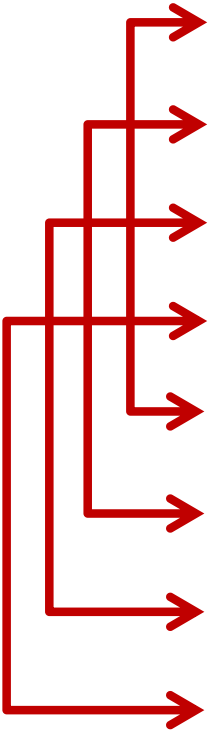
Blank indicates F

Deriving Determination Predicates, Using Tabular Approach

Identifying Determination Using Truth Table

$$p = a \wedge (b \vee c)$$

Major clause: **a**



row	a	b	c	p	p _a	p _b	p _c
1	T	T	T	T	T		
2	T	T		T	T		
3	T		T	T	T		
4	T						
5		T	T		T		
6		T			T		
7			T		T		
8							

Blank indicates F

Identifying Determination Using Truth Table

$$p = a \wedge (b \vee c)$$

Major clause: **b**

row	a	b	c	p	p _a	p _b	p _c
1	T	T	T	T	T		
2	T	T		T	T	T	
3	T		T	T	T		
4	T					T	
5		T	T		T		
6		T			T		
7			T		T		
8							

Blank indicates F

Identifying Determination Using Truth Table

$$p = a \wedge (b \vee c)$$

Major clause: **c**

row	a	b	c	p	p _a	p _b	p _c
1	T	T	T	T	T		
2	T	T		T	T	T	
3	T		T	T	T		T
4	T					T	T
5		T	T		T		
6		T			T		
7			T		T		
8							

Blank indicates F

What's next?

- Use determination
- Apply logic coverage criteria to derive test requirements and design test cases
 - Active Clause Coverage (ACC)
 - General Active Clause Coverage (GACC)
 - Correlated Active Clause Coverage (CACC)
 - Restricted Active Clause Coverage (RACC)