

Circus in Isabelle/UTP

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1 Introduction

This document contains a mechanisation in Isabelle/UTP [1] of Circus [2].

2 Circus Trace Merge

```
theory utp-circus-traces
  imports UTP-Stateful-Failures.utp-sf-rdes
begin
```

2.1 Function Definition

```

fun tr-par ::
  '∅ set ⇒ '∅ list ⇒ '∅ list ⇒ '∅ list set where
tr-par cs [] = {} |
tr-par cs (e # t) [] = (if e ∈ cs then {} else {[e]} ∩ (tr-par cs t [])) |
tr-par cs [] (e # t) = (if e ∈ cs then {} else {[e]} ∩ (tr-par cs [] t)) |
tr-par cs (e1 # t1) (e2 # t2) =
  (if e1 = e2
   then
    if e1 ∈ cs
    then {[e1]} ∩ (tr-par cs t1 t2)
    else
      ({[e1]} ∩ (tr-par cs t1 (e2 # t2))) ∪
      ({[e2]} ∩ (tr-par cs (e1 # t1) t2))
   else
    if e1 ∈ cs then
      if e2 ∈ cs then {}
      else
        {[e2]} ∩ (tr-par cs (e1 # t1) t2)
    else
      if e2 ∈ cs then
        {[e1]} ∩ (tr-par cs t1 (e2 # t2))
      else
        {[e1]} ∩ (tr-par cs t1 (e2 # t2)) ∪
        {[e2]} ∩ (tr-par cs (e1 # t1) t2)
  )

```

abbreviation *tr-inter* :: '∅ list ⇒ '∅ list ⇒ '∅ list set (**infixr** |||_t 100) **where**
x |||_t *y* ≡ *tr-par* {} *x y*

2.2 Lifted Trace Merge

```

syntax -utr-par ::
  logic ⇒ logic ⇒ logic ⇒ logic ((- ★-/ -) [100, 0, 101] 100)

```

The function *trop* is used to lift ternary operators.

translations

t1 ★_{cs} *t2* == (CONST *bop*) (CONST *tr-par* cs) *t1 t2*

2.3 Trace Merge Lemmas

```

lemma tr-par-empty:
tr-par cs t1 [] = {takeWhile (λx. x ∉ cs) t1}
tr-par cs [] t2 = {takeWhile (λx. x ∉ cs) t2}
— Subgoal 1
apply (induct t1; simp)
— Subgoal 2
apply (induct t2; simp)
done

```

```

lemma tr-par-sym:
tr-par cs t1 t2 = tr-par cs t2 t1
apply (induct t1 arbitrary: t2)
— Subgoal 1
apply (simp add: tr-par-empty)
— Subgoal 2

```

apply (*induct-tac* *t2*)

— Subgoal 2.1

apply (*clarsimp*)

— Subgoal 2.2

apply (*clarsimp*)

apply (*blast*)

done

lemma *tr-inter-sym*: $x \parallel_t y = y \parallel_t x$

by (*simp* *add*: *tr-par-sym*)

lemma *trace-merge-nil* [*simp*]: $x \star_{\{\}} \langle \rangle = \{x\}_u$

by (*pred-auto*, *simp-all* *add*: *tr-par-empty*, *metis* *takeWhile-eq-all-conv*)

lemma *trace-merge-empty* [*simp*]:

$(\langle \rangle \star_{cs} \langle \rangle) = \{\langle \rangle\}_u$

by (*rel-auto*)

lemma *trace-merge-single-empty* [*simp*]:

$a \in cs \implies \langle \ll a \gg \rangle \star_{cs} \langle \rangle = \{\langle \rangle\}_u$

by (*rel-auto*)

lemma *trace-merge-empty-single* [*simp*]:

$a \in cs \implies \langle \rangle \star_{cs} \langle \ll a \gg \rangle = \{\langle \rangle\}_u$

by (*rel-auto*)

lemma *trace-merge-commute*: $t_1 \star_{cs} t_2 = t_2 \star_{cs} t_1$

by (*rel-simp*, *simp* *add*: *tr-par-sym*)

lemma *csp-trace-simps* [*simp*]:

$v + \langle \rangle = v \langle \rangle + v = v$

$bop (\#) x xs \hat{^}_u ys = bop (\#) x (xs \hat{^}_u ys)$

by (*rel-auto*)⁺

Alternative characterisation of traces, adapted from CSP-Prover

inductive-set

parx :: 'a set => ('a list * 'a list * 'a list) set

for *X* :: 'a set

where

parx-nil-nil [*intro*]:

$([], [], []) \in parx X \mid$

parx-Ev-nil [*intro*]:

$[] (u, s, []) \in parx X ; a \notin X []$

$\implies (a \# u, a \# s, []) \in parx X \mid$

parx-nil-Ev [*intro*]:

$[] (u, [], t) \in parx X ; a \notin X []$

$\implies (a \# u, [], a \# t) \in parx X \mid$

parx-Ev-sync [*intro*]:

$[] (u, s, t) \in parx X ; a \in X []$

$\implies (a \# u, a \# s, a \# t) \in parx X \mid$

parx-Ev-left [intro]:

$\llbracket (u, s, t) \in \text{parx } X ; a \notin X \rrbracket$
 $\implies (a \# u, a \# s, t) \in \text{parx } X \mid$

parx-Ev-right [intro]:

$\llbracket (u, s, t) \in \text{parx } X ; a \notin X \rrbracket$
 $\implies (a \# u, s, a \# t) \in \text{parx } X$

lemma *parx-implies-tr-par*: $(t, t_1, t_2) \in \text{parx } cs \implies t \in \text{tr-par } cs \ t_1 \ t_2$

apply (*induct rule*: *parx.induct*)

apply (*auto*)

apply (*case-tac t*)

apply (*auto*)

apply (*case-tac s*)

apply (*auto*)

done

end

3 Syntax and Translations for Event Prefix

theory *utp-circus-prefix*

imports *UTP-Stateful-Failures.utp-sf-rdes*

begin

syntax

-simple-prefix :: *logic* \Rightarrow *logic* \Rightarrow *logic* $(- \rightarrow - [63, 62] \ 62)$

translations

$a \rightarrow P == \text{CONST PrefixCSP} \ll a \gg P$

We next configure a syntax for mixed prefixes.

nonterminal *prefix-elem'* **and** *mixed-prefix'*

syntax *-end-prefix* :: *prefix-elem'* \Rightarrow *mixed-prefix'* $(-)$

Input Prefix: $\dots?(x)$

syntax *-simple-input-prefix* :: *id* \Rightarrow *prefix-elem'* $(?'(-'))$

Input Prefix with Constraint: $\dots?(x : P)$

syntax *-input-prefix* :: *id* \Rightarrow $('\sigma, '\varepsilon)$ *action* \Rightarrow *prefix-elem'* $(?'(- : / -'))$

Output Prefix: $\dots![v]e$

A variable name must currently be provided for outputs, too. Fix?!

syntax *-output-prefix* :: *uexp* \Rightarrow *prefix-elem'* $(!'(-'))$

syntax *-output-prefix* :: *uexp* \Rightarrow *prefix-elem'* $(.'(-'))$

syntax (**output**) *-output-prefix-pp* :: *uexp* \Rightarrow *prefix-elem'* $(!'(-'))$

syntax

-prefix-aux :: *pttrn* \Rightarrow *logic* \Rightarrow *prefix-elem'*

Mixed-Prefix Action: $c \dots (\text{prefix}) \rightarrow A$

syntax *-mixed-prefix* :: *prefix-elim'* \Rightarrow *mixed-prefix'* \Rightarrow *mixed-prefix'* (--)

syntax

-prefix-action ::
 ('a, 'ε) *chan* \Rightarrow *mixed-prefix'* \Rightarrow ('σ, 'ε) *action* \Rightarrow ('σ, 'ε) *action*
 ((-- \rightarrow / -) [63, 63, 62] 62)

Syntax translations

definition *lconj* :: ('a \Rightarrow 'α *upred*) \Rightarrow ('b \Rightarrow 'α *upred*) \Rightarrow ('a \times 'b \Rightarrow 'α *upred*) (**infixr** \wedge_l 35)
where [*upred-defs*]: (*P* \wedge_l *Q*) \equiv ($\lambda (x,y). P\ x \wedge Q\ y$)

definition *outp-constraint* (**infix** =_o 60) **where**

[*upred-defs*]: *outp-constraint* *v* \equiv ($\lambda x. \ll x \gg =_u v$)

translations

-simple-input-prefix *x* \equiv *-input-prefix* *x* *true*
-mixed-prefix (*-input-prefix* *x* *P*) (*-prefix-aux* *y* *Q*) \rightarrow
-prefix-aux (*-pattern* *x* *y*) (($\lambda x. P$) \wedge_l *Q*)
-mixed-prefix (*-output-prefix* *P*) (*-prefix-aux* *y* *Q*) \rightarrow
-prefix-aux (*-pattern* *-idtdummy* *y*) ((*CONST* *outp-constraint* *P*) \wedge_l *Q*)
-end-prefix (*-input-prefix* *x* *P*) \rightarrow *-prefix-aux* *x* ($\lambda x. P$)
-end-prefix (*-output-prefix* *P*) \rightarrow *-prefix-aux* *-idtdummy* (*CONST* *outp-constraint* *P*)
-prefix-action *c* (*-prefix-aux* *x* *P*) *A* == (*CONST* *InputCSP*) *c* *P* ($\lambda x. A$)

Basic print translations; more work needed

translations

-simple-input-prefix *x* <= *-input-prefix* *x* *true*
-output-prefix *v* <= *-prefix-aux* *p* (*CONST* *outp-constraint* *v*)
-output-prefix *u* (*-output-prefix* *v*)
 <= *-prefix-aux* *p* ($\lambda(x1, y1). \text{CONST outp-constraint } u\ x2 \wedge \text{CONST outp-constraint } v\ y2$)
-input-prefix *x* *P* <= *-prefix-aux* *v* ($\lambda x. P$)
x!(*v*) $\rightarrow P$ <= *CONST* *OutputCSP* *x* *v* *P*

term *x*!(1)!(*y*) $\rightarrow P$

term *x*?(*v*) $\rightarrow P$

term *x*?(*v*:*false*) $\rightarrow P$

term *x*!(⟨1⟩) $\rightarrow P$

term *x*?(*v*)!(1) $\rightarrow P$

term *x*!(⟨1⟩)!(2)?(*v*:*true*) $\rightarrow P$

Basic translations for state variable communications

syntax

-csp-input-var :: *logic* \Rightarrow *id* \Rightarrow *logic* \Rightarrow *logic* (-?'(-:')) [63, 0, 0] 62)
-csp-inputu-var :: *logic* \Rightarrow *id* \Rightarrow *logic* (-?'(-')) [63, 0] 62)
-csp-output-var :: *logic* \Rightarrow *uexp* \Rightarrow *logic* (-!'(-')) [63, 0] 62)

translations

c?(*x*:*A*) \rightarrow *CONST* *InputVarCSP* *c* *x* *A*
c?(*x*) \rightarrow *CONST* *InputVarCSP* *c* *x* ($\lambda x. \text{true}$)
c?(*x*:*A*) <= *CONST* *InputVarCSP* *c* *x* ($\lambda x'. A$)
c?(*x*) <= *c*?(*x*:*true*)
-csp-output-var *c* *e* == *CONST* *DoCSP* (*c*·*e*)_u

lemma *outp-constraint-prod*:

(*outp-constraint* $\ll a \gg x \wedge \text{outp-constraint } \ll b \gg y$) =

outp-constraint $\ll(a, b)\gg (x, y)$
by (*simp add: outp-constraint-def, pred-auto*)

lemma *subst-outp-constraint* [*usubst*]:
 $\sigma \uparrow (v =_o x) = (\sigma \uparrow v =_o x)$
by (*rel-auto*)

lemma *UINF-one-point-simp* [*rpred*]:
 $\ll \bigwedge i. P\ i\ is\ R1 \gg \implies (\bigcap x \cdot \ll i \gg =_o x)_{S<} \wedge P(x) = P(i)$
by (*rel-blast*)

lemma *USUP-one-point-simp* [*rpred*]:
 $\ll \bigwedge i. P\ i\ is\ R1 \gg \implies (\bigcup x \cdot \ll i \gg =_o x)_{S<} \Rightarrow_r P(x) = P(i)$
by (*rel-blast*)

lemma *USUP-eq-event-eq* [*rpred*]:
assumes $\bigwedge y. P(y)\ is\ RR$
shows $(\bigcup y \cdot [v =_o y]_{S<} \Rightarrow_r P(y)) = P(y) \ll y \rightarrow [v]_{S\leftarrow} \gg$
proof –
have $(\bigcup y \cdot [v =_o y]_{S<} \Rightarrow_r RR(P(y))) = RR(P(y)) \ll y \rightarrow [v]_{S\leftarrow} \gg$
apply (*rel-simp, safe*)
apply *metis*
apply *blast*
apply *simp*
done
thus *?thesis*
by (*simp add: Healthy-if assms*)
qed

lemma *UINF-eq-event-eq* [*rpred*]:
assumes $\bigwedge y. P(y)\ is\ RR$
shows $(\bigcap y \cdot [v =_o y]_{S<} \wedge P(y)) = P(y) \ll y \rightarrow [v]_{S\leftarrow} \gg$
proof –
have $(\bigcap y \cdot [v =_o y]_{S<} \wedge RR(P(y))) = RR(P(y)) \ll y \rightarrow [v]_{S\leftarrow} \gg$
by (*rel-simp, safe, metis*)
thus *?thesis*
by (*simp add: Healthy-if assms*)
qed

Proofs that the input constrained parser versions of output is the same as the regular definition.

lemma *output-prefix-is-OutputCSP* [*simp*]:
assumes *A is NCSP*
shows $x!(P) \rightarrow A = OutputCSP\ x\ P\ A\ (\text{is } ?lhs = ?rhs)$
by (*rdes-eq cls: assms*)

lemma *OutputCSP-pair-simp* [*simp*]:
 $P\ is\ NCSP \implies a.(\ll i \gg).(\ll j \gg) \rightarrow P = OutputCSP\ a\ \ll (i, j) \gg P$
using *output-prefix-is-OutputCSP[of P a]*
by (*simp add: outp-constraint-prod lconj-def InputCSP-def closure del: output-prefix-is-OutputCSP*)

lemma *OutputCSP-triple-simp* [*simp*]:
 $P\ is\ NCSP \implies a.(\ll i \gg).(\ll j \gg).(\ll k \gg) \rightarrow P = OutputCSP\ a\ \ll (i, j, k) \gg P$
using *output-prefix-is-OutputCSP[of P a]*
by (*simp add: outp-constraint-prod lconj-def InputCSP-def closure del: output-prefix-is-OutputCSP*)

end

4 Circus Parallel Composition

theory *utp-circus-parallel*

imports

utp-circus-prefix

utp-circus-traces

begin

4.1 Merge predicates

definition *CSPInnerMerge* :: $('α \Longrightarrow 'σ) \Rightarrow 'ψ \text{ set} \Rightarrow ('β \Longrightarrow 'σ) \Rightarrow (('σ, 'ψ) \text{ sfrd}) \text{ merge } (N_C)$ **where**
 $[upred-defs]:$

$CSPInnerMerge \ ns1 \ cs \ ns2 = ($
 $\$ref' \subseteq_u ((\$0:ref \cup_u \$1:ref) \cap_u \ll cs \gg) \cup_u ((\$0:ref \cap_u \$1:ref) - \ll cs \gg) \wedge$
 $\$<:tr \leq_u \$tr' \wedge$
 $(\$tr' - \$<:tr) \in_u (\$0:tr - \$<:tr) \star_{CS} (\$1:tr - \$<:tr) \wedge$
 $(\$0:tr - \$<:tr) \downarrow_u \ll cs \gg =_u (\$1:tr - \$<:tr) \downarrow_u \ll cs \gg \wedge$
 $\$st' =_u (\$<:st \oplus \$0:st \text{ on } \&ns1) \oplus \$1:st \text{ on } \&ns2)$

definition *CSPInnerInterleave* :: $('α \Longrightarrow 'σ) \Rightarrow ('β \Longrightarrow 'σ) \Rightarrow (('σ, 'ψ) \text{ sfrd}) \text{ merge } (N_I)$ **where**
 $[upred-defs]:$

$N_I \ ns1 \ ns2 = ($
 $\$ref' \subseteq_u (\$0:ref \cap_u \$1:ref) \wedge$
 $\$<:tr \leq_u \$tr' \wedge$
 $(\$tr' - \$<:tr) \in_u (\$0:tr - \$<:tr) \star_{\{\}} (\$1:tr - \$<:tr) \wedge$
 $\$st' =_u (\$<:st \oplus \$0:st \text{ on } \&ns1) \oplus \$1:st \text{ on } \&ns2)$

An intermediate merge hides the state, whilst a final merge hides the refusals.

definition *CSPInterMerge* **where**

$[upred-defs]: CSPInterMerge \ P \ cs \ Q = (P \parallel_{(\exists \ \$st' \cdot N_C \ 0_L \ cs \ 0_L)} Q)$

definition *CSPFinalMerge* **where**

$[upred-defs]: CSPFinalMerge \ P \ ns1 \ cs \ ns2 \ Q = (P \parallel_{(\exists \ \$ref' \cdot N_C \ ns1 \ cs \ ns2)} Q)$

syntax

$-cinter\text{-}merge :: logic \Rightarrow logic \Rightarrow logic \Rightarrow logic \ (- \ll _ \rrbracket^I - [85,0,86] \ 86)$
 $-cfinal\text{-}merge :: logic \Rightarrow salpha \Rightarrow logic \Rightarrow salpha \Rightarrow logic \Rightarrow logic \ (- \ll _ \rrbracket^F - [85,0,0,86] \ 86)$
 $-wrC :: logic \Rightarrow logic \Rightarrow logic \Rightarrow logic \ (- \text{wr}[_]_C - [85,0,86] \ 86)$

translations

$-cinter\text{-}merge \ P \ cs \ Q == CONST \ CSPInterMerge \ P \ cs \ Q$
 $-cfinal\text{-}merge \ P \ ns1 \ cs \ ns2 \ Q == CONST \ CSPFinalMerge \ P \ ns1 \ cs \ ns2 \ Q$
 $-wrC \ P \ cs \ Q == P \ wr_R(N_C \ 0_L \ cs \ 0_L) \ Q$

lemma *CSPInnerMerge-R2m* $[closure]: N_C \ ns1 \ cs \ ns2 \text{ is } R2m$
by (rel-auto)

lemma *CSPInnerMerge-RDM* $[closure]: N_C \ ns1 \ cs \ ns2 \text{ is } RDM$
by (rule RDM-intro, simp add: closure, simp-all add: CSPInnerMerge-def unrest)

lemma *ex-ref'-R2m-closed* $[closure]:$

assumes $P \text{ is } R2m$

shows $(\exists \ \$ref' \cdot P) \text{ is } R2m$

proof –

have $R2m(\exists \$ref' \cdot R2m P) = (\exists \$ref' \cdot R2m P)$
 by (rel-auto)
 thus ?thesis
 by (metis Healthy-def' assms)

qed

lemma *CSPInnerMerge-unrests* [unrest]:

$\$<:ok \# N_C ns1 cs ns2$
 $\$<:wait \# N_C ns1 cs ns2$
 by (rel-auto)+

lemma *CSPInterMerge-RR-closed* [closure]:

assumes P is RR Q is RR
 shows $P \llbracket cs \rrbracket^I Q$ is RR
 by (simp add: CSPInterMerge-def parallel-RR-closed assms closure unrest)

lemma *CSPInterMerge-unrest-ref* [unrest]:

assumes P is CRR Q is CRR
 shows $\$ref \# P \llbracket cs \rrbracket^I Q$

proof –

have $\$ref \# CRR(P) \llbracket cs \rrbracket^I CRR(Q)$
 by (rel-blast)
 thus ?thesis
 by (simp add: Healthy-if assms)

qed

lemma *CSPInterMerge-unrest-st'* [unrest]:

$\$st' \# P \llbracket cs \rrbracket^I Q$
 by (rel-auto)

lemma *CSPInterMerge-CRR-closed* [closure]:

assumes P is CRR Q is CRR
 shows $P \llbracket cs \rrbracket^I Q$ is CRR
 by (simp add: CRR-implies-RR CRR-intro CSPInterMerge-RR-closed CSPInterMerge-unrest-ref assms)

lemma *CSPFinalMerge-RR-closed* [closure]:

assumes P is RR Q is RR
 shows $P \llbracket ns1|cs|ns2 \rrbracket^F Q$ is RR
 by (simp add: CSPFinalMerge-def parallel-RR-closed assms closure unrest)

lemma *CSPFinalMerge-unrest-ref* [unrest]:

assumes P is CRR Q is CRR
 shows $\$ref \# P \llbracket ns1|cs|ns2 \rrbracket^F Q$

proof –

have $\$ref \# CRR(P) \llbracket ns1|cs|ns2 \rrbracket^F CRR(Q)$
 by (rel-blast)
 thus ?thesis
 by (simp add: Healthy-if assms)

qed

lemma *CSPFinalMerge-CRR-closed* [closure]:

assumes P is CRR Q is CRR
 shows $P \llbracket ns1|cs|ns2 \rrbracket^F Q$ is CRR
 by (simp add: CRR-implies-RR CRR-intro CSPFinalMerge-RR-closed CSPFinalMerge-unrest-ref assms)

lemma *CSPFinalMerge-unrest-ref'* [unrest]:

assumes P is CRR Q is CRR
 shows $\$ref' \# P \llbracket ns1 | cs | ns2 \rrbracket^F Q$

proof –

have $\$ref' \# CRR(P) \llbracket ns1 | cs | ns2 \rrbracket^F CRR(Q)$
 by (rel-blast)
 thus ?thesis
 by (simp add: Healthy-if assms)

qed

lemma *CSPFinalMerge-CRF-closed* [closure]:

assumes P is CRF Q is CRF
 shows $P \llbracket ns1 | cs | ns2 \rrbracket^F Q$ is CRF
 by (rule CRF-intro, simp-all add: assms unrest closure)

lemma *CSPInnerMerge-empty-Interleave*:

$N_C ns1 \{ \} ns2 = N_I ns1 ns2$
 by (rel-auto)

definition *CSPMerge* :: $('α \implies 'σ) \Rightarrow 'ψ \text{ set} \Rightarrow ('β \implies 'σ) \Rightarrow (('σ, 'ψ) \text{ sfrd}) \text{ merge } (M_C)$ **where**
 $[upred-defs]: M_C ns1 cs ns2 = M_R(N_C ns1 cs ns2) ;; Skip$

definition *CSPInterleave* :: $('α \implies 'σ) \Rightarrow ('β \implies 'σ) \Rightarrow (('σ, 'ψ) \text{ sfrd}) \text{ merge } (M_I)$ **where**
 $[upred-defs]: M_I ns1 ns2 = M_R(N_I ns1 ns2) ;; Skip$

lemma *swap-CSPInnerMerge*:

$ns1 \bowtie ns2 \implies swap_m ;; (N_C ns1 cs ns2) = (N_C ns2 cs ns1)$
 apply (rel-auto)
 using tr-par-sym apply blast
 apply (simp add: lens-indep-comm)
 using tr-par-sym apply blast
 apply (simp add: lens-indep-comm)

done

lemma *SymMerge-CSPInnerMerge-NS* [closure]: $N_C 0_L cs 0_L$ is SymMerge
 by (simp add: Healthy-def swap-CSPInnerMerge)

lemma *SymMerge-CSPInnerInterleave* [closure]:

$N_I 0_L 0_L$ is SymMerge
 by (metis CSPInnerMerge-empty-Interleave SymMerge-CSPInnerMerge-NS)

lemma *SymMerge-CSPInnerInterleave* [closure]:

AssocMerge $(N_I 0_L 0_L)$
 apply (rel-auto)
 apply (rename-tac tr tr₂' ref₀ tr₀' ref₀' tr₁' ref₁' tr' ref₂' tr_i' ref₃)

oops

lemma *CSPInterMerge-right-false* [rpred]: $P \llbracket cs \rrbracket^I \text{ false} = \text{false}$
 by (simp add: CSPInterMerge-def)

lemma *CSPInterMerge-left-false* [rpred]: $\text{false} \llbracket cs \rrbracket^I P = \text{false}$
 by (rel-auto)

lemma *CSPFinalMerge-right-false* [rpred]: $P \llbracket ns1 | cs | ns2 \rrbracket^F \text{ false} = \text{false}$

by (simp add: CSPFinalMerge-def)

lemma CSPFinalMerge-left-false [rpred]: false $\llbracket ns1|cs|ns2 \rrbracket^F P = false$
by (simp add: CSPFinalMerge-def)

lemma CSPInnerMerge-commute:

assumes $ns1 \bowtie ns2$

shows $P \parallel_{N_C} ns1 \ cs \ ns2 \ Q = Q \parallel_{N_C} ns2 \ cs \ ns1 \ P$

proof –

have $P \parallel_{N_C} ns1 \ cs \ ns2 \ Q = P \parallel_{swap_m} ;; N_C \ ns2 \ cs \ ns1 \ Q$
by (simp add: assms lens-indep-sym swap-CSPInnerMerge)

also have $\dots = Q \parallel_{N_C} ns2 \ cs \ ns1 \ P$

by (metis par-by-merge-commute-swap)

finally show ?thesis .

qed

lemma CSPInterMerge-commute:

$P \llbracket cs \rrbracket^I Q = Q \llbracket cs \rrbracket^I P$

proof –

have $P \llbracket cs \rrbracket^I Q = P \parallel_{\exists \$st' \cdot N_C \ 0_L \ cs \ 0_L} Q$

by (simp add: CSPInterMerge-def)

also have $\dots = P \parallel_{\exists \$st' \cdot (swap_m ;; N_C \ 0_L \ cs \ 0_L)} Q$

by (simp add: swap-CSPInnerMerge lens-indep-sym)

also have $\dots = P \parallel_{swap_m} ;; (\exists \$st' \cdot N_C \ 0_L \ cs \ 0_L) \ Q$

by (simp add: seqr-exists-right)

also have $\dots = Q \parallel_{(\exists \$st' \cdot N_C \ 0_L \ cs \ 0_L)} P$

by (simp add: par-by-merge-commute-swap[THEN sym])

also have $\dots = Q \llbracket cs \rrbracket^I P$

by (simp add: CSPInterMerge-def)

finally show ?thesis .

qed

lemma CSPFinalMerge-commute:

assumes $ns1 \bowtie ns2$

shows $P \llbracket ns1|cs|ns2 \rrbracket^F Q = Q \llbracket ns2|cs|ns1 \rrbracket^F P$

proof –

have $P \llbracket ns1|cs|ns2 \rrbracket^F Q = P \parallel_{\exists \$ref' \cdot N_C \ ns1 \ cs \ ns2} Q$

by (simp add: CSPFinalMerge-def)

also have $\dots = P \parallel_{\exists \$ref' \cdot (swap_m ;; N_C \ ns2 \ cs \ ns1)} Q$

by (simp add: swap-CSPInnerMerge lens-indep-sym assms)

also have $\dots = P \parallel_{swap_m} ;; (\exists \$ref' \cdot N_C \ ns2 \ cs \ ns1) \ Q$

by (simp add: seqr-exists-right)

also have $\dots = Q \parallel_{(\exists \$ref' \cdot N_C \ ns2 \ cs \ ns1)} P$

by (simp add: par-by-merge-commute-swap[THEN sym])

also have $\dots = Q \llbracket ns2|cs|ns1 \rrbracket^F P$

by (simp add: CSPFinalMerge-def)

finally show ?thesis .

qed

Important theorem that shows the form of a parallel process

lemma CSPInnerMerge-form:

fixes $P \ Q :: ('s, 'v) \text{ action}$

assumes $vwb\text{-}lens \ ns1 \ vwb\text{-}lens \ ns2 \ P \text{ is } RR \ Q \text{ is } RR$

shows

$P \parallel_{N_C}^{ns1\ cs\ ns2} Q =$
 $(\exists (ref_0, ref_1, st_0, st_1, tt_0, tt_1) \cdot$
 $P[\langle\langle ref_0 \rangle\rangle, \langle\langle st_0 \rangle\rangle, \langle\langle tt_0 \rangle\rangle / \$ref', \$st', \$tr, \$tr'] \wedge Q[\langle\langle ref_1 \rangle\rangle, \langle\langle st_1 \rangle\rangle, \langle\langle tt_1 \rangle\rangle / \$ref', \$st', \$tr, \$tr']$
 $\wedge \$ref' \subseteq_u ((\langle\langle ref_0 \rangle\rangle \cup_u \langle\langle ref_1 \rangle\rangle) \cap_u \langle\langle cs \rangle\rangle) \cup_u ((\langle\langle ref_0 \rangle\rangle \cap_u \langle\langle ref_1 \rangle\rangle) - \langle\langle cs \rangle\rangle)$
 $\wedge \$tr \leq_u \tr'
 $\wedge \&tt \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} \langle\langle tt_1 \rangle\rangle$
 $\wedge \langle\langle tt_0 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle =_u \langle\langle tt_1 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle$
 $\wedge \$st' =_u (\$st \oplus \langle\langle st_0 \rangle\rangle\ on\ \&ns1) \oplus \langle\langle st_1 \rangle\rangle\ on\ \&ns2)$
(is ?lhs = ?rhs)
proof –
have $P: (\exists \{ \$ok', \$wait' \} \cdot R2(P)) = P$ **(is ?P' = -)**
by (*simp add: ex-unrest ex-plus Healthy-if assms RR-implies-R2 unrest closure*)
have $Q: (\exists \{ \$ok', \$wait' \} \cdot R2(Q)) = Q$ **(is ?Q' = -)**
by (*simp add: ex-unrest ex-plus Healthy-if assms RR-implies-R2 unrest closure*)
from *assms(1,2)*
have $?P' \parallel_{N_C}^{ns1\ cs\ ns2} ?Q' =$
 $(\exists (ref_0, ref_1, st_0, st_1, tt_0, tt_1) \cdot$
 $?P'[\langle\langle ref_0 \rangle\rangle, \langle\langle st_0 \rangle\rangle, \langle\langle tt_0 \rangle\rangle / \$ref', \$st', \$tr, \$tr'] \wedge ?Q'[\langle\langle ref_1 \rangle\rangle, \langle\langle st_1 \rangle\rangle, \langle\langle tt_1 \rangle\rangle / \$ref', \$st', \$tr, \$tr']$
 $\wedge \$ref' \subseteq_u ((\langle\langle ref_0 \rangle\rangle \cup_u \langle\langle ref_1 \rangle\rangle) \cap_u \langle\langle cs \rangle\rangle) \cup_u ((\langle\langle ref_0 \rangle\rangle \cap_u \langle\langle ref_1 \rangle\rangle) - \langle\langle cs \rangle\rangle)$
 $\wedge \$tr \leq_u \tr'
 $\wedge \&tt \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} \langle\langle tt_1 \rangle\rangle$
 $\wedge \langle\langle tt_0 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle =_u \langle\langle tt_1 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle$
 $\wedge \$st' =_u (\$st \oplus \langle\langle st_0 \rangle\rangle\ on\ \&ns1) \oplus \langle\langle st_1 \rangle\rangle\ on\ \&ns2)$
apply (*simp add: par-by-merge-alt-def, rel-auto, blast*)
apply (*rename-tac ok wait tr st ref tr' ref' ref_0 ref_1 st_0 st_1 tr_0 ok_0 tr_1 wait_0 ok_1 wait_1*)
apply (*rule-tac x=ok in exI*)
apply (*rule-tac x=wait in exI*)
apply (*rule-tac x=tr in exI*)
apply (*rule-tac x=st in exI*)
apply (*rule-tac x=ref in exI*)
apply (*rule-tac x=tr @ tr_0 in exI*)
apply (*rule-tac x=st_0 in exI*)
apply (*rule-tac x=ref_0 in exI*)
apply (*auto*)
apply (*metis Prefix-Order.prefixI append-minus*)
done
thus *?thesis*
by (*simp add: P Q*)
qed

lemma *CSPInterMerge-form:*

fixes $P\ Q :: ('σ, 'φ)\ action$

assumes $P\ is\ RR\ Q\ is\ RR$

shows

$P \llbracket cs \rrbracket^I Q =$

$(\exists (ref_0, ref_1, st_0, st_1, tt_0, tt_1) \cdot$
 $P[\langle\langle ref_0 \rangle\rangle, \langle\langle st_0 \rangle\rangle, \langle\langle tt_0 \rangle\rangle / \$ref', \$st', \$tr, \$tr'] \wedge Q[\langle\langle ref_1 \rangle\rangle, \langle\langle st_1 \rangle\rangle, \langle\langle tt_1 \rangle\rangle / \$ref', \$st', \$tr, \$tr']$
 $\wedge \$ref' \subseteq_u ((\langle\langle ref_0 \rangle\rangle \cup_u \langle\langle ref_1 \rangle\rangle) \cap_u \langle\langle cs \rangle\rangle) \cup_u ((\langle\langle ref_0 \rangle\rangle \cap_u \langle\langle ref_1 \rangle\rangle) - \langle\langle cs \rangle\rangle)$
 $\wedge \$tr \leq_u \tr'
 $\wedge \&tt \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} \langle\langle tt_1 \rangle\rangle$
 $\wedge \langle\langle tt_0 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle =_u \langle\langle tt_1 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle)$
(is ?lhs = ?rhs)

proof –

have $?lhs = (\exists \$st' \cdot P \parallel_{N_C}^{0_L\ cs\ 0_L} Q)$

by (*simp add: CSPInterMerge-def par-by-merge-def seqr-exists-right*)

also have ... =

(\exists $\$st'$.
 (\exists ($ref_0, ref_1, st_0, st_1, tt_0, tt_1$) .
 $P[\langle\langle ref_0 \rangle\rangle, \langle\langle st_0 \rangle\rangle, \langle\langle tt_0 \rangle\rangle / \$ref', \$st', \$tr, \$tr'] \wedge Q[\langle\langle ref_1 \rangle\rangle, \langle\langle st_1 \rangle\rangle, \langle\langle tt_1 \rangle\rangle / \$ref', \$st', \$tr, \$tr']$
 $\wedge \$ref' \subseteq_u ((\langle\langle ref_0 \rangle\rangle \cup_u \langle\langle ref_1 \rangle\rangle) \cap_u \langle\langle cs \rangle\rangle) \cup_u ((\langle\langle ref_0 \rangle\rangle \cap_u \langle\langle ref_1 \rangle\rangle) - \langle\langle cs \rangle\rangle)$
 $\wedge \$tr \leq_u \tr'
 $\wedge \&tt \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} \langle\langle tt_1 \rangle\rangle$
 $\wedge \langle\langle tt_0 \rangle\rangle \downarrow_u \langle\langle cs \rangle\rangle =_u \langle\langle tt_1 \rangle\rangle \downarrow_u \langle\langle cs \rangle\rangle$
 $\wedge \$st' =_u (\$st \oplus \langle\langle st_0 \rangle\rangle \text{ on } \emptyset) \oplus \langle\langle st_1 \rangle\rangle \text{ on } \emptyset)$)

by (simp add: CSPInnerMerge-form pr-var-def assms)

also have ... = ?rhs

by (rel-blast)

finally show ?thesis .

qed

lemma CSPFinalMerge-form:

fixes $P \ Q :: ('s, 'v) \text{ action}$

assumes $vwb\text{-lens } ns1 \ vwb\text{-lens } ns2 \ P \text{ is } RR \ Q \text{ is } RR \ \$ref' \ \# \ P \ \$ref' \ \# \ Q$

shows

($P \llbracket ns1 | cs | ns2 \rrbracket^F Q$) =
 (\exists (st_0, st_1, tt_0, tt_1) .
 $P[\langle\langle st_0 \rangle\rangle, \langle\langle st_1 \rangle\rangle, \langle\langle tt_0 \rangle\rangle / \$st', \$tr, \$tr'] \wedge Q[\langle\langle st_1 \rangle\rangle, \langle\langle st_1 \rangle\rangle, \langle\langle tt_1 \rangle\rangle / \$st', \$tr, \$tr']$
 $\wedge \$tr \leq_u \tr'
 $\wedge \&tt \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} \langle\langle tt_1 \rangle\rangle$
 $\wedge \langle\langle tt_0 \rangle\rangle \downarrow_u \langle\langle cs \rangle\rangle =_u \langle\langle tt_1 \rangle\rangle \downarrow_u \langle\langle cs \rangle\rangle$
 $\wedge \$st' =_u (\$st \oplus \langle\langle st_0 \rangle\rangle \text{ on } \&ns1) \oplus \langle\langle st_1 \rangle\rangle \text{ on } \&ns2)$)

(is ?lhs = ?rhs)

proof -

have ?lhs = ($\exists \ \$ref' \cdot P \parallel_{N_C \ ns1 \ cs \ ns2} Q$)

by (simp add: CSPFinalMerge-def par-by-merge-def seqr-exists-right)

also have ... =

($\exists \ \$ref'$.
 (\exists ($ref_0, ref_1, st_0, st_1, tt_0, tt_1$) .
 $P[\langle\langle ref_0 \rangle\rangle, \langle\langle st_0 \rangle\rangle, \langle\langle tt_0 \rangle\rangle / \$ref', \$st', \$tr, \$tr'] \wedge Q[\langle\langle ref_1 \rangle\rangle, \langle\langle st_1 \rangle\rangle, \langle\langle tt_1 \rangle\rangle / \$ref', \$st', \$tr, \$tr']$
 $\wedge \$ref' \subseteq_u ((\langle\langle ref_0 \rangle\rangle \cup_u \langle\langle ref_1 \rangle\rangle) \cap_u \langle\langle cs \rangle\rangle) \cup_u ((\langle\langle ref_0 \rangle\rangle \cap_u \langle\langle ref_1 \rangle\rangle) - \langle\langle cs \rangle\rangle)$
 $\wedge \$tr \leq_u \tr'
 $\wedge \&tt \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} \langle\langle tt_1 \rangle\rangle$
 $\wedge \langle\langle tt_0 \rangle\rangle \downarrow_u \langle\langle cs \rangle\rangle =_u \langle\langle tt_1 \rangle\rangle \downarrow_u \langle\langle cs \rangle\rangle$
 $\wedge \$st' =_u (\$st \oplus \langle\langle st_0 \rangle\rangle \text{ on } \&ns1) \oplus \langle\langle st_1 \rangle\rangle \text{ on } \&ns2)$)

by (simp add: CSPInnerMerge-form assms)

also have ... =

($\exists \ \$ref'$.
 (\exists ($ref_0, ref_1, st_0, st_1, tt_0, tt_1$) .
 ($\exists \ \$ref' \cdot P$) $[\langle\langle ref_0 \rangle\rangle, \langle\langle st_0 \rangle\rangle, \langle\langle tt_0 \rangle\rangle / \$ref', \$st', \$tr, \$tr'] \wedge (\exists \ \$ref' \cdot Q)[\langle\langle ref_1 \rangle\rangle, \langle\langle st_1 \rangle\rangle, \langle\langle tt_1 \rangle\rangle / \$ref', \$st', \$tr, \$tr']$
 $\wedge \$ref' \subseteq_u ((\langle\langle ref_0 \rangle\rangle \cup_u \langle\langle ref_1 \rangle\rangle) \cap_u \langle\langle cs \rangle\rangle) \cup_u ((\langle\langle ref_0 \rangle\rangle \cap_u \langle\langle ref_1 \rangle\rangle) - \langle\langle cs \rangle\rangle)$
 $\wedge \$tr \leq_u \tr'
 $\wedge \&tt \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} \langle\langle tt_1 \rangle\rangle$
 $\wedge \langle\langle tt_0 \rangle\rangle \downarrow_u \langle\langle cs \rangle\rangle =_u \langle\langle tt_1 \rangle\rangle \downarrow_u \langle\langle cs \rangle\rangle$
 $\wedge \$st' =_u (\$st \oplus \langle\langle st_0 \rangle\rangle \text{ on } \&ns1) \oplus \langle\langle st_1 \rangle\rangle \text{ on } \&ns2)$)

by (simp add: ex-unrest assms)

also have ... =

(\exists (st_0, st_1, tt_0, tt_1) .
 ($\exists \ \$ref' \cdot P$) $[\langle\langle st_0 \rangle\rangle, \langle\langle st_1 \rangle\rangle, \langle\langle tt_0 \rangle\rangle / \$st', \$tr, \$tr'] \wedge (\exists \ \$ref' \cdot Q)[\langle\langle st_1 \rangle\rangle, \langle\langle st_1 \rangle\rangle, \langle\langle tt_1 \rangle\rangle / \$st', \$tr, \$tr']$
 $\wedge \$tr \leq_u \tr'
 $\wedge \&tt \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} \langle\langle tt_1 \rangle\rangle$)

$\wedge \ll tt_0 \gg \upharpoonright_u \ll cs \gg =_u \ll tt_1 \gg \upharpoonright_u \ll cs \gg$
 $\wedge \$st' =_u (\$st \oplus \ll st_0 \gg \text{ on } \&ns1) \oplus \ll st_1 \gg \text{ on } \&ns2)$
 by (rel-blast)
 also have ... = ?rhs
 by (simp add: ex-unrest assms)
 finally show ?thesis .
 qed

lemma CSPInterleave-merge: $M_I \ ns1 \ ns2 = M_C \ ns1 \ \{\} \ ns2$
 by (rel-auto)

lemma csp-wrR-def:
 $P \ wr[cs]_C \ Q = (\neg_r ((\neg_r \ Q) ;; \ U0 \wedge P ;; \ U1 \wedge \$<:st' =_u \$st \wedge \$<:tr' =_u \$tr) ;; \ N_C \ 0_L \ cs \ 0_L ;; \ R1 \ true)$
 by (rel-auto, metis+)

lemma csp-wrR-ns-irr:
 $(P \ wr_R(N_C \ ns1 \ cs \ ns2) \ Q) = (P \ wr[cs]_C \ Q)$
 by (rel-auto)

lemma csp-wrR-CRC-closed [closure]:
 assumes $P \text{ is } CRR \ Q \text{ is } CRR$
 shows $P \ wr[cs]_C \ Q \text{ is } CRC$
proof –
 have $\$ref \ \# \ P \ wr[cs]_C \ Q$
 by (simp add: csp-wrR-def rpred closure assms unrest)
 thus ?thesis
 by (rule CRC-intro, simp-all add: closure assms)
 qed

lemma ref'-unrest-final-merge [unrest]:
 $\$ref' \ \# \ P \ \ll ns1 | cs | ns2 \gg^F \ Q$
 by (rel-auto)

lemma inter-merge-CDC-closed [closure]:
 $P \ \ll cs \gg^I \ Q \text{ is } CDC$
 using le-less-trans by (rel-blast)

lemma CSPInterMerge-alt-def:
 $P \ \ll cs \gg^I \ Q = (\exists \ \$st' \cdot P \ \parallel_{N_C \ 0_L \ cs \ 0_L} Q)$
 by (simp add: par-by-merge-def CSPInterMerge-def seqr-exists-right)

lemma CSPFinalMerge-alt-def:
 $P \ \ll ns1 | cs | ns2 \gg^F \ Q = (\exists \ \$ref' \cdot P \ \parallel_{N_C \ ns1 \ cs \ ns2} Q)$
 by (simp add: par-by-merge-def CSPFinalMerge-def seqr-exists-right)

lemma merge-csp-do-left:
 assumes $vwb\text{-}lens \ ns1 \ vwb\text{-}lens \ ns2 \ ns1 \ \bowtie \ ns2 \ P \text{ is } RR$
 shows $\Phi(s_0, \sigma_0, t_0) \parallel_{N_C \ ns1 \ cs \ ns2} P =$
 $(\exists \ (ref_1, st_1, tt_1) \cdot$
 $\ [s_0]_{S<} \wedge$
 $\ [\$ref' \mapsto_s \ll ref_1 \gg, \$st' \mapsto_s \ll st_1 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll tt_1 \gg] \ \dagger \ P \wedge$
 $\ \$ref' \subseteq_u \ll cs \gg \cup_u (\ll ref_1 \gg - \ll cs \gg) \wedge$
 $\ [\ll trace \gg \in_u t_0 \star_{cs} \ll tt_1 \gg \wedge t_0 \upharpoonright_u \ll cs \gg =_u \ll tt_1 \gg \upharpoonright_u \ll cs \gg]_t \wedge$
 $\ \$st' =_u \$st \oplus \ll \sigma_0 \gg (\$st)_a \text{ on } \&ns1 \oplus \ll st_1 \gg \text{ on } \&ns2)$

(is ?lhs = ?rhs)

proof –

have ?lhs =

$(\exists (ref_0, ref_1, st_0, st_1, tt_0, tt_1) \cdot$

$[\$ref' \mapsto_s \langle\langle ref_0 \rangle\rangle, \$st' \mapsto_s \langle\langle st_0 \rangle\rangle, \$tr \mapsto_s \langle\rangle, \$tr' \mapsto_s \langle\langle tt_0 \rangle\rangle] \dagger \Phi(s_0, \sigma_0, t_0) \wedge$

$[\$ref' \mapsto_s \langle\langle ref_1 \rangle\rangle, \$st' \mapsto_s \langle\langle st_1 \rangle\rangle, \$tr \mapsto_s \langle\rangle, \$tr' \mapsto_s \langle\langle tt_1 \rangle\rangle] \dagger P \wedge$

$\$ref' \subseteq_u (\langle\langle ref_0 \rangle\rangle \cup_u \langle\langle ref_1 \rangle\rangle) \cap_u \langle\langle cs \rangle\rangle \cup_u (\langle\langle ref_0 \rangle\rangle \cap_u \langle\langle ref_1 \rangle\rangle - \langle\langle cs \rangle\rangle) \wedge$

$\$tr \leq_u \$tr' \wedge$

$\&tt \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} \langle\langle tt_1 \rangle\rangle \wedge \langle\langle tt_0 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle =_u \langle\langle tt_1 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle \wedge \$st' =_u \$st \oplus \langle\langle st_0 \rangle\rangle \text{ on } \&ns1$

$\oplus \langle\langle st_1 \rangle\rangle \text{ on } \&ns2)$

by (simp add: CSPInnerMerge-form assms closure)

also have ... =

$(\exists (ref_1, st_1, tt_1) \cdot$

$[s_0]_{S<} \wedge$

$[\$ref' \mapsto_s \langle\langle ref_1 \rangle\rangle, \$st' \mapsto_s \langle\langle st_1 \rangle\rangle, \$tr \mapsto_s \langle\rangle, \$tr' \mapsto_s \langle\langle tt_1 \rangle\rangle] \dagger P \wedge$

$\$ref' \subseteq_u \langle\langle cs \rangle\rangle \cup_u (\langle\langle ref_1 \rangle\rangle - \langle\langle cs \rangle\rangle) \wedge$

$[\langle\langle trace \rangle\rangle \in_u t_0 \star_{cs} \langle\langle tt_1 \rangle\rangle \wedge t_0 \upharpoonright_u \langle\langle cs \rangle\rangle =_u \langle\langle tt_1 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle]_t \wedge$

$\$st' =_u \$st \oplus \langle\langle \sigma_0 \rangle\rangle (\$st)_a \text{ on } \&ns1 \oplus \langle\langle st_1 \rangle\rangle \text{ on } \&ns2)$

by (rel-blast)

finally show ?thesis .

qed

lemma merge-csp-do-right:

assumes vwb-lens ns1 vwb-lens ns2 ns1 \bowtie ns2 *P* is RR

shows $P \parallel_{N_C} ns1 \text{ cs } ns2 \Phi(s_1, \sigma_1, t_1) =$

$(\exists (ref_0, st_0, tt_0) \cdot$

$[\$ref' \mapsto_s \langle\langle ref_0 \rangle\rangle, \$st' \mapsto_s \langle\langle st_0 \rangle\rangle, \$tr \mapsto_s \langle\rangle, \$tr' \mapsto_s \langle\langle tt_0 \rangle\rangle] \dagger P \wedge$

$[s_1]_{S<} \wedge$

$\$ref' \subseteq_u \langle\langle cs \rangle\rangle \cup_u (\langle\langle ref_0 \rangle\rangle - \langle\langle cs \rangle\rangle) \wedge$

$[\langle\langle trace \rangle\rangle \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} t_1 \wedge \langle\langle tt_0 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle =_u t_1 \upharpoonright_u \langle\langle cs \rangle\rangle]_t \wedge$

$\$st' =_u \$st \oplus \langle\langle st_0 \rangle\rangle \text{ on } \&ns1 \oplus \langle\langle \sigma_1 \rangle\rangle (\$st)_a \text{ on } \&ns2)$

(is ?lhs = ?rhs)

proof –

have ?lhs = $\Phi(s_1, \sigma_1, t_1) \parallel_{N_C} ns2 \text{ cs } ns1 P$

by (simp add: CSPInnerMerge-commute assms)

also from assms **have** ... = ?rhs

apply (simp add: assms merge-csp-do-left lens-indep-sym)

apply (rel-auto)

using assms(3) lens-indep-comm tr-par-sym **apply** fastforce

using assms(3) lens-indep.lens-put-comm tr-par-sym **apply** fastforce

done

finally show ?thesis .

qed

lemma merge-csp-enable-right:

assumes vwb-lens ns1 vwb-lens ns2 ns1 \bowtie ns2 *P* is RR

shows $P \parallel_{N_C} ns1 \text{ cs } ns2 \mathcal{E}(s_0, t_0, E_0) =$

$(\exists (ref_0, ref_1, st_0, st_1, tt_0) \cdot$

$[s_0]_{S<} \wedge$

$[\$ref' \mapsto_s \langle\langle ref_0 \rangle\rangle, \$st' \mapsto_s \langle\langle st_0 \rangle\rangle, \$tr \mapsto_s \langle\rangle, \$tr' \mapsto_s \langle\langle tt_0 \rangle\rangle] \dagger P \wedge$

$(\forall e \cdot \langle\langle e \rangle\rangle \in_u [E_0]_{S<} \Rightarrow \langle\langle e \rangle\rangle \notin_u \langle\langle ref_1 \rangle\rangle) \wedge$

$\$ref' \subseteq_u (\langle\langle ref_0 \rangle\rangle \cup_u \langle\langle ref_1 \rangle\rangle) \cap_u \langle\langle cs \rangle\rangle \cup_u (\langle\langle ref_0 \rangle\rangle \cap_u \langle\langle ref_1 \rangle\rangle - \langle\langle cs \rangle\rangle) \wedge$

$[\langle\langle trace \rangle\rangle \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} t_0 \wedge \langle\langle tt_0 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle =_u t_0 \upharpoonright_u \langle\langle cs \rangle\rangle]_t \wedge$

$\$st' =_u \$st \oplus \langle\langle st_0 \rangle\rangle \text{ on } \&ns1 \oplus \langle\langle st_1 \rangle\rangle \text{ on } \&ns2)$

(is ?lhs = ?rhs)

proof –

have $?lhs = (\exists (ref_0, ref_1, st_0, st_1, tt_0, tt_1) \cdot$
 $[\$ref' \mapsto_s \langle ref_0 \rangle, \$st' \mapsto_s \langle st_0 \rangle, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \langle tt_0 \rangle] \dagger P \wedge$
 $[\$ref' \mapsto_s \langle ref_1 \rangle, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \langle tt_1 \rangle] \dagger \mathcal{E}(s_0, t_0, E_0) \wedge$
 $\$ref' \subseteq_u (\langle ref_0 \rangle \cup_u \langle ref_1 \rangle) \cap_u \langle cs \rangle \cup_u (\langle ref_0 \rangle \cap_u \langle ref_1 \rangle - \langle cs \rangle) \wedge$
 $\$tr \leq_u \$tr' \wedge \&tt \in_u \langle tt_0 \rangle \star_{cs} \langle tt_1 \rangle \wedge \langle tt_0 \rangle \upharpoonright_u \langle cs \rangle =_u \langle tt_1 \rangle \upharpoonright_u \langle cs \rangle \wedge \$st' =_u \$st$
 $\oplus \langle st_0 \rangle \text{ on } \&ns1 \oplus \langle st_1 \rangle \text{ on } \&ns2)$
by (*simp add: CSPInnerMerge-form assms closure unrest usubst*)
also have $\dots = (\exists (ref_0, ref_1, st_0, st_1, tt_0, tt_1) \cdot [\$ref' \mapsto_s \langle ref_0 \rangle, \$st' \mapsto_s \langle st_0 \rangle, \$tr \mapsto_s \langle \rangle, \tr'
 $\mapsto_s \langle tt_0 \rangle] \dagger P \wedge$
 $([s_0]_{S<} \wedge \langle tt_1 \rangle =_u [t_0]_{S<} \wedge (\forall e \cdot \langle e \rangle \in_u [E_0]_{S<} \Rightarrow \langle e \rangle \notin_u \langle ref_1 \rangle)) \wedge$
 $\$ref' \subseteq_u (\langle ref_0 \rangle \cup_u \langle ref_1 \rangle) \cap_u \langle cs \rangle \cup_u (\langle ref_0 \rangle \cap_u \langle ref_1 \rangle - \langle cs \rangle) \wedge$
 $\$tr \leq_u \$tr' \wedge \&tt \in_u \langle tt_0 \rangle \star_{cs} \langle tt_1 \rangle \wedge \langle tt_0 \rangle \upharpoonright_u \langle cs \rangle =_u \langle tt_1 \rangle \upharpoonright_u \langle cs \rangle \wedge \$st' =_u \$st$
 $\oplus \langle st_0 \rangle \text{ on } \&ns1 \oplus \langle st_1 \rangle \text{ on } \&ns2)$
by (*simp add: csp-enable-def usubst unrest*)
also have $\dots = (\exists (ref_0, ref_1, st_0, st_1, tt_0) \cdot$
 $[s_0]_{S<} \wedge$
 $[\$ref' \mapsto_s \langle ref_0 \rangle, \$st' \mapsto_s \langle st_0 \rangle, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \langle tt_0 \rangle] \dagger P \wedge$
 $(\forall e \cdot \langle e \rangle \in_u [E_0]_{S<} \Rightarrow \langle e \rangle \notin_u \langle ref_1 \rangle) \wedge$
 $\$ref' \subseteq_u (\langle ref_0 \rangle \cup_u \langle ref_1 \rangle) \cap_u \langle cs \rangle \cup_u (\langle ref_0 \rangle \cap_u \langle ref_1 \rangle - \langle cs \rangle) \wedge$
 $[\langle trace \rangle \in_u \langle tt_0 \rangle \star_{cs} t_0 \wedge \langle tt_0 \rangle \upharpoonright_u \langle cs \rangle =_u t_0 \upharpoonright_u \langle cs \rangle]_t \wedge$
 $\$st' =_u \$st \oplus \langle st_0 \rangle \text{ on } \&ns1 \oplus \langle st_1 \rangle \text{ on } \&ns2)$
by (*rel-blast*)
finally show $?thesis$.
qed

lemma *merge-csp-enable-left:*

assumes *vwb-lens ns1 vwb-lens ns2 ns1 \bowtie ns2 P is RR*

shows $\mathcal{E}(s_0, t_0, E_0) \parallel_{N_C} ns1 \text{ cs } ns2 P =$

$(\exists (ref_0, ref_1, st_0, st_1, tt_0) \cdot$
 $[s_0]_{S<} \wedge$
 $[\$ref' \mapsto_s \langle ref_0 \rangle, \$st' \mapsto_s \langle st_1 \rangle, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \langle tt_0 \rangle] \dagger P \wedge$
 $(\forall e \cdot \langle e \rangle \in_u [E_0]_{S<} \Rightarrow \langle e \rangle \notin_u \langle ref_1 \rangle) \wedge$
 $\$ref' \subseteq_u (\langle ref_0 \rangle \cup_u \langle ref_1 \rangle) \cap_u \langle cs \rangle \cup_u (\langle ref_0 \rangle \cap_u \langle ref_1 \rangle - \langle cs \rangle) \wedge$
 $[\langle trace \rangle \in_u t_0 \star_{cs} \langle tt_0 \rangle \wedge \langle tt_0 \rangle \upharpoonright_u \langle cs \rangle =_u t_0 \upharpoonright_u \langle cs \rangle]_t \wedge$
 $\$st' =_u \$st \oplus \langle st_0 \rangle \text{ on } \&ns1 \oplus \langle st_1 \rangle \text{ on } \&ns2)$

(**is** $?lhs = ?rhs$)

proof –

have $?lhs = P \parallel_{N_C} ns2 \text{ cs } ns1 \mathcal{E}(s_0, t_0, E_0)$

by (*simp add: CSPInnerMerge-commute assms*)

also from *assms* **have** $\dots = ?rhs$

apply (*simp add: merge-csp-enable-right assms(4) lens-indep-sym*)

apply (*rel-auto*)

oops

The result of merge two terminated stateful traces is to (1) require both state preconditions hold, (2) merge the traces using, and (3) merge the state using a parallel assignment.

lemma *FinalMerge-csp-do-left:*

assumes *vwb-lens ns1 vwb-lens ns2 ns1 \bowtie ns2 P is RR* $\$ref' \nmid P$

shows $\Phi(s_0, \sigma_0, t_0) \llbracket ns1 | cs | ns2 \rrbracket^F P =$

$(\exists (st_1, t_1) \cdot$
 $[s_0]_{S<} \wedge$
 $[\$st' \mapsto_s \langle st_1 \rangle, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \langle t_1 \rangle] \dagger P \wedge$
 $[\langle trace \rangle \in_u t_0 \star_{cs} \langle t_1 \rangle \wedge t_0 \upharpoonright_u \langle cs \rangle =_u \langle t_1 \rangle \upharpoonright_u \langle cs \rangle]_t \wedge$
 $\$st' =_u \$st \oplus \langle \sigma_0 \rangle (\$st)_a \text{ on } \&ns1 \oplus \langle st_1 \rangle \text{ on } \&ns2)$

(is ?lhs = ?rhs)

proof –

have ?lhs =

$(\exists (st_0, st_1, tt_0, tt_1) \cdot$

$[\$st' \mapsto_s \langle st_0 \rangle, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \langle tt_0 \rangle] \dagger \Phi(s_0, \sigma_0, t_0) \wedge$

$[\$st' \mapsto_s \langle st_1 \rangle, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \langle tt_1 \rangle] \dagger RR(\exists \$ref' \cdot P) \wedge$

$\$tr \leq_u \$tr' \wedge \&tt \in_u \langle tt_0 \rangle \star_{cs} \langle tt_1 \rangle \wedge \langle tt_0 \rangle \upharpoonright_u \langle cs \rangle =_u \langle tt_1 \rangle \upharpoonright_u \langle cs \rangle \wedge$

$\$st' =_u \$st \oplus \langle st_0 \rangle \text{ on } \&ns1 \oplus \langle st_1 \rangle \text{ on } \&ns2)$

by (simp add: CSPFinalMerge-form ex-unrest Healthy-if unrest closure assms)

also have ... =

$(\exists (st_1, tt_1) \cdot$

$[s_0]_{S<} \wedge$

$[\$st' \mapsto_s \langle st_1 \rangle, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \langle tt_1 \rangle] \dagger RR(\exists \$ref' \cdot P) \wedge$

$\langle \text{trace} \rangle \in_u t_0 \star_{cs} \langle tt_1 \rangle \wedge t_0 \upharpoonright_u \langle cs \rangle =_u \langle tt_1 \rangle \upharpoonright_u \langle cs \rangle]_t \wedge$

$\$st' =_u \$st \oplus \langle \sigma_0 \rangle (\$st)_a \text{ on } \&ns1 \oplus \langle st_1 \rangle \text{ on } \&ns2)$

by (rel-blast)

also have ... =

$(\exists (st_1, t_1) \cdot$

$[s_0]_{S<} \wedge$

$[\$st' \mapsto_s \langle st_1 \rangle, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \langle t_1 \rangle] \dagger P \wedge$

$\langle \text{trace} \rangle \in_u t_0 \star_{cs} \langle t_1 \rangle \wedge t_0 \upharpoonright_u \langle cs \rangle =_u \langle t_1 \rangle \upharpoonright_u \langle cs \rangle]_t \wedge$

$\$st' =_u \$st \oplus \langle \sigma_0 \rangle (\$st)_a \text{ on } \&ns1 \oplus \langle st_1 \rangle \text{ on } \&ns2)$

by (simp add: ex-unrest Healthy-if unrest closure assms)

finally show ?thesis .

qed

lemma FinalMerge-csp-do-right:

assumes vwb-lens ns1 vwb-lens ns2 ns1 \bowtie ns2 P is RR \$ref' \nmid P

shows P $\llbracket ns1 | cs | ns2 \rrbracket^F \Phi(s_1, \sigma_1, t_1) =$

$(\exists (st_0, t_0) \cdot$

$[\$st' \mapsto_s \langle st_0 \rangle, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \langle t_0 \rangle] \dagger P \wedge$

$[s_1]_{S<} \wedge$

$\langle \text{trace} \rangle \in_u \langle t_0 \rangle \star_{cs} t_1 \wedge \langle t_0 \rangle \upharpoonright_u \langle cs \rangle =_u t_1 \upharpoonright_u \langle cs \rangle]_t \wedge$

$\$st' =_u \$st \oplus \langle st_0 \rangle \text{ on } \&ns1 \oplus \langle \sigma_1 \rangle (\$st)_a \text{ on } \&ns2)$

 (is ?lhs = ?rhs)

proof –

have P $\llbracket ns1 | cs | ns2 \rrbracket^F \Phi(s_1, \sigma_1, t_1) = \Phi(s_1, \sigma_1, t_1) \llbracket ns2 | cs | ns1 \rrbracket^F P$

by (simp add: assms CSPFinalMerge-commute)

also have ... = ?rhs

apply (simp add: FinalMerge-csp-do-left assms lens-indep-sym conj-comm)

apply (rel-auto)

using assms(3) lens-indep.lens-put-comm tr-par-sym **apply** fastforce+

done

finally show ?thesis .

qed

lemma FinalMerge-csp-do:

assumes vwb-lens ns1 vwb-lens ns2 ns1 \bowtie ns2

shows $\Phi(s_1, \sigma_1, t_1) \llbracket ns1 | cs | ns2 \rrbracket^F \Phi(s_2, \sigma_2, t_2) =$

$([s_1 \wedge s_2]_{S<} \wedge \langle \text{trace} \rangle \in_u t_1 \star_{cs} t_2 \wedge t_1 \upharpoonright_u \langle cs \rangle =_u t_2 \upharpoonright_u \langle cs \rangle]_t \wedge [(\sigma_1 [\&ns1 | \&ns2]_s \sigma_2)_a]_{S'})$

 (is ?lhs = ?rhs)

proof –

have ?lhs =

$(\exists (st_0, st_1, tt_0, tt_1) \cdot$

$$\begin{aligned}
& [\$st' \mapsto_s \langle st_0 \rangle, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \langle tt_0 \rangle] \dagger \Phi(s_1, \sigma_1, t_1) \wedge \\
& [\$st' \mapsto_s \langle st_1 \rangle, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \langle tt_1 \rangle] \dagger \Phi(s_2, \sigma_2, t_2) \wedge \\
& \$tr \leq_u \$tr' \wedge \langle tt \rangle \in_u \langle tt_0 \rangle \star_{cs} \langle tt_1 \rangle \wedge \langle tt_0 \rangle \upharpoonright_u \langle cs \rangle =_u \langle tt_1 \rangle \upharpoonright_u \langle cs \rangle \wedge \\
& \$st' =_u \$st \oplus \langle st_0 \rangle \text{ on } \&ns1 \oplus \langle st_1 \rangle \text{ on } \&ns2) \\
& \text{by (simp add: CSPFinalMerge-form unrest closure assms)} \\
& \text{also have ... =} \\
& ([s_1 \wedge s_2]_{S<} \wedge [\langle trace \rangle \in_u t_1 \star_{cs} t_2 \wedge t_1 \upharpoonright_u \langle cs \rangle =_u t_2 \upharpoonright_u \langle cs \rangle]_t \wedge [\langle \sigma_1 [\&ns1 | \&ns2]_s \sigma_2 \rangle_a]_{S'}) \\
& \text{by (rel-auto)} \\
& \text{finally show ?thesis .} \\
& \text{qed}
\end{aligned}$$

lemma *FinalMerge-csp-do'* [rpred]:
assumes *vwb-lens ns1 vwb-lens ns2 ns1* \bowtie *ns2*
shows $\Phi(s_1, \sigma_1, t_1) \llbracket ns1 | cs | ns2 \rrbracket^F \Phi(s_2, \sigma_2, t_2) =$
 $(\exists \text{ trace} \cdot \Phi(s_1 \wedge s_2 \wedge \langle trace \rangle \in_u t_1 \star_{cs} t_2 \wedge t_1 \upharpoonright_u \langle cs \rangle =_u t_2 \upharpoonright_u \langle cs \rangle, \sigma_1 [\&ns1 | \&ns2]_s \sigma_2,$
 $\langle trace \rangle))$
by (simp add: FinalMerge-csp-do assms, rel-auto)

lemma *CSPFinalMerge-UINF-mem-left* [rpred]:
 $(\bigcap i \in A \cdot P(i)) \llbracket ns1 | cs | ns2 \rrbracket^F Q = (\bigcap i \in A \cdot P(i) \llbracket ns1 | cs | ns2 \rrbracket^F Q)$
by (simp add: CSPFinalMerge-def par-by-merge-USUP-mem-left)

lemma *CSPFinalMerge-UINF-ind-left* [rpred]:
 $(\bigcap i \cdot P(i)) \llbracket ns1 | cs | ns2 \rrbracket^F Q = (\bigcap i \cdot P(i) \llbracket ns1 | cs | ns2 \rrbracket^F Q)$
by (simp add: CSPFinalMerge-def par-by-merge-USUP-ind-left)

lemma *CSPFinalMerge-UINF-mem-right* [rpred]:
 $P \llbracket ns1 | cs | ns2 \rrbracket^F (\bigcap i \in A \cdot Q(i)) = (\bigcap i \in A \cdot P \llbracket ns1 | cs | ns2 \rrbracket^F Q(i))$
by (simp add: CSPFinalMerge-def par-by-merge-USUP-mem-right)

lemma *CSPFinalMerge-UINF-ind-right* [rpred]:
 $P \llbracket ns1 | cs | ns2 \rrbracket^F (\bigcap i \cdot Q(i)) = (\bigcap i \cdot P \llbracket ns1 | cs | ns2 \rrbracket^F Q(i))$
by (simp add: CSPFinalMerge-def par-by-merge-USUP-ind-right)

lemma *InterMerge-csp-enable-left*:
assumes *P is RR \$st' \# P*
shows $\mathcal{E}(s_0, t_0, E_0) \llbracket cs \rrbracket^I P =$
 $(\exists (\text{ref}_0, \text{ref}_1, t_1) \cdot$
 $[s_0]_{S<} \wedge (\forall e \cdot \langle e \rangle \in_u [E_0]_{S<} \Rightarrow \langle e \rangle \notin_u \langle \text{ref}_0 \rangle) \wedge$
 $[\$ref' \mapsto_s \langle \text{ref}_1 \rangle, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \langle t_1 \rangle] \dagger P \wedge$
 $\$ref' \subseteq_u (\langle \text{ref}_0 \rangle \cup_u \langle \text{ref}_1 \rangle) \cap_u \langle cs \rangle \cup_u (\langle \text{ref}_0 \rangle \cap_u \langle \text{ref}_1 \rangle - \langle cs \rangle) \wedge$
 $[\langle trace \rangle \in_u t_0 \star_{cs} \langle t_1 \rangle \wedge t_0 \upharpoonright_u \langle cs \rangle =_u \langle t_1 \rangle \upharpoonright_u \langle cs \rangle]_t)$
(is ?lhs = ?rhs)
apply (simp add: CSPInterMerge-form ex-unrest Healthy-if unrest closure assms usubst)
apply (simp add: csp-enable-def usubst unrest assms closure)
apply (rel-auto)
done

lemma *InterMerge-csp-enable*:
 $\mathcal{E}(s_1, t_1, E_1) \llbracket cs \rrbracket^I \mathcal{E}(s_2, t_2, E_2) =$
 $([s_1 \wedge s_2]_{S<} \wedge$
 $(\forall e \in [(E_1 \cap_u E_2 \cap_u \langle cs \rangle) \cup_u ((E_1 \cup_u E_2) - \langle cs \rangle)]_{S<} \cdot \langle e \rangle \notin_u \$ref') \wedge$
 $[\langle trace \rangle \in_u t_1 \star_{cs} t_2 \wedge t_1 \upharpoonright_u \langle cs \rangle =_u t_2 \upharpoonright_u \langle cs \rangle]_t)$

(is ?lhs = ?rhs)

proof –

have ?lhs =

$(\exists (ref_0, ref_1, st_0, st_1, tt_0, tt_1) \cdot$

$[\$ref' \mapsto_s \langle\langle ref_0 \rangle\rangle, \$st' \mapsto_s \langle\langle st_0 \rangle\rangle, \$tr \mapsto_s \langle\rangle, \$tr' \mapsto_s \langle\langle tt_0 \rangle\rangle] \dagger \mathcal{E}(s_1, t_1, E_1) \wedge$

$[\$ref' \mapsto_s \langle\langle ref_1 \rangle\rangle, \$st' \mapsto_s \langle\langle st_1 \rangle\rangle, \$tr \mapsto_s \langle\rangle, \$tr' \mapsto_s \langle\langle tt_1 \rangle\rangle] \dagger \mathcal{E}(s_2, t_2, E_2) \wedge$

$\$ref' \subseteq_u (\langle\langle ref_0 \rangle\rangle \cup_u \langle\langle ref_1 \rangle\rangle) \cap_u \langle\langle cs \rangle\rangle \cup_u (\langle\langle ref_0 \rangle\rangle \cap_u \langle\langle ref_1 \rangle\rangle - \langle\langle cs \rangle\rangle) \wedge$

$\$tr \leq_u \$tr' \wedge \&tt \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} \langle\langle tt_1 \rangle\rangle \wedge \langle\langle tt_0 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle =_u \langle\langle tt_1 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle)$

by (simp add: CSPInterMerge-form unrest closure)

also have ... =

$(\exists (ref_0, ref_1, tt_0, tt_1) \cdot$

$[\$ref' \mapsto_s \langle\langle ref_0 \rangle\rangle, \$tr \mapsto_s \langle\rangle, \$tr' \mapsto_s \langle\langle tt_0 \rangle\rangle] \dagger \mathcal{E}(s_1, t_1, E_1) \wedge$

$[\$ref' \mapsto_s \langle\langle ref_1 \rangle\rangle, \$tr \mapsto_s \langle\rangle, \$tr' \mapsto_s \langle\langle tt_1 \rangle\rangle] \dagger \mathcal{E}(s_2, t_2, E_2) \wedge$

$\$ref' \subseteq_u (\langle\langle ref_0 \rangle\rangle \cup_u \langle\langle ref_1 \rangle\rangle) \cap_u \langle\langle cs \rangle\rangle \cup_u (\langle\langle ref_0 \rangle\rangle \cap_u \langle\langle ref_1 \rangle\rangle - \langle\langle cs \rangle\rangle) \wedge$

$\$tr \leq_u \$tr' \wedge \&tt \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} \langle\langle tt_1 \rangle\rangle \wedge \langle\langle tt_0 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle =_u \langle\langle tt_1 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle)$

by (rel-auto)

also have ... =

$([s_1 \wedge s_2]_{S<} \wedge$

$(\forall e \in [(E_1 \cap_u E_2 \cap_u \langle\langle cs \rangle\rangle) \cup_u ((E_1 \cup_u E_2) - \langle\langle cs \rangle\rangle)]_{S<} \cdot \langle\langle e \rangle\rangle \notin_u \$ref') \wedge$

$[\langle\langle trace \rangle\rangle \in_u t_1 \star_{cs} t_2 \wedge t_1 \upharpoonright_u \langle\langle cs \rangle\rangle =_u t_2 \upharpoonright_u \langle\langle cs \rangle\rangle]_t$

)

apply (rel-auto)

apply (rename-tac tr st tr' ref')

apply (rule-tac x=– $\llbracket E_1 \rrbracket_e$ st in exI)

apply (simp)

apply (rule-tac x=– $\llbracket E_2 \rrbracket_e$ st in exI)

apply (auto)

done

finally show ?thesis .

qed

lemma InterMerge-csp-enable' [rpred]:

$\mathcal{E}(s_1, t_1, E_1) \llbracket cs \rrbracket^I \mathcal{E}(s_2, t_2, E_2) =$

$(\exists trace \cdot \mathcal{E}(s_1 \wedge s_2 \wedge \langle\langle trace \rangle\rangle \in_u t_1 \star_{cs} t_2 \wedge t_1 \upharpoonright_u \langle\langle cs \rangle\rangle =_u t_2 \upharpoonright_u \langle\langle cs \rangle\rangle$

$, \langle\langle trace \rangle\rangle$

$, (E_1 \cap_u E_2 \cap_u \langle\langle cs \rangle\rangle) \cup_u ((E_1 \cup_u E_2) - \langle\langle cs \rangle\rangle)))$

by (simp add: InterMerge-csp-enable, rel-auto)

lemma InterMerge-csp-enable-csp-do [rpred]:

$\mathcal{E}(s_1, t_1, E_1) \llbracket cs \rrbracket^I \Phi(s_2, \sigma_2, t_2) =$

$(\exists trace \cdot \mathcal{E}(s_1 \wedge s_2 \wedge \langle\langle trace \rangle\rangle \in_u t_1 \star_{cs} t_2 \wedge t_1 \upharpoonright_u \langle\langle cs \rangle\rangle =_u t_2 \upharpoonright_u \langle\langle cs \rangle\rangle, \langle\langle trace \rangle\rangle, E_1 - \langle\langle cs \rangle\rangle))$

(is ?lhs = ?rhs)

proof –

have ?lhs =

$(\exists (ref_0, ref_1, st_0, st_1, tt_0, tt_1) \cdot$

$[\$ref' \mapsto_s \langle\langle ref_0 \rangle\rangle, \$st' \mapsto_s \langle\langle st_0 \rangle\rangle, \$tr \mapsto_s \langle\rangle, \$tr' \mapsto_s \langle\langle tt_0 \rangle\rangle] \dagger \mathcal{E}(s_1, t_1, E_1) \wedge$

$[\$ref' \mapsto_s \langle\langle ref_1 \rangle\rangle, \$st' \mapsto_s \langle\langle st_1 \rangle\rangle, \$tr \mapsto_s \langle\rangle, \$tr' \mapsto_s \langle\langle tt_1 \rangle\rangle] \dagger \Phi(s_2, \sigma_2, t_2) \wedge$

$\$ref' \subseteq_u (\langle\langle ref_0 \rangle\rangle \cup_u \langle\langle ref_1 \rangle\rangle) \cap_u \langle\langle cs \rangle\rangle \cup_u (\langle\langle ref_0 \rangle\rangle \cap_u \langle\langle ref_1 \rangle\rangle - \langle\langle cs \rangle\rangle) \wedge$

$\$tr \leq_u \$tr' \wedge \&tt \in_u \langle\langle tt_0 \rangle\rangle \star_{cs} \langle\langle tt_1 \rangle\rangle \wedge \langle\langle tt_0 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle =_u \langle\langle tt_1 \rangle\rangle \upharpoonright_u \langle\langle cs \rangle\rangle)$

by (simp add: CSPInterMerge-form unrest closure)

also have ... =

$(\exists (ref_0, ref_1, tt_0) \cdot$

$[\$ref' \mapsto_s \langle\langle ref_0 \rangle\rangle, \$tr \mapsto_s \langle\rangle, \$tr' \mapsto_s \langle\langle tt_0 \rangle\rangle] \dagger \mathcal{E}(s_1, t_1, E_1) \wedge$

$[s_2]_{S<} \wedge$

$\$ref' \subseteq_u (\langle\langle ref_0 \rangle\rangle \cup_u \langle\langle ref_1 \rangle\rangle) \cap_u \langle\langle cs \rangle\rangle \cup_u (\langle\langle ref_0 \rangle\rangle \cap_u \langle\langle ref_1 \rangle\rangle - \langle\langle cs \rangle\rangle) \wedge$

$$[\llbracket \text{trace} \rrbracket \in_u t_1 \star_{cs} t_2 \wedge t_1 \downarrow_u \llbracket cs \rrbracket =_u t_2 \downarrow_u \llbracket cs \rrbracket]_t)$$
 by (*rel-auto*)
 also have ... = $([s_1 \wedge s_2]_{S<} \wedge (\forall e \in [(E_1 - \llbracket cs \rrbracket)]_{S<} \cdot \llbracket e \rrbracket \notin_u \$ref') \wedge$

$$[\llbracket \text{trace} \rrbracket \in_u t_1 \star_{cs} t_2 \wedge t_1 \downarrow_u \llbracket cs \rrbracket =_u t_2 \downarrow_u \llbracket cs \rrbracket]_t)$$

 by (*rel-auto*)
 also have ... = $(\exists \text{ trace} \cdot \mathcal{E}(s_1 \wedge s_2 \wedge \llbracket \text{trace} \rrbracket \in_u t_1 \star_{cs} t_2 \wedge t_1 \downarrow_u \llbracket cs \rrbracket =_u t_2 \downarrow_u \llbracket cs \rrbracket, \llbracket \text{trace} \rrbracket, E_1 - \llbracket cs \rrbracket))$
 by (*rel-auto*)
 finally show *?thesis* .
 qed

lemma *InterMerge-csp-do-csp-enable* [*rpred*]:

$$\Phi(s_1, \sigma_1, t_1) \llbracket cs \rrbracket^I \mathcal{E}(s_2, t_2, E_2) =$$

$$(\exists \text{ trace} \cdot \mathcal{E}(s_1 \wedge s_2 \wedge \llbracket \text{trace} \rrbracket \in_u t_1 \star_{cs} t_2 \wedge t_1 \downarrow_u \llbracket cs \rrbracket =_u t_2 \downarrow_u \llbracket cs \rrbracket, \llbracket \text{trace} \rrbracket, E_2 - \llbracket cs \rrbracket))$$

 (is *?lhs* = *?rhs*)

proof –

have $\Phi(s_1, \sigma_1, t_1) \llbracket cs \rrbracket^I \mathcal{E}(s_2, t_2, E_2) = \mathcal{E}(s_2, t_2, E_2) \llbracket cs \rrbracket^I \Phi(s_1, \sigma_1, t_1)$
 by (*simp add: CSPInterMerge-commute*)
 also have ... = *?rhs*
 by (*simp add: rpred trace-merge-commute eq-upred-sym, rel-auto*)
 finally show *?thesis* .

qed

lemma *CSPInterMerge-or-left* [*rpred*]:

$$(P \vee Q) \llbracket cs \rrbracket^I R = (P \llbracket cs \rrbracket^I R \vee Q \llbracket cs \rrbracket^I R)$$

 by (*simp add: CSPInterMerge-def par-by-merge-or-left*)

lemma *CSPInterMerge-or-right* [*rpred*]:

$$P \llbracket cs \rrbracket^I (Q \vee R) = (P \llbracket cs \rrbracket^I Q \vee P \llbracket cs \rrbracket^I R)$$

 by (*simp add: CSPInterMerge-def par-by-merge-or-right*)

lemma *CSPFinalMerge-or-left* [*rpred*]:

$$(P \vee Q) \llbracket ns1 | cs | ns2 \rrbracket^F R = (P \llbracket ns1 | cs | ns2 \rrbracket^F R \vee Q \llbracket ns1 | cs | ns2 \rrbracket^F R)$$

 by (*simp add: CSPFinalMerge-def par-by-merge-or-left*)

lemma *CSPFinalMerge-or-right* [*rpred*]:

$$P \llbracket ns1 | cs | ns2 \rrbracket^F (Q \vee R) = (P \llbracket ns1 | cs | ns2 \rrbracket^F Q \vee P \llbracket ns1 | cs | ns2 \rrbracket^F R)$$

 by (*simp add: CSPFinalMerge-def par-by-merge-or-right*)

lemma *CSPInterMerge-UINF-mem-left* [*rpred*]:

$$(\bigcap_{i \in A} i \cdot P(i)) \llbracket cs \rrbracket^I Q = (\bigcap_{i \in A} i \cdot P(i) \llbracket cs \rrbracket^I Q)$$

 by (*simp add: CSPInterMerge-def par-by-merge-USUP-mem-left*)

lemma *CSPInterMerge-UINF-ind-left* [*rpred*]:

$$(\bigcap i \cdot P(i)) \llbracket cs \rrbracket^I Q = (\bigcap i \cdot P(i) \llbracket cs \rrbracket^I Q)$$

 by (*simp add: CSPInterMerge-def par-by-merge-USUP-ind-left*)

lemma *CSPInterMerge-UINF-mem-right* [*rpred*]:

$$P \llbracket cs \rrbracket^I (\bigcap_{i \in A} i \cdot Q(i)) = (\bigcap_{i \in A} i \cdot P \llbracket cs \rrbracket^I Q(i))$$

 by (*simp add: CSPInterMerge-def par-by-merge-USUP-mem-right*)

lemma *CSPInterMerge-UINF-ind-right* [*rpred*]:

$$P \llbracket cs \rrbracket^I (\bigcap i \cdot Q(i)) = (\bigcap i \cdot P \llbracket cs \rrbracket^I Q(i))$$

 by (*simp add: CSPInterMerge-def par-by-merge-USUP-ind-right*)

lemma *CSPInterMerge-shEx-left* [*rpred*]:
 $(\exists i \cdot P(i)) \llbracket cs \rrbracket^I Q = (\exists i \cdot P(i) \llbracket cs \rrbracket^I Q)$
using *CSPInterMerge-UINF-ind-left*[*of P cs Q*]
by (*simp add: UINF-is-exists*)

lemma *CSPInterMerge-shEx-right* [*rpred*]:
 $P \llbracket cs \rrbracket^I (\exists i \cdot Q(i)) = (\exists i \cdot P \llbracket cs \rrbracket^I Q(i))$
using *CSPInterMerge-UINF-ind-right*[*of P cs Q*]
by (*simp add: UINF-is-exists*)

lemma *par-by-merge-seq-remove*: $(P \parallel_M \text{;;} R \text{ } Q) = (P \parallel_M Q) \text{;;} R$
by (*simp add: par-by-merge-seq-add[THEN sym]*)

lemma *utrace-leg*: $(x \leq_u y) = (\exists z \cdot y =_u x \hat{\cdot}_u \ll z \gg)$
by (*rel-auto*)

lemma *trace-pred-R1-true*: $[P(\text{trace})]_t \text{;;} R1 \text{ true} = [(\exists tt_0 \cdot \ll tt_0 \gg \leq_u \ll \text{trace} \gg \wedge P(tt_0))]_t$
apply (*rel-auto*)
using *minus-cancel-le* **apply** *blast*
apply (*metis diff-add-cancel-left' le-add trace-class.add-diff-cancel-left trace-class.add-left-mono*)
done

lemma *wrC-csp-do-init* [*wp*]:
 $\Phi(s_1, \sigma_1, t_1) \text{ wr}[cs]_C \mathcal{I}(s_2, t_2) =$
 $(\forall (tt_0, tt_1) \cdot \mathcal{I}(s_1 \wedge s_2 \wedge \ll tt_1 \gg \in_u (t_2 \hat{\cdot}_u \ll tt_0 \gg) \star_{cs} t_1 \wedge t_2 \hat{\cdot}_u \ll tt_0 \gg \downarrow_u \ll cs \gg =_u t_1 \downarrow_u \ll cs \gg,$
 $\ll tt_1 \gg))$
(is ?lhs = ?rhs)

proof –

have *?lhs* =

$(\neg_r (\exists (ref_0, st_0, tt_0) \cdot$
 $[\$ref' \mapsto_s \ll ref_0 \gg, \$st' \mapsto_s \ll st_0 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll tt_0 \gg] \dagger (\neg_r \mathcal{I}(s_2, t_2)) \wedge$
 $[s_1]_{S<} \wedge$
 $\$ref' \subseteq_u \ll cs \gg \cup_u (\ll ref_0 \gg - \ll cs \gg) \wedge$
 $[\ll \text{trace} \gg \in_u \ll tt_0 \gg \star_{cs} t_1 \wedge \ll tt_0 \gg \downarrow_u \ll cs \gg =_u t_1 \downarrow_u \ll cs \gg]_t \wedge$
 $\$st' =_u \$st) \text{;;} R1 \text{ true})$

by (*simp add: wrR-def par-by-merge-seq-remove merge-csp-do-right pr-var-def closure Healthy-if rpred*)

also have ... =

$(\neg_r (\exists tt_0 \cdot ([s_2]_{S<} \wedge [t_2]_{S<} \leq_u \ll tt_0 \gg) \wedge [s_1]_{S<} \wedge$
 $[\ll \text{trace} \gg \in_u \ll tt_0 \gg \star_{cs} t_1 \wedge \ll tt_0 \gg \downarrow_u \ll cs \gg =_u t_1 \downarrow_u \ll cs \gg]_t) \text{;;} R1 \text{ true})$

by (*rel-auto*)

also have ... =

$(\neg_r (\exists tt_0 \cdot ([s_2]_{S<} \wedge (\exists tt_1 \cdot \ll tt_0 \gg =_u [t_2]_{S<} \hat{\cdot}_u \ll tt_1 \gg)) \wedge [s_1]_{S<} \wedge$
 $[\ll \text{trace} \gg \in_u \ll tt_0 \gg \star_{cs} t_1 \wedge \ll tt_0 \gg \downarrow_u \ll cs \gg =_u t_1 \downarrow_u \ll cs \gg]_t) \text{;;} R1 \text{ true})$

by (*simp add: utrace-leg*)

also have ... =

$(\neg_r (\exists tt_1 \cdot [s_1 \wedge s_2 \wedge \ll \text{trace} \gg \in_u (t_2 \hat{\cdot}_u \ll tt_1 \gg) \star_{cs} t_1 \wedge t_2 \hat{\cdot}_u \ll tt_1 \gg \downarrow_u \ll cs \gg =_u t_1 \downarrow_u \ll cs \gg]_t)$

;; *R1 true*)

by (*rel-auto*)

also have ... =

$(\forall tt_1 \cdot \neg_r ([s_1 \wedge s_2 \wedge \ll \text{trace} \gg \in_u (t_2 \hat{\cdot}_u \ll tt_1 \gg) \star_{cs} t_1 \wedge t_2 \hat{\cdot}_u \ll tt_1 \gg \downarrow_u \ll cs \gg =_u t_1 \downarrow_u \ll cs \gg]_t$

;; *R1 true*)

by (*rel-auto*)

also have ... =

$(\forall (tt_0, tt_1) \cdot \neg_r ([s_1 \wedge s_2 \wedge \ll tt_0 \gg \leq_u \ll \text{trace} \gg \wedge \ll tt_0 \gg \in_u (t_2 \hat{\cdot}_u \ll tt_1 \gg) \star_{cs} t_1 \wedge t_2 \hat{\cdot}_u \ll tt_1 \gg \downarrow_u$

$\ll cs \gg =_u t_1 \downarrow_u \ll cs \gg_t$)
 by (simp add: trace-pred-R1-true, rel-auto)
 also have ... = ?rhs
 by (rel-auto)
 finally show ?thesis .
 qed

lemma wrC-csp-do-false [wp]:
 $\Phi(s_1, \sigma_1, t_1) \text{ wr}[cs]_C \text{ false} =$
 $(\forall (tt_0, tt_1) \cdot \mathcal{I}(s_1 \wedge \ll tt_1 \gg \in_u \ll tt_0 \gg \star_{cs} t_1 \wedge \ll tt_0 \gg \downarrow_u \ll cs \gg =_u t_1 \downarrow_u \ll cs \gg, \ll tt_1 \gg))$
 (is ?lhs = ?rhs)

proof –
 have ?lhs = $\Phi(s_1, \sigma_1, t_1) \text{ wr}[cs]_C \mathcal{I}(\text{true}, \langle \rangle)$
 by (simp add: rpred)
 also have ... = ?rhs
 by (simp add: wp)
 finally show ?thesis .
 qed

lemma wrC-csp-enable-init [wp]:
 fixes $t_1 \ t_2 :: ('a \text{ list}, 'b) \text{ uexpr}$
 shows
 $\mathcal{E}(s_1, t_1, E_1) \text{ wr}[cs]_C \mathcal{I}(s_2, t_2) =$
 $(\forall (tt_0, tt_1) \cdot \mathcal{I}(s_1 \wedge s_2 \wedge \ll tt_1 \gg \in_u (t_2 \hat{ }_u \ll tt_0 \gg) \star_{cs} t_1 \wedge t_2 \hat{ }_u \ll tt_0 \gg \downarrow_u \ll cs \gg =_u t_1 \downarrow_u \ll cs \gg, \ll tt_1 \gg))$
 (is ?lhs = ?rhs)

proof –
 have ?lhs =
 $(\neg_r (\exists (ref_0, ref_1, st_0, st_1 :: 'b,$
 $tt_0) \cdot [s_1]_{S<} \wedge$
 $[\$ref' \mapsto_s \ll ref_0 \gg, \$st' \mapsto_s \ll st_0 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll tt_0 \gg] \dagger (\neg_r \mathcal{I}(s_2, t_2)) \wedge$
 $(\forall e \cdot \ll e \gg \in_u [E_1]_{S<} \Rightarrow \ll e \gg \notin_u \ll ref_1 \gg) \wedge$
 $\$ref' \subseteq_u (\ll ref_0 \gg \cup_u \ll ref_1 \gg) \cap_u \ll cs \gg \cup_u (\ll ref_0 \gg \cap_u \ll ref_1 \gg - \ll cs \gg) \wedge$
 $[\ll trace \gg \in_u \ll tt_0 \gg \star_{cs} t_1 \wedge \ll tt_0 \gg \downarrow_u \ll cs \gg =_u t_1 \downarrow_u \ll cs \gg]_t \wedge \$st' =_u \$st) ;;_h$
 $R1 \text{ true})$
 by (simp add: wrR-def par-by-merge-seq-remove merge-csp-enable-right pr-var-def closure Healthy-if
 rpred)
 also have ... =
 $(\neg_r (\exists tt_0 \cdot ([s_2]_{S<} \wedge [t_2]_{S<} \leq_u \ll tt_0 \gg) \wedge [s_1]_{S<} \wedge$
 $[\ll trace \gg \in_u \ll tt_0 \gg \star_{cs} t_1 \wedge \ll tt_0 \gg \downarrow_u \ll cs \gg =_u t_1 \downarrow_u \ll cs \gg]_t) ;; R1 \text{ true})$
 by (rel-blast)
 also have ... =
 $(\neg_r (\exists tt_0 \cdot ([s_2]_{S<} \wedge (\exists tt_1 \cdot \ll tt_0 \gg =_u [t_2]_{S<} \hat{ }_u \ll tt_1 \gg)) \wedge [s_1]_{S<} \wedge$
 $[\ll trace \gg \in_u \ll tt_0 \gg \star_{cs} t_1 \wedge \ll tt_0 \gg \downarrow_u \ll cs \gg =_u t_1 \downarrow_u \ll cs \gg]_t) ;; R1 \text{ true})$
 by (simp add: utrace-leg)
 also have ... =
 $(\neg_r (\exists tt_1 \cdot [s_1 \wedge s_2 \wedge \ll trace \gg \in_u (t_2 \hat{ }_u \ll tt_1 \gg) \star_{cs} t_1 \wedge t_2 \hat{ }_u \ll tt_1 \gg \downarrow_u \ll cs \gg =_u t_1 \downarrow_u \ll cs \gg]_t)$
 $; R1 \text{ true})$
 by (rel-auto)
 also have ... =
 $(\forall tt_1 \cdot \neg_r ([s_1 \wedge s_2 \wedge \ll trace \gg \in_u (t_2 \hat{ }_u \ll tt_1 \gg) \star_{cs} t_1 \wedge t_2 \hat{ }_u \ll tt_1 \gg \downarrow_u \ll cs \gg =_u t_1 \downarrow_u \ll cs \gg]_t$
 $; R1 \text{ true}))$
 by (rel-auto)
 also have ... =
 $(\forall (tt_0, tt_1) \cdot \neg_r ([s_1 \wedge s_2 \wedge \ll tt_0 \gg \leq_u \ll trace \gg \wedge \ll tt_0 \gg \in_u (t_2 \hat{ }_u \ll tt_1 \gg) \star_{cs} t_1 \wedge t_2 \hat{ }_u \ll tt_1 \gg \downarrow_u$

$\ll cs \gg =_u t_1 \upharpoonright_u \ll cs \gg_t))$
 by (*simp add: trace-pred-R1-true, rel-auto*)
 also have ... = ?rhs
 by (*rel-auto*)
 finally show ?thesis .
 qed

lemma *wrC-csp-enable-false* [wp]:
 $\mathcal{E}(s_1, t_1, E) \text{ wr}[cs]_C \text{ false} =$
 $(\forall (tt_0, tt_1) \cdot \mathcal{I}(s_1 \wedge \ll tt_1 \gg \in_u \ll tt_0 \gg \star_{cs} t_1 \wedge \ll tt_0 \gg \upharpoonright_u \ll cs \gg =_u t_1 \upharpoonright_u \ll cs \gg, \ll tt_1 \gg))$
 (is ?lhs = ?rhs)
proof –
 have ?lhs = $\mathcal{E}(s_1, t_1, E) \text{ wr}[cs]_C \mathcal{I}(\text{true}, \langle \rangle)$
 by (*simp add: rpred*)
 also have ... = ?rhs
 by (*simp add: wp*)
 finally show ?thesis .
 qed

4.2 Parallel operator

syntax

-par-circus :: $\text{logic} \Rightarrow \text{salpha} \Rightarrow \text{logic} \Rightarrow \text{salpha} \Rightarrow \text{logic} \Rightarrow \text{logic} \quad (- \ll - \parallel - \gg - [75, 0, 0, 0, 76] \ 76)$
-par-csp :: $\text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \quad (- \ll - \parallel_C - [75, 0, 76] \ 76)$
-inter-circus :: $\text{logic} \Rightarrow \text{salpha} \Rightarrow \text{salpha} \Rightarrow \text{logic} \Rightarrow \text{logic} \quad (- \ll - \parallel - \gg - [75, 0, 0, 76] \ 76)$

translations

-par-circus $P \text{ ns1 } cs \text{ ns2 } Q == P \parallel_{M_C} \text{ ns1 } cs \text{ ns2 } Q$
-par-csp $P \text{ cs } Q == \text{-par-circus } P \ 0_L \text{ cs } 0_L \ Q$
-inter-circus $P \text{ ns1 } ns2 \ Q == \text{-par-circus } P \text{ ns1 } \{\} \text{ ns2 } Q$

abbreviation *InterleaveCSP* :: $('s, 'e) \text{ action} \Rightarrow ('s, 'e) \text{ action} \Rightarrow ('s, 'e) \text{ action} \quad (\text{infixr } \parallel 75)$
where $P \parallel Q \equiv P \ll \emptyset \parallel \emptyset \gg Q$

abbreviation *SynchroniseCSP* :: $('s, 'e) \text{ action} \Rightarrow ('s, 'e) \text{ action} \Rightarrow ('s, 'e) \text{ action} \quad (\text{infixr } \parallel 75)$
where $P \parallel Q \equiv P \ll UNIV \parallel_C Q$

definition *CSP5* :: $'\varphi \text{ process} \Rightarrow '\varphi \text{ process}$ **where**
 $[upred-defs]: \text{CSP5}(P) = (P \parallel \text{Skip})$

definition *C2* :: $('s, 'e) \text{ action} \Rightarrow ('s, 'e) \text{ action}$ **where**
 $[upred-defs]: \text{C2}(P) = (P \ll \Sigma \parallel \{\} \parallel \emptyset \gg \text{Skip})$

definition *CACT* :: $('s, 'e) \text{ action} \Rightarrow ('s, 'e) \text{ action}$ **where**
 $[upred-defs]: \text{CACT}(P) = \text{C2}(\text{NCSP}(P))$

abbreviation *CPROC* :: $'e \text{ process} \Rightarrow 'e \text{ process}$ **where**
 $\text{CPROC}(P) \equiv \text{CACT}(P)$

lemma *Skip-right-form*:

assumes $P_1 \text{ is } RC \ P_2 \text{ is } RR \ P_3 \text{ is } RR \ \$st' \ \# \ P_2$
 shows $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) ;; \text{Skip} = \mathbf{R}_s(P_1 \vdash P_2 \diamond (\exists \$ref' \cdot P_3))$

proof –

have $1:RR(P_3) ;; \Phi(\text{true}, id, \langle \rangle) = (\exists \$ref' \cdot RR(P_3))$
 by (*rel-auto*)
 show ?thesis

by (*rdes-simp cls: assms, metis 1 Healthy-if assms(3)*)
qed

lemma *ParCSP-rdes-def* [*rdes-def*]:

fixes $P_1 :: ('s, 'e)$ *action*

assumes P_1 is CRC Q_1 is CRC P_2 is CRR Q_2 is CRR P_3 is CRR Q_3 is CRR

$\$st' \# P_2 \$st' \# Q_2$

$ns1 \bowtie ns2$

shows $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \llbracket ns1 \parallel cs \parallel ns2 \rrbracket \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3) =$
 $\mathbf{R}_s(((Q_1 \Rightarrow_r Q_2) \text{ wr}[cs]_C P_1 \wedge (Q_1 \Rightarrow_r Q_3) \text{ wr}[cs]_C P_1 \wedge$
 $(P_1 \Rightarrow_r P_2) \text{ wr}[cs]_C Q_1 \wedge (P_1 \Rightarrow_r P_3) \text{ wr}[cs]_C Q_1) \vdash$
 $(P_2 \llbracket cs \rrbracket^I Q_2 \vee P_3 \llbracket cs \rrbracket^I Q_2 \vee P_2 \llbracket cs \rrbracket^I Q_3) \diamond$
 $(P_3 \llbracket ns1 \parallel cs \parallel ns2 \rrbracket^F Q_3))$

(**is** $?P \llbracket ns1 \parallel cs \parallel ns2 \rrbracket ?Q = ?rhs$)

proof –

have 1: $\bigwedge P Q. P \text{ wr}_R(N_C \ ns1 \ cs \ ns2) \ Q = P \text{ wr}[cs]_C \ Q \bigwedge P Q. P \text{ wr}_R(N_C \ ns2 \ cs \ ns1) \ Q = P \text{ wr}[cs]_C \ Q$

by (*rel-auto*) +

have 2: $(\exists \$st' \cdot N_C \ ns1 \ cs \ ns2) = (\exists \$st' \cdot N_C \ 0_L \ cs \ 0_L)$

by (*rel-auto*)

have $?P \llbracket ns1 \parallel cs \parallel ns2 \rrbracket ?Q = (?P \parallel_{M_R(N_C \ ns1 \ cs \ ns2)} ?Q) ;;_h \text{Skip}$

by (*simp add: CSPMerge-def par-by-merge-seq-add*)

also

have ... = $\mathbf{R}_s(((Q_1 \Rightarrow_r Q_2) \text{ wr}[cs]_C P_1 \wedge$
 $(Q_1 \Rightarrow_r Q_3) \text{ wr}[cs]_C P_1 \wedge$
 $(P_1 \Rightarrow_r P_2) \text{ wr}[cs]_C Q_1 \wedge$
 $(P_1 \Rightarrow_r P_3) \text{ wr}[cs]_C Q_1) \vdash$
 $(P_2 \llbracket cs \rrbracket^I Q_2 \vee$
 $P_3 \llbracket cs \rrbracket^I Q_2 \vee$
 $P_2 \llbracket cs \rrbracket^I Q_3) \diamond$
 $P_3 \parallel_{N_C \ ns1 \ cs \ ns2} Q_3) ;;_h \text{Skip}$

by (*simp add: parallel-rdes-def swap-CSPInnerMerge CSPInterMerge-def closure assms 1 2*)

also

have ... = $\mathbf{R}_s(((Q_1 \Rightarrow_r Q_2) \text{ wr}[cs]_C P_1 \wedge$
 $(Q_1 \Rightarrow_r Q_3) \text{ wr}[cs]_C P_1 \wedge$
 $(P_1 \Rightarrow_r P_2) \text{ wr}[cs]_C Q_1 \wedge$
 $(P_1 \Rightarrow_r P_3) \text{ wr}[cs]_C Q_1) \vdash$
 $(P_2 \llbracket cs \rrbracket^I Q_2 \vee$
 $P_3 \llbracket cs \rrbracket^I Q_2 \vee$
 $P_2 \llbracket cs \rrbracket^I Q_3) \diamond$
 $(\exists \$ref' \cdot (P_3 \parallel_{N_C \ ns1 \ cs \ ns2} Q_3)))$

by (*simp add: Skip-right-form closure parallel-RR-closed assms unrest*)

also

have ... = $\mathbf{R}_s(((Q_1 \Rightarrow_r Q_2) \text{ wr}[cs]_C P_1 \wedge$
 $(Q_1 \Rightarrow_r Q_3) \text{ wr}[cs]_C P_1 \wedge$
 $(P_1 \Rightarrow_r P_2) \text{ wr}[cs]_C Q_1 \wedge$
 $(P_1 \Rightarrow_r P_3) \text{ wr}[cs]_C Q_1) \vdash$
 $(P_2 \llbracket cs \rrbracket^I Q_2 \vee$
 $P_3 \llbracket cs \rrbracket^I Q_2 \vee$
 $P_2 \llbracket cs \rrbracket^I Q_3) \diamond$
 $(P_3 \llbracket ns1 \parallel cs \parallel ns2 \rrbracket^F Q_3))$

proof –

have $(\exists \$ref' \cdot (P_3 \parallel_{N_C \ ns1 \ cs \ ns2} Q_3)) = (P_3 \llbracket ns1 \parallel cs \parallel ns2 \rrbracket^F Q_3)$

by (*rel-blast*)

thus *?thesis* by *simp*

qed
 finally show ?thesis .
 qed

4.3 Parallel Laws

lemma *ParCSP-expand*:

$P \llbracket ns1 \parallel cs \parallel ns2 \rrbracket Q = (P \parallel_{RN_C} ns1 \ cs \ ns2 \ Q) \ ; \ ; \ Skip$
 by (simp add: CSPMerge-def par-by-merge-seq-add)

lemma *parallel-is-CSP* [closure]:

assumes P is CSP Q is CSP
 shows $(P \llbracket ns1 \parallel cs \parallel ns2 \rrbracket Q)$ is CSP

proof –

have $(P \parallel_{M_R(N_C \ ns1 \ cs \ ns2)} Q)$ is CSP
 by (simp add: closure assms)
 hence $(P \parallel_{M_R(N_C \ ns1 \ cs \ ns2)} Q) \ ; \ ; \ Skip$ is CSP
 by (simp add: closure)
 thus ?thesis
 by (simp add: CSPMerge-def par-by-merge-seq-add)

qed

lemma *parallel-is-NCSP* [closure]:

assumes $ns1 \bowtie ns2$ P is NCSP Q is NCSP
 shows $(P \llbracket ns1 \parallel cs \parallel ns2 \rrbracket Q)$ is NCSP

proof –

have $(P \llbracket ns1 \parallel cs \parallel ns2 \rrbracket Q) = (\mathbf{R}_s(pre_R \ P \vdash \ peri_R \ P \diamond \ post_R \ P) \llbracket ns1 \parallel cs \parallel ns2 \rrbracket \mathbf{R}_s(pre_R \ Q \vdash \ peri_R \ Q \diamond \ post_R \ Q))$
 $\diamond \ post_R \ Q))$
 by (metis NCSP-implies-NSRD NSRD-is-SRD SRD-reactive-design-alt assms wait'-cond-peri-post-cmt)
 also have ... is NCSP
 by (simp add: ParCSP-rdes-def assms closure unrest)
 finally show ?thesis .

qed

theorem *parallel-commutative*:

assumes $ns1 \bowtie ns2$
 shows $(P \llbracket ns1 \parallel cs \parallel ns2 \rrbracket Q) = (Q \llbracket ns2 \parallel cs \parallel ns1 \rrbracket P)$

proof –

have $(P \llbracket ns1 \parallel cs \parallel ns2 \rrbracket Q) = P \parallel_{swap_m} \ ; \ ; \ (M_C \ ns2 \ cs \ ns1) \ Q$
 by (simp add: CSPMerge-def segr-assoc[THEN sym] swap-merge-rd swap-CSPInnerMerge lens-indep-sym assms)
 also have ... = $Q \llbracket ns2 \parallel cs \parallel ns1 \rrbracket P$
 by (metis par-by-merge-commute-swap)
 finally show ?thesis .

qed

CSP5 is precisely *C2* when applied to a process

lemma *CSP5-is-C2*:

fixes $P :: 'e \text{ process}$
 assumes P is NCSP
 shows $CSP5(P) = C2(P)$
 unfolding CSP5-def C2-def by (rdes-eq cls: assms)

The form of C2 tells us that a normal CSP process has a downward closed set of refusals

lemma *csp-do-triv-merge*:

assumes P is CRF
shows $P \llbracket \Sigma[\{\}|\emptyset] \rrbracket^F \Phi(\text{true}, \text{id}, \langle \rangle) = P$ (**is** ?lhs = ?rhs)
proof –
 have ?lhs = $(\exists (st_0, t_0) \cdot [\$st' \mapsto_s \ll st_0 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_0 \gg] \dagger \text{CRF}(P) \wedge [\text{true}]_{S<} \wedge [\ll \text{trace} \gg =_u \ll t_0 \gg]_t \wedge \$st' =_u \$st \oplus \ll st_0 \gg \text{ on } \&\mathbf{v} \oplus \ll \text{id} \gg (\$st)_a \text{ on } \emptyset)$
 by (simp add: FinalMerge-csp-do-right assms closure unrest Healthy-if, rel-auto)
 also have ... = $\text{CRF}(P)$
 by (rel-auto)
 finally show ?thesis
 by (simp add: assms Healthy-if)
qed

lemma *csp-do-triv-wr*:

assumes P is CRC
shows $\Phi(\text{true}, \text{id}, \langle \rangle) \text{ wr}[\{\}]_C P = P$ (**is** ?lhs = ?rhs)
proof –
 have ?lhs = $(\neg_r (\exists (\text{ref}_0, st_0, tt_0) \cdot [\$ref' \mapsto_s \ll \text{ref}_0 \gg, \$st' \mapsto_s \ll st_0 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll tt_0 \gg] \dagger (\exists \$ref'; \$st' \cdot \text{RR}(\neg_r P)) \wedge \$ref' \subseteq_u \ll \text{ref}_0 \gg \wedge [\ll \text{trace} \gg =_u \ll tt_0 \gg]_t \wedge \$st' =_u \$st) ;; R1 \text{ true})$
 by (simp add: wrR-def par-by-merge-seq-remove rpred merge-csp-do-right ex-unrest Healthy-if pr-var-def closure assms unrest usubst)
 also have ... = $(\neg_r (\exists \$ref'; \$st' \cdot \text{RR}(\neg_r P)) ;; R1 \text{ true})$
 by (rel-auto, meson order-refl)
 also have ... = $(\neg_r (\neg_r P) ;; R1 \text{ true})$
 by (simp add: Healthy-if closure ex-unrest unrest assms)
 also have ... = P
 by (metis CRC-implies-RC Healthy-def RC1-def RC-implies-RC1 assms)
 finally show ?thesis .
qed

lemma *C2-form*:

assumes P is NCSP
shows $C2(P) = \mathbf{R}_s (\text{pre}_R P \vdash (\exists \text{ref}_0 \cdot \text{peri}_R P \llbracket \ll \text{ref}_0 \gg / \$ref' \rrbracket \wedge \$ref' \subseteq_u \ll \text{ref}_0 \gg) \diamond \text{post}_R P)$
proof –
 have 1: $\Phi(\text{true}, \text{id}, \langle \rangle) \text{ wr}[\{\}]_C \text{pre}_R P = \text{pre}_R P$ (**is** ?lhs = ?rhs)
 by (simp add: csp-do-triv-wr closure assms)
 have 2: $(\text{pre}_R P \Rightarrow_r \text{peri}_R P) \llbracket \{\} \rrbracket^I \Phi(\text{true}, \text{id}, \langle \rangle) = (\exists \text{ref}_0 \cdot (\text{peri}_R P) \llbracket \ll \text{ref}_0 \gg / \$ref' \rrbracket \wedge \$ref' \subseteq_u \ll \text{ref}_0 \gg) (\text{is } ?lhs = ?rhs)$
proof –
 have ?lhs = $\text{peri}_R P \llbracket \{\} \rrbracket^I \Phi(\text{true}, \text{id}, \langle \rangle)$
 by (simp add: SRD-peri-under-pre closure assms unrest)
 also have ... = $(\exists \$st' \cdot (\text{peri}_R P \parallel_{N_C} \emptyset_L \{\} \emptyset_L \Phi(\text{true}, \text{id}, \langle \rangle)))$
 by (simp add: CSPInterMerge-def par-by-merge-def segr-exists-right)
 also have ... = $(\exists \$st' \cdot \exists (\text{ref}_0, st_0, tt_0) \cdot [\$ref' \mapsto_s \ll \text{ref}_0 \gg, \$st' \mapsto_s \ll st_0 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll tt_0 \gg] \dagger (\exists \$st' \cdot \text{RR}(\text{peri}_R P)) \wedge \$ref' \subseteq_u \ll \text{ref}_0 \gg \wedge [\ll \text{trace} \gg =_u \ll tt_0 \gg]_t \wedge \$st' =_u \$st)$
 by (simp add: merge-csp-do-right pr-var-def assms Healthy-if closure rpred unrest ex-unrest)
 also have ... = $(\exists \text{ref}_0 \cdot (\exists \$st' \cdot \text{RR}(\text{peri}_R P)) \llbracket \ll \text{ref}_0 \gg / \$ref' \rrbracket \wedge \$ref' \subseteq_u \ll \text{ref}_0 \gg)$
 by (rel-auto)
 also have ... = ?rhs
 by (simp add: closure ex-unrest Healthy-if unrest assms)

finally show *?thesis* .
qed
have $3: (pre_R P \Rightarrow_r post_R P) \llbracket \Sigma | \{\} | \emptyset \rrbracket^F \Phi(true, id, \langle \rangle) = post_R(P)$ (**is** *?lhs = ?rhs*)
by (*simp add: csp-do-triv-merge SRD-post-under-pre unrest assms closure*)
show *?thesis*
proof –
have $C2(P) = \mathbf{R}_s (\Phi(true, id, \langle \rangle) wr[\{\}]_C pre_R P \vdash$
 $(pre_R P \Rightarrow_r peri_R P) \llbracket \{\} \rrbracket^I \Phi(true, id, \langle \rangle) \diamond (pre_R P \Rightarrow_r post_R P) \llbracket \Sigma | \{\} | \emptyset \rrbracket^F \Phi(true, id, \langle \rangle))$
by (*simp add: C2-def, rdes-simp cls: assms, simp add: id-def pr-var-def*)
also have $\dots = \mathbf{R}_s (pre_R P \vdash (\exists ref_0 \cdot peri_R P \llbracket \llbracket ref_0 \rrbracket / \$ref' \rrbracket \wedge \$ref' \subseteq_u \llbracket ref_0 \rrbracket) \diamond post_R P)$
by (*simp add: 1 2 3*)
finally show *?thesis* .
qed
qed

lemma *C2-CDC-form*:
assumes P *is NCSP*
shows $C2(P) = \mathbf{R}_s (pre_R P \vdash CDC(peri_R P) \diamond post_R P)$
by (*simp add: C2-form assms CDC-def*)

lemma *C2-rdes-def*:
assumes P_1 *is CRC* P_2 *is CRR* P_3 *is CRR* $\$st' \# P_2 \$ref' \# P_3$
shows $C2(\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3)) = \mathbf{R}_s(P_1 \vdash CDC(P_2) \diamond P_3)$
by (*simp add: C2-form assms closure rdes unrest usubst, rel-auto*)

lemma *C2-NCSP-intro*:
assumes P *is NCSP* $peri_R(P)$ *is CDC*
shows P *is C2*
proof –
have $C2(P) = \mathbf{R}_s (pre_R P \vdash CDC(peri_R P) \diamond post_R P)$
by (*simp add: C2-CDC-form assms(1)*)
also have $\dots = \mathbf{R}_s (pre_R P \vdash peri_R P \diamond post_R P)$
by (*simp add: Healthy-if assms*)
also have $\dots = P$
by (*simp add: NCSP-implies-CSP SRD-reactive-tri-design assms(1)*)
finally show *?thesis*
by (*simp add: Healthy-def*)
qed

lemma *C2-rdes-intro*:
assumes P_1 *is CRC* P_2 *is CRR* P_2 *is CDC* P_3 *is CRR* $\$st' \# P_2 \$ref' \# P_3$
shows $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3)$ *is C2*
unfolding *Healthy-def*
by (*simp add: C2-rdes-def assms unrest closure Healthy-if*)

lemma *C2-implies-CDC-peri [closure]*:
assumes P *is NCSP* P *is C2*
shows $peri_R(P)$ *is CDC*
proof –
have $peri_R(P) = peri_R (\mathbf{R}_s (pre_R P \vdash CDC(peri_R P) \diamond post_R P))$
by (*metis C2-CDC-form Healthy-if assms(1) assms(2)*)
also have $\dots = CDC (pre_R P \Rightarrow_r peri_R P)$
by (*simp add: rdes rpred assms closure unrest del: NSRD-peri-under-pre*)
also have $\dots = CDC (peri_R P)$
by (*simp add: SRD-peri-under-pre closure unrest assms*)

finally show ?thesis
 by (simp add: Healthy-def)
 qed

lemma CACT-intro:
 assumes P is NCSP P is C2
 shows P is CACT
 by (metis CACT-def Healthy-def assms(1) assms(2))

lemma CACT-rdes-intro:
 assumes P_1 is CRC P_2 is CRR P_2 is CDC P_3 is CRR $\$st' \# P_2 \$ref' \# P_3$
 shows $\mathbf{R}_s (P_1 \vdash P_2 \diamond P_3)$ is CACT
 by (rule CACT-intro, simp add: closure assms, rule C2-rdes-intro, simp-all add: assms)

lemma C2-NCSP-quasi-commute:
 assumes P is NCSP
 shows $C2(\text{NCSP}(P)) = \text{NCSP}(C2(P))$
 proof –
 have 1: $C2(\text{NCSP}(P)) = C2(P)$
 by (simp add: assms Healthy-if)
 also have ... = $\mathbf{R}_s (pre_R P \vdash CDC(per_i_R P) \diamond post_R P)$
 by (simp add: C2-CDC-form assms)
 also have ... is NCSP
 by (rule NCSP-rdes-intro, simp-all add: closure assms unrest)
 finally show ?thesis
 by (simp add: Healthy-if 1)
 qed

lemma C2-quasi-idem:
 assumes P is NCSP
 shows $C2(C2(P)) = C2(P)$
 proof –
 have $C2(C2(P)) = C2(C2(\mathbf{R}_s(pre_R(P) \vdash per_i_R(P) \diamond post_R(P))))$
 by (simp add: NCSP-implies-NSRD NSRD-is-SRD SRD-reactive-tri-design assms)
 also have ... = $\mathbf{R}_s (pre_R P \vdash CDC(per_i_R P) \diamond post_R P)$
 by (simp add: C2-rdes-def closure assms unrest CDC-idem)
 also have ... = $C2(P)$
 by (simp add: C2-CDC-form assms)
 finally show ?thesis .
 qed

lemma CACT-implies-NCSP [closure]:
 assumes P is CACT
 shows P is NCSP
 proof –
 have $P = C2(\text{NCSP}(\text{NCSP}(P)))$
 by (metis CACT-def Healthy-Idempotent Healthy-if NCSP-Idempotent assms)
 also have ... = $\text{NCSP}(C2(\text{NCSP}(P)))$
 by (simp add: C2-NCSP-quasi-commute Healthy-Idempotent NCSP-Idempotent)
 also have ... is NCSP
 by (metis CACT-def Healthy-def assms calculation)
 finally show ?thesis .
 qed

lemma CACT-implies-C2 [closure]:

assumes P is CACT
shows P is C2
by (metis CACT-def CACT-implies-NCSP Healthy-def assms)

lemma CACT-idem: $\text{CACT}(\text{CACT}(P)) = \text{CACT}(P)$
by (simp add: CACT-def C2-NCSP-quasi-commute[THEN sym] C2-quasi-idem Healthy-Idempotent Healthy-if NCSP-Idempotent)

lemma CACT-Idempotent: Idempotent CACT
by (simp add: CACT-idem Idempotent-def)

lemma PACT-elim [RD-elim]:
 $\llbracket X \text{ is CACT}; P(\mathbf{R}_s(\text{pre}_R(X) \vdash \text{peri}_R(X) \diamond \text{post}_R(X))) \rrbracket \implies P(X)$
using CACT-implies-NCSP NCSP-elim **by** blast

lemma Miracle-C2-closed [closure]: Miracle is C2
by (rdes-simp, rule C2-rdes-intro, simp-all add: closure unrest)

lemma Chaos-C2-closed [closure]: Chaos is C2
by (rdes-simp, rule C2-rdes-intro, simp-all add: closure unrest)

lemma Skip-C2-closed [closure]: Skip is C2
by (rdes-simp, rule C2-rdes-intro, simp-all add: closure unrest)

lemma Stop-C2-closed [closure]: Stop is C2
by (rdes-simp, rule C2-rdes-intro, simp-all add: closure unrest)

lemma Miracle-CACT-closed [closure]: Miracle is CACT
by (simp add: CACT-intro Miracle-C2-closed csp-theory.top-closed)

lemma Chaos-CACT-closed [closure]: Chaos is CACT
by (simp add: CACT-intro closure)

lemma Skip-CACT-closed [closure]: Skip is CACT
by (simp add: CACT-intro closure)

lemma Stop-CACT-closed [closure]: Stop is CACT
by (simp add: CACT-intro closure)

lemma seq-C2-closed [closure]:
assumes P is NCSP P is C2 Q is NCSP Q is C2
shows $P ;; Q$ is C2
by (rdes-simp cls: assms(1,3), rule C2-rdes-intro, simp-all add: closure assms unrest)

lemma seq-CACT-closed [closure]:
assumes P is CACT Q is CACT
shows $P ;; Q$ is CACT
by (meson CACT-implies-C2 CACT-implies-NCSP CACT-intro assms csp-theory.Healthy-Sequence seq-C2-closed)

lemma AssignsCSP-C2 [closure]: $\langle \sigma \rangle_C$ is C2
by (rdes-simp, rule C2-rdes-intro, simp-all add: closure unrest)

lemma AssignsCSP-CACT [closure]: $\langle \sigma \rangle_C$ is CACT
by (simp add: CACT-intro closure)

lemma *map-st-ext-CDC-closed* [closure]:

assumes P is CDC

shows $P \oplus_r \text{map-st}_L[a]$ is CDC

proof –

have $\text{CDC } P \oplus_r \text{map-st}_L[a]$ is CDC

by (*rel-auto*)

thus ?thesis

by (*simp add: assms Healthy-if*)

qed

lemma *rdes-frame-ext-C2-closed* [closure]:

assumes P is NCSP P is C2

shows $a:[P]_R^+$ is C2

by (*rdes-simp cls:assms(2), rule C2-rdes-intro, simp-all add: closure assms unrest*)

lemma *rdes-frame-ext-CACT-closed* [closure]:

assumes *vwb-lens* a P is CACT

shows $a:[P]_R^+$ is CACT

by (*rule CACT-intro, simp-all add: closure assms*)

lemma *UINF-C2-closed* [closure]:

assumes $A \neq \{\}$ $\bigwedge i. i \in A \implies P(i)$ is NCSP $\bigwedge i. i \in A \implies P(i)$ is C2

shows $(\bigcap i \in A \cdot P(i))$ is C2

proof –

have $(\bigcap i \in A \cdot P(i)) = (\bigcap i \in A \cdot \mathbf{R}_s(\text{pre}_R(P(i)) \vdash \text{peri}_R(P(i)) \diamond \text{post}_R(P(i))))$

by (*simp add: closure SRD-reactive-tri-design assms cong: UINF-cong*)

also have ... is C2

by (*rdes-simp cls: assms, rule C2-rdes-intro, simp-all add: closure unrest assms*)

finally show ?thesis .

qed

lemma *UINF-CACT-closed* [closure]:

assumes $A \neq \{\}$ $\bigwedge i. i \in A \implies P(i)$ is CACT

shows $(\bigcap i \in A \cdot P(i))$ is CACT

by (*rule CACT-intro, simp-all add: assms closure*)

lemma *inf-C2-closed* [closure]:

assumes P is NCSP Q is NCSP P is C2 Q is C2

shows $P \sqcap Q$ is C2

by (*rdes-simp cls: assms, rule C2-rdes-intro, simp-all add: closure unrest assms*)

lemma *cond-CDC-closed* [closure]:

assumes P is CDC Q is CDC

shows $P \triangleleft b \triangleright_R Q$ is CDC

proof –

have $\text{CDC } P \triangleleft b \triangleright_R \text{CDC } Q$ is CDC

by (*rel-auto*)

thus ?thesis

by (*simp add: Healthy-if assms*)

qed

lemma *cond-C2-closed* [closure]:

assumes P is NCSP Q is NCSP P is C2 Q is C2

shows $P \triangleleft b \triangleright_R Q$ is C2

by (rdes-simp cls: assms, rule C2-rdes-intro, simp-all add: closure unrest assms)

lemma cond-CACT-closed [closure]:
 assumes P is CACT Q is CACT
 shows $P \triangleleft b \triangleright_R Q$ is CACT
 by (rule CACT-intro, simp-all add: assms closure)

lemma gcomm-C2-closed [closure]:
 assumes P is NCSP P is C2
 shows $b \rightarrow_R P$ is C2
 by (rdes-simp cls: assms, rule C2-rdes-intro, simp-all add: closure unrest assms)

lemma AssumeCircus-CACT [closure]: $[b]_C$ is CACT
 by (metis AssumeCircus-NCSP AssumeCircus-def CACT-intro NCSP-Skip Skip-C2-closed gcomm-C2-closed)

lemma StateInvR-CACT [closure]: $\text{inv}_R(b)$ is CACT
 by (simp add: CACT-rdes-intro rdes-def closure unrest)

lemma AlternateR-C2-closed [closure]:
 assumes
 $\bigwedge i. i \in A \implies P(i)$ is NCSP Q is NCSP
 $\bigwedge i. i \in A \implies P(i)$ is C2 Q is C2
 shows $(\text{if}_R i \in A \cdot g(i) \rightarrow P(i) \text{ else } Q \text{ fi})$ is C2
proof (cases $A = \{\}$)
 case True
 then show ?thesis
 by (simp add: assms(4))
 next
 case False
 then show ?thesis
 by (simp add: AlternateR-def closure assms)
qed

lemma AlternateR-CACT-closed [closure]:
 assumes $\bigwedge i. i \in A \implies P(i)$ is CACT Q is CACT
 shows $(\text{if}_R i \in A \cdot g(i) \rightarrow P(i) \text{ else } Q \text{ fi})$ is CACT
 by (rule CACT-intro, simp-all add: closure assms)

lemma AlternateR-list-C2-closed [closure]:
 assumes
 $\bigwedge b P. (b, P) \in \text{set } A \implies P$ is NCSP Q is NCSP
 $\bigwedge b P. (b, P) \in \text{set } A \implies P$ is C2 Q is C2
 shows $(\text{AlternateR-list } A \ Q)$ is C2
 apply (simp add: AlternateR-list-def)
 apply (rule AlternateR-C2-closed)
 apply (auto simp add: assms closure)
 apply (metis assms nth-mem prod.collapse)+
 done

lemma AlternateR-list-CACT-closed [closure]:
 assumes $\bigwedge b P. (b, P) \in \text{set } A \implies P$ is CACT Q is CACT
 shows $(\text{AlternateR-list } A \ Q)$ is CACT
 by (rule CACT-intro, simp-all add: closure assms)

lemma R4-CRR-closed [closure]: P is CRR $\implies R4(P)$ is CRR

by (rule CRR-intro, simp-all add: closure unrest R4-def)

lemma *WhileC-C2-closed* [closure]:

assumes *P* is NCSP *P* is Productive *P* is C2

shows *while_C b do P od* is C2

proof –

have *while_C b do P od* = *while_C b do Productive(R_s (pre_R P ⊢ peri_R P ◊ post_R P)) od*

by (simp add: assms Healthy-if SRD-reactive-tri-design closure)

also have ... = *while_C b do R_s (pre_R P ⊢ peri_R P ◊ R4(post_R P)) od*

by (simp add: Productive-RHS-design-form unrest assms rdes closure R4-def)

also have ... is C2

by (simp add: WhileC-def, simp add: closure assms unrest rdes-def C2-rdes-intro)

finally show ?thesis .

qed

lemma *WhileC-CACT-closed* [closure]:

assumes *P* is CACT *P* is Productive

shows *while_C b do P od* is CACT

using CACT-implies-C2 CACT-implies-NCSP CACT-intro WhileC-C2-closed WhileC-NCSP-closed
assms by blast

lemma *IterateC-C2-closed* [closure]:

assumes

$\bigwedge i. i \in A \implies P(i) \text{ is NCSP } \bigwedge i. i \in A \implies P(i) \text{ is Productive } \bigwedge i. i \in A \implies P(i) \text{ is C2}$

shows (*do_C i ∈ A · g(i) → P(i) od*) is C2

unfolding *IterateC-def* by (simp add: closure assms)

lemma *IterateC-CACT-closed* [closure]:

assumes

$\bigwedge i. i \in A \implies P(i) \text{ is CACT } \bigwedge i. i \in A \implies P(i) \text{ is Productive}$

shows (*do_C i ∈ A · g(i) → P(i) od*) is CACT

by (metis CACT-implies-C2 CACT-implies-NCSP CACT-intro IterateC-C2-closed IterateC-NCSP-closed
assms)

lemma *IterateC-list-C2-closed* [closure]:

assumes

$\bigwedge b P. (b, P) \in \text{set } A \implies P \text{ is NCSP}$

$\bigwedge b P. (b, P) \in \text{set } A \implies P \text{ is Productive}$

$\bigwedge b P. (b, P) \in \text{set } A \implies P \text{ is C2}$

shows (*IterateC-list A*) is C2

unfolding *IterateC-list-def*

by (rule IterateC-C2-closed, (metis assms atLeastLessThan-iff nth-map nth-mem prod.collapse)+)

lemma *IterateC-list-CACT-closed* [closure]:

assumes

$\bigwedge b P. (b, P) \in \text{set } A \implies P \text{ is CACT}$

$\bigwedge b P. (b, P) \in \text{set } A \implies P \text{ is Productive}$

shows (*IterateC-list A*) is CACT

by (metis CACT-implies-C2 CACT-implies-NCSP CACT-intro IterateC-list-C2-closed IterateC-list-NCSP-closed
assms)

lemma *GuardCSP-C2-closed* [closure]:

assumes *P* is NCSP *P* is C2

shows *g &_C P* is C2

by (rdes-simp cls: assms(1), rule C2-rdes-intro, simp-all add: closure assms unrest)

lemma *GuardCSP-CACT-closed* [closure]:
assumes P is CACT
shows $g \ \&_C \ P$ is CACT
by (rule CACT-intro, simp-all add: closure assms)

lemma *DoCSP-C2* [closure]:
 $do_C(a)$ is C2
by (rdes-simp, rule C2-rdes-intro, simp-all add: closure unrest)

lemma *DoCSP-CACT* [closure]:
 $do_C(a)$ is CACT
by (rule CACT-intro, simp-all add: closure)

lemma *PrefixCSP-C2-closed* [closure]:
assumes P is NCSP P is C2
shows $a \rightarrow_C P$ is C2
unfolding PrefixCSP-def **by** (metis DoCSP-C2 Healthy-def NCSP-DoCSP NCSP-implies-CSP assms seq-C2-closed)

lemma *PrefixCSP-CACT-closed* [closure]:
assumes P is CACT
shows $a \rightarrow_C P$ is CACT
using CACT-implies-C2 CACT-implies-NCSP CACT-intro NCSP-PrefixCSP PrefixCSP-C2-closed assms **by** blast

lemma *ExtChoice-C2-closed* [closure]:
assumes $\bigwedge i. i \in I \implies P(i)$ is NCSP $\bigwedge i. i \in I \implies P(i)$ is C2
shows $(\square i \in I \cdot P(i))$ is C2
proof (cases $I = \{\}$)
case True
then show ?thesis **by** (simp add: closure ExtChoice-empty)
next
case False
show ?thesis
by (rule C2-NCSP-intro, simp-all add: assms closure unrest periR-ExtChoice-ind' False)
qed

lemma *ExtChoice-CACT-closed* [closure]:
assumes $\bigwedge i. i \in I \implies P(i)$ is CACT
shows $(\square i \in I \cdot P(i))$ is CACT
by (rule CACT-intro, simp-all add: closure assms)

lemma *extChoice-C2-closed* [closure]:
assumes P is NCSP P is C2 Q is NCSP Q is C2
shows $P \sqcap Q$ is C2
proof –
have $P \sqcap Q = (\square I \in \{P, Q\} \cdot I)$
by (simp add: extChoice-def)
also have ... is C2
by (rule ExtChoice-C2-closed, auto simp add: assms)
finally show ?thesis .
qed

lemma *extChoice-CACT-closed* [closure]:

assumes P is CACT Q is CACT
shows $P \sqcap Q$ is CACT
by (rule CACT-intro, simp-all add: closure assms)

lemma *state-srea-C2-closed* [closure]:
assumes P is NCSP P is C2
shows $\text{state } 'a \cdot P$ is C2
by (rule C2-NCSP-intro, simp-all add: closure rdes assms)

lemma *state-srea-CACT-closed* [closure]:
assumes P is CACT
shows $\text{state } 'a \cdot P$ is CACT
by (rule CACT-intro, simp-all add: closure assms)

lemma *parallel-C2-closed* [closure]:
assumes $ns1 \bowtie ns2$ P is NCSP Q is NCSP P is C2 Q is C2
shows $(P \llbracket ns1 \parallel cs \parallel ns2 \rrbracket Q)$ is C2
proof –
have $(P \llbracket ns1 \parallel cs \parallel ns2 \rrbracket Q) = (\mathbf{R}_s(\text{pre}_R P \vdash \text{peri}_R P \diamond \text{post}_R P) \llbracket ns1 \parallel cs \parallel ns2 \rrbracket \mathbf{R}_s(\text{pre}_R Q \vdash \text{peri}_R Q \diamond \text{post}_R Q))$
by (metis NCSP-implies-NSRD NSRD-is-SRD SRD-reactive-design-alt assms wait'-cond-peri-post-cmt)
also have ... is C2
by (simp add: ParCSP-rdes-def C2-rdes-intro assms closure unrest)
finally show ?thesis .
qed

lemma *parallel-CACT-closed* [closure]:
assumes $ns1 \bowtie ns2$ P is CACT Q is CACT
shows $(P \llbracket ns1 \parallel cs \parallel ns2 \rrbracket Q)$ is CACT
by (meson CACT-implies-C2 CACT-implies-NCSP CACT-intro assms parallel-C2-closed parallel-is-NCSP)

lemma *RenameCSP-C2-closed* [closure]:
assumes P is NCSP P is C2
shows $P \langle f \rangle_C$ is C2
by (simp add: RenameCSP-def C2-rdes-intro RenameCSP-pre-CRC-closed closure assms unrest)

lemma *RenameCSP-CACT-closed* [closure]:
assumes P is CACT
shows $P \langle f \rangle_C$ is CACT
by (rule CACT-intro, simp-all add: closure assms)

This property depends on downward closure of the refusals

lemma *rename-extChoice-pre*:
assumes $\text{inj } f$ P is NCSP Q is NCSP P is C2 Q is C2
shows $(P \sqcap Q) \langle f \rangle_C = (P \langle f \rangle_C \sqcap Q \langle f \rangle_C)$
by (rdes-eq-split cls: assms)

lemma *rename-extChoice*:
assumes $\text{inj } f$ P is CACT Q is CACT
shows $(P \sqcap Q) \langle f \rangle_C = (P \langle f \rangle_C \sqcap Q \langle f \rangle_C)$
by (simp add: CACT-implies-C2 CACT-implies-NCSP assms rename-extChoice-pre)

lemma *interleave-commute*:
 $P \parallel Q = Q \parallel P$
by (auto intro: parallel-commutative zero-lens-indep)

lemma *interleave-unit*:
assumes P is CPROC
shows $P \parallel \text{Skip} = P$
by (*metis* CACT-implies-C2 CACT-implies-NCSP CSP5-def CSP5-is-C2 Healthy-if *assms*)

lemma *parallel-miracle*:
 P is NCSP \implies Miracle $\llbracket ns1 \parallel cs \parallel ns2 \rrbracket P = \text{Miracle}$
by (*simp add*: CSPMerge-def *par-by-merge-seq-add*[THEN *sym*] Miracle-parallel-left-zero Skip-right-unit closure)

lemma *parallel-assigns*:
assumes *vwb-lens* $ns1$ *vwb-lens* $ns2$ $ns1 \bowtie ns2$ $x \subseteq_L ns1$ $y \subseteq_L ns2$
shows $(x :=_C u) \llbracket ns1 \parallel cs \parallel ns2 \rrbracket (y :=_C v) = x, y :=_C u, v$
using *assms* **by** (*rdes-eq*)

definition *Accept* :: ('s, 'e) action **where**
 $[upred-defs, rdes-def]: \text{Accept} = \mathbf{R}_s(\text{true}_r \vdash \mathcal{E}(\text{true}, \langle \rangle, \ll UNIV \gg) \diamond \text{false})$

definition [*upred-defs*, *rdes-def*]: $\text{CACC}(P) = (P \vee \text{Accept})$

lemma *CACC-form*:
assumes P_1 is RC P_2 is RR $\$st' \# P_2$ P_3 is RR
shows $\text{CACC}(\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3)) = \mathbf{R}_s(P_1 \vdash (\mathcal{E}(\text{true}, \langle \rangle, \ll UNIV \gg) \vee P_2) \diamond P_3)$
by (*rdes-eq cls: assms*)

lemma *CACC-refines-Accept*:
assumes P is CACC
shows $P \sqsubseteq \text{Accept}$
proof –
have $\text{CACC}(P) \sqsubseteq \text{Accept}$ **by** *rel-auto*
thus ?thesis **by** (*simp add: Healthy-if assms*)
qed

lemma *DoCSP-CACC* [*closure*]: $\text{do}_C(e)$ is CACC
unfolding *Healthy-def* **by** (*rdes-eq*)

lemma *CACC-seq-closure-left* [*closure*]:
assumes P is NCSP P is CACC Q is NCSP
shows $(P ;; Q)$ is CACC
proof –
have 1: $(P ;; Q) = \text{CACC}(P) ;; Q$
by (*simp add: Healthy-if assms*(2))
also have 2: $\dots = \mathbf{R}_s((pre_R P \wedge post_R P \text{ wp}_r pre_R Q) \vdash (peri_R P \vee \mathcal{E}(\text{true}, \langle \rangle, \ll UNIV \gg) \vee post_R P ;; peri_R Q) \diamond post_R P ;; post_R Q)$
by (*rdes-simp cls: assms*)
also have $\dots = \text{CACC}(\dots)$
by (*rdes-eq cls: assms*)
also have $\dots = \text{CACC}(P ;; Q)$
by (*simp add: 1 2*)
finally show ?thesis
by (*simp add: Healthy-def*)
qed

lemma *CACC-seq-closure-right*:

assumes P is NCSP $P \parallel Chaos = Chaos$ Q is NCSP Q is CACC
shows $(P \parallel Q)$ is CACC
oops

lemma *Chaos-is-CACC [closure]*: $Chaos$ is CACC

unfolding *Healthy-def* **by** (*rdes-eq*)

lemma *intChoice-is-CACC [closure]*:

assumes P is NCSP P is CACC Q is NCSP Q is CACC
shows $P \sqcap Q$ is CACC

proof –

have $CACC(P) \sqcap CACC(Q)$ is CACC
unfolding *Healthy-def* **by** (*rdes-eq cls: assms*)
thus *?thesis*
by (*simp add: Healthy-if assms(2) assms(4)*)

qed

lemma *extChoice-is-CACC [closure]*:

assumes P is NCSP P is CACC Q is NCSP Q is CACC
shows $P \sqcup Q$ is CACC

proof –

have $CACC(P) \sqcup CACC(Q)$ is CACC
unfolding *Healthy-def* **by** (*rdes-eq cls: assms*)
thus *?thesis*
by (*simp add: Healthy-if assms(2) assms(4)*)

qed

lemma *Chaos-par-zero*:

assumes P is NCSP P is CACC $ns1 \bowtie ns2$
shows $Chaos \llbracket ns1 \parallel cs \parallel ns2 \rrbracket P = Chaos$

proof –

have *pprop*: $(\forall (tt_0, tt_1) \cdot \mathcal{I}(\langle \langle tt_1 \rangle \rangle \in_u \langle \langle tt_0 \rangle \rangle \star_{cs} \langle \rangle \wedge \langle \langle tt_0 \rangle \rangle \upharpoonright_u \langle \langle cs \rangle \rangle =_u \langle \rangle \upharpoonright_u \langle \langle cs \rangle \rangle, \langle \langle tt_1 \rangle \rangle)) = false$
by (*rel-simp, auto simp add: tr-par-empty*)
(metis append-Nil2 seq-filter-Nil takeWhile.simps(1))

have $1:P = \mathbf{R}_s(pre_R(P) \vdash peri_R(P) \diamond post_R(P))$
by (*simp add: NCSP-implies-NSRD NSRD-is-SRD SRD-reactive-tri-design assms(1)*)

have $\dots \sqsubseteq \mathbf{R}_s(true_r \vdash \mathcal{E}(true, \langle \rangle, \langle \langle UNIV \rangle \rangle) \diamond false)$
by (*metis 1 Accept-def CACC-refines-Accept assms(2)*)

hence $peri_R P \sqsubseteq (pre_R P \wedge \mathcal{E}(true, \langle \rangle, \langle \langle UNIV \rangle \rangle))$
by (*auto simp add: RHS-tri-design-refine' closure assms*)

hence $peri_R(P) = ((pre_R P \wedge \mathcal{E}(true, \langle \rangle, \langle \langle UNIV \rangle \rangle)) \vee peri_R(P))$
by (*simp add: assms(2) utp-pred-laws.sup.absorb2*)

with 1 have $P = \mathbf{R}_s(pre_R(P) \vdash (pre_R(P) \wedge \mathcal{E}(true, \langle \rangle, \langle \langle UNIV \rangle \rangle) \vee peri_R(P)) \diamond post_R(P))$
by (*simp add: NCSP-implies-CSP SRD-reactive-tri-design assms(1)*)

also have $\dots = \mathbf{R}_s(pre_R(P) \vdash (\mathcal{E}(true, \langle \rangle, \langle \langle UNIV \rangle \rangle) \vee peri_R(P)) \diamond post_R(P))$
by (*rel-auto*)

also have $Chaos \llbracket ns1 \parallel cs \parallel ns2 \rrbracket \dots = Chaos$
 by (*rdes-simp cls: assms, simp add: pprop*)

finally show *?thesis* .

qed

lemma *Chaos-inter-zero*:

assumes *P is NCSP P is CACC*

shows $Chaos \parallel P = Chaos$

by (*simp add: Chaos-par-zero assms(1) assms(2)*)

end

5 Hiding

theory *utp-circus-hiding*

imports *utp-circus-parallel*

begin

5.1 Hiding in peri- and postconditions

definition *hide-rea* (*hide_r*) where

[*upred-defs*]: $hide_r P E = (\exists s \cdot (P \llbracket str^u \ll s \gg, (\ll E \gg \cup_u \$ref') / \$tr', \$ref' \rrbracket \wedge \$tr' =_u str^u (\ll s \gg \downarrow_u \ll -E \gg)))$

lemma *hide-rea-CRR-closed* [*closure*]:

assumes *P is CRR*

shows *hide_r P E is CRR*

proof –

have $CRR(hide_r (CRR P) E) = hide_r (CRR P) E$

by (*rel-auto, fastforce+*)

thus *?thesis*

by (*metis Healthy-def' assms*)

qed

lemma *hide-rea-CDC* [*closure*]:

assumes *P is CDC*

shows *hide_r P E is CDC*

proof –

have $CDC(hide_r (CDC P) E) = hide_r (CDC P) E$

by (*rel-blast*)

thus *?thesis*

by (*simp add: Healthy-if Healthy-intro assms*)

qed

lemma *hide-rea-false* [*rpred*]: $hide_r false E = false$

by (*rel-auto*)

lemma *hide-rea-disj* [*rpred*]: $hide_r (P \vee Q) E = (hide_r P E \vee hide_r Q E)$

by (*rel-auto*)

lemma *hide-rea-csp-enable* [*rpred*]:

$hide_r \mathcal{E}(s, t, E) F = \mathcal{E}(s \wedge E - \ll F \gg =_u E, t \downarrow_u \ll -F \gg, E)$

by (*rel-auto*)

lemma *hide-rea-csp-do* [*rpred*]: $hide_r \Phi(s, \sigma, t) E = \Phi(s, \sigma, t \downarrow_u \ll -E \gg)$

by (rel-auto)

lemma filter-eval [simp]:

(bop Cons x xs) \downarrow_u E = (bop Cons x (xs \downarrow_u E) \triangleleft x \in_u E \triangleright xs \downarrow_u E)

by (rel-simp)

lemma hide-rea-seq [rpred]:

assumes P is CRR \$ref' $\#$ P Q is CRR

shows $hide_r (P ;; Q) E = hide_r P E ;; hide_r Q E$

proof –

have $hide_r (CRR(\exists \$ref' \cdot P) ;; CRR(Q)) E = hide_r (CRR(\exists \$ref' \cdot P)) E ;; hide_r (CRR Q) E$

apply (simp add: hide-rea-def usubst unrest CRR-seqr-form)

apply (simp add: CRR-form)

apply (rel-auto)

using seq-filter-append apply fastforce

apply (metis seq-filter-append)

done

thus ?thesis

by (simp add: Healthy-if assms ex-unrest)

qed

lemma hide-rea-true-R1-true [rpred]:

$hide_r (R1 \text{ true}) A ;; R1 \text{ true} = R1 \text{ true}$

by (rel-auto, metis append-Nil2 seq-filter-Nil)

lemma hide-rea-shEx [rpred]: $hide_r (\exists i \cdot P(i)) cs = (\exists i \cdot hide_r (P i) cs)$

by (rel-auto)

lemma hide-rea-empty [rpred]:

assumes P is RR

shows $hide_r P \{\} = P$

proof –

have $hide_r (RR P) \{\} = (RR P)$

by (rel-auto; force)

thus ?thesis

by (simp add: Healthy-if assms)

qed

lemma hide-rea-twice [rpred]: $hide_r (hide_r P A) B = hide_r P (A \cup B)$

apply (rel-auto)

apply (metis (no-types, hide-lams) semilattice-sup-class.sup-assoc)

apply (metis (no-types, lifting) semilattice-sup-class.sup-assoc seq-filter-twice)

done

lemma st'-unrest-hide-rea [unrest]: $\$st' \# P \implies \$st' \# hide_r P E$

by (simp add: hide-rea-def unrest)

lemma ref'-unrest-hide-rea [unrest]: $\$ref' \# P \implies \$ref' \# hide_r P E$

by (simp add: hide-rea-def unrest usubst)

5.2 Hiding in preconditions

definition abs-rea :: ('s, 'e) action \Rightarrow 'e set \Rightarrow ('s, 'e) action (abs_r) **where**

[upred-defs]: $abs_r P E = (\neg_r (hide_r (\neg_r P) E ;; true_r))$

lemma abs-rea-false [rpred]: $abs_r \text{ false } E = \text{ false }$

by (*rel-simp*, *metis append.right-neutral seq-filter-Nil*)

lemma *abs-rea-conj* [*rpred*]: $abs_r (P \wedge Q) E = (abs_r P E \wedge abs_r Q E)$
 by (*rel-blast*)

lemma *abs-rea-true* [*rpred*]: $abs_r true_r E = true_r$
 by (*rel-auto*)

lemma *abs-rea-RC-closed* [*closure*]:
 assumes *P is CRR*
 shows $abs_r P E$ is *CRC*
proof –
 have *RC1* ($abs_r (CRR P) E$) = $abs_r (CRR P) E$
 apply (*rel-auto*)
 apply (*metis order-refl*)
 apply *blast*
 done
 hence $abs_r P E$ is *RC1*
 by (*simp add: assms Healthy-if Healthy-intro closure*)
 thus ?thesis
 by (*rule-tac CRC-intro'', simp-all add: abs-rea-def closure assms unrest*)
qed

lemma *hide-rea-impl-under-abs*:
 assumes *P is CRC Q is CRR*
 shows $(abs_r P A \Rightarrow_r hide_r (P \Rightarrow_r Q) A) = (abs_r P A \Rightarrow_r hide_r Q A)$
 by (*simp add: RC1-def abs-rea-def rea-impl-def rpred closure assms unrest*)
 (*rel-auto, metis order-refl*)

lemma *abs-rea-not-CRR*: P is *CRR* $\implies abs_r (\neg_r P) E = (\neg_r hide_r P E ;; R1 true)$
 by (*simp add: abs-rea-def rpred closure*)

lemma *abs-rea-wpR* [*rpred*]:
 assumes *P is CRR \$ref' \# P Q is CRC*
 shows $abs_r (P wp_r Q) A = (hide_r P A) wp_r (abs_r Q A)$
 by (*simp add: wp-rea-def abs-rea-not-CRR hide-rea-seq assms closure*)
 (*simp add: abs-rea-def rpred closure assms seqr-assoc*)

lemma *abs-rea-empty* [*rpred*]:
 assumes *P is RC*
 shows $abs_r P \{\} = P$
proof –
 have $abs_r (RC P) \{\} = (RC P)$
 apply (*rel-auto*)
 apply (*metis diff-add-cancel-left' order-refl plus-list-def*)
 using *dual-order.trans* apply *blast*
 done
 thus ?thesis
 by (*simp add: Healthy-if assms*)
qed

lemma *abs-rea-twice* [*rpred*]:
 assumes *P is CRC*
 shows $abs_r (abs_r P A) B = abs_r P (A \cup B)$ (*is ?lhs = ?rhs*)

proof –
have $?lhs = \text{abs}_r (\neg_r \text{hide}_r (\neg_r P) A ;; R1 \text{ true}) B$
by (*simp add: abs-rea-def*)
thus $?thesis$
by (*simp add: abs-rea-def rpred closure unrest segr-assoc assms*)
qed

5.3 Hiding Operator

In common with the UTP book definition of hiding, this definition does not introduce divergence if there is an infinite sequence of events that are hidden. For this, we would need a more complex precondition which is left for future work.

definition *HideCSP* :: $(s, e) \text{ action} \Rightarrow e \text{ set} \Rightarrow (s, e) \text{ action}$ (**infixl** \setminus_C 80) **where**
 $[upred-defs]:$
 $\text{HideCSP } P \ E = \mathbf{R}_s(\text{abs}_r(\text{pre}_R(P)) \ E \vdash \text{hide}_r(\text{peri}_R(P)) \ E \diamond \text{hide}_r(\text{post}_R(P)) \ E)$

lemma *HideCSP-rdes-def* [*rdes-def*]:
assumes $P \text{ is CRC } Q \text{ is CRR } R \text{ is CRR}$
shows $\mathbf{R}_s(P \vdash Q \diamond R) \setminus_C A = \mathbf{R}_s(\text{abs}_r(P) \ A \vdash \text{hide}_r \ Q \ A \diamond \text{hide}_r \ R \ A)$ (**is** $?lhs = ?rhs$)
proof –
have $?lhs = \mathbf{R}_s(\text{abs}_r \ P \ A \vdash \text{hide}_r \ (P \Rightarrow_r \ Q) \ A \diamond \text{hide}_r \ (P \Rightarrow_r \ R) \ A)$
by (*simp add: HideCSP-def rdes assms closure*)
also have $\dots = \mathbf{R}_s(\text{abs}_r \ P \ A \vdash (\text{abs}_r \ P \ A \Rightarrow_r \text{hide}_r \ (P \Rightarrow_r \ Q) \ A) \diamond (\text{abs}_r \ P \ A \Rightarrow_r \text{hide}_r \ (P \Rightarrow_r \ R) \ A))$
by (*metis RHS-tri-design-conj conj-idem utp-pred-laws.sup.idem*)
also have $\dots = ?rhs$
by (*metis RHS-tri-design-conj assms conj-idem hide-rea-impl-under-abs utp-pred-laws.sup.idem*)
finally show $?thesis$.
qed

lemma *HideCSP-NCSP-closed* [*closure*]: $P \text{ is NCSP} \Longrightarrow P \setminus_C E \text{ is NCSP}$
by (*simp add: HideCSP-def closure unrest*)

lemma *HideCSP-C2-closed* [*closure*]:
assumes $P \text{ is NCSP } P \text{ is C2}$
shows $P \setminus_C E \text{ is C2}$
by (*rdes-simp cls: assms, simp add: C2-rdes-intro closure unrest assms*)

lemma *HideCSP-CACT-closed* [*closure*]:
assumes $P \text{ is CACT}$
shows $P \setminus_C E \text{ is CACT}$
by (*rule CACT-intro, simp-all add: closure assms*)

lemma *HideCSP-Chaos*: $\text{Chaos} \setminus_C E = \text{Chaos}$
by (*rdes-simp*)

lemma *HideCSP-Miracle*: $\text{Miracle} \setminus_C A = \text{Miracle}$
by (*rdes-eq*)

lemma *HideCSP-AssignsCSP*:
 $\langle \sigma \rangle_C \setminus_C A = \langle \sigma \rangle_C$
by (*rdes-eq*)

lemma *HideCSP-cond*:
assumes $P \text{ is NCSP } Q \text{ is NCSP}$

shows $(P \triangleleft b \triangleright_R Q) \setminus_C A = (P \setminus_C A \triangleleft b \triangleright_R Q \setminus_C A)$
by (*rdes-eq cls: assms*)

lemma *HideCSP-int-choice*:

assumes P is NCSP Q is NCSP
shows $(P \sqcap Q) \setminus_C A = (P \setminus_C A \sqcap Q \setminus_C A)$
by (*rdes-eq cls: assms*)

lemma *HideCSP-guard*:

assumes P is NCSP
shows $(b \&_C P) \setminus_C A = b \&_C (P \setminus_C A)$
by (*rdes-eq cls: assms*)

lemma *HideCSP-seq*:

assumes P is NCSP Q is NCSP
shows $(P ;; Q) \setminus_C A = (P \setminus_C A ;; Q \setminus_C A)$
by (*rdes-eq-split cls: assms*)

lemma *HideCSP-DoCSP* [*rdes-def*]:

$do_C(a) \setminus_C A = (Skip \triangleleft (a \in_u \ll A \gg) \triangleright_R do_C(a))$
by (*rdes-eq*)

lemma *HideCSP-PrefixCSP*:

assumes P is NCSP
shows $(a \rightarrow_C P) \setminus_C A = ((P \setminus_C A) \triangleleft (a \in_u \ll A \gg) \triangleright_R (a \rightarrow_C (P \setminus_C A)))$
apply (*simp add: PrefixCSP-def Healthy-if HideCSP-seq HideCSP-DoCSP closure assms rdes rpred*)
apply (*simp add: HideCSP-NCSP-closed Skip-left-unit assms cond-st-distr*)
done

lemma *HideCSP-empty*:

assumes P is NCSP
shows $P \setminus_C \{\} = P$
by (*rdes-eq cls: assms*)

lemma *HideCSP-twice*:

assumes P is NCSP
shows $P \setminus_C A \setminus_C B = P \setminus_C (A \cup B)$
by (*rdes-simp cls: assms*)

lemma *HideCSP-Skip*: $Skip \setminus_C A = Skip$

by (*rdes-eq*)

lemma *HideCSP-Stop*: $Stop \setminus_C A = Stop$

by (*rdes-eq*)

end

6 Meta theory for Circus

theory *utp-circus*

imports

utp-circus-traces

utp-circus-parallel

utp-circus-hiding

begin end

7 Easy to use Circus-M parser

```
theory utp-circus-easy-parser
imports utp-circus UTP.utp-easy-parser
begin recall-syntax
```

We change := so that it refers to the Circus operator

```
no-adhoc-overloading
  uassigns assigns-r
```

```
adhoc-overloading
  uassigns AssignsCSP
```

```
syntax
  -GuardCSP :: uexp  $\Rightarrow$  logic  $\Rightarrow$  logic (infixr && 60)
```

```
no-translations
  -uwhile-top b P == CONST while-top b P
```

```
translations
  -uwhile-top b P == CONST WhileC b P
```

```
end
```

References

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