

# Kleene Algebra in Unifying Theories of Programming

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## Abstract

This development links Isabelle/UTP to the mechanised Kleene Algebra (KA) hierarchy for Isabelle/HOL. We substantiate the required KA laws, and provides a large body of additional theorems for alphabetised relations which are provided by the KA library. Additionally, we show how such theorems can be lifted to a subclass of UTP theories, provided certain conditions hold.

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## 1 Kleene Algebra and UTP

```
theory utp-kleene
  imports
    KAT-and-DRA.KAT
    UTP.utp
begin
```

This theory instantiates the Kleene Algebra [6] (KA) hierarchy, mechanised in Isabelle/HOL by Armstrong, Gomes, Struth et al [1, 4, 2]., for Isabelle/UTP alphabetised relations [3, 5]. Specifically, we substantiate the required dioid and KA laws in the type class hierarchy, which allows us to make use of all theorems proved in the former work. Moreover, we also prove an important result that a subclass of UTP theories, which we call “Kleene UTP theories”, always form Kleene algebras. The proof of the latter is obtained by lifting laws from the KA hierarchy.

### 1.1 Syntax setup

It is necessary to replace parts of the KA syntax to ensure compatibility with UTP. We therefore delete various bits of notation, and hide some constants.

```
purge-notation star (-* [101] 100)
```

```
recall-syntax
```

```
purge-notation n-op (n - [90] 91)
```

```
purge-notation ts-ord (infix  $\sqsubseteq$  50)
```

**notation**  $n\text{-op}$  ( $\mathbf{n}[-]$ )  
**notation**  $t$  ( $\mathbf{n}^2[-]$ )  
**notation**  $ts\text{-ord}$  (**infix**  $\sqsubseteq_t$  50)

**hide-const**  $t$

## 1.2 Kleene Algebra Instantiations

Next, import the laws of Kleene Algebra into the UTP relational calculus. We show that relations form a dioid and a Kleene algebra via two locales, the interpretation of which exports a large library of algebraic laws.

**interpretation** *urel-dioid*: *dioid*

**where**  $plus = op \sqcap$  **and**  $times = op ;;_h$  **and**  $less\text{-eq} = less\text{-eq}$  **and**  $less = less$

**proof**

**fix**  $P Q R :: 'a \text{ hrel}$   
**show**  $(P \sqcap Q) ;; R = P ;; R \sqcap Q ;; R$   
**by** (*simp add: upred-semiring.distrib-right*)  
**show**  $(Q \sqsubseteq P) = (P \sqcap Q = Q)$   
**by** (*simp add: semilattice-sup-class.le-iff-sup*)  
**show**  $(P < Q) = (Q \sqsubseteq P \wedge \neg P = Q)$   
**by** (*simp add: less-le*)  
**show**  $P \sqcap P = P$   
**by** *simp*

**qed**

**interpretation** *urel-ka*: *kleene-algebra*

**where**  $plus = op \sqcap$  **and**  $times = op ;;_h$  **and**  $one = skip\text{-}r$  **and**  $zero = false_h$  **and**  $less\text{-eq} = less\text{-eq}$   
**and**  $less = less$  **and**  $star = ustar$

**proof**

**fix**  $P Q R :: 'a \text{ hrel}$   
**show**  $II ;; P = P$  **by** *simp*  
**show**  $P ;; II = P$  **by** *simp*  
**show**  $false \sqcap P = P$  **by** *simp*  
**show**  $false ;; P = false$  **by** *simp*  
**show**  $P ;; false = false$  **by** *simp*  
**show**  $P^* \sqsubseteq II \sqcap P ;; P^*$   
**using** *ustar-sub-unfoldl* **by** *blast*  
**show**  $Q \sqsubseteq R \sqcap P ;; Q \implies Q \sqsubseteq P^* ;; R$   
**by** (*simp add: ustar-inductl*)  
**show**  $Q \sqsubseteq R \sqcap Q ;; P \implies Q \sqsubseteq R ;; P^*$   
**by** (*simp add: ustar-inductr*)

**qed**

We also show that UTP relations form a Kleene Algebra with Tests [7, 4] (KAT).

**interpretation** *urel-kat*: *kat*

**where**  $plus = op \sqcap$  **and**  $times = op ;;_h$  **and**  $one = skip\text{-}r$  **and**  $zero = false_h$  **and**  $less\text{-eq} = less\text{-eq}$   
**and**  $less = less$  **and**  $star = ustar$  **and**  $n\text{-op} = \lambda x. II \wedge (\neg x)$   
**by** (*unfold-locales, rel-auto+*)

We can now access the laws of KA and KAT for UTP relations as below.

**thm** *urel-ka.star-inductr-var*  
**thm** *urel-ka.star-trans*  
**thm** *urel-ka.star-square*  
**thm** *urel-ka.independence1*

### 1.3 Derived Laws

We prove that UTP assumptions are tests.

**lemma** *test-rassume* [simp]: *urel-kat.test*  $[b]^\top$   
 by (simp add: *urel-kat.test-def*, *rel-auto*)

The KAT laws can be used to prove results like the one below.

**lemma** *while-kat-form*:  
 $\text{while } b \text{ do } P \text{ od} = ([b]^\top ;; P)^* ;; [\neg b]^\top$  (is ?lhs = ?rhs)  
**proof** –  
 have  $1:(II::'a \text{ hrel}) \sqcap (II::'a \text{ hrel}) ;; [\neg b]^\top = II$   
 by (metis *assume-true test-rassume urel-kat.test-absorb1*)  
 have ?lhs =  $([b]^\top ;; P \sqcap [\neg b]^\top ;; II)^* ;; [\neg b]^\top$   
 by (simp add: *while-star-form rcond-rassume-expand*)  
 also have  $\dots = (([b]^\top ;; P)^* ;; [\neg b]^\top)^* ;; [\neg b]^\top$   
 by (metis *segr-right-unit urel-ka.star-denest*)  
 also have  $\dots = (([b]^\top ;; P)^* ;; (II \sqcap [\neg b]^\top)^*)^* ;; [\neg b]^\top$   
 by (metis *urel-ka.star2*)  
 also have  $\dots = (([b]^\top ;; P)^* ;; (II)^*)^* ;; [\neg b]^\top$   
 by (metis *1 segr-left-unit*)  
 also have  $\dots = (([b]^\top ;; P)^*)^* ;; [\neg b]^\top$   
 by (metis *urel-ka.mult-oner urel-ka.star-one*)  
 also have  $\dots = ?rhs$   
 by (metis *urel-ka.star-invol*)  
**finally show** ?thesis .  
**qed**

**lemma** *uplus-invol* [simp]:  $(P^+)^+ = P^+$   
 by (metis *RA1 uplus-def urel-ka.conway.dagger-trans-eq urel-ka.star-denest-var-2 urel-ka.star-invol*)

**lemma** *uplus-alt-def*:  $P^+ = P^* ;; P$   
 by (simp add: *uplus-def urel-ka.star-slide-var*)

### 1.4 UTP Theories with Kleene Algebra

A Kleene UTP theory is continuous UTP theory with a left unit. The star in such a context has already been defined by lifting the relational Kleene star. Here, we use the KA theorems obtained above to provide corresponding theorems for a Kleene UTP theory.

**locale** *utp-theory-kleene* = *utp-theory-cont-unital-zero1*  
**begin**

**lemma** *Star-def*:  $P\star = P^* ;; \mathcal{I}\mathcal{I}$   
 by (simp add: *utp-star-def*)

**lemma** *Star-alt-def*:  
 assumes  $P \text{ is } \mathcal{H}$   
 shows  $P\star = \mathcal{I}\mathcal{I} \sqcap P^+$   
**proof** –  
 from *assms* have  $P^+ = P^* ;; P ;; \mathcal{I}\mathcal{I}$   
 by (simp add: *Unit-Right uplus-alt-def*)  
 then show ?thesis  
 by (simp add: *RA1 utp-star-def*)  
**qed**

**lemma** *Star-Healthy* [closure]:

assumes  $P$  is  $\mathcal{H}$

shows  $P^\star$  is  $\mathcal{H}$

by (simp add: assms closure Star-alt-def)

**lemma** *Star-unfoldl*:

$P^\star \sqsubseteq \mathcal{II} \sqcap P$  ;;  $P^\star$

by (simp add: RA1 utp-star-def)

**lemma** *Star-inductl*:

assumes  $R$  is  $\mathcal{H}$   $Q \sqsubseteq P$  ;;  $Q \sqcap R$

shows  $Q \sqsubseteq P^\star$ ;; $R$

**proof** –

from assms(2) have  $Q \sqsubseteq R$   $Q \sqsubseteq P$  ;;  $Q$

by auto

thus ?thesis

by (simp add: Unit-Left assms(1) upred-semiring.mult-assoc urel-ka.star-inductl utp-star-def)

qed

**lemma** *Star-invol*:

assumes  $P$  is  $\mathcal{H}$

shows  $P^{\star\star} = P^\star$

by (metis (no-types) RA1 Unit-Left Unit-self assms urel-ka.star-invol urel-ka.star-sim3 utp-star-def)

**lemma** *Star-test*:

assumes  $P$  is  $\mathcal{H}$  utest  $\mathcal{T}$   $P$

shows  $P^\star = \mathcal{II}$

by (metis utp-star-def Star-alt-def Unit-Right Unit-self assms semilattice-sup-class.sup.absorb1 semilattice-sup-class.sup-urel-ka.star-inductr-var-eq2 urel-ka.star-sim1 utest-def)

**lemma** *Star-lemma-1*:

$P$  is  $\mathcal{H} \implies \mathcal{II}$  ;;  $P^\star$  ;;  $\mathcal{II} = P^\star$  ;;  $\mathcal{II}$

by (metis utp-star-def Star-Healthy Unit-Left)

**lemma** *Star-lemma-2*:

assumes  $P$  is  $\mathcal{H}$   $Q$  is  $\mathcal{H}$

shows  $(P^\star ;; Q^\star ;; \mathcal{II})^\star ;; \mathcal{II} = (P^\star ;; Q^\star)^\star ;; \mathcal{II}$

by (metis (no-types) assms RA1 Star-lemma-1 Unit-self urel-ka.star-sim3)

**lemma** *Star-denest*:

assumes  $P$  is  $\mathcal{H}$   $Q$  is  $\mathcal{H}$

shows  $(P \sqcap Q)^\star = (P^\star ;; Q^\star)^\star$

by (metis (no-types, lifting) RA1 utp-star-def Star-lemma-1 Star-lemma-2 assms urel-ka.star-denest)

**lemma** *Star-denest-disj*:

assumes  $P$  is  $\mathcal{H}$   $Q$  is  $\mathcal{H}$

shows  $(P \vee Q)^\star = (P^\star ;; Q^\star)^\star$

by (simp add: disj-upred-def Star-denest assms)

**lemma** *Star-unfoldl-eq*:

assumes  $P$  is  $\mathcal{H}$

shows  $\mathcal{II} \sqcap P$  ;;  $P^\star = P^\star$

by (simp add: RA1 utp-star-def)

```

lemma uplus-Star-def:
  assumes  $P$  is  $\mathcal{H}$ 
  shows  $P^+ = (P ;; P^\star)$ 
  by (metis (full-types) RA1 utp-star-def Unit-Left Unit-Right assms uplus-def urel-ka.conway.dagger-slide)

lemma Star-trade-skip:
   $P$  is  $\mathcal{H} \implies \mathcal{II} ;; P^\star = P^\star ;; \mathcal{II}$ 
  by (simp add: Unit-Left Unit-Right urel-ka.star-sim3)

lemma Star-slide:
  assumes  $P$  is  $\mathcal{H}$ 
  shows  $(P ;; P^\star) = (P^\star ;; P)$  (is  $?lhs = ?rhs$ )
proof –
  have  $?lhs = P ;; P^\star ;; \mathcal{II}$ 
    by (simp add: utp-star-def)
  also have  $\dots = P ;; \mathcal{II} ;; P^\star$ 
    by (simp add: Star-trade-skip assms)
  also have  $\dots = P ;; P^\star$ 
    by (simp add: RA1 Unit-Right assms)
  also have  $\dots = P^\star ;; P$ 
    by (simp add: urel-ka.star-slide-var)
  also have  $\dots = ?rhs$ 
    by (metis RA1 utp-star-def Unit-Left assms)
  finally show  $?thesis$  .
qed

lemma Star-unfoldr-eq:
  assumes  $P$  is  $\mathcal{H}$ 
  shows  $\mathcal{II} \sqcap P^\star ;; P = P^\star$ 
  using Star-slide Star-unfoldl-eq assms by auto

lemma Star-inductr:
  assumes  $P$  is  $\mathcal{H}$   $R$  is  $\mathcal{H}$   $Q \sqsubseteq P \sqcap Q ;; R$ 
  shows  $Q \sqsubseteq P ;; R^\star$ 
  by (metis (full-types) RA1 Star-def Star-trade-skip Unit-Right assms urel-ka.star-inductr')

lemma Star-Top:  $\top^\star = \mathcal{II}$ 
  by (simp add: Star-test top-healthy utest-Top)

end

end

```

## References

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