

# Isabelle/UTP: Mechanised reasoning for the UTP

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# 1 Parser Utilities

**theory** *utp-parser-utils*

**imports**

*Main*

**begin**

**syntax**

*-id-string*     :: *id*  $\Rightarrow$  *string* (*IDSTR'*(-))

**ML**  $\ll$

*signature* *UTP-PARSER-UTILS* =

*sig*

*val* *mk-nib* : *int*  $\rightarrow$  *Ast.ast*

*val* *mk-char* : *string*  $\rightarrow$  *Ast.ast*

*val* *mk-string* : *string list*  $\rightarrow$  *Ast.ast*

*val* *string-ast-tr* : *Ast.ast list*  $\rightarrow$  *Ast.ast*

*end*;

*structure* *Utp-Parser-Utils* : *UTP-PARSER-UTILS* =

*struct*

*val* *mk-nib* =

*Ast.Constant* *o* *Lexicon.mark-const* *o*

*fst* *o* *Term.dest-Const* *o* *HOLogic.mk-char*;

*fun* *mk-char* *s* =

*if* *Symbol.is-ascii* *s* *then*

*Ast.Appl* [*Ast.Constant* @{*const-syntax* *Char*}, *mk-nib* (*ord* *s* *div* 16), *mk-nib* (*ord* *s* *mod* 16)]

*else* *error* (*Non-ASCII symbol: ^ quote* *s*);

*fun* *mk-string* [] = *Ast.Constant* @{*const-syntax* *Nil*}

| *mk-string* (*c* :: *cs*) =

*Ast.Appl* [*Ast.Constant* @{*const-syntax* *List.Cons*}, *mk-char* *c*, *mk-string* *cs*];

*fun* *string-ast-tr* [*Ast.Variable* *str*] =

(*case* *Lexicon.explode-str* (*str*, *Position.none*) *of*

[] =>

*Ast.Appl*

[*Ast.Constant* @{*syntax-const* -*constrain*},

*Ast.Constant* @{*const-syntax* *Nil*}, *Ast.Constant* @{*type-syntax* *string*}]

| *ss* => *mk-string* (*map* *Symbol-Pos.symbol* *ss*))

```

| string-ast-tr [Ast.Appl [Ast.Constant @{syntax-const -constrain}, ast1, ast2]] =
  Ast.Appl [Ast.Constant @{syntax-const -constrain}, string-ast-tr [ast1], ast2]
| string-ast-tr asts = raise Ast.AST (string-tr, asts);

end

signature NAME-UTILS =
sig
  val deep-unmark-const : term -> term
  val right-crop-by : int -> string -> string
  val last-char-str : string -> string
  val repeat-char : char -> int -> string
  val mk-id : string -> term
end;

structure Name-Utils : NAME-UTILS =
struct
  fun unmark-const-term (Const (name, typ)) =
    Const (Lexicon.unmark-const name, typ)
  | unmark-const-term term = term;

  val deep-unmark-const =
    (map-aterms unmark-const-term);

  fun right-crop-by n s =
    String.substring (s, 0, (String.size s) - n);

  fun last-char-str s =
    String.str (String.sub (s, (String.size s) - 1));

  fun repeat-char c n =
    if n > 0 then (String.str c) ^ (repeat-char c (n - 1)) else ;

  fun mk-id name = Free (name, dummyT);
end;
>>

parse-translation <<
let
  fun id-string-tr [Free (full-name, -)] = HOLLogic.mk-string full-name
  | id-string-tr [Const (full-name, -)] = HOLLogic.mk-string full-name
  | id-string-tr - = raise Match;
in
  [(@{syntax-const -id-string}, K id-string-tr)]
end
>>
end

```

## 2 UTP variables

```

theory utp-var
imports
  ../utils/utp-imports
  utp-parser-utils
begin

```

In this first UTP theory we set up variable, which are built on lenses. A large part of this theory is setting up the parser for UTP variable syntax.

## 2.1 Initial syntax setup

We will overload the square order relation with refinement and also the lattice operators so we will turn off these notations.

### purge-notation

```
Order.le (infixl  $\sqsubseteq_1$  50) and
Lattice.sup ( $\sqcup_1$ - [90] 90) and
Lattice.inf ( $\sqcap_1$ - [90] 90) and
Lattice.join (infixl  $\sqcup_1$  65) and
Lattice.meet (infixl  $\sqcap_1$  70) and
LFP ( $\mu$ ) and
GFP ( $\nu$ ) and
Set.member (op :) and
Set.member ((-/ : -) [51, 51] 50) and
disj (infixr | 30)
```

We hide HOL's built-in relation type since we will replace it with our own

### hide-type rel

```
type-synonym 'a relation = ('a  $\times$  'a) set
```

```
declare fst-vwb-lens [simp]
declare snd-vwb-lens [simp]
declare comp-vwb-lens [simp]
declare lens-indep-left-ext [simp]
declare lens-indep-right-ext [simp]
```

## 2.2 Variable foundations

This theory describes the foundational structure of UTP variables, upon which the rest of our model rests. We start by defining alphabets, which following [3, 4] in this shallow model are simply represented as types  $'\alpha$ , though by convention usually a record type where each field corresponds to a variable. UTP variables in this frame are simply modelled as lenses  $'a \Longrightarrow '\alpha$ , where the view type  $'a$  is the variable type, and the source type  $'\alpha$  is the alphabet or state-space type.

We define some lifting functions for variables to create input and output variables. These simply lift the alphabet to a tuple type since relations will ultimately be defined by a tuple alphabet.

**definition**  $in-var :: ('a \Longrightarrow '\alpha) \Rightarrow ('a \Longrightarrow '\alpha \times '\beta)$  **where**

$[lens-defs]: in-var\ x = x ;_L\ fst_L$

**definition**  $out-var :: ('a \Longrightarrow '\beta) \Rightarrow ('a \Longrightarrow '\alpha \times '\beta)$  **where**

$[lens-defs]: out-var\ x = x ;_L\ snd_L$

Variables can also be used to effectively define sets of variables. Here we define the the universal alphabet ( $\Sigma$ ) to be the bijective lens  $1_L$ . This characterises the whole of the source type, and thus is effectively the set of all alphabet variables.

**abbreviation**  $(input)\ univ-alpha :: ('\alpha \Longrightarrow '\alpha)\ (\Sigma)$  **where**

$univ-alpha \equiv 1_L$

The next construct is vacuous and simply exists to help the parser distinguish predicate variables from input and output variables.

**definition**  $pr\text{-}var :: ('a \Longrightarrow 'b) \Rightarrow ('a \Longrightarrow 'b)$  **where**  
 $[lens\text{-}defs]: pr\text{-}var\ x = x$

### 2.3 Variable lens properties

We can now easily show that our UTP variable construction are various classes of well-behaved lens .

**lemma**  $in\text{-}var\text{-}semi\text{-}uvar$   $[simp]:$   
 $mwb\text{-}lens\ x \Longrightarrow mwb\text{-}lens\ (in\text{-}var\ x)$   
**by**  $(simp\ add: comp\text{-}mwb\text{-}lens\ in\text{-}var\text{-}def)$

**lemma**  $pr\text{-}var\text{-}mwb\text{-}lens$   $[simp]:$   
 $mwb\text{-}lens\ x \Longrightarrow mwb\text{-}lens\ (pr\text{-}var\ x)$   
**by**  $(simp\ add: pr\text{-}var\text{-}def)$

**lemma**  $pr\text{-}var\text{-}vwb\text{-}lens$   $[simp]:$   
 $vwb\text{-}lens\ x \Longrightarrow vwb\text{-}lens\ (pr\text{-}var\ x)$   
**by**  $(simp\ add: pr\text{-}var\text{-}def)$

**lemma**  $in\text{-}var\text{-}uvar$   $[simp]:$   
 $vwb\text{-}lens\ x \Longrightarrow vwb\text{-}lens\ (in\text{-}var\ x)$   
**by**  $(simp\ add: in\text{-}var\text{-}def)$

**lemma**  $out\text{-}var\text{-}semi\text{-}uvar$   $[simp]:$   
 $mwb\text{-}lens\ x \Longrightarrow mwb\text{-}lens\ (out\text{-}var\ x)$   
**by**  $(simp\ add: comp\text{-}mwb\text{-}lens\ out\text{-}var\text{-}def)$

**lemma**  $out\text{-}var\text{-}uvar$   $[simp]:$   
 $vwb\text{-}lens\ x \Longrightarrow vwb\text{-}lens\ (out\text{-}var\ x)$   
**by**  $(simp\ add: out\text{-}var\text{-}def)$

Moreover, we can show that input and output variables are independent, since they refer to different sections of the alphabet.

**lemma**  $in\text{-}out\text{-}indep$   $[simp]:$   
 $in\text{-}var\ x \bowtie out\text{-}var\ y$   
**by**  $(simp\ add: lens\text{-}indep\text{-}def\ in\text{-}var\text{-}def\ out\text{-}var\text{-}def\ fst\text{-}lens\text{-}def\ snd\text{-}lens\text{-}def\ lens\text{-}comp\text{-}def)$

**lemma**  $out\text{-}in\text{-}indep$   $[simp]:$   
 $out\text{-}var\ x \bowtie in\text{-}var\ y$   
**by**  $(simp\ add: lens\text{-}indep\text{-}def\ in\text{-}var\text{-}def\ out\text{-}var\text{-}def\ fst\text{-}lens\text{-}def\ snd\text{-}lens\text{-}def\ lens\text{-}comp\text{-}def)$

**lemma**  $in\text{-}var\text{-}indep$   $[simp]:$   
 $x \bowtie y \Longrightarrow in\text{-}var\ x \bowtie in\text{-}var\ y$   
**by**  $(simp\ add: in\text{-}var\text{-}def\ out\text{-}var\text{-}def)$

**lemma**  $out\text{-}var\text{-}indep$   $[simp]:$   
 $x \bowtie y \Longrightarrow out\text{-}var\ x \bowtie out\text{-}var\ y$   
**by**  $(simp\ add: out\text{-}var\text{-}def)$

**lemma**  $pr\text{-}var\text{-}indeps$   $[simp]:$   
 $x \bowtie y \Longrightarrow pr\text{-}var\ x \bowtie y$   
 $x \bowtie y \Longrightarrow x \bowtie pr\text{-}var\ y$   
**by**  $(simp\text{-}all\ add: pr\text{-}var\text{-}def)$

**lemma**  $prod\text{-}lens\text{-}indep\text{-}in\text{-}var$   $[simp]:$

$a \bowtie x \implies a \times_L b \bowtie \text{in-var } x$   
**by** (*metis in-var-def in-var-indep out-in-indep out-var-def plus-pres-lens-indep prod-as-plus*)

**lemma** *prod-lens-indep-out-var* [*simp*]:  
 $b \bowtie x \implies a \times_L b \bowtie \text{out-var } x$   
**by** (*metis in-out-indep in-var-def out-var-def out-var-indep plus-pres-lens-indep prod-as-plus*)

**lemma** *in-var-pr-var* [*simp*]:  
 $\text{in-var } (\text{pr-var } x) = \text{in-var } x$   
**by** (*simp add: pr-var-def*)

**lemma** *out-var-pr-var* [*simp*]:  
 $\text{out-var } (\text{pr-var } x) = \text{out-var } x$   
**by** (*simp add: pr-var-def*)

Similar properties follow for sublens

**lemma** *in-var-sublens* [*simp*]:  
 $y \subseteq_L x \implies \text{in-var } y \subseteq_L \text{in-var } x$   
**by** (*metis (no-types, hide-lams) in-var-def lens-comp-assoc sublens-def*)

**lemma** *out-var-sublens* [*simp*]:  
 $y \subseteq_L x \implies \text{out-var } y \subseteq_L \text{out-var } x$   
**by** (*metis (no-types, hide-lams) out-var-def lens-comp-assoc sublens-def*)

**lemma** *pr-var-sublens* [*simp*]:  
 $y \subseteq_L x \implies \text{pr-var } y \subseteq_L \text{pr-var } x$   
**by** (*simp add: pr-var-def*)

## 2.4 Lens simplifications

We also define some lookup abstraction simplifications.

**lemma** *var-lookup-in* [*simp*]:  $\text{lens-get } (\text{in-var } x) (A, A') = \text{lens-get } x A$   
**by** (*simp add: in-var-def fst-lens-def lens-comp-def*)

**lemma** *var-lookup-out* [*simp*]:  $\text{lens-get } (\text{out-var } x) (A, A') = \text{lens-get } x A'$   
**by** (*simp add: out-var-def snd-lens-def lens-comp-def*)

**lemma** *var-update-in* [*simp*]:  $\text{lens-put } (\text{in-var } x) (A, A') v = (\text{lens-put } x A v, A')$   
**by** (*simp add: in-var-def fst-lens-def lens-comp-def*)

**lemma** *var-update-out* [*simp*]:  $\text{lens-put } (\text{out-var } x) (A, A') v = (A, \text{lens-put } x A' v)$   
**by** (*simp add: out-var-def snd-lens-def lens-comp-def*)

## 2.5 Syntax translations

In order to support nice syntax for variables, we here set up some translations. The first step is to introduce a collection of non-terminals.

**nonterminal** *svid* **and** *svids* **and** *svar* **and** *svars* **and** *salpha*

These non-terminals correspond to the following syntactic entities. Non-terminal *svid* is an atomic variable identifier, and *svids* is a list of identifier. *svar* is a decorated variable, such as an input or output variable, and *svars* is a list of decorated variables. *salpha* is an alphabet or set of variables. Such sets can be constructed only through lens composition due to typing restrictions. Next we introduce some syntax constructors.



### **syntax** — Identifiers

*-svid* ::  $id \Rightarrow svid$  (- [999] 999)  
*-svid-unit* ::  $svid \Rightarrow svids$  (-)  
*-svid-list* ::  $svid \Rightarrow svids \Rightarrow svids$  (-, / -)  
*-svid-alpha* ::  $svid$  (**v**)  
*-svid-dot* ::  $svid \Rightarrow svid \Rightarrow svid$  (:- [999,998] 999)

A variable identifier can either be a HOL identifier, the complete set of variables in the alphabet **v**, or a composite identifier separated by colons, which corresponds to a sort of qualification. The final option is effectively a lens composition.

### **syntax** — Decorations

*-spvar* ::  $svid \Rightarrow svar$  (&- [998] 998)  
*-sinvar* ::  $svid \Rightarrow svar$  (\$- [998] 998)  
*-soutvar* ::  $svid \Rightarrow svar$  (\$-' [998] 998)

A variable can be decorated with an ampersand, to indicate it is a predicate variable, with a dollar to indicate its an unprimed relational variable, or a dollar and “acute” symbol to indicate its a primed relational variable. Isabelle’s parser is extensible so additional decorations can be and are added later.

### **syntax** — Variable sets

*-salphaid* ::  $id \Rightarrow salpha$  (- [998] 998)  
*-salphavar* ::  $svar \Rightarrow salpha$  (- [998] 998)  
*-salphaparen* ::  $salpha \Rightarrow salpha$  ('(-'))  
*-salphacomp* ::  $salpha \Rightarrow salpha \Rightarrow salpha$  (**infixr** ; 75)  
*-salphaprod* ::  $salpha \Rightarrow salpha \Rightarrow salpha$  (**infixr**  $\times$  85)  
*-salpha-all* ::  $salpha$  ( $\Sigma$ )  
*-salpha-none* ::  $salpha$  ( $\emptyset$ )  
*-svar-nil* ::  $svar \Rightarrow svars$  (-)  
*-svar-cons* ::  $svar \Rightarrow svars \Rightarrow svars$  (-, / -)  
*-salphaset* ::  $svars \Rightarrow salpha$  ({-})  
*-salphamk* ::  $logic \Rightarrow salpha$

The terminals of an alphabet are either HOL identifiers or UTP variable identifiers. We support two ways of constructing alphabets; by composition of smaller alphabets using a semi-colon or by a set-style construction  $\{a, b, c\}$  with a list of UTP variables.

### **syntax** — Quotations

*-ualpha-set* ::  $svars \Rightarrow logic$  ({-} <sub>$\alpha$</sub> )  
*-svar* ::  $svar \Rightarrow logic$  ('(-') <sub>$v$</sub> )

For various reasons, the syntax constructors above all yield specific grammar categories and will not parser at the HOL top level (basically this is to do with us wanting to reuse the syntax for expressions). As a result we provide some quotation constructors above.

Next we need to construct the syntax translations rules. First we need a few polymorphic constants.

### **consts**

*svar* ::  $'v \Rightarrow 'e$   
*ivar* ::  $'v \Rightarrow 'e$   
*ovar* ::  $'v \Rightarrow 'e$

### **ad hoc-overloading**

*svar* *pr-var* **and** *ivar* *in-var* **and** *ovar* *out-var*

The functions above turn a representation of a variable (type  $'v$ ), including its name and type, into some lens type  $'e$ . *svar* constructs a predicate variable, *ivar* and input variables, and *ovar*

and output variable. The functions bridge between the model and encoding of the variable and its interpretation as a lens in order to integrate it into the general lens-based framework. Overriding these functions is then all we need to make use of any kind of variables in terms of interfacing it with the system. Although in core UTP variables are always modelled using record field, we can overload these constants to allow other kinds of variables, such as deep variables with explicit syntax and type information.

Finally, we set up the translations rules.

#### translations

— Identifiers

-svid  $x \rightarrow x$

-svid-alpha  $\Rightarrow \Sigma$

-svid-dot  $x y \rightarrow y ;_L x$

— Decorations

-spvar  $\Sigma \leftarrow \text{CONST svar } \text{CONST id-lens}$

-sinvar  $\Sigma \leftarrow \text{CONST ivar } 1_L$

-soutvar  $\Sigma \leftarrow \text{CONST ovar } 1_L$

-spvar  $(\text{-svid-dot } x y) \leftarrow \text{CONST svar } (\text{CONST lens-comp } y x)$

-sinvar  $(\text{-svid-dot } x y) \leftarrow \text{CONST ivar } (\text{CONST lens-comp } y x)$

-soutvar  $(\text{-svid-dot } x y) \leftarrow \text{CONST ovar } (\text{CONST lens-comp } y x)$

-spvar  $x \Rightarrow \text{CONST svar } x$

-sinvar  $x \Rightarrow \text{CONST ivar } x$

-soutvar  $x \Rightarrow \text{CONST ovar } x$

— Alphabets

-salphaparen  $a \rightarrow a$

-salphaid  $x \rightarrow x$

-salphacomp  $x y \rightarrow x +_L y$

-salphaprod  $a b \Rightarrow a \times_L b$

-salphavar  $x \rightarrow x$

-svar-nil  $x \rightarrow x$

-svar-cons  $x xs \rightarrow x +_L xs$

-salphaset  $A \rightarrow A$

$(\text{-svar-cons } x (\text{-salphamk } y)) \leftarrow \text{-salphamk } (x +_L y)$

$x \leftarrow \text{-salphamk } x$

-salpha-all  $\Rightarrow 1_L$

-salpha-none  $\Rightarrow 0_L$

— Quotations

-ualpha-set  $A \rightarrow A$

-svar  $x \rightarrow x$

The translation rules mainly convert syntax into lens constructions, using a mixture of lens operators and the bespoke variable definitions. Notably, a colon variable identifier qualification becomes a lens composition, and variable sets are constructed using len sum. The translation rules are carefully crafted to ensure both parsing and pretty printing.

Finally we create the following useful utility translation function that allows us to construct a UTP variable (lens) type given a return and alphabet type.

#### syntax

-uvar-ty  $:: \text{type} \Rightarrow \text{type} \Rightarrow \text{type}$

#### parse-translation (

let

```

  fun uvar-ty-tr [ty] = Syntax.const @{type-syntax lens} $ ty $ Syntax.const @{type-syntax dummy}
    | uvar-ty-tr ts = raise TERM (uvar-ty-tr, ts);
in [(@{syntax-const -uvar-ty}, K uvar-ty-tr)] end
>

```

end

### 3 UTP expressions

```

theory utp-expr
imports
  utp-var
begin

```

#### 3.1 Expression type

**purge-notation** *BNF-Def.convolve* ((-, / -))

Before building the predicate model, we will build a model of expressions that generalise alphabetised predicates. Expressions are represented semantically as mapping from the alphabet  $'\alpha$  to the expression's type  $'a$ . This general model will allow us to unify all constructions under one type. The majority definitions in the file are given using the *lifting* package [6], which allows us to reuse much of the existing library of HOL functions.

**typedef**  $('t, 'a) \text{ uexpr} = \text{UNIV} :: ('a \Rightarrow 't) \text{ set} ..$

**setup-lifting** *type-definition-uexpr*

**notation** *Rep-uexpr* ( $\llbracket \cdot \rrbracket_e$ )

**lemma** *uexpr-eq-iff*:

$e = f \longleftrightarrow (\forall b. \llbracket e \rrbracket_e b = \llbracket f \rrbracket_e b)$   
**using** *Rep-uexpr-inject*[*of e f, THEN sym*] **by** (*auto*)

The term  $\llbracket e \rrbracket_e b$  effectively refers to the semantic interpretation of the expression under the state-space valuation (or variables binding)  $b$ . It can be used, in concert with the *lifting* package, to interpret UTP constructs to their HOL equivalents. We create some theorem sets to store such transfer theorems.

**named-theorems** *ueval* and *lit-simps*

#### 3.2 Core expression constructs

A variable expression corresponds to the lens *get* function associated with a variable. Specifically, given a lens the expression always returns that portion of the state-space referred to by the lens.

**lift-definition**  $\text{var} :: ('t \Rightarrow 'a) \Rightarrow ('t, 'a) \text{ uexpr}$  **is** *lens-get* .

A literal is simply a constant function expression, always returning the same value for any binding.

**lift-definition**  $\text{lit} :: 't \Rightarrow ('t, 'a) \text{ uexpr}$  **is**  $\lambda v b. v$  .

We define lifting for unary, binary, ternary, and quaternary expression constructs, that simply take a HOL function with correct number of arguments and apply it function to all possible results of the expressions.

**lift-definition**  $uop :: ('a \Rightarrow 'b) \Rightarrow ('a, 'α) uexpr \Rightarrow ('b, 'α) uexpr$   
**is**  $\lambda f e b. f (e b) .$

**lift-definition**  $bop ::$   
 $('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow ('a, 'α) uexpr \Rightarrow ('b, 'α) uexpr \Rightarrow ('c, 'α) uexpr$   
**is**  $\lambda f u v b. f (u b) (v b) .$

**lift-definition**  $trop ::$   
 $('a \Rightarrow 'b \Rightarrow 'c \Rightarrow 'd) \Rightarrow ('a, 'α) uexpr \Rightarrow ('b, 'α) uexpr \Rightarrow ('c, 'α) uexpr \Rightarrow ('d, 'α) uexpr$   
**is**  $\lambda f u v w b. f (u b) (v b) (w b) .$

**lift-definition**  $qtop ::$   
 $('a \Rightarrow 'b \Rightarrow 'c \Rightarrow 'd \Rightarrow 'e) \Rightarrow$   
 $('a, 'α) uexpr \Rightarrow ('b, 'α) uexpr \Rightarrow ('c, 'α) uexpr \Rightarrow ('d, 'α) uexpr \Rightarrow$   
 $('e, 'α) uexpr$   
**is**  $\lambda f u v w x b. f (u b) (v b) (w b) (x b) .$

We also define a UTP expression version of function ( $\lambda$ ) abstraction, that takes a function producing an expression and produces an expression producing a function.

**lift-definition**  $ulambda :: ('a \Rightarrow ('b, 'α) uexpr) \Rightarrow ('a \Rightarrow 'b, 'α) uexpr$   
**is**  $\lambda f A x. f x A .$

UTP expression is equality is simply HOL equality lifted using the *bop* binary expression constructor.

**definition**  $eq-upred :: ('a, 'α) uexpr \Rightarrow ('a, 'α) uexpr \Rightarrow (bool, 'α) uexpr$   
**where**  $eq-upred x y = bop HOL.eq x y$

We define syntax for expressions using adhoc-overloading – this allows us to later define operators on different types if necessary (e.g. when adding types for new UTP theories).

**consts**

$ulit \quad :: 't \Rightarrow 'e \ (\ll-\gg)$   
 $ueq \quad :: 'a \Rightarrow 'a \Rightarrow 'b \ (\text{infixl} =_u 50)$

**adhoc-overloading**

$ulit \ lit \ \mathbf{and}$   
 $ueq \ eq-upred$

A literal is the expression  $\ll v \gg$ , where  $v$  is any HOL term. Actually, the literal construct is very versatile and also allows us to refer to HOL variables within UTP expressions, and has a variety of other uses. It can therefore also be considered as a kind of quotation mechanism.

We also set up syntax for UTP variable expressions.

**syntax**

$-uuvar :: svar \Rightarrow logic \ (-)$

**translations**

$-uuvar x == CONST var x$

Since we already have a parser for variables, we can directly reuse it and simply apply the *var* expression construct to lift the resulting variable to an expression.

### 3.3 Type class instantiations

Isabelle/HOL of course provides a large hierarchy of type classes that provide constructs such as numerals and the arithmetic operators. Fortunately we can directly make use of these for UTP expressions, and thus we now perform a long list of appropriate instantiations. We first

lift the core arithmetic constants and operators using a mixture of literals, unary, and binary expression constructors.

```
instantiation uexpr :: (zero, type) zero
begin
  definition zero-uexpr-def:  $0 = \text{lit } 0$ 
instance ..
end
```

```
instantiation uexpr :: (one, type) one
begin
  definition one-uexpr-def:  $1 = \text{lit } 1$ 
instance ..
```

```
end
```

```
instantiation uexpr :: (plus, type) plus
begin
  definition plus-uexpr-def:  $u + v = \text{bop } (op +) u v$ 
instance ..
end
```

It should be noted that instantiating the unary minus class, *uminus*, will also provide negation UTP predicates later.

```
instantiation uexpr :: (uminus, type) uminus
begin
  definition uminus-uexpr-def:  $- u = \text{uop } \text{uminus } u$ 
instance ..
end
```

```
instantiation uexpr :: (minus, type) minus
begin
  definition minus-uexpr-def:  $u - v = \text{bop } (op -) u v$ 
instance ..
end
```

```
instantiation uexpr :: (times, type) times
begin
  definition times-uexpr-def:  $u * v = \text{bop } (op *) u v$ 
instance ..
end
```

```
instance uexpr :: (Rings.dvd, type) Rings.dvd ..
```

```
instantiation uexpr :: (divide, type) divide
begin
  definition divide-uexpr :: ('a, 'b) uexpr  $\Rightarrow$  ('a, 'b) uexpr  $\Rightarrow$  ('a, 'b) uexpr where
    divide-uexpr u v = bop divide u v
instance ..
end
```

```
instantiation uexpr :: (inverse, type) inverse
begin
  definition inverse-uexpr :: ('a, 'b) uexpr  $\Rightarrow$  ('a, 'b) uexpr
    where inverse-uexpr u = uop inverse u
instance ..
```

**end**

**instantiation** *uexpr* :: (*modulo*, *type*) *modulo*

**begin**

**definition** *mod-uexpr-def*:  $u \text{ mod } v = \text{bop } (\text{op mod}) \ u \ v$

**instance** ..

**end**

**instantiation** *uexpr* :: (*sgn*, *type*) *sgn*

**begin**

**definition** *sgn-uexpr-def*:  $\text{sgn } u = \text{uop sgn } u$

**instance** ..

**end**

**instantiation** *uexpr* :: (*abs*, *type*) *abs*

**begin**

**definition** *abs-uexpr-def*:  $\text{abs } u = \text{uop abs } u$

**instance** ..

**end**

Once we've set up all the core constructs for arithmetic, we can also instantiate the type classes for various algebras, including groups and rings. The proofs are done by definitional expansion, the *transfer* tactic, and then finally the theorems of the underlying HOL operators. This is mainly routine, so we don't comment further.

**instance** *uexpr* :: (*semigroup-mult*, *type*) *semigroup-mult*

**by** (*intro-classes*) (*simp add: times-uexpr-def one-uexpr-def, transfer, simp add: mult.assoc*)+

**instance** *uexpr* :: (*monoid-mult*, *type*) *monoid-mult*

**by** (*intro-classes*) (*simp add: times-uexpr-def one-uexpr-def, transfer, simp*)+

**instance** *uexpr* :: (*semigroup-add*, *type*) *semigroup-add*

**by** (*intro-classes*) (*simp add: plus-uexpr-def zero-uexpr-def, transfer, simp add: add.assoc*)+

**instance** *uexpr* :: (*monoid-add*, *type*) *monoid-add*

**by** (*intro-classes*) (*simp add: plus-uexpr-def zero-uexpr-def, transfer, simp*)+

**instance** *uexpr* :: (*ab-semigroup-add*, *type*) *ab-semigroup-add*

**by** (*intro-classes*) (*simp add: plus-uexpr-def, transfer, simp add: add.commute*)+

**instance** *uexpr* :: (*cancel-semigroup-add*, *type*) *cancel-semigroup-add*

**by** (*intro-classes*) (*simp add: plus-uexpr-def, transfer, simp add: fun-eq-iff*)+

**instance** *uexpr* :: (*cancel-ab-semigroup-add*, *type*) *cancel-ab-semigroup-add*

**by** (*intro-classes*, (*simp add: plus-uexpr-def minus-uexpr-def, transfer, simp add: fun-eq-iff add.commute cancel-ab-semigroup-add-class.diff-diff-add*)

**instance** *uexpr* :: (*group-add*, *type*) *group-add*

**by** (*intro-classes*)

    (*simp add: plus-uexpr-def uminus-uexpr-def minus-uexpr-def zero-uexpr-def, transfer, simp*)+

**instance** *uexpr* :: (*ab-group-add*, *type*) *ab-group-add*

**by** (*intro-classes*)

    (*simp add: plus-uexpr-def uminus-uexpr-def minus-uexpr-def zero-uexpr-def, transfer, simp*)+

**instance** *uexpr* :: (*semiring*, *type*) *semiring*

**by** (*intro-classes*) (*simp add: plus-uepr-def times-uepr-def, transfer, simp add: fun-eq-iff add.commute semiring-class.distrib-right semiring-class.distrib-left*)+

**instance** *uepr* :: (*ring-1*, *type*) *ring-1*

**by** (*intro-classes*) (*simp add: plus-uepr-def uminus-uepr-def minus-uepr-def times-uepr-def zero-uepr-def one-uepr-def, transfer, simp add: fun-eq-iff*)+

We can also define the order relation on expressions. Now, unlike the previous group and ring constructs, the order relations  $op \leq$  and  $op \leq$  return a *bool* type. This order is not therefore the lifted order which allows us to compare the valuation of two expressions, but rather the order on expressions themselves. Notably, this instantiation will later allow us to talk about predicate refinements and complete lattices.

**instantiation** *uepr* :: (*ord*, *type*) *ord*

**begin**

**lift-definition** *less-eq-uepr* :: (*'a*, *'b*) *uepr*  $\Rightarrow$  (*'a*, *'b*) *uepr*  $\Rightarrow$  *bool*

**is**  $\lambda P Q. (\forall A. P A \leq Q A) .$

**definition** *less-uepr* :: (*'a*, *'b*) *uepr*  $\Rightarrow$  (*'a*, *'b*) *uepr*  $\Rightarrow$  *bool*

**where** *less-uepr* *P Q* = (*P*  $\leq$  *Q*  $\wedge$   $\neg$  *Q*  $\leq$  *P*)

**instance** ..

**end**

UTP expressions whose return type is a partial ordered type, are also partially ordered as the following instantiation demonstrates.

**instance** *uepr* :: (*order*, *type*) *order*

**proof**

**fix** *x y z* :: (*'a*, *'b*) *uepr*

**show** (*x* < *y*) = (*x*  $\leq$  *y*  $\wedge$   $\neg$  *y*  $\leq$  *x*) **by** (*simp add: less-uepr-def*)

**show** *x*  $\leq$  *x* **by** (*transfer, auto*)

**show** *x*  $\leq$  *y*  $\implies$  *y*  $\leq$  *z*  $\implies$  *x*  $\leq$  *z*

**by** (*transfer, blast intro:order.trans*)

**show** *x*  $\leq$  *y*  $\implies$  *y*  $\leq$  *x*  $\implies$  *x* = *y*

**by** (*transfer, rule ext, simp add: eq-iff*)

**qed**

We also lift the properties from certain ordered groups.

**instance** *uepr* :: (*ordered-ab-group-add*, *type*) *ordered-ab-group-add*

**by** (*intro-classes*) (*simp add: plus-uepr-def, transfer, simp*)

**instance** *uepr* :: (*ordered-ab-group-add-abs*, *type*) *ordered-ab-group-add-abs*

**apply** (*intro-classes*)

**apply** (*simp add: abs-uepr-def zero-uepr-def plus-uepr-def uminus-uepr-def, transfer, simp add:*

*abs-ge-self abs-le-iff abs-triangle-ineq*)+

**apply** (*metis ab-group-add-class.ab-diff-conv-add-uminus abs-ge-minus-self abs-ge-self add-mono-thms-linordered-semiring*)  
**done**

The following instantiation sets up numerals. This will allow us to have Isabelle number representations (i.e. 3,7,42,198 etc.) to UTP expressions directly.

**instance** *uepr* :: (*numeral*, *type*) *numeral*

**by** (*intro-classes, simp add: plus-uepr-def, transfer, simp add: add.assoc*)

The following two theorems also set up interpretation of numerals, meaning a UTP numeral can always be converted to a HOL numeral.

**lemma** *numeral-uepr-rep-eq*:  $\llbracket \text{numeral } x \rrbracket_e b = \text{numeral } x$

```

apply (induct x)
apply (simp add: lit.rep-eq one-uepr-def)
apply (simp add: bop.rep-eq numeral-Bit0 plus-uepr-def)
apply (simp add: bop.rep-eq lit.rep-eq numeral-code(3) one-uepr-def plus-uepr-def)
done

```

```

lemma numeral-uepr-simp: numeral x = <<numeral x>>
by (simp add: uepr-eq-iff numeral-uepr-rep-eq lit.rep-eq)

```

We can also lift a few arithmetic properties from the class instantiations above using *transfer*.

```

lemma uepr-diff-zero [simp]:
  fixes a :: ('α::trace, 'a) uepr
  shows a - 0 = a
by (simp add: minus-uepr-def zero-uepr-def, transfer, auto)

```

```

lemma uepr-add-diff-cancel-left [simp]:
  fixes a b :: ('α::trace, 'a) uepr
  shows (a + b) - a = b
by (simp add: minus-uepr-def plus-uepr-def, transfer, auto)

```

### 3.4 Overloaded expression constructors

For convenience, we often want to utilise the same expression syntax for multiple constructs. This can be achieved using ad-hoc overloading. We create a number of polymorphic constants and then overload their definitions using appropriate implementations. In order for this to work, each collection must have its own unique type. Thus we do not use the HOL map type directly, but rather our own partial function type, for example.

```

consts
  — Empty elements, for example empty set, nil list, 0...
  uempty    :: 'f
  — Function application, map application, list application...
  uapply    :: 'f ⇒ 'k ⇒ 'v
  — Function update, map update, list update...
  upd       :: 'f ⇒ 'k ⇒ 'v ⇒ 'f
  — Domain of maps, lists...
  udom      :: 'f ⇒ 'a set
  — Range of maps, lists...
  uran      :: 'f ⇒ 'b set
  — Domain restriction
  udomres   :: 'a set ⇒ 'f ⇒ 'f
  — Range restriction
  uranres   :: 'f ⇒ 'b set ⇒ 'f
  — Collection cardinality
  ucard     :: 'f ⇒ nat
  — Collection summation
  usums     :: 'f ⇒ 'a
  — Construct a collection from a list of entries
  uentries  :: 'k set ⇒ ('k ⇒ 'v) ⇒ 'f

```

We need a function corresponding to function application in order to overload.

```

definition fun-apply :: ('a ⇒ 'b) ⇒ ('a ⇒ 'b)
where fun-apply f x = f x

```

```

declare fun-apply-def [simp]

```



**definition** *ffun-entries* :: 'k set  $\Rightarrow$  ('k  $\Rightarrow$  'v)  $\Rightarrow$  ('k, 'v) *ffun* **where**  
*ffun-entries* d f = *graph-ffun* {(k, f k) | k. k  $\in$  d}

We then set up the overloading for a number of useful constructs for various collections.

#### ad hoc-overloading

*uempty* 0 **and**  
*uapply* *fun-apply* **and** *uapply* *nth* **and** *uapply* *pfun-app* **and**  
*uapply* *ffun-app* **and**  
*upd* *pfun-upd* **and** *upd* *ffun-upd* **and** *upd* *list-augment* **and**  
*uom* *Domain* **and** *uom* *pdom* **and** *uom* *fdom* **and** *uom* *seq-dom* **and**  
*uom* *Range* **and** *uran* *pran* **and** *uran* *fran* **and** *uran* *set* **and**  
*uomres* *pdom-res* **and** *uomres* *fdom-res* **and**  
*uranres* *pran-res* **and** *uomres* *fran-res* **and**  
*ucard* *card* **and** *ucard* *pcard* **and** *ucard* *length* **and**  
*usums* *list-sum* **and** *usums* *Sum* **and** *usums* *pfun-sum* **and**  
*uentries* *pfun-entries* **and** *uentries* *ffun-entries*

### 3.5 Syntax translations

The follows a large number of translations that lift HOL functions to UTP expressions using the various expression constructors defined above. Much of the time we try to keep the HOL syntax but add a "u" subscript.

**abbreviation** (*input*) *ulens-override* x f g  $\equiv$  *lens-override* f g x

#### translations

0  $\leq$  CONST *uempty* — We have to do this so we don't see *uempty*. Is there a better way of printing?

We add new non-terminals for UTP tuples and maplets.

**nonterminal** *utuple-args* **and** *umaplet* **and** *umaplets*

**syntax** — Core expression constructs

-*ucoerce* :: *logic*  $\Rightarrow$  *type*  $\Rightarrow$  *logic* (**infix** :<sub>u</sub> 50)  
-*ulambda* :: *pttrn*  $\Rightarrow$  *logic*  $\Rightarrow$  *logic* ( $\lambda$  - - - [0, 10] 10)  
-*ulens-ovrd* :: *logic*  $\Rightarrow$  *logic*  $\Rightarrow$  *salpha*  $\Rightarrow$  *logic* (-  $\oplus$  - on - [85, 0, 86] 86)

#### translations

$\lambda x \cdot p ==$  CONST *ulambda* ( $\lambda x. p$ )  
 $x :_u 'a ==$  x :: ('a, -) *uexpr*  
-*ulens-ovrd* f g a  $\Rightarrow$  CONST *bop* (CONST *ulens-override* a) f g  
-*ulens-ovrd* f g a  $\leq$  CONST *bop* ( $\lambda x y. CONST$  *lens-override* x1 y1 a) f g

**syntax** — Tuples

-*utuple* :: ('a, 'α) *uexpr*  $\Rightarrow$  *utuple-args*  $\Rightarrow$  ('a \* 'b, 'α) *uexpr* ((1'(-, -')<sub>u</sub>))  
-*utuple-arg* :: ('a, 'α) *uexpr*  $\Rightarrow$  *utuple-args* (-)  
-*utuple-args* :: ('a, 'α) *uexpr*  $\Rightarrow$  *utuple-args*  $\Rightarrow$  *utuple-args* (-, / -)  
-*uunit* :: ('a, 'α) *uexpr* ('()<sub>u</sub>)  
-*ufst* :: ('a  $\times$  'b, 'α) *uexpr*  $\Rightarrow$  ('a, 'α) *uexpr* ( $\pi_1$ '(-))  
-*usnd* :: ('a  $\times$  'b, 'α) *uexpr*  $\Rightarrow$  ('b, 'α) *uexpr* ( $\pi_2$ '(-))

#### translations

()<sub>u</sub> ==  $\ll() \gg$   
(x, y)<sub>u</sub> == CONST *bop* (CONST *Pair*) x y  
-*utuple* x (-*utuple-args* y z) == -*utuple* x (-*utuple-arg* (-*utuple* y z))

$\pi_1(x) == \text{CONST uop CONST fst } x$   
 $\pi_2(x) == \text{CONST uop CONST snd } x$

**syntax** — Polymorphic constructs

$\text{-uundef} :: \text{logic } (\perp_u)$   
 $\text{-umap-empty} :: \text{logic } ([\ ]_u)$   
 $\text{-uapply} :: ('a \Rightarrow 'b, 'a) \text{ uexpr} \Rightarrow \text{utuple-args} \Rightarrow ('b, 'a) \text{ uexpr } (-'(-')_a [999,0] 999)$   
 $\text{-umaplet} :: [\text{logic}, \text{logic}] \Rightarrow \text{umaplet } (- \text{ /}\mapsto\text{ / } -)$   
 $\quad \quad \quad :: \text{umaplet} \Rightarrow \text{umaplets} \quad (-)$   
 $\text{-UMaplets} :: [\text{umaplet}, \text{umaplets}] \Rightarrow \text{umaplets } (-, / -)$   
 $\text{-UMapUpd} :: [\text{logic}, \text{umaplets}] \Rightarrow \text{logic } (-'(-')_u [900,0] 900)$   
 $\text{-UMap} :: \text{umaplets} \Rightarrow \text{logic } ((1[\ ]_u))$   
 $\text{-ucard} :: \text{logic} \Rightarrow \text{logic } (\#_u '(-'))$   
 $\text{-ules} :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } (\mathbf{infix} <_u 50)$   
 $\text{-uleq} :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } (\mathbf{infix} \leq_u 50)$   
 $\text{-ugreat} :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } (\mathbf{infix} >_u 50)$   
 $\text{-ugeq} :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } (\mathbf{infix} \geq_u 50)$   
 $\text{-uceil} :: \text{logic} \Rightarrow \text{logic } (\lceil - \rceil_u)$   
 $\text{-ufloor} :: \text{logic} \Rightarrow \text{logic } (\lfloor - \rfloor_u)$   
 $\text{-udom} :: \text{logic} \Rightarrow \text{logic } (\text{dom}_u '(-'))$   
 $\text{-uran} :: \text{logic} \Rightarrow \text{logic } (\text{ran}_u '(-'))$   
 $\text{-usum} :: \text{logic} \Rightarrow \text{logic } (\text{sum}_u '(-'))$   
 $\text{-udom-res} :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } (\mathbf{infixl} \triangleleft_u 85)$   
 $\text{-uran-res} :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } (\mathbf{infixl} \triangleright_u 85)$   
 $\text{-umin} :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } (\text{min}_u '(-, -'))$   
 $\text{-umax} :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } (\text{max}_u '(-, -'))$   
 $\text{-ugcd} :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } (\text{gcd}_u '(-, -'))$   
 $\text{-uentries} :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } (\text{entr}_u '(-, -'))$

**translations**

— Pretty printing for adhoc-overloaded constructs

$f(x)_a <= \text{CONST uapply } f \ x$   
 $\text{dom}_u(f) <= \text{CONST udom } f$   
 $\text{ran}_u(f) <= \text{CONST uran } f$   
 $A \triangleleft_u f <= \text{CONST udomres } A \ f$   
 $f \triangleright_u A <= \text{CONST uranres } f \ A$   
 $\#_u(f) <= \text{CONST ucard } f$   
 $f(k \mapsto v)_u <= \text{CONST uupd } f \ k \ v$

— Overloaded construct translations

$f(x,y,z,u)_a == \text{CONST bop CONST uapply } f \ (x,y,z,u)_u$   
 $f(x,y,z)_a == \text{CONST bop CONST uapply } f \ (x,y,z)_u$   
 $f(x,y)_a == \text{CONST bop CONST uapply } f \ (x,y)_u$   
 $f(x)_a == \text{CONST bop CONST uapply } f \ x$   
 $\#_u(xs) == \text{CONST uop CONST ucard } xs$   
 $\text{sum}_u(A) == \text{CONST uop CONST usums } A$   
 $\text{dom}_u(f) == \text{CONST uop CONST udom } f$   
 $\text{ran}_u(f) == \text{CONST uop CONST uran } f$   
 $[\ ]_u == \ll \text{CONST uempty} \gg$   
 $\perp_u == \ll \text{CONST undefined} \gg$   
 $A \triangleleft_u f == \text{CONST bop } (\text{CONST udomres}) \ A \ f$   
 $f \triangleright_u A == \text{CONST bop } (\text{CONST uranres}) \ f \ A$   
 $\text{entr}_u(d,f) == \text{CONST bop CONST uentries } d \ \ll f \gg$   
 $\text{-UMapUpd } m \ (\text{-UMaplets } xy \ ms) == \text{-UMapUpd } (\text{-UMapUpd } m \ xy) \ ms$   
 $\text{-UMapUpd } m \ (\text{-umaplet } x \ y) == \text{CONST trop CONST uupd } m \ x \ y$

$-UMap\ ms \quad == \ -UMapUpd\ []_u\ ms$   
 $-UMap\ (-UMaplets\ ms1\ ms2) \quad <= \ -UMapUpd\ (-UMap\ ms1)\ ms2$   
 $-UMaplets\ ms1\ (-UMaplets\ ms2\ ms3) <= \ -UMaplets\ (-UMaplets\ ms1\ ms2)\ ms3$

— Type-class polymorphic constructs

$x <_u y \quad == \ CONST\ bop\ (op\ <)\ x\ y$   
 $x \leq_u y \quad == \ CONST\ bop\ (op\ \leq)\ x\ y$   
 $x >_u y \quad ==> \ y <_u x$   
 $x \geq_u y \quad ==> \ y \leq_u x$   
 $min_u(x, y) \quad == \ CONST\ bop\ (CONST\ min)\ x\ y$   
 $max_u(x, y) \quad == \ CONST\ bop\ (CONST\ max)\ x\ y$   
 $gcd_u(x, y) \quad == \ CONST\ bop\ (CONST\ gcd)\ x\ y$   
 $[x]_u \quad == \ CONST\ uop\ CONST\ ceiling\ x$   
 $[x]_u \quad == \ CONST\ uop\ CONST\ floor\ x$

**syntax** — Lists / Sequences

$-unil \quad :: \ ('a\ list, '\alpha)\ uexpr\ (\langle \rangle)$   
 $-ulist \quad :: \ args \Rightarrow ('a\ list, '\alpha)\ uexpr\ (\langle (-) \rangle)$   
 $-uappend \quad :: \ ('a\ list, '\alpha)\ uexpr \Rightarrow ('a\ list, '\alpha)\ uexpr \Rightarrow ('a\ list, '\alpha)\ uexpr\ (\mathbf{infixr}\ \hat{\_}_u\ 80)$   
 $-ulast \quad :: \ ('a\ list, '\alpha)\ uexpr \Rightarrow ('a, '\alpha)\ uexpr\ (last_u'(-))$   
 $-ufront \quad :: \ ('a\ list, '\alpha)\ uexpr \Rightarrow ('a\ list, '\alpha)\ uexpr\ (front_u'(-))$   
 $-uhead \quad :: \ ('a\ list, '\alpha)\ uexpr \Rightarrow ('a, '\alpha)\ uexpr\ (head_u'(-))$   
 $-utail \quad :: \ ('a\ list, '\alpha)\ uexpr \Rightarrow ('a\ list, '\alpha)\ uexpr\ (tail_u'(-))$   
 $-utake \quad :: \ (nat, '\alpha)\ uexpr \Rightarrow ('a\ list, '\alpha)\ uexpr \Rightarrow ('a\ list, '\alpha)\ uexpr\ (take_u'(-, / -))$   
 $-udrop \quad :: \ (nat, '\alpha)\ uexpr \Rightarrow ('a\ list, '\alpha)\ uexpr \Rightarrow ('a\ list, '\alpha)\ uexpr\ (drop_u'(-, / -))$   
 $-ufilter \quad :: \ ('a\ list, '\alpha)\ uexpr \Rightarrow ('a\ set, '\alpha)\ uexpr \Rightarrow ('a\ list, '\alpha)\ uexpr\ (\mathbf{infixl}\ \downarrow_u\ 75)$   
 $-uextract \quad :: \ ('a\ set, '\alpha)\ uexpr \Rightarrow ('a\ list, '\alpha)\ uexpr \Rightarrow ('a\ list, '\alpha)\ uexpr\ (\mathbf{infixl}\ \uparrow_u\ 75)$   
 $-uelems \quad :: \ ('a\ list, '\alpha)\ uexpr \Rightarrow ('a\ set, '\alpha)\ uexpr\ (elems_u'(-))$   
 $-usorted \quad :: \ ('a\ list, '\alpha)\ uexpr \Rightarrow (bool, '\alpha)\ uexpr\ (sorted_u'(-))$   
 $-udistinct \quad :: \ ('a\ list, '\alpha)\ uexpr \Rightarrow (bool, '\alpha)\ uexpr\ (distinct_u'(-))$   
 $-uupto \quad :: \ logic \Rightarrow logic \Rightarrow logic\ (\langle \dots \rangle)$   
 $-uupt \quad :: \ logic \Rightarrow logic \Rightarrow logic\ (\langle \dots < \rangle)$   
 $-umap \quad :: \ logic \Rightarrow logic \Rightarrow logic\ (map_u)$

**translations**

$\langle \rangle \quad == \ \llbracket \rrbracket$   
 $\langle x, xs \rangle \quad == \ CONST\ bop\ (op\ \#)\ x\ \langle xs \rangle$   
 $\langle x \rangle \quad == \ CONST\ bop\ (op\ \#)\ x\ \llbracket \rrbracket$   
 $x \hat{\_}_u y \quad == \ CONST\ bop\ (op\ @)\ x\ y$   
 $last_u(xs) \quad == \ CONST\ uop\ CONST\ last\ xs$   
 $front_u(xs) \quad == \ CONST\ uop\ CONST\ butlast\ xs$   
 $head_u(xs) \quad == \ CONST\ uop\ CONST\ hd\ xs$   
 $tail_u(xs) \quad == \ CONST\ uop\ CONST\ tl\ xs$   
 $drop_u(n, xs) \quad == \ CONST\ bop\ CONST\ drop\ n\ xs$   
 $take_u(n, xs) \quad == \ CONST\ bop\ CONST\ take\ n\ xs$   
 $elems_u(xs) \quad == \ CONST\ uop\ CONST\ set\ xs$   
 $sorted_u(xs) \quad == \ CONST\ uop\ CONST\ sorted\ xs$   
 $distinct_u(xs) \quad == \ CONST\ uop\ CONST\ distinct\ xs$   
 $xs \downarrow_u A \quad == \ CONST\ bop\ CONST\ seq-filter\ xs\ A$   
 $A \uparrow_u xs \quad == \ CONST\ bop\ (op\ \uparrow_l)\ A\ xs$   
 $\langle n..k \rangle \quad == \ CONST\ bop\ CONST\ upto\ n\ k$   
 $\langle n..<k \rangle \quad == \ CONST\ bop\ CONST\ upt\ n\ k$   
 $map_u\ f\ xs \quad == \ CONST\ bop\ CONST\ map\ f\ xs$

**syntax** — Sets

$-ufinite \quad :: \text{logic} \Rightarrow \text{logic} \ (\text{finite}_u '(-))$   
 $-uempset \quad :: ('a \text{ set}, ' \alpha) \text{ uexpr} \ (\{\}_u)$   
 $-uset \quad :: \text{args} \Rightarrow ('a \text{ set}, ' \alpha) \text{ uexpr} \ (\{(-)\}_u)$   
 $-uunion \quad :: ('a \text{ set}, ' \alpha) \text{ uexpr} \Rightarrow ('a \text{ set}, ' \alpha) \text{ uexpr} \Rightarrow ('a \text{ set}, ' \alpha) \text{ uexpr} \ (\text{infixl} \cup_u 65)$   
 $-uinter \quad :: ('a \text{ set}, ' \alpha) \text{ uexpr} \Rightarrow ('a \text{ set}, ' \alpha) \text{ uexpr} \Rightarrow ('a \text{ set}, ' \alpha) \text{ uexpr} \ (\text{infixl} \cap_u 70)$   
 $-umem \quad :: ('a, ' \alpha) \text{ uexpr} \Rightarrow ('a \text{ set}, ' \alpha) \text{ uexpr} \Rightarrow (\text{bool}, ' \alpha) \text{ uexpr} \ (\text{infix} \in_u 50)$   
 $-usubset \quad :: ('a \text{ set}, ' \alpha) \text{ uexpr} \Rightarrow ('a \text{ set}, ' \alpha) \text{ uexpr} \Rightarrow (\text{bool}, ' \alpha) \text{ uexpr} \ (\text{infix} \subset_u 50)$   
 $-usubseteq \quad :: ('a \text{ set}, ' \alpha) \text{ uexpr} \Rightarrow ('a \text{ set}, ' \alpha) \text{ uexpr} \Rightarrow (\text{bool}, ' \alpha) \text{ uexpr} \ (\text{infix} \subseteq_u 50)$

#### translations

$\text{finite}_u(x) == \text{CONST } uop \ (\text{CONST } finite) \ x$   
 $\{\}_u == \ll\{\}\gg$   
 $\{x, xs\}_u == \text{CONST } bop \ (\text{CONST } insert) \ x \ \{xs\}_u$   
 $\{x\}_u == \text{CONST } bop \ (\text{CONST } insert) \ x \ \ll\{\}\gg$   
 $A \cup_u B == \text{CONST } bop \ (op \cup) \ A \ B$   
 $A \cap_u B == \text{CONST } bop \ (op \cap) \ A \ B$   
 $x \in_u A == \text{CONST } bop \ (op \in) \ x \ A$   
 $A \subset_u B == \text{CONST } bop \ (op \subset) \ A \ B$   
 $f \subset_u g <= \text{CONST } bop \ (op \subset_p) \ f \ g$   
 $f \subset_u g <= \text{CONST } bop \ (op \subset_f) \ f \ g$   
 $A \subseteq_u B == \text{CONST } bop \ (op \subseteq) \ A \ B$   
 $f \subseteq_u g <= \text{CONST } bop \ (op \subseteq_p) \ f \ g$   
 $f \subseteq_u g <= \text{CONST } bop \ (op \subseteq_f) \ f \ g$

#### syntax — Partial functions

$-umap-plus \quad :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \ (\text{infixl} \oplus_u 85)$   
 $-umap-minus \quad :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \ (\text{infixl} \ominus_u 85)$

#### translations

$f \oplus_u g \Rightarrow (f :: ((-, -) \text{ pfun}, -) \text{ uexpr}) + g$   
 $f \ominus_u g \Rightarrow (f :: ((-, -) \text{ pfun}, -) \text{ uexpr}) - g$

#### syntax — Sum types

$-uinl \quad :: \text{logic} \Rightarrow \text{logic} \ (\text{inl}_u '(-))$   
 $-uinr \quad :: \text{logic} \Rightarrow \text{logic} \ (\text{inr}_u '(-))$

#### translations

$\text{inl}_u(x) == \text{CONST } uop \ \text{CONST } Inl \ x$   
 $\text{inr}_u(x) == \text{CONST } uop \ \text{CONST } Inr \ x$

### 3.6 Lifting set collectors

We provide syntax for various types of set collectors, including intervals and the Z-style set comprehension which is purpose built as a new lifted definition.

#### syntax

$-uset-atLeastAtMost \quad :: ('a, ' \alpha) \text{ uexpr} \Rightarrow ('a, ' \alpha) \text{ uexpr} \Rightarrow ('a \text{ set}, ' \alpha) \text{ uexpr} \ ((1\{-..\}_u))$   
 $-uset-atLeastLessThan \quad :: ('a, ' \alpha) \text{ uexpr} \Rightarrow ('a, ' \alpha) \text{ uexpr} \Rightarrow ('a \text{ set}, ' \alpha) \text{ uexpr} \ ((1\{--<}_u))$   
 $-uset-compr \quad :: \text{pttrn} \Rightarrow ('a \text{ set}, ' \alpha) \text{ uexpr} \Rightarrow (\text{bool}, ' \alpha) \text{ uexpr} \Rightarrow ('b, ' \alpha) \text{ uexpr} \Rightarrow ('b \text{ set}, ' \alpha) \text{ uexpr} \ ((1\{- :/ - \mid - \cdot / -\}_u))$   
 $-uset-compr-nset \quad :: \text{pttrn} \Rightarrow (\text{bool}, ' \alpha) \text{ uexpr} \Rightarrow ('b, ' \alpha) \text{ uexpr} \Rightarrow ('b \text{ set}, ' \alpha) \text{ uexpr} \ ((1\{- \mid - \cdot / -\}_u))$

#### lift-definition ZedSetCompr ::

$('a \text{ set}, ' \alpha) \text{ uexpr} \Rightarrow ('a \Rightarrow (\text{bool}, ' \alpha) \text{ uexpr} \times ('b, ' \alpha) \text{ uexpr}) \Rightarrow ('b \text{ set}, ' \alpha) \text{ uexpr}$   
**is**  $\lambda A \ PF \ b. \ \{ \text{snd} \ (PF \ x) \ b \mid x. \ x \in A \ b \wedge \text{fst} \ (PF \ x) \ b \} .$

### translations

$\{x..y\}_u == \text{CONST } \text{bop } \text{CONST } \text{atLeastAtMost } x \ y$   
 $\{x..<y\}_u == \text{CONST } \text{bop } \text{CONST } \text{atLeastLessThan } x \ y$   
 $\{x \mid P \cdot F\}_u == \text{CONST } \text{ZedSetCompr } (\text{CONST } \text{ulit } \text{CONST } \text{UNIV}) (\lambda x. (P, F))$   
 $\{x : A \mid P \cdot F\}_u == \text{CONST } \text{ZedSetCompr } A (\lambda x. (P, F))$

### 3.7 Lifting limits

We also lift the following functions on topological spaces for taking function limits, and describing continuity.

**definition**  $\text{ulim-left} :: 'a::\text{order-topology} \Rightarrow ('a \Rightarrow 'b) \Rightarrow 'b::t2\text{-space}$  **where**  
 $\text{ulim-left} = (\lambda p \ f. \text{Lim } (\text{at-left } p) \ f)$

**definition**  $\text{ulim-right} :: 'a::\text{order-topology} \Rightarrow ('a \Rightarrow 'b) \Rightarrow 'b::t2\text{-space}$  **where**  
 $\text{ulim-right} = (\lambda p \ f. \text{Lim } (\text{at-right } p) \ f)$

**definition**  $\text{ucont-on} :: ('a::\text{topological-space} \Rightarrow 'b::\text{topological-space}) \Rightarrow 'a \text{ set} \Rightarrow \text{bool}$  **where**  
 $\text{ucont-on} = (\lambda f \ A. \text{continuous-on } A \ f)$

### syntax

$\text{-ulim-left} :: \text{id} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } (\lim_u '(- \rightarrow -^-)'(-))$   
 $\text{-ulim-right} :: \text{id} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } (\lim_u '(- \rightarrow -^+)'(-))$   
 $\text{-ucont-on} :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } (\mathbf{infix } \text{cont-on}_u \ 90)$

### translations

$\lim_u(x \rightarrow p^-)(e) == \text{CONST } \text{bop } \text{CONST } \text{ulim-left } p (\lambda x \cdot e)$   
 $\lim_u(x \rightarrow p^+)(e) == \text{CONST } \text{bop } \text{CONST } \text{ulim-right } p (\lambda x \cdot e)$   
 $f \text{ cont-on}_u A == \text{CONST } \text{bop } \text{CONST } \text{continuous-on } A \ f$

### 3.8 Evaluation laws for expressions

We now collect together all the definitional theorems for expression constructs, and use them to build an evaluation strategy for expressions that we will later use to construct proof tactics for UTP predicates.

**lemmas**  $\text{ueexpr-defs} =$

$\text{zero-ueexpr-def}$   
 $\text{one-ueexpr-def}$   
 $\text{plus-ueexpr-def}$   
 $\text{uminus-ueexpr-def}$   
 $\text{minus-ueexpr-def}$   
 $\text{times-ueexpr-def}$   
 $\text{inverse-ueexpr-def}$   
 $\text{divide-ueexpr-def}$   
 $\text{sgn-ueexpr-def}$   
 $\text{abs-ueexpr-def}$   
 $\text{mod-ueexpr-def}$   
 $\text{eq-upred-def}$   
 $\text{numeral-ueexpr-simp}$   
 $\text{ulim-left-def}$   
 $\text{ulim-right-def}$   
 $\text{ucont-on-def}$   
 $\text{plus-list-def}$

The following laws show how to evaluate the core expressions constructs in terms of which the

above definitions are defined. Thus, using these theorems together, we can convert any UTP expression into a pure HOL expression. All these theorems are marked as *ueval* theorems which can be used for evaluation.

**lemma** *lit-ueval* [*ueval*]:  $\llbracket \langle x \rangle \rrbracket_e b = x$   
**by** (*transfer*, *simp*)

**lemma** *var-ueval* [*ueval*]:  $\llbracket \text{var } x \rrbracket_e b = \text{get}_x b$   
**by** (*transfer*, *simp*)

**lemma** *uop-ueval* [*ueval*]:  $\llbracket \text{uop } f \ x \rrbracket_e b = f (\llbracket x \rrbracket_e b)$   
**by** (*transfer*, *simp*)

**lemma** *bop-ueval* [*ueval*]:  $\llbracket \text{bop } f \ x \ y \rrbracket_e b = f (\llbracket x \rrbracket_e b) (\llbracket y \rrbracket_e b)$   
**by** (*transfer*, *simp*)

**lemma** *trop-ueval* [*ueval*]:  $\llbracket \text{trop } f \ x \ y \ z \rrbracket_e b = f (\llbracket x \rrbracket_e b) (\llbracket y \rrbracket_e b) (\llbracket z \rrbracket_e b)$   
**by** (*transfer*, *simp*)

**lemma** *qtop-ueval* [*ueval*]:  $\llbracket \text{qtop } f \ x \ y \ z \ w \rrbracket_e b = f (\llbracket x \rrbracket_e b) (\llbracket y \rrbracket_e b) (\llbracket z \rrbracket_e b) (\llbracket w \rrbracket_e b)$   
**by** (*transfer*, *simp*)

We also add all the definitional expressions to the evaluation theorem set.

**declare** *ueval-defs* [*ueval*]

### 3.9 Misc laws

We also prove a few useful algebraic and expansion laws for expressions.

**lemma** *uop-const* [*simp*]:  $\text{uop } \text{id } u = u$   
**by** (*transfer*, *simp*)

**lemma** *bop-const-1* [*simp*]:  $\text{bop } (\lambda x \ y. \ y) \ u \ v = v$   
**by** (*transfer*, *simp*)

**lemma** *bop-const-2* [*simp*]:  $\text{bop } (\lambda x \ y. \ x) \ u \ v = u$   
**by** (*transfer*, *simp*)

**lemma** *uinter-empty-1* [*simp*]:  $x \cap_u \{\} = \{\}$   
**by** (*transfer*, *simp*)

**lemma** *uinter-empty-2* [*simp*]:  $\{\} \cap_u x = \{\}$   
**by** (*transfer*, *simp*)

**lemma** *union-empty-1* [*simp*]:  $\{\} \cup_u x = x$   
**by** (*transfer*, *simp*)

**lemma** *uset-minus-empty* [*simp*]:  $x - \{\} = x$   
**by** (*simp add: ueval-defs*, *transfer*, *simp*)

**lemma** *ulist-filter-empty* [*simp*]:  $x \upharpoonright_u \{\} = \langle \rangle$   
**by** (*transfer*, *simp*)

**lemma** *tail-cons* [*simp*]:  $\text{tail}_u (\langle x \rangle \hat{\ }_u xs) = xs$   
**by** (*transfer*, *simp*)

**lemma** *ufun-apply-lit* [*simp*]:  
 $\llbracket f \rrbracket (\llbracket x \rrbracket)_a = \llbracket f(x) \rrbracket$   
**by** (*transfer*, *simp*)

### 3.10 Literalise tactics

The following tactic converts literal HOL expressions to UTP expressions and vice-versa via a collection of simplification rules. The two tactics are called "literalise", which converts UTP to expressions to HOL expressions – i.e. it pushes them into literals – and *unliteralise* that reverses this. We collect the equations in a theorem attribute called "lit\_simps".

**lemma** *lit-num-simps* [*lit\_simps*]:  $\llbracket 0 \rrbracket = 0$   $\llbracket 1 \rrbracket = 1$   $\llbracket \text{numeral } n \rrbracket = \text{numeral } n$   $\llbracket -x \rrbracket = - \llbracket x \rrbracket$   
**by** (*simp-all* add: *ueval*, *transfer*, *simp*)

**lemma** *lit-arith-simps* [*lit\_simps*]:  
 $\llbracket -x \rrbracket = - \llbracket x \rrbracket$   
 $\llbracket x + y \rrbracket = \llbracket x \rrbracket + \llbracket y \rrbracket$   $\llbracket x - y \rrbracket = \llbracket x \rrbracket - \llbracket y \rrbracket$   
 $\llbracket x * y \rrbracket = \llbracket x \rrbracket * \llbracket y \rrbracket$   $\llbracket x / y \rrbracket = \llbracket x \rrbracket / \llbracket y \rrbracket$   
 $\llbracket x \text{ div } y \rrbracket = \llbracket x \rrbracket \text{ div } \llbracket y \rrbracket$   
**by** (*simp* add: *uexpr-defs*, *transfer*, *simp*) +

**lemma** *lit-fun-simps* [*lit\_simps*]:  
 $\llbracket i \ x \ y \ z \ u \rrbracket = \text{qtop } i \ \llbracket x \rrbracket \ \llbracket y \rrbracket \ \llbracket z \rrbracket \ \llbracket u \rrbracket$   
 $\llbracket h \ x \ y \ z \rrbracket = \text{trop } h \ \llbracket x \rrbracket \ \llbracket y \rrbracket \ \llbracket z \rrbracket$   
 $\llbracket g \ x \ y \rrbracket = \text{bop } g \ \llbracket x \rrbracket \ \llbracket y \rrbracket$   
 $\llbracket f \ x \rrbracket = \text{uop } f \ \llbracket x \rrbracket$   
**by** (*transfer*, *simp*) +

In general *unliteralise* converts function applications to corresponding expression liftings. Since some operators, like + and \*, have specific operators we also have to use  $0 = \llbracket \_ \rrbracket_u$

$1 = \llbracket 1 :: ?'a \rrbracket$   
 $?u + ?v = \text{bop } op + ?u \ ?v$   
 $- ?u = \text{uop } uminus \ ?u$   
 $?u - ?v = \text{bop } op - ?u \ ?v$   
 $?u * ?v = \text{bop } op * ?u \ ?v$   
 $\text{inverse } ?u = \text{uop } inverse \ ?u$   
 $?u \text{ div } ?v = \text{bop } op \text{ div } ?u \ ?v$   
 $\text{sgn } ?u = \text{uop } sgn \ ?u$   
 $|?u| = \text{uop } abs \ ?u$   
 $?u \text{ mod } ?v = \text{bop } op \text{ mod } ?u \ ?v$   
 $(?x =_u ?y) = \text{bop } op = ?x \ ?y$   
 $\text{numeral } ?x = \llbracket \text{numeral } ?x \rrbracket$   
 $\text{ulim-left} = (\lambda p. \text{Lim } (\text{at-left } p))$   
 $\text{ulim-right} = (\lambda p. \text{Lim } (\text{at-right } p))$   
 $\text{ucont-on} = (\lambda f \ A. \text{continuous-on } A \ f)$

$op + = op @$  in reverse to correctly interpret these. Moreover, numerals must be handled separately by first simplifying them and then converting them into UTP expression numerals; hence the following two simplification rules.

**lemma** *lit-numeral-1*:  $\text{uop numeral } x = \text{Abs-uexpr } (\lambda b. \text{numeral } (\llbracket x \rrbracket_e \ b))$   
**by** (*simp* add: *uop-def*)

```

lemma lit-numeral-2: Abs-uepr ( $\lambda b. \text{numeral } v$ ) = numeral v
  by (metis lit.abs-eq lit-num-simps(3))

method literalise = (unfold lit-simps[THEN sym])
method unliteralise = (unfold lit-simps uepr-defs[THEN sym];
  (unfold lit-numeral-1 ; (unfold ueval); (unfold lit-numeral-2)))?+
end

```

## 4 Unrestriction

```

theory utp-unrest
  imports utp-expr
begin

```

### 4.1 Definitions and Core Syntax

Unrestriction is an encoding of semantic freshness that allows us to reason about the presence of variables in predicates without being concerned with abstract syntax trees. An expression  $p$  is unrestricted by lens  $x$ , written  $x \# p$ , if altering the value of  $x$  has no effect on the valuation of  $p$ . This is a sufficient notion to prove many laws that would ordinarily rely on an  $fv$  function.

Unrestriction was first defined in the work of Marcel Oliveira [10, 9] in his UTP mechanisation in *ProofPowerZ*. Our definition modifies his in that our variables are semantically characterised as lenses, and supported by the lens laws, rather than named syntactic entities. We effectively fuse the ideas from both Feliachi [3] and Oliveira's [9] mechanisations of the UTP, the former being also purely semantic in nature.

We first set up overloaded syntax for unrestricted, as several concepts will have this defined.

```

consts
  unrest :: 'a  $\Rightarrow$  'b  $\Rightarrow$  bool

syntax
  -unrest :: salpha  $\Rightarrow$  logic  $\Rightarrow$  logic  $\Rightarrow$  logic (infix  $\#$  20)

translations
  -unrest x p == CONST unrest x p
  -unrest (-salphaset (-salphamk (x +L y))) P <= -unrest (x +L y) P

```

Our syntax translations support both variables and variable sets such that we can write down predicates like  $\&x \# P$  and also  $\{\&x, \&y, \&z\} \# P$ .

We set up a simple tactic for discharging unrestricted conjectures using a simplification set.

```

named-theorems unrest
method unrest-tac = (simp add: unrest)?

```

Unrestriction for expressions is defined as a lifted construct using the underlying lens operations. It states that lens  $x$  is unrestricted by expression  $e$  provided that, for any state-space binding  $b$  and variable valuation  $v$ , the value which the expression evaluates to is unaltered if we set  $x$  to  $v$  in  $b$ . In other words, we cannot effect the behaviour of  $e$  by changing  $x$ . Thus  $e$  does not observe the portion of state-space characterised by  $x$ . We add this definition to our overloaded constant.

```

lift-definition unrest-uepr :: ('a  $\Rightarrow$  'α)  $\Rightarrow$  ('b, 'α) uepr  $\Rightarrow$  bool
is  $\lambda x e. \forall b v. e \text{ (put}_x b v) = e b$  .

```



## adhoc-overloading

*unrest unrest-uepr*

**lemma** *unrest-expr-alt-def*:

*weak-lens*  $x \implies (x \# P) = (\forall b \ b'. \llbracket P \rrbracket_e (b \oplus_L b' \text{ on } x) = \llbracket P \rrbracket_e b)$   
**by** (*transfer*, *metis lens-override-def weak-lens.put-get*)

## 4.2 Unrestriction laws

We now prove unrestricted laws for the key constructs of our expression model. Many of these depend on lens properties and so variously employ the assumptions *mwb-lens* and *vwb-lens*, depending on the number of assumptions from the lenses theory is required.

Firstly, we prove a general property – if  $x$  and  $y$  are both unrestricted in  $P$ , then their composition is also unrestricted in  $P$ . One can interpret the composition here as a union – if the two sets of variables  $x$  and  $y$  are unrestricted, then so is their union.

**lemma** *unrest-var-comp* [*unrest*]:

$\llbracket x \# P; y \# P \rrbracket \implies x; y \# P$   
**by** (*transfer*, *simp add: lens-defs*)

**lemma** *unrest-svar* [*unrest*]:  $x \# P \implies \&x \# P$

**by** (*transfer*, *simp add: lens-defs*)

No lens is restricted by a literal, since it returns the same value for any state binding.

**lemma** *unrest-lit* [*unrest*]:  $x \# \llbracket v \rrbracket$

**by** (*transfer*, *simp*)

If one lens is smaller than another, then any unrestricted on the larger lens implies unrestricted on the smaller.

**lemma** *unrest-sublens*:

**fixes**  $P :: ('a, 'a) \text{ uepr}$   
**assumes**  $x \# P \ y \subseteq_L x$   
**shows**  $y \# P$   
**using** *assms*  
**by** (*transfer*, *metis (no-types, lifting) lens.select-convs(2) lens-comp-def sublens-def*)

If two lenses are equivalent, and thus they characterise the same state-space regions, then clearly unrestrictedions over them are equivalent.

**lemma** *unrest-equiv*:

**fixes**  $P :: ('a, 'a) \text{ uepr}$   
**assumes** *mwb-lens*  $y \ x \approx_L y \ x \# P$   
**shows**  $y \# P$   
**by** (*metis assms lens-equiv-def sublens-pres-mwb sublens-put-put unrest-uepr.rep-eq*)

The following laws demonstrate the primary motivation for lens independence: a variable expression is unrestricted by another variable only when the two variables are independent. Lens independence thus effectively allows us to semantically characterise when two variables, or sets of variables, are different.

**lemma** *unrest-var* [*unrest*]:  $\llbracket \text{mwb-lens } x; x \bowtie y \rrbracket \implies y \# \text{var } x$

**by** (*transfer*, *auto*)

**lemma** *unrest-iuvar* [*unrest*]:  $\llbracket \text{mwb-lens } x; x \bowtie y \rrbracket \implies \$y \# \$x$

by (simp add: unrest-var)

**lemma** *unrest-ouvar* [unrest]:  $\llbracket \text{mwb-lens } x; x \bowtie y \rrbracket \Longrightarrow \$y' \# \$x'$   
 by (simp add: unrest-var)

The following laws follow automatically from independence of input and output variables.

**lemma** *unrest-iuvar-ouvar* [unrest]:  
 fixes  $x :: ('a \Longrightarrow 'a)$   
 assumes *mwb-lens*  $y$   
 shows  $\$x \# \$y'$   
 by (metis prod.collapse unrest-uexpr.rep-eq var.rep-eq var-lookup-out var-update-in)

**lemma** *unrest-ouvar-iuvar* [unrest]:  
 fixes  $x :: ('a \Longrightarrow 'a)$   
 assumes *mwb-lens*  $y$   
 shows  $\$x' \# \$y$   
 by (metis prod.collapse unrest-uexpr.rep-eq var.rep-eq var-lookup-in var-update-out)

Unrestriction distributes through the various function lifting expression constructs; this allows us to prove unrestrictions for the majority of the expression language.

**lemma** *unrest-uop* [unrest]:  $x \# e \Longrightarrow x \# \text{uop } f e$   
 by (transfer, simp)

**lemma** *unrest-bop* [unrest]:  $\llbracket x \# u; x \# v \rrbracket \Longrightarrow x \# \text{bop } f u v$   
 by (transfer, simp)

**lemma** *unrest-trop* [unrest]:  $\llbracket x \# u; x \# v; x \# w \rrbracket \Longrightarrow x \# \text{trop } f u v w$   
 by (transfer, simp)

**lemma** *unrest-qtop* [unrest]:  $\llbracket x \# u; x \# v; x \# w; x \# y \rrbracket \Longrightarrow x \# \text{qtop } f u v w y$   
 by (transfer, simp)

For convenience, we also prove unrestriction rules for the bespoke operators on equality, numbers, arithmetic etc.

**lemma** *unrest-eq* [unrest]:  $\llbracket x \# u; x \# v \rrbracket \Longrightarrow x \# u =_u v$   
 by (simp add: eq-upred-def, transfer, simp)

**lemma** *unrest-zero* [unrest]:  $x \# 0$   
 by (simp add: unrest-lit zero-uexpr-def)

**lemma** *unrest-one* [unrest]:  $x \# 1$   
 by (simp add: one-uexpr-def unrest-lit)

**lemma** *unrest-numeral* [unrest]:  $x \# (\text{numeral } n)$   
 by (simp add: numeral-uexpr-simp unrest-lit)

**lemma** *unrest-sgn* [unrest]:  $x \# u \Longrightarrow x \# \text{sgn } u$   
 by (simp add: sgn-uexpr-def unrest-uop)

**lemma** *unrest-abs* [unrest]:  $x \# u \Longrightarrow x \# \text{abs } u$   
 by (simp add: abs-uexpr-def unrest-uop)

**lemma** *unrest-plus* [unrest]:  $\llbracket x \# u; x \# v \rrbracket \Longrightarrow x \# u + v$   
 by (simp add: plus-uexpr-def unrest)

```

lemma unrest-uminus [unrest]:  $x \# u \implies x \# - u$ 
  by (simp add: uminus-uepr-def unrest)

lemma unrest-minus [unrest]:  $\llbracket x \# u; x \# v \rrbracket \implies x \# u - v$ 
  by (simp add: minus-uepr-def unrest)

lemma unrest-times [unrest]:  $\llbracket x \# u; x \# v \rrbracket \implies x \# u * v$ 
  by (simp add: times-uepr-def unrest)

lemma unrest-divide [unrest]:  $\llbracket x \# u; x \# v \rrbracket \implies x \# u / v$ 
  by (simp add: divide-uepr-def unrest)

```

For a  $\lambda$ -term we need to show that the characteristic function expression does not restrict  $v$  for any input value  $x$ .

```

lemma unrest-ulambda [unrest]:
   $\llbracket \bigwedge x. v \# F x \rrbracket \implies v \# (\lambda x. F x)$ 
  by (transfer, simp)

```

**end**

## 5 Used-by

```

theory utp-usedby
  imports utp-unrest
begin

```

The used-by predicate is the dual of unrestriction. It states that the given lens is an upper-bound on the size of state space the given expression depends on. It is similar to stating that the lens is a valid alphabet for the predicate. For convenience, and because the predicate uses a similar form, we will reuse much of unrestriction's infrastructure.

```

consts
  usedBy :: 'a  $\Rightarrow$  'b  $\Rightarrow$  bool

```

```

syntax
  -usedBy :: salpha  $\Rightarrow$  logic  $\Rightarrow$  logic  $\Rightarrow$  logic (infix  $\#$  20)

```

```

translations
  -usedBy x p == CONST usedBy x p
  -usedBy (-salphaset (-salphamk (x +L y))) P <= -usedBy (x +L y) P

```

```

lift-definition usedBy-uepr :: ('b  $\implies$  'α)  $\Rightarrow$  ('a, 'α) uepr  $\Rightarrow$  bool
is  $\lambda x e. (\forall b b'. e (b' \oplus_L b \text{ on } x) = e b)$  .

```

```

adhoc-overloading usedBy usedBy-uepr

```

```

lemma usedBy-lit [unrest]:  $x \# \langle v \rangle$ 
  by (transfer, simp)

```

```

lemma usedBy-sublens:
  fixes P :: ('a, 'α) uepr
  assumes  $x \# P x \subseteq_L y \text{ vwb-lens } y$ 
  shows  $y \# P$ 
  using assms
  by (transfer, auto,metis lens-override-def lens-override-idem sublens-obs-get vwb-lens-mwb)

```

**lemma** *usedBy-svar* [unrest]:  $x \Vdash P \implies \&x \Vdash P$   
**by** (*transfer*, *simp* add: *lens-defs*)

**lemma** *usedBy-lens-plus-1* [unrest]:  $x \Vdash P \implies x;y \Vdash P$   
**by** (*transfer*, *simp* add: *lens-defs*)

**lemma** *usedBy-lens-plus-2* [unrest]:  $\llbracket x \bowtie y; y \Vdash P \rrbracket \implies x;y \Vdash P$   
**by** (*transfer*, *auto* *simp* add: *lens-defs* *lens-indep-comm*)

Linking used-by to unrestricted: if  $x$  is used-by  $P$ , and  $x$  is independent of  $y$ , then  $P$  cannot depend on any variable in  $y$ .

**lemma** *usedBy-indep-uses*:  
**fixes**  $P :: ('a, 'α) \text{ uexpr}$   
**assumes**  $x \Vdash P \bowtie y$   
**shows**  $y \nVdash P$   
**using** *assms* **by** (*transfer*, *auto*, *metis* *lens-indep-get* *lens-override-def*)

**lemma** *usedBy-var* [unrest]:  
**assumes** *vwb-lens*  $x \ y \subseteq_L x$   
**shows**  $x \Vdash \text{var } y$   
**using** *assms*  
**by** (*transfer*, *simp* add: *uexpr-defs* *pr-var-def*)  
(*metis* *lens-override-def* *lens-override-idem* *sublens-obs-get* *vwb-lens-mwb*)

**lemma** *usedBy-uop* [unrest]:  $x \Vdash e \implies x \Vdash \text{uop } f \ e$   
**by** (*transfer*, *simp*)

**lemma** *usedBy-bop* [unrest]:  $\llbracket x \Vdash u; x \Vdash v \rrbracket \implies x \Vdash \text{bop } f \ u \ v$   
**by** (*transfer*, *simp*)

**lemma** *usedBy-trop* [unrest]:  $\llbracket x \Vdash u; x \Vdash v; x \Vdash w \rrbracket \implies x \Vdash \text{trop } f \ u \ v \ w$   
**by** (*transfer*, *simp*)

**lemma** *usedBy-qtop* [unrest]:  $\llbracket x \Vdash u; x \Vdash v; x \Vdash w; x \Vdash y \rrbracket \implies x \Vdash \text{qtop } f \ u \ v \ w \ y$   
**by** (*transfer*, *simp*)

For convenience, we also prove used-by rules for the bespoke operators on equality, numbers, arithmetic etc.

**lemma** *usedBy-eq* [unrest]:  $\llbracket x \Vdash u; x \Vdash v \rrbracket \implies x \Vdash u =_u v$   
**by** (*simp* add: *eq-upred-def*, *transfer*, *simp*)

**lemma** *usedBy-zero* [unrest]:  $x \Vdash 0$   
**by** (*simp* add: *usedBy-lit* *zero-uexpr-def*)

**lemma** *usedBy-one* [unrest]:  $x \Vdash 1$   
**by** (*simp* add: *one-uexpr-def* *usedBy-lit*)

**lemma** *usedBy-numeral* [unrest]:  $x \Vdash (\text{numeral } n)$   
**by** (*simp* add: *numeral-uexpr-simp* *usedBy-lit*)

**lemma** *usedBy-sgn* [unrest]:  $x \Vdash u \implies x \Vdash \text{sgn } u$   
**by** (*simp* add: *sgn-uexpr-def* *usedBy-uop*)

**lemma** *usedBy-abs* [unrest]:  $x \Vdash u \implies x \Vdash \text{abs } u$

```

by (simp add: abs-ueexpr-def usedBy-uop)

lemma usedBy-plus [unrest]:  $\llbracket x \Downarrow u; x \Downarrow v \rrbracket \implies x \Downarrow u + v$ 
  by (simp add: plus-ueexpr-def unrest)

lemma usedBy-uminus [unrest]:  $x \Downarrow u \implies x \Downarrow - u$ 
  by (simp add: uminus-ueexpr-def unrest)

lemma usedBy-minus [unrest]:  $\llbracket x \Downarrow u; x \Downarrow v \rrbracket \implies x \Downarrow u - v$ 
  by (simp add: minus-ueexpr-def unrest)

lemma usedBy-times [unrest]:  $\llbracket x \Downarrow u; x \Downarrow v \rrbracket \implies x \Downarrow u * v$ 
  by (simp add: times-ueexpr-def unrest)

lemma usedBy-divide [unrest]:  $\llbracket x \Downarrow u; x \Downarrow v \rrbracket \implies x \Downarrow u / v$ 
  by (simp add: divide-ueexpr-def unrest)

lemma usedBy-ulambda [unrest]:
   $\llbracket \bigwedge x. v \Downarrow F x \rrbracket \implies v \Downarrow (\lambda x. F x)$ 
  by (transfer, simp)

end

```

## 6 Substitution

```

theory utp-subst
imports
  utp-expr
  utp-unrest
begin

```

### 6.1 Substitution definitions

Variable substitution, like unrestriction, will be characterised semantically using lenses and state-spaces. Effectively a substitution  $\sigma$  is simply a function on the state-space which can be applied to an expression  $e$  using the syntax  $\sigma \dagger e$ . We introduce a polymorphic constant that will be used to represent application of a substitution, and also a set of theorems to represent laws.

```

consts
  usubst :: 's  $\Rightarrow$  'a  $\Rightarrow$  'b (infixr  $\dagger$  80)

```

```

named-theorems usubst

```

A substitution is simply a transformation on the alphabet; it shows how variables should be mapped to different values. Most of the time these will be homogeneous functions but for flexibility we also allow some operations to be heterogeneous.

```

type-synonym (' $\alpha$ , ' $\beta$ ) psubst = ' $\alpha \Rightarrow$  ' $\beta$ 
type-synonym ' $\alpha$  usubst = ' $\alpha \Rightarrow$  ' $\alpha$ 

```

Application of a substitution simply applies the function  $\sigma$  to the state binding  $b$  before it is handed to  $e$  as an input. This effectively ensures all variables are updated in  $e$ .

```

lift-definition subst :: (' $\alpha$ , ' $\beta$ ) psubst  $\Rightarrow$  ('a, ' $\alpha$ ) ueexpr  $\Rightarrow$  ('a, ' $\alpha$ ) ueexpr is
 $\lambda \sigma e b. e (\sigma b)$  .

```

### ad hoc-overloading

*usubst subst*

Substitutions can be updated by associating variables with expressions. We thus create an additional polymorphic constant to represent updating the value of a variable to an expression in a substitution, where the variable is modelled by type  $'v$ . This again allows us to support different notions of variables, such as deep variables, later.

**consts** *subst-upd* ::  $('α, 'β) psubst \Rightarrow 'v \Rightarrow ('a, 'α) uexpr \Rightarrow ('α, 'β) psubst$

The following function takes a substitution from state-space  $'α$  to  $'β$ , a lens with source  $'β$  and view  $'a$ , and an expression over  $'α$  and returning a value of type  $'a$ , and produces an updated substitution. It does this by constructing a substitution function that takes state binding  $b$ , and updates the state first by applying the original substitution  $\sigma$ , and then updating the part of the state associated with lens  $x$  with expression evaluated in the context of  $b$ . This effectively means that  $x$  is now associated with expression  $v$ . We add this definition to our overloaded constant.

**definition** *subst-upd-uvar* ::  $('α, 'β) psubst \Rightarrow ('a \Rightarrow 'β) \Rightarrow ('a, 'α) uexpr \Rightarrow ('α, 'β) psubst$  **where**  
*subst-upd-uvar*  $\sigma x v = (\lambda b. put_x (\sigma b) (\llbracket v \rrbracket_e b))$

### ad hoc-overloading

*subst-upd subst-upd-uvar*

The next function looks up the expression associated with a variable in a substitution by use of the *get* lens function.

**lift-definition** *usubst-lookup* ::  $('α, 'β) psubst \Rightarrow ('a \Rightarrow 'β) \Rightarrow ('a, 'α) uexpr (\langle - \rangle_s)$   
**is**  $\lambda \sigma x b. get_x (\sigma b)$ .

Substitutions also exhibit a natural notion of unrestriction which states that  $\sigma$  does not restrict  $x$  if application of  $\sigma$  to an arbitrary state  $\rho$  will not effect the valuation of  $x$ . Put another way, it requires that *put* and the substitution commute.

**definition** *unrest-usubst* ::  $('a \Rightarrow 'α) \Rightarrow 'α usubst \Rightarrow bool$   
**where** *unrest-usubst*  $x \sigma = (\forall \rho v. \sigma (put_x \rho v) = put_x (\sigma \rho) v)$

### ad hoc-overloading

*unrest unrest-usubst*

Parallel substitutions allow us to divide the state space into three segments using two lens, A and B. They correspond to the part of the state that should be updated by the respective substitution. The two lenses should be independent. If any part of the state is not covered by either lenses then this area is left unchanged (framed).

**definition** *par-subst* ::  $'α usubst \Rightarrow ('a \Rightarrow 'α) \Rightarrow ('b \Rightarrow 'α) \Rightarrow 'α usubst \Rightarrow 'α usubst$  **where**  
*par-subst*  $\sigma_1 A B \sigma_2 = (\lambda s. (s \oplus_L (\sigma_1 s) \text{ on } A) \oplus_L (\sigma_2 s) \text{ on } B)$

## 6.2 Syntax translations

We support two kinds of syntax for substitutions, one where we construct a substitution using a maplet-style syntax, with variables mapping to expressions. Such a constructed substitution can be applied to an expression. Alternatively, we support the more traditional notation,  $P\llbracket v/x \rrbracket$ , which also support multiple simultaneous substitutions. We have to use double square brackets as the single ones are already well used.

We set up non-terminals to represent a single substitution maplet, a sequence of maplets, a list of expressions, and a list of alphabets. The parser effectively uses *subst-upd* to construct substitutions from multiple variables.

**nonterminal** *smaplet* **and** *smaplets* **and** *uexprs* **and** *salphas*

**syntax**

```
-smaplet :: [salpha, 'a] => smaplet      (- /↦s/ -)
           :: smaplet => smaplets        (-)
-SMaplets :: [smaplet, smaplets] => smaplets (-,/ -)
-SubstUpd :: ['m usubst, smaplets] => 'm usubst (-/'(-) [900,0] 900)
-Subst    :: smaplets => 'a ↦ 'b         ((1[-]))
-psubst   :: [logic, svars, uexprs] ⇒ logic
-subst    :: logic ⇒ uexprs ⇒ salphas ⇒ logic ((-['/-]) [990,0,0] 991)
-uexprs   :: [logic, uexprs] => uexprs (-,/ -)
           :: logic => uexprs (-)
-salphas  :: [salpha, salphas] => salphas (-,/ -)
           :: salpha => salphas (-)
-par-subst :: logic ⇒ salpha ⇒ salpha ⇒ logic ⇒ logic (- [-]ₛ - [100,0,0,101] 101)
```

**translations**

```
-SubstUpd m (-SMaplets xy ms) == -SubstUpd (-SubstUpd m xy) ms
-SubstUpd m (-smaplet x y)    == CONST subst-upd m x y
-Subst ms                     == -SubstUpd (CONST id) ms
-Subst (-SMaplets ms1 ms2)    <= -SubstUpd (-Subst ms1) ms2
-SMaplets ms1 (-SMaplets ms2 ms3) <= -SMaplets (-SMaplets ms1 ms2) ms3
-subst P es vs => CONST subst (-psubst (CONST id) vs es) P
-psubst m (-salphas x xs) (-uexprs v vs) => -psubst (-psubst m x v) xs vs
-psubst m x v => CONST subst-upd m x v
P[v/$x] <= CONST usubst (CONST subst-upd (CONST id) (CONST ivar x) v) P
P[v/$x ^] <= CONST usubst (CONST subst-upd (CONST id) (CONST ovar x) v) P
P[v/x] <= CONST usubst (CONST subst-upd (CONST id) x v) P
-par-subst σ₁ A B σ₂ == CONST par-subst σ₁ A B σ₂
```

Thus we can write things like  $\sigma(x \mapsto_s v)$  to update a variable  $x$  in  $\sigma$  with expression  $v$ ,  $[x \mapsto_s e, y \mapsto_s f]$  to construct a substitution with two variables, and finally  $P[v/x]$ , the traditional syntax.

We can now express deletion of a substitution maplet.

**definition** *subst-del* :: 'α usubst ⇒ ('a ⇒ 'α) ⇒ 'α usubst (**infix**  $-_s$  85) **where**  
*subst-del* σ  $x = \sigma(x \mapsto_s \&x)$

### 6.3 Substitution application laws

We set up a simple substitution tactic that applies substitution and unrestriction laws

**method** *subst-tac* = (*simp add: usubst unrest*)?

Evaluation of a substitution expression involves application of the substitution to different variables. Thus we first prove laws for these cases. The simplest substitution, *id*, when applied to any variable  $x$  simply returns the variable expression, since *id* has no effect.

**lemma** *usubst-lookup-id* [*usubst*]:  $\langle id \rangle_s x = \text{var } x$   
**by** (*transfer, simp*)

**lemma** *subst-upd-id-lam* [*usubst*]:  $\text{subst-upd } (\lambda x. x) x v = \text{subst-upd } id x v$   
**by** (*simp add: id-def*)

A substitution update naturally yields the given expression.

**lemma** *usubst-lookup-upd* [*usubst*]:  
**assumes** *mwb-lens* *x*  
**shows**  $\langle \sigma(x \mapsto_s v) \rangle_s x = v$   
**using** *assms*  
**by** (*simp add: subst-upd-uvar-def, transfer*) (*simp*)

**lemma** *usubst-lookup-upd-pr-var* [*usubst*]:  
**assumes** *mwb-lens* *x*  
**shows**  $\langle \sigma(x \mapsto_s v) \rangle_s (pr\text{-}var\ x) = v$   
**using** *assms*  
**by** (*simp add: subst-upd-uvar-def pr-var-def, transfer*) (*simp*)

Substitution update is idempotent.

**lemma** *usubst-upd-idem* [*usubst*]:  
**assumes** *mwb-lens* *x*  
**shows**  $\sigma(x \mapsto_s u, x \mapsto_s v) = \sigma(x \mapsto_s v)$   
**by** (*simp add: subst-upd-uvar-def assms comp-def*)

Substitution updates commute when the lenses are independent.

**lemma** *usubst-upd-comm*:  
**assumes**  $x \bowtie y$   
**shows**  $\sigma(x \mapsto_s u, y \mapsto_s v) = \sigma(y \mapsto_s v, x \mapsto_s u)$   
**using** *assms*  
**by** (*rule-tac ext, auto simp add: subst-upd-uvar-def assms comp-def lens-indep-comm*)

**lemma** *usubst-upd-comm2*:  
**assumes**  $z \bowtie y$  **and** *mwb-lens* *x*  
**shows**  $\sigma(x \mapsto_s u, y \mapsto_s v, z \mapsto_s s) = \sigma(x \mapsto_s u, z \mapsto_s s, y \mapsto_s v)$   
**using** *assms*  
**by** (*rule-tac ext, auto simp add: subst-upd-uvar-def assms comp-def lens-indep-comm*)

**lemma** *subst-upd-pr-var* [*usubst*]:  $s(\&x \mapsto_s v) = s(x \mapsto_s v)$   
**by** (*simp add: pr-var-def*)

A substitution which swaps two independent variables is an injective function.

**lemma** *swap-usubst-inj*:  
**fixes**  $x\ y :: ('a \implies 'a)$   
**assumes** *vwb-lens* *x* *vwb-lens* *y*  $x \bowtie y$   
**shows** *inj*  $[x \mapsto_s \&y, y \mapsto_s \&x]$   
**proof** (*rule injI*)  
**fix**  $b_1 :: 'a$  **and**  $b_2 :: 'a$   
**assume**  $[x \mapsto_s \&y, y \mapsto_s \&x] b_1 = [x \mapsto_s \&y, y \mapsto_s \&x] b_2$   
**hence**  $a: put_y (put_x b_1 (\llbracket \&y \rrbracket_e b_1)) (\llbracket \&x \rrbracket_e b_1) = put_y (put_x b_2 (\llbracket \&y \rrbracket_e b_2)) (\llbracket \&x \rrbracket_e b_2)$   
**by** (*auto simp add: subst-upd-uvar-def*)  
**then have**  $(\forall a\ b\ c. put_x (put_y a\ b)\ c = put_y (put_x a\ c)\ b) \wedge$   
 $(\forall a\ b. get_x (put_y a\ b) = get_x a) \wedge (\forall a\ b. get_y (put_x a\ b) = get_y a)$   
**by** (*simp add: assms(3) lens-indep.lens-put-irr2 lens-indep-comm*)  
**then show**  $b_1 = b_2$   
**by** (*metis a assms(1) assms(2) pr-var-def var.rep-eq vwb-lens.source-determination vwb-lens-def*  
*wb-lens-def weak-lens.put-get*)  
**qed**

**lemma** *usubst-upd-var-id* [*usubst*]:



```

vwb-lens x  $\implies$  [x  $\mapsto_s$  var x] = id
apply (simp add: subst-upd-uvar-def)
apply (transfer)
apply (rule ext)
apply (auto)
done

```

```

lemma usubst-upd-pr-var-id [usubst]:
  vwb-lens x  $\implies$  [x  $\mapsto_s$  var (pr-var x)] = id
  apply (simp add: subst-upd-uvar-def pr-var-def)
  apply (transfer)
  apply (rule ext)
  apply (auto)
done

```

```

lemma usubst-upd-comm-dash [usubst]:
  fixes x :: ('a  $\implies$  'α)
  shows  $\sigma(\$x' \mapsto_s v, \$x \mapsto_s u) = \sigma(\$x \mapsto_s u, \$x' \mapsto_s v)$ 
  using out-in-indep usubst-upd-comm by blast

```

```

lemma subst-upd-lens-plus [usubst]:
  subst-upd  $\sigma$  (x +L y)  $\ll(u,v)\gg = \sigma(y \mapsto_s \ll v \gg, x \mapsto_s \ll u \gg)$ 
  by (simp add: lens-defs uexpr-defs subst-upd-uvar-def, transfer, auto)

```

```

lemma subst-upd-in-lens-plus [usubst]:
  subst-upd  $\sigma$  (ivar (x +L y))  $\ll(u,v)\gg = \sigma(\$y \mapsto_s \ll v \gg, \$x \mapsto_s \ll u \gg)$ 
  by (simp add: lens-defs uexpr-defs subst-upd-uvar-def, transfer, auto simp add: prod.case-eq-if)

```

```

lemma subst-upd-out-lens-plus [usubst]:
  subst-upd  $\sigma$  (ovar (x +L y))  $\ll(u,v)\gg = \sigma(\$y' \mapsto_s \ll v \gg, \$x' \mapsto_s \ll u \gg)$ 
  by (simp add: lens-defs uexpr-defs subst-upd-uvar-def, transfer, auto simp add: prod.case-eq-if)

```

```

lemma usubst-lookup-upd-indep [usubst]:
  assumes mwb-lens x x  $\boxtimes$  y
  shows  $\langle \sigma(y \mapsto_s v) \rangle_s x = \langle \sigma \rangle_s x$ 
  using assms
  by (simp add: subst-upd-uvar-def, transfer, simp)

```

If a variable is unrestricted in a substitution then it's application has no effect.

```

lemma usubst-apply-unrest [usubst]:
   $\ll vwb-lens x; x \# \sigma \gg \implies \langle \sigma \rangle_s x = \text{var } x$ 
  by (simp add: unrest-usubst-def, transfer, auto simp add: fun-eq-iff, metis vwb-lens-wb wb-lens.get-put
  wb-lens-weak weak-lens.put-get)

```

There follows various laws about deleting variables from a substitution.

```

lemma subst-del-id [usubst]:
  vwb-lens x  $\implies$  id -s x = id
  by (simp add: subst-del-def subst-upd-uvar-def pr-var-def, transfer, auto)

```

```

lemma subst-del-upd-same [usubst]:
  mwb-lens x  $\implies$   $\sigma(x \mapsto_s v) -_s x = \sigma -_s x$ 
  by (simp add: subst-del-def subst-upd-uvar-def)

```

```

lemma subst-del-upd-diff [usubst]:
  x  $\boxtimes$  y  $\implies$   $\sigma(y \mapsto_s v) -_s x = (\sigma -_s x)(y \mapsto_s v)$ 

```

by (simp add: subst-del-def subst-upd-uvar-def lens-indep-comm)

If a variable is unrestricted in an expression, then any substitution of that variable has no effect on the expression .

**lemma** *subst-unrest* [usubst]:  $x \# P \implies \sigma(x \mapsto_s v) \dagger P = \sigma \dagger P$   
 by (simp add: subst-upd-uvar-def, transfer, auto)

**lemma** *subst-compose-upd* [usubst]:  $x \# \sigma \implies \sigma \circ \varrho(x \mapsto_s v) = (\sigma \circ \varrho)(x \mapsto_s v)$   
 by (simp add: subst-upd-uvar-def, transfer, auto simp add: unrest-usubst-def)

Any substitution is a monotonic function.

**lemma** *subst-mono*: *mono* (subst  $\sigma$ )  
 by (simp add: less-eq-uexpr.rep-eq mono-def subst.rep-eq)

## 6.4 Substitution laws

We now prove the key laws that show how a substitution should be performed for every expression operator, including the core function operators, literals, variables, and the arithmetic operators. They are all added to the *usubst* theorem attribute so that we can apply them using the substitution tactic.

**lemma** *id-subst* [usubst]:  $id \dagger v = v$   
 by (transfer, simp)

**lemma** *subst-lit* [usubst]:  $\sigma \dagger \langle\!\langle v \rangle\!\rangle = \langle\!\langle v \rangle\!\rangle$   
 by (transfer, simp)

**lemma** *subst-var* [usubst]:  $\sigma \dagger \text{var } x = \langle\sigma\rangle_s x$   
 by (transfer, simp)

**lemma** *usubst-ulambda* [usubst]:  $\sigma \dagger (\lambda x \cdot P(x)) = (\lambda x \cdot \sigma \dagger P(x))$   
 by (transfer, simp)

**lemma** *unrest-usubst-del* [unrest]:  $\llbracket \text{vwb-lens } x; x \# (\langle\sigma\rangle_s x); x \# \sigma -_s x \rrbracket \implies x \# (\sigma \dagger P)$   
 by (simp add: subst-del-def subst-upd-uvar-def unrest-uexpr-def unrest-usubst-def subst.rep-eq usubst-lookup.rep-eq) (metis vwb-lens.put-eq)

We add the symmetric definition of input and output variables to substitution laws so that the variables are correctly normalised after substitution.

**lemma** *subst-uop* [usubst]:  $\sigma \dagger \text{uop } f v = \text{uop } f (\sigma \dagger v)$   
 by (transfer, simp)

**lemma** *subst-bop* [usubst]:  $\sigma \dagger \text{bop } f u v = \text{bop } f (\sigma \dagger u) (\sigma \dagger v)$   
 by (transfer, simp)

**lemma** *subst-trop* [usubst]:  $\sigma \dagger \text{trop } f u v w = \text{trop } f (\sigma \dagger u) (\sigma \dagger v) (\sigma \dagger w)$   
 by (transfer, simp)

**lemma** *subst-qtop* [usubst]:  $\sigma \dagger \text{qtop } f u v w x = \text{qtop } f (\sigma \dagger u) (\sigma \dagger v) (\sigma \dagger w) (\sigma \dagger x)$   
 by (transfer, simp)

**lemma** *subst-plus* [usubst]:  $\sigma \dagger (x + y) = \sigma \dagger x + \sigma \dagger y$   
 by (simp add: plus-uexpr-def subst-bop)

**lemma** *subst-times* [usubst]:  $\sigma \dagger (x * y) = \sigma \dagger x * \sigma \dagger y$

by (simp add: times-ueexpr-def subst-bop)

**lemma** subst-mod [usubst]:  $\sigma \dagger (x \text{ mod } y) = \sigma \dagger x \text{ mod } \sigma \dagger y$   
 by (simp add: mod-ueexpr-def usubst)

**lemma** subst-div [usubst]:  $\sigma \dagger (x \text{ div } y) = \sigma \dagger x \text{ div } \sigma \dagger y$   
 by (simp add: divide-ueexpr-def usubst)

**lemma** subst-minus [usubst]:  $\sigma \dagger (x - y) = \sigma \dagger x - \sigma \dagger y$   
 by (simp add: minus-ueexpr-def subst-bop)

**lemma** subst-uminus [usubst]:  $\sigma \dagger (- x) = - (\sigma \dagger x)$   
 by (simp add: uminus-ueexpr-def subst-uop)

**lemma** usubst-sgn [usubst]:  $\sigma \dagger \text{sgn } x = \text{sgn } (\sigma \dagger x)$   
 by (simp add: sgn-ueexpr-def subst-uop)

**lemma** usubst-abs [usubst]:  $\sigma \dagger \text{abs } x = \text{abs } (\sigma \dagger x)$   
 by (simp add: abs-ueexpr-def subst-uop)

**lemma** subst-zero [usubst]:  $\sigma \dagger 0 = 0$   
 by (simp add: zero-ueexpr-def subst-lit)

**lemma** subst-one [usubst]:  $\sigma \dagger 1 = 1$   
 by (simp add: one-ueexpr-def subst-lit)

**lemma** subst-eq-upred [usubst]:  $\sigma \dagger (x =_u y) = (\sigma \dagger x =_u \sigma \dagger y)$   
 by (simp add: eq-upred-def usubst)

This laws shows the effect of applying one substitution after another – we simply use function composition to compose them.

**lemma** subst-subst [usubst]:  $\sigma \dagger \varrho \dagger e = (\varrho \circ \sigma) \dagger e$   
 by (transfer, simp)

The next law is similar, but shows how such a substitution is to be applied to every updated variable additionally.

**lemma** subst-upd-comp [usubst]:  
 fixes  $x :: ('a \Rightarrow 'a)$   
 shows  $\varrho(x \mapsto_s v) \circ \sigma = (\varrho \circ \sigma)(x \mapsto_s \sigma \dagger v)$   
 by (rule ext, simp add: ueexpr-defs subst-upd-uvar-def, transfer, simp)

**lemma** subst-singleton:  
 fixes  $x :: ('a \Rightarrow 'a)$   
 assumes  $x \# \sigma$   
 shows  $\sigma(x \mapsto_s v) \dagger P = (\sigma \dagger P)[[v/x]]$   
 using assms  
 by (simp add: usubst)

**lemmas** subst-to-singleton = subst-singleton id-subst

## 6.5 Ordering substitutions

We set up a purely syntactic order on variable lenses which is useful for the substitution normal form.

**definition** *var-name-ord* :: (*'a*  $\Rightarrow$  *'α*)  $\Rightarrow$  (*'b*  $\Rightarrow$  *'α*)  $\Rightarrow$  *bool* **where**  
*[no-atp]*: *var-name-ord* *x y* = *True*

**syntax**

*-var-name-ord* :: *salpha*  $\Rightarrow$  *salpha*  $\Rightarrow$  *bool* (**infix**  $\prec_v$  65)

**translations**

*-var-name-ord* *x y* == *CONST* *var-name-ord* *x y*

A fact of the form  $x \prec_v y$  has no logical information; it simply exists to define a total order on named lenses that is useful for normalisation. The following theorem is simply an instance of the commutativity law for substitutions. However, that law could not be a simplification law as it would cause the simplifier to loop. Assuming that the variable order is a total order then this theorem will not loop.

**lemma** *usubst-upd-comm-ord* [*usubst*]:

**assumes**  $x \bowtie y \ y \prec_v x$

**shows**  $\sigma(x \mapsto_s u, y \mapsto_s v) = \sigma(y \mapsto_s v, x \mapsto_s u)$

**by** (*simp* *add*: *assms*(1) *usubst-upd-comm*)

## 6.6 Unrestriction laws

These are the key unrestricted theorems for substitutions and expressions involving substitutions.

**lemma** *unrest-usubst-single* [*unrest*]:

$\llbracket \text{mwb-lens } x; x \# v \rrbracket \Longrightarrow x \# P[v/x]$

**by** (*transfer*, *auto* *simp* *add*: *subst-upd-uvar-def* *unrest-ueexpr-def*)

**lemma** *unrest-usubst-id* [*unrest*]:

*mwb-lens* *x*  $\Longrightarrow x \# id$

**by** (*simp* *add*: *unrest-usubst-def*)

**lemma** *unrest-usubst-upd* [*unrest*]:

$\llbracket x \bowtie y; x \# \sigma; x \# v \rrbracket \Longrightarrow x \# \sigma(y \mapsto_s v)$

**by** (*simp* *add*: *subst-upd-uvar-def* *unrest-usubst-def* *unrest-ueexpr.rep-eq* *lens-indep-comm*)

**lemma** *unrest-subst* [*unrest*]:

$\llbracket x \# P; x \# \sigma \rrbracket \Longrightarrow x \# (\sigma \uparrow P)$

**by** (*transfer*, *simp* *add*: *unrest-usubst-def*)

## 6.7 Parallel Substitution Laws

**lemma** *par-subst-id* [*usubst*]:

$\llbracket \text{vwb-lens } A; \text{vwb-lens } B \rrbracket \Longrightarrow id [A|B]_s id = id$

**by** (*simp* *add*: *par-subst-def* *lens-override-idem* *id-def*)

**lemma** *par-subst-left-empty* [*usubst*]:

$\llbracket \text{vwb-lens } A \rrbracket \Longrightarrow \sigma [\emptyset|A]_s \varrho = id [\emptyset|A]_s \varrho$

**by** (*simp* *add*: *par-subst-def* *pr-var-def*)

**lemma** *par-subst-right-empty* [*usubst*]:

$\llbracket \text{vwb-lens } A \rrbracket \Longrightarrow \sigma [A|\emptyset]_s \varrho = \sigma [A|\emptyset]_s id$

**by** (*simp* *add*: *par-subst-def* *pr-var-def*)

**lemma** *par-subst-comm*:

$\llbracket A \bowtie B \rrbracket \implies \sigma [A|B]_s \varrho = \varrho [B|A]_s \sigma$   
**by** (*simp add: par-subst-def lens-override-def lens-indep-comm*)

**lemma** *par-subst-upd-left-in* [*usubst*]:  
 $\llbracket \text{vwb-lens } A; A \bowtie B; x \subseteq_L A \rrbracket \implies \sigma(x \mapsto_s v) [A|B]_s \varrho = (\sigma [A|B]_s \varrho)(x \mapsto_s v)$   
**by** (*simp add: par-subst-def subst-upd-uvar-def lens-override-put-right-in*)  
*(simp add: lens-indep-comm lens-override-def sublens-pres-indep)*

**lemma** *par-subst-upd-left-out* [*usubst*]:  
 $\llbracket \text{vwb-lens } A; x \bowtie A \rrbracket \implies \sigma(x \mapsto_s v) [A|B]_s \varrho = (\sigma [A|B]_s \varrho)$   
**by** (*simp add: par-subst-def subst-upd-uvar-def lens-override-put-right-out*)

**lemma** *par-subst-upd-right-in* [*usubst*]:  
 $\llbracket \text{vwb-lens } B; A \bowtie B; x \subseteq_L B \rrbracket \implies \sigma [A|B]_s \varrho(x \mapsto_s v) = (\sigma [A|B]_s \varrho)(x \mapsto_s v)$   
**using** *lens-indep-sym par-subst-comm par-subst-upd-left-in* **by** *fastforce*

**lemma** *par-subst-upd-right-out* [*usubst*]:  
 $\llbracket \text{vwb-lens } B; A \bowtie B; x \bowtie B \rrbracket \implies \sigma [A|B]_s \varrho(x \mapsto_s v) = (\sigma [A|B]_s \varrho)$   
**by** (*simp add: par-subst-comm par-subst-upd-left-out*)

**end**

## 7 UTP Tactics

```
theory utp-tactics
imports Eisbach Lenses Interp utp-expr utp-unrest utp-usedby
keywords update-ueexpr-rep-eq-thms :: thy-decl
begin
```

In this theory, we define several automatic proof tactics that use transfer techniques to re-interpret proof goals about UTP predicates and relations in terms of pure HOL conjectures. The fundamental tactics to achieve this are *pred-simp* and *rel-simp*; a more detailed explanation of their behaviour is given below. The tactics can be given optional arguments to fine-tune their behaviour. By default, they use a weaker but faster form of transfer using rewriting; the option *robust*, however, forces them to use the slower but more powerful transfer of Isabelle’s lifting package. A second option *no-interp* suppresses the re-interpretation of state spaces in order to eradicate record for tuple types prior to automatic proof.

In addition to *pred-simp* and *rel-simp*, we also provide the tactics *pred-auto* and *rel-auto*, as well as *pred-blast* and *rel-blast*; they, in essence, sequence the simplification tactics with the methods *auto* and *blast*, respectively.

### 7.1 Theorem Attributes

The following named attributes have to be introduced already here since our tactics must be able to see them. Note that we do not want to import the theories *utp-pred* and *utp-rel* here, so that both can potentially already make use of the tactics we define in this theory.

```
named-theorems upred-defs upred definitional theorems
named-theorems urel-defs urel definitional theorems
```

### 7.2 Generic Methods

We set up several automatic tactics that recast theorems on UTP predicates into equivalent HOL predicates, eliminating artefacts of the mechanisation as much as this is possible. Our approach is first to unfold all relevant definition of the UTP predicate model, then perform a transfer, and finally simplify by using lens and variable definitions, the split laws of alphabet records, and interpretation laws to convert record-based state spaces into products. The definition of the respective methods is facilitated by the Eisbach tool: we define generic methods that are parametrised by the tactics used for transfer, interpretation and subsequent automatic proof. Note that the tactics only apply to the head goal.

#### Generic Predicate Tactics

```
method gen-pred-tac methods transfer-tac interp-tac prove-tac = (
  ((unfold upred-defs) [1])?;
  (transfer-tac),
  (simp add: fun-eq-iff
    lens-defs upred-defs alpha-splits Product-Type.split-beta)?,
  (interp-tac)?);
(prove-tac)
```

#### Generic Relational Tactics

```
method gen-rel-tac methods transfer-tac interp-tac prove-tac = (
  ((unfold upred-defs urel-defs) [1])?;
  (transfer-tac),
```

```
(simp add: fun-eq-iff relcomp-unfold OO-def
  lens-defs upred-defs alpha-splits Product-Type.split-beta)?,
(interp-tac)?);
(prove-tac)
```

### 7.3 Transfer Tactics

Next, we define the component tactics used for transfer.

#### 7.3.1 Robust Transfer

Robust transfer uses the transfer method of the lifting package.

```
method slow-uexpr-transfer = (transfer)
```

#### 7.3.2 Faster Transfer

Fast transfer side-steps the use of the (*transfer*) method in favour of plain rewriting with the underlying *rep-eq*... laws of lifted definitions. For moderately complex terms, surprisingly, the transfer step turned out to be a bottle-neck in some proofs; we observed that faster transfer resulted in a speed-up of approximately 30% when building the UTP theory heaps. On the downside, tactics using faster transfer do not always work but merely in about 95% of the cases. The approach typically works well when proving predicate equalities and refinements conjectures.

A known limitation is that the faster tactic, unlike lifting transfer, does not turn free variables into meta-quantified ones. This can, in some cases, interfere with the interpretation step and cause subsequent application of automatic proof tactics to fail. A fix is in progress [TODO].

**Attribute Setup** We first configure a dynamic attribute *uexpr-rep-eq-thms* to automatically collect all *rep-eq*- laws of lifted definitions on the *uexpr* type.

**ML-file** *uexpr-rep-eq.ML*

```
setup ⟨⟨
  Global-Theory.add-thms-dynamic (@{binding uexpr-rep-eq-thms},
    uexpr-rep-eq.get-uexpr-rep-eq-thms o Context.theory-of)
  ⟩⟩
```

We next configure a command **update-uexpr-rep-eq-thms** in order to update the content of the *uexpr-rep-eq-thms* attribute. Although the relevant theorems are collected automatically, for efficiency reasons, the user has to manually trigger the update process. The command must hence be executed whenever new lifted definitions for type *uexpr* are created. The updating mechanism uses **find-theorems** under the hood.

```
ML ⟨⟨
  Outer-Syntax.command @{command-keyword update-uexpr-rep-eq-thms}
    reread and update content of the uexpr-rep-eq-thms attribute
    (Scan.succeed (Toplevel.theory uexpr-rep-eq.read-uexpr-rep-eq-thms));
  ⟩⟩
```

**update-uexpr-rep-eq-thms** — Read *uexpr-rep-eq-thms* here.

Lastly, we require several named-theorem attributes to record the manual transfer laws and extra simplifications, so that the user can dynamically extend them in child theories.

**named-theorems** *uexpr-transfer-laws uexpr transfer laws*

**declare** *uexpr-eq-iff* [*uexpr-transfer-laws*]

**named-theorems** *uexpr-transfer-extra extra simplifications for uexpr transfer*

**declare** *unrest-uexpr.rep-eq* [*uexpr-transfer-extra*]

*usedBy-uexpr.rep-eq* [*uexpr-transfer-extra*]

*utp-expr.numeral-uexpr.rep-eq* [*uexpr-transfer-extra*]

*utp-expr.less-eq-uexpr.rep-eq* [*uexpr-transfer-extra*]

*Abs-uexpr-inverse* [*simplified, uexpr-transfer-extra*]

*Rep-uexpr-inverse* [*uexpr-transfer-extra*]

**Tactic Definition** We have all ingredients now to define the fast transfer tactic as a single simplification step.

**method** *fast-uexpr-transfer* =

(*simp add: uexpr-transfer-laws uexpr-rep-eq-thms uexpr-transfer-extra*)

## 7.4 Interpretation

The interpretation of record state spaces as products is done using the laws provided by the utility theory *Interp*. Note that this step can be suppressed by using the *no-interp* option.

**method** *uexpr-interp-tac* = (*simp add: lens-interp-laws*)?

## 7.5 User Tactics

In this section, we finally set-up the six user tactics: *pred-simp*, *rel-simp*, *pred-auto*, *rel-auto*, *pred-blast* and *rel-blast*. For this, we first define the proof strategies that are to be applied *after* the transfer steps.

**method** *utp-simp-tac* = (*clarsimp*)?

**method** *utp-auto-tac* = ((*clarsimp*)?; *auto*)

**method** *utp-blast-tac* = ((*clarsimp*)?; *blast*)

The ML file below provides ML constructor functions for tactics that process arguments suitable and invoke the generic methods *gen-pred-tac* and *gen-rel-tac* with suitable arguments.

**ML-file** *utp-tactics.ML*

Finally, we execute the relevant outer commands for method setup. Sadly, this cannot be done at the level of Eisbach since the latter does not provide a convenient mechanism to process symbolic flags as arguments. It may be worth to put in a feature request with the developers of the Eisbach tool.

```
method-setup pred-simp = <<
  (Scan.lift UTP-Tactics.scan-args) >>
  (fn args => fn ctx =>
    let val prove-tac = Basic-Tactics.utp-simp-tac in
      (UTP-Tactics.inst-gen-pred-tac args prove-tac ctx)
    end);
>>
```

```
method-setup rel-simp = <<
  (Scan.lift UTP-Tactics.scan-args) >>
  (fn args => fn ctx =>
    let val prove-tac = Basic-Tactics.utp-simp-tac in
```



```

      (UTP-Tactics.inst-gen-rel-tac args prove-tac ctx)
    end);
  >>

method-setup pred-auto = <<
  (Scan.lift UTP-Tactics.scan-args) >>
  (fn args => fn ctx =>
    let val prove-tac = Basic-Tactics.utp-auto-tac in
    (UTP-Tactics.inst-gen-pred-tac args prove-tac ctx)
  end);
  >>

method-setup rel-auto = <<
  (Scan.lift UTP-Tactics.scan-args) >>
  (fn args => fn ctx =>
    let val prove-tac = Basic-Tactics.utp-auto-tac in
    (UTP-Tactics.inst-gen-rel-tac args prove-tac ctx)
  end);
  >>

method-setup pred-blast = <<
  (Scan.lift UTP-Tactics.scan-args) >>
  (fn args => fn ctx =>
    let val prove-tac = Basic-Tactics.utp-blast-tac in
    (UTP-Tactics.inst-gen-pred-tac args prove-tac ctx)
  end);
  >>

method-setup rel-blast = <<
  (Scan.lift UTP-Tactics.scan-args) >>
  (fn args => fn ctx =>
    let val prove-tac = Basic-Tactics.utp-blast-tac in
    (UTP-Tactics.inst-gen-rel-tac args prove-tac ctx)
  end);
  >>

```

end

## 8 Alphabetised Predicates

```
theory utp-pred
imports
  utp-expr
  utp-subst
  utp-tactics
begin
```

In this theory we begin to create an Isabelle version of the alphabetised predicate calculus that is described in Chapter 1 of the UTP book [5].

### 8.1 Predicate type and syntax

An alphabetised predicate is simply a boolean valued expression.

```
type-synonym 'α upred = (bool, 'α) uexpr
```

#### translations

```
(type) 'α upred <= (type) (bool, 'α) uexpr
```

We want to remain as close as possible to the mathematical UTP syntax, but also want to be conservative with HOL. For this reason we chose not to steal syntax from HOL, but where possible use polymorphism to allow selection of the appropriate operator (UTP vs. HOL). Thus we will first remove the standard syntax for conjunction, disjunction, and negation, and replace these with adhoc overloaded definitions. We similarly use polymorphic constants for the other predicate calculus operators.

#### purge-notation

```
conj (infixr  $\wedge$  35) and
disj (infixr  $\vee$  30) and
Not ( $\neg$  - [40] 40)
```

#### consts

```
uttrue :: 'a (true)
ufalse :: 'a (false)
uconj :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr  $\wedge$  35)
udisj :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr  $\vee$  30)
wimpl :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr  $\Rightarrow$  25)
wiff :: 'a  $\Rightarrow$  'a  $\Rightarrow$  'a (infixr  $\Leftrightarrow$  25)
unot :: 'a  $\Rightarrow$  'a ( $\neg$  - [40] 40)
uex :: ('a  $\Rightarrow$  'α)  $\Rightarrow$  'p  $\Rightarrow$  'p
uall :: ('a  $\Rightarrow$  'α)  $\Rightarrow$  'p  $\Rightarrow$  'p
ushEx :: ['a  $\Rightarrow$  'p]  $\Rightarrow$  'p
ushAll :: ['a  $\Rightarrow$  'p]  $\Rightarrow$  'p
```

#### adhoc-overloading

```
uconj conj and
udisj disj and
unot Not
```

We set up two versions of each of the quantifiers: *uex* / *uall* and *ushEx* / *ushAll*. The former pair allows quantification of UTP variables, whilst the latter allows quantification of HOL variables

in concert with the literal expression constructor  $\llbracket x \rrbracket$ . Both varieties will be needed at various points. Syntactically they are distinguished by a boldface quantifier for the HOL versions (achieved by the "bold" escape in Isabelle).

**nonterminal** *idt-list*

**syntax**

```
-idt-el  :: idt ⇒ idt-list (-)
-idt-list :: idt ⇒ idt-list ⇒ idt-list ((-, / -) [0, 1])
-uex     :: salpha ⇒ logic ⇒ logic (∃ - · - [0, 10] 10)
-uall    :: salpha ⇒ logic ⇒ logic (∀ - · - [0, 10] 10)
-ushEx   :: pttrn ⇒ logic ⇒ logic (∃ - · - [0, 10] 10)
-ushAll  :: pttrn ⇒ logic ⇒ logic (∀ - · - [0, 10] 10)
-ushBEx  :: pttrn ⇒ logic ⇒ logic ⇒ logic (∃ - ∈ - · - [0, 0, 10] 10)
-ushBAll :: pttrn ⇒ logic ⇒ logic ⇒ logic (∀ - ∈ - · - [0, 0, 10] 10)
-ushGAll :: pttrn ⇒ logic ⇒ logic ⇒ logic (∀ - | - · - [0, 0, 10] 10)
-ushGtAll :: idt ⇒ logic ⇒ logic ⇒ logic (∀ - > - · - [0, 0, 10] 10)
-ushLtAll :: idt ⇒ logic ⇒ logic ⇒ logic (∀ - < - · - [0, 0, 10] 10)
-uvar-res :: logic ⇒ salpha ⇒ logic (infixl  $\downarrow_v$  90)
```

**translations**

```
-uex x P          == CONST uex x P
-uex (-salphaset (-salphamk (x +L y))) P <= -uex (x +L y) P
-uall x P          == CONST uall x P
-uall (-salphaset (-salphamk (x +L y))) P <= -uall (x +L y) P
-ushEx x P         == CONST ushEx (λ x. P)
∃ x ∈ A · P        => ∃ x ·  $\llbracket x \rrbracket \in_u A \wedge P$ 
-ushAll x P        == CONST ushAll (λ x. P)
∀ x ∈ A · P        => ∀ x ·  $\llbracket x \rrbracket \in_u A \Rightarrow P$ 
∀ x | P · Q        => ∀ x · P ⇒ Q
∀ x > y · P        => ∀ x ·  $\llbracket x \rrbracket >_u y \Rightarrow P$ 
∀ x < y · P        => ∀ x ·  $\llbracket x \rrbracket <_u y \Rightarrow P$ 
```

## 8.2 Predicate operators

We chose to maximally reuse definitions and laws built into HOL. For this reason, when introducing the core operators we proceed by lifting operators from the polymorphic algebraic hierarchy of HOL. Thus the initial definitions take place in the context of type class instantiations. We first introduce our own class called *refine* that will add the refinement operator syntax to the HOL partial order class.

**class** *refine* = *order*

**abbreviation** *refineBy* :: 'a::*refine* ⇒ 'a ⇒ bool (**infix**  $\sqsubseteq$  50) **where**  
*P*  $\sqsubseteq$  *Q*  $\equiv$  *less-eq* *Q* *P*

Since, on the whole, lattices in UTP are the opposite way up to the standard definitions in HOL, we syntactically invert the lattice operators. This is the one exception where we do steal HOL syntax, but I think it makes sense for UTP. Indeed we make this inversion for all of the lattice operators.

**purge-notation** *Lattices.inf* (**infixl**  $\sqcap$  70)  
**notation** *Lattices.inf* (**infixl**  $\sqcap$  70)  
**purge-notation** *Lattices.sup* (**infixl**  $\sqcup$  65)  
**notation** *Lattices.sup* (**infixl**  $\sqcup$  65)

**purge-notation** *Inf* ( $\sqcap$  - [900] 900)  
**notation** *Inf* ( $\sqcup$  - [900] 900)  
**purge-notation** *Sup* ( $\sqcup$  - [900] 900)  
**notation** *Sup* ( $\sqcap$  - [900] 900)

**purge-notation** *Orderings.bot* ( $\perp$ )  
**notation** *Orderings.bot* ( $\top$ )  
**purge-notation** *Orderings.top* ( $\top$ )  
**notation** *Orderings.top* ( $\perp$ )

**purge-syntax**

-INF1    :: *pttrns*  $\Rightarrow$  'b  $\Rightarrow$  'b    (( $\exists \sqcap$  -./ -) [0, 10] 10)  
 -INF    :: *pttrn*  $\Rightarrow$  'a set  $\Rightarrow$  'b  $\Rightarrow$  'b (( $\exists \sqcap$  - $\in$  -./ -) [0, 0, 10] 10)  
 -SUP1    :: *pttrns*  $\Rightarrow$  'b  $\Rightarrow$  'b    (( $\exists \sqcup$  -./ -) [0, 10] 10)  
 -SUP    :: *pttrn*  $\Rightarrow$  'a set  $\Rightarrow$  'b  $\Rightarrow$  'b (( $\exists \sqcup$  - $\in$  -./ -) [0, 0, 10] 10)

**syntax**

-INF1    :: *pttrns*  $\Rightarrow$  'b  $\Rightarrow$  'b    (( $\exists \sqcup$  -./ -) [0, 10] 10)  
 -INF    :: *pttrn*  $\Rightarrow$  'a set  $\Rightarrow$  'b  $\Rightarrow$  'b (( $\exists \sqcup$  - $\in$  -./ -) [0, 0, 10] 10)  
 -SUP1    :: *pttrns*  $\Rightarrow$  'b  $\Rightarrow$  'b    (( $\exists \sqcap$  -./ -) [0, 10] 10)  
 -SUP    :: *pttrn*  $\Rightarrow$  'a set  $\Rightarrow$  'b  $\Rightarrow$  'b (( $\exists \sqcap$  - $\in$  -./ -) [0, 0, 10] 10)

We trivially instantiate our refinement class

**instance** *uexpr* :: (*order*, *type*) *refine* ..

— Configure transfer law for refinement for the fast relational tactics.

**theorem** *upred-ref-iff* [*uexpr-transfer-laws*]:

$(P \sqsubseteq Q) = (\forall b. \llbracket Q \rrbracket_e b \longrightarrow \llbracket P \rrbracket_e b)$

**apply** (*transfer*)

**apply** (*clarsimp*)

**done**

Next we introduce the lattice operators, which is again done by lifting.

**instantiation** *uexpr* :: (*lattice*, *type*) *lattice*

**begin**

**lift-definition** *sup-uexpr* :: ('a, 'b) *uexpr*  $\Rightarrow$  ('a, 'b) *uexpr*  $\Rightarrow$  ('a, 'b) *uexpr*  
**is**  $\lambda P Q A. \text{Lattices.sup } (P A) (Q A)$  .

**lift-definition** *inf-uexpr* :: ('a, 'b) *uexpr*  $\Rightarrow$  ('a, 'b) *uexpr*  $\Rightarrow$  ('a, 'b) *uexpr*  
**is**  $\lambda P Q A. \text{Lattices.inf } (P A) (Q A)$  .

**instance**

**by** (*intro-classes*) (*transfer*, *auto*)+

**end**

**instantiation** *uexpr* :: (*bounded-lattice*, *type*) *bounded-lattice*

**begin**

**lift-definition** *bot-uexpr* :: ('a, 'b) *uexpr* **is**  $\lambda A. \text{Orderings.bot}$  .

**lift-definition** *top-uexpr* :: ('a, 'b) *uexpr* **is**  $\lambda A. \text{Orderings.top}$  .

**instance**

**by** (*intro-classes*) (*transfer*, *auto*)+

**end**

**lemma** *top-uexpr-rep-eq* [*simp*]:

$\llbracket \text{Orderings.bot} \rrbracket_e b = \text{False}$

**by** (*transfer*, *auto*)

```

lemma bot-uepr-rep-eq [simp]:
   $\llbracket \text{Orderings.top} \rrbracket_e b = \text{True}$ 
  by (transfer, auto)

```

```

instance uepr :: (distrib-lattice, type) distrib-lattice
  by (intro-classes) (transfer, rule ext, auto simp add: sup-inf-distrib1)

```

Finally we show that predicates form a Boolean algebra (under the lattice operators), a complete lattice, a completely distribute lattice, and a complete boolean algebra. This equip us with a very complete theory for basic logical propositions.

```

instance uepr :: (boolean-algebra, type) boolean-algebra
apply (intro-classes, unfold uepr-defs; transfer, rule ext)
apply (simp-all add: sup-inf-distrib1 diff-eq)
done

```

```

instantiation uepr :: (complete-lattice, type) complete-lattice
begin
  lift-definition Inf-uepr :: ('a, 'b) uepr set  $\Rightarrow$  ('a, 'b) uepr
  is  $\lambda PS A. \text{INF } P:PS. P(A)$  .
  lift-definition Sup-uepr :: ('a, 'b) uepr set  $\Rightarrow$  ('a, 'b) uepr
  is  $\lambda PS A. \text{SUP } P:PS. P(A)$  .
instance
  by (intro-classes)
  (transfer, auto intro: INF-lower SUP-upper simp add: INF-greatest SUP-least)+
end

```

```

instance uepr :: (complete-distrib-lattice, type) complete-distrib-lattice
  apply (intro-classes)
  apply (transfer, rule ext, auto)
  using sup-INF apply fastforce
  apply (transfer, rule ext, auto)
  using inf-SUP apply fastforce
done

```

```

instance uepr :: (complete-boolean-algebra, type) complete-boolean-algebra ..

```

From the complete lattice, we can also define and give syntax for the fixed-point operators. Like the lattice operators, these are reversed in UTP.

```

syntax
  -mu :: pttrn  $\Rightarrow$  logic  $\Rightarrow$  logic ( $\mu$  - - [0, 10] 10)
  -nu :: pttrn  $\Rightarrow$  logic  $\Rightarrow$  logic ( $\nu$  - - [0, 10] 10)

```

```

notation gfp ( $\mu$ )
notation lfp ( $\nu$ )

```

```

translations
   $\nu X \cdot P == \text{CONST lfp } (\lambda X. P)$ 
   $\mu X \cdot P == \text{CONST gfp } (\lambda X. P)$ 

```

With the lattice operators defined, we can proceed to give definitions for the standard predicate operators in terms of them.

```

definition true-upred = (Orderings.top :: 'a upred)
definition false-upred = (Orderings.bot :: 'a upred)

```

**definition** *conj-upred* = (*Lattices.inf* :: 'α upred ⇒ 'α upred ⇒ 'α upred)

**definition** *disj-upred* = (*Lattices.sup* :: 'α upred ⇒ 'α upred ⇒ 'α upred)

**definition** *not-upred* = (*uminus* :: 'α upred ⇒ 'α upred)

**definition** *diff-upred* = (*minus* :: 'α upred ⇒ 'α upred ⇒ 'α upred)

**abbreviation** *Conj-upred* :: 'α upred set ⇒ 'α upred (Λ- [900] 900) **where**  
 $\bigwedge A \equiv \bigcap A$

**abbreviation** *Disj-upred* :: 'α upred set ⇒ 'α upred (∨- [900] 900) **where**  
 $\bigvee A \equiv \bigcup A$

**notation**

*conj-upred* (**infixr**  $\wedge_p$  35) **and**

*disj-upred* (**infixr**  $\vee_p$  30)

Perhaps slightly confusingly, the UTP infimum is the HOL supremum and vice-versa. This is because, again, in UTP the lattice is inverted due to the definition of refinement and a desire to have miracle at the top, and abort at the bottom.

**lift-definition** *UINF* :: ('a ⇒ 'α upred) ⇒ ('a ⇒ ('b::complete-lattice, 'α) uexpr) ⇒ ('b, 'α) uexpr  
**is** λ P F b. *Sup* {⌊F x⌋<sub>e</sub> b | x. ⌊P x⌋<sub>e</sub> b} .

**lift-definition** *USUP* :: ('a ⇒ 'α upred) ⇒ ('a ⇒ ('b::complete-lattice, 'α) uexpr) ⇒ ('b, 'α) uexpr  
**is** λ P F b. *Inf* {⌊F x⌋<sub>e</sub> b | x. ⌊P x⌋<sub>e</sub> b} .

**declare** *UINF-def* [*upred-defs*]

**declare** *USUP-def* [*upred-defs*]

**syntax**

-*USup* :: *pttrn* ⇒ *logic* ⇒ *logic* (Λ - · - [0, 10] 10)  
-*USup* :: *pttrn* ⇒ *logic* ⇒ *logic* (⊓ - · - [0, 10] 10)  
-*USup-mem* :: *pttrn* ⇒ *logic* ⇒ *logic* ⇒ *logic* (Λ - ∈ · - [0, 10] 10)  
-*USup-mem* :: *pttrn* ⇒ *logic* ⇒ *logic* ⇒ *logic* (⊓ - ∈ · - [0, 10] 10)  
-*USUP* :: *pttrn* ⇒ *logic* ⇒ *logic* ⇒ *logic* (Λ - | · - [0, 0, 10] 10)  
-*USUP* :: *pttrn* ⇒ *logic* ⇒ *logic* ⇒ *logic* (⊓ - | · - [0, 0, 10] 10)  
-*UInf* :: *pttrn* ⇒ *logic* ⇒ *logic* (∨ - · - [0, 10] 10)  
-*UInf* :: *pttrn* ⇒ *logic* ⇒ *logic* (⊔ - · - [0, 10] 10)  
-*UInf-mem* :: *pttrn* ⇒ *logic* ⇒ *logic* ⇒ *logic* (∨ - ∈ · - [0, 10] 10)  
-*UInf-mem* :: *pttrn* ⇒ *logic* ⇒ *logic* ⇒ *logic* (⊔ - ∈ · - [0, 10] 10)  
-*UINF* :: *pttrn* ⇒ *logic* ⇒ *logic* ⇒ *logic* (∨ - | · - [0, 10] 10)  
-*UINF* :: *pttrn* ⇒ *logic* ⇒ *logic* ⇒ *logic* (⊔ - | · - [0, 10] 10)

**translations**

⊓ x | P · F => *CONST UINF* (λ x. P) (λ x. F)  
⊓ x · F == ⊓ x | true · F  
⊓ x · F == ⊓ x | true · F  
⊓ x ∈ A · F => ⊓ x | «x» ∈<sub>u</sub> «A» · F  
⊓ x ∈ A · F <= ⊓ x | «y» ∈<sub>u</sub> «A» · F  
⊓ x | P · F <= *CONST UINF* (λ y. P) (λ x. F)  
⊓ x | P · F(x) <= *CONST UINF* (λ x. P) F  
⊓ x | P · F => *CONST USUP* (λ x. P) (λ x. F)  
⊓ x · F == ⊓ x | true · F  
⊓ x ∈ A · F => ⊓ x | «x» ∈<sub>u</sub> «A» · F  
⊓ x ∈ A · F <= ⊓ x | «y» ∈<sub>u</sub> «A» · F  
⊓ x | P · F <= *CONST USUP* (λ y. P) (λ x. F)  
⊓ x | P · F(x) <= *CONST USUP* (λ x. P) F

We also define the other predicate operators

**lift-definition** *impl* ::  $'\alpha \text{ upred} \Rightarrow '\alpha \text{ upred} \Rightarrow '\alpha \text{ upred}$  **is**  
 $\lambda P Q A. P A \longrightarrow Q A$  .

**lift-definition** *iff-upred* ::  $'\alpha \text{ upred} \Rightarrow '\alpha \text{ upred} \Rightarrow '\alpha \text{ upred}$  **is**  
 $\lambda P Q A. P A \longleftrightarrow Q A$  .

**lift-definition** *ex* ::  $('a \Longrightarrow '\alpha) \Rightarrow '\alpha \text{ upred} \Rightarrow '\alpha \text{ upred}$  **is**  
 $\lambda x P b. (\exists v. P(\text{put}_x b v))$  .

**lift-definition** *shEx* ::  $['\beta \Rightarrow '\alpha \text{ upred}] \Rightarrow '\alpha \text{ upred}$  **is**  
 $\lambda P A. \exists x. (P x) A$  .

**lift-definition** *all* ::  $('a \Longrightarrow '\alpha) \Rightarrow '\alpha \text{ upred} \Rightarrow '\alpha \text{ upred}$  **is**  
 $\lambda x P b. (\forall v. P(\text{put}_x b v))$  .

**lift-definition** *shAll* ::  $['\beta \Rightarrow '\alpha \text{ upred}] \Rightarrow '\alpha \text{ upred}$  **is**  
 $\lambda P A. \forall x. (P x) A$  .

We define the following operator which is dual of existential quantification. It hides the valuation of variables other than  $x$  through existential quantification.

**lift-definition** *var-res* ::  $'\alpha \text{ upred} \Rightarrow ('a \Longrightarrow '\alpha) \Rightarrow '\alpha \text{ upred}$  **is**  
 $\lambda P x b. \exists b'. P(b' \oplus_L b \text{ on } x)$  .

#### translations

*-uvar-res*  $P a \equiv \text{CONST } \text{var-res } P a$

We have to add a  $u$  subscript to the closure operator as I don't want to override the syntax for HOL lists (we'll be using them later).

**lift-definition** *closure* ::  $'\alpha \text{ upred} \Rightarrow '\alpha \text{ upred}$  ( $[-]_u$ ) **is**  
 $\lambda P A. \forall A'. P A'$  .

**lift-definition** *taut* ::  $'\alpha \text{ upred} \Rightarrow \text{bool}$  ( $'\text{'}$ )  
**is**  $\lambda P. \forall A. P A$  .

— Configuration for UTP tactics (see *utp-tactics*).

**update-uexpr-rep-eq-thms** — Reread *rep-eq* theorems.

**declare** *utp-pred.taut.rep-eq* [*upred-defs*]

#### ad hoc overloading

*uttrue* *true-upred* **and**  
*ufalse* *false-upred* **and**  
*unot* *not-upred* **and**  
*uconj* *conj-upred* **and**  
*udisj* *disj-upred* **and**  
*uimpl* *impl* **and**  
*uiff* *iff-upred* **and**  
*uex* *ex* **and**  
*uall* *all* **and**  
*ushEx* *shEx* **and**  
*ushAll* *shAll*

#### syntax

$-uneq \quad :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \text{ (infixl } \neq_u 50)$   
 $-unmem \quad :: ('a, 'α) \text{ uexpr} \Rightarrow ('a \text{ set}, 'α) \text{ uexpr} \Rightarrow (\text{bool}, 'α) \text{ uexpr} \text{ (infix } \notin_u 50)$

#### translations

$x \neq_u y == \text{CONST unot } (x =_u y)$   
 $x \notin_u A == \text{CONST unot } (\text{CONST bop } (op \in) x A)$

**declare** *true-upred-def* [*upred-defs*]  
**declare** *false-upred-def* [*upred-defs*]  
**declare** *conj-upred-def* [*upred-defs*]  
**declare** *disj-upred-def* [*upred-defs*]  
**declare** *not-upred-def* [*upred-defs*]  
**declare** *diff-upred-def* [*upred-defs*]  
**declare** *subst-upd-uvar-def* [*upred-defs*]  
**declare** *par-subst-def* [*upred-defs*]  
**declare** *unrest-usubst-def* [*upred-defs*]  
**declare** *uexpr-defs* [*upred-defs*]

**lemma** *true-alt-def*: *true* =  $\ll \text{True} \gg$   
**by** (*pred-auto*)

**lemma** *false-alt-def*: *false* =  $\ll \text{False} \gg$   
**by** (*pred-auto*)

**declare** *true-alt-def* [*THEN sym, lit-simps*]  
**declare** *false-alt-def* [*THEN sym, lit-simps*]

We set up syntax for the conditional. This is effectively an infix version of if-then-else where the condition is in the middle.

**abbreviation** *cond* ::  
 $('a, 'α) \text{ uexpr} \Rightarrow 'α \text{ upred} \Rightarrow ('a, 'α) \text{ uexpr} \Rightarrow ('a, 'α) \text{ uexpr}$   
 $((\beta \triangleleft - \triangleright / -) [52, 0, 53] 52)$   
**where**  $P \triangleleft b \triangleright Q \equiv \text{trop If } b P Q$

### 8.3 Unrestriction Laws

**lemma** *unrest-allE*:  
 $\ll \Sigma \# P; P = \text{true} \Longrightarrow Q; P = \text{false} \Longrightarrow Q \gg \Longrightarrow Q$   
**by** (*pred-auto*)

**lemma** *unrest-true* [*unrest*]:  $x \# \text{true}$   
**by** (*pred-auto*)

**lemma** *unrest-false* [*unrest*]:  $x \# \text{false}$   
**by** (*pred-auto*)

**lemma** *unrest-conj* [*unrest*]:  $\ll x \# (P :: 'α \text{ upred}); x \# Q \gg \Longrightarrow x \# P \wedge Q$   
**by** (*pred-auto*)

**lemma** *unrest-disj* [*unrest*]:  $\ll x \# (P :: 'α \text{ upred}); x \# Q \gg \Longrightarrow x \# P \vee Q$   
**by** (*pred-auto*)

**lemma** *unrest-UNIF* [*unrest*]:  
 $\ll (\bigwedge i. x \# P(i)); (\bigwedge i. x \# Q(i)) \gg \Longrightarrow x \# (\bigcap i \mid P(i) \cdot Q(i))$   
**by** (*pred-auto*)



**lemma** *unrest-USUP* [*unrest*]:  
 $\llbracket (\bigwedge i. x \# P(i)); (\bigwedge i. x \# Q(i)) \rrbracket \Longrightarrow x \# (\bigsqcup i \mid P(i) \cdot Q(i))$   
**by** (*pred-auto*)

**lemma** *unrest-UINF-mem* [*unrest*]:  
 $\llbracket (\bigwedge i. i \in A \Longrightarrow x \# P(i)) \rrbracket \Longrightarrow x \# (\prod_{i \in A} P(i))$   
**by** (*pred-simp*, *metis*)

**lemma** *unrest-USUP-mem* [*unrest*]:  
 $\llbracket (\bigwedge i. i \in A \Longrightarrow x \# P(i)) \rrbracket \Longrightarrow x \# (\bigsqcup_{i \in A} P(i))$   
**by** (*pred-simp*, *metis*)

**lemma** *unrest-impl* [*unrest*]:  $\llbracket x \# P; x \# Q \rrbracket \Longrightarrow x \# P \Rightarrow Q$   
**by** (*pred-auto*)

**lemma** *unrest-iff* [*unrest*]:  $\llbracket x \# P; x \# Q \rrbracket \Longrightarrow x \# P \Leftrightarrow Q$   
**by** (*pred-auto*)

**lemma** *unrest-not* [*unrest*]:  $x \# (P :: 'a \text{ upred}) \Longrightarrow x \# (\neg P)$   
**by** (*pred-auto*)

The sublens proviso can be thought of as membership below.

**lemma** *unrest-ex-in* [*unrest*]:  
 $\llbracket \text{mwb-lens } y; x \subseteq_L y \rrbracket \Longrightarrow x \# (\exists y \cdot P)$   
**by** (*pred-auto*)

**declare** *sublens-refl* [*simp*]  
**declare** *lens-plus-ub* [*simp*]  
**declare** *lens-plus-right-sublens* [*simp*]  
**declare** *comp-wb-lens* [*simp*]  
**declare** *comp-mwb-lens* [*simp*]  
**declare** *plus-mwb-lens* [*simp*]

**lemma** *unrest-ex-diff* [*unrest*]:  
**assumes**  $x \bowtie y \ y \# P$   
**shows**  $y \# (\exists x \cdot P)$   
**using** *assms*  
**apply** (*pred-auto*)  
**using** *lens-indep-comm* **apply** *fastforce* +  
**done**

**lemma** *unrest-all-in* [*unrest*]:  
 $\llbracket \text{mwb-lens } y; x \subseteq_L y \rrbracket \Longrightarrow x \# (\forall y \cdot P)$   
**by** (*pred-auto*)

**lemma** *unrest-all-diff* [*unrest*]:  
**assumes**  $x \bowtie y \ y \# P$   
**shows**  $y \# (\forall x \cdot P)$   
**using** *assms*  
**by** (*pred-simp*, *simp-all* *add: lens-indep-comm*)

**lemma** *unrest-var-res-diff* [*unrest*]:  
**assumes**  $x \bowtie y$   
**shows**  $y \# (P \upharpoonright_v x)$

**using** *assms* **by** (*pred-auto*)

**lemma** *unrest-var-res-in* [*unrest*]:  
**assumes** *mwb-lens*  $x\ y \subseteq_L x\ y \# P$   
**shows**  $y \# (P \upharpoonright_v x)$   
**using** *assms*  
**apply** (*pred-auto*)  
**apply** *fastforce*  
**apply** (*metis* (*no-types*, *lifting*) *mwb-lens-weak weak-lens.put-get*)  
**done**

**lemma** *unrest-shEx* [*unrest*]:  
**assumes**  $\bigwedge y. x \# P(y)$   
**shows**  $x \# (\exists y. P(y))$   
**using** *assms* **by** (*pred-auto*)

**lemma** *unrest-shAll* [*unrest*]:  
**assumes**  $\bigwedge y. x \# P(y)$   
**shows**  $x \# (\forall y. P(y))$   
**using** *assms* **by** (*pred-auto*)

**lemma** *unrest-closure* [*unrest*]:  
 $x \# [P]_u$   
**by** (*pred-auto*)

## 8.4 Used-by laws

**lemma** *usedBy-not* [*unrest*]:  
 $\llbracket x \# P \rrbracket \implies x \# (\neg P)$   
**by** (*pred-simp*)

**lemma** *usedBy-conj* [*unrest*]:  
 $\llbracket x \# P; x \# Q \rrbracket \implies x \# (P \wedge Q)$   
**by** (*pred-simp*)

**lemma** *usedBy-disj* [*unrest*]:  
 $\llbracket x \# P; x \# Q \rrbracket \implies x \# (P \vee Q)$   
**by** (*pred-simp*)

**lemma** *usedBy-impl* [*unrest*]:  
 $\llbracket x \# P; x \# Q \rrbracket \implies x \# (P \Rightarrow Q)$   
**by** (*pred-simp*)

**lemma** *usedBy-iff* [*unrest*]:  
 $\llbracket x \# P; x \# Q \rrbracket \implies x \# (P \Leftrightarrow Q)$   
**by** (*pred-simp*)

## 8.5 Substitution Laws

Substitution is monotone

**lemma** *subst-mono*:  $P \sqsubseteq Q \implies (\sigma \dagger P) \sqsubseteq (\sigma \dagger Q)$   
**by** (*pred-auto*)

**lemma** *subst-true* [*usubst*]:  $\sigma \dagger \text{true} = \text{true}$   
**by** (*pred-auto*)

**lemma** *subst-false* [*usubst*]:  $\sigma \dagger \text{false} = \text{false}$   
**by** (*pred-auto*)

**lemma** *subst-not* [*usubst*]:  $\sigma \dagger (\neg P) = (\neg \sigma \dagger P)$   
**by** (*pred-auto*)

**lemma** *subst-impl* [*usubst*]:  $\sigma \dagger (P \Rightarrow Q) = (\sigma \dagger P \Rightarrow \sigma \dagger Q)$   
**by** (*pred-auto*)

**lemma** *subst-iff* [*usubst*]:  $\sigma \dagger (P \Leftrightarrow Q) = (\sigma \dagger P \Leftrightarrow \sigma \dagger Q)$   
**by** (*pred-auto*)

**lemma** *subst-disj* [*usubst*]:  $\sigma \dagger (P \vee Q) = (\sigma \dagger P \vee \sigma \dagger Q)$   
**by** (*pred-auto*)

**lemma** *subst-conj* [*usubst*]:  $\sigma \dagger (P \wedge Q) = (\sigma \dagger P \wedge \sigma \dagger Q)$   
**by** (*pred-auto*)

**lemma** *subst-sup* [*usubst*]:  $\sigma \dagger (P \sqcap Q) = (\sigma \dagger P \sqcap \sigma \dagger Q)$   
**by** (*pred-auto*)

**lemma** *subst-inf* [*usubst*]:  $\sigma \dagger (P \sqcup Q) = (\sigma \dagger P \sqcup \sigma \dagger Q)$   
**by** (*pred-auto*)

**lemma** *subst-UNIF* [*usubst*]:  $\sigma \dagger (\prod i \mid P(i) \cdot Q(i)) = (\prod i \mid (\sigma \dagger P(i)) \cdot (\sigma \dagger Q(i)))$   
**by** (*pred-auto*)

**lemma** *subst-USUP* [*usubst*]:  $\sigma \dagger (\bigsqcup i \mid P(i) \cdot Q(i)) = (\bigsqcup i \mid (\sigma \dagger P(i)) \cdot (\sigma \dagger Q(i)))$   
**by** (*pred-auto*)

**lemma** *subst-closure* [*usubst*]:  $\sigma \dagger [P]_u = [P]_u$   
**by** (*pred-auto*)

**lemma** *subst-shEx* [*usubst*]:  $\sigma \dagger (\exists x \cdot P(x)) = (\exists x \cdot \sigma \dagger P(x))$   
**by** (*pred-auto*)

**lemma** *subst-shAll* [*usubst*]:  $\sigma \dagger (\forall x \cdot P(x)) = (\forall x \cdot \sigma \dagger P(x))$   
**by** (*pred-auto*)

TODO: Generalise the quantifier substitution laws to n-ary substitutions

**lemma** *subst-ex-same* [*usubst*]:  
*mbw-lens*  $x \Rightarrow \sigma(x \mapsto_s v) \dagger (\exists x \cdot P) = \sigma \dagger (\exists x \cdot P)$   
**by** (*pred-auto*)

**lemma** *subst-ex-same'* [*usubst*]:  
*mbw-lens*  $x \Rightarrow \sigma(x \mapsto_s v) \dagger (\exists \&x \cdot P) = \sigma \dagger (\exists \&x \cdot P)$   
**by** (*pred-auto*)

**lemma** *subst-ex-indep* [*usubst*]:  
**assumes**  $x \bowtie y \ y \nparallel v$   
**shows**  $(\exists y \cdot P) \llbracket v/x \rrbracket = (\exists y \cdot P \llbracket v/x \rrbracket)$   
**using** *assms*  
**apply** (*pred-auto*)  
**using** *lens-indep-comm* **apply** *fastforce*+

done

**lemma** *subst-ex-unrest* [*usubst*]:  
 $x \# \sigma \implies \sigma \uparrow (\exists x \cdot P) = (\exists x \cdot \sigma \uparrow P)$   
 by (*pred-auto*)

**lemma** *subst-all-same* [*usubst*]:  
 $mwb\text{-}lens\ x \implies \sigma(x \mapsto_s v) \uparrow (\forall x \cdot P) = \sigma \uparrow (\forall x \cdot P)$   
 by (*simp add: id-subst subst-unrest unrest-all-in*)

**lemma** *subst-all-indep* [*usubst*]:  
 assumes  $x \bowtie y \ y \# v$   
 shows  $(\forall y \cdot P) \llbracket v/x \rrbracket = (\forall y \cdot P \llbracket v/x \rrbracket)$   
 using *assms*  
 by (*pred-simp, simp-all add: lens-indep-comm*)

end

## 9 Predicate Calculus Laws

**theory** *utp-pred-laws*  
 imports *utp-pred*  
 begin

### 9.1 Propositional Logic

Showing that predicates form a Boolean Algebra (under the predicate operators as opposed to the lattice operators) gives us many useful laws.

**interpretation** *boolean-algebra diff-upred not-upred conj-upred op ≤ op <*  
*disj-upred false-upred true-upred*  
 by (*unfold-locales; pred-auto*)

**lemma** *taut-true* [*simp*]: ‘true’  
 by (*pred-auto*)

**lemma** *taut-false* [*simp*]: ‘false’ = *False*  
 by (*pred-auto*)

**lemma** *upred-eval-taut*:  
 $\llbracket P \llbracket \llbracket b \rrbracket / \&\mathbf{v} \rrbracket \rrbracket = \llbracket P \rrbracket_e b$   
 by (*pred-auto*)

**lemma** *refBy-order*:  $P \sqsubseteq Q = \llbracket Q \Rightarrow P \rrbracket$   
 by (*pred-auto*)

**lemma** *conj-idem* [*simp*]:  $((P::'\alpha\ \text{upred}) \wedge P) = P$   
 by (*pred-auto*)

**lemma** *disj-idem* [*simp*]:  $((P::'\alpha\ \text{upred}) \vee P) = P$   
 by (*pred-auto*)

**lemma** *conj-comm*:  $((P::'\alpha\ \text{upred}) \wedge Q) = (Q \wedge P)$   
 by (*pred-auto*)

**lemma** *disj-comm*:  $((P::'\alpha \text{ upred}) \vee Q) = (Q \vee P)$   
**by** (*pred-auto*)

**lemma** *conj-subst*:  $P = R \implies ((P::'\alpha \text{ upred}) \wedge Q) = (R \wedge Q)$   
**by** (*pred-auto*)

**lemma** *disj-subst*:  $P = R \implies ((P::'\alpha \text{ upred}) \vee Q) = (R \vee Q)$   
**by** (*pred-auto*)

**lemma** *conj-assoc*:  $((P::'\alpha \text{ upred}) \wedge Q) \wedge S = (P \wedge (Q \wedge S))$   
**by** (*pred-auto*)

**lemma** *disj-assoc*:  $((P::'\alpha \text{ upred}) \vee Q) \vee S = (P \vee (Q \vee S))$   
**by** (*pred-auto*)

**lemma** *conj-disj-abs*:  $((P::'\alpha \text{ upred}) \wedge (P \vee Q)) = P$   
**by** (*pred-auto*)

**lemma** *disj-conj-abs*:  $((P::'\alpha \text{ upred}) \vee (P \wedge Q)) = P$   
**by** (*pred-auto*)

**lemma** *conj-disj-distr*:  $((P::'\alpha \text{ upred}) \wedge (Q \vee R)) = ((P \wedge Q) \vee (P \wedge R))$   
**by** (*pred-auto*)

**lemma** *disj-conj-distr*:  $((P::'\alpha \text{ upred}) \vee (Q \wedge R)) = ((P \vee Q) \wedge (P \vee R))$   
**by** (*pred-auto*)

**lemma** *true-disj-zero* [*simp*]:  
 $(P \vee \text{true}) = \text{true} \quad (\text{true} \vee P) = \text{true}$   
**by** (*pred-auto*)+

**lemma** *true-conj-zero* [*simp*]:  
 $(P \wedge \text{false}) = \text{false} \quad (\text{false} \wedge P) = \text{false}$   
**by** (*pred-auto*)+

**lemma** *imp-vacuous* [*simp*]:  $(\text{false} \Rightarrow u) = \text{true}$   
**by** (*pred-auto*)

**lemma** *imp-true* [*simp*]:  $(p \Rightarrow \text{true}) = \text{true}$   
**by** (*pred-auto*)

**lemma** *true-imp* [*simp*]:  $(\text{true} \Rightarrow p) = p$   
**by** (*pred-auto*)

**lemma** *impl-mp1* [*simp*]:  $(P \wedge (P \Rightarrow Q)) = (P \wedge Q)$   
**by** (*pred-auto*)

**lemma** *impl-mp2* [*simp*]:  $((P \Rightarrow Q) \wedge P) = (Q \wedge P)$   
**by** (*pred-auto*)

**lemma** *impl-adjoin*:  $((P \Rightarrow Q) \wedge R) = ((P \wedge R \Rightarrow Q \wedge R) \wedge R)$   
**by** (*pred-auto*)

**lemma** *impl-refine-intro*:  
 $\llbracket Q_1 \sqsubseteq P_1; P_2 \sqsubseteq (P_1 \wedge Q_2) \rrbracket \implies (P_1 \Rightarrow P_2) \sqsubseteq (Q_1 \Rightarrow Q_2)$

by (pred-auto)

**lemma** spec-refine:

$Q \sqsubseteq (P \wedge R) \implies (P \Rightarrow Q) \sqsubseteq R$

by (rel-auto)

**lemma** impl-disjI:  $\llbracket 'P \Rightarrow R'; 'Q \Rightarrow R' \rrbracket \implies '(P \vee Q) \Rightarrow R'$

by (rel-auto)

**lemma** conditional-iff:

$(P \Rightarrow Q) = (P \Rightarrow R) \longleftrightarrow 'P \Rightarrow (Q \Leftrightarrow R)'$

by (pred-auto)

**lemma** p-and-not-p [simp]:  $(P \wedge \neg P) = \text{false}$

by (pred-auto)

**lemma** p-or-not-p [simp]:  $(P \vee \neg P) = \text{true}$

by (pred-auto)

**lemma** p-imp-p [simp]:  $(P \Rightarrow P) = \text{true}$

by (pred-auto)

**lemma** p-iff-p [simp]:  $(P \Leftrightarrow P) = \text{true}$

by (pred-auto)

**lemma** p-imp-false [simp]:  $(P \Rightarrow \text{false}) = (\neg P)$

by (pred-auto)

**lemma** not-conj-deMorgans [simp]:  $(\neg ((P::'\alpha \text{ upred}) \wedge Q)) = ((\neg P) \vee (\neg Q))$

by (pred-auto)

**lemma** not-disj-deMorgans [simp]:  $(\neg ((P::'\alpha \text{ upred}) \vee Q)) = ((\neg P) \wedge (\neg Q))$

by (pred-auto)

**lemma** conj-disj-not-abs [simp]:  $((P::'\alpha \text{ upred}) \wedge ((\neg P) \vee Q)) = (P \wedge Q)$

by (pred-auto)

**lemma** subsumption1:

$'P \Rightarrow Q' \implies (P \vee Q) = Q$

by (pred-auto)

**lemma** subsumption2:

$'Q \Rightarrow P' \implies (P \vee Q) = P$

by (pred-auto)

**lemma** neg-conj-cancel1:  $(\neg P \wedge (P \vee Q)) = (\neg P \wedge Q :: '\alpha \text{ upred})$

by (pred-auto)

**lemma** neg-conj-cancel2:  $(\neg Q \wedge (P \vee Q)) = (\neg Q \wedge P :: '\alpha \text{ upred})$

by (pred-auto)

**lemma** double-negation [simp]:  $(\neg \neg (P::'\alpha \text{ upred})) = P$

by (pred-auto)

**lemma** true-not-false [simp]:  $\text{true} \neq \text{false} \text{ false} \neq \text{true}$

by (pred-auto)+

**lemma** *closure-conj-distr*:  $([P]_u \wedge [Q]_u) = [P \wedge Q]_u$   
by (pred-auto)

**lemma** *closure-imp-distr*:  $'[P \Rightarrow Q]_u \Rightarrow [P]_u \Rightarrow [Q]_u'$   
by (pred-auto)

**lemma** *true-iff [simp]*:  $(P \Leftrightarrow \text{true}) = P$   
by (pred-auto)

**lemma** *taut-iff-eq*:  
 $'P \Leftrightarrow Q' \longleftrightarrow (P = Q)$   
by (pred-auto)

**lemma** *impl-alt-def*:  $(P \Rightarrow Q) = (\neg P \vee Q)$   
by (pred-auto)

## 9.2 Lattice laws

**lemma** *uinf-or*:  
fixes  $P Q :: 'a \text{ upred}$   
shows  $(P \sqcap Q) = (P \vee Q)$   
by (pred-auto)

**lemma** *usup-and*:  
fixes  $P Q :: 'a \text{ upred}$   
shows  $(P \sqcup Q) = (P \wedge Q)$   
by (pred-auto)

**lemma** *UINF-alt-def*:  
 $(\bigcap i \mid A(i) \cdot P(i)) = (\bigcap i \cdot A(i) \wedge P(i))$   
by (rel-auto)

**lemma** *USUP-true [simp]*:  $(\bigcup P \mid F(P) \cdot \text{true}) = \text{true}$   
by (pred-auto)

**lemma** *UINF-mem-UNIV [simp]*:  $(\bigcap x \in \text{UNIV} \cdot P(x)) = (\bigcap x \cdot P(x))$   
by (pred-auto)

**lemma** *USUP-mem-UNIV [simp]*:  $(\bigcup x \in \text{UNIV} \cdot P(x)) = (\bigcup x \cdot P(x))$   
by (pred-auto)

**lemma** *USUP-false [simp]*:  $(\bigcup i \cdot \text{false}) = \text{false}$   
by (pred-simp)

**lemma** *UINF-true [simp]*:  $(\bigcap i \cdot \text{true}) = \text{true}$   
by (pred-simp)

**lemma** *UINF-mem-true [simp]*:  $A \neq \{\} \implies (\bigcap i \in A \cdot \text{true}) = \text{true}$   
by (pred-auto)

**lemma** *UINF-false [simp]*:  $(\bigcap i \mid P(i) \cdot \text{false}) = \text{false}$   
by (pred-auto)

**lemma** *UINF-cong-eq*:

$$\llbracket \bigwedge x. P_1(x) = P_2(x); \bigwedge x. 'P_1(x) \Rightarrow Q_1(x) =_u Q_2(x)' \rrbracket \implies$$

$$(\bigcap x \mid P_1(x) \cdot Q_1(x)) = (\bigcap x \mid P_2(x) \cdot Q_2(x))$$
 by (unfold UINF-def, pred-simp, metis)

**lemma** UINF-as-Sup:  $(\bigcap P \in \mathcal{P} \cdot P) = \bigcap \mathcal{P}$   
 apply (simp add: upred-defs bop.rep-eq lit.rep-eq Sup-uexpr-def)  
 apply (pred-simp)  
 apply (rule cong[of Sup])  
 apply (auto)  
 done

**lemma** UINF-as-Sup-collect:  $(\bigcap P \in A \cdot f(P)) = (\bigcap P \in A. f(P))$   
 apply (simp add: upred-defs bop.rep-eq lit.rep-eq Sup-uexpr-def)  
 apply (pred-simp)  
 apply (simp add: Setcompr-eq-image)  
 done

**lemma** UINF-as-Sup-collect':  $(\bigcap P \cdot f(P)) = (\bigcap P. f(P))$   
 apply (simp add: upred-defs bop.rep-eq lit.rep-eq Sup-uexpr-def)  
 apply (pred-simp)  
 apply (simp add: full-SetCompr-eq)  
 done

**lemma** UINF-as-Sup-image:  $(\bigcap P \mid \ll P \gg \in_u \ll A \gg \cdot f(P)) = \bigcap (f ' A)$   
 apply (simp add: upred-defs bop.rep-eq lit.rep-eq Sup-uexpr-def)  
 apply (pred-simp)  
 apply (rule cong[of Sup])  
 apply (auto)  
 done

**lemma** USUP-as-Inf:  $(\bigsqcup P \in \mathcal{P} \cdot P) = \bigsqcup \mathcal{P}$   
 apply (simp add: upred-defs bop.rep-eq lit.rep-eq Inf-uexpr-def)  
 apply (pred-simp)  
 apply (rule cong[of Inf])  
 apply (auto)  
 done

**lemma** USUP-as-Inf-collect:  $(\bigsqcup P \in A \cdot f(P)) = (\bigsqcup P \in A. f(P))$   
 apply (simp add: upred-defs bop.rep-eq lit.rep-eq Sup-uexpr-def)  
 apply (pred-simp)  
 apply (simp add: Setcompr-eq-image)  
 done

**lemma** USUP-as-Inf-collect':  $(\bigsqcup P \cdot f(P)) = (\bigsqcup P. f(P))$   
 apply (simp add: upred-defs bop.rep-eq lit.rep-eq Sup-uexpr-def)  
 apply (pred-simp)  
 apply (simp add: full-SetCompr-eq)  
 done

**lemma** USUP-as-Inf-image:  $(\bigsqcup P \in \mathcal{P} \cdot f(P)) = \bigsqcup (f ' \mathcal{P})$   
 apply (simp add: upred-defs bop.rep-eq lit.rep-eq Inf-uexpr-def)  
 apply (pred-simp)  
 apply (rule cong[of Inf])  
 apply (auto)  
 done



**lemma** *USUP-image-eq* [simp]:  $USUP (\lambda i. \llbracket i \rrbracket \in_u \llbracket f \cdot A \rrbracket) g = (\bigsqcup_{i \in A} \cdot g(f(i)))$   
**by** (*pred-simp*, *rule-tac cong[of Inf Inf]*, *auto*)

**lemma** *UINF-image-eq* [simp]:  $UINF (\lambda i. \llbracket i \rrbracket \in_u \llbracket f \cdot A \rrbracket) g = (\bigsqcap_{i \in A} \cdot g(f(i)))$   
**by** (*pred-simp*, *rule-tac cong[of Sup Sup]*, *auto*)

**lemma** *subst-continuous* [usubst]:  $\sigma \dagger (\bigsqcap A) = (\bigsqcap \{\sigma \dagger P \mid P. P \in A\})$   
**by** (*simp add: UINF-as-Sup[THEN sym]* *usubst setcompr-eq-image*)

**lemma** *not-UINF*:  $(\neg (\bigsqcap_{i \in A} \cdot P(i))) = (\bigsqcup_{i \in A} \cdot \neg P(i))$   
**by** (*pred-auto*)

**lemma** *not-USUP*:  $(\neg (\bigsqcup_{i \in A} \cdot P(i))) = (\bigsqcap_{i \in A} \cdot \neg P(i))$   
**by** (*pred-auto*)

**lemma** *UINF-empty* [simp]:  $(\bigsqcap_{i \in \{\}} \cdot P(i)) = \text{false}$   
**by** (*pred-auto*)

**lemma** *UINF-insert* [simp]:  $(\bigsqcap_{i \in \text{insert } x \text{ } xs} \cdot P(i)) = (P(x) \sqcap (\bigsqcap_{i \in xs} \cdot P(i)))$   
**apply** (*pred-simp*)  
**apply** (*subst Sup-insert[THEN sym]*)  
**apply** (*rule-tac cong[of Sup Sup]*)  
**apply** (*auto*)  
**done**

**lemma** *USUP-empty* [simp]:  $(\bigsqcup_{i \in \{\}} \cdot P(i)) = \text{true}$   
**by** (*pred-auto*)

**lemma** *USUP-insert* [simp]:  $(\bigsqcup_{i \in \text{insert } x \text{ } xs} \cdot P(i)) = (P(x) \sqcup (\bigsqcup_{i \in xs} \cdot P(i)))$   
**apply** (*pred-simp*)  
**apply** (*subst Inf-insert[THEN sym]*)  
**apply** (*rule-tac cong[of Inf Inf]*)  
**apply** (*auto*)  
**done**

**lemma** *conj-UINF-dist*:  
 $(P \wedge (\bigsqcap_{Q \in S} \cdot F(Q))) = (\bigsqcap_{Q \in S} \cdot P \wedge F(Q))$   
**by** (*simp add: upred-defs bop.rep-eq lit.rep-eq, pred-auto*)

**lemma** *disj-UINF-dist*:  
 $S \neq \{\} \implies (P \vee (\bigsqcap_{Q \in S} \cdot F(Q))) = (\bigsqcap_{Q \in S} \cdot P \vee F(Q))$   
**by** (*simp add: upred-defs bop.rep-eq lit.rep-eq, pred-auto*)

**lemma** *conj-USUP-dist*:  
 $S \neq \{\} \implies (P \wedge (\bigsqcup_{Q \in S} \cdot F(Q))) = (\bigsqcup_{Q \in S} \cdot P \wedge F(Q))$   
**by** (*subst uexpr-eq-iff, auto simp add: conj-upred-def USUP.rep-eq inf-uexpr.rep-eq bop.rep-eq lit.rep-eq*)

**lemma** *USUP-conj-USUP*:  $((\bigsqcup_{P \in A} \cdot F(P)) \wedge (\bigsqcup_{P \in A} \cdot G(P))) = (\bigsqcup_{P \in A} \cdot F(P) \wedge G(P))$   
**by** (*simp add: upred-defs bop.rep-eq lit.rep-eq, pred-auto*)

**lemma** *UINF-all-cong*:  
**assumes**  $\bigwedge P. F(P) = G(P)$   
**shows**  $(\bigsqcap P \cdot F(P)) = (\bigsqcap P \cdot G(P))$   
**by** (*simp add: UINF-as-Sup-collect assms*)

**lemma** *UINF-cong*:

**assumes**  $\bigwedge P. P \in A \implies F(P) = G(P)$   
**shows**  $(\bigcap P \in A \cdot F(P)) = (\bigcap P \in A \cdot G(P))$   
**by** (*simp add: UINF-as-Sup-collect assms*)

**lemma** *USUP-all-cong*:

**assumes**  $\bigwedge P. F(P) = G(P)$   
**shows**  $(\bigsqcup P \cdot F(P)) = (\bigsqcup P \cdot G(P))$   
**by** (*simp add: assms*)

**lemma** *USUP-cong*:

**assumes**  $\bigwedge P. P \in A \implies F(P) = G(P)$   
**shows**  $(\bigsqcup P \in A \cdot F(P)) = (\bigsqcup P \in A \cdot G(P))$   
**by** (*simp add: USUP-as-Inf-collect assms*)

**lemma** *UINF-subset-mono*:  $A \subseteq B \implies (\bigcap P \in B \cdot F(P)) \sqsubseteq (\bigcap P \in A \cdot F(P))$   
**by** (*simp add: SUP-subset-mono UINF-as-Sup-collect*)

**lemma** *USUP-subset-mono*:  $A \subseteq B \implies (\bigsqcup P \in A \cdot F(P)) \sqsubseteq (\bigsqcup P \in B \cdot F(P))$   
**by** (*simp add: INF-superset-mono USUP-as-Inf-collect*)

**lemma** *UINF-impl*:  $(\bigcap P \in A \cdot F(P) \Rightarrow G(P)) = ((\bigsqcup P \in A \cdot F(P)) \Rightarrow (\bigcap P \in A \cdot G(P)))$   
**by** (*pred-auto*)

**lemma** *UINF-all-nats* [*simp*]:

**fixes**  $P :: \text{nat} \Rightarrow 'a \text{ upred}$   
**shows**  $(\bigcap n \cdot \bigcap i \in \{0..n\} \cdot P(i)) = (\bigcap i \in \{0..\} \cdot P(i))$   
**by** (*pred-auto*)

**lemma** *UINF-refines'*:

**assumes**  $\bigwedge i. P \sqsubseteq Q(i)$   
**shows**  $P \sqsubseteq (\bigcap i \cdot Q(i))$   
**using** *assms*  
**apply** (*rel-auto*) **using** *Sup-le-iff* **by** *fastforce*

### 9.3 Equality laws

**lemma** *eq-upred-refl* [*simp*]:  $(x =_u x) = \text{true}$   
**by** (*pred-auto*)

**lemma** *eq-upred-sym*:  $(x =_u y) = (y =_u x)$   
**by** (*pred-auto*)

**lemma** *eq-cong-left*:

**assumes**  $\text{vwb-lens } x \ \$x \ \# \ Q \ \$x' \ \# \ Q \ \$x \ \# \ R \ \$x' \ \# \ R$   
**shows**  $((\$x' =_u \$x \wedge Q) = (\$x' =_u \$x \wedge R)) \longleftrightarrow (Q = R)$   
**using** *assms*  
**by** (*pred-simp, (meson mwb-lens-def vwb-lens-mwb weak-lens-def)+*)

**lemma** *conj-eq-in-var-subst*:

**fixes**  $x :: ('a \implies 'a)$   
**assumes**  $\text{vwb-lens } x$   
**shows**  $(P \wedge \$x =_u v) = (P[v/\$x] \wedge \$x =_u v)$   
**using** *assms*  
**by** (*pred-simp, (metis vwb-lens-wb wb-lens.get-put)+*)

**lemma** *conj-eq-out-var-subst*:

**fixes**  $x :: ('a \Rightarrow 'a)$   
**assumes** *vwb-lens*  $x$   
**shows**  $(P \wedge \$x' =_u v) = (P[v/\$x'] \wedge \$x' =_u v)$   
**using** *assms*  
**by** (*pred-simp*, (*metis vwb-lens-wb wb-lens.get-put*)+)

**lemma** *conj-pos-var-subst*:

**assumes** *vwb-lens*  $x$   
**shows**  $(\$x \wedge Q) = (\$x \wedge Q[true/\$x])$   
**using** *assms*  
**by** (*pred-auto*, *metis (full-types) vwb-lens-wb wb-lens.get-put*, *metis (full-types) vwb-lens-wb wb-lens.get-put*)

**lemma** *conj-neg-var-subst*:

**assumes** *vwb-lens*  $x$   
**shows**  $(\neg \$x \wedge Q) = (\neg \$x \wedge Q[false/\$x])$   
**using** *assms*  
**by** (*pred-auto*, *metis (full-types) vwb-lens-wb wb-lens.get-put*, *metis (full-types) vwb-lens-wb wb-lens.get-put*)

**lemma** *upred-eq-true [simp]*:  $(p =_u true) = p$

**by** (*pred-auto*)

**lemma** *upred-eq-false [simp]*:  $(p =_u false) = (\neg p)$

**by** (*pred-auto*)

**lemma** *upred-true-eq [simp]*:  $(true =_u p) = p$

**by** (*pred-auto*)

**lemma** *upred-false-eq [simp]*:  $(false =_u p) = (\neg p)$

**by** (*pred-auto*)

**lemma** *conj-var-subst*:

**assumes** *vwb-lens*  $x$   
**shows**  $(P \wedge \text{var } x =_u v) = (P[v/x] \wedge \text{var } x =_u v)$   
**using** *assms*  
**by** (*pred-simp*, (*metis (full-types) vwb-lens-def wb-lens.get-put*)+)

## 9.4 HOL Variable Quantifiers

**lemma** *shEx-unbound [simp]*:  $(\exists x \cdot P) = P$

**by** (*pred-auto*)

**lemma** *shEx-bool [simp]*:  $\text{shEx } P = (P \text{ True} \vee P \text{ False})$

**by** (*pred-simp*, *metis (full-types)*)

**lemma** *shEx-commute*:  $(\exists x \cdot \exists y \cdot P x y) = (\exists y \cdot \exists x \cdot P x y)$

**by** (*pred-auto*)

**lemma** *shEx-cong*:  $[\bigwedge x. P x = Q x] \Longrightarrow \text{shEx } P = \text{shEx } Q$

**by** (*pred-auto*)

**lemma** *shAll-unbound [simp]*:  $(\forall x \cdot P) = P$

**by** (*pred-auto*)

**lemma** *shAll-bool [simp]*:  $\text{shAll } P = (P \text{ True} \wedge P \text{ False})$

by (*pred-simp*, *metis* (*full-types*))

**lemma** *shAll-cong*:  $\llbracket \bigwedge x. P\ x = Q\ x \rrbracket \implies \text{shAll}\ P = \text{shAll}\ Q$   
 by (*pred-auto*)

Quantifier lifting

**named-theorems** *uquant-lift*

**lemma** *shEx-lift-conj-1* [*uquant-lift*]:  
 $((\exists x \cdot P(x)) \wedge Q) = (\exists x \cdot P(x) \wedge Q)$   
 by (*pred-auto*)

**lemma** *shEx-lift-conj-2* [*uquant-lift*]:  
 $(P \wedge (\exists x \cdot Q(x))) = (\exists x \cdot P \wedge Q(x))$   
 by (*pred-auto*)

## 9.5 Case Splitting

**lemma** *eq-split-subst*:  
 assumes *vwb-lens* *x*  
 shows  $(P = Q) \longleftrightarrow (\forall v. P\llbracket \llbracket v \rrbracket / x \rrbracket = Q\llbracket \llbracket v \rrbracket / x \rrbracket)$   
 using *assms*  
 by (*pred-auto*, *metis* *vwb-lens-wb* *wb-lens.source-stability*)

**lemma** *eq-split-substI*:  
 assumes *vwb-lens* *x*  $\bigwedge v. P\llbracket \llbracket v \rrbracket / x \rrbracket = Q\llbracket \llbracket v \rrbracket / x \rrbracket$   
 shows  $P = Q$   
 using *assms*(1) *assms*(2) *eq-split-subst* by *blast*

**lemma** *taut-split-subst*:  
 assumes *vwb-lens* *x*  
 shows  $\langle P \rangle \longleftrightarrow (\forall v. \langle P\llbracket \llbracket v \rrbracket / x \rrbracket \rangle)$   
 using *assms*  
 by (*pred-auto*, *metis* *vwb-lens-wb* *wb-lens.source-stability*)

**lemma** *eq-split*:  
 assumes  $\langle P \Rightarrow Q \rangle$   $\langle Q \Rightarrow P \rangle$   
 shows  $P = Q$   
 using *assms*  
 by (*pred-auto*)

**lemma** *bool-eq-splitI*:  
 assumes *vwb-lens* *x*  $P\llbracket \text{true}/x \rrbracket = Q\llbracket \text{true}/x \rrbracket$   $P\llbracket \text{false}/x \rrbracket = Q\llbracket \text{false}/x \rrbracket$   
 shows  $P = Q$   
 by (*metis* (*full-types*) *assms* *eq-split-subst* *false-alt-def* *true-alt-def*)

**lemma** *subst-bool-split*:  
 assumes *vwb-lens* *x*  
 shows  $\langle P \rangle = \langle (P\llbracket \text{false}/x \rrbracket \wedge P\llbracket \text{true}/x \rrbracket) \rangle$   
**proof** –  
 from *assms* have  $\langle P \rangle = (\forall v. \langle P\llbracket \llbracket v \rrbracket / x \rrbracket \rangle)$   
 by (*subst* *taut-split-subst*[*of* *x*], *auto*)  
 also have  $\dots = \langle P\llbracket \llbracket \text{True} \rrbracket / x \rrbracket \wedge P\llbracket \llbracket \text{False} \rrbracket / x \rrbracket \rangle$   
 by (*metis* (*mono-tags*, *lifting*))  
 also have  $\dots = \langle (P\llbracket \text{false}/x \rrbracket \wedge P\llbracket \text{true}/x \rrbracket) \rangle$   
 by (*pred-auto*)

finally show *?thesis* .  
qed

**lemma** *subst-eq-replace*:  
fixes  $x :: ('a \implies 'a)$   
shows  $(p[u/x] \wedge u =_u v) = (p[v/x] \wedge u =_u v)$   
by (*pred-auto*)

## 9.6 UTP Quantifiers

**lemma** *one-point*:  
assumes *mbw-lens*  $x \not\# v$   
shows  $(\exists x \cdot P \wedge \text{var } x =_u v) = P[v/x]$   
using *assms*  
by (*pred-auto*)

**lemma** *exists-twice*: *mbw-lens*  $x \implies (\exists x \cdot \exists x \cdot P) = (\exists x \cdot P)$   
by (*pred-auto*)

**lemma** *all-twice*: *mbw-lens*  $x \implies (\forall x \cdot \forall x \cdot P) = (\forall x \cdot P)$   
by (*pred-auto*)

**lemma** *exists-sub*:  $\llbracket \text{mbw-lens } y; x \subseteq_L y \rrbracket \implies (\exists x \cdot \exists y \cdot P) = (\exists y \cdot P)$   
by (*pred-auto*)

**lemma** *all-sub*:  $\llbracket \text{mbw-lens } y; x \subseteq_L y \rrbracket \implies (\forall x \cdot \forall y \cdot P) = (\forall y \cdot P)$   
by (*pred-auto*)

**lemma** *ex-commute*:  
assumes  $x \bowtie y$   
shows  $(\exists x \cdot \exists y \cdot P) = (\exists y \cdot \exists x \cdot P)$   
using *assms*  
apply (*pred-auto*)  
using *lens-indep-comm* apply *fastforce* +  
done

**lemma** *all-commute*:  
assumes  $x \bowtie y$   
shows  $(\forall x \cdot \forall y \cdot P) = (\forall y \cdot \forall x \cdot P)$   
using *assms*  
apply (*pred-auto*)  
using *lens-indep-comm* apply *fastforce* +  
done

**lemma** *ex-equiv*:  
assumes  $x \approx_L y$   
shows  $(\exists x \cdot P) = (\exists y \cdot P)$   
using *assms*  
by (*pred-simp*, *metis* (*no-types*, *lifting*) *lens.select-convs*(2))

**lemma** *all-equiv*:  
assumes  $x \approx_L y$   
shows  $(\forall x \cdot P) = (\forall y \cdot P)$   
using *assms*  
by (*pred-simp*, *metis* (*no-types*, *lifting*) *lens.select-convs*(2))

**lemma** *ex-zero*:

$(\exists \emptyset \cdot P) = P$   
**by** (*pred-auto*)

**lemma** *all-zero*:

$(\forall \emptyset \cdot P) = P$   
**by** (*pred-auto*)

**lemma** *ex-plus*:

$(\exists y; x \cdot P) = (\exists x \cdot \exists y \cdot P)$   
**by** (*pred-auto*)

**lemma** *all-plus*:

$(\forall y; x \cdot P) = (\forall x \cdot \forall y \cdot P)$   
**by** (*pred-auto*)

**lemma** *closure-all*:

$[P]_u = (\forall \Sigma \cdot P)$   
**by** (*pred-auto*)

**lemma** *unrest-as-exists*:

$wb\text{-}lens\ x \implies (x \# P) \longleftrightarrow ((\exists x \cdot P) = P)$   
**by** (*pred-simp*, *metis wb-lens.put-eq*)

**lemma** *ex-mono*:  $P \sqsubseteq Q \implies (\exists x \cdot P) \sqsubseteq (\exists x \cdot Q)$

**by** (*pred-auto*)

**lemma** *ex-weakens*:  $wb\text{-}lens\ x \implies (\exists x \cdot P) \sqsubseteq P$

**by** (*pred-simp*, *metis wb-lens.get-put*)

**lemma** *all-mono*:  $P \sqsubseteq Q \implies (\forall x \cdot P) \sqsubseteq (\forall x \cdot Q)$

**by** (*pred-auto*)

**lemma** *all-strengthens*:  $wb\text{-}lens\ x \implies P \sqsubseteq (\forall x \cdot P)$

**by** (*pred-simp*, *metis wb-lens.get-put*)

**lemma** *ex-unrest*:  $x \# P \implies (\exists x \cdot P) = P$

**by** (*pred-auto*)

**lemma** *all-unrest*:  $x \# P \implies (\forall x \cdot P) = P$

**by** (*pred-auto*)

**lemma** *not-ex-not*:  $\neg (\exists x \cdot \neg P) = (\forall x \cdot P)$

**by** (*pred-auto*)

**lemma** *not-all-not*:  $\neg (\forall x \cdot \neg P) = (\exists x \cdot P)$

**by** (*pred-auto*)

## 9.7 Variable Restriction

**lemma** *var-res-all*:

$P \upharpoonright_v \Sigma = P$

**by** (*rel-auto*)

**lemma** *var-res-twice*:

$mwb\text{-}lens\ x \implies P \upharpoonright_v x \upharpoonright_v x = P \upharpoonright_v x$

by (*pred-auto*)

## 9.8 Conditional laws

**lemma** *cond-def*:

$$(P \triangleleft b \triangleright Q) = ((b \wedge P) \vee ((\neg b) \wedge Q))$$

by (*pred-auto*)

**lemma** *cond-idem*:  $(P \triangleleft b \triangleright P) = P$  by (*pred-auto*)

**lemma** *cond-symm*:  $(P \triangleleft b \triangleright Q) = (Q \triangleleft \neg b \triangleright P)$  by (*pred-auto*)

**lemma** *cond-assoc*:  $((P \triangleleft b \triangleright Q) \triangleleft c \triangleright R) = (P \triangleleft b \wedge c \triangleright (Q \triangleleft c \triangleright R))$  by (*pred-auto*)

**lemma** *cond-distr*:  $(P \triangleleft b \triangleright (Q \triangleleft c \triangleright R)) = ((P \triangleleft b \triangleright Q) \triangleleft c \triangleright (P \triangleleft b \triangleright R))$  by (*pred-auto*)

**lemma** *cond-unit-T* [*simp*]:  $(P \triangleleft \text{true} \triangleright Q) = P$  by (*pred-auto*)

**lemma** *cond-unit-F* [*simp*]:  $(P \triangleleft \text{false} \triangleright Q) = Q$  by (*pred-auto*)

**lemma** *cond-conj-not*:  $((P \triangleleft b \triangleright Q) \wedge (\neg b)) = (Q \wedge (\neg b))$

by (*rel-auto*)

**lemma** *cond-and-T-integrate*:

$$((P \wedge b) \vee (Q \triangleleft b \triangleright R)) = ((P \vee Q) \triangleleft b \triangleright R)$$

by (*pred-auto*)

**lemma** *cond-L6*:  $(P \triangleleft b \triangleright (Q \triangleleft b \triangleright R)) = (P \triangleleft b \triangleright R)$  by (*pred-auto*)

**lemma** *cond-L7*:  $(P \triangleleft b \triangleright (P \triangleleft c \triangleright Q)) = (P \triangleleft b \vee c \triangleright Q)$  by (*pred-auto*)

**lemma** *cond-and-distr*:  $((P \wedge Q) \triangleleft b \triangleright (R \wedge S)) = ((P \triangleleft b \triangleright R) \wedge (Q \triangleleft b \triangleright S))$  by (*pred-auto*)

**lemma** *cond-or-distr*:  $((P \vee Q) \triangleleft b \triangleright (R \vee S)) = ((P \triangleleft b \triangleright R) \vee (Q \triangleleft b \triangleright S))$  by (*pred-auto*)

**lemma** *cond-imp-distr*:

$$((P \Rightarrow Q) \triangleleft b \triangleright (R \Rightarrow S)) = ((P \triangleleft b \triangleright R) \Rightarrow (Q \triangleleft b \triangleright S))$$

by (*pred-auto*)

**lemma** *cond-eq-distr*:

$$((P \Leftrightarrow Q) \triangleleft b \triangleright (R \Leftrightarrow S)) = ((P \triangleleft b \triangleright R) \Leftrightarrow (Q \triangleleft b \triangleright S))$$

by (*pred-auto*)

**lemma** *cond-conj-distr*:  $(P \wedge (Q \triangleleft b \triangleright S)) = ((P \wedge Q) \triangleleft b \triangleright (P \wedge S))$  by (*pred-auto*)

**lemma** *cond-disj-distr*:  $(P \vee (Q \triangleleft b \triangleright S)) = ((P \vee Q) \triangleleft b \triangleright (P \vee S))$  by (*pred-auto*)

**lemma** *cond-neg*:  $\neg (P \triangleleft b \triangleright Q) = ((\neg P) \triangleleft b \triangleright (\neg Q))$  by (*pred-auto*)

**lemma** *cond-conj*:  $P \triangleleft b \wedge c \triangleright Q = (P \triangleleft c \triangleright Q) \triangleleft b \triangleright Q$

by (*pred-auto*)

**lemma** *spec-cond-dist*:  $(P \Rightarrow (Q \triangleleft b \triangleright R)) = ((P \Rightarrow Q) \triangleleft b \triangleright (P \Rightarrow R))$

by (*pred-auto*)

**lemma** *cond-USUP-dist*:  $(\bigsqcup_{P \in S} P \cdot F(P)) \triangleleft b \triangleright (\bigsqcup_{P \in S} P \cdot G(P)) = (\bigsqcup_{P \in S} P \cdot F(P) \triangleleft b \triangleright G(P))$

by (*pred-auto*)

**lemma** *cond-UNIF-dist*:  $(\prod P \in S \cdot F(P)) \triangleleft b \triangleright (\prod P \in S \cdot G(P)) = (\prod P \in S \cdot F(P) \triangleleft b \triangleright G(P))$   
**by** (*pred-auto*)

**lemma** *cond-var-subst-left*:

**assumes** *vwb-lens x*

**shows**  $(P \llbracket \text{true}/x \rrbracket \triangleleft \text{var } x \triangleright Q) = (P \triangleleft \text{var } x \triangleright Q)$

**using** *assms* **by** (*pred-auto*, *metis (full-types) vwb-lens-wb wb-lens.get-put*)

**lemma** *cond-var-subst-right*:

**assumes** *vwb-lens x*

**shows**  $(P \triangleleft \text{var } x \triangleright Q \llbracket \text{false}/x \rrbracket) = (P \triangleleft \text{var } x \triangleright Q)$

**using** *assms* **by** (*pred-auto*, *metis (full-types) vwb-lens.put-eq*)

**lemma** *cond-var-split*:

*vwb-lens x*  $\implies (P \llbracket \text{true}/x \rrbracket \triangleleft \text{var } x \triangleright P \llbracket \text{false}/x \rrbracket) = P$

**by** (*rel-simp*, (*metis (full-types) vwb-lens.put-eq*)+)

**lemma** *cond-assign-subst*:

*vwb-lens x*  $\implies (P \triangleleft \text{utp-expr.var } x =_u v \triangleright Q) = (P \llbracket v/x \rrbracket \triangleleft \text{utp-expr.var } x =_u v \triangleright Q)$

**apply** (*rel-simp*) **using** *vwb-lens.put-eq* **by** *force*

**lemma** *conj-conds*:

$(P1 \triangleleft b \triangleright Q1 \wedge P2 \triangleleft b \triangleright Q2) = (P1 \wedge P2) \triangleleft b \triangleright (Q1 \wedge Q2)$

**by** *pred-auto*

**lemma** *disj-conds*:

$(P1 \triangleleft b \triangleright Q1 \vee P2 \triangleleft b \triangleright Q2) = (P1 \vee P2) \triangleleft b \triangleright (Q1 \vee Q2)$

**by** *pred-auto*

## 9.9 Additional Expression Laws

**lemma** *le-pred-refl [simp]*:

**fixes** *x :: ('a::preorder, 'α) uexpr*

**shows**  $(x \leq_u x) = \text{true}$

**by** (*pred-auto*)

**lemma** *uzero-le-laws [simp]*:

$(0 :: ('a::\{\text{linordered-semidom}\}, 'α) uexpr) \leq_u \text{numeral } x = \text{true}$

$(1 :: ('a::\{\text{linordered-semidom}\}, 'α) uexpr) \leq_u \text{numeral } x = \text{true}$

$(0 :: ('a::\{\text{linordered-semidom}\}, 'α) uexpr) \leq_u 1 = \text{true}$

**by** (*pred-simp*)+

**lemma** *unumeral-le-1 [simp]*:

**assumes**  $(\text{numeral } i :: 'a::\{\text{numeral,ord}\}) \leq \text{numeral } j$

**shows**  $(\text{numeral } i :: ('a, 'α) uexpr) \leq_u \text{numeral } j = \text{true}$

**using** *assms* **by** (*pred-auto*)

**lemma** *unumeral-le-2 [simp]*:

**assumes**  $(\text{numeral } i :: 'a::\{\text{numeral,linorder}\}) > \text{numeral } j$

**shows**  $(\text{numeral } i :: ('a, 'α) uexpr) \leq_u \text{numeral } j = \text{false}$

**using** *assms* **by** (*pred-auto*)

**lemma** *uset-laws [simp]*:

$x \in_u \{\} = \text{false}$

$x \in_u \{m..n\}_u = (m \leq_u x \wedge x \leq_u n)$

**by** (*pred-auto*)+



**lemma** *pfun-entries-apply* [simp]:  
 $(entr_u(d, f) :: (('k, 'v) pfun, 'α) uepr)(i)_a = ((\ll f \gg)(i)_a) \triangleleft i \in_u d \triangleright \perp_u$   
**by** (*pred-auto*)

**lemma** *uupdate-pfun* [simp]:  
**fixes**  $m :: (('k, 'v) pfun, 'α) uepr$   
**shows**  $dom_u(m(k \mapsto v)_u) = \{k\}_u \cup_u dom_u(m)$   
**by** (*rel-auto*)

**lemma** *uapply-uupdate-pfun* [simp]:  
**fixes**  $m :: (('k, 'v) pfun, 'α) uepr$   
**shows**  $(m(k \mapsto v)_u)(i)_a = v \triangleleft i =_u k \triangleright m(i)_a$   
**by** (*rel-auto*)

**lemma** *ulit-eq* [simp]:  $x = y \implies (\ll x \gg =_u \ll y \gg) = true$   
**by** (*rel-auto*)

**lemma** *ulit-neg* [simp]:  $x \neq y \implies (\ll x \gg =_u \ll y \gg) = false$   
**by** (*rel-auto*)

**lemma** *uset-mems* [simp]:  
 $x \in_u \{y\}_u = (x =_u y)$   
 $x \in_u A \cup_u B = (x \in_u A \vee x \in_u B)$   
 $x \in_u A \cap_u B = (x \in_u A \wedge x \in_u B)$   
**by** (*rel-auto*)<sup>+</sup>

## 9.10 Refinement By Observation

Function to obtain the set of observations of a predicate

**definition** *obs-upred* ::  $'α \text{ upred} \Rightarrow 'α \text{ set } ([\![\cdot]\!]_o)$   
**where** [*upred-defs*]:  $[P]_o = \{b. [P]_e b\}$

**lemma** *obs-upred-refine-iff*:  
 $P \sqsubseteq Q \longleftrightarrow [Q]_o \subseteq [P]_o$   
**by** (*pred-auto*)

A refinement can be demonstrated by considering only the observations of the predicates which are relevant, i.e. not unrestricted, for them. In other words, if the alphabet can be split into two disjoint segments,  $x$  and  $y$ , and neither predicate refers to  $y$  then only  $x$  need be considered when checking for observations.

**lemma** *refine-by-obs*:  
**assumes**  $x \bowtie y \text{ bij-lens } (x +_L y) \ y \nmid P \ y \nmid Q \ \{v. 'P[\ll v \gg/x]\} \subseteq \{v. 'Q[\ll v \gg/x]\}$   
**shows**  $Q \sqsubseteq P$   
**using** *assms*(3-5)  
**apply** (*simp add: obs-upred-refine-iff subset-eq*)  
**apply** (*pred-simp*)  
**apply** (*rename-tac b*)  
**apply** (*drule-tac x=get\_x b in spec*)  
**apply** (*auto simp add: assms*)  
**apply** (*metis assms*(1) *assms*(2) *bij-lens.axioms*(2) *bij-lens-axioms-def lens-override-def lens-override-plus*)<sup>+</sup>  
**done**

## 9.11 Cylindric Algebra

**lemma** *C1*:  $(\exists x \cdot \text{false}) = \text{false}$   
**by** (*pred-auto*)

**lemma** *C2*:  $\text{wb-lens } x \implies 'P \Rightarrow (\exists x \cdot P)'$   
**by** (*pred-simp, metis wb-lens.get-put*)

**lemma** *C3*:  $\text{mwb-lens } x \implies (\exists x \cdot (P \wedge (\exists x \cdot Q))) = ((\exists x \cdot P) \wedge (\exists x \cdot Q))$   
**by** (*pred-auto*)

**lemma** *C4a*:  $x \approx_L y \implies (\exists x \cdot \exists y \cdot P) = (\exists y \cdot \exists x \cdot P)$   
**by** (*pred-simp, metis (no-types, lifting) lens.select-convs(2)*)<sup>+</sup>

**lemma** *C4b*:  $x \bowtie y \implies (\exists x \cdot \exists y \cdot P) = (\exists y \cdot \exists x \cdot P)$   
**using** *ex-commute* **by** *blast*

**lemma** *C5*:  
**fixes**  $x :: ('a \implies 'a)$   
**shows**  $(\&x =_u \&x) = \text{true}$   
**by** (*pred-auto*)

**lemma** *C6*:  
**assumes**  $\text{wb-lens } x \bowtie y \bowtie z$   
**shows**  $(\&y =_u \&z) = (\exists x \cdot \&y =_u \&x \wedge \&x =_u \&z)$   
**using** *assms*  
**by** (*pred-simp, (metis lens-indep-def)*)<sup>+</sup>

**lemma** *C7*:  
**assumes**  $\text{weak-lens } x \bowtie y$   
**shows**  $((\exists x \cdot \&x =_u \&y \wedge P) \wedge (\exists x \cdot \&x =_u \&y \wedge \neg P)) = \text{false}$   
**using** *assms*  
**by** (*pred-simp, simp add: lens-indep-sym*)

**end**

## 10 Fixed-points and Recursion

**theory** *utp-recursion*  
**imports** *utp-pred-laws*  
**begin**

### 10.1 Fixed-point Laws

**lemma** *mu-id*:  $(\mu X \cdot X) = \text{true}$   
**by** (*simp add: antisym gfp-upperbound*)

**lemma** *mu-const*:  $(\mu X \cdot P) = P$   
**by** (*simp add: gfp-const*)

**lemma** *nu-id*:  $(\nu X \cdot X) = \text{false}$   
**by** (*meson lfp-lowerbound utp-pred-laws.bot.extremum-unique*)

**lemma** *nu-const*:  $(\nu X \cdot P) = P$   
**by** (*simp add: lfp-const*)

```

lemma mu-refine-intro:
  assumes  $(C \Rightarrow S) \sqsubseteq F(C \Rightarrow S)$  ' $C \Rightarrow (\mu F \Leftrightarrow \nu F)$ '
  shows  $(C \Rightarrow S) \sqsubseteq \mu F$ 
proof –
  from assms have  $(C \Rightarrow S) \sqsubseteq \nu F$ 
    by (simp add: lfp-lowerbound)
  with assms show ?thesis
    by (pred-auto)
qed

```

## 10.2 Obtaining Unique Fixed-points

Obtaining termination proofs via approximation chains. Theorems and proofs adapted from Chapter 2, page 63 of the UTP book [5].

**type-synonym** '*a chain* = *nat*  $\Rightarrow$  '*a upred*

**definition** *chain* :: '*a chain*  $\Rightarrow$  *bool* **where**  
*chain* *Y* =  $((Y\ 0 = \text{false}) \wedge (\forall\ i. Y\ (\text{Suc } i) \sqsubseteq Y\ i))$

**lemma** *chain0* [*simp*]: *chain* *Y*  $\Longrightarrow$  *Y* 0 = *false*  
**by** (*simp add: chain-def*)

**lemma** *chainI*:  
**assumes**  $Y\ 0 = \text{false} \wedge i. Y\ (\text{Suc } i) \sqsubseteq Y\ i$   
**shows** *chain* *Y*  
**using** *assms* **by** (*auto simp add: chain-def*)

**lemma** *chainE*:  
**assumes** *chain* *Y*  $\wedge i. \llbracket Y\ 0 = \text{false}; Y\ (\text{Suc } i) \sqsubseteq Y\ i \rrbracket \Longrightarrow P$   
**shows** *P*  
**using** *assms* **by** (*simp add: chain-def*)

**lemma** *L274*:  
**assumes**  $\forall\ n. ((E\ n \wedge_p X) = (E\ n \wedge Y))$   
**shows**  $(\bigcap (\text{range } E) \wedge X) = (\bigcap (\text{range } E) \wedge Y)$   
**using** *assms* **by** (*pred-auto*)

Constructive chains

**definition** *constr* ::  
 $(\text{'a upred} \Rightarrow \text{'a upred}) \Rightarrow \text{'a chain} \Rightarrow \text{bool}$  **where**  
*constr* *F* *E*  $\longleftrightarrow \text{chain } E \wedge (\forall\ X\ n. ((F(X) \wedge E(n+1)) = (F(X \wedge E(n)) \wedge E(n+1))))$

This lemma gives a way of showing that there is a unique fixed-point when the predicate function can be built using a constructive function F over an approximation chain E

**lemma** *chain-pred-terminates*:  
**assumes** *constr* *F* *E* *mono* *F*  
**shows**  $(\bigcap (\text{range } E) \wedge \mu F) = (\bigcap (\text{range } E) \wedge \nu F)$   
**proof** –  
**from** *assms* **have**  $\forall\ n. (E\ n \wedge \mu F) = (E\ n \wedge \nu F)$   
**proof** (*rule-tac allI*)  
**fix** *n*  
**from** *assms* **show**  $(E\ n \wedge \mu F) = (E\ n \wedge \nu F)$   
**proof** (*induct n*)  
**case** 0 **thus** ?case **by** (*simp add: constr-def*)

```

next
  case (Suc n)
  note hyp = this
  thus ?case
  proof -
    have (E (n + 1)  $\wedge$   $\mu$  F) = (E (n + 1)  $\wedge$  F ( $\mu$  F))
      using gfp-unfold[OF hyp(3), THEN sym] by (simp add: constr-def)
    also from hyp have ... = (E (n + 1)  $\wedge$  F (E n  $\wedge$   $\mu$  F))
      by (metis conj-comm constr-def)
    also from hyp have ... = (E (n + 1)  $\wedge$  F (E n  $\wedge$   $\nu$  F))
      by simp
    also from hyp have ... = (E (n + 1)  $\wedge$   $\nu$  F)
      by (metis (no-types, lifting) conj-comm constr-def lfp-unfold)
    ultimately show ?thesis
      by simp
  qed
qed
qed
qed
thus ?thesis
  by (auto intro: L274)
qed

```

```

theorem constr-fp-uniq:
  assumes constr F E mono F  $\sqcap$  (range E) = C
  shows (C  $\wedge$   $\mu$  F) = (C  $\wedge$   $\nu$  F)
  using assms(1) assms(2) assms(3) chain-pred-terminates by blast

```

end

## 11 UTP Events

```

theory utp-event
imports utp-pred
begin

```

### 11.1 Events

Events of some type  $'\vartheta$  are just the elements of that type.

```
type-synonym ' $\vartheta$  event = ' $\vartheta$ 
```

### 11.2 Channels

Typed channels are modelled as functions. Below,  $'a$  determines the channel type and  $'\vartheta$  the underlying event type. As with values, it is difficult to introduce channels as monomorphic types due to the fact that they can have arbitrary parametrisations in term of  $'a$ . Applying a channel to an element of its type yields an event, as we may expect. Though this is not formalised here, we may also sensibly assume that all channel- representing functions are injective. Note: is there benefit in formalising this here?

```
type-synonym ('a, ' $\vartheta$ ) chan = ' $a \Rightarrow '\vartheta$  event
```

A downside of the approach is that the event type  $'\vartheta$  must be able to encode *all* events of a process model, and hence cannot be fixed upfront for a single channel or channel set. To do so, we actually require a notion of ‘extensible’ datatypes, in analogy to extensible record types.

Another solution is to encode a notion of channel scoping that namely uses *sum* types to lift channel types into extensible ones, that is using channel-set specific scoping operators. This is a current work in progress.

### 11.2.1 Operators

The  $Z$  type of a channel corresponds to the entire carrier of the underlying HOL type of that channel. Strictly, the function is redundant but was added to mirror the mathematical account in [?]. (TODO: Ask Simon Foster for [?])

**definition** *chan-type* :: ('a, 'v) chan  $\Rightarrow$  'a set ( $\delta_u$ ) **where**  
 [upred-defs]:  $\delta_u c = UNIV$

The next lifted function creates an expression that yields a channel event, from an expression on the channel type 'a.

**definition** *chan-apply* ::  
 ('a, 'v) chan  $\Rightarrow$  ('a, 'α) uexpr  $\Rightarrow$  ('v event, 'α) uexpr (('·/-')<sub>u</sub>) **where**  
 [upred-defs]:  $(c \cdot e)_u = \ll c \gg (e)_a$

**lemma** *unrest-chan-apply* [unrest]:  $x \# e \Longrightarrow x \# (c \cdot e)_u$   
 by (rel-auto)

**lemma** *usubst-chan-apply* [usubst]:  $\sigma \dagger (c \cdot v)_u = (c \cdot \sigma \dagger v)_u$   
 by (rel-auto)

end

## 12 Alphabet Manipulation

**theory** *utp-alphabet*  
**imports**  
*utp-pred utp-event*  
**begin**

### 12.1 Preliminaries

Alphabets are simply types that characterise the state-space of an expression. Thus the Isabelle type system ensures that predicates cannot refer to variables not in the alphabet as this would be a type error. Often one would like to add or remove additional variables, for example if we wish to have a predicate which ranges only a smaller state-space, and then lift it into a predicate over a larger one. This is useful, for example, when dealing with relations which refer only to undashed variables (conditions) since we can use the type system to ensure well-formedness.

In this theory we will set up operators for extending and contracting and alphabet. We first set up a theorem attribute for alphabet laws and a tactic.

**named-theorems** *alpha*

**method** *alpha-tac* = (simp add: alpha unrest)?

### 12.2 Alphabet Extrusion

Alter an alphabet by application of a lens that demonstrates how the smaller alphabet ( $\beta$ ) injects into the larger alphabet ( $\alpha$ ). This changes the type of the expression so it is parametrised over

the large alphabet. We do this by using the lens *get* function to extract the smaller state binding, and then apply this to the expression.

We call this "extrusion" rather than "extension" because if the extension lens is bijective then it does not extend the alphabet. Nevertheless, it does have an effect because the type will be different which can be useful when converting predicates with equivalent alphabets.

**lift-definition**  $aext :: ('a, 'β) uexpr \Rightarrow ('β, 'α) lens \Rightarrow ('a, 'α) uexpr$  (**infixr**  $\oplus_p$  95)  
**is**  $\lambda P x b. P (get_x b)$  .

#### update-uexpr-rep-eq-thms

Next we prove some of the key laws. Extending an alphabet twice is equivalent to extending by the composition of the two lenses.

**lemma** *aext-twice*:  $(P \oplus_p a) \oplus_p b = P \oplus_p (a ;_L b)$   
**by** (*pred-auto*)

The bijective  $\Sigma$  lens identifies the source and view types. Thus an alphabet extension using this has no effect.

**lemma** *aext-id* [*alpha*]:  $P \oplus_p 1_L = P$   
**by** (*pred-auto*)

Literals do not depend on any variables, and thus applying an alphabet extension only alters the predicate's type, and not its valuation .

**lemma** *aext-lit* [*alpha*]:  $\ll v \gg \oplus_p a = \ll v \gg$   
**by** (*pred-auto*)

**lemma** *aext-zero* [*alpha*]:  $0 \oplus_p a = 0$   
**by** (*pred-auto*)

**lemma** *aext-one* [*alpha*]:  $1 \oplus_p a = 1$   
**by** (*pred-auto*)

**lemma** *aext-numeral* [*alpha*]:  $\text{numeral } n \oplus_p a = \text{numeral } n$   
**by** (*pred-auto*)

**lemma** *aext-true* [*alpha*]:  $\text{true} \oplus_p a = \text{true}$   
**by** (*pred-auto*)

**lemma** *aext-false* [*alpha*]:  $\text{false} \oplus_p a = \text{false}$   
**by** (*pred-auto*)

**lemma** *aext-not* [*alpha*]:  $(\neg P) \oplus_p x = (\neg (P \oplus_p x))$   
**by** (*pred-auto*)

**lemma** *aext-and* [*alpha*]:  $(P \wedge Q) \oplus_p x = (P \oplus_p x \wedge Q \oplus_p x)$   
**by** (*pred-auto*)

**lemma** *aext-or* [*alpha*]:  $(P \vee Q) \oplus_p x = (P \oplus_p x \vee Q \oplus_p x)$   
**by** (*pred-auto*)

**lemma** *aext-imp* [*alpha*]:  $(P \Rightarrow Q) \oplus_p x = (P \oplus_p x \Rightarrow Q \oplus_p x)$   
**by** (*pred-auto*)

**lemma** *aext-iff* [*alpha*]:  $(P \Leftrightarrow Q) \oplus_p x = (P \oplus_p x \Leftrightarrow Q \oplus_p x)$

**by** (*pred-auto*)

**lemma** *aext-shAll* [*alpha*]:  $(\forall x \cdot P(x)) \oplus_p a = (\forall x \cdot P(x) \oplus_p a)$   
**by** (*pred-auto*)

**lemma** *aext-event* [*alpha*]:  $(c \cdot v)_u \oplus_p a = (c \cdot v \oplus_p a)_u$   
**by** (*pred-auto*)

Alphabet extension distributes through the function liftings.

**lemma** *aext-uop* [*alpha*]:  $uop\ f\ u \oplus_p a = uop\ f\ (u \oplus_p a)$   
**by** (*pred-auto*)

**lemma** *aext-bop* [*alpha*]:  $bop\ f\ u\ v \oplus_p a = bop\ f\ (u \oplus_p a)\ (v \oplus_p a)$   
**by** (*pred-auto*)

**lemma** *aext-trop* [*alpha*]:  $trop\ f\ u\ v\ w \oplus_p a = trop\ f\ (u \oplus_p a)\ (v \oplus_p a)\ (w \oplus_p a)$   
**by** (*pred-auto*)

**lemma** *aext-qtop* [*alpha*]:  $qtop\ f\ u\ v\ w\ x \oplus_p a = qtop\ f\ (u \oplus_p a)\ (v \oplus_p a)\ (w \oplus_p a)\ (x \oplus_p a)$   
**by** (*pred-auto*)

**lemma** *aext-plus* [*alpha*]:  
 $(x + y) \oplus_p a = (x \oplus_p a) + (y \oplus_p a)$   
**by** (*pred-auto*)

**lemma** *aext-minus* [*alpha*]:  
 $(x - y) \oplus_p a = (x \oplus_p a) - (y \oplus_p a)$   
**by** (*pred-auto*)

**lemma** *aext-uminus* [*simp*]:  
 $(-x) \oplus_p a = -(x \oplus_p a)$   
**by** (*pred-auto*)

**lemma** *aext-times* [*alpha*]:  
 $(x * y) \oplus_p a = (x \oplus_p a) * (y \oplus_p a)$   
**by** (*pred-auto*)

**lemma** *aext-divide* [*alpha*]:  
 $(x / y) \oplus_p a = (x \oplus_p a) / (y \oplus_p a)$   
**by** (*pred-auto*)

Extending a variable expression over  $x$  is equivalent to composing  $x$  with the alphabet, thus effectively yielding a variable whose source is the large alphabet.

**lemma** *aext-var* [*alpha*]:  
 $var\ x \oplus_p a = var\ (x ;_L a)$   
**by** (*pred-auto*)

**lemma** *aext-ulambda* [*alpha*]:  $((\lambda x \cdot P(x)) \oplus_p a) = (\lambda x \cdot P(x) \oplus_p a)$   
**by** (*pred-auto*)

Alphabet extension is monotonic and continuous.

**lemma** *aext-mono*:  $P \sqsubseteq Q \implies P \oplus_p a \sqsubseteq Q \oplus_p a$   
**by** (*pred-auto*)

**lemma** *aext-cont* [*alpha*]:  $vwb\ lens\ a \implies (\bigsqcap A) \oplus_p a = (\bigsqcap P \in A. P \oplus_p a)$

by (*pred-simp*)

If a variable is unrestricted in a predicate, then the extended variable is unrestricted in the predicate with an alphabet extension.

**lemma** *unrest-aext* [*unrest*]:

$\llbracket \text{mwb-lens } a; x \# p \rrbracket \implies \text{unrest } (x ;_L a) (p \oplus_p a)$   
 by (*transfer*, *simp add: lens-comp-def*)

If a given variable (or alphabet)  $b$  is independent of the extension lens  $a$ , that is, it is outside the original state-space of  $p$ , then it follows that once  $p$  is extended by  $a$  then  $b$  cannot be restricted.

**lemma** *unrest-aext-indep* [*unrest*]:

$a \bowtie b \implies b \# (p \oplus_p a)$   
 by *pred-auto*

### 12.3 Alphabet Restriction

Restrict an alphabet by application of a lens that demonstrates how the smaller alphabet ( $\beta$ ) injects into the larger alphabet ( $\alpha$ ). Unlike extension, this operation can lose information if the expressions refers to variables in the larger alphabet.

**lift-definition** *arestr* :: ( $'a, 'a$ ) *uexpr*  $\Rightarrow$  ( $'\beta, 'a$ ) *lens*  $\Rightarrow$  ( $'a, 'a$ ) *uexpr* (**infixr**  $\downarrow_p$  90)  
 is  $\lambda P x b. P (\text{create}_x b)$ .

**update-uexpr-rep-eq-thms**

**lemma** *arestr-id* [*alpha*]:  $P \downarrow_p 1_L = P$   
 by (*pred-auto*)

**lemma** *arestr-aext* [*simp*]:  $\text{mwb-lens } a \implies (P \oplus_p a) \downarrow_p a = P$   
 by (*pred-auto*)

If an expression's alphabet can be divided into two disjoint sections and the expression does not depend on the second half then restricting the expression to the first half is loss-less.

**lemma** *aext-arestr* [*alpha*]:

**assumes**  $\text{mwb-lens } a \text{ bij-lens } (a +_L b) \ a \bowtie b \ b \# P$   
**shows**  $(P \downarrow_p a) \oplus_p a = P$

**proof** –

**from** *assms*(2) **have**  $1_L \subseteq_L a +_L b$

**by** (*simp add: bij-lens-equiv-id lens-equiv-def*)

**with** *assms*(1,3,4) **show** *?thesis*

**apply** (*auto simp add: id-lens-def lens-plus-def sublens-def lens-comp-def prod.case-eq-if*)

**apply** (*pred-simp*)

**apply** (*metis lens-indep-comm mwb-lens-weak weak-lens.put-get*)

**done**

**qed**

**lemma** *arestr-lit* [*alpha*]:  $\llbracket v \rrbracket \downarrow_p a = \llbracket v \rrbracket$   
 by (*pred-auto*)

**lemma** *arestr-zero* [*alpha*]:  $0 \downarrow_p a = 0$   
 by (*pred-auto*)

**lemma** *arestr-one* [*alpha*]:  $1 \downarrow_p a = 1$   
 by (*pred-auto*)



**lemma** *arestr-numeral* [*alpha*]: *numeral n*  $\vdash_p a = \text{numeral } n$   
**by** (*pred-auto*)

**lemma** *arestr-var* [*alpha*]:  
 $\text{var } x \vdash_p a = \text{var } (x \text{ /}_L a)$   
**by** (*pred-auto*)

**lemma** *arestr-true* [*alpha*]: *true*  $\vdash_p a = \text{true}$   
**by** (*pred-auto*)

**lemma** *arestr-false* [*alpha*]: *false*  $\vdash_p a = \text{false}$   
**by** (*pred-auto*)

**lemma** *arestr-not* [*alpha*]:  $(\neg P) \vdash_p a = (\neg (P \vdash_p a))$   
**by** (*pred-auto*)

**lemma** *arestr-and* [*alpha*]:  $(P \wedge Q) \vdash_p x = (P \vdash_p x \wedge Q \vdash_p x)$   
**by** (*pred-auto*)

**lemma** *arestr-or* [*alpha*]:  $(P \vee Q) \vdash_p x = (P \vdash_p x \vee Q \vdash_p x)$   
**by** (*pred-auto*)

**lemma** *arestr-imp* [*alpha*]:  $(P \Rightarrow Q) \vdash_p x = (P \vdash_p x \Rightarrow Q \vdash_p x)$   
**by** (*pred-auto*)

## 12.4 Alphabet Lens Laws

**lemma** *alpha-in-var* [*alpha*]:  $x ;_L \text{fst}_L = \text{in-var } x$   
**by** (*simp add: in-var-def*)

**lemma** *alpha-out-var* [*alpha*]:  $x ;_L \text{snd}_L = \text{out-var } x$   
**by** (*simp add: out-var-def*)

**lemma** *in-var-prod-lens* [*alpha*]:  
 $\text{wb-lens } Y \Rightarrow \text{in-var } x ;_L (X \times_L Y) = \text{in-var } (x ;_L X)$   
**by** (*simp add: in-var-def prod-as-plus lens-comp-assoc fst-lens-plus*)

**lemma** *out-var-prod-lens* [*alpha*]:  
 $\text{wb-lens } X \Rightarrow \text{out-var } x ;_L (X \times_L Y) = \text{out-var } (x ;_L Y)$   
**apply** (*simp add: out-var-def prod-as-plus lens-comp-assoc*)  
**apply** (*subst snd-lens-plus*)  
**using** *comp-wb-lens fst-vwb-lens vwb-lens-wb* **apply** *blast*  
**apply** (*simp add: alpha-in-var alpha-out-var*)  
**apply** (*simp*)  
**done**

## 12.5 Substitution Alphabet Extension

This allows us to extend the alphabet of a substitution, in a similar way to expressions.

**definition** *subst-ext* ::  $'\alpha \text{ usubst} \Rightarrow ('\alpha \Rightarrow '\beta) \Rightarrow '\beta \text{ usubst}$  (**infix**  $\oplus_s$  65) **where**  
 $[\text{upred-defs}]: \sigma \oplus_s x = (\lambda s. \text{put}_x s (\sigma (\text{get}_x s)))$

**lemma** *id-subst-ext* [*usubst*]:  
 $\text{wb-lens } x \Rightarrow \text{id} \oplus_s x = \text{id}$   
**by** *pred-auto*

**lemma** *upd-subst-ext* [*alpha*]:

$vwb\text{-}lens\ x \implies \sigma(y \mapsto_s v) \oplus_s x = (\sigma \oplus_s x)(\&x:y \mapsto_s v \oplus_p x)$

**by** *pred-auto*

**lemma** *apply-subst-ext* [*alpha*]:

$vwb\text{-}lens\ x \implies (\sigma \dagger e) \oplus_p x = (\sigma \oplus_s x) \dagger (e \oplus_p x)$

**by** (*pred-auto*)

**lemma** *axt-upred-eq* [*alpha*]:

$((e =_u f) \oplus_p a) = ((e \oplus_p a) =_u (f \oplus_p a))$

**by** (*pred-auto*)

## 12.6 Substitution Alphabet Restriction

This allows us to reduce the alphabet of a substitution, in a similar way to expressions.

**definition** *subst-res* ::  $'\alpha\ usubst \Rightarrow (' \beta \implies '\alpha) \Rightarrow '\beta\ usubst$  (**infix**  $\downarrow_s$  65) **where**

[*upred-defs*]:  $\sigma \downarrow_s x = (\lambda s. get_x(\sigma(create_x s)))$

**lemma** *id-subst-res* [*usubst*]:

$mwb\text{-}lens\ x \implies id \downarrow_s x = id$

**by** *pred-auto*

**lemma** *upd-subst-res* [*alpha*]:

$mwb\text{-}lens\ x \implies \sigma(\&x:y \mapsto_s v) \downarrow_s x = (\sigma \downarrow_s x)(\&y \mapsto_s v \downarrow_p x)$

**by** (*pred-auto*)

**lemma** *subst-ext-res* [*usubst*]:

$mwb\text{-}lens\ x \implies (\sigma \oplus_s x) \downarrow_s x = \sigma$

**by** (*pred-auto*)

**lemma** *unrest-subst-alpha-ext* [*unrest*]:

$x \bowtie y \implies x \# (P \oplus_s y)$

**by** (*pred-simp robust, metis lens-indep-def*)

**end**

## 13 Lifting expressions

**theory** *utp-lift*

**imports**

*utp-alphabet*

**begin**

### 13.1 Lifting definitions

We define operators for converting an expression to and from a relational state space with the help of alphabet extrusion and restriction. In general throughout Isabelle/UTP we adopt the notation  $\lceil P \rceil$  with some subscript to denote lifting an expression into a larger alphabet, and  $\lfloor P \rfloor$  for dropping into a smaller alphabet.

The following two functions lift and drop an expression, respectively, whose alphabet is  $'\alpha$ , into a product alphabet  $'\alpha \times '\beta$ . This allows us to deal with expressions which refer only to undashed variables, and use the type-system to ensure this.

**abbreviation** *lift-pre* ::  $('a, '\alpha)\ uexpr \Rightarrow ('a, '\alpha \times '\beta)\ uexpr$  ( $\lceil \_ \rceil_{<}$ )

**where**  $\lceil P \rceil_{<} \equiv P \oplus_p fst_L$

**abbreviation**  $drop\text{-}pre :: ('a, 'α \times 'β) \text{ ue}xpr \Rightarrow ('a, 'α) \text{ ue}xpr (\lceil - \rceil_{<})$

**where**  $\lceil P \rceil_{<} \equiv P \downarrow_p fst_L$

The following two functions lift and drop an expression, respectively, whose alphabet is  $'β$ , into a product alphabet  $'α \times 'β$ . This allows us to deal with expressions which refer only to dashed variables.

**abbreviation**  $lift\text{-}post :: ('a, 'β) \text{ ue}xpr \Rightarrow ('a, 'α \times 'β) \text{ ue}xpr (\lceil - \rceil_{>})$

**where**  $\lceil P \rceil_{>} \equiv P \oplus_p snd_L$

**abbreviation**  $drop\text{-}post :: ('a, 'α \times 'β) \text{ ue}xpr \Rightarrow ('a, 'β) \text{ ue}xpr (\lceil - \rceil_{>})$

**where**  $\lceil P \rceil_{>} \equiv P \downarrow_p snd_L$

**abbreviation**  $lift\text{-}cond\text{-}pre (\lceil - \rceil_{\leftarrow})$  **where**  $\lceil P \rceil_{\leftarrow} \equiv P \oplus_p (1_L \times_L 0_L)$

**abbreviation**  $lift\text{-}cond\text{-}post (\lceil - \rceil_{\rightarrow})$  **where**  $\lceil P \rceil_{\rightarrow} \equiv P \oplus_p (0_L \times_L 1_L)$

**abbreviation**  $drop\text{-}cond\text{-}pre (\lceil - \rceil_{\leftarrow})$  **where**  $\lceil P \rceil_{\leftarrow} \equiv P \downarrow_p (1_L \times_L 0_L)$

**abbreviation**  $drop\text{-}cond\text{-}post (\lceil - \rceil_{\rightarrow})$  **where**  $\lceil P \rceil_{\rightarrow} \equiv P \downarrow_p (0_L \times_L 1_L)$

## 13.2 Lifting Laws

With the help of our alphabet laws, we can prove some intuitive laws about alphabet lifting. For example, lifting variables yields an unprimed or primed relational variable expression, respectively.

**lemma**  $lift\text{-}pre\text{-}var [simp]$ :

$\lceil var\ x \rceil_{<} = \$x$

**by** ( $alpha\text{-}tac$ )

**lemma**  $lift\text{-}post\text{-}var [simp]$ :

$\lceil var\ x \rceil_{>} = \$x'$

**by** ( $alpha\text{-}tac$ )

**lemma**  $lift\text{-}cond\text{-}pre\text{-}var [simp]$ :

$\lceil \$x \rceil_{\leftarrow} = \$x$

**by** ( $pred\text{-}auto$ )

**lemma**  $lift\text{-}cond\text{-}post\text{-}var [simp]$ :

$\lceil \$x' \rceil_{\rightarrow} = \$x'$

**by** ( $pred\text{-}auto$ )

## 13.3 Substitution Laws

**lemma**  $pre\text{-}var\text{-}subst [usubst]$ :

$\sigma(\$x \mapsto_s \ll v \gg) \dagger \lceil P \rceil_{<} = \sigma \dagger \lceil P[\ll v \gg / \&x] \rceil_{<}$

**by** ( $pred\text{-}simp$ )

## 13.4 Unrestriction laws

Crucially, the lifting operators allow us to demonstrate unrestricted properties. For example, we can show that no primed variable is restricted in an expression over only the first element of the state-space product type.

**lemma**  $unrest\text{-}dash\text{-}var\text{-}pre [unrest]$ :

```

fixes  $x :: ('a \Longrightarrow 'α)$ 
shows  $\$x' \# \lceil p \rceil <$ 
by (pred-auto)

```

```

lemma unrest-dash-var-cond-pre [unrest]:
  fixes  $x :: ('a \Longrightarrow 'α)$ 
  shows  $\$x' \# \lceil P \rceil \leftarrow$ 
  by (pred-auto)
end

```

## 14 Alphabetised Relations

```

theory utp-rel
imports
  utp-pred-laws
  utp-recursion
  utp-lift
  utp-tactics
begin

```

An alphabetised relation is simply a predicate whose state-space is a product type. In this theory we construct the core operators of the relational calculus, and prove a library of associated theorems, based on Chapters 2 and 5 of the UTP book [5].

### 14.1 Relational Alphabets

We set up convenient syntax to refer to the input and output parts of the alphabet, as is common in UTP. Since we are in a product space, these are simply the lenses  $fst_L$  and  $snd_L$ .

```

definition  $inα :: ('α \Longrightarrow 'α \times 'β)$  where
[lens-defs]:  $inα = fst_L$ 

```

```

definition  $outα :: ('β \Longrightarrow 'α \times 'β)$  where
[lens-defs]:  $outα = snd_L$ 

```

```

lemma inα-uvar [simp]: vwb-lens  $inα$ 
by (unfold-locales, auto simp add: inα-def)

```

```

lemma outα-uvar [simp]: vwb-lens  $outα$ 
by (unfold-locales, auto simp add: outα-def)

```

```

lemma var-in-alpha [simp]:  $x ;_L inα = ivar\ x$ 
by (simp add: fst-lens-def inα-def in-var-def)

```

```

lemma var-out-alpha [simp]:  $x ;_L outα = ovar\ x$ 
by (simp add: outα-def out-var-def snd-lens-def)

```

```

lemma drop-pre-inv [simp]:  $\llbracket outα \# p \rrbracket \Longrightarrow \lceil \lfloor p \rfloor \rfloor < = p$ 
by (pred-simp)

```

```

lemma usubst-lookup-ivar-unrest [usubst]:
   $inα \# \sigma \Longrightarrow \langle \sigma \rangle_s (ivar\ x) = \$x$ 
by (rel-simp, metis fstI)

```

```

lemma usubst-lookup-ovar-unrest [usubst]:

```

$out\alpha \# \sigma \implies \langle \sigma \rangle_s (ovar\ x) = \$x'$   
**by** (*rel-simp*, *metis sndI*)

**lemma** *out-alpha-in-indep* [*simp*]:  
 $out\alpha \bowtie in-var\ x\ in-var\ x \bowtie out\alpha$   
**by** (*simp-all add: in-var-def out $\alpha$ -def lens-indep-def fst-lens-def snd-lens-def lens-comp-def*)

**lemma** *in-alpha-out-indep* [*simp*]:  
 $in\alpha \bowtie out-var\ x\ out-var\ x \bowtie in\alpha$   
**by** (*simp-all add: in-var-def in $\alpha$ -def lens-indep-def fst-lens-def lens-comp-def*)

The following two functions lift a predicate substitution to a relational one.

**abbreviation** *usubst-rel-lift* ::  $'\alpha\ usubst \Rightarrow ('\alpha \times '\beta)\ usubst\ ([\_ ]_s)$  **where**  
 $[\sigma]_s \equiv \sigma \oplus_s in\alpha$

**abbreviation** *usubst-rel-drop* ::  $(''\alpha \times '\alpha)\ usubst \Rightarrow '\alpha\ usubst\ ([\_ ]_s)$  **where**  
 $[\sigma]_s \equiv \sigma \upharpoonright_s in\alpha$

The alphabet of a relation then consists wholly of the input and output portions.

**lemma** *alpha-in-out*:  
 $\Sigma \approx_L in\alpha +_L out\alpha$   
**by** (*simp add: fst-snd-id-lens in $\alpha$ -def lens-equiv-refl out $\alpha$ -def*)

## 14.2 Relational Types and Operators

We create type synonyms for conditions (which are simply predicates) – i.e. relations without dashed variables –, alphabetised relations where the input and output alphabet can be different, and finally homogeneous relations.

**type-synonym**  $'\alpha\ cond = '\alpha\ upred$   
**type-synonym**  $(''\alpha, '\beta)\ rel = (''\alpha \times '\beta)\ upred$   
**type-synonym**  $'\alpha\ hrel = (''\alpha \times '\alpha)\ upred$   
**type-synonym**  $('a, '\alpha)\ hexpr = ('a, '\alpha \times '\alpha)\ uexpr$

### translations

$(type)\ (''\alpha, '\beta)\ rel \leq (type)\ (''\alpha \times '\beta)\ upred$

We set up some overloaded constants for sequential composition and the identity in case we want to overload their definitions later.

### consts

*useq* ::  $'a \Rightarrow 'b \Rightarrow 'c$  (**infixr** ;; 71)  
*uskip* ::  $'a$  (*II*)

We define a specialised version of the conditional where the condition can refer only to undashed variables, as is usually the case in programs, but not universally in UTP models. We implement this by lifting the condition predicate into the relational state-space with construction  $[b]_{<}$ .

### abbreviation

*rcond* ::  $(''\alpha, '\beta)\ rel \Rightarrow '\alpha\ cond \Rightarrow (''\alpha, '\beta)\ rel \Rightarrow (''\alpha, '\beta)\ rel$   
 $((\beta \triangleleft \_ \triangleright_r \_) [52, 0, 53] 52)$   
**where**  $(P \triangleleft b \triangleright_r Q) \equiv (P \triangleleft [b]_{<} \triangleright Q)$

Sequential composition is heterogeneous, and simply requires that the output alphabet of the first matches then input alphabet of the second. We define it by lifting HOL's built-in relational composition operator (*op O*). Since this returns a set, the definition states that the state binding *b* is an element of this set.

**lift-definition**  $seqr :: ('\alpha, '\beta) \text{rel} \Rightarrow (''\beta, '\gamma) \text{rel} \Rightarrow (''\alpha \times '\gamma) \text{upred}$   
**is**  $\lambda P Q b. b \in (\{p. P p\} \cap \{q. Q q\})$  .

#### ad hoc-overloading

*useq seqr*

We also set up a homogeneous sequential composition operator, and versions of *true* and *false* that are explicitly typed by a homogeneous alphabet.

**abbreviation**  $seqh :: '\alpha \text{hrel} \Rightarrow '\alpha \text{hrel} \Rightarrow '\alpha \text{hrel}$  (**infixr**  $::_h$  71) **where**  
 $seqh P Q \equiv (P :: Q)$

**abbreviation**  $truer :: '\alpha \text{hrel}$  ( $true_h$ ) **where**  
 $truer \equiv true$

**abbreviation**  $falsr :: '\alpha \text{hrel}$  ( $false_h$ ) **where**  
 $falsr \equiv false$

We define the relational converse operator as an alphabet extrusion on the bijective lens  $swap_L$  that swaps the elements of the product state-space.

**abbreviation**  $conv-r :: ('a, '\alpha \times '\beta) \text{uepr} \Rightarrow ('a, '\beta \times '\alpha) \text{uepr}$  ( $-^{\sim}$  [999] 999)  
**where**  $conv-r e \equiv e \oplus_p swap_L$

Assignment is defined using substitutions, where latter defines what each variable should map to. The definition of the operator identifies the after state binding,  $b'$ , with the substitution function applied to the before state binding  $b$ .

**lift-definition**  $assigns-r :: '\alpha \text{usubst} \Rightarrow '\alpha \text{hrel}$  ( $\langle \cdot \rangle_a$ )  
**is**  $\lambda \sigma (b, b'). b' = \sigma(b)$  .

Relational identity, or skip, is then simply an assignment with the identity substitution: it simply identifies all variables.

**definition**  $skip-r :: '\alpha \text{hrel}$  **where**  
 $[urel-defs]: skip-r = assigns-r id$

#### ad hoc-overloading

*uskip skip-r*

We set up iterated sequential composition which iterates an indexed predicate over the elements of a list.

**definition**  $seqr-iter :: 'a \text{list} \Rightarrow ('a \Rightarrow '\alpha \text{hrel}) \Rightarrow '\alpha \text{hrel}$  **where**  
 $[urel-defs]: seqr-iter xs P = foldr (\lambda i Q. P(i) :: Q) xs II$

A singleton assignment simply applies a singleton substitution function, and similarly for a double assignment.

**abbreviation**  $assign-r :: ('t \Rightarrow '\alpha) \Rightarrow ('t, '\alpha) \text{uepr} \Rightarrow '\alpha \text{hrel}$   
**where**  $assign-r x v \equiv \langle [x \mapsto_s v] \rangle_a$

**abbreviation**  $assign-2-r ::$   
 $(t1 \Rightarrow '\alpha) \Rightarrow (t2 \Rightarrow '\alpha) \Rightarrow (t1, '\alpha) \text{uepr} \Rightarrow (t2, '\alpha) \text{uepr} \Rightarrow '\alpha \text{hrel}$   
**where**  $assign-2-r x y u v \equiv assigns-r [x \mapsto_s u, y \mapsto_s v]$

We also define the alphabetised skip operator that identifies all input and output variables in the given alphabet lens. All other variables are unrestricted. We also set up syntax for it.

**definition**  $skip-ra :: ('\beta, '\alpha) \text{lens} \Rightarrow '\alpha \text{hrel}$  **where**

[urel-defs]:  $\text{skip-ra } v = (\$v' =_u \$v)$

Similarly, we define the alphabetised assignment operator.

**definition**  $\text{assigns-ra} :: 'a \text{ usubst} \Rightarrow ('b, 'a) \text{ lens} \Rightarrow 'a \text{ hrel } (\langle \cdot \rangle_-)$  **where**  
 $\langle \sigma \rangle_a = ([\sigma]_s \uparrow \text{skip-ra } a)$

Assumptions ( $c^\top$ ) and assertions ( $c_\perp$ ) are encoded as conditionals. An assumption behaves like skip if the condition is true, and otherwise behaves like *false* (miracle). An assertion is the same, but yields *true*, which is an abort.

**definition**  $\text{rassume} :: 'a \text{ upred} \Rightarrow 'a \text{ hrel } ([\cdot]^\top)$  **where**  
 $[\text{urel-defs}]: \text{rassume } c = II \triangleleft c \triangleright_r \text{false}$

**definition**  $\text{rassert} :: 'a \text{ upred} \Rightarrow 'a \text{ hrel } (\{\cdot\}_\perp)$  **where**  
 $[\text{urel-defs}]: \text{rassert } c = II \triangleleft c \triangleright_r \text{true}$

A test is like a precondition, except that it identifies to the postcondition, and is thus a refinement of  $II$ . It forms the basis for Kleene Algebra with Tests [7, 1] (KAT), which embeds a Boolean algebra into a Kleene algebra to represent conditions.

**definition**  $\text{lift-test} :: 'a \text{ cond} \Rightarrow 'a \text{ hrel } ([\cdot]_t)$   
**where**  $[\text{urel-defs}]: [\cdot]_t = ([\cdot]_< \wedge II)$

We define two variants of while loops based on strongest and weakest fixed points. The former is *false* for an infinite loop, and the latter is *true*.

**definition**  $\text{while} :: 'a \text{ cond} \Rightarrow 'a \text{ hrel} \Rightarrow 'a \text{ hrel } (\text{while}^\top - \text{do} - \text{od})$  **where**  
 $[\text{urel-defs}]: \text{while}^\top b \text{ do } P \text{ od} = (\nu X \cdot (P ;; X) \triangleleft b \triangleright_r II)$

**abbreviation**  $\text{while-top} :: 'a \text{ cond} \Rightarrow 'a \text{ hrel} \Rightarrow 'a \text{ hrel } (\text{while} - \text{do} - \text{od})$  **where**  
 $\text{while } b \text{ do } P \text{ od} \equiv \text{while}^\top b \text{ do } P \text{ od}$

**definition**  $\text{while-bot} :: 'a \text{ cond} \Rightarrow 'a \text{ hrel} \Rightarrow 'a \text{ hrel } (\text{while}_\perp - \text{do} - \text{od})$  **where**  
 $[\text{urel-defs}]: \text{while}_\perp b \text{ do } P \text{ od} = (\mu X \cdot (P ;; X) \triangleleft b \triangleright_r II)$

While loops with invariant decoration (cf. [1]).

**definition**  $\text{while-inv} :: 'a \text{ cond} \Rightarrow 'a \text{ cond} \Rightarrow 'a \text{ hrel} \Rightarrow 'a \text{ hrel } (\text{while} - \text{invar} - \text{do} - \text{od})$  **where**  
 $[\text{urel-defs}]: \text{while } b \text{ invar } p \text{ do } S \text{ od} = \text{while } b \text{ do } S \text{ od}$

We implement a poor man's version of alphabet restriction that hides a variable within a relation.

**definition**  $\text{rel-var-res} :: 'a \text{ hrel} \Rightarrow ('a \Rightarrow 'a) \Rightarrow 'a \text{ hrel } (\text{infix } \downarrow_\alpha \text{ } 80)$  **where**  
 $[\text{urel-defs}]: P \downarrow_\alpha x = (\exists \$x \cdot \exists \$x' \cdot P)$

We next describe frames and antiframes with the help of lenses. A frame states that  $P$  defines the behaviour of all variables not in  $a$ , and all those in  $a$  remain the same. An antiframe describes the converse: all variables in  $a$  are specified by  $P$ , and all other variables remain the same. For more information please see [8].

**definition**  $\text{frame} :: ('a \Rightarrow 'a) \Rightarrow 'a \text{ hrel} \Rightarrow 'a \text{ hrel}$  **where**  
 $[\text{urel-defs}]: \text{frame } a P = (\exists st \cdot P[\ll st \gg / \$\mathbf{v}' ] \wedge \$\mathbf{v}' =_u \ll st \gg \oplus \$\mathbf{v} \text{ on } \&a)$

**definition**  $\text{antiframe} :: ('a \Rightarrow 'a) \Rightarrow 'a \text{ hrel} \Rightarrow 'a \text{ hrel}$  **where**  
 $[\text{urel-defs}]: \text{antiframe } a P = (\exists st \cdot P[\ll st \gg / \$\mathbf{v}' ] \wedge \$\mathbf{v}' =_u \$\mathbf{v} \oplus \ll st \gg \text{ on } \&a)$

The nameset operator can be used to hide a portion of the after-state that lies outside the lens  $a$ . It can be useful to partition a relation's variables in order to conjoin it with another relation.

**definition** *nameset* :: ('a  $\Rightarrow$  'α)  $\Rightarrow$  'α hrel  $\Rightarrow$  'α hrel **where**  
[urel-defs]: *nameset* a P = (P  $\vdash_v$  {'v,\$a'})

### 14.3 Syntax Translations

#### syntax

- Alternative traditional conditional syntax
- utp-if :: *logic*  $\Rightarrow$  *logic*  $\Rightarrow$  *logic*  $\Rightarrow$  *logic* ((if<sub>u</sub> (-)/ then (-)/ else (-)) [0, 0, 71] 71)
- Iterated sequential composition
- seqr-iter :: *pttrn*  $\Rightarrow$  'a list  $\Rightarrow$  'σ hrel  $\Rightarrow$  'σ hrel ((3;; - : - •/ -) [0, 0, 10] 10)
- Single and multiple assignement
- assignment :: *svids*  $\Rightarrow$  *uexprs*  $\Rightarrow$  'α hrel ('(-) := '(-))
- assignment :: *svids*  $\Rightarrow$  *uexprs*  $\Rightarrow$  'α hrel (**infixr** := 72)
- Indexed assignment
- assignment-upd :: *svid*  $\Rightarrow$  *logic*  $\Rightarrow$  *logic*  $\Rightarrow$  *logic* (([-] :=/ -) [73, 0, 0] 72)
- Substitution constructor
- mk-usubst :: *svids*  $\Rightarrow$  *uexprs*  $\Rightarrow$  'α usubst
- Alphabetised skip
- skip-ra :: *salpha*  $\Rightarrow$  *logic* (II.)
- Frame
- frame :: *salpha*  $\Rightarrow$  *logic*  $\Rightarrow$  *logic* (:-[] [79,0] 80)
- Antiframe
- antiframe :: *salpha*  $\Rightarrow$  *logic*  $\Rightarrow$  *logic* (:-[] [99,0] 100)
- Nameset
- nameset :: *salpha*  $\Rightarrow$  *logic*  $\Rightarrow$  *logic* (ns - - - [0,999] 999)

#### translations

- utp-if b P Q  $\Rightarrow$  P  $\triangleleft$  b  $\triangleright_r$  Q
- ;; x : l • P  $\Rightarrow$  (CONST seqr-iter) l (λx. P)
- mk-usubst σ (-svid-unit x) v  $\Rightarrow$  σ(&x  $\mapsto_s$  v)
- mk-usubst σ (-svid-list x xs) (-uexprs v vs)  $\Rightarrow$  (-mk-usubst (σ(&x  $\mapsto_s$  v)) xs vs)
- assignment xs vs  $\Rightarrow$  CONST assigns-r (-mk-usubst (CONST id) xs vs)
- x := v <= CONST assigns-r (CONST subst-upd (CONST id) (CONST svar x) v)
- x := v <= CONST assigns-r (CONST subst-upd (CONST id) x v)
- x,y := u,v <= CONST assigns-r (CONST subst-upd (CONST subst-upd (CONST id) (CONST svar x) u) (CONST svar y) v)
- Indexed assignment uses the overloaded collection update function *uupd*.
- x [k] := v  $\Rightarrow$  x := &x(k  $\mapsto$  v)<sub>u</sub>
- skip-ra v  $\Rightarrow$  CONST skip-ra v
- frame x P  $\Rightarrow$  CONST frame x P
- antiframe x P  $\Rightarrow$  CONST antiframe x P
- nameset x P  $\Rightarrow$  CONST nameset x P

The following code sets up pretty-printing for homogeneous relational expressions. We cannot do this via the “translations” command as we only want the rule to apply when the input and output alphabet types are the same. The code has to deconstruct a ('a, 'α) *uexpr* type, determine that it is relational (product alphabet), and then checks if the types *alpha* and *beta* are the same. If they are, the type is printed as a *hexpr*. Otherwise, we have no match. We then set up a regular translation for the *hrel* type that uses this.

#### print-translation <<

```
let
fun tr' ctx [ a
    , Const (@{type-syntax prod},-) $ alpha $ beta ] =
  if (alpha = beta)
  then Syntax.const @ {type-syntax hexpr} $ a $ alpha
```



```

    else raise Match;
in [(@{type-syntax uexpr},tr')]
end
>>

```

#### translations

(type)  $'\alpha$  hrel  $\leq$  (type) (bool,  $'\alpha$ ) hexpr

### 14.4 Relation Properties

We describe some properties of relations, including functional and injective relations.

**definition** *ufunctional* :: ( $'a$ ,  $'b$ ) rel  $\Rightarrow$  bool  
**where** [*urel-defs*]: *ufunctional*  $R \longleftrightarrow II \sqsubseteq R^-$  ;;  $R$

**definition** *uinj* :: ( $'a$ ,  $'b$ ) rel  $\Rightarrow$  bool  
**where** [*urel-defs*]: *uinj*  $R \longleftrightarrow II \sqsubseteq R$  ;;  $R^-$

— Configuration for UTP tactics (see *utp-tactics*).

**update-uexpr-rep-eq-thms** — Reread *rep-eq* theorems.

### 14.5 Unrestriction Laws

**lemma** *unrest-iuvar* [*unrest*]:  $out\alpha \# \$x$   
**by** (*metis fst-snd-lens-indep lift-pre-var out $\alpha$ -def unrest-aext-indep*)

**lemma** *unrest-ouvar* [*unrest*]:  $in\alpha \# \$x'$   
**by** (*metis in $\alpha$ -def lift-post-var snd-fst-lens-indep unrest-aext-indep*)

**lemma** *unrest-semir-undash* [*unrest*]:  
**fixes**  $x :: ('a \Rightarrow '\alpha)$   
**assumes**  $\$x \# P$   
**shows**  $\$x \# P$  ;;  $Q$   
**using** *assms* **by** (*rel-auto*)

**lemma** *unrest-semir-dash* [*unrest*]:  
**fixes**  $x :: ('a \Rightarrow '\alpha)$   
**assumes**  $\$x' \# Q$   
**shows**  $\$x' \# P$  ;;  $Q$   
**using** *assms* **by** (*rel-auto*)

**lemma** *unrest-cond* [*unrest*]:  
 $\llbracket x \# P; x \# b; x \# Q \rrbracket \Longrightarrow x \# P \triangleleft b \triangleright Q$   
**by** (*rel-auto*)

**lemma** *unrest-in $\alpha$ -var* [*unrest*]:  
 $\llbracket mwb\text{-}lens\ x; in\alpha \# (P :: ('a, ('\alpha \times '\beta))\ uexpr) \rrbracket \Longrightarrow \$x \# P$   
**by** (*rel-auto*)

**lemma** *unrest-out $\alpha$ -var* [*unrest*]:  
 $\llbracket mwb\text{-}lens\ x; out\alpha \# (P :: ('a, ('\alpha \times '\beta))\ uexpr) \rrbracket \Longrightarrow \$x' \# P$   
**by** (*rel-auto*)

**lemma** *unrest-pre-out $\alpha$*  [*unrest*]:  $out\alpha \# \lceil b \rceil_<$   
**by** (*transfer, auto simp add: out $\alpha$ -def*)

**lemma** *unrest-post-in $\alpha$*  [*unrest*]:  $in\alpha \# [b]_>$   
**by** (*transfer*, *auto simp add: in $\alpha$ -def*)

**lemma** *unrest-pre-in-var* [*unrest*]:  
 $x \# p1 \implies \$x \# [p1]_<$   
**by** (*transfer*, *simp*)

**lemma** *unrest-post-out-var* [*unrest*]:  
 $x \# p1 \implies \$x' \# [p1]_>$   
**by** (*transfer*, *simp*)

**lemma** *unrest-convr-out $\alpha$*  [*unrest*]:  
 $in\alpha \# p \implies out\alpha \# p^-$   
**by** (*transfer*, *auto simp add: lens-defs*)

**lemma** *unrest-convr-in $\alpha$*  [*unrest*]:  
 $out\alpha \# p \implies in\alpha \# p^-$   
**by** (*transfer*, *auto simp add: lens-defs*)

**lemma** *unrest-in-rel-var-res* [*unrest*]:  
 $vwb\text{-}lens\ x \implies \$x \# (P \upharpoonright_\alpha x)$   
**by** (*simp add: rel-var-res-def unrest*)

**lemma** *unrest-out-rel-var-res* [*unrest*]:  
 $vwb\text{-}lens\ x \implies \$x' \# (P \upharpoonright_\alpha x)$   
**by** (*simp add: rel-var-res-def unrest*)

**lemma** *unrest-out-alpha-usubst-rel-lift* [*unrest*]:  
 $out\alpha \# [\sigma]_s$   
**by** (*rel-auto*)

## 14.6 Substitution laws

**lemma** *subst-seq-left* [*usubst*]:  
 $out\alpha \# \sigma \implies \sigma \dagger (P ;; Q) = (\sigma \dagger P) ;; Q$   
**by** (*rel-simp*, (*metis* (*no-types*, *lifting*) *Pair-inject surjective-pairing*)+)

**lemma** *subst-seq-right* [*usubst*]:  
 $in\alpha \# \sigma \implies \sigma \dagger (P ;; Q) = P ;; (\sigma \dagger Q)$   
**by** (*rel-simp*, (*metis* (*no-types*, *lifting*) *Pair-inject surjective-pairing*)+)

The following laws support substitution in heterogeneous relations for polymorphically typed literal expressions. These cannot be supported more generically due to limitations in HOL's type system. The laws are presented in a slightly strange way so as to be as general as possible.

**lemma** *bool-seqr-laws* [*usubst*]:  
**fixes**  $x :: (bool \implies 'a)$   
**shows**  
 $\bigwedge P\ Q\ \sigma. \sigma(\$x \mapsto_s true) \dagger (P ;; Q) = \sigma \dagger (P[true/\$x] ;; Q)$   
 $\bigwedge P\ Q\ \sigma. \sigma(\$x \mapsto_s false) \dagger (P ;; Q) = \sigma \dagger (P[false/\$x] ;; Q)$   
 $\bigwedge P\ Q\ \sigma. \sigma(\$x' \mapsto_s true) \dagger (P ;; Q) = \sigma \dagger (P ;; Q[true/\$x'])$   
 $\bigwedge P\ Q\ \sigma. \sigma(\$x' \mapsto_s false) \dagger (P ;; Q) = \sigma \dagger (P ;; Q[false/\$x'])$   
**by** (*rel-auto*)+

**lemma** *zero-one-seqr-laws* [*usubst*]:

**fixes**  $x :: (- \Longrightarrow ' \alpha)$

**shows**

$\bigwedge P Q \sigma. \sigma(\$x \mapsto_s 0) \dagger (P ;; Q) = \sigma \dagger (P \llbracket 0/\$x \rrbracket ;; Q)$   
 $\bigwedge P Q \sigma. \sigma(\$x \mapsto_s 1) \dagger (P ;; Q) = \sigma \dagger (P \llbracket 1/\$x \rrbracket ;; Q)$   
 $\bigwedge P Q \sigma. \sigma(\$x' \mapsto_s 0) \dagger (P ;; Q) = \sigma \dagger (P ;; Q \llbracket 0/\$x' \rrbracket)$   
 $\bigwedge P Q \sigma. \sigma(\$x' \mapsto_s 1) \dagger (P ;; Q) = \sigma \dagger (P ;; Q \llbracket 1/\$x' \rrbracket)$   
**by** (*rel-auto*)<sup>+</sup>

**lemma** *numeral-seqr-laws* [*usubst*]:

**fixes**  $x :: (- \Longrightarrow ' \alpha)$

**shows**

$\bigwedge P Q \sigma. \sigma(\$x \mapsto_s \text{numeral } n) \dagger (P ;; Q) = \sigma \dagger (P \llbracket \text{numeral } n/\$x \rrbracket ;; Q)$   
 $\bigwedge P Q \sigma. \sigma(\$x' \mapsto_s \text{numeral } n) \dagger (P ;; Q) = \sigma \dagger (P ;; Q \llbracket \text{numeral } n/\$x' \rrbracket)$   
**by** (*rel-auto*)<sup>+</sup>

**lemma** *usubst-condr* [*usubst*]:

$\sigma \dagger (P \triangleleft b \triangleright Q) = (\sigma \dagger P \triangleleft \sigma \dagger b \triangleright \sigma \dagger Q)$

**by** (*rel-auto*)

**lemma** *subst-skip-r* [*usubst*]:

$\text{out}\alpha \# \sigma \Longrightarrow \sigma \dagger II = \langle \lfloor \sigma \rfloor_s \rangle_a$

**by** (*rel-simp*, (*metis* (*mono-tags*, *lifting*) *prod.sel*(1) *sndI* *surjective-pairing*))<sup>+</sup>

**lemma** *subst-pre-skip* [*usubst*]:  $\lceil \sigma \rceil_s \dagger II = \langle \sigma \rangle_a$

**by** (*rel-auto*)

**lemma** *usubst-upd-in-comp* [*usubst*]:

$\sigma(\&\text{in}\alpha:x \mapsto_s v) = \sigma(\$x \mapsto_s v)$

**by** (*simp* *add*: *pr-var-def fst-lens-def in $\alpha$ -def in-var-def*)

**lemma** *usubst-upd-out-comp* [*usubst*]:

$\sigma(\&\text{out}\alpha:x \mapsto_s v) = \sigma(\$x' \mapsto_s v)$

**by** (*simp* *add*: *pr-var-def out $\alpha$ -def out-var-def snd-lens-def*)

**lemma** *subst-lift-upd* [*alpha*]:

**fixes**  $x :: ('a \Longrightarrow ' \alpha)$

**shows**  $\lceil \sigma(x \mapsto_s v) \rceil_s = \lceil \sigma \rceil_s(\$x \mapsto_s \lceil v \rceil_<)$

**by** (*simp* *add*: *alpha usubst*, *simp* *add*: *pr-var-def fst-lens-def in $\alpha$ -def in-var-def*)

**lemma** *subst-drop-upd* [*alpha*]:

**fixes**  $x :: ('a \Longrightarrow ' \alpha)$

**shows**  $\lfloor \sigma(\$x \mapsto_s v) \rfloor_s = \lfloor \sigma \rfloor_s(x \mapsto_s \lfloor v \rfloor_<)$

**by** *pred-simp*

**lemma** *subst-lift-pre* [*usubst*]:  $\lceil \sigma \rceil_s \dagger \lceil b \rceil_< = \lceil \sigma \dagger b \rceil_<$

**by** (*metis* *apply-subst-ext fst-vwb-lens in $\alpha$ -def*)

**lemma** *unrest-usubst-lift-in* [*unrest*]:

$x \# P \Longrightarrow \$x \# \lceil P \rceil_s$

**by** *pred-simp*

**lemma** *unrest-usubst-lift-out* [*unrest*]:

**fixes**  $x :: ('a \Longrightarrow ' \alpha)$

**shows**  $\$x' \# \lfloor P \rfloor_s$

**by** *pred-simp*

## 14.7 Alphabet laws

**lemma** *aext-cond* [*alpha*]:

$(P \triangleleft b \triangleright Q) \oplus_p a = ((P \oplus_p a) \triangleleft (b \oplus_p a) \triangleright (Q \oplus_p a))$   
**by** (*rel-auto*)

**lemma** *aext-seq* [*alpha*]:

$wb\text{-}lens\ a \implies ((P ;; Q) \oplus_p (a \times_L a)) = ((P \oplus_p (a \times_L a)) ;; (Q \oplus_p (a \times_L a)))$   
**by** (*rel-simp*, *metis wb-lens-weak weak-lens.put-get*)

## 14.8 Relational unrestriction

Relational unrestriction states that a variable is unchanged by a relation. Eventually I'd also like to have it state that the relation also does not depend on the variable's initial value, but I'm not sure how to state that yet. For now we represent this by the parametric healthiness condition *RID*.

**definition** *RID* ::  $('a \implies 'a) \Rightarrow 'a \text{ hrel} \Rightarrow 'a \text{ hrel}$

**where**  $RID\ x\ P = ((\exists\ \$x \cdot \exists\ \$x' \cdot P) \wedge \$x' =_u \$x)$

**declare** *RID-def* [*urel-defs*]

**lemma** *RID-idem*:

$mwb\text{-}lens\ x \implies RID(x)(RID(x)(P)) = RID(x)(P)$   
**by** (*rel-auto*)

**lemma** *RID-mono*:

$P \sqsubseteq Q \implies RID(x)(P) \sqsubseteq RID(x)(Q)$   
**by** (*rel-auto*)

**lemma** *RID-skip-r*:

$vwb\text{-}lens\ x \implies RID(x)(II) = II$   
**apply** (*rel-auto*) **using** *vwb-lens.put-eq* **by** *fastforce*

**lemma** *RID-disj*:

$RID(x)(P \vee Q) = (RID(x)(P) \vee RID(x)(Q))$   
**by** (*rel-auto*)

**lemma** *RID-conj*:

$vwb\text{-}lens\ x \implies RID(x)(RID(x)(P) \wedge RID(x)(Q)) = (RID(x)(P) \wedge RID(x)(Q))$   
**by** (*rel-auto*)

**lemma** *RID-assigns-r-diff*:

$\llbracket vwb\text{-}lens\ x; x \# \sigma \rrbracket \implies RID(x)(\langle \sigma \rangle_a) = \langle \sigma \rangle_a$   
**apply** (*rel-auto*)  
**apply** (*metis vwb-lens.put-eq*)  
**apply** (*metis vwb-lens-wb wb-lens.get-put wb-lens-weak weak-lens.put-get*)

**done**

**lemma** *RID-assign-r-same*:

$vwb\text{-}lens\ x \implies RID(x)(x := v) = II$   
**apply** (*rel-auto*)  
**using** *vwb-lens.put-eq* **apply** *fastforce*

**done**

**lemma** *RID-seq-left*:

```

assumes vwb-lens x
shows  $RID(x)(RID(x)(P) ;; Q) = (RID(x)(P) ;; RID(x)(Q))$ 
proof –
  have  $RID(x)(RID(x)(P) ;; Q) = ((\exists \$x \cdot \exists \$x' \cdot ((\exists \$x \cdot \exists \$x' \cdot P) \wedge \$x' =_u \$x) ;; Q) \wedge \$x' =_u \$x)$ 
    by (simp add: RID-def usubst)
  also from assms have  $\dots = (((\exists \$x \cdot \exists \$x' \cdot P) \wedge (\exists \$x \cdot \$x' =_u \$x)) ;; (\exists \$x' \cdot Q)) \wedge \$x' =_u \$x)$ 
    by (rel-auto)
  also from assms have  $\dots = (((\exists \$x \cdot \exists \$x' \cdot P) ;; (\exists \$x \cdot \exists \$x' \cdot Q)) \wedge \$x' =_u \$x)$ 
    apply (rel-auto)
    apply (metis vwb-lens.put-eq)
    apply (metis mwb-lens.put-put vwb-lens-mwb)
  done
  also from assms have  $\dots = (((\exists \$x \cdot \exists \$x' \cdot P) \wedge \$x' =_u \$x) ;; (\exists \$x \cdot \exists \$x' \cdot Q)) \wedge \$x' =_u \$x)$ 
    by (rel-simp, metis (full-types) mwb-lens.put-put vwb-lens-def wb-lens-weak weak-lens.put-get)
  also have  $\dots = (((\exists \$x \cdot \exists \$x' \cdot P) \wedge \$x' =_u \$x) ;; ((\exists \$x \cdot \exists \$x' \cdot Q) \wedge \$x' =_u \$x)) \wedge \$x' =_u \$x)$ 
    by (rel-simp, fastforce)
  also have  $\dots = (((\exists \$x \cdot \exists \$x' \cdot P) \wedge \$x' =_u \$x) ;; ((\exists \$x \cdot \exists \$x' \cdot Q) \wedge \$x' =_u \$x))$ 
    by (rel-auto)
  also have  $\dots = (RID(x)(P) ;; RID(x)(Q))$ 
    by (rel-auto)
  finally show ?thesis .
qed

```

**lemma** *RID-seq-right*:

```

assumes vwb-lens x
shows  $RID(x)(P ;; RID(x)(Q)) = (RID(x)(P) ;; RID(x)(Q))$ 
proof –
  have  $RID(x)(P ;; RID(x)(Q)) = ((\exists \$x \cdot \exists \$x' \cdot P ;; ((\exists \$x \cdot \exists \$x' \cdot Q) \wedge \$x' =_u \$x)) \wedge \$x' =_u \$x)$ 
    by (simp add: RID-def usubst)
  also from assms have  $\dots = (((\exists \$x \cdot P) ;; (\exists \$x \cdot \exists \$x' \cdot Q) \wedge (\exists \$x' \cdot \$x' =_u \$x)) \wedge \$x' =_u \$x)$ 
    by (rel-auto)
  also from assms have  $\dots = (((\exists \$x \cdot \exists \$x' \cdot P) ;; (\exists \$x \cdot \exists \$x' \cdot Q)) \wedge \$x' =_u \$x)$ 
    apply (rel-auto)
    apply (metis vwb-lens.put-eq)
    apply (metis mwb-lens.put-put vwb-lens-mwb)
  done
  also from assms have  $\dots = (((\exists \$x \cdot \exists \$x' \cdot P) \wedge \$x' =_u \$x) ;; (\exists \$x \cdot \exists \$x' \cdot Q)) \wedge \$x' =_u \$x)$ 
    by (rel-simp robust, metis (full-types) mwb-lens.put-put vwb-lens-def wb-lens-weak weak-lens.put-get)
  also have  $\dots = (((\exists \$x \cdot \exists \$x' \cdot P) \wedge \$x' =_u \$x) ;; ((\exists \$x \cdot \exists \$x' \cdot Q) \wedge \$x' =_u \$x)) \wedge \$x' =_u \$x)$ 
    by (rel-simp, fastforce)
  also have  $\dots = (((\exists \$x \cdot \exists \$x' \cdot P) \wedge \$x' =_u \$x) ;; ((\exists \$x \cdot \exists \$x' \cdot Q) \wedge \$x' =_u \$x))$ 
    by (rel-auto)
  also have  $\dots = (RID(x)(P) ;; RID(x)(Q))$ 
    by (rel-auto)
  finally show ?thesis .
qed

```

**definition** *unrest-relation* ::  $('a \implies 'a) \Rightarrow 'a \text{ hrel} \Rightarrow \text{bool}$  (**infix**  $\#\#$  20)

**where**  $(x \#\# P) \longleftrightarrow (P = RID(x)(P))$

**declare** *unrest-relation-def* [*urel-defs*]

**lemma** *skip-r-runrest* [*unrest*]:

*vwb-lens*  $x \implies x \# \# I$   
**by** (*simp add: RID-skip-r unrest-relation-def*)

**lemma** *assigns-r-runrest*:

$\llbracket \text{vwb-lens } x; x \# \sigma \rrbracket \implies x \# \# \langle \sigma \rangle_a$   
**by** (*simp add: RID-assigns-r-diff unrest-relation-def*)

**lemma** *seq-r-runrest* [*unrest*]:

**assumes** *vwb-lens*  $x \# \# P \# \# Q$   
**shows**  $x \# \# (P ;; Q)$   
**by** (*metis RID-seq-left assms unrest-relation-def*)

**lemma** *false-runrest* [*unrest*]:  $x \# \# \text{false}$

**by** (*rel-auto*)

**lemma** *and-runrest* [*unrest*]:  $\llbracket \text{vwb-lens } x; x \# \# P; x \# \# Q \rrbracket \implies x \# \# (P \wedge Q)$

**by** (*metis RID-conj unrest-relation-def*)

**lemma** *or-runrest* [*unrest*]:  $\llbracket x \# \# P; x \# \# Q \rrbracket \implies x \# \# (P \vee Q)$

**by** (*simp add: RID-disj unrest-relation-def*)

## 14.9 Relational alphabet extension

**lift-definition** *rel-alpha-ext* ::  $'\beta \text{ hrel} \Rightarrow (' \beta \implies ' \alpha) \Rightarrow ' \alpha \text{ hrel}$  (**infix**  $\oplus_R$  65)  
**is**  $\lambda P \ x \ (b1, b2). P \ (\text{get}_x \ b1, \text{get}_x \ b2) \wedge (\forall \ b. b1 \oplus_L b \text{ on } x = b2 \oplus_L b \text{ on } x)$  .

**lemma** *rel-alpha-ext-alt-def*:

**assumes** *vwb-lens*  $y \ x \ +_L y \approx_L 1_L \ x \bowtie y$   
**shows**  $P \oplus_R x = (P \oplus_P (x \times_L x) \wedge \$y' =_u \$y)$   
**using** *assms*  
**apply** (*rel-auto robust, simp-all add: lens-override-def*)  
**apply** (*metis lens-indep-get lens-indep-sym*)  
**apply** (*metis vwb-lens-def wb-lens.get-put wb-lens-def weak-lens.put-get*)

**done**

**end**

## 15 Meta-level substitution

**theory** *utp-meta-subst*

**imports** *utp-rel utp-event*

**begin**

Meta substitution substituted a HOL variable in a UTP expression for another UTP expression. It is analogous to UTP substitution, but acts on functions.

**lift-definition** *msubst* ::  $('b \Rightarrow ('a, ' \alpha) \text{ uexpr}) \Rightarrow ('b, ' \alpha) \text{ uexpr} \Rightarrow ('a, ' \alpha) \text{ uexpr}$   
**is**  $\lambda F \ v \ b. F \ (v \ b) \ b$  .

**update-uexpr-rep-eq-thms** — Reread *rep-eq* theorems.

**syntax**

*-msubst* :: *logic*  $\Rightarrow$  *pttrn*  $\Rightarrow$  *logic*  $\Rightarrow$  *logic* (( $\neg$ →)) [990,0,0] 991)

#### translations

*-msubst* *P x v* == *CONST msubst* ( $\lambda x. P$ ) *v*

**lemma** *msubst-true* [*usubst*]:  $\text{true} \llbracket x \rightarrow v \rrbracket = \text{true}$   
by (*pred-auto*)

**lemma** *msubst-false* [*usubst*]:  $\text{false} \llbracket x \rightarrow v \rrbracket = \text{false}$   
by (*pred-auto*)

**lemma** *msubst-lit* [*usubst*]:  $\llbracket x \rrbracket \llbracket x \rightarrow v \rrbracket = v$   
by (*pred-auto*)

**lemma** *msubst-lit-2-1* [*usubst*]:  $\llbracket x \rrbracket \llbracket (x, y) \rightarrow (u, v)_u \rrbracket = u$   
by (*pred-auto*)

**lemma** *msubst-lit-2-2* [*usubst*]:  $\llbracket y \rrbracket \llbracket (x, y) \rightarrow (u, v)_u \rrbracket = v$   
by (*pred-auto*)

**lemma** *msubst-lit'* [*usubst*]:  $\llbracket y \rrbracket \llbracket x \rightarrow v \rrbracket = \llbracket y \rrbracket$   
by (*pred-auto*)

**lemma** *msubst-lit'-2* [*usubst*]:  $\llbracket z \rrbracket \llbracket (x, y) \rightarrow v \rrbracket = \llbracket z \rrbracket$   
by (*pred-auto*)

**lemma** *msubst-uop* [*usubst*]:  $(\text{uop } f \ (v \ x)) \llbracket x \rightarrow u \rrbracket = \text{uop } f \ ((v \ x) \llbracket x \rightarrow u \rrbracket)$   
by (*rel-auto*)

**lemma** *msubst-uop-2* [*usubst*]:  $(\text{uop } f \ (v \ x \ y)) \llbracket (x, y) \rightarrow u \rrbracket = \text{uop } f \ ((v \ x \ y) \llbracket (x, y) \rightarrow u \rrbracket)$   
by (*pred-simp*, *pred-simp*)

**lemma** *msubst-bop* [*usubst*]:  $(\text{bop } f \ (v \ x) \ (w \ x)) \llbracket x \rightarrow u \rrbracket = \text{bop } f \ ((v \ x) \llbracket x \rightarrow u \rrbracket) \ ((w \ x) \llbracket x \rightarrow u \rrbracket)$   
by (*rel-auto*)

**lemma** *msubst-bop-2* [*usubst*]:  $(\text{bop } f \ (v \ x \ y) \ (w \ x \ y)) \llbracket (x, y) \rightarrow u \rrbracket = \text{bop } f \ ((v \ x \ y) \llbracket (x, y) \rightarrow u \rrbracket) \ ((w \ x \ y) \llbracket (x, y) \rightarrow u \rrbracket)$   
by (*pred-simp*, *pred-simp*)

**lemma** *msubst-not* [*usubst*]:  $(\neg P(x)) \llbracket x \rightarrow v \rrbracket = (\neg ((P \ x) \llbracket x \rightarrow v \rrbracket))$   
by (*pred-auto*)

**lemma** *msubst-not-2* [*usubst*]:  $(\neg P \ x \ y) \llbracket (x, y) \rightarrow v \rrbracket = (\neg ((P \ x \ y) \llbracket (x, y) \rightarrow v \rrbracket))$   
by (*pred-auto*) +

**lemma** *msubst-disj* [*usubst*]:  $(P(x) \vee Q(x)) \llbracket x \rightarrow v \rrbracket = ((P(x)) \llbracket x \rightarrow v \rrbracket \vee (Q(x)) \llbracket x \rightarrow v \rrbracket)$   
by (*pred-auto*)

**lemma** *msubst-disj-2* [*usubst*]:  $(P \ x \ y \vee Q \ x \ y) \llbracket (x, y) \rightarrow v \rrbracket = ((P \ x \ y) \llbracket (x, y) \rightarrow v \rrbracket \vee (Q \ x \ y) \llbracket (x, y) \rightarrow v \rrbracket)$   
by (*pred-auto*) +

**lemma** *msubst-conj* [*usubst*]:  $(P(x) \wedge Q(x)) \llbracket x \rightarrow v \rrbracket = ((P(x)) \llbracket x \rightarrow v \rrbracket \wedge (Q(x)) \llbracket x \rightarrow v \rrbracket)$   
by (*pred-auto*)

**lemma** *msubst-conj-2* [*usubst*]:  $(P \ x \ y \wedge Q \ x \ y) \llbracket (x, y) \rightarrow v \rrbracket = ((P \ x \ y) \llbracket (x, y) \rightarrow v \rrbracket \wedge (Q \ x \ y) \llbracket (x, y) \rightarrow v \rrbracket)$

by (pred-auto)+

**lemma** *msubst-implies* [usubst]:

$(P\ x \Rightarrow Q\ x) \llbracket x \rightarrow v \rrbracket = ((P\ x) \llbracket x \rightarrow v \rrbracket \Rightarrow (Q\ x) \llbracket x \rightarrow v \rrbracket)$

by (pred-auto)

**lemma** *msubst-implies-2* [usubst]:

$(P\ x\ y \Rightarrow Q\ x\ y) \llbracket (x, y) \rightarrow v \rrbracket = ((P\ x\ y) \llbracket (x, y) \rightarrow v \rrbracket \Rightarrow (Q\ x\ y) \llbracket (x, y) \rightarrow v \rrbracket)$

by (pred-auto)+

**lemma** *msubst-shAll* [usubst]:

$(\forall\ x \cdot P\ x\ y) \llbracket y \rightarrow v \rrbracket = (\forall\ x \cdot (P\ x\ y) \llbracket y \rightarrow v \rrbracket)$

by (pred-auto)

**lemma** *msubst-shAll-2* [usubst]:

$(\forall\ x \cdot P\ x\ y\ z) \llbracket (y, z) \rightarrow v \rrbracket = (\forall\ x \cdot (P\ x\ y\ z) \llbracket (y, z) \rightarrow v \rrbracket)$

by (pred-auto)+

**lemma** *msubst-event* [usubst]:

$(c \cdot v\ x)_u \llbracket x \rightarrow u \rrbracket = (c \cdot (v\ x) \llbracket x \rightarrow u \rrbracket)_u$

by (pred-simp)

**lemma** *msubst-event-2* [usubst]:

$(c \cdot v\ x\ y)_u \llbracket (x, y) \rightarrow u \rrbracket = (c \cdot (v\ x\ y) \llbracket (x, y) \rightarrow u \rrbracket)_u$

by (pred-simp)+

**lemma** *msubst-var* [usubst]:

$(\text{utp-expr.var}\ x) \llbracket y \rightarrow u \rrbracket = \text{utp-expr.var}\ x$

by (pred-simp)

**lemma** *msubst-var-2* [usubst]:

$(\text{utp-expr.var}\ x) \llbracket (y, z) \rightarrow u \rrbracket = \text{utp-expr.var}\ x$

by (pred-simp)+

**lemma** *msubst-seq* [usubst]:  $(P(x) ;; Q(x)) \llbracket x \rightarrow \llbracket v \rrbracket \rrbracket = ((P(x)) \llbracket x \rightarrow \llbracket v \rrbracket \rrbracket ;; (Q(x)) \llbracket x \rightarrow \llbracket v \rrbracket \rrbracket)$

by (rel-auto)

**lemma** *msubst-unrest* [unrest]:  $\llbracket \bigwedge v. x \# P(v); x \# k \rrbracket \Longrightarrow x \# P(v) \llbracket v \rightarrow k \rrbracket$

by (pred-auto)

end

## 16 UTP Deduction Tactic

**theory** *utp-deduct*

**imports** *utp-pred*

**begin**

**named-theorems** *uintro*

**named-theorems** *uelim*

**named-theorems** *udest*

**lemma** *uttrueI* [uintro]:  $\llbracket \text{true} \rrbracket_e b$

by (pred-auto)



**lemma** *uopI* [*uintro*]:  $f \llbracket x \rrbracket_e b \Longrightarrow \llbracket uop\ f\ x \rrbracket_e b$   
**by** (*pred-auto*)

**lemma** *bopI* [*uintro*]:  $f \llbracket x \rrbracket_e b \llbracket y \rrbracket_e b \Longrightarrow \llbracket bop\ f\ x\ y \rrbracket_e b$   
**by** (*pred-auto*)

**lemma** *tropI* [*uintro*]:  $f \llbracket x \rrbracket_e b \llbracket y \rrbracket_e b \llbracket z \rrbracket_e b \Longrightarrow \llbracket trop\ f\ x\ y\ z \rrbracket_e b$   
**by** (*pred-auto*)

**lemma** *uconjI* [*uintro*]:  $\llbracket \llbracket p \rrbracket_e b ; \llbracket q \rrbracket_e b \rrbracket \Longrightarrow \llbracket p \wedge q \rrbracket_e b$   
**by** (*pred-auto*)

**lemma** *uconjE* [*uelim*]:  $\llbracket \llbracket p \wedge q \rrbracket_e b ; \llbracket \llbracket p \rrbracket_e b ; \llbracket q \rrbracket_e b \rrbracket \Longrightarrow P \rrbracket \Longrightarrow P$   
**by** (*pred-auto*)

**lemma** *uimpI* [*uintro*]:  $\llbracket \llbracket p \rrbracket_e b \Longrightarrow \llbracket q \rrbracket_e b \rrbracket \Longrightarrow \llbracket p \Rightarrow q \rrbracket_e b$   
**by** (*pred-auto*)

**lemma** *uimpE* [*elim*]:  $\llbracket \llbracket p \Rightarrow q \rrbracket_e b ; (\llbracket p \rrbracket_e b \Longrightarrow \llbracket q \rrbracket_e b) \Longrightarrow P \rrbracket \Longrightarrow P$   
**by** (*pred-auto*)

**lemma** *ushAllI* [*uintro*]:  $\llbracket \bigwedge x. \llbracket p(x) \rrbracket_e b \rrbracket \Longrightarrow \llbracket \bigvee x. p(x) \rrbracket_e b$   
**by** *pred-auto*

**lemma** *ushExI* [*uintro*]:  $\llbracket \llbracket p(x) \rrbracket_e b \rrbracket \Longrightarrow \llbracket \exists x. p(x) \rrbracket_e b$   
**by** *pred-auto*

**lemma** *udeduct-tautI* [*uintro*]:  $\llbracket \bigwedge b. \llbracket p \rrbracket_e b \rrbracket \Longrightarrow 'p'$   
**using** *taut.rep-eq* **by** *blast*

**lemma** *udeduct-refineI* [*uintro*]:  $\llbracket \bigwedge b. \llbracket q \rrbracket_e b \Longrightarrow \llbracket p \rrbracket_e b \rrbracket \Longrightarrow p \sqsubseteq q$   
**by** *pred-auto*

**lemma** *udeduct-eqI* [*uintro*]:  $\llbracket \bigwedge b. \llbracket p \rrbracket_e b \Longrightarrow \llbracket q \rrbracket_e b ; \bigwedge b. \llbracket q \rrbracket_e b \Longrightarrow \llbracket p \rrbracket_e b \rrbracket \Longrightarrow p = q$   
**by** (*pred-auto*)

Some of the following lemmas help backward reasoning with bindings

**lemma** *conj-implies*:  $\llbracket \llbracket P \wedge Q \rrbracket_e b \rrbracket \Longrightarrow \llbracket P \rrbracket_e b \wedge \llbracket Q \rrbracket_e b$   
**by** *pred-auto*

**lemma** *conj-implies2*:  $\llbracket \llbracket P \rrbracket_e b \wedge \llbracket Q \rrbracket_e b \rrbracket \Longrightarrow \llbracket P \wedge Q \rrbracket_e b$   
**by** *pred-auto*

**lemma** *disj-eq*:  $\llbracket \llbracket P \rrbracket_e b \vee \llbracket Q \rrbracket_e b \rrbracket \Longrightarrow \llbracket P \vee Q \rrbracket_e b$   
**by** *pred-auto*

**lemma** *disj-eq2*:  $\llbracket \llbracket P \vee Q \rrbracket_e b \rrbracket \Longrightarrow \llbracket P \rrbracket_e b \vee \llbracket Q \rrbracket_e b$   
**by** *pred-auto*

**lemma** *conj-eq-subst*:  $(\llbracket P \wedge Q \rrbracket_e b \wedge \llbracket P \rrbracket_e b = \llbracket R \rrbracket_e b) = (\llbracket R \wedge Q \rrbracket_e b \wedge \llbracket P \rrbracket_e b = \llbracket R \rrbracket_e b)$   
**by** *pred-auto*

**lemma** *conj-imp-subst*:  $(\llbracket P \wedge Q \rrbracket_e b \wedge (\llbracket Q \rrbracket_e b \longrightarrow (\llbracket P \rrbracket_e b = \llbracket R \rrbracket_e b))) = (\llbracket R \wedge Q \rrbracket_e b \wedge (\llbracket Q \rrbracket_e b \longrightarrow (\llbracket P \rrbracket_e b = \llbracket R \rrbracket_e b)))$   
**by** *pred-auto*

**lemma** *disj-imp-subst*:  $(\llbracket Q \wedge (P \vee S) \rrbracket_e b \wedge (\llbracket Q \rrbracket_e b \longrightarrow (\llbracket P \rrbracket_e b = \llbracket R \rrbracket_e b))) = (\llbracket Q \wedge (R \vee S) \rrbracket_e b \wedge (\llbracket Q \rrbracket_e b \longrightarrow (\llbracket P \rrbracket_e b = \llbracket R \rrbracket_e b)))$   
**by** *pred-auto*

Simplifications on value equality

**lemma** *ueqpr-eq*:  $(\llbracket e_0 \rrbracket_e b = \llbracket e_1 \rrbracket_e b) = \llbracket e_0 =_u e_1 \rrbracket_e b$   
**by** *pred-auto*

**lemma** *ueqpr-trans*:  $(\llbracket P \wedge e_0 =_u e_1 \rrbracket_e b \wedge (\llbracket P \rrbracket_e b \longrightarrow \llbracket e_1 =_u e_2 \rrbracket_e b)) = (\llbracket P \wedge e_0 =_u e_2 \rrbracket_e b \wedge (\llbracket P \rrbracket_e b \longrightarrow \llbracket e_1 =_u e_2 \rrbracket_e b))$   
**by** (*pred-auto*)

**lemma** *ueqpr-trans2*:  $(\llbracket P \wedge Q \wedge e_0 =_u e_1 \rrbracket_e b \wedge (\llbracket Q \rrbracket_e b \longrightarrow \llbracket e_1 =_u e_2 \rrbracket_e b)) = (\llbracket P \wedge Q \wedge e_0 =_u e_2 \rrbracket_e b \wedge (\llbracket P \rrbracket_e b \longrightarrow \llbracket e_1 =_u e_2 \rrbracket_e b))$   
**by** (*pred-auto*)

**lemma** *uequality*:  $(\llbracket \llbracket R \rrbracket_e b = \llbracket Q \rrbracket_e b \rrbracket \implies \llbracket P \wedge R \rrbracket_e b = \llbracket P \wedge Q \rrbracket_e b$   
**by** *pred-auto*

**lemma** *ueqe1*:  $\llbracket \llbracket P \rrbracket_e b \implies (\llbracket R \rrbracket_e b = \llbracket Q \rrbracket_e b) \rrbracket \implies (\llbracket P \wedge R \rrbracket_e b \implies \llbracket P \wedge Q \rrbracket_e b)$   
**by** *pred-auto*

**lemma** *ueqe2*:  $(\llbracket P \rrbracket_e b \implies (\llbracket Q \rrbracket_e b = \llbracket R \rrbracket_e b) \wedge \llbracket Q \wedge P \rrbracket_e b = \llbracket R \wedge P \rrbracket_e b) \implies$   
 $(\llbracket P \rrbracket_e b \implies (\llbracket Q \rrbracket_e b = \llbracket R \rrbracket_e b))$   
**by** *pred-auto*

**lemma** *ueqe3*:  $\llbracket \llbracket P \rrbracket_e b \implies (\llbracket Q \rrbracket_e b = \llbracket R \rrbracket_e b) \rrbracket \implies (\llbracket R \wedge P \rrbracket_e b = \llbracket Q \wedge P \rrbracket_e b)$   
**by** *pred-auto*

The following allows simplifying the equality if  $P \Rightarrow Q = R$

**lemma** *ueqe3-imp*:  $(\bigwedge b. \llbracket P \rrbracket_e b \implies (\llbracket Q \rrbracket_e b = \llbracket R \rrbracket_e b)) \implies ((R \wedge P) = (Q \wedge P))$   
**by** *pred-auto*

**lemma** *ueqe3-imp3*:  $(\bigwedge b. \llbracket P \rrbracket_e b \implies (\llbracket Q \rrbracket_e b = \llbracket R \rrbracket_e b)) \implies ((P \wedge Q) = (P \wedge R))$   
**by** *pred-auto*

**lemma** *ueqe3-imp2*:  $\llbracket (\bigwedge b. \llbracket P0 \wedge P1 \rrbracket_e b \implies \llbracket Q \rrbracket_e b \implies \llbracket R \rrbracket_e b = \llbracket S \rrbracket_e b) \rrbracket \implies ((P0 \wedge P1 \wedge (Q \Rightarrow R)) = (P0 \wedge P1 \wedge (Q \Rightarrow S)))$   
**by** *pred-auto*

The following can introduce the binding notation into predicates

**lemma** *conj-bind-dist*:  $\llbracket P \wedge Q \rrbracket_e b = (\llbracket P \rrbracket_e b \wedge \llbracket Q \rrbracket_e b)$   
**by** *pred-auto*

**lemma** *disj-bind-dist*:  $\llbracket P \vee Q \rrbracket_e b = (\llbracket P \rrbracket_e b \vee \llbracket Q \rrbracket_e b)$   
**by** *pred-auto*

**lemma** *imp-bind-dist*:  $\llbracket P \Rightarrow Q \rrbracket_e b = (\llbracket P \rrbracket_e b \longrightarrow \llbracket Q \rrbracket_e b)$   
**by** *pred-auto*  
**end**

## 17 Relational Calculus Laws

**theory** *utp-rel-laws*  
**imports** *utp-rel*  
**begin**

### 17.1 Conditional Laws

**lemma** *comp-cond-left-distr*:  
 $((P \triangleleft b \triangleright_r Q) ;; R) = ((P ;; R) \triangleleft b \triangleright_r (Q ;; R))$   
**by** (*rel-auto*)

**lemma** *cond-seq-left-distr*:  
 $out\alpha \# b \implies ((P \triangleleft b \triangleright Q) ;; R) = ((P ;; R) \triangleleft b \triangleright (Q ;; R))$   
**by** (*rel-auto*)

**lemma** *cond-seq-right-distr*:  
 $in\alpha \# b \implies (P ;; (Q \triangleleft b \triangleright R)) = ((P ;; Q) \triangleleft b \triangleright (P ;; R))$   
**by** (*rel-auto*)

**lemma** *cond-mono*:  
 $\llbracket P_1 \sqsubseteq P_2; Q_1 \sqsubseteq Q_2 \rrbracket \implies (P_1 \triangleleft b \triangleright Q_1) \sqsubseteq (P_2 \triangleleft b \triangleright Q_2)$   
**by** (*rel-auto*)

**lemma** *cond-monotonic*:  
 $\llbracket mono\ P; mono\ Q \rrbracket \implies mono\ (\lambda X. P\ X \triangleleft b \triangleright Q\ X)$   
**by** (*simp add: mono-def, rel-blast*)

### 17.2 Precondition and Postcondition Laws

**theorem** *precond-equiv*:  
 $P = (P ;; true) \longleftrightarrow (out\alpha \# P)$   
**by** (*rel-auto*)

**theorem** *postcond-equiv*:  
 $P = (true ;; P) \longleftrightarrow (in\alpha \# P)$   
**by** (*rel-auto*)

**lemma** *precond-right-unit*:  $out\alpha \# p \implies (p ;; true) = p$   
**by** (*metis precond-equiv*)

**lemma** *postcond-left-unit*:  $in\alpha \# p \implies (true ;; p) = p$   
**by** (*metis postcond-equiv*)

**theorem** *precond-left-zero*:  
**assumes**  $out\alpha \# p \neq false$   
**shows**  $(true ;; p) = true$   
**using** *assms* **by** (*rel-auto*)

**theorem** *feasibile-iff-true-right-zero*:  
 $P ;; true = true \longleftrightarrow \exists out\alpha \cdot P$   
**by** (*rel-auto*)

### 17.3 Sequential Composition Laws

**lemma** *seqr-assoc*:  $P ;; (Q ;; R) = (P ;; Q) ;; R$

**by** (*rel-auto*)

**lemma** *seqr-left-unit* [*simp*]:  
 $II ;; P = P$   
**by** (*rel-auto*)

**lemma** *seqr-right-unit* [*simp*]:  
 $P ;; II = P$   
**by** (*rel-auto*)

**lemma** *seqr-left-zero* [*simp*]:  
 $false ;; P = false$   
**by** *pred-auto*

**lemma** *seqr-right-zero* [*simp*]:  
 $P ;; false = false$   
**by** *pred-auto*

**lemma** *impl-seqr-mono*:  $\llbracket 'P \Rightarrow Q'; 'R \Rightarrow S' \rrbracket \Longrightarrow '(P ;; R) \Rightarrow (Q ;; S)'$   
**by** (*pred-blast*)

**lemma** *seqr-mono*:  
 $\llbracket P_1 \sqsubseteq P_2; Q_1 \sqsubseteq Q_2 \rrbracket \Longrightarrow (P_1 ;; Q_1) \sqsubseteq (P_2 ;; Q_2)$   
**by** (*rel-blast*)

**lemma** *seqr-monotonic*:  
 $\llbracket mono\ P; mono\ Q \rrbracket \Longrightarrow mono\ (\lambda X. P\ X ;; Q\ X)$   
**by** (*simp add: mono-def, rel-blast*)

**lemma** *seqr-exists-left*:  
 $((\exists \$x \cdot P) ;; Q) = (\exists \$x \cdot (P ;; Q))$   
**by** (*rel-auto*)

**lemma** *seqr-exists-right*:  
 $(P ;; (\exists \$x' \cdot Q)) = (\exists \$x' \cdot (P ;; Q))$   
**by** (*rel-auto*)

**lemma** *seqr-or-distl*:  
 $((P \vee Q) ;; R) = ((P ;; R) \vee (Q ;; R))$   
**by** (*rel-auto*)

**lemma** *seqr-or-distr*:  
 $(P ;; (Q \vee R)) = ((P ;; Q) \vee (P ;; R))$   
**by** (*rel-auto*)

**lemma** *seqr-and-distr-ufunc*:  
 $ufunctional\ P \Longrightarrow (P ;; (Q \wedge R)) = ((P ;; Q) \wedge (P ;; R))$   
**by** (*rel-auto*)

**lemma** *seqr-and-distl-ujnj*:  
 $ujnj\ R \Longrightarrow ((P \wedge Q) ;; R) = ((P ;; R) \wedge (Q ;; R))$   
**by** (*rel-auto*)

**lemma** *seqr-unfold*:  
 $(P ;; Q) = (\exists v \cdot P \llbracket \llbracket v \rrbracket / \$\mathbf{v}' \rrbracket \wedge Q \llbracket \llbracket v \rrbracket / \$\mathbf{v} \rrbracket)$

by (rel-auto)

lemma *seqr-middle*:

assumes *vwb-lens*  $x$   
shows  $(P ;; Q) = (\exists v \cdot P[\llbracket v \rrbracket / \$x'] ;; Q[\llbracket v \rrbracket / \$x])$   
using *assms*  
apply (rel-auto robust)  
apply (rename-tac  $xa$   $P$   $Q$   $a$   $b$   $y$ )  
apply (rule-tac  $x = \text{get}_{xa} y$  in  $exI$ )  
apply (rule-tac  $x = y$  in  $exI$ )  
apply (*simp*)  
done

lemma *seqr-left-one-point*:

assumes *vwb-lens*  $x$   
shows  $((P \wedge \$x' =_u \llbracket v \rrbracket) ;; Q) = (P[\llbracket v \rrbracket / \$x'] ;; Q[\llbracket v \rrbracket / \$x])$   
using *assms*  
by (rel-auto, metis *vwb-lens-wb wb-lens.get-put*)

lemma *seqr-right-one-point*:

assumes *vwb-lens*  $x$   
shows  $(P ;; (\$x =_u \llbracket v \rrbracket \wedge Q)) = (P[\llbracket v \rrbracket / \$x'] ;; Q[\llbracket v \rrbracket / \$x])$   
using *assms*  
by (rel-auto, metis *vwb-lens-wb wb-lens.get-put*)

lemma *seqr-left-one-point-true*:

assumes *vwb-lens*  $x$   
shows  $((P \wedge \$x') ;; Q) = (P[\llbracket \text{true} \rrbracket / \$x'] ;; Q[\llbracket \text{true} \rrbracket / \$x])$   
by (metis *assms seqr-left-one-point true-alt-def upred-eq-true*)

lemma *seqr-left-one-point-false*:

assumes *vwb-lens*  $x$   
shows  $((P \wedge \neg \$x') ;; Q) = (P[\llbracket \text{false} \rrbracket / \$x'] ;; Q[\llbracket \text{false} \rrbracket / \$x])$   
by (metis *assms false-alt-def seqr-left-one-point upred-eq-false*)

lemma *seqr-right-one-point-true*:

assumes *vwb-lens*  $x$   
shows  $(P ;; (\$x \wedge Q)) = (P[\llbracket \text{true} \rrbracket / \$x'] ;; Q[\llbracket \text{true} \rrbracket / \$x])$   
by (metis *assms seqr-right-one-point true-alt-def upred-eq-true*)

lemma *seqr-right-one-point-false*:

assumes *vwb-lens*  $x$   
shows  $(P ;; (\neg \$x \wedge Q)) = (P[\llbracket \text{false} \rrbracket / \$x'] ;; Q[\llbracket \text{false} \rrbracket / \$x])$   
by (metis *assms false-alt-def seqr-right-one-point upred-eq-false*)

lemma *seqr-insert-ident-left*:

assumes *vwb-lens*  $x$   $\$x' \# P$   $\$x \# Q$   
shows  $((\$x' =_u \$x \wedge P) ;; Q) = (P ;; Q)$   
using *assms*  
by (rel-simp, meson *vwb-lens-wb wb-lens-weak weak-lens.put-get*)

lemma *seqr-insert-ident-right*:

assumes *vwb-lens*  $x$   $\$x' \# P$   $\$x \# Q$   
shows  $(P ;; (\$x' =_u \$x \wedge Q)) = (P ;; Q)$   
using *assms*

by (rel-simp, metis (no-types, hide-lams) vwb-lens-def wb-lens-def weak-lens.put-get)

**lemma** seq-var-ident-lift:

assumes vwb-lens  $x$   $\$x'$   $\#$   $P$   $\$x$   $\#$   $Q$   
 shows  $((\$x' =_u \$x \wedge P) ;; (\$x' =_u \$x \wedge Q)) = (\$x' =_u \$x \wedge (P ;; Q))$   
 using assms apply (rel-auto)  
 by (metis (no-types, lifting) vwb-lens-wb wb-lens-weak weak-lens.put-get)

**lemma** seqr-bool-split:

assumes vwb-lens  $x$   
 shows  $P ;; Q = (P \llbracket \text{true}/\$x \rrbracket ;; Q \llbracket \text{true}/\$x \rrbracket \vee P \llbracket \text{false}/\$x' \rrbracket ;; Q \llbracket \text{false}/\$x \rrbracket)$   
 using assms  
 by (subst seqr-middle[of  $x$ ], simp-all add: true-alt-def false-alt-def)

**lemma** cond-inter-var-split:

assumes vwb-lens  $x$   
 shows  $(P \triangleleft \$x' \triangleright Q) ;; R = (P \llbracket \text{true}/\$x' \rrbracket ;; R \llbracket \text{true}/\$x \rrbracket \vee Q \llbracket \text{false}/\$x' \rrbracket ;; R \llbracket \text{false}/\$x \rrbracket)$

**proof** –

have  $(P \triangleleft \$x' \triangleright Q) ;; R = ((\$x' \wedge P) ;; R \vee (\neg \$x' \wedge Q) ;; R)$   
 by (simp add: cond-def seqr-or-distl)  
 also have  $\dots = ((P \wedge \$x') ;; R \vee (Q \wedge \neg \$x') ;; R)$   
 by (rel-auto)  
 also have  $\dots = (P \llbracket \text{true}/\$x' \rrbracket ;; R \llbracket \text{true}/\$x \rrbracket \vee Q \llbracket \text{false}/\$x' \rrbracket ;; R \llbracket \text{false}/\$x \rrbracket)$   
 by (simp add: seqr-left-one-point-true seqr-left-one-point-false assms)  
 finally show ?thesis .

**qed**

**theorem** seqr-pre-transfer:  $\text{in}\alpha \# q \implies ((P \wedge q) ;; R) = (P ;; (q^- \wedge R))$

by (rel-auto)

**theorem** seqr-pre-transfer':

$((P \wedge \lceil q \rceil_>) ;; R) = (P ;; (\lceil q \rceil_< \wedge R))$   
 by (rel-auto)

**theorem** seqr-post-out:  $\text{in}\alpha \# r \implies (P ;; (Q \wedge r)) = ((P ;; Q) \wedge r)$

by (rel-blast)

**lemma** seqr-post-var-out:

fixes  $x :: (\text{bool} \implies 'a)$   
 shows  $(P ;; (Q \wedge \$x')) = ((P ;; Q) \wedge \$x')$   
 by (rel-auto)

**theorem** seqr-post-transfer:  $\text{out}\alpha \# q \implies (P ;; (q \wedge R)) = ((P \wedge q^-) ;; R)$

by (rel-auto)

**lemma** seqr-pre-out:  $\text{out}\alpha \# p \implies ((p \wedge Q) ;; R) = (p \wedge (Q ;; R))$

by (rel-blast)

**lemma** seqr-pre-var-out:

fixes  $x :: (\text{bool} \implies 'a)$   
 shows  $((\$x \wedge P) ;; Q) = (\$x \wedge (P ;; Q))$   
 by (rel-auto)

**lemma** seqr-true-lemma:

$(P = (\neg ((\neg P) ;; \text{true}))) = (P = (P ;; \text{true}))$

by (rel-auto)

**lemma** *seqr-to-conj*:  $\llbracket \text{out}\alpha \# P; \text{in}\alpha \# Q \rrbracket \implies (P ;; Q) = (P \wedge Q)$   
 by (metis postcond-left-unit seqr-pre-out utp-pred-laws.inf-top.right-neutral)

**lemma** *shEx-lift-seq-1* [uquant-lift]:  
 $((\exists x \cdot P x) ;; Q) = (\exists x \cdot (P x ;; Q))$   
 by pred-auto

**lemma** *shEx-lift-seq-2* [uquant-lift]:  
 $(P ;; (\exists x \cdot Q x)) = (\exists x \cdot (P ;; Q x))$   
 by pred-auto

## 17.4 Iterated Sequential Composition Laws

**lemma** *iter-seqr-nil* [simp]:  $(;; i : [] \cdot P(i)) = II$   
 by (simp add: seqr-iter-def)

**lemma** *iter-seqr-cons* [simp]:  $(;; i : (x \# xs) \cdot P(i)) = P(x) ;; (;; i : xs \cdot P(i))$   
 by (simp add: seqr-iter-def)

## 17.5 Quantale Laws

**lemma** *seq-Sup-distl*:  $P ;; (\bigcap A) = (\bigcap_{Q \in A} P ;; Q)$   
 by (transfer, auto)

**lemma** *seq-Sup-distr*:  $(\bigcap A) ;; Q = (\bigcap_{P \in A} P ;; Q)$   
 by (transfer, auto)

**lemma** *seq-UNIF-distl*:  $P ;; (\bigcap_{Q \in A} F(Q)) = (\bigcap_{Q \in A} P ;; F(Q))$   
 by (simp add: UNIF-as-Sup-collect seq-Sup-distl)

**lemma** *seq-UNIF-distl'*:  $P ;; (\bigcap Q \cdot F(Q)) = (\bigcap Q \cdot P ;; F(Q))$   
 by (metis UNIF-mem-UNIV seq-UNIF-distl)

**lemma** *seq-UNIF-distr*:  $(\bigcap_{P \in A} F(P)) ;; Q = (\bigcap_{P \in A} P \cdot F(P) ;; Q)$   
 by (simp add: UNIF-as-Sup-collect seq-Sup-distr)

**lemma** *seq-UNIF-distr'*:  $(\bigcap P \cdot F(P)) ;; Q = (\bigcap P \cdot F(P) ;; Q)$   
 by (metis UNIF-mem-UNIV seq-UNIF-distr)

**lemma** *seq-SUP-distl*:  $P ;; (\bigcap_{i \in A} Q(i)) = (\bigcap_{i \in A} P ;; Q(i))$   
 by (metis image-image seq-Sup-distl)

**lemma** *seq-SUP-distr*:  $(\bigcap_{i \in A} P(i)) ;; Q = (\bigcap_{i \in A} P(i) ;; Q)$   
 by (simp add: seq-Sup-distr)

## 17.6 Skip Laws

**lemma** *cond-skip*:  $\text{out}\alpha \# b \implies (b \wedge II) = (II \wedge b^-)$   
 by (rel-auto)

**lemma** *pre-skip-post*:  $(\lceil b \rceil_{<} \wedge II) = (II \wedge \lceil b \rceil_{>})$   
 by (rel-auto)

**lemma** *skip-var*:

**fixes**  $x :: (bool \implies 'a)$   
**shows**  $(\$x \wedge II) = (II \wedge \$x')$   
**by**  $(rel-auto)$

**lemma** *skip-r-unfold*:

$vwb-lens\ x \implies II = (\$x' =_u \$x \wedge II \upharpoonright_\alpha x)$   
**by**  $(rel-simp, metis\ mwb-lens.put-put\ vwb-lens-mwb\ vwb-lens-wb\ wb-lens.get-put)$

**lemma** *skip-r-alpha-eq*:

$II = (\$v' =_u \$v)$   
**by**  $(rel-auto)$

**lemma** *skip-ra-unfold*:

$II_{x;y} = (\$x' =_u \$x \wedge II_y)$   
**by**  $(rel-auto)$

**lemma** *skip-res-as-ra*:

$\llbracket vwb-lens\ y; x +_L y \approx_L 1_L; x \bowtie y \rrbracket \implies II \upharpoonright_\alpha x = II_y$   
**apply**  $(rel-auto)$   
**apply**  $(metis\ (no-types,\ lifting)\ lens-indep-def)$   
**apply**  $(metis\ vwb-lens.put-eq)$

**done**

## 17.7 Assignment Laws

**lemma** *assigns-subst* [*usubst*]:

$\llbracket \sigma \rrbracket_s \upharpoonright \langle \varrho \rangle_a = \langle \varrho \circ \sigma \rangle_a$   
**by**  $(rel-auto)$

**lemma** *assigns-r-comp*:  $(\langle \sigma \rangle_a ;; P) = (\llbracket \sigma \rrbracket_s \upharpoonright P)$

**by**  $(rel-auto)$

**lemma** *assigns-r-feasible*:

$(\langle \sigma \rangle_a ;; true) = true$   
**by**  $(rel-auto)$

**lemma** *assign-subst* [*usubst*]:

$\llbracket mwb-lens\ x; mwb-lens\ y \rrbracket \implies [\$x \mapsto_s [u]_<] \upharpoonright (y := v) = (x, y := u, [x \mapsto_s u] \upharpoonright v)$   
**by**  $(rel-auto)$

**lemma** *assigns-idem*:  $mwb-lens\ x \implies (x, x := u, v) = (x := v)$

**by**  $(simp\ add:\ usubst)$

**lemma** *assigns-comp*:  $(\langle f \rangle_a ;; \langle g \rangle_a) = \langle g \circ f \rangle_a$

**by**  $(simp\ add:\ assigns-r-comp\ usubst)$

**lemma** *assigns-r-conv*:

$bij\ f \implies \langle f \rangle_a^- = \langle inv\ f \rangle_a$   
**by**  $(rel-auto,\ simp-all\ add:\ bij-is-inj\ bij-is-surj\ surj-f-inv-f)$

**lemma** *assign-pred-transfer*:

**fixes**  $x :: ('a \implies 'a)$   
**assumes**  $\$x \# b\ out\alpha \# b$   
**shows**  $(b \wedge x := v) = (x := v \wedge b^-)$   
**using** *assms* **by**  $(rel-blast)$



**lemma** *assign-r-comp*:  $x := u ;; P = P[[u]_{<}/\$x]$   
 by (*simp add: assigns-r-comp usubst alpha*)

**lemma** *assign-test*:  $\text{mwb-lens } x \implies (x := \llbracket u \rrbracket ;; x := \llbracket v \rrbracket) = (x := \llbracket v \rrbracket)$   
 by (*simp add: assigns-comp usubst*)

**lemma** *assign-twice*:  $\llbracket \text{mwb-lens } x; x \# f \rrbracket \implies (x := e ;; x := f) = (x := f)$   
 by (*simp add: assigns-comp usubst*)

**lemma** *assign-commute*:  
 assumes  $x \bowtie y \ x \# f \ y \# e$   
 shows  $(x := e ;; y := f) = (y := f ;; x := e)$   
 using *assms*  
 by (*rel-simp, simp-all add: lens-indep-comm*)

**lemma** *assign-cond*:  
 fixes  $x :: ('a \implies 'c)$   
 assumes  $\text{out}\alpha \# b$   
 shows  $(x := e ;; (P \triangleleft b \triangleright Q)) = ((x := e ;; P) \triangleleft (b[[e]_{<}/\$x]) \triangleright (x := e ;; Q))$   
 by (*rel-auto*)

**lemma** *assign-rcond*:  
 fixes  $x :: ('a \implies 'c)$   
 shows  $(x := e ;; (P \triangleleft b \triangleright_r Q)) = ((x := e ;; P) \triangleleft (b[[e/x]]) \triangleright_r (x := e ;; Q))$   
 by (*rel-auto*)

**lemma** *assign-r-alt-def*:  
 fixes  $x :: ('a \implies 'c)$   
 shows  $x := v = H[[v]_{<}/\$x]$   
 by (*rel-auto*)

**lemma** *assigns-r-ufunc*: *ufunctional*  $\langle f \rangle_a$   
 by (*rel-auto*)

**lemma** *assigns-r-uinj*:  $\text{inj } f \implies \text{uinj } \langle f \rangle_a$   
 by (*rel-simp, simp add: inj-eq*)

**lemma** *assigns-r-swap-uinj*:  
 $\llbracket \text{vwb-lens } x; \text{vwb-lens } y; x \bowtie y \rrbracket \implies \text{uinj } (x, y := \&y, \&x)$   
 using *assigns-r-uinj swap-usubst-inj*  
 by (*simp add: assigns-r-uinj swap-usubst-inj subst-upd-pr-var*)

**lemma** *assign-unfold*:  
 $\text{vwb-lens } x \implies (x := v) = (\$x' =_u [v]_{<} \wedge H|_{\alpha} x)$   
 apply (*rel-auto, auto simp add: comp-def*)  
 using *vwb-lens.put-eq* by *fastforce*

## 17.8 Converse Laws

**lemma** *convr-invol* [*simp*]:  $p^{--} = p$   
 by *pred-auto*

**lemma** *lit-convr* [*simp*]:  $\llbracket v \rrbracket^{\neg} = \llbracket v \rrbracket$   
 by *pred-auto*

**lemma** *uivar-convr* [*simp*]:

**fixes**  $x :: ('a \implies 'a)$   
**shows**  $(\$x)^- = \$x'$   
**by** *pred-auto*

**lemma** *uovar-convr* [*simp*]:  
**fixes**  $x :: ('a \implies 'a)$   
**shows**  $(\$x')^- = \$x$   
**by** *pred-auto*

**lemma** *uop-convr* [*simp*]:  $(uop\ f\ u)^- = uop\ f\ (u^-)$   
**by** (*pred-auto*)

**lemma** *bop-convr* [*simp*]:  $(bop\ f\ u\ v)^- = bop\ f\ (u^-)\ (v^-)$   
**by** (*pred-auto*)

**lemma** *eq-convr* [*simp*]:  $(p =_u q)^- = (p^- =_u q^-)$   
**by** (*pred-auto*)

**lemma** *not-convr* [*simp*]:  $(\neg p)^- = (\neg p^-)$   
**by** (*pred-auto*)

**lemma** *disj-convr* [*simp*]:  $(p \vee q)^- = (q^- \vee p^-)$   
**by** (*pred-auto*)

**lemma** *conj-convr* [*simp*]:  $(p \wedge q)^- = (q^- \wedge p^-)$   
**by** (*pred-auto*)

**lemma** *seqr-convr* [*simp*]:  $(p ;; q)^- = (q^- ;; p^-)$   
**by** (*rel-auto*)

**lemma** *pre-convr* [*simp*]:  $[p]_{<}^- = [p]_{>}$   
**by** (*rel-auto*)

**lemma** *post-convr* [*simp*]:  $[p]_{>}^- = [p]_{<}$   
**by** (*rel-auto*)

## 17.9 Assertion and Assumption Laws

**declare** *sublens-def* [*lens-defs del*]

**lemma** *assume-false*:  $[false]^\top = false$   
**by** (*rel-auto*)

**lemma** *assume-true*:  $[true]^\top = II$   
**by** (*rel-auto*)

**lemma** *assume-seq*:  $[b]^\top ;; [c]^\top = [b \wedge c]^\top$   
**by** (*rel-auto*)

**lemma** *assert-false*:  $\{false\}_\perp = true$   
**by** (*rel-auto*)

**lemma** *assert-true*:  $\{true\}_\perp = II$   
**by** (*rel-auto*)

**lemma** *assert-seq*:  $\{b\}_\perp ;; \{c\}_\perp = \{b \wedge c\}_\perp$

by (rel-auto)

**lemma** *frame-disj*:  $(x:\llbracket P \rrbracket \vee x:\llbracket Q \rrbracket) = x:\llbracket P \vee Q \rrbracket$   
 by (rel-auto)

**lemma** *frame-seq*:  
 $\llbracket \text{vwb-lens } x; \$x' \# P; \$x \# Q \rrbracket \implies (x:\llbracket P \rrbracket ;; x:\llbracket Q \rrbracket) = x:\llbracket P ;; Q \rrbracket$   
 by (rel-simp, metis vwb-lens-def wb-lens-weak weak-lens.put-get)

**lemma** *antiframe-to-frame*:  
 $\llbracket x \bowtie y; x +_L y = 1_L \rrbracket \implies x:\llbracket P \rrbracket = y:\llbracket P \rrbracket$   
 by (rel-auto, metis lens-indep-def, metis lens-indep-def surj-pair)

**lemma** *antiframe-skip* [simp]:  
 $\text{vwb-lens } x \implies x:\llbracket II \rrbracket = II$   
 by (rel-auto)

**lemma** *antiframe-assign-in*:  
 $\llbracket \text{vwb-lens } a; x \subseteq_L a \rrbracket \implies a:[x := v] = x := v$   
 by (rel-auto, simp-all add: lens-get-put-quasi-commute lens-put-of-quotient)

**lemma** *nameset-skip*:  $\text{vwb-lens } x \implies (ns \ x \cdot II) = II_x$   
 by (rel-auto, meson vwb-lens-wb wb-lens.get-put)

**lemma** *nameset-skip-ra*:  $\text{vwb-lens } x \implies (ns \ x \cdot II_x) = II_x$   
 by (rel-auto)

**declare** *sublens-def* [lens-defs]

## 17.10 While Loop Laws

**theorem** *while-unfold*:  
 $\text{while } b \text{ do } P \text{ od} = ((P ;; \text{while } b \text{ do } P \text{ od}) \triangleleft b \triangleright_r II)$   
**proof** –  
 have  $m:\text{mono } (\lambda X. (P ;; X) \triangleleft b \triangleright_r II)$   
 by (auto intro: monoI segr-mono cond-mono)  
 have  $(\text{while } b \text{ do } P \text{ od}) = (\nu \ X \cdot (P ;; X) \triangleleft b \triangleright_r II)$   
 by (simp add: while-def)  
 also have  $\dots = ((P ;; (\nu \ X \cdot (P ;; X) \triangleleft b \triangleright_r II)) \triangleleft b \triangleright_r II)$   
 by (subst lfp-unfold, simp-all add: m)  
 also have  $\dots = ((P ;; \text{while } b \text{ do } P \text{ od}) \triangleleft b \triangleright_r II)$   
 by (simp add: while-def)  
 finally show ?thesis .  
**qed**

**theorem** *while-false*:  $\text{while false do } P \text{ od} = II$   
 by (subst while-unfold, simp add: aext-false)

**theorem** *while-true*:  $\text{while true do } P \text{ od} = \text{false}$   
 apply (simp add: while-def alpha)  
 apply (rule antisym)  
 apply (simp-all)  
 apply (rule lfp-lowerbound)  
 apply (simp)  
**done**

**theorem** *while-bot-unfold*:  
 $while_{\perp} b \text{ do } P \text{ od} = ((P ;; while_{\perp} b \text{ do } P \text{ od}) \triangleleft b \triangleright_r II)$   
**proof** –  
 have  $m:mono (\lambda X. (P ;; X) \triangleleft b \triangleright_r II)$   
 by (*auto intro: monoI seqr-mono cond-mono*)  
 have  $(while_{\perp} b \text{ do } P \text{ od}) = (\mu X \cdot (P ;; X) \triangleleft b \triangleright_r II)$   
 by (*simp add: while-bot-def*)  
 also have  $\dots = ((P ;; (\mu X \cdot (P ;; X) \triangleleft b \triangleright_r II)) \triangleleft b \triangleright_r II)$   
 by (*subst gfp-unfold, simp-all add: m*)  
 also have  $\dots = ((P ;; while_{\perp} b \text{ do } P \text{ od}) \triangleleft b \triangleright_r II)$   
 by (*simp add: while-bot-def*)  
 finally show *?thesis* .  
**qed**

**theorem** *while-bot-false*:  $while_{\perp} false \text{ do } P \text{ od} = II$   
 by (*simp add: while-bot-def mu-const alpha*)

**theorem** *while-bot-true*:  $while_{\perp} true \text{ do } P \text{ od} = (\mu X \cdot P ;; X)$   
 by (*simp add: while-bot-def alpha*)

An infinite loop with a feasible body corresponds to a program error (non-termination).

**theorem** *while-infinite*:  $P ;; true_h = true \implies while_{\perp} true \text{ do } P \text{ od} = true$   
 apply (*simp add: while-bot-true*)  
 apply (*rule antisym*)  
 apply (*simp*)  
 apply (*rule gfp-upperbound*)  
 apply (*simp*)  
**done**

## 17.11 Algebraic Properties

**interpretation** *upred-semiring*: *semiring-1*  
 where *times* = *seqr* and *one* = *skip-r* and *zero* = *false<sub>h</sub>* and *plus* = *Lattices.sup*  
 by (*unfold-locales, (rel-auto)+*)

We introduce the power syntax derived from semirings

**abbreviation** *upower* :: ' $\alpha \text{ hrel} \Rightarrow \text{nat} \Rightarrow \alpha \text{ hrel}$  (*infixr* ^ 80) **where**  
 $upower P n \equiv upred-semiring.power P n$

**translations**

$P \wedge i \leq CONST power.power II op ;; P i$   
 $P \wedge i \leq (CONST power.power II op ;; P) i$

Set up transfer tactic for powers

**lemma** *upower-rep-eq*:  
 $\llbracket P \wedge i \rrbracket_e = (\lambda b. b \in (\{p. \llbracket P \rrbracket_e p\} \wedge i))$   
**proof** (*induct i arbitrary: P*)  
 case 0  
 then show *?case*  
 by (*auto, rel-auto*)  
**next**  
 case (*Suc i*)  
 show *?case*  
 by (*simp add: Suc seqr.rep-eq relpow-commute*)  
**qed**

```

lemma upower-rep-eq-alt:
   $\llbracket \text{power.power } \langle \text{id} \rangle_a \text{ op } ;; P \ i \rrbracket_e = (\lambda b. b \in (\{p. \llbracket P \rrbracket_e p\} \hat{\ } i))$ 
  by (metis skip-r-def upower-rep-eq)

```

**update-uexpr-rep-eq-thms**

```

lemma Sup-power-expand:
  fixes P :: nat  $\Rightarrow$  'a::complete-lattice
  shows  $P(0) \sqcap (\bigsqcap i. P(i+1)) = (\bigsqcap i. P(i))$ 
proof -
  have UNIV = insert (0::nat) {1..}
    by auto
  moreover have  $(\bigsqcap i. P(i)) = \bigsqcap (P \text{ ` } UNIV)$ 
    by (blast)
  moreover have  $\bigsqcap (P \text{ ` } \text{insert } 0 \text{ } \{1..\}) = P(0) \sqcap SUPREMUM \{1..\} P$ 
    by (simp)
  moreover have  $SUPREMUM \{1..\} P = (\bigsqcap i. P(i+1))$ 
    by (simp add: atLeast-Suc-greaterThan)
  ultimately show ?thesis
    by (simp only:)
qed

```

```

lemma Sup-upto-Suc:  $(\bigsqcap i \in \{0..Suc\ n\}. P \hat{\ } i) = (\bigsqcap i \in \{0..n\}. P \hat{\ } i) \sqcap P \hat{\ } Suc\ n$ 
proof -
  have  $(\bigsqcap i \in \{0..Suc\ n\}. P \hat{\ } i) = (\bigsqcap i \in \text{insert } (Suc\ n) \text{ } \{0..n\}. P \hat{\ } i)$ 
    by (simp add: atLeast0-atMost-Suc)
  also have  $\dots = P \hat{\ } Suc\ n \sqcap (\bigsqcap i \in \{0..n\}. P \hat{\ } i)$ 
    by (simp)
  finally show ?thesis
    by (simp add: Lattices.sup-commute)
qed

```

The following two proofs are adapted from the AFP entry [Kleene Algebra](#). See also [2, 1].

```

lemma upower-inductl:  $Q \sqsubseteq (P ;; Q \sqcap R) \Longrightarrow Q \sqsubseteq P \hat{\ } n ;; R$ 
proof (induct n)
  case 0
  then show ?case by (auto)
next
  case (Suc n)
  then show ?case
    by (auto, metis (no-types, hide-lams) dual-order.trans order-refl seqr-assoc seqr-mono)
qed

```

```

lemma upower-inductr:
  assumes  $Q \sqsubseteq (R \sqcap Q ;; P)$ 
  shows  $Q \sqsubseteq R ;; (P \hat{\ } n)$ 
using assms proof (induct n)
  case 0
  then show ?case by auto
next
  case (Suc n)
  have  $R ;; P \hat{\ } Suc\ n = (R ;; P \hat{\ } n) ;; P$ 
    by (metis seqr-assoc upred-semiring.power-Suc2)
  also have  $Q ;; P \sqsubseteq \dots$ 

```

using *Suc.hyps assms seqr-mono* by *auto*  
 also have  $Q \sqsubseteq \dots$   
 using *assms* by *auto*  
 finally show *?case* .  
 qed

**lemma** *SUP-atLeastAtMost-first*:  
 fixes  $P :: \text{nat} \Rightarrow 'a::\text{complete-lattice}$   
 assumes  $m \leq n$   
 shows  $(\bigcap_{i \in \{m..n\}}. P(i)) = P(m) \sqcap (\bigcap_{i \in \{Suc\ m..n\}}. P(i))$   
 by (*metis SUP-insert assms atLeastAtMost-insertL*)

**lemma** *upower-seqr-iter*:  $P \hat{\ } n = (;; Q : \text{replicate } n\ P \cdot Q)$   
 by (*induct n, simp-all*)

### 17.11.1 Kleene Star

**definition** *ustar* ::  $'\alpha\ hrel \Rightarrow '\alpha\ hrel\ (-^* [999]\ 999)$  **where**  
 $P^* = (\bigcap_{i \in \{0..\}} \cdot P^i)$

**lemma** *ustar-rep-eq*:  
 $\llbracket P^* \rrbracket_e = (\lambda b. b \in (\{p. \llbracket P \rrbracket_e\ p\}^*))$   
 by (*simp add: ustar-def, rel-auto, simp-all add: relpow-imp-rtrancl rtrancl-imp-relpow*)

**update-uexpr-rep-eq-thms**

### 17.11.2 Omega

**definition** *uomega* ::  $'\alpha\ hrel \Rightarrow '\alpha\ hrel\ (-^\omega [999]\ 999)$  **where**  
 $P^\omega = (\mu\ X \cdot P ;; X)$

## 17.12 Relation Algebra Laws

**theorem** *RA1*:  $(P ;; (Q ;; R)) = ((P ;; Q) ;; R)$   
 using *seqr-assoc* by *auto*

**theorem** *RA2*:  $(P ;; II) = P\ (II ;; P) = P$   
 by *simp-all*

**theorem** *RA3*:  $P^{--} = P$   
 by *simp*

**theorem** *RA4*:  $(P ;; Q)^- = (Q^- ;; P^-)$   
 by *simp*

**theorem** *RA5*:  $(P \vee Q)^- = (P^- \vee Q^-)$   
 by (*rel-auto*)

**theorem** *RA6*:  $((P \vee Q) ;; R) = (P ;; R \vee Q ;; R)$   
 using *seqr-or-distl* by *blast*

**theorem** *RA7*:  $((P^- ;; (\neg(P ;; Q))) \vee (\neg Q)) = (\neg Q)$   
 by (*rel-auto*)

### 17.13 Kleene Algebra Laws

**theorem** *ustar-unfoldl*:  $P^* \sqsubseteq II \sqcap P;;P^*$

by (*rel-simp*, *simp add: rtrancl-into-trancl2 trancl-into-rtrancl*)

**theorem** *ustar-inductl*:

assumes  $Q \sqsubseteq R$   $Q \sqsubseteq P;;Q$

shows  $Q \sqsubseteq P^*;;R$

**proof** –

have  $P^*;;R = (\bigsqcap i. P \wedge i;;R)$

by (*simp add: ustar-def UINF-as-Sup-collect' seq-SUP-distr*)

also have  $Q \sqsubseteq \dots$

by (*simp add: SUP-least assms upower-inductl*)

finally show *?thesis* .

qed

**theorem** *ustar-inductr*:

assumes  $Q \sqsubseteq R$   $Q \sqsubseteq Q;;P$

shows  $Q \sqsubseteq R;;P^*$

**proof** –

have  $R;;P^* = (\bigsqcap i. R;;P \wedge i)$

by (*simp add: ustar-def UINF-as-Sup-collect' seq-SUP-distl*)

also have  $Q \sqsubseteq \dots$

by (*simp add: SUP-least assms upower-inductr*)

finally show *?thesis* .

qed

**lemma** *ustar-refines-nu*:  $(\nu X \cdot P;;X \sqcap II) \sqsubseteq P^*$

by (*metis (no-types, lifting) lfp-greatest semilattice-sup-class.le-sup-iff  
semilattice-sup-class.sup-idem upred-semiring.mult-2-right  
upred-semiring.one-add-one ustar-inductl*)

**lemma** *ustar-as-nu*:  $P^* = (\nu X \cdot P;;X \sqcap II)$

**proof** (*rule antisym*)

show  $(\nu X \cdot P;;X \sqcap II) \sqsubseteq P^*$

by (*simp add: ustar-refines-nu*)

show  $P^* \sqsubseteq (\nu X \cdot P;;X \sqcap II)$

by (*metis lfp-lowerbound upred-semiring.add-commute ustar-unfoldl*)

qed

### 17.14 Omega Algebra Laws

**lemma** *uomega-induct*:

$P;;P^\omega \sqsubseteq P^\omega$

by (*simp add: uomega-def, metis eq-refl gfp-unfold monoI seqr-mono*)

end

### 17.15 Relational Hoare calculus

**theory** *utp-hoare*

**imports** *utp-rel-laws*

**begin**

**named-theorems** *hoare* **and** *hoare-safe*

**method** *hoare-split* **uses** *hr* =  
 ((*simp add: assigns-r-comp usubst unrest*)?, — Eliminate assignments where possible  
 (*auto*  
   *intro: hoare intro!: hoare-safe hr*  
   *simp add: assigns-r-comp conj-comm conj-assoc usubst unrest*))[1] — Apply Hoare logic laws

**method** *hoare-auto* **uses** *hr* = (*hoare-split hr: hr; rel-auto?*)

**definition** *hoare-r* :: ' $\alpha$  cond  $\Rightarrow$  ' $\alpha$  hrel  $\Rightarrow$  ' $\alpha$  cond  $\Rightarrow$  bool ( $\{-\}/ -/ \{-\}_u$ ) **where**  
 $\{p\}Q\{r\}_u = (([p]_< \Rightarrow [r]_>) \sqsubseteq Q)$

**declare** *hoare-r-def* [*upred-defs*]

**lemma** *hoare-r-conj* [*hoare-safe*]:  $\llbracket \{p\}Q\{r\}_u; \{p\}Q\{s\}_u \rrbracket \Longrightarrow \{p\}Q\{r \wedge s\}_u$   
**by** *rel-auto*

**lemma** *hoare-r-weaken-pre* [*hoare*]:  
 $\{p\}Q\{r\}_u \Longrightarrow \{p \wedge q\}Q\{r\}_u$   
 $\{q\}Q\{r\}_u \Longrightarrow \{p \wedge q\}Q\{r\}_u$   
**by** *rel-auto+*

**lemma** *hoare-r-conseq*:  $\llbracket 'p_1 \Rightarrow p_2'; \{p_2\}S\{q_2\}_u; 'q_2 \Rightarrow q_1' \rrbracket \Longrightarrow \{p_1\}S\{q_1\}_u$   
**by** *rel-auto*

**lemma** *assigns-hoare-r* [*hoare-safe*]: ' $p \Rightarrow \sigma \dagger q$ '  $\Longrightarrow \{p\}\langle\sigma\rangle_a\{q\}_u$   
**by** *rel-auto*

**lemma** *skip-hoare-r* [*hoare-safe*]:  $\{p\}II\{p\}_u$   
**by** *rel-auto*

**lemma** *seq-hoare-r*:  $\llbracket \{p\}Q_1\{s\}_u; \{s\}Q_2\{r\}_u \rrbracket \Longrightarrow \{p\}Q_1;; Q_2\{r\}_u$   
**by** *rel-auto*

**lemma** *seq-hoare-invariant* [*hoare-safe*]:  $\llbracket \{p\}Q_1\{p\}_u; \{p\}Q_2\{p\}_u \rrbracket \Longrightarrow \{p\}Q_1;; Q_2\{p\}_u$   
**by** *rel-auto*

**lemma** *seq-hoare-stronger-pre-1* [*hoare-safe*]:  
 $\llbracket \{p \wedge q\}Q_1\{p \wedge q\}_u; \{p \wedge q\}Q_2\{q\}_u \rrbracket \Longrightarrow \{p \wedge q\}Q_1;; Q_2\{q\}_u$   
**by** *rel-auto*

**lemma** *seq-hoare-stronger-pre-2* [*hoare-safe*]:  
 $\llbracket \{p \wedge q\}Q_1\{p \wedge q\}_u; \{p \wedge q\}Q_2\{p\}_u \rrbracket \Longrightarrow \{p \wedge q\}Q_1;; Q_2\{p\}_u$   
**by** *rel-auto*

**lemma** *seq-hoare-inv-r-2* [*hoare*]:  $\llbracket \{p\}Q_1\{q\}_u; \{q\}Q_2\{q\}_u \rrbracket \Longrightarrow \{p\}Q_1;; Q_2\{q\}_u$   
**by** *rel-auto*

**lemma** *seq-hoare-inv-r-3* [*hoare*]:  $\llbracket \{p\}Q_1\{p\}_u; \{p\}Q_2\{q\}_u \rrbracket \Longrightarrow \{p\}Q_1;; Q_2\{q\}_u$   
**by** *rel-auto*

**lemma** *cond-hoare-r* [*hoare-safe*]:  $\llbracket \{b \wedge p\}S\{q\}_u; \{\neg b \wedge p\}T\{q\}_u \rrbracket \Longrightarrow \{p\}S \triangleleft b \triangleright_r T\{q\}_u$   
**by** *rel-auto*

Frame rule: If starting  $S$  in a state satisfying *peestablishesq* in the final state, then we can insert an invariant predicate  $r$  when  $S$  is framed by  $a$ , provided that  $r$  does not refer to variables in the frame, and  $q$  does not refer to variables outside the frame.



```

lemma frame-hoare-r [hoare-safe]:
  assumes vwb-lens a a  $\#$  r a  $\Vdash$  q  $\llbracket p \wedge r \rrbracket S \llbracket q \rrbracket_u$ 
  shows  $\llbracket p \wedge r \rrbracket a.[S] \llbracket q \wedge r \rrbracket_u$ 
  using assms by (rel-simp)

lemma frame-hoare-r' [hoare-safe]:
  assumes vwb-lens a a  $\#$  r a  $\Vdash$  q  $\llbracket r \wedge p \rrbracket S \llbracket q \rrbracket_u$ 
  shows  $\llbracket r \wedge p \rrbracket a.[S] \llbracket r \wedge q \rrbracket_u$ 
  using assms
  by (simp add: frame-hoare-r utp-pred-laws.inf commute)

lemma while-hoare-r [hoare-safe]:
  assumes  $\llbracket p \wedge b \rrbracket S \llbracket p \rrbracket_u$ 
  shows  $\llbracket p \rrbracket \text{while } b \text{ do } S \text{ od } \llbracket \neg b \wedge p \rrbracket_u$ 
  using assms
  by (simp add: while-def hoare-r-def, rule-tac lfp-lowerbound) (rel-auto)

lemma while-invr-hoare-r [hoare-safe]:
  assumes  $\llbracket p \wedge b \rrbracket S \llbracket p \rrbracket_u$  ‘pre  $\Rightarrow$  p’ ‘ $(\neg b \wedge p) \Rightarrow$  post’
  shows  $\llbracket \text{pre} \rrbracket \text{while } b \text{ invr } p \text{ do } S \text{ od } \llbracket \text{post} \rrbracket_u$ 
  by (metis assms hoare-r-conseq while-hoare-r while-inv-def)
end

```

## 17.16 Weakest precondition calculus

```

theory utp-wp
imports utp-hoare
begin

```

A very quick implementation of wp – more laws still needed!

**named-theorems** *wp*

**method** *wp-tac* = (*simp add: wp*)

```

consts
  uwp :: 'a  $\Rightarrow$  'b  $\Rightarrow$  'c (infix wp 60)

```

**definition** *wp-upred* :: (' $\alpha$ , ' $\beta$ ) *rel*  $\Rightarrow$  ' $\beta$  *cond*  $\Rightarrow$  ' $\alpha$  *cond* **where**  
*wp-upred* *Q* *r* =  $\lfloor \neg (Q \;; (\neg \lceil r \rceil_{<})) \rfloor :: ('\alpha, '\beta) \text{ rel} \rfloor_{<}$

```

adhoc-overloading
  uwp wp-upred

```

**declare** *wp-upred-def* [*urel-defs*]

```

lemma wp-true [wp]: p wp true = true
  by (rel-simp)

```

```

theorem wp-assigns-r [wp]:
   $\langle \sigma \rangle_a \text{ wp } r = \sigma \uparrow r$ 
  by rel-auto

```

```

theorem wp-skip-r [wp]:
  II wp r = r
  by rel-auto

```

**theorem** *wp-abort* [wp]:  
 $r \neq \text{true} \implies \text{true wp } r = \text{false}$   
**by** *rel-auto*

**theorem** *wp-conj* [wp]:  
 $P \text{ wp } (q \wedge r) = (P \text{ wp } q \wedge P \text{ wp } r)$   
**by** *rel-auto*

**theorem** *wp-seq-r* [wp]:  $(P ;; Q) \text{ wp } r = P \text{ wp } (Q \text{ wp } r)$   
**by** *rel-auto*

**theorem** *wp-cond* [wp]:  $(P \triangleleft b \triangleright_r Q) \text{ wp } r = ((b \Rightarrow P \text{ wp } r) \wedge ((\neg b) \Rightarrow Q \text{ wp } r))$   
**by** *rel-auto*

**theorem** *wp-hoare-link*:  
 $\{p\} Q \{r\}_u \longleftrightarrow (Q \text{ wp } r \sqsubseteq p)$   
**by** *rel-auto*

If two programs have the same weakest precondition for any postcondition then the programs are the same.

**theorem** *wp-eq-intro*:  $\llbracket \bigwedge r. P \text{ wp } r = Q \text{ wp } r \rrbracket \implies P = Q$   
**by** (*rel-auto robust, fastforce+*)  
**end**

## 18 UTP Theories

**theory** *utp-theory*  
**imports** *utp-rel-laws*  
**begin**

Closure laws for theories

**named-theorems** *closure*

### 18.1 Complete lattice of predicates

**definition** *upred-lattice* ::  $('a \text{ upred}) \text{ gorder } (\mathcal{P})$  **where**  
 $\text{upred-lattice} = (\text{carrier} = \text{UNIV}, \text{eq} = (\text{op} =), \text{le} = \text{op} \sqsubseteq)$

$\mathcal{P}$  is the complete lattice of alphabetised predicates. All other theories will be defined relative to it.

**interpretation** *upred-lattice*: *complete-lattice*  $\mathcal{P}$   
**proof** (*unfold-locales, simp-all add: upred-lattice-def*)  
**fix**  $A :: 'a \text{ upred set}$   
**show**  $\exists s. \text{is-lub } (\text{carrier} = \text{UNIV}, \text{eq} = \text{op} =, \text{le} = \text{op} \sqsubseteq) s A$   
**apply** (*rule-tac x= $\sqcup$  A in exI*)  
**apply** (*rule least-UpperI*)  
**apply** (*auto intro: Inf-greatest simp add: Inf-lower Upper-def*)  
**done**  
**show**  $\exists i. \text{is-glb } (\text{carrier} = \text{UNIV}, \text{eq} = \text{op} =, \text{le} = \text{op} \sqsubseteq) i A$   
**apply** (*rule-tac x= $\sqcap$  A in exI*)  
**apply** (*rule greatest-LowerI*)  
**apply** (*auto intro: Sup-least simp add: Sup-upper Lower-def*)  
**done**  
**qed**

**lemma** *upred-weak-complete-lattice* [simp]: *weak-complete-lattice*  $\mathcal{P}$   
**by** (simp add: upred-lattice.weak.weak-complete-lattice-axioms)

**lemma** *upred-lattice-eq* [simp]:  
 $op \text{ .}=\mathcal{P} = op =$   
**by** (simp add: upred-lattice-def)

**lemma** *upred-lattice-le* [simp]:  
 $le \mathcal{P} P Q = (P \sqsubseteq Q)$   
**by** (simp add: upred-lattice-def)

**lemma** *upred-lattice-carrier* [simp]:  
 $carrier \mathcal{P} = UNIV$   
**by** (simp add: upred-lattice-def)

## 18.2 Healthiness conditions

**type-synonym**  $'\alpha \text{ health} = '\alpha \text{ upred} \Rightarrow '\alpha \text{ upred}$

**definition**  
 $Healthy::'\alpha \text{ upred} \Rightarrow '\alpha \text{ health} \Rightarrow \text{bool}$  (**infix** is 30)  
**where**  $P \text{ is } H \equiv (H P = P)$

**lemma** *Healthy-def'*:  $P \text{ is } H \longleftrightarrow (H P = P)$   
**unfolding** *Healthy-def* **by** auto

**lemma** *Healthy-if*:  $P \text{ is } H \Longrightarrow (H P = P)$   
**unfolding** *Healthy-def* **by** auto

**declare** *Healthy-def'* [upred-defs]

**abbreviation** *Healthy-carrier* ::  $'\alpha \text{ health} \Rightarrow '\alpha \text{ upred set}$  ( $\llbracket - \rrbracket_H$ )  
**where**  $\llbracket H \rrbracket_H \equiv \{P. P \text{ is } H\}$

**lemma** *Healthy-carrier-image*:  
 $A \subseteq \llbracket \mathcal{H} \rrbracket_H \Longrightarrow \mathcal{H} \text{ ' } A = A$   
**by** (auto simp add: image-def, (metis *Healthy-if mem-Collect-eq subsetCE*)+)

**lemma** *Healthy-carrier-Collect*:  $A \subseteq \llbracket H \rrbracket_H \Longrightarrow A = \{H(P) \mid P. P \in A\}$   
**by** (simp add: *Healthy-carrier-image Setcompr-eq-image*)

**lemma** *Healthy-func*:  
 $\llbracket F \in \llbracket \mathcal{H}_1 \rrbracket_H \rightarrow \llbracket \mathcal{H}_2 \rrbracket_H; P \text{ is } \mathcal{H}_1 \rrbracket \Longrightarrow \mathcal{H}_2(F(P)) = F(P)$   
**using** *Healthy-if* **by** blast

**lemma** *Healthy-apply-closed*:  
**assumes**  $F \in \llbracket H \rrbracket_H \rightarrow \llbracket H \rrbracket_H P \text{ is } H$   
**shows**  $F(P) \text{ is } H$   
**using** *assms(1) assms(2)* **by** auto

**lemma** *Healthy-set-image-member*:  
 $\llbracket P \in F \text{ ' } A; \bigwedge x. F x \text{ is } H \rrbracket \Longrightarrow P \text{ is } H$   
**by** blast

**lemma** *Healthy-SUPREMUM*:

$A \subseteq \llbracket H \rrbracket_H \implies \text{SUPREMUM } A \ H = \bigcap A$   
**by** (*drule Healthy-carrier-image, presburger*)

**lemma** *Healthy-INFIMUM*:

$A \subseteq \llbracket H \rrbracket_H \implies \text{INFIMUM } A \ H = \bigcup A$   
**by** (*drule Healthy-carrier-image, presburger*)

**lemma** *Healthy-nu [closure]*:

**assumes** *mono F F*  $F \in \llbracket id \rrbracket_H \rightarrow \llbracket H \rrbracket_H$   
**shows**  $\nu F$  *is H*  
**by** (*metis (mono-tags) Healthy-def Healthy-func assms eq-id-iff lfp-unfold*)

**lemma** *Healthy-mu [closure]*:

**assumes** *mono F F*  $F \in \llbracket id \rrbracket_H \rightarrow \llbracket H \rrbracket_H$   
**shows**  $\mu F$  *is H*  
**by** (*metis (mono-tags) Healthy-def Healthy-func assms eq-id-iff gfp-unfold*)

**lemma** *Healthy-subset-member*:  $\llbracket A \subseteq \llbracket H \rrbracket_H; P \in A \rrbracket \implies H(P) = P$

**by** (*meson Ball-Collect Healthy-if*)

**lemma** *is-Healthy-subset-member*:  $\llbracket A \subseteq \llbracket H \rrbracket_H; P \in A \rrbracket \implies P$  *is H*

**by** *blast*

### 18.3 Properties of healthiness conditions

**definition** *Idempotent* ::  $'\alpha$  *health*  $\Rightarrow$  *bool* **where**

$\text{Idempotent}(H) \longleftrightarrow (\forall P. H(H(P)) = H(P))$

**abbreviation** *Monotonic* ::  $'\alpha$  *health*  $\Rightarrow$  *bool* **where**

$\text{Monotonic}(H) \equiv \text{mono } H$

**definition** *IMH* ::  $'\alpha$  *health*  $\Rightarrow$  *bool* **where**

$\text{IMH}(H) \longleftrightarrow \text{Idempotent}(H) \wedge \text{Monotonic}(H)$

**definition** *Antitone* ::  $'\alpha$  *health*  $\Rightarrow$  *bool* **where**

$\text{Antitone}(H) \longleftrightarrow (\forall P \ Q. Q \sqsubseteq P \longrightarrow (H(P) \sqsubseteq H(Q)))$

**definition** *Conjunctive* ::  $'\alpha$  *health*  $\Rightarrow$  *bool* **where**

$\text{Conjunctive}(H) \longleftrightarrow (\exists Q. \forall P. H(P) = (P \wedge Q))$

**definition** *FunctionalConjunctive* ::  $'\alpha$  *health*  $\Rightarrow$  *bool* **where**

$\text{FunctionalConjunctive}(H) \longleftrightarrow (\exists F. \forall P. H(P) = (P \wedge F(P)) \wedge \text{Monotonic}(F))$

**definition** *WeakConjunctive* ::  $'\alpha$  *health*  $\Rightarrow$  *bool* **where**

$\text{WeakConjunctive}(H) \longleftrightarrow (\forall P. \exists Q. H(P) = (P \wedge Q))$

**definition** *Disjunctuous* ::  $'\alpha$  *health*  $\Rightarrow$  *bool* **where**

[*upred-defs*]:  $\text{Disjunctuous } H = (\forall P \ Q. H(P \sqcap Q) = (H(P) \sqcap H(Q)))$

**definition** *Continuous* ::  $'\alpha$  *health*  $\Rightarrow$  *bool* **where**

[*upred-defs*]:  $\text{Continuous } H = (\forall A. A \neq \{\} \longrightarrow H(\bigcap A) = \bigcap (H \restriction A))$

**lemma** *Healthy-Idempotent [closure]*:

$\text{Idempotent } H \implies H(P)$  *is H*

**by** (*simp add: Healthy-def Idempotent-def*)

**lemma** *Healthy-range*:  $\text{Idempotent } H \implies \text{range } H = \llbracket H \rrbracket_H$   
 by (auto simp add: image-def Healthy-if Healthy-Idempotent, metis Healthy-if)

**lemma** *Idempotent-id* [simp]: *Idempotent id*  
 by (simp add: Idempotent-def)

**lemma** *Idempotent-comp* [intro]:  
 $\llbracket \text{Idempotent } f; \text{Idempotent } g; f \circ g = g \circ f \rrbracket \implies \text{Idempotent } (f \circ g)$   
 by (auto simp add: Idempotent-def comp-def, metis)

**lemma** *Idempotent-image*:  $\text{Idempotent } f \implies f \text{ ' } f \text{ ' } A = f \text{ ' } A$   
 by (metis (mono-tags, lifting) Idempotent-def image-cong image-image)

**lemma** *Monotonic-id* [simp]: *Monotonic id*  
 by (simp add: monoI)

**lemma** *Monotonic-comp* [intro]:  
 $\llbracket \text{Monotonic } f; \text{Monotonic } g \rrbracket \implies \text{Monotonic } (f \circ g)$   
 by (simp add: mono-def)

**lemma** *Conjunctive-Idempotent*:  
 $\text{Conjunctive}(H) \implies \text{Idempotent}(H)$   
 by (auto simp add: Conjunctive-def Idempotent-def)

**lemma** *Conjunctive-Monotonic*:  
 $\text{Conjunctive}(H) \implies \text{Monotonic}(H)$   
**unfolding** *Conjunctive-def* *mono-def*  
**using** *dual-order.trans* **by** *fastforce*

**lemma** *Conjunctive-conj*:  
**assumes** *Conjunctive*(*HC*)  
**shows**  $\text{HC}(P \wedge Q) = (\text{HC}(P) \wedge Q)$   
**using** *assms* **unfolding** *Conjunctive-def*  
**by** (metis *utp-pred-laws.inf.assoc utp-pred-laws.inf.commute*)

**lemma** *Conjunctive-distr-conj*:  
**assumes** *Conjunctive*(*HC*)  
**shows**  $\text{HC}(P \wedge Q) = (\text{HC}(P) \wedge \text{HC}(Q))$   
**using** *assms* **unfolding** *Conjunctive-def*  
**by** (metis *Conjunctive-conj assms utp-pred-laws.inf.assoc utp-pred-laws.inf-right-idem*)

**lemma** *Conjunctive-distr-disj*:  
**assumes** *Conjunctive*(*HC*)  
**shows**  $\text{HC}(P \vee Q) = (\text{HC}(P) \vee \text{HC}(Q))$   
**using** *assms* **unfolding** *Conjunctive-def*  
**using** *utp-pred-laws.inf-sup-distrib2* **by** *fastforce*

**lemma** *Conjunctive-distr-cond*:  
**assumes** *Conjunctive*(*HC*)  
**shows**  $\text{HC}(P \triangleleft b \triangleright Q) = (\text{HC}(P) \triangleleft b \triangleright \text{HC}(Q))$   
**using** *assms* **unfolding** *Conjunctive-def*  
**by** (metis *cond-conj-distr utp-pred-laws.inf-commute*)

**lemma** *FunctionalConjunctive-Monotonic*:  
 $\text{FunctionalConjunctive}(H) \implies \text{Monotonic}(H)$

**unfolding** *FunctionalConjunctive-def* **by** (*metis mono-def utp-pred-laws.inf-mono*)

**lemma** *WeakConjunctive-Refinement*:  
**assumes** *WeakConjunctive(HC)*  
**shows**  $P \sqsubseteq HC(P)$   
**using** *assms unfolding WeakConjunctive-def* **by** (*metis utp-pred-laws.inf.cobounded1*)

**lemma** *WeakConjunctive-Healthy-Refinement*:  
**assumes** *WeakConjunctive(HC)* **and** *P is HC*  
**shows**  $HC(P) \sqsubseteq P$   
**using** *assms unfolding WeakConjunctive-def Healthy-def* **by** *simp*

**lemma** *WeakConjunctive-implies-WeakConjunctive*:  
*Conjunctive(H)  $\implies$  WeakConjunctive(H)*  
**unfolding** *WeakConjunctive-def Conjunctive-def* **by** *pred-auto*

**declare** *Conjunctive-def* [*upred-defs*]  
**declare** *mono-def* [*upred-defs*]

**lemma** *Disjunctuous-Monotonic*: *Disjunctuous H  $\implies$  Monotonic H*  
**by** (*metis Disjunctuous-def mono-def semilattice-sup-class.le-iff-sup*)

**lemma** *ContinuousD* [*dest*]:  $\llbracket \text{Continuous } H; A \neq \{\} \rrbracket \implies H (\bigsqcap A) = (\bigsqcap P \in A. H(P))$   
**by** (*simp add: Continuous-def*)

**lemma** *Continuous-Disjunctuous*: *Continuous H  $\implies$  Disjunctuous H*  
**apply** (*auto simp add: Continuous-def Disjunctuous-def*)  
**apply** (*rename-tac P Q*)  
**apply** (*drule-tac x={P,Q} in spec*)  
**apply** (*simp*)  
**done**

**lemma** *Continuous-Monotonic* [*closure*]: *Continuous H  $\implies$  Monotonic H*  
**by** (*simp add: Continuous-Disjunctuous Disjunctuous-Monotonic*)

**lemma** *Continuous-comp* [*intro*]:  
 $\llbracket \text{Continuous } f; \text{Continuous } g \rrbracket \implies \text{Continuous } (f \circ g)$   
**by** (*simp add: Continuous-def*)

Closure laws derived from continuity

**lemma** *Sup-Continuous-closed* [*closure*]:  
 $\llbracket \text{Continuous } H; \bigwedge i. i \in A \implies P(i) \text{ is } H; A \neq \{\} \rrbracket \implies (\bigsqcap i \in A. P(i)) \text{ is } H$   
**by** (*drule ContinuousD[of H P 'A], simp add: UINF-mem-UNIV[THEN sym] UINF-as-Sup[THEN sym]*)  
*(metis (no-types, lifting) Healthy-def' SUP-cong image-image)*

**lemma** *UINF-mem-Continuous-closed* [*closure*]:  
 $\llbracket \text{Continuous } H; \bigwedge i. i \in A \implies P(i) \text{ is } H; A \neq \{\} \rrbracket \implies (\bigsqcap i \in A. P(i)) \text{ is } H$   
**by** (*simp add: Sup-Continuous-closed UINF-as-Sup-collect*)

**lemma** *UINF-mem-Continuous-closed-pair* [*closure*]:  
**assumes** *Continuous H  $\bigwedge i j. (i, j) \in A \implies P i j \text{ is } H A \neq \{\}$*   
**shows**  $(\bigsqcap (i,j) \in A. P i j) \text{ is } H$   
**proof** –  
**have**  $(\bigsqcap (i,j) \in A. P i j) = (\bigsqcap x \in A. P (\text{fst } x) (\text{snd } x))$

```

    by (rel-auto)
  also have ... is H
    by (metis (mono-tags) UINF-mem-Continuous-closed assms(1) assms(2) assms(3) prod.collapse)
  finally show ?thesis .
qed

```

```

lemma UINF-mem-Continuous-closed-triple [closure]:
  assumes Continuous H  $\bigwedge i j k. (i, j, k) \in A \implies P i j k \text{ is } H$  A  $\neq \{\}$ 
  shows  $(\bigcap (i,j,k) \in A \cdot P i j k) \text{ is } H$ 
proof -
  have  $(\bigcap (i,j,k) \in A \cdot P i j k) = (\bigcap x \in A \cdot P (fst x) (fst (snd x)) (snd (snd x)))$ 
    by (rel-auto)
  also have ... is H
    by (metis (mono-tags) UINF-mem-Continuous-closed assms(1) assms(2) assms(3) prod.collapse)
  finally show ?thesis .
qed

```

```

lemma UINF-Continuous-closed [closure]:
   $\llbracket \text{Continuous } H; \bigwedge i. P(i) \text{ is } H \rrbracket \implies (\bigcap i \cdot P(i)) \text{ is } H$ 
  using UINF-mem-Continuous-closed[of H UNIV P]
  by (simp add: UINF-mem-UNIV)

```

All continuous functions are also Scott-continuous

```

lemma sup-continuous-Continuous [closure]: Continuous F  $\implies$  sup-continuous F
  by (simp add: Continuous-def sup-continuous-def)

```

```

lemma Healthy-fixed-points [simp]: fps  $\mathcal{P}$  H =  $\llbracket H \rrbracket_H$ 
  by (simp add: fps-def upred-lattice-def Healthy-def)

```

```

lemma USUP-healthy:  $A \subseteq \llbracket H \rrbracket_H \implies (\bigsqcup P \in A \cdot F(P)) = (\bigsqcup P \in A \cdot F(H(P)))$ 
  by (rule USUP-cong, simp add: Healthy-subset-member)

```

```

lemma UINF-healthy:  $A \subseteq \llbracket H \rrbracket_H \implies (\bigcap P \in A \cdot F(P)) = (\bigcap P \in A \cdot F(H(P)))$ 
  by (rule UINF-cong, simp add: Healthy-subset-member)

```

```

lemma upred-lattice-Idempotent [simp]: Idem $\mathcal{P}$  H = Idempotent H
  using upred-lattice.weak-partial-order-axioms by (auto simp add: idempotent-def Idempotent-def)

```

```

lemma upred-lattice-Monotonic [simp]: Mono $\mathcal{P}$  H = Monotonic H
  using upred-lattice.weak-partial-order-axioms by (auto simp add: isotone-def mono-def)

```

## 18.4 UTP theories hierarchy

```

typedef (' $\mathcal{T}$ , ' $\alpha$ ) uthy = UNIV :: unit set
  by auto

```

We create a unitary parametric type to represent UTP theories. These are merely tags and contain no data other than to help the type-system resolve polymorphic definitions. The two parameters denote the name of the UTP theory – as a unique type – and the minimal alphabet that the UTP theory requires. We will then use Isabelle’s ad-hoc overloading mechanism to associate theory constructs, like healthiness conditions and units, with each of these types. This will allow the type system to retrieve definitions based on a particular theory context.

```

definition uthy :: ('a, 'b) uthy where
  uthy = Abs-uthy ()

```

**lemma** *uthy-eq* [intro]:  
**fixes**  $x\ y :: ('a, 'b)\ \text{uthy}$   
**shows**  $x = y$   
**by** (*cases x, cases y, simp*)

**syntax**  
 $-UTHY :: \text{type} \Rightarrow \text{type} \Rightarrow \text{logic}\ (UTHY'(-, -))$

**translations**  
 $UTHY('T, 'a) == \text{CONST}\ \text{uthy} :: ('T, 'a)\ \text{uthy}$

We set up polymorphic constants to denote the healthiness conditions associated with a UTP theory. Unfortunately we can currently only characterise UTP theories of homogeneous relations; this is due to restrictions in the instantiation of Isabelle's polymorphic constants which apparently cannot specialise types in this way.

**consts**  
 $\text{utp-hcond} :: ('T, 'a)\ \text{uthy} \Rightarrow ('a \times 'a)\ \text{health}\ (\mathcal{H}_1)$

**definition** *utp-order* ::  $('a \times 'a)\ \text{health} \Rightarrow 'a\ \text{hrel}\ \text{gorder}$  **where**  
 $\text{utp-order}\ H = (\mid \text{carrier} = \{P.\ P\ \text{is}\ H\}, \text{eq} = (\text{op} =), \text{le} = \text{op} \sqsubseteq \mid)$

**abbreviation** *uthy-order*  $T \equiv \text{utp-order}\ \mathcal{H}_T$

Constant *utp-order* obtains the order structure associated with a UTP theory. Its carrier is the set of healthy predicates, equality is HOL equality, and the order is refinement.

**lemma** *utp-order-carrier* [simp]:  
 $\text{carrier}\ (\text{utp-order}\ H) = \llbracket H \rrbracket_H$   
**by** (*simp add: utp-order-def*)

**lemma** *utp-order-eq* [simp]:  
 $\text{eq}\ (\text{utp-order}\ T) = \text{op} =$   
**by** (*simp add: utp-order-def*)

**lemma** *utp-order-le* [simp]:  
 $\text{le}\ (\text{utp-order}\ T) = \text{op} \sqsubseteq$   
**by** (*simp add: utp-order-def*)

**lemma** *utp-partial-order*:  $\text{partial-order}\ (\text{utp-order}\ T)$   
**by** (*unfold-locales, simp-all add: utp-order-def*)

**lemma** *utp-weak-partial-order*:  $\text{weak-partial-order}\ (\text{utp-order}\ T)$   
**by** (*unfold-locales, simp-all add: utp-order-def*)

**lemma** *mono-Monotone-utp-order*:  
 $\text{mono}\ f \Longrightarrow \text{Monotone}\ (\text{utp-order}\ T)\ f$   
**apply** (*auto simp add: isotone-def*)  
**apply** (*metis partial-order-def utp-partial-order*)  
**apply** (*metis monoD*)  
**done**

**lemma** *isotone-utp-orderI*:  $\text{Monotonic}\ H \Longrightarrow \text{isotone}\ (\text{utp-order}\ X)\ (\text{utp-order}\ Y)\ H$   
**by** (*auto simp add: mono-def isotone-def utp-weak-partial-order*)

**lemma** *Mono-utp-orderI*:  
 $\llbracket \bigwedge P\ Q.\ \llbracket P \sqsubseteq Q; P\ \text{is}\ H; Q\ \text{is}\ H \rrbracket \Longrightarrow F(P) \sqsubseteq F(Q) \rrbracket \Longrightarrow \text{Mono}_{\text{utp-order}\ H}\ F$



by (auto simp add: isotone-def utp-weak-partial-order)

The UTP order can equivalently be characterised as the fixed point lattice, *fpl*.

**lemma** *utp-order-fpl*:  $utp\text{-}order\ H = fpl\ \mathcal{P}\ H$   
 by (auto simp add: utp-order-def upred-lattice-def fps-def Healthy-def)

**definition** *uth-eq* ::  $('T_1, 'α) uthy \Rightarrow ('T_2, 'α) uthy \Rightarrow bool$  (**infix**  $\approx_T$  50) **where**  
 $T_1 \approx_T T_2 \longleftrightarrow \llbracket \mathcal{H}_{T_1} \rrbracket_H = \llbracket \mathcal{H}_{T_2} \rrbracket_H$

**lemma** *uth-eq-refl*:  $T \approx_T T$   
 by (simp add: uth-eq-def)

**lemma** *uth-eq-sym*:  $T_1 \approx_T T_2 \longleftrightarrow T_2 \approx_T T_1$   
 by (auto simp add: uth-eq-def)

**lemma** *uth-eq-trans*:  $\llbracket T_1 \approx_T T_2; T_2 \approx_T T_3 \rrbracket \Longrightarrow T_1 \approx_T T_3$   
 by (auto simp add: uth-eq-def)

**definition** *uthy-plus* ::  $('T_1, 'α) uthy \Rightarrow ('T_2, 'α) uthy \Rightarrow ('T_1 \times 'T_2, 'α) uthy$  (**infixl**  $+_T$  65) **where**  
 $uthy\text{-}plus\ T_1\ T_2 = uthy$

**overloading**

*prod-hcond* == *utp-hcond* ::  $('T_1 \times 'T_2, 'α) uthy \Rightarrow ('α \times 'α) health$   
**begin**

The healthiness condition of a relation is simply identity, since every alphabetised relation is healthy.

**definition** *prod-hcond* ::  $('T_1 \times 'T_2, 'α) uthy \Rightarrow ('α \times 'α) upred \Rightarrow ('α \times 'α) upred$  **where**  
 $prod\text{-}hcond\ T = \mathcal{H}_{UTHY}('T_1, 'α) \circ \mathcal{H}_{UTHY}('T_2, 'α)$

**end**

## 18.5 UTP theory hierarchy

We next define a hierarchy of locales that characterise different classes of UTP theory. Minimally we require that a UTP theory's healthiness condition is idempotent.

**locale** *utp-theory* =  
 fixes  $\mathcal{T} :: ('T, 'α) uthy$  (**structure**)  
 assumes *HCond-Idem*:  $\mathcal{H}(\mathcal{H}(P)) = \mathcal{H}(P)$   
**begin**

**lemma** *uthy-simp*:  
 $uthy = \mathcal{T}$   
 by blast

A UTP theory fixes  $\mathcal{T}$ , the structural element denoting the UTP theory. All constants associated with UTP theories can then be resolved by the type system.

**lemma** *HCond-Idempotent* [*closure,intro*]: *Idempotent*  $\mathcal{H}$   
 by (simp add: Idempotent-def HCond-Idem)

**sublocale** *partial-order uthy-order*  $\mathcal{T}$   
 by (unfold-locales, simp-all add: utp-order-def)  
**end**

Theory summation is commutative provided the healthiness conditions commute.

**lemma** *uthy-plus-comm*:

**assumes**  $\mathcal{H}_{T_1} \circ \mathcal{H}_{T_2} = \mathcal{H}_{T_2} \circ \mathcal{H}_{T_1}$

**shows**  $T_1 +_T T_2 \approx_T T_2 +_T T_1$

**proof** –

**have**  $T_1 = \text{uthy } T_2 = \text{uthy}$

**by** *blast+*

**thus** *?thesis*

**using** *assms* **by** (*simp add: uth-eq-def prod-hcond-def*)

**qed**

**lemma** *uthy-plus-assoc*:  $T_1 +_T (T_2 +_T T_3) \approx_T (T_1 +_T T_2) +_T T_3$

**by** (*simp add: uth-eq-def prod-hcond-def comp-def*)

**lemma** *uthy-plus-idem*:  $\text{utp-theory } T \implies T +_T T \approx_T T$

**by** (*simp add: uth-eq-def prod-hcond-def Healthy-def utp-theory.HCond-Idem utp-theory.uthy-simp*)

**locale** *utp-theory-lattice* = *utp-theory*  $\mathcal{T}$  + *complete-lattice* *uthy-order*  $\mathcal{T}$  **for**  $\mathcal{T} :: ('T, 'a) \text{uthy}$  (**structure**)

The healthiness conditions of a UTP theory lattice form a complete lattice, and allows us to make use of complete lattice results from HOL-Algebra, such as the Knaster-Tarski theorem. We can also retrieve lattice operators as below.

**abbreviation** *utp-top* ( $\top_1$ )

**where** *utp-top*  $\mathcal{T} \equiv \text{top } (\text{uthy-order } \mathcal{T})$

**abbreviation** *utp-bottom* ( $\perp_1$ )

**where** *utp-bottom*  $\mathcal{T} \equiv \text{bottom } (\text{uthy-order } \mathcal{T})$

**abbreviation** *utp-join* (**infixl**  $\sqcup_1$  65) **where**

*utp-join*  $\mathcal{T} \equiv \text{join } (\text{uthy-order } \mathcal{T})$

**abbreviation** *utp-meet* (**infixl**  $\sqcap_1$  70) **where**

*utp-meet*  $\mathcal{T} \equiv \text{meet } (\text{uthy-order } \mathcal{T})$

**abbreviation** *utp-sup* ( $\bigsqcup_1$ - [90] 90) **where**

*utp-sup*  $\mathcal{T} \equiv \text{Lattice.sup } (\text{uthy-order } \mathcal{T})$

**abbreviation** *utp-inf* ( $\bigsqcap_1$ - [90] 90) **where**

*utp-inf*  $\mathcal{T} \equiv \text{Lattice.inf } (\text{uthy-order } \mathcal{T})$

**abbreviation** *utp-gfp* ( $\nu_1$ ) **where**

*utp-gfp*  $\mathcal{T} \equiv \text{GFP } (\text{uthy-order } \mathcal{T})$

**abbreviation** *utp-lfp* ( $\mu_1$ ) **where**

*utp-lfp*  $\mathcal{T} \equiv \text{LFP } (\text{uthy-order } \mathcal{T})$

**syntax**

*-tmu* :: *logic*  $\Rightarrow$  *pttrn*  $\Rightarrow$  *logic*  $\Rightarrow$  *logic* ( $\mu_1$  - - [0, 10] 10)

*-tnu* :: *logic*  $\Rightarrow$  *pttrn*  $\Rightarrow$  *logic*  $\Rightarrow$  *logic* ( $\nu_1$  - - [0, 10] 10)

**notation** *gfp* ( $\mu$ )

**notation** *lfp* ( $\nu$ )

**translations**

$\nu_T X \cdot P == \text{CONST } \text{utp-lfp } T (\lambda X. P)$

$\mu_T X \cdot P == \text{CONST } \text{utp-gfp } T (\lambda X. P)$

**lemma** *upred-lattice-inf*:

*Lattice.inf*  $\mathcal{P} A = \bigsqcap A$

**by** (*metis Sup-least Sup-upper UNIV-I antisym-conv subsetI upred-lattice.weak.inf-greatest upred-lattice.weak.inf-lower upred-lattice-carrier upred-lattice-le*)

We can then derive a number of properties about these operators, as below.

**context** *utp-theory-lattice*

**begin**

**lemma** *LFP-healthy-comp*:  $\mu F = \mu (F \circ \mathcal{H})$

**proof** –

**have**  $\{P. (P \text{ is } \mathcal{H}) \wedge F P \sqsubseteq P\} = \{P. (P \text{ is } \mathcal{H}) \wedge F (\mathcal{H} P) \sqsubseteq P\}$

**by** (*auto simp add: Healthy-def*)

**thus** *?thesis*

**by** (*simp add: LFP-def*)

**qed**

**lemma** *GFP-healthy-comp*:  $\nu F = \nu (F \circ \mathcal{H})$

**proof** –

**have**  $\{P. (P \text{ is } \mathcal{H}) \wedge P \sqsubseteq F P\} = \{P. (P \text{ is } \mathcal{H}) \wedge P \sqsubseteq F (\mathcal{H} P)\}$

**by** (*auto simp add: Healthy-def*)

**thus** *?thesis*

**by** (*simp add: GFP-def*)

**qed**

**lemma** *top-healthy [closure]*:  $\top \text{ is } \mathcal{H}$

**using** *weak.top-closed* **by** *auto*

**lemma** *bottom-healthy [closure]*:  $\perp \text{ is } \mathcal{H}$

**using** *weak.bottom-closed* **by** *auto*

**lemma** *utp-top*:  $P \text{ is } \mathcal{H} \implies P \sqsubseteq \top$

**using** *weak.top-higher* **by** *auto*

**lemma** *utp-bottom*:  $P \text{ is } \mathcal{H} \implies \perp \sqsubseteq P$

**using** *weak.bottom-lower* **by** *auto*

**end**

**lemma** *upred-top*:  $\top_{\mathcal{P}} = \text{false}$

**using** *ball-UNIV greatest-def* **by** *fastforce*

**lemma** *upred-bottom*:  $\perp_{\mathcal{P}} = \text{true}$

**by** *fastforce*

One way of obtaining a complete lattice is showing that the healthiness conditions are monotone, which the below locale characterises.

**locale** *utp-theory-mono* = *utp-theory* +

**assumes** *HCond-Mono [closure,intro]*: *Monotonic*  $\mathcal{H}$

**sublocale** *utp-theory-mono*  $\subseteq$  *utp-theory-lattice*

**proof** –

We can then use the Knaster-Tarski theorem to obtain a complete lattice, and thus provide all

the usual properties.

```

interpret weak-complete-lattice fpl  $\mathcal{P} \mathcal{H}$ 
  by (rule Knaster-Tarski, auto simp add: upred-lattice.weak.weak-complete-lattice-axioms)

have complete-lattice (fpl  $\mathcal{P} \mathcal{H}$ )
  by (unfold-locales, simp add: fps-def sup-exists, (blast intro: sup-exists inf-exists)+)

hence complete-lattice (uthy-order  $\mathcal{T}$ )
  by (simp add: utp-order-def, simp add: upred-lattice-def)

thus utp-theory-lattice  $\mathcal{T}$ 
  by (simp add: utp-theory-axioms utp-theory-lattice-def)
qed

```

```

context utp-theory-mono
begin

```

In a monotone theory, the top and bottom can always be obtained by applying the healthiness condition to the predicate top and bottom, respectively.

```

lemma healthy-top:  $\top = \mathcal{H}(\text{false})$ 
proof -
  have  $\top = \top_{fpl \mathcal{P} \mathcal{H}}$ 
    by (simp add: utp-order-fpl)
  also have  $\dots = \mathcal{H} \top_{\mathcal{P}}$ 
    using Knaster-Tarski-idem-extremes(1)[of  $\mathcal{P} \mathcal{H}$ ]
    by (simp add: HCond-Idempotent HCond-Mono)
  also have  $\dots = \mathcal{H} \text{false}$ 
    by (simp add: upred-top)
  finally show ?thesis .
qed

```

```

lemma healthy-bottom:  $\perp = \mathcal{H}(\text{true})$ 
proof -
  have  $\perp = \perp_{fpl \mathcal{P} \mathcal{H}}$ 
    by (simp add: utp-order-fpl)
  also have  $\dots = \mathcal{H} \perp_{\mathcal{P}}$ 
    using Knaster-Tarski-idem-extremes(2)[of  $\mathcal{P} \mathcal{H}$ ]
    by (simp add: HCond-Idempotent HCond-Mono)
  also have  $\dots = \mathcal{H} \text{true}$ 
    by (simp add: upred-bottom)
  finally show ?thesis .
qed

```

```

lemma healthy-inf:
  assumes  $A \subseteq \llbracket \mathcal{H} \rrbracket_H$ 
  shows  $\bigcap A = \mathcal{H} (\bigcap A)$ 
proof -
  have 1: weak-complete-lattice (uthy-order  $\mathcal{T}$ )
    by (simp add: weak.weak-complete-lattice-axioms)
  have 2: Monouthy-order  $\mathcal{T}$   $\mathcal{H}$ 
    by (simp add: HCond-Mono isotone-utp-orderI)
  have 3: Idemuthy-order  $\mathcal{T}$   $\mathcal{H}$ 
    by (simp add: HCond-Idem idempotent-def)
  show ?thesis
    using Knaster-Tarski-idem-inf-eq[OF upred-weak-complete-lattice, of  $\mathcal{H}$ ]

```

by (simp, metis HCond-Idempotent HCond-Mono assms partial-object.simps(3) upred-lattice-def  
upred-lattice-inf utp-order-def)  
qed

end

locale utp-theory-continuous = utp-theory +  
assumes HCond-Cont [closure,intro]: Continuous  $\mathcal{H}$

sublocale utp-theory-continuous  $\subseteq$  utp-theory-mono

proof

show Monotonic  $\mathcal{H}$

by (simp add: Continuous-Monotonic HCond-Cont)

qed

context utp-theory-continuous

begin

lemma healthy-inf-cont:

assumes  $A \subseteq \llbracket \mathcal{H} \rrbracket_H$   $A \neq \{\}$

shows  $\sqcap A = \sqcap A$

proof -

have  $\sqcap A = \sqcap (\mathcal{H}'A)$

using Continuous-def HCond-Cont assms(1) assms(2) healthy-inf by auto

also have  $\dots = \sqcap A$

by (unfold Healthy-carrier-image[OF assms(1)], simp)

finally show ?thesis .

qed

lemma healthy-inf-def:

assumes  $A \subseteq \llbracket \mathcal{H} \rrbracket_H$

shows  $\sqcap A = (\text{if } (A = \{\}) \text{ then } \top \text{ else } (\sqcap A))$

using assms healthy-inf-cont weak.weak-inf-empty by auto

lemma healthy-meet-cont:

assumes  $P$  is  $\mathcal{H}$   $Q$  is  $\mathcal{H}$

shows  $P \sqcap Q = P \sqcap Q$

using healthy-inf-cont[of  $\{P, Q\}$ ] assms

by (simp add: Healthy-if meet-def)

lemma meet-is-healthy [closure]:

assumes  $P$  is  $\mathcal{H}$   $Q$  is  $\mathcal{H}$

shows  $P \sqcap Q$  is  $\mathcal{H}$

by (metis Continuous-Disjunctous Disjunctuous-def HCond-Cont Healthy-def' assms(1) assms(2))

lemma meet-bottom [simp]:

assumes  $P$  is  $\mathcal{H}$

shows  $P \sqcap \perp = \perp$

by (simp add: assms semilattice-sup-class.sup-absorb2 utp-bottom)

lemma meet-top [simp]:

assumes  $P$  is  $\mathcal{H}$

shows  $P \sqcap \top = P$

by (simp add: assms semilattice-sup-class.sup-absorb1 utp-top)

The UTP theory lfp operator can be rewritten to the alphabetised predicate lfp when in a

continuous context.

```

theorem utp-lfp-def:
  assumes Monotonic  $F \in \llbracket \mathcal{H} \rrbracket_H \rightarrow \llbracket \mathcal{H} \rrbracket_H$ 
  shows  $\mu F = (\mu X \cdot F(\mathcal{H}(X)))$ 
proof (rule antisym)
  have ne:  $\{P. (P \text{ is } \mathcal{H}) \wedge F P \sqsubseteq P\} \neq \{\}$ 
  proof –
    have  $F \top \sqsubseteq \top$ 
    using assms(2) utp-top weak.top-closed by force
    thus ?thesis
    by (auto, rule-tac x= $\top$  in exI, auto simp add: top-healthy)
  qed
show  $\mu F \sqsubseteq (\mu X \cdot F(\mathcal{H} X))$ 
proof –
  have  $\sqcap \{P. (P \text{ is } \mathcal{H}) \wedge F(P) \sqsubseteq P\} \sqsubseteq \sqcap \{P. F(\mathcal{H}(P)) \sqsubseteq P\}$ 
  proof –
    have  $1: \bigwedge P. F(\mathcal{H}(P)) = \mathcal{H}(F(\mathcal{H}(P)))$ 
    by (metis HCond-Idem Healthy-def assms(2) funcset-mem mem-Collect-eq)
    show ?thesis
    proof (rule Sup-least, auto)
      fix  $P$ 
      assume  $a: F(\mathcal{H} P) \sqsubseteq P$ 
      hence  $F: (F(\mathcal{H} P)) \sqsubseteq (\mathcal{H} P)$ 
      by (metis 1 HCond-Mono mono-def)
      show  $\sqcap \{P. (P \text{ is } \mathcal{H}) \wedge F P \sqsubseteq P\} \sqsubseteq P$ 
      proof (rule Sup-upper2[of F ( $\mathcal{H} P$ )])
        show  $F(\mathcal{H} P) \in \{P. (P \text{ is } \mathcal{H}) \wedge F P \sqsubseteq P\}$ 
        proof (auto)
          show  $F(\mathcal{H} P) \text{ is } \mathcal{H}$ 
          by (metis 1 Healthy-def)
          show  $F(F(\mathcal{H} P)) \sqsubseteq F(\mathcal{H} P)$ 
          using F mono-def assms(1) by blast
        qed
        show  $F(\mathcal{H} P) \sqsubseteq P$ 
        by (simp add: a)
      qed
    qed
  qed

with ne show ?thesis
  by (simp add: LFP-def gfp-def, subst healthy-inf-cont, auto simp add: lfp-def)
qed
from ne show  $(\mu X \cdot F(\mathcal{H} X)) \sqsubseteq \mu F$ 
  apply (simp add: LFP-def gfp-def, subst healthy-inf-cont, auto simp add: lfp-def)
  apply (rule Sup-least)
  apply (auto simp add: Healthy-def Sup-upper)
done
qed

```

**end**

In another direction, we can also characterise UTP theories that are relational. Minimally this requires that the healthiness condition is closed under sequential composition.

```

locale utp-theory-rel =
  utp-theory +

```

**assumes** *Healthy-Sequence* [closure]:  $\llbracket P \text{ is } \mathcal{H}; Q \text{ is } \mathcal{H} \rrbracket \implies (P ;; Q) \text{ is } \mathcal{H}$

**locale** *utp-theory-cont-rel* = *utp-theory-continuous* + *utp-theory-rel*  
**begin**

**lemma** *seq-cont-Sup-distl*:

**assumes**  $P \text{ is } \mathcal{H} \ A \subseteq \llbracket \mathcal{H} \rrbracket_H \ A \neq \{\}$

**shows**  $P ;; (\bigcap A) = \bigcap \{P ;; Q \mid Q. Q \in A\}$

**proof** –

**have**  $\{P ;; Q \mid Q. Q \in A\} \subseteq \llbracket \mathcal{H} \rrbracket_H$

**using** *Healthy-Sequence* *assms*(1) *assms*(2) **by** (*auto*)

**thus** *?thesis*

**by** (*simp add: healthy-inf-cont seq-Sup-distl setcompr-eq-image assms*)

**qed**

**lemma** *seq-cont-Sup-distr*:

**assumes**  $Q \text{ is } \mathcal{H} \ A \subseteq \llbracket \mathcal{H} \rrbracket_H \ A \neq \{\}$

**shows**  $(\bigcap A) ;; Q = \bigcap \{P ;; Q \mid P. P \in A\}$

**proof** –

**have**  $\{P ;; Q \mid P. P \in A\} \subseteq \llbracket \mathcal{H} \rrbracket_H$

**using** *Healthy-Sequence* *assms*(1) *assms*(2) **by** (*auto*)

**thus** *?thesis*

**by** (*simp add: healthy-inf-cont seq-Sup-distr setcompr-eq-image assms*)

**qed**

**end**

There also exist UTP theories with units, and the following operator is a theory specific operator for them.

**consts**

*utp-unit* ::  $(\mathcal{T}, \alpha) \text{ uty} \Rightarrow \alpha \text{ hrel } (\mathcal{I}\mathcal{I}_1)$

Not all theories have both a left and a right unit (e.g. H1-H2 designs) and so we split up the locale into two cases.

**locale** *utp-theory-left-unital* =

*utp-theory-rel* +

**assumes** *Healthy-Left-Unit* [closure]:  $\mathcal{I}\mathcal{I} \text{ is } \mathcal{H}$

**and** *Left-Unit*:  $P \text{ is } \mathcal{H} \implies (\mathcal{I}\mathcal{I} ;; P) = P$

**locale** *utp-theory-right-unital* =

*utp-theory-rel* +

**assumes** *Healthy-Right-Unit* [closure]:  $\mathcal{I}\mathcal{I} \text{ is } \mathcal{H}$

**and** *Right-Unit*:  $P \text{ is } \mathcal{H} \implies (P ;; \mathcal{I}\mathcal{I}) = P$

**locale** *utp-theory-unital* =

*utp-theory-rel* +

**assumes** *Healthy-Unit* [closure]:  $\mathcal{I}\mathcal{I} \text{ is } \mathcal{H}$

**and** *Unit-Left*:  $P \text{ is } \mathcal{H} \implies (\mathcal{I}\mathcal{I} ;; P) = P$

**and** *Unit-Right*:  $P \text{ is } \mathcal{H} \implies (P ;; \mathcal{I}\mathcal{I}) = P$

**locale** *utp-theory-mono-unital* = *utp-theory-mono* + *utp-theory-unital*

**definition** *utp-star* ( $\star_1$  [999] 999) **where**

*utp-star*  $\mathcal{T} \ P = (\nu_{\mathcal{T}} (\lambda X. (P ;; X) \sqcap_{\mathcal{T}} \mathcal{I}\mathcal{I}_{\mathcal{T}}))$

**definition** *utp-omega* ( $\omega_1$  [999] 999) **where**

*utp-omega*  $\mathcal{T} P = (\mu_{\mathcal{T}} (\lambda X. (P ;; X)))$

**locale** *utp-pre-left-quantale* = *utp-theory-continuous* + *utp-theory-left-unital*  
**begin**

**lemma** *star-healthy* [closure]:  $P\star$  is  $\mathcal{H}$

**by** (*metis mem-Collect-eq utp-order-carrier utp-star-def weak.GFP-closed*)

**lemma** *star-unfold*:  $P$  is  $\mathcal{H} \implies P\star = (P;;P\star) \sqcap \mathcal{II}$

**apply** (*simp add: utp-star-def healthy-meet-cont*)

**apply** (*subst GFP-unfold*)

**apply** (*rule Mono-utp-orderI*)

**apply** (*simp add: healthy-meet-cont closure semilattice-sup-class.le-supI1 segr-mono*)

**apply** (*auto intro: funcsetI*)

**apply** (*simp add: Healthy-Left-Unit Healthy-Sequence healthy-meet-cont meet-is-healthy*)

**using** *Healthy-Left-Unit Healthy-Sequence healthy-meet-cont weak.GFP-closed* **apply** *auto*

**done**

**end**

**sublocale** *utp-theory-unital*  $\subseteq$  *utp-theory-left-unital*

**by** (*simp add: Healthy-Unit Unit-Left Healthy-Sequence utp-theory-rel-def utp-theory-axioms utp-theory-rel-axioms-def utp-theory-left-unital-axioms-def utp-theory-left-unital-def*)

**sublocale** *utp-theory-unital*  $\subseteq$  *utp-theory-right-unital*

**by** (*simp add: Healthy-Unit Unit-Right Healthy-Sequence utp-theory-rel-def utp-theory-axioms utp-theory-rel-axioms-def utp-theory-right-unital-axioms-def utp-theory-right-unital-def*)

## 18.6 Theory of relations

We can exemplify the creation of a UTP theory with the theory of relations, a trivial theory.

**typeddecl** *REL*

**abbreviation** *REL*  $\equiv$  *UTHY*(*REL*, ' $\alpha$ )

We declare the type *REL* to be the tag for this theory. We need know nothing about this type (other than it's non-empty), since it is merely a name. We also create the corresponding constant to refer to the theory. Then we can use it to instantiate the relevant polymorphic constants.

**overloading**

*rel-hcond* == *utp-hcond* :: (*REL*, ' $\alpha$ ) *uthy*  $\Rightarrow$  (' $\alpha$   $\times$  ' $\alpha$ ) *health*

*rel-unit* == *utp-unit* :: (*REL*, ' $\alpha$ ) *uthy*  $\Rightarrow$  ' $\alpha$  *hrel*

**begin**

The healthiness condition of a relation is simply identity, since every alphabetised relation is healthy.

**definition** *rel-hcond* :: (*REL*, ' $\alpha$ ) *uthy*  $\Rightarrow$  (' $\alpha$   $\times$  ' $\alpha$ ) *upred*  $\Rightarrow$  (' $\alpha$   $\times$  ' $\alpha$ ) *upred* **where**

*rel-hcond* *T* = *id*

The unit of the theory is simply the relational unit.

**definition** *rel-unit* :: (*REL*, ' $\alpha$ ) *uthy*  $\Rightarrow$  ' $\alpha$  *hrel* **where**

*rel-unit* *T* = *II*

**end**



Finally we can show that relations are a monotone and unital theory using a locale interpretation, which requires that we prove all the relevant properties. It's convenient to rewrite some of the theorems so that the provisos are more UTP like; e.g. that the carrier is the set of healthy predicates.

**interpretation** *rel-theory: utp-theory-mono-unital REL*  
**rewrites** *carrier (uthy-order REL) =  $\llbracket id \rrbracket_H$*   
**by** (*unfold-locales, simp-all add: rel-hcond-def rel-unit-def Healthy-def*)

We can then, for instance, determine what the top and bottom of our new theory is.

**lemma** *REL-top:  $\top_{REL} = false$*   
**by** (*simp add: rel-theory.healthy-top, simp add: rel-hcond-def*)

**lemma** *REL-bottom:  $\perp_{REL} = true$*   
**by** (*simp add: rel-theory.healthy-bottom, simp add: rel-hcond-def*)

A number of theorems have been exported, such as the fixed point unfolding laws.

**thm** *rel-theory.GFP-unfold*

## 18.7 Theory links

We can also describe links between theories, such as Galois connections and retractions, using the following notation.

**definition** *mk-conn* ( $- \Leftarrow \langle -, - \rangle \Rightarrow - [90, 0, 0, 91] \ 91$ ) **where**  
 $H1 \Leftarrow \langle \mathcal{H}_1, \mathcal{H}_2 \rangle \Rightarrow H2 \equiv () \text{ order}A = \text{utp-order } H1, \text{ order}B = \text{utp-order } H2, \text{ lower} = \mathcal{H}_2, \text{ upper} = \mathcal{H}_1 \ ()$

**abbreviation** *mk-conn'* ( $- \Leftarrow \langle -, - \rangle \rightarrow - [90, 0, 0, 91] \ 91$ ) **where**  
 $T1 \Leftarrow \langle \mathcal{H}_1, \mathcal{H}_2 \rangle \rightarrow T2 \equiv \mathcal{H}_{T1} \Leftarrow \langle \mathcal{H}_1, \mathcal{H}_2 \rangle \Rightarrow \mathcal{H}_{T2}$

**lemma** *mk-conn-orderA [simp]:  $\mathcal{X}_{H1} \Leftarrow \langle \mathcal{H}_1, \mathcal{H}_2 \rangle \Rightarrow H2 = \text{utp-order } H1$*   
**by** (*simp add: mk-conn-def*)

**lemma** *mk-conn-orderB [simp]:  $\mathcal{Y}_{H1} \Leftarrow \langle \mathcal{H}_1, \mathcal{H}_2 \rangle \Rightarrow H2 = \text{utp-order } H2$*   
**by** (*simp add: mk-conn-def*)

**lemma** *mk-conn-lower [simp]:  $\pi_* H1 \Leftarrow \langle \mathcal{H}_1, \mathcal{H}_2 \rangle \Rightarrow H2 = \mathcal{H}_1$*   
**by** (*simp add: mk-conn-def*)

**lemma** *mk-conn-upper [simp]:  $\pi^* H1 \Leftarrow \langle \mathcal{H}_1, \mathcal{H}_2 \rangle \Rightarrow H2 = \mathcal{H}_2$*   
**by** (*simp add: mk-conn-def*)

**lemma** *galois-comp:  $(H2 \Leftarrow \langle \mathcal{H}_3, \mathcal{H}_4 \rangle \Rightarrow H3) \circ_g (H1 \Leftarrow \langle \mathcal{H}_1, \mathcal{H}_2 \rangle \Rightarrow H2) = H1 \Leftarrow \langle \mathcal{H}_1 \circ \mathcal{H}_3, \mathcal{H}_4 \circ \mathcal{H}_2 \rangle \Rightarrow H3$*   
**by** (*simp add: comp-galcon-def mk-conn-def*)

Example Galois connection / retract: Existential quantification

**lemma** *Idempotent-ex: mwb-lens  $x \Longrightarrow \text{Idempotent } (ex \ x)$*   
**by** (*simp add: Idempotent-def exists-twice*)

**lemma** *Monotonic-ex: mwb-lens  $x \Longrightarrow \text{Monotonic } (ex \ x)$*   
**by** (*simp add: mono-def ex-mono*)

**lemma** *ex-closed-unrest:*  
*mwb-lens  $x \Longrightarrow \llbracket ex \ x \rrbracket_H = \{P. x \# P\}$*   
**by** (*simp add: Healthy-def unrest-as-exists*)

Any theory can be composed with an existential quantification to produce a Galois connection

**theorem** *ex-retract*:

**assumes** *vwb-lens*  $x$  *Idempotent*  $H$   $ex\ x \circ H = H \circ ex\ x$

**shows** *retract*  $((ex\ x \circ H) \Leftarrow \langle ex\ x, H \rangle \Rightarrow H)$

**proof** (*unfold-locales*, *simp-all*)

**show**  $H \in \llbracket ex\ x \circ H \rrbracket_H \rightarrow \llbracket H \rrbracket_H$

**using** *Healthy-Idempotent* *assms* **by** *blast*

**from** *assms*(1) *assms*(3)[*THEN sym*] **show**  $ex\ x \in \llbracket H \rrbracket_H \rightarrow \llbracket ex\ x \circ H \rrbracket_H$

**by** (*simp add: Pi-iff Healthy-def fun-eq-iff exists-twice*)

**fix**  $P\ Q$

**assume**  $P$  is  $(ex\ x \circ H)$   $Q$  is  $H$

**thus**  $(H\ P \sqsubseteq Q) = (P \sqsubseteq (\exists\ x \cdot Q))$

**by** (*metis* (*no-types*, *lifting*) *Healthy-Idempotent Healthy-if* *assms comp-apply dual-order.trans ex-weakens* *utp-pred-laws.ex-mono vwb-lens-wb*)

**next**

**fix**  $P$

**assume**  $P$  is  $(ex\ x \circ H)$

**thus**  $(\exists\ x \cdot H\ P) \sqsubseteq P$

**by** (*simp add: Healthy-def*)

**qed**

**corollary** *ex-retract-id*:

**assumes** *vwb-lens*  $x$

**shows** *retract*  $(ex\ x \Leftarrow \langle ex\ x, id \rangle \Rightarrow id)$

**using** *assms ex-retract*[**where**  $H=id$ ] **by** (*auto*)

**end**

## 19 Concurrent Programming

**theory** *utp-concurrency*

**imports**

*utp-hoare*

*utp-rel*

*utp-tactics*

*utp-theory*

**begin**

In this theory we describe the UTP scheme for concurrency, *parallel-by-merge*, which provides a general parallel operator parametrised by a “merge predicate” that explains how to merge the after states of the composed predicates. It can thus be applied to many languages and concurrency schemes, with this theory providing a number of generic laws. The operator is explained in more detail in Chapter 7 of the UTP book [5].

### 19.1 Variable Renamings

In parallel-by-merge constructions, a merge predicate defines the behaviour following execution of of parallel processes,  $P \parallel Q$ , as a relation that merges the output of  $P$  and  $Q$ . In order to achieve this we need to separate the variable values output from  $P$  and  $Q$ , and in addition the variable values before execution. The following three constructs do these separations. The initial state-space before execution is  $'\alpha$ , the final state-space after the first parallel process is  $'\beta_0$ , and the final state-space for the second is  $'\beta_1$ . These three functions lift variables on these three state-spaces, respectively.

**alphabet**  $(' \alpha, ' \beta_0, ' \beta_1)$  *mrg* =

$mrg\text{-}prior :: 'α$   
 $mrg\text{-}left :: 'β_0$   
 $mrg\text{-}right :: 'β_1$

**definition**  $pre\text{-}uvar :: ('a \Rightarrow 'α) \Rightarrow ('a \Rightarrow ('α, 'β_0, 'β_1) \text{ mrg})$  **where**  
 $[upred\text{-}defs]: pre\text{-}uvar\ x = x ;_L mrg\text{-}prior$

**definition**  $left\text{-}uvar :: ('a \Rightarrow 'β_0) \Rightarrow ('a \Rightarrow ('α, 'β_0, 'β_1) \text{ mrg})$  **where**  
 $[upred\text{-}defs]: left\text{-}uvar\ x = x ;_L mrg\text{-}left$

**definition**  $right\text{-}uvar :: ('a \Rightarrow 'β_1) \Rightarrow ('a \Rightarrow ('α, 'β_0, 'β_1) \text{ mrg})$  **where**  
 $[upred\text{-}defs]: right\text{-}uvar\ x = x ;_L mrg\text{-}right$

We set up syntax for the three variable classes using a subscript  $<$ , 0- $x$ , and 1- $x$ , respectively.

**syntax**

$\text{-}svarpre :: svid \Rightarrow svid\ (-< [999] 999)$   
 $\text{-}svarleft :: svid \Rightarrow svid\ (0-- [999] 999)$   
 $\text{-}svarright :: svid \Rightarrow svid\ (1-- [999] 999)$

**translations**

$\text{-}svarpre\ x == CONST\ pre\text{-}uvar\ x$   
 $\text{-}svarleft\ x == CONST\ left\text{-}uvar\ x$   
 $\text{-}svarright\ x == CONST\ right\text{-}uvar\ x$   
 $\text{-}svarpre\ \Sigma <= CONST\ pre\text{-}uvar\ 1_L$   
 $\text{-}svarleft\ \Sigma <= CONST\ left\text{-}uvar\ 1_L$   
 $\text{-}svarright\ \Sigma <= CONST\ right\text{-}uvar\ 1_L$

We proved behavedness closure properties about the lenses.

**lemma**  $left\text{-}uvar\ [simp]: vwb\text{-}lens\ x \Longrightarrow vwb\text{-}lens\ (left\text{-}uvar\ x)$   
**by** ( $simp\ add: left\text{-}uvar\text{-}def$ )

**lemma**  $right\text{-}uvar\ [simp]: vwb\text{-}lens\ x \Longrightarrow vwb\text{-}lens\ (right\text{-}uvar\ x)$   
**by** ( $simp\ add: right\text{-}uvar\text{-}def$ )

**lemma**  $pre\text{-}uvar\ [simp]: vwb\text{-}lens\ x \Longrightarrow vwb\text{-}lens\ (pre\text{-}uvar\ x)$   
**by** ( $simp\ add: pre\text{-}uvar\text{-}def$ )

**lemma**  $left\text{-}uvar\text{-}mwb\ [simp]: mwb\text{-}lens\ x \Longrightarrow mwb\text{-}lens\ (left\text{-}uvar\ x)$   
**by** ( $simp\ add: left\text{-}uvar\text{-}def$ )

**lemma**  $right\text{-}uvar\text{-}mwb\ [simp]: mwb\text{-}lens\ x \Longrightarrow mwb\text{-}lens\ (right\text{-}uvar\ x)$   
**by** ( $simp\ add: right\text{-}uvar\text{-}def$ )

**lemma**  $pre\text{-}uvar\text{-}mwb\ [simp]: mwb\text{-}lens\ x \Longrightarrow mwb\text{-}lens\ (pre\text{-}uvar\ x)$   
**by** ( $simp\ add: pre\text{-}uvar\text{-}def$ )

We prove various independence laws about the variable classes.

**lemma**  $left\text{-}uvar\text{-}indep\text{-}right\text{-}uvar\ [simp]:$   
 $left\text{-}uvar\ x \bowtie right\text{-}uvar\ y$   
**by** ( $simp\ add: left\text{-}uvar\text{-}def\ right\text{-}uvar\text{-}def\ lens\text{-}comp\text{-}assoc[THEN\ sym]$ )

**lemma**  $left\text{-}uvar\text{-}indep\text{-}pre\text{-}uvar\ [simp]:$   
 $left\text{-}uvar\ x \bowtie pre\text{-}uvar\ y$   
**by** ( $simp\ add: left\text{-}uvar\text{-}def\ pre\text{-}uvar\text{-}def$ )

**lemma** *left-uvar-indep-left-uvar* [simp]:  
 $x \bowtie y \implies \text{left-uvar } x \bowtie \text{left-uvar } y$   
**by** (simp add: left-uvar-def)

**lemma** *right-uvar-indep-left-uvar* [simp]:  
 $\text{right-uvar } x \bowtie \text{left-uvar } y$   
**by** (simp add: lens-indep-sym)

**lemma** *right-uvar-indep-pre-uvar* [simp]:  
 $\text{right-uvar } x \bowtie \text{pre-uvar } y$   
**by** (simp add: right-uvar-def pre-uvar-def)

**lemma** *right-uvar-indep-right-uvar* [simp]:  
 $x \bowtie y \implies \text{right-uvar } x \bowtie \text{right-uvar } y$   
**by** (simp add: right-uvar-def)

**lemma** *pre-uvar-indep-left-uvar* [simp]:  
 $\text{pre-uvar } x \bowtie \text{left-uvar } y$   
**by** (simp add: lens-indep-sym)

**lemma** *pre-uvar-indep-right-uvar* [simp]:  
 $\text{pre-uvar } x \bowtie \text{right-uvar } y$   
**by** (simp add: lens-indep-sym)

**lemma** *pre-uvar-indep-pre-uvar* [simp]:  
 $x \bowtie y \implies \text{pre-uvar } x \bowtie \text{pre-uvar } y$   
**by** (simp add: pre-uvar-def)

## 19.2 Merge Predicates

A merge predicate is a relation whose input has three parts: the prior variables, the output variables of the left predicate, and the output of the right predicate.

**type-synonym**  $'\alpha \text{ merge} = ((' \alpha, ' \alpha, ' \alpha) \text{ mrg}, ' \alpha) \text{ rel}$

skip is the merge predicate which ignores the output of both parallel predicates

**definition**  $\text{skip}_m :: ' \alpha \text{ merge}$  **where**  
[upred-defs]:  $\text{skip}_m = (\$ \mathbf{v}' =_u \$ \mathbf{v}_<)$

swap is a predicate that the swaps the left and right indices; it is used to specify commutativity of the parallel operator

**definition**  $\text{swap}_m :: ((' \alpha, ' \beta, ' \beta) \text{ mrg}) \text{ hrel}$  **where**  
[upred-defs]:  $\text{swap}_m = (0 - \mathbf{v}, 1 - \mathbf{v}) := (\& 1 - \mathbf{v}, \& 0 - \mathbf{v})$

A symmetric merge is one for which swapping the order of the merged concurrent predicates has no effect. We represent this by the following healthiness condition that states that  $\text{swap}_m$  is a left-unit.

**abbreviation**  $\text{SymMerge} :: ' \alpha \text{ merge} \Rightarrow ' \alpha \text{ merge}$  **where**  
 $\text{SymMerge}(M) \equiv (\text{swap}_m ;; M)$

## 19.3 Separating Simulations

U0 and U1 are relations modify the variables of the input state-space such that they become indexed with 0 and 1, respectively.

**definition**  $U0 :: ('β_0, ('α, 'β_0, 'β_1) \text{ mrg}) \text{ rel}$  **where**  
 $[upred-defs]: U0 = (\$0 - \mathbf{v}' =_u \$\mathbf{v})$

**definition**  $U1 :: ('β_1, ('α, 'β_0, 'β_1) \text{ mrg}) \text{ rel}$  **where**  
 $[upred-defs]: U1 = (\$1 - \mathbf{v}' =_u \$\mathbf{v})$

**lemma**  $U0\text{-swap}: (U0 ;; \text{swap}_m) = U1$   
**by**  $(\text{rel-auto})$

**lemma**  $U1\text{-swap}: (U1 ;; \text{swap}_m) = U0$   
**by**  $(\text{rel-auto})$

As shown below, separating simulations can also be expressed using the following two alphabet extrusions

**definition**  $U0\alpha$  **where**  $[upred-defs]: U0\alpha = (1_L \times_L \text{mrg-left})$

**definition**  $U1\alpha$  **where**  $[upred-defs]: U1\alpha = (1_L \times_L \text{mrg-right})$

We then create the following intuitive syntax for separating simulations.

**abbreviation**  $U0\text{-alpha-lift } (\lceil - \rceil_0)$  **where**  $\lceil P \rceil_0 \equiv P \oplus_p U0\alpha$

**abbreviation**  $U1\text{-alpha-lift } (\lceil - \rceil_1)$  **where**  $\lceil P \rceil_1 \equiv P \oplus_p U1\alpha$

$\lceil P \rceil_0$  is predicate  $P$  where all variables are indexed by 0, and  $\lceil P \rceil_1$  is where all variables are indexed by 1. We can thus equivalently express separating simulations using alphabet extrusion.

**lemma**  $U0\text{-as-alpha}: (P ;; U0) = \lceil P \rceil_0$   
**by**  $(\text{rel-auto})$

**lemma**  $U1\text{-as-alpha}: (P ;; U1) = \lceil P \rceil_1$   
**by**  $(\text{rel-auto})$

**lemma**  $U0\alpha\text{-vwb-lens } [simp]: \text{vwb-lens } U0\alpha$   
**by**  $(\text{simp add: } U0\alpha\text{-def id-vwb-lens prod-vwb-lens})$

**lemma**  $U1\alpha\text{-vwb-lens } [simp]: \text{vwb-lens } U1\alpha$   
**by**  $(\text{simp add: } U1\alpha\text{-def id-vwb-lens prod-vwb-lens})$

**lemma**  $U0\alpha\text{-indep-right-uvar } [simp]: \text{vwb-lens } x \implies U0\alpha \bowtie \text{out-var } (\text{right-uvar } x)$   
**by**  $(\text{force intro: plus-pres-lens-indep fst-snd-lens-indep lens-indep-left-comp}$   
 $\text{simp add: } U0\alpha\text{-def right-uvar-def out-var-def prod-as-plus lens-comp-assoc}[\text{THEN sym}])$

**lemma**  $U1\alpha\text{-indep-left-uvar } [simp]: \text{vwb-lens } x \implies U1\alpha \bowtie \text{out-var } (\text{left-uvar } x)$   
**by**  $(\text{force intro: plus-pres-lens-indep fst-snd-lens-indep lens-indep-left-comp}$   
 $\text{simp add: } U1\alpha\text{-def left-uvar-def out-var-def prod-as-plus lens-comp-assoc}[\text{THEN sym}])$

**lemma**  $U0\text{-alpha-lift-bool-subst } [usubst]:$   
 $\sigma(\$0 - x' \mapsto_s \text{true}) \dagger \lceil P \rceil_0 = \sigma \dagger \lceil P \llbracket \text{true} / \$x' \rrbracket \rceil_0$   
 $\sigma(\$0 - x' \mapsto_s \text{false}) \dagger \lceil P \rceil_0 = \sigma \dagger \lceil P \llbracket \text{false} / \$x' \rrbracket \rceil_0$   
**by**  $(\text{pred-auto+})$

**lemma**  $U1\text{-alpha-lift-bool-subst } [usubst]:$   
 $\sigma(\$1 - x' \mapsto_s \text{true}) \dagger \lceil P \rceil_1 = \sigma \dagger \lceil P \llbracket \text{true} / \$x' \rrbracket \rceil_1$   
 $\sigma(\$1 - x' \mapsto_s \text{false}) \dagger \lceil P \rceil_1 = \sigma \dagger \lceil P \llbracket \text{false} / \$x' \rrbracket \rceil_1$   
**by**  $(\text{pred-auto+})$

**lemma** *U0-alpha-out-var* [*alpha*]:  $\lceil \$x' \rceil_0 = \$0 - x'$   
**by** (*rel-auto*)

**lemma** *U1-alpha-out-var* [*alpha*]:  $\lceil \$x' \rceil_1 = \$1 - x'$   
**by** (*rel-auto*)

**lemma** *U0-skip* [*alpha*]:  $\lceil II \rceil_0 = (\$0 - \mathbf{v}' =_u \$\mathbf{v})$   
**by** (*rel-auto*)

**lemma** *U1-skip* [*alpha*]:  $\lceil II \rceil_1 = (\$1 - \mathbf{v}' =_u \$\mathbf{v})$   
**by** (*rel-auto*)

**lemma** *U0-seqr* [*alpha*]:  $\lceil P ;; Q \rceil_0 = P ;; \lceil Q \rceil_0$   
**by** (*rel-auto*)

**lemma** *U1-seqr* [*alpha*]:  $\lceil P ;; Q \rceil_1 = P ;; \lceil Q \rceil_1$   
**by** (*rel-auto*)

**lemma** *U0 $\alpha$ -comp-in-var* [*alpha*]:  $(in-var\ x) ;_L U0\alpha = in-var\ x$   
**by** (*simp add: U0 $\alpha$ -def alpha-in-var in-var-prod-lens pre-uvar-def*)

**lemma** *U0 $\alpha$ -comp-out-var* [*alpha*]:  $(out-var\ x) ;_L U0\alpha = out-var\ (left-uvar\ x)$   
**by** (*simp add: U0 $\alpha$ -def alpha-out-var id-wb-lens left-uvar-def out-var-prod-lens*)

**lemma** *U1 $\alpha$ -comp-in-var* [*alpha*]:  $(in-var\ x) ;_L U1\alpha = in-var\ x$   
**by** (*simp add: U1 $\alpha$ -def alpha-in-var in-var-prod-lens pre-uvar-def*)

**lemma** *U1 $\alpha$ -comp-out-var* [*alpha*]:  $(out-var\ x) ;_L U1\alpha = out-var\ (right-uvar\ x)$   
**by** (*simp add: U1 $\alpha$ -def alpha-out-var id-wb-lens right-uvar-def out-var-prod-lens*)

## 19.4 Associative Merges

Associativity of a merge means that if we construct a three way merge from a two way merge and then rotate the three inputs of the merge to the left, then we get exactly the same three way merge back.

We first construct the operator that constructs the three way merge by effectively wiring up the two way merge in an appropriate way.

**definition** *ThreeWayMerge* ::  $'\alpha \text{ merge} \Rightarrow (('\alpha, '\alpha, (' \alpha, '\alpha, '\alpha) \text{ mrg}) \text{ mrg}, '\alpha) \text{ rel } (\mathbf{M3}'(-'))$  **where**  
 $[upred-defs]: \text{ThreeWayMerge } M = ((\$0 - \mathbf{v}' =_u \$0 - \mathbf{v} \wedge \$1 - \mathbf{v}' =_u \$1 - \mathbf{v} \wedge \$\mathbf{v}_{<} =_u \$\mathbf{v}_{<}) ;; M ;; U0 \wedge \$1 - \mathbf{v}' =_u \$1 - \mathbf{v} \wedge \$\mathbf{v}_{<} =_u \$\mathbf{v}_{<}) ;; M$

The next definition rotates the inputs to a three way merge to the left one place.

**abbreviation** *rotate<sub>m</sub>* **where**  $rotate_m \equiv (0 - \mathbf{v}, 1 - 0 - \mathbf{v}, 1 - 1 - \mathbf{v}) := (\&1 - 0 - \mathbf{v}, \&1 - 1 - \mathbf{v}, \&0 - \mathbf{v})$

Finally, a merge is associative if rotating the inputs does not effect the output.

**definition** *AssocMerge* ::  $'\alpha \text{ merge} \Rightarrow \text{bool}$  **where**  
 $[upred-defs]: \text{AssocMerge } M = (rotate_m ;; \mathbf{M3}(M) = \mathbf{M3}(M))$

## 19.5 Parallel Operators

We implement the following useful abbreviation for separating of two parallel processes and copying of the before variables, all to act as input to the merge predicate.

**abbreviation** *par-sep* (**infixr**  $\parallel_s$  85) **where**

$$P \parallel_s Q \equiv (P ;; U0) \wedge (Q ;; U1) \wedge \$\mathbf{v}_{<}' =_u \$\mathbf{v}$$

The following implementation of parallel by merge is less general than the book version, in that it does not properly partition the alphabet into two disjoint segments. We could actually achieve this specifying lenses into the larger alphabet, but this would complicate the definition of programs. May reconsider later.

**definition** *par-by-merge*  $(- \parallel - [85, 0, 86] \ 85)$   
**where**  $[upred-defs]: P \parallel_M Q = (P \parallel_s Q ;; M)$

**lemma** *par-by-merge-alt-def*:  $P \parallel_M Q = ([P]_0 \wedge [Q]_1 \wedge \$\mathbf{v}_{<}' =_u \$\mathbf{v}) ;; M$   
**by** (*simp add: par-by-merge-def U0-as-alpha U1-as-alpha*)

**lemma** *shEx-pbm-left*:  $((\exists x \cdot P x) \parallel_M Q) = (\exists x \cdot (P x \parallel_M Q))$   
**by** (*rel-auto*)

**lemma** *shEx-pbm-right*:  $(P \parallel_M (\exists x \cdot Q x)) = (\exists x \cdot (P \parallel_M Q x))$   
**by** (*rel-auto*)

## 19.6 Unrestriction Laws

**lemma** *unrest-in-par-by-merge*  $[unrest]:$   
 $\llbracket \$x \# P; \$x_{<} \# M; \$x \# Q \rrbracket \implies \$x \# P \parallel_M Q$   
**by** (*rel-auto, fastforce+*)

**lemma** *unrest-out-par-by-merge*  $[unrest]:$   
 $\llbracket \$x' \# M \rrbracket \implies \$x' \# P \parallel_M Q$   
**by** (*rel-auto*)

## 19.7 Substitution laws

Substitution is a little tricky because when we push the expression through the composition operator the alphabet of the expression must also change. Consequently for now we only support literal substitution, though this could be generalised with suitable alphabet coercsions. We need quite a number of variants to support this which are below.

**lemma** *U0-seq-subst*:  $(P ;; U0) \llbracket \llbracket v \rrbracket / \$0 - x' \rrbracket = (P \llbracket \llbracket v \rrbracket / \$x' \rrbracket ;; U0)$   
**by** (*rel-auto*)

**lemma** *U1-seq-subst*:  $(P ;; U1) \llbracket \llbracket v \rrbracket / \$1 - x' \rrbracket = (P \llbracket \llbracket v \rrbracket / \$x' \rrbracket ;; U1)$   
**by** (*rel-auto*)

**lemma** *lit-pbm-subst*  $[usubst]:$   
**fixes**  $x :: (- \implies 'a)$   
**shows**  
 $\bigwedge P Q M \sigma. \sigma(\$x \mapsto_s \llbracket v \rrbracket) \dagger (P \parallel_M Q) = \sigma \dagger ((P \llbracket \llbracket v \rrbracket / \$x \rrbracket) \parallel_{M \llbracket \llbracket v \rrbracket / \$x_{<} \rrbracket} (Q \llbracket \llbracket v \rrbracket / \$x \rrbracket))$   
 $\bigwedge P Q M \sigma. \sigma(\$x' \mapsto_s \llbracket v \rrbracket) \dagger (P \parallel_M Q) = \sigma \dagger (P \parallel_{M \llbracket \llbracket v \rrbracket / \$x' \rrbracket} Q)$   
**by** (*rel-auto*) $+$

**lemma** *bool-pbm-subst*  $[usubst]:$   
**fixes**  $x :: (- \implies 'a)$   
**shows**  
 $\bigwedge P Q M \sigma. \sigma(\$x \mapsto_s \text{false}) \dagger (P \parallel_M Q) = \sigma \dagger ((P \llbracket \text{false} / \$x \rrbracket) \parallel_{M \llbracket \text{false} / \$x_{<} \rrbracket} (Q \llbracket \text{false} / \$x \rrbracket))$   
 $\bigwedge P Q M \sigma. \sigma(\$x \mapsto_s \text{true}) \dagger (P \parallel_M Q) = \sigma \dagger ((P \llbracket \text{true} / \$x \rrbracket) \parallel_{M \llbracket \text{true} / \$x_{<} \rrbracket} (Q \llbracket \text{true} / \$x \rrbracket))$   
 $\bigwedge P Q M \sigma. \sigma(\$x' \mapsto_s \text{false}) \dagger (P \parallel_M Q) = \sigma \dagger (P \parallel_{M \llbracket \text{false} / \$x' \rrbracket} Q)$

$\bigwedge P Q M \sigma. \sigma(\$x' \mapsto_s \text{true}) \dagger (P \parallel_M Q) = \sigma \dagger (P \parallel_M \llbracket \text{true}/\$x' \rrbracket Q)$   
**by** (rel-auto)+

**lemma** *zero-one-pbm-subst* [usubst]:

**fixes**  $x :: (- \implies 'a)$

**shows**

$\bigwedge P Q M \sigma. \sigma(\$x \mapsto_s 0) \dagger (P \parallel_M Q) = \sigma \dagger ((P \llbracket 0/\$x \rrbracket) \parallel_M \llbracket 0/\$x \rrbracket (Q \llbracket 0/\$x \rrbracket))$

$\bigwedge P Q M \sigma. \sigma(\$x \mapsto_s 1) \dagger (P \parallel_M Q) = \sigma \dagger ((P \llbracket 1/\$x \rrbracket) \parallel_M \llbracket 1/\$x \rrbracket (Q \llbracket 1/\$x \rrbracket))$

$\bigwedge P Q M \sigma. \sigma(\$x' \mapsto_s 0) \dagger (P \parallel_M Q) = \sigma \dagger (P \parallel_M \llbracket 0/\$x' \rrbracket Q)$

$\bigwedge P Q M \sigma. \sigma(\$x' \mapsto_s 1) \dagger (P \parallel_M Q) = \sigma \dagger (P \parallel_M \llbracket 1/\$x' \rrbracket Q)$

**by** (rel-auto)+

**lemma** *numeral-pbm-subst* [usubst]:

**fixes**  $x :: (- \implies 'a)$

**shows**

$\bigwedge P Q M \sigma. \sigma(\$x \mapsto_s \text{numeral } n) \dagger (P \parallel_M Q) = \sigma \dagger ((P \llbracket \text{numeral } n/\$x \rrbracket) \parallel_M \llbracket \text{numeral } n/\$x \rrbracket (Q \llbracket \text{numeral } n/\$x \rrbracket))$

$\bigwedge P Q M \sigma. \sigma(\$x' \mapsto_s \text{numeral } n) \dagger (P \parallel_M Q) = \sigma \dagger (P \parallel_M \llbracket \text{numeral } n/\$x' \rrbracket Q)$

**by** (rel-auto)+

## 19.8 Parallel-by-merge laws

**lemma** *par-by-merge-false* [simp]:

$P \parallel_{\text{false}} Q = \text{false}$

**by** (rel-auto)

**lemma** *par-by-merge-left-false* [simp]:

$\text{false} \parallel_M Q = \text{false}$

**by** (rel-auto)

**lemma** *par-by-merge-right-false* [simp]:

$P \parallel_M \text{false} = \text{false}$

**by** (rel-auto)

**lemma** *par-by-merge-seq-add*:  $(P \parallel_M Q) ;; R = (P \parallel_M ;; R Q)$

**by** (simp add: par-by-merge-def seqr-assoc)

A skip parallel-by-merge yields a skip whenever the parallel predicates are both feasible.

**lemma** *par-by-merge-skip*:

**assumes**  $P ;; \text{true} = \text{true} Q ;; \text{true} = \text{true}$

**shows**  $P \parallel_{\text{skip}_m} Q = \text{II}$

**using** *assms* **by** (rel-auto)

**lemma** *skip-merge-swap*:  $\text{swap}_m ;; \text{skip}_m = \text{skip}_m$

**by** (rel-auto)

**lemma** *par-sep-swap*:  $P \parallel_s Q ;; \text{swap}_m = Q \parallel_s P$

**by** (rel-auto)

Parallel-by-merge commutes when the merge predicate is unchanged by swap

**lemma** *par-by-merge-commute-swap*:

**shows**  $P \parallel_M Q = Q \parallel_{\text{swap}_m} ;; M P$

**proof** –

**have**  $Q \parallel_{\text{swap}_m} ;; M P = (((Q ;; U0) \wedge (P ;; U1) \wedge \$\mathbf{v}_{<}' =_u \$\mathbf{v}) ;; \text{swap}_m) ;; M$



by (simp add: par-by-merge-def seqr-assoc)  
 also have ... = (((Q ;; U0 ;; swap<sub>m</sub>) ∧ (P ;; U1 ;; swap<sub>m</sub>) ∧ \$v\_{<}' =\_u \$v) ;; M)  
 by (rel-auto)  
 also have ... = (((Q ;; U1) ∧ (P ;; U0) ∧ \$v\_{<}' =\_u \$v) ;; M)  
 by (simp add: U0-swap U1-swap)  
 also have ... = P ∥<sub>M</sub> Q  
 by (simp add: par-by-merge-def utp-pred-laws.inf.left-commute)  
 finally show ?thesis ..  
 qed

**theorem** par-by-merge-commute:  
 assumes *M is SymMerge*  
 shows  $P \parallel_M Q = Q \parallel_M P$   
 by (metis Healthy-if assms par-by-merge-commute-swap)

**lemma** par-by-merge-mono-1:  
 assumes  $P_1 \sqsubseteq P_2$   
 shows  $P_1 \parallel_M Q \sqsubseteq P_2 \parallel_M Q$   
 using assms by (rel-auto)

**lemma** par-by-merge-mono-2:  
 assumes  $Q_1 \sqsubseteq Q_2$   
 shows  $(P \parallel_M Q_1) \sqsubseteq (P \parallel_M Q_2)$   
 using assms by (rel-blast)

**lemma** par-by-merge-mono:  
 assumes  $P_1 \sqsubseteq P_2$   $Q_1 \sqsubseteq Q_2$   
 shows  $P_1 \parallel_M Q_1 \sqsubseteq P_2 \parallel_M Q_2$   
 by (meson assms dual-order.trans par-by-merge-mono-1 par-by-merge-mono-2)

**theorem** par-by-merge-assoc:  
 assumes *M is SymMerge AssocMerge M*  
 shows  $(P \parallel_M Q) \parallel_M R = P \parallel_M (Q \parallel_M R)$

**proof** –  
 have  $(P \parallel_M Q) \parallel_M R = ((P ;; U0) \wedge (Q ;; U0 ;; U1) \wedge (R ;; U1 ;; U1) \wedge $v_{<}' =_u $v) ;; \mathbf{M3}(M)$   
 by (rel-blast)  
 also have ... =  $((P ;; U0) \wedge (Q ;; U0 ;; U1) \wedge (R ;; U1 ;; U1) \wedge $v_{<}' =_u $v) ;; rotate_m ;; \mathbf{M3}(M)$   
 using AssocMerge-def assms(2) by force  
 also have ... =  $((Q ;; U0) \wedge (R ;; U0 ;; U1) \wedge (P ;; U1 ;; U1) \wedge $v_{<}' =_u $v) ;; \mathbf{M3}(M)$   
 by (rel-blast)  
 also have ... =  $(Q \parallel_M R) \parallel_M P$   
 by (rel-blast)  
 also have ... =  $P \parallel_M (Q \parallel_M R)$   
 by (simp add: assms(1) par-by-merge-commute)  
 finally show ?thesis .  
 qed

**theorem** par-by-merge-choice-left:  
 $(P \sqcap Q) \parallel_M R = (P \parallel_M R) \sqcap (Q \parallel_M R)$   
 by (rel-auto)

**theorem** par-by-merge-choice-right:  
 $P \parallel_M (Q \sqcap R) = (P \parallel_M Q) \sqcap (P \parallel_M R)$   
 by (rel-auto)

**theorem** *par-by-merge-USUP-mem-left*:  
 $(\prod_{i \in I} i \cdot P(i)) \parallel_M Q = (\prod_{i \in I} i \cdot P(i)) \parallel_M Q$   
**by** (*rel-auto*)

**theorem** *par-by-merge-USUP-ind-left*:  
 $(\prod_{i \cdot P(i)} i) \parallel_M Q = (\prod_{i \cdot P(i)} i) \parallel_M Q$   
**by** (*rel-auto*)

**theorem** *par-by-merge-USUP-mem-right*:  
 $P \parallel_M (\prod_{i \in I} i \cdot Q(i)) = (\prod_{i \in I} i \cdot P \parallel_M Q(i))$   
**by** (*rel-auto*)

**theorem** *par-by-merge-USUP-ind-right*:  
 $P \parallel_M (\prod_{i \cdot Q(i)} i) = (\prod_{i \cdot P \parallel_M Q(i)} i)$   
**by** (*rel-auto*)

## 19.9 Example: Simple State-Space Division

The following merge predicate divides the state space using a pair of independent lenses.

**definition** *StateMerge* ::  $('a \Rightarrow 'a) \Rightarrow ('b \Rightarrow 'a) \Rightarrow 'a \text{ merge } (M[-]_{\sigma})$  **where**  
 $[upred-defs]: M[a|b]_{\sigma} = (\$v' =_u (\$v_{<} \oplus \$0 - v \text{ on } \&a) \oplus \$1 - v \text{ on } \&b)$

**lemma** *swap-StateMerge*:  $a \bowtie b \Rightarrow (swap_m ;; M[a|b]_{\sigma}) = M[b|a]_{\sigma}$   
**by** (*rel-auto*, *simp-all add: lens-indep-comm*)

**abbreviation** *StateParallel* ::  $'a \text{ hrel} \Rightarrow ('a \Rightarrow 'a) \Rightarrow ('b \Rightarrow 'a) \Rightarrow 'a \text{ hrel} \Rightarrow 'a \text{ hrel } (-|-|_{\sigma} -$   
 $[85,0,0,86] \ 86)$

**where**  $P \mid a|b|_{\sigma} Q \equiv P \parallel_{M[a|b]_{\sigma}} Q$

**lemma** *StateParallel-commute*:  $a \bowtie b \Rightarrow P \mid a|b|_{\sigma} Q = Q \mid b|a|_{\sigma} P$   
**by** (*metis par-by-merge-commute-swap swap-StateMerge*)

**lemma** *StateParallel-form*:

$P \mid a|b|_{\sigma} Q = (\exists (st_0, st_1) \cdot P[\ll st_0 \gg / \$v'] \wedge Q[\ll st_1 \gg / \$v'] \wedge \$v' =_u (\$v \oplus \ll st_0 \gg \text{ on } \&a) \oplus \ll st_1 \gg \text{ on } \&b)$   
**by** (*rel-auto*)

**lemma** *StateParallel-frame-left*:

**assumes** *vwb-lens a*  
**shows**  $a:[P] \mid a|b|_{\sigma} Q = P \mid a|b|_{\sigma} Q$   
**using** *assms by (simp add: StateParallel-form, rel-auto, blast, fastforce)*

**lemma** *StateParallel-frame-right*:

**assumes** *vwb-lens b*  
**shows**  $P \mid a|b|_{\sigma} b:[Q] = P \mid a|b|_{\sigma} Q$   
**using** *assms by (simp add: StateParallel-form, rel-auto, blast, fastforce)*

We can frame all the variables that the parallel operator refers to

**lemma** *StateParallel-frame*:

**assumes** *vwb-lens a vwb-lens b a \bowtie b*  
**shows**  $\{\&a, \&b\}:[P \mid a|b|_{\sigma} Q] = P \mid a|b|_{\sigma} Q$   
**using** *assms*  
**apply** (*simp add: StateParallel-form, rel-auto*)  
**using** *lens-indep-comm apply fastforce+*

**done**

**lemma** *StateParallel-skip*:

**assumes** *vwb-lens a vwb-lens b*  $a \bowtie b$   
**shows**  $\llbracket a \rrbracket_\sigma P = b \cdot \llbracket P \rrbracket$   
**using** *assms* **by** (*rel-auto*)

Parallel Hoare logic rule. This employs something similar to separating conjunction in the postcondition, but we explicitly require that the two conjuncts only refer to variables on the left and right of the parallel composition explicitly.

**theorem** *StateParallel-hoare* [*hoare*]:

**assumes**  $\llbracket c \rrbracket P \llbracket d_1 \rrbracket_u \llbracket c \rrbracket Q \llbracket d_2 \rrbracket_u$   $a \bowtie b$   $a \nmid d_1$   $b \nmid d_2$   
**shows**  $\llbracket c \rrbracket P \llbracket a \rrbracket_\sigma Q \llbracket d_1 \wedge d_2 \rrbracket_u$

**proof** –

— Parallelise the specification

**from** *assms*(4,5)

**have**  $1: (\llbracket c \rrbracket < \Rightarrow \llbracket d_1 \wedge d_2 \rrbracket >) \sqsubseteq (\llbracket c \rrbracket < \Rightarrow \llbracket d_1 \rrbracket >) \llbracket a \rrbracket_\sigma (\llbracket c \rrbracket < \Rightarrow \llbracket d_2 \rrbracket >)$  (**is** *?lhs*  $\sqsubseteq$  *?rhs*)  
**by** (*simp add: StateParallel-form, rel-auto,metis assms(3) lens-indep-comm*)

— Prove Hoare rule by monotonicity of parallelism

**have**  $2: ?rhs \sqsubseteq P \llbracket a \rrbracket_\sigma Q$

**proof** (*rule par-by-merge-mono*)

**show**  $(\llbracket c \rrbracket < \Rightarrow \llbracket d_1 \rrbracket >) \sqsubseteq P$

**using** *assms*(1) *hoare-r-def* **by** *auto*

**show**  $(\llbracket c \rrbracket < \Rightarrow \llbracket d_2 \rrbracket >) \sqsubseteq Q$

**using** *assms*(2) *hoare-r-def* **by** *auto*

**qed**

**show** *?thesis*

**unfolding** *hoare-r-def* **using** 1 2 *order-trans* **by** *auto*

**qed**

Specialised version of the above law where an invariant expression referring to variables outside the frame is preserved.

**theorem** *StateParallel-frame-hoare* [*hoare*]:

**assumes** *vwb-lens a vwb-lens b*  $a \bowtie b$   $a \nmid d_1$   $b \nmid d_2$   $a \# c_1$   $b \# c_1$   $\llbracket c_1 \wedge c_2 \rrbracket P \llbracket d_1 \rrbracket_u \llbracket c_1 \wedge c_2 \rrbracket Q \llbracket d_2 \rrbracket_u$   
**shows**  $\llbracket c_1 \wedge c_2 \rrbracket P \llbracket a \rrbracket_\sigma Q \llbracket c_1 \wedge d_1 \wedge d_2 \rrbracket_u$

**proof** –

**have**  $\llbracket c_1 \wedge c_2 \rrbracket \{ \&a, \&b \} : \llbracket P \rrbracket_\sigma Q \llbracket c_1 \wedge d_1 \wedge d_2 \rrbracket_u$

**by** (*auto intro!: frame-hoare-r' StateParallel-hoare simp add: assms unrest plus-vwb-lens*)

**thus** *?thesis*

**by** (*simp add: StateParallel-frame assms*)

**qed**

**end**

## 20 Relational Operational Semantics

**theory** *utp-rel-opsem*

**imports**

*utp-rel-laws*

*utp-hoare*

**begin**

This theory uses the laws of relational calculus to create a basic operational semantics. It is based on Chapter 10 of the UTP book [5].

**fun** *trel* ::  $'\alpha$  *usubst*  $\times$   $'\alpha$  *hrel*  $\Rightarrow$   $'\alpha$  *usubst*  $\times$   $'\alpha$  *hrel*  $\Rightarrow$  *bool* (**infix**  $\rightarrow_u$  85) **where**

$$(\sigma, P) \rightarrow_u (\varrho, Q) \longleftrightarrow (\langle \sigma \rangle_a ;; P) \sqsubseteq (\langle \varrho \rangle_a ;; Q)$$

**lemma** *trans-trel*:

$\llbracket (\sigma, P) \rightarrow_u (\varrho, Q); (\varrho, Q) \rightarrow_u (\varphi, R) \rrbracket \implies (\sigma, P) \rightarrow_u (\varphi, R)$   
**by** *auto*

**lemma** *skip-trel*:  $(\sigma, II) \rightarrow_u (\sigma, II)$

**by** *simp*

**lemma** *assigns-trel*:  $(\sigma, \langle \varrho \rangle_a) \rightarrow_u (\varrho \circ \sigma, II)$

**by** (*simp add: assigns-comp*)

**lemma** *assign-trel*:

$(\sigma, x := v) \rightarrow_u (\sigma(x \mapsto_s \sigma \dagger v), II)$   
**by** (*simp add: assigns-comp usubst*)

**lemma** *seq-trel*:

**assumes**  $(\sigma, P) \rightarrow_u (\varrho, Q)$   
**shows**  $(\sigma, P ;; R) \rightarrow_u (\varrho, Q ;; R)$   
**by** (*metis (no-types, lifting) assms order-refl seqr-assoc seqr-mono trel.simps*)

**lemma** *seq-skip-trel*:

$(\sigma, II ;; P) \rightarrow_u (\sigma, P)$   
**by** *simp*

**lemma** *nondet-left-trel*:

$(\sigma, P \sqcap Q) \rightarrow_u (\sigma, P)$

**by** (*metis (no-types, hide-lams) disj-comm disj-upred-def semilattice-sup-class.sup.absorb-iff1 semilattice-sup-class.sup.l*  
*seqr-or-distr trel.simps*)

**lemma** *nondet-right-trel*:

$(\sigma, P \sqcap Q) \rightarrow_u (\sigma, Q)$

**by** (*simp add: seqr-mono*)

**lemma** *rcond-true-trel*:

**assumes**  $\sigma \dagger b = \text{true}$   
**shows**  $(\sigma, P \triangleleft b \triangleright_r Q) \rightarrow_u (\sigma, P)$   
**using** *assms*  
**by** (*simp add: assigns-r-comp usubst aext-true cond-unit-T*)

**lemma** *rcond-false-trel*:

**assumes**  $\sigma \dagger b = \text{false}$   
**shows**  $(\sigma, P \triangleleft b \triangleright_r Q) \rightarrow_u (\sigma, Q)$   
**using** *assms*  
**by** (*simp add: assigns-r-comp usubst aext-false cond-unit-F*)

**lemma** *while-true-trel*:

**assumes**  $\sigma \dagger b = \text{true}$   
**shows**  $(\sigma, \text{while } b \text{ do } P \text{ od}) \rightarrow_u (\sigma, P ;; \text{while } b \text{ do } P \text{ od})$   
**by** (*metis assms rcond-true-trel while-unfold*)

**lemma** *while-false-trel*:

**assumes**  $\sigma \dagger b = \text{false}$   
**shows**  $(\sigma, \text{while } b \text{ do } P \text{ od}) \rightarrow_u (\sigma, II)$   
**by** (*metis assms rcond-false-trel while-unfold*)

Theorem linking Hoare calculus and operational semantics. If we start  $Q$  in a state  $\sigma_0$  satisfying  $p$ , and  $Q$  reaches final state  $\sigma_1$  then  $r$  holds in this final state.

**theorem** *hoare-opsem-link*:

```

  {p}Q{r}_u = (∀ σ0 σ1. 'σ0 † p' ∧ (σ0, Q) →u (σ1, II) → 'σ1 † r')
  apply (rel-auto)
  apply (rename-tac a b)
  apply (drule-tac x=λ -. a in spec, simp)
  apply (drule-tac x=λ -. b in spec, simp)
done

```

**declare** *trel.simps* [*simp del*]

**end**

## 20.1 Variable blocks

**theory** *utp-local*

**imports** *utp-theory*

**begin**

Local variables are represented as lenses whose view type is a list of values. A variable therefore effectively records the stack of values that variable has had, if any. This allows us to denote variable scopes using assignments that push and pop this stack to add or delete a particular local variable.

**type-synonym** ('a, 'α) *lvar* = ('a list ⇒ 'α)

Different UTP theories have different assignment operators; consequently in order to generically characterise variable blocks we need to abstractly characterise assignments. We first create two polymorphic constants that characterise the underlying program state model of a UTP theory.

**consts**

```

  pvar          :: ('T, 'α) uthy ⇒ 'β ⇒ 'α (s1)
  pvar-assigns :: ('T, 'α) uthy ⇒ 'β usubst ⇒ 'α hrel ((-)1)

```

*pvar* is a lens from the program state, 'β, to the overall global state 'α, which also contains none user-space information, such as observational variables. *pvar-assigns* takes as parameter a UTP theory and returns an assignment operator which maps a substitution over the program state to a homogeneous relation on the global state. We now set up some syntax translations for these operators.

**syntax**

```

  -svid-pvar :: ('T, 'α) uthy ⇒ svid (s1)
  -thy-asgn  :: ('T, 'α) uthy ⇒ svids ⇒ uexprs ⇒ logic (infixr ::= 72)

```

**translations**

```

  -svid-pvar T => CONST pvar T
  -thy-asgn T xs vs => CONST pvar-assigns T (-mk-usubst (CONST id) xs vs)

```

Next, we define constants to represent the top most variable on the local variable stack, and the remainder after this. We define these in terms of the list lens, and so for each another lens is produced.

**definition** *top-var* :: ('a::two, 'α) *lvar* ⇒ ('a ⇒ 'α) **where**

[*upred-defs*]: *top-var*  $x = (list-lens\ 0\ ;_L\ x)$

The remainder of the local variable stack (the tail)

**definition** *rest-var* :: ('a::two, 'α) lvar ⇒ ('a list ⇒ 'α) **where**  
[upred-defs]: *rest-var* x = (tl-lens ;<sub>L</sub> x)

We can show that the top variable is a mainly well-behaved lense, and that the top most variable lens is independent of the rest of the stack.

**lemma** *top-mwb-lens* [simp]: *mwb-lens* x ⇒ *mwb-lens* (top-var x)  
**by** (simp add: list-mwb-lens top-var-def)

**lemma** *top-rest-var-indep* [simp]:  
*mwb-lens* x ⇒ top-var x ⋈ rest-var x  
**by** (simp add: lens-indep-left-comp rest-var-def top-var-def)

**lemma** *top-var-pres-indep* [simp]:  
x ⋈ y ⇒ top-var x ⋈ y  
**by** (simp add: lens-indep-left-ext top-var-def)

**syntax**  
-*top-var* :: svid ⇒ svid (@- [999] 999)  
-*rest-var* :: svid ⇒ svid (↓- [999] 999)

**translations**  
-*top-var* x == CONST top-var x  
-*rest-var* x == CONST rest-var x

With operators to represent local variables, assignments, and stack manipulation defined, we can go about defining variable blocks themselves.

**definition** *var-begin* :: ('T, 'α) uthy ⇒ ('a, 'β) lvar ⇒ 'α hrel **where**  
[urel-defs]: *var-begin* T x = x ::<sub>T</sub> ⟨⋈ undefined⟩<sub>u</sub> &x

**definition** *var-end* :: ('T, 'α) uthy ⇒ ('a, 'β) lvar ⇒ 'α hrel **where**  
[urel-defs]: *var-end* T x = (x ::<sub>T</sub> tail<sub>u</sub>(&x))

*var-begin* takes as parameters a UTP theory and a local variable, and uses the theory assignment operator to push and undefined value onto the variable stack. *var-end* removes the top most variable from the stack in a similar way.

**definition** *var-vlet* :: ('T, 'α) uthy ⇒ ('a, 'α) lvar ⇒ 'α hrel **where**  
[urel-defs]: *var-vlet* T x = ((\$x ≠<sub>u</sub> ⟨⟩) ∧ II<sub>T</sub>)

Next we set up the typical UTP variable block syntax, though with a suitable subscript index to represent the UTP theory parameter.

**syntax**  
-*var-begin* :: logic ⇒ svid ⇒ logic (var<sub>1</sub> - [100] 100)  
-*var-begin-asn* :: logic ⇒ svid ⇒ logic ⇒ logic (var<sub>1</sub> - := -)  
-*var-end* :: logic ⇒ svid ⇒ logic (end<sub>1</sub> - [100] 100)  
-*var-vlet* :: logic ⇒ svid ⇒ logic (vlet<sub>1</sub> - [100] 100)  
-*var-scope* :: logic ⇒ svid ⇒ logic ⇒ logic (var<sub>1</sub> - ⋅ - [0,10] 10)  
-*var-scope-ty* :: logic ⇒ svid ⇒ type ⇒ logic ⇒ logic (var<sub>1</sub> - :: ⋅ - [0,0,10] 10)  
-*var-scope-ty-assign* :: logic ⇒ svid ⇒ type ⇒ logic ⇒ logic ⇒ logic (var<sub>1</sub> - :: ⋅ - := ⋅ - [0,0,0,10] 10)

**translations**  
-*var-begin* T x == CONST var-begin T x  
-*var-begin-asn* T x e ==> var<sub>T</sub> x ;; @x ::<sub>T</sub> e  
-*var-end* T x == CONST var-end T x  
-*var-vlet* T x == CONST var-vlet T x

$$\begin{aligned} \text{var}_T x \cdot P &\Rightarrow \text{var}_T x ;; ((\lambda x. P) (CONST \text{top-var } x)) ;; \text{end}_T x \\ \text{var}_T x \cdot P &\Rightarrow \text{var}_T x ;; ((\lambda x. P) (CONST \text{top-var } x)) ;; \text{end}_T x \end{aligned}$$

In order to substantiate standard variable block laws, we need some underlying laws about assignments, which is the purpose of the following locales.

**locale** *utp-prog-var* = *utp-theory*  $\mathcal{T}$  **for**  $\mathcal{T} :: ('T, 'a) \text{thy} (\text{structure}) +$   
**fixes**  $\mathcal{VT} :: 'a \text{ itself}$   
**assumes** *pvar-uvar*: *vwb-lens* ( $s :: 'a \Rightarrow 'a$ )  
**and** *Healthy-pvar-assigns* [*closure*]:  $\langle \sigma :: 'a \text{ usubst} \rangle$  is  $\mathcal{H}$   
**and** *pvar-assigns-comp*:  $(\langle \sigma \rangle ;; \langle \varrho \rangle) = \langle \varrho \circ \sigma \rangle$

We require that (1) the user-space variable is a very well-behaved lens, (2) that the assignment operator is healthy, and (3) that composing two assignments is equivalent to composing their substitutions. The next locale extends this with a left unit.

**locale** *utp-local-var* = *utp-prog-var*  $\mathcal{T}$   $V + \text{utp-theory-left-unital } \mathcal{T}$  **for**  $\mathcal{T} :: ('T, 'a) \text{thy} (\text{structure})$   
**and**  $V :: 'a \text{ itself} +$   
**assumes** *pvar-assign-unit*:  $\langle \text{id} :: 'a \text{ usubst} \rangle = \mathcal{II}$   
**begin**

If a left unit exists then an assignment with an identity substitution should yield the identity relation, as the above assumption requires. With these laws available, we can prove the main laws of variable blocks.

**lemma** *var-begin-healthy* [*closure*]:  
**fixes**  $x :: ('a, 'a) \text{ lvar}$   
**shows** *var*  $x$  is  $\mathcal{H}$   
**by** (*simp add: var-begin-def Healthy-pvar-assigns*)

**lemma** *var-end-healthy* [*closure*]:  
**fixes**  $x :: ('a, 'a) \text{ lvar}$   
**shows** *end*  $x$  is  $\mathcal{H}$   
**by** (*simp add: var-end-def Healthy-pvar-assigns*)

The beginning and end of a variable block are both healthy theory elements.

**lemma** *var-open-close*:  
**fixes**  $x :: ('a, 'a) \text{ lvar}$   
**assumes** *vwb-lens*  $x$   
**shows**  $(\text{var } x ;; \text{end } x) = \mathcal{II}$   
**by** (*simp add: var-begin-def var-end-def shEx-lift-seq-1 Healthy-pvar-assigns pvar-assigns-comp pvar-assign-unit usubst assms*)

Opening and then immediately closing a variable blocks yields a skip.

**lemma** *var-open-close-commute*:  
**fixes**  $x :: ('a, 'a) \text{ lvar}$  **and**  $y :: ('b, 'b) \text{ lvar}$   
**assumes** *vwb-lens*  $x$  *vwb-lens*  $y$   $x \bowtie y$   
**shows**  $(\text{var } x ;; \text{end } y) = (\text{end } y ;; \text{var } x)$   
**by** (*simp add: var-begin-def var-end-def shEx-lift-seq-1 shEx-lift-seq-2 Healthy-pvar-assigns pvar-assigns-comp assms usubst unrest lens-indep-sym, simp add: assms usubst-upd-comm*)

The beginning and end of variable blocks from different variables commute.

**lemma** *var-block-vacuous*:  
**fixes**  $x :: ('a::\text{two}, 'a) \text{ lvar}$   
**assumes** *vwb-lens*  $x$

```

shows (var x ·  $\mathcal{II}$ ) =  $\mathcal{II}$ 
by (simp add: Left-Unit assms var-end-healthy var-open-close)

```

A variable block with a skip inside results in a skip.

**end**

Example instantiation for the theory of relations

**overloading**

```

rel-pvar == pvar :: (REL, 'α) uthy ⇒ 'α ⇒ 'α
rel-pvar-assigns == pvar-assigns :: (REL, 'α) uthy ⇒ 'α usubst ⇒ 'α hrel

```

**begin**

```

definition rel-pvar :: (REL, 'α) uthy ⇒ 'α ⇒ 'α where

```

```

[upred-defs]: rel-pvar T = 1L

```

```

definition rel-pvar-assigns :: (REL, 'α) uthy ⇒ 'α usubst ⇒ 'α hrel where

```

```

[upred-defs]: rel-pvar-assigns T σ = ⟨σ⟩a

```

**end**

**interpretation** rel-local-var: utp-local-var UTHY(REL, 'α) TYPE('α)

**proof** –

```

interpret vw: vwb-lens pvar REL :: 'α ⇒ 'α

```

```

by (simp add: rel-pvar-def id-vwb-lens)

```

```

show utp-local-var TYPE('α) UTHY(REL, 'α)

```

**proof**

```

show  $\bigwedge \sigma :: 'α \Rightarrow 'α. \langle \sigma \rangle_{REL} \text{ is } \mathcal{H}_{REL}$ 

```

```

by (simp add: rel-pvar-assigns-def rel-hcond-def Healthy-def)

```

```

show  $\bigwedge (\sigma :: 'α \Rightarrow 'α) \varrho. \langle \sigma \rangle_{UTHY(REL, 'α)} \;; \langle \varrho \rangle_{REL} = \langle \varrho \circ \sigma \rangle_{REL}$ 

```

```

by (simp add: rel-pvar-assigns-def assigns-comp)

```

```

show  $\langle id :: 'α \Rightarrow 'α \rangle_{UTHY(REL, 'α)} = \mathcal{II}_{REL}$ 

```

```

by (simp add: rel-pvar-assigns-def rel-unit-def skip-r-def)

```

**qed**

**qed**

**end**

## 21 Meta-theory for the Standard Core

**theory** utp

**imports**

utp-var

utp-expr

utp-unrest

utp-usedby

utp-subst

utp-meta-subst

utp-alphabet

utp-lift

utp-pred

utp-pred-laws

utp-recursion

utp-deduct

utp-rel

utp-rel-laws

utp-tactics

utp-hoare

utp-wp



*utp-theory*  
*utp-concurrency*  
*utp-rel-opsem*  
*utp-local*  
*utp-event*  
**begin end**

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