

Stateful-Failure Reactive Designs in Isabelle/UTP

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Abstract

Stateful-Failure Reactive Designs specialise reactive design contracts with failures traces, as present in languages like CSP and Circus. A failure trace consists of a sequence of events and a refusal set. It intuitively represents a quiescent observation, where certain events have previously occurred, and others are currently being accepted. Following the UTP book, we add an observational variable to represent refusal sets, and healthiness conditions that ensure their well-formedness. Using these, we also specialise our theory of reactive relations with operators to characterise both completed and quiescent interactions, and an accompanying equational theory. We use these to define the core operators — including assignment, event occurrence, and external choice — and specialise our proof strategy to support these. We also demonstrate a link with the CSP failures-divergences semantic model.

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1 Introduction

This document contains a mechanisation in Isabelle/UTP [1] of an specialisation of stateful reactive designs with refusal information, as present in languages like Circus [2].

2 Stateful-Failure Core Types

```
theory utp-sfrd-core
  imports UTP-Reactive-Designs.utp-rea-designs
begin
```

2.1 SFRD Alphabet

alphabet φ *csp-vars* = σ *rsp-vars* +
 $ref :: \varphi$ *set*

declare *csp-vars.defs* [*lens-defs*]
declare *csp-vars.splits* [*alpha-splits*]

The following two locale interpretations are a technicality to improve the behaviour of the automatic tactics. They enable (re)interpretation of state spaces in order to remove any occurrences of lens types, replacing them by tuple types after the tactics *pred-simp* and *rel-simp* are applied. Eventually, it would be desirable to automate preform these interpretations automatically as part of the **alphabet** command.

interpretation *alphabet-csp-prd*:
 $lens_interp \lambda(ok, wait, tr, m). (ok, wait, tr, ref_v m, more m)$
apply (*unfold-locales*)
apply (*rule injI*)
apply (*clarsimp*)
done

interpretation *alphabet-csp-rel*:
 $lens_interp \lambda(ok, ok', wait, wait', tr, tr', m, m').$
 $(ok, ok', wait, wait', tr, tr', ref_v m, ref_v m', more m, more m')$
apply (*unfold-locales*)
apply (*rule injI*)
apply (*clarsimp*)
done

type-synonym (σ, φ) *st-csp* = $(\sigma, \varphi$ *list*, $(\varphi, unit)$ *csp-vars-scheme*) *rsp*
type-synonym (σ, φ) *action* = (σ, φ) *st-csp hrel*
type-synonym φ *csp* = $(unit, \varphi)$ *st-csp*
type-synonym φ *process* = φ *csp hrel*

There is some slight imprecision with the translations, in that we don't bother to check if the trace event type and refusal set event types are the same. Essentially this is because its very difficult to construct processes where this would be the case. However, it may be better to add a proper ML print translation in the future.

translations

$(type) (\sigma, \varphi) st-csp \leq (type) (\sigma, \varphi$ *list*, $\varphi1$ *csp-vars*) *rsp*
 $(type) (\sigma, \varphi) action \leq (type) (\sigma, \varphi) st-csp hrel$
 $(type) \varphi process \leq (type) (unit, \varphi) action$

notation *csp-vars-child-lens_a* (Σ_c)

notation *csp-vars-child-lens* (Σ_C)

2.2 Basic laws

lemma *R2c-tr-ext*: $R2c (\$tr' =_u \$tr \hat{^}_u \langle [a]_{S<} \rangle) = (\$tr' =_u \$tr \hat{^}_u \langle [a]_{S<} \rangle)$
by (*rel-auto*)

lemma *circus-alpha-bij-lens*:

$bij_lens (\{\$ok, \$ok', \$wait, \$wait', \$tr, \$tr', \$st, \$st', \$ref, \$ref'\}_{\alpha} :: - \implies (s, e) st-csp \times (s, e) st-csp)$
by (*unfold-locales, lens-simp+*)

2.3 Unrestriction laws

lemma *pre-unrest-ref* [*unrest*]: $\$ref \# P \implies \$ref \# pre_R(P)$
by (*simp add: pre_R-def unrest*)

lemma *peri-unrest-ref* [*unrest*]: $\$ref \# P \implies \$ref \# peri_R(P)$
by (*simp add: peri_R-def unrest*)

lemma *post-unrest-ref* [*unrest*]: $\$ref \# P \implies \$ref \# post_R(P)$
by (*simp add: post_R-def unrest*)

lemma *cmt-unrest-ref* [*unrest*]: $\$ref \# P \implies \$ref \# cmt_R(P)$
by (*simp add: cmt_R-def unrest*)

lemma *st-lift-unrest-ref'* [*unrest*]: $\$ref' \# \lceil b \rceil_{S<} \implies$
by (*rel-auto*)

lemma *RHS-design-ref-unrest* [*unrest*]:
 $\llbracket \$ref \# P; \$ref \# Q \rrbracket \implies \$ref \# (\mathbf{R}_s(P \vdash Q)) \llbracket false/\$wait \rrbracket$
by (*simp add: RHS-def R1-def R2c-def R2s-def R3h-def design-def usubst unrest*)

lemma *R1-ref-unrest* [*unrest*]: $\$ref \# P \implies \$ref \# R1(P)$
by (*simp add: R1-def unrest*)

lemma *R2c-ref-unrest* [*unrest*]: $\$ref \# P \implies \$ref \# R2c(P)$
by (*simp add: R2c-def unrest*)

lemma *R1-ref'-unrest* [*unrest*]: $\$ref' \# P \implies \$ref' \# R1(P)$
by (*simp add: R1-def unrest*)

lemma *R2c-ref'-unrest* [*unrest*]: $\$ref' \# P \implies \$ref' \# R2c(P)$
by (*simp add: R2c-def unrest*)

lemma *R2s-notin-ref'*: $R2s(\lceil \ll x \gg \rceil_{S<} \notin_u \$ref') = (\lceil \ll x \gg \rceil_{S<} \notin_u \$ref')$
by (*pred-auto*)

lemma *unrest-circus-alpha*:
fixes $P :: ('e, 't) \text{ action}$
assumes
 $\$ok \# P \ \$ok' \# P \ \$wait \# P \ \$wait' \# P \ \$tr \# P$
 $\$tr' \# P \ \$st \# P \ \$st' \# P \ \$ref \# P \ \$ref' \# P$
shows $\Sigma \# P$
by (*rule bij-lens-unrest-all[OF circus-alpha-bij-lens], simp add: unrest assms*)

lemma *unrest-all-circus-vars*:
fixes $P :: ('s, 'e) \text{ action}$
assumes $\$ok \# P \ \$ok' \# P \ \$wait \# P \ \$wait' \# P \ \$ref \# P \ \Sigma \# r' \ \Sigma \# s \ \Sigma \# s' \ \Sigma \# t \ \Sigma \# t'$
shows $\Sigma \# [\$ref' \mapsto_s r', \$st \mapsto_s s, \$st' \mapsto_s s', \$tr \mapsto_s t, \$tr' \mapsto_s t'] \dagger P$
using *assms*
by (*simp add: bij-lens-unrest-all-eq[OF circus-alpha-bij-lens] unrest-plus-split plus-vwb-lens*)
(simp add: unrest usubst closure)

lemma *unrest-all-circus-vars-st-st'*:
fixes $P :: ('s, 'e) \text{ action}$
assumes $\$ok \# P \ \$ok' \# P \ \$wait \# P \ \$wait' \# P \ \$ref \# P \ \$ref' \# P \ \Sigma \# s \ \Sigma \# s' \ \Sigma \# t \ \Sigma \# t'$
shows $\Sigma \# [\$st \mapsto_s s, \$st' \mapsto_s s', \$tr \mapsto_s t, \$tr' \mapsto_s t'] \dagger P$

```

using assms
by (simp add: bij-lens-unrest-all-eq[OF circus-alpha-bij-lens] unrest-plus-split plus-vwb-lens)
    (simp add: unrest usubst closure)

lemma unrest-all-circus-vars-st:
  fixes  $P :: ('s, 'e) \text{ action}$ 
  assumes  $\$ok \# P \$ok' \# P \$wait \# P \$wait' \# P \$ref \# P \$ref' \# P \$st' \# P \Sigma \# s \Sigma \# t \Sigma \# t'$ 
  shows  $\Sigma \# [\$st \mapsto_s s, \$tr \mapsto_s t, \$tr' \mapsto_s t'] \dagger P$ 
  using assms
  by (simp add: bij-lens-unrest-all-eq[OF circus-alpha-bij-lens] unrest-plus-split plus-vwb-lens)
    (simp add: unrest usubst closure)

lemma unrest-any-circus-var:
  fixes  $P :: ('s, 'e) \text{ action}$ 
  assumes  $\$ok \# P \$ok' \# P \$wait \# P \$wait' \# P \$ref \# P \$ref' \# P \Sigma \# s \Sigma \# s' \Sigma \# t \Sigma \# t'$ 
  shows  $x \# [\$st \mapsto_s s, \$st' \mapsto_s s', \$tr \mapsto_s t, \$tr' \mapsto_s t'] \dagger P$ 
  by (simp add: unrest-all-var unrest-all-circus-vars-st-st' assms)

lemma unrest-any-circus-var-st:
  fixes  $P :: ('s, 'e) \text{ action}$ 
  assumes  $\$ok \# P \$ok' \# P \$wait \# P \$wait' \# P \$ref \# P \$ref' \# P \$st' \# P \Sigma \# s \Sigma \# t \Sigma \# t'$ 
  shows  $x \# [\$st \mapsto_s s, \$tr \mapsto_s t, \$tr' \mapsto_s t'] \dagger P$ 
  by (simp add: unrest-all-var unrest-all-circus-vars-st assms)

end

```

3 Stateful-Failure Reactive Relations

```

theory utp-sfrd-rel
  imports utp-sfrd-core
begin

```

3.1 Healthiness Conditions

CSP Reactive Relations

definition $CRR :: ('s, 'e) \text{ action} \Rightarrow ('s, 'e) \text{ action}$ **where**
[upred-defs]: $CRR(P) = (\exists \$ref \cdot RR(P))$

lemma *CRR-idem*: $CRR(CRR(P)) = CRR(P)$
by (*rel-auto*)

lemma *Idempotent-CRR* [*closure*]: *Idempotent CRR*
by (*simp add: CRR-idem Idempotent-def*)

lemma *Continuous-CRR* [*closure*]: *Continuous CRR*
by (*rel-blast*)

lemma *CRR-intro*:
assumes $\$ref \# P$ *P is RR*
shows *P is CRR*
by (*simp add: CRR-def Healthy-def, simp add: Healthy-if assms ex-unrest*)

lemma *CRR-form*: $CRR(P) = (\exists \{\$ok, \$ok', \$wait, \$wait', \$ref\} \cdot (\exists tt_0 \cdot P[\langle \rangle / \$tr][\langle \langle tt_0 \rangle \rangle / \$tr']) \wedge \$tr' =_u \$tr \hat{^}_u \langle \langle tt_0 \rangle \rangle)$

by (rel-auto; fastforce)

lemma *CRR-seqr-form*:

$CRR(P) ;; CRR(Q) =$

$(\exists tt_1 \cdot \exists tt_2 \cdot ((\exists \{ \$ok, \$ok', \$wait, \$wait', \$ref \} \cdot P) [\langle \rangle / \$tr] [\ll tt_1 \gg / \$tr'] ;;$
 $(\exists \{ \$ok, \$ok', \$wait, \$wait', \$ref \} \cdot Q) [\langle \rangle / \$tr] [\ll tt_2 \gg / \$tr'] \wedge \$tr' =_u \$tr \wedge_u$
 $\ll tt_1 \gg \wedge_u \ll tt_2 \gg))$

apply (rel-auto)

apply (metis (no-types, hide-lams) Prefix-Order.prefixE append.assoc plus-list-def trace-class.add-diff-cancel-left)

apply (metis append.assoc le-add plus-list-def trace-class.add-diff-cancel-left)

done

CSP Reactive Conditions

definition *CRC* :: ('s,'e) action \Rightarrow ('s,'e) action **where**

[upred-defs]: $CRC(P) = (\exists \$ref \cdot RC(P))$

lemma *CRC-intro*:

assumes $\$ref \# P$ *P is RC*

shows *P is CRC*

by (simp add: CRC-def Healthy-def, simp add: Healthy-if assms ex-unrest)

lemma *CRC-intro'*:

assumes *P is CRR* *P is RC*

shows *P is CRC*

by (metis CRC-def CRR-def Healthy-def RC-implies-RR assms)

lemma *ref-unrest-RR* [unrest]: $\$ref \# P \Longrightarrow \$ref \# RR P$

by (rel-auto, blast+)

lemma *ref-unrest-RC1* [unrest]: $\$ref \# P \Longrightarrow \$ref \# RC1 P$

by (rel-auto, blast+)

lemma *ref-unrest-RC* [unrest]: $\$ref \# P \Longrightarrow \$ref \# RC P$

by (simp add: RC-R2-def ref-unrest-RC1 ref-unrest-RR)

lemma *RR-ex-ref*: $RR (\exists \$ref \cdot RR P) = (\exists \$ref \cdot RR P)$

by (rel-auto)

lemma *RC1-ex-ref*: $RC1 (\exists \$ref \cdot RC1 P) = (\exists \$ref \cdot RC1 P)$

by (rel-auto, meson dual-order.trans)

lemma *ex-ref'-RR-closed* [closure]:

assumes *P is RR*

shows $(\exists \$ref' \cdot P)$ *is RR*

proof –

have $RR (\exists \$ref' \cdot RR(P)) = (\exists \$ref' \cdot RR(P))$

by (rel-auto)

thus ?thesis

by (metis Healthy-def assms)

qed

lemma *CRC-idem*: $CRC(CRC(P)) = CRC(P)$

apply (simp add: CRC-def ex-unrest unrest)

apply (simp add: RC-def RR-ex-ref)

apply (*metis* (*no-types*, *hide-lams*) *Healthy-def RC1-RR-closed RC1-ex-ref RR-ex-ref RR-idem*)
done

lemma *Idempotent-CRC* [*closure*]: *Idempotent CRC*
by (*simp add: CRC-idem Idempotent-def*)

3.2 Closure Properties

lemma *CRR-implies-RR* [*closure*]:
assumes *P is CRR*
shows *P is RR*

proof –
have $RR(CRR(P)) = CRR(P)$
by (*rel-auto*)
thus *?thesis*
by (*metis Healthy-def' assms*)
qed

lemma *CRC-intro''*:
assumes *P is CRR P is RC1*
shows *P is CRC*
by (*simp add: CRC-intro' CRR-implies-RR RC-intro' assms*)

lemma *CRC-implies-RR* [*closure*]:
assumes *P is CRC*
shows *P is RR*
proof –
have $RR(CRC(P)) = CRC(P)$
by (*rel-auto*)
(*metis* (*no-types*, *lifting*) *Prefix-Order.prefixE Prefix-Order.prefixI append.assoc append-minus*) +
thus *?thesis*
by (*metis Healthy-def assms*)
qed

lemma *CRC-implies-RC* [*closure*]:
assumes *P is CRC*
shows *P is RC*
proof –
have $RC1(CRC(P)) = CRC(P)$
by (*rel-auto, meson dual-order.trans*)
thus *?thesis*
by (*simp add: CRC-implies-RR Healthy-if RC1-def RC-intro assms*)
qed

lemma *CRR-unrest-ref* [*unrest*]: $P \text{ is } CRR \implies \$ref \# P$
by (*metis CRR-def CRR-implies-RR Healthy-def in-var-uvar ref-vwb-lens unrest-as-exists*)

lemma *CRC-implies-CRR* [*closure*]:
assumes *P is CRC*
shows *P is CRR*
apply (*rule CRR-intro*)
apply (*simp-all add: unrest assms closure*)
apply (*metis CRC-def CRC-implies-RC Healthy-def assms in-var-uvar ref-vwb-lens unrest-as-exists*)
done

lemma *unrest-ref'-neg-RC* [*unrest*]:

assumes P is RR P is RC
shows $\$ref' \# P$
proof –
 have $P = (\neg_r \neg_r P)$
 by (*simp add: closure rpred assms*)
 also have $\dots = (\neg_r (\neg_r P) ;; true_r)$
 by (*metis Healthy-if RC1-def RC-implies-RC1 assms(2) calculation*)
 also have $\$ref' \# \dots$
 by (*rel-auto*)
 finally show *?thesis* .
qed

lemma *rea-true-CRR* [closure]: $true_r$ is CRR
 by (*rel-auto*)

lemma *rea-true-CRC* [closure]: $true_r$ is CRC
 by (*rel-auto*)

lemma *false-CRR* [closure]: $false$ is CRR
 by (*rel-auto*)

lemma *false-CRC* [closure]: $false$ is CRC
 by (*rel-auto*)

lemma *st-pred-CRR* [closure]: $[P]_{S<}$ is CRR
 by (*rel-auto*)

lemma *st-post-unrest-ref'* [unrest]: $\$ref' \# [b]_{S>}$
 by (*rel-auto*)

lemma *st-post-CRR* [closure]: $[b]_{S>}$ is CRR
 by (*rel-auto*)

lemma *st-cond-CRC* [closure]: $[P]_{S<}$ is CRC
 by (*rel-auto*)

lemma *rea-rename-CRR-closed* [closure]:
 assumes P is CRR
 shows $P(\lfloor f \rfloor_r)$ is CRR
proof –
 have $\$ref \# (CRR P)(\lfloor f \rfloor_r)$
 by (*rel-auto*)
 thus *?thesis*
 by (*rule-tac CRR-intro, simp-all add: closure Healthy-if assms*)
qed

lemma *st-subst-CRR-closed* [closure]:
 assumes P is CRR
 shows $(\sigma \uparrow_S P)$ is CRR
 by (*rule CRR-intro, simp-all add: unrest closure assms*)

lemma *st-subst-CRC-closed* [closure]:
 assumes P is CRC
 shows $(\sigma \uparrow_S P)$ is CRC
 by (*rule CRC-intro, simp-all add: closure assms unrest*)

lemma *conj-CRC-closed* [closure]:

$\llbracket P \text{ is CRC}; Q \text{ is CRC} \rrbracket \implies (P \wedge Q) \text{ is CRC}$

by (rule CRC-intro, simp-all add: unrest closure)

lemma *disj-CRC-closed* [closure]:

$\llbracket P \text{ is CRC}; Q \text{ is CRC} \rrbracket \implies (P \vee Q) \text{ is CRC}$

by (rule CRC-intro, simp-all add: unrest closure)

lemma *st-cond-left-impl-CRC-closed* [closure]:

$P \text{ is CRC} \implies ([b]_{S<} \Rightarrow_r P) \text{ is CRC}$

by (rule CRC-intro, simp-all add: unrest closure)

lemma *unrest-ref-map-st* [unrest]: $\$ref \# P \implies \$ref \# P \oplus_r \text{map-st}_L[a]$

by (rel-auto)

lemma *unrest-ref'-map-st* [unrest]: $\$ref' \# P \implies \$ref' \# P \oplus_r \text{map-st}_L[a]$

by (rel-auto)

lemma *unrest-ref-rdes-frame-ext* [unrest]:

$\$ref \# P \implies \$ref \# a:[P]_r^+$

by (rel-blast)

lemma *unrest-ref'-rdes-frame-ext* [unrest]:

$\$ref' \# P \implies \$ref' \# a:[P]_r^+$

by (rel-blast)

lemma *map-st-ext-CRR-closed* [closure]:

assumes $P \text{ is CRR}$

shows $P \oplus_r \text{map-st}_L[a] \text{ is CRR}$

by (rule CRR-intro, simp-all add: closure unrest assms)

lemma *map-st-ext-CRC-closed* [closure]:

assumes $P \text{ is CRC}$

shows $P \oplus_r \text{map-st}_L[a] \text{ is CRC}$

by (rule CRC-intro, simp-all add: closure unrest assms)

lemma *rdes-frame-ext-CRR-closed* [closure]:

assumes $P \text{ is CRR}$

shows $a:[P]_r^+ \text{ is CRR}$

by (rule CRR-intro, simp-all add: closure unrest assms)

lemma *USUP-CRC-closed* [closure]: $\llbracket A \neq \{\}; \bigwedge i. i \in A \implies P i \text{ is CRC} \rrbracket \implies (\bigsqcup i \in A \cdot P i) \text{ is CRC}$

by (rule CRC-intro, simp-all add: unrest closure)

lemma *UINF-CRR-closed* [closure]: $\llbracket \bigwedge i. i \in A \implies P i \text{ is CRR} \rrbracket \implies (\bigsqcap i \in A \cdot P i) \text{ is CRR}$

by (rule CRR-intro, simp-all add: unrest closure)

lemma *cond-CRC-closed* [closure]:

assumes $P \text{ is CRC } Q \text{ is CRC}$

shows $P \triangleleft b \triangleright_R Q \text{ is CRC}$

by (rule CRC-intro, simp-all add: closure assms unrest)

lemma *shEx-CRR-closed* [closure]:

```

assumes  $\bigwedge x. P\ x\ is\ CRR$ 
shows  $(\exists x \cdot P(x))\ is\ CRR$ 
proof –
  have  $CRR(\exists x \cdot CRR(P(x))) = (\exists x \cdot CRR(P(x)))$ 
    by (rel-auto)
  thus ?thesis
    by (metis Healthy-def assms shEx-cong)
qed

lemma USUP-ind-CRR-closed [closure]:
  assumes  $\bigwedge i. P\ i\ is\ CRR$ 
  shows  $(\bigsqcup i \cdot P(i))\ is\ CRR$ 
  by (rule CRR-intro, simp-all add: assms unrest closure)

lemma UINF-ind-CRR-closed [closure]:
  assumes  $\bigwedge i. P\ i\ is\ CRR$ 
  shows  $(\bigsqcap i \cdot P(i))\ is\ CRR$ 
  by (rule CRR-intro, simp-all add: assms unrest closure)

lemma cond-tt-CRR-closed [closure]:
  assumes  $P\ is\ CRR\ Q\ is\ CRR$ 
  shows  $P \triangleleft \$tr' =_u \$tr \triangleright Q\ is\ CRR$ 
  by (rule CRR-intro, simp-all add: unrest assms closure)

lemma rea-implies-CRR-closed [closure]:
   $\llbracket P\ is\ CRR; Q\ is\ CRR \rrbracket \implies (P \Rightarrow_r Q)\ is\ CRR$ 
  by (simp-all add: CRR-intro closure unrest)

lemma conj-CRR-closed [closure]:
   $\llbracket P\ is\ CRR; Q\ is\ CRR \rrbracket \implies (P \wedge Q)\ is\ CRR$ 
  by (simp-all add: CRR-intro closure unrest)

lemma disj-CRR-closed [closure]:
   $\llbracket P\ is\ CRR; Q\ is\ CRR \rrbracket \implies (P \vee Q)\ is\ CRR$ 
  by (rule CRR-intro, simp-all add: unrest closure)

lemma rea-not-CRR-closed [closure]:
   $P\ is\ CRR \implies (\neg_r P)\ is\ CRR$ 
  using false-CRR rea-implies-CRR-closed by fastforce

lemma disj-R1-closed [closure]:  $\llbracket P\ is\ R1; Q\ is\ R1 \rrbracket \implies (P \vee Q)\ is\ R1$ 
  by (rel-blast)

lemma st-cond-R1-closed [closure]:  $\llbracket P\ is\ R1; Q\ is\ R1 \rrbracket \implies (P \triangleleft b \triangleright_R Q)\ is\ R1$ 
  by (rel-blast)

lemma cond-st-RR-closed [closure]:
  assumes  $P\ is\ RR\ Q\ is\ RR$ 
  shows  $(P \triangleleft b \triangleright_R Q)\ is\ RR$ 
  apply (rule RR-intro, simp-all add: unrest closure assms, simp add: Healthy-def R2c-condr)
  apply (simp add: Healthy-if assms RR-implies-R2c)
  apply (rel-auto)
done

lemma cond-st-CRR-closed [closure]:

```

$\llbracket P \text{ is CRR}; Q \text{ is CRR} \rrbracket \implies (P \triangleleft b \triangleright_R Q) \text{ is CRR}$
by (*simp-all add: CRR-intro closure unrest*)

lemma *seq-CRR-closed* [*closure*]:
assumes $P \text{ is CRR } Q \text{ is RR}$
shows $(P ;; Q) \text{ is CRR}$
by (*rule CRR-intro, simp-all add: unrest assms closure*)

lemma *wp-rea-CRC* [*closure*]: $\llbracket P \text{ is CRR}; Q \text{ is RC} \rrbracket \implies P \text{ wp}_r Q \text{ is CRC}$
by (*rule CRC-intro, simp-all add: unrest closure*)

lemma *USUP-ind-CRC-closed* [*closure*]:
 $\llbracket \bigwedge i. P i \text{ is CRC} \rrbracket \implies (\bigsqcup i. P i) \text{ is CRC}$
by (*metis CRC-implies-CRR CRC-implies-RC USUP-ind-CRR-closed USUP-ind-RC-closed false-CRC rea-not-CRR-closed wp-rea-CRC wp-rea-RC-false*)

lemma *tr-extend-seqr-lit* [*rdes*]:
fixes $P :: ('s, 'e) \text{ action}$
assumes $\$ok \nmid P \$wait \nmid P \$ref \nmid P$
shows $(\$tr' =_u \$tr \hat{\ }_u \langle \ll a \gg \rangle \wedge \$st' =_u \$st) ;; P = P[\$tr \hat{\ }_u \langle \ll a \gg \rangle / \$tr]$
using *assms* **by** (*rel-auto, meson*)

lemma *tr-assign-comp* [*rdes*]:
fixes $P :: ('s, 'e) \text{ action}$
assumes $\$ok \nmid P \$wait \nmid P \$ref \nmid P$
shows $(\$tr' =_u \$tr \wedge \lceil \langle \sigma \rangle_a \rceil_s) ;; P = \lceil \sigma \rceil_{s\sigma} \dagger P$
using *assms* **by** (*rel-auto, meson*)

lemma *RR-msubst-tt*: $RR((P t) \llbracket t \rightarrow \& tt \rrbracket) = (RR (P t)) \llbracket t \rightarrow \& tt \rrbracket$
by (*rel-auto*)

lemma *RR-msubst-ref'*: $RR((P r) \llbracket r \rightarrow \$ref' \rrbracket) = (RR (P r)) \llbracket r \rightarrow \$ref' \rrbracket$
by (*rel-auto*)

lemma *msubst-tt-RR* [*closure*]: $\llbracket \bigwedge t. P t \text{ is RR} \rrbracket \implies (P t) \llbracket t \rightarrow \& tt \rrbracket \text{ is RR}$
by (*simp add: Healthy-def RR-msubst-tt*)

lemma *msubst-ref'-RR* [*closure*]: $\llbracket \bigwedge r. P r \text{ is RR} \rrbracket \implies (P r) \llbracket r \rightarrow \$ref' \rrbracket \text{ is RR}$
by (*simp add: Healthy-def RR-msubst-ref'*)

lemma *conj-less-tr-RR-closed* [*closure*]:
assumes $P \text{ is CRR}$
shows $(P \wedge \$tr <_u \$tr') \text{ is CRR}$
proof –
have $CRR(CRR(P) \wedge \$tr <_u \$tr') = (CRR(P) \wedge \$tr <_u \$tr')$
apply (*rel-auto, blast+*)
using *less-le* **apply** *fastforce+*
done
thus *?thesis*
by (*metis Healthy-def assms*)
qed

lemma *R4-CRR-closed* [*closure*]: $P \text{ is CRR} \implies R4(P) \text{ is CRR}$
by (*simp add: R4-def conj-less-tr-RR-closed*)

lemma *R5-CRR-closed* [closure]:

assumes P is CRR

shows $R5(P)$ is CRR

proof –

have $R5(CRR(P))$ is CRR

by (rel-auto; blast)

thus ?thesis

by (simp add: assms Healthy-if)

qed

lemma *conj-eq-tr-RR-closed* [closure]:

assumes P is CRR

shows $(P \wedge \$tr' =_u \$tr)$ is CRR

proof –

have $CRR(CRR(P) \wedge \$tr' =_u \$tr) = (CRR(P) \wedge \$tr' =_u \$tr)$

by (rel-auto, blast+)

thus ?thesis

by (metis Healthy-def assms)

qed

lemma *all-ref-CRC-closed* [closure]:

P is CRC $\implies (\forall \$ref \cdot P)$ is CRC

by (simp add: CRC-implies-CRR CRR-unrest-ref all-unrest)

lemma *ex-ref-CRR-closed* [closure]:

P is CRR $\implies (\exists \$ref \cdot P)$ is CRR

by (simp add: CRR-unrest-ref ex-unrest)

lemma *ex-st'-CRR-closed* [closure]:

P is CRR $\implies (\exists \$st' \cdot P)$ is CRR

by (rule CRR-intro, simp-all add: closure unrest)

lemma *ex-ref'-CRR-closed* [closure]:

P is CRR $\implies (\exists \$ref' \cdot P)$ is CRR

using CRR-implies-RR CRR-intro CRR-unrest-ref ex-ref'-RR-closed out-in-indep unrest-ex-diff by blast

3.3 Introduction laws

Extensionality principles for introducing refinement and equality of Circus reactive relations. It is necessary only to consider a subset of the variables that are present.

lemma *CRR-refine-ext*:

assumes

P is CRR Q is CRR

$\bigwedge t s s' r'. P[\langle \rangle, \langle t \rangle, \langle s \rangle, \langle s' \rangle, \langle r' \rangle / \$tr, \$tr', \$st, \$st', \$ref'] \sqsubseteq Q[\langle \rangle, \langle t \rangle, \langle s \rangle, \langle s' \rangle, \langle r' \rangle / \$tr, \$tr', \$st, \$st', \$ref']$

shows $P \sqsubseteq Q$

proof –

have $\bigwedge t s s' r'. (CRR P)[\langle \rangle, \langle t \rangle, \langle s \rangle, \langle s' \rangle, \langle r' \rangle / \$tr, \$tr', \$st, \$st', \$ref']$

$\sqsubseteq (CRR Q)[\langle \rangle, \langle t \rangle, \langle s \rangle, \langle s' \rangle, \langle r' \rangle / \$tr, \$tr', \$st, \$st', \$ref']$

using assms by (simp add: Healthy-if)

hence $CRR P \sqsubseteq CRR Q$

by (rel-auto)

thus ?thesis

by (metis Healthy-if assms(1) assms(2))

qed

lemma *CRR-eq-ext*:

assumes

P is CRR Q is CRR

$\bigwedge t s s' r'. P[\langle \rangle, \langle t \rangle, \langle s \rangle, \langle s' \rangle, \langle r' \rangle / \$tr, \$tr', \$st, \$st', \$ref'] = Q[\langle \rangle, \langle t \rangle, \langle s \rangle, \langle s' \rangle, \langle r' \rangle / \$tr, \$tr', \$st, \$st', \$ref']$

shows $P = Q$

proof –

have $\bigwedge t s s' r'. (CRR P)[\langle \rangle, \langle t \rangle, \langle s \rangle, \langle s' \rangle, \langle r' \rangle / \$tr, \$tr', \$st, \$st', \$ref']$
 $= (CRR Q)[\langle \rangle, \langle t \rangle, \langle s \rangle, \langle s' \rangle, \langle r' \rangle / \$tr, \$tr', \$st, \$st', \$ref']$

using *assms* **by** (*simp add: Healthy-if*)

hence $CRR P = CRR Q$

by (*rel-auto*)

thus *?thesis*

by (*metis Healthy-if assms(1) assms(2)*)

qed

lemma *CRR-refine-impl-prop*:

assumes P is CRR Q is CRR

$\bigwedge t s s' r'. 'Q[\langle r' \rangle, \langle s \rangle, \langle s' \rangle, \langle \rangle, \langle t \rangle / \$ref', \$st, \$st', \$tr, \$tr']' \implies 'P[\langle r' \rangle, \langle s \rangle, \langle s' \rangle, \langle \rangle, \langle t \rangle / \$ref', \$st, \$st', \$tr, \$tr']'$

shows $P \sqsubseteq Q$

by (*rule CRR-refine-ext, simp-all add: assms closure unrest usubst*)

(*rule refine-prop-intro, simp-all add: unrest unrest-all-circus-vars closure assms*)

3.4 Weakest Precondition

lemma *nil-least [simp]*:

$\langle \rangle \leq_u x = \text{true}$ **by** *rel-auto*

lemma *minus-nil [simp]*:

$xs - \langle \rangle = xs$ **by** *rel-auto*

lemma *wp-rea-circus-lemma-1*:

assumes P is CRR $\$ref' \# P$

shows $\text{out}\alpha \# P[\langle s_0 \rangle, \langle t_0 \rangle / \$st', \$tr']$

proof –

have $\text{out}\alpha \# (CRR (\exists \$ref' \cdot P))[\langle s_0 \rangle, \langle t_0 \rangle / \$st', \$tr']$

by (*rel-auto*)

thus *?thesis*

by (*simp add: Healthy-if assms(1) assms(2) ex-unrest*)

qed

lemma *wp-rea-circus-lemma-2*:

assumes P is CRR

shows $\text{in}\alpha \# P[\langle s_0 \rangle, \langle t_0 \rangle / \$st, \$tr]$

proof –

have $\text{in}\alpha \# (CRR P)[\langle s_0 \rangle, \langle t_0 \rangle / \$st, \$tr]$

by (*rel-auto*)

thus *?thesis*

by (*simp add: Healthy-if assms ex-unrest*)

qed

The meaning of reactive weakest precondition for Circus. $P \text{ wp}_r Q$ means that, whenever P terminates in a state s_0 having done the interaction trace t_0 , which is a prefix of the overall trace, then Q must be satisfied. This in particular means that the remainder of the trace after t_0 must not be a divergent behaviour of Q .

lemma *wp-rea-circus-form*:

assumes P is CRR $\$ref' \# P$ Q is CRC

shows $(P \text{ wp}_r Q) = (\forall (s_0, t_0) \cdot \ll t_0 \gg \leq_u \$tr' \wedge P[\ll s_0 \gg, \ll t_0 \gg / \$st', \$tr']) \Rightarrow_r Q[\ll s_0 \gg, \ll t_0 \gg / \$st, \$tr])$

proof –

have $(P \text{ wp}_r Q) = (\neg_r (\exists t_0 \cdot P[\ll t_0 \gg / \$tr']) ; (\neg_r Q)[\ll t_0 \gg / \$tr'] \wedge \ll t_0 \gg \leq_u \$tr')$

by (*simp-all add: wp-rea-def R2-tr-middle closure assms*)

also have $\dots = (\neg_r (\exists (s_0, t_0) \cdot P[\ll s_0 \gg, \ll t_0 \gg / \$st', \$tr']) ; (\neg_r Q)[\ll s_0 \gg, \ll t_0 \gg / \$st, \$tr'] \wedge \ll t_0 \gg \leq_u \$tr')$

by (*rel-blast*)

also have $\dots = (\neg_r (\exists (s_0, t_0) \cdot P[\ll s_0 \gg, \ll t_0 \gg / \$st', \$tr'] \wedge (\neg_r Q)[\ll s_0 \gg, \ll t_0 \gg / \$st, \$tr'] \wedge \ll t_0 \gg \leq_u \$tr'))$

by (*simp add: seqr-to-conj add: wp-rea-circus-lemma-1 wp-rea-circus-lemma-2 assms closure conj-assoc*)

also have $\dots = (\forall (s_0, t_0) \cdot \neg_r P[\ll s_0 \gg, \ll t_0 \gg / \$st', \$tr'] \vee \neg_r (\neg_r Q)[\ll s_0 \gg, \ll t_0 \gg / \$st, \$tr'] \vee \neg_r \ll t_0 \gg \leq_u \$tr')$

by (*rel-auto*)

also have $\dots = (\forall (s_0, t_0) \cdot \neg_r P[\ll s_0 \gg, \ll t_0 \gg / \$st', \$tr'] \vee \neg_r (\neg_r RR Q)[\ll s_0 \gg, \ll t_0 \gg / \$st, \$tr'] \vee \neg_r \ll t_0 \gg \leq_u \$tr')$

by (*simp add: Healthy-if assms closure*)

also have $\dots = (\forall (s_0, t_0) \cdot \neg_r P[\ll s_0 \gg, \ll t_0 \gg / \$st', \$tr'] \vee (RR Q)[\ll s_0 \gg, \ll t_0 \gg / \$st, \$tr'] \vee \neg_r \ll t_0 \gg \leq_u \$tr')$

by (*rel-auto*)

also have $\dots = (\forall (s_0, t_0) \cdot \ll t_0 \gg \leq_u \$tr' \wedge P[\ll s_0 \gg, \ll t_0 \gg / \$st', \$tr']) \Rightarrow_r (RR Q)[\ll s_0 \gg, \ll t_0 \gg / \$st, \$tr])$

by (*rel-auto*)

also have $\dots = (\forall (s_0, t_0) \cdot \ll t_0 \gg \leq_u \$tr' \wedge P[\ll s_0 \gg, \ll t_0 \gg / \$st', \$tr']) \Rightarrow_r Q[\ll s_0 \gg, \ll t_0 \gg / \$st, \$tr])$

by (*simp add: Healthy-if assms closure*)

finally show *?thesis* .

qed

lemma *wp-rea-circus-form-alt*:

assumes P is CRR $\$ref' \# P$ Q is CRC

shows $(P \text{ wp}_r Q) = (\forall (s_0, t_0) \cdot \$tr \hat{=}^u \ll t_0 \gg \leq_u \$tr' \wedge P[\ll s_0 \gg, \langle \rangle, \ll t_0 \gg / \$st', \$tr, \$tr']) \Rightarrow_r R1(Q[\ll s_0 \gg, \langle \rangle, (\&tt - \ll t_0 \gg) / \$st, \$tr, \$tr'])$

proof –

have $(P \text{ wp}_r Q) = R2(P \text{ wp}_r Q)$

by (*simp add: CRC-implies-RR CRR-implies-RR Healthy-if RR-implies-R2 assms wp-rea-R2-closed*)

also have $\dots = R2(\forall (s_0, tr_0) \cdot \ll tr_0 \gg \leq_u \$tr' \wedge (RR P)[\ll s_0 \gg, \ll tr_0 \gg / \$st', \$tr'] \Rightarrow_r (RR Q)[\ll s_0 \gg, \ll tr_0 \gg / \$st, \$tr])$

by (*simp add: wp-rea-circus-form assms closure Healthy-if*)

also have $\dots = (\exists tt_0 \cdot (\forall (s_0, tr_0) \cdot \ll tr_0 \gg \leq_u \ll tt_0 \gg \wedge (RR P)[\ll s_0 \gg, \langle \rangle, \ll tr_0 \gg / \$st', \$tr, \$tr']) \Rightarrow_r (RR Q)[\ll s_0 \gg, \ll tr_0 \gg, \ll tt_0 \gg / \$st, \$tr, \$tr']) \wedge \$tr' =_u \$tr \hat{=}^u \ll tt_0 \gg)$

by (*simp add: R2-form, rel-auto*)

also have $\dots = (\exists tt_0 \cdot (\forall (s_0, tr_0) \cdot \ll tr_0 \gg \leq_u \ll tt_0 \gg \wedge (RR P)[\ll s_0 \gg, \langle \rangle, \ll tr_0 \gg / \$st', \$tr, \$tr']) \Rightarrow_r (RR Q)[\ll s_0 \gg, \langle \rangle, \ll tt_0 - tr_0 \gg / \$st, \$tr, \$tr']) \wedge \$tr' =_u \$tr \hat{=}^u \ll tt_0 \gg)$

by (*rel-auto*)

also have $\dots = (\exists tt_0 \cdot (\forall (s_0, tr_0) \cdot \$tr \hat{=}^u \ll tr_0 \gg \leq_u \$tr' \wedge (RR P)[\ll s_0 \gg, \langle \rangle, \ll tr_0 \gg / \$st', \$tr, \$tr']) \Rightarrow_r (RR Q)[\ll s_0 \gg, \langle \rangle, (\&tt - \ll tr_0 \gg) / \$st, \$tr, \$tr']) \wedge \$tr' =_u \$tr \hat{=}^u \ll tt_0 \gg)$

by (*rel-auto, (metis list-concat-minus-list-concat)+*)

also have $\dots = (\forall (s_0, tr_0) \cdot \$tr \hat{=}^u \ll tr_0 \gg \leq_u \$tr' \wedge (RR P)[\ll s_0 \gg, \langle \rangle, \ll tr_0 \gg / \$st', \$tr, \$tr']) \Rightarrow_r R1((RR Q)[\ll s_0 \gg, \langle \rangle, (\&tt - \ll tr_0 \gg) / \$st, \$tr, \$tr'])$

by (*rel-auto, blast+*)

also have $\dots = (\forall (s_0, t_0) \cdot \$tr \hat{=}^u \ll t_0 \gg \leq_u \$tr' \wedge P[\ll s_0 \gg, \langle \rangle, \ll t_0 \gg / \$st', \$tr, \$tr']) \Rightarrow_r R1(Q[\ll s_0 \gg, \langle \rangle, (\&tt - \ll t_0 \gg) / \$st, \$tr, \$tr'])$

by (*simp add: Healthy-if assms closure*)

finally show *?thesis* .
qed

lemma *wp-rea-circus-form-alt*:

assumes P is CRR $\$ref' \# P$ Q is CRC

shows $(P \text{ wp}_r Q) = (\forall (s_0, t_0) \cdot \$tr \hat{^}_u \ll t_0 \gg \leq_u \$tr' \wedge P[\ll s_0 \gg, \langle \rangle, \ll t_0 \gg / \$st', \$tr, \$tr'])$
 $\Rightarrow_r R1(Q[\ll s_0 \gg, \langle \rangle, (\&tt - \ll t_0 \gg) / \$st, \$tr, \$tr'])$

oops

3.5 Trace Substitution

definition *trace-subst* $(-\llbracket \cdot \rrbracket_t [999, 0] 999)$

where [*upred-defs*]: $P\llbracket v \rrbracket_t = (P[(\&tt - \lceil v \rceil_{S<}) / \&tt] \wedge \$tr + \lceil v \rceil_{S<} \leq_u \$tr')$

lemma *unrest-trace-subst* [*unrest*]:

$\llbracket \text{mwb-lens } x; x \bowtie (\$tr)_v; x \bowtie (\$tr')_v; x \bowtie (\$st)_v; x \# P \rrbracket \Longrightarrow x \# P\llbracket v \rrbracket_t$

by (*simp add: trace-subst-def lens-indep-sym unrest*)

lemma *trace-subst-RR-closed* [*closure*]:

assumes P is RR

shows $P\llbracket v \rrbracket_t$ is RR

proof –

have $(RR P)\llbracket v \rrbracket_t$ is RR

apply (*rel-auto*)

apply (*metis diff-add-cancel-left' trace-class.add-left-mono*)

apply (*metis le-add minus-cancel-le trace-class.add-diff-cancel-left*)

using *le-add order-trans* apply *blast*

done

thus *?thesis*

by (*simp add: Healthy-if assms*)

qed

lemma *trace-subst-CRR-closed* [*closure*]:

assumes P is CRR

shows $P\llbracket v \rrbracket_t$ is CRR

by (*rule CRR-intro, simp-all add: closure assms unrest*)

lemma *tsubst-nil* [*usubst*]:

assumes P is CRR

shows $P[\langle \rangle]_t = P$

proof –

have $(CRR P)[\langle \rangle]_t = CRR P$

by (*rel-auto*)

thus *?thesis*

by (*simp add: Healthy-if assms*)

qed

lemma *tsubst-false* [*usubst*]: $false\llbracket y \rrbracket_t = false$

by *rel-auto*

lemma *cond-rea-tt-subst* [*usubst*]:

$(P \triangleleft b \triangleright_R Q)\llbracket v \rrbracket_t = (P\llbracket v \rrbracket_t \triangleleft b \triangleright_R Q\llbracket v \rrbracket_t)$

by (*rel-auto*)

lemma *tsubst-conj* [*usubst*]: $(P \wedge Q)\llbracket v \rrbracket_t = (P\llbracket v \rrbracket_t \wedge Q\llbracket v \rrbracket_t)$

by (*rel-auto*)

lemma *tsubst-disj* [*usubst*]: $(P \vee Q)[v]_t = (P[v]_t \vee Q[v]_t)$
by (*rel-auto*)

lemma *rea-subst-R1-closed* [*closure*]: $P[v]_t$ is *R1*
apply (*rel-auto*) **using** *le-add order.trans* **by** *blast*

lemma *tsubst-UINF-ind* [*usubst*]: $(\prod i \cdot P(i))[v]_t = (\prod i \cdot (P(i))[v]_t)$
by (*rel-auto*)

3.6 Initial Interaction

definition *rea-init* :: $'s \text{ upred} \Rightarrow ('t::\text{trace}, 's) \text{ uexpr} \Rightarrow ('s, 't, 'a, 'b) \text{ rel-rsp } (\mathcal{I}'(-, -))$ **where**
 $[upred-defs]: \mathcal{I}(s, t) = (\lceil s \rceil_{S<} \wedge \$tr + \lceil t \rceil_{S<} \leq_u \$tr')$

$\mathcal{I}(s, t)$ is a predicate stating that, if the initial state satisfies state predicate s , then the trace t is an initial trace.

lemma *unrest-rea-init* [*unrest*]:
 $\llbracket x \bowtie (\$tr)_v; x \bowtie (\$tr')_v; x \bowtie (\$st)_v \rrbracket \Longrightarrow x \# \mathcal{I}(s, t)$
by (*simp add: rea-init-def unrest lens-indep-sym*)

lemma *rea-init-R1* [*closure*]: $\mathcal{I}(s, t)$ is *R1*
apply (*rel-auto*) **using** *dual-order.trans le-add* **by** *blast*

lemma *rea-init-R2c* [*closure*]: $\mathcal{I}(s, t)$ is *R2c*
apply (*rel-auto*)
apply (*metis diff-add-cancel-left' trace-class.add-left-mono*)
apply (*metis le-add minus-cancel-le trace-class.add-diff-cancel-left*)
done

lemma *rea-init-R2* [*closure*]: $\mathcal{I}(s, t)$ is *R2*
by (*metis Healthy-def R1-R2c-is-R2 rea-init-R1 rea-init-R2c*)

lemma *csp-init-RR* [*closure*]: $\mathcal{I}(s, t)$ is *RR*
apply (*rel-auto*)
apply (*metis diff-add-cancel-left' trace-class.add-left-mono*)
apply (*metis le-add minus-cancel-le trace-class.add-diff-cancel-left*)
apply (*metis le-add less-le less-le-trans*)
done

lemma *csp-init-CRR* [*closure*]: $\mathcal{I}(s, t)$ is *CRR*
by (*rule CRR-intro, simp-all add: unrest closure*)

lemma *rea-init-impl-st* [*closure*]: $(\mathcal{I}(b, t) \Rightarrow_r [c]_{S<})$ is *RC*
apply (*rule RC-intro*)
apply (*simp add: closure*)
apply (*rel-auto*)
using *order-trans* **by** *auto*

lemma *rea-init-RC1*:
 $\neg_r \mathcal{I}(P, t)$ is *RC1*
apply (*rel-auto*) **using** *dual-order.trans* **by** *blast*

lemma *init-acts-empty* [*rpred*]: $\mathcal{I}(\text{true}, \langle \rangle) = \text{true}_r$
by (*rel-auto*)

lemma *rea-not-init* [*rpred*]:

$(\neg_r \mathcal{I}(P, \langle \rangle)) = \mathcal{I}(\neg P, \langle \rangle)$
by (*rel-auto*)

lemma *rea-init-conj* [*rpred*]:

$(\mathcal{I}(P, t) \wedge \mathcal{I}(Q, t)) = \mathcal{I}(P \wedge Q, t)$
by (*rel-auto*)

lemma *rea-init-empty-trace* [*rpred*]: $\mathcal{I}(s, \langle \rangle) = [s]_{s <}$

by (*rel-auto*)

lemma *rea-init-disj-same* [*rpred*]: $(\mathcal{I}(s_1, t) \vee \mathcal{I}(s_2, t)) = \mathcal{I}(s_1 \vee s_2, t)$

by (*rel-auto*)

lemma *rea-init-impl-same* [*rpred*]: $(\mathcal{I}(s_1, t) \Rightarrow_r \mathcal{I}(s_2, t)) = (\mathcal{I}(s_1, t) \Rightarrow_r [s_2]_{s <})$

apply (*rel-auto*) **using** *dual-order.trans le-add* **by** *blast+*

lemma *tsubst-st-cond* [*usubst*]: $[P]_{s <} \llbracket t \rrbracket_t = \mathcal{I}(P, t)$

by (*rel-auto*)

lemma *tsubst-rea-init* [*usubst*]: $(\mathcal{I}(s, x)) \llbracket y \rrbracket_t = \mathcal{I}(s, y + x)$

apply (*rel-auto*)

apply (*metis add.assoc diff-add-cancel-left' trace-class.add-le-imp-le-left trace-class.add-left-mono*)

apply (*metis add.assoc diff-add-cancel-left' le-add trace-class.add-le-imp-le-left trace-class.add-left-mono*)

done

lemma *tsubst-rea-not* [*usubst*]: $(\neg_r P) \llbracket v \rrbracket_t = ((\neg_r P \llbracket v \rrbracket_t) \wedge \mathcal{I}(\text{true}, v))$

apply (*rel-auto*)

using *le-add order-trans* **by** *blast*

lemma *tsubst-true* [*usubst*]: $\text{true}_r \llbracket v \rrbracket_t = \mathcal{I}(\text{true}, v)$

by (*rel-auto*)

lemma *R4-csp-init* [*rpred*]: $R4(\mathcal{I}(s, \text{bop Cons } x \text{ } xs)) = \mathcal{I}(s, \text{bop Cons } x \text{ } xs)$

using *less-list-def* **by** (*rel-blast*)

lemma *R5-csp-init* [*rpred*]: $R5(\mathcal{I}(s, \text{bop Cons } x \text{ } xs)) = \text{false}$

by (*rel-auto*)

lemma *R4-trace-subst* [*rpred*]:

$R4(P \llbracket \text{bop Cons } x \text{ } xs \rrbracket_t) = P \llbracket \text{bop Cons } x \text{ } xs \rrbracket_t$

using *le-imp-less-or-eq* **by** (*rel-blast*)

lemma *R5-trace-subst* [*rpred*]:

$R5(P \llbracket \text{bop Cons } x \text{ } xs \rrbracket_t) = \text{false}$

by (*rel-auto*)

3.7 Enabled Events

definition *csp-enable* :: $'s \text{ upred} \Rightarrow ('e \text{ list}, 's) \text{ uexpr} \Rightarrow ('e \text{ set}, 's) \text{ uexpr} \Rightarrow ('s, 'e) \text{ action } (\mathcal{E}'(-, -, -))$

where

$[upred-defs]: \mathcal{E}(s, t, E) = ([s]_{s <} \wedge \text{\$tr}' =_u \text{\$tr} \hat{_}_u [t]_{s <} \wedge (\forall e \in [E]_{s <} \cdot \llbracket e \rrbracket \notin_u \text{\$ref}'))$

Predicate $\mathcal{E}(s, t, E)$ states that, if the initial state satisfies predicate s , then t is a possible (failure) trace, such that the events in the set E are enabled after the given interaction.

lemma *csp-enable-R1-closed* [closure]: $\mathcal{E}(s, t, E)$ is R1
 by (rel-auto)

lemma *csp-enable-R2-closed* [closure]: $\mathcal{E}(s, t, E)$ is R2c
 by (rel-auto)

lemma *csp-enable-RR* [closure]: $\mathcal{E}(s, t, E)$ is CRR
 by (rel-auto)

lemma *tsubst-csp-enable* [usubst]: $\mathcal{E}(s, t_2, e) \llbracket t_1 \rrbracket_t = \mathcal{E}(s, t_1 \hat{=} t_2, e)$
 apply (rel-auto)
 apply (metis append.assoc less-eq-list-def prefix-concat-minus)
 apply (simp add: list-concat-minus-list-concat)
 done

lemma *csp-enable-unrests* [unrest]:
 $\llbracket x \bowtie (\$tr)_v; x \bowtie (\$tr')_v; x \bowtie (\$st)_v; x \bowtie (\$ref')_v \rrbracket \implies x \# \mathcal{E}(s, t, e)$
 by (simp add: csp-enable-def R1-def lens-indep-sym unrest)

lemma *st-unrest-csp-enable* [unrest]: $\llbracket \&\mathbf{v} \# s; \&\mathbf{v} \# t; \&\mathbf{v} \# E \rrbracket \implies \$st \# \mathcal{E}(s, t, E)$
 by (simp add: csp-enable-def unrest)

lemma *csp-enable-tr'-eq-tr* [rpred]:
 $\mathcal{E}(s, \langle \rangle, r) \triangleleft \$tr' =_u \$tr \triangleright false = \mathcal{E}(s, \langle \rangle, r)$
 by (rel-auto)

lemma *csp-enable-st-pred* [rpred]:
 $([s_1]_{S<} \wedge \mathcal{E}(s_2, t, E)) = \mathcal{E}(s_1 \wedge s_2, t, E)$
 by (rel-auto)

lemma *csp-enable-conj* [rpred]:
 $(\mathcal{E}(s, t, E_1) \wedge \mathcal{E}(s, t, E_2)) = \mathcal{E}(s, t, E_1 \cup_u E_2)$
 by (rel-auto)

lemma *csp-enable-cond* [rpred]:
 $\mathcal{E}(s_1, t_1, E_1) \triangleleft b \triangleright_R \mathcal{E}(s_2, t_2, E_2) = \mathcal{E}(s_1 \triangleleft b \triangleright s_2, t_1 \triangleleft b \triangleright t_2, E_1 \triangleleft b \triangleright E_2)$
 by (rel-auto)

lemma *csp-enable-rea-assm* [rpred]:
 $[b]^\top_r ;; \mathcal{E}(s, t, E) = \mathcal{E}(b \wedge s, t, E)$
 by (rel-auto)

lemma *csp-enable-tr-empty*: $\mathcal{E}(true, \langle \rangle, \{v\}_u) = (\$tr' =_u \$tr \wedge [v]_{S<} \notin_u \$ref')$
 by (rel-auto)

lemma *csp-enable-nothing*: $\mathcal{E}(true, \langle \rangle, \{\}_u) = (\$tr' =_u \$tr)$
 by (rel-auto)

lemma *msubst-nil-csp-enable* [usubst]:
 $\mathcal{E}(s(x), t(x), E(x)) \llbracket x \rightarrow \langle \rangle \rrbracket = \mathcal{E}(s(x) \llbracket x \rightarrow \langle \rangle \rrbracket, t(x) \llbracket x \rightarrow \langle \rangle \rrbracket, E(x) \llbracket x \rightarrow \langle \rangle \rrbracket)$
 by (pred-auto)

lemma *msubst-csp-enable* [usubst]:
 $\mathcal{E}(s(x), t(x), E(x)) \llbracket x \rightarrow [v]_{S\leftarrow} \rrbracket = \mathcal{E}(s(x) \llbracket x \rightarrow v \rrbracket, t(x) \llbracket x \rightarrow v \rrbracket, E(x) \llbracket x \rightarrow v \rrbracket)$
 by (rel-auto)

lemma *csp-enable-false* [*rpred*]: $\mathcal{E}(\text{false}, t, E) = \text{false}$
by (*rel-auto*)

lemma *conj-csp-enable* [*rpred*]: $(\mathcal{E}(b_1, t, E_1) \wedge \mathcal{E}(b_2, t, E_2)) = \mathcal{E}(b_1 \wedge b_2, t, E_1 \cup_u E_2)$
by (*rel-auto*)

lemma *USUP-csp-enable* [*rpred*]:
 $(\bigsqcup x \cdot \mathcal{E}(s, t, A(x))) = \mathcal{E}(s, t, (\bigvee x \cdot A(x)))$
by (*rel-auto*)

lemma *R4-csp-enable-nil* [*rpred*]:
 $R4(\mathcal{E}(s, \langle \rangle, E)) = \text{false}$
by (*rel-auto*)

lemma *R5-csp-enable-nil* [*rpred*]:
 $R5(\mathcal{E}(s, \langle \rangle, E)) = \mathcal{E}(s, \langle \rangle, E)$
by (*rel-auto*)

lemma *R4-csp-enable-Cons* [*rpred*]:
 $R4(\mathcal{E}(s, \text{bop Cons } x \text{ } xs, E)) = \mathcal{E}(s, \text{bop Cons } x \text{ } xs, E)$
by (*rel-auto*, *simp add: Prefix-Order.strict-prefixI'*)

lemma *R5-csp-enable-Cons* [*rpred*]:
 $R5(\mathcal{E}(s, \text{bop Cons } x \text{ } xs, E)) = \text{false}$
by (*rel-auto*)

lemma *rel-aext-csp-enable* [*alpha*]:
 $\text{vwb-lens } a \implies \mathcal{E}(s, t, E) \oplus_r \text{map-st}_L[a] = \mathcal{E}(s \oplus_p a, t \oplus_p a, E \oplus_p a)$
by (*rel-auto*)

3.8 Completed Trace Interaction

definition *csp-do* :: $'s \text{ upred} \Rightarrow ('s \Rightarrow 's) \Rightarrow ('e \text{ list}, 's) \text{ uexpr} \Rightarrow ('s, 'e) \text{ action } (\Phi'(-, -, -))$ **where**
 $[\text{upred-defs}]: \Phi(s, \sigma, t) = ([s]_{S<} \wedge \text{\$tr}' =_u \text{\$tr} \hat{_}_u [t]_{S<} \wedge [\langle \sigma \rangle_a]_S)$

Predicate $\Phi(s, \sigma, t)$ states that if the initial state satisfies s , and the trace t is performed, then afterwards the state update σ is executed.

lemma *unrest-csp-do* [*unrest*]:
 $\llbracket x \bowtie (\text{\$tr})_v; x \bowtie (\text{\$tr}')_v; x \bowtie (\text{\$st})_v; x \bowtie (\text{\$st}')_v \rrbracket \implies x \sharp \Phi(s, \sigma, t)$
by (*simp-all add: csp-do-def alpha-in-var alpha-out-var prod-as-plus unrest lens-indep-sym*)

lemma *csp-do-CRR* [*closure*]: $\Phi(s, \sigma, t)$ *is CRR*
by (*rel-auto*)

lemma *csp-do-R4-closed* [*closure*]:
 $\Phi(b, \sigma, \text{bop Cons } x \text{ } xs)$ *is R4*
by (*rel-auto*, *simp add: Prefix-Order.strict-prefixI'*)

lemma *st-pred-conj-csp-do* [*rpred*]:
 $([b]_{S<} \wedge \Phi(s, \sigma, t)) = \Phi(b \wedge s, \sigma, t)$
by (*rel-auto*)

lemma *trea-subst-csp-do* [*usubst*]:
 $(\Phi(s, \sigma, t_2)) \llbracket t_1 \rrbracket_t = \Phi(s, \sigma, t_1 \hat{_}_u t_2)$

apply (*rel-auto*)
apply (*metis append.assoc less-eq-list-def prefix-concat-minus*)
apply (*simp add: list-concat-minus-list-concat*)
done

lemma *st-subst-csp-do* [*usubst*]:
 $\lceil \sigma \rceil_{S\sigma} \uparrow \Phi(s, \varrho, t) = \Phi(\sigma \uparrow s, \varrho \circ \sigma, \sigma \uparrow t)$
by (*rel-auto*)

lemma *csp-init-do* [*rpred*]: $(\mathcal{I}(s1, t) \wedge \Phi(s2, \sigma, t)) = \Phi(s1 \wedge s2, \sigma, t)$
by (*rel-auto*)

lemma *csp-do-false* [*rpred*]: $\Phi(\text{false}, s, t) = \text{false}$
by (*rel-auto*)

lemma *csp-do-assign* [*rpred*]:
assumes *P is CRR*
shows $\Phi(s, \sigma, t) ;; P = ([s]_{S<} \wedge (\lceil \sigma \rceil_{S\sigma} \uparrow P)) \llbracket t \rrbracket_t$
proof –
have $\Phi(s, \sigma, t) ;; CRR(P) = ([s]_{S<} \wedge (\lceil \sigma \rceil_{S\sigma} \uparrow CRR(P))) \llbracket t \rrbracket_t$
by (*rel-blast*)
thus *?thesis*
by (*simp add: Healthy-if assms*)
qed

lemma *subst-state-csp-enable* [*usubst*]:
 $\lceil \sigma \rceil_{S\sigma} \uparrow \mathcal{E}(s, t_2, e) = \mathcal{E}(\sigma \uparrow s, \sigma \uparrow t_2, \sigma \uparrow e)$
by (*rel-auto*)

lemma *csp-do-assign-enable* [*rpred*]:
 $\Phi(s_1, \sigma, t_1) ;; \mathcal{E}(s_2, t_2, e) = \mathcal{E}(s_1 \wedge \sigma \uparrow s_2, t_1 \hat{^}_u(\sigma \uparrow t_2), (\sigma \uparrow e))$
by (*simp add: rpred closure usubst*)

lemma *csp-do-assign-do* [*rpred*]:
 $\Phi(s_1, \sigma, t_1) ;; \Phi(s_2, \varrho, t_2) = \Phi(s_1 \wedge (\sigma \uparrow s_2), \varrho \circ \sigma, t_1 \hat{^}_u(\sigma \uparrow t_2))$
by (*rel-auto*)

lemma *csp-do-cond* [*rpred*]:
 $\Phi(s_1, \sigma, t_1) \triangleleft b \triangleright_R \Phi(s_2, \varrho, t_2) = \Phi(s_1 \triangleleft b \triangleright s_2, \sigma \triangleleft b \triangleright_s \varrho, t_1 \triangleleft b \triangleright t_2)$
by (*rel-auto*)

lemma *rea-assm-csp-do* [*rpred*]:
 $[b]^\top_r ;; \Phi(s, \sigma, t) = \Phi(b \wedge s, \sigma, t)$
by (*rel-auto*)

lemma *csp-do-skip* [*rpred*]:
assumes *P is CRR*
shows $\Phi(\text{true}, id, t) ;; P = P \llbracket t \rrbracket_t$
proof –
have $\Phi(\text{true}, id, t) ;; CRR(P) = (CRR P) \llbracket t \rrbracket_t$
by (*rel-auto*)
thus *?thesis*
by (*simp add: Healthy-if assms*)
qed

lemma *wp-rea-csp-do-lemma*:

fixes $P :: ('σ, 'φ) \text{ action}$
assumes $\$ok \# P \$wait \# P \$ref \# P$
shows $(\llbracket \langle \sigma \rangle_a \rrbracket_s \wedge \$tr' =_u \$tr \hat{^}_u \llbracket t \rrbracket_{s<} \rrbracket ; P = (\llbracket \sigma \rrbracket_{s\sigma} \dagger P) \llbracket \$tr \hat{^}_u \llbracket t \rrbracket_{s<} / \$tr \rrbracket$
using *assms* **by** (*rel-auto*, *meson*)

lemma *wp-rea-csp-do* [*wp*]:

fixes $P :: ('σ, 'φ) \text{ action}$
assumes $P \text{ is } CRR$
shows $\Phi(s, \sigma, t) \text{ wp}_r P = (\mathcal{I}(s, t) \Rightarrow_r (\llbracket \sigma \rrbracket_{s\sigma} \dagger P) \llbracket t \rrbracket_t)$

proof –

have $\Phi(s, \sigma, t) \text{ wp}_r CRR(P) = (\mathcal{I}(s, t) \Rightarrow_r (\llbracket \sigma \rrbracket_{s\sigma} \dagger CRR(P)) \llbracket t \rrbracket_t)$
by (*rel-blast*)
thus *?thesis*
by (*simp add: assms Healthy-if*)

qed

lemma *csp-do-power-Suc* [*rpred*]:

$\Phi(\text{true}, \text{id}, t) \hat{^} (\text{Suc } i) = \Phi(\text{true}, \text{id}, \text{iter}[\text{Suc } i](t))$
by (*induct i*, (*rel-auto*)+)

lemma *csp-power-do-comp* [*rpred*]:

assumes $P \text{ is } CRR$
shows $\Phi(\text{true}, \text{id}, t) \hat{^} i ; P = \Phi(\text{true}, \text{id}, \text{iter}[i](t)) ; P$
apply (*cases i*)
apply (*simp-all add: rpred usubst assms closure*)
done

lemma *wp-rea-csp-do-skip* [*wp*]:

fixes $Q :: ('σ, 'φ) \text{ action}$
assumes $P \text{ is } CRR$
shows $\Phi(s, \text{id}, t) \text{ wp}_r P = (\mathcal{I}(s, t) \Rightarrow_r P \llbracket t \rrbracket_t)$

proof –

have $\Phi(s, \text{id}, t) \text{ wp}_r P = \Phi(s, \text{id}, t) \text{ wp}_r P$
by (*simp add: skip-r-def*)
thus *?thesis* **by** (*simp add: wp assms usubst alpha*)

qed

lemma *msubst-csp-do* [*usubst*]:

$\Phi(s(x), \sigma, t(x)) \llbracket x \rightarrow \llbracket v \rrbracket_{s\leftarrow} \rrbracket = \Phi(s(x) \llbracket x \rightarrow v \rrbracket, \sigma, t(x) \llbracket x \rightarrow v \rrbracket)$
by (*rel-auto*)

lemma *rea-frame-ext-csp-do* [*frame*]:

$vwb\text{-}lens \ a \Longrightarrow a : [\Phi(s, \sigma, t)]_r^+ = \Phi(s \oplus_p a, \sigma \oplus_s a, t \oplus_p a)$
by (*rel-auto*)

3.9 Downward closure of refusals

We define downward closure of the pericondition by the following healthiness condition

definition *CDC* :: $('s, 'e) \text{ action} \Rightarrow ('s, 'e) \text{ action}$ **where**

[*upred-defs*]: $CDC(P) = (\exists \text{ ref}_0 \cdot P \llbracket \llbracket \text{ref}_0 \rrbracket / \$ref' \rrbracket \wedge \$ref' \subseteq_u \llbracket \text{ref}_0 \rrbracket \gg)$

lemma *CDC-idem*: $CDC(CDC(P)) = CDC(P)$

by (*rel-auto*)

lemma *CDC-RR-commute*: $CDC(RR(P)) = RR(CDC(P))$
by (*rel-blast*)

lemma *CDC-RR-closed* [*closure*]: P is $RR \implies CDC(P)$ is RR
by (*metis CDC-RR-commute Healthy-def*)

lemma *CDC-CRR-commute*: $CDC(CRR P) = CRR(CDC P)$
by (*rel-blast*)

lemma *CDC-CRR-closed* [*closure*]:
assumes P is CRR
shows $CDC(P)$ is CRR
by (*rule CRR-intro, simp add: CDC-def unrest assms closure, simp add: unrest assms closure*)

lemma *CDC-unrest* [*unrest*]: $\llbracket vwb\text{-}lens\ x; (\$ref')_v \bowtie x; x \# P \rrbracket \implies x \# CDC(P)$
by (*simp add: CDC-def unrest usubst lens-indep-sym*)

lemma *CDC-R4-commute*: $CDC(R4(P)) = R4(CDC(P))$
by (*rel-auto*)

lemma *R4-CDC-closed* [*closure*]: P is $CDC \implies R4(P)$ is CDC
by (*simp add: CDC-R4-commute Healthy-def*)

lemma *CDC-R5-commute*: $CDC(R5(P)) = R5(CDC(P))$
by (*rel-auto*)

lemma *R5-CDC-closed* [*closure*]: P is $CDC \implies R5(P)$ is CDC
by (*simp add: CDC-R5-commute Healthy-def*)

lemma *rea-true-CDC* [*closure*]: $true_r$ is CDC
by (*rel-auto*)

lemma *false-CDC* [*closure*]: $false$ is CDC
by (*rel-auto*)

lemma *CDC-UINF-closed* [*closure*]:
assumes $\bigwedge i. i \in I \implies P\ i$ is CDC
shows $(\bigsqcap i \in I \cdot P\ i)$ is CDC
using *assms* **by** (*rel-blast*)

lemma *CDC-disj-closed* [*closure*]:
assumes P is CDC Q is CDC
shows $(P \vee Q)$ is CDC
proof –
have $CDC(P \vee Q) = (CDC(P) \vee CDC(Q))$
by (*rel-auto*)
thus *?thesis*
by (*metis Healthy-def assms(1) assms(2)*)
qed

lemma *CDC-USUP-closed* [*closure*]:
assumes $\bigwedge i. i \in I \implies P\ i$ is CDC
shows $(\bigsqcup i \in I \cdot P\ i)$ is CDC
using *assms* **by** (*rel-blast*)

lemma *CDC-conj-closed* [closure]:
assumes P is CDC Q is CDC
shows $(P \wedge Q)$ is CDC
using *assms* **by** (*rel-auto*, *blast*, *meson*)

lemma *CDC-rea-impl* [*rpred*]:
 $\$ref' \# P \implies CDC(P \Rightarrow_r Q) = (P \Rightarrow_r CDC(Q))$
by (*rel-auto*)

lemma *rea-impl-CDC-closed* [closure]:
assumes $\$ref' \# P$ Q is CDC
shows $(P \Rightarrow_r Q)$ is CDC
using *assms* **by** (*simp add: CDC-rea-impl Healthy-def*)

lemma *seq-CDC-closed* [closure]:
assumes Q is CDC
shows $(P ;; Q)$ is CDC
proof –
have $CDC(P ;; Q) = P ;; CDC(Q)$
by (*rel-blast*)
thus ?thesis
by (*metis Healthy-def assms*)
qed

lemma *st-subst-CDC-closed* [closure]:
assumes P is CDC
shows $(\sigma \uparrow_S P)$ is CDC
proof –
have $(\sigma \uparrow_S CDC P)$ is CDC
by (*rel-auto*)
thus ?thesis
by (*simp add: assms Healthy-if*)
qed

lemma *rea-st-cond-CDC* [closure]: $[g]_{S<}$ is CDC
by (*rel-auto*)

lemma *csp-enable-CDC* [closure]: $\mathcal{E}(s, t, E)$ is CDC
by (*rel-auto*)

lemma *state-srea-CDC-closed* [closure]:
assumes P is CDC
shows *state* $'a \cdot P$ is CDC
proof –
have *state* $'a \cdot CDC(P)$ is CDC
by (*rel-blast*)
thus ?thesis
by (*simp add: Healthy-if assms*)
qed

3.10 Renaming

abbreviation *pre-image* $f B \equiv \{x. f(x) \in B\}$

definition *csp-rename* :: (s, e) action $\Rightarrow (e \Rightarrow f) \Rightarrow (s, f)$ action $((-)\llbracket - \rrbracket_c [999, 0] 999)$ **where**
 $[upred-defs]: P \llbracket f \rrbracket_c = R2((\$tr' =_u \langle \rangle \wedge \$st' =_u \$st) ;; P ;; (\$tr' =_u \text{map}_u \llbracket f \rrbracket \$tr \wedge \$st' =_u \$st \wedge$

$uop \ (pre\text{-}image \ f) \ \$ref' \subseteq_u \ \$ref))$

lemma *csp-rename-CRR-closed* [closure]:

assumes P is CRR
shows $P \llbracket f \rrbracket_c$ is CRR

proof –

have $(CRR \ P) \llbracket f \rrbracket_c$ is CRR
by (rel-auto)
thus ?thesis by (simp add: assms Healthy-if)

qed

lemma *csp-rename-disj* [rpred]: $(P \vee Q) \llbracket f \rrbracket_c = (P \llbracket f \rrbracket_c \vee Q \llbracket f \rrbracket_c)$

by (rel-blast)

lemma *csp-rename-UINF-ind* [rpred]: $(\bigcap i \cdot P \ i) \llbracket f \rrbracket_c = (\bigcap i \cdot (P \ i) \llbracket f \rrbracket_c)$

by (rel-blast)

lemma *csp-rename-UINF-mem* [rpred]: $(\bigcap i \in A \cdot P \ i) \llbracket f \rrbracket_c = (\bigcap i \in A \cdot (P \ i) \llbracket f \rrbracket_c)$

by (rel-blast)

Renaming distributes through conjunction only when both sides are downward closed

lemma *csp-rename-conj* [rpred]:

assumes $inj \ f$ P is CRR Q is CRR P is CDC Q is CDC
shows $(P \wedge Q) \llbracket f \rrbracket_c = (P \llbracket f \rrbracket_c \wedge Q \llbracket f \rrbracket_c)$

proof –

from assms(1) have $((CDC \ (CRR \ P)) \wedge (CDC \ (CRR \ Q))) \llbracket f \rrbracket_c = ((CDC \ (CRR \ P)) \llbracket f \rrbracket_c \wedge (CDC \ (CRR \ Q)) \llbracket f \rrbracket_c)$

apply (rel-auto)

apply blast

apply blast

apply (meson order-refl order-trans)

done

thus ?thesis

by (simp add: assms Healthy-if)

qed

lemma *csp-rename-seq* [rpred]:

assumes P is CRR Q is CRR

shows $(P ;; Q) \llbracket f \rrbracket_c = P \llbracket f \rrbracket_c ;; Q \llbracket f \rrbracket_c$

oops

lemma *csp-rename-R4* [rpred]:

$(R4(P)) \llbracket f \rrbracket_c = R4(P \llbracket f \rrbracket_c)$

apply (rel-auto, blast)

using less-le apply fastforce

apply (metis (mono-tags, lifting) Prefix-Order.Nil-prefix append-Nil2 diff-add-cancel-left' less-le list.simps(8) plus-list-def)

done

lemma *csp-rename-R5* [rpred]:

$(R5(P)) \llbracket f \rrbracket_c = R5(P \llbracket f \rrbracket_c)$

apply (rel-auto, blast)

using less-le apply fastforce

done

lemma *csp-rename-do* [*rpred*]: $\Phi(s, \sigma, t) \llbracket f \rrbracket_c = \Phi(s, \sigma, \text{map}_u \llbracket f \rrbracket t)$
by (*rel-auto*)

lemma *csp-rename-enable* [*rpred*]: $\mathcal{E}(s, t, E) \llbracket f \rrbracket_c = \mathcal{E}(s, \text{map}_u \llbracket f \rrbracket t, \text{uop}(\text{image } f) E)$
by (*rel-auto*)

lemma *st'-unrest-csp-rename* [*unrest*]: $\$st' \# P \implies \$st' \# P \llbracket f \rrbracket_c$
by (*rel-blast*)

lemma *ref'-unrest-csp-rename* [*unrest*]: $\$ref' \# P \implies \$ref' \# P \llbracket f \rrbracket_c$
by (*rel-blast*)

lemma *csp-rename-CDC-closed* [*closure*]:
 $P \text{ is CDC} \implies P \llbracket f \rrbracket_c \text{ is CDC}$
by (*rel-blast*)

lemma *csp-do-CDC* [*closure*]: $\Phi(s, \sigma, t) \text{ is CDC}$
by (*rel-auto*)

end

4 Stateful-Failure Healthiness Conditions

theory *utp-sfrd-healths*
imports *utp-sfrd-rel*
begin

5 Definitions

We here define extra healthiness conditions for stateful-failure reactive designs.

abbreviation *CSP1* :: $((\sigma', \varphi) \text{ st-csp} \times (\sigma', \varphi) \text{ st-csp}) \text{ health}$
where $CSP1(P) \equiv RD1(P)$

abbreviation *CSP2* :: $((\sigma', \varphi) \text{ st-csp} \times (\sigma', \varphi) \text{ st-csp}) \text{ health}$
where $CSP2(P) \equiv RD2(P)$

abbreviation *CSP* :: $((\sigma', \varphi) \text{ st-csp} \times (\sigma', \varphi) \text{ st-csp}) \text{ health}$
where $CSP(P) \equiv SRD(P)$

definition *STOP* :: $\varphi \text{ process}$ **where**
[*upred-defs*]: $STOP = CSP1(\$ok' \wedge R3c(\$tr' =_u \$tr \wedge \$wait'))$

definition *SKIP* :: $\varphi \text{ process}$ **where**
[*upred-defs*]: $SKIP = \mathbf{R}_s(\exists \$ref \cdot CSP1(II))$

definition *Stop* :: $(\sigma', \varphi) \text{ action}$ **where**
[*upred-defs*]: $Stop = \mathbf{R}_s(true \vdash (\$tr' =_u \$tr \wedge \$wait'))$

definition *Skip* :: $(\sigma', \varphi) \text{ action}$ **where**
[*upred-defs*]: $Skip = \mathbf{R}_s(true \vdash (\$tr' =_u \$tr \wedge \neg \$wait' \wedge \$st' =_u \$st))$

definition *CSP3* :: $((\sigma', \varphi) \text{ st-csp} \times (\sigma', \varphi) \text{ st-csp}) \text{ health}$ **where**
[*upred-defs*]: $CSP3(P) = (Skip ;; P)$

definition $CSP_4 :: ((\sigma, \varphi) \text{ st-csp} \times (\sigma, \varphi) \text{ st-csp}) \text{ health}$ **where**
 $[upred-defs]: CSP_4(P) = (P ;; Skip)$

definition $NCSP :: ((\sigma, \varphi) \text{ st-csp} \times (\sigma, \varphi) \text{ st-csp}) \text{ health}$ **where**
 $[upred-defs]: NCSP = CSP_3 \circ CSP_4 \circ CSP$

Productive and normal processes

abbreviation $PCSP \equiv Productive \circ NCSP$

Instantaneous and normal processes

abbreviation $ICSP \equiv ISRD1 \circ NCSP$

5.1 Healthiness condition properties

$SKIP$ is the same as $Skip$, and $STOP$ is the same as $Stop$, when we consider stateless CSP processes. This is because any reference to the st variable degenerates when the alphabet type coerces its type to be empty. We therefore need not consider $SKIP$ and $STOP$ actions.

theorem $SKIP\text{-is-Skip}$ $[simp]: SKIP = Skip$
by $(rel\text{-auto})$

theorem $STOP\text{-is-Stop}$ $[simp]: STOP = Stop$
by $(rel\text{-auto})$

theorem $Skip\text{-UTP-form}$: $Skip = \mathbf{R}_s(\exists \$ref \cdot CSP1(II))$
by $(rel\text{-auto})$

lemma $Skip\text{-is-CSP}$ $[closure]$:
 $Skip \text{ is CSP}$
by $(simp \text{ add: } Skip\text{-def RHS-design-is-SRD unrest})$

lemma $Skip\text{-RHS-tri-design}$:
 $Skip = \mathbf{R}_s(true \vdash (false \diamond (\$tr' =_u \$tr \wedge \$st' =_u \$st)))$
by $(rel\text{-auto})$

lemma $Skip\text{-RHS-tri-design}'$ $[rdes-def]$:
 $Skip = \mathbf{R}_s(true_r \vdash (false \diamond \Phi(true, id, \langle \rangle)))$
by $(rel\text{-auto})$

lemma $Skip\text{-frame}$ $[frame]$: $vwb\text{-lens } a \implies a:[Skip]_R^+ = Skip$
by $(rdes\text{-eq})$

lemma $Stop\text{-is-CSP}$ $[closure]$:
 $Stop \text{ is CSP}$
by $(simp \text{ add: } Stop\text{-def RHS-design-is-SRD unrest})$

lemma $Stop\text{-RHS-tri-design}$: $Stop = \mathbf{R}_s(true \vdash (\$tr' =_u \$tr) \diamond false)$
by $(rel\text{-auto})$

lemma $Stop\text{-RHS-rdes-def}$ $[rdes-def]$: $Stop = \mathbf{R}_s(true_r \vdash \mathcal{E}(true, \langle \rangle, \{\}_u) \diamond false)$
by $(rel\text{-auto})$

lemma $preR\text{-Skip}$ $[rdes]$: $pre_R(Skip) = true_r$
by $(rel\text{-auto})$

lemma *periR-Skip* [rdes]: $\text{peri}_R(\text{Skip}) = \text{false}$
by (*rel-auto*)

lemma *postR-Skip* [rdes]: $\text{post}_R(\text{Skip}) = \Phi(\text{true}, \text{id}, \langle \rangle)$
by (*rel-auto*)

lemma *Productive-Stop* [closure]:
Stop is Productive
by (*simp add: Stop-RHS-tri-design Healthy-def Productive-RHS-design-form unrest*)

lemma *Skip-left-lemma*:
assumes *P is CSP*
shows $\text{Skip} ;; P = \mathbf{R}_s ((\forall \$\text{ref} \cdot \text{pre}_R P) \vdash (\exists \$\text{ref} \cdot \text{cmt}_R P))$
proof –
have $\text{Skip} ;; P =$
 $\mathbf{R}_s ((\$tr' =_u \$tr \wedge \$st' =_u \$st) \text{wp}_r \text{pre}_R P \vdash$
 $(\$tr' =_u \$tr \wedge \$st' =_u \$st) ;; \text{peri}_R P \diamond$
 $(\$tr' =_u \$tr \wedge \$st' =_u \$st) ;; \text{post}_R P)$
by (*simp add: SRD-composition-wp alpha rdes closure wp assms rpred C1, rel-auto*)
also have $\dots = \mathbf{R}_s ((\forall \$\text{ref} \cdot \text{pre}_R P) \vdash$
 $(\$tr' =_u \$tr \wedge \neg \$\text{wait}' \wedge \$st' =_u \$st) ;; ((\exists \$st \cdot [II]_D) \triangleleft \$\text{wait} \triangleright \text{cmt}_R P))$
by (*rule cong[of $\mathbf{R}_s \mathbf{R}_s$], simp, rel-auto*)
also have $\dots = \mathbf{R}_s ((\forall \$\text{ref} \cdot \text{pre}_R P) \vdash (\exists \$\text{ref} \cdot \text{cmt}_R P))$
by (*rule cong[of $\mathbf{R}_s \mathbf{R}_s$], simp, rel-auto*)
finally show *?thesis* .
qed

lemma *Skip-left-unit-ref-unrest*:
assumes *P is CSP* $\$ref \# P \llbracket \text{false} / \$\text{wait} \rrbracket$
shows $\text{Skip} ;; P = P$
using *assms*
by (*simp add: Skip-left-lemma*)
(metis SRD-reactive-design-alt all-unrest cmt-unrest-ref cmt-wait-false ex-unrest pre-unrest-ref pre-wait-false)

lemma *CSP3-intro*:
 $\llbracket P \text{ is CSP}; \$ref \# P \llbracket \text{false} / \$\text{wait} \rrbracket \rrbracket \implies P \text{ is CSP3}$
by (*simp add: CSP3-def Healthy-def' Skip-left-unit-ref-unrest*)

lemma *ref-unrest-RHS-design*:
assumes $\$ref \# P \ \$ref \# Q_1 \ \$ref \# Q_2$
shows $\$ref \# (\mathbf{R}_s(P \vdash Q_1 \diamond Q_2)) \text{f}$
by (*simp add: RHS-def R1-def R2c-def R2s-def R3h-def design-def unrest usubst assms*)

lemma *CSP3-SRD-intro*:
assumes *P is CSP* $\$ref \# \text{pre}_R(P) \ \$ref \# \text{peri}_R(P) \ \$ref \# \text{post}_R(P)$
shows *P is CSP3*
proof –
have $P: \mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P)) = P$
by (*simp add: SRD-reactive-design-alt assms(1) wait'-cond-peri-post-cmt[THEN sym]*)
have $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P)) \text{ is CSP3}$
by (*rule CSP3-intro, simp add: assms P, simp add: ref-unrest-RHS-design assms*)
thus *?thesis*
by (*simp add: P*)
qed

lemma *Skip-unrest-ref* [unrest]: $\$ref \# Skip\llbracket false/\$wait \rrbracket$
 by (simp add: Skip-def RHS-def R1-def R2c-def R2s-def R3h-def design-def usubst unrest)

lemma *Skip-unrest-ref'* [unrest]: $\$ref' \# Skip\llbracket false/\$wait \rrbracket$
 by (simp add: Skip-def RHS-def R1-def R2c-def R2s-def R3h-def design-def usubst unrest)

lemma *CSP3-iff*:
 assumes *P is CSP*
 shows $P \text{ is } CSP3 \longleftrightarrow (\$ref \# P\llbracket false/\$wait \rrbracket)$

proof

assume 1: *P is CSP3*
 have $\$ref \# (Skip ;; P)\llbracket false/\$wait \rrbracket$
 by (simp add: usubst unrest)
 with 1 show $\$ref \# P\llbracket false/\$wait \rrbracket$
 by (metis CSP3-def Healthy-def)

next

assume 1: $\$ref \# P\llbracket false/\$wait \rrbracket$
 show *P is CSP3*
 by (simp add: 1 CSP3-intro assms)

qed

lemma *CSP3-unrest-ref* [unrest]:
 assumes *P is CSP* *P is CSP3*
 shows $\$ref \# pre_R(P) \ \$ref \# peri_R(P) \ \$ref \# post_R(P)$

proof –

have $a: (\$ref \# P\llbracket false/\$wait \rrbracket)$
 using *CSP3-iff assms* **by** blast
 from *a* show $\$ref \# pre_R(P)$
 by (rel-blast)
 from *a* show $\$ref \# peri_R(P)$
 by (rel-blast)
 from *a* show $\$ref \# post_R(P)$
 by (rel-blast)

qed

lemma *CSP3-rdes*:
 assumes *P is RR* *Q is RR* *R is RR*
 shows $CSP3(\mathbf{R}_s(P \vdash Q \diamond R)) = \mathbf{R}_s((\forall \$ref \cdot P) \vdash (\exists \$ref \cdot Q) \diamond (\exists \$ref \cdot R))$
 by (simp add: CSP3-def Skip-left-lemma closure assms rdes, rel-auto)

lemma *CSP3-form*:
 assumes *P is CSP*
 shows $CSP3(P) = \mathbf{R}_s((\forall \$ref \cdot pre_R(P)) \vdash (\exists \$ref \cdot peri_R(P)) \diamond (\exists \$ref \cdot post_R(P)))$
 by (simp add: CSP3-def Skip-left-lemma assms, rel-auto)

lemma *CSP3-Skip* [closure]:
Skip is CSP3
 by (rule CSP3-intro, simp add: Skip-is-CSP, simp add: Skip-def unrest)

lemma *CSP3-Stop* [closure]:
Stop is CSP3
 by (rule CSP3-intro, simp add: Stop-is-CSP, simp add: Stop-def unrest)

lemma *CSP3-Idempotent* [closure]: *Idempotent CSP3*
 by (metis (no-types, lifting) CSP3-Skip CSP3-def Healthy-if Idempotent-def seqr-assoc)

lemma *CSP3-Continuous: Continuous CSP3*

by (simp add: Continuous-def CSP3-def seq-Sup-distl)

lemma *Skip-right-lemma:*

assumes *P is CSP*

shows $P \;; \text{Skip} = \mathbf{R}_s ((\neg_r \text{pre}_R P) \text{wp}_r \text{false} \vdash ((\exists \$st' \cdot \text{cmt}_R P) \triangleleft \$wait' \triangleright (\exists \$ref' \cdot \text{cmt}_R P)))$

proof –

have $P \;; \text{Skip} = \mathbf{R}_s ((\neg_r \text{pre}_R P) \text{wp}_r \text{false} \vdash (\exists \$st' \cdot \text{peri}_R P) \diamond \text{post}_R P \;; (\$tr' =_u \$tr \wedge \$st' =_u \$st))$

by (simp add: SRD-composition-wp closure assms wp rdes rpred, rel-auto)

also have $\dots = \mathbf{R}_s ((\neg_r \text{pre}_R P) \text{wp}_r \text{false} \vdash ((\text{cmt}_R P \;; (\exists \$st \cdot [II]_D)) \triangleleft \$wait' \triangleright (\text{cmt}_R P \;; (\$tr' =_u \$tr \wedge \neg \$wait \wedge \$st' =_u \$st))))$

by (rule cong[of $\mathbf{R}_s \mathbf{R}_s$], simp, rel-auto)

also have $\dots = \mathbf{R}_s ((\neg_r \text{pre}_R P) \text{wp}_r \text{false} \vdash ((\exists \$st' \cdot \text{cmt}_R P) \triangleleft \$wait' \triangleright (\text{cmt}_R P \;; (\$tr' =_u \$tr \wedge \neg \$wait \wedge \$st' =_u \$st))))$

by (rule cong[of $\mathbf{R}_s \mathbf{R}_s$], simp, rel-auto)

also have $\dots = \mathbf{R}_s ((\neg_r \text{pre}_R P) \text{wp}_r \text{false} \vdash ((\exists \$st' \cdot \text{cmt}_R P) \triangleleft \$wait' \triangleright (\exists \$ref' \cdot \text{cmt}_R P)))$

by (rule cong[of $\mathbf{R}_s \mathbf{R}_s$], simp, rel-auto)

finally show ?thesis .

qed

lemma *Skip-right-tri-lemma:*

assumes *P is CSP*

shows $P \;; \text{Skip} = \mathbf{R}_s ((\neg_r \text{pre}_R P) \text{wp}_r \text{false} \vdash ((\exists \$st' \cdot \text{peri}_R P) \diamond (\exists \$ref' \cdot \text{post}_R P)))$

proof –

have $((\exists \$st' \cdot \text{cmt}_R P) \triangleleft \$wait' \triangleright (\exists \$ref' \cdot \text{cmt}_R P)) = ((\exists \$st' \cdot \text{peri}_R P) \diamond (\exists \$ref' \cdot \text{post}_R P))$

by (rel-auto)

thus ?thesis by (simp add: Skip-right-lemma[OF assms])

qed

lemma *CSP4-intro:*

assumes *P is CSP* $(\neg_r \text{pre}_R(P)) \;; R1(\text{true}) = (\neg_r \text{pre}_R(P))$

$\$st' \# (\text{cmt}_R P) \llbracket \text{true}/\$wait' \rrbracket \$ref' \# (\text{cmt}_R P) \llbracket \text{false}/\$wait' \rrbracket$

shows *P is CSP4*

proof –

have $\text{CSP4}(P) = \mathbf{R}_s ((\neg_r \text{pre}_R P) \text{wp}_r \text{false} \vdash ((\exists \$st' \cdot \text{cmt}_R P) \triangleleft \$wait' \triangleright (\exists \$ref' \cdot \text{cmt}_R P)))$

by (simp add: CSP4-def Skip-right-lemma assms(1))

also have $\dots = \mathbf{R}_s (\text{pre}_R(P) \vdash ((\exists \$st' \cdot \text{cmt}_R P) \llbracket \text{true}/\$wait' \rrbracket \triangleleft \$wait' \triangleright (\exists \$ref' \cdot \text{cmt}_R P) \llbracket \text{false}/\$wait' \rrbracket))$

by (simp add: wp-rea-def assms(2) rpred closure cond-var-subst-left cond-var-subst-right)

also have $\dots = \mathbf{R}_s (\text{pre}_R(P) \vdash ((\exists \$st' \cdot (\text{cmt}_R P) \llbracket \text{true}/\$wait' \rrbracket) \triangleleft \$wait' \triangleright (\exists \$ref' \cdot (\text{cmt}_R P) \llbracket \text{false}/\$wait' \rrbracket)))$

by (simp add: usubst unrest)

also have $\dots = \mathbf{R}_s (\text{pre}_R P \vdash ((\text{cmt}_R P) \llbracket \text{true}/\$wait' \rrbracket \triangleleft \$wait' \triangleright (\text{cmt}_R P) \llbracket \text{false}/\$wait' \rrbracket))$

by (simp add: ex-unrest assms)

also have $\dots = \mathbf{R}_s (\text{pre}_R P \vdash \text{cmt}_R P)$

by (simp add: cond-var-split)

also have $\dots = P$

by (simp add: SRD-reactive-design-alt assms(1))

finally show ?thesis

by (simp add: Healthy-def')

qed

lemma *CSP4-RC-intro:*

assumes P is CSP $pre_R(P)$ is RC
 $\$st' \# (cmt_R P) \llbracket true/\$wait' \rrbracket \$ref' \# (cmt_R P) \llbracket false/\$wait' \rrbracket$
shows P is CSP4

proof –

have $(\neg_r pre_R(P)) \;; R1(true) = (\neg_r pre_R(P))$
by (metis (no-types, lifting) R1-seqr-closure assms(2) rea-not-R1 rea-not-false rea-not-not wp-rea-RC-false wp-rea-def)
thus ?thesis
by (simp add: CSP4-intro assms)

qed

lemma *CSP4-rdes:*

assumes P is RR Q is RR R is RR
shows $CSP4(\mathbf{R}_s(P \vdash Q \diamond R)) = \mathbf{R}_s((\neg_r P) wp_r false \vdash ((\exists \$st' \cdot Q) \diamond (\exists \$ref' \cdot R)))$
by (simp add: CSP4-def Skip-right-lemma closure assms rdes, rel-auto, blast+)

lemma *CSP4-form:*

assumes P is CSP
shows $CSP4(P) = \mathbf{R}_s((\neg_r pre_R P) wp_r false \vdash ((\exists \$st' \cdot peri_R P) \diamond (\exists \$ref' \cdot post_R P)))$
by (simp add: CSP4-def Skip-right-tri-lemma assms)

lemma *Skip-srdes-right-unit:*

$(Skip \:: (\sigma, \varphi) \text{ action}) \;; II_R = Skip$
by (rdes-simp)

lemma *Skip-srdes-left-unit:*

$II_R \;; (Skip \:: (\sigma, \varphi) \text{ action}) = Skip$
by (rdes-eq)

lemma *CSP4-right-subsumes-RD3:* $RD3(CSP4(P)) = CSP4(P)$

by (metis (no-types, hide-lams) CSP4-def RD3-def Skip-srdes-right-unit seqr-assoc)

lemma *CSP4-implies-RD3:* P is CSP4 $\implies P$ is RD3

by (metis CSP4-right-subsumes-RD3 Healthy-def)

lemma *CSP4-tri-intro:*

assumes P is CSP $(\neg_r pre_R(P)) \;; R1(true) = (\neg_r pre_R(P)) \$st' \# peri_R(P) \$ref' \# post_R(P)$
shows P is CSP4
using assms
by (rule-tac CSP4-intro, simp-all add: pre_R-def peri_R-def post_R-def usubst cmt_R-def)

lemma *CSP4-NSRD-intro:*

assumes P is NSRD $\$ref' \# post_R(P)$
shows P is CSP4
by (simp add: CSP4-tri-intro NSRD-is-SRD NSRD-neg-pre-unit NSRD-st'-unrest-peri assms)

lemma *CSP3-commutes-CSP4:* $CSP3(CSP4(P)) = CSP4(CSP3(P))$

by (simp add: CSP3-def CSP4-def seqr-assoc)

lemma *NCSP-implies-CSP [closure]:* P is NCSP $\implies P$ is CSP

by (metis (no-types, hide-lams) CSP3-def CSP4-def Healthy-def NCSP-def SRD-idem SRD-seqr-closure Skip-is-CSP comp-apply)

lemma *NCSP-elim* [*RD-elim*]:

$\llbracket X \text{ is NCSP}; P(\mathbf{R}_s(\text{pre}_R(X) \vdash \text{peri}_R(X) \diamond \text{post}_R(X))) \rrbracket \implies P(X)$
by (*simp add: SRD-reactive-tri-design closure*)

lemma *NCSP-implies-CSP3* [*closure*]:

$P \text{ is NCSP} \implies P \text{ is CSP3}$

by (*metis (no-types, lifting) CSP3-def Healthy-def' NCSP-def Skip-is-CSP Skip-left-unit-ref-unrest Skip-unrest-ref comp-apply segr-assoc*)

lemma *NCSP-implies-CSP4* [*closure*]:

$P \text{ is NCSP} \implies P \text{ is CSP4}$

by (*metis (no-types, hide-lams) CSP3-commutes-CSP4 Healthy-def NCSP-def NCSP-implies-CSP NCSP-implies-CSP3 comp-apply*)

lemma *NCSP-implies-RD3* [*closure*]: $P \text{ is NCSP} \implies P \text{ is RD3}$

by (*metis CSP3-commutes-CSP4 CSP4-right-subsumes-RD3 Healthy-def NCSP-def comp-apply*)

lemma *NCSP-implies-NSRD* [*closure*]: $P \text{ is NCSP} \implies P \text{ is NSRD}$

by (*simp add: NCSP-implies-CSP NCSP-implies-RD3 SRD-RD3-implies-NSRD*)

lemma *NCSP-subset-implies-CSP* [*closure*]:

$A \subseteq \llbracket \text{NCSP} \rrbracket_H \implies A \subseteq \llbracket \text{CSP} \rrbracket_H$

using *NCSP-implies-CSP* **by** *blast*

lemma *NCSP-subset-implies-NSRD* [*closure*]:

$A \subseteq \llbracket \text{NCSP} \rrbracket_H \implies A \subseteq \llbracket \text{NSRD} \rrbracket_H$

using *NCSP-implies-NSRD* **by** *blast*

lemma *CSP-Healthy-subset-member*: $\llbracket P \in A; A \subseteq \llbracket \text{CSP} \rrbracket_H \rrbracket \implies P \text{ is CSP}$

by (*simp add: is-Healthy-subset-member*)

lemma *CSP3-Healthy-subset-member*: $\llbracket P \in A; A \subseteq \llbracket \text{CSP3} \rrbracket_H \rrbracket \implies P \text{ is CSP3}$

by (*simp add: is-Healthy-subset-member*)

lemma *CSP4-Healthy-subset-member*: $\llbracket P \in A; A \subseteq \llbracket \text{CSP4} \rrbracket_H \rrbracket \implies P \text{ is CSP4}$

by (*simp add: is-Healthy-subset-member*)

lemma *NCSP-Healthy-subset-member*: $\llbracket P \in A; A \subseteq \llbracket \text{NCSP} \rrbracket_H \rrbracket \implies P \text{ is NCSP}$

by (*simp add: is-Healthy-subset-member*)

lemma *NCSP-intro*:

assumes $P \text{ is CSP } P \text{ is CSP3 } P \text{ is CSP4}$

shows $P \text{ is NCSP}$

by (*metis Healthy-def NCSP-def assms comp-eq-dest-lhs*)

lemma *Skip-left-unit*: $P \text{ is NCSP} \implies \text{Skip} ;; P = P$

by (*metis (full-types) CSP3-def Healthy-if NCSP-implies-CSP3*)

lemma *Skip-right-unit*: $P \text{ is NCSP} \implies P ;; \text{Skip} = P$

by (*metis (full-types) CSP4-def Healthy-if NCSP-implies-CSP4*)

lemma *NCSP-NSRD-intro*:

assumes $P \text{ is NSRD } \$\text{ref} \# \text{pre}_R(P) \$\text{ref} \# \text{peri}_R(P) \$\text{ref} \# \text{post}_R(P) \$\text{ref}' \# \text{post}_R(P)$

shows $P \text{ is NCSP}$

by (*simp add: CSP3-SRD-intro CSP4-NSRD-intro NCSP-intro NSRD-is-SRD assms*)

lemma *CSP4-neg-pre-unit*:

assumes P is CSP P is CSP4

shows $(\neg_r \text{pre}_R(P)) \;; R1(\text{true}) = (\neg_r \text{pre}_R(P))$

by (*simp add: CSP4-implies-RD3 NSRD-neg-pre-unit SRD-RD3-implies-NSRD assms(1) assms(2)*)

lemma *NSRD-CSP4-intro*:

assumes P is CSP P is CSP4

shows P is NSRD

by (*simp add: CSP4-implies-RD3 SRD-RD3-implies-NSRD assms(1) assms(2)*)

lemma *NCSP-form*:

$NCSP\ P = \mathbf{R}_s((\forall \$ref \cdot (\neg_r \text{pre}_R(P))\ wp_r\ \text{false}) \vdash ((\exists \$ref \cdot \exists \$st' \cdot \text{peri}_R(P)) \diamond (\exists \$ref \cdot \exists \$ref' \cdot \text{post}_R(P))))$

proof –

have $NCSP\ P = CSP3\ (CSP4\ (NSRD\ P))$

by (*metis (no-types, hide-lams) CSP4-def NCSP-def NSRD-alt-def RA1 RD3-def Skip-srdes-left-unit o-apply*)

also

have $\dots = \mathbf{R}_s((\forall \$ref \cdot (\neg_r \text{pre}_R\ (NSRD\ P))\ wp_r\ \text{false}) \vdash ((\exists \$ref \cdot \exists \$st' \cdot \text{peri}_R\ (NSRD\ P)) \diamond (\exists \$ref \cdot \exists \$ref' \cdot \text{post}_R\ (NSRD\ P))))$

by (*simp add: CSP3-form CSP4-form closure unrest rdes, rel-auto*)

also have $\dots = \mathbf{R}_s((\forall \$ref \cdot (\neg_r \text{pre}_R(P))\ wp_r\ \text{false}) \vdash ((\exists \$ref \cdot \exists \$st' \cdot \text{peri}_R(P)) \diamond (\exists \$ref \cdot \exists \$ref' \cdot \text{post}_R(P))))$

by (*simp add: NSRD-form rdes closure, rel-blast*)

finally show *?thesis* .

qed

lemma *CSP4-st'-unrest-peri* [*unrest*]:

assumes P is CSP P is CSP4

shows $\$st' \not\# \text{peri}_R(P)$

by (*simp add: NSRD-CSP4-intro NSRD-st'-unrest-peri assms*)

lemma *CSP4-healthy-form*:

assumes P is CSP P is CSP4

shows $P = \mathbf{R}_s((\neg_r \text{pre}_R\ P)\ wp_r\ \text{false} \vdash ((\exists \$st' \cdot \text{peri}_R(P)) \diamond (\exists \$ref' \cdot \text{post}_R(P))))$

proof –

have $P = \mathbf{R}_s((\neg_r \text{pre}_R\ P)\ wp_r\ \text{false} \vdash ((\exists \$st' \cdot \text{cmt}_R\ P) \triangleleft \$wait' \triangleright (\exists \$ref' \cdot \text{cmt}_R\ P)))$

by (*metis CSP4-def Healthy-def Skip-right-lemma assms(1) assms(2)*)

also have $\dots = \mathbf{R}_s((\neg_r \text{pre}_R\ P)\ wp_r\ \text{false} \vdash ((\exists \$st' \cdot \text{cmt}_R\ P) \llbracket \text{true}/\$wait' \rrbracket \triangleleft \$wait' \triangleright (\exists \$ref' \cdot \text{cmt}_R\ P) \llbracket \text{false}/\$wait' \rrbracket))$

by (*metis (no-types, hide-lams) subst-wait'-left-subst subst-wait'-right-subst wait'-cond-def*)

also have $\dots = \mathbf{R}_s((\neg_r \text{pre}_R\ P)\ wp_r\ \text{false} \vdash ((\exists \$st' \cdot \text{peri}_R(P)) \diamond (\exists \$ref' \cdot \text{post}_R(P))))$

by (*simp add: wait'-cond-def usubst peri_R-def post_R-def cmt_R-def unrest*)

finally show *?thesis* .

qed

lemma *CSP4-ref'-unrest-pre* [*unrest*]:

assumes P is CSP P is CSP4

shows $\$ref' \not\# \text{pre}_R(P)$

proof –

have $\text{pre}_R(P) = \text{pre}_R(\mathbf{R}_s((\neg_r \text{pre}_R\ P)\ wp_r\ \text{false} \vdash ((\exists \$st' \cdot \text{peri}_R(P)) \diamond (\exists \$ref' \cdot \text{post}_R(P))))$

using *CSP4-healthy-form assms(1) assms(2)* **by** *fastforce*

also have $\dots = (\neg_r \text{pre}_R\ P)\ wp_r\ \text{false}$

by (simp add: rea-pre-RHS-design wp-rea-def usubst unrest
 CSP4-neg-pre-unit R1-rea-not R2c-preR R2c-rea-not assms)
 also have $\$ref' \# \dots$
 by (simp add: wp-rea-def unrest)
 finally show ?thesis .
 qed

lemma NCSP-set-unrest-pre-wait':
 assumes $A \subseteq \llbracket NCSP \rrbracket_H$
 shows $\bigwedge P. P \in A \implies \$wait' \# pre_R(P)$
proof –
 fix P
 assume $P \in A$
 hence P is NSRD
 using NCSP-implies-NSRD assms by auto
 thus $\$wait' \# pre_R(P)$
 using NSRD-wait'-unrest-pre by blast
 qed

lemma CSP4-set-unrest-pre-st':
 assumes $A \subseteq \llbracket CSP \rrbracket_H$ $A \subseteq \llbracket CSP4 \rrbracket_H$
 shows $\bigwedge P. P \in A \implies \$st' \# pre_R(P)$
proof –
 fix P
 assume $P \in A$
 hence P is NSRD
 using NSRD-CSP4-intro assms(1) assms(2) by blast
 thus $\$st' \# pre_R(P)$
 using NSRD-st'-unrest-pre by blast
 qed

lemma CSP4-ref'-unrest-post [unrest]:
 assumes P is CSP P is CSP4
 shows $\$ref' \# post_R(P)$
proof –
 have $post_R(P) = post_R(\mathbf{R}_s((\neg_r pre_R P) wp_r false \vdash ((\exists \$st' \cdot peri_R(P)) \diamond (\exists \$ref' \cdot post_R(P))))$
 using CSP4-healthy-form assms(1) assms(2) by fastforce
 also have $\dots = R1 (R2c ((\neg_r pre_R P) wp_r false \Rightarrow_r (\exists \$ref' \cdot post_R P)))$
 by (simp add: rea-post-RHS-design usubst unrest wp-rea-def)
 also have $\$ref' \# \dots$
 by (simp add: R1-def R2c-def wp-rea-def unrest)
 finally show ?thesis .
 qed

lemma CSP3-Chaos [closure]: Chaos is CSP3
 by (simp add: Chaos-def, rule CSP3-intro, simp-all add: RHS-design-is-SRD unrest)

lemma CSP4-Chaos [closure]: Chaos is CSP4
 by (rule CSP4-tri-intro, simp-all add: closure rdes unrest)

lemma NCSP-Chaos [closure]: Chaos is NCSP
 by (simp add: NCSP-intro closure)

lemma CSP3-Miracle [closure]: Miracle is CSP3
 by (simp add: Miracle-def, rule CSP3-intro, simp-all add: RHS-design-is-SRD unrest)

lemma *CSP4-Miracle* [closure]: *Miracle is CSP4*
 by (rule *CSP4-tri-intro*, simp-all add: closure rdes unrest)

lemma *NCSP-Miracle* [closure]: *Miracle is NCSP*
 by (simp add: NCSP-intro closure)

lemma *NCSP-seqr-closure* [closure]:
 assumes *P is NCSP Q is NCSP*
 shows *P ;; Q is NCSP*
 by (metis (no-types, lifting) *CSP3-def CSP4-def Healthy-def' NCSP-implies-CSP NCSP-implies-CSP3*
NCSP-implies-CSP4 NCSP-intro SRD-seqr-closure assms(1) assms(2) seqr-assoc)

lemma *CSP4-Skip* [closure]: *Skip is CSP4*
 apply (rule *CSP4-intro*, simp-all add: *Skip-is-CSP*)
 apply (simp-all add: *Skip-def rea-pre-RHS-design rea-cmt-RHS-design usubst unrest R2c-true*)
 done

lemma *NCSP-Skip* [closure]: *Skip is NCSP*
 by (metis *CSP3-Skip CSP4-Skip Healthy-def NCSP-def Skip-is-CSP comp-apply*)

lemma *CSP4-Stop* [closure]: *Stop is CSP4*
 apply (rule *CSP4-intro*, simp-all add: *Stop-is-CSP*)
 apply (simp-all add: *Stop-def rea-pre-RHS-design rea-cmt-RHS-design usubst unrest R2c-true*)
 done

lemma *NCSP-Stop* [closure]: *Stop is NCSP*
 by (metis *CSP3-Stop CSP4-Stop Healthy-def NCSP-def Stop-is-CSP comp-apply*)

lemma *CSP4-Idempotent*: *Idempotent CSP4*
 by (metis (no-types, lifting) *CSP3-Skip CSP3-def CSP4-def Healthy-if Idempotent-def seqr-assoc*)

lemma *CSP4-Continuous*: *Continuous CSP4*
 by (simp add: *Continuous-def CSP4-def seq-Sup-distr*)

lemma *rdes-frame-ext-NCSP-closed* [closure]:
 assumes *vwb-lens a P is NCSP*
 shows *a:[P]_R⁺ is NCSP*
 by (metis (no-types, lifting) *CSP3-def CSP4-def Healthy-intro NCSP-Skip NCSP-implies-NSRD NCSP-intro*
NSRD-is-SRD Skip-frame Skip-left-unit Skip-right-unit assms(1) assms(2) rdes-frame-ext-NSRD-closed
seq-srea-frame)

lemma *preR-Stop* [rdes]: *pre_R(Stop) = true_r*
 by (simp add: *Stop-def Stop-is-CSP rea-pre-RHS-design unrest usubst R2c-true*)

lemma *periR-Stop* [rdes]: *peri_R(Stop) = $\mathcal{E}(\text{true}, \langle \rangle, \{\}_u)$*
 by (rel-auto)

lemma *postR-Stop* [rdes]: *post_R(Stop) = false*
 by (rel-auto)

lemma *cmtR-Stop* [rdes]: *cmt_R(Stop) = ($\$tr' =_u \$tr \wedge \$wait'$)*
 by (rel-auto)

lemma *NCSP-Idempotent* [closure]: *Idempotent NCSP*

by (clarsimp simp add: NCSP-def Idempotent-def)
 (metis (no-types, hide-lams) CSP3-Idempotent CSP3-def CSP4-Idempotent CSP4-def Healthy-def
 Idempotent-def SRD-idem SRD-seqr-closure Skip-is-CSP seqr-assoc)

lemma NCSP-Continuous [closure]: Continuous NCSP
 by (simp add: CSP3-Continuous CSP4-Continuous Continuous-comp NCSP-def SRD-Continuous)

lemma preR-CRR [closure]: P is NCSP \implies $pre_R(P)$ is CRR
 by (rule CRR-intro, simp-all add: closure unrest)

lemma periR-CRR [closure]: P is NCSP \implies $peri_R(P)$ is CRR
 by (rule CRR-intro, simp-all add: closure unrest)

lemma postR-CRR [closure]: P is NCSP \implies $post_R(P)$ is CRR
 by (rule CRR-intro, simp-all add: closure unrest)

lemma NCSP-rdes-intro [closure]:
 assumes P is CRC Q is CRR R is CRR
 $\$st' \# Q \$ref' \# R$
 shows $\mathbf{R}_s(P \vdash Q \diamond R)$ is NCSP
 apply (rule NCSP-intro)
 apply (simp-all add: closure assms)
 apply (rule CSP3-SRD-intro)
 apply (simp-all add: rdes closure assms unrest)
 apply (rule CSP4-tri-intro)
 apply (simp-all add: rdes closure assms unrest)
 apply (metis (no-types, lifting) CRC-implies-RC R1-seqr-closure assms(1) rea-not-R1 rea-not-false
 rea-not-not wp-rea-RC-false wp-rea-def)
 done

lemma NCSP-preR-CRC [closure]:
 assumes P is NCSP
 shows $pre_R(P)$ is CRC
 by (rule CRC-intro, simp-all add: closure assms unrest)

lemma CSP3-Sup-closure [closure]:
 $A \subseteq \llbracket CSP3 \rrbracket_H \implies (\bigcap A)$ is CSP3
 apply (auto simp add: CSP3-def Healthy-def seq-Sup-distl)
 apply (rule cong[of Sup])
 apply (simp)
 using image-iff apply force
 done

lemma CSP4-Sup-closure [closure]:
 $A \subseteq \llbracket CSP4 \rrbracket_H \implies (\bigcap A)$ is CSP4
 apply (auto simp add: CSP4-def Healthy-def seq-Sup-distr)
 apply (rule cong[of Sup])
 apply (simp)
 using image-iff apply force
 done

lemma NCSP-Sup-closure [closure]: $\llbracket A \subseteq \llbracket NCSP \rrbracket_H; A \neq \{\} \rrbracket \implies (\bigcap A)$ is NCSP
 apply (rule NCSP-intro, simp-all add: closure)
 apply (metis (no-types, lifting) Ball-Collect CSP3-Sup-closure NCSP-implies-CSP3)
 apply (metis (no-types, lifting) Ball-Collect CSP4-Sup-closure NCSP-implies-CSP4)

done

lemma *NCSP-SUP-closure* [closure]: $\llbracket \bigwedge i. P(i) \text{ is NCSP}; A \neq \{\} \rrbracket \implies (\bigcap_{i \in A}. P(i)) \text{ is NCSP}$
 by (metis (mono-tags, lifting) Ball-Collect NCSP-Sup-closure image-iff image-is-empty)

lemma *PCSP-implies-NCSP* [closure]:

assumes *P* is PCSP

shows *P* is NCSP

proof –

have $P = \text{Productive}(\text{NCSP}(\text{NCSP } P))$

by (metis (no-types, hide-lams) Healthy-def' Idempotent-def NCSP-Idempotent assms comp-apply)

also have $\dots = \mathbf{R}_s ((\forall \$ref \cdot (\neg_r \text{pre}_R(\text{NCSP } P)) \text{wp}_r \text{false}) \vdash$
 $(\exists \$ref \cdot \exists \$st' \cdot \text{peri}_R(\text{NCSP } P)) \diamond$
 $((\exists \$ref \cdot \exists \$ref' \cdot \text{post}_R(\text{NCSP } P)) \wedge \$tr <_u \$tr'))$

by (simp add: NCSP-form Productive-RHS-design-form unrest closure)

also have \dots is NCSP

apply (rule NCSP-rdes-intro)

apply (rule CRC-intro)

apply (simp-all add: unrest ex-unrest all-unrest closure)

done

finally show ?thesis .

qed

lemma *PCSP-elim* [RD-elim]:

assumes *X* is PCSP $P (\mathbf{R}_s ((\text{pre}_R X) \vdash \text{peri}_R X \diamond (R_4(\text{post}_R X))))$

shows $P X$

by (metis R4-def Healthy-if NCSP-implies-CSP PCSP-implies-NCSP Productive-form assms comp-apply)

lemma *ICSP-implies-NCSP* [closure]:

assumes *P* is ICSP

shows *P* is NCSP

proof –

have $P = \text{ISRDI}(\text{NCSP}(\text{NCSP } P))$

by (metis (no-types, hide-lams) Healthy-def' Idempotent-def NCSP-Idempotent assms comp-apply)

also have $\dots = \text{ISRDI} (\mathbf{R}_s ((\forall \$ref \cdot (\neg_r \text{pre}_R(\text{NCSP } P)) \text{wp}_r \text{false}) \vdash$
 $(\exists \$ref \cdot \exists \$st' \cdot \text{peri}_R(\text{NCSP } P)) \diamond$
 $(\exists \$ref \cdot \exists \$ref' \cdot \text{post}_R(\text{NCSP } P))))$

by (simp add: NCSP-form)

also have $\dots = \mathbf{R}_s ((\forall \$ref \cdot (\neg_r \text{pre}_R(\text{NCSP } P)) \text{wp}_r \text{false}) \vdash$
 $\text{false} \diamond$
 $((\exists \$ref \cdot \exists \$ref' \cdot \text{post}_R(\text{NCSP } P)) \wedge \$tr' =_u \$tr))$

by (simp-all add: ISRDI-RHS-design-form closure rdes unrest)

also have \dots is NCSP

apply (rule NCSP-rdes-intro)

apply (rule CRC-intro)

apply (simp-all add: unrest ex-unrest all-unrest closure)

done

finally show ?thesis .

qed

lemma *ICSP-implies-ISRDI* [closure]:

assumes *P* is ICSP

shows *P* is ISRDI

by (metis (no-types, hide-lams) Healthy-def ICSP-implies-NCSP ISRDI-def NCSP-implies-ISRDI assms)

comp-apply)

lemma *ICSP-elim* [*RD-elim*]:

assumes *X is ICSP* *P* ($\mathbf{R}_s ((pre_R X) \vdash false \diamond (post_R X \wedge \$tr' =_u \$tr))$)

shows *P X*

by (*metis Healthy-if NCSP-implies-CSP ICSP-implies-NCSP ISRD1-form assms comp-apply*)

lemma *ICSP-Stop-right-zero-lemma*:

$(P \wedge (\$tr' =_u \$tr)) ;; true_r = true_r \implies (P \wedge (\$tr' =_u \$tr)) ;; (\$tr' =_u \$tr) = (\$tr' =_u \$tr)$

by (*rel-blast*)

lemma *ICSP-Stop-right-zero*:

assumes *P is ICSP* $pre_R(P) = true_r$ $post_R(P) ;; true_r = true_r$

shows *P ;; Stop = Stop*

proof –

from *assms(3)* **have** $1:(post_R P \wedge \$tr' =_u \$tr) ;; true_r = true_r$

by (*rel-auto, metis (full-types, hide-lams) dual-order.antisym order-refl*)

show *?thesis*

by (*rdes-simp cls: assms(1), simp add: csp-enable-nothing assms(2) ICSP-Stop-right-zero-lemma[OF 1]*)

qed

lemma *ICSP-intro*: $\llbracket P \text{ is NCSP}; P \text{ is ISRD1} \rrbracket \implies P \text{ is ICSP}$

using *Healthy-comp* **by** *blast*

lemma *seq-ICSP-closed* [*closure*]:

assumes *P is ICSP* *Q is ICSP*

shows *P ;; Q is ICSP*

by (*meson ICSP-implies-ISRD ICSP-implies-NCSP ICSP-intro ISRD-implies-ISRD1 NCSP-seqr-closure assms seq-ISRD-closed*)

lemma *Miracle-ICSP* [*closure*]: *Miracle is ICSP*

by (*rule ICSP-intro, simp add: closure, simp add: ISRD1-rdes-intro rdes-def closure*)

5.2 CSP theories

typeddecl *TCSP*

abbreviation $TCSP \equiv UTHY(TCSP, ('\sigma, '\varphi) \text{ st-csp})$

overloading

tcspl-hcond == *utpl-hcond* :: $(TCSP, ('\sigma, '\varphi) \text{ st-csp}) \text{ uthy} \Rightarrow ((' \sigma, '\varphi) \text{ st-csp} \times (' \sigma, '\varphi) \text{ st-csp}) \text{ health}$

tcspl-unit == *utpl-unit* :: $(TCSP, ('\sigma, '\varphi) \text{ st-csp}) \text{ uthy} \Rightarrow (' \sigma, '\varphi) \text{ action}$

begin

definition *tcspl-hcond* :: $(TCSP, ('\sigma, '\varphi) \text{ st-csp}) \text{ uthy} \Rightarrow ((' \sigma, '\varphi) \text{ st-csp} \times (' \sigma, '\varphi) \text{ st-csp}) \text{ health}$ **where**

[*upred-defs*]: *tcspl-hcond* *T* = *NCSP*

definition *tcspl-unit* :: $(TCSP, ('\sigma, '\varphi) \text{ st-csp}) \text{ uthy} \Rightarrow (' \sigma, '\varphi) \text{ action}$ **where**

[*upred-defs*]: *tcspl-unit* *T* = *Skip*

end

interpretation *csp-theory*: *utpl-theory-kleene* $UTHY(TCSP, ('\sigma, '\varphi) \text{ st-csp})$

rewrites $\bigwedge P. P \in \text{carrier (uthy-order } TCSP) \longleftrightarrow P \text{ is NCSP}$

and $P \text{ is } \mathcal{H}_{TCSP} \longleftrightarrow P \text{ is NCSP}$

and $\mathcal{II}_{TCSP} = \text{Skip}$

and $\top_{TCSP} = \text{Miracle}$

and $\text{carrier (uthy-order } TCSP) \rightarrow \text{carrier (uthy-order } TCSP) \equiv \llbracket NCSP \rrbracket_H \rightarrow \llbracket NCSP \rrbracket_H$

and $A \subseteq \text{carrier } (\text{uthy-order } TCSP) \longleftrightarrow A \subseteq \llbracket NCSP \rrbracket_H$
and $le \text{ (uthy-order } TCSP) = (\sqsubseteq)$
proof –
interpret lat : $\text{utp-theory-continuous } UTHY(TCSP, ('σ, 'φ) \text{ st-csp})$
by ($\text{unfold-locales, simp-all add: tcsp-hcond-def closure Healthy-if}$)
show 1 : $\top_{TCSP} = (\text{Miracle} :: ('σ, 'φ) \text{ action})$
by ($\text{metis NCSP-Miracle NCSP-implies-CSP lat.top-healthy lat.utp-theory-continuous-axioms srdes-theory-continuous.}$
 $\text{tcsp-hcond-def upred-semiring.add-commute utp-theory-continuous.meet-top})$

thus $\text{utp-theory-kleene } UTHY(TCSP, ('σ, 'φ) \text{ st-csp})$
by ($\text{unfold-locales, simp-all add: tcsp-hcond-def tcsp-unit-def Skip-left-unit Skip-right-unit closure}$
 $\text{Healthy-if Miracle-left-zero }$)
qed ($\text{simp-all add: tcsp-hcond-def tcsp-unit-def closure Healthy-if}$)

declare $\text{csp-theory.top-healthy [simp del]}$
declare $\text{csp-theory.bottom-healthy [simp del]}$

abbreviation $\text{TestC } (test_C) \text{ where}$
 $\text{test}_C P \equiv \text{utest } TCSP P$

abbreviation $\text{StarC} :: ('σ, 'φ) \text{ action} \Rightarrow ('σ, 'φ) \text{ action } (-^{*C} [999] 999) \text{ where}$
 $\text{StarC } P \equiv P \star_{TCSP}$

lemma csp-bottom-Chaos : $\perp_{TCSP} = \text{Chaos}$
using $\text{NCSP-Chaos NCSP-implies-CSP}$ **by** auto

lemma csp-top-Miracle : $\top_{TCSP} = \text{Miracle}$
by ($\text{simp add: csp-theory.healthy-top csp-theory.utp-theory-mono-axioms utp-theory-mono.healthy-top}$)

5.3 Algebraic laws

lemma Stop-left-zero :
assumes $P \text{ is CSP}$
shows $\text{Stop} ;; P = \text{Stop}$
by ($\text{simp add: NSRD-seq-post-false assms NCSP-implies-NSRD NCSP-Stop postR-Stop}$)

end

6 Stateful-Failure Reactive Contracts

theory $\text{utp-sfrd-contracts}$
imports utp-sfrd-healths
begin

definition $\text{mk-CRD} :: 's \text{ upred} \Rightarrow ('e \text{ list} \Rightarrow 'e \text{ set} \Rightarrow 's \text{ upred}) \Rightarrow ('e \text{ list} \Rightarrow 's \text{ hrel}) \Rightarrow ('s, 'e) \text{ action}$
where
 $[\text{rdes-def}]: \text{mk-CRD } P Q R = \mathbf{R}_s([P]_{S<} \vdash [Q \ x \ r]_{S<} \llbracket x \rightarrow \&tt \rrbracket \llbracket r \rightarrow \$ref \ ' \rrbracket \diamond [R(x)]_S \llbracket x \rightarrow \&tt \rrbracket)$

syntax
 $\text{-ref-var} :: \text{logic}$
 $\text{-mk-CRD} :: \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic } ([\text{-} / \vdash \text{-} / \mid \text{-}]_C)$

parse-translation \ll
 let
 $\text{fun ref-var-tr } [] = \text{Syntax.free refs}$

$| \text{ref-var-tr} - = \text{raise Match};$
in
 $[(\text{@}\{\text{syntax-const -ref-var}\}, K \text{ref-var-tr})]$
end
 \gg

translations

$[P \vdash Q \mid R]_C \Rightarrow \text{CONST mk-CRD } P (\lambda \text{-trace-var -ref-var. } Q) (\lambda \text{-trace-var. } R)$
 $[P \vdash Q \mid R]_C \Leftarrow \text{CONST mk-CRD } P (\lambda x r. Q) (\lambda y. R)$

lemma *CSP-mk-CRD* [*closure*]: $[P \vdash Q \text{ trace refs} \mid R(\text{trace})]_C$ is CSP
by (*simp add: mk-CRD-def closure unrest*)

lemma *preR-mk-CRD* [*rdes*]: $\text{pre}_R([P \vdash Q \text{ trace refs} \mid R(\text{trace})]_C) = [P]_{S<}$
by (*simp add: mk-CRD-def rea-pre-RHS-design usubst unrest R2c-not R2c-lift-state-pre rea-st-cond-def, rel-auto*)

lemma *periR-mk-CRD* [*rdes*]: $\text{peri}_R([P \vdash Q \text{ trace refs} \mid R(\text{trace})]_C) = ([P]_{S<} \Rightarrow_r ([Q \text{ trace refs}]_{S<})) \llbracket (\text{trace}, \text{refs}) \rightarrow (\&tt, \$n) \rrbracket$
by (*simp add: mk-CRD-def rea-peri-RHS-design usubst unrest R2c-not R2c-lift-state-pre impl-alt-def R2c-disj R2c-msubst-tt R1-disj, rel-auto*)

lemma *postR-mk-CRD* [*rdes*]: $\text{post}_R([P \vdash Q \text{ trace refs} \mid R(\text{trace})]_C) = ([P]_{S<} \Rightarrow_r ([R(\text{trace})]_S) \llbracket \text{trace} \rightarrow \&tt \rrbracket)$
by (*simp add: mk-CRD-def rea-post-RHS-design usubst unrest R2c-not R2c-lift-state-pre impl-alt-def R2c-disj R2c-msubst-tt R1-disj, rel-auto*)

Refinement introduction law for contracts

lemma *CRD-contract-refine*:

assumes
 $Q \text{ is CSP } '[P_1]_{S<} \Rightarrow \text{pre}_R Q'$
 $'[P_1]_{S<} \wedge \text{peri}_R Q \Rightarrow [P_2 \text{ } t \text{ } r]_{S<} \llbracket t \rightarrow \&tt \rrbracket \llbracket r \rightarrow \$ref' \rrbracket'$
 $'[P_1]_{S<} \wedge \text{post}_R Q \Rightarrow [P_3 \text{ } x]_S \llbracket x \rightarrow \&tt \rrbracket'$
shows $[P_1 \vdash P_2 \text{ trace refs} \mid P_3(\text{trace})]_C \sqsubseteq Q$

proof –

have $[P_1 \vdash P_2 \text{ trace refs} \mid P_3(\text{trace})]_C \sqsubseteq \mathbf{R}_s(\text{pre}_R(Q) \vdash \text{peri}_R(Q) \diamond \text{post}_R(Q))$
using *assms by (simp add: mk-CRD-def, rule-tac sdes-tri-refine-intro, rel-auto+)*
thus *thesis*
by (*simp add: SRD-reactive-tri-design assms(1)*)

qed

lemma *CRD-contract-refine'*:

assumes
 $Q \text{ is CSP } '[P_1]_{S<} \Rightarrow \text{pre}_R Q'$
 $[P_2 \text{ } t \text{ } r]_{S<} \llbracket t \rightarrow \&tt \rrbracket \llbracket r \rightarrow \$ref' \rrbracket \sqsubseteq ([P_1]_{S<} \wedge \text{peri}_R Q)$
 $[P_3 \text{ } x]_S \llbracket x \rightarrow \&tt \rrbracket \sqsubseteq ([P_1]_{S<} \wedge \text{post}_R Q)$
shows $[P_1 \vdash P_2 \text{ trace refs} \mid P_3(\text{trace})]_C \sqsubseteq Q$
using *assms by (rule-tac CRD-contract-refine, simp-all add: refBy-order)*

lemma *CRD-refine-CRD*:

assumes
 $'[P_1]_{S<} \Rightarrow ([Q_1]_{S<} :: ('e, 's) \text{ action})'$
 $([P_2 \text{ } x \text{ } r]_{S<} \llbracket x \rightarrow \&tt \rrbracket \llbracket r \rightarrow \$ref' \rrbracket) \sqsubseteq ([P_1]_{S<} \wedge [Q_2 \text{ } x \text{ } r]_{S<} \llbracket x \rightarrow \&tt \rrbracket \llbracket r \rightarrow \$ref' \rrbracket :: ('e, 's) \text{ action})$
 $[P_3 \text{ } x]_S \llbracket x \rightarrow \&tt \rrbracket \sqsubseteq ([P_1]_{S<} \wedge [Q_3 \text{ } x]_S \llbracket x \rightarrow \&tt \rrbracket :: ('e, 's) \text{ action})$
shows $([P_1 \vdash P_2 \text{ trace refs} \mid P_3 \text{ trace}]_C :: ('e, 's) \text{ action}) \sqsubseteq [Q_1 \vdash Q_2 \text{ trace refs} \mid Q_3 \text{ trace}]_C$
using *assms*
by (*simp add: mk-CRD-def, rule-tac sdes-tri-refine-intro, rel-auto+*)

lemma *CRD-refine-rdes*:

assumes

‘ $[P_1]_{S<} \Rightarrow Q_1$ ’

$([P_2 \ x \ r]_{S<} \llbracket x \rightarrow \&tt \rrbracket [r \rightarrow \$ref']) \sqsubseteq ([P_1]_{S<} \wedge Q_2)$

$[P_3 \ x]_{S'} \llbracket x \rightarrow \&tt \rrbracket \sqsubseteq ([P_1]_{S<} \wedge Q_3)$

shows $([P_1 \vdash P_2 \text{ trace refs} \mid P_3 \text{ trace}]_C :: ('e, 's) \text{ action}) \sqsubseteq$

$\mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3)$

using *assms*

by (*simp add: mk-CRD-def, rule-tac srdes-tri-refine-intro, rel-auto+*)

lemma *CRD-refine-rdes'*:

assumes

$Q_2 \text{ is } RR$

$Q_3 \text{ is } RR$

‘ $[P_1]_{S<} \Rightarrow Q_1$ ’

$\bigwedge t. ([P_2 \ t \ r]_{S<} \llbracket r \rightarrow \$ref' \rrbracket) \sqsubseteq ([P_1]_{S<} \wedge Q_2 \llbracket \langle \rangle, \ll t \gg / \$tr, \$tr' \rrbracket)$

$\bigwedge t. [P_3 \ t]_{S'} \sqsubseteq ([P_1]_{S<} \wedge Q_3 \llbracket \langle \rangle, \ll t \gg / \$tr, \$tr' \rrbracket)$

shows $([P_1 \vdash P_2 \text{ trace refs} \mid P_3 \text{ trace}]_C :: ('e, 's) \text{ action}) \sqsubseteq$

$\mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3)$

proof (*simp add: mk-CRD-def, rule srdes-tri-refine-intro*)

show ‘ $[P_1]_{S<} \Rightarrow Q_1$ ’ **by** (*fact assms(3)*)

have $\bigwedge t. ([P_2 \ t \ r]_{S<} \llbracket r \rightarrow \$ref' \rrbracket) \sqsubseteq ([P_1]_{S<} \wedge (RR \ Q_2) \llbracket \langle \rangle, \ll t \gg / \$tr, \$tr' \rrbracket)$

by (*simp add: assms Healthy-if*)

hence ‘ $[P_1]_{S<} \wedge RR(Q_2) \Rightarrow [P_2 \ x \ r]_{S<} \llbracket x \rightarrow \&tt \rrbracket [r \rightarrow \$ref']$ ’

by (*rel-simp; meson*)

thus ‘ $[P_1]_{S<} \wedge Q_2 \Rightarrow [P_2 \ x \ r]_{S<} \llbracket x \rightarrow \&tt \rrbracket [r \rightarrow \$ref']$ ’

by (*simp add: Healthy-if assms*)

have $\bigwedge t. [P_3 \ t]_{S'} \sqsubseteq ([P_1]_{S<} \wedge (RR \ Q_3) \llbracket \langle \rangle, \ll t \gg / \$tr, \$tr' \rrbracket)$

by (*simp add: assms Healthy-if*)

hence ‘ $[P_1]_{S<} \wedge (RR \ Q_3) \Rightarrow [P_3 \ x]_{S'} \llbracket x \rightarrow \&tt \rrbracket$ ’

by (*rel-simp; meson*)

thus ‘ $[P_1]_{S<} \wedge Q_3 \Rightarrow [P_3 \ x]_{S'} \llbracket x \rightarrow \&tt \rrbracket$ ’

by (*simp add: Healthy-if assms*)

qed

end

7 External Choice

theory *utp-sfrd-extchoice*

imports

utp-sfrd-healths

utp-sfrd-rel

begin

7.1 Definitions and syntax

definition *ExtChoice* ::

$('\sigma, ' \varphi) \text{ action set} \Rightarrow (' \sigma, ' \varphi) \text{ action}$ **where**

$[upred-defs]: \text{ExtChoice } A = \mathbf{R}_s((\bigsqcup P \in A \cdot pre_R(P)) \vdash ((\bigsqcup P \in A \cdot cmt_R(P)) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (\prod P \in A \cdot cmt_R(P))))$

syntax

-*ExtChoice* :: *pttrn* \Rightarrow '*a set* \Rightarrow '*b* \Rightarrow '*b* $((\exists \square \text{ -- } \cdot / \text{ --}) [0, 0, 10] 10)$
-*ExtChoice-simp* :: *pttrn* \Rightarrow '*b* \Rightarrow '*b* $((\exists \square \text{ -- } \cdot / \text{ --}) [0, 10] 10)$

translations

$\square P \in A \cdot B \Rightarrow \text{CONST } \text{ExtChoice } ((\lambda P. B) \text{ ' } A)$
 $\square P \cdot B \Rightarrow \text{CONST } \text{ExtChoice } (\text{CONST } \text{range } (\lambda P. B))$

definition *extChoice* ::

$(\text{'}\sigma, \text{'}\varphi) \text{ action} \Rightarrow (\text{'}\sigma, \text{'}\varphi) \text{ action} \Rightarrow (\text{'}\sigma, \text{'}\varphi) \text{ action}$ (**infixl** \square 59) **where**
[*upred-defs*]: $P \square Q \equiv \text{ExtChoice } \{P, Q\}$

Small external choice as an indexed big external choice.

lemma *extChoice-alt-def*:

$P \square Q = (\square i :: \text{nat} \in \{0, 1\} \cdot P \triangleleft \ll i = 0 \gg \triangleright Q)$
by (*simp add: extChoice-def ExtChoice-def*)

7.2 Basic laws

7.3 Algebraic laws

lemma *ExtChoice-empty*: $\text{ExtChoice } \{\} = \text{Stop}$

by (*simp add: ExtChoice-def cond-def Stop-def*)

lemma *ExtChoice-single*:

$P \text{ is CSP} \Rightarrow \text{ExtChoice } \{P\} = P$
by (*simp add: ExtChoice-def usup-and uinf-or SRD-reactive-design-alt*)

7.4 Reactive design calculations

lemma *ExtChoice-rdes*:

assumes $\bigwedge i. \$ok' \nmid P(i) \ A \neq \{\}$

shows $(\square i \in A \cdot \mathbf{R}_s(P(i) \vdash Q(i))) = \mathbf{R}_s((\square i \in A \cdot P(i)) \vdash ((\square i \in A \cdot Q(i)) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (\square i \in A \cdot Q(i))))$

proof –

have $(\square i \in A \cdot \mathbf{R}_s(P(i) \vdash Q(i))) =$
 $\mathbf{R}_s((\square i \in A \cdot \text{pre}_R(\mathbf{R}_s(P(i) \vdash Q(i)))) \vdash$
 $((\square i \in A \cdot \text{cmt}_R(\mathbf{R}_s(P(i) \vdash Q(i))))$
 $\triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright$
 $(\square i \in A \cdot \text{cmt}_R(\mathbf{R}_s(P(i) \vdash Q(i))))))$

by (*simp add: ExtChoice-def*)

also have ... =

$\mathbf{R}_s((\square i \in A \cdot R1(R2c(\text{pre}_s \uparrow P(i)))) \vdash$
 $((\square i \in A \cdot R1(R2c(\text{cmt}_s \uparrow (P(i) \Rightarrow Q(i))))$
 $\triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright$
 $(\square i \in A \cdot R1(R2c(\text{cmt}_s \uparrow (P(i) \Rightarrow Q(i))))))$

by (*simp add: rea-pre-RHS-design rea-cmt-RHS-design*)

also have ... =

$\mathbf{R}_s((\square i \in A \cdot R1(R2c(\text{pre}_s \uparrow P(i)))) \vdash$
 $R1(R2c$
 $((\square i \in A \cdot R1(R2c(\text{cmt}_s \uparrow (P(i) \Rightarrow Q(i))))$
 $\triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright$
 $(\square i \in A \cdot R1(R2c(\text{cmt}_s \uparrow (P(i) \Rightarrow Q(i))))))$

by (*metis (no-types, lifting) RHS-design-export-R1 RHS-design-export-R2c*)

also have ... =

$\mathbf{R}_s((\square i \in A \cdot R1(R2c(\text{pre}_s \uparrow P(i)))) \vdash$

$R1(R2c$
 $((\sqcup i \in A \cdot (cmt_s \dagger (P(i) \Rightarrow Q(i))))$
 $\triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright$
 $(\sqcap i \in A \cdot (cmt_s \dagger (P(i) \Rightarrow Q(i))))))$
by (*simp add: R2c-UNIF R2c-condr R1-cond R1-idem R1-R2c-commute R2c-idem R1-UNIF assms*
R1-USUP R2c-USUP)
also have ... =
 $\mathbf{R}_s ((\sqcup i \in A \cdot R1 (R2c (pre_s \dagger P(i)))) \vdash$
 $cmt_s \dagger$
 $((\sqcup i \in A \cdot (cmt_s \dagger (P(i) \Rightarrow Q(i))))$
 $\triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright$
 $(\sqcap i \in A \cdot (cmt_s \dagger (P(i) \Rightarrow Q(i))))))$
by (*metis (no-types, lifting) RHS-design-export-R1 RHS-design-export-R2c rdes-export-cmt*)
also have ... =
 $\mathbf{R}_s ((\sqcup i \in A \cdot R1 (R2c (pre_s \dagger P(i)))) \vdash$
 $cmt_s \dagger$
 $((\sqcup i \in A \cdot (P(i) \Rightarrow Q(i)))$
 $\triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright$
 $(\sqcap i \in A \cdot (P(i) \Rightarrow Q(i))))$
by (*simp add: usubst*)
also have ... =
 $\mathbf{R}_s ((\sqcup i \in A \cdot R1 (R2c (pre_s \dagger P(i)))) \vdash$
 $((\sqcup i \in A \cdot (P(i) \Rightarrow Q(i))) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (\sqcap i \in A \cdot (P(i) \Rightarrow Q(i))))$
by (*simp add: rdes-export-cmt*)
also have ... =
 $\mathbf{R}_s ((R1(R2c(\sqcup i \in A \cdot (pre_s \dagger P(i)))) \vdash$
 $((\sqcup i \in A \cdot (P(i) \Rightarrow Q(i))) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (\sqcap i \in A \cdot (P(i) \Rightarrow Q(i))))$
by (*simp add: not-UNIF R1-UNIF R2c-UNIF assms*)
also have ... =
 $\mathbf{R}_s ((R2c(\sqcup i \in A \cdot (pre_s \dagger P(i)))) \vdash$
 $((\sqcup i \in A \cdot (P(i) \Rightarrow Q(i))) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (\sqcap i \in A \cdot (P(i) \Rightarrow Q(i))))$
by (*simp add: R1-design-R1-pre*)
also have ... =
 $\mathbf{R}_s (((\sqcup i \in A \cdot (pre_s \dagger P(i)))) \vdash$
 $((\sqcup i \in A \cdot (P(i) \Rightarrow Q(i))) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (\sqcap i \in A \cdot (P(i) \Rightarrow Q(i))))$
by (*metis (no-types, lifting) RHS-design-R2c-pre*)
also have ... =
 $\mathbf{R}_s (([\$ok \mapsto_s true, \$wait \mapsto_s false] \dagger (\sqcup i \in A \cdot P(i))) \vdash$
 $((\sqcup i \in A \cdot (P(i) \Rightarrow Q(i))) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (\sqcap i \in A \cdot (P(i) \Rightarrow Q(i))))$
proof –
from *assms* **have** $\bigwedge i. pre_s \dagger P(i) = [\$ok \mapsto_s true, \$wait \mapsto_s false] \dagger P(i)$
by (*rel-auto*)
thus *?thesis*
by (*simp add: usubst*)
qed
also have ... =
 $\mathbf{R}_s ((\sqcup i \in A \cdot P(i)) \vdash ((\sqcup i \in A \cdot (P(i) \Rightarrow Q(i))) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (\sqcap i \in A \cdot (P(i) \Rightarrow$
 $Q(i))))))$
by (*simp add: rdes-export-pre not-UNIF*)
also have ... = $\mathbf{R}_s ((\sqcup i \in A \cdot P(i)) \vdash ((\sqcup i \in A \cdot Q(i)) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (\sqcap i \in A \cdot Q(i))))$
by (*rule cong[of $\mathbf{R}_s \mathbf{R}_s$], simp, rel-auto, blast+*)

finally show *?thesis* .
qed

lemma *ExtChoice-tri-rdes*:

assumes $\bigwedge i . \$ok' \# P_1(i) \ A \neq \{\}$

shows $(\Box_{i \in A} \cdot \mathbf{R}_s(P_1(i) \vdash P_2(i) \diamond P_3(i))) =$

$\mathbf{R}_s((\Box_{i \in A} \cdot P_1(i)) \vdash (((\Box_{i \in A} \cdot P_2(i)) \triangleleft \$tr' =_u \$tr \triangleright (\Box_{i \in A} \cdot P_2(i))) \diamond (\Box_{i \in A} \cdot P_3(i))))$

proof –

have $(\Box_{i \in A} \cdot \mathbf{R}_s(P_1(i) \vdash P_2(i) \diamond P_3(i))) =$

$\mathbf{R}_s((\Box_{i \in A} \cdot P_1(i)) \vdash ((\Box_{i \in A} \cdot P_2(i) \diamond P_3(i)) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (\Box_{i \in A} \cdot P_2(i) \diamond P_3(i))))$

by (*simp add: ExtChoice-rdes assms*)

also

have ... =

$\mathbf{R}_s((\Box_{i \in A} \cdot P_1(i)) \vdash ((\Box_{i \in A} \cdot P_2(i) \diamond P_3(i)) \triangleleft \$wait' \wedge \$tr' =_u \$tr \triangleright (\Box_{i \in A} \cdot P_2(i) \diamond P_3(i))))$

by (*simp add: conj-comm*)

also

have ... =

$\mathbf{R}_s((\Box_{i \in A} \cdot P_1(i)) \vdash (((\Box_{i \in A} \cdot P_2(i) \diamond P_3(i)) \triangleleft \$tr' =_u \$tr \triangleright (\Box_{i \in A} \cdot P_2(i) \diamond P_3(i))) \diamond (\Box_{i \in A} \cdot P_2(i) \diamond P_3(i))))$

by (*simp add: cond-conj wait'-cond-def*)

also

have ... = $\mathbf{R}_s((\Box_{i \in A} \cdot P_1(i)) \vdash (((\Box_{i \in A} \cdot P_2(i)) \triangleleft \$tr' =_u \$tr \triangleright (\Box_{i \in A} \cdot P_2(i))) \diamond (\Box_{i \in A} \cdot P_3(i))))$

by (*rule cong[of $\mathbf{R}_s \ \mathbf{R}_s$], simp, rel-auto*)

finally show *?thesis* .

qed

lemma *ExtChoice-tri-rdes' [rdes-def]*:

assumes $\bigwedge i . \$ok' \# P_1(i) \ A \neq \{\}$

shows $(\Box_{i \in A} \cdot \mathbf{R}_s(P_1(i) \vdash P_2(i) \diamond P_3(i))) =$

$\mathbf{R}_s((\Box_{i \in A} \cdot P_1(i)) \vdash (((\Box_{i \in A} \cdot R5(P_2(i))) \vee (\Box_{i \in A} \cdot R4(P_2(i)))) \diamond (\Box_{i \in A} \cdot P_3(i))))$

by (*simp add: ExtChoice-tri-rdes assms, rel-auto, simp-all add: less-le assms*)

lemma *ExtChoice-tri-rdes-def [rdes-def]*:

assumes $A \subseteq \llbracket CSP \rrbracket_H$

shows $ExtChoice \ A = \mathbf{R}_s((\Box_{P \in A} \cdot pre_R \ P) \vdash (((\Box_{P \in A} \cdot peri_R \ P) \triangleleft \$tr' =_u \$tr \triangleright (\Box_{P \in A} \cdot peri_R \ P)) \diamond (\Box_{P \in A} \cdot post_R \ P)))$

proof –

have $((\Box_{P \in A} \cdot cmt_R \ P) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (\Box_{P \in A} \cdot cmt_R \ P)) =$

$((\Box_{P \in A} \cdot cmt_R \ P) \triangleleft \$tr' =_u \$tr \triangleright (\Box_{P \in A} \cdot cmt_R \ P)) \diamond (\Box_{P \in A} \cdot cmt_R \ P)$

by (*rel-auto*)

also have ... = $((\Box_{P \in A} \cdot peri_R \ P) \triangleleft \$tr' =_u \$tr \triangleright (\Box_{P \in A} \cdot peri_R \ P)) \diamond (\Box_{P \in A} \cdot post_R \ P)$

by (*rel-auto*)

finally show *?thesis*

by (*simp add: ExtChoice-def*)

qed

lemma *extChoice-rdes*:

assumes $\$ok' \# P_1 \ \$ok' \# Q_1$

shows $\mathbf{R}_s(P_1 \vdash P_2) \Box \mathbf{R}_s(Q_1 \vdash Q_2) = \mathbf{R}_s((P_1 \wedge Q_1) \vdash ((P_2 \wedge Q_2) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (P_2 \vee Q_2)))$

proof –

have $(\Box_{i::nat \in \{0, 1\}} \cdot \mathbf{R}_s(P_1 \vdash P_2) \triangleleft \ll i = 0 \gg \triangleright \mathbf{R}_s(Q_1 \vdash Q_2)) = (\Box_{i::nat \in \{0, 1\}} \cdot \mathbf{R}_s((P_1 \vdash P_2) \triangleleft \ll i = 0 \gg \triangleright (Q_1 \vdash Q_2)))$

by (*simp only: RHS-cond R2c-lit*)

also have ... = $(\Box i :: nat \in \{0, 1\} \cdot \mathbf{R}_s ((P_1 \triangleleft \ll i = 0 \gg \triangleright Q_1) \vdash (P_2 \triangleleft \ll i = 0 \gg \triangleright Q_2)))$
 by (simp add: design-condr)
 also have ... = $\mathbf{R}_s ((P_1 \wedge Q_1) \vdash ((P_2 \wedge Q_2) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (P_2 \vee Q_2)))$
 by (subst ExtChoice-rdes, simp-all add: assms unrest uinf-or usup-and)
 finally show ?thesis by (simp add: extChoice-alt-def)
 qed

lemma extChoice-tri-rdes:

assumes $\$ok' \# P_1 \$ok' \# Q_1$

shows $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \Box \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3) =$

$\mathbf{R}_s ((P_1 \wedge Q_1) \vdash (((P_2 \wedge Q_2) \triangleleft \$tr' =_u \$tr \triangleright (P_2 \vee Q_2)) \diamond (P_3 \vee Q_3)))$

proof –

have $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \Box \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3) =$

$\mathbf{R}_s ((P_1 \wedge Q_1) \vdash ((P_2 \diamond P_3 \wedge Q_2 \diamond Q_3) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (P_2 \diamond P_3 \vee Q_2 \diamond Q_3)))$

by (simp add: extChoice-rdes assms)

also

have ... = $\mathbf{R}_s ((P_1 \wedge Q_1) \vdash ((P_2 \diamond P_3 \wedge Q_2 \diamond Q_3) \triangleleft \$wait' \wedge \$tr' =_u \$tr \triangleright (P_2 \diamond P_3 \vee Q_2 \diamond Q_3)))$

by (simp add: conj-comm)

also

have ... = $\mathbf{R}_s ((P_1 \wedge Q_1) \vdash$

$((P_2 \diamond P_3 \wedge Q_2 \diamond Q_3) \triangleleft \$tr' =_u \$tr \triangleright (P_2 \diamond P_3 \vee Q_2 \diamond Q_3)) \diamond (P_2 \diamond P_3 \vee Q_2 \diamond Q_3)))$

by (simp add: cond-conj wait'-cond-def)

also

have ... = $\mathbf{R}_s ((P_1 \wedge Q_1) \vdash (((P_2 \wedge Q_2) \triangleleft \$tr' =_u \$tr \triangleright (P_2 \vee Q_2)) \diamond (P_3 \vee Q_3)))$

by (rule cong[of $\mathbf{R}_s \mathbf{R}_s$], simp, rel-auto)

finally show ?thesis .

qed

lemma extChoice-rdes-def:

assumes P_1 is RR Q_1 is RR

shows $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \Box \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3) =$

$\mathbf{R}_s ((P_1 \wedge Q_1) \vdash (((P_2 \wedge Q_2) \triangleleft \$tr' =_u \$tr \triangleright (P_2 \vee Q_2)) \diamond (P_3 \vee Q_3)))$

by (subst extChoice-tri-rdes, simp-all add: assms unrest)

lemma extChoice-rdes-def' [rdes-def]:

assumes P_1 is RR Q_1 is RR

shows $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \Box \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3) =$

$\mathbf{R}_s ((P_1 \wedge Q_1) \vdash ((R5(P_2 \wedge Q_2) \vee R4(P_2 \vee Q_2)) \diamond (P_3 \vee Q_3)))$

by (simp add: extChoice-rdes-def assms, rel-auto, simp-all add: less-le)

lemma CSP-ExtChoice [closure]:

ExtChoice A is CSP

by (simp add: ExtChoice-def RHS-design-is-SRD unrest)

lemma CSP-extChoice [closure]:

$P \Box Q$ is CSP

by (simp add: CSP-ExtChoice extChoice-def)

lemma preR-ExtChoice [rdes]:

assumes $A \neq \{\}$ $A \subseteq \llbracket CSP \rrbracket_H$

shows $pre_R(ExtChoice A) = (\bigsqcup P \in A \cdot pre_R(P))$

proof –

have $pre_R(ExtChoice A) = (R1 (R2c ((\bigsqcup P \in A \cdot pre_R P))))$

by (simp add: ExtChoice-def rea-pre-RHS-design usubst unrest)

also from assms have ... = $(R1 (R2c (\bigsqcup P \in A \cdot (pre_R(CSP(P)))))$

by (*metis USUP-healthy*)
 also from *assms* have ... = $(\bigsqcup P \in A \cdot (\text{pre}_R(\text{CSP}(P))))$
 by (*rel-auto*)
 also from *assms* have ... = $(\bigsqcup P \in A \cdot (\text{pre}_R(P)))$
 by (*metis USUP-healthy*)
 finally show *?thesis* .
 qed

lemma *preR-ExtChoice-ind* [*rdes*]:

assumes $A \neq \{\}$ \wedge $P. P \in A \implies F(P)$ is CSP
 shows $\text{pre}_R(\bigsqcup P \in A \cdot F(P)) = (\bigsqcup P \in A \cdot \text{pre}_R(F(P)))$
 using *assms* by (*subst preR-ExtChoice, auto*)

lemma *periR-ExtChoice* [*rdes*]:

assumes $A \subseteq \llbracket \text{NCSP} \rrbracket_H$ $A \neq \{\}$
 shows $\text{peri}_R(\text{ExtChoice } A) = ((\bigsqcup P \in A \cdot \text{pre}_R(P)) \Rightarrow_r (\bigsqcup P \in A \cdot \text{peri}_R P)) \triangleleft \$tr' =_u \$tr \triangleright (\bigsqcap P \in A \cdot \text{peri}_R P)$

proof –

have $\text{peri}_R(\text{ExtChoice } A) = \text{peri}_R(\mathbf{R}_s((\bigsqcup P \in A \cdot \text{pre}_R P)) \vdash$
 $(\bigsqcup P \in A \cdot \text{peri}_R P) \triangleleft \$tr' =_u \$tr \triangleright (\bigsqcap P \in A \cdot \text{peri}_R P)) \diamond$
 $(\bigsqcap P \in A \cdot \text{post}_R P)))$
 by (*simp add: ExtChoice-tri-rdes-def assms closure*)

also have ... = $\text{peri}_R(\mathbf{R}_s((\bigsqcup P \in A \cdot \text{pre}_R(\text{NCSP } P)) \vdash$
 $(\bigsqcup P \in A \cdot \text{peri}_R(\text{NCSP } P)) \triangleleft \$tr' =_u \$tr \triangleright (\bigsqcap P \in A \cdot \text{peri}_R(\text{NCSP } P))) \diamond$
 $(\bigsqcap P \in A \cdot \text{post}_R P)))$
 by (*simp add: UINF-healthy[OF assms(1), THEN sym] USUP-healthy[OF assms(1), THEN sym]*)

also have ... = $R1(R2c((\bigsqcup P \in A \cdot \text{pre}_R(\text{NCSP } P)) \Rightarrow_r$
 $(\bigsqcup P \in A \cdot \text{peri}_R(\text{NCSP } P))$
 $\triangleleft \$tr' =_u \$tr \triangleright$
 $(\bigsqcap P \in A \cdot \text{peri}_R(\text{NCSP } P))))$

proof –

have $(\bigsqcup P \in A \cdot [\$ok \mapsto_s \text{true}, \$ok' \mapsto_s \text{true}, \$wait \mapsto_s \text{false}, \$wait' \mapsto_s \text{true}] \dagger \text{pre}_R(\text{NCSP } P))$
 = $(\bigsqcup P \in A \cdot \text{pre}_R(\text{NCSP } P))$
 by (*rule USUP-cong, simp add: closure usubst unrest assms*)
 thus *?thesis*
 by (*simp add: rea-peri-RHS-design Healthy-Idempotent SRD-Idempotent usubst unrest assms*)

qed

also have ... = $R1((\bigsqcup P \in A \cdot \text{pre}_R(\text{NCSP } P)) \Rightarrow_r$
 $(\bigsqcup P \in A \cdot \text{peri}_R(\text{NCSP } P))$
 $\triangleleft \$tr' =_u \$tr \triangleright$
 $(\bigsqcap P \in A \cdot \text{peri}_R(\text{NCSP } P)))$

by (*simp add: R2c-rea-impl R2c-condr R2c-UINF R2c-preR R2c-periR R2c-tr'-minus-tr R2c-USUP closure*)

also have ... = $((\bigsqcup P \in A \cdot \text{pre}_R(\text{NCSP } P)) \Rightarrow_r (\bigsqcup P \in A \cdot \text{peri}_R(\text{NCSP } P)))$
 $\triangleleft \$tr' =_u \$tr \triangleright$
 $((\bigsqcup P \in A \cdot \text{pre}_R(\text{NCSP } P)) \Rightarrow_r (\bigsqcap P \in A \cdot \text{peri}_R(\text{NCSP } P)))$

by (*simp add: R1-rea-impl R1-cond R1-USUP R1-UINF assms Healthy-if closure, rel-auto*)

also have ... = $((\bigsqcup P \in A \cdot \text{pre}_R(\text{NCSP } P)) \Rightarrow_r (\bigsqcup P \in A \cdot \text{peri}_R(\text{NCSP } P)))$
 $\triangleleft \$tr' =_u \$tr \triangleright$
 $((\bigsqcap P \in A \cdot \text{pre}_R(\text{NCSP } P)) \Rightarrow_r \text{peri}_R(\text{NCSP } P)))$

by (*simp add: UINF-rea-impl[THEN sym]*)

also have ... = $((\bigsqcup P \in A \cdot \text{pre}_R(\text{NCSP } P)) \Rightarrow_r (\bigsqcup P \in A \cdot \text{peri}_R(\text{NCSP } P)))$
 $\triangleleft \$tr' =_u \$tr \triangleright$

$((\prod P \in A \cdot \text{peri}_R (\text{NCSP } P)))$
 by (simp add: SRD-*peri-under-pre* closure assms unrest)
 also have ... = $((\prod P \in A \cdot \text{pre}_R P) \Rightarrow_r (\prod P \in A \cdot \text{peri}_R P))$
 $\triangleleft \$tr' =_u \$tr \triangleright$
 $((\prod P \in A \cdot \text{peri}_R P))$
 by (simp add: UINF-healthy[OF assms(1), THEN sym] USUP-healthy[OF assms(1), THEN sym])
 finally show ?thesis .
 qed

lemma *periR-ExtChoice'*:

assumes $A \subseteq \llbracket \text{NCSP} \rrbracket_H A \neq \{\}$
 shows $\text{peri}_R(\text{ExtChoice } A) = (R5((\prod P \in A \cdot \text{pre}_R(P)) \Rightarrow_r (\prod P \in A \cdot \text{peri}_R P)) \vee (\prod P \in A \cdot R4(\text{peri}_R P)))$
 using assms(2)
 by (simp add: *periR-ExtChoice* assms(1), rel-auto)

lemma *periR-ExtChoice-ind* [rdes]:

assumes $\bigwedge P. P \in A \implies F(P)$ is NCSP $A \neq \{\}$
 shows $\text{peri}_R(\prod P \in A \cdot F(P)) = ((\prod P \in A \cdot \text{pre}_R(F P)) \Rightarrow_r (\prod P \in A \cdot \text{peri}_R (F P))) \triangleleft \$tr' =_u \$tr \triangleright (\prod P \in A \cdot \text{peri}_R (F P))$
 using assms by (subst *periR-ExtChoice*, auto simp add: closure unrest)

lemma *periR-ExtChoice-ind'*:

assumes $\bigwedge P. P \in A \implies F(P)$ is NCSP $A \neq \{\}$
 shows $\text{peri}_R(\prod P \in A \cdot F(P)) = (R5((\prod P \in A \cdot \text{pre}_R(F P)) \Rightarrow_r (\prod P \in A \cdot \text{peri}_R (F P))) \vee (\prod P \in A \cdot R4(\text{peri}_R (F P))))$
 using assms by (subst *periR-ExtChoice'*, auto simp add: closure unrest)

lemma *postR-ExtChoice* [rdes]:

assumes $A \subseteq \llbracket \text{NCSP} \rrbracket_H A \neq \{\}$
 shows $\text{post}_R(\text{ExtChoice } A) = (\prod P \in A \cdot \text{post}_R P)$

proof –

have $\text{post}_R (\text{ExtChoice } A) = \text{post}_R (\mathbf{R}_s ((\prod P \in A \cdot \text{pre}_R P) \vdash ((\prod P \in A \cdot \text{peri}_R P) \triangleleft \$tr' =_u \$tr \triangleright (\prod P \in A \cdot \text{peri}_R P)) \diamond (\prod P \in A \cdot \text{post}_R P)))$
 by (simp add: *ExtChoice-tri-rdes-def* closure assms)

also have ... = $\text{post}_R (\mathbf{R}_s ((\prod P \in A \cdot \text{pre}_R (\text{NCSP } P)) \vdash ((\prod P \in A \cdot \text{peri}_R P) \triangleleft \$tr' =_u \$tr \triangleright (\prod P \in A \cdot \text{peri}_R P)) \diamond (\prod P \in A \cdot \text{post}_R (\text{NCSP } P))))$
 by (simp add: UINF-healthy[OF assms(1), THEN sym] USUP-healthy[OF assms(1), THEN sym])

also have ... = $R1 (R2c ((\prod P \in A \cdot \text{pre}_R (\text{NCSP } P)) \Rightarrow_r (\prod P \in A \cdot \text{post}_R (\text{NCSP } P))))$

proof –

have $(\prod P \in A \cdot [\$ok \mapsto_s \text{true}, \$ok' \mapsto_s \text{true}, \$wait \mapsto_s \text{false}, \$wait' \mapsto_s \text{false}] \dagger \text{pre}_R (\text{NCSP } P)) = (\prod P \in A \cdot \text{pre}_R (\text{NCSP } P))$
 by (rule USUP-cong, simp add: usubst closure unrest assms)
 thus ?thesis
 by (simp add: rea-post-RHS-design Healthy-Idempotent SRD-Idempotent usubst unrest assms)

qed

also have ... = $R1 ((\prod P \in A \cdot \text{pre}_R (\text{NCSP } P)) \Rightarrow_r (\prod P \in A \cdot \text{post}_R (\text{NCSP } P)))$

by (simp add: R2c-rea-impl R2c-condr R2c-UINF R2c-preR R2c-postR
 R2c-tr'-minus-tr R2c-USUP closure)

also from assms(2) have ... = $((\prod P \in A \cdot \text{pre}_R (\text{NCSP } P)) \Rightarrow_r (\prod P \in A \cdot \text{post}_R (\text{NCSP } P)))$

by (simp add: R1-rea-impl R1-cond R1-USUP R1-UINF assms Healthy-if closure)

also have ... = $(\bigwedge P \in A \cdot \text{pre}_R(\text{NCSP } P) \Rightarrow_r \text{post}_R(\text{NCSP } P))$
 by (simp add: UINF-rea-impl)
 also have ... = $(\bigwedge P \in A \cdot \text{post}_R(\text{NCSP } P))$
 by (simp add: SRD-post-under-pre closure assms unrest)
 finally show ?thesis
 by (simp add: UINF-healthy[OF assms(1), THEN sym] USUP-healthy[OF assms(1), THEN sym])
 qed

lemma *postR-ExtChoice-ind* [rdes]:
 assumes $\bigwedge P. P \in A \implies F(P) \text{ is NCSP } A \neq \{\}$
 shows $\text{post}_R(\bigwedge P \in A \cdot F(P)) = (\bigwedge P \in A \cdot \text{post}_R(F(P)))$
 using assms by (subst postR-ExtChoice, auto simp add: closure unrest)

lemma *preR-extChoice*:
 assumes $P \text{ is CSP } Q \text{ is CSP } \$wait' \# \text{pre}_R(P) \$wait' \# \text{pre}_R(Q)$
 shows $\text{pre}_R(P \sqcap Q) = (\text{pre}_R(P) \wedge \text{pre}_R(Q))$
 by (simp add: extChoice-def preR-ExtChoice assms usup-and)

lemma *preR-extChoice'* [rdes]:
 assumes $P \text{ is NCSP } Q \text{ is NCSP}$
 shows $\text{pre}_R(P \sqcap Q) = (\text{pre}_R(P) \wedge \text{pre}_R(Q))$
 by (simp add: preR-extChoice closure assms unrest)

lemma *periR-extChoice* [rdes]:
 assumes $P \text{ is NCSP } Q \text{ is NCSP}$
 shows $\text{peri}_R(P \sqcap Q) = ((\text{pre}_R(P) \wedge \text{pre}_R(Q) \Rightarrow_r \text{peri}_R(P) \wedge \text{peri}_R(Q)) \triangleleft \$tr' =_u \$tr \triangleright (\text{peri}_R(P) \vee \text{peri}_R(Q)))$
 using assms
 by (simp add: extChoice-def, subst periR-ExtChoice, auto simp add: usup-and uinf-or)

lemma *postR-extChoice* [rdes]:
 assumes $P \text{ is NCSP } Q \text{ is NCSP}$
 shows $\text{post}_R(P \sqcap Q) = (\text{post}_R(P) \vee \text{post}_R(Q))$
 using assms
 by (simp add: extChoice-def, subst postR-ExtChoice, auto simp add: usup-and uinf-or)

lemma *ExtChoice-cong*:
 assumes $\bigwedge P. P \in A \implies F(P) = G(P)$
 shows $(\bigwedge P \in A \cdot F(P)) = (\bigwedge P \in A \cdot G(P))$
 using assms image-cong by force

lemma *ref-unrest-ExtChoice*:
 assumes
 $\bigwedge P. P \in A \implies \$ref \# \text{pre}_R(P)$
 $\bigwedge P. P \in A \implies \$ref \# \text{cmt}_R(P)$
 shows $\$ref \# (\text{ExtChoice } A) \llbracket \text{false} / \$wait \rrbracket$
proof –
 have $\bigwedge P. P \in A \implies \$ref \# \text{pre}_R(P \llbracket 0 / \$tr \rrbracket)$
 using assms by (rel-blast)
 with assms show ?thesis
 by (simp add: ExtChoice-def RHS-def R1-def R2c-def R2s-def R3h-def design-def usubst unrest)
 qed

lemma *CSP4-ExtChoice*:
 assumes $A \subseteq \llbracket \text{NCSP} \rrbracket_H$

shows *ExtChoice A is CSP4*
proof (*cases A = {}*)
 case *True* **thus** *?thesis*
 by (*simp add: ExtChoice-empty Healthy-def CSP4-def, simp add: Skip-is-CSP Stop-left-zero*)
next
 case *False*
 have 1: $(\neg_r (\neg_r \text{pre}_R (\text{ExtChoice } A)) \text{ ;; }_h R1 \text{ true}) = \text{pre}_R (\text{ExtChoice } A)$
 proof –
 have $\bigwedge P. P \in A \implies (\neg_r \text{pre}_R(P)) \text{ ;; } R1 \text{ true} = (\neg_r \text{pre}_R(P))$
 by (*simp add: NCSP-Healthy-subset-member NCSP-implies-NSRD NSRD-neg-pre-unit assms*)
 thus *?thesis*
 apply (*simp add: False preR-ExtChoice closure NCSP-set-unrest-pre-wait' assms not-UINF seq-UINF-distr not-USUP*)
 apply (*rule USUP-cong*)
 apply (*simp add: rpred assms closure*)
 done
 qed
 have 2: $\$st' \# \text{peri}_R (\text{ExtChoice } A)$
 proof –
 have $a: \bigwedge P. P \in A \implies \$st' \# \text{pre}_R(P)$
 by (*simp add: NCSP-Healthy-subset-member NCSP-implies-NSRD NSRD-st'-unrest-pre assms*)
 have $b: \bigwedge P. P \in A \implies \$st' \# \text{peri}_R(P)$
 by (*simp add: NCSP-Healthy-subset-member NCSP-implies-NSRD NSRD-st'-unrest-peri assms*)
 from $a \ b$ **show** *?thesis*
 apply (*subst periR-ExtChoice*)
 apply (*simp-all add: assms closure unrest CSP4-set-unrest-pre-st' NCSP-set-unrest-pre-wait'*
False)
 done
 qed
 have 3: $\$ref' \# \text{post}_R (\text{ExtChoice } A)$
 proof –
 have $a: \bigwedge P. P \in A \implies \$ref' \# \text{pre}_R(P)$
 by (*simp add: CSP4-ref'-unrest-pre CSP-Healthy-subset-member NCSP-Healthy-subset-member NCSP-implies-CSP4 NCSP-subset-implies-CSP assms*)
 have $b: \bigwedge P. P \in A \implies \$ref' \# \text{post}_R(P)$
 by (*simp add: CSP4-ref'-unrest-post CSP-Healthy-subset-member NCSP-Healthy-subset-member NCSP-implies-CSP4 NCSP-subset-implies-CSP assms*)
 from $a \ b$ **show** *?thesis*
 by (*subst postR-ExtChoice, simp-all add: assms CSP4-set-unrest-pre-st' NCSP-set-unrest-pre-wait'*
unrest False)
 qed
 show *?thesis*
 by (*rule CSP4-tri-intro, simp-all add: 1 2 3 assms closure*)
 (*metis 1 R1-seqr-closure rea-not-R1 rea-not-not rea-true-R1*)
qed

lemma *CSP4-extChoice [closure]*:
 assumes *P is NCSP Q is NCSP*
 shows $P \sqsubseteq Q$ *is CSP4*
 by (*simp add: extChoice-def, rule CSP4-ExtChoice, simp-all add: assms*)

lemma *NCSP-ExtChoice [closure]*:
 assumes $A \subseteq \llbracket \text{NCSP} \rrbracket_H$
 shows *ExtChoice A is NCSP*
proof (*cases A = {}*)


```

case True
then show ?thesis by (simp add: ExtChoice-empty closure)
next
case False
show ?thesis
proof (rule NCSP-intro)
  from assms have cls:  $A \subseteq \llbracket CSP \rrbracket_H$   $A \subseteq \llbracket CSP^3 \rrbracket_H$   $A \subseteq \llbracket CSP^4 \rrbracket_H$ 
  using NCSP-implies-CSP NCSP-implies-CSP3 NCSP-implies-CSP4 by blast+
  have wu:  $\bigwedge P. P \in A \implies \$wait' \nmid pre_R(P)$ 
  using NCSP-implies-NSRD NSRD-wait'-unrest-pre assms by force
  show 1:ExtChoice A is CSP
  by (metis (mono-tags) Ball-Collect CSP-ExtChoice NCSP-implies-CSP assms)
  from cls show ExtChoice A is CSP3
  by (rule-tac CSP3-SRD-intro, simp-all add: CSP-Healthy-subset-member CSP3-Healthy-subset-member
  closure rdes unrest wu assms 1 False)
  from cls show ExtChoice A is CSP4
  by (simp add: CSP4-ExtChoice assms)
qed
qed

```

lemma *ExtChoice-NCSP-closed* [closure]:
 assumes $\bigwedge i. i \in I \implies P(i)$ is NCSP
 shows $(\bigwedge i \in I. P(i))$ is NCSP
 by (simp add: NCSP-ExtChoice assms image-subset-iff)

lemma *NCSP-extChoice* [closure]:
 assumes P is NCSP Q is NCSP
 shows $P \sqcap Q$ is NCSP
 by (simp add: NCSP-ExtChoice assms extChoice-def)

7.5 Productivity and Guardedness

lemma *Productive-ExtChoice* [closure]:
 assumes $A \neq \{\}$ $A \subseteq \llbracket NCSP \rrbracket_H$ $A \subseteq \llbracket Productive \rrbracket_H$
 shows ExtChoice A is Productive

proof –

have 1: $\bigwedge P. P \in A \implies \$wait' \nmid pre_R(P)$
 using NCSP-implies-NSRD NSRD-wait'-unrest-pre assms(2) by blast
 show ?thesis

proof (rule Productive-intro, simp-all add: assms closure rdes 1 unrest)

have $((\bigwedge P \in A. pre_R P) \wedge (\bigwedge P \in A. post_R P)) =$
 $((\bigwedge P \in A. pre_R P) \wedge (\bigwedge P \in A. (pre_R P \wedge post_R P)))$
 by (rel-auto)

moreover have $(\bigwedge P \in A. (pre_R P \wedge post_R P)) = (\bigwedge P \in A. ((pre_R P \wedge post_R P) \wedge \$tr <_u$
 $\$tr'))$

by (rule UINF-cong, metis (no-types, lifting) 1 Ball-Collect NCSP-implies-CSP Productive-post-refines-tr-increase
 assms utp-pred-laws.inf.absorb1)

ultimately show $(\$tr' >_u \$tr) \sqsubseteq ((\bigwedge P \in A. pre_R P) \wedge (\bigwedge P \in A. post_R P))$
 by (rel-auto)

qed
 qed

lemma *Productive-extChoice* [closure]:
 assumes P is NCSP Q is NCSP P is Productive Q is Productive
 shows $P \sqcap Q$ is Productive

by (simp add: extChoice-def Productive-ExtChoice assms)

lemma *ExtChoice-Guarded* [closure]:

assumes $\bigwedge P. P \in A \implies \text{Guarded } P$

shows *Guarded* $(\lambda X. \Box P \in A \cdot P(X))$

proof (rule *GuardedI*)

fix X n

have $\bigwedge Y. ((\Box P \in A \cdot P \ Y) \wedge \text{gvr}(n+1)) = ((\Box P \in A \cdot (P \ Y \wedge \text{gvr}(n+1))) \wedge \text{gvr}(n+1))$

proof –

fix Y

let $?lhs = ((\Box P \in A \cdot P \ Y) \wedge \text{gvr}(n+1))$ and $?rhs = ((\Box P \in A \cdot (P \ Y \wedge \text{gvr}(n+1))) \wedge \text{gvr}(n+1))$

have $a: ?lhs \llbracket \text{false}/\$ok \rrbracket = ?rhs \llbracket \text{false}/\$ok \rrbracket$

by (rel-auto)

have $b: ?lhs \llbracket \text{true}/\$ok \rrbracket \llbracket \text{true}/\$wait \rrbracket = ?rhs \llbracket \text{true}/\$ok \rrbracket \llbracket \text{true}/\$wait \rrbracket$

by (rel-auto)

have $c: ?lhs \llbracket \text{true}/\$ok \rrbracket \llbracket \text{false}/\$wait \rrbracket = ?rhs \llbracket \text{true}/\$ok \rrbracket \llbracket \text{false}/\$wait \rrbracket$

by (simp add: *ExtChoice-def RHS-def R1-def R2c-def R2s-def R3h-def design-def usubst unrest, rel-blast*)

show $?lhs = ?rhs$

using a b c

by (rule-tac *bool-eq-splitI*[of *in-var ok*], simp, rule-tac *bool-eq-splitI*[of *in-var wait*], simp-all)

qed

moreover have $((\Box P \in A \cdot (P \ X \wedge \text{gvr}(n+1))) \wedge \text{gvr}(n+1)) = ((\Box P \in A \cdot (P \ (X \wedge \text{gvr}(n))) \wedge \text{gvr}(n+1))) \wedge \text{gvr}(n+1))$

proof –

have $(\Box P \in A \cdot (P \ X \wedge \text{gvr}(n+1))) = (\Box P \in A \cdot (P \ (X \wedge \text{gvr}(n)) \wedge \text{gvr}(n+1)))$

proof (rule *ExtChoice-cong*)

fix P assume $P \in A$

thus $(P \ X \wedge \text{gvr}(n+1)) = (P \ (X \wedge \text{gvr}(n)) \wedge \text{gvr}(n+1))$

using *Guarded-def assms* by blast

qed

thus *?thesis* by simp

qed

ultimately show $((\Box P \in A \cdot P \ X) \wedge \text{gvr}(n+1)) = ((\Box P \in A \cdot (P \ (X \wedge \text{gvr}(n)))) \wedge \text{gvr}(n+1))$

by simp

qed

lemma *extChoice-Guarded* [closure]:

assumes *Guarded* P *Guarded* Q

shows *Guarded* $(\lambda X. P(X) \Box Q(X))$

proof –

have *Guarded* $(\lambda X. \Box F \in \{P, Q\} \cdot F(X))$

by (rule *ExtChoice-Guarded*, auto simp add: *assms*)

thus *?thesis*

by (simp add: *extChoice-def*)

qed

7.6 Algebraic laws

lemma *extChoice-comm*:

$P \Box Q = Q \Box P$

by (unfold *extChoice-def*, simp add: *insert-commute*)

lemma *extChoice-idem*:

$P \text{ is CSP} \implies P \Box P = P$

by (unfold *extChoice-def*, simp add: *ExtChoice-single*)

lemma *extChoice-assoc*:

assumes P is CSP Q is CSP R is CSP

shows $P \sqcap Q \sqcap R = P \sqcap (Q \sqcap R)$

proof –

have $P \sqcap Q \sqcap R = \mathbf{R}_s(\text{pre}_R(P) \vdash \text{cmt}_R(P)) \sqcap \mathbf{R}_s(\text{pre}_R(Q) \vdash \text{cmt}_R(Q)) \sqcap \mathbf{R}_s(\text{pre}_R(R) \vdash \text{cmt}_R(R))$

by (*simp add: SRD-reactive-design-alt assms(1) assms(2) assms(3)*)

also have ... =

$\mathbf{R}_s(((\text{pre}_R P \wedge \text{pre}_R Q) \wedge \text{pre}_R R) \vdash$
 $((\text{cmt}_R P \wedge \text{cmt}_R Q) \triangleleft \text{\$tr}' =_u \text{\$tr} \wedge \text{\$wait}' \triangleright (\text{cmt}_R P \vee \text{cmt}_R Q) \wedge \text{cmt}_R R)$
 $\triangleleft \text{\$tr}' =_u \text{\$tr} \wedge \text{\$wait}' \triangleright$
 $((\text{cmt}_R P \wedge \text{cmt}_R Q) \triangleleft \text{\$tr}' =_u \text{\$tr} \wedge \text{\$wait}' \triangleright (\text{cmt}_R P \vee \text{cmt}_R Q) \vee \text{cmt}_R R)))$

by (*simp add: extChoice-rdes unrest*)

also have ... =

$\mathbf{R}_s(((\text{pre}_R P \wedge \text{pre}_R Q) \wedge \text{pre}_R R) \vdash$
 $((\text{cmt}_R P \wedge \text{cmt}_R Q) \wedge \text{cmt}_R R)$
 $\triangleleft \text{\$tr}' =_u \text{\$tr} \wedge \text{\$wait}' \triangleright$
 $((\text{cmt}_R P \vee \text{cmt}_R Q) \vee \text{cmt}_R R)))$

by (*rule cong[of $\mathbf{R}_s \mathbf{R}_s$], simp, rel-auto*)

also have ... =

$\mathbf{R}_s((\text{pre}_R P \wedge \text{pre}_R Q \wedge \text{pre}_R R) \vdash$
 $((\text{cmt}_R P \wedge (\text{cmt}_R Q \wedge \text{cmt}_R R))$
 $\triangleleft \text{\$tr}' =_u \text{\$tr} \wedge \text{\$wait}' \triangleright$
 $(\text{cmt}_R P \vee (\text{cmt}_R Q \vee \text{cmt}_R R))))$

by (*simp add: conj-assoc disj-assoc*)

also have ... =

$\mathbf{R}_s((\text{pre}_R P \wedge \text{pre}_R Q \wedge \text{pre}_R R) \vdash$
 $((\text{cmt}_R P \wedge (\text{cmt}_R Q \wedge \text{cmt}_R R) \triangleleft \text{\$tr}' =_u \text{\$tr} \wedge \text{\$wait}' \triangleright (\text{cmt}_R Q \vee \text{cmt}_R R))$
 $\triangleleft \text{\$tr}' =_u \text{\$tr} \wedge \text{\$wait}' \triangleright$
 $(\text{cmt}_R P \vee (\text{cmt}_R Q \wedge \text{cmt}_R R) \triangleleft \text{\$tr}' =_u \text{\$tr} \wedge \text{\$wait}' \triangleright (\text{cmt}_R Q \vee \text{cmt}_R R))))$

by (*rule cong[of $\mathbf{R}_s \mathbf{R}_s$], simp, rel-auto*)

also have ... = $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{cmt}_R(P)) \sqcap (\mathbf{R}_s(\text{pre}_R(Q) \vdash \text{cmt}_R(Q)) \sqcap \mathbf{R}_s(\text{pre}_R(R) \vdash \text{cmt}_R(R)))$

by (*simp add: extChoice-rdes unrest*)

also have ... = $P \sqcap (Q \sqcap R)$

by (*simp add: SRD-reactive-design-alt assms(1) assms(2) assms(3)*)

finally show *?thesis* .

qed

lemma *extChoice-Stop*:

assumes Q is CSP

shows $\text{Stop} \sqcap Q = Q$

using *assms*

proof –

have $\text{Stop} \sqcap Q = \mathbf{R}_s(\text{true} \vdash (\text{\$tr}' =_u \text{\$tr} \wedge \text{\$wait}')) \sqcap \mathbf{R}_s(\text{pre}_R(Q) \vdash \text{cmt}_R(Q))$

by (*simp add: Stop-def SRD-reactive-design-alt assms*)

also have ... = $\mathbf{R}_s(\text{pre}_R Q \vdash (((\text{\$tr}' =_u \text{\$tr} \wedge \text{\$wait}') \wedge \text{cmt}_R Q) \triangleleft \text{\$tr}' =_u \text{\$tr} \wedge \text{\$wait}' \triangleright (\text{\$tr}' =_u \text{\$tr} \wedge \text{\$wait}' \vee \text{cmt}_R Q)))$

by (*simp add: extChoice-rdes unrest*)

also have ... = $\mathbf{R}_s(\text{pre}_R Q \vdash (\text{cmt}_R Q \triangleleft \text{\$tr}' =_u \text{\$tr} \wedge \text{\$wait}' \triangleright \text{cmt}_R Q))$

by (*metis (no-types, lifting) cond-def eq-upred-sym neg-conj-cancel1 utp-pred-laws.inf.left-idem*)

also have ... = $\mathbf{R}_s(\text{pre}_R Q \vdash \text{cmt}_R Q)$

by (*simp add: cond-idem*)

also have ... = Q

by (*simp add: SRD-reactive-design-alt assms*)

finally show *?thesis* .

qed

lemma *extChoice-Chaos*:

assumes *Q is CSP*

shows $\text{Chaos} \sqcap Q = \text{Chaos}$

proof –

have $\text{Chaos} \sqcap Q = \mathbf{R}_s(\text{false} \vdash \text{true}) \sqcap \mathbf{R}_s(\text{pre}_R(Q) \vdash \text{cmt}_R(Q))$

by (*simp add: Chaos-def SRD-reactive-design-alt assms*)

also have $\dots = \mathbf{R}_s(\text{false} \vdash (\text{cmt}_R Q \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright \text{true}))$

by (*simp add: extChoice-rdes unrest*)

also have $\dots = \mathbf{R}_s(\text{false} \vdash \text{true})$

by (*rule cong[of $\mathbf{R}_s \mathbf{R}_s$], simp, rel-auto*)

also have $\dots = \text{Chaos}$

by (*simp add: Chaos-def*)

finally show *?thesis* .

qed

lemma *extChoice-Dist*:

assumes *P is CSP* $S \subseteq \llbracket \text{CSP} \rrbracket_H$ $S \neq \{\}$

shows $P \sqcap (\bigsqcup S) = (\bigsqcup_{Q \in S} P \sqcap Q)$

proof –

let $?S1 = \text{pre}_R S$ and $?S2 = \text{cmt}_R S$

have $P \sqcap (\bigsqcup S) = P \sqcap (\bigsqcup_{Q \in S} \mathbf{R}_s(\text{pre}_R(Q) \vdash \text{cmt}_R(Q)))$

by (*simp add: SRD-as-reactive-design[THEN sym] Healthy-SUPREMUM UINF-as-Sup-collect assms*)

also have $\dots = \mathbf{R}_s(\text{pre}_R(P) \vdash \text{cmt}_R(P)) \sqcap \mathbf{R}_s((\bigsqcup_{Q \in S} \text{pre}_R(Q)) \vdash (\bigsqcup_{Q \in S} \text{cmt}_R(Q)))$

by (*simp add: RHS-design-USUP SRD-reactive-design-alt assms*)

also have $\dots = \mathbf{R}_s((\text{pre}_R(P) \wedge (\bigsqcup_{Q \in S} \text{pre}_R(Q))) \vdash$

$((\text{cmt}_R(P) \wedge (\bigsqcup_{Q \in S} \text{cmt}_R(Q)))$

$\triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright$

$(\text{cmt}_R(P) \vee (\bigsqcup_{Q \in S} \text{cmt}_R(Q))))$

by (*simp add: extChoice-rdes unrest*)

also have $\dots = \mathbf{R}_s((\bigsqcup_{Q \in S} \text{pre}_R P \wedge \text{pre}_R Q) \vdash$

$(\bigsqcup_{Q \in S} (\text{cmt}_R P \wedge \text{cmt}_R Q) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (\text{cmt}_R P \vee \text{cmt}_R Q)))$

by (*simp add: conj-USUP-dist conj-UINF-dist disj-UINF-dist cond-UINF-dist assms*)

also have $\dots = (\bigsqcup_{Q \in S} \mathbf{R}_s((\text{pre}_R P \wedge \text{pre}_R Q) \vdash$

$((\text{cmt}_R P \wedge \text{cmt}_R Q) \triangleleft \$tr' =_u \$tr \wedge \$wait' \triangleright (\text{cmt}_R P \vee \text{cmt}_R Q))))$

by (*simp add: assms RHS-design-USUP*)

also have $\dots = (\bigsqcup_{Q \in S} \mathbf{R}_s(\text{pre}_R(P) \vdash \text{cmt}_R(P)) \sqcap \mathbf{R}_s(\text{pre}_R(Q) \vdash \text{cmt}_R(Q)))$

by (*simp add: extChoice-rdes unrest*)

also have $\dots = (\bigsqcup_{Q \in S} P \sqcap \text{CSP}(Q))$

by (*simp add: UINF-as-Sup-collect, metis (no-types, lifting) Healthy-if SRD-as-reactive-design assms(1)*)

also have $\dots = (\bigsqcup_{Q \in S} P \sqcap Q)$

by (*rule SUP-cong, simp-all add: Healthy-subset-member[OF assms(2)]*)

finally show *?thesis* .

qed

lemma *extChoice-dist*:

assumes *P is CSP* *Q is CSP* *R is CSP*

shows $P \sqcap (Q \sqcap R) = (P \sqcap Q) \sqcap (P \sqcap R)$

using *assms extChoice-Dist*[of $P \{Q, R\}$] by *simp*

lemma *ExtChoice-seq-distr*:

assumes $\bigwedge i. i \in A \implies P i \text{ is PCSP}$ *Q is NCSP*

shows $(\bigsqcup_{i \in A} P i) ;; Q = (\bigsqcup_{i \in A} P i ;; Q)$

```

proof (cases A = {})
  case True
  then show ?thesis
    by (simp add: ExtChoice-empty NCSP-implies-CSP Stop-left-zero assms(2))
next
  case False
  show ?thesis
  proof –
    have 1: ( $\Box i \in A \cdot P\ i$ ) = ( $\Box i \in A \cdot (\mathbf{R}_s ((pre_R (P\ i)) \vdash peri_R (P\ i) \diamond (R4(post_R (P\ i)))))$ )
      (is ?X = ?Y)
    by (rule ExtChoice-cong, metis (no-types, hide-lams) R4-def Healthy-if NCSP-implies-CSP PCSP-implies-NCSP
      Productive-form assms(1) comp-apply)
    have 2: ( $\Box i \in A \cdot P\ i \;;\ Q$ ) = ( $\Box i \in A \cdot (\mathbf{R}_s ((pre_R (P\ i)) \vdash peri_R (P\ i) \diamond (R4(post_R (P\ i))))) \;;\ Q$ )
      (is ?X = ?Y)
    by (rule ExtChoice-cong, metis (no-types, hide-lams) R4-def Healthy-if NCSP-implies-CSP PCSP-implies-NCSP
      Productive-form assms(1) comp-apply)
    show ?thesis
    by (simp add: 1 2, rdes-eq cls: assms False cong: ExtChoice-cong USUP-cong)
  qed
qed

lemma extChoice-seq-distr:
  assumes P is PCSP Q is PCSP R is NCSP
  shows (P  $\Box$  Q) ;; R = (P ;; R  $\Box$  Q ;; R)
  by (rdes-eq cls: assms)

lemma extChoice-seq-distl:
  assumes P is ICSP Q is ICSP R is NCSP
  shows P ;; (Q  $\Box$  R) = (P ;; Q  $\Box$  P ;; R)
  by (rdes-eq cls: assms)

lemma extchoice-StateInvR-refine:
  assumes
    P is NCSP Q is NCSP
     $sinv_R(b) \sqsubseteq P\ sinv_R(b) \sqsubseteq Q$ 
  shows  $sinv_R(b) \sqsubseteq P\ \Box\ Q$ 
proof –
  have 1:
     $pre_R\ P \sqsubseteq [b]_{S<} [b]_{S>} \sqsubseteq ([b]_{S<} \wedge post_R\ P)$ 
     $pre_R\ Q \sqsubseteq [b]_{S<} [b]_{S>} \sqsubseteq ([b]_{S<} \wedge post_R\ Q)$ 
  by (metis (no-types, lifting) CRR-implies-RR NCSP-implies-CSP RHS-tri-design-refine SRD-reactive-tri-design
    StateInvR-def assms periR-RR postR-RR preR-CRR rea-st-cond-RR rea-true-RR refBy-order st-post-CRR)+
  show ?thesis
  by (rdes-refine-split cls: assms(1–2), simp-all add: 1 closure assms truer-bottom-rpred utp-pred-laws.inf-sup-distrib1)
qed

end

```

8 Stateful-Failure Programs

```

theory utp-sfrd-prog
imports
  UTP.utp-full
  utp-sfrd-extchoice

```

begin

8.1 Conditionals

lemma *NCSP-cond-srea* [closure]:
assumes P is NCSP Q is NCSP
shows $P \triangleleft b \triangleright_R Q$ is NCSP
by (rule *NCSP-NSRD-intro*, simp-all add: closure rdes assms unrest)

8.2 Guarded commands

lemma *GuardedCommR-NCSP-closed* [closure]:
assumes P is NCSP
shows $g \rightarrow_R P$ is NCSP
by (simp add: gcmd-def closure assms)

8.3 Alternation

lemma *AlternateR-NCSP-closed* [closure]:
assumes $\bigwedge i. i \in A \implies P(i)$ is NCSP Q is NCSP
shows (if_R $i \in A \cdot g(i) \rightarrow P(i)$ else Q fi) is NCSP
proof (cases $A = \{\}$)
 case *True*
 then show ?thesis
 by (simp add: assms)
next
 case *False*
 then show ?thesis
 by (simp add: AlternateR-def closure assms)
qed

lemma *AlternateR-list-NCSP-closed* [closure]:
assumes $\bigwedge b P. (b, P) \in \text{set } A \implies P$ is NCSP Q is NCSP
shows (AlternateR-list A Q) is NCSP
apply (simp add: AlternateR-list-def)
apply (rule AlternateR-NCSP-closed)
apply (auto simp add: assms)
apply (metis assms(1) eq-snd-iff nth-mem)
done

8.4 While Loops

lemma *NSRD-coerce-NCSP*:
 P is NSRD $\implies \text{Skip} ;; P ;; \text{Skip}$ is NCSP
by (metis (no-types, hide-lams) CSP3-Skip CSP3-def CSP4-def Healthy-def NCSP-Skip NCSP-implies-CSP NCSP-intro NSRD-is-SRD RA1 SRD-seqr-closure)

definition *WhileC* :: 's upred \Rightarrow ('s, 'e) action \Rightarrow ('s, 'e) action (while_C - do - od) **where**
[rdes-def]: while_C b do P od = Skip ;; while_R b do P od ;; Skip

lemma *WhileC-NCSP-closed* [closure]:
assumes P is NCSP P is Productive
shows while_C b do P od is NCSP
by (simp add: WhileC-def NSRD-coerce-NCSP assms closure)

lemma *WhileC-false*:

P is NCSP \implies WhileC false $P = \text{Skip}$
 by (simp add: NCSP-implies-NSRD Skip-srdes-left-unit WhileC-def WhileR-false)

8.5 Iteration Construction

definition $\text{IterateC} :: 'a \text{ set} \Rightarrow ('a \Rightarrow 's \text{ upred}) \Rightarrow ('a \Rightarrow ('s, 'e) \text{ action}) \Rightarrow ('s, 'e) \text{ action}$
where [upred-defs, ndes-simp]: $\text{IterateC } A \ g \ P = \text{while}_C (\bigvee i \in A \cdot g(i)) \text{ do } (if_R i \in A \cdot g(i) \rightarrow P(i) \text{ fi})$
 od

lemma $\text{IterateC-IterateR-def}$: $\text{IterateC } A \ g \ P = \text{Skip} ;; \text{IterateR } A \ g \ P ;; \text{Skip}$
 by (simp add: IterateC-def IterateR-def WhileC-def)

definition $\text{IterateC-list} :: ('s \text{ upred} \times ('s, 'e) \text{ action}) \text{ list} \Rightarrow ('s, 'e) \text{ action}$ **where**
 [upred-defs, ndes-simp]:
 $\text{IterateC-list } xs = \text{IterateC } \{0..<\text{length } xs\} (\lambda i. \text{map fst } xs ! i) (\lambda i. \text{map snd } xs ! i)$

syntax

-iter-C $:: \text{pttrn} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{do}_C \text{ -} \in \cdot \text{ -} \rightarrow \text{ -} \text{ od}$
 -iter-gcommC $:: \text{gcomms} \Rightarrow \text{logic} \Rightarrow \text{do}_C / \text{ -} / \text{od}$

translations

-iter-C $x \ A \ g \ P \Rightarrow \text{CONST } \text{IterateC } A \ (\lambda x. g) (\lambda x. P)$
 -iter-C $x \ A \ g \ P \Leftarrow \text{CONST } \text{IterateC } A \ (\lambda x. g) (\lambda x'. P)$
 -iter-gcommC $cs \rightarrow \text{CONST } \text{IterateC-list } cs$
 -iter-gcommC $(\text{-gcomm-show } cs) \leftarrow \text{CONST } \text{IterateC-list } cs$

lemma $\text{IterateC-NCSP-closed}$ [closure]:

assumes

$\bigwedge i. i \in I \implies P(i) \text{ is NCSP}$

$\bigwedge i. i \in I \implies P(i) \text{ is Productive}$

shows $\text{do}_C i \in I \cdot g(i) \rightarrow P(i) \text{ od is NCSP}$

by (simp add: IterateC-IterateR-def IterateR-NSRD-closed NCSP-implies-NSRD NSRD-coerce-NCSP
 assms(1) assms(2))

lemma $\text{IterateC-list-NCSP-closed}$ [closure]:

assumes

$\bigwedge b \ P. (b, P) \in \text{set } A \implies P \text{ is NCSP}$

$\bigwedge b \ P. (b, P) \in \text{set } A \implies P \text{ is Productive}$

shows $\text{IterateC-list } A \text{ is NCSP}$

apply (simp add: IterateC-list-def, rule IterateC-NCSP-closed)

apply (metis assms atLeastLessThan-iff nth-map nth-mem prod.collapse)+

done

lemma $\text{IterateC-list-alt-def}$:

$\text{IterateC-list } xs = \text{while}_C (\bigvee b \in \text{set}(\text{map fst } xs) \cdot b) \text{ do } \text{AlternateR-list } xs \text{ Chaos od}$

proof –

have $(\bigvee i \in \{0..<\text{length}(xs)\} \cdot (\text{map fst } xs) ! i) = (\bigvee b \in \text{set}(\text{map fst } xs) \cdot b)$

by (rel-auto, metis nth-mem prod.collapse, metis fst-conv in-set-conv-nth nth-map)

thus ?thesis

by (simp add: IterateC-list-def IterateC-def AlternateR-list-def)

qed

lemma IterateC-empty :

$\text{do}_C i \in \{\} \cdot g(i) \rightarrow P(i) \text{ od} = \text{Skip}$

by (simp add: IterateC-IterateR-def IterateR-empty closure Skip-srdes-left-unit)

lemma *IterateC-singleton:*

assumes $P\ k$ is NCSP $P\ k$ is Productive

shows $do_C\ i \in \{k\} \cdot g(i) \rightarrow P(i)$ od = while_C $g(k)$ do $P(k)$ od (**is** ?lhs = ?rhs)

by (simp add: IterateC-IterateR-def IterateR-singleton NCSP-implies-NSRD WhileC-def assms)

lemma *IterateC-outer-refine-intro:*

assumes $I \neq \{\}$ $\bigwedge i. i \in I \implies P\ i$ is NCSP $\bigwedge i. i \in I \implies P\ i$ is Productive

$\bigwedge i. i \in I \implies S \sqsubseteq (b\ i \rightarrow_R P\ i ;; S)$ S is NCSP

$S \sqsubseteq [\neg (\bigcap i \in I \cdot b\ i)]^\top_R$

shows $S \sqsubseteq do_C\ i \in I \cdot b(i) \rightarrow P(i)$ od

proof –

have $S \sqsubseteq do_R\ i \in I \cdot b(i) \rightarrow P(i)$ od

by (simp add: IterateR-outer-refine-intro NCSP-implies-NSRD assms)

thus ?thesis

unfolding IterateC-IterateR-def

by (metis (full-types) Skip-left-unit Skip-right-unit assms(5) urel-dioid.mult-isol urel-dioid.mult-isor)

qed

lemma *IterateC-outer-refine-init-intro:*

assumes

$\bigwedge i. i \in A \implies P\ i$ is NCSP

$\bigwedge i. i \in A \implies P\ i$ is Productive

S is NCSP I is NCSP

$S \sqsubseteq I ;; [\neg (\bigcap i \in A \cdot b\ i)]^\top_R$

$\bigwedge i. i \in A \implies S \sqsubseteq S ;; b\ i \rightarrow_R P\ i$

$\bigwedge i. i \in A \implies S \sqsubseteq I ;; b\ i \rightarrow_R P\ i$

shows $S \sqsubseteq I ;; do_C\ i \in A \cdot b(i) \rightarrow P(i)$ od

proof (cases $A = \{\}$)

case True

with assms(5) **show** ?thesis

by (simp add: IterateC-empty assms closure Skip-right-unit AssumeR-true NSRD-right-unit)

next

case False

have $S \sqsubseteq I ;; do_R\ i \in A \cdot b(i) \rightarrow P(i)$ od

by (simp add: IterateR-outer-refine-init-intro NCSP-implies-NSRD assms False)

thus ?thesis

unfolding IterateC-IterateR-def

by (metis (no-types, hide-lams) RA1 Skip-right-unit assms(3) assms(4) urel-dioid.mult-isor)

qed

lemma *IterateC-list-outer-refine-intro:*

assumes

$A \neq []$ S is NCSP

$\bigwedge b\ P. (b, P) \in \text{set } A \implies P$ is NCSP

$\bigwedge b\ P. (b, P) \in \text{set } A \implies P$ is Productive

$\bigwedge b\ P. (b, P) \in \text{set } A \implies S \sqsubseteq (b \rightarrow_R P ;; S)$

$S \sqsubseteq [\neg (\bigcap (b, P) \in \text{set } A \cdot b)]^\top_R$

shows $S \sqsubseteq \text{IterateC-list } A$

proof –

have $(\bigcap i \in \{0..<\text{length}(A)\} \cdot (\text{map fst } A) ! i) = (\bigcap (b, P) \in \text{set } A \cdot b)$

by (rel-auto, metis nth-mem prod.exhaust-sel, metis fst-conv in-set-conv-nth nth-map)

thus ?thesis

apply (simp add: IterateC-list-def)

apply (rule IterateC-outer-refine-intro)


```

  apply (simp-all add: closure assms)
  apply (metis assms(3) nth-mem prod.collapse)
  apply (metis assms(4) nth-mem prod.collapse)
done
qed

```

lemma *IterateC-list-outer-refine-init-intro*:

```

assumes
  S is NCSP I is NCSP
   $\bigwedge b P. (b, P) \in \text{set } A \implies P \text{ is NCSP}$ 
   $\bigwedge b P. (b, P) \in \text{set } A \implies P \text{ is Productive}$ 
   $S \sqsubseteq I ;; [\neg (\bigwedge (b, P) \in \text{set } A \cdot b)]^\top_R$ 
   $\bigwedge b P. (b, P) \in \text{set } A \implies S \sqsubseteq S ;; b \rightarrow_R P$ 
   $\bigwedge b P. (b, P) \in \text{set } A \implies S \sqsubseteq I ;; b \rightarrow_R P$ 
shows  $S \sqsubseteq I ;; \text{IterateC-list } A$ 
proof -
  have  $(\bigwedge i \in \{0..<\text{length}(A)\} \cdot (\text{map fst } A) ! i) = (\bigwedge (b, P) \in \text{set } A \cdot b)$ 
  by (rel-auto, metis nth-mem prod.exhaust-sel, metis fst-conv in-set-conv-nth nth-map)
  thus ?thesis
  apply (simp add: IterateC-list-def)
  apply (rule IterateC-outer-refine-init-intro)
  apply (simp-all add: closure assms)
  apply (metis assms(3) nth-mem prod.collapse)
  apply (metis assms(4) nth-mem prod.collapse)
done
qed

```

8.6 Assignment

definition *AssignsCSP* :: $'\sigma \text{ usubst} \Rightarrow (' \sigma, ' \varphi) \text{ action } (\langle \cdot \rangle_C)$ **where**
 $[\text{upred-defs}]: \text{AssignsCSP } \sigma = \mathbf{R}_s(\text{true} \vdash \text{false} \diamond (\$tr' =_u \$tr \wedge [\langle \sigma \rangle_a]_S))$

syntax

```

-assigns-csp :: svids  $\Rightarrow$  uexprs  $\Rightarrow$  logic  $(\langle \cdot \rangle :=_C \langle \cdot \rangle)$ 
-assigns-csp :: svids  $\Rightarrow$  uexprs  $\Rightarrow$  logic (infixr  $:=_C$  64)

```

translations

```

-assigns-csp xs vs => CONST AssignsCSP (-mk-usubst (CONST id) xs vs)
-assigns-csp x v <= CONST AssignsCSP (CONST subst-upd (CONST id) x v)
-assigns-csp x v <= -assigns-csp (-spvar x) v
x,y :=C u,v <= CONST AssignsCSP (CONST subst-upd (CONST subst-upd (CONST id) (CONST svar x) u) (CONST svar y) v)

```

lemma *preR-AssignsCSP* [rdes]: $\text{pre}_R(\langle \sigma \rangle_C) = \text{true}_r$
by (rel-auto)

lemma *periR-AssignsCSP* [rdes]: $\text{peri}_R(\langle \sigma \rangle_C) = \text{false}$
by (rel-auto)

lemma *postR-AssignsCSP* [rdes]: $\text{post}_R(\langle \sigma \rangle_C) = \Phi(\text{true}, \sigma, \langle \rangle)$
by (rel-auto)

lemma *AssignsCSP-rdes-def* [rdes-def] : $\langle \sigma \rangle_C = \mathbf{R}_s(\text{true}_r \vdash \text{false} \diamond \Phi(\text{true}, \sigma, \langle \rangle))$
by (rel-auto)

lemma *AssignsCSP-CSP* [closure]: $\langle \sigma \rangle_C$ is CSP
 by (simp add: AssignsCSP-def RHS-tri-design-is-SRD unrest)

lemma *AssignsCSP-CSP3* [closure]: $\langle \sigma \rangle_C$ is CSP3
 by (rule CSP3-intro, simp add: closure, rel-auto)

lemma *AssignsCSP-CSP4* [closure]: $\langle \sigma \rangle_C$ is CSP4
 by (rule CSP4-intro, simp add: closure, rel-auto+)

lemma *AssignsCSP-NCSP* [closure]: $\langle \sigma \rangle_C$ is NCSP
 by (simp add: AssignsCSP-CSP AssignsCSP-CSP3 AssignsCSP-CSP4 NCSP-intro)

lemma *AssignsCSP-ICSP* [closure]: $\langle \sigma \rangle_C$ is ICSP
 apply (rule ICSP-intro, simp add: closure, simp add: rdes-def)
 apply (rule ISRD1-rdes-intro)
 apply (simp-all add: closure)
 apply (rel-auto)
 done

lemma *AssignsCSP-as-AssignsR*: $\langle \sigma \rangle_R ;; \text{Skip} = \langle \sigma \rangle_C$
 by (rdes-eq)

lemma *AssignC-init-refine-intro*:
 assumes
 $\text{vwb-lens } x \ \$st:x \ \# \ P_2 \ \$st:x \ \# \ P_3$
 $P_2 \text{ is } RR \ P_3 \text{ is } RR \ Q \text{ is } NCSP$
 $\mathbf{R}_s([\&x =_u \ll k \gg]_{S<} \vdash P_2 \diamond P_3) \sqsubseteq Q$
 shows $\mathbf{R}_s(\text{true}_r \vdash P_2 \diamond P_3) \sqsubseteq (x :=_C \ll k \gg) ;; Q$
 by (simp add: AssignsCSP-as-AssignsR[THEN sym] assms seqr-assoc Skip-left-unit AssignR-init-refine-intro closure)

lemma *AssignsCSP-refines-sinv*:
 assumes ' $\sigma \vdash b$ '
 shows $\text{sinv}_R(b) \sqsubseteq \langle \sigma \rangle_C$
 apply (rdes-refine-split)
 apply (simp-all)
 apply (metis rea-st-cond-true st-cond-conj utp-pred-laws.inf.absorb-iff2 utp-pred-laws.inf-top-left)
 using assms apply (rel-auto)
 done

8.7 Assignment with update

There are different collections that we would like to assign to, but they all have different types and perhaps more importantly different conditions on the update being well defined. For example, for a list well-definedness equates to the index being less than the length etc. Thus we here set up a polymorphic constant for CSP assignment updates that can be specialised to different types.

definition *AssignCSP-update* ::
 $(f \Rightarrow 'k \text{ set}) \Rightarrow (f \Rightarrow 'k \Rightarrow 'v \Rightarrow 'f) \Rightarrow (f \Rightarrow ' \sigma) \Rightarrow$
 $('k, ' \sigma) \text{ uexpr} \Rightarrow ('v, ' \sigma) \text{ uexpr} \Rightarrow (' \sigma, ' \varphi) \text{ action} \text{ where}$
 $[upred-defs, rdes-def]: \text{AssignCSP-update domf updatef } x \ k \ v =$
 $\mathbf{R}_s([k \in_u \text{uop domf } (\&x)]_{S<} \vdash \text{false} \diamond \Phi(\text{true}, [x \mapsto_s \text{trop updatef } (\&x) \ k \ v], \langle \rangle))$

All different assignment updates have the same syntax; the type resolves which implementation to use.

syntax

$-csp\text{-}assign\text{-}upd :: svid \Rightarrow uexp \Rightarrow uexp \Rightarrow logic \ (-[-] :=_C \ - [61,0,62] \ 62)$

translations

$-csp\text{-}assign\text{-}upd \ x \ k \ v == CONST \ AssignCSP\text{-}update \ (CONST \ udom) \ (CONST \ uupd) \ x \ k \ v$

lemma *AssignCSP-update-CSP* [closure]:

AssignCSP-update domf updatef x k v is CSP

by (simp add: *AssignCSP-update-def RHS-tri-design-is-SRD unrest*)

lemma *preR-AssignCSP-update* [rdes]:

$pre_R(AssignCSP\text{-}update \ domf \ updatef \ x \ k \ v) = [k \in_u \ uop \ domf \ (\&x)]_{S<}$

by (rel-auto)

lemma *periR-AssignCSP-update* [rdes]:

$peri_R(AssignCSP\text{-}update \ domf \ updatef \ x \ k \ v) = [k \notin_u \ uop \ domf \ (\&x)]_{S<}$

by (rel-simp)

lemma *post-AssignCSP-update* [rdes]:

$post_R(AssignCSP\text{-}update \ domf \ updatef \ x \ k \ v) =$

$(\Phi(true, [x \mapsto_s \ trop \ updatef \ (\&x) \ k \ v], \langle \rangle) \triangleleft (k \in_u \ uop \ domf \ (\&x)) \triangleright_R \ R1(true))$

by (rel-auto)

lemma *AssignCSP-update-NCSP* [closure]:

(AssignCSP-update domf updatef x k v) is NCSP

proof (rule *NCSP-intro*)

show *(AssignCSP-update domf updatef x k v) is CSP*

by (simp add: *closure*)

show *(AssignCSP-update domf updatef x k v) is CSP3*

by (rule *CSP3-SRD-intro*, simp-all add: *csp-do-def closure rdes unrest*)

show *(AssignCSP-update domf updatef x k v) is CSP4*

by (rule *CSP4-tri-intro*, simp-all add: *csp-do-def closure rdes unrest, rel-auto*)

qed

8.8 State abstraction

lemma *ref-unrest-abs-st* [unrest]:

$\$ref \ \# \ P \Longrightarrow \$ref \ \# \ \langle P \rangle_S$

$\$ref' \ \# \ P \Longrightarrow \$ref' \ \# \ \langle P \rangle_S$

by (rel-simp)+

lemma *NCSP-state-srea* [closure]: *P is NCSP \Longrightarrow state 'a \cdot P is NCSP*

apply (rule *NCSP-NSRD-intro*)

apply (simp-all add: *closure rdes*)

apply (simp-all add: *state-srea-def unrest closure*)

done

8.9 Assumptions

definition *AssumeCircus* ($[-]_C$) where

[rdes-def]: $[b]_C = b \rightarrow_R \text{Skip}$

lemma *AssumeCircus-NCSP* [closure]: $[b]_C$ is NCSP

by (simp add: *AssumeCircus-def GuardedCommR-NCSP-closed NCSP-Skip*)

lemma *AssumeCircus-AssumeR*: $\text{Skip} ;; [b]^\top_R = [b]_C \ [b]^\top_R ;; \text{Skip} = [b]_C$

by (rdes-eq)+

lemma *AssumeR-comp-AssumeCircus*: P is NCSP $\implies P \;; [b]^\top_R = P \;; [b]_C$
 by (metis (no-types, hide-lams) AssumeCircus-AssumeR(1) RA1 Skip-right-unit)

lemma *gcmd-AssumeCircus*:
 P is NCSP $\implies b \rightarrow_R P = [b]_C \;; P$
 by (simp add: AssumeCircus-def NCSP-implies-NSRD Skip-left-unit gcmd-seq-distr)

lemma *rdes-assume-pre-refine*:
 assumes P is NCSP
 shows $P \sqsubseteq [b]_C \;; P$
 by (rdes-refine cls: assms)

8.10 Guards

definition *GuardCSP* ::

$'\sigma \text{ cond} \Rightarrow$
 $('\sigma, '\varphi) \text{ action} \Rightarrow$
 $('\sigma, '\varphi) \text{ action where}$

$[upred-defs]$: $\text{GuardCSP } g \ A = \mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r \text{pre}_R(A)) \vdash ((\lceil g \rceil_{S<} \wedge \text{cmt}_R(A)) \vee (\lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$

syntax

-GuardCSP :: $uexp \Rightarrow logic \Rightarrow logic$ (**infixr** $\&_u$ 60)

translations

-GuardCSP $b \ P == \text{CONST GuardCSP } b \ P$

lemma *Guard-tri-design*:

$g \&_u P = \mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r \text{pre}_R P) \vdash (\text{peri}_R(P) \triangleleft \lceil g \rceil_{S<} \triangleright (\$tr' =_u \$tr)) \diamond (\lceil g \rceil_{S<} \wedge \text{post}_R(P)))$

proof –

have $(\lceil g \rceil_{S<} \wedge \text{cmt}_R P \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait') = (\text{peri}_R(P) \triangleleft \lceil g \rceil_{S<} \triangleright (\$tr' =_u \$tr)) \diamond (\lceil g \rceil_{S<} \wedge \text{post}_R(P))$

by (rel-auto)

thus ?thesis by (simp add: GuardCSP-def)

qed

lemma *csp-do-cond-conj*:

assumes P is CRR

shows $(\lceil b \rceil_{S<} \wedge P) = \Phi(b, id, \langle \rangle) \;; P$

proof –

have $(\lceil b \rceil_{S<} \wedge \text{CRR}(P)) = \Phi(b, id, \langle \rangle) \;; \text{CRR}(P)$

by (rel-auto)

thus ?thesis

by (simp add: Healthy-if assms)

qed

lemma *Guard-rdes-def* [rdes-def]:

assumes P is RR Q is CRR R is CRR

shows $g \&_u \mathbf{R}_s(P \vdash Q \diamond R) = \mathbf{R}_s((\mathcal{I}(g, \langle \rangle) \Rightarrow_r P) \vdash ((\Phi(g, id, \langle \rangle) \;; Q) \vee \mathcal{E}(\neg g, \langle \rangle, \{ \}_u)) \diamond (\Phi(g, id, \langle \rangle) \;; R))$

(is ?lhs = ?rhs)

proof –

have ?lhs = $\mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r P) \vdash ((P \Rightarrow_r Q) \triangleleft \lceil g \rceil_{S<} \triangleright (\$tr' =_u \$tr)) \diamond (\lceil g \rceil_{S<} \wedge (P \Rightarrow_r R)))$

by (simp add: Guard-tri-design rdes assms closure)

also have ... = $\mathbf{R}_s ((\mathcal{I}(g, \langle \rangle) \Rightarrow_r P) \vdash (([g]_{S<} \wedge Q) \vee \mathcal{E}(\neg g, \langle \rangle, \{ \}_u)) \diamond ([g]_{S<} \wedge R))$
 by (rel-auto)
 also have ... = $\mathbf{R}_s ((\mathcal{I}(g, \langle \rangle) \Rightarrow_r P) \vdash ((\Phi(g, id, \langle \rangle) ;; Q) \vee \mathcal{E}(\neg g, \langle \rangle, \{ \}_u)) \diamond (\Phi(g, id, \langle \rangle) ;; R))$
 by (simp add: assms(2) assms(3) csp-do-cond-conj)
 finally show ?thesis .
 qed

lemma Guard-rdes-def':

assumes \$ok' $\#$ P
 shows $g \&_u (\mathbf{R}_s(P \vdash Q)) = \mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r P) \vdash (\lceil g \rceil_{S<} \wedge Q \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 proof -
 have $g \&_u (\mathbf{R}_s(P \vdash Q)) = \mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r pre_R (\mathbf{R}_s(P \vdash Q))) \vdash (\lceil g \rceil_{S<} \wedge cmt_R (\mathbf{R}_s(P \vdash Q)) \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 by (simp add: GuardCSP-def)
 also have ... = $\mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r R1(R2c(pre_s \dagger P))) \vdash (\lceil g \rceil_{S<} \wedge R1(R2c(cmt_s \dagger (P \Rightarrow Q))) \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 by (simp add: rea-pre-RHS-design rea-cmt-RHS-design)
 also have ... = $\mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r R1(R2c(pre_s \dagger P))) \vdash R1(R2c(\lceil g \rceil_{S<} \wedge R1(R2c(cmt_s \dagger (P \Rightarrow Q)))) \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 by (metis (no-types, lifting) RHS-design-export-R1 RHS-design-export-R2c)
 also have ... = $\mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r R1(R2c(pre_s \dagger P))) \vdash R1(R2c(\lceil g \rceil_{S<} \wedge (cmt_s \dagger (P \Rightarrow Q)) \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 by (simp add: R1-R2c-commute R1-disj R1-extend-conj' R1-idem R2c-and R2c-disj R2c-idem)
 also have ... = $\mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r R1(R2c(pre_s \dagger P))) \vdash (\lceil g \rceil_{S<} \wedge (cmt_s \dagger (P \Rightarrow Q)) \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 by (metis (no-types, lifting) RHS-design-export-R1 RHS-design-export-R2c)
 also have ... = $\mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r R1(R2c(pre_s \dagger P))) \vdash cmt_s \dagger (\lceil g \rceil_{S<} \wedge (cmt_s \dagger (P \Rightarrow Q)) \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 by (simp add: rdes-export-cmt)
 also have ... = $\mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r R1(R2c(pre_s \dagger P))) \vdash cmt_s \dagger (\lceil g \rceil_{S<} \wedge (P \Rightarrow Q) \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 by (simp add: usubst)
 also have ... = $\mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r R1(R2c(pre_s \dagger P))) \vdash (\lceil g \rceil_{S<} \wedge (P \Rightarrow Q) \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 by (simp add: rdes-export-cmt)
 also from assms have ... = $\mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r (pre_s \dagger P)) \vdash (\lceil g \rceil_{S<} \wedge (P \Rightarrow Q) \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 by (rel-auto)
 also have ... = $\mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r pre_s \dagger P) \llbracket true, false / \$ok, \$wait \rrbracket \vdash (\lceil g \rceil_{S<} \wedge (P \Rightarrow Q) \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 by (simp add: rdes-export-pre)
 also from assms have ... = $\mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r P) \llbracket true, false / \$ok, \$wait \rrbracket \vdash (\lceil g \rceil_{S<} \wedge (P \Rightarrow Q) \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 by (rel-auto)
 also from assms have ... = $\mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r P) \vdash (\lceil g \rceil_{S<} \wedge (P \Rightarrow Q) \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 by (simp add: rdes-export-pre)
 also have ... = $\mathbf{R}_s((\lceil g \rceil_{S<} \Rightarrow_r P) \vdash (\lceil g \rceil_{S<} \wedge Q \vee \lceil \neg g \rceil_{S<} \wedge \$tr' =_u \$tr \wedge \$wait'))$
 by (rule cong[of $\mathbf{R}_s \mathbf{R}_s$], simp, rel-auto)
 finally show ?thesis .
 qed

lemma CSP-Guard [closure]: $b \&_u P$ is CSP

by (simp add: GuardCSP-def, rule RHS-design-is-SRD, simp-all add: unrest)

lemma *preR-Guard* [rdes]: P is CSP $\implies \text{pre}_R(b \&_u P) = ([b]_{S<} \Rightarrow_r \text{pre}_R P)$
 by (simp add: Guard-tri-design rea-pre-RHS-design usubst unrest R2c-preR R2c-lift-state-pre R2c-rea-impl R1-rea-impl R1-preR Healthy-if, rel-auto)

lemma *periR-Guard* [rdes]:
 assumes P is NCSP
 shows $\text{peri}_R(b \&_u P) = (\text{peri}_R P \triangleleft b \triangleright_R \mathcal{E}(\text{true}, \langle \rangle, \{\}_u))$

proof –
 have $\text{peri}_R(b \&_u P) = (([b]_{S<} \Rightarrow_r \text{pre}_R P) \Rightarrow_r (\text{peri}_R P \triangleleft [b]_{S<} \triangleright (\$tr' =_u \$tr)))$
 by (simp add: asms Guard-tri-design rea-peri-RHS-design usubst unrest R1-rea-impl R2c-rea-not R2c-rea-impl R2c-preR R2c-periR R2c-tr'-minus-tr R2c-lift-state-pre R2c-condr closure Healthy-if R1-cond R1-tr'-eq-tr)
 also have $\dots = ((\text{pre}_R P \Rightarrow_r \text{peri}_R P) \triangleleft [b]_{S<} \triangleright (\$tr' =_u \$tr))$
 by (rel-auto)
 also have $\dots = (\text{peri}_R P \triangleleft [b]_{S<} \triangleright (\$tr' =_u \$tr))$
 by (simp add: SRD-peri-under-pre add: unrest closure asms)
 finally show ?thesis
 by rel-auto
qed

lemma *postR-Guard* [rdes]:
 assumes P is NCSP
 shows $\text{post}_R(b \&_u P) = ([b]_{S<} \wedge \text{post}_R P)$

proof –
 have $\text{post}_R(b \&_u P) = (([b]_{S<} \Rightarrow_r \text{pre}_R P) \Rightarrow_r ([b]_{S<} \wedge \text{post}_R P))$
 by (simp add: Guard-tri-design rea-post-RHS-design usubst unrest R2c-rea-not R2c-and R2c-rea-impl R2c-preR R2c-postR R2c-tr'-minus-tr R2c-lift-state-pre R2c-condr R1-rea-impl R1-extend-conj' R1-post-SRD closure asms)
 also have $\dots = ([b]_{S<} \wedge (\text{pre}_R P \Rightarrow_r \text{post}_R P))$
 by (rel-auto)
 also have $\dots = ([b]_{S<} \wedge \text{post}_R P)$
 by (simp add: SRD-post-under-pre add: unrest closure asms)
 also have $\dots = ([b]_{S<} \wedge \text{post}_R P)$
 by (metis CSP-Guard R1-extend-conj R1-post-SRD calculation rea-st-cond-def)
 finally show ?thesis .
qed

lemma *CSP3-Guard* [closure]:
 assumes P is CSP P is CSP3
 shows $b \&_u P$ is CSP3

proof –
 from asms have $1:\$ref \# P \llbracket \text{false}/\$wait \rrbracket$
 by (simp add: CSP-Guard CSP3-iff)
 hence $\$ref \# \text{pre}_R (P \llbracket 0/\$tr \rrbracket) \ \$ref \# \text{pre}_R P \ \$ref \# \text{cmt}_R P$
 by (pred-blast)+
 hence $\$ref \# (b \&_u P) \llbracket \text{false}/\$wait \rrbracket$
 by (simp add: CSP3-iff GuardCSP-def RHS-def R1-def R2c-def R2s-def R3h-def design-def unrest usubst)
 thus ?thesis
 by (metis CSP3-intro CSP-Guard)
qed

lemma *CSP4-Guard* [closure]:
 assumes P is NCSP
 shows $b \&_u P$ is CSP4

proof (*rule CSP4-tri-intro*[*OF CSP-Guard*])
show $(\neg_r \text{pre}_R (b \&_u P)) \;; R1 \text{ true} = (\neg_r \text{pre}_R (b \&_u P))$
proof –
have $a: (\neg_r \text{pre}_R P) \;; R1 \text{ true} = (\neg_r \text{pre}_R P)$
by (*simp add: CSP4-neg-pre-unit assms closure*)
have $(\neg_r ([b]_{S<} \Rightarrow_r \text{pre}_R P)) \;; R1 \text{ true} = (\neg_r ([b]_{S<} \Rightarrow_r \text{pre}_R P))$
proof –
have $1: (\neg_r ([b]_{S<} \Rightarrow_r \text{pre}_R P)) = ([b]_{S<} \wedge (\neg_r \text{pre}_R P))$
by (*rel-auto*)
also have $2: \dots = ([b]_{S<} \wedge ((\neg_r \text{pre}_R P) \;; R1 \text{ true}))$
by (*simp add: a*)
also have $3: \dots = (\neg_r ([b]_{S<} \Rightarrow_r \text{pre}_R P)) \;; R1 \text{ true}$
by (*rel-auto*)
finally show *?thesis* ..
qed
thus *?thesis*
by (*simp add: preR-Guard periR-Guard NSRD-CSP4-intro closure assms unrest*)
qed
show $\$st' \# \text{peri}_R (b \&_u P)$
by (*simp add: preR-Guard periR-Guard NSRD-CSP4-intro closure assms unrest*)
show $\$ref' \# \text{post}_R (b \&_u P)$
by (*simp add: preR-Guard postR-Guard NSRD-CSP4-intro closure assms unrest*)
qed

lemma *NCSP-Guard* [*closure*]:

assumes *P is NCSP*
shows *b &_u P is NCSP*

proof –

have *P is CSP*
using *NCSP-implies-CSP assms* **by** *blast*

then show *?thesis*

by (*metis* (*no-types*) *CSP3-Guard CSP3-commutes-CSP4 CSP4-Guard CSP4-Idempotent CSP-Guard Healthy-Idempotent Healthy-def NCSP-def assms comp-apply*)

qed

lemma *Productive-Guard* [*closure*]:

assumes *P is CSP P is Productive \$wait' \# pre_R(P)*
shows *b &_u P is Productive*

proof –

have $b \&_u P = b \&_u \mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond (\text{post}_R(P) \wedge \$tr <_u \$tr'))$
by (*metis Healthy-def Productive-form assms(1) assms(2)*)

also have $\dots =$

$\mathbf{R}_s (([b]_{S<} \Rightarrow_r \text{pre}_R P) \vdash ((\text{pre}_R P \Rightarrow_r \text{peri}_R P) \triangleleft [b]_{S<} \triangleright (\$tr' =_u \$tr)) \diamond ([b]_{S<} \wedge (\text{pre}_R P \Rightarrow_r \text{post}_R P \wedge \$tr' >_u \$tr)))$

by (*simp add: Guard-tri-design rea-pre-RHS-design rea-peri-RHS-design rea-post-RHS-design unrest assms*)

usubst R1-preR Healthy-if R1-rea-impl R1-peri-SRD R1-extend-conj' R2c-preR R2c-not R2c-rea-impl

R2c-periR R2c-postR R2c-and R2c-tr-less-tr' R1-tr-less-tr'

also have $\dots = \mathbf{R}_s (([b]_{S<} \Rightarrow_r \text{pre}_R P) \vdash (\text{peri}_R P \triangleleft [b]_{S<} \triangleright (\$tr' =_u \$tr)) \diamond (([b]_{S<} \wedge \text{post}_R P) \wedge \$tr' >_u \$tr))$

by (*rel-auto*)

also have $\dots = \text{Productive}(b \&_u P)$

by (*simp add: Productive-def Guard-tri-design RHS-tri-design-par unrest*)

finally show ?thesis
 by (simp add: Healthy-def')
 qed

lemma *Guard-refines-sinv*:
 assumes P is NCSP $\text{sinv}_R(b) \sqsubseteq P$
 shows $\text{sinv}_R(b) \sqsubseteq g \ \&_u \ P$
proof –
 from *assms*
 have $\mathbf{R}_s([b]_{S<} \vdash R1 \text{ true} \diamond [b]_{S>}) \sqsubseteq \mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P))$
 by (simp add: rdes-def NCSP-implies-CSP SRD-reactive-tri-design)
 thus ?thesis
 apply (simp add: RHS-tri-design-refine' closure unrest *assms*)
 apply (safe)
 apply (rdes-refine cls: *assms*(1))
 done
 qed

8.11 Basic events

definition $\text{do}_u ::$
 $(\text{'}\varphi, \text{'}\sigma) \text{ uexpr} \Rightarrow (\text{'}\sigma, \text{'}\varphi) \text{ action}$ **where**
 $[\text{upred-defs}]: \text{do}_u e = ((\text{'}\sigma' =_u \text{'}\sigma \wedge \lceil e \rceil_{S<} \notin_u \text{'}\sigma \text{ ref}') \triangleleft \text{'}\sigma' \triangleright (\text{'}\sigma' =_u \text{'}\sigma \wedge \lceil e \rceil_{S<} \wedge \text{'}\sigma' =_u \text{'}\sigma))$

definition $\text{DoCSP} :: (\text{'}\varphi, \text{'}\sigma) \text{ uexpr} \Rightarrow (\text{'}\sigma, \text{'}\varphi) \text{ action}$ (do_C) **where**
 $[\text{upred-defs}]: \text{DoCSP } a = \mathbf{R}_s(\text{true} \vdash \text{do}_u a)$

lemma $R1\text{-DoAct}: R1(\text{do}_u(a)) = \text{do}_u(a)$
 by (rel-auto)

lemma $R2c\text{-DoAct}: R2c(\text{do}_u(a)) = \text{do}_u(a)$
 by (rel-auto)

lemma $\text{DoCSP-alt-def}: \text{do}_C(a) = R3h(\text{CSP1}(\text{'}\sigma \wedge \text{do}_u(a)))$
 apply (simp add: DoCSP-def RHS-def design-def impl-alt-def R1-R3h-commute R2c-R3h-commute R2c-disj
 $R2c\text{-not } R2c\text{-ok } R2c\text{-ok}' \ R2c\text{-and } R2c\text{-DoAct } R1\text{-disj } R1\text{-extend-conj}' \ R1\text{-DoAct}$)
 apply (rel-auto)
 done

lemma $\text{DoAct-unrests } [\text{unrest}]:$
 $\text{'}\sigma \not\# \text{do}_u(a) \ \text{'}\sigma' \not\# \text{do}_u(a)$
 by (pred-auto)+

lemma $\text{DoCSP-RHS-tri } [\text{rdes-def}]:$
 $\text{do}_C(a) = \mathbf{R}_s(\text{true}_r \vdash (\mathcal{E}(\text{true}, \langle \rangle, \{a\}_u) \diamond \Phi(\text{true}, \text{id}, \langle a \rangle)))$
 by (simp add: DoCSP-def $\text{do}_u\text{-def}$ wait'-cond-def, rel-auto)

lemma $\text{CSP-DoCSP } [\text{closure}]: \text{do}_C(a) \text{ is CSP}$
 by (simp add: DoCSP-def $\text{do}_u\text{-def}$ RHS-design-is-SRD unrest)

lemma $\text{preR-DoCSP } [\text{rdes}]: \text{pre}_R(\text{do}_C(a)) = \text{true}_r$
 by (simp add: DoCSP-def rea-pre-RHS-design unrest usubst R2c-true)

lemma $\text{periR-DoCSP } [\text{rdes}]: \text{peri}_R(\text{do}_C(a)) = \mathcal{E}(\text{true}, \langle \rangle, \{a\}_u)$

by (*rel-auto*)

lemma *postR-DoCSP* [*rdes*]: $\text{post}_R(\text{do}_C(a)) = \Phi(\text{true}, \text{id}, \langle a \rangle)$
 by (*rel-auto*)

lemma *CSP3-DoCSP* [*closure*]: $\text{do}_C(a)$ is *CSP3*
 by (*rule CSP3-intro*[*OF CSP-DoCSP*])
 (*simp add: DoCSP-def do_u-def RHS-def design-def R1-def R2c-def R2s-def R3h-def unrest usubst*)

lemma *CSP4-DoCSP* [*closure*]: $\text{do}_C(a)$ is *CSP4*
 by (*rule CSP4-tri-intro*[*OF CSP-DoCSP*], *simp-all add: preR-DoCSP periR-DoCSP postR-DoCSP unrest*)

lemma *NCSP-DoCSP* [*closure*]: $\text{do}_C(a)$ is *NCSP*
 by (*metis CSP3-DoCSP CSP4-DoCSP CSP-DoCSP Healthy-def NCSP-def comp-apply*)

lemma *Productive-DoCSP* [*closure*]:
 ($\text{do}_C a :: (\sigma, \psi)$ action) is *Productive*
proof –
 have $((\Phi(\text{true}, \text{id}, \langle a \rangle) \wedge \text{\$tr}' >_u \text{\$tr}) :: (\sigma, \psi)$ action)
 = $(\Phi(\text{true}, \text{id}, \langle a \rangle))$
 by (*rel-auto*, *simp add: Prefix-Order.strict-prefixI'*)
 hence $\text{Productive}(\text{do}_C a) = \text{do}_C a$
 by (*simp add: Productive-RHS-design-form DoCSP-RHS-tri unrest*)
 thus ?thesis
 by (*simp add: Healthy-def*)
qed

lemma *PCSP-DoCSP* [*closure*]:
 ($\text{do}_C a :: (\sigma, \psi)$ action) is *PCSP*
 by (*simp add: Healthy-comp NCSP-DoCSP Productive-DoCSP*)

lemma *wp-rea-DoCSP-lemma*:
 fixes $P :: (\sigma, \varphi)$ action
 assumes $\text{\$ok} \# P \text{\$wait} \# P$
 shows $(\text{\$tr}' =_u \text{\$tr} \hat{^}_u \langle \lceil a \rceil_{S<} \rangle \wedge \text{\$st}' =_u \text{\$st}) ;; P = (\exists \text{\$ref} \cdot P[\text{\$tr} \hat{^}_u \langle \lceil a \rceil_{S<} \rangle / \text{\$tr}])$
 using *assms*
 by (*rel-auto*, *meson*)

lemma *wp-rea-DoCSP*:
 assumes P is *NCSP*
 shows $(\text{\$tr}' =_u \text{\$tr} \hat{^}_u \langle \lceil a \rceil_{S<} \rangle \wedge \text{\$st}' =_u \text{\$st}) \text{wp}_r \text{pre}_R P =$
 $(\neg_r (\neg_r \text{pre}_R P)[\text{\$tr} \hat{^}_u \langle \lceil a \rceil_{S<} \rangle / \text{\$tr}])$
 by (*simp add: wp-rea-def wp-rea-DoCSP-lemma unrest usubst ex-unrest assms closure*)

lemma *wp-rea-DoCSP-alt*:
 assumes P is *NCSP*
 shows $(\text{\$tr}' =_u \text{\$tr} \hat{^}_u \langle \lceil a \rceil_{S<} \rangle \wedge \text{\$st}' =_u \text{\$st}) \text{wp}_r \text{pre}_R P =$
 $(\text{\$tr}' \geq_u \text{\$tr} \hat{^}_u \langle \lceil a \rceil_{S<} \rangle \Rightarrow_r (\text{pre}_R P)[\text{\$tr} \hat{^}_u \langle \lceil a \rceil_{S<} \rangle / \text{\$tr}])$
 by (*simp add: wp-rea-DoCSP assms rea-not-def R1-def usubst unrest, rel-auto*)

lemma *DoCSP-refine-sinv*: $\text{sinv}_R(b) \sqsubseteq \text{do}_C(a)$
 by (*rdes-refine*)

8.12 Event prefix

definition *PrefixCSP* ::

$(\sigma, \sigma) \text{ ueexpr} \Rightarrow$
 $(\sigma, \sigma) \text{ action} \Rightarrow$
 $(\sigma, \sigma) \text{ action } (- \rightarrow_C - [81, 80] 80) \text{ where}$
 $[upred-defs]: \text{PrefixCSP } a \ P = (do_C(a) ;; \text{CSP}(P))$

abbreviation *OutputCSP* $c \ v \ P \equiv \text{PrefixCSP } (c.v)_u \ P$

lemma *CSP-PrefixCSP [closure]*: *PrefixCSP* $a \ P$ is *CSP*
 by (simp add: *PrefixCSP-def closure*)

lemma *CSP3-PrefixCSP [closure]*:
PrefixCSP $a \ P$ is *CSP3*
 by (metis (no-types, hide-lams) *CSP3-DoCSP CSP3-def Healthy-def PrefixCSP-def seqr-assoc*)

lemma *CSP4-PrefixCSP [closure]*:
 assumes P is *CSP* P is *CSP4*
 shows *PrefixCSP* $a \ P$ is *CSP4*
 by (metis (no-types, hide-lams) *CSP4-def Healthy-def PrefixCSP-def assms(1) assms(2) seqr-assoc*)

lemma *NCSP-PrefixCSP [closure]*:
 assumes P is *NCSP*
 shows *PrefixCSP* $a \ P$ is *NCSP*
 by (metis (no-types, hide-lams) *CSP3-PrefixCSP CSP3-commutes-CSP4 CSP4-Idempotent CSP4-PrefixCSP CSP-PrefixCSP Healthy-Idempotent Healthy-def NCSP-def NCSP-implies-CSP assms comp-apply*)

lemma *Productive-PrefixCSP [closure]*: P is *NCSP* \implies *PrefixCSP* $a \ P$ is *Productive*
 by (simp add: *Healthy-if NCSP-DoCSP NCSP-implies-NSRD NSRD-is-SRD PrefixCSP-def Productive-DoCSP Productive-seq-1*)

lemma *PCSP-PrefixCSP [closure]*: P is *NCSP* \implies *PrefixCSP* $a \ P$ is *PCSP*
 by (simp add: *Healthy-comp NCSP-PrefixCSP Productive-PrefixCSP*)

lemma *PrefixCSP-Guarded [closure]*: *Guarded* (*PrefixCSP* a)

proof –

have *PrefixCSP* $a = (\lambda X. do_C(a) ;; \text{CSP}(X))$
 by (simp add: *fun-eq-iff PrefixCSP-def*)
 thus ?thesis
 using *Guarded-if-Productive NCSP-DoCSP NCSP-implies-NSRD Productive-DoCSP* by auto
 qed

lemma *PrefixCSP-type [closure]*: *PrefixCSP* $a \in \llbracket H \rrbracket_H \rightarrow \llbracket \text{CSP} \rrbracket_H$
 using *CSP-PrefixCSP* by blast

lemma *PrefixCSP-Continuous [closure]*: *Continuous* (*PrefixCSP* a)
 by (simp add: *Continuous-def PrefixCSP-def ContinuousD[OF SRD-Continuous] seq-Sup-distl*)

lemma *PrefixCSP-RHS-tri-lemma1*:
 $R1 \ (R2s \ (\$tr' =_u \$tr \hat{=} \langle [a]_{S<} \wedge [II]_R)) = (\$tr' =_u \$tr \hat{=} \langle [a]_{S<} \wedge [II]_R))$
 by (rel-auto)

lemma *PrefixCSP-RHS-tri-lemma2*:
 fixes $P :: (\sigma, \sigma) \text{ action}$
 assumes $\$ok \ \# \ P \ \$wait \ \# \ P$

shows $((\$tr' =_u \$tr \hat{^}_u \langle [a]_{S<} \rangle \wedge \$st' =_u \$st) \wedge \neg \$wait')$;; $P = (\exists \$ref \cdot P[\$tr \hat{^}_u \langle [a]_{S<} \rangle / \$tr])$
using *assms*
by (*rel-auto*, *meson*, *fastforce*)

lemma *tr-extend-seqr*:

fixes $P :: ('σ, 'φ) \text{ action}$
assumes $\$ok \# P \$wait \# P \$ref \# P$
shows $(\$tr' =_u \$tr \hat{^}_u \langle [a]_{S<} \rangle \wedge \$st' =_u \$st)$;; $P = P[\$tr \hat{^}_u \langle [a]_{S<} \rangle / \$tr]$
using *assms* **by** (*simp add: wp-rea-DoCSP-lemma assms unrest ex-unrest*)

lemma *trace-ext-R1-closed* [*closure*]: P is $R1 \implies P[\$tr \hat{^}_u e / \$tr]$ is $R1$
by (*rel-blast*)

lemma *preR-PrefixCSP-NCSP* [*rdes*]:

assumes P is $NCSP$
shows $pre_R(PrefixCSP\ a\ P) = (\mathcal{I}(true, \langle a \rangle) \Rightarrow_r (pre_R\ P)[\langle a \rangle]_t)$
by (*simp add: PrefixCSP-def assms closure rdes rpred Healthy-if wp usubst unrest*)

lemma *periR-PrefixCSP* [*rdes*]:

assumes P is $NCSP$
shows $peri_R(PrefixCSP\ a\ P) = (\mathcal{E}(true, \langle \rangle, \{a\}_u) \vee (peri_R\ P)[\langle a \rangle]_t)$

proof –

have $peri_R(PrefixCSP\ a\ P) = peri_R(do_C\ a\ ;;\ P)$
by (*simp add: PrefixCSP-def closure assms Healthy-if*)
also have $\dots = ((\mathcal{I}(true, \langle a \rangle) \Rightarrow_r pre_R\ P[\langle a \rangle]_t) \Rightarrow_r \$tr' =_u \$tr \wedge [a]_{S<} \notin_u \$ref' \vee peri_R\ P[\langle a \rangle]_t)$
by (*simp add: assms NSRD-CSP4-intro csp-enable-tr-empty closure rdes unrest ex-unrest usubst rpred wp*)
also have $\dots = (\mathcal{E}(true, \langle \rangle, \{a\}_u) \vee ((\mathcal{I}(true, \langle a \rangle) \Rightarrow_r pre_R\ P[\langle a \rangle]_t) \Rightarrow_r peri_R\ P[\langle a \rangle]_t))$
by (*rel-auto*)
also have $\dots = (\mathcal{E}(true, \langle \rangle, \{a\}_u) \vee ((pre_R(P) \Rightarrow_r peri_R\ P)[\langle a \rangle]_t))$
by (*rel-auto*)
also have $\dots = (\mathcal{E}(true, \langle \rangle, \{a\}_u) \vee (peri_R\ P)[\langle a \rangle]_t)$
by (*simp add: SRD-peri-under-pre assms closure unrest*)
finally show *?thesis* .

qed

lemma *postR-PrefixCSP* [*rdes*]:

assumes P is $NCSP$
shows $post_R(PrefixCSP\ a\ P) = (post_R\ P)[\langle a \rangle]_t$

proof –

have $post_R(PrefixCSP\ a\ P) = ((\mathcal{I}(true, \langle a \rangle) \Rightarrow_r (pre_R\ P)[\langle a \rangle]_t) \Rightarrow_r (post_R\ P)[\langle a \rangle]_t)$
by (*simp add: PrefixCSP-def assms Healthy-if*)
(simp add: assms Healthy-if wp closure rdes rpred wp-rea-DoCSP-lemma unrest ex-unrest usubst)
also have $\dots = (\mathcal{I}(true, \langle a \rangle) \wedge (pre_R\ P \Rightarrow_r post_R\ P)[\langle a \rangle]_t)$
by (*rel-auto*)
also have $\dots = (\mathcal{I}(true, \langle a \rangle) \wedge (post_R\ P)[\langle a \rangle]_t)$
by (*simp add: SRD-post-under-pre assms closure unrest*)
also have $\dots = (post_R\ P)[\langle a \rangle]_t$
by (*rel-auto*)
finally show *?thesis* .

qed

lemma *PrefixCSP-RHS-tri*:

assumes P is $NCSP$
shows $PrefixCSP\ a\ P = \mathbf{R}_s((\mathcal{I}(true, \langle a \rangle) \Rightarrow_r pre_R\ P[\langle a \rangle]_t) \vdash (\mathcal{E}(true, \langle \rangle, \{a\}_u) \vee peri_R\ P[\langle a \rangle]_t) \diamond$

$post_R P[\langle a \rangle]_t$

by (simp add: PrefixCSP-def Healthy-if unrest assms closure NSRD-composition-wp rdes rpred usubst wp)

For prefix, we can chose whether to propagate the assumptions or not, hence there are two laws.

lemma *PrefixCSP-rdes-def-1* [rdes-def]:

assumes P is CRC Q is CRR R is CRR
 $\$st' \# Q \$ref' \# R$
 shows $PrefixCSP\ a\ (\mathbf{R}_s(P \vdash Q \diamond R)) = \mathbf{R}_s((\mathcal{I}(true, \langle a \rangle) \Rightarrow_r P[\langle a \rangle]_t) \vdash (\mathcal{E}(true, \langle \rangle, \{a\}_u) \vee Q[\langle a \rangle]_t) \diamond R[\langle a \rangle]_t)$
 apply (subst PrefixCSP-RHS-tri)
 apply (rule NCSP-rdes-intro)
 apply (simp-all add: assms rdes closure)
 apply (rel-auto)
 done

lemma *PrefixCSP-rdes-def-2*:

assumes P is CRC Q is CRR R is CRR
 $\$st' \# Q \$ref' \# R$
 shows $PrefixCSP\ a\ (\mathbf{R}_s(P \vdash Q \diamond R)) = \mathbf{R}_s((\mathcal{I}(true, \langle a \rangle) \Rightarrow_r P[\langle a \rangle]_t) \vdash (\mathcal{E}(true, \langle \rangle, \{a\}_u) \vee (P \wedge Q)[\langle a \rangle]_t) \diamond (P \wedge R)[\langle a \rangle]_t)$
 apply (subst PrefixCSP-RHS-tri)
 apply (rule NCSP-rdes-intro)
 apply (simp-all add: assms rdes closure)
 apply (rel-auto)
 done

8.13 Guarded external choice

abbreviation *GuardedChoiceCSP* :: $'\vartheta$ set $\Rightarrow (' \vartheta \Rightarrow (' \sigma, ' \vartheta)$ action) $\Rightarrow (' \sigma, ' \vartheta)$ action **where**
GuardedChoiceCSP $A\ P \equiv (\Box x \in A \cdot PrefixCSP \ll x \gg (P(x)))$

syntax

-*GuardedChoiceCSP* :: logic \Rightarrow logic \Rightarrow logic \Rightarrow logic (\Box - \in - \rightarrow - [0,0,85] 86)

translations

$\Box x \in A \rightarrow P == CONST\ GuardedChoiceCSP\ A\ (\lambda x. P)$

lemma *GuardedChoiceCSP* [rdes-def]:

assumes $\bigwedge x. P(x)$ is NCSP $A \neq \{\}$
 shows $(\Box x \in A \rightarrow P(x)) =$
 $\mathbf{R}_s((\bigwedge x \in A \cdot \mathcal{I}(true, \langle \ll x \gg \rangle) \Rightarrow_r pre_R (P\ x)[\langle \ll x \gg \rangle]_t) \vdash$
 $((\bigwedge x \in A \cdot \mathcal{E}(true, \langle \rangle, \{\ll x \gg\}_u)) \triangleleft \$tr' =_u \$tr \triangleright (\bigwedge x \in A \cdot peri_R (P\ x)[\langle \ll x \gg \rangle]_t)) \diamond$
 $(\bigwedge x \in A \cdot post_R (P\ x)[\langle \ll x \gg \rangle]_t))$
 by (simp add: PrefixCSP-RHS-tri assms ExtChoice-tri-rdes closure unrest, rel-auto)

8.14 Input prefix

definition *InputCSP* ::

$('a, ' \vartheta)$ chan $\Rightarrow ('a \Rightarrow ' \sigma\ upred) \Rightarrow ('a \Rightarrow (' \sigma, ' \vartheta)$ action) $\Rightarrow (' \sigma, ' \vartheta)$ action **where**
 [upred-defs]: *InputCSP* $c\ A\ P = (\Box x \in UNIV \cdot A(x) \ \&_u\ PrefixCSP\ (c \cdot \ll x \gg)_u\ (P\ x))$

definition *InputVarCSP* :: $('a, ' \vartheta)$ chan $\Rightarrow ('a \Rightarrow ' \sigma) \Rightarrow ('a \Rightarrow ' \sigma\ upred) \Rightarrow (' \sigma, ' \vartheta)$ action **where**
 [upred-defs, rdes-def]: *InputVarCSP* $c\ x\ A = InputCSP\ c\ A\ (\lambda v. \langle [x \mapsto_s \ll v \gg] \rangle_C)$

definition do_I ::

$(\text{'}a, \text{'}\vartheta) \text{ chan} \Rightarrow$
 $(\text{'}a \Rightarrow (\text{'}\sigma, \text{'}\vartheta) \text{ st-csp}) \Rightarrow$
 $(\text{'}a \Rightarrow (\text{'}\sigma, \text{'}\vartheta) \text{ action}) \Rightarrow$
 $(\text{'}\sigma, \text{'}\vartheta) \text{ action where}$
 $\text{do}_I \ c \ x \ P = ($
 $(\$tr' =_u \$tr \wedge \{e : \ll \delta_u(c) \gg \mid P(e) \cdot (c \cdot \ll e \gg)_u\}_u \cap_u \$ref' =_u \{\}_u)$
 $\triangleleft \$wait' \triangleright$
 $((\$tr' - \$tr) \in_u \{e : \ll \delta_u(c) \gg \mid P(e) \cdot \langle (c \cdot \ll e \gg)_u \rangle\}_u \wedge (c \cdot \$x')_u =_u \text{last}_u(\$tr'))))$

lemma *InputCSP-CSP [closure]: InputCSP c A P is CSP*
by (*simp add: CSP-ExtChoice InputCSP-def*)

lemma *InputCSP-NCSP [closure]: $\llbracket \bigwedge v. P(v) \text{ is NCSP} \rrbracket \Rightarrow \text{InputCSP c A P is NCSP}$*
apply (*simp add: InputCSP-def*)
apply (*rule NCSP-ExtChoice*)
apply (*simp add: NCSP-Guard NCSP-PrefixCSP image-Collect-subsetI top-set-def*)
done

lemma *Productive-InputCSP [closure]:*
 $\llbracket \bigwedge v. P(v) \text{ is NCSP} \rrbracket \Rightarrow \text{InputCSP x A P is Productive}$
by (*auto simp add: InputCSP-def unrest closure intro: Productive-ExtChoice*)

lemma *preR-InputCSP [rdes]:*
assumes $\bigwedge v. P(v) \text{ is NCSP}$
shows $\text{pre}_R(\text{InputCSP a A P}) = (\bigsqcup v \cdot [A(v)]_{S<} \Rightarrow_r \mathcal{I}(\text{true}, \langle (a \cdot \ll v \gg)_u \rangle) \Rightarrow_r (\text{pre}_R(P(v))) \llbracket \langle (a \cdot \ll v \gg)_u \rangle \rrbracket_t)$
by (*simp add: InputCSP-def rdes closure assms alpha usubst unrest*)

lemma *periR-InputCSP [rdes]:*
assumes $\bigwedge v. P(v) \text{ is NCSP}$
shows $\text{peri}_R(\text{InputCSP a A P}) =$
 $((\bigsqcup x \cdot [A(x)]_{S<} \Rightarrow_r \mathcal{E}(\text{true}, \langle \rangle, \{(a \cdot \ll x \gg)_u\}_u))$
 $\triangleleft \$tr' =_u \$tr \triangleright$
 $(\bigcap x \cdot [A(x)]_{S<} \wedge (\text{peri}_R(P x)) \llbracket \langle (a \cdot \ll x \gg)_u \rangle \rrbracket_t)$
by (*simp add: InputCSP-def rdes closure assms, rel-auto*)

lemma *postR-InputCSP [rdes]:*
assumes $\bigwedge v. P(v) \text{ is NCSP}$
shows $\text{post}_R(\text{InputCSP a A P}) =$
 $(\bigcap x \cdot [A x]_{S<} \wedge \text{post}_R(P x) \llbracket \langle (a \cdot \ll x \gg)_u \rangle \rrbracket_t)$
using *assms by (simp add: InputCSP-def rdes closure assms usubst unrest)*

lemma *InputCSP-rdes-def [rdes-def]:*
assumes $\bigwedge v. P(v) \text{ is CRC} \wedge v. Q(v) \text{ is CRR} \wedge v. R(v) \text{ is CRR}$
 $\bigwedge v. \$st' \nmid Q(v) \wedge v. \$ref' \nmid R(v)$
shows $\text{InputCSP a A } (\lambda v. \mathbf{R}_s(P(v) \vdash Q(v) \diamond R(v))) =$
 $\mathbf{R}_s(\bigsqcup v \cdot ([A(v)]_{S<} \Rightarrow_r \mathcal{I}(\text{true}, \langle (a \cdot \ll v \gg)_u \rangle) \Rightarrow_r (P v) \llbracket \langle (a \cdot \ll v \gg)_u \rangle \rrbracket_t))$
 $\vdash (((\bigsqcup x \cdot \mathbf{R}_5([A(x)]_{S<} \Rightarrow_r \mathcal{E}(\text{true}, \langle \rangle, \{(a \cdot \ll x \gg)_u\}_u))))$
 \vee
 $(\bigcap x \cdot [A(x)]_{S<} \wedge (P x \wedge Q x) \llbracket \langle (a \cdot \ll x \gg)_u \rangle \rrbracket_t))$
 $\diamond (\bigcap x \cdot [A x]_{S<} \wedge (P x \wedge R x) \llbracket \langle (a \cdot \ll x \gg)_u \rangle \rrbracket_t))$ (**is** *?lhs = ?rhs*)

proof –

have $1: \text{pre}_R(?lhs) = (\bigsqcup v \cdot [A v]_{S<} \Rightarrow_r \mathcal{I}(\text{true}, \langle (a \cdot \ll v \gg)_u \rangle) \Rightarrow_r P v \llbracket \langle (a \cdot \ll v \gg)_u \rangle \rrbracket_t)$ (**is** *- = ?A*)
by (*simp add: rdes NCSP-rdes-intro assms conj-comm closure*)
have $2: \text{peri}_R(?lhs) = (\bigsqcup x \cdot [A x]_{S<} \Rightarrow_r \mathcal{E}(\text{true}, \langle \rangle, \{(a \cdot \ll x \gg)_u\}_u)) \triangleleft \$tr' =_u \$tr \triangleright (\bigcap x \cdot [A x]_{S<} \wedge (P x \Rightarrow_r Q x) \llbracket \langle (a \cdot \ll x \gg)_u \rangle \rrbracket_t)$

(is - = ?B)
 by (simp add: rdes NCSP-rdes-intro assms closure)
 have $\exists \text{post}_R(?lhs) = (\bigwedge x \cdot [A\ x]_{S<} \wedge (P\ x \Rightarrow_r R\ x) \llbracket \langle (a \cdot \langle x \rangle)_u \rangle_t \rrbracket)$
 (is - = ?C)
 by (simp add: rdes NCSP-rdes-intro assms closure)
 have ?lhs = $\mathbf{R}_s(?A \vdash ?B \diamond ?C)$
 by (subst SRD-reactive-tri-design[THEN sym], simp-all add: closure 1 2 3)
 also have ... = ?rhs
 by (rel-auto)
 finally show ?thesis .
 qed

8.15 Renaming

definition *RenameCSP* :: ('s, 'e) action \Rightarrow ('e \Rightarrow 'f) \Rightarrow ('s, 'f) action $((-)\llbracket - \rrbracket_C [999, 0] 999)$ **where**
 $[upred-defs]: \text{RenameCSP } Pf = \mathbf{R}_s((\neg_r (\neg_r \text{pre}_R(P))\llbracket f \rrbracket_c ;; \text{true}_r) \vdash ((\text{peri}_R(P))\llbracket f \rrbracket_c) \diamond ((\text{post}_R(P))\llbracket f \rrbracket_c))$

lemma *RenameCSP-rdes-def* [rdes-def]:

assumes *P is CRC Q is CRR R is CRR*
 shows $(\mathbf{R}_s(P \vdash Q \diamond R))\llbracket f \rrbracket_C = \mathbf{R}_s((\neg_r (\neg_r P)\llbracket f \rrbracket_c ;; \text{true}_r) \vdash Q\llbracket f \rrbracket_c \diamond R\llbracket f \rrbracket_c)$ (is ?lhs = ?rhs)
proof –
 have ?lhs = $\mathbf{R}_s((\neg_r (\neg_r P)\llbracket f \rrbracket_c ;; \text{true}_r) \vdash (P \Rightarrow_r Q)\llbracket f \rrbracket_c \diamond (P \Rightarrow_r R)\llbracket f \rrbracket_c)$
 by (simp add: RenameCSP-def rdes closure assms)
 also have ... = $\mathbf{R}_s((\neg_r (\neg_r \text{CRC}(P))\llbracket f \rrbracket_c ;; \text{true}_r) \vdash (\text{CRC}(P) \Rightarrow_r \text{CRR}(Q))\llbracket f \rrbracket_c \diamond (\text{CRC}(P) \Rightarrow_r \text{CRR}(R))\llbracket f \rrbracket_c)$
 by (simp add: Healthy-if assms)
 also have ... = $\mathbf{R}_s((\neg_r (\neg_r \text{CRC}(P))\llbracket f \rrbracket_c ;; \text{true}_r) \vdash (\text{CRR}(Q))\llbracket f \rrbracket_c \diamond (\text{CRR}(R))\llbracket f \rrbracket_c)$
 by (rel-auto, (metis order-refl)+)
 also have ... = ?rhs
 by (simp add: Healthy-if assms)
 finally show ?thesis .
 qed

lemma *RenameCSP-pre-CRC-closed*:

assumes *P is CRR*
 shows $\neg_r (\neg_r P)\llbracket f \rrbracket_c ;; R1 \text{ true is CRC}$
 apply (rule CRC-intro'')
 apply (simp add: unrest closure assms)
 apply (simp add: Healthy-def, simp add: RC1-def rpred closure CRC-idem assms segr-assoc)
 done

lemma *RenameCSP-NCSP-closed* [closure]:

assumes *P is NCSP*
 shows $P\llbracket f \rrbracket_C$ is NCSP
 by (simp add: RenameCSP-def RenameCSP-pre-CRC-closed closure assms unrest)

lemma *csp-rename-false* [rpred]:

$\text{false}\llbracket f \rrbracket_c = \text{false}$
 by (rel-auto)

lemma *umap-nil* [simp]: $\text{map}_u f \langle \rangle = \langle \rangle$

by (rel-auto)

lemma *rename-Skip*: $\text{Skip}\llbracket f \rrbracket_C = \text{Skip}$

by (rdes-eq)

lemma *rename-Chaos*: $\text{Chaos}(\llbracket f \rrbracket)_C = \text{Chaos}$
by (*rdes-eq-split*; *rel-simp*; *force*)

lemma *rename-Miracle*: $\text{Miracle}(\llbracket f \rrbracket)_C = \text{Miracle}$
by (*rdes-eq*)

lemma *rename-DoCSP*: $(\text{do}_C(a))(\llbracket f \rrbracket)_C = \text{do}_C(\llbracket f \rrbracket(a))$
by (*rdes-eq*)

8.16 Algebraic laws

lemma *AssignCSP-conditional*:
assumes *vwb-lens* x
shows $x :=_C e \triangleleft b \triangleright_R x :=_C f = x :=_C (e \triangleleft b \triangleright f)$
by (*rdes-eq cls: assms*)

lemma *AssignsCSP-id*: $\langle \text{id} \rangle_C = \text{Skip}$
by (*rel-auto*)

lemma *Guard-comp*:
 $g \&_u h \&_u P = (g \wedge h) \&_u P$
by (*rule antisym, rel-blast, rel-blast*)

lemma *Guard-false* [*simp*]: $\text{false} \&_u P = \text{Stop}$
by (*simp add: GuardCSP-def Stop-def rpred closure alpha R1-design-R1-pre*)

lemma *Guard-true* [*simp*]:
 $P \text{ is CSP} \implies \text{true} \&_u P = P$
by (*simp add: GuardCSP-def alpha SRD-reactive-design-alt closure rpred*)

lemma *Guard-conditional*:
assumes $P \text{ is NCSP}$
shows $b \&_u P = P \triangleleft b \triangleright_R \text{Stop}$
by (*rdes-eq cls: assms*)

lemma *Guard-expansion*:
 $(g_1 \vee g_2) \&_u P = (g_1 \&_u P) \sqcap (g_2 \&_u P)$
by (*rel-auto*)

lemma *Conditional-as-Guard*:
assumes $P \text{ is NCSP}$ $Q \text{ is NCSP}$
shows $P \triangleleft b \triangleright_R Q = b \&_u P \sqcap (\neg b) \&_u Q$
by (*rdes-eq cls: assms; simp add: le-less*)

lemma *PrefixCSP-dist*:
 $\text{PrefixCSP } a (P \sqcap Q) = (\text{PrefixCSP } a P) \sqcap (\text{PrefixCSP } a Q)$
using *Continuous-Disjunctous Disjunctuous-def PrefixCSP-Continuous* **by** *auto*

lemma *DoCSP-is-Prefix*:
 $\text{do}_C(a) = \text{PrefixCSP } a \text{ Skip}$
by (*simp add: PrefixCSP-def Healthy-if closure, metis CSP4-DoCSP CSP4-def Healthy-def*)

lemma *PrefixCSP-seq*:
assumes $P \text{ is CSP}$ $Q \text{ is CSP}$
shows $(\text{PrefixCSP } a P) ;; Q = (\text{PrefixCSP } a (P ;; Q))$
by (*simp add: PrefixCSP-def seqr-assoc Healthy-if assms closure*)

lemma *PrefixCSP-extChoice-dist*:

assumes P is NCSP Q is NCSP R is NCSP

shows $((a \rightarrow_C P) \sqcap (b \rightarrow_C Q)) ;; R = (a \rightarrow_C P ;; R) \sqcap (b \rightarrow_C Q ;; R)$

by (*simp add: PCSP-PrefixCSP assms(1) assms(2) assms(3) extChoice-seq-distr*)

lemma *GuardedChoiceCSP-dist*:

assumes $\bigwedge i. i \in A \implies P(i)$ is NCSP Q is NCSP

shows $\square x \in A \rightarrow P(x) ;; Q = \square x \in A \rightarrow (P(x) ;; Q)$

by (*simp add: ExtChoice-seq-distr PrefixCSP-seq closure assms cong: ExtChoice-cong*)

Alternation can be re-expressed as an external choice when the guards are disjoint

declare *ExtChoice-tri-rdes* [*rdes-def*]

declare *ExtChoice-tri-rdes'* [*rdes-def del*]

declare *extChoice-rdes-def* [*rdes-def*]

declare *extChoice-rdes-def'* [*rdes-def del*]

lemma *AlternateR-as-ExtChoice*:

assumes

$\bigwedge i. i \in A \implies P(i)$ is NCSP Q is NCSP

$\bigwedge i j. \llbracket i \in A; j \in A; i \neq j \rrbracket \implies (g\ i \wedge g\ j) = \text{false}$

shows $(\text{if}_R\ i \in A \cdot g(i) \rightarrow P(i) \text{ else } Q\ fi) =$

$(\square i \in A \cdot g(i) \ \&_u\ P(i)) \sqcap (\bigwedge i \in A \cdot \neg g(i)) \ \&_u\ Q$

proof (*cases A = {}*)

case *True*

then show *?thesis* **by** (*simp add: ExtChoice-empty extChoice-Stop closure assms*)

next

case *False*

show *?thesis*

proof –

have $1: (\bigwedge i \in A \cdot g\ i \rightarrow_R P\ i) = (\bigwedge i \in A \cdot g\ i \rightarrow_R \mathbf{R}_s(\text{pre}_R(P\ i) \vdash \text{peri}_R(P\ i) \diamond \text{post}_R(P\ i)))$

by (*rule UINF-cong, simp add: NCSP-implies-CSP SRD-reactive-tri-design assms(1)*)

have $2: (\square i \in A \cdot g(i) \ \&_u\ P(i)) = (\square i \in A \cdot g(i) \ \&_u\ \mathbf{R}_s(\text{pre}_R(P\ i) \vdash \text{peri}_R(P\ i) \diamond \text{post}_R(P\ i)))$

by (*rule ExtChoice-cong, simp add: NCSP-implies-NSRD NSRD-is-SRD SRD-reactive-tri-design assms(1)*)

from *assms(3)* **show** *?thesis*

by (*simp add: AlternateR-def 1 2*)

(*rdes-eq cls: assms(1–2)_simps: False cong: UINF-cong ExtChoice-cong*)

qed

qed

declare *ExtChoice-tri-rdes* [*rdes-def del*]

declare *ExtChoice-tri-rdes'* [*rdes-def*]

declare *extChoice-rdes-def* [*rdes-def del*]

declare *extChoice-rdes-def'* [*rdes-def*]

end

9 Recursion in Stateful-Failures

theory *utp-sfrd-recursion*

imports *utp-sfrd-contracts utp-sfrd-prog*

begin

9.1 Fixed-points

The CSP weakest fixed-point is obtained simply by precomposing the body with the CSP healthiness condition.

abbreviation *mu-CSP* :: $((\sigma, \varphi) \text{ action} \Rightarrow (\sigma, \varphi) \text{ action}) \Rightarrow (\sigma, \varphi) \text{ action } (\mu_C)$ **where**
 $\mu_C F \equiv \mu (F \circ \text{CSP})$

syntax

-mu-CSP :: *pttrn* \Rightarrow *logic* \Rightarrow *logic* $(\mu_C \cdot \cdot \cdot [\theta, 10] \ 10)$

translations

$\mu_C X \cdot P == \text{CONST } \text{mu-CSP } (\lambda X. P)$

lemma *mu-CSP-equiv*:

assumes *Monotonic F* $F \in \llbracket \text{CSP} \rrbracket_H \rightarrow \llbracket \text{CSP} \rrbracket_H$

shows $(\mu_R F) = (\mu_C F)$

by (*simp add: srd-mu-equiv assms comp-def*)

lemma *mu-CSP-unfold*:

P is CSP $\implies (\mu_C X \cdot P ;; X) = P ;; (\mu_C X \cdot P ;; X)$

apply (*subst gfp-unfold*)

apply (*simp-all add: closure Healthy-if*)

done

lemma *mu-csp-expand* [*rdes*]: $(\mu_C ((;;) Q)) = (\mu X \cdot Q ;; \text{CSP } X)$

by (*simp add: comp-def*)

lemma *mu-csp-basic-refine*:

assumes

P is CSP Q is NCSP Q is Productive $\text{pre}_R(P) = \text{true}_r \text{pre}_R(Q) = \text{true}_r$

$\text{peri}_R P \sqsubseteq \text{peri}_R Q$

$\text{peri}_R P \sqsubseteq \text{post}_R Q ;; \text{peri}_R P$

shows $P \sqsubseteq (\mu_C X \cdot Q ;; X)$

proof (*rule SRD-refine-intro', simp-all add: closure usubst alpha rpred rdes unrest wp seq-UINF-distr assms*)

show $\text{peri}_R P \sqsubseteq (\bigcap i \cdot \text{post}_R Q \wedge i ;; \text{peri}_R Q)$

proof (*rule UINF-refines'*)

fix *i*

show $\text{peri}_R P \sqsubseteq \text{post}_R Q \wedge i ;; \text{peri}_R Q$

proof (*induct i*)

case 0

then show ?*case* **by** (*simp add: assms*)

next

case (*Suc i*)

then show ?*case*

by (*meson assms(6) assms(7) semilattice-sup-class.le-sup-iff upower-inductl*)

qed

qed

qed

lemma *CRD-mu-basic-refine*:

fixes *P* :: 'e list \Rightarrow 'e set \Rightarrow 's upred

assumes

Q is NCSP Q is Productive $pre_R(Q) = true_r$
 $[P \ t \ r]_{S<} \llbracket (t, r) \rightarrow (\&tt, \$ref')_u \rrbracket \sqsubseteq peri_R \ Q$
 $[P \ t \ r]_{S<} \llbracket (t, r) \rightarrow (\&tt, \$ref')_u \rrbracket \sqsubseteq post_R \ Q \ ;;_h [P \ t \ r]_{S<} \llbracket (t, r) \rightarrow (\&tt, \$ref')_u \rrbracket$
shows $[true \vdash P \ trace \ refs \mid R]_C \sqsubseteq (\mu_C \ X \cdot Q \ ;; \ X)$
proof (rule *mu-csp-basic-refine*, *simp-all* add: *msubst-pair* *assms* *closure* *alpha* *rdes* *rpred* *Healthy-if* *R1-false*)
show $[P \ trace \ refs]_{S<} \llbracket trace \rightarrow \&tt \rrbracket \llbracket refs \rightarrow \$ref' \rrbracket \sqsubseteq peri_R \ Q$
using *assms* **by** (*simp* add: *msubst-pair*)
show $[P \ trace \ refs]_{S<} \llbracket trace \rightarrow \&tt \rrbracket \llbracket refs \rightarrow \$ref' \rrbracket \sqsubseteq post_R \ Q \ ;; [P \ trace \ refs]_{S<} \llbracket trace \rightarrow \&tt \rrbracket \llbracket refs \rightarrow \$ref' \rrbracket$
using *assms* **by** (*simp* add: *msubst-pair*)
qed

9.2 Example action expansion

lemma *mu-example1*: $(\mu \ X \cdot \ll a \gg \rightarrow_C \ X) = (\bigcap i \cdot do_C(\ll a \gg) \wedge (i+1)) \ ;; \text{Miracle}$
by (*simp* add: *PrefixCSP-def* *mu-csp-form-1* *closure*)

lemma *preR-mu-example1* [*rdes*]: $pre_R(\mu \ X \cdot \ll a \gg \rightarrow_C \ X) = true_r$
by (*simp* add: *mu-example1* *rdes* *closure* *unrest* *wp*)

lemma *periR-mu-example1* [*rdes*]:
 $peri_R(\mu \ X \cdot \ll a \gg \rightarrow_C \ X) = (\bigcap i \cdot \mathcal{E}(true, iter[i](\ll a \gg), \{\ll a \gg\}_u))$
by (*simp* add: *mu-example1* *rdes* *rpred* *closure* *unrest* *wp* *seq-UINF-distr* *alpha* *usubst*)

lemma *postR-mu-example1* [*rdes*]:
 $post_R(\mu \ X \cdot \ll a \gg \rightarrow_C \ X) = false$
by (*simp* add: *mu-example1* *rdes* *closure* *unrest* *wp*)

end

10 Linking to the Failures-Divergences Model

theory *utp-sfrd-fdsem*
imports *utp-sfrd-recursion*
begin

10.1 Failures-Divergences Semantics

The following functions play a similar role to those in Roscoe's CSP semantics, and are calculated from the Circus reactive design semantics. A major difference is that these three functions account for state. Each divergence, trace, and failure is subject to an initial state. Moreover, the traces are terminating traces, and therefore also provide a final state following the given interaction. A more subtle difference from the Roscoe semantics is that the set of traces do not include the divergences. The same semantic information is present, but we construct a direct analogy with the pre-, peri- and postconditions of our reactive designs.

definition *divergences* :: $(\sigma, \varphi) \text{ action} \Rightarrow \sigma \Rightarrow \varphi \text{ list set } (dv[-] - [0,100] \ 100)$ **where**
 [*upred-defs*]: $divergences \ P \ s = \{t \mid t. '(\neg_r \ pre_R(P)) \llbracket \ll s \gg, \langle \rangle, \ll t \gg / \$st, \$tr, \$tr' \rrbracket '\}$

definition *traces* :: $(\sigma, \varphi) \text{ action} \Rightarrow \sigma \Rightarrow (\varphi \text{ list} \times \sigma) \text{ set } (tr[-] - [0,100] \ 100)$ **where**
 [*upred-defs*]: $traces \ P \ s = \{(t, s') \mid t \ s'. '(pre_R(P) \wedge post_R(P)) \llbracket \ll s \gg, \ll s' \gg, \langle \rangle, \ll t \gg / \$st, \$st', \$tr, \$tr' \rrbracket '\}$

definition *failures* :: $(\sigma, \varphi) \text{ action} \Rightarrow \sigma \Rightarrow (\varphi \text{ list} \times \sigma) \text{ set } (fl[-] - [0,100] \ 100)$ **where**
 [*upred-defs*]: $failures \ P \ s = \{(t, r) \mid t \ r. '(pre_R(P) \wedge peri_R(P)) \llbracket \ll r \gg, \ll s \gg, \langle \rangle, \ll t \gg / \$ref', \$st, \$tr, \$tr' \rrbracket '\}$

lemma *trace-divergence-disj*:

assumes P is NCSP $(t, s') \in tr[P]s \ t \in dv[P]s$
shows *False*
using *assms(2,3)*
by (*simp add: traces-def divergences-def, rdes-simp cls:assms, rel-auto*)

lemma *preR-refine-divergences*:

assumes P is NCSP Q is NCSP $\bigwedge s. dv[P]s \subseteq dv[Q]s$
shows $pre_R(P) \sqsubseteq pre_R(Q)$

proof (*rule CRR-refine-impl-prop, simp-all add: assms closure usubst unrest*)

fix $t s$

assume a : $[\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \dagger pre_R Q$

with a **show** $[\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \dagger pre_R P$

proof (*rule-tac ccontr*)

from *assms(3)[of s]* **have** b : $t \in dv[P]s \implies t \in dv[Q]s$

by (*auto*)

assume \neg $[\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \dagger pre_R P$

hence \neg $[\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \dagger CRC(pre_R P)$

by (*simp add: assms closure Healthy-if*)

hence $[\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \dagger (\neg_r CRC(pre_R P))$

by (*rel-auto*)

hence $[\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \dagger (\neg_r pre_R P)$

by (*simp add: assms closure Healthy-if*)

with $a b$ **show** *False*

by (*rel-auto*)

qed

qed

lemma *preR-eq-divergences*:

assumes P is NCSP Q is NCSP $\bigwedge s. dv[P]s = dv[Q]s$

shows $pre_R(P) = pre_R(Q)$

by (*metis assms dual-order.antisym order-refl preR-refine-divergences*)

lemma *periR-refine-failures*:

assumes P is NCSP Q is NCSP $\bigwedge s. fl[Q]s \subseteq fl[P]s$

shows $(pre_R(P) \wedge peri_R(P)) \sqsubseteq (pre_R(Q) \wedge peri_R(Q))$

proof (*rule CRR-refine-impl-prop, simp-all add: assms closure unrest subst-unrest-3*)

fix $t s r'$

assume a : $[\$ref' \mapsto_s \ll r' \gg, \$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \dagger (pre_R Q \wedge peri_R Q)$

from *assms(3)[of s]* **have** b : $(t, r') \in fl[Q]s \implies (t, r') \in fl[P]s$

by (*auto*)

with a **show** $[\$ref' \mapsto_s \ll r' \gg, \$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \dagger (pre_R P \wedge peri_R P)$

by (*simp add: failures-def*)

qed

lemma *periR-eq-failures*:

assumes P is NCSP Q is NCSP $\bigwedge s. fl[P]s = fl[Q]s$

shows $(pre_R(P) \wedge peri_R(P)) = (pre_R(Q) \wedge peri_R(Q))$

by (*metis (full-types) assms dual-order.antisym order-refl periR-refine-failures*)

lemma *postR-refine-traces*:

assumes P is NCSP Q is NCSP $\bigwedge s. tr[Q]s \subseteq tr[P]s$

shows $(pre_R(P) \wedge post_R(P)) \sqsubseteq (pre_R(Q) \wedge post_R(Q))$

proof (*rule CRR-refine-impl-prop, simp-all add: assms closure unrest subst-unrest-5*)

fix $t s s'$

assume a : $[\$st \mapsto_s \ll s \gg, \$st' \mapsto_s \ll s' \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \uparrow (pre_R Q \wedge post_R Q)$
from $assms(3)[of\ s]$ **have** b : $(t, s') \in tr\llbracket Q \rrbracket s \implies (t, s') \in tr\llbracket P \rrbracket s$
by (*auto*)
with a **show** $[\$st \mapsto_s \ll s \gg, \$st' \mapsto_s \ll s' \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \uparrow (pre_R P \wedge post_R P)$
by (*simp add: traces-def*)
qed

lemma *postR-eq-traces*:
assumes P is NCSP Q is NCSP $\wedge s. tr\llbracket P \rrbracket s = tr\llbracket Q \rrbracket s$
shows $(pre_R(P) \wedge post_R(P)) = (pre_R(Q) \wedge post_R(Q))$
by (*metis assms dual-order.antisym order-refl postR-refine-traces*)

lemma *circus-fd-refine-intro*:
assumes P is NCSP Q is NCSP $\wedge s. dv\llbracket Q \rrbracket s \subseteq dv\llbracket P \rrbracket s \wedge s. fl\llbracket Q \rrbracket s \subseteq fl\llbracket P \rrbracket s \wedge s. tr\llbracket Q \rrbracket s \subseteq tr\llbracket P \rrbracket s$
shows $P \sqsubseteq Q$
proof (*rule SRD-refine-intro', simp-all add: closure assms*)
show a : $pre_R P \Rightarrow pre_R Q$
using $assms(1)$ $assms(2)$ $assms(3)$ *preR-refine-divergences refBy-order* **by** *blast*
show $peri_R P \sqsubseteq (pre_R P \wedge peri_R Q)$
proof –
have $peri_R P \sqsubseteq (pre_R Q \wedge peri_R Q)$
by (*metis (no-types) assms(1) assms(2) assms(4) periR-refine-failures utp-pred-laws.le-inf-iff*)
then show *?thesis*
by (*metis a refBy-order utp-pred-laws.inf.order-iff utp-pred-laws.inf-assoc*)
qed
show $post_R P \sqsubseteq (pre_R P \wedge post_R Q)$
proof –
have $post_R P \sqsubseteq (pre_R Q \wedge post_R Q)$
by (*meson assms(1) assms(2) assms(5) postR-refine-traces utp-pred-laws.le-inf-iff*)
then show *?thesis*
by (*metis a refBy-order utp-pred-laws.inf.absorb-iff1 utp-pred-laws.inf-assoc*)
qed
qed

10.2 Circus Operators

lemma *traces-Skip*:
 $tr\llbracket Skip \rrbracket s = \{(\llbracket, s\rrbracket)\}$
by (*simp add: traces-def rdes alpha closure, rel-simp*)

lemma *failures-Skip*:
 $fl\llbracket Skip \rrbracket s = \{\}$
by (*simp add: failures-def, rdes-calc*)

lemma *divergences-Skip*:
 $dv\llbracket Skip \rrbracket s = \{\}$
by (*simp add: divergences-def, rdes-calc*)

lemma *traces-Stop*:
 $tr\llbracket Stop \rrbracket s = \{\}$
by (*simp add: traces-def, rdes-calc*)

lemma *failures-Stop*:
 $fl\llbracket Stop \rrbracket s = \{(\llbracket, E) \mid E. True\}$
by (*simp add: failures-def, rdes-calc, rel-auto*)

lemma *divergences-Stop*:
 $dv\llbracket Stop \rrbracket s = \{\}$
by (*simp add: divergences-def, rdes-calc*)

lemma *traces-AssignsCSP*:
 $tr\llbracket \langle \sigma \rangle_C \rrbracket s = \{(\llbracket \cdot \rrbracket, \sigma(s))\}$
by (*simp add: traces-def rdes closure usubst alpha, rel-auto*)

lemma *failures-AssignsCSP*:
 $fl\llbracket \langle \sigma \rangle_C \rrbracket s = \{\}$
by (*simp add: failures-def, rdes-calc*)

lemma *divergences-AssignsCSP*:
 $dv\llbracket \langle \sigma \rangle_C \rrbracket s = \{\}$
by (*simp add: divergences-def, rdes-calc*)

lemma *failures-Miracle*: $fl\llbracket Miracle \rrbracket s = \{\}$
by (*simp add: failures-def rdes closure usubst*)

lemma *divergences-Miracle*: $dv\llbracket Miracle \rrbracket s = \{\}$
by (*simp add: divergences-def rdes closure usubst*)

lemma *failures-Chaos*: $fl\llbracket Chaos \rrbracket s = \{\}$
by (*simp add: failures-def rdes, rel-auto*)

lemma *divergences-Chaos*: $dv\llbracket Chaos \rrbracket s = UNIV$
by (*simp add: divergences-def rdes, rel-auto*)

lemma *traces-Chaos*: $tr\llbracket Chaos \rrbracket s = \{\}$
by (*simp add: traces-def rdes closure usubst*)

lemma *divergences-cond*:
assumes P is NCSP Q is NCSP
shows $dv\llbracket P \triangleleft b \triangleright_R Q \rrbracket s = (if (\llbracket b \rrbracket_e s) then dv\llbracket P \rrbracket s else dv\llbracket Q \rrbracket s)$
by (*rdes-simp cls: assms, simp add: divergences-def traces-def rdes closure rpred assms, rel-auto*)

lemma *traces-cond*:
assumes P is NCSP Q is NCSP
shows $tr\llbracket P \triangleleft b \triangleright_R Q \rrbracket s = (if (\llbracket b \rrbracket_e s) then tr\llbracket P \rrbracket s else tr\llbracket Q \rrbracket s)$
by (*rdes-simp cls: assms, simp add: divergences-def traces-def rdes closure rpred assms, rel-auto*)

lemma *failures-cond*:
assumes P is NCSP Q is NCSP
shows $fl\llbracket P \triangleleft b \triangleright_R Q \rrbracket s = (if (\llbracket b \rrbracket_e s) then fl\llbracket P \rrbracket s else fl\llbracket Q \rrbracket s)$
by (*rdes-simp cls: assms, simp add: divergences-def failures-def rdes closure rpred assms, rel-auto*)

lemma *divergences-guard*:
assumes P is NCSP
shows $dv\llbracket g \&_u P \rrbracket s = (if (\llbracket g \rrbracket_e s) then dv\llbracket g \&_u P \rrbracket s else \{\})$
by (*rdes-simp cls: assms, simp add: divergences-def traces-def rdes closure rpred assms, rel-auto*)

lemma *traces-do*: $tr\llbracket do_C(e) \rrbracket s = \{(\llbracket e \rrbracket_e s, s)\}$
by (*rdes-simp, simp add: traces-def rdes closure rpred, rel-auto*)

lemma *failures-do*: $fl\llbracket do_C(e) \rrbracket s = \{(\llbracket \cdot \rrbracket, E) \mid E. \llbracket e \rrbracket_e s \notin E\}$

by (rdes-simp, simp add: failures-def rdes closure rpred usubst, rel-auto)

lemma *divergences-do*: $dv\llbracket do_C(e) \rrbracket s = \{\}$
 by (rel-auto)

lemma *divergences-seq*:
 fixes $P :: ('s, 'e) \text{ action}$
 assumes $P \text{ is NCSP } Q \text{ is NCSP}$
 shows $dv\llbracket P ;; Q \rrbracket s = dv\llbracket P \rrbracket s \cup \{t_1 @ t_2 \mid t_1 \ t_2 \ s_0. (t_1, s_0) \in tr\llbracket P \rrbracket s \wedge t_2 \in dv\llbracket Q \rrbracket s_0\}$
 (is ?lhs = ?rhs)
 oops

lemma *traces-seq*:
 fixes $P :: ('s, 'e) \text{ action}$
 assumes $P \text{ is NCSP } Q \text{ is NCSP}$
 shows $tr\llbracket P ;; Q \rrbracket s =$

$$\{(t_1 @ t_2, s') \mid t_1 \ t_2 \ s_0 \ s'. (t_1, s_0) \in tr\llbracket P \rrbracket s \wedge (t_2, s') \in tr\llbracket Q \rrbracket s_0$$

$$\wedge (t_1 @ t_2) \notin dv\llbracket P \rrbracket s$$

$$\wedge (\forall (t, s_1) \in tr\llbracket P \rrbracket s. t \leq t_1 @ t_2 \longrightarrow (t_1 @ t_2) - t \notin dv\llbracket Q \rrbracket s_1) \}$$

 (is ?lhs = ?rhs)

proof

show ?lhs \subseteq ?rhs

proof (rdes-expand cls: assms, simp add: traces-def divergences-def rdes closure assms rdes-def unrest rpred usubst, auto)

fix $t :: 'e \text{ list}$ and $s' :: 's$

let $\sigma = [\$st \mapsto_s \llbracket s \rrbracket, \$st' \mapsto_s \llbracket s' \rrbracket, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \llbracket t \rrbracket]$

assume

a1: $\sigma \vdash (post_R P ;; post_R Q)'$ and

a2: $[\$st \mapsto_s \llbracket s \rrbracket, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \llbracket t \rrbracket] \vdash pre_R P'$ and

a3: $[\$st \mapsto_s \llbracket s \rrbracket, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \llbracket t \rrbracket] \vdash (post_R P \text{ wp}_r pre_R Q)'$

from a1 have $\sigma \vdash (\exists tr_0. ((post_R P) \llbracket \llbracket tr_0 \rrbracket / \$tr' \rrbracket ;; (post_R Q) \llbracket \llbracket tr_0 \rrbracket / \$tr \rrbracket) \wedge \llbracket tr_0 \rrbracket \leq_u \$tr')$
 by (simp add: R2-tr-middle assms closure)

then obtain tr_0 where $p1: \sigma \vdash ((post_R P) \llbracket \llbracket tr_0 \rrbracket / \$tr' \rrbracket ;; (post_R Q) \llbracket \llbracket tr_0 \rrbracket / \$tr \rrbracket)'$ and $tr_0 \leq t$

apply (simp add: usubst)

apply (erule taut-shEx-elim)

apply (simp add: unrest-all-circus-vars-st-st' closure unrest assms)

apply (rel-auto)

done

from p1 have $\sigma \vdash (\exists st_0. (post_R P) \llbracket \llbracket tr_0 \rrbracket / \$tr' \rrbracket \llbracket \llbracket st_0 \rrbracket / \$st' \rrbracket ;; (post_R Q) \llbracket \llbracket tr_0 \rrbracket / \$tr \rrbracket \llbracket \llbracket st_0 \rrbracket / \$st \rrbracket)'$
 by (simp add: seqr-middle[of st, THEN sym])

then obtain s_0 where $\sigma \vdash ((post_R P) \llbracket \llbracket s_0 \rrbracket, \llbracket tr_0 \rrbracket / \$st', \$tr' \rrbracket ;; (post_R Q) \llbracket \llbracket s_0 \rrbracket, \llbracket tr_0 \rrbracket / \$st, \$tr \rrbracket)'$

apply (simp add: usubst)

apply (erule taut-shEx-elim)

apply (simp add: unrest-all-circus-vars-st-st' closure unrest assms)

apply (rel-auto)

done

hence $(([\$st \mapsto_s \llbracket s \rrbracket, \$st' \mapsto_s \llbracket s_0 \rrbracket, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \llbracket tr_0 \rrbracket] \vdash post_R P) ;;$
 $([\$st \mapsto_s \llbracket s_0 \rrbracket, \$st' \mapsto_s \llbracket s' \rrbracket, \$tr \mapsto_s \llbracket tr_0 \rrbracket, \$tr' \mapsto_s \llbracket t \rrbracket] \vdash post_R Q))'$

by (rel-auto)

hence $((([\$st \mapsto_s \llbracket s \rrbracket, \$st' \mapsto_s \llbracket s_0 \rrbracket, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \llbracket tr_0 \rrbracket] \vdash post_R P) \wedge$
 $([\$st \mapsto_s \llbracket s_0 \rrbracket, \$st' \mapsto_s \llbracket s' \rrbracket, \$tr \mapsto_s \llbracket tr_0 \rrbracket, \$tr' \mapsto_s \llbracket t \rrbracket] \vdash post_R Q))'$

by (simp add: seqr-to-conj unrest-any-circus-var assms closure unrest)

hence $postP: ([\$st \mapsto_s \llbracket s \rrbracket, \$st' \mapsto_s \llbracket s_0 \rrbracket, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \llbracket tr_0 \rrbracket] \vdash post_R P)'$ and
 $postQ: ([\$st \mapsto_s \llbracket s_0 \rrbracket, \$st' \mapsto_s \llbracket s' \rrbracket, \$tr \mapsto_s \llbracket tr_0 \rrbracket, \$tr' \mapsto_s \llbracket t \rrbracket] \vdash post_R Q)'$

by (rel-auto)+
 from postQ' have '[$\$st \mapsto_s \ll s_0 \gg$, $\$st' \mapsto_s \ll s' \gg$] \dagger [$\$tr \mapsto_s \ll tr_0 \gg$, $\$tr' \mapsto_s \ll tr_0 \gg$] + ($\ll t \gg$ - $\ll tr_0 \gg$) \dagger post_R Q'
 using tr0 by (rel-auto)
 hence '[$\$st \mapsto_s \ll s_0 \gg$, $\$st' \mapsto_s \ll s' \gg$] \dagger [$\$tr \mapsto_s 0$, $\$tr' \mapsto_s \ll t \gg$ - $\ll tr_0 \gg$] \dagger post_R Q'
 by (simp add: R2-subst-tr closure assms)
 hence postQ: '[$\$st \mapsto_s \ll s_0 \gg$, $\$st' \mapsto_s \ll s' \gg$, $\$tr \mapsto_s \langle \rangle$, $\$tr' \mapsto_s \ll t - tr_0 \gg$] \dagger post_R Q'
 by (rel-auto)
 have preP: '[$\$st \mapsto_s \ll s \gg$, $\$tr \mapsto_s \langle \rangle$, $\$tr' \mapsto_s \ll tr_0 \gg$] \dagger pre_R P'
 proof -
 have (pre_R P)[$0, \ll tr_0 \gg / \$tr, \tr'] \sqsubseteq (pre_R P)[$0, \ll t \gg / \$tr, \tr']
 by (simp add: RC-prefix-refine closure assms tr0)
 hence [$\$st \mapsto_s \ll s \gg$, $\$tr \mapsto_s \langle \rangle$, $\$tr' \mapsto_s \ll tr_0 \gg$] \dagger pre_R P \sqsubseteq [$\$st \mapsto_s \ll s \gg$, $\$tr \mapsto_s \langle \rangle$, $\$tr' \mapsto_s \ll t \gg$] \dagger pre_R P
 by (rel-auto)
 thus ?thesis
 by (simp add: taut-refine-impl a2)
 qed

have preQ: '[$\$st \mapsto_s \ll s_0 \gg$, $\$tr \mapsto_s \langle \rangle$, $\$tr' \mapsto_s \ll t - tr_0 \gg$] \dagger pre_R Q'

proof -

from postP a3 have '[$\$st \mapsto_s \ll s_0 \gg$, $\$tr \mapsto_s \ll tr_0 \gg$, $\$tr' \mapsto_s \ll t \gg$] \dagger pre_R Q'

apply (simp add: wp-rea-def)

apply (rel-auto)

using tr0 apply blast+

done

hence '[$\$st \mapsto_s \ll s_0 \gg$] \dagger [$\$tr \mapsto_s \ll tr_0 \gg$, $\$tr' \mapsto_s \ll tr_0 \gg$] + ($\ll t \gg$ - $\ll tr_0 \gg$) \dagger pre_R Q'

by (rel-auto)

hence '[$\$st \mapsto_s \ll s_0 \gg$] \dagger [$\$tr \mapsto_s 0$, $\$tr' \mapsto_s \ll t \gg$ - $\ll tr_0 \gg$] \dagger pre_R Q'

by (simp add: R2-subst-tr closure assms)

thus ?thesis

by (rel-auto)

qed

from a2 have ndiv: \neg '[$\$st \mapsto_s \ll s \gg$, $\$tr \mapsto_s \langle \rangle$, $\$tr' \mapsto_s \ll t \gg$] \dagger (\neg_r pre_R P)'

by (rel-auto)

have t-minus-tr0: $tr_0 @ (t - tr_0) = t$

using append-minus tr0 by blast

from a3

have wpr: $\bigwedge t_0 s_1.$

'[$\$st \mapsto_s \ll s \gg$, $\$tr \mapsto_s \langle \rangle$, $\$tr' \mapsto_s \ll t_0 \gg$] \dagger pre_R P' \implies

'[$\$st \mapsto_s \ll s \gg$, $\$st' \mapsto_s \ll s_1 \gg$, $\$tr \mapsto_s \langle \rangle$, $\$tr' \mapsto_s \ll t_0 \gg$] \dagger post_R P' \implies

$t_0 \leq t \implies$ '[$\$st \mapsto_s \ll s_1 \gg$, $\$tr \mapsto_s \langle \rangle$, $\$tr' \mapsto_s \ll t - t_0 \gg$] \dagger (\neg_r pre_R Q)' \implies False

proof -

fix $t_0 s_1$

assume b:

'[$\$st \mapsto_s \ll s \gg$, $\$tr \mapsto_s \langle \rangle$, $\$tr' \mapsto_s \ll t_0 \gg$] \dagger pre_R P'

'[$\$st \mapsto_s \ll s \gg$, $\$st' \mapsto_s \ll s_1 \gg$, $\$tr \mapsto_s \langle \rangle$, $\$tr' \mapsto_s \ll t_0 \gg$] \dagger post_R P'

$t_0 \leq t$

'[$\$st \mapsto_s \ll s_1 \gg$, $\$tr \mapsto_s \langle \rangle$, $\$tr' \mapsto_s \ll t - t_0 \gg$] \dagger (\neg_r pre_R Q)'

from a3 have c: $\forall (s_0, t_0) \cdot \ll t_0 \gg \leq_u \ll t \gg$

$$\wedge [\$st \mapsto_s \ll s \gg, \$st' \mapsto_s \ll s_0 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_0 \gg] \dagger post_R P$$

$$\Rightarrow [\$st \mapsto_s \ll s_0 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg - \ll t_0 \gg] \dagger pre_R Q'$$

by (simp add: wp-rea-circus-form-alt[of post_R P pre_R Q] closure assms unrest usubst)
 (rel-simp)

from c b(2-4) show False
 by (rel-auto)
 qed

show $\exists t_1 t_2.$
 $t = t_1 @ t_2 \wedge$
 $(\exists s_0. ' [\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_1 \gg] \dagger pre_R P \wedge$
 $[\$st \mapsto_s \ll s \gg, \$st' \mapsto_s \ll s_0 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_1 \gg] \dagger post_R P' \wedge$
 $' [\$st \mapsto_s \ll s_0 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_2 \gg] \dagger pre_R Q \wedge$
 $[\$st \mapsto_s \ll s_0 \gg, \$st' \mapsto_s \ll s' \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_2 \gg] \dagger post_R Q' \wedge$
 $\neg ' [\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_1 @ t_2 \gg] \dagger (\neg_r pre_R P)' \wedge$
 $(\forall t_0 s_1. ' [\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_0 \gg] \dagger pre_R P \wedge$
 $[\$st \mapsto_s \ll s \gg, \$st' \mapsto_s \ll s_1 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_0 \gg] \dagger post_R P' \longrightarrow$
 $t_0 \leq t_1 @ t_2 \longrightarrow \neg ' [\$st \mapsto_s \ll s_1 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll (t_1 @ t_2) - t_0 \gg] \dagger (\neg_r$
 $pre_R Q)')$
 apply (rule-tac x=tr₀ in exI)
 apply (rule-tac x=(t - tr₀) in exI)
 apply (auto)
 using tr0 apply auto[1]
 apply (rule-tac x=s₀ in exI)
 apply (auto intro:wpr simp add: taut-conj preP preQ postP postQ ndiv wpr t-minus-tr0)
 done
 qed

show ?rhs \subseteq ?lhs
 proof (rdes-expand cls: assms, simp add: traces-def divergences-def rdes closure assms rdes-def unrest
 rpred usubst, auto)
 fix t₁ t₂ :: 'e list and s₀ s' :: 's
 assume
 a1: $\neg ' [\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_1 @ t_2 \gg] \dagger (\neg_r pre_R P)'$ and
 a2: $' [\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_1 \gg] \dagger pre_R P'$ and
 a3: $' [\$st \mapsto_s \ll s \gg, \$st' \mapsto_s \ll s_0 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_1 \gg] \dagger post_R P'$ and
 a4: $' [\$st \mapsto_s \ll s_0 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_2 \gg] \dagger pre_R Q'$ and
 a5: $' [\$st \mapsto_s \ll s_0 \gg, \$st' \mapsto_s \ll s' \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_2 \gg] \dagger post_R Q'$ and
 a6: $\forall t s_1. ' [\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \dagger pre_R P \wedge$
 $[\$st \mapsto_s \ll s \gg, \$st' \mapsto_s \ll s_1 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \dagger post_R P' \longrightarrow$
 $t \leq t_1 @ t_2 \longrightarrow \neg ' [\$st \mapsto_s \ll s_1 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll (t_1 @ t_2) - t \gg] \dagger (\neg_r pre_R Q)'$

from a1 have preP: $' [\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_1 @ t_2 \gg] \dagger (pre_R P)'$
 by (simp add: taut-not unrest-all-circus-vars-st assms closure unrest, rel-auto)

have $' [\$st \mapsto_s \ll s_0 \gg, \$st' \mapsto_s \ll s' \gg, \$tr \mapsto_s \ll t_1 \gg, \$tr' \mapsto_s \ll t_1 \gg + \ll t_2 \gg] \dagger post_R Q'$
 proof -
 have $[\$st \mapsto_s \ll s_0 \gg, \$st' \mapsto_s \ll s' \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_2 \gg] \dagger post_R Q =$
 $[\$st \mapsto_s \ll s_0 \gg, \$st' \mapsto_s \ll s' \gg] \dagger [\$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_2 \gg] \dagger post_R Q$
 by rel-auto
 also have ... = $[\$st \mapsto_s \ll s_0 \gg, \$st' \mapsto_s \ll s' \gg] \dagger [\$tr \mapsto_s \ll t_1 \gg, \$tr' \mapsto_s \ll t_1 \gg + \ll t_2 \gg] \dagger post_R Q$
 by (simp add: R2-subst-tr assms closure, rel-auto)
 finally show ?thesis using a5
 by (rel-auto)

qed
with *a3*
have *postPQ*: ‘ $[\$st \mapsto_s \ll s \gg, \$st' \mapsto_s \ll s' \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_1 @ t_2 \gg] \dagger (post_R P ;; post_R Q)'$ ’
by (*rel-auto*, *meson Prefix-Order.prefixI*)

have ‘ $[\$st \mapsto_s \ll s_0 \gg, \$tr \mapsto_s \ll t_1 \gg, \$tr' \mapsto_s \ll t_1 \gg + \ll t_2 \gg] \dagger pre_R Q'$ ’
proof –
have $[\$st \mapsto_s \ll s_0 \gg, \$tr \mapsto_s \ll t_1 \gg, \$tr' \mapsto_s \ll t_1 \gg + \ll t_2 \gg] \dagger pre_R Q =$
 $[\$st \mapsto_s \ll s_0 \gg] \dagger [\$tr \mapsto_s \ll t_1 \gg, \$tr' \mapsto_s \ll t_1 \gg + \ll t_2 \gg] \dagger pre_R Q$
by *rel-auto*
also have $\dots = [\$st \mapsto_s \ll s_0 \gg] \dagger [\$tr \mapsto_s 0, \$tr' \mapsto_s \ll t_2 \gg] \dagger pre_R Q$
by (*simp add: R2-subst-tr assms closure*)
finally show *?thesis using a4*
by (*rel-auto*)
qed

from *a6*
have *a6'*: $\bigwedge t s_1. \llbracket t \leq t_1 @ t_2; '[\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \dagger pre_R P'; '[\$st \mapsto_s \ll s \gg,$
 $\$st' \mapsto_s \ll s_1 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t \gg] \dagger post_R P' \rrbracket \implies$
 $'[\$st \mapsto_s \ll s_1 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll (t_1 @ t_2) - t \gg] \dagger pre_R Q'$
apply (*subst (asm) taut-not*)
apply (*simp add: unrest-all-circus-vars-st assms closure unrest*)
apply (*rel-auto*)
done

have *wpR*: ‘ $[\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_1 @ t_2 \gg] \dagger (post_R P wp_r pre_R Q)'$ ’
proof –
have $\bigwedge s_1 t_0. \llbracket t_0 \leq t_1 @ t_2; '[\$st \mapsto_s \ll s \gg, \$st' \mapsto_s \ll s_1 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_0 \gg] \dagger post_R P'$
 \rrbracket
 $\implies '[\$st \mapsto_s \ll s_1 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll (t_1 @ t_2) - t_0 \gg] \dagger pre_R Q'$
proof –
fix *s1 t0*
assume *c*: $t_0 \leq t_1 @ t_2$ ‘ $[\$st \mapsto_s \ll s \gg, \$st' \mapsto_s \ll s_1 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_0 \gg] \dagger post_R P'$ ’

have *preP'*: ‘ $[\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_0 \gg] \dagger pre_R P'$ ’
proof –
have $(pre_R P) \llbracket 0, \ll t_0 \gg / \$tr, \$tr' \rrbracket \sqsubseteq (pre_R P) \llbracket 0, \ll t_1 @ t_2 \gg / \$tr, \$tr' \rrbracket$
by (*simp add: RC-prefix-refine closure assms c*)
hence $[\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll t_0 \gg] \dagger pre_R P \sqsubseteq [\$st \mapsto_s \ll s \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s$
 $\ll t_1 @ t_2 \gg] \dagger pre_R P$
by (*rel-auto*)
thus *?thesis*
by (*simp add: taut-refine-impl preP*)
qed

with *c a3 preP a6'* **show** ‘ $[\$st \mapsto_s \ll s_1 \gg, \$tr \mapsto_s \langle \rangle, \$tr' \mapsto_s \ll (t_1 @ t_2) - t_0 \gg] \dagger pre_R$
 Q' ’
by (*simp*)
qed

thus *?thesis*
apply (*simp-all add: wp-rea-circus-form-alt assms closure unrest usubst rea-impl-alt-def*)
apply (*simp add: R1-def usubst tcontr-alt-def*)

```

  apply (auto intro!: taut-shAll-intro-2)
  apply (rule taut-impl-intro)
  apply (simp add: unrest-all-circus-vars-st-st' unrest closure assms)
  apply (rel-simp)
done
qed
show '([ $\$st \mapsto_s \ll s \gg$ ,  $\$tr \mapsto_s \langle \rangle$ ,  $\$tr' \mapsto_s \ll t_1 @ t_2 \gg$ ]  $\dagger$   $pre_R P \wedge$ 
  [ $\$st \mapsto_s \ll s \gg$ ,  $\$tr \mapsto_s \langle \rangle$ ,  $\$tr' \mapsto_s \ll t_1 @ t_2 \gg$ ]  $\dagger$  ( $post_R P \wp_r pre_R Q$ ))  $\wedge$ 
  [ $\$st \mapsto_s \ll s \gg$ ,  $\$st' \mapsto_s \ll s' \gg$ ,  $\$tr \mapsto_s \langle \rangle$ ,  $\$tr' \mapsto_s \ll t_1 @ t_2 \gg$ ]  $\dagger$  ( $post_R P ;; post_R Q$ )'
  by (auto simp add: taut-conj preP postPQ wpR)
qed
qed

```

lemma *Cons-minus* [*simp*]: $(a \# t) - [a] = t$
 by (metis append-Cons append-Nil append-minus)

lemma *traces-prefix*:
 assumes P is NCSP
 shows $tr[\ll a \gg \rightarrow_C P]s = \{(a \# t, s') \mid t s'. (t, s') \in tr[P]s\}$
 apply (auto simp add: PrefixCSP-def traces-seq traces-do divergences-do lit.rep-eq assms closure
 Healthy-if trace-divergence-disj)
 apply (meson assms trace-divergence-disj)
 done

10.3 Deadlock Freedom

The following is a specification for deadlock free actions. In any intermediate observation, there must be at least one enabled event.

definition *CDF* :: $('s, 'e)$ action where
 $[rdes-def]$: $CDF = \mathbf{R}_s(true_r \vdash (\prod (s, t, E, e) \cdot \mathcal{E}(\ll s \gg, \ll t \gg, \ll insert\ e\ E \gg)) \diamond true_r)$

lemma *CDF-NCSP* [*closure*]: *CDF* is NCSP
 apply (simp add: CDF-def)
 apply (rule NCSP-rdes-intro)
 apply (simp-all add: closure unrest)
 done

lemma *CDF-seq-idem*: $CDF ;; CDF = CDF$
 by (rdes-eq)

lemma *CDF-refine-intro*: $CDF \sqsubseteq P \implies CDF \sqsubseteq (CDF ;; P)$
 by (metis CDF-seq-idem urel-diod.mult-isol)

lemma *Skip-deadlock-free*: $CDF \sqsubseteq Skip$
 by (rdes-refine)

lemma *CDF-ext-st* [*alpha*]: $CDF \oplus_p abs-st_L = CDF$
 by (rdes-eq)

end

11 Meta-theory for Stateful-Failure Reactive Designs

theory *utp-sf-rdes*

```
imports  
  utp-sfrd-core  
  utp-sfrd-rel  
  utp-sfrd-healths  
  utp-sfrd-contracts  
  utp-sfrd-extchoice  
  utp-sfrd-prog  
  utp-sfrd-recursion  
  utp-sfrd-fdsem  
begin end
```

References

- [1] S. Foster, F. Zeyda, and J. Woodcock. Unifying heterogeneous state-spaces with lenses. In *ICTAC*, LNCS 9965. Springer, 2016.
- [2] M. V. M. Oliveira. *Formal Derivation of State-Rich Reactive Programs using Circus*. PhD thesis, Department of Computer Science - University of York, UK, 2006. YCST-2006-02.