

Reactive Designs in Isabelle/UTP

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Abstract

Reactive designs combine the UTP theories of reactive processes and designs to characterise reactive programs. Whereas sequential imperative programs are expected to run until termination, reactive programs pause at instances to allow interaction with the environment using abstract events, and often do not terminate at all. Thus, whereas a design describes the precondition and postcondition for a program, to characterise initial and final states, a reactive design also has a “pericondition”, which characterises intermediate quiescent observations. This gives rise to a notion of “reactive contract”, which specifies the assumptions a program makes of its environment, and the guarantees it will make of its own behaviour in both intermediate and final observations. This Isabelle/UTP document mechanises the UTP theory of reactive designs, including its healthiness conditions, signature, and a large library of algebraic laws of reactive programming.

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1 Introduction

This document contains a mechanisation in Isabelle/UTP [3] of our theory of reactive designs. Reactive designs form an important semantic foundation for reactive modelling languages such as Circus [5]. For more details of this work, please see our recent paper [2].

2 Reactive Designs Healthiness Conditions

```
theory utp-rdes-healths
  imports UTP-Reactive.utp-reactive
begin
```

2.1 Preliminaries

```
named-theorems rdes and rdes-def and RD-elim
```

```
type-synonym ('s,'t) rdes = ('s,'t,unit) hrel-rsp
```

```
translations
```

```
(type) ('s,'t) rdes <= (type) ('s,'t, unit) hrel-rsp
```

```
lemma R2-st-ex: R2 ( $\exists$  $st  $\cdot$  P) = ( $\exists$  $st  $\cdot$  R2(P))
  by (rel-auto)
```

```
lemma R2s-st'-eq-st:
  R2s($st' =u $st) = ($st' =u $st)
  by (rel-auto)
```

```
lemma R2c-st'-eq-st:
  R2c($st' =u $st) = ($st' =u $st)
  by (rel-auto)
```

```
lemma R1-des-lift-skip: R1( $\lceil$  II  $\rceil_D$ ) =  $\lceil$  II  $\rceil_D$ 
  by (rel-auto)
```

```
lemma R2-des-lift-skip:
  R2( $\lceil$  II  $\rceil_D$ ) =  $\lceil$  II  $\rceil_D$ 
  apply (rel-auto) using minus-zero-eq by blast
```

```
lemma R1-R2c-ex-st: R1 (R2c ( $\exists$  $st'  $\cdot$  Q1)) = ( $\exists$  $st'  $\cdot$  R1 (R2c Q1))
  by (rel-auto)
```

2.2 Identities

We define two identities for reactive designs, which correspond to the regular and state-sensitive versions of reactive designs, respectively. The former is the one used in the UTP book and related publications for CSP.

```
definition skip-rea :: ('t::trace, 'α) hrel-rp (IIc) where
skip-rea-def [urel-defs]: IIc = (II  $\vee$  ( $\neg$  $ok  $\wedge$  $tr  $\leq_u$  $tr'))
```

definition *skip-srea* :: ('s, 't::trace, 'α) hrel-rsp (*II*_R) **where**
skip-srea-def [*urel-defs*]: *II*_R = ((\exists \$st • *II*_c) \triangleleft \$wait \triangleright *II*_c)

lemma *skip-rea-R1-lemma*: *II*_c = *R1*(\$ok \Rightarrow *II*)
by (*rel-auto*)

lemma *skip-rea-form*: *II*_c = (*II* \triangleleft \$ok \triangleright *R1*(*true*))
by (*rel-auto*)

lemma *skip-srea-form*: *II*_R = ((\exists \$st • *II*) \triangleleft \$wait \triangleright *II*) \triangleleft \$ok \triangleright *R1*(*true*)
by (*rel-auto*)

lemma *R1-skip-rea*: *R1*(*II*_c) = *II*_c
by (*rel-auto*)

lemma *R2c-skip-rea*: *R2c* *II*_c = *II*_c
by (*simp add: skip-rea-def R2c-and R2c-disj R2c-skip-r R2c-not R2c-ok R2c-tr'-ge-tr*)

lemma *R2-skip-rea*: *R2*(*II*_c) = *II*_c
by (*metis R1-R2c-is-R2 R1-skip-rea R2c-skip-rea*)

lemma *R2c-skip-srea*: *R2c*(*II*_R) = *II*_R
apply (*rel-auto*) **using** *minus-zero-eq* **by** *blast+*

lemma *skip-srea-R1 [closure]*: *II*_R *is* *R1*
by (*rel-auto*)

lemma *skip-srea-R2c [closure]*: *II*_R *is* *R2c*
by (*simp add: Healthy-def R2c-skip-srea*)

lemma *skip-srea-R2 [closure]*: *II*_R *is* *R2*
by (*metis Healthy-def' R1-R2c-is-R2 R2c-skip-srea skip-srea-R1*)

2.3 RD1: Divergence yields arbitrary traces

definition *RD1* :: ('t::trace, 'α, 'β) rel-rp \Rightarrow ('t, 'α, 'β) rel-rp **where**
[*upred-defs*]: *RD1*(*P*) = (*P* \vee (\neg \$ok \wedge \$tr \leq_u \$tr'))

RD1 is essentially *H1* from the theory of designs, but viewed through the prism of reactive processes.

lemma *RD1-idem*: *RD1*(*RD1*(*P*)) = *RD1*(*P*)
by (*rel-auto*)

lemma *RD1-Idempotent*: *Idempotent* *RD1*
by (*simp add: Idempotent-def RD1-idem*)

lemma *RD1-mono*: *P* \sqsubseteq *Q* \Longrightarrow *RD1*(*P*) \sqsubseteq *RD1*(*Q*)
by (*rel-auto*)

lemma *RD1-Monotonic*: *Monotonic* *RD1*
using *mono-def RD1-mono* **by** *blast*

lemma *RD1-Continuous*: *Continuous* *RD1*
by (*rel-auto*)

lemma *R1-true-RD1-closed* [closure]: $R1(true)$ is *RD1*
by (*rel-auto*)

lemma *RD1-wait-false* [closure]: P is *RD1* $\implies P \llbracket false/\$wait \rrbracket$ is *RD1*
by (*rel-auto*)

lemma *RD1-wait'-false* [closure]: P is *RD1* $\implies P \llbracket false/\$wait' \rrbracket$ is *RD1*
by (*rel-auto*)

lemma *RD1-seq*: $RD1(RD1(P) ;; RD1(Q)) = RD1(P) ;; RD1(Q)$
by (*rel-auto*)

lemma *RD1-seq-closure* [closure]: $\llbracket P \text{ is } RD1; Q \text{ is } RD1 \rrbracket \implies P ;; Q \text{ is } RD1$
by (*metis Healthy-def' RD1-seq*)

lemma *RD1-R1-commute*: $RD1(R1(P)) = R1(RD1(P))$
by (*rel-auto*)

lemma *RD1-R2c-commute*: $RD1(R2c(P)) = R2c(RD1(P))$
by (*rel-auto*)

lemma *RD1-via-R1*: $R1(H1(P)) = RD1(R1(P))$
by (*rel-auto*)

lemma *RD1-R1-cases*: $RD1(R1(P)) = (R1(P) \triangleleft \$ok \triangleright R1(true))$
by (*rel-auto*)

lemma *skip-rea-RD1-skip*: $II_c = RD1(II)$
by (*rel-auto*)

lemma *skip-srea-RD1* [closure]: II_R is *RD1*
by (*rel-auto*)

lemma *RD1-algebraic-intro*:

assumes

P is *R1* ($R1(true_h) ;; P) = R1(true_h) (II_c ;; P) = P$

shows P is *RD1*

proof –

have $P = (II_c ;; P)$

by (*simp add: asms(3)*)

also have $\dots = (R1(\$ok \Rightarrow II) ;; P)$

by (*simp add: skip-rea-R1-lemma*)

also have $\dots = (((\neg \$ok \wedge R1(true)) ;; P) \vee P)$

by (*metis (no-types, lifting) R1-def seqr-left-unit seqr-or-distl skip-rea-R1-lemma skip-rea-def utp-pred-laws.inf-top-left utp-pred-laws.sup-commute*)

also have $\dots = ((R1(\neg \$ok) ;; (R1(true_h) ;; P)) \vee P)$

using *dual-order.trans* **by** (*rel-blast*)

also have $\dots = ((R1(\neg \$ok) ;; R1(true_h)) \vee P)$

by (*simp add: asms(2)*)

also have $\dots = (R1(\neg \$ok) \vee P)$

by (*rel-auto*)

also have $\dots = RD1(P)$

by (*rel-auto*)

finally show *?thesis*

by (simp add: Healthy-def)
qed

theorem *RD1-left-zero*:

assumes P is $R1$ P is $RD1$
shows $(R1(true) ;; P) = R1(true)$

proof –

have $(R1(true) ;; R1(RD1(P))) = R1(true)$
by (rel-auto)
thus ?thesis
by (simp add: Healthy-if assms(1) assms(2))

qed

theorem *RD1-left-unit*:

assumes P is $R1$ P is $RD1$
shows $(II_c ;; P) = P$

proof –

have $(II_c ;; R1(RD1(P))) = R1(RD1(P))$
by (rel-auto)
thus ?thesis
by (simp add: Healthy-if assms(1) assms(2))

qed

lemma *RD1-alt-def*:

assumes P is $R1$
shows $RD1(P) = (P \triangleleft \$ok \triangleright R1(true))$

proof –

have $RD1(R1(P)) = (R1(P) \triangleleft \$ok \triangleright R1(true))$
by (rel-auto)
thus ?thesis
by (simp add: Healthy-if assms)

qed

theorem *RD1-algebraic*:

assumes P is $R1$
shows P is $RD1 \iff (R1(true_h) ;; P) = R1(true_h) \wedge (II_c ;; P) = P$
using *RD1-algebraic-intro RD1-left-unit RD1-left-zero assms* by blast

2.4 R3c and R3h: Reactive design versions of R3

definition $R3c :: ('t::trace, 'α) hrel-rp \Rightarrow ('t, 'α) hrel-rp$ **where**
[upred-defs]: $R3c(P) = (II_c \triangleleft \$wait \triangleright P)$

definition $R3h :: ('s, 't::trace, 'α) hrel-rsp \Rightarrow ('s, 't, 'α) hrel-rsp$ **where**
R3h-def [upred-defs]: $R3h(P) = ((\exists \$st \cdot II_c) \triangleleft \$wait \triangleright P)$

lemma *R3c-idem*: $R3c(R3c(P)) = R3c(P)$
by (rel-auto)

lemma *R3c-Idempotent*: *Idempotent R3c*
by (simp add: Idempotent-def R3c-idem)

lemma *R3c-mono*: $P \sqsubseteq Q \implies R3c(P) \sqsubseteq R3c(Q)$
by (rel-auto)

lemma *R3c-Monotonic*: *Monotonic R3c*

by (simp add: mono-def R3c-mono)

lemma *R3c-Continuous: Continuous R3c*
by (rel-auto)

lemma *R3h-idem: R3h(R3h(P)) = R3h(P)*
by (rel-auto)

lemma *R3h-Idempotent: Idempotent R3h*
by (simp add: Idempotent-def R3h-idem)

lemma *R3h-mono: $P \sqsubseteq Q \implies R3h(P) \sqsubseteq R3h(Q)$*
by (rel-auto)

lemma *R3h-Monotonic: Monotonic R3h*
by (simp add: mono-def R3h-mono)

lemma *R3h-Continuous: Continuous R3h*
by (rel-auto)

lemma *R3h-inf: $R3h(P \sqcap Q) = R3h(P) \sqcap R3h(Q)$*
by (rel-auto)

lemma *R3h-UINF:*
 $A \neq \{\} \implies R3h(\bigsqcap i \in A \cdot P(i)) = (\bigsqcap i \in A \cdot R3h(P(i)))$
by (rel-auto)

lemma *R3h-cond: $R3h(P \triangleleft b \triangleright Q) = (R3h(P) \triangleleft b \triangleright R3h(Q))$*
by (rel-auto)

lemma *R3c-via-RD1-R3: $RD1(R3(P)) = R3c(RD1(P))$*
by (rel-auto)

lemma *R3c-RD1-def: P is RD1 $\implies R3c(P) = RD1(R3(P))$*
by (simp add: Healthy-if R3c-via-RD1-R3)

lemma *RD1-R3c-commute: $RD1(R3c(P)) = R3c(RD1(P))$*
by (rel-auto)

lemma *R1-R3c-commute: $R1(R3c(P)) = R3c(R1(P))$*
by (rel-auto)

lemma *R2c-R3c-commute: $R2c(R3c(P)) = R3c(R2c(P))$*
apply (rel-auto) using minus-zero-eq by blast+

lemma *R1-R3h-commute: $R1(R3h(P)) = R3h(R1(P))$*
by (rel-auto)

lemma *R2c-R3h-commute: $R2c(R3h(P)) = R3h(R2c(P))$*
apply (rel-auto) using minus-zero-eq by blast+

lemma *RD1-R3h-commute: $RD1(R3h(P)) = R3h(RD1(P))$*
by (rel-auto)

lemma *R3c-cancels-R3: $R3c(R3(P)) = R3c(P)$*

by (rel-auto)

lemma *R3-cancels-R3c*: $R3(R3c(P)) = R3(P)$

by (rel-auto)

lemma *R3h-cancels-R3c*: $R3h(R3c(P)) = R3h(P)$

by (rel-auto)

lemma *R3c-semir-form*:

$(R3c(P) ;; R3c(R1(Q))) = R3c(P ;; R3c(R1(Q)))$

by (rel-simp, safe, auto intro: order-trans)

lemma *R3h-semir-form*:

$(R3h(P) ;; R3h(R1(Q))) = R3h(P ;; R3h(R1(Q)))$

by (rel-simp, safe, auto intro: order-trans, blast+)

lemma *R3c-seq-closure*:

assumes P is $R3c$ Q is $R3c$ Q is $R1$

shows $(P ;; Q)$ is $R3c$

by (metis Healthy-def' R3c-semir-form assms)

lemma *R3h-seq-closure* [closure]:

assumes P is $R3h$ Q is $R3h$ Q is $R1$

shows $(P ;; Q)$ is $R3h$

by (metis Healthy-def' R3h-semir-form assms)

lemma *R3c-R3-left-seq-closure*:

assumes P is $R3$ Q is $R3c$

shows $(P ;; Q)$ is $R3c$

proof –

have $(P ;; Q) = ((P ;; Q) \llbracket \text{true}/\$wait \rrbracket \triangleleft \$wait \triangleright (P ;; Q))$

by (metis cond-var-split cond-var-subst-right in-var-uvar wait-vwb-lens)

also have $\dots = (((II \triangleleft \$wait \triangleright P) ;; Q) \llbracket \text{true}/\$wait \rrbracket \triangleleft \$wait \triangleright (P ;; Q))$

by (metis Healthy-def' R3-def assms(1))

also have $\dots = ((II \llbracket \text{true}/\$wait \rrbracket ;; Q) \triangleleft \$wait \triangleright (P ;; Q))$

by (subst-tac)

also have $\dots = (((II \wedge \$wait') ;; Q) \triangleleft \$wait \triangleright (P ;; Q))$

by (metis (no-types, lifting) cond-def conj-pos-var-subst seqr-pre-var-out skip-var utp-pred-laws.inf-left-idem wait-vwb-lens)

also have $\dots = ((II \llbracket \text{true}/\$wait \rrbracket' ;; Q \llbracket \text{true}/\$wait \rrbracket) \triangleleft \$wait \triangleright (P ;; Q))$

by (metis seqr-pre-transfer seqr-right-one-point true-alt-def uovar-convr upred-eq-true utp-rel.unrest-ouvar vwb-lens-mwb wait-vwb-lens)

also have $\dots = ((II \llbracket \text{true}/\$wait \rrbracket' ;; (II_c \triangleleft \$wait \triangleright Q) \llbracket \text{true}/\$wait \rrbracket) \triangleleft \$wait \triangleright (P ;; Q))$

by (metis Healthy-def' R3c-def assms(2))

also have $\dots = ((II \llbracket \text{true}/\$wait \rrbracket' ;; II_c \llbracket \text{true}/\$wait \rrbracket) \triangleleft \$wait \triangleright (P ;; Q))$

by (subst-tac)

also have $\dots = (((II \wedge \$wait') ;; II_c) \triangleleft \$wait \triangleright (P ;; Q))$

by (metis seqr-pre-transfer seqr-right-one-point true-alt-def uovar-convr upred-eq-true utp-rel.unrest-ouvar vwb-lens-mwb wait-vwb-lens)

also have $\dots = ((II ;; II_c) \triangleleft \$wait \triangleright (P ;; Q))$

by (simp add: cond-def seqr-pre-transfer utp-rel.unrest-ouvar)

also have $\dots = (II_c \triangleleft \$wait \triangleright (P ;; Q))$

by simp

also have $\dots = R3c(P ;; Q)$

by (simp add: R3c-def)

finally show ?thesis
 by (simp add: Healthy-def)
 qed

lemma R3c-cases: $R3c(P) = ((II \triangleleft \$ok \triangleright R1(true)) \triangleleft \$wait \triangleright P)$
 by (rel-auto)

lemma R3h-cases: $R3h(P) = (((\exists \$st \cdot II) \triangleleft \$ok \triangleright R1(true)) \triangleleft \$wait \triangleright P)$
 by (rel-auto)

lemma R3h-form: $R3h(P) = II_R \triangleleft \$wait \triangleright P$
 by (rel-auto)

lemma R3c-subst-wait: $R3c(P) = R3c(P_f)$
 by (simp add: R3c-def cond-var-subst-right)

lemma R3h-subst-wait: $R3h(P) = R3h(P_f)$
 by (simp add: R3h-cases cond-var-subst-right)

lemma skip-srea-R3h [closure]: II_R is R3h
 by (rel-auto)

lemma R3h-wait-true:
 assumes P is R3h
 shows $P_t = II_R_t$

proof –
 have $P_t = (II_R \triangleleft \$wait \triangleright P)_t$
 by (metis Healthy-if R3h-form assms)
 also have $\dots = II_R_t$
 by (simp add: usubst)
 finally show ?thesis .
 qed

2.5 RD2: A reactive specification cannot require non-termination

definition RD2 where
 [upred-defs]: $RD2(P) = H2(P)$

RD2 is just H2 since the type system will automatically have J identifying the reactive variables as required.

lemma RD2-idem: $RD2(RD2(P)) = RD2(P)$
 by (simp add: H2-idem RD2-def)

lemma RD2-Idempotent: Idempotent RD2
 by (simp add: Idempotent-def RD2-idem)

lemma RD2-mono: $P \sqsubseteq Q \implies RD2(P) \sqsubseteq RD2(Q)$
 by (simp add: H2-def RD2-def segr-mono)

lemma RD2-Monotonic: Monotonic RD2
 using mono-def RD2-mono by blast

lemma RD2-Continuous: Continuous RD2
 by (rel-auto)

lemma *RD1-RD2-commute*: $RD1(RD2(P)) = RD2(RD1(P))$
 by (*rel-auto*)

lemma *RD2-R3c-commute*: $RD2(R3c(P)) = R3c(RD2(P))$
 by (*rel-auto*)

lemma *RD2-R3h-commute*: $RD2(R3h(P)) = R3h(RD2(P))$
 by (*rel-auto*)

2.6 Major healthiness conditions

definition *RH* :: $(t::trace, 'α) \text{ hrel-rp} \Rightarrow (t, 'α) \text{ hrel-rp } (\mathbf{R})$
where [*upred-defs*]: $RH(P) = R1(R2c(R3c(P)))$

definition *RHS* :: $(s, t::trace, 'α) \text{ hrel-rsp} \Rightarrow (s, t, 'α) \text{ hrel-rsp } (\mathbf{R}_s)$
where [*upred-defs*]: $RHS(P) = R1(R2c(R3h(P)))$

definition *RD* :: $(t::trace, 'α) \text{ hrel-rp} \Rightarrow (t, 'α) \text{ hrel-rp}$
where [*upred-defs*]: $RD(P) = RD1(RD2(RP(P)))$

definition *SRD* :: $(s, t::trace, 'α) \text{ hrel-rsp} \Rightarrow (s, t, 'α) \text{ hrel-rsp}$
where [*upred-defs*]: $SRD(P) = RD1(RD2(RHS(P)))$

lemma *RH-comp*: $RH = R1 \circ R2c \circ R3c$
 by (*auto simp add: RH-def*)

lemma *RHS-comp*: $RHS = R1 \circ R2c \circ R3h$
 by (*auto simp add: RHS-def*)

lemma *RD-comp*: $RD = RD1 \circ RD2 \circ RP$
 by (*auto simp add: RD-def*)

lemma *SRD-comp*: $SRD = RD1 \circ RD2 \circ RHS$
 by (*auto simp add: SRD-def*)

lemma *RH-idem*: $\mathbf{R}(\mathbf{R}(P)) = \mathbf{R}(P)$
 by (*simp add: R1-R2c-commute R1-R3c-commute R1-idem R2c-R3c-commute R2c-idem R3c-idem RH-def*)

lemma *RH-Idempotent*: *Idempotent* \mathbf{R}
 by (*simp add: Idempotent-def RH-idem*)

lemma *RH-Monotonic*: *Monotonic* \mathbf{R}
 by (*metis (no-types, lifting) R1-Monotonic R2c-Monotonic R3c-mono RH-def mono-def*)

lemma *RH-Continuous*: *Continuous* \mathbf{R}
 by (*simp add: Continuous-comp R1-Continuous R2c-Continuous R3c-Continuous RH-comp*)

lemma *RHS-idem*: $\mathbf{R}_s(\mathbf{R}_s(P)) = \mathbf{R}_s(P)$
 by (*simp add: R1-R2c-is-R2 R1-R3h-commute R2-idem R2c-R3h-commute R3h-idem RHS-def*)

lemma *RHS-Idempotent* [*closure*]: *Idempotent* \mathbf{R}_s
 by (*simp add: Idempotent-def RHS-idem*)

lemma *RHS-Monotonic*: *Monotonic* \mathbf{R}_s
 by (*simp add: mono-def R1-R2c-is-R2 R2-mono R3h-mono RHS-def*)

lemma *RHS-mono*: $P \sqsubseteq Q \implies \mathbf{R}_s(P) \sqsubseteq \mathbf{R}_s(Q)$
using *mono-def RHS-Monotonic* **by** *blast*

lemma *RHS-Continuous* [*closure*]: *Continuous* \mathbf{R}_s
by (*simp add: Continuous-comp R1-Continuous R2c-Continuous R3h-Continuous RHS-comp*)

lemma *RHS-inf*: $\mathbf{R}_s(P \sqcap Q) = \mathbf{R}_s(P) \sqcap \mathbf{R}_s(Q)$
using *Continuous-Disjunctous Disjunctuous-def RHS-Continuous* **by** *auto*

lemma *RHS-INF*:
 $A \neq \{\} \implies \mathbf{R}_s(\bigsqcap i \in A \cdot P(i)) = (\bigsqcap i \in A \cdot \mathbf{R}_s(P(i)))$
by (*simp add: RHS-def R3h-UNIF R2c-USUP R1-USUP*)

lemma *RHS-sup*: $\mathbf{R}_s(P \sqcup Q) = \mathbf{R}_s(P) \sqcup \mathbf{R}_s(Q)$
by (*rel-auto*)

lemma *RHS-SUP*:
 $A \neq \{\} \implies \mathbf{R}_s(\bigsqcup i \in A \cdot P(i)) = (\bigsqcup i \in A \cdot \mathbf{R}_s(P(i)))$
by (*rel-auto*)

lemma *RHS-cond*: $\mathbf{R}_s(P \triangleleft b \triangleright Q) = (\mathbf{R}_s(P) \triangleleft R2c\ b \triangleright \mathbf{R}_s(Q))$
by (*simp add: RHS-def R3h-cond R2c-condr R1-cond*)

lemma *RD-alt-def*: $RD(P) = RD1(RD2(\mathbf{R}(P)))$
by (*simp add: R3c-via-RD1-R3 RD1-R1-commute RD1-R2c-commute RD1-R3c-commute RD1-RD2-commute RH-def RD-def RP-def*)

lemma *RD1-RH-commute*: $RD1(\mathbf{R}(P)) = \mathbf{R}(RD1(P))$
by (*simp add: RD1-R1-commute RD1-R2c-commute RD1-R3c-commute RH-def*)

lemma *RD2-RH-commute*: $RD2(\mathbf{R}(P)) = \mathbf{R}(RD2(P))$
by (*metis R1-H2-commute R2c-H2-commute RD2-R3c-commute RD2-def RH-def*)

lemma *RD-idem*: $RD(RD(P)) = RD(P)$
by (*simp add: RD-alt-def RD1-RH-commute RD2-RH-commute RD1-RD2-commute RD2-idem RD1-idem RH-idem*)

lemma *RD-Monotonic*: *Monotonic* RD
by (*simp add: Monotonic-comp RD1-Monotonic RD2-Monotonic RD-comp RP-Monotonic*)

lemma *RD-Continuous*: *Continuous* RD
by (*simp add: Continuous-comp RD1-Continuous RD2-Continuous RD-comp RP-Continuous*)

lemma *R3-RD-RP*: $R3(RD(P)) = RP(RD1(RD2(P)))$
by (*metis (no-types, lifting) R1-R2c-is-R2 R2-R3-commute R3-cancels-R3c RD1-RH-commute RD2-RH-commute RD-alt-def RH-def RP-def*)

lemma *RD1-RHS-commute*: $RD1(\mathbf{R}_s(P)) = \mathbf{R}_s(RD1(P))$
by (*simp add: RD1-R1-commute RD1-R2c-commute RD1-R3h-commute RHS-def*)

lemma *RD2-RHS-commute*: $RD2(\mathbf{R}_s(P)) = \mathbf{R}_s(RD2(P))$
by (*metis R1-H2-commute R2c-H2-commute RD2-R3h-commute RD2-def RHS-def*)

lemma *SRD-idem*: $SRD(SRD(P)) = SRD(P)$

by (*simp add: RD1-RD2-commute RD1-RHS-commute RD1-idem RD2-RHS-commute RD2-idem RHS-idem SRD-def*)

lemma *SRD-Idempotent* [*closure*]: *Idempotent SRD*
by (*simp add: Idempotent-def SRD-idem*)

lemma *SRD-Monotonic*: *Monotonic SRD*
by (*simp add: Monotonic-comp RD1-Monotonic RD2-Monotonic RHS-Monotonic SRD-comp*)

lemma *SRD-Continuous* [*closure*]: *Continuous SRD*
by (*simp add: Continuous-comp RD1-Continuous RD2-Continuous RHS-Continuous SRD-comp*)

lemma *SRD-RHS-H1-H2*: $SRD(P) = \mathbf{R}_s(\mathbf{H}(P))$
by (*rel-auto*)

lemma *SRD-healths* [*closure*]:
assumes *P is SRD*
shows *P is R1 P is R2 P is R3h P is RD1 P is RD2*
apply (*metis Healthy-def R1-idem RD1-RHS-commute RD2-RHS-commute RHS-def SRD-def assms*)
apply (*metis Healthy-def R1-R2c-is-R2 R2-idem RD1-RHS-commute RD2-RHS-commute RHS-def SRD-def assms*)
apply (*metis Healthy-def R1-R3h-commute R2c-R3h-commute R3h-idem RD1-R3h-commute RD2-R3h-commute RHS-def SRD-def assms*)
apply (*metis Healthy-def' RD1-idem SRD-def assms*)
apply (*metis Healthy-def' RD1-RD2-commute RD2-idem SRD-def assms*)
done

lemma *SRD-intro*:
assumes *P is R1 P is R2 P is R3h P is RD1 P is RD2*
shows *P is SRD*
by (*metis Healthy-def R1-R2c-is-R2 RHS-def SRD-def assms(2) assms(3) assms(4) assms(5)*)

lemma *SRD-ok-false* [*usubst*]: $P \text{ is } SRD \implies P \llbracket \text{false}/\$ok \rrbracket = R1(\text{true})$
by (*metis (no-types, hide-lams) H1-H2-eq-design Healthy-def R1-ok-false RD1-R1-commute RD1-via-R1 RD2-def SRD-def SRD-healths(1) design-ok-false*)

lemma *SRD-ok-true-wait-true* [*usubst*]:
assumes *P is SRD*
shows $P \llbracket \text{true}, \text{true}/\$ok, \$wait \rrbracket = (\exists \$st \cdot II) \llbracket \text{true}, \text{true}/\$ok, \$wait \rrbracket$
proof –
have $P = (\exists \$st \cdot II) \triangleleft \$ok \triangleright R1 \text{ true} \triangleleft \$wait \triangleright P$
by (*metis Healthy-def R3h-cases SRD-healths(3) assms*)
moreover have $((\exists \$st \cdot II) \triangleleft \$ok \triangleright R1 \text{ true} \triangleleft \$wait \triangleright P) \llbracket \text{true}, \text{true}/\$ok, \$wait \rrbracket = (\exists \$st \cdot II) \llbracket \text{true}, \text{true}/\$ok, \$wait \rrbracket$
by (*simp add: usubst*)
ultimately show *?thesis*
by (*simp*)
qed

lemma *SRD-left-zero-1*: $P \text{ is } SRD \implies R1(\text{true}) ;; P = R1(\text{true})$
by (*simp add: RD1-left-zero SRD-healths(1) SRD-healths(4)*)

lemma *SRD-left-zero-2*:
assumes *P is SRD*
shows $(\exists \$st \cdot II) \llbracket \text{true}, \text{true}/\$ok, \$wait \rrbracket ;; P = (\exists \$st \cdot II) \llbracket \text{true}, \text{true}/\$ok, \$wait \rrbracket$

proof –

have $(\exists \$st \cdot II) \llbracket true, true / \$ok, \$wait \rrbracket \;; R3h(P) = (\exists \$st \cdot II) \llbracket true, true / \$ok, \$wait \rrbracket$
by *(rel-auto)*
thus *?thesis*
by *(simp add: Healthy-if SRD-healths(3) assms)*
qed

2.7 UTP theories

We create two theory objects: one for reactive designs and one for stateful reactive designs.

typeddecl *RDES*
typeddecl *SRDES*

abbreviation *RDES* $\equiv UTHY(RDES, ('t::trace, 'α) rp)$
abbreviation *SRDES* $\equiv UTHY(SRDES, ('s, 't::trace, 'α) rsp)$

overloading

rdes-hcond $\equiv utp-hcond :: (RDES, ('t::trace, 'α) rp) uthy \Rightarrow (('t, 'α) rp \times ('t, 'α) rp) health$
srdes-hcond $\equiv utp-hcond :: (SRDES, ('s, 't::trace, 'α) rsp) uthy \Rightarrow (('s, 't, 'α) rsp \times ('s, 't, 'α) rsp) health$
begin
definition *rdes-hcond* $:: (RDES, ('t::trace, 'α) rp) uthy \Rightarrow (('t, 'α) rp \times ('t, 'α) rp) health$ **where**
[upred-defs]: rdes-hcond T = RD
definition *srdes-hcond* $:: (SRDES, ('s, 't::trace, 'α) rsp) uthy \Rightarrow (('s, 't, 'α) rsp \times ('s, 't, 'α) rsp) health$
where
[upred-defs]: srdes-hcond T = SRD
end

interpretation *rdes-theory*: *utp-theory* *UTHY*(*RDES*, $('t::trace, 'α) rp$)
by *(unfold-locales, simp-all add: rdes-hcond-def RD-idem)*

interpretation *rdes-theory-continuous*: *utp-theory-continuous* *UTHY*(*RDES*, $('t::trace, 'α) rp$)
rewrites $\bigwedge P. P \in carrier (uthy-order RDES) \longleftrightarrow P \text{ is } RD$
and *carrier* $(uthy-order RDES) \rightarrow carrier (uthy-order RDES) \equiv \llbracket RD \rrbracket_H \rightarrow \llbracket RD \rrbracket_H$
and *le* $(uthy-order RDES) = (\sqsubseteq)$
and *eq* $(uthy-order RDES) = (=)$
by *(unfold-locales, simp-all add: rdes-hcond-def RD-Continuous)*

interpretation *rdes-rea-galois*:

galois-connection $(RDES \leftarrow \langle RD1 \circ RD2, R3 \rangle \rightarrow REA)$
proof *(simp add: mk-conn-def, rule galois-connectionI', simp-all add: utp-partial-order rdes-hcond-def rea-hcond-def)*
show $R3 \in \llbracket RD \rrbracket_H \rightarrow \llbracket RP \rrbracket_H$
by *(metis (no-types, lifting) Healthy-def' Pi-I R3-RD-RP RP-idem mem-Collect-eq)*
show $RD1 \circ RD2 \in \llbracket RP \rrbracket_H \rightarrow \llbracket RD \rrbracket_H$
by *(simp add: Pi-iff Healthy-def, metis RD-def RD-idem)*
show *isotone* $(utp-order RD) (utp-order RP) R3$
by *(simp add: R3-Monotonic isotone-utp-orderI)*
show *isotone* $(utp-order RP) (utp-order RD) (RD1 \circ RD2)$
by *(simp add: Monotonic-comp RD1-Monotonic RD2-Monotonic isotone-utp-orderI)*
fix $P :: ('a, 'b) hrel-rp$
assume *P is RD*
thus $P \sqsubseteq RD1 (RD2 (R3 P))$
by *(metis Healthy-if R3-RD-RP RD-def RP-idem eq-iff)*
next

```

fix  $P :: ('a, 'b) \text{ hrel-rp}$ 
assume  $a: P \text{ is } RP$ 
thus  $R3 (RD1 (RD2 P)) \sqsubseteq P$ 
proof –
  have  $R3 (RD1 (RD2 P)) = RP (RD1 (RD2(P)))$ 
    by (metis Healthy-if R3-RD-RP RD-def a)
  moreover have  $RD1(RD2(P)) \sqsubseteq P$ 
    by (rel-auto)
  ultimately show ?thesis
    by (metis Healthy-if RP-mono a)
qed
qed

interpretation rdes-rea-retract:
  retract (RDES  $\leftarrow \langle RD1 \circ RD2, R3 \rangle \rightarrow$  REA)
  by (unfold-locales, simp-all add: mk-conn-def utp-partial-order rdes-hcond-def rea-hcond-def)
    (metis Healthy-if R3-RD-RP RD-def RP-idem eq-refl)

interpretation srdes-theory: utp-theory UTHY(SRDES, ('s, 't::trace, 'α) rsp)
  by (unfold-locales, simp-all add: srdes-hcond-def SRD-idem)

interpretation srdes-theory-continuous: utp-theory-continuous UTHY(SRDES, ('s, 't::trace, 'α) rsp)
  rewrites  $\bigwedge P. P \in \text{carrier (uthy-order SRDES)} \longleftrightarrow P \text{ is SRD}$ 
  and  $P \text{ is } \mathcal{H}_{SRDES} \longleftrightarrow P \text{ is SRD}$ 
  and  $(\mu X \cdot F (\mathcal{H}_{SRDES} X)) = (\mu X \cdot F (SRD X))$ 
  and  $\text{carrier (uthy-order SRDES)} \rightarrow \text{carrier (uthy-order SRDES)} \equiv \llbracket SRD \rrbracket_H \rightarrow \llbracket SRD \rrbracket_H$ 
  and  $\llbracket \mathcal{H}_{SRDES} \rrbracket_H \rightarrow \llbracket \mathcal{H}_{SRDES} \rrbracket_H \equiv \llbracket SRD \rrbracket_H \rightarrow \llbracket SRD \rrbracket_H$ 
  and  $le \text{ (uthy-order SRDES)} = (\sqsubseteq)$ 
  and  $eq \text{ (uthy-order SRDES)} = (=)$ 
  by (unfold-locales, simp-all add: srdes-hcond-def SRD-Continuous)

declare srdes-theory-continuous.top-healthy [simp del]
declare srdes-theory-continuous.bottom-healthy [simp del]

abbreviation Chaos :: ('s, 't::trace, 'α) hrel-rsp where
  Chaos  $\equiv \perp_{SRDES}$ 

abbreviation Miracle :: ('s, 't::trace, 'α) hrel-rsp where
  Miracle  $\equiv \top_{SRDES}$ 

thm srdes-theory-continuous.weak.bottom-lower
thm srdes-theory-continuous.weak.top-higher
thm srdes-theory-continuous.meet-bottom
thm srdes-theory-continuous.meet-top

abbreviation srd-lfp ( $\mu_R$ ) where  $\mu_R F \equiv \mu_{SRDES} F$ 

abbreviation srd-gfp ( $\nu_R$ ) where  $\nu_R F \equiv \nu_{SRDES} F$ 

syntax
  -srd-mu :: pttrn  $\Rightarrow$  logic  $\Rightarrow$  logic ( $\mu_R \cdot \cdot \cdot [0, 10] 10$ )
  -srd-nu :: pttrn  $\Rightarrow$  logic  $\Rightarrow$  logic ( $\nu_R \cdot \cdot \cdot [0, 10] 10$ )

translations
   $\mu_R X \cdot P == \mu_R (\lambda X. P)$ 

```

$$\nu_R X \cdot P == \mu_R (\lambda X. P)$$

The reactive design weakest fixed-point can be defined in terms of relational calculus one.

lemma *srd-mu-equiv*:

assumes *Monotonic* $F F \in \llbracket SRD \rrbracket_H \rightarrow \llbracket SRD \rrbracket_H$
shows $(\mu_R X \cdot F(X)) = (\mu X \cdot F(SRD(X)))$
by (*metis assms srdes-hcond-def srdes-theory-continuous.utp-lfp-def*)

end

3 Reactive Design Specifications

theory *utp-rdes-designs*

imports *utp-rdes-healths*

begin

3.1 Reactive design forms

lemma *srdes-skip-def*: $II_R = \mathbf{R}_s(\text{true} \vdash (\$tr' =_u \$tr \wedge \neg \$wait' \wedge \lceil II \rceil_R))$
apply (*rel-auto*) **using** *minus-zero-eq* **by** *blast+*

lemma *Chaos-def*: $\text{Chaos} = \mathbf{R}_s(\text{false} \vdash \text{true})$

proof –

have $\text{Chaos} = SRD(\text{true})$
by (*metis srdes-hcond-def srdes-theory-continuous.healthy-bottom*)
also have $\dots = \mathbf{R}_s(\mathbf{H}(\text{true}))$
by (*simp add: SRD-RHS-H1-H2*)
also have $\dots = \mathbf{R}_s(\text{false} \vdash \text{true})$
by (*metis H1-design H2-true design-false-pre*)
finally show *?thesis* .

qed

lemma *Miracle-def*: $\text{Miracle} = \mathbf{R}_s(\text{true} \vdash \text{false})$

proof –

have $\text{Miracle} = SRD(\text{false})$
by (*metis srdes-hcond-def srdes-theory-continuous.healthy-top*)
also have $\dots = \mathbf{R}_s(\mathbf{H}(\text{false}))$
by (*simp add: SRD-RHS-H1-H2*)
also have $\dots = \mathbf{R}_s(\text{true} \vdash \text{false})$
by (*metis (no-types, lifting) H1-H2-eq-design p-imp-p subst-impl subst-not utp-pred-laws.compl-bot-eq utp-pred-laws.compl-top-eq*)
finally show *?thesis* .

qed

lemma *RD1-reactive-design*: $RD1(\mathbf{R}(P \vdash Q)) = \mathbf{R}(P \vdash Q)$

by (*rel-auto*)

lemma *RD2-reactive-design*:

assumes $\$ok' \nmid P \ \$ok' \nmid Q$
shows $RD2(\mathbf{R}(P \vdash Q)) = \mathbf{R}(P \vdash Q)$
using *assms*
by (*metis H2-design RD2-RH-commute RD2-def*)

lemma *RD1-st-reactive-design*: $RD1(\mathbf{R}_s(P \vdash Q)) = \mathbf{R}_s(P \vdash Q)$

by (*rel-auto*)

lemma *RD2-st-reactive-design*:
 assumes $\$ok' \# P \ \$ok' \# Q$
 shows $RD2(\mathbf{R}_s(P \vdash Q)) = \mathbf{R}_s(P \vdash Q)$
 using *assms*
 by (metis *H2-design RD2-RHS-commute RD2-def*)

lemma *wait-false-design*:
 $(P \vdash Q)_f = ((P_f) \vdash (Q_f))$
 by (rel-auto)

lemma *RD-RH-design-form*:
 $RD(P) = \mathbf{R}((\neg P^f_f) \vdash P^t_f)$

proof –

have $RD(P) = RD1(RD2(R1(R2c(R3c(P)))))$
 by (simp add: *RD-alt-def RH-def*)
 also have $\dots = RD1(H2(R1(R2s(R3c(P)))))$
 by (simp add: *R1-R2s-R2c RD2-def*)
 also have $\dots = RD1(R1(H2(R2s(R3c(P)))))$
 by (simp add: *R1-H2-commute*)
 also have $\dots = R1(H1(R1(H2(R2s(R3c(P))))))$
 by (simp add: *R1-idem RD1-via-R1*)
 also have $\dots = R1(H1(H2(R2s(R3c(R1(P))))))$
 by (simp add: *R1-H2-commute R1-R2c-commute R1-R2s-R2c R1-R3c-commute RD1-via-R1*)
 also have $\dots = R1(R2s(H1(H2(R3c(R1(P))))))$
 by (simp add: *R2s-H1-commute R2s-H2-commute*)
 also have $\dots = R1(R2s(H1(R3c(H2(R1(P))))))$
 by (metis *RD2-R3c-commute RD2-def*)
 also have $\dots = R2(R1(H1(R3c(H2(R1(P))))))$
 by (metis *R1-R2-commute R1-idem R2-def*)
 also have $\dots = R2(R3c(R1(\mathbf{H}(R1(P)))))$
 by (simp add: *R1-R3c-commute RD1-R3c-commute RD1-via-R1*)
 also have $\dots = RH(\mathbf{H}(R1(P)))$
 by (metis *R1-R2s-R2c R1-R3c-commute R2-R1-form RH-def*)
 also have $\dots = RH(\mathbf{H}(P))$
 by (simp add: *R1-H2-commute R1-R2c-commute R1-R3c-commute R1-idem RD1-via-R1 RH-def*)
 also have $\dots = RH((\neg P^f_f) \vdash P^t_f)$
 by (simp add: *H1-H2-eq-design*)
 also have $\dots = \mathbf{R}((\neg P^f_f) \vdash P^t_f)$
 by (metis (no-types, lifting) *R3c-subst-wait RH-def subst-not wait-false-design*)
 finally show ?thesis .

qed

lemma *RD-reactive-design*:
 assumes *P is RD*
 shows $\mathbf{R}((\neg P^f_f) \vdash P^t_f) = P$
 by (metis *RD-RH-design-form Healthy-def' assms*)

lemma *RD-RH-design*:
 assumes $\$ok' \# P \ \$ok' \# Q$
 shows $RD(\mathbf{R}(P \vdash Q)) = \mathbf{R}(P \vdash Q)$
 by (simp add: *RD1-reactive-design RD2-reactive-design RD-alt-def RH-idem assms(1) assms(2)*)

lemma *RH-design-is-RD*:
 assumes $\$ok' \# P \ \$ok' \# Q$

shows $\mathbf{R}(P \vdash Q)$ is RD
 by (simp add: $RD\text{-}RH\text{-}design$ $Healthy\text{-}def'$ $assms(1)$ $assms(2)$)

lemma $SRD\text{-}RH\text{-}design\text{-}form$:

$SRD(P) = \mathbf{R}_s((\neg P^f_f) \vdash P^t_f)$

proof –

have $SRD(P) = R1(R2c(R3h(RD1(RD2(R1(P)))))$

by (metis (no-types, lifting) $R1\text{-}H2\text{-}commute$ $R1\text{-}R2c\text{-}commute$ $R1\text{-}R3h\text{-}commute$ $R1\text{-}idem$ $R2c\text{-}H2\text{-}commute$ $RD1\text{-}R1\text{-}commute$ $RD1\text{-}R2c\text{-}commute$ $RD1\text{-}R3h\text{-}commute$ $RD2\text{-}R3h\text{-}commute$ $RD2\text{-}def$ $RHS\text{-}def$ $SRD\text{-}def$)

also have $\dots = R1(R2s(R3h(\mathbf{H}(P))))$

by (metis (no-types, lifting) $R1\text{-}H2\text{-}commute$ $R1\text{-}R2c\text{-}is\text{-}R2$ $R1\text{-}R3h\text{-}commute$ $R2\text{-}R1\text{-}form$ $RD1\text{-}via\text{-}R1$ $RD2\text{-}def$)

also have $\dots = \mathbf{R}_s(\mathbf{H}(P))$

by (simp add: $R1\text{-}R2s\text{-}R2c$ $RHS\text{-}def$)

also have $\dots = \mathbf{R}_s((\neg P^f_f) \vdash P^t_f)$

by (simp add: $H1\text{-}H2\text{-}eq\text{-}design$)

also have $\dots = \mathbf{R}_s((\neg P^f_f) \vdash P^t_f)$

by (metis (no-types, lifting) $R3h\text{-}subst\text{-}wait$ $RHS\text{-}def$ $subst\text{-}not$ $wait\text{-}false\text{-}design$)

finally show ?thesis .

qed

lemma $SRD\text{-}reactive\text{-}design$:

assumes P is SRD

shows $\mathbf{R}_s((\neg P^f_f) \vdash P^t_f) = P$

by (metis $SRD\text{-}RH\text{-}design\text{-}form$ $Healthy\text{-}def'$ $assms$)

lemma $SRD\text{-}RH\text{-}design$:

assumes $\$ok' \# P$ $\$ok' \# Q$

shows $SRD(\mathbf{R}_s(P \vdash Q)) = \mathbf{R}_s(P \vdash Q)$

by (simp add: $RD1\text{-}st\text{-}reactive\text{-}design$ $RD2\text{-}st\text{-}reactive\text{-}design$ $RHS\text{-}idem$ $SRD\text{-}def$ $assms(1)$ $assms(2)$)

lemma $RHS\text{-}design\text{-}is\text{-}SRD$:

assumes $\$ok' \# P$ $\$ok' \# Q$

shows $\mathbf{R}_s(P \vdash Q)$ is SRD

by (simp add: $Healthy\text{-}def'$ $SRD\text{-}RH\text{-}design$ $assms(1)$ $assms(2)$)

lemma $SRD\text{-}RHS\text{-}H1\text{-}H2$: $SRD(P) = \mathbf{R}_s(\mathbf{H}(P))$

by (metis (no-types, lifting) $H1\text{-}H2\text{-}eq\text{-}design$ $R3h\text{-}subst\text{-}wait$ $RHS\text{-}def$ $SRD\text{-}RH\text{-}design\text{-}form$ $subst\text{-}not$ $wait\text{-}false\text{-}design$)

3.2 Auxiliary healthiness conditions

definition [upred-defs]: $R3c\text{-}pre(P) = (true \triangleleft \$wait \triangleright P)$

definition [upred-defs]: $R3c\text{-}post(P) = (\lceil II \rceil_D \triangleleft \$wait \triangleright P)$

definition [upred-defs]: $R3h\text{-}post(P) = ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright P)$

lemma $R3c\text{-}pre\text{-}conj$: $R3c\text{-}pre(P \wedge Q) = (R3c\text{-}pre(P) \wedge R3c\text{-}pre(Q))$

by (rel-auto)

lemma $R3c\text{-}pre\text{-}seq$:

$(true ;; Q) = true \implies R3c\text{-}pre(P ;; Q) = (R3c\text{-}pre(P) ;; Q)$

by (rel-auto)

lemma $unrest\text{-}ok\text{-}R3c\text{-}pre$ [unrest]: $\$ok \# P \implies \$ok \# R3c\text{-}pre(P)$

by (simp add: R3c-pre-def cond-def unrest)

lemma *unrest-ok'-R3c-pre* [unrest]: $\$ok' \# P \implies \$ok' \# R3c\text{-pre}(P)$
 by (simp add: R3c-pre-def cond-def unrest)

lemma *unrest-ok-R3c-post* [unrest]: $\$ok \# P \implies \$ok \# R3c\text{-post}(P)$
 by (simp add: R3c-post-def cond-def unrest)

lemma *unrest-ok-R3c-post'* [unrest]: $\$ok' \# P \implies \$ok' \# R3c\text{-post}(P)$
 by (simp add: R3c-post-def cond-def unrest)

lemma *unrest-ok-R3h-post* [unrest]: $\$ok \# P \implies \$ok \# R3h\text{-post}(P)$
 by (simp add: R3h-post-def cond-def unrest)

lemma *unrest-ok-R3h-post'* [unrest]: $\$ok' \# P \implies \$ok' \# R3h\text{-post}(P)$
 by (simp add: R3h-post-def cond-def unrest)

3.3 Composition laws

theorem *R1-design-composition*:

fixes $P\ Q :: ('t::\text{trace}, 'a, 'b) \text{ rel-rp}$

and $R\ S :: ('t, 'b, 'c) \text{ rel-rp}$

assumes $\$ok' \# P\ \$ok' \# Q\ \$ok \# R\ \$ok \# S$

shows

$(R1(P \vdash Q) ;; R1(R \vdash S)) =$
 $R1((\neg (R1(\neg P) ;; R1(true)) \wedge \neg (R1(Q) ;; R1(\neg R))) \vdash (R1(Q) ;; R1(S)))$

proof –

have $(R1(P \vdash Q) ;; R1(R \vdash S)) = ((R1(P \vdash Q))^t ;; R1(R \vdash S) \llbracket true/\$ok \rrbracket \vee (R1(P \vdash Q))^f ;; R1(R \vdash S) \llbracket false/\$ok \rrbracket)$

by (rule segr-bool-split[of ok], simp)

also from *assms* have $\dots = ((R1((\$ok \wedge P) \Rightarrow (true \wedge Q)) ;; R1((true \wedge R) \Rightarrow (\$ok' \wedge S)))$
 $\vee (R1((\$ok \wedge P) \Rightarrow (false \wedge Q)) ;; R1((false \wedge R) \Rightarrow (\$ok' \wedge S)))$

by (simp add: design-def usubst R1-def)

also from *assms* have $\dots = ((R1((\$ok \wedge P) \Rightarrow Q) ;; R1(R \Rightarrow (\$ok' \wedge S)))$
 $\vee (R1(\neg (\$ok \wedge P)) ;; R1(true)))$

by *simp*

also from *assms* have $\dots = ((R1(\neg \$ok \vee \neg P \vee Q) ;; R1(\neg R \vee (\$ok' \wedge S)))$
 $\vee (R1(\neg \$ok \vee \neg P) ;; R1(true)))$

by (simp add: impl-alt-def utp-pred-laws.sup.assoc)

also from *assms* have $\dots = (((R1(\neg \$ok \vee \neg P) \vee R1(Q)) ;; R1(\neg R \vee (\$ok' \wedge S)))$
 $\vee (R1(\neg \$ok \vee \neg P) ;; R1(true)))$

by (simp add: R1-disj utp-pred-laws.disj-assoc)

also from *assms* have $\dots = ((R1(\neg \$ok \vee \neg P) ;; R1(\neg R \vee (\$ok' \wedge S)))$
 $\vee (R1(Q) ;; R1(\neg R \vee (\$ok' \wedge S)))$
 $\vee (R1(\neg \$ok \vee \neg P) ;; R1(true)))$

by (simp add: segr-or-distl utp-pred-laws.sup.assoc)

also from *assms* have $\dots = ((R1(Q) ;; R1(\neg R \vee (\$ok' \wedge S)))$
 $\vee (R1(\neg \$ok \vee \neg P) ;; R1(true)))$

by (rel-blast)

also from *assms* have $\dots = ((R1(Q) ;; (R1(\neg R) \vee R1(S) \wedge \$ok'))$
 $\vee (R1(\neg \$ok \vee \neg P) ;; R1(true)))$

by (simp add: R1-disj R1-extend-conj utp-pred-laws.inf-commute)

also have $\dots = ((R1(Q) ;; (R1(\neg R) \vee R1(S) \wedge \$ok'))$
 $\vee ((R1(\neg \$ok) :: ('t, 'a, 'b) \text{ rel-rp}) ;; R1(true)))$
 $\vee (R1(\neg P) ;; R1(true)))$

by (simp add: R1-disj segr-or-distl)

```

also have ... = ((R1(Q) ;; (R1(¬ R) ∨ R1(S) ∧ $ok'))
  ∨ (R1(¬ $ok))
  ∨ (R1(¬ P) ;; R1(true)))
proof -
  have ((R1(¬ $ok) :: ('t, 'α, 'β) rel-rp) ;; R1(true)) =
    (R1(¬ $ok) :: ('t, 'α, 'γ) rel-rp)
  by (rel-auto)
  thus ?thesis
  by simp
qed
also have ... = ((R1(Q) ;; (R1(¬ R) ∨ (R1(S ∧ $ok')))))
  ∨ R1(¬ $ok)
  ∨ (R1(¬ P) ;; R1(true)))
  by (simp add: R1-extend-conj)
also have ... = ( (R1(Q) ;; (R1(¬ R)))
  ∨ (R1(Q) ;; (R1(S ∧ $ok'))))
  ∨ R1(¬ $ok)
  ∨ (R1(¬ P) ;; R1(true)))
  by (simp add: seqr-or-distr utp-pred-laws.sup.assoc)
also have ... = R1( (R1(Q) ;; (R1(¬ R)))
  ∨ (R1(Q) ;; (R1(S ∧ $ok'))))
  ∨ (¬ $ok)
  ∨ (R1(¬ P) ;; R1(true)))
  by (simp add: R1-disj R1-seqr)
also have ... = R1( (R1(Q) ;; (R1(¬ R)))
  ∨ ((R1(Q) ;; R1(S)) ∧ $ok'))
  ∨ (¬ $ok)
  ∨ (R1(¬ P) ;; R1(true)))
  by (rel-blast)
also have ... = R1(¬($ok ∧ ¬ (R1(¬ P) ;; R1(true)) ∧ ¬ (R1(Q) ;; (R1(¬ R)))))
  ∨ ((R1(Q) ;; R1(S)) ∧ $ok'))
  by (rel-blast)
also have ... = R1(($ok ∧ ¬ (R1(¬ P) ;; R1(true)) ∧ ¬ (R1(Q) ;; (R1(¬ R)))))
  ⇒ ($ok' ∧ (R1(Q) ;; R1(S))))
  by (simp add: impl-alt-def utp-pred-laws.inf-commute)
also have ... = R1((¬ (R1(¬ P) ;; R1(true)) ∧ ¬ (R1(Q) ;; R1(¬ R))) ⊢ (R1(Q) ;; R1(S)))
  by (simp add: design-def)
finally show ?thesis .
qed

```

theorem *R1-design-composition-RR:*
assumes *P is RR Q is RR R is RR S is RR*
shows
 $(R1(P ⊢ Q) ;; R1(R ⊢ S)) = R1(((¬_r P) wp_r false ∧ Q wp_r R) ⊢ (Q ;; S))$
apply (*subst R1-design-composition*)
apply (*simp-all add: assms unrest wp-rea-def Healthy-if closure*)
apply (*rel-auto*)
done

theorem *R1-design-composition-RC:*
assumes *P is RC Q is RR R is RR S is RR*
shows
 $(R1(P ⊢ Q) ;; R1(R ⊢ S)) = R1((P ∧ Q wp_r R) ⊢ (Q ;; S))$
by (*simp add: R1-design-composition-RR assms unrest Healthy-if closure wp*)

lemma *R2s-design*: $R2s(P \vdash Q) = (R2s(P) \vdash R2s(Q))$

by (*simp add: R2s-def design-def usubst*)

lemma *R2c-design*: $R2c(P \vdash Q) = (R2c(P) \vdash R2c(Q))$

by (*simp add: design-def impl-alt-def R2c-disj R2c-not R2c-ok R2c-and R2c-ok'*)

lemma *R1-R3c-design*:

$R1(R3c(P \vdash Q)) = R1(R3c\text{-pre}(P) \vdash R3c\text{-post}(Q))$

by (*rel-auto*)

lemma *R1-R3h-design*:

$R1(R3h(P \vdash Q)) = R1(R3c\text{-pre}(P) \vdash R3h\text{-post}(Q))$

by (*rel-auto*)

lemma *R3c-R1-design-composition*:

assumes $\$ok' \# P \$ok' \# Q \$ok \# R \$ok \# S$

shows $(R3c(R1(P \vdash Q)) ;; R3c(R1(R \vdash S))) =$

$R3c(R1((\neg (R1(\neg P) ;; R1(true)) \wedge \neg ((R1(Q) \wedge \neg \$wait') ;; R1(\neg R)))$
 $\vdash (R1(Q) ;; (\llbracket II \rrbracket_D \triangleleft \$wait \triangleright R1(S))))$

proof –

have 1: $(\neg (R1(\neg R3c\text{-pre } P) ;; R1 true)) = (R3c\text{-pre } (\neg (R1(\neg P) ;; R1 true)))$

by (*rel-auto*)

have 2: $(\neg (R1(R3c\text{-post } Q) ;; R1(\neg R3c\text{-pre } R))) = R3c\text{-pre}(\neg ((R1 Q \wedge \neg \$wait') ;; R1(\neg R)))$

by (*rel-auto, blast+*)

have 3: $(R1(R3c\text{-post } Q) ;; R1(R3c\text{-post } S)) = R3c\text{-post } (R1 Q ;; (\llbracket II \rrbracket_D \triangleleft \$wait \triangleright R1 S))$

by (*rel-auto*)

show ?thesis

apply (*simp add: R3c-semir-form R1-R3c-commute[THEN sym] R1-R3c-design unrest*)

apply (*subst R1-design-composition*)

apply (*simp-all add: unrest assms R3c-pre-conj 1 2 3*)

done

qed

lemma *R3h-R1-design-composition*:

assumes $\$ok' \# P \$ok' \# Q \$ok \# R \$ok \# S$

shows $(R3h(R1(P \vdash Q)) ;; R3h(R1(R \vdash S))) =$

$R3h(R1((\neg (R1(\neg P) ;; R1(true)) \wedge \neg ((R1(Q) \wedge \neg \$wait') ;; R1(\neg R)))$
 $\vdash (R1(Q) ;; ((\exists \$st \cdot \llbracket II \rrbracket_D) \triangleleft \$wait \triangleright R1(S))))$

proof –

have 1: $(\neg (R1(\neg R3c\text{-pre } P) ;; R1 true)) = (R3c\text{-pre } (\neg (R1(\neg P) ;; R1 true)))$

by (*rel-auto*)

have 2: $(\neg (R1(R3h\text{-post } Q) ;; R1(\neg R3c\text{-pre } R))) = R3c\text{-pre}(\neg ((R1 Q \wedge \neg \$wait') ;; R1(\neg R)))$

by (*rel-auto, blast+*)

have 3: $(R1(R3h\text{-post } Q) ;; R1(R3h\text{-post } S)) = R3h\text{-post } (R1 Q ;; ((\exists \$st \cdot \llbracket II \rrbracket_D) \triangleleft \$wait \triangleright R1 S))$

by (*rel-auto, blast+*)

show ?thesis

apply (*simp add: R3h-semir-form R1-R3h-commute[THEN sym] R1-R3h-design unrest*)

apply (*subst R1-design-composition*)

apply (*simp-all add: unrest assms R3c-pre-conj 1 2 3*)

done

qed

lemma *R2-design-composition*:

assumes $\$ok' \# P \$ok' \# Q \$ok \# R \$ok \# S$

shows $(R2(P \vdash Q) ;; R2(R \vdash S)) =$

$R2((\neg (R1 (\neg R2c P) ;; R1 true) \wedge \neg (R1 (R2c Q) ;; R1 (\neg R2c R))) \vdash (R1 (R2c Q) ;; R1 (R2c S)))$
apply (*simp add: R2-R2c-def R2c-design R1-design-composition assms unrest R2c-not R2c-and R2c-disj R1-R2c-commute*[*THEN sym*] *R2c-idem R2c-R1-seq*)
apply (*metis* (*no-types*, *lifting*) *R2c-R1-seq R2c-not R2c-true*)
done

lemma *RH-design-composition:*

assumes $\$ok' \# P \$ok' \# Q \$ok \# R \$ok \# S$

shows $(RH(P \vdash Q) ;; RH(R \vdash S)) =$

$RH((\neg (R1 (\neg R2s P) ;; R1 true) \wedge \neg ((R1 (R2s Q) \wedge (\neg \$wait')) ;; R1 (\neg R2s R))) \vdash$
 $(R1 (R2s Q) ;; (\lceil II \rceil_D \triangleleft \$wait \triangleright R1 (R2s S))))$

proof –

have 1: $R2c (R1 (\neg R2s P) ;; R1 true) = (R1 (\neg R2s P) ;; R1 true)$

proof –

have 1: $(R1 (\neg R2s P) ;; R1 true) = (R1(R2 (\neg P) ;; R2 true))$

by (*rel-auto*)

have $R2c(R1(R2 (\neg P) ;; R2 true)) = R2c(R1(R2 (\neg P) ;; R2 true))$

using *R2c-not* **by** *blast*

also have $\dots = R2(R2 (\neg P) ;; R2 true)$

by (*metis R1-R2c-commute R1-R2c-is-R2*)

also have $\dots = (R2 (\neg P) ;; R2 true)$

by (*simp add: R2-seqr-distribute*)

also have $\dots = (R1 (\neg R2s P) ;; R1 true)$

by (*simp add: R2-def R2s-not R2s-true*)

finally show *?thesis*

by (*simp add: 1*)

qed

have 2: $R2c ((R1 (R2s Q) \wedge \neg \$wait') ;; R1 (\neg R2s R)) = ((R1 (R2s Q) \wedge \neg \$wait') ;; R1 (\neg R2s R))$

proof –

have $((R1 (R2s Q) \wedge \neg \$wait') ;; R1 (\neg R2s R)) = R1 (R2 (Q \wedge \neg \$wait') ;; R2 (\neg R))$

by (*rel-auto*)

hence $R2c ((R1 (R2s Q) \wedge \neg \$wait') ;; R1 (\neg R2s R)) = (R2 (Q \wedge \neg \$wait') ;; R2 (\neg R))$

by (*metis R1-R2c-commute R1-R2c-is-R2 R2-seqr-distribute*)

also have $\dots = ((R1 (R2s Q) \wedge \neg \$wait') ;; R1 (\neg R2s R))$

by (*rel-auto*)

finally show *?thesis* .

qed

have 3: $R2c((R1 (R2s Q) ;; (\lceil II \rceil_D \triangleleft \$wait \triangleright R1 (R2s S)))) = (R1 (R2s Q) ;; (\lceil II \rceil_D \triangleleft \$wait \triangleright R1 (R2s S)))$

proof –

have $R2c(((R1 (R2s Q))\llbracket true/\$wait' \rrbracket ;; (\lceil II \rceil_D \triangleleft \$wait \triangleright R1 (R2s S))\llbracket true/\$wait \rrbracket))$

$= ((R1 (R2s Q))\llbracket true/\$wait' \rrbracket ;; (\lceil II \rceil_D \triangleleft \$wait \triangleright R1 (R2s S))\llbracket true/\$wait \rrbracket)$

proof –

have $R2c(((R1 (R2s Q))\llbracket true/\$wait' \rrbracket ;; (\lceil II \rceil_D \triangleleft \$wait \triangleright R1 (R2s S))\llbracket true/\$wait \rrbracket)) =$

$R2c(R1 (R2s (Q\llbracket true/\$wait' \rrbracket)) ;; \lceil II \rceil_D \llbracket true/\$wait \rrbracket)$

by (*simp add: usubst cond-unit-T R1-def R2s-def*)

also have $\dots = R2c(R2(Q\llbracket true/\$wait' \rrbracket) ;; R2(\lceil II \rceil_D \llbracket true/\$wait \rrbracket))$

by (*metis R2-def R2-des-lift-skip R2-subst-wait-true*)

also have $\dots = (R2(Q\llbracket true/\$wait' \rrbracket) ;; R2(\lceil II \rceil_D \llbracket true/\$wait \rrbracket))$

using *R2c-seq* **by** *blast*

also have $\dots = ((R1 (R2s Q))\llbracket true/\$wait' \rrbracket ;; (\lceil II \rceil_D \triangleleft \$wait \triangleright R1 (R2s S))\llbracket true/\$wait \rrbracket)$

```

    apply (simp add: usubst R2-des-lift-skip)
    apply (metis R2-def R2-des-lift-skip R2-subst-wait'-true R2-subst-wait-true)
  done
  finally show ?thesis .
qed
moreover have R2c(((R1 (R2s Q))[false/$wait'] ;; ([II]D < $wait > R1 (R2s S))[false/$wait]))
  = ((R1 (R2s Q))[false/$wait'] ;; ([II]D < $wait > R1 (R2s S))[false/$wait])
  by (simp add: usubst cond-unit-F)
  (metis (no-types, hide-lams) R1-wait'-false R1-wait-false R2-def R2-subst-wait'-false R2-subst-wait-false
R2c-seq)
ultimately show ?thesis
proof -
  have [II]D < $wait > R1 (R2s S) = R2 ([II]D < $wait > S)
  by (simp add: R1-R2c-is-R2 R1-R2s-R2c R2-condr' R2-des-lift-skip R2s-wait)
  then show ?thesis
  by (simp add: R1-R2c-is-R2 R1-R2s-R2c R2c-seq)
qed
qed

have (R1(R2s(R3c(P ⊢ Q))) ;; R1(R2s(R3c(R ⊢ S)))) =
  ((R3c(R1(R2s(P) ⊢ R2s(Q)))) ;; R3c(R1(R2s(R) ⊢ R2s(S))))
  by (metis (no-types, hide-lams) R1-R2s-R2c R1-R3c-commute R2c-R3c-commute R2s-design)
also have ... = R3c (R1 ((¬ (R1 (¬ R2s P) ;; R1 true) ∧ ¬ ((R1 (R2s Q) ∧ ¬ $wait') ;; R1 (¬ R2s
R)))) ⊢
  (R1 (R2s Q) ;; ([II]D < $wait > R1 (R2s S))))
  by (simp add: R3c-R1-design-composition assms unrest)
also have ... = R3c(R1(R2c((¬ (R1 (¬ R2s P) ;; R1 true) ∧ ¬ ((R1 (R2s Q) ∧ ¬ $wait') ;; R1 (¬
R2s R)))) ⊢
  (R1 (R2s Q) ;; ([II]D < $wait > R1 (R2s S))))))
  by (simp add: R2c-design R2c-and R2c-not 1 2 3)
finally show ?thesis
  by (simp add: R1-R2s-R2c R1-R3c-commute R2c-R3c-commute RH-def)
qed

lemma RHS-design-composition:
  assumes $ok' # P $ok' # Q $ok # R $ok # S
  shows (Rs(P ⊢ Q) ;; Rs(R ⊢ S)) =
    Rs((¬ (R1 (¬ R2s P) ;; R1 true) ∧ ¬ ((R1 (R2s Q) ∧ ¬ $wait') ;; R1 (¬ R2s R))) ⊢
      (R1 (R2s Q) ;; ((∃ $st • [II]D < $wait > R1 (R2s S))))
proof -
  have 1: R2c (R1 (¬ R2s P) ;; R1 true) = (R1 (¬ R2s P) ;; R1 true)
  proof -
    have 1:(R1 (¬ R2s P) ;; R1 true) = (R1(R2 (¬ P) ;; R2 true))
    by (rel-auto, blast)
    have R2c(R1(R2 (¬ P) ;; R2 true)) = R2c(R1(R2 (¬ P) ;; R2 true))
    using R2c-not by blast
    also have ... = R2(R2 (¬ P) ;; R2 true)
    by (metis R1-R2c-commute R1-R2c-is-R2)
    also have ... = (R2 (¬ P) ;; R2 true)
    by (simp add: R2-seqr-distribute)
    also have ... = (R1 (¬ R2s P) ;; R1 true)
    by (simp add: R2-def R2s-not R2s-true)
    finally show ?thesis
    by (simp add: 1)
  qed

```

have 2:R2c ((R1 (R2s Q) \wedge \neg \$wait') ;; R1 (\neg R2s R)) = ((R1 (R2s Q) \wedge \neg \$wait') ;; R1 (\neg R2s R))

proof –

have ((R1 (R2s Q) \wedge \neg \$wait') ;; R1 (\neg R2s R)) = R1 (R2 (Q \wedge \neg \$wait') ;; R2 (\neg R))
by (rel-auto, blast+)

hence R2c ((R1 (R2s Q) \wedge \neg \$wait') ;; R1 (\neg R2s R)) = (R2 (Q \wedge \neg \$wait') ;; R2 (\neg R))
by (metis (no-types, lifting) R1-R2c-commute R1-R2c-is-R2 R2-seqr-distribute)

also have ... = ((R1 (R2s Q) \wedge \neg \$wait') ;; R1 (\neg R2s R))
by (rel-auto, blast+)

finally show ?thesis .

qed

have 3:R2c((R1 (R2s Q) ;; ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 (R2s S)))) =
(R1 (R2s Q) ;; ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 (R2s S)))

proof –

have R2c(((R1 (R2s Q)) \llbracket true/\$wait' \rrbracket ;; ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 (R2s S)) \llbracket true/\$wait \rrbracket))
= ((R1 (R2s Q)) \llbracket true/\$wait' \rrbracket ;; ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 (R2s S)) \llbracket true/\$wait \rrbracket)

proof –

have R2c(((R1 (R2s Q)) \llbracket true/\$wait' \rrbracket ;; ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 (R2s S)) \llbracket true/\$wait \rrbracket)) =
R2c(R1 (R2s (Q \llbracket true/\$wait' \rrbracket)) ;; (\exists \$st \cdot \lceil II \rceil_D) \llbracket true/\$wait \rrbracket)

by (simp add: usubst cond-unit-T R1-def R2s-def)

also have ... = R2c(R2(Q \llbracket true/\$wait' \rrbracket) ;; R2((\exists \$st \cdot \lceil II \rceil_D) \llbracket true/\$wait \rrbracket))

by (metis (no-types, lifting) R2-def R2-des-lift-skip R2-subst-wait-true R2-st-ex)

also have ... = (R2(Q \llbracket true/\$wait' \rrbracket) ;; R2((\exists \$st \cdot \lceil II \rceil_D) \llbracket true/\$wait \rrbracket))

using R2c-seq by blast

also have ... = ((R1 (R2s Q)) \llbracket true/\$wait' \rrbracket ;; ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 (R2s S)) \llbracket true/\$wait \rrbracket)

apply (simp add: usubst R2-des-lift-skip)

apply (metis (no-types) R2-def R2-des-lift-skip R2-st-ex R2-subst-wait'-true R2-subst-wait-true)

done

finally show ?thesis .

qed

moreover have R2c(((R1 (R2s Q)) \llbracket false/\$wait' \rrbracket ;; ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 (R2s S)) \llbracket false/\$wait \rrbracket))
= ((R1 (R2s Q)) \llbracket false/\$wait' \rrbracket ;; ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 (R2s S)) \llbracket false/\$wait \rrbracket)

by (simp add: usubst)

(metis (no-types, lifting) R1-wait'-false R1-wait-false R2-R1-form R2-subst-wait'-false R2-subst-wait-false

R2c-seq)

ultimately show ?thesis

by (smt R2-R1-form R2-condr' R2-des-lift-skip R2-st-ex R2c-seq R2s-wait)

qed

have (R1(R2s(R3h(P \vdash Q))) ;; R1(R2s(R3h(R \vdash S)))) =

((R3h(R1(R2s(P) \vdash R2s(Q)))) ;; R3h(R1(R2s(R) \vdash R2s(S))))

by (metis (no-types, hide-lams) R1-R2s-R2c R1-R3h-commute R2c-R3h-commute R2s-design)

also have ... = R3h (R1 ((\neg (R1 (\neg R2s P) ;; R1 true) \wedge \neg ((R1 (R2s Q) \wedge \neg \$wait') ;; R1 (\neg R2s R))) \vdash

(R1 (R2s Q) ;; ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 (R2s S))))

by (simp add: R3h-R1-design-composition asms unrest)

also have ... = R3h(R1(R2c((\neg (R1 (\neg R2s P) ;; R1 true) \wedge \neg ((R1 (R2s Q) \wedge \neg \$wait') ;; R1 (\neg R2s R))) \vdash

(R1 (R2s Q) ;; ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 (R2s S))))

by (simp add: R2c-design R2c-and R2c-not 1 2 3)

finally show ?thesis

by (simp add: R1-R2s-R2c R1-R3h-commute R2c-R3h-commute RHS-def)

qed

lemma *RHS-R2s-design-composition:*

assumes

$\$ok' \# P \ \$ok' \# Q \ \$ok \# R \ \$ok \# S$
 $P \text{ is } R2s \ Q \text{ is } R2s \ R \text{ is } R2s \ S \text{ is } R2s$

shows $(\mathbf{R}_s(P \vdash Q) ;; \mathbf{R}_s(R \vdash S)) =$

$\mathbf{R}_s((\neg (R1 \ (\neg P) ;; R1 \ true) \wedge \neg ((R1 \ Q \wedge \neg \$wait') ;; R1 \ (\neg R))) \vdash$
 $(R1 \ Q ;; ((\exists \$st \cdot [II]_D) \triangleleft \$wait \triangleright R1 \ S)))$

proof –

have $f1: R2s \ P = P$

by (*meson Healthy-def assms(5)*)

have $f2: R2s \ Q = Q$

by (*meson Healthy-def assms(6)*)

have $f3: R2s \ R = R$

by (*meson Healthy-def assms(7)*)

have $R2s \ S = S$

by (*meson Healthy-def assms(8)*)

then show *?thesis*

using $f3 \ f2 \ f1$ **by** (*simp add: RHS-design-composition assms(1) assms(2) assms(3) assms(4)*)

qed

lemma *RH-design-export-R1:* $\mathbf{R}(P \vdash Q) = \mathbf{R}(P \vdash R1(Q))$

by (*rel-auto*)

lemma *RH-design-export-R2s:* $\mathbf{R}(P \vdash Q) = \mathbf{R}(P \vdash R2s(Q))$

by (*rel-auto*)

lemma *RH-design-export-R2c:* $\mathbf{R}(P \vdash Q) = \mathbf{R}(P \vdash R2c(Q))$

by (*rel-auto*)

lemma *RHS-design-export-R1:* $\mathbf{R}_s(P \vdash Q) = \mathbf{R}_s(P \vdash R1(Q))$

by (*rel-auto*)

lemma *RHS-design-export-R2s:* $\mathbf{R}_s(P \vdash Q) = \mathbf{R}_s(P \vdash R2s(Q))$

by (*rel-auto*)

lemma *RHS-design-export-R2c:* $\mathbf{R}_s(P \vdash Q) = \mathbf{R}_s(P \vdash R2c(Q))$

by (*rel-auto*)

lemma *RHS-design-export-R2:* $\mathbf{R}_s(P \vdash Q) = \mathbf{R}_s(P \vdash R2(Q))$

by (*rel-auto*)

lemma *R1-design-R1-pre:*

$\mathbf{R}_s(R1(P) \vdash Q) = \mathbf{R}_s(P \vdash Q)$

by (*rel-auto*)

lemma *RHS-design-ok-wait:* $\mathbf{R}_s(P \llbracket true, false / \$ok, \$wait \rrbracket \vdash Q \llbracket true, false / \$ok, \$wait \rrbracket) = \mathbf{R}_s(P \vdash Q)$

by (*rel-auto*)

lemma *RHS-design-neg-R1-pre:*

$\mathbf{R}_s((\neg R1 \ P) \vdash R) = \mathbf{R}_s((\neg P) \vdash R)$

by (*rel-auto*)

lemma *RHS-design-conj-neg-R1-pre:*

$\mathbf{R}_s(((\neg R1 \ P) \wedge Q) \vdash R) = \mathbf{R}_s(((\neg P) \wedge Q) \vdash R)$

by (rel-auto)

lemma *RHS-pre-lemma*: $(\mathbf{R}_s P)^f_f = R1(R2c(P^f_f))$
 by (rel-auto)

lemma *RHS-design-R2c-pre*:
 $\mathbf{R}_s(R2c(P) \vdash Q) = \mathbf{R}_s(P \vdash Q)$
 by (rel-auto)

3.4 Refinement introduction laws

lemma *R1-design-refine*:

assumes

P_1 is R1 P_2 is R1 Q_1 is R1 Q_2 is R1
 $\$ok \# P_1 \$ok' \# P_1 \$ok \# P_2 \$ok' \# P_2$
 $\$ok \# Q_1 \$ok' \# Q_1 \$ok \# Q_2 \$ok' \# Q_2$

shows $R1(P_1 \vdash P_2) \sqsubseteq R1(Q_1 \vdash Q_2) \longleftrightarrow 'P_1 \Rightarrow Q_1' \wedge 'P_1 \wedge Q_2 \Rightarrow P_2'$

proof –

have $R1((\exists \$ok; \$ok' \cdot P_1) \vdash (\exists \$ok; \$ok' \cdot P_2)) \sqsubseteq R1((\exists \$ok; \$ok' \cdot Q_1) \vdash (\exists \$ok; \$ok' \cdot Q_2))$
 $\longleftrightarrow 'R1(\exists \$ok; \$ok' \cdot P_1) \Rightarrow R1(\exists \$ok; \$ok' \cdot Q_1)' \wedge 'R1(\exists \$ok; \$ok' \cdot P_1) \wedge R1(\exists \$ok; \$ok'$

$\cdot Q_2) \Rightarrow R1(\exists \$ok; \$ok' \cdot P_2)'$

by (rel-auto, meson+)

thus ?thesis

by (simp-all add: ex-unrest ex-plus Healthy-if assms)

qed

lemma *R1-design-refine-RR*:

assumes P_1 is RR P_2 is RR Q_1 is RR Q_2 is RR

shows $R1(P_1 \vdash P_2) \sqsubseteq R1(Q_1 \vdash Q_2) \longleftrightarrow 'P_1 \Rightarrow Q_1' \wedge 'P_1 \wedge Q_2 \Rightarrow P_2'$

by (simp add: R1-design-refine assms unrest closure)

lemma *RHS-design-refine*:

assumes

P_1 is R1 P_2 is R1 Q_1 is R1 Q_2 is R1
 P_1 is R2c P_2 is R2c Q_1 is R2c Q_2 is R2c
 $\$ok \# P_1 \$ok' \# P_1 \$ok \# P_2 \$ok' \# P_2$
 $\$ok \# Q_1 \$ok' \# Q_1 \$ok \# Q_2 \$ok' \# Q_2$
 $\$wait \# P_1 \$wait \# P_2 \$wait \# Q_1 \$wait \# Q_2$

shows $\mathbf{R}_s(P_1 \vdash P_2) \sqsubseteq \mathbf{R}_s(Q_1 \vdash Q_2) \longleftrightarrow 'P_1 \Rightarrow Q_1' \wedge 'P_1 \wedge Q_2 \Rightarrow P_2'$

proof –

have $\mathbf{R}_s(P_1 \vdash P_2) \sqsubseteq \mathbf{R}_s(Q_1 \vdash Q_2) \longleftrightarrow R1(R3h(R2c(P_1 \vdash P_2))) \sqsubseteq R1(R3h(R2c(Q_1 \vdash Q_2)))$
 by (simp add: R2c-R3h-commute RHS-def)

also have $\dots \longleftrightarrow R1(R3h(P_1 \vdash P_2)) \sqsubseteq R1(R3h(Q_1 \vdash Q_2))$

by (simp add: Healthy-if R2c-design assms)

also have $\dots \longleftrightarrow R1(R3h(P_1 \vdash P_2)) \llbracket false/\$wait \rrbracket \sqsubseteq R1(R3h(Q_1 \vdash Q_2)) \llbracket false/\$wait \rrbracket$

by (rel-auto, metis+)

also have $\dots \longleftrightarrow R1(P_1 \vdash P_2) \llbracket false/\$wait \rrbracket \sqsubseteq R1(Q_1 \vdash Q_2) \llbracket false/\$wait \rrbracket$

by (rel-auto)

also have $\dots \longleftrightarrow R1(P_1 \vdash P_2) \sqsubseteq R1(Q_1 \vdash Q_2)$

by (simp add: usubst assms closure unrest)

also have $\dots \longleftrightarrow 'P_1 \Rightarrow Q_1' \wedge 'P_1 \wedge Q_2 \Rightarrow P_2'$

by (simp add: R1-design-refine assms)

finally show ?thesis .

qed

lemma *srdes-refine-intro*:

assumes $\text{'}P_1 \Rightarrow P_2\text{'}$ $\text{'}P_1 \wedge Q_2 \Rightarrow Q_1\text{'}$
shows $\mathbf{R}_s(P_1 \vdash Q_1) \sqsubseteq \mathbf{R}_s(P_2 \vdash Q_2)$
by (*simp add: RHS-mono assms design-refine-intro*)

3.5 Distribution laws

lemma *RHS-design-choice*: $\mathbf{R}_s(P_1 \vdash Q_1) \sqcap \mathbf{R}_s(P_2 \vdash Q_2) = \mathbf{R}_s((P_1 \wedge P_2) \vdash (Q_1 \vee Q_2))$
by (*metis RHS-inf design-choice*)

lemma *RHS-design-sup*: $\mathbf{R}_s(P_1 \vdash Q_1) \sqcup \mathbf{R}_s(P_2 \vdash Q_2) = \mathbf{R}_s((P_1 \vee P_2) \vdash ((P_1 \Rightarrow Q_1) \wedge (P_2 \Rightarrow Q_2)))$
by (*metis RHS-sup design-inf*)

lemma *RHS-design-USUP*:
assumes $A \neq \{\}$
shows $(\prod i \in A \cdot \mathbf{R}_s(P(i) \vdash Q(i))) = \mathbf{R}_s((\prod i \in A \cdot P(i)) \vdash (\prod i \in A \cdot Q(i)))$
by (*subst RHS-INF[OF assms, THEN sym], simp add: design-UINF-mem assms*)

end

4 Reactive Design Triples

theory *utp-rdes-triples*
imports *utp-rdes-designs*
begin

4.1 Diamond notation

term ($::$)

definition *wait'-cond* ::
 $(t::\text{trace}, 'a, 'b) \text{ rel-rp} \Rightarrow (t, 'a, 'b) \text{ rel-rp} \Rightarrow (t, 'a, 'b) \text{ rel-rp}$ (**infixr** \diamond 60) **where**
 $[\text{upred-defs}]: P \diamond Q = (P \triangleleft \$\text{wait}' \triangleright Q)$

lemma *wait'-cond-unrest* [*unrest*]:
 $\llbracket \text{out-var wait} \bowtie x; x \# P; x \# Q \rrbracket \Longrightarrow x \# (P \diamond Q)$
by (*simp add: wait'-cond-def unrest*)

lemma *wait'-cond-subst* [*usubst*]:
 $\$ \text{wait}' \# \sigma \Longrightarrow \sigma \dagger (P \diamond Q) = (\sigma \dagger P) \diamond (\sigma \dagger Q)$
by (*simp add: wait'-cond-def usubst unrest*)

lemma *wait'-cond-left-false*: $\text{false} \diamond P = (\neg \$ \text{wait}' \wedge P)$
by (*rel-auto*)

lemma *wait'-cond-seq*: $((P \diamond Q) ;; R) = ((P ;; (\$ \text{wait}' \wedge R)) \vee (Q ;; (\neg \$ \text{wait}' \wedge R)))$
by (*simp add: wait'-cond-def cond-def seqr-or-distl, rel-blast*)

lemma *wait'-cond-true*: $(P \diamond Q \wedge \$ \text{wait}') = (P \wedge \$ \text{wait}')$
by (*rel-auto*)

lemma *wait'-cond-false*: $(P \diamond Q \wedge (\neg \$ \text{wait}')) = (Q \wedge (\neg \$ \text{wait}'))$
by (*rel-auto*)

lemma *wait'-cond-idem*: $P \diamond P = P$
by (*rel-auto*)

lemma *wait'-cond-conj-exchange*:

$$((P \diamond Q) \wedge (R \diamond S)) = (P \wedge R) \diamond (Q \wedge S)$$

by (*rel-auto*)

lemma *subst-wait'-cond-true* [*usubst*]: $(P \diamond Q) \llbracket \text{true}/\$wait' \rrbracket = P \llbracket \text{true}/\$wait' \rrbracket$

by (*rel-auto*)

lemma *subst-wait'-cond-false* [*usubst*]: $(P \diamond Q) \llbracket \text{false}/\$wait' \rrbracket = Q \llbracket \text{false}/\$wait' \rrbracket$

by (*rel-auto*)

lemma *subst-wait'-left-subst*: $(P \llbracket \text{true}/\$wait' \rrbracket \diamond Q) = (P \diamond Q)$

by (*rel-auto*)

lemma *subst-wait'-right-subst*: $(P \diamond Q \llbracket \text{false}/\$wait' \rrbracket) = (P \diamond Q)$

by (*rel-auto*)

lemma *wait'-cond-split*: $P \llbracket \text{true}/\$wait' \rrbracket \diamond P \llbracket \text{false}/\$wait' \rrbracket = P$

by (*simp add: wait'-cond-def cond-var-split*)

lemma *wait-cond'-assoc* [*simp*]: $P \diamond Q \diamond R = P \diamond R$

by (*rel-auto*)

lemma *wait-cond'-shadow*: $(P \diamond Q) \diamond R = P \diamond Q \diamond R$

by (*rel-auto*)

lemma *wait-cond'-conj* [*simp*]: $P \diamond (Q \wedge (R \diamond S)) = P \diamond (Q \wedge S)$

by (*rel-auto*)

lemma *R1-wait'-cond*: $R1(P \diamond Q) = R1(P) \diamond R1(Q)$

by (*rel-auto*)

lemma *R2s-wait'-cond*: $R2s(P \diamond Q) = R2s(P) \diamond R2s(Q)$

by (*simp add: wait'-cond-def R2s-def R2s-def usubst*)

lemma *R2-wait'-cond*: $R2(P \diamond Q) = R2(P) \diamond R2(Q)$

by (*simp add: R2-def R2s-wait'-cond R1-wait'-cond*)

lemma *wait'-cond-R1-closed* [*closure*]:

$$\llbracket P \text{ is } R1; Q \text{ is } R1 \rrbracket \implies P \diamond Q \text{ is } R1$$

by (*simp add: Healthy-def R1-wait'-cond*)

lemma *wait'-cond-R2c-closed* [*closure*]: $\llbracket P \text{ is } R2c; Q \text{ is } R2c \rrbracket \implies P \diamond Q \text{ is } R2c$

by (*simp add: R2c-condr wait'-cond-def Healthy-def, rel-auto*)

4.2 Export laws

lemma *RH-design-peri-R1*: $\mathbf{R}(P \vdash R1(Q) \diamond R) = \mathbf{R}(P \vdash Q \diamond R)$

by (*metis (no-types, lifting) R1-idem R1-wait'-cond RH-design-export-R1*)

lemma *RH-design-post-R1*: $\mathbf{R}(P \vdash Q \diamond R1(R)) = \mathbf{R}(P \vdash Q \diamond R)$

by (*metis R1-wait'-cond RH-design-export-R1 RH-design-peri-R1*)

lemma *RH-design-peri-R2s*: $\mathbf{R}(P \vdash R2s(Q) \diamond R) = \mathbf{R}(P \vdash Q \diamond R)$

by (*metis (no-types, lifting) R2s-idem R2s-wait'-cond RH-design-export-R2s*)

lemma *RH-design-post-R2s*: $\mathbf{R}(P \vdash Q \diamond R2s(R)) = \mathbf{R}(P \vdash Q \diamond R)$
 by (metis (no-types, lifting) *R2s-idem R2s-wait'-cond RH-design-export-R2s*)

lemma *RH-design-peri-R2c*: $\mathbf{R}(P \vdash R2c(Q) \diamond R) = \mathbf{R}(P \vdash Q \diamond R)$
 by (metis *R1-R2s-R2c RH-design-peri-R1 RH-design-peri-R2s*)

lemma *RHS-design-peri-R1*: $\mathbf{R}_s(P \vdash R1(Q) \diamond R) = \mathbf{R}_s(P \vdash Q \diamond R)$
 by (metis (no-types, lifting) *R1-idem R1-wait'-cond RHS-design-export-R1*)

lemma *RHS-design-post-R1*: $\mathbf{R}_s(P \vdash Q \diamond R1(R)) = \mathbf{R}_s(P \vdash Q \diamond R)$
 by (metis *R1-wait'-cond RHS-design-export-R1 RHS-design-peri-R1*)

lemma *RHS-design-peri-R2s*: $\mathbf{R}_s(P \vdash R2s(Q) \diamond R) = \mathbf{R}_s(P \vdash Q \diamond R)$
 by (metis (no-types, lifting) *R2s-idem R2s-wait'-cond RHS-design-export-R2s*)

lemma *RHS-design-post-R2s*: $\mathbf{R}_s(P \vdash Q \diamond R2s(R)) = \mathbf{R}_s(P \vdash Q \diamond R)$
 by (metis *R2s-wait'-cond RHS-design-export-R2s RHS-design-peri-R2s*)

lemma *RHS-design-peri-R2c*: $\mathbf{R}_s(P \vdash R2c(Q) \diamond R) = \mathbf{R}_s(P \vdash Q \diamond R)$
 by (metis *R1-R2s-R2c RHS-design-peri-R1 RHS-design-peri-R2s*)

lemma *RH-design-lemma1*:
 $RH(P \vdash (R1(R2c(Q)) \vee R) \diamond S) = RH(P \vdash (Q \vee R) \diamond S)$
 by (metis (no-types, lifting) *R1-R2c-is-R2 R1-R2s-R2c R2-R1-form R2-disj R2c-idem RH-design-peri-R1 RH-design-peri-R2s*)

lemma *RHS-design-lemma1*:
 $RHS(P \vdash (R1(R2c(Q)) \vee R) \diamond S) = RHS(P \vdash (Q \vee R) \diamond S)$
 by (metis (no-types, lifting) *R1-R2c-is-R2 R1-R2s-R2c R2-R1-form R2-disj R2c-idem RHS-design-peri-R1 RHS-design-peri-R2s*)

4.3 Pre-, peri-, and postconditions

4.3.1 Definitions

abbreviation $pre_s \equiv [\$ok \mapsto_s true, \$ok' \mapsto_s false, \$wait \mapsto_s false]$

abbreviation $cmt_s \equiv [\$ok \mapsto_s true, \$ok' \mapsto_s true, \$wait \mapsto_s false]$

abbreviation $peri_s \equiv [\$ok \mapsto_s true, \$ok' \mapsto_s true, \$wait \mapsto_s false, \$wait' \mapsto_s true]$

abbreviation $post_s \equiv [\$ok \mapsto_s true, \$ok' \mapsto_s true, \$wait \mapsto_s false, \$wait' \mapsto_s false]$

abbreviation $npre_R(P) \equiv pre_s \dagger P$

definition [*upred-defs*]: $pre_R(P) = (\neg_r npre_R(P))$

definition [*upred-defs*]: $cmt_R(P) = R1(cmt_s \dagger P)$

definition [*upred-defs*]: $peri_R(P) = R1(peri_s \dagger P)$

definition [*upred-defs*]: $post_R(P) = R1(post_s \dagger P)$

4.3.2 Unrestriction laws

lemma *ok-pre-unrest* [*unrest*]: $\$ok \# pre_R P$
 by (simp add: *pre_R-def unrest usubst*)

lemma *ok-peri-unrest* [*unrest*]: $\$ok \# peri_R P$
 by (simp add: *peri_R-def unrest usubst*)

lemma *ok-post-unrest* [*unrest*]: $\$ok \# post_R P$

by (simp add: post_R-def unrest usubst)

lemma ok-cmt-unrest [unrest]: \$ok # cmt_R P
by (simp add: cmt_R-def unrest usubst)

lemma ok'-pre-unrest [unrest]: \$ok' # pre_R P
by (simp add: pre_R-def unrest usubst)

lemma ok'-peri-unrest [unrest]: \$ok' # peri_R P
by (simp add: peri_R-def unrest usubst)

lemma ok'-post-unrest [unrest]: \$ok' # post_R P
by (simp add: post_R-def unrest usubst)

lemma ok'-cmt-unrest [unrest]: \$ok' # cmt_R P
by (simp add: cmt_R-def unrest usubst)

lemma wait-pre-unrest [unrest]: \$wait # pre_R P
by (simp add: pre_R-def unrest usubst)

lemma wait-peri-unrest [unrest]: \$wait # peri_R P
by (simp add: peri_R-def unrest usubst)

lemma wait-post-unrest [unrest]: \$wait # post_R P
by (simp add: post_R-def unrest usubst)

lemma wait-cmt-unrest [unrest]: \$wait # cmt_R P
by (simp add: cmt_R-def unrest usubst)

lemma wait'-peri-unrest [unrest]: \$wait' # peri_R P
by (simp add: peri_R-def unrest usubst)

lemma wait'-post-unrest [unrest]: \$wait' # post_R P
by (simp add: post_R-def unrest usubst)

4.3.3 Substitution laws

lemma pre_s-design: pre_s † (P ⊢ Q) = (¬ pre_s † P)
by (simp add: design-def pre_R-def usubst)

lemma peri_s-design: peri_s † (P ⊢ Q ◊ R) = peri_s † (P ⇒ Q)
by (simp add: design-def usubst wait'-cond-def)

lemma post_s-design: post_s † (P ⊢ Q ◊ R) = post_s † (P ⇒ R)
by (simp add: design-def usubst wait'-cond-def)

lemma cmt_s-design: cmt_s † (P ⊢ Q) = cmt_s † (P ⇒ Q)
by (simp add: design-def usubst wait'-cond-def)

lemma pre_s-R1 [usubst]: pre_s † R1(P) = R1(pre_s † P)
by (simp add: R1-def usubst)

lemma pre_s-R2c [usubst]: pre_s † R2c(P) = R2c(pre_s † P)
by (simp add: R2c-def R2s-def usubst)

lemma peri_s-R1 [usubst]: peri_s † R1(P) = R1(peri_s † P)

by (simp add: R1-def usubst)

lemma *peri_s-R2c* [usubst]: $peri_s \dagger R2c(P) = R2c(peri_s \dagger P)$
 by (simp add: R2c-def R2s-def usubst)

lemma *post_s-R1* [usubst]: $post_s \dagger R1(P) = R1(post_s \dagger P)$
 by (simp add: R1-def usubst)

lemma *post_s-R2c* [usubst]: $post_s \dagger R2c(P) = R2c(post_s \dagger P)$
 by (simp add: R2c-def R2s-def usubst)

lemma *cmt_s-R1* [usubst]: $cmt_s \dagger R1(P) = R1(cmt_s \dagger P)$
 by (simp add: R1-def usubst)

lemma *cmt_s-R2c* [usubst]: $cmt_s \dagger R2c(P) = R2c(cmt_s \dagger P)$
 by (simp add: R2c-def R2s-def usubst)

lemma *pre-wait-false*:
 $pre_R(P \llbracket false/\$wait \rrbracket) = pre_R(P)$
 by (rel-auto)

lemma *cmt-wait-false*:
 $cmt_R(P \llbracket false/\$wait \rrbracket) = cmt_R(P)$
 by (rel-auto)

lemma *rea-pre-RHS-design*: $pre_R(\mathbf{R}_s(P \vdash Q)) = R1(R2c(pre_s \dagger P))$
 by (simp add: RHS-def usubst R3h-def pre_R-def pre_s-design R1-negate-R1 R2c-not rea-not-def)

lemma *rea-cmt-RHS-design*: $cmt_R(\mathbf{R}_s(P \vdash Q)) = R1(R2c(cmt_s \dagger (P \Rightarrow Q)))$
 by (simp add: RHS-def usubst R3h-def cmt_R-def cmt_s-design R1-idem)

lemma *rea-peri-RHS-design*: $peri_R(\mathbf{R}_s(P \vdash Q \diamond R)) = R1(R2c(peri_s \dagger (P \Rightarrow_r Q)))$
 by (simp add: RHS-def usubst peri_R-def R3h-def peri_s-design, rel-auto)

lemma *rea-post-RHS-design*: $post_R(\mathbf{R}_s(P \vdash Q \diamond R)) = R1(R2c(post_s \dagger (P \Rightarrow_r R)))$
 by (simp add: RHS-def usubst post_R-def R3h-def post_s-design, rel-auto)

lemma *peri-cmt-def*: $peri_R(P) = (cmt_R(P)) \llbracket true/\$wait \rrbracket$
 by (rel-auto)

lemma *post-cmt-def*: $post_R(P) = (cmt_R(P)) \llbracket false/\$wait \rrbracket$
 by (rel-auto)

lemma *rdes-export-cmt*: $\mathbf{R}_s(P \vdash cmt_s \dagger Q) = \mathbf{R}_s(P \vdash Q)$
 by (rel-auto)

lemma *rdes-export-pre*: $\mathbf{R}_s((P \llbracket true, false/\$ok, \$wait \rrbracket) \vdash Q) = \mathbf{R}_s(P \vdash Q)$
 by (rel-auto)

4.3.4 Healthiness laws

lemma *wait'-unrest-pre-SRD* [unrest]:
 $\$wait' \# pre_R(P) \Longrightarrow \$wait' \# pre_R(SRD\ P)$
 apply (rel-auto)
 using least-zero apply blast+
 done

lemma *R1-R2s-cmt-SRD*:

assumes *P is SRD*

shows $R1(R2s(cmt_R(P))) = cmt_R(P)$

by (*metis (no-types, lifting) R1-R2c-commute R1-R2s-R2c R1-idem R2c-idem SRD-reactive-design assms rea-cmt-RHS-design*)

lemma *R1-R2s-peri-SRD*:

assumes *P is SRD*

shows $R1(R2s(peri_R(P))) = peri_R(P)$

by (*metis (no-types, hide-lams) Healthy-def R1-R2s-R2c R2-def R2-idem RHS-def SRD-RH-design-form assms R1-idem peri_R-def peri_s-R1 peri_s-R2c*)

lemma *R1-peri-SRD*:

assumes *P is SRD*

shows $R1(peri_R(P)) = peri_R(P)$

proof –

have $R1(peri_R(P)) = R1(R1(R2s(peri_R(P))))$

by (*simp add: R1-R2s-peri-SRD assms*)

also have $\dots = peri_R(P)$

by (*simp add: R1-idem, simp add: R1-R2s-peri-SRD assms*)

finally show *?thesis* .

qed

lemma *periR-SRD-R1 [closure]*: $P \text{ is SRD} \implies peri_R(P) \text{ is R1}$

by (*simp add: Healthy-def' R1-peri-SRD*)

lemma *R1-R2c-peri-RHS*:

assumes *P is SRD*

shows $R1(R2c(peri_R(P))) = peri_R(P)$

by (*metis R1-R2s-R2c R1-R2s-peri-SRD assms*)

lemma *R1-R2s-post-SRD*:

assumes *P is SRD*

shows $R1(R2s(post_R(P))) = post_R(P)$

by (*metis (no-types, hide-lams) Healthy-def R1-R2s-R2c R1-idem R2-def R2-idem RHS-def SRD-RH-design-form assms post_R-def post_s-R1 post_s-R2c*)

lemma *R2c-peri-SRD*:

assumes *P is SRD*

shows $R2c(peri_R(P)) = peri_R(P)$

by (*metis R1-R2c-commute R1-R2c-peri-RHS R1-peri-SRD assms*)

lemma *R1-post-SRD*:

assumes *P is SRD*

shows $R1(post_R(P)) = post_R(P)$

proof –

have $R1(post_R(P)) = R1(R1(R2s(post_R(P))))$

by (*simp add: R1-R2s-post-SRD assms*)

also have $\dots = post_R(P)$

by (*simp add: R1-idem, simp add: R1-R2s-post-SRD assms*)

finally show *?thesis* .

qed

lemma *R2c-post-SRD*:

assumes P is SRD
shows $R2c(post_R(P)) = post_R(P)$
by (metis $R1-R2c-commute$ $R1-R2s-R2c$ $R1-R2s-post-SRD$ $R1-post-SRD$ $assms$)

lemma $postR-SRD-R1$ [closure]: P is $SRD \implies post_R(P)$ is $R1$
by (simp add: $Healthy-def'$ $R1-post-SRD$)

lemma $R1-R2c-post-RHS$:
assumes P is SRD
shows $R1(R2c(post_R(P))) = post_R(P)$
by (metis $R1-R2s-R2c$ $R1-R2s-post-SRD$ $assms$)

lemma $R2-cmt-conj-wait'$:
 P is $SRD \implies R2(cmt_R P \wedge \neg \$wait') = (cmt_R P \wedge \neg \$wait')$
by (simp add: $R2-def$ $R2s-conj$ $R2s-not$ $R2s-wait'$ $R1-extend-conj$ $R1-R2s-cmt-SRD$)

lemma $R2c-preR$:
 P is $SRD \implies R2c(pre_R(P)) = pre_R(P)$
by (metis (no-types, lifting) $R1-R2c-commute$ $R2c-idem$ $SRD-reactive-design$ $rea-pre-RHS-design$)

lemma $preR-R2c-closed$ [closure]: P is $SRD \implies pre_R(P)$ is $R2c$
by (simp add: $Healthy-def'$ $R2c-preR$)

lemma $R2c-periR$:
 P is $SRD \implies R2c(per_iR(P)) = per_iR(P)$
by (metis (no-types, lifting) $R1-R2c-commute$ $R1-R2s-R2c$ $R1-R2s-peri-SRD$ $R2c-idem$)

lemma $periR-R2c-closed$ [closure]: P is $SRD \implies per_iR(P)$ is $R2c$
by (simp add: $Healthy-def$ $R2c-peri-SRD$)

lemma $R2c-postR$:
 P is $SRD \implies R2c(post_R(P)) = post_R(P)$
by (metis (no-types, hide-lams) $R1-R2c-commute$ $R1-R2c-is-R2$ $R1-R2s-post-SRD$ $R2-def$ $R2s-idem$)

lemma $postR-R2c-closed$ [closure]: P is $SRD \implies post_R(P)$ is $R2c$
by (simp add: $Healthy-def$ $R2c-post-SRD$)

lemma $periR-RR$ [closure]: P is $SRD \implies per_iR(P)$ is RR
by (rule $RR-intro$, simp-all add: $closure$ $unrest$)

lemma $postR-RR$ [closure]: P is $SRD \implies post_R(P)$ is RR
by (rule $RR-intro$, simp-all add: $closure$ $unrest$)

lemma $wpR-trace-ident-pre$ [wp]:
 $(\$tr' =_u \$tr \wedge \lceil II \rceil_R) \wp_r pre_R P = pre_R P$
by (rel-auto)

lemma $R1-preR$ [closure]:
 $pre_R(P)$ is $R1$
by (rel-auto)

lemma $trace-ident-left-periR$:
 $(\$tr' =_u \$tr \wedge \lceil II \rceil_R) ;; per_iR(P) = per_iR(P)$
by (rel-auto)

lemma *trace-ident-left-postR*:

$(\$tr' =_u \$tr \wedge \lceil II \rceil_R) ;; post_R(P) = post_R(P)$
by (*rel-auto*)

lemma *trace-ident-right-postR*:

$post_R(P) ;; (\$tr' =_u \$tr \wedge \lceil II \rceil_R) = post_R(P)$
by (*rel-auto*)

lemma *preR-R2-closed* [*closure*]: P is SRD $\implies pre_R(P)$ is R2

by (*simp add: R2-comp-def Healthy-comp closure*)

lemma *periR-R2-closed* [*closure*]: P is SRD $\implies peri_R(P)$ is R2

by (*simp add: Healthy-def' R1-R2c-peri-RHS R2-R2c-def*)

lemma *postR-R2-closed* [*closure*]: P is SRD $\implies post_R(P)$ is R2

by (*simp add: Healthy-def' R1-R2c-post-RHS R2-R2c-def*)

4.3.5 Calculation laws

lemma *wait'-cond-peri-post-cmt* [*rdes*]:

$cmt_R P = peri_R P \diamond post_R P$

by (*rel-auto*)

lemma *preR-rdes* [*rdes*]:

assumes P is RR

shows $pre_R(\mathbf{R}_s(P \vdash Q \diamond R)) = P$

by (*simp add: rea-pre-RHS-design unrest usubst assms Healthy-if RR-implies-R2c RR-implies-R1*)

lemma *periR-rdes* [*rdes*]:

assumes P is RR Q is RR

shows $peri_R(\mathbf{R}_s(P \vdash Q \diamond R)) = (P \Rightarrow_r Q)$

by (*simp add: rea-peri-RHS-design unrest usubst assms Healthy-if RR-implies-R2c closure*)

lemma *postR-rdes* [*rdes*]:

assumes P is RR R is RR

shows $post_R(\mathbf{R}_s(P \vdash Q \diamond R)) = (P \Rightarrow_r R)$

by (*simp add: rea-post-RHS-design unrest usubst assms Healthy-if RR-implies-R2c closure*)

lemma *preR-Chaos* [*rdes*]: $pre_R(Chaos) = false$

by (*simp add: Chaos-def, rel-simp*)

lemma *periR-Chaos* [*rdes*]: $peri_R(Chaos) = true_r$

by (*simp add: Chaos-def, rel-simp*)

lemma *postR-Chaos* [*rdes*]: $post_R(Chaos) = true_r$

by (*simp add: Chaos-def, rel-simp*)

lemma *preR-Miracle* [*rdes*]: $pre_R(Miracle) = true_r$

by (*simp add: Miracle-def, rel-auto*)

lemma *periR-Miracle* [*rdes*]: $peri_R(Miracle) = false$

by (*simp add: Miracle-def, rel-auto*)

lemma *postR-Miracle* [*rdes*]: $post_R(Miracle) = false$

by (*simp add: Miracle-def, rel-auto*)

lemma *preR-srdes-skip* [rdes]: $pre_R(II_R) = true_r$
by (*rel-auto*)

lemma *periR-srdes-skip* [rdes]: $peri_R(II_R) = false$
by (*rel-auto*)

lemma *postR-srdes-skip* [rdes]: $post_R(II_R) = (\$tr' =_u \$tr \wedge [II]_R)$
by (*rel-auto*)

lemma *preR-INF* [rdes]: $A \neq \{\} \implies pre_R(\sqcap A) = (\bigwedge P \in A \cdot pre_R(P))$
by (*rel-auto*)

lemma *periR-INF* [rdes]: $peri_R(\sqcap A) = (\bigvee P \in A \cdot peri_R(P))$
by (*rel-auto*)

lemma *postR-INF* [rdes]: $post_R(\sqcap A) = (\bigvee P \in A \cdot post_R(P))$
by (*rel-auto*)

lemma *preR-UINF* [rdes]: $pre_R(\sqcap i \cdot P(i)) = (\sqcap i \cdot pre_R(P(i)))$
by (*rel-auto*)

lemma *periR-UINF* [rdes]: $peri_R(\sqcap i \cdot P(i)) = (\sqcap i \cdot peri_R(P(i)))$
by (*rel-auto*)

lemma *postR-UINF* [rdes]: $post_R(\sqcap i \cdot P(i)) = (\sqcap i \cdot post_R(P(i)))$
by (*rel-auto*)

lemma *preR-UINF-member* [rdes]: $A \neq \{\} \implies pre_R(\sqcap i \in A \cdot P(i)) = (\sqcap i \in A \cdot pre_R(P(i)))$
by (*rel-auto*)

lemma *preR-UINF-member-2* [rdes]: $A \neq \{\} \implies pre_R(\sqcap (i,j) \in A \cdot P i j) = (\sqcap (i,j) \in A \cdot pre_R(P i j))$
by (*rel-auto*)

lemma *preR-UINF-member-3* [rdes]: $A \neq \{\} \implies pre_R(\sqcap (i,j,k) \in A \cdot P i j k) = (\sqcap (i,j,k) \in A \cdot pre_R(P i j k))$
by (*rel-auto*)

lemma *periR-UINF-member* [rdes]: $peri_R(\sqcap i \in A \cdot P(i)) = (\sqcap i \in A \cdot peri_R(P(i)))$
by (*rel-auto*)

lemma *periR-UINF-member-2* [rdes]: $peri_R(\sqcap (i,j) \in A \cdot P i j) = (\sqcap (i,j) \in A \cdot peri_R(P i j))$
by (*rel-auto*)

lemma *periR-UINF-member-3* [rdes]: $peri_R(\sqcap (i,j,k) \in A \cdot P i j k) = (\sqcap (i,j,k) \in A \cdot peri_R(P i j k))$
by (*rel-auto*)

lemma *postR-UINF-member* [rdes]: $post_R(\sqcap i \in A \cdot P(i)) = (\sqcap i \in A \cdot post_R(P(i)))$
by (*rel-auto*)

lemma *postR-UINF-member-2* [rdes]: $post_R(\sqcap (i,j) \in A \cdot P i j) = (\sqcap (i,j) \in A \cdot post_R(P i j))$
by (*rel-auto*)

lemma *postR-UINF-member-3* [rdes]: $post_R(\sqcap (i,j,k) \in A \cdot P i j k) = (\sqcap (i,j,k) \in A \cdot post_R(P i j k))$
by (*rel-auto*)

lemma *preR-inf* [rdes]: $pre_R(P \sqcap Q) = (pre_R(P) \wedge pre_R(Q))$
by (*rel-auto*)

lemma *periR-inf* [rdes]: $peri_R(P \sqcap Q) = (peri_R(P) \vee peri_R(Q))$
by (*rel-auto*)

lemma *postR-inf* [rdes]: $post_R(P \sqcap Q) = (post_R(P) \vee post_R(Q))$
by (*rel-auto*)

lemma *preR-SUP* [rdes]: $pre_R(\bigsqcup A) = (\bigvee P \in A \cdot pre_R(P))$
by (*rel-auto*)

lemma *periR-SUP* [rdes]: $A \neq \{\} \implies peri_R(\bigsqcup A) = (\bigwedge P \in A \cdot peri_R(P))$
by (*rel-auto*)

lemma *postR-SUP* [rdes]: $A \neq \{\} \implies post_R(\bigsqcup A) = (\bigwedge P \in A \cdot post_R(P))$
by (*rel-auto*)

4.4 Formation laws

lemma *srdes-skip-tri-design* [rdes-def]: $II_R = \mathbf{R}_s(true_r \vdash false \diamond II_r)$
by (*simp add: srdes-skip-def, rel-auto*)

lemma *Chaos-tri-def* [rdes-def]: $Chaos = \mathbf{R}_s(false \vdash false \diamond false)$
by (*simp add: Chaos-def design-false-pre*)

lemma *Miracle-tri-def* [rdes-def]: $Miracle = \mathbf{R}_s(true_r \vdash false \diamond false)$
by (*simp add: Miracle-def R1-design-R1-pre wait'-cond-idem*)

lemma *RHS-tri-design-form*:

assumes P_1 is RR P_2 is RR P_3 is RR

shows $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) = (II_R \triangleleft \$wait \triangleright ((\$ok \wedge P_1) \Rightarrow_r (\$ok' \wedge (P_2 \diamond P_3))))$

proof –

have $\mathbf{R}_s(RR(P_1) \vdash RR(P_2) \diamond RR(P_3)) = (II_R \triangleleft \$wait \triangleright ((\$ok \wedge RR(P_1)) \Rightarrow_r (\$ok' \wedge (RR(P_2) \diamond RR(P_3)))))$

apply (*rel-auto*) **using** *minus-zero-eq* **by** *blast*

thus *?thesis*

by (*simp add: Healthy-if assms*)

qed

lemma *RHS-design-pre-post-form*:

$\mathbf{R}_s((\neg P^f_f) \vdash P^t_f) = \mathbf{R}_s(pre_R(P) \vdash cmt_R(P))$

proof –

have $\mathbf{R}_s((\neg P^f_f) \vdash P^t_f) = \mathbf{R}_s((\neg P^f_f)[\![true/\$ok]\!] \vdash P^t_f[\![true/\$ok]\!])$

by (*simp add: design-subst-ok*)

also have $\dots = \mathbf{R}_s(pre_R(P) \vdash cmt_R(P))$

by (*simp add: pre_R-def cmt_R-def usubst, rel-auto*)

finally show *?thesis* .

qed

lemma *SRD-as-reactive-design*:

$SRD(P) = \mathbf{R}_s(pre_R(P) \vdash cmt_R(P))$

by (*simp add: RHS-design-pre-post-form SRD-RH-design-form*)

lemma *SRD-reactive-design-alt*:

assumes P is *SRD*

shows $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{cmt}_R(P)) = P$
proof –
 have $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{cmt}_R(P)) = \mathbf{R}_s((\neg P^f_f) \vdash P^t_f)$
 by (*simp add: RHS-design-pre-post-form*)
 thus *?thesis*
 by (*simp add: SRD-reactive-design assms*)
qed

lemma *SRD-reactive-tri-design-lemma*:
 $\text{SRD}(P) = \mathbf{R}_s((\neg P^f_f) \vdash P^t_f \llbracket \text{true}/\$wait' \rrbracket \diamond P^t_f \llbracket \text{false}/\$wait' \rrbracket)$
 by (*simp add: SRD-RH-design-form wait'-cond-split*)

lemma *SRD-as-reactive-tri-design*:
 $\text{SRD}(P) = \mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P))$
proof –
 have $\text{SRD}(P) = \mathbf{R}_s((\neg P^f_f) \vdash P^t_f \llbracket \text{true}/\$wait' \rrbracket \diamond P^t_f \llbracket \text{false}/\$wait' \rrbracket)$
 by (*simp add: SRD-RH-design-form wait'-cond-split*)
 also have $\dots = \mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P))$
 apply (*simp add: usubst*)
 apply (*subst design-subst-ok-ok'[THEN sym]*)
 apply (*simp add: pre_R-def peri_R-def post_R-def usubst unrest*)
 apply (*rel-auto*)
 done
 finally show *?thesis* .
qed

lemma *SRD-reactive-tri-design*:
 assumes *P is SRD*
 shows $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P)) = P$
 by (*metis Healthy-if SRD-as-reactive-tri-design assms*)

lemma *SRD-elim* [*RD-elim*]: $\llbracket P \text{ is SRD}; Q(\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P))) \rrbracket \implies Q(P)$
 by (*simp add: SRD-reactive-tri-design*)

lemma *RHS-tri-design-is-SRD* [*closure*]:
 assumes $\$ok' \# P \ \$ok' \# Q \ \$ok' \# R$
 shows $\mathbf{R}_s(P \vdash Q \diamond R)$ *is SRD*
 by (*rule RHS-design-is-SRD, simp-all add: unrest assms*)

lemma *SRD-rdes-intro* [*closure*]:
 assumes *P is RR Q is RR R is RR*
 shows $\mathbf{R}_s(P \vdash Q \diamond R)$ *is SRD*
 by (*rule RHS-tri-design-is-SRD, simp-all add: unrest closure assms*)

lemma *USUP-R1-R2s-cmt-SRD*:
 assumes $A \subseteq \llbracket \text{SRD} \rrbracket_H$
 shows $(\bigsqcup P \in A \cdot R1 (R2s (\text{cmt}_R P))) = (\bigsqcup P \in A \cdot \text{cmt}_R P)$
 by (*rule USUP-cong[of A], metis (mono-tags, lifting) Ball-Collect R1-R2s-cmt-SRD assms*)

lemma *UINF-R1-R2s-cmt-SRD*:
 assumes $A \subseteq \llbracket \text{SRD} \rrbracket_H$
 shows $(\bigsqcap P \in A \cdot R1 (R2s (\text{cmt}_R P))) = (\bigsqcap P \in A \cdot \text{cmt}_R P)$
 by (*rule UINF-cong[of A], metis (mono-tags, lifting) Ball-Collect R1-R2s-cmt-SRD assms*)

4.4.1 Order laws

lemma *preR-antitone*: $P \sqsubseteq Q \implies \text{pre}_R(Q) \sqsubseteq \text{pre}_R(P)$
by (*rel-auto*)

lemma *periR-monotone*: $P \sqsubseteq Q \implies \text{peri}_R(P) \sqsubseteq \text{peri}_R(Q)$
by (*rel-auto*)

lemma *postR-monotone*: $P \sqsubseteq Q \implies \text{post}_R(P) \sqsubseteq \text{post}_R(Q)$
by (*rel-auto*)

4.5 Composition laws

theorem *RH-tri-design-composition*:

assumes $\$ok' \# P \ \$ok' \# Q_1 \ \$ok' \# Q_2 \ \$ok \# R \ \$ok \# S_1 \ \$ok \# S_2$
 $\$wait' \# Q_2 \ \$wait \# S_1 \ \$wait \# S_2$

shows $(RH(P \vdash Q_1 \diamond Q_2) ;; RH(R \vdash S_1 \diamond S_2)) =$
 $RH((\neg (R1 (\neg R2s P) ;; R1 \text{ true}) \wedge \neg ((R1 (R2s Q_2) \wedge \neg \$wait') ;; R1 (\neg R2s R))) \vdash$
 $((Q_1 \vee (R1 (R2s Q_2) ;; R1 (R2s S_1))) \diamond ((R1 (R2s Q_2) ;; R1 (R2s S_2))))$

proof –

have $1: (\neg ((R1 (R2s (Q_1 \diamond Q_2)) \wedge \neg \$wait') ;; R1 (\neg R2s R))) =$
 $(\neg ((R1 (R2s Q_2) \wedge \neg \$wait') ;; R1 (\neg R2s R)))$

by (*metis* (*no-types*, *hide-lams*) *R1-extend-conj* *R2s-conj* *R2s-not* *R2s-wait'* *wait'-cond-false*)

have $2: (R1 (R2s (Q_1 \diamond Q_2)) ;; ([II]_D \triangleleft \$wait \triangleright R1 (R2s (S_1 \diamond S_2)))) =$
 $((R1 (R2s Q_1) \vee (R1 (R2s Q_2) ;; R1 (R2s S_1))) \diamond (R1 (R2s Q_2) ;; R1 (R2s S_2)))$

proof –

have $(R1 (R2s Q_1) ;; (\$wait \wedge ([II]_D \triangleleft \$wait \triangleright R1 (R2s S_1) \diamond R1 (R2s S_2))))$
 $= (R1 (R2s Q_1) \wedge \$wait')$

proof –

have $(R1 (R2s Q_1) ;; (\$wait \wedge ([II]_D \triangleleft \$wait \triangleright R1 (R2s S_1) \diamond R1 (R2s S_2))))$
 $= (R1 (R2s Q_1) ;; (\$wait \wedge [II]_D))$

by (*rel-auto*)

also have $\dots = ((R1 (R2s Q_1) ;; [II]_D) \wedge \$wait')$

by (*rel-auto*)

also from *assms*(2) **have** $\dots = ((R1 (R2s Q_1)) \wedge \$wait')$

by (*simp* *add: lift-des-skip-dr-unit-unrest unrest*)

finally show *?thesis* .

qed

moreover have $(R1 (R2s Q_2) ;; (\neg \$wait \wedge ([II]_D \triangleleft \$wait \triangleright R1 (R2s S_1) \diamond R1 (R2s S_2))))$
 $= ((R1 (R2s Q_2) ;; (R1 (R2s S_1) \diamond R1 (R2s S_2)))$

proof –

have $(R1 (R2s Q_2) ;; (\neg \$wait \wedge ([II]_D \triangleleft \$wait \triangleright R1 (R2s S_1) \diamond R1 (R2s S_2))))$
 $= (R1 (R2s Q_2) ;; (\neg \$wait \wedge (R1 (R2s S_1) \diamond R1 (R2s S_2))))$

by (*metis* (*no-types*, *lifting*) *cond-def* *conj-disj-not-abs* *utp-pred-laws.double-compl* *utp-pred-laws.inf.left-idem* *utp-pred-laws.sup-assoc* *utp-pred-laws.sup-inf-absorb*)

also have $\dots = ((R1 (R2s Q_2)) \llbracket \text{false} / \$wait' \rrbracket ;; (R1 (R2s S_1) \diamond R1 (R2s S_2)) \llbracket \text{false} / \$wait \rrbracket)$
by (*metis* *false-alt-def* *seqr-right-one-point upred-eq-false wait-vwb-lens*)

also have $\dots = ((R1 (R2s Q_2) ;; (R1 (R2s S_1) \diamond R1 (R2s S_2)))$
by (*simp* *add: wait'-cond-def usubst unrest assms*)

finally show *?thesis* .

qed

moreover
have $((R1 \ (R2s \ Q_1) \wedge \$wait') \vee ((R1 \ (R2s \ Q_2)) ;; (R1 \ (R2s \ S_1) \diamond R1 \ (R2s \ S_2))))$
 $= (R1 \ (R2s \ Q_1) \vee (R1 \ (R2s \ Q_2) ;; R1 \ (R2s \ S_1))) \diamond ((R1 \ (R2s \ Q_2) ;; R1 \ (R2s \ S_2)))$
by (*simp add: wait'-cond-def cond-seq-right-distr cond-and-T-integrate unrest*)

ultimately show *?thesis*
by (*simp add: R2s-wait'-cond R1-wait'-cond wait'-cond-seq*)
qed

show *?thesis*
apply (*subst RH-design-composition*)
apply (*simp-all add: assms*)
apply (*simp add: assms wait'-cond-def unrest*)
apply (*simp add: assms wait'-cond-def unrest*)
apply (*simp add: 1 2*)
apply (*simp add: R1-R2s-R2c RH-design-lemma1*)
done
qed

theorem *R1-design-composition-RR:*
assumes *P is RR Q is RR R is RR S is RR*
shows
 $(R1(P \vdash Q) ;; R1(R \vdash S)) = R1(((\neg_r P) \text{ wp}_r \text{ false} \wedge Q \text{ wp}_r R) \vdash (Q ;; S))$
apply (*subst R1-design-composition*)
apply (*simp-all add: assms unrest wp-rea-def Healthy-if closure*)
apply (*rel-auto*)
done

theorem *R1-design-composition-RC:*
assumes *P is RC Q is RR R is RR S is RR*
shows
 $(R1(P \vdash Q) ;; R1(R \vdash S)) = R1((P \wedge Q \text{ wp}_r R) \vdash (Q ;; S))$
by (*simp add: R1-design-composition-RR assms unrest Healthy-if closure wp*)

theorem *RHS-tri-design-composition:*
assumes $\$ok' \# P \ \$ok' \# Q_1 \ \$ok' \# Q_2 \ \$ok \# R \ \$ok \# S_1 \ \$ok \# S_2$
 $\$wait \# R \ \$wait' \# Q_2 \ \$wait \# S_1 \ \$wait \# S_2$
shows $(\mathbf{R}_s(P \vdash Q_1 \diamond Q_2) ;; \mathbf{R}_s(R \vdash S_1 \diamond S_2)) =$
 $\mathbf{R}_s((\neg (R1 \ (\neg R2s \ P) ;; R1 \ true) \wedge \neg (R1(R2s \ Q_2) ;; R1 \ (\neg R2s \ R))) \vdash$
 $((\exists \$st' \cdot Q_1) \vee (R1 \ (R2s \ Q_2) ;; R1 \ (R2s \ S_1))) \diamond ((R1 \ (R2s \ Q_2) ;; R1 \ (R2s \ S_2))))$

proof –
have $1: (\neg ((R1 \ (R2s \ (Q_1 \diamond Q_2)) \wedge \neg \$wait') ;; R1 \ (\neg R2s \ R))) =$
 $(\neg ((R1 \ (R2s \ Q_2) \wedge \neg \$wait') ;; R1 \ (\neg R2s \ R)))$
by (*metis (no-types, hide-lams) R1-extend-conj R2s-conj R2s-not R2s-wait' wait'-cond-false*)
have $2: (R1 \ (R2s \ (Q_1 \diamond Q_2)) ;; ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 \ (R2s \ (S_1 \diamond S_2)))) =$
 $((\exists \$st' \cdot R1 \ (R2s \ Q_1)) \vee (R1 \ (R2s \ Q_2) ;; R1 \ (R2s \ S_1))) \diamond (R1 \ (R2s \ Q_2) ;; R1 \ (R2s \ S_2)))$

proof –
have $(R1 \ (R2s \ Q_1) ;; (\$wait \wedge ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 \ (R2s \ S_1) \diamond R1 \ (R2s \ S_2))))$
 $= (\exists \$st' \cdot ((R1 \ (R2s \ Q_1)) \wedge \$wait'))$

proof –
have $(R1 \ (R2s \ Q_1) ;; (\$wait \wedge ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 \ (R2s \ S_1) \diamond R1 \ (R2s \ S_2))))$
 $= (R1 \ (R2s \ Q_1) ;; (\$wait \wedge (\exists \$st \cdot \lceil II \rceil_D)))$
by (*rel-auto, blast+*)
also have $\dots = ((R1 \ (R2s \ Q_1) ;; (\exists \$st \cdot \lceil II \rceil_D)) \wedge \$wait')$

by (rel-auto)
 also from *assms*(2) have ... = $(\exists \$st' \cdot ((R1 (R2s Q_1)) \wedge \$wait'))$
 by (rel-auto, blast)
 finally show ?thesis .
 qed

moreover have $(R1 (R2s Q_2) ;; (\neg \$wait \wedge ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 (R2s S_1) \diamond R1 (R2s S_2))))$
 $= ((R1 (R2s Q_2)) ;; (R1 (R2s S_1) \diamond R1 (R2s S_2)))$
 proof –
 have $(R1 (R2s Q_2) ;; (\neg \$wait \wedge ((\exists \$st \cdot \lceil II \rceil_D) \triangleleft \$wait \triangleright R1 (R2s S_1) \diamond R1 (R2s S_2))))$
 $= (R1 (R2s Q_2) ;; (\neg \$wait \wedge (R1 (R2s S_1) \diamond R1 (R2s S_2))))$
 by (metis (no-types, lifting) cond-def conj-disj-not-abs utp-pred-laws.double-compl utp-pred-laws.inf.left-idem
 utp-pred-laws.sup-assoc utp-pred-laws.sup-inf-absorb)
 also have ... = $((R1 (R2s Q_2))\llbracket false/\$wait \rrbracket ;; (R1 (R2s S_1) \diamond R1 (R2s S_2))\llbracket false/\$wait \rrbracket)$
 by (metis false-alt-def seqr-right-one-point upred-eq-false wait-vwb-lens)
 also have ... = $((R1 (R2s Q_2)) ;; (R1 (R2s S_1) \diamond R1 (R2s S_2)))$
 by (simp add: wait'-cond-def usubst unrest assms)
 finally show ?thesis .
 qed

moreover
 have $((R1 (R2s Q_1) \wedge \$wait') \vee ((R1 (R2s Q_2)) ;; (R1 (R2s S_1) \diamond R1 (R2s S_2))))$
 $= (R1 (R2s Q_1) \vee (R1 (R2s Q_2) ;; R1 (R2s S_1))) \diamond ((R1 (R2s Q_2) ;; R1 (R2s S_2)))$
 by (simp add: wait'-cond-def cond-seq-right-distr cond-and-T-integrate unrest)
 ultimately show ?thesis
 by (simp add: R2s-wait'-cond R1-wait'-cond wait'-cond-seq ex-conj-contr-right unrest)
 (simp add: cond-and-T-integrate cond-seq-right-distr unrest-var wait'-cond-def)
 qed

from *assms*(7,8) have $\exists: (R1 (R2s Q_2) \wedge \neg \$wait') ;; R1 (\neg R2s R) = R1 (R2s Q_2) ;; R1 (\neg R2s R)$
 by (rel-auto, blast, meson)
 show ?thesis
 apply (subst RHS-design-composition)
 apply (simp-all add: assms)
 apply (simp add: assms wait'-cond-def unrest)
 apply (simp add: assms wait'-cond-def unrest)
 apply (simp add: 1 2 3)
 apply (simp add: R1-R2s-R2c RHS-design-lemma1)
 apply (metis R1-R2c-ex-st RHS-design-lemma1)
 done
 qed

theorem *RHS-tri-design-composition-wp*:
 assumes $\$ok' \# P \$ok' \# Q_1 \$ok' \# Q_2 \$ok \# R \$ok \# S_1 \$ok \# S_2$
 $\$wait \# R \$wait' \# Q_2 \$wait \# S_1 \$wait \# S_2$
 $P \text{ is } R2c \ Q_1 \text{ is } R1 \ Q_1 \text{ is } R2c \ Q_2 \text{ is } R1 \ Q_2 \text{ is } R2c$
 $R \text{ is } R2c \ S_1 \text{ is } R1 \ S_1 \text{ is } R2c \ S_2 \text{ is } R1 \ S_2 \text{ is } R2c$
 shows $\mathbf{R}_s(P \vdash Q_1 \diamond Q_2) ;; \mathbf{R}_s(R \vdash S_1 \diamond S_2) =$

$\mathbf{R}_s(((\neg_r P) \text{ wp}_r \text{ false} \wedge Q_2 \text{ wp}_r R) \vdash (((\exists \$st' \cdot Q_1) \sqcap (Q_2 ;; S_1)) \diamond (Q_2 ;; S_2)))$ (is ?lhs =
 ?rhs)
proof –
 have ?lhs = $\mathbf{R}_s(((\neg R1 (\neg P) ;; R1 \text{ true} \wedge \neg Q_2 ;; R1 (\neg R)) \vdash (((\exists \$st' \cdot Q_1) \sqcap (Q_2 ;; S_1)) \diamond (Q_2 ;; S_2)))$
 by (simp add: RHS-tri-design-composition assms Healthy-if R2c-healthy-R2s disj-upred-def)
 (metis (no-types, hide-lams) R1-negate-R1 R2c-healthy-R2s assms(11,16))
 also have ... = ?rhs
 by (rel-auto)
 finally show ?thesis .
qed

theorem *RHS-tri-design-composition-RR-wp*:

assumes P is RR Q_1 is RR Q_2 is RR

R is RR S_1 is RR S_2 is RR

shows $\mathbf{R}_s(P \vdash Q_1 \diamond Q_2) ;; \mathbf{R}_s(R \vdash S_1 \diamond S_2) =$

$\mathbf{R}_s(((\neg_r P) \text{ wp}_r \text{ false} \wedge Q_2 \text{ wp}_r R) \vdash (((\exists \$st' \cdot Q_1) \sqcap (Q_2 ;; S_1)) \diamond (Q_2 ;; S_2)))$ (is ?lhs =
 ?rhs)

by (simp add: RHS-tri-design-composition-wp add: closure assms unrest RR-implies-R2c)

lemma *RHS-tri-normal-design-composition*:

assumes

$\$ok' \# P \$ok' \# Q_1 \$ok' \# Q_2 \$ok \# R \$ok \# S_1 \$ok \# S_2$

$\$wait \# R \$wait' \# Q_2 \$wait \# S_1 \$wait \# S_2$

P is R2c Q_1 is R1 Q_1 is R2c Q_2 is R1 Q_2 is R2c

R is R2c S_1 is R1 S_1 is R2c S_2 is R1 S_2 is R2c

$R1 (\neg P) ;; R1(\text{true}) = R1(\neg P) \$st' \# Q_1$

shows $\mathbf{R}_s(P \vdash Q_1 \diamond Q_2) ;; \mathbf{R}_s(R \vdash S_1 \diamond S_2)$

$= \mathbf{R}_s((P \wedge Q_2 \text{ wp}_r R) \vdash (Q_1 \vee (Q_2 ;; S_1)) \diamond (Q_2 ;; S_2))$

proof –

have $\mathbf{R}_s(P \vdash Q_1 \diamond Q_2) ;; \mathbf{R}_s(R \vdash S_1 \diamond S_2) =$

$\mathbf{R}_s((R1 (\neg P) \text{ wp}_r \text{ false} \wedge Q_2 \text{ wp}_r R) \vdash (((\exists \$st' \cdot Q_1) \sqcap (Q_2 ;; S_1)) \diamond (Q_2 ;; S_2)))$

by (simp-all add: RHS-tri-design-composition-wp rea-not-def assms unrest)

also have ... = $\mathbf{R}_s((P \wedge Q_2 \text{ wp}_r R) \vdash (Q_1 \vee (Q_2 ;; S_1)) \diamond (Q_2 ;; S_2))$

by (simp add: assms wp-rea-def ex-unrest, rel-auto)

finally show ?thesis .

qed

lemma *RHS-tri-normal-design-composition' [rdes-def]*:

assumes P is RC Q_1 is RR $\$st' \# Q_1$ Q_2 is RR R is RR S_1 is RR S_2 is RR

shows $\mathbf{R}_s(P \vdash Q_1 \diamond Q_2) ;; \mathbf{R}_s(R \vdash S_1 \diamond S_2)$

$= \mathbf{R}_s((P \wedge Q_2 \text{ wp}_r R) \vdash (Q_1 \vee (Q_2 ;; S_1)) \diamond (Q_2 ;; S_2))$

proof –

have $R1 (\neg P) ;; R1 \text{ true} = R1(\neg P)$

using RC-implies-RC1[OF assms(1)]

by (simp add: Healthy-def RC1-def rea-not-def)

(metis R1-negate-R1 R1-seqr utp-pred-laws.double-compl)

thus ?thesis

by (simp add: RHS-tri-normal-design-composition assms closure unrest RR-implies-R2c)

qed

lemma *RHS-tri-design-right-unit-lemma*:

assumes $\$ok' \# P \$ok' \# Q \$ok' \# R \$wait' \# R$

shows $\mathbf{R}_s(P \vdash Q \diamond R) ;; \Pi_R = \mathbf{R}_s((\neg_r (\neg_r P) ;; \text{true}_r) \vdash ((\exists \$st' \cdot Q) \diamond R))$

proof –

have $\mathbf{R}_s(P \vdash Q \diamond R) ;; II_R = \mathbf{R}_s(P \vdash Q \diamond R) ;; \mathbf{R}_s(\text{true} \vdash \text{false} \diamond (\$tr' =_u \$tr \wedge [II]_R))$
by (*simp add: srdes-skip-tri-design, rel-auto*)
also have $\dots = \mathbf{R}_s((\neg R1 \ (\neg R2s \ P) ;; R1 \ \text{true}) \vdash (\exists \$st' \cdot Q) \diamond (R1 \ (R2s \ R) ;; R1 \ (R2s \ (\$tr' =_u \$tr \wedge [II]_R))))$
by (*simp-all add: RHS-tri-design-composition assms unrest R2s-true R1-false R2s-false*)
also have $\dots = \mathbf{R}_s((\neg R1 \ (\neg R2s \ P) ;; R1 \ \text{true}) \vdash (\exists \$st' \cdot Q) \diamond R1 \ (R2s \ R))$
proof –
from *assms(3,4)* **have** $(R1 \ (R2s \ R) ;; R1 \ (R2s \ (\$tr' =_u \$tr \wedge [II]_R))) = R1 \ (R2s \ R)$
by (*rel-auto, metis (no-types, lifting) minus-zero-eq, meson order-refl trace-class.diff-cancel*)
thus *?thesis*
by *simp*
qed
also have $\dots = \mathbf{R}_s((\neg (\neg P) ;; R1 \ \text{true}) \vdash ((\exists \$st' \cdot Q) \diamond R))$
by (*metis (no-types, lifting) R1-R2s-R1-true-lemma R1-R2s-R2c R2c-not RHS-design-R2c-pre RHS-design-neg-R1-pre RHS-design-post-R1 RHS-design-post-R2s*)
also have $\dots = \mathbf{R}_s((\neg_r (\neg_r P) ;; \text{true}_r) \vdash ((\exists \$st' \cdot Q) \diamond R))$
by (*rel-auto*)
finally show *?thesis* .
qed

lemma *SRD-composition-wp*:

assumes *P is SRD Q is SRD*

shows $(P ;; Q) = \mathbf{R}_s(((\neg_r \text{pre}_R P) \text{wp}_r \text{false} \wedge \text{post}_R P \text{wp}_r \text{pre}_R Q) \vdash ((\exists \$st' \cdot \text{peri}_R P) \vee (\text{post}_R P ;; \text{peri}_R Q)) \diamond (\text{post}_R P ;; \text{post}_R Q))$

(*is ?lhs = ?rhs*)

proof –

have $(P ;; Q) = (\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P)) ;; \mathbf{R}_s(\text{pre}_R(Q) \vdash \text{peri}_R(Q) \diamond \text{post}_R(Q)))$

by (*simp add: SRD-reactive-tri-design assms(1) assms(2)*)

also from *assms*

have $\dots = ?rhs$

by (*simp add: RHS-tri-design-composition-wp disj-upred-def unrest assms closure*)

finally show *?thesis* .

qed

4.6 Refinement introduction laws

lemma *RHS-tri-design-refine*:

assumes *P₁ is RR P₂ is RR P₃ is RR Q₁ is RR Q₂ is RR Q₃ is RR*

shows $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \sqsubseteq \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3) \longleftrightarrow 'P_1 \Rightarrow Q_1' \wedge 'P_1 \wedge Q_2 \Rightarrow P_2' \wedge 'P_1 \wedge Q_3 \Rightarrow P_3'$

(*is ?lhs = ?rhs*)

proof –

have $?lhs \longleftrightarrow 'P_1 \Rightarrow Q_1' \wedge 'P_1 \wedge Q_2 \diamond Q_3 \Rightarrow P_2 \diamond P_3'$

by (*simp add: RHS-design-refine assms closure RR-implies-R2c unrest ex-unrest*)

also have $\dots \longleftrightarrow 'P_1 \Rightarrow Q_1' \wedge '(P_1 \wedge Q_2) \diamond (P_1 \wedge Q_3) \Rightarrow P_2 \diamond P_3'$

by (*rel-auto*)

also have $\dots \longleftrightarrow 'P_1 \Rightarrow Q_1' \wedge '((P_1 \wedge Q_2) \diamond (P_1 \wedge Q_3) \Rightarrow P_2 \diamond P_3)[[true/\$wait]]' \wedge '((P_1 \wedge Q_2) \diamond (P_1 \wedge Q_3) \Rightarrow P_2 \diamond P_3)[[false/\$wait]]'$

by (*rel-auto, metis*)

also have $\dots \longleftrightarrow ?rhs$

by (*simp add: usubst unrest assms*)

finally show *?thesis* .

qed

lemma *RHS-tri-design-refine'*:

assumes *P₁ is RR P₂ is RR P₃ is RR Q₁ is RR Q₂ is RR Q₃ is RR*

shows $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \sqsubseteq \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3) \longleftrightarrow (Q_1 \sqsubseteq P_1) \wedge (P_2 \sqsubseteq (P_1 \wedge Q_2)) \wedge (P_3 \sqsubseteq (P_1 \wedge Q_3))$
by (*simp add: RHS-tri-design-refine assms, rel-auto*)

lemma *srdes-tri-refine-intro*:

assumes $\text{'}P_1 \Rightarrow P_2\text{'}$ $\text{'}P_1 \wedge Q_2 \Rightarrow Q_1\text{'}$ $\text{'}P_1 \wedge R_2 \Rightarrow R_1\text{'}$
shows $\mathbf{R}_s(P_1 \vdash Q_1 \diamond R_1) \sqsubseteq \mathbf{R}_s(P_2 \vdash Q_2 \diamond R_2)$
using *assms*
by (*rule-tac srdes-refine-intro, simp-all, rel-auto*)

lemma *srdes-tri-eq-intro*:

assumes $P_1 = Q_1$ $P_2 = Q_2$ $P_3 = Q_3$
shows $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) = \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3)$
using *assms* **by** (*simp*)

lemma *srdes-tri-refine-intro'*:

assumes $P_2 \sqsubseteq P_1$ $Q_1 \sqsubseteq (P_1 \wedge Q_2)$ $R_1 \sqsubseteq (P_1 \wedge R_2)$
shows $\mathbf{R}_s(P_1 \vdash Q_1 \diamond R_1) \sqsubseteq \mathbf{R}_s(P_2 \vdash Q_2 \diamond R_2)$
using *assms*
by (*rule-tac srdes-tri-refine-intro, simp-all add: refBy-order*)

lemma *SRD-peri-under-pre*:

assumes P is SRD $\$wait' \nmid pre_R(P)$
shows $(pre_R(P) \Rightarrow_r peri_R(P)) = peri_R(P)$

proof –

have $peri_R(P) =$
 $peri_R(\mathbf{R}_s(pre_R(P) \vdash peri_R(P) \diamond post_R(P)))$
by (*simp add: SRD-reactive-tri-design assms*)
also have $\dots = (pre_R P \Rightarrow_r peri_R P)$
by (*simp add: rea-pre-RHS-design rea-peri-RHS-design assms*
 $unrest usubst R1-peri-SRD R2c-preR R1-rea-impl R2c-rea-impl R2c-periR$)
finally show *?thesis* ..

qed

lemma *SRD-post-under-pre*:

assumes P is SRD $\$wait' \nmid pre_R(P)$
shows $(pre_R(P) \Rightarrow_r post_R(P)) = post_R(P)$

proof –

have $post_R(P) =$
 $post_R(\mathbf{R}_s(pre_R(P) \vdash peri_R(P) \diamond post_R(P)))$
by (*simp add: SRD-reactive-tri-design assms*)
also have $\dots = (pre_R P \Rightarrow_r post_R P)$
by (*simp add: rea-pre-RHS-design rea-post-RHS-design assms*
 $unrest usubst R1-post-SRD R2c-preR R1-rea-impl R2c-rea-impl R2c-postR$)
finally show *?thesis* ..

qed

lemma *SRD-refine-intro*:

assumes
 P is SRD Q is SRD
 $\text{'}pre_R(P) \Rightarrow pre_R(Q)\text{'}$ $\text{'}pre_R(P) \wedge peri_R(Q) \Rightarrow peri_R(P)\text{'}$ $\text{'}pre_R(P) \wedge post_R(Q) \Rightarrow post_R(P)\text{'}$
shows $P \sqsubseteq Q$
by (*metis SRD-reactive-tri-design assms(1) assms(2) assms(3) assms(4) assms(5) srdes-tri-refine-intro*)

lemma *SRD-refine-intro'*:

assumes
 P is SRD Q is SRD
 $\text{'pre}_R(P) \Rightarrow \text{pre}_R(Q) \text{'}$ $\text{peri}_R(P) \sqsubseteq (\text{pre}_R(P) \wedge \text{peri}_R(Q))$ $\text{post}_R(P) \sqsubseteq (\text{pre}_R(P) \wedge \text{post}_R(Q))$
shows $P \sqsubseteq Q$
using *assms* **by** (*rule-tac SRD-refine-intro, simp-all add: refBy-order*)

lemma *SRD-eq-intro*:

assumes
 P is SRD Q is SRD $\text{pre}_R(P) = \text{pre}_R(Q)$ $\text{peri}_R(P) = \text{peri}_R(Q)$ $\text{post}_R(P) = \text{post}_R(Q)$
shows $P = Q$
by (*metis SRD-reactive-tri-design assms*)

4.7 Closure laws

lemma *SRD-srdes-skip [closure]*: Π_R is SRD

by (*simp add: srdes-skip-def RHS-design-is-SRD unrest*)

lemma *SRD-seqr-closure [closure]*:

assumes P is SRD Q is SRD

shows $(P ;; Q)$ is SRD

proof –

have $(P ;; Q) = \mathbf{R}_s (((\neg_r \text{pre}_R P) \text{wp}_r \text{false} \wedge \text{post}_R P \text{wp}_r \text{pre}_R Q) \vdash$
 $((\exists \$st' \cdot \text{peri}_R P) \vee (\text{post}_R P ;; \text{peri}_R Q)) \diamond (\text{post}_R P ;; \text{post}_R Q))$

by (*simp add: SRD-composition-wp assms(1) assms(2)*)

also have ... is SRD

by (*rule RHS-design-is-SRD, simp-all add: wp-rea-def unrest*)

finally show ?thesis .

qed

lemma *SRD-power-Suc [closure]*: P is SRD $\implies P^\wedge(\text{Suc } n)$ is SRD

proof (*induct n*)

case 0

then show ?case

by (*simp*)

next

case (*Suc n*)

then show ?case

using *SRD-seqr-closure* **by** (*simp add: SRD-seqr-closure upred-semiring.power-Suc*)

qed

lemma *SRD-power-comp [closure]*: P is SRD $\implies P ;; P^\wedge n$ is SRD

by (*metis SRD-power-Suc upred-semiring.power-Suc*)

lemma *uplus-SRD-closed [closure]*: P is SRD $\implies P^+$ is SRD

by (*simp add: uplus-power-def closure*)

lemma *SRD-Sup-closure [closure]*:

assumes $A \subseteq \llbracket \text{SRD} \rrbracket_H$ $A \neq \{\}$

shows $(\bigcap A)$ is SRD

proof –

have $\text{SRD } (\bigcap A) = (\bigcap (\text{SRD } A))$

by (*simp add: ContinuousD SRD-Continuous assms(2)*)

also have ... = $(\bigcap A)$

by (*simp only: Healthy-carrier-image assms*)

finally show ?thesis **by** (*simp add: Healthy-def*)

qed

4.8 Distribution laws

lemma *RHS-tri-design-choice* [rdes-def]:

$\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \sqcap \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3) = \mathbf{R}_s((P_1 \wedge Q_1) \vdash (P_2 \vee Q_2) \diamond (P_3 \vee Q_3))$
apply (*simp add: RHS-design-choice*)
apply (*rule cong[of \mathbf{R}_s \mathbf{R}_s]*)
apply (*simp*)
apply (*rel-auto*)
done

lemma *RHS-tri-design-disj* [rdes-def]:

$(\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \vee \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3)) = \mathbf{R}_s((P_1 \wedge Q_1) \vdash (P_2 \vee Q_2) \diamond (P_3 \vee Q_3))$
by (*simp add: RHS-tri-design-choice disj-upred-def*)

lemma *RHS-tri-design-sup* [rdes-def]:

$\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \sqcup \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3) = \mathbf{R}_s((P_1 \vee Q_1) \vdash ((P_1 \Rightarrow_r P_2) \wedge (Q_1 \Rightarrow_r Q_2)) \diamond ((P_1 \Rightarrow_r P_3) \wedge (Q_1 \Rightarrow_r Q_3)))$
by (*simp add: RHS-design-sup, rel-auto*)

lemma *RHS-tri-design-conj* [rdes-def]:

$(\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \wedge \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3)) = \mathbf{R}_s((P_1 \vee Q_1) \vdash ((P_1 \Rightarrow_r P_2) \wedge (Q_1 \Rightarrow_r Q_2)) \diamond ((P_1 \Rightarrow_r P_3) \wedge (Q_1 \Rightarrow_r Q_3)))$
by (*simp add: RHS-tri-design-sup conj-upred-def*)

lemma *SRD-UINF* [rdes-def]:

assumes $A \neq \{\}$ $A \subseteq \llbracket \text{SRD} \rrbracket_H$
shows $\sqcap A = \mathbf{R}_s((\bigwedge P \in A \cdot \text{pre}_R(P)) \vdash (\bigvee P \in A \cdot \text{peri}_R(P)) \diamond (\bigvee P \in A \cdot \text{post}_R(P)))$
proof –
have $\sqcap A = \mathbf{R}_s(\text{pre}_R(\sqcap A) \vdash \text{peri}_R(\sqcap A) \diamond \text{post}_R(\sqcap A))$
by (*metis SRD-as-reactive-tri-design assms srdes-hcond-def*
srdes-theory-continuous.healthy-inf srdes-theory-continuous.healthy-inf-def)
also have $\dots = \mathbf{R}_s((\bigwedge P \in A \cdot \text{pre}_R(P)) \vdash (\bigvee P \in A \cdot \text{peri}_R(P)) \diamond (\bigvee P \in A \cdot \text{post}_R(P)))$
by (*simp add: preR-INF periR-INF postR-INF assms*)
finally show *?thesis* .
qed

lemma *RHS-tri-design-USUP* [rdes-def]:

assumes $A \neq \{\}$
shows $(\sqcap i \in A \cdot \mathbf{R}_s(P(i) \vdash Q(i) \diamond R(i))) = \mathbf{R}_s((\bigwedge i \in A \cdot P(i)) \vdash (\bigwedge i \in A \cdot Q(i)) \diamond (\bigwedge i \in A \cdot R(i)))$
by (*subst RHS-INF[OF assms, THEN sym], simp add: design-UINF-mem assms, rel-auto*)

lemma *SRD-UINF-mem*:

assumes $A \neq \{\}$ $\bigwedge i. P \ i \text{ is } \text{SRD}$
shows $(\sqcap i \in A \cdot P \ i) = \mathbf{R}_s((\bigwedge i \in A \cdot \text{pre}_R(P \ i)) \vdash (\bigvee i \in A \cdot \text{peri}_R(P \ i)) \diamond (\bigvee i \in A \cdot \text{post}_R(P \ i)))$
(is ?lhs = ?rhs)
proof –
have $?lhs = (\sqcap (P \ ' A))$
by (*rel-auto*)
also have $\dots = \mathbf{R}_s((\bigwedge Pa \in P \ ' A \cdot \text{pre}_R \ Pa) \vdash (\bigwedge Pa \in P \ ' A \cdot \text{peri}_R \ Pa) \diamond (\bigwedge Pa \in P \ ' A \cdot \text{post}_R \ Pa))$
by (*subst rdes-def, simp-all add: assms image-subsetI*)
also have $\dots = ?rhs$
by (*rel-auto*)
finally show *?thesis* .
qed

lemma *RHS-tri-design-UINF-ind* [*rdes-def*]:

$(\prod i \cdot \mathbf{R}_s(P_1(i) \vdash P_2(i) \diamond P_3(i))) = \mathbf{R}_s((\bigwedge i \cdot P_1 i) \vdash (\bigvee i \cdot P_2(i)) \diamond (\bigvee i \cdot P_3(i)))$
 by (*rel-auto*)

lemma *cond-srea-form* [*rdes-def*]:

$\mathbf{R}_s(P \vdash Q_1 \diamond Q_2) \triangleleft b \triangleright_R \mathbf{R}_s(R \vdash S_1 \diamond S_2) =$
 $\mathbf{R}_s((P \triangleleft b \triangleright_R R) \vdash (Q_1 \triangleleft b \triangleright_R S_1) \diamond (Q_2 \triangleleft b \triangleright_R S_2))$

proof –

have $\mathbf{R}_s(P \vdash Q_1 \diamond Q_2) \triangleleft b \triangleright_R \mathbf{R}_s(R \vdash S_1 \diamond S_2) = \mathbf{R}_s(P \vdash Q_1 \diamond Q_2) \triangleleft R2c(\lceil b \rceil_{S<}) \triangleright \mathbf{R}_s(R \vdash S_1 \diamond S_2)$

by (*pred-auto*)

also have $\dots = \mathbf{R}_s(P \vdash Q_1 \diamond Q_2 \triangleleft b \triangleright_R R \vdash S_1 \diamond S_2)$

by (*simp add: RHS-cond lift-cond-srea-def*)

also have $\dots = \mathbf{R}_s((P \triangleleft b \triangleright_R R) \vdash (Q_1 \diamond Q_2 \triangleleft b \triangleright_R S_1 \diamond S_2))$

by (*simp add: design-condr lift-cond-srea-def*)

also have $\dots = \mathbf{R}_s((P \triangleleft b \triangleright_R R) \vdash (Q_1 \triangleleft b \triangleright_R S_1) \diamond (Q_2 \triangleleft b \triangleright_R S_2))$

by (*rule cong[of \mathbf{R}_s \mathbf{R}_s], simp, rel-auto*)

finally show *?thesis* .

qed

lemma *SRD-cond-srea* [*closure*]:

assumes *P is SRD Q is SRD*

shows $P \triangleleft b \triangleright_R Q$ is *SRD*

proof –

have $P \triangleleft b \triangleright_R Q = \mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P)) \triangleleft b \triangleright_R \mathbf{R}_s(\text{pre}_R(Q) \vdash \text{peri}_R(Q) \diamond \text{post}_R(Q))$

by (*simp add: SRD-reactive-tri-design assms*)

also have $\dots = \mathbf{R}_s((\text{pre}_R P \triangleleft b \triangleright_R \text{pre}_R Q) \vdash (\text{peri}_R P \triangleleft b \triangleright_R \text{peri}_R Q) \diamond (\text{post}_R P \triangleleft b \triangleright_R \text{post}_R Q))$

by (*simp add: cond-srea-form*)

also have \dots is *SRD*

by (*simp add: RHS-tri-design-is-SRD lift-cond-srea-def unrest*)

finally show *?thesis* .

qed

4.9 Algebraic laws

lemma *SRD-left-unit*:

assumes *P is SRD*

shows $\Pi_R ;; P = P$

by (*simp add: SRD-composition-wp closure rdes wp C1 R1-negate-R1 R1-false*

rpred trace-ident-left-periR trace-ident-left-postR SRD-reactive-tri-design assms)

lemma *skip-srea-self-unit* [*simp*]:

$\Pi_R ;; \Pi_R = \Pi_R$

by (*simp add: SRD-left-unit closure*)

lemma *SRD-right-unit-tri-lemma*:

assumes *P is SRD*

shows $P ;; \Pi_R = \mathbf{R}_s((\neg_r \text{pre}_R P) \text{wp}_r \text{false} \vdash (\exists \$st' \cdot \text{peri}_R P) \diamond \text{post}_R P)$

by (*simp add: SRD-composition-wp closure rdes wp rpred trace-ident-right-postR assms*)

lemma *Miracle-left-zero*:

assumes *P is SRD*

shows *Miracle* ;; $P = \text{Miracle}$

proof –

have *Miracle* ;; $P = \mathbf{R}_s(\text{true} \vdash \text{false})$;; $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{cmt}_R(P))$
by (*simp add: Miracle-def SRD-reactive-design-alt assms*)
also have ... = $\mathbf{R}_s(\text{true} \vdash \text{false})$
by (*simp add: RHS-design-composition unrest R1-false R2s-false R2s-true*)
also have ... = *Miracle*
by (*simp add: Miracle-def*)
finally show ?thesis .
qed

lemma *Chaos-left-zero*:
assumes P is *SRD*
shows ($\text{Chaos} ;; P$) = *Chaos*

proof –

have $\text{Chaos} ;; P = \mathbf{R}_s(\text{false} \vdash \text{true})$;; $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{cmt}_R(P))$
by (*simp add: Chaos-def SRD-reactive-design-alt assms*)
also have ... = $\mathbf{R}_s((\neg R1 \text{ true} \wedge \neg (R1 \text{ true} \wedge \neg \$\text{wait}')) ;; R1 (\neg R2s (\text{pre}_R P))) \vdash$
 $R1 \text{ true} ;; ((\exists \$st \cdot [II]_D) \triangleleft \$\text{wait} \triangleright R1 (R2s (\text{cmt}_R P)))$
by (*simp add: RHS-design-composition unrest R2s-false R2s-true R1-false*)
also have ... = $\mathbf{R}_s((\text{false} \wedge \neg (R1 \text{ true} \wedge \neg \$\text{wait}')) ;; R1 (\neg R2s (\text{pre}_R P))) \vdash$
 $R1 \text{ true} ;; ((\exists \$st \cdot [II]_D) \triangleleft \$\text{wait} \triangleright R1 (R2s (\text{cmt}_R P)))$
by (*simp add: RHS-design-conj-neg-R1-pre*)
also have ... = $\mathbf{R}_s(\text{true})$
by (*simp add: design-false-pre*)
also have ... = $\mathbf{R}_s(\text{false} \vdash \text{true})$
by (*simp add: design-def*)
also have ... = *Chaos*
by (*simp add: Chaos-def*)
finally show ?thesis .
qed

lemma *SRD-right-Chaos-tri-lemma*:
assumes P is *SRD*
shows $P ;; \text{Chaos} = \mathbf{R}_s(((\neg_r \text{pre}_R P) \text{wp}_r \text{false} \wedge \text{post}_R P \text{wp}_r \text{false}) \vdash (\exists \$st' \cdot \text{peri}_R P) \diamond \text{false})$
by (*simp add: SRD-composition-wp closure rdes assms wp, rel-auto*)

lemma *SRD-right-Miracle-tri-lemma*:
assumes P is *SRD*
shows $P ;; \text{Miracle} = \mathbf{R}_s((\neg_r \text{pre}_R P) \text{wp}_r \text{false} \vdash (\exists \$st' \cdot \text{peri}_R P) \diamond \text{false})$
by (*simp add: SRD-composition-wp closure rdes assms wp, rel-auto*)

Stateful reactive designs are left unital

overloading

$\text{srdes-unit} == \text{utp-unit} :: (\text{SRDES}, ('s, 't :: \text{trace}, ' \alpha) \text{rsp}) \text{uthy} \Rightarrow ('s, 't, ' \alpha) \text{hrel-rsp}$

begin

definition $\text{srdes-unit} :: (\text{SRDES}, ('s, 't :: \text{trace}, ' \alpha) \text{rsp}) \text{uthy} \Rightarrow ('s, 't, ' \alpha) \text{hrel-rsp}$ **where**
 $\text{srdes-unit } T = II_R$

end

interpretation *srdes-left-unital*: *utp-theory-left-unital SRDES*

by (*unfold-locales, simp-all add: srdes-hcond-def srdes-unit-def SRD-seqr-closure SRD-srdes-skip SRD-left-unit*)

4.10 Recursion laws

lemma *mono-srd-iter*:

assumes $\text{mono } F$ $F \in \llbracket \text{SRD} \rrbracket_H \rightarrow \llbracket \text{SRD} \rrbracket_H$
shows $\text{mono } (\lambda X. \mathbf{R}_s(\text{pre}_R(F X) \vdash \text{peri}_R(F X) \diamond \text{post}_R(F X)))$

```

apply (rule monoI)
apply (rule srdes-tri-refine-intro')
apply (meson assms(1) monoE preR-antitone utp-pred-laws.le-infI2)
apply (meson assms(1) monoE periR-monotone utp-pred-laws.le-infI2)
apply (meson assms(1) monoE postR-monotone utp-pred-laws.le-infI2)
done

```

lemma *mu-srd-SRD*:

```

assumes mono F F ∈  $\llbracket \text{SRD} \rrbracket_H \rightarrow \llbracket \text{SRD} \rrbracket_H$ 
shows  $(\mu X \cdot \mathbf{R}_s(\text{pre}_R(F X) \vdash \text{peri}_R(F X) \diamond \text{post}_R(F X)))$  is SRD
apply (subst gfp-unfold)
apply (simp add: mono-srd-iter assms)
apply (rule RHS-tri-design-is-SRD)
apply (simp-all add: unrest)
done

```

lemma *mu-srd-iter*:

```

assumes mono F F ∈  $\llbracket \text{SRD} \rrbracket_H \rightarrow \llbracket \text{SRD} \rrbracket_H$ 
shows  $(\mu X \cdot \mathbf{R}_s(\text{pre}_R(F(X)) \vdash \text{peri}_R(F(X)) \diamond \text{post}_R(F(X)))) = F(\mu X \cdot \mathbf{R}_s(\text{pre}_R(F(X)) \vdash$ 
 $\text{peri}_R(F(X)) \diamond \text{post}_R(F(X))))$ 
apply (subst gfp-unfold)
apply (simp add: mono-srd-iter assms)
apply (subst SRD-as-reactive-tri-design[THEN sym])
using Healthy-func assms(1) assms(2) mu-srd-SRD apply blast
done

```

lemma *mu-srd-form*:

```

assumes mono F F ∈  $\llbracket \text{SRD} \rrbracket_H \rightarrow \llbracket \text{SRD} \rrbracket_H$ 
shows  $\mu_R F = (\mu X \cdot \mathbf{R}_s(\text{pre}_R(F(X)) \vdash \text{peri}_R(F(X)) \diamond \text{post}_R(F(X))))$ 
proof –
have 1:  $F(\mu X \cdot \mathbf{R}_s(\text{pre}_R(F X) \vdash \text{peri}_R(F X) \diamond \text{post}_R(F X)))$  is SRD
by (simp add: Healthy-apply-closed assms(1) assms(2) mu-srd-SRD)
have 2:  $\text{Mono}_{\text{uthy-order}} \text{SRDES } F$ 
by (simp add: assms(1) mono-Monotone-utp-order)
hence 3:  $\mu_R F = F(\mu_R F)$ 
by (simp add: srdes-theory-continuous.LFP-unfold[THEN sym] assms)
hence  $\mathbf{R}_s(\text{pre}_R(F(F(\mu_R F))) \vdash \text{peri}_R(F(F(\mu_R F))) \diamond \text{post}_R(F(F(\mu_R F)))) = \mu_R F$ 
using SRD-reactive-tri-design by force
hence  $(\mu X \cdot \mathbf{R}_s(\text{pre}_R(F X) \vdash \text{peri}_R(F X) \diamond \text{post}_R(F X))) \sqsubseteq F(\mu_R F)$ 
by (simp add: 2 srdes-theory-continuous.weak.LFP-lemma3 gfp-upperbound assms)
thus ?thesis
using assms 1 3 srdes-theory-continuous.weak.LFP-lowerbound eq-iff mu-srd-iter
by (metis (mono-tags, lifting))
qed

```

lemma *Monotonic-SRD-comp [closure]*: *Monotonic ((;;) P ◦ SRD)*

```

by (simp add: mono-def R1-R2c-is-R2 R2-mono R3h-mono RD1-mono RD2-mono RHS-def SRD-def
seqr-mono)

```

end

5 Normal Reactive Designs

theory *utp-rdes-normal*

imports

$utp-rdes-triples$
 $UTP-KAT.utp-kleene$
begin

This additional healthiness condition is analogous to H3

definition $RD3$ **where**

$[upred-defs]: RD3(P) = P ;; II_R$

lemma $RD3-idem: RD3(RD3(P)) = RD3(P)$

proof –

have $a: II_R ;; II_R = II_R$
by ($simp$ $add: SRD-left-unit SRD-srdes-skip$)
show $?thesis$
by ($simp$ $add: RD3-def seqr-assoc a$)

qed

lemma $RD3-Idempotent [closure]: Idempotent RD3$

by ($simp$ $add: Idempotent-def RD3-idem$)

lemma $RD3-continuous: RD3(\sqcap A) = (\sqcap P \in A. RD3(P))$

by ($simp$ $add: RD3-def seq-Sup-distr$)

lemma $RD3-Continuous [closure]: Continuous RD3$

by ($simp$ $add: Continuous-def RD3-continuous$)

lemma $RD3-right-subsumes-RD2: RD2(RD3(P)) = RD3(P)$

proof –

have $a: II_R ;; J = II_R$
by ($rel-auto$)
show $?thesis$
by ($metis$ ($no-types$, $hide-lams$) $H2-def RD2-def RD3-def a seqr-assoc$)

qed

lemma $RD3-left-subsumes-RD2: RD3(RD2(P)) = RD3(P)$

proof –

have $a: J ;; II_R = II_R$
by ($rel-simp$, $safe$, $blast+$)
show $?thesis$
by ($metis$ ($no-types$, $hide-lams$) $H2-def RD2-def RD3-def a seqr-assoc$)

qed

lemma $RD3-implies-RD2: P \text{ is } RD3 \implies P \text{ is } RD2$

by ($metis$ $Healthy-def RD3-right-subsumes-RD2$)

lemma $RD3-intro-pre:$

assumes $P \text{ is } SRD (\neg_r pre_R(P)) ;; true_r = (\neg_r pre_R(P)) \$st' \# peri_R(P)$
shows $P \text{ is } RD3$

proof –

have $RD3(P) = \mathbf{R}_s ((\neg_r pre_R P) wp_r false \vdash (\exists \$st' \cdot peri_R P) \diamond post_R P)$
by ($simp$ $add: RD3-def SRD-right-unit-tri-lemma assms$)
also have $\dots = \mathbf{R}_s ((\neg_r pre_R P) wp_r false \vdash peri_R P \diamond post_R P)$
by ($simp$ $add: assms(3) ex-unrest$)
also have $\dots = \mathbf{R}_s ((\neg_r pre_R P) wp_r false \vdash cmt_R P)$
by ($simp$ $add: wait'-cond-peri-post-cmt$)
also have $\dots = \mathbf{R}_s (pre_R P \vdash cmt_R P)$

by (simp add: assms(2) rpred wp-rea-def R1-preR)
 finally show ?thesis
 by (metis Healthy-def SRD-as-reactive-design assms(1))
 qed

lemma *RHS-tri-design-right-unit-lemma:*

assumes $\$ok' \# P \$ok' \# Q \$ok' \# R \$wait' \# R$
 shows $\mathbf{R}_s(P \vdash Q \diamond R) ;; II_R = \mathbf{R}_s((\neg_r (\neg_r P) ;; true_r) \vdash ((\exists \$st' \cdot Q) \diamond R))$
proof –
 have $\mathbf{R}_s(P \vdash Q \diamond R) ;; II_R = \mathbf{R}_s(P \vdash Q \diamond R) ;; \mathbf{R}_s(true \vdash false \diamond (\$tr' =_u \$tr \wedge [II]_R))$
 by (simp add: srdes-skip-tri-design, rel-auto)
 also have $\dots = \mathbf{R}_s((\neg R1 (\neg R2s P) ;; R1 true) \vdash (\exists \$st' \cdot Q) \diamond (R1 (R2s R) ;; R1 (R2s (\$tr' =_u \$tr \wedge [II]_R))))$
 by (simp-all add: RHS-tri-design-composition assms unrest R2s-true R1-false R2s-false)
 also have $\dots = \mathbf{R}_s((\neg R1 (\neg R2s P) ;; R1 true) \vdash (\exists \$st' \cdot Q) \diamond R1 (R2s R))$
proof –
 from assms(3,4) have $(R1 (R2s R) ;; R1 (R2s (\$tr' =_u \$tr \wedge [II]_R))) = R1 (R2s R)$
 by (rel-auto, metis (no-types, lifting) minus-zero-eq, meson order-refl trace-class.diff-cancel)
 thus ?thesis
 by simp
 qed
 also have $\dots = \mathbf{R}_s((\neg (\neg P) ;; R1 true) \vdash ((\exists \$st' \cdot Q) \diamond R))$
 by (metis (no-types, lifting) R1-R2s-R1-true-lemma R1-R2s-R2c R2c-not RHS-design-R2c-pre RHS-design-neg-R1-pre RHS-design-post-R1 RHS-design-post-R2s)
 also have $\dots = \mathbf{R}_s((\neg_r (\neg_r P) ;; true_r) \vdash ((\exists \$st' \cdot Q) \diamond R))$
 by (rel-auto)
 finally show ?thesis .
 qed

lemma *RHS-tri-design-RD3-intro:*

assumes
 $\$ok' \# P \$ok' \# Q \$ok' \# R \$st' \# Q \$wait' \# R$
 $P \text{ is } R1 (\neg_r P) ;; true_r = (\neg_r P)$
 shows $\mathbf{R}_s(P \vdash Q \diamond R) \text{ is } RD3$
apply (simp add: Healthy-def RD3-def)
apply (subst RHS-tri-design-right-unit-lemma)
apply (simp-all add: assms ex-unrest rpred)
 done

RD3 reactive designs are those whose assumption can be written as a conjunction of a precondition on (undashed) program variables, and a negated statement about the trace. The latter allows us to state that certain events must not occur in the trace – which are effectively safety properties.

lemma *R1-right-unit-lemma:*

$\llbracket out\alpha \# b; out\alpha \# e \rrbracket \implies (\neg_r b \vee \$tr \hat{^}_u e \leq_u \$tr') ;; R1(true) = (\neg_r b \vee \$tr \hat{^}_u e \leq_u \$tr')$
by (rel-auto, blast, metis (no-types, lifting) dual-order.trans)

lemma *RHS-tri-design-RD3-intro-form:*

assumes
 $out\alpha \# b \text{ out}\alpha \# e \$ok' \# Q \$st' \# Q \$ok' \# R \$wait' \# R$
 shows $\mathbf{R}_s((b \wedge \neg_r \$tr \hat{^}_u e \leq_u \$tr') \vdash Q \diamond R) \text{ is } RD3$
apply (rule RHS-tri-design-RD3-intro)
apply (simp-all add: assms unrest closure rpred)
apply (subst R1-right-unit-lemma)
apply (simp-all add: assms unrest)

done

definition $NSRD :: ('s, 't :: trace, 'α) hrel-rsp \Rightarrow ('s, 't, 'α) hrel-rsp$
where $[upred-defs]: NSRD = RD1 \circ RD3 \circ RHS$

lemma $RD1\text{-}RD3\text{-}commute$: $RD1(RD3(P)) = RD3(RD1(P))$
by $(rel\text{-}auto, blast+)$

lemma $NSRD\text{-}is\text{-}SRD$ $[closure]$: $P \text{ is } NSRD \implies P \text{ is } SRD$
by $(simp \text{ add: } Healthy\text{-}def \ NSRD\text{-}def \ SRD\text{-}def, \text{metis } Healthy\text{-}def \ RD1\text{-}RD3\text{-}commute \ RD2\text{-}RHS\text{-}commute \ RD3\text{-}def \ RD3\text{-}right\text{-}subsumes\text{-}RD2 \ SRD\text{-}def \ SRD\text{-}idem \ SRD\text{-}segr\text{-}closure \ SRD\text{-}srdes\text{-}skip)$

lemma $NSRD\text{-}elim$ $[RD\text{-}elim]$:
 $\llbracket P \text{ is } NSRD; Q(\mathbf{R}_s(pre_R(P) \vdash peri_R(P) \diamond post_R(P))) \rrbracket \implies Q(P)$
by $(simp \text{ add: } RD\text{-}elim \ closure)$

lemma $NSRD\text{-}Idempotent$ $[closure]$: $Idempotent \ NSRD$
by $(clarsimp \ simp \text{ add: } Idempotent\text{-}def \ NSRD\text{-}def, \text{metis } (no\text{-}types, \text{hide-lams}) \ Healthy\text{-}def \ RD1\text{-}RD3\text{-}commute \ RD3\text{-}def \ RD3\text{-}idem \ RD3\text{-}left\text{-}subsumes\text{-}RD2 \ SRD\text{-}def \ SRD\text{-}idem \ SRD\text{-}segr\text{-}closure \ SRD\text{-}srdes\text{-}skip)$

lemma $NSRD\text{-}Continuous$ $[closure]$: $Continuous \ NSRD$
by $(simp \text{ add: } Continuous\text{-}comp \ NSRD\text{-}def \ RD1\text{-}Continuous \ RD3\text{-}Continuous \ RHS\text{-}Continuous)$

lemma $NSRD\text{-}form$:
 $NSRD(P) = \mathbf{R}_s((\neg_r (\neg_r pre_R(P)) ;; R1 \text{ true}) \vdash ((\exists \$st' \cdot peri_R(P)) \diamond post_R(P)))$
proof –
have $NSRD(P) = RD3(SRD(P))$
by $(metis (no\text{-}types, \text{lifting}) \ NSRD\text{-}def \ RD1\text{-}RD3\text{-}commute \ RD3\text{-}left\text{-}subsumes\text{-}RD2 \ SRD\text{-}def \ comp\text{-}def)$
also have $\dots = RD3(\mathbf{R}_s(pre_R(P) \vdash peri_R(P) \diamond post_R(P)))$
by $(simp \text{ add: } SRD\text{-}as\text{-}reactive\text{-}tri\text{-}design)$
also have $\dots = \mathbf{R}_s(pre_R(P) \vdash peri_R(P) \diamond post_R(P)) ;; II_R$
by $(simp \text{ add: } RD3\text{-}def)$
also have $\dots = \mathbf{R}_s((\neg_r (\neg_r pre_R(P)) ;; R1 \text{ true}) \vdash ((\exists \$st' \cdot peri_R(P)) \diamond post_R(P)))$
by $(simp \text{ add: } RHS\text{-}tri\text{-}design\text{-}right\text{-}unit\text{-}lemma \ unrest)$
finally show $?thesis$.

qed

lemma $NSRD\text{-}healthy\text{-}form$:
assumes $P \text{ is } NSRD$
shows $\mathbf{R}_s((\neg_r (\neg_r pre_R(P)) ;; R1 \text{ true}) \vdash ((\exists \$st' \cdot peri_R(P)) \diamond post_R(P))) = P$
by $(metis \ Healthy\text{-}def \ NSRD\text{-}form \ assms)$

lemma $NSRD\text{-}Sup\text{-}closure$ $[closure]$:
assumes $A \subseteq \llbracket NSRD \rrbracket_H \ A \neq \{\}$
shows $\sqcap A \text{ is } NSRD$

proof –
have $NSRD (\sqcap A) = (\sqcap (NSRD `A))$
by $(simp \text{ add: } ContinuousD \ NSRD\text{-}Continuous \ assms(2))$
also have $\dots = (\sqcap A)$
by $(simp \text{ only: } Healthy\text{-}carrier\text{-}image \ assms)$
finally show $?thesis$ **by** $(simp \text{ add: } Healthy\text{-}def)$

qed

lemma $intChoice\text{-}NSRD\text{-}closed$ $[closure]$:
assumes $P \text{ is } NSRD \ Q \text{ is } NSRD$

shows $P \sqcap Q$ *is NSRD*
using *NSRD-Sup-closure*[of $\{P, Q\}$] **by** (*simp add: assms*)

lemma *NSRD-SUP-closure* [*closure*]:
 $\llbracket \bigwedge i. i \in A \implies P(i) \text{ is NSRD}; A \neq \{\} \rrbracket \implies (\bigwedge i \in A. P(i)) \text{ is NSRD}$
by (*rule NSRD-Sup-closure, auto*)

lemma *NSRD-neg-pre-unit*:
assumes P *is NSRD*
shows $(\neg_r \text{pre}_R(P)) ;; \text{true}_r = (\neg_r \text{pre}_R(P))$
proof –
have $(\neg_r \text{pre}_R(P)) = (\neg_r \text{pre}_R(\mathbf{R}_s((\neg_r (\neg_r \text{pre}_R(P)) ;; R1 \text{true}) \vdash ((\exists \$st' \cdot \text{peri}_R(P)) \diamond \text{post}_R(P))))$
by (*simp add: NSRD-healthy-form assms*)
also have $\dots = R1 (R2c ((\neg_r \text{pre}_R P) ;; R1 \text{true}))$
by (*simp add: rea-pre-RHS-design R1-negate-R1 R1-idem R1-rea-not' R2c-rea-not usubst rpred unrest closure*)
also have $\dots = (\neg_r \text{pre}_R P) ;; R1 \text{true}$
by (*simp add: R1-R2c-seqr-distribute closure assms*)
finally show *?thesis*
by (*simp add: rea-not-def*)
qed

lemma *NSRD-neg-pre-left-zero*:
assumes P *is NSRD* Q *is R1* Q *is RD1*
shows $(\neg_r \text{pre}_R(P)) ;; Q = (\neg_r \text{pre}_R(P))$
by (*metis (no-types, hide-lams) NSRD-neg-pre-unit RD1-left-zero assms(1) assms(2) assms(3) seqr-assoc*)

lemma *NSRD-st'-unrest-peri* [*unrest*]:
assumes P *is NSRD*
shows $\$st' \# \text{peri}_R(P)$
proof –
have $\text{peri}_R(P) = \text{peri}_R(\mathbf{R}_s((\neg_r (\neg_r \text{pre}_R(P)) ;; R1 \text{true}) \vdash ((\exists \$st' \cdot \text{peri}_R(P)) \diamond \text{post}_R(P))))$
by (*simp add: NSRD-healthy-form assms*)
also have $\dots = R1 (R2c (\neg_r (\neg_r \text{pre}_R P) ;; R1 \text{true} \Rightarrow_r (\exists \$st' \cdot \text{peri}_R P)))$
by (*simp add: rea-peri-RHS-design usubst unrest*)
also have $\$st' \# \dots$
by (*simp add: R1-def R2c-def unrest*)
finally show *?thesis* .
qed

lemma *NSRD-wait'-unrest-pre* [*unrest*]:
assumes P *is NSRD*
shows $\$wait' \# \text{pre}_R(P)$
proof –
have $\text{pre}_R(P) = \text{pre}_R(\mathbf{R}_s((\neg_r (\neg_r \text{pre}_R(P)) ;; R1 \text{true}) \vdash ((\exists \$st' \cdot \text{peri}_R(P)) \diamond \text{post}_R(P))))$
by (*simp add: NSRD-healthy-form assms*)
also have $\dots = (R1 (R2c (\neg_r (\neg_r \text{pre}_R P) ;; R1 \text{true})))$
by (*simp add: rea-pre-RHS-design usubst unrest*)
also have $\$wait' \# \dots$
by (*simp add: R1-def R2c-def unrest*)
finally show *?thesis* .
qed

lemma *NSRD-st'-unrest-pre* [*unrest*]:
assumes P *is NSRD*

shows $\$st' \# pre_R(P)$
proof –
 have $pre_R(P) = pre_R(\mathbf{R}_s((\neg_r (\neg_r pre_R(P)) ;; R1 \text{ true}) \vdash ((\exists \$st' \cdot peri_R(P)) \diamond post_R(P))))$
 by (*simp add: NSRD-healthy-form assms*)
 also have $\dots = R1 (R2c (\neg_r (\neg_r pre_R P) ;; R1 \text{ true}))$
 by (*simp add: rea-pre-RHS-design usubst unrest*)
 also have $\$st' \# \dots$
 by (*simp add: R1-def R2c-def unrest*)
 finally show *?thesis* .
qed

lemma *NSRD-peri-under-pre* [*rpred*]:
 $P \text{ is NSRD} \implies (pre_R P \Rightarrow_r peri_R P) = peri_R P$
 by (*simp add: SRD-peri-under-pre unrest closure*)

lemma *NSRD-post-under-pre* [*rpred*]:
 $P \text{ is NSRD} \implies (pre_R P \Rightarrow_r post_R P) = post_R P$
 by (*simp add: SRD-post-under-pre unrest closure*)

lemma *NSRD-peri-seq-under-pre*:
 assumes $P \text{ is NSRD } Q \text{ is NSRD}$
 shows $(pre_R P \Rightarrow_r peri_R P \vee post_R P ;; peri_R Q) = (peri_R P \vee post_R P ;; peri_R Q)$
 by (*metis NSRD-peri-under-pre assms(1) rea-impl-def utp-pred-laws.disj-assoc*)

lemma *NSRD-postR-seq-periR-impl*:
 assumes $P \text{ is NSRD } Q \text{ is NSRD}$
 shows $(post_R P wp_r pre_R Q \Rightarrow_r (post_R P ;; peri_R Q)) = (post_R P ;; peri_R Q)$
 by (*metis NSRD-is-SRD NSRD-peri-under-pre assms postR-RR wpR-impl-post-spec*)

lemma *NSRD-postR-seq-postR-impl*:
 assumes $P \text{ is NSRD } Q \text{ is NSRD}$
 shows $(post_R P wp_r pre_R Q \Rightarrow_r (post_R P ;; post_R Q)) = (post_R P ;; post_R Q)$
 by (*metis NSRD-is-SRD NSRD-post-under-pre assms postR-RR wpR-impl-post-spec*)

lemma *NSRD-peri-under-assms*:
 assumes $P \text{ is NSRD } Q \text{ is NSRD}$
 shows $(pre_R P \wedge post_R P wp_r pre_R Q \Rightarrow_r peri_R P \vee post_R P ;; peri_R Q) = (peri_R P \vee post_R P ;; peri_R Q)$
 by (*metis (no-types, lifting) NSRD-peri-seq-under-pre assms NSRD-postR-seq-periR-impl rea-impl-conj rea-impl-disj*)

lemma *NSRD-peri-under-assms'*:
 assumes $P \text{ is NSRD } Q \text{ is NSRD}$
 shows $(post_R P wp_r pre_R Q \Rightarrow_r peri_R P \vee post_R P ;; peri_R Q) = (peri_R P \vee post_R P ;; peri_R Q)$
 by (*simp add: NSRD-postR-seq-periR-impl assms rea-impl-disj*)

lemma *NSRD-post-under-assms*:
 assumes $P \text{ is NSRD } Q \text{ is NSRD}$
 shows $(pre_R P \wedge post_R P wp_r pre_R Q \Rightarrow_r post_R P ;; post_R Q) = (pre_R P \Rightarrow_r (post_R P ;; post_R Q))$
 by (*metis NSRD-postR-seq-postR-impl assms(1) assms(2) rea-impl-conj*)

lemma *NSRD-alt-def*: $NSRD(P) = RD3(SRD(P))$
 by (*metis NSRD-def RD1-RD3-commute RD3-left-subsumes-RD2 SRD-def comp-eq-dest-lhs*)

lemma *preR-RR* [*closure*]: $P \text{ is NSRD} \implies pre_R(P) \text{ is RR}$

by (rule *RR-intro*, *simp-all add: closure unrest*)

lemma *NSRD-neg-pre-RC* [closure]:
 assumes *P is NSRD*
 shows *pre_R(P) is RC*
 by (rule *RC-intro*, *simp-all add: closure assms NSRD-neg-pre-unit rpred*)

lemma *NSRD-intro*:
 assumes *P is SRD* ($\neg_r \text{pre}_R(P)$) ;; *true_r* = ($\neg_r \text{pre}_R(P)$) $\$st' \# \text{peri}_R(P)$
 shows *P is NSRD*
proof –
 have *NSRD(P) = R_s(($\neg_r (\neg_r \text{pre}_R(P))$) ;; *R1 true*)* $\vdash ((\exists \$st' \cdot \text{peri}_R(P)) \diamond \text{post}_R(P))$
 by (*simp add: NSRD-form*)
 also have ... = *R_s(pre_R P $\vdash \text{peri}_R P \diamond \text{post}_R P$)*
 by (*simp add: assms ex-unrest rpred closure*)
 also have ... = *P*
 by (*simp add: SRD-reactive-tri-design assms(1)*)
 finally show ?thesis
 using *Healthy-def* by blast
qed

lemma *NSRD-intro'*:
 assumes *P is R2 P is R3h P is RD1 P is RD3*
 shows *P is NSRD*
 by (*metis (no-types, hide-lams) Healthy-def NSRD-def R1-R2c-is-R2 RHS-def assms comp-apply*)

lemma *NSRD-RC-intro*:
 assumes *P is SRD pre_R(P) is RC* $\$st' \# \text{peri}_R(P)$
 shows *P is NSRD*
 by (*metis Healthy-def NSRD-form SRD-reactive-tri-design assms(1) assms(2) assms(3) ex-unrest rea-not-false wp-rea-RC-false wp-rea-def*)

lemma *NSRD-rdes-intro* [closure]:
 assumes *P is RC Q is RR R is RR* $\$st' \# Q$
 shows *R_s(P $\vdash Q \diamond R$) is NSRD*
 by (rule *NSRD-RC-intro*, *simp-all add: rdes closure assms unrest*)

lemma *SRD-RD3-implies-NSRD*:
 $\llbracket P \text{ is SRD}; P \text{ is RD3} \rrbracket \implies P \text{ is NSRD}$
 by (*metis (no-types, lifting) Healthy-def NSRD-def RHS-idem SRD-healths(4) SRD-reactive-design comp-apply*)

lemma *NSRD-iff*:
 $P \text{ is NSRD} \longleftrightarrow ((P \text{ is SRD}) \wedge (\neg_r \text{pre}_R(P)) \text{ ;; } R1(\text{true}) = (\neg_r \text{pre}_R(P)) \wedge (\$st' \# \text{peri}_R(P)))$
 by (*meson NSRD-intro NSRD-is-SRD NSRD-neg-pre-unit NSRD-st'-unrest-peri*)

lemma *NSRD-is-RD3* [closure]:
 assumes *P is NSRD*
 shows *P is RD3*
 by (*simp add: NSRD-is-SRD NSRD-neg-pre-unit NSRD-st'-unrest-peri RD3-intro-pre assms*)

lemma *NSRD-refine-elim*:
 assumes
 $P \sqsubseteq Q$ *P is NSRD Q is NSRD*

$\llbracket \text{'pre}_R(P) \Rightarrow \text{pre}_R(Q) \text{'}; \text{'pre}_R(P) \wedge \text{peri}_R(Q) \Rightarrow \text{peri}_R(P) \text{'}; \text{'pre}_R(P) \wedge \text{post}_R(Q) \Rightarrow \text{post}_R(P) \text{' } \rrbracket$
 $\Rightarrow R$

shows R

proof –

have $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P)) \sqsubseteq \mathbf{R}_s(\text{pre}_R(Q) \vdash \text{peri}_R(Q) \diamond \text{post}_R(Q))$

by (simp add: NSRD-is-SRD SRD-reactive-tri-design assms(1) assms(2) assms(3))

hence 1: $\text{'pre}_R P \Rightarrow \text{pre}_R Q \text{'}$ and 2: $\text{'pre}_R P \wedge \text{peri}_R Q \Rightarrow \text{peri}_R P \text{'}$ and 3: $\text{'pre}_R P \wedge \text{post}_R Q \Rightarrow \text{post}_R P \text{'}$

by (simp-all add: RHS-tri-design-refine assms closure)

with assms(4) show ?thesis

by simp

qed

lemma NSRD-right-unit: P is NSRD $\Rightarrow P ;; II_R = P$

by (metis Healthy-if NSRD-is-RD3 RD3-def)

lemma NSRD-composition-wp:

assumes P is NSRD Q is SRD

shows $P ;; Q =$

$\mathbf{R}_s((\text{pre}_R P \wedge \text{post}_R P \text{ wp}_r \text{pre}_R Q) \vdash (\text{peri}_R P \vee (\text{post}_R P ;; \text{peri}_R Q)) \diamond (\text{post}_R P ;; \text{post}_R Q))$

by (simp add: SRD-composition-wp assms NSRD-is-SRD wp-rea-def NSRD-neg-pre-unit NSRD-st'-unrest-peri R1-negate-R1 R1-preR ex-unrest rpred)

lemma preR-NSRD-seq-lemma:

assumes P is NSRD Q is SRD

shows $R1 (R2c (\text{post}_R P ;; (\neg_r \text{pre}_R Q))) = \text{post}_R P ;; (\neg_r \text{pre}_R Q)$

proof –

have $\text{post}_R P ;; (\neg_r \text{pre}_R Q) = R1(R2c(\text{post}_R P)) ;; R1(R2c(\neg_r \text{pre}_R Q))$

by (simp add: NSRD-is-SRD R1-R2c-post-RHS R1-rea-not R2c-preR R2c-rea-not assms(1) assms(2))

also have $\dots = R1 (R2c (\text{post}_R P ;; (\neg_r \text{pre}_R Q)))$

by (simp add: R1-seqr R2c-R1-seq calculation)

finally show ?thesis ..

qed

lemma preR-NSRD-seq [rdes]:

assumes P is NSRD Q is SRD

shows $\text{pre}_R(P ;; Q) = (\text{pre}_R P \wedge \text{post}_R P \text{ wp}_r \text{pre}_R Q)$

by (simp add: NSRD-composition-wp assms rea-pre-RHS-design usubst unrest wp-rea-def R2c-disj R1-disj R2c-and R2c-preR R1-R2c-commute[THEN sym] R1-extend-conj' R1-idem R2c-not closure)

(metis (no-types, lifting) Healthy-def Healthy-if NSRD-is-SRD R1-R2c-commute

R1-R2c-seqr-distribute R1-seqr-closure assms(1) assms(2) postR-R2c-closed postR-SRD-R1

preR-R2c-closed rea-not-R1 rea-not-R2c)

lemma periR-NSRD-seq [rdes]:

assumes P is NSRD Q is NSRD

shows $\text{peri}_R(P ;; Q) = ((\text{pre}_R P \wedge \text{post}_R P \text{ wp}_r \text{pre}_R Q) \Rightarrow_r (\text{peri}_R P \vee (\text{post}_R P ;; \text{peri}_R Q)))$

by (simp add: NSRD-composition-wp assms closure rea-peri-RHS-design usubst unrest wp-rea-def

R1-extend-conj' R1-disj R1-R2c-seqr-distribute R2c-disj R2c-and R2c-rea-impl R1-rea-impl'

R2c-preR R2c-periR R1-rea-not' R2c-rea-not R1-peri-SRD)

lemma postR-NSRD-seq [rdes]:

assumes P is NSRD Q is NSRD

shows $\text{post}_R(P ;; Q) = ((\text{pre}_R P \wedge \text{post}_R P \text{ wp}_r \text{pre}_R Q) \Rightarrow_r (\text{post}_R P ;; \text{post}_R Q))$

by (simp add: NSRD-composition-wp assms closure rea-post-RHS-design usubst unrest wp-rea-def

*R1-extend-conj' R1-disj R1-R2c-seqr-distribute R2c-disj R2c-and R2c-rea-impl R1-rea-impl'
R2c-preR R2c-periR R1-rea-not' R2c-rea-not)*

lemma *NSRD-seqr-closure* [closure]:

assumes *P is NSRD Q is NSRD*

shows *(P ;; Q) is NSRD*

proof –

have $(\neg_r \text{post}_R P \text{wp}_r \text{pre}_R Q) ;; \text{true}_r = (\neg_r \text{post}_R P \text{wp}_r \text{pre}_R Q)$

by (*simp add: wp-rea-def rpred assms closure seqr-assoc NSRD-neg-pre-unit*)

moreover have $\$st' \# \text{pre}_R P \wedge \text{post}_R P \text{wp}_r \text{pre}_R Q \Rightarrow_r \text{peri}_R P \vee \text{post}_R P ;; \text{peri}_R Q$

by (*simp add: unrest assms wp-rea-def*)

ultimately show *?thesis*

by (*rule-tac NSRD-intro, simp-all add: seqr-or-distl NSRD-neg-pre-unit assms closure rdes unrest*)

qed

lemma *RHS-tri-normal-design-composition:*

assumes

$\$ok' \# P \$ok' \# Q_1 \$ok' \# Q_2 \$ok \# R \$ok \# S_1 \$ok \# S_2$

$\$wait \# R \$wait' \# Q_2 \$wait \# S_1 \$wait \# S_2$

P is R2c Q₁ is R1 Q₁ is R2c Q₂ is R1 Q₂ is R2c

R is R2c S₁ is R1 S₁ is R2c S₂ is R1 S₂ is R2c

R1 ($\neg P$) ;; R1(true) = R1($\neg P$) \$st' # Q₁

shows $\mathbf{R}_s(P \vdash Q_1 \diamond Q_2) ;; \mathbf{R}_s(R \vdash S_1 \diamond S_2)$

$= \mathbf{R}_s((P \wedge Q_2 \text{wp}_r R) \vdash (Q_1 \vee (Q_2 ;; S_1)) \diamond (Q_2 ;; S_2))$

proof –

have $\mathbf{R}_s(P \vdash Q_1 \diamond Q_2) ;; \mathbf{R}_s(R \vdash S_1 \diamond S_2) =$

$\mathbf{R}_s((R1(\neg P) \text{wp}_r \text{false} \wedge Q_2 \text{wp}_r R) \vdash ((\exists \$st' \cdot Q_1) \sqcap (Q_2 ;; S_1)) \diamond (Q_2 ;; S_2))$

by (*simp-all add: RHS-tri-design-composition-wp rea-not-def assms unrest*)

also have $\dots = \mathbf{R}_s((P \wedge Q_2 \text{wp}_r R) \vdash (Q_1 \vee (Q_2 ;; S_1)) \diamond (Q_2 ;; S_2))$

by (*simp add: assms wp-rea-def ex-unrest, rel-auto*)

finally show *?thesis* .

qed

lemma *RHS-tri-normal-design-composition'* [rdes-def]:

assumes *P is RC Q₁ is RR \$st' # Q₁ Q₂ is RR R is RR S₁ is RR S₂ is RR*

shows $\mathbf{R}_s(P \vdash Q_1 \diamond Q_2) ;; \mathbf{R}_s(R \vdash S_1 \diamond S_2)$

$= \mathbf{R}_s((P \wedge Q_2 \text{wp}_r R) \vdash (Q_1 \vee (Q_2 ;; S_1)) \diamond (Q_2 ;; S_2))$

proof –

have *R1 ($\neg P$) ;; R1 true = R1($\neg P$)*

using *RC-implies-RC1[OF assms(1)]*

by (*simp add: Healthy-def RC1-def rea-not-def*)

(*metis R1-negate-R1 R1-seqr utp-pred-laws.double-compl*)

thus *?thesis*

by (*simp add: RHS-tri-normal-design-composition assms closure unrest RR-implies-R2c*)

qed

If a normal reactive design has postcondition false, then it is a left zero for sequential composition.

lemma *NSRD-seq-post-false:*

assumes *P is NSRD Q is SRD post_R(P) = false*

shows *P ;; Q = P*

apply (*simp add: NSRD-composition-wp assms wp rpred closure*)

using *NSRD-is-SRD SRD-reactive-tri-design assms(1,3)* **apply** *fastforce*

done

lemma *NSRD-srd-skip* [closure]: *II_R is NSRD*
 by (rule *NSRD-intro*, simp-all add: rdes closure unrest)

lemma *NSRD-Chaos* [closure]: *Chaos is NSRD*
 by (rule *NSRD-intro*, simp-all add: closure rdes unrest)

lemma *NSRD-Miracle* [closure]: *Miracle is NSRD*
 by (rule *NSRD-intro*, simp-all add: closure rdes unrest)

Post-composing a miracle filters out the non-terminating behaviours

lemma *NSRD-right-Miracle-tri-lemma*:
 assumes *P is NSRD*
 shows *P ;; Miracle = R_s (pre_R P ⊢ peri_R P ◊ false)*
 by (simp add: *NSRD-composition-wp closure assms rdes wp rpred*)

The set of non-terminating behaviours is a subset

lemma *NSRD-right-Miracle-refines*:
 assumes *P is NSRD*
 shows *P ⊆ P ;; Miracle*
proof –
 have *R_s (pre_R P ⊢ peri_R P ◊ post_R P) ⊆ R_s (pre_R P ⊢ peri_R P ◊ false)*
 by (rule *srdes-tri-refine-intro*, rel-auto+)
 thus ?thesis
 by (simp add: *NSRD-elim NSRD-right-Miracle-tri-lemma assms*)
qed

lemma *upower-Suc-NSRD-closed* [closure]:
P is NSRD ⇒ P ^ Suc n is NSRD
proof (induct n)
 case 0
 then show ?case
 by (simp)
next
 case (Suc n)
 then show ?case
 by (simp add: *NSRD-seqr-closure upred-semiring.power-Suc*)
qed

lemma *NSRD-power-Suc* [closure]:
P is NSRD ⇒ P ;; P ^ n is NSRD
 by (metis *upower-Suc-NSRD-closed upred-semiring.power-Suc*)

lemma *uplus-NSRD-closed* [closure]: *P is NSRD ⇒ P⁺ is NSRD*
 by (simp add: *uplus-power-def closure*)

lemma *preR-power*:
 assumes *P is NSRD*
 shows *pre_R(P ;; Pⁿ) = (⊔_{i∈{0..n}}. (post_R(P) ^ i) wp_r (pre_R(P)))*
proof (induct n)
 case 0
 then show ?case
 by (simp add: *wp closure*)
next
 case (Suc n) **note** *hyp = this*
 have *pre_R (P ^ (Suc n + 1)) = pre_R (P ;; P ^ (n+1))*


```

  by (simp add: upred-semiring.power-Suc)
also have ... = (preR P ∧ postR P wpr preR (P ^ (Suc n)))
  using NSRD-iff assms preR-NSRD-seq upower-Suc-NSRD-closed by fastforce
also have ... = (preR P ∧ postR P wpr (⋂ i∈{0..n}. postR P ^ i wpr preR P))
  by (simp add: hyp upred-semiring.power-Suc)
also have ... = (preR P ∧ (⋂ i∈{0..n}. postR P wpr (postR P ^ i wpr preR P)))
  by (simp add: wp)
also have ... = (preR P ∧ (⋂ i∈{0..n}. (postR P ^ (i+1) wpr preR P)))
proof -
  have ∧ i. R1 (postR P ^ i ;; (¬r preR P)) = (postR P ^ i ;; (¬r preR P))
    by (induct-tac i, simp-all add: closure Healthy-if assms)
  thus ?thesis
    by (simp add: wp-rea-def upred-semiring.power-Suc seqr-assoc rpred closure assms)
qed
also have ... = (postR P ^ 0 wpr preR P ∧ (⋂ i∈{0..n}. (postR P ^ (i+1) wpr preR P)))
  by (simp add: wp assms closure)
also have ... = (postR P ^ 0 wpr preR P ∧ (⋂ i∈{1..Suc n}. (postR P ^ i wpr preR P)))
proof -
  have (⋂ i∈{0..n}. (postR P ^ (i+1) wpr preR P)) = (⋂ i∈{1..Suc n}. (postR P ^ i wpr preR P))
    by (rule cong[of Inf], simp-all add: fun-eq-iff)
    (metis (no-types, lifting) image-Suc-atLeastAtMost image-cong image-image)
  thus ?thesis by simp
qed
also have ... = (⋂ i∈insert 0 {1..Suc n}. (postR P ^ i wpr preR P))
  by (simp add: conj-upred-def)
also have ... = (⋂ i∈{0..Suc n}. postR P ^ i wpr preR P)
  by (simp add: atLeast0-atMost-Suc-eq-insert-0)
finally show ?case by (simp add: upred-semiring.power-Suc)
qed

lemma preR-power' [rdes]:
  assumes P is NSRD
  shows preR(P ;; P^n) = (⋂ i∈{0..n}. (postR(P) ^ i) wpr (preR(P)))
  by (simp add: preR-power assms USUP-as-Inf[THEN sym])

lemma preR-power-Suc [rdes]:
  assumes P is NSRD
  shows preR(P^(Suc n)) = (⋂ i∈{0..n}. (postR(P) ^ i) wpr (preR(P)))
  by (simp add: upred-semiring.power-Suc rdes assms)

declare upred-semiring.power-Suc [simp]

lemma periR-power:
  assumes P is NSRD
  shows periR(P ;; P^n) = (preR(P^(Suc n)) ⇒r (⋂ i∈{0..n}. postR(P) ^ i ;; periR(P)))
proof (induct n)
  case 0
  then show ?case
    by (simp add: NSRD-is-SRD NSRD-wait'-unrest-pre SRD-peri-under-pre assms)
next
  case (Suc n) note hyp = this
  have periR (P ^ (Suc n + 1)) = periR (P ;; P ^ (n+1))
    by (simp)
  also have ... = (preR(P ^ (Suc n + 1)) ⇒r (periR P ∨ postR P ;; periR (P ;; P ^ n)))
    by (simp add: closure assms rdes)

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also have ... = (preR(P ^ (Suc n + 1)) ⇒r (periR P ∨ postR P ;; (preR (P ^ (Suc n)) ⇒r (⋀ i ∈ {0..n}.
postR P ^ i) ;; periR P)))
  by (simp only: hyp)
also
have ... = (preR P ⇒r periR P ∨ (postR P wpr preR (P ;; P ^ n) ⇒r postR P ;; (preR (P ;; P ^ n)
⇒r (⋀ i ∈ {0..n}. postR P ^ i) ;; periR P)))
  by (simp add: rdes closure assms, rel-blast)
also
have ... = (preR P ⇒r periR P ∨ (postR P wpr preR (P ;; P ^ n) ⇒r postR P ;; (⋀ i ∈ {0..n}. postR
P ^ i) ;; periR P)))
proof –
  have (⋀ i ∈ {0..n}. postR P ^ i) is R1
  by (simp add: NSRD-is-SRD R1-Continuous R1-power Sup-Continuous-closed assms postR-SRD-R1)
  hence 1:(⋀ i ∈ {0..n}. postR P ^ i) ;; periR P) is R1
  by (simp add: closure assms)
  hence (preR (P ;; P ^ n) ⇒r (⋀ i ∈ {0..n}. postR P ^ i) ;; periR P) is R1
  by (simp add: closure)
  hence (postR P wpr preR (P ;; P ^ n) ⇒r postR P ;; (preR (P ;; P ^ n) ⇒r (⋀ i ∈ {0..n}. postR P
^ i) ;; periR P))
    = (postR P wpr preR (P ;; P ^ n) ⇒r R1(postR P) ;; R1(preR (P ;; P ^ n) ⇒r (⋀ i ∈ {0..n}.
postR P ^ i) ;; periR P))
  by (simp add: Healthy-if R1-post-SRD assms closure)
  thus ?thesis
  by (simp only: wp-rea-impl-lemma, simp add: Healthy-if 1, simp add: R1-post-SRD assms closure)
qed
also
have ... = (preR P ∧ postR P wpr preR (P ;; P ^ n) ⇒r periR P ∨ postR P ;; (⋀ i ∈ {0..n}. postR
P ^ i) ;; periR P))
  by (pred-auto)
also
have ... = (preR P ∧ postR P wpr preR (P ;; P ^ n) ⇒r periR P ∨ ((⋀ i ∈ {0..n}. postR P ^ (Suc
i)) ;; periR P))
  by (simp add: seq-Sup-distl seqr-assoc[THEN sym])
also
have ... = (preR P ∧ postR P wpr preR (P ;; P ^ n) ⇒r periR P ∨ ((⋀ i ∈ {1..Suc n}. postR P ^ i)
;; periR P))
proof –
  have (⋀ i ∈ {0..n}. postR P ^ Suc i) = (⋀ i ∈ {1..Suc n}. postR P ^ i)
  apply (rule cong[of Sup], auto)
  apply (metis atLeast0AtMost atMost-iff image-Suc-atLeastAtMost rev-image-eqI upred-semiring.power-Suc)
  using Suc-le-D apply fastforce
done
thus ?thesis by simp
qed
also
have ... = (preR P ∧ postR P wpr preR (P ;; P ^ n) ⇒r ((⋀ i ∈ {0..Suc n}. postR P ^ i) ;; periR P)
  by (simp add: SUP-atLeastAtMost-first uinf-or seqr-or-distl seqr-or-distr)
also
have ... = (preR(P ^ (Suc (Suc n)))) ⇒r ((⋀ i ∈ {0..Suc n}. postR P ^ i) ;; periR P)
  by (simp add: rdes closure assms)
finally show ?case by (simp)
qed

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lemma periR-power' [rdes]:
  assumes P is NSRD

```

shows $\text{peri}_R(P ;; P^\wedge n) = (\text{pre}_R(P^\wedge(\text{Suc } n)) \Rightarrow_r (\bigwedge i \in \{0..n\} \cdot \text{post}_R(P) \wedge i) ;; \text{peri}_R(P))$
by (*simp add: periR-power assms UINF-as-Sup[THEN sym]*)

lemma *periR-power-Suc* [*rdes*]:

assumes *P is NSRD*

shows $\text{peri}_R(P^\wedge(\text{Suc } n)) = (\text{pre}_R(P^\wedge(\text{Suc } n)) \Rightarrow_r (\bigwedge i \in \{0..n\} \cdot \text{post}_R(P) \wedge i) ;; \text{peri}_R(P))$

by (*simp add: rdes assms*)

lemma *postR-power* [*rdes*]:

assumes *P is NSRD*

shows $\text{post}_R(P ;; P^\wedge n) = (\text{pre}_R(P^\wedge(\text{Suc } n)) \Rightarrow_r \text{post}_R(P) \wedge \text{Suc } n)$

proof (*induct n*)

case 0

then show ?*case*

by (*simp add: NSRD-is-SRD NSRD-wait'-unrest-pre SRD-post-under-pre assms*)

next

case (*Suc n*) **note** *hyp = this*

have $\text{post}_R(P \wedge (\text{Suc } n + 1)) = \text{post}_R(P ;; P^\wedge(n+1))$

by (*simp*)

also have $\dots = (\text{pre}_R(P \wedge (\text{Suc } n + 1)) \Rightarrow_r (\text{post}_R P ;; \text{post}_R(P ;; P^\wedge n)))$

by (*simp add: closure assms rdes*)

also have $\dots = (\text{pre}_R(P \wedge (\text{Suc } n + 1)) \Rightarrow_r (\text{post}_R P ;; (\text{pre}_R(P \wedge \text{Suc } n) \Rightarrow_r \text{post}_R P \wedge \text{Suc } n)))$

by (*simp only: hyp*)

also

have $\dots = (\text{pre}_R P \Rightarrow_r (\text{post}_R P \text{ wp}_r \text{pre}_R(P \wedge \text{Suc } n) \Rightarrow_r \text{post}_R P ;; (\text{pre}_R(P \wedge \text{Suc } n) \Rightarrow_r \text{post}_R P \wedge \text{Suc } n)))$

by (*simp add: rdes closure assms, pred-auto*)

also

have $\dots = (\text{pre}_R P \Rightarrow_r (\text{post}_R P \text{ wp}_r \text{pre}_R(P \wedge \text{Suc } n) \Rightarrow_r \text{post}_R P ;; \text{post}_R P \wedge \text{Suc } n))$

by (*metis (no-types, lifting) Healthy-if NSRD-is-SRD NSRD-power-Suc R1-power assms hyp postR-SRD-R1 upred-semiring.power-Suc wp-rea-impl-lemma*)

also

have $\dots = (\text{pre}_R P \wedge \text{post}_R P \text{ wp}_r \text{pre}_R(P \wedge \text{Suc } n) \Rightarrow_r \text{post}_R P \wedge \text{Suc } (\text{Suc } n))$

by (*pred-auto*)

also have $\dots = (\text{pre}_R(P^\wedge(\text{Suc } (\text{Suc } n))) \Rightarrow_r \text{post}_R P \wedge \text{Suc } (\text{Suc } n))$

by (*simp add: rdes closure assms*)

finally show ?*case* **by** (*simp*)

qed

lemma *postR-power-Suc* [*rdes*]:

assumes *P is NSRD*

shows $\text{post}_R(P^\wedge(\text{Suc } n)) = (\text{pre}_R(P^\wedge(\text{Suc } n)) \Rightarrow_r \text{post}_R(P) \wedge \text{Suc } n)$

by (*simp add: rdes assms*)

lemma *power-rdes-def* [*rdes-def*]:

assumes *P is RC Q is RR R is RR \$st' \# Q*

shows $(\mathbf{R}_s(P \vdash Q \diamond R))^\wedge(\text{Suc } n)$

$= \mathbf{R}_s((\bigwedge i \in \{0..n\} \cdot (R \wedge i) \text{ wp}_r P) \vdash ((\bigwedge i \in \{0..n\} \cdot R \wedge i) ;; Q) \diamond (R \wedge \text{Suc } n))$

proof (*induct n*)

case 0

then show ?*case*

by (*simp add: wp assms closure*)

next

case (*Suc n*)

have 1: $(P \wedge (\bigsqcup i \in \{0..n\} \cdot R \text{ wp}_r (R \hat{~} i \text{ wp}_r P))) = (\bigsqcup i \in \{0..Suc\ n\} \cdot R \hat{~} i \text{ wp}_r P)$
(is ?lhs = ?rhs)
proof –
have ?lhs = $(P \wedge (\bigsqcup i \in \{0..n\} \cdot (R \hat{~} Suc\ i \text{ wp}_r P)))$
by (simp add: wp closure assms)
also have ... = $(P \wedge (\bigsqcup i \in \{0..n\} \cdot (R \hat{~} Suc\ i \text{ wp}_r P)))$
by (simp only: USUP-as-Inf-collect)
also have ... = $(P \wedge (\bigsqcup i \in \{1..Suc\ n\} \cdot (R \hat{~} i \text{ wp}_r P)))$
by (metis (no-types, lifting) INF-cong One-nat-def image-Suc-atLeastAtMost image-image)
also have ... = $(\bigsqcup i \in \text{insert } 0\ \{1..Suc\ n\} \cdot (R \hat{~} i \text{ wp}_r P))$
by (simp add: wp assms closure conj-upred-def)
also have ... = $(\bigsqcup i \in \{0..Suc\ n\} \cdot (R \hat{~} i \text{ wp}_r P))$
by (simp add: atLeastAtMost-insertL)
finally show ?thesis
by (simp add: USUP-as-Inf-collect)
qed

have 2: $(Q \vee R ;; (\prod i \in \{0..n\} \cdot R \hat{~} i) ;; Q) = (\prod i \in \{0..Suc\ n\} \cdot R \hat{~} i) ;; Q$
(is ?lhs = ?rhs)
proof –
have ?lhs = $(Q \vee (\prod i \in \{0..n\} \cdot R \hat{~} Suc\ i) ;; Q)$
by (simp add: seqr-assoc[THEN sym] seq-UINF-distl)
also have ... = $(Q \vee (\prod i \in \{0..n\} \cdot R \hat{~} Suc\ i) ;; Q)$
by (simp only: UINF-as-Sup-collect)
also have ... = $(Q \vee (\prod i \in \{1..Suc\ n\} \cdot R \hat{~} i) ;; Q)$
by (metis One-nat-def image-Suc-atLeastAtMost image-image)
also have ... = $((\prod i \in \text{insert } 0\ \{1..Suc\ n\} \cdot R \hat{~} i) ;; Q)$
by (simp add: disj-upred-def[THEN sym] seqr-or-distl)
also have ... = $((\prod i \in \{0..Suc\ n\} \cdot R \hat{~} i) ;; Q)$
by (simp add: atLeastAtMost-insertL)
finally show ?thesis
by (simp add: UINF-as-Sup-collect)
qed

have 3: $(\prod i \in \{0..n\} \cdot R \hat{~} i) ;; Q$ is RR
proof –
have $(\prod i \in \{0..n\} \cdot R \hat{~} i) ;; Q = (\prod i \in \{0..n\} \cdot R \hat{~} i) ;; Q$
by (simp add: UINF-as-Sup-collect)
also have ... = $(\prod i \in \text{insert } 0\ \{1..n\} \cdot R \hat{~} i) ;; Q$
by (simp add: atLeastAtMost-insertL)
also have ... = $(Q \vee (\prod i \in \{1..n\} \cdot R \hat{~} i) ;; Q)$
by (metis (no-types, lifting) SUP-insert disj-upred-def seqr-left-unit seqr-or-distl upred-semiring.power-0)
also have ... = $(Q \vee (\prod i \in \{0..<n\} \cdot R \hat{~} Suc\ i) ;; Q)$
by (metis One-nat-def atLeastLessThanSuc-atLeastAtMost image-Suc-atLeastLessThan image-image)
also have ... = $(Q \vee (\prod i \in \{0..<n\} \cdot R \hat{~} Suc\ i) ;; Q)$
by (simp add: UINF-as-Sup-collect)
also have ... is RR
by (simp-all add: closure assms)
finally show ?thesis .
qed

from 1 2 3 Suc **show** ?case
by (simp add: Suc RHS-tri-normal-design-composition' closure assms wp)
qed

declare upred-semiring.power-Suc [simp del]

theorem *uplus-rdes-def* [*rdes-def*]:
assumes P is RC Q is RR R is RR $\$st'$ $\#$ Q
shows $(\mathbf{R}_s(P \vdash Q \diamond R))^+ = \mathbf{R}_s(R^{*r} \text{ wp}_r P \vdash (R^{*r} ;; Q) \diamond R^+)$
proof –
have $1: (\bigcap i \cdot R \hat{=} i) ;; Q = R^{*r} ;; Q$
by (*metis* (*no-types*) *RA1* *assms*(2) *rea-skip-unit*(2) *rrel-thy.Star-def* *ustar-alt-def*)
show *?thesis*
by (*simp* *add: uplus-power-def seq-UINF-distr wp closure assms rdes-def*)
(metis 1 seq-UINF-distr')
qed

5.1 UTP theory

typedec *NSRDES*

abbreviation $NSRDES \equiv UTHY(NSRDES, ('s, 't::trace, 'α) \text{ rsp})$

overloading

nsrdes-hcond == *utp-hcond* :: $(NSRDES, ('s, 't::trace, 'α) \text{ rsp}) \text{ uthy} \Rightarrow ((('s, 't, 'α) \text{ rsp} \times ('s, 't, 'α) \text{ rsp}) \text{ health}$

nsrdes-unit == *utp-unit* :: $(NSRDES, ('s, 't::trace, 'α) \text{ rsp}) \text{ uthy} \Rightarrow ('s, 't, 'α) \text{ hrel-rsp}$

begin

definition *nsrdes-hcond* :: $(NSRDES, ('s, 't::trace, 'α) \text{ rsp}) \text{ uthy} \Rightarrow ((('s, 't, 'α) \text{ rsp} \times ('s, 't, 'α) \text{ rsp}) \text{ health}$ **where**

[*upred-defs*]: *nsrdes-hcond* $T = NSRD$

definition *nsrdes-unit* :: $(NSRDES, ('s, 't::trace, 'α) \text{ rsp}) \text{ uthy} \Rightarrow ('s, 't, 'α) \text{ hrel-rsp}$ **where**

[*upred-defs*]: *nsrdes-unit* $T = II_R$

end

Here, we show that normal stateful reactive designs form a Kleene UTP theory, and thus a Kleene algebra [4, 1]. This provides the basis for reasoning about iterative reactive contracts.

interpretation *nsrd-thy: utp-theory-kleene* $UTHY(NSRDES, ('s, 't::trace, 'α) \text{ rsp})$

rewrites $\bigwedge P. P \in \text{carrier (uthy-order } NSRDES) \longleftrightarrow P \text{ is } NSRD$

and $P \text{ is } \mathcal{H}_{NSRDES} \longleftrightarrow P \text{ is } NSRD$

and $(\mu X \cdot F (\mathcal{H}_{NSRDES} X)) = (\mu X \cdot F (NSRD X))$

and $\text{carrier (uthy-order } NSRDES) \rightarrow \text{carrier (uthy-order } NSRD) \equiv \llbracket NSRD \rrbracket_H \rightarrow \llbracket NSRD \rrbracket_H$

and $\llbracket \mathcal{H}_{NSRDES} \rrbracket_H \rightarrow \llbracket \mathcal{H}_{NSRDES} \rrbracket_H \equiv \llbracket NSRD \rrbracket_H \rightarrow \llbracket NSRD \rrbracket_H$

and $\top_{NSRDES} = \text{Miracle}$

and $\mathcal{II}_{NSRDES} = II_R$

and $le \text{ (uthy-order } NSRDES) = (\sqsubseteq)$

proof –

interpret *lat: utp-theory-continuous* $UTHY(NSRDES, ('s, 't, 'α) \text{ rsp})$

by (*unfold-locales*, *simp-all* *add: nsrdes-hcond-def nsrdes-unit-def closure Healthy-if*)

show $1: \top_{NSRDES} = (\text{Miracle} :: ('s, 't, 'α) \text{ hrel-rsp})$

by (*metis* *NSRD-Miracle* *NSRD-is-SRD* *lat.top-healthy* *lat.utp-theory-continuous-axioms* *nsrdes-hcond-def* *srdes-theory-continuous.meet-top* *upred-semiring.add-commute* *utp-theory-continuous.meet-top*)

thus *utp-theory-kleene* $UTHY(NSRDES, ('s, 't, 'α) \text{ rsp})$

by (*unfold-locales*, *simp-all* *add: nsrdes-hcond-def nsrdes-unit-def closure Healthy-if* *Miracle-left-zero* *SRD-left-unit* *NSRD-right-unit*)

qed (*simp-all* *add: nsrdes-hcond-def nsrdes-unit-def closure Healthy-if*)

declare *nsrd-thy.top-healthy* [*simp* *del*]

declare *nsrd-thy.bottom-healthy* [*simp* *del*]

abbreviation $TestR$ ($test_R$) **where**
 $test_R P \equiv utest\ NSRDES\ P$

abbreviation $StarR :: ('s, 't::trace, 'α) hrel-rsp \Rightarrow ('s, 't, 'α) hrel-rsp\ (-^{*R}\ [999]\ 999)$ **where**
 $StarR\ P \equiv P \star NSRDES$

We also show how to calculate the Kleene closure of a reactive design.

lemma $StarR-rdes-def$ [$rdes-def$]:

assumes P is RC Q is RR R is RR $\$st' \# Q$

shows $(\mathbf{R}_s(P \vdash Q \diamond R))^{*R} = \mathbf{R}_s((R^{*r}\ wp_r\ P) \vdash (R^{*r} ;; Q) \diamond R^{*r})$

by ($simp\ add: rrel-thy.Star-alt-def\ nsrd-thy.Star-alt-def\ assms\ closure\ rdes-def\ unrest\ rpred\ disj-upred-def$)

end

6 Syntax for reactive design contracts

theory $utp-rdes-contracts$

imports $utp-rdes-normal$

begin

We give an experimental syntax for reactive design contracts $[P \vdash Q | R]_R$, where P is a precondition on undashed state variables only, Q is a periccondition that can refer to the trace and before state but not the after state, and R is a postcondition. Both Q and R can refer only to the trace contribution through a HOL variable $trace$ which is bound to $\&tt$.

definition $mk-RD :: 's\ upred \Rightarrow ('t::trace \Rightarrow 's\ upred) \Rightarrow ('t \Rightarrow 's\ hrel) \Rightarrow ('s, 't, 'a) hrel-rsp$ **where**
 $mk-RD\ P\ Q\ R = \mathbf{R}_s([P]_{s<} \vdash [Q(x)]_{s<} \llbracket x \rightarrow \&tt \rrbracket \diamond [R(x)]_s \llbracket x \rightarrow \&tt \rrbracket)$

definition $trace-pred :: ('t::trace \Rightarrow 's\ upred) \Rightarrow ('s, 't, 'α) hrel-rsp$ **where**
 $[upred-defs]: trace-pred\ P = [(P\ x)]_{s<} \llbracket x \rightarrow \&tt \rrbracket$

syntax

$-trace-var :: logic$

$-mk-RD :: logic \Rightarrow logic \Rightarrow logic \Rightarrow logic\ ([-/ \vdash -/ \mid -]_R)$

$-trace-pred :: logic \Rightarrow logic\ ([_]_t)$

parse-translation \ll

let

$fun\ trace-var-tr\ [] = Syntax.free\ trace$

$\mid trace-var-tr\ - = raise\ Match;$

in

$[(\@ \{syntax-const\ -trace-var\}, K\ trace-var-tr)]$

end

\gg

translations

$[P \vdash Q \mid R]_R \Rightarrow CONST\ mk-RD\ P\ (\lambda\ -trace-var.\ Q)\ (\lambda\ -trace-var.\ R)$

$[P \vdash Q \mid R]_R \Leftarrow CONST\ mk-RD\ P\ (\lambda\ x.\ Q)\ (\lambda\ y.\ R)$

$[P]_t \Rightarrow CONST\ trace-pred\ (\lambda\ -trace-var.\ P)$

$[P]_t \Leftarrow CONST\ trace-pred\ (\lambda\ t.\ P)$

lemma $SRD-mk-RD$ [$closure$]: $[P \vdash Q(trace) \mid R(trace)]_R$ is SRD

by ($simp\ add: mk-RD-def\ closure\ unrest$)

lemma *preR-mk-RD* [rdes]: $\text{pre}_R([P \vdash Q(\text{trace}) \mid R(\text{trace})]_R) = R1(\lceil P \rceil_{S<})$
by (*simp add: mk-RD-def rea-pre-RHS-design usubst unrest R2c-not R2c-lift-state-pre*)

lemma *trace-pred-RR-closed* [closure]:
 $\lceil P \text{ trace} \rceil_t$ is *RR*
by (*rel-auto*)

lemma *unrest-trace-pred-st'* [unrest]:
 $\$st' \# \lceil P \text{ trace} \rceil_t$
by (*rel-auto*)

lemma *R2c-msubst-tt*: $R2c(\text{msubst}(\lambda x. \lceil Q \ x \rceil_S) \ \&tt) = (\text{msubst}(\lambda x. \lceil Q \ x \rceil_S) \ \&tt)$
by (*rel-auto*)

lemma *periR-mk-RD* [rdes]: $\text{peri}_R([P \vdash Q(\text{trace}) \mid R(\text{trace})]_R) = (\lceil P \rceil_{S<} \Rightarrow_r R1((\lceil Q(\text{trace}) \rceil_{S<}) \llbracket \text{trace} \rightarrow \&tt \rrbracket))$
by (*simp add: mk-RD-def rea-peri-RHS-design usubst unrest R2c-not R2c-lift-state-pre R2c-disj R2c-msubst-tt R1-disj R2c-rea-impl R1-rea-impl*)

lemma *postR-mk-RD* [rdes]: $\text{post}_R([P \vdash Q(\text{trace}) \mid R(\text{trace})]_R) = (\lceil P \rceil_{S<} \Rightarrow_r R1((\lceil R(\text{trace}) \rceil_S) \llbracket \text{trace} \rightarrow \&tt \rrbracket))$
by (*simp add: mk-RD-def rea-post-RHS-design usubst unrest R2c-not R2c-lift-state-pre impl-alt-def R2c-disj R2c-msubst-tt R2c-rea-impl R1-rea-impl*)

Refinement introduction law for contracts

lemma *RD-contract-refine*:

assumes

Q is *SRD* ‘ $\lceil P_1 \rceil_{S<} \Rightarrow \text{pre}_R Q$ ’,
‘ $\lceil P_1 \rceil_{S<} \wedge \text{peri}_R Q \Rightarrow \lceil P_2 \ x \rceil_{S<} \llbracket x \rightarrow \&tt \rrbracket$ ’,
‘ $\lceil P_1 \rceil_{S<} \wedge \text{post}_R Q \Rightarrow \lceil P_3 \ x \rceil_S \llbracket x \rightarrow \&tt \rrbracket$ ’,

shows $[P_1 \vdash P_2(\text{trace}) \mid P_3(\text{trace})]_R \sqsubseteq Q$

proof –

have $[P_1 \vdash P_2(\text{trace}) \mid P_3(\text{trace})]_R \sqsubseteq \mathbf{R}_s(\text{pre}_R(Q) \vdash \text{peri}_R(Q) \diamond \text{post}_R(Q))$

using *assms*

by (*simp add: mk-RD-def, rule-tac srdes-tri-refine-intro, simp-all*)

thus *?thesis*

by (*simp add: SRD-reactive-tri-design assms(1)*)

qed

end

7 Reactive design tactics

theory *utp-rdes-tactics*

imports *utp-rdes-triples*

begin

Theorems for normalisation

lemmas *rdes-rel-norms* =

prod.case-eq-if

conj-assoc

disj-assoc

utp-pred-laws.distrib(3,4)

conj-UINF-dist

conj-UINF-ind-dist

seqr-or-distl

seq-or-distr
seq-UINF-distl
seq-UINF-distl'
seq-UINF-distr
seq-UINF-distr'

The following tactic can be used to simply and evaluate reactive predicates.

method *rpred-simp* = (*uepr-simp* *simps*: *rpred* *usubst* *closure* *unrest*)

Tactic to expand out healthy reactive design predicates into the syntactic triple form.

method *rdes-expand* **uses** *cls* = (*insert* *cls*, (*erule* *RD-elim*)*+*)

Tactic to simplify the definition of a reactive design

method *rdes-simp* **uses** *cls* *cong* *simps* =
 ((*rdes-expand* *cls*: *cls*)?, (*simp* *add*: *closure*)?, (*simp* *add*: *rdes-def* *rdes-rel-norms* *rdes* *rpred* *cls* *closure*
alpha *frame* *usubst* *unrest* *wp* *simps* *cong*: *cong*))

Tactic to split a refinement conjecture into three POs

method *rdes-refine-split* **uses** *cls* *cong* *simps* =
 (*rdes-simp* *cls*: *cls* *cong*: *cong* *simps*: *simps*; *rule-tac* *srdes-tri-refine-intro*)

Tactic to split an equality conjecture into three POs

method *rdes-eq-split* **uses** *cls* *cong* *simps* =
 (*rdes-simp* *cls*: *cls* *cong*: *cong* *simps*: *simps*; (*rule-tac* *srdes-tri-eq-intro*))

Tactic to prove a refinement

method *rdes-refine* **uses** *cls* *cong* *simps* =
 (*rdes-refine-split* *cls*: *cls* *cong*: *cong* *simps*: *simps*; (*insert* *cls*; *rel-auto*))

Tactics to prove an equality

method *rdes-eq* **uses** *cls* *cong* *simps* =
 (*rdes-eq-split* *cls*: *cls* *cong*: *cong* *simps*: *simps*; *rel-auto*)

Via antisymmetry

method *rdes-eq-anti* **uses** *cls* *cong* *simps* =
 (*rdes-simp* *cls*: *cls* *cong*: *cong* *simps*: *simps*; (*rule-tac* *antisym*; (*rule-tac* *srdes-tri-refine-intro*; *rel-auto*)))

Tactic to calculate pre/peri/postconditions from reactive designs

method *rdes-calc* = (*simp* *add*: *rdes* *rpred* *closure* *alpha* *usubst* *unrest* *wp* *prod.case-eq-if*)

The following tactic attempts to prove a reactive design refinement by calculation of the pre-, peri-, and postconditions and then showing three implications between them using *rel-blast*.

method *rdspl-refine* =
 (*rule-tac* *SRD-refine-intro*; (*simp* *add*: *closure* *rdes* *unrest* *usubst* ; *rel-blast*?))

The following tactic combines antisymmetry with the previous tactic to prove an equality.

method *rdspl-eq* =
 (*rule-tac* *antisym*, *rdes-refine*, *rdes-refine*)

end

8 Reactive design parallel-by-merge

```

theory utp-rdes-parallel
  imports
    utp-rdes-normal
    utp-rdes-tactics
begin

```

R3h implicitly depends on RD1, and therefore it requires that both sides be RD1. We also require that both sides are R3c, and that $wait_m$ is a quasi-unit, and div_m yields divergence.

```

lemma st-U0-alpha:  $\exists \$st \cdot II \rfloor_0 = (\exists \$st \cdot \lceil II \rceil_0)$ 
  by (rel-auto)

```

```

lemma st-U1-alpha:  $\exists \$st \cdot II \rfloor_1 = (\exists \$st \cdot \lceil II \rceil_1)$ 
  by (rel-auto)

```

```

definition skip-rm :: ('s, 't::trace, 'α) rsp merge (IIRM) where
  [upred-defs]: IIRM =  $(\exists \$st_{<} \cdot skip_m \vee (\neg \$ok_{<} \wedge \$tr_{<} \leq_u \$tr'))$ 

```

```

definition [upred-defs]: R3hm(M) = (IIRM ◁ $wait_{<} ▷ M)

```

```

lemma R3hm-idem: R3hm(R3hm(P)) = R3hm(P)
  by (rel-auto)

```

```

lemma R3h-par-by-merge [closure]:
  assumes P is R3h Q is R3h M is R3hm
  shows (P ||M Q) is R3h

```

proof –

```

  have (P ||M Q) = (((P ||M Q) ⟦true/$ok⟧ ◁ $ok ▷ (P ||M Q) ⟦false/$ok⟧) ⟦true/$wait⟧ ◁ $wait ▷ (P ||M Q))

```

```

    by (simp add: cond-var-subst-left cond-var-subst-right)

```

```

  also have ... = (((P ||M Q) ⟦true,true/$ok,$wait⟧ ◁ $ok ▷ (P ||M Q) ⟦false,true/$ok,$wait⟧) ◁ $wait ▷ (P ||M Q))

```

```

    by (rel-auto)

```

```

  also have ... = (((∃ $st · II) ⟦true,true/$ok,$wait⟧ ◁ $ok ▷ (P ||M Q) ⟦false,true/$ok,$wait⟧) ◁ $wait ▷ (P ||M Q))

```

proof –

```

  have (P ||M Q) ⟦true,true/$ok,$wait⟧ = ((⟦P⟧0 ∧ ⟦Q⟧1 ∧ $v_{<}′ =u $v) ;; R3hm(M)) ⟦true,true/$ok,$wait⟧
    by (simp add: par-by-merge-def U0-as-alpha U1-as-alpha assms Healthy-if)

```

```

  also have ... = ((⟦P⟧0 ∧ ⟦Q⟧1 ∧ $v_{<}′ =u $v) ;; (∃ $st_{<} · $v′ =u $v_{<})) ⟦true,true/$ok,$wait⟧

```

```

    by (rel-blast)

```

```

  also have ... = ((⟦R3h(P)⟧0 ∧ ⟦R3h(Q)⟧1 ∧ $v_{<}′ =u $v) ;; (∃ $st_{<} · $v′ =u $v_{<})) ⟦true,true/$ok,$wait⟧

```

```

    by (simp add: assms Healthy-if)

```

```

  also have ... = (∃ $st · II) ⟦true,true/$ok,$wait⟧

```

```

    by (rel-auto)

```

```

  finally show ?thesis by (simp add: closure assms unrest)

```

qed

```

  also have ... = (((∃ $st · II) ⟦true,true/$ok,$wait⟧ ◁ $ok ▷ (R1(true)) ⟦false,true/$ok,$wait⟧) ◁ $wait ▷ (P ||M Q))

```

proof –

```

  have (P ||M Q) ⟦false,true/$ok,$wait⟧ = ((⟦P⟧0 ∧ ⟦Q⟧1 ∧ $v_{<}′ =u $v) ;; R3hm(M)) ⟦false,true/$ok,$wait⟧
    by (simp add: par-by-merge-def U0-as-alpha U1-as-alpha assms Healthy-if)

```

```

  also have ... = ((⟦P⟧0 ∧ ⟦Q⟧1 ∧ $v_{<}′ =u $v) ;; ($tr_{<} ≤u $tr′)) ⟦false,true/$ok,$wait⟧

```

```

    by (rel-blast)

```

```

  also have ... = ((⟦R3h(P)⟧0 ∧ ⟦R3h(Q)⟧1 ∧ $v_{<}′ =u $v) ;; ($tr_{<} ≤u $tr′)) ⟦false,true/$ok,$wait⟧

```

```

    by (simp add: assms Healthy-if)

```

also have ... = $(R1(true))\llbracket false, true / \$ok, \$wait \rrbracket$
 by (rel-blast)
 finally show ?thesis by simp
 qed
 also have ... = $((\exists \$st \cdot II) \triangleleft \$ok \triangleright R1(true)) \triangleleft \$wait \triangleright (P \parallel_M Q)$
 by (rel-auto)
 also have ... = $R3h(P \parallel_M Q)$
 by (simp add: R3h-cases)
 finally show ?thesis
 by (simp add: Healthy-def)
 qed

definition [upred-defs]: $RD1m(M) = (M \vee \neg \$ok_{<} \wedge \$tr_{<} \leq_u \$tr')$

lemma *RD1-par-by-merge* [closure]:

assumes P is $R1$ Q is $R1$ M is $R1m$ P is $RD1$ Q is $RD1$ M is $RD1m$
 shows $(P \parallel_M Q)$ is $RD1$

proof –

have 1: $(RD1(R1(P)) \parallel_{RD1m(R1m(M))} RD1(R1(Q)))\llbracket false / \$ok \rrbracket = R1(true)$
 by (rel-blast)
 have $(P \parallel_M Q) = (P \parallel_M Q)\llbracket true / \$ok \rrbracket \triangleleft \$ok \triangleright (P \parallel_M Q)\llbracket false / \$ok \rrbracket$
 by (simp add: cond-var-split)
 also have ... = $R1(P \parallel_M Q) \triangleleft \$ok \triangleright R1(true)$
 by (metis 1 Healthy-if R1-par-by-merge assms calculation
 cond-idem cond-var-subst-right in-var-uvar ok-vwb-lens)
 also have ... = $RD1(P \parallel_M Q)$
 by (simp add: Healthy-if R1-par-by-merge RD1-alt-def assms(3))
 finally show ?thesis
 by (simp add: Healthy-def)

qed

lemma *RD2-par-by-merge* [closure]:

assumes M is $RD2$
 shows $(P \parallel_M Q)$ is $RD2$

proof –

have $(P \parallel_M Q) = ((P \parallel_s Q) ;; M)$
 by (simp add: par-by-merge-def)
 also from assms have ... = $((P \parallel_s Q) ;; (M ;; J))$
 by (simp add: Healthy-def' RD2-def H2-def)
 also from assms have ... = $((P \parallel_s Q) ;; M) ;; J$
 by (simp add: seqr-assoc)
 also from assms have ... = $RD2(P \parallel_M Q)$
 by (simp add: RD2-def H2-def par-by-merge-def)
 finally show ?thesis
 by (simp add: Healthy-def')

qed

lemma *SRD-par-by-merge*:

assumes P is SRD Q is SRD M is $R1m$ M is $R2m$ M is $R3hm$ M is $RD1m$ M is $RD2$
 shows $(P \parallel_M Q)$ is SRD
 by (rule SRD-intro, simp-all add: assms closure SRD-healths)

definition *nmerge-rd0* (N_0) **where**

[upred-defs]: $N_0(M) = (\$wait' =_u (\$0 - wait \vee \$1 - wait) \wedge \$tr_{<} \leq_u \$tr' \wedge (\exists \$0 - ok; \$1 - ok; \$ok_{<} \$ok'; \$0 - wait; \$1 - wait; \$wait_{<} \$wait' \cdot M))$

definition *nmerge-rd1* (N_1) **where**

[*upred-defs*]: $N_1(M) = (\$ok' =_u (\$0-ok \wedge \$1-ok) \wedge N_0(M))$

definition *nmerge-rd* (N_R) **where**

[*upred-defs*]: $N_R(M) = ((\exists \$st_< \cdot \$v' =_u \$v_<) \triangleleft \$wait_< \triangleright N_1(M)) \triangleleft \$ok_< \triangleright (\$tr_< \leq_u \$tr')$

definition *merge-rd1* (M_1) **where**

[*upred-defs*]: $M_1(M) = (N_1(M) ;; II_R)$

definition *merge-rd* (M_R) **where**

[*upred-defs*]: $M_R(M) = N_R(M) ;; II_R$

abbreviation *rdes-par* ($- \parallel_R -$ - [85,0,86] 85) **where**

$P \parallel_{RM} Q \equiv P \parallel_{M_R(M)} Q$

Healthiness condition for reactive design merge predicates

definition [*upred-defs*]: $RDM(M) = R2m(\exists \$0-ok; \$1-ok; \$ok_<; \$ok'; \$0-wait; \$1-wait; \$wait_<; \$wait' \cdot M)$

lemma *nmerge-rd-is-R1m* [*closure*]:

$N_R(M)$ *is* *R1m*

by (*rel-blast*)

lemma *R2m-nmerge-rd*: $R2m(N_R(R2m(M))) = N_R(R2m(M))$

apply (*rel-auto*) **using** *minus-zero-eq* **by** *blast+*

lemma *nmerge-rd-is-R2m* [*closure*]:

M *is* *R2m* $\implies N_R(M)$ *is* *R2m*

by (*metis Healthy-def' R2m-nmerge-rd*)

lemma *nmerge-rd-is-R3hm* [*closure*]: $N_R(M)$ *is* *R3hm*

by (*rel-blast*)

lemma *nmerge-rd-is-RD1m* [*closure*]: $N_R(M)$ *is* *RD1m*

by (*rel-blast*)

lemma *merge-rd-is-RD3*: $M_R(M)$ *is* *RD3*

by (*metis Healthy-Idempotent RD3-Idempotent RD3-def merge-rd-def*)

lemma *merge-rd-is-RD2*: $M_R(M)$ *is* *RD2*

by (*simp add: RD3-implies-RD2 merge-rd-is-RD3*)

lemma *par-rdes-NSRD* [*closure*]:

assumes P *is* *SRD* Q *is* *SRD* M *is* *RDM*

shows $P \parallel_{RM} Q$ *is* *NSRD*

proof –

have $(P \parallel_{N_R} M Q ;; II_R)$ *is* *NSRD*

by (*rule NSRD-intro', simp-all add: SRD-healths closure assms*)

(*metis (no-types, lifting) Healthy-def R2-par-by-merge R2-seqr-closure R2m-nmerge-rd RDM-def SRD-healths(2) assms skip-srea-R2*

, *metis Healthy-Idempotent RD3-Idempotent RD3-def*)

thus *?thesis*

by (*simp add: merge-rd-def par-by-merge-def seqr-assoc*)

qed

lemma *RDM-intro*:

assumes M is *R2m* $\$0-ok \# M \$1-ok \# M \$ok_{<} \# M \$ok' \# M$
 $\$0-wait \# M \$1-wait \# M \$wait_{<} \# M \$wait' \# M$
shows M is *RDM*
using *assms*
by (*simp add: Healthy-def RDM-def ex-unrest unrest*)

lemma *RDM-unrests* [*unrest*]:

assumes M is *RDM*
shows $\$0-ok \# M \$1-ok \# M \$ok_{<} \# M \$ok' \# M$
 $\$0-wait \# M \$1-wait \# M \$wait_{<} \# M \$wait' \# M$
by (*subst Healthy-if[OF assms, THEN sym], simp-all add: RDM-def unrest, rel-auto*) $+$

lemma *RDM-R1m* [*closure*]: M is *RDM* $\implies M$ is *R1m*

by (*metis (no-types, hide-lams) Healthy-def R1m-idem R2m-def RDM-def*)

lemma *RDM-R2m* [*closure*]: M is *RDM* $\implies M$ is *R2m*

by (*metis (no-types, hide-lams) Healthy-def R2m-idem RDM-def*)

lemma *ex-st'-R2m-closed* [*closure*]:

assumes P is *R2m*
shows $(\exists \$st' \cdot P)$ is *R2m*

proof $-$

have $R2m(\exists \$st' \cdot R2m P) = (\exists \$st' \cdot R2m P)$
by (*rel-auto*)
thus *?thesis*
by (*metis Healthy-def' assms*)

qed

lemma *parallel-RR-closed*:

assumes P is *RR* Q is *RR* M is *R2m*
 $\$ok_{<} \# M \$wait_{<} \# M \$ok' \# M \$wait' \# M$
shows $P \parallel_M Q$ is *RR*
by (*rule RR-R2-intro, simp-all add: unrest assms RR-implies-R2 closure*)

lemma *parallel-ok-cases*:

$((P \parallel_s Q) ;; M) = ($
 $((P^t \parallel_s Q^t) ;; (M \llbracket true, true / \$0-ok, \$1-ok \rrbracket)) \vee$
 $((P^f \parallel_s Q^t) ;; (M \llbracket false, true / \$0-ok, \$1-ok \rrbracket)) \vee$
 $((P^t \parallel_s Q^f) ;; (M \llbracket true, false / \$0-ok, \$1-ok \rrbracket)) \vee$
 $((P^f \parallel_s Q^f) ;; (M \llbracket false, false / \$0-ok, \$1-ok \rrbracket)))$

proof $-$

have $((P \parallel_s Q) ;; M) = (\exists ok_0 \cdot (P \parallel_s Q) \llbracket \llbracket ok_0 \rrbracket / \$0-ok \rrbracket ;; M \llbracket \llbracket ok_0 \rrbracket / \$0-ok \rrbracket)$
by (*subst seqr-middle[of left-uvar ok], simp-all*)
also have $\dots = (\exists ok_0 \cdot \exists ok_1 \cdot ((P \parallel_s Q) \llbracket \llbracket ok_0 \rrbracket / \$0-ok \rrbracket \llbracket \llbracket ok_1 \rrbracket / \$1-ok \rrbracket ;; (M \llbracket \llbracket ok_0 \rrbracket / \$0-ok \rrbracket \llbracket \llbracket ok_1 \rrbracket / \$1-ok \rrbracket))$
by (*subst seqr-middle[of right-uvar ok], simp-all*)
also have $\dots = (\exists ok_0 \cdot \exists ok_1 \cdot (P \llbracket \llbracket ok_0 \rrbracket / \$ok' \rrbracket \parallel_s Q \llbracket \llbracket ok_1 \rrbracket / \$ok' \rrbracket) ;; (M \llbracket \llbracket ok_0 \rrbracket, \llbracket ok_1 \rrbracket / \$0-ok, \$1-ok \rrbracket))$
by (*rel-auto robust*)
also have $\dots = ($
 $((P^t \parallel_s Q^t) ;; (M \llbracket true, true / \$0-ok, \$1-ok \rrbracket)) \vee$
 $((P^f \parallel_s Q^t) ;; (M \llbracket false, true / \$0-ok, \$1-ok \rrbracket)) \vee$
 $((P^t \parallel_s Q^f) ;; (M \llbracket true, false / \$0-ok, \$1-ok \rrbracket)) \vee$
 $((P^f \parallel_s Q^f) ;; (M \llbracket false, false / \$0-ok, \$1-ok \rrbracket)))$
by (*simp add: true-alt-def[THEN sym] false-alt-def[THEN sym] disj-assoc*)

```

    utp-pred-laws.sup.left-commute utp-pred-laws.sup-commute usubst)
  finally show ?thesis .
qed

lemma skip-srea-ok-f [usubst]:
  
$$II_R^f = R1(\neg \$ok)$$

  by (rel-auto)

lemma nmerge0-rd-unrest [unrest]:
  
$$\$0-ok \# N_0 \ M \ \$1-ok \# N_0 \ M$$

  by (pred-auto)+

lemma parallel-assm-lemma:
  assumes  $P$  is RD2
  shows  $pre_s \uparrow (P \parallel_{M_R(M)} Q) = ((pre_s \uparrow P) \parallel_{N_0(M)} ;; R1(true) (cmt_s \uparrow Q))$ 
     $\vee ((cmt_s \uparrow P) \parallel_{N_0(M)} ;; R1(true) (pre_s \uparrow Q))$ 
proof -
  have  $pre_s \uparrow (P \parallel_{M_R(M)} Q) = pre_s \uparrow ((P \parallel_s Q) ;; M_R(M))$ 
    by (simp add: par-by-merge-def)
  also have ... =  $((P \parallel_s Q) \llbracket true, false / \$ok, \$wait \rrbracket ;; N_R \ M ;; R1(\neg \$ok))$ 
    by (simp add: merge-rd-def usubst, rel-auto)
  also have ... =  $((P \llbracket true, false / \$ok, \$wait \rrbracket \parallel_s Q \llbracket true, false / \$ok, \$wait \rrbracket) ;; N_1(M) ;; R1(\neg \$ok))$ 
    by (rel-auto robust, (metis)+)
  also have ... = (
     $((P \llbracket true, false / \$ok, \$wait \rrbracket)^t \parallel_s (Q \llbracket true, false / \$ok, \$wait \rrbracket)^t) ;; ((N_1 \ M) \llbracket true, true / \$0-ok, \$1-ok \rrbracket$ 
     $;; R1(\neg \$ok)) \vee$ 
     $((P \llbracket true, false / \$ok, \$wait \rrbracket)^f \parallel_s (Q \llbracket true, false / \$ok, \$wait \rrbracket)^t) ;; ((N_1 \ M) \llbracket false, true / \$0-ok, \$1-ok \rrbracket$ 
     $;; R1(\neg \$ok)) \vee$ 
     $((P \llbracket true, false / \$ok, \$wait \rrbracket)^t \parallel_s (Q \llbracket true, false / \$ok, \$wait \rrbracket)^f) ;; ((N_1 \ M) \llbracket true, false / \$0-ok, \$1-ok \rrbracket$ 
     $;; R1(\neg \$ok)) \vee$ 
     $((P \llbracket true, false / \$ok, \$wait \rrbracket)^f \parallel_s (Q \llbracket true, false / \$ok, \$wait \rrbracket)^f) ;; ((N_1 \ M) \llbracket false, false / \$0-ok, \$1-ok \rrbracket$ 
     $;; R1(\neg \$ok))$ 
  )
  (is - =  $(?C1 \vee_p ?C2 \vee_p ?C3 \vee_p ?C4)$ )
  by (subst parallel-ok-cases, subst-tac)
  also have ... =  $(?C2 \vee ?C3)$ 
proof -
  have  $?C1 = false$ 
    by (rel-auto)
  moreover have  $'?C4 \Rightarrow ?C3'$  (is  $'(?A ;; ?B) \Rightarrow (?C ;; ?D)'$ )
proof -
  from assms have  $'P^f \Rightarrow P^t'$ 
    by (metis RD2-def H2-equivalence Healthy-def)
  hence  $P: 'P^f_f \Rightarrow P^t_f'$ 
    by (rel-auto)
  have  $'?A \Rightarrow ?C'$ 
    using  $P$  by (rel-auto)
  moreover have  $'?B \Rightarrow ?D'$ 
    by (rel-auto)
  ultimately show ?thesis
    by (simp add: impl-seqr-mono)
qed
ultimately show ?thesis
  by (simp add: subsumption2)
qed
also have ... = (

```

$((pre_s \uparrow P) \parallel_s (cmt_s \uparrow Q)) \parallel (N_0 M \parallel R1(true)) \vee$
 $((cmt_s \uparrow P) \parallel_s (pre_s \uparrow Q)) \parallel (N_0 M \parallel R1(true))$
by (*rel-auto, metis+*)
also have ... = (
 $((pre_s \uparrow P) \parallel_{N_0 M \parallel R1(true)} (cmt_s \uparrow Q)) \vee$
 $((cmt_s \uparrow P) \parallel_{N_0 M \parallel R1(true)} (pre_s \uparrow Q))$
by (*simp add: par-by-merge-def*)
finally show *?thesis* .
qed

lemma *pre_s-SRD*:
assumes *P is SRD*
shows $pre_s \uparrow P = (\neg_r pre_R(P))$
proof –
have $pre_s \uparrow P = pre_s \uparrow \mathbf{R}_s(pre_R P \vdash peri_R P \diamond post_R P)$
by (*simp add: SRD-reactive-tri-design assms*)
also have ... = $R1(R2c(\neg pre_s \uparrow pre_R P))$
by (*simp add: RHS-def usubst R3h-def pre_s-design*)
also have ... = $R1(R2c(\neg pre_R P))$
by (*rel-auto*)
also have ... = $(\neg_r pre_R P)$
by (*simp add: R2c-not R2c-preR assms rea-not-def*)
finally show *?thesis* .
qed

lemma *parallel-assm*:
assumes *P is SRD Q is SRD*
shows $pre_R(P \parallel_{M_R(M)} Q) = (\neg_r ((\neg_r pre_R(P)) \parallel_{N_0(M)} R1(true) cmt_R(Q)) \wedge$
 $\neg_r (cmt_R(P) \parallel_{N_0(M)} R1(true) (\neg_r pre_R(Q))))$
(is ?lhs = ?rhs)
proof –
have $pre_R(P \parallel_{M_R(M)} Q) = (\neg_r (pre_s \uparrow P) \parallel_{N_0 M} R1 true (cmt_s \uparrow Q) \wedge$
 $\neg_r (cmt_s \uparrow P) \parallel_{N_0 M} R1 true (pre_s \uparrow Q))$
by (*simp add: pre_R-def parallel-assm-lemma assms SRD-healths R1-conj rea-not-def [THEN sym]*)
also have ... = *?rhs*
by (*simp add: pre_s-SRD assms cmt_R-def Healthy-if closure unrest*)
finally show *?thesis* .
qed

lemma *parallel-assm-unrest-wait' [unrest]*:
 $\llbracket P \text{ is SRD}; Q \text{ is SRD} \rrbracket \implies \$wait' \# pre_R(P \parallel_{M_R(M)} Q)$
by (*simp add: parallel-assm, simp add: par-by-merge-def unrest*)

lemma *JL1*: $(M_1 M)^t \llbracket false, true / \$0-ok, \$1-ok \rrbracket = N_0(M) \parallel R1(true)$
by (*rel-blast*)

lemma *JL2*: $(M_1 M)^t \llbracket true, false / \$0-ok, \$1-ok \rrbracket = N_0(M) \parallel R1(true)$
by (*rel-blast*)

lemma *JL3*: $(M_1 M)^t \llbracket false, false / \$0-ok, \$1-ok \rrbracket = N_0(M) \parallel R1(true)$
by (*rel-blast*)

lemma *JL4*: $(M_1 \ M)^t \llbracket \text{true}, \text{true} / \$0\text{--}ok, \$1\text{--}ok \rrbracket = (\$ok' \wedge N_0 \ M) ;; II_R^t$
 by (*simp add: merge-rd1-def usubst nmerge-rd1-def unrest*)

lemma *parallel-commitment-lemma-1*:

assumes *P* is *RD2*

shows $cmt_s \dagger (P \parallel_{M_R(M)} Q) =$

$((cmt_s \dagger P) \parallel_{(\$ok' \wedge N_0 \ M) ;; II_R^t (cmt_s \dagger Q)}) \vee$
 $((pre_s \dagger P) \parallel_{N_0(M) ;; R1(true) (cmt_s \dagger Q)}) \vee$
 $((cmt_s \dagger P) \parallel_{N_0(M) ;; R1(true) (pre_s \dagger Q)})$

proof –

have $cmt_s \dagger (P \parallel_{M_R(M)} Q) = (P \llbracket \text{true}, \text{false} / \$ok, \$wait \rrbracket \parallel_{(M_1(M))^t} Q \llbracket \text{true}, \text{false} / \$ok, \$wait \rrbracket)$

by (*simp add: usubst, rel-auto*)

also have $\dots = ((P \llbracket \text{true}, \text{false} / \$ok, \$wait \rrbracket \parallel_s Q \llbracket \text{true}, \text{false} / \$ok, \$wait \rrbracket) ;; (M_1 \ M)^t)$

by (*simp add: par-by-merge-def*)

also have $\dots =$

$((cmt_s \dagger P) \parallel_s (cmt_s \dagger Q)) ;; ((M_1 \ M)^t \llbracket \text{true}, \text{true} / \$0\text{--}ok, \$1\text{--}ok \rrbracket) \vee$
 $((pre_s \dagger P) \parallel_s (cmt_s \dagger Q)) ;; ((M_1 \ M)^t \llbracket \text{false}, \text{true} / \$0\text{--}ok, \$1\text{--}ok \rrbracket) \vee$
 $((cmt_s \dagger P) \parallel_s (pre_s \dagger Q)) ;; ((M_1 \ M)^t \llbracket \text{true}, \text{false} / \$0\text{--}ok, \$1\text{--}ok \rrbracket) \vee$
 $((pre_s \dagger P) \parallel_s (pre_s \dagger Q)) ;; ((M_1 \ M)^t \llbracket \text{false}, \text{false} / \$0\text{--}ok, \$1\text{--}ok \rrbracket)$

by (*subst parallel-ok-cases, subst-tac*)

also have $\dots =$

$((cmt_s \dagger P) \parallel_s (cmt_s \dagger Q)) ;; ((M_1 \ M)^t \llbracket \text{true}, \text{true} / \$0\text{--}ok, \$1\text{--}ok \rrbracket) \vee$
 $((pre_s \dagger P) \parallel_s (cmt_s \dagger Q)) ;; (N_0(M) ;; R1(true)) \vee$
 $((cmt_s \dagger P) \parallel_s (pre_s \dagger Q)) ;; (N_0(M) ;; R1(true)) \vee$
 $((pre_s \dagger P) \parallel_s (pre_s \dagger Q)) ;; (N_0(M) ;; R1(true))$
 (**is** $= (?C1 \vee_p ?C2 \vee_p ?C3 \vee_p ?C4)$)

by (*simp add: JL1 JL2 JL3*)

also have $\dots =$

$((cmt_s \dagger P) \parallel_s (cmt_s \dagger Q)) ;; ((M_1(M))^t \llbracket \text{true}, \text{true} / \$0\text{--}ok, \$1\text{--}ok \rrbracket) \vee$
 $((pre_s \dagger P) \parallel_s (cmt_s \dagger Q)) ;; (N_0(M) ;; R1(true)) \vee$
 $((cmt_s \dagger P) \parallel_s (pre_s \dagger Q)) ;; (N_0(M) ;; R1(true))$

proof –

from *assms* **have** $P^f \Rightarrow P^t$

by (*metis RD2-def H2-equivalence Healthy-def*)

hence $P: P^f_f \Rightarrow P^t_f$

by (*rel-auto*)

have $?C4 \Rightarrow ?C3$ (**is** $(?A ;; ?B) \Rightarrow (?C ;; ?D)$)

proof –

have $?A \Rightarrow ?C$

using *P* by (*rel-auto*)

thus *?thesis*

by (*simp add: impl-seqr-mono*)

qed

thus *?thesis*

by (*simp add: subsumption2*)

qed

finally show *?thesis*

by (*simp add: par-by-merge-def JL4*)

qed

lemma *parallel-commitment-lemma-2*:

assumes *P* is *RD2*

shows $cmt_s \dagger (P \parallel_{M_R(M)} Q) =$

$((cmt_s \dagger P) \parallel_{(\$ok' \wedge N_0 \ M) ;; II_R^t (cmt_s \dagger Q)}) \vee pre_s \dagger (P \parallel_{M_R(M)} Q)$

by (simp add: parallel-commitment-lemma-1 assms parallel-assm-lemma)

lemma parallel-commitment-lemma-3:

M is $R1m \implies (\$ok' \wedge N_0 M) ;; II_{R^t}$ is $R1m$

by (rel-simp, safe, metis+)

lemma parallel-commitment:

assumes P is SRD Q is SRD M is RDM

shows $cmt_R(P \parallel_{M_R(M)} Q) = (pre_R(P \parallel_{M_R(M)} Q) \Rightarrow_r cmt_R(P) \parallel_{(\$ok' \wedge N_0 M) ;; II_{R^t}} cmt_R(Q))$

by (simp add: parallel-commitment-lemma-2 parallel-commitment-lemma-3 Healthy-if assms cmt_R-def pre_s-SRD closure rea-impl-def disj-comm unrest)

theorem parallel-reactive-design:

assumes P is SRD Q is SRD M is RDM

shows $(P \parallel_{M_R(M)} Q) = \mathbf{R}_s($

$(\neg_r ((\neg_r pre_R(P)) \parallel_{N_0(M) ;; R1(true)} cmt_R(Q)) \wedge$

$\neg_r (cmt_R(P) \parallel_{N_0(M) ;; R1(true)} (\neg_r pre_R(Q)))) \vdash$

$(cmt_R(P) \parallel_{(\$ok' \wedge N_0 M) ;; II_{R^t}} cmt_R(Q)))$ (is ?lhs = ?rhs)

proof –

have $(P \parallel_{M_R(M)} Q) = \mathbf{R}_s(pre_R(P \parallel_{M_R(M)} Q) \vdash cmt_R(P \parallel_{M_R(M)} Q))$

by (metis Healthy-def NSRD-is-SRD SRD-as-reactive-design assms(1) assms(2) assms(3) par-rdes-NSRD)

also have ... = ?rhs

by (simp add: parallel-assm parallel-commitment design-export-spec assms, rel-auto)

finally show ?thesis .

qed

lemma parallel-periccondition-lemma1:

$(\$ok' \wedge P) ;; II_R \llbracket true, true / \$ok', \$wait' \rrbracket = (\exists \$st' \cdot P) \llbracket true, true / \$ok', \$wait' \rrbracket$

(is ?lhs = ?rhs)

proof –

have ?lhs = $(\$ok' \wedge P) ;; (\exists \$st \cdot II) \llbracket true, true / \$ok', \$wait' \rrbracket$

by (rel-blast)

also have ... = ?rhs

by (rel-auto)

finally show ?thesis .

qed

lemma parallel-periccondition-lemma2:

assumes M is RDM

shows $(\exists \$st' \cdot N_0(M)) \llbracket true, true / \$ok', \$wait' \rrbracket = ((\$0-wait \vee \$1-wait) \wedge (\exists \$st' \cdot M))$

proof –

have $(\exists \$st' \cdot N_0(M)) \llbracket true, true / \$ok', \$wait' \rrbracket = (\exists \$st' \cdot (\$0-wait \vee \$1-wait) \wedge \$tr' \geq_u \$tr_{<} \wedge M)$

by (simp add: usubst unrest nmerge-rd0-def ex-unrest Healthy-if R1m-def assms)

also have ... = $(\exists \$st' \cdot (\$0-wait \vee \$1-wait) \wedge M)$

by (metis (no-types, hide-lams) Healthy-if R1m-def R1m-idem R2m-def RDM-def assms utp-pred-laws.inf-commute)

also have ... = $(\$0-wait \vee \$1-wait) \wedge (\exists \$st' \cdot M)$

by (rel-auto)

finally show ?thesis .

qed

lemma parallel-periccondition-lemma3:

$((\$0-wait \vee \$1-wait) \wedge (\exists \$st' \cdot M)) = ((\$0-wait \wedge \$1-wait \wedge (\exists \$st' \cdot M)) \vee (\neg \$0-wait \wedge \$1-wait \wedge (\exists \$st' \cdot M)))$

by (rel-auto)

lemma *parallel-pericondition* [rdes]:

fixes $M :: ('s, 't :: \text{trace}, 'a) \text{ rsp merge}$

assumes P is SRD Q is SRD M is RDM

shows $\text{peri}_R(P \parallel_{M_R(M)} Q) = (\text{pre}_R(P \parallel_{M_R} M Q) \Rightarrow_r \text{peri}_R(P) \parallel_{\exists \$st' \cdot M} \text{peri}_R(Q)$
 $\vee \text{post}_R(P) \parallel_{\exists \$st' \cdot M} \text{peri}_R(Q)$
 $\vee \text{peri}_R(P) \parallel_{\exists \$st' \cdot M} \text{post}_R(Q))$

proof –

have $\text{peri}_R(P \parallel_{M_R(M)} Q) =$

$(\text{pre}_R(P \parallel_{M_R} M Q) \Rightarrow_r \text{cmt}_R P \parallel_{(\$ok' \wedge N_0 M) ;; II_R[\text{true}, \text{true}/\$ok', \$wait']} \text{cmt}_R Q)$

by (simp add: peri-cmt-def parallel-commitment SRD-healths assms usubst unrest assms)

also have $\dots = (\text{pre}_R(P \parallel_{M_R} M Q) \Rightarrow_r \text{cmt}_R P \parallel_{(\exists \$st' \cdot N_0 M)[\text{true}, \text{true}/\$ok', \$wait']} \text{cmt}_R Q)$

by (simp add: parallel-pericondition-lemma1)

also have $\dots = (\text{pre}_R(P \parallel_{M_R} M Q) \Rightarrow_r \text{cmt}_R P \parallel_{(\$0\text{-wait} \vee \$1\text{-wait}) \wedge (\exists \$st' \cdot M)} \text{cmt}_R Q)$

by (simp add: parallel-pericondition-lemma2 assms)

also have $\dots = (\text{pre}_R(P \parallel_{M_R} M Q) \Rightarrow_r ((\lceil \text{cmt}_R P \rceil_0 \wedge \lceil \text{cmt}_R Q \rceil_1 \wedge \$\mathbf{v}_{<}' =_u \$\mathbf{v}) ;; (\$0\text{-wait} \wedge$
 $\$1\text{-wait} \wedge (\exists \$st' \cdot M)))$
 $\vee (\lceil \text{cmt}_R P \rceil_0 \wedge \lceil \text{cmt}_R Q \rceil_1 \wedge \$\mathbf{v}_{<}' =_u \$\mathbf{v}) ;; (\neg \$0\text{-wait} \wedge \$1\text{-wait}$
 $\wedge (\exists \$st' \cdot M)))$

$\vee (\lceil \text{cmt}_R P \rceil_0 \wedge \lceil \text{cmt}_R Q \rceil_1 \wedge \$\mathbf{v}_{<}' =_u \$\mathbf{v}) ;; (\$0\text{-wait} \wedge \neg \1-wait
 $\wedge (\exists \$st' \cdot M)))$

by (simp add: par-by-merge-alt-def parallel-pericondition-lemma3 seqr-or-distr)

also have $\dots = (\text{pre}_R(P \parallel_{M_R} M Q) \Rightarrow_r ((\lceil \text{peri}_R P \rceil_0 \wedge \lceil \text{peri}_R Q \rceil_1 \wedge \$\mathbf{v}_{<}' =_u \$\mathbf{v}) ;; (\exists \$st' \cdot M)$

$\vee (\lceil \text{post}_R P \rceil_0 \wedge \lceil \text{peri}_R Q \rceil_1 \wedge \$\mathbf{v}_{<}' =_u \$\mathbf{v}) ;; (\exists \$st' \cdot M)$

$\vee (\lceil \text{peri}_R P \rceil_0 \wedge \lceil \text{post}_R Q \rceil_1 \wedge \$\mathbf{v}_{<}' =_u \$\mathbf{v}) ;; (\exists \$st' \cdot M)))$

by (simp add: seqr-right-one-point-true seqr-right-one-point-false cmt_R-def post_R-def peri_R-def usubst unrest assms)

also have $\dots = (\text{pre}_R(P \parallel_{M_R} M Q) \Rightarrow_r \text{peri}_R(P) \parallel_{\exists \$st' \cdot M} \text{peri}_R(Q)$

$\vee \text{post}_R(P) \parallel_{\exists \$st' \cdot M} \text{peri}_R(Q)$

$\vee \text{peri}_R(P) \parallel_{\exists \$st' \cdot M} \text{post}_R(Q))$

by (simp add: par-by-merge-alt-def)

finally show ?thesis .

qed

lemma *parallel-postcondition-lemma1*:

$(\$ok' \wedge P) ;; II_R[\text{true}, \text{false}/\$ok', \$wait'] = P[\text{true}, \text{false}/\$ok', \$wait']$

(is ?lhs = ?rhs)

proof –

have ?lhs = $(\$ok' \wedge P) ;; II[\text{true}, \text{false}/\$ok', \$wait']$

by (rel-blast)

also have $\dots = ?rhs$

by (rel-auto)

finally show ?thesis .

qed

lemma *parallel-postcondition-lemma2*:

assumes M is RDM

shows $(N_0(M))[\text{true}, \text{false}/\$ok', \$wait'] = ((\neg \$0\text{-wait} \wedge \neg \$1\text{-wait}) \wedge M)$

proof –

have $(N_0(M))[\text{true}, \text{false}/\$ok', \$wait'] = ((\neg \$0\text{-wait} \wedge \neg \$1\text{-wait}) \wedge \$tr' \geq_u \$tr_{<} \wedge M)$

by (simp add: usubst unrest nmerge-rd0-def ex-unrest Healthy-if R1m-def assms)

also have $\dots = ((\neg \$0\text{-wait} \wedge \neg \$1\text{-wait}) \wedge M)$

by (metis Healthy-if R1m-def RDM-R1m assms utp-pred-laws.inf-commute)

finally show ?thesis .

qed

lemma *parallel-postcondition* [rdes]:

fixes $M :: ('s, 't :: \text{trace}, 'a) \text{ rsp merge}$

assumes P is SRD Q is SRD M is RDM

shows $\text{post}_R(P \parallel_{M_R(M)} Q) = (\text{pre}_R(P \parallel_{M_R(M)} Q) \Rightarrow_r \text{post}_R(P) \parallel_M \text{post}_R(Q))$

proof –

have $\text{post}_R(P \parallel_{M_R(M)} Q) =$

$(\text{pre}_R(P \parallel_{M_R(M)} Q) \Rightarrow_r \text{cmt}_R P \parallel_{(\$ok' \wedge N_0 M)} ;; II_R[\text{true}, \text{false}/\$ok', \$wait'] \text{cmt}_R Q)$

by (simp add: post-cmt-def parallel-commitment assms usubst unrest SRD-healths)

also have $\dots = (\text{pre}_R(P \parallel_{M_R(M)} Q) \Rightarrow_r \text{cmt}_R P \parallel_{(\neg \$0\text{-wait} \wedge \neg \$1\text{-wait} \wedge M)} \text{cmt}_R Q)$

by (simp add: parallel-postcondition-lemma1 parallel-postcondition-lemma2 assms,
simp add: utp-pred-laws.inf-commute utp-pred-laws.inf-left-commute)

also have $\dots = (\text{pre}_R(P \parallel_{M_R(M)} Q) \Rightarrow_r \text{post}_R P \parallel_M \text{post}_R Q)$

by (simp add: par-by-merge-alt-def seqr-right-one-point-false usubst unrest cmt_R-def post_R-def assms)

finally show ?thesis .

qed

lemma *parallel-precondition-lemma*:

fixes $M :: ('s, 't :: \text{trace}, 'a) \text{ rsp merge}$

assumes P is NSRD Q is NSRD M is RDM

shows $(\neg_r \text{pre}_R(P)) \parallel_{N_0(M)} ;; R1(\text{true}) \text{cmt}_R(Q) =$

$((\neg_r \text{pre}_R P) \parallel_M ;; R1(\text{true}) \text{peri}_R Q \vee (\neg_r \text{pre}_R P) \parallel_M ;; R1(\text{true}) \text{post}_R Q)$

proof –

have $((\neg_r \text{pre}_R(P)) \parallel_{N_0(M)} ;; R1(\text{true}) \text{cmt}_R(Q)) =$

$((\neg_r \text{pre}_R(P)) \parallel_{N_0(M)} ;; R1(\text{true}) (\text{peri}_R(Q) \diamond \text{post}_R(Q)))$

by (simp add: wait'-cond-peri-post-cmt)

also have $\dots = (([\neg_r \text{pre}_R(P)]_0 \wedge [\text{peri}_R(Q) \diamond \text{post}_R(Q)]_1 \wedge \$v_{<}' =_u \$v) ;; N_0(M) ;; R1(\text{true}))$

by (simp add: par-by-merge-alt-def)

also have $\dots = (([\neg_r \text{pre}_R(P)]_0 \wedge [\text{peri}_R(Q)]_1 \triangleleft \$1\text{-wait}' \triangleright [\text{post}_R(Q)]_1 \wedge \$v_{<}' =_u \$v) ;; N_0(M) ;; R1(\text{true}))$

by (simp add: wait'-cond-def alpha)

also have $\dots = ((([\neg_r \text{pre}_R(P)]_0 \wedge [\text{peri}_R(Q)]_1 \wedge \$v_{<}' =_u \$v) \triangleleft \$1\text{-wait}' \triangleright ([\neg_r \text{pre}_R(P)]_0 \wedge [\text{post}_R(Q)]_1 \wedge \$v_{<}' =_u \$v)) ;; N_0(M) ;; R1(\text{true}))$

(is (?P ;; -) = (?Q ;; -))

proof –

have $?P = ?Q$

by (rel-auto)

thus ?thesis by simp

qed

also have $\dots = (([\neg_r \text{pre}_R P]_0 \wedge [\text{peri}_R Q]_1 \wedge \$v_{<}' =_u \$v) [\text{true}/\$1\text{-wait}'] ;; (N_0 M ;; R1 \text{true}) [\text{true}/\$1\text{-wait}'] \vee$

$([\neg_r \text{pre}_R P]_0 \wedge [\text{post}_R Q]_1 \wedge \$v_{<}' =_u \$v) [\text{false}/\$1\text{-wait}'] ;; (N_0 M ;; R1 \text{true}) [\text{false}/\$1\text{-wait}']$

by (simp add: cond-inter-var-split)

also have $\dots = (([\neg_r \text{pre}_R P]_0 \wedge [\text{peri}_R Q]_1 \wedge \$v_{<}' =_u \$v) ;; N_0 M [\text{true}/\$1\text{-wait}'] ;; R1 \text{true} \vee$

$([\neg_r \text{pre}_R P]_0 \wedge [\text{post}_R Q]_1 \wedge \$v_{<}' =_u \$v) ;; N_0 M [\text{false}/\$1\text{-wait}'] ;; R1 \text{true})$

by (simp add: usubst unrest)

also have $\dots = (([\neg_r \text{pre}_R P]_0 \wedge [\text{peri}_R Q]_1 \wedge \$v_{<}' =_u \$v) ;; (\$wait' \wedge M) ;; R1 \text{true} \vee$

$([\neg_r \text{pre}_R P]_0 \wedge [\text{post}_R Q]_1 \wedge \$v_{<}' =_u \$v) ;; (\$wait' =_u \$0\text{-wait} \wedge M) ;; R1 \text{true})$

proof –

have $(\$tr' \geq_u \$tr_{<} \wedge M) = M$

using RDM-R1m[OF assms(3)]

by (simp add: Healthy-def R1m-def conj-comm)

thus ?thesis
 by (simp add: nmerge-rd0-def unrest assms closure ex-unrest usubst)
 qed
 also have ... = (($\lceil \neg_r \text{pre}_R P \rceil_0 \wedge \lceil \text{peri}_R Q \rceil_1 \wedge \$\mathbf{v}_{<} =_u \$\mathbf{v}$) ;; M ;; R1 true \vee
 ($\lceil \neg_r \text{pre}_R P \rceil_0 \wedge \lceil \text{post}_R Q \rceil_1 \wedge \$\mathbf{v}_{<} =_u \$\mathbf{v}$) ;; M ;; R1 true)
 (is ($?P_1 \vee_p ?P_2$) = ($?Q_1 \vee ?Q_2$))
 proof –
 have $?P_1 = (\lceil \neg_r \text{pre}_R P \rceil_0 \wedge \lceil \text{peri}_R Q \rceil_1 \wedge \$\mathbf{v}_{<} =_u \$\mathbf{v}) ;; (M \wedge \$\text{wait}') ;; R1 \text{ true}$
 by (simp add: conj-comm)
 hence 1: $?P_1 = ?Q_1$
 by (simp add: segr-left-one-point-true segr-left-one-point-false add: unrest usubst closure assms)
 have $?P_2 = ((\lceil \neg_r \text{pre}_R P \rceil_0 \wedge \lceil \text{post}_R Q \rceil_1 \wedge \$\mathbf{v}_{<} =_u \$\mathbf{v}) ;; (M \wedge \$\text{wait}') ;; R1 \text{ true} \vee$
 ($\lceil \neg_r \text{pre}_R P \rceil_0 \wedge \lceil \text{post}_R Q \rceil_1 \wedge \$\mathbf{v}_{<} =_u \$\mathbf{v}$) ;; (M $\wedge \neg \$\text{wait}'$) ;; R1 true)
 by (subst segr-bool-split[of left-uvar wait], simp-all add: usubst unrest assms closure conj-comm)
 hence 2: $?P_2 = ?Q_2$
 by (simp add: segr-left-one-point-true segr-left-one-point-false unrest usubst closure assms)
 from 1 2 show ?thesis by simp
 qed
 also have ... = (($\neg_r \text{pre}_R P$) \parallel_M ;; R1(true) $\text{peri}_R Q \vee (\neg_r \text{pre}_R P)$ \parallel_M ;; R1(true) $\text{post}_R Q$)
 by (simp add: par-by-merge-alt-def)
 finally show ?thesis .
 qed

lemma swap-nmerge-rd0:
 $\text{swap}_m ;; N_0(M) = N_0(\text{swap}_m ;; M)$
 by (rel-auto, meson+)

lemma SymMerge-nmerge-rd0 [closure]:
 $M \text{ is SymMerge} \implies N_0(M) \text{ is SymMerge}$
 by (rel-auto, meson+)

lemma swap-merge-rd':
 $\text{swap}_m ;; N_R(M) = N_R(\text{swap}_m ;; M)$
 by (rel-blast)

lemma swap-merge-rd:
 $\text{swap}_m ;; M_R(M) = M_R(\text{swap}_m ;; M)$
 by (simp add: merge-rd-def segr-assoc[THEN sym] swap-merge-rd')

lemma SymMerge-merge-rd [closure]:
 $M \text{ is SymMerge} \implies M_R(M) \text{ is SymMerge}$
 by (simp add: Healthy-def swap-merge-rd)

lemma nmerge-rd1-merge3:
 assumes $M \text{ is RDM}$
 shows $\mathbf{M3}(N_1(M)) = (\$ok' =_u (\$0-ok \wedge \$1-0-ok \wedge \$1-1-ok) \wedge$
 $\$wait' =_u (\$0-wait \vee \$1-0-wait \vee \$1-1-wait) \wedge$
 $\mathbf{M3}(M))$

proof –
 have $\mathbf{M3}(N_1(M)) = \mathbf{M3}(\$ok' =_u (\$0-ok \wedge \$1-ok) \wedge$
 $\$wait' =_u (\$0-wait \vee \$1-wait) \wedge$
 $\$tr_{<} \leq_u \$tr' \wedge$
 $(\exists \{\$0-ok, \$1-ok, \$ok_{<}, \$ok', \$0-wait, \$1-wait, \$wait_{<}, \$wait'\} \cdot \text{RDM}(M)))$
 by (simp add: nmerge-rd1-def nmerge-rd0-def assms Healthy-if)
 also have ... = $\mathbf{M3}(\$ok' =_u (\$0-ok \wedge \$1-ok) \wedge \$wait' =_u (\$0-wait \vee \$1-wait) \wedge \text{RDM}(M))$

by (rel-blast)
 also have ... = $(\$ok' =_u (\$0-ok \wedge \$1-0-ok \wedge \$1-1-ok) \wedge \$wait' =_u (\$0-wait \vee \$1-0-wait \vee \$1-1-wait) \wedge \mathbf{M3}(RDM(M)))$
 by (rel-blast)
 also have ... = $(\$ok' =_u (\$0-ok \wedge \$1-0-ok \wedge \$1-1-ok) \wedge \$wait' =_u (\$0-wait \vee \$1-0-wait \vee \$1-1-wait) \wedge \mathbf{M3}(M))$
 by (simp add: assms Healthy-if)
 finally show ?thesis .
 qed

lemma nmerge-rd-merge3:

$\mathbf{M3}(N_R(M)) = (\exists \$st_{<} \cdot \$\mathbf{v}' =_u \$\mathbf{v}_{<}) \triangleleft \$wait_{<} \triangleright \mathbf{M3}(N_1 M) \triangleleft \$ok_{<} \triangleright (\$tr_{<} \leq_u \$tr')$
 by (rel-blast)

lemma AssocMerge-nmerge-rd:

assumes M is RDM AssocMerge M
 shows AssocMerge($N_R(M)$)

proof –

have 1: $\mathbf{M3}(M) = rotate_m ;; \mathbf{M3}(M)$
 using assms by (simp add: AssocMerge-def)
 have $rotate_m ;; (\mathbf{M3}(N_R(M))) =$
 $rotate_m ;;$
 $((\exists \$st_{<} \cdot \$\mathbf{v}' =_u \$\mathbf{v}_{<}) \triangleleft \$wait_{<} \triangleright$
 $(\$ok' =_u (\$0-ok \wedge \$1-0-ok \wedge \$1-1-ok) \wedge \$wait' =_u (\$0-wait \vee \$1-0-wait \vee$
 $\$1-1-wait) \wedge \mathbf{M3}(M)) \triangleleft \$ok_{<} \triangleright$
 $(\$tr_{<} \leq_u \$tr'))$
 by (simp add: AssocMerge-def nmerge-rd-merge3 nmerge-rd1-merge3 assms)
 also have ... =
 $((\exists \$st_{<} \cdot \$\mathbf{v}' =_u \$\mathbf{v}_{<}) \triangleleft \$wait_{<} \triangleright$
 $(\$ok' =_u (\$0-ok \wedge \$1-0-ok \wedge \$1-1-ok) \wedge \$wait' =_u (\$0-wait \vee \$1-0-wait \vee$
 $\$1-1-wait) \wedge (rotate_m ;; \mathbf{M3}(M))) \triangleleft \$ok_{<} \triangleright$
 $(\$tr_{<} \leq_u \$tr'))$
 by (rel-blast)
 also have ... =
 $((\exists \$st_{<} \cdot \$\mathbf{v}' =_u \$\mathbf{v}_{<}) \triangleleft \$wait_{<} \triangleright$
 $(\$ok' =_u (\$0-ok \wedge \$1-0-ok \wedge \$1-1-ok) \wedge \$wait' =_u (\$0-wait \vee \$1-0-wait \vee$
 $\$1-1-wait) \wedge \mathbf{M3}(M)) \triangleleft \$ok_{<} \triangleright$
 $(\$tr_{<} \leq_u \$tr'))$
 using 1 by auto
 also have ... = $\mathbf{M3}(N_R(M))$
 by (simp add: AssocMerge-def nmerge-rd-merge3 nmerge-rd1-merge3 assms)
 finally show ?thesis
 using AssocMerge-def by blast
 qed

lemma swap-merge-RDM-closed [closure]:

assumes M is RDM
 shows $swap_m ;; M$ is RDM

proof –

have $RDM(swap_m ;; RDM(M)) = (swap_m ;; RDM(M))$
 by (rel-auto)
 thus ?thesis
 by (metis Healthy-def' assms)
 qed

lemma *parallel-precondition*:

fixes $M :: ('s, 't :: \text{trace}, 'a) \text{ rsp merge}$

assumes P is NSRD Q is NSRD M is RDM

shows $\text{pre}_R(P \parallel_{M_R(M)} Q) =$

$$\begin{aligned} & (\neg_r ((\neg_r \text{pre}_R P) \parallel_M ;; R1(\text{true}) \text{peri}_R Q) \wedge \\ & \neg_r ((\neg_r \text{pre}_R P) \parallel_M ;; R1(\text{true}) \text{post}_R Q) \wedge \\ & \neg_r ((\neg_r \text{pre}_R Q) \parallel_{(\text{swap}_m ;; M) ;; R1(\text{true}) \text{peri}_R P} P) \wedge \\ & \neg_r ((\neg_r \text{pre}_R Q) \parallel_{(\text{swap}_m ;; M) ;; R1(\text{true}) \text{post}_R P} P)) \end{aligned}$$

proof –

have $a: (\neg_r \text{pre}_R(P)) \parallel_{N_0(M) ;; R1(\text{true}) \text{cmt}_R(Q)} =$

$$((\neg_r \text{pre}_R P) \parallel_M ;; R1(\text{true}) \text{peri}_R Q \vee (\neg_r \text{pre}_R P) \parallel_M ;; R1(\text{true}) \text{post}_R Q)$$

by (*simp add: parallel-precondition-lemma assms*)

have $b: (\neg_r \text{cmt}_R P \parallel_{N_0 M} ;; R1 \text{ true } (\neg_r \text{pre}_R Q)) =$

$$(\neg_r (\neg_r \text{pre}_R(Q)) \parallel_{N_0(\text{swap}_m ;; M) ;; R1(\text{true}) \text{cmt}_R(P)})$$

by (*simp add: swap-nmerge-rd0[THEN sym] seqr-assoc[THEN sym] par-by-merge-def par-sep-swap*)

have $c: (\neg_r \text{pre}_R(Q)) \parallel_{N_0(\text{swap}_m ;; M) ;; R1(\text{true}) \text{cmt}_R(P)} =$

$$((\neg_r \text{pre}_R Q) \parallel_{(\text{swap}_m ;; M) ;; R1(\text{true}) \text{peri}_R P} P \vee (\neg_r \text{pre}_R Q) \parallel_{(\text{swap}_m ;; M) ;; R1(\text{true}) \text{post}_R P} P)$$

$P)$

by (*simp add: parallel-precondition-lemma closure assms*)

show *?thesis*

by (*simp add: parallel-assm closure assms a b c, rel-auto*)

qed

Weakest Parallel Precondition

definition $\text{wr}_R ::$

$('t :: \text{trace}, 'a) \text{ hrel-rp} \Rightarrow$

$('t :: \text{trace}, 'a) \text{ rp merge} \Rightarrow$

$('t, 'a) \text{ hrel-rp} \Rightarrow$

$('t, 'a) \text{ hrel-rp } (- \text{ wr}_R'(-) - [60, 0, 61] \ 61)$

where [*upred-defs*]: $Q \text{ wr}_R(M) P = (\neg_r ((\neg_r P) \parallel_M ;; R1(\text{true}) Q))$

lemma $\text{wr}_R\text{-R1}$ [*closure*]:

$M \text{ is R1m} \Rightarrow Q \text{ wr}_R(M) P \text{ is R1}$

by (*simp add: wrR-def closure*)

lemma $R2\text{-rea-not}$: $R2(\neg_r P) = (\neg_r R2(P))$

by (*rel-auto*)

lemma $\text{wr}_R\text{-R2-lemma}$:

assumes $P \text{ is R2 } Q \text{ is R2 } M \text{ is R2m}$

shows $((\neg_r P) \parallel_M Q) ;; R1(\text{true}_h) \text{ is R2}$

proof –

have $(\neg_r P) \parallel_M Q \text{ is R2}$

by (*simp add: closure assms*)

thus *?thesis*

by (*simp add: closure*)

qed

lemma $\text{wr}_R\text{-R2}$ [*closure*]:

assumes $P \text{ is R2 } Q \text{ is R2 } M \text{ is R2m}$

shows $Q \text{ wr}_R(M) P \text{ is R2}$

proof –

have $((\neg_r P) \parallel_M Q) ;; R1(true_h) \text{ is } R2$
by (*simp add: wrR-R2-lemma assms*)
thus *?thesis*
by (*simp add: wrR-def wrR-R2-lemma par-by-merge-seq-add closure*)
qed

lemma *wrR-RR [closure]*:
assumes $P \text{ is } RR \ Q \text{ is } RR \ M \text{ is } RDM$
shows $Q \text{ wr}_R(M) \ P \text{ is } RR$
apply (*rule RR-intro*)
apply (*simp-all add: unrest assms closure wrR-def rpred*)
apply (*metis (no-types, lifting) Healthy-def' R1-R2c-commute R1-R2c-is-R2 R1-rea-not RDM-R2m*
RR-implies-R2 assms(1) assms(2) assms(3) par-by-merge-seq-add rea-not-R2-closed
wrR-R2-lemma)
done

lemma *wrR-RC [closure]*:
assumes $P \text{ is } RR \ Q \text{ is } RR \ M \text{ is } RDM$
shows $(Q \text{ wr}_R(M) \ P) \text{ is } RC$
apply (*rule RC-intro*)
apply (*simp add: closure assms*)
apply (*simp add: wrR-def rpred closure assms*)
apply (*simp add: par-by-merge-def seqr-assoc*)
done

lemma *wppR-choice [wp]*: $(P \vee Q) \text{ wr}_R(M) \ R = (P \text{ wr}_R(M) \ R \wedge Q \text{ wr}_R(M) \ R)$
proof –
have $(P \vee Q) \text{ wr}_R(M) \ R =$
 $(\neg_r ((\neg_r R) ;; U0 \wedge (P ;; U1 \vee Q ;; U1) \wedge \$\mathbf{v}_{<}' =_u \$\mathbf{v}) ;; M ;; true_r)$
by (*simp add: wrR-def par-by-merge-def seqr-or-distl*)
also have $\dots = (\neg_r ((\neg_r R) ;; U0 \wedge P ;; U1 \wedge \$\mathbf{v}_{<}' =_u \$\mathbf{v} \vee (\neg_r R) ;; U0 \wedge Q ;; U1 \wedge \$\mathbf{v}_{<}' =_u$
 $\$ \mathbf{v}) ;; M ;; true_r)$
by (*simp add: conj-disj-distr utp-pred-laws.inf-sup-distrib2*)
also have $\dots = (\neg_r (((\neg_r R) ;; U0 \wedge P ;; U1 \wedge \$\mathbf{v}_{<}' =_u \$\mathbf{v}) ;; M ;; true_r \vee$
 $((\neg_r R) ;; U0 \wedge Q ;; U1 \wedge \$\mathbf{v}_{<}' =_u \$\mathbf{v}) ;; M ;; true_r))$
by (*simp add: seqr-or-distl*)
also have $\dots = (P \text{ wr}_R(M) \ R \wedge Q \text{ wr}_R(M) \ R)$
by (*simp add: wrR-def par-by-merge-def*)
finally show *?thesis* .
qed

lemma *wppR-miracle [wp]*: $false \text{ wr}_R(M) \ P = true_r$
by (*simp add: wrR-def*)

lemma *wppR-true [wp]*: $P \text{ wr}_R(M) \ true_r = true_r$
by (*simp add: wrR-def*)

lemma *parallel-precondition-wr [rdes]*:
assumes $P \text{ is } NSRD \ Q \text{ is } NSRD \ M \text{ is } RDM$
shows $pre_R(P \parallel_{M_R(M)} Q) = (peri_R(Q) \text{ wr}_R(M) \ pre_R(P) \wedge post_R(Q) \text{ wr}_R(M) \ pre_R(P) \wedge$
 $peri_R(P) \text{ wr}_R(\text{swap}_m ;; M) \ pre_R(Q) \wedge post_R(P) \text{ wr}_R(\text{swap}_m ;; M) \ pre_R(Q))$
by (*simp add: assms parallel-precondition wrR-def*)

lemma *parallel-rdes-def [rdes-def]*:
assumes $P_1 \text{ is } RC \ P_2 \text{ is } RR \ P_3 \text{ is } RR \ Q_1 \text{ is } RC \ Q_2 \text{ is } RR \ Q_3 \text{ is } RR$

$\$st' \# P_2 \ \$st' \# Q_2$
 $M \text{ is } RDM$
shows $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \parallel_{M_R(M)} \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3) =$
 $\mathbf{R}_s(((Q_1 \Rightarrow_r Q_2) \text{ wr}_R(M) P_1 \wedge (Q_1 \Rightarrow_r Q_3) \text{ wr}_R(M) P_1 \wedge$
 $(P_1 \Rightarrow_r P_2) \text{ wr}_R(\text{swap}_m ;; M) Q_1 \wedge (P_1 \Rightarrow_r P_3) \text{ wr}_R(\text{swap}_m ;; M) Q_1) \vdash$
 $((P_1 \Rightarrow_r P_2) \parallel_{\exists \$st' . M} (Q_1 \Rightarrow_r Q_2) \vee$
 $(P_1 \Rightarrow_r P_3) \parallel_{\exists \$st' . M} (Q_1 \Rightarrow_r Q_2) \vee (P_1 \Rightarrow_r P_2) \parallel_{\exists \$st' . M} (Q_1 \Rightarrow_r Q_3)) \diamond$
 $((P_1 \Rightarrow_r P_3) \parallel_M (Q_1 \Rightarrow_r Q_3))) \text{ (is } ?lhs = ?rhs)$
proof –
have $?lhs = \mathbf{R}_s(\text{pre}_R ?lhs \vdash \text{peri}_R ?lhs \diamond \text{post}_R ?lhs)$
by (*simp add: SRD-reactive-tri-design assms closure*)
also have $\dots = ?rhs$
by (*simp add: rdes closure unrest assms, rel-auto*)
finally show $?thesis$.
qed

lemma *Miracle-parallel-left-zero:*
assumes $P \text{ is } SRD \ M \text{ is } RDM$
shows $\text{Miracle} \parallel_{RM} P = \text{Miracle}$
proof –
have $\text{pre}_R(\text{Miracle} \parallel_{RM} P) = \text{true}_r$
by (*simp add: parallel-assm wait'-cond-idem rdes closure assms*)
moreover hence $\text{cmt}_R(\text{Miracle} \parallel_{RM} P) = \text{false}$
by (*simp add: rdes closure wait'-cond-idem SRD-healths assms*)
ultimately have $\text{Miracle} \parallel_{RM} P = \mathbf{R}_s(\text{true}_r \vdash \text{false})$
by (*metis NSRD-iff SRD-reactive-design-alt assms par-rdes-NSRD sdes-theory-continuous.weak.top-closed*)
thus $?thesis$
by (*simp add: Miracle-def R1-design-R1-pre*)
qed

lemma *Miracle-parallel-right-zero:*
assumes $P \text{ is } SRD \ M \text{ is } RDM$
shows $P \parallel_{RM} \text{Miracle} = \text{Miracle}$
proof –
have $\text{pre}_R(P \parallel_{RM} \text{Miracle}) = \text{true}_r$
by (*simp add: wait'-cond-idem parallel-assm rdes closure assms*)
moreover hence $\text{cmt}_R(P \parallel_{RM} \text{Miracle}) = \text{false}$
by (*simp add: wait'-cond-idem rdes closure SRD-healths assms*)
ultimately have $P \parallel_{RM} \text{Miracle} = \mathbf{R}_s(\text{true}_r \vdash \text{false})$
by (*metis NSRD-iff SRD-reactive-design-alt assms par-rdes-NSRD sdes-theory-continuous.weak.top-closed*)
thus $?thesis$
by (*simp add: Miracle-def R1-design-R1-pre*)
qed

8.1 Example basic merge

definition $\text{BasicMerge} :: (('s, 't :: \text{trace}, \text{unit}) \text{ rsp}) \text{ merge } (N_B) \text{ where}$
 $[\text{upred-defs}]: \text{BasicMerge} = (\$tr_{<} \leq_u \$tr' \wedge \$tr' - \$tr_{<} =_u \$0 - tr - \$tr_{<} \wedge \$tr' - \$tr_{<} =_u \$1 - tr$
 $- \$tr_{<} \wedge \$st' =_u \$st_{<})$

abbreviation $\text{rbasic-par } (- \parallel_B - [85, 86] \ 85) \text{ where}$
 $P \parallel_B Q \equiv P \parallel_{M_R(N_B)} Q$

lemma *BasicMerge-RDM [closure]:* $N_B \text{ is } RDM$
by (*rule RDM-intro, (rel-auto)+*)

lemma *BasicMerge-SymMerge* [closure]:
 N_B is *SymMerge*
 by (rel-auto)

lemma *BasicMerge'-calc*:
 assumes $\$ok' \# P \$wait' \# P \$ok' \# Q \$wait' \# Q$ P is $R2$ Q is $R2$
 shows $P \parallel_{N_B} Q = ((\exists \$st' \cdot P) \wedge (\exists \$st' \cdot Q) \wedge \$st' =_u \$st)$
 using *assms*

proof –

have $P: (\exists \{ \$ok', \$wait' \} \cdot R2(P)) = P$ (**is** $?P' = -$)
 by (simp add: ex-unrest ex-plus Healthy-if *assms*)
 have $Q: (\exists \{ \$ok', \$wait' \} \cdot R2(Q)) = Q$ (**is** $?Q' = -$)
 by (simp add: ex-unrest ex-plus Healthy-if *assms*)
 have $?P' \parallel_{N_B} ?Q' = ((\exists \$st' \cdot ?P') \wedge (\exists \$st' \cdot ?Q') \wedge \$st' =_u \$st)$
 by (simp add: par-by-merge-alt-def, rel-auto, blast+)
 thus *?thesis*
 by (simp add: P Q)

qed

8.2 Simple parallel composition

definition *rea-design-par* ::

$(\text{'s}, \text{'t}::\text{trace}, \text{'}\alpha) \text{ hrel-rsp} \Rightarrow (\text{'s}, \text{'t}, \text{'}\alpha) \text{ hrel-rsp} \Rightarrow (\text{'s}, \text{'t}, \text{'}\alpha) \text{ hrel-rsp}$ (**infixr** \parallel_R 85)
 where [*upred-defs*]: $P \parallel_R Q = \mathbf{R}_s((\text{pre}_R(P) \wedge \text{pre}_R(Q)) \vdash (\text{cmt}_R(P) \wedge \text{cmt}_R(Q)))$

lemma *RHS-design-par*:

assumes
 $\$ok' \# P_1 \$ok' \# P_2$
 shows $\mathbf{R}_s(P_1 \vdash Q_1) \parallel_R \mathbf{R}_s(P_2 \vdash Q_2) = \mathbf{R}_s((P_1 \wedge P_2) \vdash (Q_1 \wedge Q_2))$

proof –

have $\mathbf{R}_s(P_1 \vdash Q_1) \parallel_R \mathbf{R}_s(P_2 \vdash Q_2) =$
 $\mathbf{R}_s(P_1 \llbracket \text{true}, \text{false} / \$ok, \$wait \rrbracket \vdash Q_1 \llbracket \text{true}, \text{false} / \$ok, \$wait \rrbracket) \parallel_R \mathbf{R}_s(P_2 \llbracket \text{true}, \text{false} / \$ok, \$wait \rrbracket \vdash$
 $Q_2 \llbracket \text{true}, \text{false} / \$ok, \$wait \rrbracket)$
 by (simp add: *RHS-design-ok-wait*)

also from *assms*

have ... =

$\mathbf{R}_s((R1 (R2c (P_1)) \wedge R1 (R2c (P_2))) \llbracket \text{true}, \text{false} / \$ok, \$wait \rrbracket \vdash$
 $(R1 (R2c (P_1 \Rightarrow Q_1)) \wedge R1 (R2c (P_2 \Rightarrow Q_2))) \llbracket \text{true}, \text{false} / \$ok, \$wait \rrbracket)$
apply (simp add: *rea-design-par-def* *rea-pre-RHS-design* *rea-cmt-RHS-design* *usubst* *unrest* *assms*)
apply (rule *cong*[of \mathbf{R}_s \mathbf{R}_s], simp)
 using *assms* **apply** (rel-auto)

done

also have ... =

$\mathbf{R}_s((R2c(P_1) \wedge R2c(P_2)) \vdash$
 $(R1 (R2s (P_1 \Rightarrow Q_1)) \wedge R1 (R2s (P_2 \Rightarrow Q_2))))$
 by (*metis* (*no-types*, *hide-lams*) *R1-R2s-R2c* *R1-conj* *R1-design-R1-pre* *RHS-design-ok-wait*)

also have ... =

$\mathbf{R}_s((P_1 \wedge P_2) \vdash (R1 (R2s (P_1 \Rightarrow Q_1)) \wedge R1 (R2s (P_2 \Rightarrow Q_2))))$
 by (simp add: *R2c-R3h-commute* *R2c-and* *R2c-design* *R2c-idem* *R2c-not* *RHS-def*)

also have ... = $\mathbf{R}_s((P_1 \wedge P_2) \vdash ((P_1 \Rightarrow Q_1) \wedge (P_2 \Rightarrow Q_2)))$

by (*metis* (*no-types*, *lifting*) *R1-conj* *R2s-conj* *RHS-design-export-R1* *RHS-design-export-R2s*)

also have ... = $\mathbf{R}_s((P_1 \wedge P_2) \vdash (Q_1 \wedge Q_2))$

by (rule *cong*[of \mathbf{R}_s \mathbf{R}_s], simp, rel-auto)

finally show *?thesis* .

qed

lemma *RHS-tri-design-par*:

assumes $\$ok' \# P_1 \$ok' \# P_2$

shows $\mathbf{R}_s(P_1 \vdash Q_1 \diamond R_1) \parallel_R \mathbf{R}_s(P_2 \vdash Q_2 \diamond R_2) = \mathbf{R}_s((P_1 \wedge P_2) \vdash (Q_1 \wedge Q_2) \diamond (R_1 \wedge R_2))$

by (*simp add: RHS-design-par assms unrest wait'-cond-conj-exchange*)

lemma *RHS-tri-design-par-RR* [*rdes-def*]:

assumes P_1 is RR P_2 is RR

shows $\mathbf{R}_s(P_1 \vdash Q_1 \diamond R_1) \parallel_R \mathbf{R}_s(P_2 \vdash Q_2 \diamond R_2) = \mathbf{R}_s((P_1 \wedge P_2) \vdash (Q_1 \wedge Q_2) \diamond (R_1 \wedge R_2))$

by (*simp add: RHS-tri-design-par unrest assms*)

lemma *RHS-comp-assoc*:

assumes P is NSRD Q is NSRD R is NSRD

shows $(P \parallel_R Q) \parallel_R R = P \parallel_R Q \parallel_R R$

by (*rdes-eq cls: assms*)

end

9 Productive Reactive Designs

theory *utp-rdes-productive*

imports *utp-rdes-parallel*

begin

9.1 Healthiness condition

A reactive design is productive if it strictly increases the trace, whenever it terminates. If it does not terminate, it is also classed as productive.

definition *Productive* :: $(\text{'s}, \text{'t}::\text{trace}, \text{'}\alpha) \text{ hrel-rsp} \Rightarrow (\text{'s}, \text{'t}, \text{'}\alpha) \text{ hrel-rsp}$ **where**

[*upred-defs*]: *Productive*(P) = $P \parallel_R \mathbf{R}_s(\text{true} \vdash \text{true} \diamond (\$tr <_u \$tr'))$

lemma *Productive-RHS-design-form*:

assumes $\$ok' \# P \$ok' \# Q \$ok' \# R$

shows *Productive*($\mathbf{R}_s(P \vdash Q \diamond R)$) = $\mathbf{R}_s(P \vdash Q \diamond (R \wedge \$tr <_u \$tr'))$

using *assms* **by** (*simp add: Productive-def RHS-tri-design-par unrest*)

lemma *Productive-form*:

Productive($\text{SRD}(P)$) = $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond (\text{post}_R(P) \wedge \$tr <_u \$tr'))$

proof –

have *Productive*($\text{SRD}(P)$) = $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P)) \parallel_R \mathbf{R}_s(\text{true} \vdash \text{true} \diamond (\$tr <_u \$tr'))$

by (*simp add: Productive-def SRD-as-reactive-tri-design*)

also have ... = $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond (\text{post}_R(P) \wedge \$tr <_u \$tr'))$

by (*simp add: RHS-tri-design-par unrest*)

finally show *?thesis* .

qed

A reactive design is productive provided that the postcondition, under the precondition, strictly increases the trace.

lemma *Productive-intro*:

assumes P is SRD $(\$tr <_u \$tr') \sqsubseteq (\text{pre}_R(P) \wedge \text{post}_R(P)) \$wait' \# \text{pre}_R(P)$

shows P is *Productive*

proof –

have $P:\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond (\text{post}_R(P) \wedge \$tr <_u \$tr')) = P$

proof –
 have $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P)) = \mathbf{R}_s(\text{pre}_R(P) \vdash (\text{pre}_R(P) \wedge \text{peri}_R(P)) \diamond (\text{pre}_R(P) \wedge \text{post}_R(P)))$
 by (metis (no-types, hide-lams) design-export-pre wait'-cond-conj-exchange wait'-cond-idem)
 also have $\dots = \mathbf{R}_s(\text{pre}_R(P) \vdash (\text{pre}_R(P) \wedge \text{peri}_R(P)) \diamond (\text{pre}_R(P) \wedge (\text{post}_R(P) \wedge \$tr <_u \$tr')))$
 by (metis assms(2) utp-pred-laws.inf.absorb1 utp-pred-laws.inf.assoc)
 also have $\dots = \mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond (\text{post}_R(P) \wedge \$tr <_u \$tr'))$
 by (metis (no-types, hide-lams) design-export-pre wait'-cond-conj-exchange wait'-cond-idem)
 finally show ?thesis
 by (simp add: SRD-reactive-tri-design assms(1))
qed
 thus ?thesis
 by (metis Healthy-def RHS-tri-design-par Productive-def ok'-pre-unrest unrest-true utp-pred-laws.inf-right-idem utp-pred-laws.inf-top-right)
qed

lemma *Productive-post-refines-tr-increase*:

assumes P is SRD P is Productive $\$wait' \nmid \text{pre}_R(P)$
 shows $(\$tr <_u \$tr') \sqsubseteq (\text{pre}_R(P) \wedge \text{post}_R(P))$
proof –
 have $\text{post}_R(P) = \text{post}_R(\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond (\text{post}_R(P) \wedge \$tr <_u \$tr')))$
 by (metis Healthy-def Productive-form assms(1) assms(2))
 also have $\dots = R1(R2c(\text{pre}_R(P) \Rightarrow (\text{post}_R(P) \wedge \$tr <_u \$tr')))$
 by (simp add: rea-post-RHS-design unrest usubst assms, rel-auto)
 also have $\dots = R1((\text{pre}_R(P) \Rightarrow (\text{post}_R(P) \wedge \$tr <_u \$tr')))$
 by (simp add: R2c-impl R2c-preR R2c-postR R2c-and R2c-tr-less-tr' assms)
 also have $(\$tr <_u \$tr') \sqsubseteq (\text{pre}_R(P) \wedge \dots)$
 by (rel-auto)
 finally show ?thesis .
qed

lemma *Continuous-Productive [closure]: Continuous Productive*

by (simp add: Continuous-def Productive-def, rel-auto)

9.2 Reactive design calculations

lemma *preR-Productive [rdes]*:

assumes P is SRD
 shows $\text{pre}_R(\text{Productive}(P)) = \text{pre}_R(P)$
proof –
 have $\text{pre}_R(\text{Productive}(P)) = \text{pre}_R(\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond (\text{post}_R(P) \wedge \$tr <_u \$tr')))$
 by (metis Healthy-def Productive-form assms)
 thus ?thesis
 by (simp add: rea-pre-RHS-design usubst unrest R2c-not R2c-preR R1-preR Healthy-if assms)
qed

lemma *periR-Productive [rdes]*:

assumes P is NSRD
 shows $\text{peri}_R(\text{Productive}(P)) = \text{peri}_R(P)$
proof –
 have $\text{peri}_R(\text{Productive}(P)) = \text{peri}_R(\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond (\text{post}_R(P) \wedge \$tr <_u \$tr')))$
 by (metis Healthy-def NSRD-is-SRD Productive-form assms)
 also have $\dots = R1(R2c(\text{pre}_R P \Rightarrow_r \text{peri}_R P))$
 by (simp add: rea-peri-RHS-design usubst unrest R2c-not assms closure)
 also have $\dots = (\text{pre}_R P \Rightarrow_r \text{peri}_R P)$
 by (simp add: R1-rea-impl R2c-rea-impl R2c-preR R2c-peri-SRD)
qed

R1-peri-SRD assms closure R1-tr-less-tr' R2c-tr-less-tr')

finally show *?thesis*
 by (*simp add: SRD-peri-under-pre assms unrest closure*)
qed

lemma *postR-Productive [rdes]:*
 assumes *P is NSRD*
 shows $\text{post}_R(\text{Productive}(P)) = (\text{pre}_R(P) \Rightarrow_r \text{post}_R(P) \wedge \$tr <_u \$tr')$
proof –
 have $\text{post}_R(\text{Productive}(P)) = \text{post}_R(\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond (\text{post}_R(P) \wedge \$tr <_u \$tr')))$
 by (*metis Healthy-def NSRD-is-SRD Productive-form assms*)
 also have $\dots = R1 (R2c (\text{pre}_R P \Rightarrow_r \text{post}_R P \wedge \$tr' >_u \$tr))$
 by (*simp add: rea-post-RHS-design usubst unrest assms closure*)
 also have $\dots = (\text{pre}_R P \Rightarrow_r \text{post}_R P \wedge \$tr' >_u \$tr)$
 by (*simp add: R1-rea-impl R2c-rea-impl R2c-preR R2c-and R1-extend-conj' R2c-post-SRD R1-post-SRD assms closure R1-tr-less-tr' R2c-tr-less-tr'*)
finally show *?thesis* .
qed

lemma *preR-frame-seq-export:*
 assumes *P is NSRD P is Productive Q is NSRD*
 shows $(\text{pre}_R P \wedge (\text{pre}_R P \wedge \text{post}_R P) ;; Q) = (\text{pre}_R P \wedge (\text{post}_R P ;; Q))$
proof –
 have $(\text{pre}_R P \wedge (\text{post}_R P ;; Q)) = (\text{pre}_R P \wedge ((\text{pre}_R P \Rightarrow_r \text{post}_R P) ;; Q))$
 by (*simp add: SRD-post-under-pre assms closure unrest*)
 also have $\dots = (\text{pre}_R P \wedge (((\neg_r \text{pre}_R P) ;; Q \vee (\text{pre}_R P \Rightarrow_r R1(\text{post}_R P)) ;; Q)))$
 by (*simp add: NSRD-is-SRD R1-post-SRD assms(1) rea-impl-def seqr-or-distl R1-preR Healthy-if*)
 also have $\dots = (\text{pre}_R P \wedge (((\neg_r \text{pre}_R P) ;; Q \vee (\text{pre}_R P \wedge \text{post}_R P) ;; Q)))$
proof –
 have $(\text{pre}_R P \vee \neg_r \text{pre}_R P) = R1 \text{ true}$
 by (*simp add: R1-preR rea-not-or*)
 then show *?thesis*
 by (*metis (no-types, lifting) R1-def conj-comm disj-comm disj-conj-distr rea-impl-def seqr-or-distl utp-pred-laws.inf-top-left utp-pred-laws.sup.left-idem*)
qed
 also have $\dots = (\text{pre}_R P \wedge (((\neg_r \text{pre}_R P) \vee (\text{pre}_R P \wedge \text{post}_R P) ;; Q)))$
 by (*simp add: NSRD-neg-pre-left-zero assms closure SRD-healths*)
 also have $\dots = (\text{pre}_R P \wedge (\text{pre}_R P \wedge \text{post}_R P) ;; Q)$
 by (*rel-blast*)
finally show *?thesis ..*
qed

9.3 Closure laws

lemma *Productive-rdes-intro:*
 assumes $(\$tr <_u \$tr') \sqsubseteq R \$ok' \# P \$ok' \# Q \$ok' \# R \$wait \# P \$wait' \# P$
 shows $(\mathbf{R}_s(P \vdash Q \diamond R))$ is *Productive*
proof (*rule Productive-intro*)
 show $\mathbf{R}_s(P \vdash Q \diamond R)$ is *SRD*
 by (*simp add: RHS-tri-design-is-SRD assms*)

from *assms(1)* **show** $(\$tr' >_u \$tr) \sqsubseteq (\text{pre}_R(\mathbf{R}_s(P \vdash Q \diamond R)) \wedge \text{post}_R(\mathbf{R}_s(P \vdash Q \diamond R)))$
apply (*simp add: rea-pre-RHS-design rea-post-RHS-design usubst assms unrest*)
using *assms(1)* **apply** (*rel-auto*)
apply *fastforce*
done

show $\$wait' \# pre_R (\mathbf{R}_s (P \vdash Q \diamond R))$
by (*simp add: rea-pre-RHS-design rea-post-RHS-design usubst R1-def R2c-def R2s-def assms unrest*)
qed

We use the $R4$ healthiness condition to characterise that the postcondition must extend the trace for a reactive design to be productive.

lemma *Productive-rdes-RR-intro*:
assumes P is RR Q is RR R is RR R is $R4$
shows $(\mathbf{R}_s(P \vdash Q \diamond R))$ is *Productive*
using *assms* **by** (*simp add: Productive-rdes-intro R4-iff-refine unrest*)

lemma *Productive-Miracle* [closure]: *Miracle is Productive*
unfolding *Miracle-tri-def Healthy-def*
by (*subst Productive-RHS-design-form, simp-all add: unrest*)

lemma *Productive-Chaos* [closure]: *Chaos is Productive*
unfolding *Chaos-tri-def Healthy-def*
by (*subst Productive-RHS-design-form, simp-all add: unrest*)

lemma *Productive-intChoice* [closure]:
assumes P is SRD P is *Productive* Q is SRD Q is *Productive*
shows $P \sqcap Q$ is *Productive*
proof –
have $P \sqcap Q =$
 $\mathbf{R}_s(pre_R(P) \vdash peri_R(P) \diamond (post_R(P) \wedge \$tr <_u \$tr')) \sqcap \mathbf{R}_s(pre_R(Q) \vdash peri_R(Q) \diamond (post_R(Q) \wedge$
 $\$tr <_u \$tr'))$
by (*metis Healthy-if Productive-form assms*)
also have $\dots = \mathbf{R}_s((pre_R P \wedge pre_R Q) \vdash (peri_R P \vee peri_R Q) \diamond ((post_R P \wedge \$tr' >_u \$tr) \vee (post_R$
 $Q \wedge \$tr' >_u \$tr)))$
by (*simp add: RHS-tri-design-choice*)
also have $\dots = \mathbf{R}_s((pre_R P \wedge pre_R Q) \vdash (peri_R P \vee peri_R Q) \diamond (((post_R P) \vee (post_R Q)) \wedge \$tr' >_u$
 $\$tr))$
by (*rule cong[of \mathbf{R}_s \mathbf{R}_s], simp, rel-auto*)
also have \dots is *Productive*
by (*simp add: Healthy-def Productive-RHS-design-form unrest*)
finally show *?thesis* .
qed

lemma *Productive-cond-rea* [closure]:
assumes P is SRD P is *Productive* Q is SRD Q is *Productive*
shows $P \triangleleft b \triangleright_R Q$ is *Productive*
proof –
have $P \triangleleft b \triangleright_R Q =$
 $\mathbf{R}_s(pre_R(P) \vdash peri_R(P) \diamond (post_R(P) \wedge \$tr <_u \$tr')) \triangleleft b \triangleright_R \mathbf{R}_s(pre_R(Q) \vdash peri_R(Q) \diamond (post_R(Q)$
 $\wedge \$tr <_u \$tr'))$
by (*metis Healthy-if Productive-form assms*)
also have $\dots = \mathbf{R}_s((pre_R P \triangleleft b \triangleright_R pre_R Q) \vdash (peri_R P \triangleleft b \triangleright_R peri_R Q) \diamond ((post_R P \wedge \$tr' >_u$
 $\$tr) \triangleleft b \triangleright_R (post_R Q \wedge \$tr' >_u \$tr)))$
by (*simp add: cond-srea-form*)
also have $\dots = \mathbf{R}_s((pre_R P \triangleleft b \triangleright_R pre_R Q) \vdash (peri_R P \triangleleft b \triangleright_R peri_R Q) \diamond (((post_R P) \triangleleft b \triangleright_R (post_R$
 $Q)) \wedge \$tr' >_u \$tr))$
by (*rule cong[of \mathbf{R}_s \mathbf{R}_s], simp, rel-auto*)
also have \dots is *Productive*
by (*simp add: Healthy-def Productive-RHS-design-form unrest*)

finally show ?thesis .

qed

lemma *Productive-seq-1* [closure]:

assumes *P is NSRD P is Productive Q is NSRD*

shows *P ;; Q is Productive*

proof –

have $P ;; Q = \mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond (\text{post}_R(P) \wedge \$tr <_u \$tr')) ;; \mathbf{R}_s(\text{pre}_R(Q) \vdash \text{peri}_R(Q) \diamond (\text{post}_R(Q)))$

by (metis *Healthy-def NSRD-is-SRD SRD-reactive-tri-design Productive-form assms(1) assms(2) assms(3)*)

also have $\dots = \mathbf{R}_s((\text{pre}_R P \wedge (\text{post}_R P \wedge \$tr' >_u \$tr) \text{wp}_r \text{pre}_R Q) \vdash (\text{peri}_R P \vee ((\text{post}_R P \wedge \$tr' >_u \$tr) ;; \text{peri}_R Q)) \diamond ((\text{post}_R P \wedge \$tr' >_u \$tr) ;; \text{post}_R Q))$

by (simp add: *RHS-tri-design-composition-wp rpred unrest closure assms wp NSRD-neg-pre-left-zero SRD-healths ex-unrest wp-rea-def disj-upred-def*)

also have $\dots = \mathbf{R}_s((\text{pre}_R P \wedge (\text{post}_R P \wedge \$tr' >_u \$tr) \text{wp}_r \text{pre}_R Q) \vdash (\text{peri}_R P \vee ((\text{post}_R P \wedge \$tr' >_u \$tr) ;; \text{peri}_R Q)) \diamond ((\text{post}_R P \wedge \$tr' >_u \$tr) ;; \text{post}_R Q \wedge \$tr' >_u \$tr))$

proof –

have $((\text{post}_R P \wedge \$tr' >_u \$tr) ;; R1(\text{post}_R Q)) = ((\text{post}_R P \wedge \$tr' >_u \$tr) ;; R1(\text{post}_R Q) \wedge \$tr' >_u \$tr)$

by (rel-auto)

thus ?thesis

by (simp add: *NSRD-is-SRD R1-post-SRD assms*)

qed

also have \dots is Productive

by (rule *Productive-rdes-intro*, simp-all add: *unrest assms closure wp-rea-def*)

finally show ?thesis .

qed

lemma *Productive-seq-2* [closure]:

assumes *P is NSRD Q is NSRD Q is Productive*

shows *P ;; Q is Productive*

proof –

have $P ;; Q = \mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond (\text{post}_R(P))) ;; \mathbf{R}_s(\text{pre}_R(Q) \vdash \text{peri}_R(Q) \diamond (\text{post}_R(Q) \wedge \$tr <_u \$tr'))$

by (metis *Healthy-def NSRD-is-SRD SRD-reactive-tri-design Productive-form assms*)

also have $\dots = \mathbf{R}_s((\text{pre}_R P \wedge \text{post}_R P \text{wp}_r \text{pre}_R Q) \vdash (\text{peri}_R P \vee (\text{post}_R P ;; \text{peri}_R Q)) \diamond (\text{post}_R P ;; (\text{post}_R Q \wedge \$tr' >_u \$tr)))$

by (simp add: *RHS-tri-design-composition-wp rpred unrest closure assms wp NSRD-neg-pre-left-zero SRD-healths ex-unrest wp-rea-def disj-upred-def*)

also have $\dots = \mathbf{R}_s((\text{pre}_R P \wedge \text{post}_R P \text{wp}_r \text{pre}_R Q) \vdash (\text{peri}_R P \vee (\text{post}_R P ;; \text{peri}_R Q)) \diamond (\text{post}_R P ;; (\text{post}_R Q \wedge \$tr' >_u \$tr) \wedge \$tr' >_u \$tr))$

proof –

have $(R1(\text{post}_R P) ;; (\text{post}_R Q \wedge \$tr' >_u \$tr) \wedge \$tr' >_u \$tr) = (R1(\text{post}_R P) ;; (\text{post}_R Q \wedge \$tr' >_u \$tr))$

by (rel-auto)

thus ?thesis

by (simp add: *NSRD-is-SRD R1-post-SRD assms*)

qed

also have \dots is Productive

by (rule *Productive-rdes-intro*, simp-all add: *unrest assms closure wp-rea-def*)

finally show ?thesis .

qed

end

10 Guarded Recursion

```
theory utp-rdes-guarded
  imports utp-rdes-productive
begin
```

10.1 Traces with a size measure

Guarded recursion relies on our ability to measure the trace's size, in order to see if it is decreasing on each iteration. Thus, we here equip the trace algebra with the *size* function that provides this.

```
class size-trace = trace + size +
  assumes
```

```
    size-zero: size 0 = 0 and
    size-nzero: s > 0  $\implies$  size(s) > 0 and
    size-plus: size (s + t) = size(s) + size(t)
```

— These axioms may be stronger than necessary. In particular, $(0::'a) < ?s \implies 0 < \text{size } ?s$ requires that a non-empty trace have a positive size. But this may not be the case with all trace models and is possibly more restrictive than necessary. In future we will explore weakening.

```
begin
```

```
lemma size-mono: s  $\leq$  t  $\implies$  size(s)  $\leq$  size(t)
  by (metis le-add1 local.diff-add-cancel-left' local.size-plus)
```

```
lemma size-strict-mono: s < t  $\implies$  size(s) < size(t)
  by (metis cancel-ab-semigroup-add-class.add-diff-cancel-left' local.diff-add-cancel-left' local.less-iff local.minus-gr-zero-iff local.size-nzero local.size-plus zero-less-diff)
```

```
lemma trace-strict-prefixE: xs < ys  $\implies$  ( $\bigwedge$ zs.  $\llbracket$  ys = xs + zs; size(zs) > 0  $\rrbracket \implies$  thesis)  $\implies$  thesis
  by (metis local.diff-add-cancel-left' local.less-iff local.minus-gr-zero-iff local.size-nzero)
```

```
lemma size-minus-trace: y  $\leq$  x  $\implies$  size(x - y) = size(x) - size(y)
  by (metis diff-add-inverse local.diff-add-cancel-left' local.size-plus)
```

end

Both natural numbers and lists are measurable trace algebras.

```
instance nat :: size-trace
  by (intro-classes, simp-all)
```

```
instance list :: (type) size-trace
  by (intro-classes, simp-all add: zero-list-def less-list-def' plus-list-def prefix-length-less)
```

```
syntax
  -usize      :: logic  $\Rightarrow$  logic (sizeu'(-))
```

```
translations
  sizeu(t) == CONST uop CONST size t
```

10.2 Guardedness

definition $gvert :: ((t::size\text{-}trace, 'a) \text{ } rp \times (t, 'a) \text{ } rp) \text{ } chain \text{ } \mathbf{where}$
 $[upred\text{-}defs]: gvert(n) \equiv (\$tr \leq_u \$tr' \wedge size_u(\&tt) <_u \ll n \gg)$

lemma $gvert\text{-}chain: chain \text{ } gvert$
apply ($simp \text{ } add: chain\text{-}def, safe$)
apply ($rel\text{-}simp$)
apply ($rel\text{-}simp$)
done

lemma $gvert\text{-}limit: \sqcap (\text{range } gvert) = (\$tr \leq_u \$tr')$
by ($rel\text{-}auto$)

definition $Guarded :: ((t::size\text{-}trace, 'a) \text{ } hrel\text{-}rp \Rightarrow (t, 'a) \text{ } hrel\text{-}rp) \Rightarrow bool \text{ } \mathbf{where}$
 $[upred\text{-}defs]: Guarded(F) = (\forall X \text{ } n. (F(X) \wedge gvert(n+1)) = (F(X \wedge gvert(n)) \wedge gvert(n+1)))$

lemma $GuardedI: \llbracket \bigwedge X \text{ } n. (F(X) \wedge gvert(n+1)) = (F(X \wedge gvert(n)) \wedge gvert(n+1)) \rrbracket \Longrightarrow Guarded \text{ } F$
by ($simp \text{ } add: Guarded\text{-}def$)

Guarded reactive designs yield unique fixed-points.

theorem $guarded\text{-}fp\text{-}uniq:$
assumes $mono \text{ } F \text{ } F \in \llbracket id \rrbracket_H \rightarrow \llbracket SRD \rrbracket_H \text{ } Guarded \text{ } F$
shows $\mu \text{ } F = \nu \text{ } F$
proof –
have $constr \text{ } F \text{ } gvert$
using $assms$
by ($auto \text{ } simp \text{ } add: constr\text{-}def \text{ } gvert\text{-}chain \text{ } Guarded\text{-}def \text{ } tcontr\text{-}alt\text{-}def'$)
hence $(\$tr \leq_u \$tr' \wedge \mu \text{ } F) = (\$tr \leq_u \$tr' \wedge \nu \text{ } F)$
apply ($rule \text{ } constr\text{-}fp\text{-}uniq$)
apply ($simp \text{ } add: assms$)
using $gvert\text{-}limit$ **apply** $blast$
done
moreover **have** $(\$tr \leq_u \$tr' \wedge \mu \text{ } F) = \mu \text{ } F$
proof –
have $\mu \text{ } F \text{ } is \text{ } R1$
by ($rule \text{ } SRD\text{-}healths(1), rule \text{ } Healthy\text{-}\mu, simp\text{-}all \text{ } add: assms$)
thus $?thesis$
by ($metis \text{ } Healthy\text{-}def \text{ } R1\text{-}def \text{ } conj\text{-}comm$)
qed
moreover **have** $(\$tr \leq_u \$tr' \wedge \nu \text{ } F) = \nu \text{ } F$
proof –
have $\nu \text{ } F \text{ } is \text{ } R1$
by ($rule \text{ } SRD\text{-}healths(1), rule \text{ } Healthy\text{-}\nu, simp\text{-}all \text{ } add: assms$)
thus $?thesis$
by ($metis \text{ } Healthy\text{-}def \text{ } R1\text{-}def \text{ } conj\text{-}comm$)
qed
ultimately **show** $?thesis$
by ($simp$)
qed

lemma $Guarded\text{-}const \text{ } [closure]: Guarded \text{ } (\lambda X. P)$
by ($simp \text{ } add: Guarded\text{-}def$)

lemma $UINF\text{-}Guarded \text{ } [closure]:$
assumes $\bigwedge P. P \in A \Longrightarrow Guarded \text{ } P$

```

shows Guarded ( $\lambda X. \sqcap P \in A \cdot P(X)$ )
proof (rule GuardedI)
  fix  $X\ n$ 
  have  $\bigwedge Y. ((\sqcap P \in A \cdot P\ Y) \wedge \text{gvr}(n+1)) = ((\sqcap P \in A \cdot (P\ Y \wedge \text{gvr}(n+1))) \wedge \text{gvr}(n+1))$ 
  proof –
    fix  $Y$ 
    let  $?lhs = ((\sqcap P \in A \cdot P\ Y) \wedge \text{gvr}(n+1))$  and  $?rhs = ((\sqcap P \in A \cdot (P\ Y \wedge \text{gvr}(n+1))) \wedge \text{gvr}(n+1))$ 
    have  $a: ?lhs \llbracket \text{false} / \$ok \rrbracket = ?rhs \llbracket \text{false} / \$ok \rrbracket$ 
      by (rel-auto)
    have  $b: ?lhs \llbracket \text{true} / \$ok \rrbracket \llbracket \text{true} / \$wait \rrbracket = ?rhs \llbracket \text{true} / \$ok \rrbracket \llbracket \text{true} / \$wait \rrbracket$ 
      by (rel-auto)
    have  $c: ?lhs \llbracket \text{true} / \$ok \rrbracket \llbracket \text{false} / \$wait \rrbracket = ?rhs \llbracket \text{true} / \$ok \rrbracket \llbracket \text{false} / \$wait \rrbracket$ 
      by (rel-auto)
    show  $?lhs = ?rhs$ 
      using  $a\ b\ c$ 
      by (rule-tac bool-eq-splitI[of in-var ok], simp, rule-tac bool-eq-splitI[of in-var wait], simp-all)
  qed
  moreover have  $((\sqcap P \in A \cdot (P\ X \wedge \text{gvr}(n+1))) \wedge \text{gvr}(n+1)) = ((\sqcap P \in A \cdot (P\ (X \wedge \text{gvr}(n)) \wedge \text{gvr}(n+1))) \wedge \text{gvr}(n+1))$ 
  proof –
    have  $(\sqcap P \in A \cdot (P\ X \wedge \text{gvr}(n+1))) = (\sqcap P \in A \cdot (P\ (X \wedge \text{gvr}(n)) \wedge \text{gvr}(n+1)))$ 
    proof (rule UINF-cong)
      fix  $P$  assume  $P \in A$ 
      thus  $(P\ X \wedge \text{gvr}(n+1)) = (P\ (X \wedge \text{gvr}(n)) \wedge \text{gvr}(n+1))$ 
        using Guarded-def assms by blast
    qed
    thus  $?thesis$  by simp
  qed
  ultimately show  $((\sqcap P \in A \cdot P\ X) \wedge \text{gvr}(n+1)) = ((\sqcap P \in A \cdot (P\ (X \wedge \text{gvr}(n)))) \wedge \text{gvr}(n+1))$ 
    by simp
qed

```

```

lemma intChoice-Guarded [closure]:
  assumes Guarded  $P$  Guarded  $Q$ 
  shows Guarded ( $\lambda X. P(X) \sqcap Q(X)$ )
proof –
  have Guarded ( $\lambda X. \sqcap F \in \{P, Q\} \cdot F(X)$ )
    by (rule UINF-Guarded, auto simp add: assms)
  thus  $?thesis$ 
    by (simp)
qed

```

```

lemma cond-srea-Guarded [closure]:
  assumes Guarded  $P$  Guarded  $Q$ 
  shows Guarded ( $\lambda X. P(X) \triangleleft b \triangleright_R Q(X)$ )
  using assms by (rel-auto)

```

A tail recursive reactive design with a productive body is guarded.

```

lemma Guarded-if-Productive [closure]:
  fixes  $P :: ('s, 't::\text{size-trace}, 'a) \text{hrel-rsp}$ 
  assumes  $P$  is NSRD  $P$  is Productive
  shows Guarded ( $\lambda X. P ;; \text{SRD}(X)$ )
proof (clarsimp simp add: Guarded-def)

```

— We split the proof into three cases corresponding to valuations for *ok*, *wait*, and *wait'* respectively.

```

fix  $X\ n$ 

```


have $a:(P ;; \text{SRD}(X) \wedge \text{gvt}(\text{Suc } n))\llbracket \text{false}/\$ok \rrbracket =$
 $(P ;; \text{SRD}(X \wedge \text{gvt } n) \wedge \text{gvt}(\text{Suc } n))\llbracket \text{false}/\$ok \rrbracket$
by (*simp add: usubst closure SRD-left-zero-1 assms*)
have $b:(P ;; \text{SRD}(X) \wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}/\$ok \rrbracket\llbracket \text{true}/\$wait \rrbracket =$
 $((P ;; \text{SRD}(X \wedge \text{gvt } n) \wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}/\$ok \rrbracket)\llbracket \text{true}/\$wait \rrbracket$
by (*simp add: usubst closure SRD-left-zero-2 assms*)
have $c:(P ;; \text{SRD}(X) \wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}/\$ok \rrbracket\llbracket \text{false}/\$wait \rrbracket =$
 $((P ;; \text{SRD}(X \wedge \text{gvt } n) \wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}/\$ok \rrbracket)\llbracket \text{false}/\$wait \rrbracket$
proof –
have $1:(P\llbracket \text{true}/\$wait' \rrbracket ;; (\text{SRD } X)\llbracket \text{true}/\$wait \rrbracket \wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket =$
 $(P\llbracket \text{true}/\$wait' \rrbracket ;; (\text{SRD } (X \wedge \text{gvt } n))\llbracket \text{true}/\$wait \rrbracket \wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket$
by (*metis (no-types, lifting) Healthy-def R3h-wait-true SRD-healths(3) SRD-idem*)
have $2:(P\llbracket \text{false}/\$wait' \rrbracket ;; (\text{SRD } X)\llbracket \text{false}/\$wait \rrbracket \wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket =$
 $(P\llbracket \text{false}/\$wait' \rrbracket ;; (\text{SRD } (X \wedge \text{gvt } n))\llbracket \text{false}/\$wait \rrbracket \wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket$
proof –
have $\text{exp}:\bigwedge Y::('s, 't, 'a) \text{ hrel-rsp. } (P\llbracket \text{false}/\$wait' \rrbracket ;; (\text{SRD } Y)\llbracket \text{false}/\$wait \rrbracket \wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket$
 $=$
 $((\neg_r \text{pre}_R P) ;; (\text{SRD}(Y))\llbracket \text{false}/\$wait \rrbracket \vee (\text{post}_R P \wedge \$tr' >_u \$tr) ;; (\text{SRD}$
 $Y)\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket))$
 $\wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket$
proof –
fix $Y :: ('s, 't, 'a) \text{ hrel-rsp}$

have $(P\llbracket \text{false}/\$wait' \rrbracket ;; (\text{SRD } Y)\llbracket \text{false}/\$wait \rrbracket \wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket =$
 $((\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond (\text{post}_R(P) \wedge \$tr <_u \$tr')))\llbracket \text{false}/\$wait' \rrbracket ;; (\text{SRD } Y)\llbracket \text{false}/\$wait \rrbracket$
 $\wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket$
by (*metis (no-types) Healthy-def Productive-form assms(1) assms(2) NSRD-is-SRD*)
also have $\dots =$
 $((R1(R2c(\text{pre}_R(P) \Rightarrow (\$ok' \wedge \text{post}_R(P) \wedge \$tr <_u \$tr')))\llbracket \text{false}/\$wait' \rrbracket ;; (\text{SRD } Y)\llbracket \text{false}/\$wait \rrbracket$
 $\wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket$
by (*simp add: RHS-def R1-def R2c-def R2s-def R3h-def RD1-def RD2-def usubst unrest assms*
closure design-def)
also have $\dots =$
 $((\neg_r \text{pre}_R(P) \vee (\$ok' \wedge \text{post}_R(P) \wedge \$tr <_u \$tr'))\llbracket \text{false}/\$wait' \rrbracket ;; (\text{SRD } Y)\llbracket \text{false}/\$wait \rrbracket$
 $\wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket$
by (*simp add: impl-alt-def R2c-disj R1-disj R2c-not assms closure R2c-and*
R2c-preR rea-not-def R1-extend-conj' R2c-ok' R2c-post-SRD R1-tr-less-tr' R2c-tr-less-tr')
also have $\dots =$
 $((\neg_r \text{pre}_R P) ;; (\text{SRD}(Y))\llbracket \text{false}/\$wait \rrbracket \vee (\$ok' \wedge \text{post}_R P \wedge \$tr' >_u \$tr) ;; (\text{SRD}$
 $Y)\llbracket \text{false}/\$wait \rrbracket)) \wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket$
by (*simp add: usubst unrest assms closure seqr-or-distl NSRD-neg-pre-left-zero*)
also have $\dots =$
 $((\neg_r \text{pre}_R P) ;; (\text{SRD}(Y))\llbracket \text{false}/\$wait \rrbracket \vee (\text{post}_R P \wedge \$tr' >_u \$tr) ;; (\text{SRD } Y)\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket))$
 $\wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket$
proof –
have $(\$ok' \wedge \text{post}_R P \wedge \$tr' >_u \$tr) ;; (\text{SRD } Y)\llbracket \text{false}/\$wait \rrbracket =$
 $((\text{post}_R P \wedge \$tr' >_u \$tr) \wedge \$ok' =_u \text{true}) ;; (\text{SRD } Y)\llbracket \text{false}/\$wait \rrbracket$
by (*rel-blast*)
also have $\dots = (\text{post}_R P \wedge \$tr' >_u \$tr)\llbracket \text{true}/\$ok' \rrbracket ;; (\text{SRD } Y)\llbracket \text{false}/\$wait \rrbracket\llbracket \text{true}/\$ok \rrbracket$
using *seqr-left-one-point*[*of ok (post_R P \wedge \\$tr' >_u \\$tr) True (SRD Y)\llbracket \text{false}/\\$wait \rrbracket*]
by (*simp add: true-alt-def[THEN sym]*)
finally show *?thesis* **by** (*simp add: usubst unrest*)
qed
finally
show $(P\llbracket \text{false}/\$wait' \rrbracket ;; (\text{SRD } Y)\llbracket \text{false}/\$wait \rrbracket \wedge \text{gvt}(\text{Suc } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket =$

$$(((\neg_r \text{pre}_R P) ;; (\text{SRD}(Y))\llbracket \text{false}/\$wait \rrbracket \vee (\text{post}_R P \wedge \$tr' >_u \$tr) ;; (\text{SRD}(Y)\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket))$$

$$\wedge \text{gvt} (Suc\ n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket .$$

qed

have 1: $((\text{post}_R P \wedge \$tr' >_u \$tr) ;; (\text{SRD } X)\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket \wedge \text{gvt} (Suc\ n)) =$
 $((\text{post}_R P \wedge \$tr' >_u \$tr) ;; (\text{SRD } (X \wedge \text{gvt } n))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket \wedge \text{gvt} (Suc\ n))$

apply (*rel-auto*)

apply (*rename-tac tr st more ok wait tr' st' more' tr₀ st₀ more₀ ok'*)

apply (*rule-tac x=tr₀ in exI, rule-tac x=st₀ in exI, rule-tac x=more₀ in exI*)

apply (*simp*)

apply (*erule trace-strict-prefixE*)

apply (*rename-tac tr st ref ok wait tr' st' ref' tr₀ st₀ ref₀ ok' zs*)

apply (*rule-tac x=False in exI*)

apply (*simp add: size-minus-trace*)

apply (*subgoal-tac size(tr) < size(tr₀)*)

apply (*simp add: less-diff-conv2 size-mono*)

using *size-strict-mono* **apply** *blast*

apply (*rename-tac tr st more ok wait tr' st' more' tr₀ st₀ more₀ ok'*)

apply (*rule-tac x=tr₀ in exI, rule-tac x=st₀ in exI, rule-tac x=more₀ in exI*)

apply (*simp*)

apply (*erule trace-strict-prefixE*)

apply (*rename-tac tr st more ok wait tr' st' more' tr₀ st₀ more₀ ok' zs*)

apply (*auto simp add: size-minus-trace*)

apply (*subgoal-tac size(tr) < size(tr₀)*)

apply (*simp add: less-diff-conv2 size-mono*)

using *size-strict-mono* **apply** *blast*

done

have 2: $(\neg_r \text{pre}_R P) ;; (\text{SRD } X)\llbracket \text{false}/\$wait \rrbracket = (\neg_r \text{pre}_R P) ;; (\text{SRD}(X \wedge \text{gvt } n))\llbracket \text{false}/\$wait \rrbracket$

by (*simp add: NSRD-neg-pre-left-zero closure assms SRD-healths*)

show *?thesis*

by (*simp add: exp 1 2 utp-pred-laws.inf-sup-distrib2*)

qed

show *?thesis*

proof –

have $(P ;; (\text{SRD } X) \wedge \text{gvt } (n+1))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket =$

$((P\llbracket \text{true}/\$wait' \rrbracket ;; (\text{SRD } X)\llbracket \text{true}/\$wait \rrbracket \wedge \text{gvt } (n+1))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket \vee$

$(P\llbracket \text{false}/\$wait' \rrbracket ;; (\text{SRD } X)\llbracket \text{false}/\$wait \rrbracket \wedge \text{gvt } (n+1))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket)$

by (*subst seqr-bool-split[of wait], simp-all add: usubst utp-pred-laws.distrib(4)*)

also

have $\dots = ((P\llbracket \text{true}/\$wait' \rrbracket ;; (\text{SRD } (X \wedge \text{gvt } n))\llbracket \text{true}/\$wait \rrbracket \wedge \text{gvt } (n+1))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket$

\vee

$(P\llbracket \text{false}/\$wait' \rrbracket ;; (\text{SRD } (X \wedge \text{gvt } n))\llbracket \text{false}/\$wait \rrbracket \wedge \text{gvt } (n+1))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket)$

by (*simp add: 1 2*)

also

have $\dots = ((P\llbracket \text{true}/\$wait' \rrbracket ;; (\text{SRD } (X \wedge \text{gvt } n))\llbracket \text{true}/\$wait \rrbracket \vee$

$P\llbracket \text{false}/\$wait' \rrbracket ;; (\text{SRD } (X \wedge \text{gvt } n))\llbracket \text{false}/\$wait \rrbracket) \wedge \text{gvt } (n+1))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket$

by (*simp add: usubst utp-pred-laws.distrib(4)*)

also have $\dots = (P ;; (\text{SRD } (X \wedge \text{gvt } n)) \wedge \text{gvt } (n+1))\llbracket \text{true}, \text{false}/\$ok, \$wait \rrbracket$

by (*subst seqr-bool-split[of wait], simp-all add: usubst*)

finally show *?thesis* **by** (*simp add: usubst*)

```

qed

qed
show (P ;; SRD(X) ∧ gvirt (Suc n)) = (P ;; SRD(X ∧ gvirt n) ∧ gvirt (Suc n))
  apply (rule-tac bool-eq-splitI[of in-var ok])
  apply (simp-all add: a)
  apply (rule-tac bool-eq-splitI[of in-var wait])
  apply (simp-all add: b c)
done
qed

```

10.3 Tail recursive fixed-point calculations

declare upred-semiring.power-Suc [simp]

```

lemma mu-csp-form-1 [rdes]:
  fixes P :: ('s, 't::size-trace, 'α) hrel-rsp
  assumes P is NSRD P is Productive
  shows (μ X · P ;; SRD(X)) = (⋀ i. P ^ (i+1)) ;; Miracle
proof -
  have 1: Continuous (λX. P ;; SRD X)
    using SRD-Continuous
  by (clarsimp simp add: Continuous-def seq-SUP-distl[THEN sym], drule-tac x=A in spec, simp)
  have 2: (λX. P ;; SRD X) ∈ ⟦id⟧H → ⟦SRD⟧H
    by (blast intro: funcsetI closure assms)
  with 1 2 have (μ X · P ;; SRD(X)) = (ν X · P ;; SRD(X))
    by (simp add: guarded-fp-uniq Guarded-if-Productive[OF assms] funcsetI closure)
  also have ... = (⋀ i. ((λX. P ;; SRD X) ^^ i) false)
    by (simp add: sup-continuous-lfp 1 sup-continuous-Continuous false-upred-def)
  also have ... = ((λX. P ;; SRD X) ^^ 0) false ∧ (⋀ i. ((λX. P ;; SRD X) ^^ (i+1)) false)
    by (subst Sup-power-expand, simp)
  also have ... = (⋀ i. ((λX. P ;; SRD X) ^^ (i+1)) false)
    by (simp)
  also have ... = (⋀ i. P ^ (i+1)) ;; Miracle
proof (rule SUP-cong, simp-all)
  fix i
  show P ;; SRD (((λX. P ;; SRD X) ^^ i) false) = (P ;; P ^ i) ;; Miracle
proof (induct i)
  case 0
  then show ?case
    by (simp, metis srdes-hcond-def srdes-theory-continuous.healthy-top)
next
  case (Suc i)
  then show ?case
    by (simp add: Healthy-if NSRD-is-SRD SRD-power-comp SRD-seqr-closure assms(1) seqr-assoc[THEN sym] srdes-theory-continuous.weak.top-closed)
qed
qed
also have ... = (⋀ i. P ^ (i+1)) ;; Miracle
  by (simp add: seq-Sup-distr)
finally show ?thesis
  by (simp add: UINF-as-Sup[THEN sym])
qed

```

```

lemma mu-csp-form-NSRD [closure]:
  fixes P :: ('s, 't::size-trace, 'α) hrel-rsp

```

assumes P is NSRD P is Productive
shows $(\mu X \cdot P ;; SRD(X))$ is NSRD
by (*simp add: mu-csp-form-1 assms closure*)

lemma *mu-csp-form-1'*:
fixes $P :: ('s, 't::size-trace, 'a) hrel-rsp$
assumes P is NSRD P is Productive
shows $(\mu X \cdot P ;; SRD(X)) = (P ;; P^*) ;; Miracle$
proof –
have $(\mu X \cdot P ;; SRD(X)) = (\bigcap i \in UNIV \cdot P ;; P \hat{~} i) ;; Miracle$
by (*simp add: mu-csp-form-1 assms closure ustar-def*)
also have $\dots = (P ;; P^*) ;; Miracle$
by (*simp only: seq-UINF-distl[THEN sym], simp add: ustar-def*)
finally show ?thesis .
qed

declare *upred-semiring.power-Suc* [*simp del*]

end

11 Reactive Design Programs

theory *utp-rdes-prog*
imports
utp-rdes-normal
utp-rdes-tactics
utp-rdes-parallel
utp-rdes-guarded
UTP-KAT.utp-kleene
begin

11.1 State substitution

lemma *srd-subst-RHS-tri-design* [*usubst*]:

$$[\sigma]_{S\sigma} \dagger \mathbf{R}_s(P \vdash Q \diamond R) = \mathbf{R}_s([\sigma]_{S\sigma} \dagger P \vdash ([\sigma]_{S\sigma} \dagger Q) \diamond ([\sigma]_{S\sigma} \dagger R))$$

by (*rel-auto*)

lemma *srd-subst-SRD-closed* [*closure*]:
assumes P is SRD
shows $[\sigma]_{S\sigma} \dagger P$ is SRD
proof –
have $SRD([\sigma]_{S\sigma} \dagger (SRD P)) = [\sigma]_{S\sigma} \dagger (SRD P)$
by (*rel-auto*)
thus ?thesis
by (*metis Healthy-def assms*)
qed

lemma *preR-srd-subst* [*rdes*]:

$$pre_R([\sigma]_{S\sigma} \dagger P) = [\sigma]_{S\sigma} \dagger pre_R(P)$$

by (*rel-auto*)

lemma *periR-srd-subst* [*rdes*]:

$$peri_R([\sigma]_{S\sigma} \dagger P) = [\sigma]_{S\sigma} \dagger peri_R(P)$$

by (*rel-auto*)

lemma *postR-srd-subst* [rdes]:
 $\text{post}_R(\lceil \sigma \rceil_{S\sigma} \dagger P) = \lceil \sigma \rceil_{S\sigma} \dagger \text{post}_R(P)$
by (*rel-auto*)

lemma *srd-subst-NSRD-closed* [closure]:
assumes P is NSRD
shows $\lceil \sigma \rceil_{S\sigma} \dagger P$ is NSRD
by (*rule NSRD-RC-intro, simp-all add: closure rdes assms unrest*)

11.2 Assignment

definition *assigns-srd* :: 's usubst \Rightarrow ('s, 't::trace, 'α) hrel-rsp ($\langle \cdot \rangle_R$) **where**
[upred-defs]: *assigns-srd* $\sigma = \mathbf{R}_s(\text{true} \vdash (\$tr' =_u \$tr \wedge \neg \$wait' \wedge \lceil \langle \sigma \rangle_a \rceil_S \wedge \$\Sigma_S' =_u \$\Sigma_S))$

syntax

-*assign-srd* :: svids \Rightarrow uexprs \Rightarrow logic ($'(-)' :=_R '(-)'$)
-*assign-srd* :: svids \Rightarrow uexprs \Rightarrow logic (**infixr** :=_R 62)

translations

-*assign-srd* $xs\ vs \Rightarrow \text{CONST assigns-srd } (-\text{mk-usubst } (\text{CONST id})\ xs\ vs)$
-*assign-srd* $x\ v \leq \text{CONST assigns-srd } (\text{CONST subst-upd } (\text{CONST id})\ x\ v)$
-*assign-srd* $x\ v \leq -\text{assign-srd } (-\text{spvar } x)\ v$
 $x, y :=_R u, v \leq \text{CONST assigns-srd } (\text{CONST subst-upd } (\text{CONST subst-upd } (\text{CONST id})\ (\text{CONST svar } x)\ u)\ (\text{CONST svar } y)\ v)$

lemma *assigns-srd-RHS-tri-des* [rdes-def]:
 $\langle \sigma \rangle_R = \mathbf{R}_s(\text{true}_r \vdash \text{false} \diamond \langle \sigma \rangle_r)$
by (*rel-auto*)

lemma *assigns-srd-NSRD-closed* [closure]: $\langle \sigma \rangle_R$ is NSRD
by (*simp add: rdes-def closure unrest*)

lemma *preR-assigs-srd* [rdes]: $\text{pre}_R(\langle \sigma \rangle_R) = \text{true}_r$
by (*simp add: rdes-def rdes closure*)

lemma *periR-assigs-srd* [rdes]: $\text{peri}_R(\langle \sigma \rangle_R) = \text{false}$
by (*simp add: rdes-def rdes closure*)

lemma *postR-assigs-srd* [rdes]: $\text{post}_R(\langle \sigma \rangle_R) = \langle \sigma \rangle_r$
by (*simp add: rdes-def rdes closure rpred*)

lemma *taut-eq-impl-property*:
 $\llbracket \text{vwb-lens } x; \$x \# P \rrbracket \Rightarrow '(\$x =_u \llbracket v \rrbracket \wedge Q) \Rightarrow P' = 'Q \llbracket \llbracket v \rrbracket / \$x \rrbracket \Rightarrow P'$
by (*rel-auto, meson mwb-lens-weak vwb-lens-mwb weak-lens.put-get*)

lemma *st-subst-taut-impl*:
assumes $\text{vwb-lens } x\ \$st:x \# Q\ P$ is RR Q is RR
shows $'[\&x \mapsto_s \llbracket k \rrbracket] \dagger_S P \Rightarrow Q' = '([\&x =_u \llbracket k \rrbracket]_{S<} \wedge P \Rightarrow Q' \text{ (is ?lhs = ?rhs)})$

proof –

have ?lhs = $'P \llbracket \llbracket k \rrbracket / \$st:x \rrbracket \Rightarrow Q'$
by (*simp add: usubst-st-lift-def alpha usubst*)
also have ... = $'(\$st:x =_u \llbracket k \rrbracket) \wedge RR(P) \Rightarrow RR(Q)'$
by (*simp add: Healthy-if assms taut-eq-impl-property*)
also have ... = $'([\&x =_u \llbracket k \rrbracket]_{S<} \wedge RR(P) \Rightarrow RR(Q)'$
by (*rel-blast*)
finally show ?thesis **by** (*simp add: assms Healthy-if*)

qed

The following law explains how to refine a program Q when it is first initialised by an assignment. Would be good if it could be generalised to a more general precondition.

lemma *AssignR-init-refine-intro*:

assumes
 $vwb\text{-}lens\ x\ \$st:x\ \# P_2\ \$st:x\ \# P_3$
 $P_2\ is\ RR\ P_3\ is\ RR\ Q\ is\ NSRD$
 $\mathbf{R}_s([\&x =_u \ll k \gg]_{S<} \vdash P_2 \diamond P_3) \sqsubseteq Q$
shows $\mathbf{R}_s(true_r \vdash P_2 \diamond P_3) \sqsubseteq (x :=_R \ll k \gg) ;; Q$
proof –
have $\mathbf{R}_s([\&x =_u \ll k \gg]_{S<} \vdash P_2 \diamond P_3) \sqsubseteq \mathbf{R}_s(pre_R(Q) \vdash peri_R(Q) \diamond post_R(Q))$
by (*simp add: NSRD-is-SRD SRD-reactive-tri-design assms*)
hence $\mathbf{R}_s(true_r \vdash P_2 \diamond P_3) \sqsubseteq x :=_R \ll k \gg ;; \mathbf{R}_s(pre_R(Q) \vdash peri_R(Q) \diamond post_R(Q))$
proof (*clarsimp simp add: rdes-def assms closure unrest rpred wp RHS-tri-design-refine, safe*)
assume $a1: '[\&x =_u \ll k \gg]_{S<} \Rightarrow pre_R(Q)'$ **and** $a2: '[\&x =_u \ll k \gg]_{S<} \wedge peri_R(Q) \Rightarrow P_2'$ **and** $a3: '[\&x =_u \ll k \gg]_{S<} \wedge post_R(Q) \Rightarrow P_3'$
from $a1\ assms(1)$ **show** $'R1\ true \Rightarrow [\&x \mapsto_s \ll k \gg] \uparrow_S pre_R(Q)'$
by (*rel-simp*)
show $'[\&x \mapsto_s \ll k \gg] \uparrow_S peri_R(Q) \Rightarrow P_2'$
by (*simp add: a2 assms st-subst-taut-impl closure*)
show $'[\&x \mapsto_s \ll k \gg] \uparrow_S post_R(Q) \Rightarrow P_3'$
by (*simp add: a3 assms st-subst-taut-impl closure*)
qed
thus *?thesis*
by (*simp add: NSRD-is-SRD SRD-reactive-tri-design assms*)
qed

11.3 Conditional

lemma *preR-cond-srea [rdes]*:

$pre_R(P \triangleleft b \triangleright_R Q) = ([b]_{S<} \wedge pre_R(P) \vee [\neg b]_{S<} \wedge pre_R(Q))$
by (*rel-auto*)

lemma *periR-cond-srea [rdes]*:

assumes $P\ is\ SRD\ Q\ is\ SRD$
shows $peri_R(P \triangleleft b \triangleright_R Q) = ([b]_{S<} \wedge peri_R(P) \vee [\neg b]_{S<} \wedge peri_R(Q))$

proof –

have $peri_R(P \triangleleft b \triangleright_R Q) = peri_R(R1(P) \triangleleft b \triangleright_R R1(Q))$
by (*simp add: Healthy-if SRD-healths assms*)

thus *?thesis*

by (*rel-auto*)

qed

lemma *postR-cond-srea [rdes]*:

assumes $P\ is\ SRD\ Q\ is\ SRD$
shows $post_R(P \triangleleft b \triangleright_R Q) = ([b]_{S<} \wedge post_R(P) \vee [\neg b]_{S<} \wedge post_R(Q))$

proof –

have $post_R(P \triangleleft b \triangleright_R Q) = post_R(R1(P) \triangleleft b \triangleright_R R1(Q))$
by (*simp add: Healthy-if SRD-healths assms*)

thus *?thesis*

by (*rel-auto*)

qed

lemma *NSRD-cond-srea [closure]*:

assumes P is NSRD Q is NSRD
shows $P \triangleleft b \triangleright_R Q$ is NSRD
proof (rule NSRD-RC-intro)
show $P \triangleleft b \triangleright_R Q$ is SRD
by (simp add: closure assms)
show $\text{pre}_R (P \triangleleft b \triangleright_R Q)$ is RC
proof –
have 1: $(\lceil \neg b \rceil_{S<} \vee \neg_r \text{pre}_R P) ;; R1(\text{true}) = (\lceil \neg b \rceil_{S<} \vee \neg_r \text{pre}_R P)$
by (metis (no-types, lifting) NSRD-neg-pre-unit aext-not assms(1) seqr-or-distl st-lift-R1-true-right)
have 2: $(\lceil b \rceil_{S<} \vee \neg_r \text{pre}_R Q) ;; R1(\text{true}) = (\lceil b \rceil_{S<} \vee \neg_r \text{pre}_R Q)$
by (simp add: NSRD-neg-pre-unit assms seqr-or-distl st-lift-R1-true-right)
show ?thesis
by (simp add: rdes closure assms)
qed
show $\$st' \# \text{peri}_R (P \triangleleft b \triangleright_R Q)$
by (simp add: rdes assms closure unrest)
qed

11.4 Assumptions

definition $\text{AssumeR} :: 's \text{ cond} \Rightarrow ('s, 't::\text{trace}, 'a) \text{ hrel-rsp } ([\cdot]^\top_R)$ **where**
 $[\text{upred-defs}]: \text{AssumeR } b = II_R \triangleleft b \triangleright_R \text{Miracle}$

lemma AssumeR-rdes-def [rdes-def]:
 $[\text{b}]^\top_R = \mathbf{R}_s(\text{true}_r \vdash \text{false} \diamond [\text{b}]^\top_r)$
unfolding AssumeR-def **by** (rdes-eq)

lemma AssumeR-NSRD [closure]: $[\text{b}]^\top_R$ is NSRD
by (simp add: AssumeR-def closure)

lemma AssumeR-false : $[\text{false}]^\top_R = \text{Miracle}$
by (rel-auto)

lemma AssumeR-true : $[\text{true}]^\top_R = II_R$
by (rel-auto)

lemma AssumeR-comp : $[\text{b}]^\top_R ;; [\text{c}]^\top_R = [\text{b} \wedge \text{c}]^\top_R$
by (rdes-simp)

lemma AssumeR-choice : $[\text{b}]^\top_R \sqcap [\text{c}]^\top_R = [\text{b} \vee \text{c}]^\top_R$
by (rdes-eq)

lemma $\text{AssumeR-refine-skip}$: $II_R \sqsubseteq [\text{b}]^\top_R$
by (rdes-refine)

lemma AssumeR-test [closure]: $\text{test}_R [\text{b}]^\top_R$
by (simp add: AssumeR-refine-skip nsrd-thy.utest-intro)

lemma Star-AssumeR : $[\text{b}]^\top_R^{\star R} = II_R$
by (simp add: AssumeR-NSRD AssumeR-test nsrd-thy.Star-test)

lemma $\text{AssumeR-choice-skip}$: $II_R \sqcap [\text{b}]^\top_R = II_R$
by (rdes-eq)

lemma $\text{AssumeR-seq-refines}$:
assumes P is NSRD

shows $P \sqsubseteq P ;; [b]^\top_R$
by (*rdes-refine cls: assms*)

lemma *cond-srea-AssumeR-form*:
assumes P is NSRD Q is NSRD
shows $P \triangleleft b \triangleright_R Q = ([b]^\top_R ;; P) \sqcap ([\neg b]^\top_R ;; Q)$
by (*rdes-eq cls: assms*)

lemma *cond-srea-insert-assume*:
assumes P is NSRD Q is NSRD
shows $P \triangleleft b \triangleright_R Q = ([b]^\top_R ;; P \triangleleft b \triangleright_R [\neg b]^\top_R ;; Q)$
by (*simp add: AssumeR-NSRD AssumeR-comp NSRD-seqr-closure RA1 assms cond-srea-AssumeR-form*)

lemma *AssumeR-cond-left*:
assumes P is NSRD Q is NSRD
shows $[b]^\top_R ;; (P \triangleleft b \triangleright_R Q) = ([b]^\top_R ;; P)$
by (*rdes-eq cls: assms*)

lemma *AssumeR-cond-right*:
assumes P is NSRD Q is NSRD
shows $[\neg b]^\top_R ;; (P \triangleleft b \triangleright_R Q) = ([\neg b]^\top_R ;; Q)$
by (*rdes-eq cls: assms*)

11.5 Guarded commands

definition *GuardedCommR* :: $'s \text{ cond} \Rightarrow ('s, 't::\text{trace}, 'a) \text{ hrel-rsp} \Rightarrow ('s, 't, 'a) \text{ hrel-rsp} (- \rightarrow_R - [85, 86] \ 85)$ **where**
gcmd-def[rdes-def]: *GuardedCommR* $g \ A = A \triangleleft g \triangleright_R \text{Miracle}$

lemma *gcmd-false[simp]*: $(\text{false} \rightarrow_R A) = \text{Miracle}$
unfolding *gcmd-def* **by** (*pred-auto*)

lemma *gcmd-true[simp]*: $(\text{true} \rightarrow_R A) = A$
unfolding *gcmd-def* **by** (*pred-auto*)

lemma *gcmd-SRD*:
assumes A is SRD
shows $(g \rightarrow_R A)$ is SRD
by (*simp add: gcmd-def SRD-cond-srea assms srdes-theory-continuous.weak.top-closed*)

lemma *gcmd-NSRD [closure]*:
assumes A is NSRD
shows $(g \rightarrow_R A)$ is NSRD
by (*simp add: gcmd-def NSRD-cond-srea assms NSRD-Miracle*)

lemma *gcmd-Productive [closure]*:
assumes A is NSRD A is Productive
shows $(g \rightarrow_R A)$ is Productive
by (*simp add: gcmd-def closure assms*)

lemma *gcmd-seq-distr*:
assumes B is NSRD
shows $(g \rightarrow_R A) ;; B = (g \rightarrow_R (A ;; B))$
by (*simp add: Miracle-left-zero NSRD-is-SRD assms cond-st-distr gcmd-def*)

lemma *gcmd-nondet-distr*:

assumes A is NSRD B is NSRD
shows $(g \rightarrow_R (A \sqcap B)) = (g \rightarrow_R A) \sqcap (g \rightarrow_R B)$
by (*rdes-eq cls: assms*)

lemma *AssumeR-as-gcmd*:

$[b]^\top_R = b \rightarrow_R II_R$
by (*rdes-eq*)

lemma *AssumeR-gcomm*:

assumes P is NSRD
shows $[b]^\top_R ;; (c \rightarrow_R P) = (b \wedge c) \rightarrow_R P$
by (*rdes-eq cls: assms*)

11.6 Generalised Alternation

definition *AlternateR*

$:: 'a \text{ set} \Rightarrow ('a \Rightarrow 's \text{ upred}) \Rightarrow ('a \Rightarrow ('s, 't::\text{trace}, 'a) \text{ hrel-rsp}) \Rightarrow ('s, 't, 'a) \text{ hrel-rsp} \Rightarrow ('s, 't, 'a) \text{ hrel-rsp}$ **where**
[upred-defs, rdes-def]: $\text{AlternateR } I \ g \ A \ B = (\bigsqcap i \in I \cdot ((g \ i) \rightarrow_R (A \ i))) \sqcap ((\neg (\bigvee i \in I \cdot g \ i)) \rightarrow_R B)$

definition *AlternateR-list*

$:: ('s \text{ upred} \times ('s, 't::\text{trace}, 'a) \text{ hrel-rsp}) \text{ list} \Rightarrow ('s, 't, 'a) \text{ hrel-rsp} \Rightarrow ('s, 't, 'a) \text{ hrel-rsp}$ **where**
[upred-defs, ndes-simp]:
 $\text{AlternateR-list } xs \ P = \text{AlternateR } \{0..<\text{length } xs\} \ (\lambda i. \text{map fst } xs \ ! \ i) \ (\lambda i. \text{map snd } xs \ ! \ i) \ P$

syntax

-altindR-els $:: \text{pttrn} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \ (if_R \ - \in \ - \cdot \ - \rightarrow \ - \text{ else } - \text{ fi})$
-altindR $:: \text{pttrn} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \ (if_R \ - \in \ - \cdot \ - \rightarrow \ - \text{ fi})$
-altgcommR-els $:: \text{gcomms} \Rightarrow \text{logic} \Rightarrow \text{logic} \ (if_R / \ - \text{ else } - / \text{ fi})$
-altgcommR $:: \text{gcomms} \Rightarrow \text{logic} \ (if_R / \ - / \text{ fi})$

translations

$if_R \ i \in I \cdot g \rightarrow A \text{ else } B \text{ fi} \rightarrow \text{CONST AlternateR } I \ (\lambda i. g) \ (\lambda i. A) \ B$
 $if_R \ i \in I \cdot g \rightarrow A \text{ fi} \rightarrow \text{CONST AlternateR } I \ (\lambda i. g) \ (\lambda i. A) \ (\text{CONST Chaos})$
 $if_R \ i \in I \cdot (g \ i) \rightarrow A \text{ else } B \text{ fi} \leftarrow \text{CONST AlternateR } I \ g \ (\lambda i. A) \ B$
 $\text{-altgcommR } cs \rightarrow \text{CONST AlternateR-list } cs \ (\text{CONST Chaos})$
 $\text{-altgcommR } (-\text{gcomm-show } cs) \leftarrow \text{CONST AlternateR-list } cs \ (\text{CONST Chaos})$
 $\text{-altgcommR-els } cs \ P \rightarrow \text{CONST AlternateR-list } cs \ P$
 $\text{-altgcommR-els } (-\text{gcomm-show } cs) \ P \leftarrow \text{CONST AlternateR-list } cs \ P$

lemma *AlternateR-NSRD-closed [closure]*:

assumes $\bigwedge i. i \in I \Rightarrow A \ i \text{ is NSRD } B \text{ is NSRD}$
shows $(if_R \ i \in I \cdot g \ i \rightarrow A \ i \text{ else } B \text{ fi}) \text{ is NSRD}$

proof (*cases* $I = \{\}$)

case *True*

then show *?thesis* **by** (*simp add: AlternateR-def assms*)

next

case *False*

then show *?thesis* **by** (*simp add: AlternateR-def closure assms*)

qed

lemma *AlternateR-empty [simp]*:

$(if_R \ i \in \{\} \cdot g \ i \rightarrow A \ i \text{ else } B \text{ fi}) = B$
by (*rdes-simp*)

lemma *AlternateR-Productive* [closure]:
assumes
 $\bigwedge i. i \in I \implies A \ i \text{ is NSRD } B \text{ is NSRD}$
 $\bigwedge i. i \in I \implies A \ i \text{ is Productive } B \text{ is Productive}$
shows $(\text{if}_R i \in I \cdot g \ i \rightarrow A \ i \text{ else } B \text{ fi}) \text{ is Productive}$
proof (cases $I = \{\}$)
case *True*
then show ?thesis
by (simp add: assms(4))
next
case *False*
then show ?thesis
by (simp add: AlternateR-def closure assms)
qed

lemma *AlternateR-singleton*:
assumes $A \ k \text{ is NSRD } B \text{ is NSRD}$
shows $(\text{if}_R i \in \{k\} \cdot g \ i \rightarrow A \ i \text{ else } B \text{ fi}) = (A(k) \triangleleft g(k) \triangleright_R B)$
by (simp add: AlternateR-def, rdes-eq cls: assms)

Convert an alternation over disjoint guards into a cascading if-then-else

lemma *AlternateR-insert-cascade*:
assumes
 $\bigwedge i. i \in I \implies A \ i \text{ is NSRD}$
 $A \ k \text{ is NSRD } B \text{ is NSRD}$
 $(g(k) \wedge (\bigvee i \in I \cdot g(i))) = \text{false}$
shows $(\text{if}_R i \in \text{insert } k \ I \cdot g \ i \rightarrow A \ i \text{ else } B \text{ fi}) = (A(k) \triangleleft g(k) \triangleright_R (\text{if}_R i \in I \cdot g(i) \rightarrow A(i) \text{ else } B \text{ fi}))$
proof (cases $I = \{\}$)
case *True*
then show ?thesis **by** (simp add: AlternateR-singleton assms)
next
case *False*
have 1: $(\bigcap i \in I \cdot g \ i \rightarrow_R A \ i) = (\bigcap i \in I \cdot g \ i \rightarrow_R \mathbf{R}_s(\text{pre}_R(A \ i) \vdash \text{peri}_R(A \ i) \diamond \text{post}_R(A \ i)))$
by (simp add: NSRD-is-SRD SRD-reactive-tri-design assms(1) cong: UINF-cong)
from assms(4) **show** ?thesis
by (simp add: AlternateR-def 1 False)
 $(\text{rdes-eq cls: assms}(1-3) \text{ False cong: UINF-cong})$
qed

lemma *AlternateR-assume-branch*:
assumes $I \neq \{\}$ $\bigwedge i. i \in I \implies P \ i \text{ is NSRD } Q \text{ is NSRD}$
shows $(\bigcap i \in I \cdot b \ i]^\top_R ;; \text{AlternateR } I \ b \ P \ Q) = (\bigcap i \in I \cdot b \ i \rightarrow_R P \ i) \ (\text{is } ?lhs = ?rhs)$
proof –
have ?lhs = $\bigcap i \in I \cdot b \ i]^\top_R ;; ((\bigcap i \in I \cdot b \ i \rightarrow_R P \ i) \sqcap (\neg (\bigcap i \in I \cdot b \ i)) \rightarrow_R Q)$
by (simp add: AlternateR-def closure assms)
also have ... = $\bigcap i \in I \cdot b \ i]^\top_R ;; (\bigcap i \in I \cdot b \ i \rightarrow_R P \ i) \sqcap \text{Miracle}$
by (simp add: seqr-inf-distr AssumeR-gcomm closure assms)
also have ... = $(\bigcap i \in I \cdot ((\bigcap i \in I \cdot b \ i) \wedge b \ i) \rightarrow_R P \ i) \sqcap \text{Miracle}$
by (simp add: seq-UINF-distl AssumeR-gcomm closure assms cong: UINF-cong)
also have ... = $(\bigcap i \in I \cdot b \ i \rightarrow_R P \ i) \sqcap \text{Miracle}$
proof –
have $\bigwedge i. i \in I \implies ((\bigcap i \in I \cdot b \ i) \wedge b \ i) = b \ i$
by (rel-auto)
thus ?thesis

by (simp cong: UINF-cong)
 qed
 also have ... = ($\bigcap i \in I \cdot b\ i \rightarrow_R P\ i$)
 by (simp add: closure assms)
 finally show ?thesis .
 qed

11.7 Choose

definition choose-srd :: ('s, 't::trace, 'α) hrel-rsp (choose_R) **where**
 [upred-defs, rdes-def]: choose_R = $\mathbf{R}_s(\text{true}_r \vdash \text{false} \diamond \text{true}_r)$

lemma preR-choose [rdes]: pre_R(choose_R) = true_r
 by (rel-auto)

lemma periR-choose [rdes]: peri_R(choose_R) = false
 by (rel-auto)

lemma postR-choose [rdes]: post_R(choose_R) = true_r
 by (rel-auto)

lemma choose-srd-SRD [closure]: choose_R is SRD
 by (simp add: choose-srd-def closure unrest)

lemma NSRD-choose-srd [closure]: choose_R is NSRD
 by (rule NSRD-intro, simp-all add: closure unrest rdes)

11.8 Divergence Freedom

definition ndiv-srd :: ('s, 't::trace, 'α) hrel-rsp (ndiv_R)
where [rdes-def]: ndiv-srd = $\mathbf{R}_s(\text{true}_r \vdash \text{true}_r \diamond \text{true}_r)$

lemma ndiv-NSRD [closure]: ndiv_R is NSRD
 by (simp add: rdes-def closure unrest)

lemma ndiv-srd-refines-preR-true:
 assumes P is SRD
 shows $\text{ndiv}_R \sqsubseteq P \longleftrightarrow \text{pre}_R(P) = \text{true}_r$ (is ?lhs \longleftrightarrow ?rhs)

proof

assume ?lhs

thus ?rhs

by (metis R1-preR ndiv-srd-def preR-antitone preR-rdes rea-true-RR rea-true-disj(2) utp-pred-laws.sup.orderE)

next

assume ?rhs

hence $\text{ndiv}_R \sqsubseteq \mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P))$

by (simp add: RHS-tri-design-conj assms ndiv-srd-def periR-SRD-R1 postR-SRD-R1 rea-true-conj(1) rea-true-impl utp-pred-laws.inf.absorb-iff2)

thus ?lhs

by (simp add: SRD-reactive-tri-design assms)

qed

lemma ndiv-srd-refines-rdes-pre-true:
 assumes P_1 is RR P_2 is RR P_3 is RR
 shows $\text{ndiv}_R \sqsubseteq \mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \longleftrightarrow P_1 = \text{true}_r$ (is ?lhs \longleftrightarrow ?rhs)
 by (simp add: ndiv-srd-refines-preR-true closure assms rdes unrest)

11.9 State Abstraction

definition *state-srea* ::

's itself \Rightarrow (*'s, 't::trace, 'α, 'β*) *rel-rsp* \Rightarrow (*unit, 't, 'α, 'β*) *rel-rsp* **where**
[upred-defs]: *state-srea* *t P* = $\langle \exists \{ \$st, \$st' \} \cdot P \rangle_S$

syntax

-state-srea :: *type* \Rightarrow *logic* \Rightarrow *logic* (*state* - - - [0,200] 200)

translations

state 'a \cdot *P* == *CONST state-srea TYPE('a) P*

lemma *R1-state-srea*: *R1*(*state 'a* \cdot *P*) = (*state 'a* \cdot *R1*(*P*))
by (*rel-auto*)

lemma *R2c-state-srea*: *R2c*(*state 'a* \cdot *P*) = (*state 'a* \cdot *R2c*(*P*))
by (*rel-auto*)

lemma *R3h-state-srea*: *R3h*(*state 'a* \cdot *P*) = (*state 'a* \cdot *R3h*(*P*))
by (*rel-auto*)

lemma *RD1-state-srea*: *RD1*(*state 'a* \cdot *P*) = (*state 'a* \cdot *RD1*(*P*))
by (*rel-auto*)

lemma *RD2-state-srea*: *RD2*(*state 'a* \cdot *P*) = (*state 'a* \cdot *RD2*(*P*))
by (*rel-auto*)

lemma *RD3-state-srea*: *RD3*(*state 'a* \cdot *P*) = (*state 'a* \cdot *RD3*(*P*))
by (*rel-auto*, *blast+*)

lemma *SRD-state-srea* [*closure*]: *P* is *SRD* \Longrightarrow *state 'a* \cdot *P* is *SRD*
by (*simp add: Healthy-def R1-state-srea R2c-state-srea R3h-state-srea RD1-state-srea RD2-state-srea RHS-def SRD-def*)

lemma *NSRD-state-srea* [*closure*]: *P* is *NSRD* \Longrightarrow *state 'a* \cdot *P* is *NSRD*
by (*metis Healthy-def NSRD-is-RD3 NSRD-is-SRD RD3-state-srea SRD-RD3-implies-NSRD SRD-state-srea*)

lemma *preR-state-srea* [*rdes*]: *pre_R*(*state 'a* \cdot *P*) = $\langle \forall \{ \$st, \$st' \} \cdot \text{pre}_R(P) \rangle_S$
by (*simp add: state-srea-def, rel-auto*)

lemma *periR-state-srea* [*rdes*]: *peri_R*(*state 'a* \cdot *P*) = *state 'a* \cdot *peri_R*(*P*)
by (*rel-auto*)

lemma *postR-state-srea* [*rdes*]: *post_R*(*state 'a* \cdot *P*) = *state 'a* \cdot *post_R*(*P*)
by (*rel-auto*)

lemma *state-srea-rdes-def* [*rdes-def*]:
assumes *P* is *RC* *Q* is *RR* *R* is *RR*
shows *state 'a* \cdot **R_s**(*P* \vdash *Q* \diamond *R*) = **R_s**($\langle \forall \{ \$st, \$st' \} \cdot P \rangle_S \vdash$ (*state 'a* \cdot *Q*) \diamond (*state 'a* \cdot *R*))
(is ?lhs = ?rhs)

proof –

have *?lhs* = **R_s**(*pre_R*(*?lhs*) \vdash *peri_R*(*?lhs*) \diamond *post_R*(*?lhs*))
by (*simp add: RC-implies-RR SRD-rdes-intro SRD-reactive-tri-design SRD-state-srea assms*)
also have ... = **R_s**($\langle \forall \{ \$st, \$st' \} \cdot P \rangle_S \vdash$ *state 'a* \cdot (*P* \Rightarrow_r *Q*) \diamond *state 'a* \cdot (*P* \Rightarrow_r *R*))
by (*simp add: rdes closure assms*)
also have ... = *?rhs*

by (rel-auto)
 finally show ?thesis .
 qed

lemma ext-st-rdes-dist [rdes-def]:
 $\mathbf{R}_s(P \vdash Q \diamond R) \oplus_p \text{abs-st}_L = \mathbf{R}_s(P \oplus_p \text{abs-st}_L \vdash Q \oplus_p \text{abs-st}_L \diamond R \oplus_p \text{abs-st}_L)$
 by (rel-auto)

lemma state-srea-refine:
 $(P \oplus_p \text{abs-st}_L) \sqsubseteq Q \implies P \sqsubseteq (\text{state-srea } \text{TYPE}('s) \ Q)$
 by (rel-auto)

11.10 Reactive Frames

definition rdes-frame-ext :: $(' \alpha \implies ' \beta) \Rightarrow (' \alpha, 't::\text{trace}, 'r) \text{ hrel-rsp} \Rightarrow (' \beta, 't, 'r) \text{ hrel-rsp}$ **where**
 [upred-defs, rdes-def]: $\text{rdes-frame-ext } a \ P = \mathbf{R}_s(\text{rel-aext } (\text{pre}_R(P)) (\text{map-st}_L \ a) \vdash \text{rel-aext } (\text{peri}_R(P)) (\text{map-st}_L \ a) \diamond a : [\text{post}_R(P)]_r^+)$

syntax
 $\text{-rdes-frame-ext} :: \text{salpha} \Rightarrow \text{logic} \Rightarrow \text{logic } (:-)_{R^+} [99,0] \ 100)$

translations
 $\text{-rdes-frame-ext } x \ P \Rightarrow \text{CONST } \text{rdes-frame-ext } x \ P$
 $\text{-rdes-frame-ext } (-\text{salphaset } (-\text{salphamk } x)) \ P \Leftarrow \text{CONST } \text{rdes-frame-ext } x \ P$

lemma RC-rel-aext-st-closed [closure]:
 assumes P is RC
 shows $\text{rel-aext } P (\text{map-st}_L \ a)$ is RC
proof –
 have $\text{RC}(\text{rel-aext } (\text{RC}(P)) (\text{map-st}_L \ a)) = \text{rel-aext } (\text{RC}(P)) (\text{map-st}_L \ a)$
 by (rel-auto)
 (metis (no-types, hide-lams) diff-add-cancel-left' dual-order.trans le-add trace-class.add-diff-cancel-left trace-class.add-left-mono)
 thus ?thesis
 by (rule-tac Healthy-intro, simp add: assms Healthy-if)
 qed

lemma rdes-frame-ext-SRD-closed:
 $\llbracket P \text{ is SRD}; \$\text{wait}' \ \# \ \text{pre}_R(P) \rrbracket \implies a : [P]_{R^+} \text{ is SRD}$
 unfolding rdes-frame-ext-def
 apply (rule SRD-rdes-intro)
 apply (simp-all add: closure unrest)
 apply (simp add: RR-R2-intro ok'-pre-unrest ok-pre-unrest preR-R2-closed rea-aext-RR wait-pre-unrest)
 done

lemma preR-rdes-frame-ext:
 $P \text{ is NSRD} \implies \text{pre}_R(a : [P]_{R^+}) = \text{rel-aext } (\text{pre}_R(P)) (\text{map-st}_L \ a)$
 by (simp add: preR-RR preR-rdes rdes-frame-ext-def rea-aext-RR)

lemma unrest-rel-aext-st' [unrest]: $\$st' \ \# \ P \implies \$st' \ \# \ \text{rel-aext } P (\text{map-st}_L \ a)$
 by (rel-auto)

lemma rdes-frame-ext-NSRD-closed:
 $P \text{ is NSRD} \implies a : [P]_{R^+} \text{ is NSRD}$
 apply (rule NSRD-RC-intro)
 apply (rule rdes-frame-ext-SRD-closed)

apply (*simp-all add: closure unrest rdes*)
apply (*simp add: NSRD-neg-pre-RC RC-rel-aext-st-closed preR-RR preR-rdes rdes-frame-ext-def rea-aext-RR*)
apply (*simp add: rdes-frame-ext-def*)
apply (*simp add: rdes closure unrest*)
done

lemma *skip-srea-frame* [*frame*]:
 $vwb\text{-}lens\ a \implies a:[II_R]_{R^+} = II_R$
by (*rdes-eq*)

lemma *seq-srea-frame* [*frame*]:
assumes *vwb-lens a P is NSRD Q is NSRD*
shows $a:[P \;; Q]_{R^+} = a:[P]_{R^+} \;; a:[Q]_{R^+}$ (**is** *?lhs = ?rhs*)
proof –
have *?lhs =* $\mathbf{R}_s ((pre_R P \wedge post_R P \wp_r pre_R Q) \oplus_r map\text{-}st_L[a] \vdash$
 $((pre_R P \wedge post_R P \wp_r pre_R Q) \oplus_r map\text{-}st_L[a] \Rightarrow_r (peri_R P \vee post_R P \;; peri_R Q)$
 $\oplus_r map\text{-}st_L[a]) \diamond$
 $a:[pre_R P \wedge post_R P \wp_r pre_R Q \Rightarrow_r post_R P \;; post_R Q]_{r^+})$
using *assms(1)* **by** (*rdes-simp cls: assms(2-3)*)
also have $\dots = \mathbf{R}_s ((pre_R P \wedge post_R P \wp_r pre_R Q) \oplus_r map\text{-}st_L[a] \vdash$
 $((peri_R P \vee post_R P \;; peri_R Q) \oplus_r map\text{-}st_L[a]) \diamond$
 $a:[post_R P \;; post_R Q]_{r^+})$
by (*rel-auto*)
also from *assms(1)* **have** $\dots = ?rhs$
apply (*rdes-eq-split cls: assms(2-3)*)
apply (*rel-auto*)
apply (*metis mwb-lens-def vwb-lens-mwb weak-lens.put-get*)
apply (*rel-auto*)
apply (*metis mwb-lens-def vwb-lens-mwb weak-lens.put-get*)
apply (*simp add: rea-frame-ext-seq*)
done
finally show *?thesis* .
qed

lemma *rdes-frame-ext-Productive-closed* [*closure*]:
assumes *P is NSRD P is Productive*
shows $x:[P]_{R^+}$ *is Productive*
proof –
have $x:[Productive(P)]_{R^+}$ *is Productive*
by (*rdes-simp cls: assms, rel-auto*)
thus *?thesis*
by (*simp add: Healthy-if assms*)
qed

11.11 While Loop

definition *WhileR* :: $'s \text{ upred} \Rightarrow ('s, 't::\text{size-trace}, 'a) \text{ hrel-rsp} \Rightarrow ('s, 't, 'a) \text{ hrel-rsp}$ (*while_R - do - od*)
where
 $WhileR\ b\ P = (\mu_R X \cdot (P \;; X) \triangleleft b \triangleright_R II_R)$

lemma *Sup-power-false*:
fixes $F :: 'a \text{ upred} \Rightarrow 'a \text{ upred}$
shows $(\bigcap i. (F \hat{\wedge} i) \text{ false}) = (\bigcap i. (F \hat{\wedge} (i+1)) \text{ false})$
proof –
have $(\bigcap i. (F \hat{\wedge} i) \text{ false}) = (F \hat{\wedge} 0) \text{ false} \sqcap (\bigcap i. (F \hat{\wedge} (i+1)) \text{ false})$

```

    by (subst Sup-power-expand, simp)
  also have ... = ( $\bigwedge i. (F \wedge (i+1))$  false)
    by (simp)
  finally show ?thesis .
qed

theorem WhileR-iter-expand:
  assumes P is NSRD P is Productive
  shows  $\text{while}_R b \text{ do } P \text{ od} = (\bigwedge i. (P \triangleleft b \triangleright_R II_R) \wedge i ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R))$  (is ?lhs = ?rhs)
proof -
  have 1: Continuous ( $\lambda X. P ;; SRD X$ )
    using SRD-Continuous
    by (clarsimp simp add: Continuous-def seq-SUP-distl[THEN sym], drule-tac  $x=A$  in spec, simp)
  have 2: Continuous ( $\lambda X. P ;; SRD X \triangleleft b \triangleright_R II_R$ )
    by (simp add: 1 closure assms)
  have ?lhs = ( $\mu_R X \cdot P ;; X \triangleleft b \triangleright_R II_R$ )
    by (simp add: WhileR-def)
  also have ... = ( $\mu X \cdot P ;; SRD(X) \triangleleft b \triangleright_R II_R$ )
    by (auto simp add: srd-mu-equiv closure assms)
  also have ... = ( $\nu X \cdot P ;; SRD(X) \triangleleft b \triangleright_R II_R$ )
    by (auto simp add: guarded-fp-uniq Guarded-if-Productive[OF assms] funcsetI closure assms)
  also have ... = ( $\bigwedge i. ((\lambda X. P ;; SRD X \triangleleft b \triangleright_R II_R) \wedge i)$  false)
    by (simp add: sup-continuous-lfp 2 sup-continuous-Continuous false-upred-def)
  also have ... = ( $\bigwedge i. ((\lambda X. P ;; SRD X \triangleleft b \triangleright_R II_R) \wedge (i+1))$  false)
    by (simp add: Sup-power-false)
  also have ... = ( $\bigwedge i. (P \triangleleft b \triangleright_R II_R) \wedge i ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)$ )
  proof (rule SUP-cong, simp)
    fix i
    show  $((\lambda X. P ;; SRD X \triangleleft b \triangleright_R II_R) \wedge (i+1))$  false =  $(P \triangleleft b \triangleright_R II_R) \wedge i ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)$ 
  proof (induct i)
    case 0
    thm if-eq-cancel
    then show ?case
      by (simp, metis srdes-hcond-def srdes-theory-continuous.healthy-top)
  next
    case (Suc i)
    show ?case
    proof -
      have  $((\lambda X. P ;; SRD X \triangleleft b \triangleright_R II_R) \wedge (Suc i + 1))$  false =
         $P ;; SRD (((\lambda X. P ;; SRD X \triangleleft b \triangleright_R II_R) \wedge (i+1))$  false)  $\triangleleft b \triangleright_R II_R$ 
        by simp
      also have ... =  $P ;; SRD ((P \triangleleft b \triangleright_R II_R) \wedge i ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)) \triangleleft b \triangleright_R II_R$ 
        using Suc.hyps by auto
      also have ... =  $P ;; ((P \triangleleft b \triangleright_R II_R) \wedge i ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)) \triangleleft b \triangleright_R II_R$ 
        by (metis (no-types, lifting) Healthy-if NSRD-cond-srea NSRD-is-SRD NSRD-power-Suc
          NSRD-srd-skip SRD-cond-srea SRD-seqr-closure assms(1) power.power-eq-if seqr-left-unit srdes-theory-continuous.top-cl
          also have ... =  $(P \triangleleft b \triangleright_R II_R) \wedge Suc i ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)$ 
        proof (induct i)
          case 0
          then show ?case
            by (simp add: NSRD-is-SRD SRD-cond-srea SRD-left-unit SRD-seqr-closure SRD-srdes-skip
              assms(1) cond-L6 cond-st-distr srdes-theory-continuous.top-closed)
        next
          case (Suc i)

```

have $1: II_R ;; ((P \triangleleft b \triangleright_R II_R) ;; (P \triangleleft b \triangleright_R II_R) \wedge i) = ((P \triangleleft b \triangleright_R II_R) ;; (P \triangleleft b \triangleright_R II_R) \wedge i)$
by (*simp add: NSRD-is-SRD RA1 SRD-cond-srea SRD-left-unit SRD-srdes-skip assms(1)*)
then show ?case
proof –
have $\bigwedge u. (u ;; (P \triangleleft b \triangleright_R II_R) \wedge \text{Suc } i) ;; (P ;; (\text{Miracle} \triangleleft b \triangleright_R (II_R)) \triangleleft b \triangleright_R (II_R) =$
 $((u \triangleleft b \triangleright_R II_R) ;; (P \triangleleft b \triangleright_R II_R) \wedge \text{Suc } i) ;; (P ;; (\text{Miracle} \triangleleft b \triangleright_R (II_R)))$
by (*metis (no-types) Suc.hyps 1 cond-L6 cond-st-distr power.power.power-Suc*)
then show ?thesis
by (*simp add: RA1 upred-semiring.power-Suc*)
qed
qed
finally show ?thesis .
qed
qed
qed
also have $\dots = (\bigcap i \cdot (P \triangleleft b \triangleright_R II_R) \wedge i) ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)$
by (*simp add: UINF-as-Sup-collect'*)
finally show ?thesis .
qed

theorem *WhileR-star-expand*:

assumes *P is NSRD P is Productive*

shows $\text{while}_R b \text{ do } P \text{ od} = (P \triangleleft b \triangleright_R II_R)^{\star R} ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)$ (**is** ?lhs = ?rhs)

proof –

have ?lhs = $(\bigcap i \cdot (P \triangleleft b \triangleright_R II_R) \wedge i) ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)$

by (*simp add: WhileR-iter-expand seq-UINF-distr' assms*)

also have $\dots = (P \triangleleft b \triangleright_R II_R)^{\star} ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)$

by (*simp add: ustar-def*)

also have $\dots = ((P \triangleleft b \triangleright_R II_R)^{\star} ;; II_R) ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)$

by (*simp add: seqr-assoc SRD-left-unit closure assms*)

also have $\dots = (P \triangleleft b \triangleright_R II_R)^{\star R} ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)$

by (*simp add: nsrd-thy.Star-def*)

finally show ?thesis .

qed

lemma *WhileR-NSRD-closed [closure]*:

assumes *P is NSRD P is Productive*

shows $\text{while}_R b \text{ do } P \text{ od}$ *is NSRD*

by (*simp add: WhileR-star-expand assms closure*)

theorem *WhileR-iter-form-lemma*:

assumes *P is NSRD*

shows $(P \triangleleft b \triangleright_R II_R)^{\star R} ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R) = ([b]^{\top}_R ;; P)^{\star R} ;; [\neg b]^{\top}_R$

proof –

have $(P \triangleleft b \triangleright_R II_R)^{\star R} ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R) = ([b]^{\top}_R ;; P) \sqcap [\neg b]^{\top}_R)^{\star R} ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)$

by (*simp add: AssumeR-NSRD NSRD-right-unit NSRD-srd-skip assms(1) cond-srea-AssumeR-form*)

also have $\dots = ([b]^{\top}_R ;; P)^{\star R} ;; [\neg b]^{\top}_R)^{\star R} ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)$

by (*simp add: AssumeR-NSRD NSRD-seqr-closure nsrd-thy.Star-denest assms(1)*)

also have $\dots = ([b]^{\top}_R ;; P)^{\star R} ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)$

by (*metis (no-types, hide-lams) RD3-def RD3-idem Star-AssumeR nsrd-thy.Star-def*)

also have $\dots = ([b]^{\top}_R ;; P)^{\star R} ;; (P ;; \text{Miracle} \triangleleft b \triangleright_R II_R)$

by (*simp add: AssumeR-NSRD NSRD-seqr-closure nsrd-thy.Star-invol assms(1)*)

also have $\dots = ([b]^{\top}_R ;; P)^{\star R} ;; ([b]^{\top}_R ;; P ;; \text{Miracle}) \sqcap [\neg b]^{\top}_R$

by (*simp add: AssumeR-NSRD NSRD-Miracle NSRD-right-unit NSRD-seqr-closure NSRD-srd-skip*)

assms(1) cond-srea-AssumeR-form
also have ... = $((([b]^\top_R \;; P)^{\star R}) \;; [b]^\top_R \;; P \;; \text{Miracle}) \sqcap (([b]^\top_R \;; P)^{\star R} \;; [\neg b]^\top_R)$
by (*simp add: upred-semiring.distrib-left*)
also have ... = $([b]^\top_R \;; P)^{\star R} \;; [\neg b]^\top_R$
proof –
have $((([b]^\top_R \;; P)^{\star R}) \;; [\neg b]^\top_R = (II_R \sqcap (([b]^\top_R \;; P)^{\star R} \;; [b]^\top_R \;; P)) \;; [\neg b]^\top_R$
by (*simp add: AssumeR-NSRD NSRD-seqr-closure nsrd-thy.Star-unfoldr-eq assms(1)*)
also have ... = $[\neg b]^\top_R \sqcap ((([b]^\top_R \;; P)^{\star R} \;; [b]^\top_R \;; P) \;; [\neg b]^\top_R)$
by (*metis (no-types, lifting) AssumeR-NSRD AssumeR-as-gcmd NSRD-srd-skip Star-AssumeR nsrd-thy.Star-slide gcmd-seq-distr skip-srea-self-unit urel-diod.distrib-right'*)
also have ... = $[\neg b]^\top_R \sqcap ((([b]^\top_R \;; P)^{\star R} \;; [b]^\top_R \;; P \;; [b \vee \neg b]^\top_R) \;; [\neg b]^\top_R)$
by (*simp add: AssumeR-true NSRD-right-unit assms(1)*)
also have ... = $[\neg b]^\top_R \sqcap ((([b]^\top_R \;; P)^{\star R} \;; [b]^\top_R \;; P \;; [b]^\top_R) \;; [\neg b]^\top_R)$
 $\sqcap ((([b]^\top_R \;; P)^{\star R} \;; [b]^\top_R \;; P \;; [\neg b]^\top_R) \;; [\neg b]^\top_R)$
by (*metis (no-types, hide-lams) AssumeR-choice upred-semiring.add-assoc upred-semiring.distrib-left upred-semiring.distrib-right*)
also have ... = $[\neg b]^\top_R \sqcap ((([b]^\top_R \;; P)^{\star R} \;; [b]^\top_R \;; P \;; ([b]^\top_R \;; [\neg b]^\top_R)) \sqcap (([b]^\top_R \;; P)^{\star R} \;; [b]^\top_R \;; P \;; ([\neg b]^\top_R \;; [\neg b]^\top_R))$
by (*simp add: RA1*)
also have ... = $[\neg b]^\top_R \sqcap ((([b]^\top_R \;; P)^{\star R} \;; [b]^\top_R \;; P \;; \text{Miracle}) \sqcap (([b]^\top_R \;; P)^{\star R} \;; [b]^\top_R \;; P \;; [\neg b]^\top_R)$
by (*simp add: AssumeR-comp AssumeR-false*)
finally have $([b]^\top_R \;; P)^{\star R} \;; [\neg b]^\top_R \sqsubseteq (([b]^\top_R \;; P)^{\star R}) \;; [b]^\top_R \;; P \;; \text{Miracle}$
by (*simp add: semilattice-sup-class.le-supI1*)
thus ?thesis
by (*simp add: semilattice-sup-class.le-iff-sup*)
qed
finally show ?thesis .
qed

theorem *WhileR-iter-form:*

assumes *P is NSRD P is Productive*
shows $\text{while}_R \ b \ \text{do} \ P \ \text{od} = ([b]^\top_R \;; P)^{\star R} \;; [\neg b]^\top_R$
by (*simp add: WhileR-iter-form-lemma WhileR-star-expand assms*)

theorem *WhileR-outer-refine-intro:*

assumes
P is NSRD P is Productive
 $S \sqsubseteq ([b]^\top_R \;; P) \;; S \sqsubseteq [\neg b]^\top_R$
shows $S \sqsubseteq \text{while}_R \ b \ \text{do} \ P \ \text{od}$
apply (*simp add: assms WhileR-iter-form*)
apply (*rule nsrd-thy.Star-inductl*)
apply (*simp-all add: closure assms*)
done

theorem *WhileR-outer-refine-init-intro:*

assumes
P is NSRD I is NSRD P is Productive
 $S \sqsubseteq I \;; [\neg b]^\top_R$
 $S \sqsubseteq S \;; [b]^\top_R \;; P$
 $S \sqsubseteq I \;; [b]^\top_R \;; P$
shows $S \sqsubseteq I \;; \text{while}_R \ b \ \text{do} \ P \ \text{od}$

proof –

have $S \sqsubseteq I \;; (([b]^\top_R \;; P) \;; ([b]^\top_R \;; P)^{\star R}) \;; [\neg b]^\top_R$

proof –

have $S \sqsubseteq I \;; ([b]^\top_R \;; P) \;; ([b]^\top_R \;; P)^{\star R}$
by (*metis* (*no-types*, *hide-lams*) *AssumeR-NSRD NSRD-seqr-closure RA1 assms(1) assms(2)*
assms(5) assms(6) nsrd-thy.Star-inductr semilattice-sup-class.le-sup-iff)
thus *?thesis*
by (*metis* (*no-types*, *lifting*) *AssumeR-NSRD AssumeR-refine-skip NSRD-right-unit NSRD-seqr-closure*
assms(1) dual-order.trans nsrd-thy.Star-Healthy urel-diod.mult-isol)
qed
moreover **have** $S \sqsubseteq I \;; II_R \;; [\neg b]^\top_R$
by (*simp add: AssumeR-NSRD assms nsrd-thy.Unit-Left*)
ultimately show *?thesis*
apply (*simp add: assms WhileR-iter-form*)
apply (*subst nsrd-thy.Star-unfoldl-eq[THEN sym]*)
apply (*auto simp add: closure assms seqr-inf-distr*)
done
qed

theorem *WhileR-false*:
assumes *P is NSRD*
shows *while_R false do P od = II_R*
by (*simp add: WhileR-def rpred closure sdes-theory-continuous.LFP-const*)

theorem *WhileR-true*:
assumes *P is NSRD P is Productive*
shows *while_R true do P od = P^{★R} ;; Miracle*
by (*simp add: WhileR-iter-form AssumeR-true AssumeR-false SRD-left-unit assms closure*)

lemma *WhileR-insert-assume*:
assumes *P is NSRD P is Productive*
shows *while_R b do ([b][⊤]_R ;; P) od = while_R b do P od*
by (*simp add: AssumeR-NSRD AssumeR-comp NSRD-seqr-closure Productive-seq-2 RA1 WhileR-iter-form assms*)

theorem *WhileR-rdes-def [rdes-def]*:
assumes *P is RC Q is RR R is RR \$st' \# Q R is R4*
shows *while_R b do R_s(P ⊢ Q ◊ R) od =*

$$\mathbf{R}_s \left(([b]^\top_r \;; R)^{\star r} \text{ wp}_r ([b]_{S<} \Rightarrow_r P) \vdash ([b]^\top_r \;; R)^{\star r} \;; [b]^\top_r \;; Q \right) \diamond ([b]^\top_r \;; R)^{\star r} \;; [\neg b]^\top_r$$

(is ?lhs = ?rhs)

proof –
have *?lhs = ([b][⊤]_R ;; R_s (P ⊢ Q ◊ R))^{★R} ;; [¬ b][⊤]_R*
by (*simp add: WhileR-iter-form Productive-rdes-RR-intro assms closure*)
also have *... = ?rhs*
by (*simp add: rdes-def assms closure unrest rpred wp del: rea-star-wp*)
finally show *?thesis* .
qed

Refinement introduction law for reactive while loops

theorem *WhileR-refine-intro*:
assumes
— Closure conditions
Q₁ is RC Q₂ is RR Q₃ is RR \$st' \# Q₂ Q₃ is R4
— Refinement conditions

$$([b]^\top_r \;; Q_3)^{\star r} \text{ wp}_r ([b]_{S<} \Rightarrow_r Q_1) \sqsubseteq P_1$$

$$P_2 \sqsubseteq [b]^\top_r \;; Q_2$$

$$P_2 \sqsubseteq [b]^\top_r \;; Q_3 \;; P_2$$

$$P_3 \sqsubseteq [\neg b]^\top_r$$

$P_3 \sqsubseteq [b]^\top_r ;; Q_3 ;; P_3$
shows $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \sqsubseteq \text{while}_R b \text{ do } \mathbf{R}_s(Q_1 \vdash Q_2 \diamond Q_3) \text{ od}$
proof (*simp add: rdes-def assms, rule srdes-tri-refine-intro*)
show $([b]^\top_r ;; Q_3)^{*r} \text{wp}_r ([b]_{S<} \Rightarrow_r Q_1) \sqsubseteq P_1$
by (*simp add: assms*)
show $P_2 \sqsubseteq (P_1 \wedge ([b]^\top_r ;; Q_3)^{*r} ;; [b]^\top_r ;; Q_2)$
proof –
have $P_2 \sqsubseteq ([b]^\top_r ;; Q_3)^{*r} ;; [b]^\top_r ;; Q_2$
by (*simp add: assms rea-assume-RR rrel-thy.Star-inductl seq-RR-closed seqr-assoc*)
thus ?thesis
by (*simp add: utp-pred-laws.le-infI2*)
qed
show $P_3 \sqsubseteq (P_1 \wedge ([b]^\top_r ;; Q_3)^{*r} ;; [\neg b]^\top_r)$
proof –
have $P_3 \sqsubseteq ([b]^\top_r ;; Q_3)^{*r} ;; [\neg b]^\top_r$
by (*simp add: assms rea-assume-RR rrel-thy.Star-inductl seqr-assoc*)
thus ?thesis
by (*simp add: utp-pred-laws.le-infI2*)
qed
qed

11.12 Iteration Construction

definition *IterateR*

$:: 'a \text{ set} \Rightarrow ('a \Rightarrow 's \text{ upred}) \Rightarrow ('a \Rightarrow ('s, 't::\text{size-trace}, 'a) \text{ hrel-rsp}) \Rightarrow ('s, 't, 'a) \text{ hrel-rsp}$
where $\text{IterateR } A \ g \ P = \text{while}_R (\bigvee i \in A \cdot g(i)) \text{ do } (\text{if}_R i \in A \cdot g(i) \rightarrow P(i) \text{ fi}) \text{ od}$

definition *IterateR-list*

$:: ('s \text{ upred} \times ('s, 't::\text{size-trace}, 'a) \text{ hrel-rsp}) \text{ list} \Rightarrow ('s, 't, 'a) \text{ hrel-rsp}$ **where**
 $[\text{upred-defs}, \text{ndes-simp}]$:

$\text{IterateR-list } xs = \text{IterateR } \{0..<\text{length } xs\} (\lambda i. \text{map fst } xs ! i) (\lambda i. \text{map snd } xs ! i)$

syntax

$\text{-iter-srd} \quad :: \text{pttrn} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \Rightarrow \text{logic} \text{ (do}_R \text{ -} \cdot \text{ -} \rightarrow \text{ - od)}$
 $\text{-iter-gcommR} \quad :: \text{gcomms} \Rightarrow \text{logic} \text{ (do}_R \text{ / - /od)}$

translations

$\text{-iter-srd } x \ A \ g \ P \Rightarrow \text{CONST IterateR } A \ (\lambda x. g) \ (\lambda x. P)$
 $\text{-iter-srd } x \ A \ g \ P \Leftarrow \text{CONST IterateR } A \ (\lambda x. g) \ (\lambda x'. P)$
 $\text{-iter-gcommR } cs \rightarrow \text{CONST IterateR-list } cs$
 $\text{-iter-gcommR } (-\text{gcomm-show } cs) \leftarrow \text{CONST IterateR-list } cs$

lemma *IterateR-NSRD-closed [closure]*:

assumes
 $\bigwedge i. i \in I \Rightarrow P(i) \text{ is NSRD}$
 $\bigwedge i. i \in I \Rightarrow P(i) \text{ is Productive}$
shows $\text{do}_R i \in I \cdot g(i) \rightarrow P(i) \text{ od is NSRD}$
by (*simp add: IterateR-def closure assms*)

lemma *IterateR-empty*:

$\text{do}_R i \in \{\} \cdot g(i) \rightarrow P(i) \text{ od} = \text{II}_R$
by (*simp add: IterateR-def srd-mu-equiv closure rpred gfp-const WhileR-false*)

lemma *IterateR-singleton*:

assumes $P \ k \text{ is NSRD } P \ k \text{ is Productive}$
shows $\text{do}_R i \in \{k\} \cdot g(i) \rightarrow P(i) \text{ od} = \text{while}_R g(k) \text{ do } P(k) \text{ od}$ (**is** ?lhs = ?rhs)

proof –

have ?lhs = while_R g k do P k < g k ▷_R Chaos od
 by (simp add: IterateR-def AlternateR-singleton assms closure)
 also have ... = while_R g k do [g k][⊤]_R ;; (P k < g k ▷_R Chaos) od
 by (simp add: WhileR-insert-assume closure assms)
 also have ... = while_R g k do P k od
 by (simp add: AssumeR-cond-left NSRD-Chaos WhileR-insert-assume assms)
 finally show ?thesis .

qed

declare IterateR-list-def [rdes-def]

declare IterateR-def [rdes-def]

lemma R4-Continuous [closure]: Continuous R4

by (rel-auto)

lemma cond-rea-R4-closed [closure]:

[[P is R4; Q is R4]] ⇒ P < b ▷_R Q is R4

by (simp add: Healthy-def R4-cond)

lemma IterateR-outer-refine-intro:

assumes $I \neq \{\}$ ∧ $i. i \in I \Rightarrow P i$ is NSRD ∧ $i. i \in I \Rightarrow P i$ is Productive

∧ $i. i \in I \Rightarrow S \sqsubseteq (b i \rightarrow_R P i ;; S)$

$S \sqsubseteq [\neg (\bigwedge i \in I \cdot b i)]^\top_R$

shows $S \sqsubseteq \text{do}_R i \in I \cdot b(i) \rightarrow P(i)$ od

apply (simp add: IterateR-def)

apply (rule WhileR-outer-refine-intro)

apply (simp-all add: assms closure AlternateR-assume-branch seq-UINF-distr UINF-refines)

done

lemma IterateR-outer-refine-init-intro:

assumes

$A \neq \{\}$ ∧ $i. i \in A \Rightarrow P i$ is NSRD

∧ $i. i \in A \Rightarrow P i$ is Productive

I is NSRD

$S \sqsubseteq I ;; [\neg (\bigwedge i \in A \cdot b i)]^\top_R$

∧ $i. i \in A \Rightarrow S \sqsubseteq S ;; b i \rightarrow_R P i$

∧ $i. i \in A \Rightarrow S \sqsubseteq I ;; b i \rightarrow_R P i$

shows $S \sqsubseteq I ;; \text{do}_R i \in A \cdot b(i) \rightarrow P(i)$ od

apply (simp add: IterateR-def)

apply (rule-tac WhileR-outer-refine-init-intro)

apply (simp-all add: assms closure AlternateR-assume-branch seq-UINF-distl UINF-refines)

done

lemma IterateR-lemma1:

$[\bigwedge i \in I \cdot b i]^\top_r ;; (\bigwedge i \in I \cdot P i < b i \triangleright_R \text{false}) = (\bigwedge i \in I \cdot [b i]^\top_r ;; P i)$

by (rel-auto; fastforce)

lemma IterateR-lemma2:

assumes $I \neq \{\}$ ∧ $i. i \in I \Rightarrow P(i)$ is RR

shows $([\bigwedge i \in I \cdot b i]_{S<} \Rightarrow_r (\bigwedge i \in I \cdot (P i) < b i \triangleright_R R1 \text{ true}) \wedge \text{false} < (\neg (\bigwedge i \in I \cdot b i))) \triangleright_R R1 \text{ true})$

$= (\bigwedge i \in I \cdot (P i) < b i \triangleright_R R1 \text{ true})$

proof –

from assms(1)

have ($(\bigsqcup i \in I \cdot b\ i)_{S<} \Rightarrow_r (\bigsqcup i \in I \cdot RR(P\ i) \triangleleft b\ i \triangleright_R R1\ true) \wedge false \triangleleft (\neg (\bigsqcup i \in I \cdot b\ i)) \triangleright_R R1\ true)$)
 $= (\bigsqcup i \in I \cdot RR(P\ i) \triangleleft b\ i \triangleright_R R1\ true)$
by (*rel-auto*)
thus *?thesis*
by (*simp add: assms Healthy-if cong: USUP-cong*)
qed

lemma *IterateR-lemma3:*

assumes $\bigwedge i. i \in I \Rightarrow P(i)$ *is* *RR*
shows $(\bigsqcup i \in I \cdot P\ i \triangleleft b\ i \triangleright_R R1\ true) = (\bigsqcup i \in I \cdot [b\ i]_{S<} \Rightarrow_r P\ i)$
proof –
have $(\bigsqcup i \in I \cdot RR(P\ i) \triangleleft b\ i \triangleright_R R1\ true) = (\bigsqcup i \in I \cdot [b\ i]_{S<} \Rightarrow_r RR(P\ i))$
by (*rel-auto*)
thus *?thesis*
by (*simp add: assms Healthy-if cong: USUP-cong*)
qed

theorem *IterateR-refine-intro:*

assumes
– Closure conditions
 $\bigwedge i. i \in I \Rightarrow Q_1(i)$ *is* *RC* $\bigwedge i. i \in I \Rightarrow Q_2(i)$ *is* *RR* $\bigwedge i. i \in I \Rightarrow Q_3(i)$ *is* *RR*
 $\bigwedge i. i \in I \Rightarrow \$st' \# Q_2(i) \bigwedge i. i \in I \Rightarrow Q_3(i)$ *is* *R4* $I \neq \{\}$
 $(\bigsqcup i \in I \cdot [b\ i]_{S<}^{\top_r} ;; Q_3\ i)^{*r} wp_r (\bigsqcup i \in I \cdot [b\ i]_{S<} \Rightarrow_r Q_1\ i) \sqsubseteq P_1$
 $P_2 \sqsubseteq (\bigsqcup i \in I \cdot [b\ i]_{S<}^{\top_r} ;; Q_2\ i)$
 $P_2 \sqsubseteq (\bigsqcup i \in I \cdot [b\ i]_{S<}^{\top_r} ;; Q_3\ i) ;; P_2$
 $P_3 \sqsubseteq [\neg (\bigsqcup i \in I \cdot b\ i)]_{S<}^{\top_r}$
 $P_3 \sqsubseteq (\bigsqcup i \in I \cdot [b\ i]_{S<}^{\top_r} ;; Q_3\ i) ;; P_3$
shows $\mathbf{R}_s(P_1 \vdash P_2 \diamond P_3) \sqsubseteq do_R\ i \in I \cdot b(i) \rightarrow \mathbf{R}_s(Q_1(i) \vdash Q_2(i) \diamond Q_3(i))$ *od*
apply (*simp add: rdes-def closure assms unrest del: WhileR-rdes-def*)
apply (*rule WhileR-refine-intro*)
apply (*simp-all add: closure assms unrest IterateR-lemma1 IterateR-lemma2 segr-assoc[THEN sym]*)
apply (*simp add: IterateR-lemma3 closure assms unrest*)
done

method *unfold-iteration* = *simp add: IterateR-list-def IterateR-def AlternateR-list-def AlternateR-def UINF-upto-expand-first*

11.13 Substitution Laws

lemma *srd-subst-Chaos* [*usubst*]:

$\sigma \dagger_S Chaos = Chaos$
by (*rdes-simp*)

lemma *srd-subst-Miracle* [*usubst*]:

$\sigma \dagger_S Miracle = Miracle$
by (*rdes-simp*)

lemma *srd-subst-skip* [*usubst*]:

$\sigma \dagger_S II_R = \langle \sigma \rangle_R$
by (*rdes-eq*)

lemma *srd-subst-assigns* [*usubst*]:

$\sigma \dagger_S \langle \varrho \rangle_R = \langle \varrho \circ \sigma \rangle_R$
by (*rdes-eq*)

11.14 Algebraic Laws

theorem *assigns-srd-id*: $\langle id \rangle_R = II_R$
by (*rdes-eq*)

theorem *assigns-srd-comp*: $\langle \sigma \rangle_R ;; \langle \varrho \rangle_R = \langle \varrho \circ \sigma \rangle_R$
by (*rdes-eq*)

theorem *assigns-srd-Miracle*: $\langle \sigma \rangle_R ;; Miracle = Miracle$
by (*rdes-eq*)

theorem *assigns-srd-Chaos*: $\langle \sigma \rangle_R ;; Chaos = Chaos$
by (*rdes-eq*)

theorem *assigns-srd-cond* : $\langle \sigma \rangle_R \triangleleft b \triangleright_R \langle \varrho \rangle_R = \langle \sigma \triangleleft b \triangleright_s \varrho \rangle_R$
by (*rdes-eq*)

theorem *assigns-srd-left-seq*:
assumes P is NSRD
shows $\langle \sigma \rangle_R ;; P = \sigma \upharpoonright_S P$
by (*rdes-simp cls: assms*)

lemma *AlternateR-seq-distr*:
assumes $\bigwedge i. A\ i$ is NSRD B is NSRD C is NSRD
shows $(if_R\ i \in I \cdot g\ i \rightarrow A\ i\ else\ B\ fi) ;; C = (if_R\ i \in I \cdot g\ i \rightarrow A\ i ;; C\ else\ B ;; C\ fi)$
proof (*cases* $I = \{\}$)
case *True*
then show ?thesis **by** (*simp*)
next
case *False*
then show ?thesis
by (*simp add: AlternateR-def upred-semiring.distrib-right seq-UINF-distr gcmd-seq-distr assms(3)*)
qed

lemma *AlternateR-is-cond-srea*:
assumes A is NSRD B is NSRD
shows $(if_R\ i \in \{a\} \cdot g \rightarrow A\ else\ B\ fi) = (A \triangleleft g \triangleright_R B)$
by (*rdes-eq cls: assms*)

lemma *AlternateR-Chaos*:
 $if_R\ i \in A \cdot g(i) \rightarrow Chaos\ fi = Chaos$
by (*cases* $A = \{\}$, *simp*, *rdes-eq*)

lemma *choose-srd-par*:
 $choose_R \parallel_R choose_R = choose_R$
by (*rdes-eq*)

11.15 Lifting designs to reactive designs

definition *des-rea-lift* :: $'s\ hrel\ des \Rightarrow ('s, 't :: trace, 'a) hrel\ rsp\ (\mathbf{R}_D)$ **where**
 $[upred\ defs]: \mathbf{R}_D(P) = \mathbf{R}_s([pre_D(P)]_S \vdash (false \diamond (\$tr' =_u \$tr \wedge [post_D(P)]_S)))$

definition *des-rea-drop* :: $('s, 't :: trace, 'a) hrel\ rsp \Rightarrow 's\ hrel\ des\ (\mathbf{D}_R)$ **where**
 $[upred\ defs]: \mathbf{D}_R(P) = \lfloor (pre_R(P))[\$tr/\$tr'] \upharpoonright_v \$st \rfloor_{S<} \vdash_n \lfloor (post_R(P))[\$tr/\$tr'] \upharpoonright_v \{\$st, \$st'\} \rfloor_S$

lemma *ndesign-rea-lift-inverse*: $\mathbf{D}_R(\mathbf{R}_D(p \vdash_n Q)) = p \vdash_n Q$
apply (*simp add: des-rea-lift-def des-rea-drop-def rea-pre-RHS-design rea-post-RHS-design*)
apply (*simp add: R1-def R2c-def R2s-def usubst unrest*)
apply (*rel-auto*)
done

lemma *ndesign-rea-lift-injective*:
assumes P is \mathbf{N} Q is \mathbf{N} $\mathbf{R}_D P = \mathbf{R}_D Q$ (**is** $?RP(P) = ?RQ(Q)$)
shows $P = Q$

proof –

have $?RP(\lfloor pre_D(P) \rfloor_{<} \vdash_n post_D(P)) = ?RQ(\lfloor pre_D(Q) \rfloor_{<} \vdash_n post_D(Q))$
by (*simp add: ndesign-form assms*)
hence $\lfloor pre_D(P) \rfloor_{<} \vdash_n post_D(P) = \lfloor pre_D(Q) \rfloor_{<} \vdash_n post_D(Q)$
by (*metis ndesign-rea-lift-inverse*)
thus *?thesis*
by (*simp add: ndesign-form assms*)

qed

lemma *des-rea-lift-closure* [*closure*]: $\mathbf{R}_D(P)$ is *SRD*
by (*simp add: des-rea-lift-def RHS-design-is-SRD unrest*)

lemma *preR-des-rea-lift* [*rdes*]:
 $pre_R(\mathbf{R}_D(P)) = R1(\lfloor pre_D(P) \rfloor_S)$
by (*rel-auto*)

lemma *periR-des-rea-lift* [*rdes*]:
 $peri_R(\mathbf{R}_D(P)) = (false \triangleleft \lfloor pre_D(P) \rfloor_S \triangleright (\$tr \leq_u \$tr'))$
by (*rel-auto*)

lemma *postR-des-rea-lift* [*rdes*]:
 $post_R(\mathbf{R}_D(P)) = ((true \triangleleft \lfloor pre_D(P) \rfloor_S \triangleright (\neg \$tr \leq_u \$tr')) \Rightarrow (\$tr' =_u \$tr \wedge \lfloor post_D(P) \rfloor_S))$
apply (*rel-auto*) **using** *minus-zero-eq* **by** *blast*

lemma *ndes-rea-lift-closure* [*closure*]:

assumes P is \mathbf{N}
shows $\mathbf{R}_D(P)$ is *NSRD*

proof –

obtain $p Q$ **where** $P = (p \vdash_n Q)$
by (*metis H1-H3-commute H1-H3-is-normal-design H1-idem Healthy-def assms*)
show *?thesis*
apply (*rule NSRD-intro*)
apply (*simp-all add: closure rdes unrest P*)
apply (*rel-auto*)
done

qed

lemma *R-D-mono*:

assumes P is \mathbf{H} Q is \mathbf{H} $P \sqsubseteq Q$
shows $\mathbf{R}_D(P) \sqsubseteq \mathbf{R}_D(Q)$
apply (*simp add: des-rea-lift-def*)
apply (*rule srdes-tri-refine-intro'*)
apply (*auto intro: H1-H2-refines assms aext-mono*)
apply (*rel-auto*)
apply (*metis (no-types, hide-lams) aext-mono assms(3) design-post-choice*)

semilattice-sup-class.sup.orderE utp-pred-laws.inf.coboundedI1 utp-pred-laws.inf commute utp-pred-laws.sup.order-iff

done

Homomorphism laws

lemma *R-D-Miracle*:

$\mathbf{R}_D(\top_D) = \text{Miracle}$
by (*simp add: Miracle-def, rel-auto*)

lemma *R-D-Chaos*:

$\mathbf{R}_D(\perp_D) = \text{Chaos}$

proof –

have $\mathbf{R}_D(\perp_D) = \mathbf{R}_D(\text{false} \vdash_r \text{true})$
by (*rel-auto*)
also have $\dots = \mathbf{R}_s(\text{false} \vdash \text{false} \diamond (\$tr' =_u \$tr))$
by (*simp add: Chaos-def des-rea-lift-def alpha*)
also have $\dots = \mathbf{R}_s(\text{true})$
by (*rel-auto*)
also have $\dots = \text{Chaos}$
by (*simp add: Chaos-def design-false-pre*)
finally show ?thesis .

qed

lemma *R-D-inf*:

$\mathbf{R}_D(P \sqcap Q) = \mathbf{R}_D(P) \sqcap \mathbf{R}_D(Q)$
by (*rule antisym, rel-auto+*)

lemma *R-D-cond*:

$\mathbf{R}_D(P \triangleleft [b]_{D<} \triangleright Q) = \mathbf{R}_D(P) \triangleleft b \triangleright_R \mathbf{R}_D(Q)$
by (*rule antisym, rel-auto+*)

lemma *R-D-seq-ndesign*:

$\mathbf{R}_D(p_1 \vdash_n Q_1) ;; \mathbf{R}_D(p_2 \vdash_n Q_2) = \mathbf{R}_D((p_1 \vdash_n Q_1) ;; (p_2 \vdash_n Q_2))$
apply (*rule antisym*)
apply (*rule SRD-refine-intro*)
apply (*simp-all add: closure rdes ndesign-composition-wp*)
using *dual-order.trans* **apply** (*rel-blast*)
using *dual-order.trans* **apply** (*rel-blast*)
apply (*rel-auto*)
apply (*rule SRD-refine-intro*)
apply (*simp-all add: closure rdes ndesign-composition-wp*)
apply (*rel-auto*)
apply (*rel-auto*)
apply (*rel-auto*)
done

lemma *R-D-seq*:

assumes $P \text{ is } \mathbf{N} \ Q \text{ is } \mathbf{N}$
shows $\mathbf{R}_D(P) ;; \mathbf{R}_D(Q) = \mathbf{R}_D(P ;; Q)$
by (*metis R-D-seq-ndesign assms ndesign-form*)

These laws are applicable only when there is no further alphabet extension

lemma *R-D-skip*:

$\mathbf{R}_D(\text{II}_D) = (\text{II}_R :: ('s, 't::\text{trace}, \text{unit}) \text{ hrel-rsp})$
apply (*rel-auto*) **using** *minus-zero-eq* **by** *blast+*

lemma *R-D-assigns*:

$\mathbf{R}_D(\langle \sigma \rangle_D) = (\langle \sigma \rangle_R :: ('s, 't::\text{trace}, \text{unit}) \text{ hrel-rsp})$
by (*simp add: assigns-d-def des-rea-lift-def alpha assigns-srd-RHS-tri-des, rel-auto*)

11.16 State Invariants

definition *StateInvR* :: '*s* upred \Rightarrow ('*s*, '*t*::trace, ' α) hrel-rsp (*sinv_R*'(-')) **where**
[rdes-def]: *sinv_R*(*b*) = $\mathbf{R}_s([b]_{S<} \vdash \text{true}_r \diamond [b]_{S>})$

lemma *StateInvR-NSRD* [*closure*]: *sinv_R*(*b*) is NSRD
by (*simp add: StateInvR-def closure unrest*)

lemma *StateInvR-srd-skip-refine*: *sinv_R*(*b*) \sqsubseteq *II_R*
by (*rdes-refine*)

lemma *StateInvR-seq-idem*:
sinv_R(*b*) ;; *sinv_R*(*b*) = *sinv_R*(*b*)
by (*rdes-eq*)

lemma *StateInvR-seq-refine*:
assumes *sinv_R*(*b*) \sqsubseteq *P* *sinv_R*(*b*) \sqsubseteq *Q*
shows *sinv_R*(*b*) \sqsubseteq *P* ;; *Q*
by (*metis (full-types) StateInvR-seq-idem assms seqr-mono*)

lemma *ndiv-StateInvR*: *ndiv_R* = *sinv_R*(*true*)
by (*rdes-eq*)

end

12 Instantaneous Reactive Designs

theory *utp-rdes-instant*
imports *utp-rdes-prog*
begin

definition *ISRD1* :: ('*s*, '*t*::trace, ' α) hrel-rsp \Rightarrow ('*s*, '*t*, ' α) hrel-rsp **where**
[upred-defs]: *ISRD1*(*P*) = *P* $\parallel_R \mathbf{R}_s(\text{true}_r \vdash \text{false} \diamond (\$tr' =_u \$tr))$

definition *ISRD* :: ('*s*, '*t*::trace, ' α) hrel-rsp \Rightarrow ('*s*, '*t*, ' α) hrel-rsp **where**
[upred-defs]: *ISRD* = *ISRD1* \circ *NSRD*

lemma *ISRD1-idem*: *ISRD1*(*ISRD1*(*P*)) = *ISRD1*(*P*)
by (*rel-auto*)

lemma *ISRD1-monotonic*: *P* \sqsubseteq *Q* \Longrightarrow *ISRD1*(*P*) \sqsubseteq *ISRD1*(*Q*)
by (*rel-auto*)

lemma *ISRD1-RHS-design-form*:
assumes $\$ok' \# P \ \$ok' \# Q \ \$ok' \# R$
shows *ISRD1*($\mathbf{R}_s(P \vdash Q \diamond R)$) = $\mathbf{R}_s(P \vdash \text{false} \diamond (R \wedge \$tr' =_u \$tr))$
using *assms* **by** (*simp add: ISRD1-def choose-srd-def RHS-tri-design-par unrest, rel-auto*)

lemma *ISRD1-form*:
ISRD1(*SRD*(*P*)) = $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{false} \diamond (\text{post}_R(P) \wedge \$tr' =_u \$tr))$
by (*simp add: ISRD1-RHS-design-form SRD-as-reactive-tri-design unrest*)

lemma *ISR1-rdes-def* [*rdes-def*]:

$\llbracket P \text{ is } RR; R \text{ is } RR \rrbracket \implies \text{ISR1}(\mathbf{R}_s(P \vdash Q \diamond R)) = \mathbf{R}_s(P \vdash \text{false} \diamond (R \wedge \$tr' =_u \$tr))$
 by (*simp add: ISR1-def rdes-def closure rpred*)

lemma *ISR-intro*:

 assumes *P is NSRD* $\text{peri}_R(P) = (\neg_r \text{pre}_R(P)) (\$tr' =_u \$tr) \sqsubseteq \text{post}_R(P)$
 shows *P is ISR*

proof –

 have $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P))$ *is ISR1*
 apply (*simp add: Healthy-def rdes-def closure assms(1-2)*)
 using *assms(3) least-zero* apply (*rel-blast*)
 done
 hence *P is ISR1*
 by (*simp add: SRD-reactive-tri-design closure assms(1)*)
 thus ?thesis
 by (*simp add: ISR-def Healthy-comp assms(1)*)

qed

lemma *ISR1-rdes-intro*:

 assumes *P is RR Q is RR* $(\$tr' =_u \$tr) \sqsubseteq Q$
 shows $\mathbf{R}_s(P \vdash \text{false} \diamond Q)$ *is ISR1*
 unfolding *Healthy-def*
 by (*simp add: ISR1-rdes-def assms closure unrest utp-pred-laws.inf.absorb1*)

lemma *ISR-rdes-intro* [*closure*]:

 assumes *P is RC Q is RR* $(\$tr' =_u \$tr) \sqsubseteq Q$
 shows $\mathbf{R}_s(P \vdash \text{false} \diamond Q)$ *is ISR*
 unfolding *Healthy-def*
 by (*simp add: ISR-def closure Healthy-if ISR1-rdes-def assms unrest utp-pred-laws.inf.absorb1*)

lemma *ISR-implies-ISR1*:

 assumes *P is ISR*
 shows *P is ISR1*

proof –

 have *ISR(P) is ISR1*
 by (*simp add: ISR-def Healthy-def ISR1-idem*)
 thus ?thesis
 by (*simp add: assms Healthy-if*)

qed

lemma *ISR-implies-SRD*:

 assumes *P is ISR*
 shows *P is SRD*

proof –

 have $1:\text{ISR}(P) = \mathbf{R}_s((\neg_r (\neg_r \text{pre}_R P) ;; R1 \text{ true} \wedge R1 \text{ true}) \vdash \text{false} \diamond (\text{post}_R P \wedge \$tr' =_u \$tr))$
 by (*simp add: NSRD-form ISR1-def ISR-def RHS-tri-design-par rdes-def unrest closure*)
 moreover have ... *is SRD*
 by (*simp add: closure unrest*)
 ultimately have *ISR(P) is SRD*
 by (*simp*)
 with *assms* show ?thesis
 by (*simp add: Healthy-def*)

qed

lemma *ISR-implies-NSRD* [*closure*]:

assumes P is ISRD
shows P is NSRD
proof –
 have $1: \text{ISRD}(P) = \text{ISRD1}(\text{RD3}(\text{SRD}(P)))$
 by (simp add: ISRD-def NSRD-def SRD-def, metis RD1-RD3-commute RD3-left-subsumes-RD2)
 also have $\dots = \text{ISRD1}(\text{RD3}(P))$
 by (simp add: assms ISRD-implies-SRD Healthy-if)
 also have $\dots = \text{ISRD1}(\mathbf{R}_s((\neg_r \text{pre}_R P) \text{wp}_r \text{false}_h \vdash (\exists \$st' \cdot \text{peri}_R P) \diamond \text{post}_R P))$
 by (simp add: RD3-def, subst SRD-right-unit-tri-lemma, simp-all add: assms ISRD-implies-SRD)
 also have $\dots = \mathbf{R}_s((\neg_r \text{pre}_R P) \text{wp}_r \text{false}_h \vdash \text{false} \diamond (\text{post}_R P \wedge \$tr' =_u \$tr))$
 by (simp add: RHS-tri-design-par ISRD1-def unrest choose-srd-def rpred closure ISRD-implies-SRD
 assms)
 also have $\dots = (\dots ;; II_R)$
 by (rdes-simp, simp add: RHS-tri-normal-design-composition' closure assms unrest ISRD-implies-SRD
 wp rpred wp-rea-false-RC)
 also have \dots is RD3
 by (simp add: Healthy-def RD3-def segr-assoc)
 finally show ?thesis
 by (simp add: SRD-RD3-implies-NSRD Healthy-if assms ISRD-implies-SRD)
qed

lemma *ISRD-form*:

assumes P is ISRD
shows $\mathbf{R}_s(\text{pre}_R(P) \vdash \text{false} \diamond (\text{post}_R(P) \wedge \$tr' =_u \$tr)) = P$
proof –

have $P = \text{ISRD1}(P)$
 by (simp add: ISRD-implies-ISRD1 assms Healthy-if)
 also have $\dots = \text{ISRD1}(\mathbf{R}_s(\text{pre}_R(P) \vdash \text{peri}_R(P) \diamond \text{post}_R(P)))$
 by (simp add: SRD-reactive-tri-design ISRD-implies-SRD assms)
 also have $\dots = \mathbf{R}_s(\text{pre}_R(P) \vdash \text{false} \diamond (\text{post}_R(P) \wedge \$tr' =_u \$tr))$
 by (simp add: ISRD1-rdes-def closure assms)
 finally show ?thesis ..
qed

lemma *ISRD-elim* [RD-elim]:

$\llbracket P \text{ is ISRD}; Q(\mathbf{R}_s(\text{pre}_R(P) \vdash \text{false} \diamond (\text{post}_R(P) \wedge \$tr' =_u \$tr))) \rrbracket \implies Q(P)$
 by (simp add: ISRD-form)

lemma *skip-srd-ISRD* [closure]: II_R is ISRD

by (rule ISRD-intro, simp-all add: rdes closure)

lemma *assigns-srd-ISRD* [closure]: $\langle \sigma \rangle_R$ is ISRD

by (rule ISRD-intro, simp-all add: rdes closure, rel-auto)

lemma *seq-ISRD-closed*:

assumes P is ISRD Q is ISRD
shows $P ;; Q$ is ISRD
apply (insert assms)
apply (erule ISRD-elim)+
apply (simp add: rdes-def closure assms unrest)
apply (rule ISRD-rdes-intro)
 apply (simp-all add: rdes-def closure assms unrest)
apply (rel-auto)
done

```

lemma ISRD-Miracle-right-zero:
  assumes  $P$  is ISRD  $\text{pre}_R(P) = \text{true}_r$ 
  shows  $P \;; \text{Miracle} = \text{Miracle}$ 
  by (rdes-simp cls: assms)

```

A recursion whose body does not extend the trace results in divergence

```

lemma ISRD-recurse-Chaos:
  assumes  $P$  is ISRD  $\text{post}_R P \;; \text{true}_r = \text{true}_r$ 
  shows  $(\mu_R X \cdot P \;; X) = \text{Chaos}$ 
proof –
  have 1:  $(\mu_R X \cdot P \;; X) = (\mu X \cdot P \;; \text{SRD}(X))$ 
    by (auto simp add: sdes-theory-continuous.utp-lfp-def closure assms)
  have  $(\mu X \cdot P \;; \text{SRD}(X)) \sqsubseteq \text{Chaos}$ 
proof (rule gfp-upperbound)
  have  $P \;; \text{Chaos} \sqsubseteq \text{Chaos}$ 
    apply (rdes-refine-split cls: assms)
    using assms(2) apply (rel-auto, metis (no-types, lifting) dual-order.antisym order-refl)
    apply (rel-auto) +
    done
  thus  $P \;; \text{SRD Chaos} \sqsubseteq \text{Chaos}$ 
    by (simp add: Healthy-if sdes-theory-continuous.bottom-closed)
qed
thus ?thesis
  by (metis 1 dual-order.antisym sdes-theory-continuous.LFP-closed sdes-theory-continuous.bottom-lower)
qed

```

```

lemma recursive-assign-Chaos:
   $(\mu_R X \cdot (\sigma)_R \;; X) = \text{Chaos}$ 
  by (rule ISRD-recurse-Chaos, simp-all add: closure rdes, rel-auto)

```

end

13 Meta-theory for Reactive Designs

```

theory utp-rea-designs
imports
  utp-rdes-healths
  utp-rdes-designs
  utp-rdes-triples
  utp-rdes-normal
  utp-rdes-contracts
  utp-rdes-tactics
  utp-rdes-parallel
  utp-rdes-prog
  utp-rdes-instant
  utp-rdes-guarded
begin end

```

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