

# Cmpe-300

Ahmet Hacıoğlu

November 2023

## 1 Introduction

### Q2

The pseudocode describes a set of nested loops, and we'll analyze the number of iterations in each loop.

1. The outer loop runs  $n$  times ( $i$  goes from 1 to  $n$ ). 2. The middle loop runs  $i^2$  times for each iteration of the outer loop ( $j$  goes from 1 to  $i^2$ ). 3. The innermost loop runs  $i$  times for each iteration of the middle loop ( $k$  goes from 1 to  $i$ ).

So, the total number of iterations ( $a$ ) is given by the sum:

$$a = \sum_{i=1}^n \sum_{j=1}^{i^2} \sum_{k=1}^i 1$$

$$a = \sum_{i=1}^n \sum_{j=1}^{i^2} i$$

$$a = \sum_{i=1}^n i \cdot i^2$$

$$a = \sum_{i=1}^n i^3$$

$$a = \left( \frac{n(n+1)}{2} \right)^2$$

Now, let's find the worst-case time complexity. The dominant term in the formula is  $n^4$ , so the worst-case time complexity is  $O(n^4)$ .  $B(n) = W(n) = A(n) = f(n) \in O(n^4)$

$$\lim_{n \rightarrow \infty} \frac{n^4 \cdot 4}{n^2 \cdot (n+1)^2} = \lim_{n \rightarrow \infty} \frac{4}{\left(1 + \frac{1}{n}\right)^2} = 4(\text{constant})$$