

inzva Algorithm Study Group

Week 8

Dynamic Programming

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Today's Contents

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3.	<u>Do it once: Suffix Query</u>	12.40 - 13.00	<u>full solution here</u>
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4.	<u>Memoization: Basic-Unbounded-Knapsack</u>	13.15 - 13.35	<u>full solution here</u>
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What is Dynamic Programming?

Dynamic Programming is mainly an **optimization** over plain [recursion](#). Wherever we see a recursive solution that has repeated calls for same inputs, we can optimize it using Dynamic Programming.[1]

Steps to solve a DP

- 1) Identify if it is a DP problem
- 2) Decide a state expression with least parameters
- 3) Formulate state relationship
- 4) Do tabulation (or add memoization)[1]

We will also cover Greedy Approaches this week, which is [different](#) from dynamic programming.

[1]: <https://www.geeksforgeeks.org/dynamic-programming/> (this is very good material)



1- Mark's Cakewalk

This is a greedy type of problem.

Complete the marcsCakewalk function below.

```
def marcsCakewalk(calorie):
    calorie.sort(reverse=True)
    mysum = 0
    for i, cal in enumerate(calorie):
        mysum += 2**i * cal
    return mysum

if __name__ == '__main__':
    fptr = open(os.environ['OUTPUT_PATH'], 'w')

    n = int(input())

    calorie = list(map(int, input().rstrip().split()))

    result = marcsCakewalk(calorie)

    fptr.write(str(result) + '\n')

    fptr.close()
```

Task: Consume cakes to get minimum calories

Important: Each cake gives $2^{\text{(position)}} * \text{calorie}$ calories

Inputs: Calorie values of cakes

Example input:

2 3 4



2- Onur and the Suffix Query

This is a do-it-once (?) type of problem. We process our input once, and note down the output, instead of for each query.

```
def count_former(recurrent_list):
    memory = []
    unique_count_list = []
    count = 0
    for item in recurrent_list[::-1]:
        if not item in memory:
            count = count + 1
            memory.append(item)
            unique_count_list.append(count)
        #print("inside ",item,count)
    return unique_count_list[::-1], count
```

```
n,m = map(int,input().split())
recurrent_list = list(map(int,input().split()))
```

```
unique_count_list, last_count = count_former(recurrent_list)
for _ in range(m):
    position = int(input())
    print(unique_count_list[position-1])
```

Task: Find number of distinct integers from position i to the end of the sequence

Inputs: The sequence, and queries consisting of beginning position i .

Example input:

```
1 5 3 4 1 5 3 4 100000 88888
1
2
9
10
```

(sequence)
(query for $i=1$)
(query for $i=2$)
(query for $i=9$)
(query for $i=10$)



3- Basic-Unbounded-Knapsack

This is a memoization problem. In a recursive process, we note down some results to cut back computation.

```
if __name__ == '__main__':  
    fptr = open(os.environ['OUTPUT_PATH'], 'w')  
  
    t = int(input())  
  
    for i in range(t):  
        nk = input().split()  
  
        n = int(nk[0])  
  
        k = int(nk[1])  
  
        arr = list(map(int, input().rstrip().split()))  
  
        memory = {}  
        result = unboundedKnapsack(k, arr)  
  
        fptr.write(str(result) + '\n')  
  
    fptr.close()
```

Task: We have a money limit to spend. We also have banknotes. Spend most money with the banknotes, without exceeding the limit.

Inputs: The sequence, and queries consisting of beginning position i.

Example input:

12	(money limit)
1 6 9	(banknotes)



3- Basic-Unbounded-Knapsack

Our recursive function:

Complete the unboundedKnapsack function below.

```
sys.setrecursionlimit(10000)
```

```
def unboundedKnapsack(k, arr):
```

```
    if k in memory:
```

```
        return memory[k]
```

```
    chosen_sum = 0
```

```
    for item in arr:
```

```
        candidate_sum = 0
```

```
        if item <= k:
```

```
            candidate_sum = unboundedKnapsack(k-item, arr) + item
```

```
        else:
```

```
            continue
```

```
    if candidate_sum > chosen_sum:
```

```
        chosen_sum = candidate_sum
```

```
    memory[k] = chosen_sum
```

```
    return chosen_sum
```



4- LCS: Longest Common Subsequence

This is a memoization problem. In a recursive process, we note down some results to cut back computation.

```
n, m = map(int, input().split())
a = list(map(int, input().rstrip().split()))
b = list(map(int, input().rstrip().split()))
```

```
memory = {}
backtrack = {}
length = longestCommonSubsequence(0, 0)
printer()
```

Task: We have two sequences. Find the longest subsequence that exists in both of the sequences.

Important: A subsequence is formed by deleting elements from the original sequence

Inputs: First sequence and second sequence

Example input:

1 2 3 4 1	(first sequence)
3 4 1 2 1 3	(second sequence)

Example to explain substrings:

Original:	1 2 3 4 1
A substring of original:	1 2 3
Another substring of original:	1 1



4- LCS: Longest Common Subsequence

Our recursive function:

```
def longestCommonSubsequence(i, j):  
    global a,b,n,m  
  
    if i>=n or j>=m:  
        return 0  
  
    if (i,j) in memory:  
        return memory[i,j]  
  
    if a[i] == b[j]:  
        memory[i,j] = 1 + longestCommonSubsequence(i+1,j+1)  
        backtrack[i,j] = 1,1  
        return memory[i,j]  
  
    left = longestCommonSubsequence(i+1,j)  
    right = longestCommonSubsequence(i, j+1)  
  
    if left>=right:  
        chosen = left  
        backtrack[i,j] = 1,0  
    elif right>left:  
        chosen = right  
        backtrack[i,j] = 0,1  
  
    memory[i,j] = chosen  
    return memory[i,j]
```



4- LCS: Longest Common Subsequence

Our print function:

```
def printer():  
    i = 0  
    j = 0  
    while(1):  
        if i+1>n or j+1>m:  
            break  
  
        if a[i] == b[j]:  
            print(a[i], end = ' ')  
  
        if (i,j) in backtrack:  
            myop = backtrack[i,j]  
            i = i + int(myop[0])  
            j = j + int(myop[1])  
        else:  
            break
```



5- LIS: Longest Increasing Subsequence

This is a memoization problem. In a recursive process, we save some results to cut back computation.

```
seq_len = int(input())
sequence = list(map(int, input().split()))
```

```
dp_memory = {}
```

```
maxi = 0
current = 0
```

```
for i in range(seq_len):
    current = LIS(i)
    if current > maxi:
        maxi = current
```

```
print(maxi)
```

In this question, we do not call all of the solution-space recursively. We also use two for loops.

This is for saving memory, and for ease of implementation.

Task: We have a sequence. Find the longest subsequence. The sequence has to be also strictly increasing.

Important: A subsequence is formed by deleting elements from the original sequence

Inputs: The sequence

Example input:

15 27 14 38 26 55 46 65 85

(the sequence)



5- LIS: Longest Increasing Subsequence

Our recursive function:

```
def LIS(start_index):  
    if start_index in dp_memory:  
        return dp_memory[start_index]
```

```
maxi_local = 0  
current_local = 0
```

```
for i in range(start_index+1, seq_len):  
    if(sequence[i]>sequence[start_index]):  
        current_local = LIS(i)  
        if current_local > maxi_local:  
            maxi_local = current_local  
dp_memory[start_index] = maxi_local + 1  
return maxi_local + 1
```

```
maxi = 0  
current = 0
```

