
MODULE *MoneyTransfer*

EXTENDS *Integers*

VARIABLES *accountA*, *accountB*, *money*, *pc*

CONSTANTS *A*, *B*, *M*

PositiveInteger(*n*) $\triangleq n \in \text{Nat} \wedge n \neq 0$

The “money” is transferred from account *A* to account *B*. The “*pc*” variable is used to control the program flow. Remember, there is no order of execution between “Subtract” and “Add”

TypeInvariant $\triangleq \wedge pc \in \{\text{“SUB”}, \text{“ADD”}, \text{“DONE”}\}$
 $\wedge accountA \in \text{Nat}$
 $\wedge accountB \in \text{Nat}$
 $\wedge PositiveInteger(money)$ no zero transfer

Init $\triangleq \wedge accountA = A$
 $\wedge accountB = B$
 $\wedge money = M$
 $\wedge pc = \text{“SUB”}$

The “money” is subtracted from account *A* if it has enough funds. After that the flow is passed to “Add” by changing “*pc*” values to “ADD”. If the balance is not enough, there is no next state, preventing from adding “money” to account *B*.

Subtract $\triangleq \wedge pc = \text{“SUB”}$
 $\wedge \text{IF } accountA \geq money$
 $\quad \text{THEN } accountA' = accountA - money \wedge pc' = \text{“ADD”}$
 $\quad \text{ELSE UNCHANGED } \langle accountA \rangle \wedge pc' = \text{“DONE”}$
 $\wedge \text{UNCHANGED } \langle accountB, money \rangle$

The “money” is added to account *B* and after that the flow is changing to “DONE”, meaning that there is no next state.

Add $\triangleq \wedge pc = \text{“ADD”}$
 $\wedge accountB' = accountB + money$
 $\wedge pc' = \text{“DONE”}$
 $\wedge \text{UNCHANGED } \langle accountA, money \rangle$

Next $\triangleq Subtract \vee Add$

vars $\triangleq \langle accountA, accountB, money, pc \rangle$

Spec $\triangleq Init \wedge \Box[Next]_{\langle vars \rangle}$
