Team notebook

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1 Data structures

1.1 Centroid decomposition

```
namespace decomposition {
 int cnt[N], depth[N], f[N]; // if depth != 0
     means was removed from the tree while
     decomposing
 int dfs (int u, int p = -1) {
   cnt[u] = 1:
   for (int v : g[u])
     if (!depth[v] && v != p)
       cnt[u] += dfs(v, u);
   return cnt[u];
 int get_centroid (int u, int r, int p = -1) {
   for (int v : g[u])
     if (!depth[v] && v != p && cnt[v] > r)
       return get_centroid(v, r, u);
   return u;
 }
 // depth of all tree's centroid is 1 and father
     is 0 (you can set up this when call this)
 int decompose (int u, int d = 1) {
   int centroid = get_centroid(u, dfs(u) >> 1);
   depth[centroid] = d; // remove this node from
       component
   /// add here magic function to count
       properties on paths
   for (int v : g[centroid])
     if (!depth[v])
       f[decompose(v, d + 1)] = centroid;
   return centroid;
```

```
}
int lca (int u, int v) { // lca on centroid tree
  for (; u != v; u = f[u])
    if (depth[v] > depth[u])
      swap(u, v);
  return u;
}
```

1.2 Fenwick tree

1.3 Heavy light decomposition

```
/// Complexity: O(|N|)
int idx; /// top is father of the chain, up is
    father of a node
vector<int> len, depth, in, out, top, up;
int dfs_len( int u, int p, int d ) {
    up[u] = p; depth[u] = d;
    int sz = 1;
    for( auto& v : g[u] ) {
        if( v == p ) continue;
        sz += dfs_len(v, u, d+1);
        if(len[ g[u][0] ] <= len[v]) swap(g[u][0], v);
    }
    return len[u] = sz;</pre>
```

```
void dfs_hld( int u, int p = 0 ) {
 in[u] = idx++;
 narr[ in[u] ] = val[u]; /// to initialize the
      segment tree
 for( auto& v : g[u] ) {
   if( v == p ) continue;
   top[v] = (v == g[u][0] ? top[u] : v);
   dfs_hld(v, u);
 out[u] = idx-1;
void update_hld( int u, int val ) {
 update_DS(in[u], val);
data query_hld( int u, int v ) {
 data val = NULL_DATA;
 while( top[u] != top[v] ) {
   if( depth[ top[u] ] < depth[ top[v] ] )</pre>
       swap(u, v);
   val = val+query_DS(in[ top[u] ], in[u]);
   u = up[ top[u] ];
 if( depth[u] > depth[v] ) swap(u, v);
 val = val+query_DS(in[u], in[v]);
 return val;
/// when updates are on edges use:
/// val[v] is cost_edge(up[v], v), mind root's
/// if(depth[u] == depth[v]) return val;
/// val = val+query_DS(in[u] + 1, in[v]);
}
void build(int n, int root) {
 top = len = in = out = up = depth =
      vector<int>(n+1):
 idx = 1; /// DS index [1, n]
 dfs_len(root, root, 0);
 top[root] = root;
 dfs_hld(root, root);
 /// initialize DS
```

1.4 Iterative Segment Tree

```
/// Complexity: O(|N|*log|N|)
struct info { int val; };
info merge(info &a, info &b) {
 return {a.val + b.val};
info NEUTRAL = {0};
struct segtree { // for point update and range
    queries, supports left to right merge
 int n; /// 0-indexed
 vector<info> t;
 segtree(int n, vector<int> &v) : n(n), t(2*n) {
   for(int i = 0; i < n; i++) t[i+n] = {v[i]};
   for(int i = n-1; i > 0; --i) t[i] =
        merge(t[i<<1], t[i<<1|1]):
 }
 info query(int 1, int r) {
   info ans_1 = NEUTRAL, ans_r = NEUTRAL;
   for(1 += n, r += n+1; 1 < r; 1 >>= 1, r >>=
       1) {
     if(1\&1) ans_1 = merge(ans_1, t[1++]);
     if(r&1) ans_r = merge(t[--r], ans_r);
   return merge(ans_1, ans_r);
 void modify(int p, int x) {
   for(t[p += n] = \{x\}; p >>= 1; ) t[p] =
        merge(t[p<<1], t[p<<1|1]);
 }
};
```

1.5 Mo's

```
/// Complexity: O(|N+Q|*sqrt(|N|)*|ADD/DEL|)
// Requires add(), delete() and get_ans()
struct query {
  int 1, r, idx;
  query (int 1, int r, int idx) : l(l), r(r),
      idx(idx) {}
};
```

```
int S; // s = sqrt(n)
bool cmp (const query &a, const query &b) {
   int A = a.1/S, B = b.1/S;
   if (A != B) return A < B;
   return A & 1 ? a.r > b.r : a.r < b.r;
}
S = sqrt(n); // n = size of array
sort(q.begin(), q.end(), cmp);
int 1 = 0, r = -1;
for (int i = 0; i < q.size(); ++i) {
   while (r < q[i].r) add(++r);
   while (1 > q[i].l) add(--l);
   while (r > q[i].r) del(r--);
   while (1 < q[i].l) del(1++);
   ans[q[i].idx] = get_ans();
}</pre>
```

1.6 Order statistics

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
typedef tree<int, null_type, less<int>,
    rb_tree_tag,
tree_order_statistics_node_update> ordered_set;
//methods
tree.find_by_order(k) //returns pointer to the
    k-th smallest element
tree.order_of_key(x) //returns how many elements
    are smaller than x
//if element does not exist
tree.end() == tree.find_by_order(k) //true
```

1.7 Persistent segment tree

```
/// Complexity: O(|N|*log|N|)
struct node {
  node *left, *right;
  int v;
  node() : left(this), right(this), v(0) {}
```

3

1.8 Rmq

```
/// Complexity: O(|N|*log|N|)
struct rmq {
 vector<vector<int>> table;
  rmg(vector<int> &v) : table(20,
      vector<int>(v.size())) {
   int n = v.size();
   for (int i = 0; i < n; i++) table[0][i] =</pre>
        v[i]:
   for (int j = 1; (1<<j) <= n; j++)
     for (int i = 0; i + (1 << j-1) < n; i++)
       table[j][i] = min(table[j-1][i],
            table[j-1][i + (1 << j-1)]);
  int query(int a, int b) {
   int j = 31 - __builtin_clz(b-a+1);
   return min(table[j][a], table[j][b-(1<<j)+1]);</pre>
};
```

1.9 Sack

```
/// Complexity: |N|*log(|N|)
int dfs(int u, int p = -1) {
  who[t] = u; fr[u] = t++;
```

```
4
```

```
pii best = \{0, -1\};
  int sz = 1;
 for(auto v : g[u])
   if(v != p) {
     int cur_sz = dfs(v, u);
     sz += cur sz:
     best = max(best, {cur_sz, v});
  to[u] = t-1;
  big[u] = best.second;
 return sz;
void add(int u, int x) { /// x == 1 add, x == -1
    delete
  cnt[u] += x;
void go(int u, int p = -1, bool keep = true) {
 for(auto v : g[u])
   if(v != p && v != big[u])
     go(v, u, 0);
  if(big[u] != -1) go(big[u], u, 1);
 /// add all small
 for(auto v : g[u])
   if(v != p && v != big[u])
     for(int i = fr[v]; i <= to[v]; i++)</pre>
       add(who[i], 1);
  add(u, 1):
  ans[u] = get(u);
  if(!keep)
   for(int i = fr[u]; i <= to[u]; i++)</pre>
      add(who[i], -1);
}
void solve(int root) {
 t = 0;
 dfs(root):
 go(root);
```

1.10 Segment tree 2D

```
/// Complexity: \log(|N|)^2 per operation struct info {
```

```
ll val:
};
info merge(info a, info b) {
 return {a.val > b.val ? a.val : b.val};
info NEUTRAL = {LLONG_MIN};
struct segtree_2d {
 int n, m; /// 0-indexed
 vector<vector<info>> t;
 segtree_2d(int n, int m) : n(n), m(m), t(2*n,
      vector<info>(2*m, NEUTRAL)) {}
  segtree_2d(int n, int m, vector<vector<info>>
      &v): n(n), m(m), t(2*n), vector<info>(2*m)
      NEUTRAL)) {
   for(int i = 0; i < n; i++)</pre>
     for(int j = 0; j < m; j++)</pre>
       t[i+n][j+m] = v[i][j];
   for(int i = 0; i < n; i++)</pre>
     for(int j = m-1; j; j--)
       t[i+n][j] = merge(t[i+n][j<<1],
           t[i+n][j<<1|1]);
   for(int i = n-1; i; i--)
     for(int j = 0; j < 2*m; j++)
       t[i][j] = merge(t[i << 1][j], t[i << 1|1][j]);
 info get(int x1, int y1, int x2, int y2) {
   info ans = NEUTRAL:
   vector<int> pos(2);
   for(x1 += n, x2 += n+1; x1 < x2; x1 >>= 1, x2
        >>= 1) {
     int q = 0;
     if(x1\&1) pos[q++] = x1++;
     if(x2\&1) pos[q++] = --x2;
     for(int i = 0; i < q; i++) {</pre>
       for(int t1 = m+y1, t2 = m+y2+1, id =
           pos[i]; t1 < t2; t1 >>= 1, t2 >>= 1) {
         if(t1\&1) ans = merge(ans, t[id][t1++]);
         if(t2\&1) ans = merge(ans, t[id][--t2]);
     }
   }
   return ans;
  void update(int x, int y, info v) {
```

```
t[x+n][y+m] = v;
for(int j = y+m; j > 1; j >>= 1)
    t[x+n][j>>1] = merge(t[x+n][j], t[x+n][j^1]);
for(int i = x+n; i > 1; i >>= 1)
    for(int j = y+m; j; j >>= 1)
     t[i>>1][j] = merge(t[i][j],t[i^1][j]);
}
};
```

1.11 Treap

```
/// Complexity: log(|N|) per operation
uniform_int_distribution<11> rnd(0, LLONG_MAX);
typedef long long T;
struct treap {
 treap *left, *right, *father;
  ll prior:
  int sz, idx;
 T value, lazy_sum, sum;
  treap(T x) {
   left = right = father = NULL;
   prior = rnd(rng);
   sz = 1;
   value = sum = x;
   lazy_sum = 0;
};
int cnt(treap* t) { return !t ? 0 : t->sz; }
T sum(treap* t) { return !t ? 0 : t->sum; }
T value(treap* t) { return !t ? 0 : t->value; }
void propagate(treap* t) {
 if(t && t->lazy_sum) {
   if(t->left) t->left->lazy_sum += t->lazy_sum;
   if(t->right) t->right->lazy_sum +=
        t->lazy_sum;
   t->sum += cnt(t) * t->lazy_sum;
   t->value += t->lazy_sum;
   t->lazy_sum = 0;
}
void update(treap *t) {
 propagate(t->left);
```

```
propagate(t->right);
 t->sz = cnt(t->left) + cnt(t->right) + 1;
 t->sum = sum(t->left) + sum(t->right) +
      value(t):
 if(t->left) t->left->father = t;
 if(t->right) t->right->father = t;
}
void add_value(treap *t, T v) {
 t->value += v;
 update(t):
void add_lazy_sum(treap *t, T v) {
 t->lazv_sum += v;
 update(t):
pair<treap*, treap*> split(treap* t, int
    left count) {
 if(!t) return {NULL, NULL};
 propagate(t);
 if(cnt(t->left) >= left_count) {
   auto got = split(t->left, left_count);
   t->left = got.second;
   update(t);
   return {got.first, t};
 } else {
   left_count = left_count-cnt(t->left)-1;
   auto got = split(t->right, left_count);
   t->right = got.first;
   update(t):
   return {t, got.second};
 }
}
treap* merge(treap *s, treap *t) {
 if(!s) return t;
 if(!t) return s:
 propagate(s);
 propagate(t);
 if(s->prior <= t->prior) {
   s->right = merge(s->right, t);
   update(s);
   return s;
 } else {
   t->left = merge(s, t->left);
   update(t);
```

```
return t:
 }
};
void print(treap *x) {
 if(!x) return;
 propagate(x);
 print(x->left);
 cout << value(x) << ", ";</pre>
 print(x->right);
int find_left_count(treap* root, treap* x) { ///
    x not inclusive
 if(!x) return 0;
 int ans = cnt(x->left):
 while(x != root) {
   treap *par = x->father;
   if(par->right == x) ans += cnt(par->left)+1;
   x = par;
 }
 return ans;
treap *root = NULL;
```

1.12 lichaor

```
typedef long long type;
struct line {
 type m, b;
 type eval (type x) {
   return m * x + b;
 }
/// M = 4 * (number of updates) * (log R) where R
    is the range of the segtree
/// oo is the max abs value that the function may
    take
const type R = 1e9+5;
const int M = 1e7;
const type oo = 2e18;
struct li chao {
 /// see comments for min
 int nodes;
```

```
line lines[M]:
int lf[M], rg[M];
void ini () {
 nodes = 0:
 memset(lf, -1, sizeof lf);
 memset(rg, -1, sizeof rg);
 lines[0] = \{0, -\infty\}; /// \text{ change to } \{0, \infty\};
int add_line (int node, type 1, type r, line
    nw) {
 if (node == -1) {
   lines[++nodes] = nw;
   return nodes;
 type m = (1 + r) >> 1;
 bool lef = nw.eval(1) > lines[node].eval(1):
      ///change > to <
 bool mid = nw.eval(m) > lines[node].eval(m);
      ///change > to <
 if (mid) swap(nw, lines[node]);
 lines[++nodes] = lines[node];
 lf[nodes] = lf[node];
 rg[nodes] = rg[node];
 int sv = nodes:
 if (r == 1) return sv;
 if (lef != mid) lf[sv] = add_line(lf[node],
      1. m. nw):
 else rg[sv] = add_line(rg[node], m + 1, r,
 return sv;
type get(int node, type 1, type r, type x) {
 if (node == -1) return -oo; //change to oo
 type m = (1 + r) / 2;
 if(1 == r) {
     return lines[node].eval(x);
 else if(x < m) 
     return max(lines[node].eval(x),
         get(lf[node], 1, m, x)); ///change max
         to min
 } else {
     return max(lines[node].eval(x),
         get(rg[node], m + 1, r, x)); ///change
         max to min
```

```
}
}
};
```

2 Dp optimization

2.1 Convex hull trick dynamic

```
/// Complexity: O(|N|*log(|N|))
typedef 11 T;
const T is_query = -(1LL<<62);</pre>
struct line {
       T m. b:
       mutable multiset<line>::iterator it, end;
       bool operator < (const line &rhs) const {</pre>
               if(rhs.b != is_query) return m <</pre>
                   rhs.m:
               auto s = next(it);
               if(s == end) return 0:
               return b - s->b < (long
                   double)(s->m - m) * rhs.m;
       }
};
struct CHT : public multiset<line> {
       bool bad(iterator y) {
               auto z = next(y);
               if(v == begin()) {
                      if(z == end()) return false;
                      return y->m == z->m && y->b
                           \leq z->b:
               auto x = prev(y);
               if(z == end()) return y->m == x->m
                   && y->b == x->b;
               return (long double)(x->b -
                   y->b)*(z->m - y->m) >= (long)
                   double) (y->b - z->b)*(y->m -
                   x->m);
       }
       void add(T m, T b) {
               auto y = insert({m, b});
```

```
y->it = y; y->end = end();
    if(bad(y)) { erase(y); return; }
    while(next(y) != end() &&
        bad(next(y))) erase(next(y));
    while(y != begin() &&
        bad(prev(y)))erase(prev(y));
}
T eval(T x) { /// for maximum
    auto 1 = *lower_bound({x,
        is_query});
    return 1.m*x+1.b;
}
}; // for minimum, you must change (b, m) to (-b,
    -m)
```

2.2 Convex hull trick

```
struct line {
 ll m, b;
 11 eval (11 x) { return m*x + b; }
struct cht {
 vector<line> lines;
 vector<lf> inter;
 int n;
 lf get_inter(line &a, line &b) { return lf(b.b
     -a.b) / (a.m - b.m); }
 inline bool ok(line &a, line &b, line &c) {
   return lf(a.b-c.b) / (c.m-a.m) > lf(a.b-b.b)
       / (b.m-a.m):
 void add(line 1) {
   n = lines.size();
   if(n && lines.back().m == 1.m &&
       lines.back().b >= 1.b) return:
   if(n == 1 && lines.back().m == 1.m &&
       lines.back().b < 1.b) lines.pop_back(),</pre>
   while (n \ge 2 \&\& !ok(lines[n-2], lines[n-1],
       1)) {
     n--:
     lines.pop_back(); inter.pop_back();
```

```
}
lines.push_back(1); n++;
if(n >= 2)
    inter.push_back(get_inter(lines[n-1],
        lines[n-2]));
}
ll get_max(lf x) {
    if(lines.size() == 0) return LLONG_MIN;
    if(lines.size() == 1) return lines[0].eval(x);
    int pos = lower_bound(inter.begin(),
        inter.end(), x) - inter.begin();
    return lines[pos].eval(x);
}
}; // for max: order slops non-decreasing, for
    min: order slops non-descending
```

2.3 Divide and conquer

```
/// Complexity: O(|N|*|K|*log|N|)
/// ****** Theory ******
/// dp[k][i]=min(dp[k1][j]+C[i][j]), j < i
/// opt[k][i] opt[k][i+1].
/// A sufficient (but not necessary) condition
    for above is
/// C[a][c] + C[b][d] C[a][d] + C[b][c]
    where a < b < c < d.
void go(int k, int l, int r, int opl, int opr) {
 if(1 > r) return;
 int mid = (1 + r) / 2, op = -1;
 11 &best = dp[mid][k];
 best = INF;
 for(int i = min(opr, mid); i >= opl; i--) {
   ll cur = dp[i][k-1] + cost(i+1, mid);
   if(best > cur) {
     best = cur;
     op = i;
 }
 go(k, l, mid-1, opl, op);
 go(k, mid+1, r, op, opr);
```

7

2.4 Knuth

```
/// Complexity: O(|N|^2)
/// ****** Theory ******
/// dp[i][j] = min(dp[i][k]+dp[k][j])+C[i][j],
    i<k<i
/// where opt[i][j1]
                          opt[i][i] opt[i+1][i].
/// sufficient (but not necessary) condition for
    above is
/// C[a][c] + C[b][d]
                            C[a][d] + C [b][c]
    and C[b][c] C[a][d] where
                abcd
for(int i = 1; i <= n; i++) {</pre>
 opt[i][i] = i;
 dp[i][i] = sum[i] - sum[i-1];
for(int len = 2; len <= n; len++)</pre>
 for(int l = 1; l+len-1 <= n; l++) {</pre>
   int r = 1 + len - 1:
   dp[1][r] = oo;
   for(int i = opt[1][r-1]; i <= opt[1+1][r];</pre>
        i++) {
     ll cur = dp[l][i-1] + dp[i+1][r] + sum[r] -
          sum[1-1];
     if(cur < dp[l][r]) {</pre>
       dp[1][r] = cur;
       opt[1][r] = i;
     }
   }
 }
```

3 Formulas

3.1 Burnside's lemma

$$\#orbitas = \frac{1}{|G|} \sum_{g \in G} |fix(g)|$$

1. **G**: Las acciones que se pueden aplicar sobre un elemento, incluyendo la identidad, eg. Shift 0 veces, Shift 1 veces...

- 2. **Fix(g)**: Es el número de elementos que al aplicar *g* vuelven a ser ser ellos mismos
- 3. Órbita: El conjunto de elementos que pueden ser iguales entre si al aplicar alguna de las acciones de G

3.2 Combinatorics

• Hockey-stick identity $\sum_{i=r}^{n} {i \choose r} = {n+1 \choose r+1}$

3.3 Law of sines and cosines

- a, b, c: lenghts, A, B, C: opposite angles, d: circumcircle
- $\frac{a}{\sin(A)} = \frac{b}{\sin(B)} = \frac{c}{\sin(C)} = d$
- $c^2 = a^2 + b^2 2ab\cos(C)$

3.4 Pythagorean triples $(a^2 + b^2 = c^2)$

- Given an arbitrary pair of integers m and n with m > n > 0: $a = m^2 n^2, b = 2mn, c = m^2 + n^2$
- The triple generated by Euclid's formula is primitive if and only if m and n are coprime and not both odd.
- To generate all Pythagorean triples uniquely: $a = k(m^2 n^2), b = k(2mn), c = k(m^2 + n^2)$

4 Geometry

4.1 2DLibrary

```
typedef long double lf;
const lf EPS = 1e-8L, INF = 5e9;
```

```
const lf E0 = 0.0L;//Keep = 0 for integer
    coordinates, otherwise = EPS
enum {OUT,IN,ON};
struct pt {
 lf x,y;
 pt(){}
 pt(lf a, lf b): x(a), y(b){}
  pt operator - (const pt &q ) const {
   return \{x - q.x, y - q.y\};
  pt operator + (const pt &q ) const {
   return \{x + q.x, y + q.y\};
  pt operator * (const lf &t ) const {
   return {x * t , y * t };
  pt operator / (const lf &t ) const {
   return {x / t , y / t };
  bool operator < ( const pt & q ) const {</pre>
   if (fabsl(x - q.x) > E0) return x < q.x;
   return y < q.y;</pre>
  }
  void normalize() {
   lf norm = hypotl( x, y );
   if( fabsl( norm ) > EPS )
     x /= norm, y /= norm;
 }
}:
pt rot90( pt p ) { return { -p.y, p.x }; }
pt rot( pt p, lf w ) {
 return { cosl(w) * p.x - sinl(w) * p.y,
      sinl(w) * p.x + cosl(w) * p.v };
lf norm2 ( pt p ) { return p.x * p.x + p.y * p.y;
lf dis2 ( pt p, pt q ) { return norm2(p-q); }
lf norm ( pt p ) { return hypotl ( p.x, p.y ); }
lf dis( pt p, pt q ) { return norm( p - q ); }
lf dot( pt p, pt q ) { return p.x * q.x + p.y *
    q.v; }
lf cross( pt p, pt q ) { return p.x * q.y - q.x *
    p.y ; }
```

```
lf orient( pt a, pt b, pt c ) { return cross( b -
    a, c - a); };
bool in_angle( pt a, pt b, pt c, pt p ) {
 //assert( fabsl( orient( a, b, c ) ) > E0 );
 if( orient( a, b, c ) < -E0 )</pre>
   return orient( a, b, p ) >= -E0 || orient( a,
        c, p) \leq E0;
 return orient(a, b, p) >= -E0 && orient(a,
      c, p) \le E0;
}
struct line {
 pt nv;
 lf c;
 line( pt _nv, lf _c ) : nv( _nv ), c( _c ) {}
 line (pt p, pt q) {
   nv = \{ p.y - q.y, q.x - p.x \};
   c = -dot(p, nv);
 lf eval( pt p ) { return dot( nv, p ) + c; }
 lf distance2( pt p ) {
   return eval( p ) / norm2( nv ) * eval( p );
 lf distance( pt p ) {
   return fabsl( eval( p ) ) / norm( nv );
 pt projection( pt p ) {
   return p - nv * ( eval( p ) / norm2( nv ) );
 }
}:
pt lines_intersection( line a, line b ) {
 lf d = cross( a.nv, b.nv );
 //assert( fabsl( d ) > E0 );
 lf dx = a.nv.v * b.c - a.c * b.nv.v;
 lf dy = a.c * b.nv.x - a.nv.x * b.c;
 return { dx / d, dy / d };
}
line bisector( pt a, pt b ) {
 pt nv = (b - a), p = (a + b) * 0.5L;
 lf c = -dot(nv, p);
 return line( nv, c );
struct Circle {
 pt center;
 lf r;
```

```
Circle( pt p, lf rad ) : center( p ), r( rad )
 Circle( pt p, pt q ) {
   center = (p + q) * 0.5L;
   r = dis(p, q) * 0.5L;
 Circle( pt a, pt b, pt c ) {
   line lb = bisector(a, b), lc = bisector(a,
       c ):
   center = lines_intersection( lb, lc );
   r = dis(a, center);
 int contains( pt &p ) {
   lf det = r * r - dis2(center, p);
   if( fabsl( det ) <= E0 ) return ON;</pre>
   return ( det > E0 ? IN : OUT );
 }
vector< pt > circle_line_intersection( Circle c,
    line 1) {
 lf h2 = c.r * c.r - 1.distance2(c.center);
 if( fabsl( h2 ) < EPS ) return { 1.projection(</pre>
      c.center ) };
 if( h2 < 0.0L ) return {};</pre>
 pt dir = rot90( 1.nv );
 pt p = 1.projection( c.center );
 lf t = sqrtl( h2 / norm2( dir ) );
 return { p + dir * t, p - dir * t };
vector< pt > circle_circle_intersection( Circle
    c1, Circle c2 ) {
 pt dir = c2.center - c1.center;
 lf d2 = dis2( c1.center, c2.center );
 if( d2 <= E0 ) { //assert( fabsl( c1.r - c2.r )</pre>
      > E0 ):
   return {};
 1f td = 0.5L * (d2 + c1.r * c1.r - c2.r * c2.r
      );
 1f h2 = c1.r * c1.r - td / d2 * td:
 pt p = c1.center + dir * (td / d2);
 if( fabsl( h2 ) < EPS ) return { p };</pre>
 if( h2 < 0.0L ) return {};</pre>
 pt dir_h = rot90(dir) * sqrtl( h2 / d2 );
```

```
return { p + dir_h, p - dir_h };
vector< pt > convex_hull( vector< pt > v ) {
 sort( v.begin(), v.end() );//remove repeated
      points if needed
 const int n = v.size();
 if(n < 3) return v;
 vector< pt > ch(2 * n);
 int k = 0:
 for( int i = 0; i < n; ++ i ) {</pre>
   while (k > 1 \&\& orient(ch[k-2], ch[k-1],
       v[i] ) <= E0 )
     --k:
   ch[k++] = v[i]:
 const int t = k:
 for( int i = n - 2; i >= 0; -- i ) {
   while (k > t \&\& orient(ch[k-2], ch[k-1],
       v[i] ) <= E0 )
     --k;
   ch[k++] = v[i];
 ch.resize( k - 1 );
 return ch:
vector<pt> minkowski( vector<pt> P, vector<pt> Q
    ) {
 rotate( P.begin(), min_element( P.begin(),
      P.end() ), P.end() ):
 rotate( Q.begin(), min_element( Q.begin(),
      Q.end() ), Q.end() );
 P.push_back(P[0]), P.push_back(P[1]);
 Q.push_back(Q[0]), Q.push_back(Q[1]);
 vector<pt> ans;
 size_t i = 0, j = 0;
 while(i < P.size() - 2 || j < Q.size() - 2) {</pre>
     ans.push_back(P[i] + Q[j]);
     lf dt = cross(P[i + 1] - P[i], Q[j + 1] -
          Q[i]);
     if(dt >= E0 && i < P.size() - 2) ++i;</pre>
     if(dt <= E0 && j < Q.size() - 2) ++j;</pre>
 return ans;
```

```
vector< pt > cut( const vector< pt > &pol, line 1
    ) {
 vector< pt > ans;
 for( int i = 0, n = pol.size(); i < n; ++ i ) {</pre>
   lf s1 = 1.eval( pol[i] ), s2 = 1.eval(
        pol[(i+1)%n]);
   if( s1 >= -EPS ) ans.push_back( pol[i] );
   if( ( s1 < -EPS && s2 > EPS ) || ( s1 > EPS
        && s2 < -EPS ) ) {
     line li = line( pol[i], pol[(i+1)%n]);
     ans.push_back( lines_intersection( 1, li ) );
 }
 return ans:
int point_in_polygon( const vector< pt > &pol,
    const pt &p ) {
 int wn = 0;
 for( int i = 0, n = pol.size(); i < n; ++ i ) {</pre>
   lf c = orient(p, pol[i], pol[(i+1)\%n]);
   if( fabsl( c ) <= E0 && dot( pol[i] - p,</pre>
        pol[(i+1)\%n] - p) \le E0) return ON;
   if( c > 0 && pol[i].y <= p.y + E0 &&</pre>
        pol[(i+1)\%n].y - p.y > E0) ++wn;
   if(c < 0 \&\& pol[(i+1)\%n].y \le p.y + E0 \&\&
        pol[i].y - p.y > E0 ) --wn;
 return wn ? IN : OUT;
int point_in_convex_polygon( const vector < pt >
    &pol, const pt &p ) {
 int low = 1, high = pol.size() - 1;
 while( high - low > 1 ) {
   int mid = ( low + high ) / 2;
   if( orient( pol[0], pol[mid], p ) >= -E0 )
       low = mid;
   else high = mid;
 if( orient( pol[0], pol[low], p ) < -E0 )</pre>
      return OUT;
 if( orient( pol[low], pol[high], p ) < -E0 )</pre>
      return OUT:
 if( orient( pol[high], pol[0], p ) < -E0 )</pre>
      return OUT;
```

```
if( low == 1 && orient( pol[0], pol[low], p )
        <= E0 ) return ON;
if( orient( pol[low], pol[high], p ) <= E0 )
        return ON;
if( high == (int) pol.size() -1 && orient(
        pol[high], pol[0], p ) <= E0 ) return ON;
return IN;
}</pre>
```

4.2 Bentley-Ottmann

```
struct seg {
   pt p, q;
   int id;
   lf get_y(double x) const {
       if (fabsl (p.x - q.x) < EPS) return p.y;</pre>
       return p.y + (q.y - p.y) * (x - p.x) /
            (q.x - p.x);
   }
};
bool operator<(const seg& a, const seg& b) {</pre>
   If x = max(min(a.p.x, a.q.x), min(b.p.x,
        b.q.x));
   return a.get_y(x) < b.get_y(x) - EPS;</pre>
} //Use where[id] = s.insert(nxt, a[id]) &
    s.erase(where[id]);
set<seg> s;
vector<set<seg>::iterator> where;
```

4.3 Closest Points

```
int mid = (1 + r) / 2;
  lf x_mid = v[mid].x;
  solve( 1, mid ):
  solve(mid + 1, r);
  vector< pt > aux;
  int p1 = 1, p2 = mid + 1;
  while ( p1 <= mid && p2 <= r ) {</pre>
   if( v[p1].y < v[p2].y ) aux.push_back(</pre>
        v[p1++]);
   else aux.push_back( v[p2++] );
  while( p1 <= mid ) aux.push_back( v[p1++] );</pre>
  while( p2 <= r ) aux.push_back( v[p2++] );</pre>
  vector < pt > nb;
  for( int i = 1; i <= r; ++ i ) {</pre>
   v[i] = aux[i-1];
   lf dx = (x_mid - v[i].x);
   if(dx * dx < d2)
     nb.push_back( v[i] );
 for( int i = 0; i < (int) nb.size(); ++ i ) {</pre>
   for( int k = i + 1; k < (int) nb.size(); ++</pre>
        k ) {
     lf dy = (nb[k].y - nb[i].y);
     if( dy * dy > d2 ) break;
     lf nd2 = dis2( nb[i], nb[k] );
     if( nd2 < d2 ) d2 = nd2, ans = {nb[i],}
          nb[k]}:
   }
 }
}:
solve( 0, v.size() -1 );
return ans;
```

4.4 Halfplane Intersection

```
struct Halfplane {
```

```
pt p, pq;
 lf angle;
 Halfplane() {}
 Halfplane(const pt& a, const pt& b) : p(a),
      pq(b - a) {
   angle = atan21(pq.y, pq.x);
  bool out(const pt& r) {
   return cross(pq, r - p) < -EPS;</pre>
 }
  bool operator < (const Halfplane& e) const {</pre>
   return angle < e.angle;</pre>
 friend pt inter(const Halfplane& s, const
      Halfplane& t) {
     lf alpha = cross((t.p - s.p), t.pq) /
          cross(s.pq, t.pq);
     return s.p + (s.pq * alpha);
 }
};
vector<pt> hp_intersect(vector<Halfplane>& H) {
 pt box[4] = { pt(INF, INF), pt(-INF, INF),
      pt(-INF, -INF), pt(INF, -INF) };
 for(int i = 0; i < 4; ++i) {
   Halfplane aux(box[i], box[(i+1) % 4]);
   H.push_back(aux);
  sort(H.begin(), H.end());
 deque<Halfplane> dq;
  int len = 0;
 for(int i = 0; i < int(H.size()); ++i ) {</pre>
   while (len > 1 && H[i].out(inter(dq[len-1],
        dq[len-2]))) {
     dq.pop_back();
     --len:
   while (len > 1 && H[i].out(inter(dq[0],
        dq[1]))) {
     dq.pop_front();
     --len:
   if (len > 0 && fabsl(cross(H[i].pq,
        dq[len-1].pq)) < EPS ) {
     if (dot(H[i].pq, dq[len-1].pq) < 0.0)</pre>
```

```
return vector<pt>();
   if (H[i].out(dq[len-1].p)) {
     dq.pop_back();
     --len:
    else continue:
  dq.push_back(H[i]);
 ++len:
}
while (len > 2 && dq[0].out(inter(dq[len-1],
    dq[len-2]))) {
 dq.pop_back();
  --len:
while (len > 2 && dq[len-1].out(inter(dq[0],
    dq[1]))) {
 dq.pop_front();
  --len:
if (len < 3) return vector<pt>();
vector<pt> ret(len);
for(int i = 0; i+1 < len; ++i )</pre>
 ret[i] = inter(dq[i], dq[i+1]);
ret.back() = inter(dq[len-1], dq[0]);
//remove repeated points if needed
return ret:
```

4.5 Min Circle

```
Circle min_circle( vector< pt > v ) {
  random_shuffle( v.begin(), v.end() );
  auto f2 = [&]( int a, int b ){
    Circle ans( v[a], v[b] );
    for( int i = 0; i < a; ++ i )
      if( ans.contains( v[i] ) == OUT )
      ans = Circle( v[i], v[a], v[b] );
  return ans;
};
auto f1 = [&]( int a ){</pre>
```

```
Circle ans( v[a], 0.0L );
  for( int i = 0; i < a; ++ i )
    if( ans.contains( v[i] ) == OUT )
      ans = f2( i, a );
  return ans;
};

Circle ans( v[0], 0.0L );
  for( int i = 1; i < (int) v.size(); ++ i )
    if( ans.contains( v[i] ) == OUT )
      ans = f1( i );

return ans;
}</pre>
```

4.6 Miscellaneous

```
lf part(pt a, pt b, T r) {
 lf l = abs(a-b);
 pt p = (b-a)/1;
 If c = dot(a, p), d = 4.0 * (c*c - dot(a, a) +
 if (d < eps) return angle(a, b) * r * r * 0.5;
 d = sqrt(d) * 0.5;
 lf s = -c - d, t = -c + d;
 if (s < 0.0) s = 0.0; else if (s > 1) s = 1;
 if (t < 0.0) t = 0.0; else if (t > 1) t = 1;
 pt u = a + p*s, v = a + p*t;
 return (cross(u, v) + (angle(a, u) + angle(v,
      b)) * r * r) * 0.5;
lf inter_cp(circle c, polygon p) {
 lf ans = 0;
 int n = p.p.size();
 for (int i = 0; i < n; i++) {</pre>
   ans += part(p[i]-c.c, p[(i+1)%4]-c.c, c.r);
 return abs(ans);
bool circumcircle_contains( triangle tr, pt D )
    {//triange CCW
```

```
pt A = tr.vert[0] - D, B = tr.vert[1] - D, C =
      tr.vert[2] - D;
 lf norm_a = norm2( tr.vert[0] ) - norm2( D );
 lf norm_b = norm2( tr.vert[1] ) - norm2( D );
 lf norm_c = norm2( tr.vert[2] ) - norm2( D );
 lf det1 = A.x * (B.y * norm_c - norm_b * C.y);
 lf det2 = B.x * (C.y * norm_a - norm_c * A.y);
 lf det3 = C.x * (A.y * norm_b - norm_a * B.y);
 return det1 + det2 + det3 > E0;
}
lf areaOfIntersectionOfTwoCircles( lf r1, lf r2,
    lf d ) {
 if(d >= r1 + r2)
   return 0.0L:
 if( d <= fabsl( r2 - r1 ) )</pre>
   return PI * ( r1 < r2 ? r1 * r1 : r2 * r2 ):
 lf alpha = safeAcos( ( r1 * r1 - r2 * r2 + d *
      d)/(2.0L * d * r1));
 lf betha = safeAcos( ( r2 * r2 - r1 * r1 + d *
      d) / (2.0L * d * r2 ));
 lf a1 = r1 * r1 * ( alpha - sinl( alpha ) *
      cosl( alpha ) );
 lf a2 = r2 * r2 * (betha - sinl(betha) *
      cosl( betha ) );
 return a1 + a2;
};
bool half(pt p) { /// true if is in (0, 180]
 assert(p.x != 0 || p.y != 0); /// the argument
      of (0.0) is undefined
 return p.y > 0 || (p.y == 0 && p.x < 0);
bool half_from(pt p, pt v = {1, 0}) {
 return cross(v,p) < 0 \mid \mid (cross(v,p) == 0 \&\&
      dot(v,p) < 0);
bool polar_cmp(const pt &a, const pt &b) {
 return make_tuple(half(a), 0) <</pre>
      make_tuple(half(b), cross(a,b));
bool in_disk(pt a, pt b, pt p) {
 return dot(a-p, b-p) <= 0;</pre>
bool on_segment(pt a, pt b, pt p) {
 return orient(a,b,p) == 0 && in_disk(a,b,p);
```

```
bool proper_inter(pt a, pt b, pt c, pt d, pt
    &out) {
 T oa = orient(c,d,a), ob = orient(c,d,b),
   oc = orient(a,b,c), od = orient(a,b,d);
 if (oa*ob < 0 && oc*od < 0) {
   out = (a*ob - b*oa) / (ob-oa);
   return true:
 return false:
set<pt> inter_ss(pt a, pt b, pt c, pt d) {
 pt out;
 if (proper_inter(a,b,c,d,out)) return {out};
 set<pt> s;
 if (on_segment(c,d,a)) s.insert(a);
 if (on_segment(c,d,b)) s.insert(b);
 if (on_segment(a,b,c)) s.insert(c);
 if (on_segment(a,b,d)) s.insert(d);
 return s;
lf pt_to_seg(pt a, pt b, pt p) {
 if(a != b) {
   line 1(a,b);
   if (1.cmp_proj(a,p) && 1.cmp_proj(p,b)) ///
        if closest to projection
     return l.dist(p); /// output distance to line
 return min(abs(p-a), abs(p-b)); /// otherwise
      distance to A or B
lf seg_to_seg(pt a, pt b, pt c, pt d) {
 pt dummy;
 if (proper_inter(a,b,c,d,dummy)) return 0;
 return min({pt_to_seg(a,b,c), pt_to_seg(a,b,d),
            pt_to_seg(c,d,a), pt_to_seg(c,d,b)});
enum {IN, OUT, ON};
struct polygon {
 vector<pt> p;
 polygon(int n) : p(n) {}
 pt centroid() {
   pt c{0, 0};
   lf scale = 6. * area(true);
```

```
for(int i = 0, n = p.size(); i < n; ++i) {</pre>
     int j = (i+1 == n ? 0 : i+1);
     c = c + (p[i] + p[j]) * cross(p[i], p[j]);
   return c / scale;
 }
 ll pick() {
   11 boundary = 0;
   for(int i = 0, n = p.size(); i < n; i++) {</pre>
     int j = (i+1 == n ? 0 : i+1);
     boundary += _{gcd}((ll)abs(p[i].x - p[j].x),
          (ll)abs(p[i].y - p[j].y));
   return area() + 1 - boundary/2;
 pt& operator[] (int i){ return p[i]; }
};
int tangents(circle c1, circle c2, bool inner,
    vector<pair<pt,pt>> &out) {
 if(inner) c2.r = -c2.r;
 pt d = c2.c-c1.c;
 double dr = c1.r-c2.r, d2 = norm(d), h2 =
      d2-dr*dr:
 if(d2 == 0 || h2 < 0) { assert(h2 != 0); return</pre>
      0; }
 for(double s : {-1,1}) {
   pt v = (d*dr + rot90ccw(d)*sqrt(h2)*s)/d2;
   out.push_back({c1.c + v*c1.r, c2.c + v*c2.r});
 return 1 + (h2 > 0);
int tangent_through_pt(pt p, circle c, pair<pt,</pre>
    pt> &out) {
 double d = abs(p - c.c);
       if(d < c.r) return 0;</pre>
 pt base = c.c-p;
 double w = sqrt(norm(base) - c.r*c.r);
 pt a = \{w, c.r\}, b = \{w, -c.r\};
 pt s = p + base*a/norm(base)*w;
 pt t = p + base*b/norm(base)*w;
 out = {s, t};
 return 1 + (abs(c.c-p) == c.r);
/// cross product 3D (VxW)
```

```
{ v.y*w.z - v.z*w.y, v.z*w.x - v.x*w.z, v.x*w.y - v.y*w.x };
```

5 Graphs

5.1 2-satisfiability

```
/// Complexity: O(|N|)
struct sat2 {
 int n:
 vector<vector<int>>> g;
 vector<int> tag;
 vector<bool> seen, value;
 stack<int> st;
 sat2(int n) : n(n), g(2,
      vector<vector<int>>(2*n)), tag(2*n),
      seen(2*n), value(2*n) { }
 int neg(int x) { return 2*n-x-1; }
 void add_or(int u, int v) { implication(neg(u),
      v); }
 void make_true(int u) { add_edge(neg(u), u); }
 void make_false(int u) { make_true(neg(u)); }
 void eq(int u, int v) {
   implication(u, v);
   implication(v, u);
 void diff(int u, int v) { eq(u, neg(v)); }
 void implication(int u, int v) {
   add_edge(u, v);
   add_edge(neg(v), neg(u));
 void add_edge(int u, int v) {
   g[0][u].push_back(v);
   g[1][v].push_back(u);
 void dfs(int id, int u, int t = 0) {
   seen[u] = true;
   for(auto& v : g[id][u])
     if(!seen[v])
       dfs(id, v, t);
   if(id == 0) st.push(u);
```

```
else tag[u] = t;
 }
 void kosaraju() {
   for(int u = 0; u < n; u++) {</pre>
     if(!seen[u]) dfs(0, u);
     if(!seen[neg(u)]) dfs(0, neg(u));
   fill(seen.begin(), seen.end(), false);
   int t = 0;
   while(!st.empty()) {
     int u = st.top(); st.pop();
     if(!seen[u]) dfs(1, u, t++);
   }
 }
 bool satisfiable() {
   kosaraju():
   for(int i = 0; i < n; i++) {</pre>
     if(tag[i] == tag[neg(i)]) return false;
     value[i] = tag[i] > tag[neg(i)];
   }
   return true;
 }
};
```

5.2 Erdos–Gallai theorem

```
/// Complexity: O(|N|*log|N|)
/// Theorem: it gives a necessary and sufficient
    condition for a finite sequence
111
            of natural numbers to be the degree
    sequence of a simple graph
bool erdos(vector<int> &d) {
 11 \text{ sum} = 0:
  for(int i = 0; i < d.size(); ++i) sum += d[i];</pre>
  if(sum & 1) return false;
  sort(d.rbegin(), d.rend());
 11 1 = 0, r = 0;
  for(int k = 1, i = d.size() - 1; k <= d.size();</pre>
      ++k) {
   1 += d[k-1]:
   if(k > i) r = d[++i];
   while (i >= k && d[i] < k+1) r += d[i--];
```

```
if(1 > 111*k*(k-1) + 111*k*(i-k+1) + r)
    return false;
}
return true;
}
```

5.3 Eulerian path

```
/// Complexity: O(|N|)
struct edge {
 int v; //list<edge>::iterator rev;
 edge(int v) : v(v) {}
void add_edge(int a, int b) {
 g[a].push_front(edge(b)); // auto ia =
      g[a].begin();
       g[b].push_front(edge(a)); // auto ib =
           g[b].begin();
       //ia->rev=ib; ib->rev=ia;
}
/// for undirected uncomment and check for path
    existance
bool eulerian(vector<int> &tour) { /// directed
  int one_in = 0, one_out = 0, start = -1;
 bool ok = true;
  for (int i = 0; i < n; i++) {</pre>
   if(out[i] && start == -1) start = i;
   if(out[i] - in[i] == 1) one_out++, start = i;
   else if(in[i] - out[i] == 1) one in++;
   else ok &= in[i] == out[i];
  ok &= one_in == one_out && one_in <= 1;
  if (ok) {
   function<void(int)> go = [&](int u) {
     while(g[u].size()) {
       int v = g[u].front().v;
       g[v].erase(g[u].front().rev);
       g[u].pop_front();
       go(v, tour);
     tour.push_back(u);
```

```
};
go(start);
reverse(tour.begin(), tour.end());
if(tour.size() == edges + 1) return true;
}
return false;
```

5.4 Number of spanning trees

```
///A -> adjacency matrix
///It is necessary to compute the D-A matrix,
    where D is a diagonal matrix
///that contains the degree of each node.
///To compute the number of spanning trees it's
    necessary to compute any
///D-A cofactor
///C(i, j) = (-1)^(i+j) * Mij
///Where Mij is the matrix determinant after
    removing row i and column j
double mat[MAX][MAX]:
///call determinant(n - 1)
double determinant(int n) {
 double det = 1.0;
 for(int k = 0; k < n; k++) {
   for(int i = k+1; i < n; i++) {</pre>
     assert(mat[k][k] != 0):
     long double factor = mat[i][k]/mat[k][k];
     for(int j = 0; j < n; j++) {</pre>
       mat[i][j] = mat[i][j] - factor*mat[k][j];
     }
    det *= mat[k][k];
 return round(det);
```

5.5 Scc

```
/// Complexity: O(|N|)
int scc(int n) {
 vector<int> dfn(n+1), low(n+1), in_stack(n+1);
 stack<int> st:
 int tag = 0;
 function<void(int, int&)> dfs = [&](int u, int
     &t) {
   dfn[u] = low[u] = ++t:
   st.push(u);
   in_stack[u] = true;
   for(auto &v : g[u]) {
     if(!dfn[v]) {
       dfs(v, t);
       low[u] = min(low[u], low[v]);
     } else if(in_stack[v])
       low[u] = min(low[u], dfn[v]);
   if (low[u] == dfn[u]) {
     int v:
     do {
       v = st.top(); st.pop();
        id[v] = tag;
      in_stack[v] = false;
     } while (v != u);
     tag++;
   }
 for(int u = 1, t; u <= n; ++u) {</pre>
   if(!dfn[u]) dfs(u, t = 0):
 return tag;
```

5.6 Tarjan tree

```
/// Complexity: O(|N|)
struct tarjan_tree {
  int n;
  vector<vector<int>> g, comps;
  vector<pii>> bridge;
  vector<int> id, art;
```

```
tarjan_tree(int n) : n(n), g(n+1), id(n+1),
    art(n+1) {}
void add_edge(vector<vector<int>> &g, int u,
    int v) { /// nodes from [1, n]
 g[u].push_back(v);
 g[v].push_back(u);
void add_edge(int u, int v) { add_edge(g, u,
    v): }
void tarjan(bool with_bridge) {
 vector<int> dfn(n+1), low(n+1);
 stack<int> st;
 comps.clear();
 function<void(int, int, int&)> dfs = [&](int
      u, int p, int &t) {
   dfn[u] = low[u] = ++t:
   st.push(u);
   int cntp = 0;
   for(int v : g[u]) {
     cntp += v == p;
     if(!dfn[v]) {
       dfs(v, u, t);
       low[u] = min(low[u], low[v]);
       if(with_bridge && low[v] > dfn[u]) {
        bridge.push_back({min(u,v), max(u,v)});
         comps.push_back({});
        for(int w = -1; w != v; )
           comps.back().push_back(w =
               st.top()), st.pop();
       if(!with_bridge && low[v] >= dfn[u]) {
         art[u] = (dfn[u] > 1 || dfn[v] > 2);
         comps.push_back({u});
         for(int w = -1; w != v; )
           comps.back().push_back(w =
               st.top()), st.pop();
      }
     else if(v != p \mid \mid cntp > 1) low[u] =
         min(low[u], dfn[v]);
   if(p == -1 && ( with_bridge || g[u].size()
       == 0 )) {
     comps.push_back({});
```

```
for(int w = -1: w != u:)
         comps.back().push_back(w = st.top()),
             st.pop();
     }
   };
   for(int u = 1, t; u \le n; ++u)
     if(!dfn[u]) dfs(u, -1, t = 0);
  vector<vector<int>> build_block_cut_tree() {
   tarjan(false);
   int t = 0;
   for(int u = 1; u \le n; ++u)
     if(art[u]) id[u] = t++;
   vector<vector<int>> tree(t+comps.size());
   for(int i = 0; i < comps.size(); ++i)</pre>
     for(int u : comps[i]) {
       if(!art[u]) id[u] = i+t;
       else add_edge(tree, i+t, id[u]);
     }
   return tree;
  vector<vector<int>> build_bridge_tree() {
   tarjan(true);
   vector<vector<int>> tree(comps.size());
   for(int i = 0; i < comps.size(); ++i)</pre>
     for(int u : comps[i]) id[u] = i;
   for(auto &b : bridge)
     add_edge(tree, id[b.first], id[b.second]);
   return tree:
 }
};
```

6 Math

6.1 Berlekamp-Massey

```
namespace linear_seq {
  int m; // a = first m terms
  vector < int > p, a;// p = recurrence, length
    is m
```

```
inline vector < int > BM(vector < int > x) { //
    finds shortest linear recurrence given
    first x terms in O(x^2)
 vector < int > ls, cur;
 int lf, ld;
 for (int i = 0; i < (int) x.size(); ++i) {</pre>
   int t = 0;
   for (int j = 0; j < (int) cur.size(); ++j)</pre>
     t = (t + 1 LL * x[i - j - 1] * cur[j]) %
          mod:
   if ((t - x[i]) \% \text{ mod} == 0) continue;
   if (!cur.size()) {
     cur.resize(i + 1);
     lf = i:
     1d = (t - x[i]) \% mod;
     continue:
   int k = 1 LL * (t - x[i]) * pw(ld, mod - 2)
        % mod:
   vector < int > c(i - lf - 1);
    c.push_back(k);
   for (int j = 0; j < (int) ls.size(); ++j)</pre>
     c.push_back(-1 LL * ls[j] * k % mod);
   if (c.size() < cur.size())</pre>
        c.resize(cur.size());
   for (int j = 0; j < (int) cur.size(); ++j)</pre>
     c[i] = (c[i] + cur[i]) \% mod;
   if (i + lf + (int) ls.size() >= (int)
        cur.size())
     ls = cur, lf = i, ld = (t - x[i]) \% mod;
    cur = c;
 }
 for (int i = 0; i < (int) cur.size(); ++i)</pre>
    cur[i] = (cur[i] % mod + mod) % mod;
 m = cur.size():
 p.resize(m), a.resize(m);
 for (int i = 0; i < m; ++i)</pre>
   p[i] = cur[i], a[i] = x[i];
 return cur;
inline vector < int > mul(vector < int > & a,
    vector < int > & b) {
 vector < int > r(2 * m);
 for (int i = 0; i < m; ++i)</pre>
```

```
if (a[i])
       for (int j = 0; j < m; ++j)
         r[i + j] = (r[i + j] + 1 LL * a[i] *
             b[i]) % mod:
   for (int i = 2 * m - 1; i >= m; --i)
     if (r[i])
       for (int j = m - 1; j \ge 0; ---j)
        r[i - j - 1] = (r[i - j - 1] + 1 LL *
             p[j] * r[i]) % mod;
   r.resize(m):
   return r;
 inline int calc(long long k) { // O(m*m*log(k))
   if (m == 0) return 0:
   vector < int > bs(m), r(m);
   if (m == 1) bs[0] = p[0];
   else bs[1] = 1:
   r[0] = 1;
   while (k) {
     if (k \& 1) r = mul(r, bs);
     bs = mul(bs, bs);
     k >>= 1;
   }
   int res = 0;
   for (int i = 0; i < m; ++i)</pre>
     res = (res + 1 LL * r[i] * a[i]) % mod;
   return res:
}
```

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6.2 Chinese remainder theorem

```
/// Complexity: |N|*log(|N|)
/// finds a suitable x that meets: x is congruent
    to a_i mod n_i
/** Works for non-coprime moduli.
Returns {-1,-1} if solution does not exist or
    input is invalid.
Otherwise, returns {x,L}, where x is the
    solution unique to mod L = LCM of mods
*/
```

```
pair<int, int> chinese_remainder_theorem(
    vector<int> A, vector<int> M ) {
 int n = A.size(), a1 = A[0], m1 = M[0];
 for(int i = 1: i < n: i++) {</pre>
   int a2 = A[i], m2 = M[i];
   int g = __gcd(m1, m2);
   if( a1 % g != a2 % g ) return {-1,-1};
   int p, q;
   eea(m1/g, m2/g, &p, &q);
   int mod = m1 / g * m2;
   q %= mod; p %= mod;
   int x = ((111*(a1\%mod)*(m2/g))\%mod*q +
        (111*(a2%mod)*(m1/g))%mod*p) % mod; // if
        WA there is overflow
   a1 = x:
   if (a1 < 0) a1 += mod:
   m1 = mod:
 return {a1, m1};
```

6.3 Constant modular inverse

6.4 Extended euclides

```
/// Complexity: O(log(|N|))
ll eea(ll a, ll b, ll& x, ll& y) {
    ll xx = y = 0; ll yy = x = 1;
    while (b) {
        ll q = a / b; ll t = b; b = a % b; a = t;
        t = xx; xx = x - q * xx; x = t;
        t = yy; yy = y - q * yy; y = t;
```

```
}
return a;
}
ll inverse(ll a, ll n) {
    ll x, y;
    ll g = eea(a, n, x, y);
    if(g > 1)
        return -1;
    return (x % n + n) % n;
}
```

6.5 FWHT

```
vector< long long > haddamard( vector< long long</pre>
    > a, bool inverse ) {
 const int n = (int) a.size();
 for( int k = 1; k < n; k <<= 1 ) {</pre>
   for( int i = 0; i < n; i += 2 * k ) {</pre>
     for( int j = 0; j < k; ++ j ) {</pre>
       long long u = a[i+j], v = a[i+j+k];
       a[i+j] = u + v;
       a[i+j+k] = u - v;
     }
   }
   if( inverse )
     for( auto &x : a )
       x >>= 1;
 }
 return a:
vector< long long > FWHT( vector< long long > a,
    vector< long long > b ) {
 auto h_a = haddamard( a, false ), h_b =
     haddamard( b, false );
 vector< long long > h_c( a.size() );
 for( int i = 0; i < (int) a.size(); ++ i )</pre>
   h_c[i] = 1LL * h_a[i] * h_b[i];
 return haddamard( h_c, true );
```

6.6 Fast Fourier transform module

```
struct FFT {
 int mod, root, root_1, root_pw;
 void fft(vector<int> & a, bool invert) {
     int n = a.size();
     for (int i = 1, j = 0; i < n; ++i) {
         int bit = n \gg 1:
         for (; j & bit; bit >>= 1)
            j ^= bit;
         j ^= bit;
         if (i < j) swap(a[i], a[j]);</pre>
     for (int len = 2; len <= n; len <<= 1) {
         int wlen = invert ? root_1 : root;
         for (int i = len; i < root_pw; i <<= 1)</pre>
             wlen = 1LL * wlen * wlen % mod;
         for (int i = 0: i < n: i += len) {
            int w = 1;
            for (int j = 0; j < len / 2; ++j) {
                int u = a[i+j], v = 1LL *
                    a[i+j+len/2] * w % mod;
                a[i+j] = u + v < mod ? u + v : u
                    + v - mod;
                a[i+j+len/2] = u - v >= 0 ? u - v
                     : u - v + mod;
                w = 1LL * w * wlen % mod;
            }
         }
     if (invert) {
         int n_1 = pw(n, mod-2, mod);
         for (int & x : a)
            x = 1LL * x * n_1 \% mod;
     }
 vector<int> multiply(vector<int> const& a,
      vector<int> const& b) {
     vector<int> fa(a.begin(), a.end()),
         fb(b.begin(), b.end());
     int n = 1:
     while (n < a.size() + b.size())</pre>
         n <<= 1;
```

```
fa.resize(n);
fb.resize(n);
fft(fa, false);
fft(fb, false);
for (int i = 0; i < n; i++)
     fa[i] = 1LL * fa[i] * fb[i] % mod;
fft(fa, true);
   return fa;
}
};
FFT A = { 998244353, 15311432, 469870224, 1<<23 };</pre>
```

6.7 Fast fourier transform

```
///Complexity: O(N log N)
#define rep(i, a, b) for(int i = a; i < (b); ++i)
#define sz(v) ((int)v.size())
#define trav(a, x) for(auto& a : x)
#define all(v) v.begin(),v.end()
typedef vector<ll> v1;
typedef vector<int> vi;
typedef complex<double> C;
typedef vector<double> vd;
void fft(vector<C>& a) {
       int n = sz(a), L = 31 - \_builtin\_clz(n);
       static vector<complex<long double>> R(2,
           1):
       static vector<C> rt(2, 1); // (^ 10%
           faster if double)
      for (static int k = 2: k < n: k *= 2) {
              R.resize(n); rt.resize(n);
              auto x = polar(1.0L, acos(-1.0L) /
                  k);
              rep(i,k,2*k) rt[i] = R[i] = i&1 ?
                  R[i/2] * x : R[i/2];
       }
       vi rev(n);
       rep(i,0,n) rev[i] = (rev[i / 2] | (i & 1)
           << L) / 2;
       rep(i,0,n) if (i < rev[i]) swap(a[i],
           a[rev[i]]):
       for (int k = 1; k < n; k *= 2)
```

```
for (int i = 0; i < n; i += 2 * k)
                  rep(j,0,k) {
                     // C z = rt[j+k] *
                          a[i+i+k]: // (25%
                          faster if hand-rolled)
                          /// include-line
                      auto x = (double)
                          *)&rt[j+k], y = (double)
                          *)&a[i+j+k]; ///
                          exclude-line
                     C z(x[0]*y[0] - x[1]*y[1],
                          x[0]*v[1] + x[1]*v[0]);
                              /// exclude-line
                      a[i + j + k] = a[i + j] - z;
                      a[i + j] += z;
              }
vl conv(const vl& a, const vl& b) {
       if (a.empty() || b.empty()) return {};
       vd res(sz(a) + sz(b) - 1);
       int L = 32 - __builtin_clz(sz(res)), n = 1
           << L:
       vector<C> in(n), out(n);
       copy(all(a), begin(in));
       rep(i,0,sz(b)) in[i].imag(b[i]);
       fft(in);
       trav(x, in) x *= x:
       rep(i,0,n) out[i] = in[-i & (n - 1)] -
           conj(in[i]);
       fft(out);
       vector<ll> resp(sz(res));
       rep(i,0,sz(res)) resp[i] =
           round(imag(out[i]) / (4.0 * n));
       return resp;
}
```

6.8 Gauss jordan

```
/// Complexity: O(|N|^3)
const int EPS = 1;
int gauss (vector<vector<int>> a, vector<int>
    &ans) {
```

```
int n = a.size(), m = a[0].size()-1:
vector<int> where(m, -1);
for(int col = 0, row = 0; col < m && row < n;</pre>
    ++col) {
  int sel = row;
 for(int i = row: i < n: ++i)</pre>
   if(abs(a[i][col]) > abs(a[sel][col])) sel =
  if(abs(a[sel][col]) < EPS) continue;</pre>
  swap(a[sel], a[row]);
  where [col] = row;
  for(int i = 0; i < n; ++i)</pre>
   if(i != row) {
     int c = divide(a[i][col], a[row][col]);
          /// precalc inverses
     for(int j = col; j <= m; ++j)</pre>
       a[i][j] = sub(a[i][j], mul(a[row][j],
            c));
   }
  ++row;
ans.assign(m, 0);
for(int i = 0; i < m; ++i)</pre>
 if(where[i] != -1) ans[i] =
      divide(a[where[i]][m], a[where[i]][i]);
for(int i = 0; i < n; ++i) {</pre>
  int sum = 0:
 for(int j = 0; j < m; ++j)
   sum = add(sum, mul(ans[j], a[i][j]));
 if(sum != a[i][m]) return 0;
for(int i = 0; i < m; ++i)</pre>
 if(where[i] == -1) return -1; /// infinite
      solutions
return 1:
```

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6.9 Lagrange Interpolation

```
/// Complexity: O(|N|^2)
vector<lf> X, F;
lf f(lf x) {
```

```
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```

```
lf answer = 0:
 for(int i = 0; i < (int)X.size(); i++) {</pre>
   lf prod = F[i];
   for(int j = 0; j < (int)X.size(); j++) {</pre>
     if(i == j) continue;
     prod = mul(prod, divide(sub(x, X[j]),
          sub(X[i], X[i])));
   answer = add(answer, prod);
 return answer;
//given y=f(x) for x [0,degree]
vector< int > interpolation( vector< int > &y ) {
 int n = (int) y.size();
 vector \leq int > u = y, ans (n), sum (n);
 ans[0] = u[0], sum[0] = 1;
 for( int i = 1; i < n; ++i ) {</pre>
   int inv = modpow( i, mod - 2 );
   for( int j = n - 1; j >= i; --j)
     u[j] = 1LL * (u[j] - u[j - 1] + mod) * inv %
   for( int j = i; j > 0; --j ) {
     sum[j] = (sum[j - 1] - 1LL * (i - 1) *
         sum[j] % mod + mod) % mod;
     ans[i] = (ans[i] + 1LL * sum[i] * u[i]) %
         mod:
   sum[0] = 1LL * (i - 1) * (mod - sum[0]) % mod:
   ans[0] = (ans[0] + 1LL * sum[0] * u[i]) % mod;
 }
 return ans;
```

6.10 Linear diophantine

```
g = eea(abs(a), abs(b), x, y);
 if(c % g) return false;
 a /= g; b /= g; c /= g;
 if(a < 0) x *= -1:
 x = (x \% b) * (c \% b) \% b;
 if(x < 0) x += b:
 y = (c - a*x) / b;
 return true:
///finds the first k \mid x + b * k / gcd(a, b) >=
ll greater_or_equal_than(ll a, ll b, ll x, ll
    val, 11 g) {
 lf got = 1.0 * (val - x) * g / b;
 return b > 0 ? ceil(got) : floor(got);
}
void get_xy (ll a, ll b, ll &x, ll &y, ll k, ll
    g) { /// if for y, change the order to b,a y,x
 x = x + b / g * k;
 y = y - a / g * k;
```

6.11 Matrix multiplication

```
const int MOD = 1e9+7;
struct matrix {
 const int N = 2:
 int m[N][N], r, c;
 matrix(int r = N, int c = N, bool iden = false)
      : r(r), c(c) {
   memset(m, 0, sizeof m);
   if (iden)
     for(int i = 0; i < r; i++) m[i][i] = 1;
 matrix operator * (const matrix &o) const {
   matrix ret(r, o.c);
   for(int i = 0; i < r; ++i)</pre>
     for(int j = 0; j < o.c; ++j) {</pre>
       ll &r = ret.m[i][j];
      for(int k = 0; k < c; ++k)
        r = (r + 111*m[i][k]*o.m[k][j]) % MOD;
```

```
return ret;
}
```

6.12 Miller rabin

```
11 mul (11 a. 11 b. 11 mod) {
 11 \text{ ret} = 0;
 for(a %= mod, b %= mod; b != 0:
   b >>= 1, a <<= 1, a = a >= mod ? a - mod : a)
   if (b & 1) {
     ret += a;
     if (ret >= mod) ret -= mod;
 return ret:
11 fpow (11 a, 11 b, 11 mod) {
 ll ans = 1:
 for (; b; b >>= 1, a = mul(a, a, mod))
   if (b & 1)
     ans = mul(ans, a, mod);
 return ans;
bool witness (ll a, ll s, ll d, ll n) {
 11 x = fpow(a, d, n);
 if (x == 1 \mid | x == n - 1) return false;
 for (int i = 0; i < s - 1; i++) {
   x = mul(x, x, n):
   if (x == 1) return true;
   if (x == n - 1) return false;
 return true;
11 \text{ test}[] = \{2, 3, 5, 7, 11, 13, 17, 19, 23, 0\};
bool is_prime (ll n) {
 if (n < 2) return false;
 if (n == 2) return true;
 if (n % 2 == 0) return false;
 11 d = n - 1, s = 0;
  while (d \% 2 == 0) ++s, d /= 2;
```

```
for (int i = 0; test[i] && test[i] < n; ++i)
  if (witness(test[i], s, d, n))
    return false;
return true;
}</pre>
```

6.13 Pollard's rho

```
11 pollard_rho(ll n, ll c) {
 11 x = 2, y = 2, i = 1, k = 2, d;
 while (true) {
   x = (mul(x, x, n) + c);
   if (x \ge n) x -= n;
   d = \_gcd(x - y, n);
   if (d > 1) return d;
   if (++i == k) y = x, k <<= 1;
 }
 return n;
void factorize(ll n, vector<ll> &f) {
 if (n == 1) return:
 if (is_prime(n)) {
   f.push_back(n);
   return:
 }
 11 d = n;
 for (int i = 2; d == n; i++)
   d = pollard_rho(n, i);
 factorize(d, f):
 factorize(n/d, f);
```

6.14 Simplex

```
/// Complexity: O(|N|^2 * |M|) N variables, N
    restrictions
const double EPS = 1e-6;
typedef vector<double> vec;
namespace simplex {
  vector<int> X, Y;
```

```
vector<vec> a:
vec b, c;
double z;
int n, m;
void pivot(int x, int y) {
  swap(X[y], Y[x]);
  b[x] /= a[x][y];
  for(int i = 0; i < m; i++)</pre>
   if(i != y)
     a[x][i] /= a[x][v];
  a[x][y] = 1 / a[x][y];
  for(int i = 0; i < n; i++)</pre>
   if(i != x && abs(a[i][v]) > EPS) {
   b[i] -= a[i][y] * b[x];
   for(int j = 0; j < m; j++)
     if(i != v)
       a[i][j] = a[i][y] * a[x][j];
   a[i][y] = -a[i][y] * a[x][y];
  }
  z += c[v] * b[x];
 for(int i = 0; i < m; i++)</pre>
   if(i != y)
     c[i] -= c[v] * a[x][i];
  c[y] = -c[y] * a[x][y];
/// A is a vector of 1 and 0. B is the limit
    restriction. C is the factors of O.F.
pair<double, vec> simplex(vector<vec> &A, vec
    &B, vec &C) {
  a = A; b = B; c = C;
  n = b.size(); m = c.size(); z = 0.0;
  X = vector<int>(m):
  Y = vector<int>(n);
 for(int i = 0; i < m; i++) X[i] = i;</pre>
  for(int i = 0; i < n; i++) Y[i] = i + m;
  while(1) {
   int x = -1, y = -1;
   double mn = -EPS;
   for(int i = 0; i < n; i++)</pre>
     if(b[i] < mn)
       mn = b[i], x = i;
   if(x < 0) break:
   for(int i = 0; i < m; i++)</pre>
     if(a[x][i] < -EPS) { y = i; break; }</pre>
```

```
assert(v \ge 0): // no sol
     pivot(x, y);
   }
    while(1) {
     double mx = EPS;
     int x = -1, y = -1;
     for(int i = 0; i < m; i++)</pre>
      if(c[i] > mx)
         mx = c[i], y = i;
     if(v < 0) break;
     double mn = 1e200;
     for(int i = 0; i < n; i++)</pre>
       if(a[i][v] > EPS && b[i] / a[i][v] < mn)</pre>
         mn = b[i] / a[i][y], x = i;
     assert(x >= 0); // unbound
     pivot(x, v):
    vec r(m);
   for(int i = 0; i < n; i++)</pre>
     if(Y[i] < m)
       r[Y[i]] = b[i];
   return make_pair(z, r);
}
```

6.15 Simpson

```
inline lf simpson(lf fl, lf fr, lf fmid, lf l, lf
    r) {
    return (fl + fr + 4.0 * fmid) * (r - 1) / 6.0;
}
lf rsimpson (lf slr, lf fl, lf fr, lf fmid, lf l,
    lf r) {
        lf mid = (l + r) * 0.5;
        lf fml = f((l + mid) * 0.5);
        lf fmr = f((mid + r) * 0.5);
        lf slm = simpson(fl, fmid, fml, l, mid);
        lf smr = simpson(fmid, fr, fmr, mid, r);
        if (fabs(slr - slm - smr) < eps) return
            slm + smr;
    return rsimpson(slm, fl, fmid, fml, l,
            mid) + rsimpson(smr, fmid, fr, fmr,</pre>
```

```
mid, r);
}

If integrate(lf l,lf r) {
    lf mid = (l + r) * .5, fl = f(l), fr =
        f(r), fmid = f(mid);
    return rsimpson(simpson(fl, fr, fmid, l,
        r), fl, fr, fmid, l, r);
}
```

7 Network flows

7.1 Blossom

```
/// Complexity: O(|E||V|^2)
struct network {
 struct struct_edge { int v; struct_edge * n; };
 typedef struct_edge* edge;
 struct_edge pool[MAXE]; ///2*n*n;
 edge top;
 vector<edge> adj;
 queue<int> q;
 vector<int> f, base, inq, inb, inp, match;
 vector<vector<int>> ed;
 network(int n) : n(n), match(n, -1), adj(n),
      top(pool), f(n), base(n),
                 inq(n), inb(n), inp(n), ed(n)
                     vector<int>(n)) {}
 void add_edge(int u, int v) {
   if(ed[u][v]) return;
   ed[u][v] = 1;
   top->v = v, top->n = adj[u], adj[u] = top++;
   top->v = u, top->n = adj[v], adj[v] = top++;
 int get_lca(int root, int u, int v) {
   fill(inp.begin(), inp.end(), 0);
   while(1) {
     inp[u = base[u]] = 1;
     if(u == root) break;
     u = f[match[u]];
```

```
while(1) {
   if(inp[v = base[v]]) return v;
   else v = f[ match[v] ];
 }
void mark(int lca, int u) {
 while(base[u] != lca) {
   int v = match[u]:
   inb[ base[u ]] = 1;
   inb[ base[v] ] = 1;
   u = f[v]:
   if(base[u] != lca) f[u] = v;
 }
}
void blossom_contraction(int s, int u, int v) {
 int lca = get_lca(s, u, v);
 fill(inb.begin(), inb.end(), 0);
 mark(lca, u); mark(lca, v);
 if(base[u] != lca) f[u] = v;
 if(base[v] != lca) f[v] = u;
 for(int u = 0; u < n; u++)
   if(inb[base[u]]) {
     base[u] = lca;
     if(!inq[u]) {
         inq[u] = 1;
         q.push(u);
     }
   }
}
int bfs(int s) {
 fill(inq.begin(), inq.end(), 0);
 fill(f.begin(), f.end(), -1);
 for(int i = 0; i < n; i++) base[i] = i;</pre>
 q = queue<int>();
 q.push(s);
 inq[s] = 1;
 while(q.size()) {
   int u = q.front(); q.pop();
   for(edge e = adi[u]; e; e = e->n) {
     int v = e \rightarrow v:
     if(base[u] != base[v] && match[u] != v) {
       if((v == s) \mid | (match[v] != -1 &&
           f[match[v]] != -1)
         blossom_contraction(s, u, v);
```

```
else if(f[v] == -1) {
          f[v] = u:
          if(match[v] == -1) return v;
          else if(!ing[match[v]]) {
            inq[match[v]] = 1;
            q.push(match[v]);
        }
   return -1;
 int doit(int u) {
   if(u == -1) return 0;
   int v = f[u]:
   doit(match[v]):
   match[v] = u; match[u] = v;
   return u != -1:
 /// (i < net.match[i]) => means match
 int maximum matching() {
   int ans = 0;
   for(int u = 0; u < n; u++)
     ans += (match[u] == -1) && doit(bfs(u));
   return ans;
 }
};
```

7.2 Dinic

```
/// Complexity: O(|E|*|V|^2)
struct edge { int v, cap, inv, flow; };
struct network {
  int n, s, t;
  vector<int> lvl;
  vector<vector<edge>> g;
  network(int n) : n(n), lvl(n), g(n) {}
  void add_edge(int u, int v, int c) {
    g[u].push_back({v, c, g[v].size(), 0});
    g[v].push_back({u, 0, g[u].size()-1, c});
}
```

```
bool bfs() {
   fill(lvl.begin(), lvl.end(), -1);
   queue<int> q;
   lvl[s] = 0:
   for(q.push(s); q.size(); q.pop()) {
     int u = q.front();
     for(auto &e : g[u]) {
      if(e.cap > 0 \&\& lvl[e.v] == -1) {
        lvl[e.v] = lvl[u]+1;
        q.push(e.v):
      }
    }
   return lvl[t] != -1:
 int dfs(int u, int nf) {
   if(u == t) return nf;
   int res = 0;
   for(auto &e : g[u]) {
     if(e.cap > 0 && lvl[e.v] == lvl[u]+1) {
       int tf = dfs(e.v, min(nf, e.cap));
      res += tf; nf -= tf; e.cap -= tf;
       g[e.v][e.inv].cap += tf;
      g[e.v][e.inv].flow -= tf;
       e.flow += tf;
       if(nf == 0) return res;
    }
   if(!res) lvl[u] = -1:
   return res;
 int max_flow(int so, int si, int res = 0) {
   s = so; t = si;
   while(bfs()) res += dfs(s, INT_MAX);
   return res:
 }
};
```

7.3 Hopcroft karp

```
/// Complexity: O(|E|*sqrt(|V|))
struct mbm {
```

```
vector<vector<int>> g;
vector<int> d, match;
int nil, l, r;
/// u \rightarrow 0 \text{ to } 1. v \rightarrow 0 \text{ to } r
mbm(int 1, int r) : 1(1), r(r), nil(1+r),
    g(1+r).
                   d(1+l+r, INF), match(l+r,
                       1+r) {}
void add_edge(int a, int b) {
  g[a].push_back(1+b);
  g[l+b].push_back(a);
bool bfs() {
  queue<int> q:
  for(int u = 0; u < 1; u++) {</pre>
   if(match[u] == nil) {
     d[u] = 0:
     q.push(u);
   } else d[u] = INF;
  }
  d[nil] = INF;
  while(q.size()) {
   int u = q.front(); q.pop();
   if(u == nil) continue;
   for(auto v : g[u]) {
     if(d[ match[v] ] == INF) {
       d[match[v]] = d[u]+1;
       q.push(match[v]);
   }
  return d[nil] != INF;
bool dfs(int u) {
  if(u == nil) return true;
  for(int v : g[u]) {
    if(d[ match[v] ] == d[u]+1 && dfs(match[v]))
     match[v] = u; match[u] = v;
     return true;
  }
  d[u] = INF;
  return false;
```

```
}
int max_matching() {
  int ans = 0;
  while(bfs()) {
    for(int u = 0; u < 1; u++) {
        ans += (match[u] == nil && dfs(u));
     }
  }
  return ans;
}</pre>
```

7.4 Maximum bipartite matching

```
/// Complexity: O(|E|*|V|)
struct mbm {
 int 1. r:
 vector<vector<int>> g;
 vector<int> match, seen;
 mbm(int l, int r) : l(l), r(r), seen(r),
      match(r), g(1) {}
 void add_edge(int 1, int r) {
      g[l].push_back(r); }
 bool dfs(int u) {
   for(auto v : g[u]) {
     if(seen[v]++) continue;
     if(match[v] == -1 || dfs(match[v])) {
       match[v] = u;
       return true;
     }
   return false;
 int max_matching() {
   int ans = 0;
   fill(match.begin(), match.end(), -1);
   for(int u = 0; u < 1; ++u) {
     fill(seen.begin(), seen.end(), 0);
     ans += dfs(u);
   return ans;
```

};

7.5 Maximum flow minimum cost

```
/// Complexity: O(|V|*|E|^2*log(|E|))
template <class type>
struct mcmf {
 struct edge { int u, v, cap, flow; type cost; };
 int n:
 vector<edge> ed;
 vector<vector<int>> g;
 vector<int> p;
 vector<type> d, phi;
 mcmf(int n) : n(n), g(n), p(n), d(n), phi(n) {}
 void add_edge(int u, int v, int cap, type cost)
   g[u].push_back(ed.size());
   ed.push_back({u, v, cap, 0, cost});
   g[v].push_back(ed.size());
   ed.push_back({v, u, 0, 0, -cost});
 bool dijkstra(int s, int t) {
   fill(d.begin(), d.end(), INF_TYPE);
   fill(p.begin(), p.end(), -1);
   set<pair<type, int>> q;
   d[s] = 0;
   for(q.insert({d[s], s}); q.size();) {
     int u = (*q.begin()).second;
         q.erase(q.begin());
     for(auto v : g[u]) {
       auto &e = ed[v];
       type nd = d[e.u]+e.cost+phi[e.u]-phi[e.v];
       if(0 < (e.cap-e.flow) && nd < d[e.v]) {
         q.erase({d[e.v], e.v});
         d[e.v] = nd; p[e.v] = v;
         q.insert({d[e.v], e.v});
       }
     }
   for(int i = 0; i < n; i++) phi[i] =</pre>
        min(INF_TYPE, phi[i]+d[i]);
   return d[t] != INF_TYPE;
```

```
pair<int, type> max_flow(int s, int t) {
   type mc = 0;
   int mf = 0:
   fill(phi.begin(), phi.end(), 0);
   while(dijkstra(s, t)) {
     int flow = INF;
     for(int v = p[t]; v != -1; v = p[ ed[v].u ])
       flow = min(flow, ed[v].cap-ed[v].flow);
     for(int v = p[t]; v != -1; v = p[ed[v].u])
       edge &e1 = ed[v];
       edge &e2 = ed[v^1];
       mc += e1.cost*flow:
       e1.flow += flow;
       e2.flow -= flow:
     }
     mf += flow;
   }
   return {mf, mc};
 }
};
```

7.6 Stoer Wagner

```
/// Complexity: O(|V|^3)
/// Tested: https://tinyurl.com/y8eu433d
struct stoer_wagner {
 int n;
 vector<vector<int>> g;
 stoer_wagner(int n) : n(n), g(n,
      vector<int>(n)) {}
 void add_edge(int a, int b, int w) { g[a][b] =
     g[b][a] = w; }
 pair<int, vector<int>> min_cut() {
   vector<int> used(n);
   vector<int> cut, best_cut;
   int best_weight = -1;
   for(int p = n-1; p >= 0; --p) {
     vector < int > w = g[0];
     vector<int> added = used;
     int prv, lst = 0;
```

```
for(int i = 0; i < p; ++i) {</pre>
      prv = lst; lst = -1;
      for(int j = 1; j < n; ++j)
        if(!added[j] && (lst == -1 || w[j] >
            w[lst]))
          lst = i:
      if(i == p-1) {
        for(int j = 0; j < n; j++)</pre>
          g[prv][i] += g[lst][i];
        for(int j = 0; j < n; j++)
          g[j][prv] = g[prv][j];
        used[lst] = true;
        cut.push_back(lst);
        if(best_weight == -1 || w[lst] <</pre>
            best_weight) {
          best_cut = cut;
          best_weight = w[lst];
      } else {
        for(int j = 0; j < n; j++)
          w[i] += g[lst][i];
        added[lst] = true;
    }
  return {best_weight, best_cut}; /// best_cut
       contains all nodes in the same set
}
}:
```

7.7 Weighted matching

```
/// Complexity: O(|V|^3)
typedef int type;
struct matching_weighted {
  int 1, r;
  vector<vector<type>> c;
  matching_weighted(int 1, int r) : 1(1), r(r),
        c(1, vector<type>(r)) {
    assert(1 <= r);
}</pre>
```

```
void add_edge(int a, int b, type cost) {
    c[a][b] = cost; }
type matching() {
  vector<type> v(r), d(r); // v: potential
  vector\langle int \rangle ml(1, -1), mr(r, -1); // matching
      pairs
  vector<int> idx(r), prev(r);
  iota(idx.begin(), idx.end(), 0);
  auto residue = [&](int i, int j) { return
      c[i][i]-v[i]; };
  for(int f = 0; f < 1; ++f) {</pre>
   for(int j = 0; j < r; ++j) {
     d[j] = residue(f, j);
     prev[j] = f;
   type w;
    int j, 1;
   for (int s = 0, t = 0;;) {
     if(s == t) {
       1 = s;
       w = d[idx[t++]];
       for(int k = t: k < r: ++k) {
         j = idx[k];
         type h = d[j];
         if (h <= w) {</pre>
           if (h < w) t = s, w = h;
           idx[k] = idx[t]:
           idx[t++] = j;
         }
       for (int k = s; k < t; ++k) {
         j = idx[k];
         if (mr[j] < 0) goto aug;</pre>
       }
     }
     int q = idx[s++], i = mr[q];
     for (int k = t; k < r; ++k) {
       j = idx[k];
       type h = residue(i, j) - residue(i, q) +
       if (h < d[i]) {</pre>
         d[j] = h;
         prev[j] = i;
         if(h == w) {
```

```
if(mr[j] < 0) goto aug;</pre>
             idx[k] = idx[t];
             idx[t++] = j;
           }
        }
       }
     aug: for (int k = 0; k < 1; ++k)
       v[idx[k]] += d[idx[k]] - w;
     int i:
     do {
       mr[j] = i = prev[j];
       swap(j, ml[i]);
     } while (i != f):
   type opt = 0;
   for (int i = 0; i < 1; ++i)</pre>
     opt += c[i][ml[i]]; // (i, ml[i]) is a
          solution
   return opt;
 }
};
```

8 Strings

8.1 Aho corasick

```
/// Complexity: O(|text|+SUM(|pattern_i|)+matches)
const static int alpha = 26;
int trie[N][alpha], fail[N], nodes;
void add(string &s, int i) {
  int cur = 0;
  for(char c : s) {
    int x = c-'a';
    if(!trie[cur][x]) trie[cur][x] = ++nodes;
    cur = trie[cur][x];
  }
  //cnt_word[cur]++;
  //end_word[cur] = i; // for i > 0
}
void build() {
```

```
queue<int> q; q.push(0);
 while(q.size()) {
   int u = q.front(); q.pop();
   for(int i = 0; i < alpha; ++i) {</pre>
     int v = trie[u][i];
     if(!v) trie[u][i] = trie[ fail[u] ][i]; //
         construir automata
     else q.push(v);
     if(!u || !v) continue;
     fail[v] = trie[ fail[u] ][i];
     //fail_out[v] = end_word[ fail[v] ] ?
         fail[v] : fail_out[ fail[v] ];
     //cnt_word[v] += cnt_word[ fail[v] ]; //
         obtener informacion del fail_padre
   }
 }
}
```

8.2 Hashing

```
/// 1000234999, 1000567999, 1000111997, 1000777121
const int MODS[] = { 1001864327, 1001265673 };
const mint BASE(256, 256), ZERO(0, 0), ONE(1, 1);
inline int add(int a, int b, const int& mod) {
    return a+b >= mod ? a+b-mod : a+b; }
inline int sbt(int a, int b, const int& mod) {
    return a-b < 0 ? a-b+mod : a-b; }</pre>
inline int mul(int a, int b, const int& mod) {
    return 1ll*a*b%mod: }
inline 11 operator ! (const mint a) { return
    (ll(a.first)<<32)|ll(a.second); }
inline mint operator + (const mint a, const mint
    b) {
 return {add(a.first, b.first, MODS[0]),
      add(a.second, b.second, MODS[1]));
inline mint operator - (const mint a, const mint
 return {sbt(a.first, b.first, MODS[0]),
      sbt(a.second, b.second, MODS[1]));
}
```

```
inline mint operator * (const mint a, const mint
    b) {
 return {mul(a.first, b.first, MODS[0]),
      mul(a.second, b.second, MODS[1])};
mint base[MAXN]:
void prepare() {
 base[0] = ONE:
 for(int i = 1; i < MAXN; i++) base[i] =</pre>
      base[i-1]*BASE:
template <class type>
struct hashing {
 vector<mint> code:
 hashing(type &t) {
   code.resize(t.size()+1);
   code[0] = ZER0:
   for (int i = 1; i < code.size(); ++i)</pre>
     code[i] = code[i-1]*BASE + mint{t[i-1],}
         t[i-1]};
 }
 mint query(int 1, int r) {
   return code[r+1] - code[1]*base[r-1+1];
 }
};
```

8.3 Kmp automaton

}

8.4 Kmp

```
/// Complexity: O(|N|)
vector<int> get_phi(string &p) {
 vector<int> phi(p.size());
 phi[0] = 0;
 for(int i = 1, j = 0; i < p.size(); ++i ) {</pre>
   while(j > 0 && p[i] != p[j] ) j = phi[j-1];
   if(p[i] == p[j]) ++j;
   phi[i] = j;
 return phi;
int get_matches(string &t, string &p) {
 vector<int> phi = get_phi(p);
 int matches = 0;
 for(int i = 0, j = 0; i < t.size(); ++i ) {</pre>
   while(j > 0 && t[i] != p[j] ) j = phi[j-1];
   if(t[i] == p[j]) ++j;
   if(j == p.size()) {
     matches++:
     j = phi[j-1];
 return matches;
```

8.5 Manacher

```
/// Complexity: O(|N|)
///to = i - from[i];
///len = to - from[i] + 1 = i - 2 * from[i] + 1;
vector<int> manacher(string &s) {
  int n = s.size(), p = 0, pr = -1;
  vector<int> from(2*n-1);
  for(int i = 0; i < 2*n-1; ++i) {
   int r = i <= 2*pr ? min(p - from[2*p - i],
        pr) : i/2;</pre>
```

8.6 Minimun expression

```
/// Complexity: O(|N|)
int minimum_expression(string s) {
    s = s+s;
    int len = s.size(), i = 0, j = 1, k = 0;
    while(i+k < len && j+k < len) {
        if(s[i+k] == s[j+k]) k++;
        else if(s[i+k] > s[j+k]) i = i+k+1, k = 0;
        else j = j+k+1, k = 0;
        if(i == j) j++;
    }
    return min(i, j);
}
```

8.7 Palindromic Tree

```
/// Complexity: O(|N|*log(|alphabet|))
struct palindromic_tree {
   struct node{
     int len, link;
     map <char,int> next;
   };
   vector <node> pt;
   int last, it;
   string s;
   palindromic_tree (string str = "") {
      pt.reserve ( s.size()+2 );
   }
}
```

```
s = '#', it = 1, last = 1: //# isn't used in
        string
   add_node (), add_node();
   pt[1].len = -1, pt[0].link = 1;
   for ( char &c: str ) add_letter ( c );
 }
 int add_node () {
   pt.push_back({});
   return pt.size()-1;
 }
 int get_link ( int v ) {
   while (s[it-pt[v].len-2] != s[it-1]) v =
       pt[v].link;
   return v:
 }
 void add letter ( char c ) {
   s += c, ++it;
   last = get_link ( last );
   if ( !pt[last].next.count (c) ) {
     int curr = add_node ();
     pt[curr].len = pt[last].len+2;
     pt[curr].link =
         pt[get_link(pt[last].link)].next[c];
     pt[last].next[c] = curr;
   last = pt[last].next[c];
 node& operator[](int i) { return pt[i]; }
 int size() { return pt.size(); }
};
```

8.8 Suffix array

```
/// Complexity: O(|N|*log(|N|))
const int alpha = 400;
struct suffix_array { // s MUST not have 0 value
  vector<int> sa, pos, lcp;
  suffix_array(string s) {
    s.push_back('$'); // always add something
       less to input, so it stays in pos 0
    int n = s.size(), mx = max(alpha, n)+2;
```

```
vector\langle int \rangle a(n), a1(n), c(n+1), c1(n+1),
        head(mx), cnt(mx);
   pos = 1cp = a;
   for(int i = 0; i < n; i++) c[i] = s[i], a[i]
        = i, cnt[c[i]]++;
   for(int i = 0; i < mx-1; i++) head[i+1] =</pre>
        head[i] + cnt[i];
   for(int k = 0; k < n; k = max(1, k << 1)) {
     for(int i = 0; i < n; i++) {</pre>
       int j = (a[i] - k + n) \% n;
       a1[head[c[i]]++]=i;
     swap(a1, a);
     for(int i = 0, x = a[0], y, col = 0; i < n;
          i++, x = a[i], y = a[i-1]) {
       c1[x] = (i \&\& c[x] == c[y] \&\& c[x+k] ==
            c[y+k]) ? col : ++col;
       if(!i || c1[x] != c1[v]) head[col] = i;
      swap(c1, c);
     if(c[ a[n-1] ] == n) break;
   }
    swap(sa, a);
   for(int i = 0; i < n; i++) pos[ sa[i] ] = i;</pre>
   for(int i = 0, k = 0, j; i < n; lcp[ pos[i++]</pre>
        ] = k) { /// lcp[i, i+1]}
     if(pos[i] == n-1) continue;
     for(k = max(0, k-1), j = sa[pos[i]+1];
          s[i+k] == s[i+k]: k++):
   }
  int& operator[] ( int i ){ return sa[i]; }
};
```

8.9 Suffix automaton

```
/// Complexity: O(|N|*log(|alphabet|))
struct suffix_automaton {
   struct node {
    int len, link, cnt, first_pos; bool end; ///
        cnt is endpos size, first_pos is minimum
        of endpos
```

```
map<char. int> next:
}:
vector<node> sa:
int last:
suffix_automaton() {}
suffix_automaton(string &s) {
 sa.reserve(s.size()*2);
 last = add node():
 sa[last].len = sa[last].cnt =
      sa[last].first_pos = sa[last].end = 0;
 sa[last].link = -1;
 for(char c : s) sa_append(c);
 mark_suffixes();
 build cnt():
int add node() {
 sa.push_back({});
 return sa.size()-1;
void mark_suffixes() {
 ///t0 is not suffix
 for(int cur = last; cur; cur = sa[cur].link)
   sa[cur].end = 1;
/// This is O(N*log(N)). Can be done O(N) by
    doing dfs and counting paths to terminal
    nodes.
void build_cnt() {
 vector<int> order(sa.size()-1):
 iota(order.begin(), order.end(), 1);
 sort(order.begin(), order.end(), [&](int a,
      int b) { return sa[a].len > sa[b].len; });
 for(auto &i : order) sa[ sa[i].link ].cnt +=
      sa[i].cnt;
 sa[0].cnt = 0; /// t0 is empty string
void sa_append(char c) {
 int cur = add_node();
 sa[cur].len = sa[last].len + 1;
 sa[cur].end = 0; sa[cur].cnt = 1;
 sa[cur].first_pos = sa[cur].len-1;
 int p = last;
 while(p != -1 && !sa[p].next[c] ){
   sa[p].next[c] = cur;
```

```
p = sa[p].link;
   if(p == -1) sa[cur].link = 0;
   else {
     int q = sa[p].next[c];
     if(sa[q].len == sa[p].len+1) sa[cur].link =
     else {
       int clone = add_node();
       sa[clone] = sa[q];
       sa[clone].len = sa[p].len+1;
       sa[clone].cnt = 0;
       sa[q].link = sa[cur].link = clone;
       while(p != -1 && sa[p].next[c] == q) {
         sa[p].next[c] = clone;
        p = sa[p].link;
     }
   }
   last = cur;
 node& operator[](int i) { return sa[i]; }
};
```

8.10 Z algorithm

```
/// Complexity: 0(|N|)
vector<int> z_algorithm (string &s) {
  int n = s.size();
  vector<int> z(n);
  int l = 0, r = 0;
  for(int i = 1; i < n; ++i) {
    z[i] = max(0, min(z[i-1], r-i+1));
    while (i+z[i] < n && s[z[i]] == s[i+z[i]])
    l = i, r = i+z[i], ++z[i];</pre>
```

```
}
return z;
```

9 Utilities

9.1 Makefile

```
compile:
    g++ -std=c++17 -static -DLOCAL -02 -Wall
         -Wshadow -Wno-unused-result -o sol
         sol.cpp
```

9.2 Pragma optimizations

```
#pragma GCC optimize ("03")
#pragma GCC target ("sse4")
#pragma GCC target ("avx,tune=native")
```

9.3 Random

```
uniform_int_distribution<T> /
    uniform_real_distribution<T> dis(fr, to);
dis(rng)
```

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9.4 gen

```
int main(int argv, char * argc[]) {
  srand( atoi( argc[1] ) );
}
```

9.5 test

```
g++ -std=c++17 gen.cpp -o gen
for((i = 0; ;++i));do
  echo $i
   ./gen $i > int
done
```

9.6 vimrc