Monster Mash A UNIX-like shell and filesystem

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1 About

Monster Mash is a UNIX-like shell and virtual filesystem. Its name was originally MASH, which stood for Mohammad and Alok's SHell. This name was taken however, so we added *monster*, referencing the hit single by Bobby Pickett. It supports various simple commands similar to basic UNIX commands. The underlying filesystem is based on UNIX's inode and data block structure.

2 Installation

The repository is available on Github (https://github.com/ahota/monstermash). To build Monster Mash, simply run make. This will create two executables, mmash and client.

3 Usage

The mmash executable will start the shell and filesystem in local mode. In this mode, the user can interact with the filesystem using the commands below.

Alternatively, Monster Mash can be started in server mode by running: mmash -s [port]. This will start the server and wait for a client to connect. The port argument is optional. If it is provided, the Monster Mash server will communicate with the client on that port. If it is not provided, the server will use the default port, 1962, which is the year *Monster Mash* was released as a single.

The Monster Mash client is used by running: client <hostname> <port>. Both the server hostname and port are required. Once the client has connected, usage is identical to running Monster Mash in local mode.

The following commands are supported:

- mkfs: Formats and (re)generates the filesystem. This is run automatically when starting Monster Mash locally or when a client connects. This creates a root directory (/) and sets it as the current directory
- open <filename> <r|w|rw>: Opens a file for reading (r), writing (w) or both (rw) depending on the flag given. If a file is opened with the w or rw flag and does not exist, it will be created. After a file is opened, a file descriptor is printed for later access
- read <fd> <size>: Reads size bytes from an open file with file descriptor fd if the file was opened for reading
- write <fd> <string>: Writes string to an open file with file descriptor fd if the file was opened for writing
- seek <fd> <offset>: Seeks to offset (bytes) in file with file descriptor fd
- close <fd>: Closes file with file descriptor fd if it was open
- mkdir <name>: Creates a directory name in the current directory. name can have spaces if it is surrounded by quotes
- rmdir <name>: Removes the directory name from the current directory. The directory must be empty
- cd <path>: Changes the current directory to path. path must be relative to the current directory. It can have spaces if surrounded by quotes
- link <target> linkname>: Creates a link linkname to file or link target. Paths to link and target can be relative

- unlink <name>: Unlinks a link name or deletes a file name. If the target would have zero links to it after this element is unlinked/deleted, the data is erased from the disk
- stat [name]: Provides the following information to the element name (or the current directory if name is not provided):
 - name on disk
 - inode ID
 - element type
 - link count
 - number of blocks allocated
 - size on disk
- 1s: Lists all elements in the current directory
- cat <name>: Prints the contents of the file name
- cp <src> <dest>: Copies the file or link src to a new file or link dest. src and dest can be relative paths with spaces if surrounded by quotes
- tree: Lists the contents of all directories hierarchically from the current directory along with their total file size
- import src> <dest>: Imports a file src from the host to a file dest on the Monster Mash filesystem
- export src> <dest>: Exports a file src from the Monster Mash filesystem to a file dest on the host
- exit: In local mode, this shuts down the shell and returns the user to the host. In server mode, this severs the connection between server and client and shuts both processes down. The Monster Mash disk is still available on the host

4 Architecture

The Monster Mash filesystem is based on concepts from a typical UNIX filesystem.

4.1 inodes

The filesystem uses inodes to keep a record of elements. Each element, whether it is a file, directory, or link, has one inode. Each inode is 8 bytes in size and has the following structure:

Type	Name	Size	Purpose
char	type	1	Type of element: 'f', 'd', or 'l'
short	id	2	Unique ID for this inode
int	first_block	4	Offset to the first block allocated for this element
char	pad	1	Padding to reach 8 bytes
	Total	8	bytes

The inode table supports up to 1024 inodes. Since each inode is 8 bytes, the inode table occupies 8 kB of the disk.

Type	Name	Size	Purpose
char *	name	32	The name of this element or 'CONTINUED'
int	$next_block$	32	Byte offset to next block or -1
short	link_count	2	Number of links to this element
char *	data	4058	Contents of this element
	Total	4096	bytes

4.2 Data Blocks

The first data block begins immediately after the inode table. Blocks are 4096 bytes each and have the following structure:

The first data block for an element (ie the block pointed to by first_block in this element's inode) contains the correct name in its name field. Names for elements are limited to 32 characters. When a block is created, its next_block pointer is set to -1 and link_count is set to 1. link_count is incremented and decremented as links to this element are created or unlinked, respectively. If a file is created and linked to, its link_count would be 2. If the file is then deleted, the data blocks associated with that file are still kept, as link_count is still 1 due to the link. If that link is later unlinked, the data blocks are erased.

Directories hold a table of their contents in the data section of their block(s). Each entry in the table contains two values: inode ID and name. inode IDs in the table are used for unique referencing in commands such as cd. Names in the table are used for printing and user-reference in commands such as 1s and tree. Files use the data section to hold their contents in plain text. Links use this section simply to hold a reference of the inode ID to which they link.

If the data in a directory or file grows beyond a single block, a new block is automatically allocated and linked to in the current block's next_block pointer. This new block's name field contains 'CONTINUED' to distinguish it as a continuation block. Block allocation and traversal is completely abstracted from the user during normal use. The number of blocks allocated for an element can be seen with the stat command.

5 Limitations

Currently mmash supports 1024 inodes as a maximum. This means there cannot be more than 1024 elements on the filesystem, including the root directory. The total inode table size is 8KB. The disk itself is 100MB total, less the inode table and block meta-data size, leaves 99.06% of the disk for data.

In the current version, the filesystem is not persistent across runs and is erased every time mmash is started.