The goal of this short note is to provide simple proofs for the "folklore facts" on the sample complexity of learning a discrete probability distribution over a known domain of size k to distance  $\varepsilon$ , with error probability  $\delta$ , can be done with  $O\left(\frac{k + \log(1/\delta)}{\varepsilon^2}\right)$ . Thanks to Gautam Kamath and John Wright for suggesting "someone should write this up as a note," and to Jiantao Jiao for discussions about the Hellinger case.

For a given distance measure d, we write  $\Phi(d, k, \varepsilon, \delta)$  for the sample complexity of learning discrete distributions over a known domain of size k, to accuracy  $\varepsilon > 0$ , with error probability  $\delta \in (0, 1]$ . As usual, asymptotics will be taken with regard to k going to infinity,  $\varepsilon$  going to 0, and  $\delta$  going to 0, in that order. Without loss of generality, we hereafter assume the domain is the set  $[k] \stackrel{\text{def}}{=} \{1, \dots, k\}$ .

## 1 Total variation distance

Recall that  $d_{\text{TV}}(p,q) = \sup_{S \subset [k]} (p(S) - q(S)) = \frac{1}{2} ||p - q||_1 \in [0,1]$  for any  $p, q \in \Delta([k])$ .

Theorem 1. 
$$\Phi(d_{\text{TV}}, k, \varepsilon, \delta) = O\left(\frac{k + \log(1/\delta)}{\varepsilon^2}\right)$$
.

First proof. Consider the empirical distribution  $\widehat{p}$  obtained by drawing n independent samples  $s_1, \ldots, s_n$  from the underlying distribution  $p \in \Delta([k])$ :

$$\widehat{p}(i) = \frac{1}{n} \sum_{j=1}^{n} \mathbb{1}_{\{s_j = i\}}, \qquad i \in [k]$$
(1)

• First, we bound the *expected* total variation distance between  $\widehat{p}$  and p, by using  $\ell_2$  distance as a proxy:

$$\mathbb{E}[\mathbf{d}_{\mathrm{TV}}(p,\widehat{p})] = \frac{1}{2}\mathbb{E}[\|p-\widehat{p}\|_1] = \frac{1}{2}\sum_{i=1}^k \mathbb{E}[|p(i)-\widehat{p}(i)|] \leq \frac{1}{2}\sum_{i=1}^k \sqrt{\mathbb{E}[(p(i)-\widehat{p}(i))^2]}$$

the last inequality by Jensen. But since, for every  $i \in [k]$ ,  $n\widehat{p}(i)$  follows a Bin(n,p(i)) distribution, we have  $\mathbb{E}\big[(p(i)-\widehat{p}(i))^2\big] = \frac{1}{n^2} \operatorname{Var}[n\widehat{p}(i)] = \frac{1}{n}p(i)(1-p(i))$ , from which

$$\mathbb{E}[\mathbf{d}_{\mathrm{TV}}(p,\widehat{p})] \le \frac{1}{2\sqrt{n}} \sum_{i=1}^{k} \sqrt{p(i)} \le \frac{1}{2} \sqrt{\frac{k}{n}}$$

the last inequality this time by Cauchy–Schwarz. Therefore, for  $n \geq \frac{k}{\varepsilon^2}$  we have  $\mathbb{E}[\mathrm{d}_{\mathrm{TV}}(p,\widehat{p})] \leq \frac{\varepsilon}{2}$ .

• Next, to convert this expected result to a high probability guarantee, we apply McDiarmid's inequality to the random variable  $f(s_1, \ldots, s_n) \stackrel{\text{def}}{=} d_{\text{TV}}(p, \widehat{p})$ , noting that changing any single sample cannot change its value by more than  $c \stackrel{\text{def}}{=} 1/n$ :

$$\Pr\left[|f(s_1,\ldots,s_n) - \mathbb{E}[f(s_1,\ldots,s_n)]| \ge \frac{\varepsilon}{2}\right] \le 2e^{-\frac{2\left(\frac{\varepsilon}{2}\right)^2}{mc^2}} = 2e^{-\frac{1}{2}n\varepsilon^2}$$

and therefore as long as  $n \geq \frac{2}{\varepsilon^2} \ln \frac{2}{\delta}$ , we have  $|f(s_1, \ldots, s_n) - \mathbb{E}[f(s_1, \ldots, s_n)]| \leq \frac{\varepsilon}{2}$  with probability at least  $1 - \delta$ .

Putting it all together, we obtain that  $d_{\text{TV}}(p, \widehat{p}) \leq \varepsilon$  with probability at least  $1 - \delta$ , as long as  $n \geq \max\left(\frac{k}{\varepsilon^2}, \frac{2}{\varepsilon^2} \ln \frac{2}{\delta}\right)$ .

Second proof – the "fun" one. Again, we will analyze the behavior of the empirical distribution  $\widehat{p}$  over n i.i.d. samples from the unknown p (cf. (1)) – because it is simple, efficiently computable, and it works. Recalling the definition of total variation distance, note that  $d_{\text{TV}}(p,\widehat{p}) > \varepsilon$  literally means there exists a subset  $S \subseteq [k]$  such that  $\widehat{p}(S) > p(S) + \varepsilon$ . There are  $2^k$  such subsets, so... let us do a union bound. 2 Fix any  $S \subseteq [k]$ . We have

$$\widehat{p}(S) = \widehat{p}(i) \stackrel{\text{(1)}}{=} \frac{1}{n} \sum_{i \in S} \sum_{j=1}^{n} \mathbb{1}_{\{s_j = i\}}$$

and so, letting  $X_j \stackrel{\text{def}}{=} \sum_{i \in S} \mathbb{1}_{\{s_j = i\}}$  for  $j \in [m]$ , we have  $\widehat{p}(S) = \frac{1}{n} \sum_{j=1}^n X_j$  where the  $X_j$ 's are i.i.d. Bernoulli random variable with parameter p(S). Here comes the Chernoff bound (actually, Hoeffding, the *other* Chernoff):

$$\Pr[\widehat{p}(S) > p(S) + \varepsilon] = \Pr\left[\frac{1}{n} \sum_{j=1}^{n} X_j > \mathbb{E}\left[\frac{1}{n} \sum_{j=1}^{n} X_j\right] + \varepsilon\right] \le e^{-2\varepsilon^2 n}$$

and therefore  $\Pr[\widehat{p}(S) > p(S) + \varepsilon] \leq \frac{\delta}{2^k}$  for any  $n \geq \frac{k \ln 2 + \log(1/\delta)}{2\varepsilon^2}$ . A union bound over these  $2^k$  possible sets S concludes the proof:

$$\Pr[\exists S \subseteq [k] \text{ s.t. } \widehat{p}(S) > p(S) + \varepsilon] \le 2^k \cdot \frac{\delta}{2^k} = \delta$$

and we are done. Badda bing badda boom, as someone would say.

## 2 Hellinger distance

Recall that  $d_H(p,q) = \frac{1}{\sqrt{2}} \sqrt{\sum_{i=1}^k (\sqrt{p(i)} - \sqrt{q(i)})^2} = \frac{1}{\sqrt{2}} \|\sqrt{p} - \sqrt{q}\|_2 \in [0,1]$  for any  $p,q \in \Delta([k])$ . The Hellinger distance has many nice properties: it is well-suited to manipulating product distributions, its square is subadditive, and is always within a quadratic factor of the total variation distance; see, e.g., [Can15, Appendix C.2].

**Theorem 2.** 
$$\Phi(d_H, k, \varepsilon, \delta) = \Theta\left(\frac{k + \log(1/\delta)}{\varepsilon^2}\right)$$

This theorem is "highly non-trivial" to establish, however; for the sake of exposition, we will show increasingly stronger bounds, starting with the easiest to establish.

**Proposition 3** (Easy bound). 
$$\Phi(d_H, k, \varepsilon, \delta) = O\left(\frac{k + \log(1/\delta)}{\varepsilon^4}\right)$$
, and  $\Phi(d_H, k, \varepsilon, \delta) = \Omega\left(\frac{k + \log(1/\delta)}{\varepsilon^2}\right)$ .

*Proof.* This is immediate from Theorem 1, recalling that  $\frac{1}{2} d_{TV}^2 \leq d_H^2 \leq d_{TV}$ .

**Proposition 4** (More involved bound). 
$$\Phi(d_H, k, \varepsilon, 1/3) = O(\frac{k}{\varepsilon^2})$$
, and  $\Phi(d_H, k, \varepsilon, \delta) = O(\frac{k}{\varepsilon^2} + \frac{\log(1/\delta)}{\varepsilon^4})$ .

*Proof.* As for total variation distance, we consider the empirical distribution  $\widehat{p}$  (cf. (1)) obtained by drawing n independent samples  $s_1, \ldots, s_n$  from  $p \in \Delta([k])$ .

• First, we bound the expected squared Hellinger distance between  $\widehat{p}$  and p: using the simple fact that  $d_H(p,q)^2 = 1 - \sum_{i=1}^k \sqrt{p(i)q(i)}$  for any  $p,q \in \Delta([k])$ ,

$$\mathbb{E}\left[\mathrm{d}_{\mathrm{H}}(p,\widehat{p})^{2}\right] = 1 - \sum_{i=1}^{k} \sqrt{p(i)} \cdot \mathbb{E}\left[\sqrt{\widehat{p}(i)}\right].$$

 $<sup>^{1}</sup>$ John Wright.

Now we would like to handle the square root inside the expectation, and of course Jensen's inequality is in the wrong direction. However, for every nonnegative r.v. X with positive expectation, letting  $Y \stackrel{\text{def}}{=} X/\mathbb{E}[X]$ , we have that

$$\begin{split} \mathbb{E}\Big[\sqrt{X}\Big] &= \sqrt{\mathbb{E}[X]} \cdot \mathbb{E}\Big[\sqrt{Y}\Big] = \sqrt{\mathbb{E}[X]} \cdot \mathbb{E}\Big[\sqrt{1 + (Y - \mathbb{E}[Y])}\Big] \\ &\geq \sqrt{\mathbb{E}[X]}\Big(1 + \frac{1}{2}\mathbb{E}[Y - \mathbb{E}[Y]] - \frac{1}{6}\mathbb{E}\big[(Y - \mathbb{E}[Y])^2\big]\,\Big) = \sqrt{\mathbb{E}[X]}\Big(1 - \frac{\operatorname{Var}X}{6\mathbb{E}[X]^2}\Big) \end{split}$$

where we use the inequality  $\sqrt{1+x} \ge 1 + \frac{x}{2} - \frac{x^2}{6}$ , which holds for  $x \ge 0$ . Since, for every  $i \in [k]$ ,  $n\widehat{p}(i)$  follows a Bin(n, p(i)) distribution, we get

$$\mathbb{E}\left[d_{H}(p,\widehat{p})^{2}\right] \geq 1 - \frac{1}{\sqrt{n}} \sum_{i=1}^{k} \sqrt{p(i)} \cdot \sqrt{np(i)} \left(1 - \frac{np(i)(1-p(i))}{6n^{2}p(i)^{2}}\right) \geq 1 - \sum_{i=1}^{k} p(i) \left(1 - \frac{1}{6np(i)}\right) = \frac{k}{6n}.$$

Therefore, for  $n \geq \frac{2k}{3\varepsilon^2}$  we have  $\mathbb{E}\Big[\mathrm{d}_{\mathrm{H}}(p,\widehat{p})^2\Big] \leq \frac{\varepsilon^2}{4}$ .

• Next, to convert this expected result to a high probability guarantee, we would like to apply McDiarmid's inequality to the random variable  $f(s_1, \ldots, s_n) \stackrel{\text{def}}{=} d_H(p, \widehat{p})^2$  as in the (first) proof of Theorem 1; unfortunately, changing a sample can change the value by up to  $c \approx 1/\sqrt{n}$ , and McDiarmid will yield only a vacuous bound.<sup>3</sup> Instead, we will use a stronger, more involved concentration inequality:

**Theorem 5** ([BLM13, Theorem 8.6]). Let  $f: \mathcal{X}^n \to \mathbb{R}$  be a measurable function and let  $X_1, \ldots, X_n$  be independent random variables taking values in  $\mathcal{X}$ . Define  $Z \stackrel{\text{def}}{=} f(X_1, \ldots, X_n)$ . Assume that there exist measurable functions  $c_i: \mathcal{X}^n \to [0, \infty)$  such that, for all  $x, y \in \mathcal{X}^n$ ,

$$f(y) - f(x) \le \sum_{i=1}^{n} c_i(x) \mathbb{1}_{\{x_i \ne y_i\}}$$

Then, setting  $v \stackrel{\text{def}}{=} \mathbb{E} \sum_{i=1}^n c_i(x)^2$  and  $v_{\infty} \stackrel{\text{def}}{=} \sup_{x \in \mathcal{X}^n} \sum_{i=1}^n c_i(x)^2$ , we have, for all t > 0,

$$\Pr[Z \ge \mathbb{E}[Z] + t] \le e^{-\frac{t^2}{2v}}, \qquad \Pr[Z \le \mathbb{E}[Z] - t] \le e^{-\frac{t^2}{2v_{\infty}}}.$$

For our f above, we have, for two any different  $x, y \in [k]^n$ , that

$$f(y) - f(x) = \frac{1}{\sqrt{n}} \sum_{i=1}^{k} \sqrt{p_i} \left( \sqrt{\sum_{j=1}^{n} \mathbb{1}_{\{x_j = i\}}} - \sqrt{\sum_{j=1}^{n} \mathbb{1}_{\{y_j = i\}}} \right)$$

$$= \frac{1}{\sqrt{n}} \sum_{i=1}^{k} \sqrt{p_i} \frac{\sum_{j=1}^{n} (\mathbb{1}_{\{x_j = i\}} - \mathbb{1}_{\{y_j = i\}})}{\sqrt{\sum_{j=1}^{n} \mathbb{1}_{\{x_j = i\}}}}$$

$$\leq \frac{1}{\sqrt{n}} \sum_{i=1}^{k} \sqrt{p_i} \frac{\sum_{j=1}^{n} \mathbb{1}_{\{x_j = i\}} \mathbb{1}_{\{x_j = i\}}}{\sqrt{\sum_{j=1}^{n} \mathbb{1}_{\{x_j = i\}}}} = \sum_{j=1}^{n} \sqrt{\frac{p_{x_j}}{n \sum_{\ell=1}^{n} \mathbb{1}_{\{x_\ell = x_j\}}}} \cdot \mathbb{1}_{\{x_j \neq y_j\}}$$

$$= c_i(x)$$

In view of Theorem 5, we then must evaluate

$$v \stackrel{\text{def}}{=} \sum_{j=1}^{n} \mathbb{E}[c_{j}(X)^{2}] = \frac{1}{n} \sum_{j=1}^{n} \sum_{i=1}^{k} p_{i}^{2} \cdot \mathbb{E}\left[\frac{1}{1 + \sum_{\ell \neq j}^{n} \mathbb{1}_{\{x_{\ell} = i\}}}\right]$$

<sup>&</sup>lt;sup>2</sup>And is inspired by the Taylor expansion  $\sqrt{1+x} \ge 1 + \frac{x}{2} - \frac{x^2}{8} + o(x^2)$ .

<sup>&</sup>lt;sup>3</sup>Try it: it's a real bummer.

where that last expectation is over  $(x_{\ell})_{\ell \neq j}$ . Since  $\sum_{\ell \neq j}^{n} \mathbb{1}_{\{x_{\ell} = i\}}$  is Binomially distributed with parameters n-1 and  $p_i$ , we can use the simple fact that

$$\mathbb{E}\left[\frac{1}{1+N}\right] = \frac{1 - (1-\rho)^{r+1}}{\rho(r+1)} \le \frac{1}{\rho(r+1)}$$

for  $N \sim \text{Bin}(r, \rho)$ , to conclude that  $v \leq \frac{1}{n^2} \sum_{j=1}^n \sum_{i=1}^k p_i = \frac{1}{n}$ . By Theorem 5, we then get

$$\Pr\left[|f(s_1,\ldots,s_n) - \mathbb{E}[f(s_1,\ldots,s_n)]| \ge \frac{\varepsilon^2}{2}\right] \le e^{-\frac{1}{8}n\varepsilon^4}$$

and therefore as long as  $n \geq \frac{8}{\varepsilon^4} \ln \frac{1}{\delta}$ , we have  $|f(s_1, \ldots, s_n) - \mathbb{E}[f(s_1, \ldots, s_n)]| \leq \frac{\varepsilon^2}{2}$  with probability at least  $1 - \delta$ .

Putting it all together, we obtain that  $d_H(p, \widehat{p})^2 \leq \varepsilon^2$  with probability at least  $1 - \delta$ , as long as  $n \geq \max\left(\frac{k}{\varepsilon^2}, \frac{8}{\varepsilon^4} \ln \frac{1}{\delta}\right)$ .

We finally get to the final, optimal bound:

Proof of Theorem 2. We will rely on a recent – and quite involved – result due to Agrawal [Agr19], analyzing the concentration of the empirical distribution  $\hat{p}$  in terms of its Kullback–Leibler (KL) divergence with regard to the true p,

$$\mathrm{KL}(\widehat{p} \| p) = \sum_{i=1}^{k} \widehat{p}_{i} \ln \frac{\widehat{p}_{i}}{p_{i}} \in [0, \infty].$$

Observing that  $d_H(p,q)^2 \leq \frac{1}{2}KL(p \parallel q)$  for any distributions p,q, the aforementioned result is actually stronger than we we need:

**Theorem 6** ([Agr19, Theorem 1.2]). Suppose  $n \ge (1 + \ln 2) \frac{k-1}{\alpha}$ . Then

$$\Pr[\operatorname{KL}(\widehat{p} \| p) > \alpha] \le e^{-n\alpha} \cdot \left(2e \cdot \left(\frac{\alpha n}{k-1} - \ln 2\right)\right)^{k-1} \le e^{-n\alpha} \cdot \left(\frac{4e \cdot \alpha n}{k}\right)^{k}.$$

In view of the above relation between Hellinger and KL, we will apply this convergence result with  $\alpha \stackrel{\text{def}}{=} 2\varepsilon^2$ , obtaining

$$\Pr[d_{\mathbf{H}}(\widehat{p}, p) > \varepsilon] \le e^{-2n\varepsilon^2 + k \ln\left(\frac{8e \cdot \varepsilon^2 n}{k}\right)}.$$

Fact 7. For  $n \ge \frac{5k}{\varepsilon^2}$ , we have  $n\varepsilon^2 \ge k \ln \frac{8e \cdot \varepsilon^2 n}{k}$ .

*Proof.* The inequality is equivalent to  $\frac{8e \cdot \varepsilon^2 n}{k} \ge 8e \cdot \ln \frac{8e \cdot \varepsilon^2 n}{k}$ , and the conclusion follows from the fact that  $x \ge 8e \ln x$  for  $x \ge 101$ .

Therefore, for  $n \ge \max\left(\frac{5k}{\varepsilon^2}, \frac{1}{\varepsilon^2} \ln \frac{1}{\delta}\right)$ , we get  $\Pr[d_H(\widehat{p}, p) > \varepsilon] \le e^{-n\varepsilon^2} \le \delta$  as desired.

## $3 - \chi^2$ and Kullback–Leibler divergences

To conclude, some remarks on Kullback–Leibler (KL) and chi-squared ( $\chi^2$ ) divergences. Recall their definition, for  $p, q \in \Delta(k)$ ,

$$KL(p \| q) = \sum_{i=1}^{k} p(i) \ln \frac{p(i)}{q(i)}, \qquad \chi^{2}(p,q) = \sum_{i=1}^{k} \frac{(p(i) - q(i))^{2}}{q(i)}$$

as well as the chain of (easily checked) inequalities

$$2d_{\text{TV}}(p,q)^2 \le \text{KL}(p \parallel q) \le \chi^2(p,q)$$

where the first one is Pinsker's. Importantly, KL and  $\chi^2$  divergences are unbounded and asymmetric, so the order of p and q matters a lot: for instance, it is easy to show that, without strong assumptions on the unknown distribution  $p \in \Delta(k)$ , the empirical estimator  $\hat{p}$  cannot achieve  $\mathrm{KL}(p \parallel \hat{p}) < \infty$  (resp.  $\chi^2(p, \hat{p}) < \infty$ ) with any finite number of samples.<sup>4</sup> So, that's uplifting. (On the other hand, other estimators than the empirical one, e.g., add-constant estimators, do provide good learning guarantees for those distance measures: see for instancr [KOPS15].)

We are going to focus here on getting  $\mathrm{KL}(\widehat{p} \parallel p)$  and  $\chi^2(\widehat{p}, p)$  down to  $\varepsilon$ . Of course, in view of the inequalities above, the latter is at least as hard as the former, and a lower bound on both follows from that on  $\mathrm{d}_{\mathrm{TV}}$ :  $\Omega((k + \log(1/\delta))/\varepsilon)$ . And, behold! The result of Agrawal [Agr19] mentioned in the proof of Theorem 2 does provide the optimal upper bound on learning in KL divergence – and it is achieved by the usual suspect, the empirical estimator:

**Theorem 8.**  $\Phi(\mathrm{KL}, k, \varepsilon, \delta) = \Theta\left(\frac{k + \log(1/\delta)}{\varepsilon}\right)$ , where by KL we refer to minimizing  $\mathrm{KL}(\widehat{p} \parallel p)$ .

The optimal sample complexity of learning in  $\chi^2$  in terms of  $k, \varepsilon, \delta$ , however, remains an open question.

## References

- [Agr19] Rohit Agrawal. Concentration of the multinomial in Kullback-Leibler divergence near the ratio of alphabet and sample sizes. CoRR, abs/1904.02291, 2019. 2, 6, 3
- [BLM13] Stéphane Boucheron, Gábor Lugosi, and Pascal Massart. Concentration Inequalities A Nonasymptotic Theory of Independence. Oxford University Press, 2013. 5
- [Can15] Clément L. Canonne. A survey on distribution testing: Your data is big. but is it blue? *Electronic Colloquium on Computational Complexity (ECCC)*, 22:63, 2015. 2
- [KOPS15] Sudeep Kamath, Alon Orlitsky, Dheeraj Pichapati, and Ananda Theertha Suresh. On learning distributions from their samples. In *Proceedings of COLT*, volume 40, pages 1066–1100. PMLR, 2015. 3

 $<sup>^4</sup>$ You can verify this: intuitively, the issue boils down to having to non-trivially learn even the elements of the support of p that have arbitrarily small probability.