

2.3

Definition

A **weighted graph** G is a graph alongside a function $w : E(G) \rightarrow \mathbb{R}^+$.

If G is a weighted graph and $H \subseteq G$, then, $w(H) := \sum_{e \in E(H)} w(e)$.

A **minimum weight spanning tree** (or MWST) is a spanning tree T such that $w(T)$ is minimized among all possible spanning trees. In other words, $w(T) \leq w(T') \forall T' \subseteq G$ where T' is a spanning tree.

Kruskal's Algorithm

INPUT Weighted graph G with n vertices

OUTPUT A MWST, T^* if G is connected, otherwise a message “ G is not connected”

STEP 1 Create a list of edges, L_E , in order from smallest weight to largest weight. Start T^* with no edges but all vertices of G .

STEP 2 If the number of edges in T^* is strictly less than $n - 1$ AND if there are still edges in L_E , examine the first edge in L_E (i.e., the edge with smallest weight), $e = \{a, b\}$

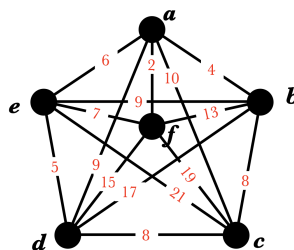
SUBSTEP 2.1 If a and b are in different components of T^* , add e to T^* and remove e from L_E . Return to STEP 2.

SUBSTEP 2.2 If a and b are in the same component, remove e from L_E and do not add to T^* . Return to STEP 2.

STEP 3 If the number of edges in T^* is $n - 1$, then output T^* , which is the MWST for G . Otherwise, the number of edges in T^* is strictly less than $n - 1$ and G was not connected.

Example

We will find a MWST for the following graph:



First, we create the following table of all the edges in G .

Edge	Cost
af	2
ab	4
de	5
ae	6
ef	7
bc	8
cd	8
ad	9
be	9
ac	10
bf	13
df	15
bd	17
cf	19
ce	21

We can read the following table describing the steps in Kruskal's algorithm from left to right (i.e., we check the edge of lowest weight, then we check the components, then we select our substep).

Edge	Components Of T'	Substep to be used
af	$\{a\}, \{b\}, \{c\}, \{d\}, \{e\}, \{f\}$	2.1
ab	$\{a, f\}, \{b\}, \{c\}, \{d\}, \{e\}$	2.1
de	$\{a, b, f\}, \{c\}, \{d\}, \{e\}$	2.1
ae	$\{a, b, f\}, \{c\}, \{d, e\}$	2.1
ef	$\{a, b, d, e, f\}, \{c\}$	2.2
bc	$\{a, b, d, e, f\}, \{c\}$	2.1
—	$\{a, b, c, d, e, f\}$	—

Proof of Kruskal's Algorithm

If G is connected, then Kruskal's Algorithm produces a minimum weight spanning tree.

Let G be a connected graph, and let T_K be the output from Kruskal's algorithm on G . It is easy to check that T_K is a spanning tree, since it is acyclic and has $n(G) - 1$ edges.

Suppose T^* is a MWST with largest edge intersection with T_K — in other words, $|E(T_K) \cap E(T^*)| \geq |E(T_K) \cap E(T')|$ for any other MWST T' .

If $T^* = T_K$, then we are done, since T_K is assumed to be a minimum weight spanning tree. Otherwise, assume toward contradiction that $T^* \neq T_K$. Let e be the first edge chosen by Kruskal's algorithm that is not in T^* . Then, by a previous result, $\exists e' \in E(T^*) - E(T_K)$ such that $T' = T^* + e - e'$ is a spanning tree.

We are assuming, however, that Kruskal's algorithm would choose e over e' . Let e_1, \dots, e_k be edges in $E(T_K)$ before e . Since e_i was selected before e , we know that $e_i \in E(T^*)$ for each $i \in [k]$.

Let $G_k = (V, \{e_1, \dots, e_k\})$. we are assuming that Kruskal's algorithm would not choose e' for two reasons:

- If e' shows up before e , then e' would connect two vertices in $G_j = (V, \{e_1, \dots, e_j\})$. Since $G_j + e' \subseteq T^*$, then T^* contains a cycle and isn't a spanning tree.
- If e' shows up after e in L_E , then $w(e') \geq w(e)$, meaning $w(T') \leq w(T^*)$, meaning that $w(T_K) \leq w(T^*)$, implying that T_K is of a lower weight than T^* .

Dijkstra's Algorithm

We want to find an algorithm to find the shortest path between two vertices in a weighted graph — Dijkstra's Algorithm can solve this.

We know that if P is the shortest u, v path, and $x \in P$, then following P from u to x will yield the shortest u, x path.

INPUT A weighted graph, G , and $u \in V(G)$

OUTPUT For each $z \in V(G)$, the distance $d(u, z)$.

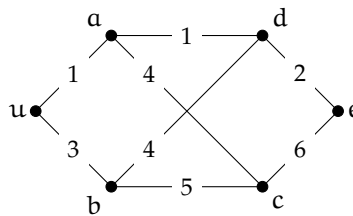
INITIALIZATION Extend the weight function such that if $xy \notin E(G)$, then $w(xy) = \infty$. Create S that contains all vertices whose distances from u are known. Let $S := \{u\}$. Let $t : V \rightarrow \mathbb{R}^+ \cup \{0, \infty\}$ which will keep track of the tentative distance between u and z . Let $t(z) := w(uz)$ for all $z \neq u$, and $t(u) := 0$.

CONDITION TO TERMINATE LOOP If $t(z) = \infty$ for all $z \notin S$ or $S = V$, then go to end.

LOOP Else, pick $v \in V - S$ such that $t(v) = \min_{z \notin S} t(z)$. Add v to S . Explore the edges from v to update tentative distances; for each edge vz with $z \notin S$, $t(z) := \min\{t(z), t(v) + w(vz)\}$.

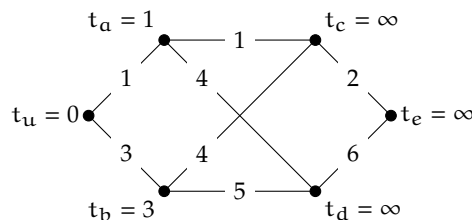
END Set $d(u, v) = t(v) \forall v \in V$.

On the following weighted graph, we can do Dijkstra's algorithm as follows:

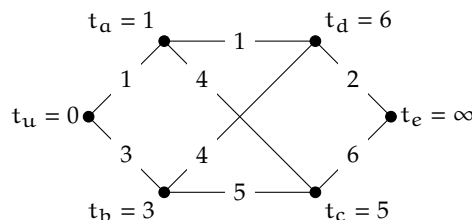


Dijkstra's Algorithm: Worked Example

We let $S = \{u\}$, and include our tentative distances for the first step.

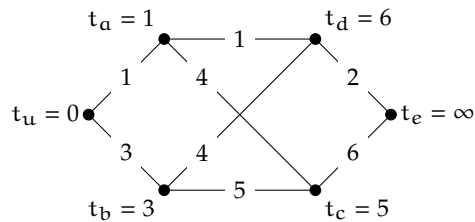


Now, $S = \{u, a\}$ because $t_a \leq t_b$. We now start from a and include our tentative distances.

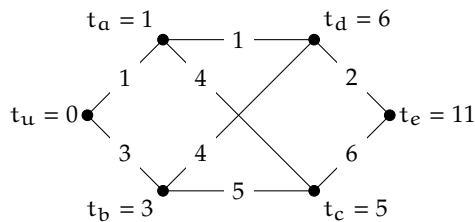


Next, we include b into S , making it $\{u, a, b\}$. We then check our tentative distances from b , where

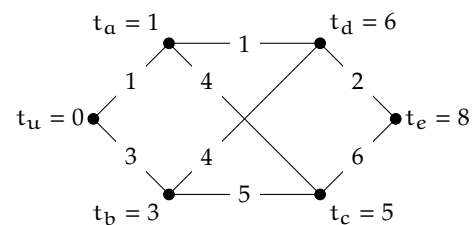
we find that they are no better than tentative distances from a , so we keep the tentative distances on c and d as what they were with a .



Next, we include c into S , making it $\{u, a, b, c\}$, and we update our tentative distances.



Finally, we include d and update tentative distances:



In order to find the direct paths, we can add arrows along the edges that were selected by Dijkstra's algorithm.

The proof can be outlined as follows:

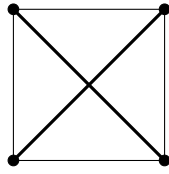
IF $z \in S$: then, $t(z) = d(u, z)$.

ELSE IF $z \notin S$: then, $t(z)$ is the length of a shortest u, z path P such that $V(P - z) \subseteq S$.

3.1

Matchings

In a simple graph, a **matching** M is a set of pairwise disjoint edges. In other words, $\forall e_i, e_j \in M$, then $e_i \cap e_j = \emptyset$. In an arbitrary graph, a matching is a set of non-loop edges with no shared endpoints. For example, the thick edges are a matching.



If v is incident to an edge, then v is **saturated** by M , otherwise it is unsaturated by M .

A **perfect matching** is a matching that saturates every vertex. We know that all graphs with perfect matchings have even number of vertices, but the alternative case may not be true (for example, we might consider a graph with an isolated vertex).

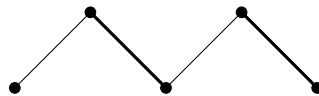
A **maximal matching** is a matching that cannot be enlarged by adding any other edges. A **maximum matching** is a matching of maximum size among all matchings. In other words, if M^* is a maximum matching, then $|M^*| \geq |M| \forall M \in G$.

- Not every maximal matching is a maximum matching.
- However, every maximum matching is a maximal matching (because you cannot extend a maximum matching by definition).

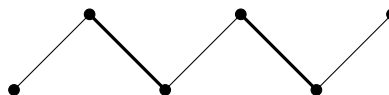
Alternating and Augmenting Paths

Given a matching M , an **M -alternating path** is a path that alternates between edges in M and edges not in M . An M -alternating path whose endpoints are unsaturated by M is an **M -augmenting path**.

An M -alternating path:



An M -augmenting path:

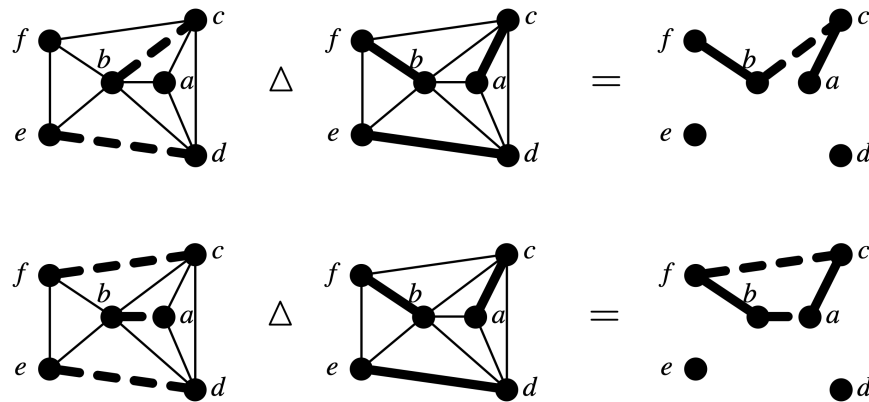


If M is a matching and there exists an M -augmenting path in G , then M is not a maximum matching (as in the path P that contains M , you can switch the matching edges as follows)



Symmetric Difference

The **symmetric difference** $G \Delta H$ is the subgraph of $G \cup H$ whose edges are the edges of $G \cup H$ that appear in exactly one of G and H . A picture of an example of a symmetric difference between matchings is shown below.



Theorems and Lemmas

Lemma 3.1.9

Every component of the symmetric difference of two matchings is a path or an even cycle.

Let M and M' be two matchings. Each vertex is incident to at most 1 edge in M and at most one edge in M' . Thus, $d(v) \leq 2$ for each vertex in M and M' . Therefore, each component is either a path or a cycle.

If a component is a cycle, then it must be even because the edges of the cycle must alternate between M and M' .

Theorem 3.1.10

A matching M in a graph G is a *maximum* matching if and only if G has no M -augmenting path.

(\Rightarrow) Suppose G has an M -augmenting path, P . Exchange the edges of P in M with the edges of P not in M . This action increases the size of the matching by 1, meaning that M was not a maximum matching initially.

(\Leftarrow) Suppose M is not a maximum matching. Let M' be a matching with more edges than M (i.e., $|M'| > |M|$). Let $H = (V(G), M \Delta M')$. By Lemma 3.1.9, we know that each component of H is either a path or an even cycle. Since $|M'| > |M|$, there is a component of H that contains more edges from M' than edges from M , which means it cannot be an even cycle — therefore, this component must be a path that alternates between M' and M . Because there are more edges from M' than M in this path, meaning this path is an M -augmenting path.

Perfect Matchings

Let $G = (X, E, Y)$ be a bipartite graph. A matching that saturates X is an **X -perfect matching**.

If there exists a set of vertices $S \subseteq X$ such that $|S| > |N(S)|$, then it is impossible for G to have an

X-perfect matching.

Theorem 3.1.11 (Hall, 1935)

A bipartite graph $G = (X, E, Y)$ has an X-perfect matching if and only if $|S| \leq |N(S)|$ for all $S \subseteq X$.

(\Rightarrow) If G has an X-perfect matching, then each vertex in X is matched to a distinct vertex in $N(S)$. Thus, $|S| \leq |N(S)|$.

(\Leftarrow) We will prove via the contrapositive. Suppose G does not have an X-perfect matching. Let M be a maximum matching in G , and let $u \in X$ be an M -unsaturated vertex. Let $S = \{x \in X : \exists M\text{-alternating } u, x \text{ path}\}$, and $T = \{y \in Y : \exists M\text{-alternating } u, y \text{ path}\}$.

We can partition S into $\{u\}, X_1, \dots, X_k$ and T into Y_1, \dots, Y_k , where $|X_i| = |Y_i|$ for $1 \leq i \leq k$. Our partition works as follows: $Y_1 = N(u)$, $X_1 = N_M(Y_1)$, $Y_2 = N(X_1) - Y_1$, $X_2 = N_M(Y_2)$, and the general form is $Y_i = N(X_{i-1}) - (Y_1 \cup \dots \cup Y_{i-1})$ and $X_i = N_M(Y_i)$.

Since M is a maximum matching, we know that there are no M -augmenting paths, meaning $|X_i| = |Y_i|$ (otherwise, if $|X_i| < |Y_i|$, we would be able to find a M -augmenting u, y path). Therefore, $|S| = |\{u\}| + |X_1| + \dots + |X_k|$ while $|T| = |Y_1| + \dots + |Y_k|$, so $|S| = |T| + 1$.

Corollary 3.1.13

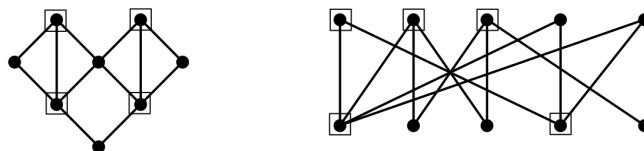
For $k > 0$, every k -regular bipartite graph has a perfect matching.

Let $k > 0$ and let $G = (X, E, Y)$ be a k -regular bipartite graph. We will show that G satisfies Hall's condition. Let $S \subseteq X$ and $[S, N(S)] = \{xy \in E(G) : x \in S \text{ and } y \in N(S)\}$, and $[X, N(S)] = \{xy \in E(G) : x \in X \text{ and } y \in N(S)\}$. Since $[S, N(S)] \subseteq [X, N(S)]$. Additionally, $\sum_{x \in S} d(x) = |[S, N(S)]|$ and $\sum_{y \in N(S)} d(y) = |[X, N(S)]|$.

Since G is k -regular, we find that $|[S, N(S)]| = k|S|$ and $|[X, N(S)]| = k|N(S)|$, meaning that $|S| \leq |N(S)|$.

Vertex Covers

A **vertex cover** is a set $Q \subseteq V(G)$ such that Q contains at least one endpoint of every edge. A picture of two vertex covers is shown below:



Of course, a very simple way to find a vertex cover is to select every vertex into Q . However, that isn't particularly useful, so we are focused on *small* vertex covers.

- The size of the **minimum vertex cover** is $\beta(G)$.

- The size of the **maximum matching** is $\alpha'(G)$ (the edge correspondent to $\alpha(G)$ for the size of the maximum independent set).

Alternatively, the minimum vertex cover is denoted $\tau(G)$ and the maximum matching is denoted $\nu(G)$.

An example of minimum vertex covers is shown below:



For any graph G , we can see that $\alpha'(G) \leq \beta(G)$ because to cover the edge of the maximum matching, we will need at least one vertex per edge in the maximum matching. Therefore, the minimum vertex cover must contain within it every vertex corresponding to an edge in the maximum matching, so $\beta(G) \geq \alpha'(G)$.

Theorem 3.1.16 (König-Egervary Theorem)

If G is a bipartite graph, then $\beta(G) = \alpha'(G)$.

Let $G = (X, E, Y)$ be a bipartite graph, and let M be a maximum matching in G . We are going to show that there exists a Q such that $|Q| = |M|$.

For each e in M , with endpoints $x \in X$ and $y \in Y$, select either x or y to be in Q as follows:

RULE: If there exists an M -alternating path that starts at an unsaturated vertex $u \in X$ and ends at y , then $y \in Q$ — else, $x \in Q$.

CLAIM: Q is a vertex cover. We will use cases to show that every edge has at least one endpoint in Q .

CASE 1: Suppose $xy \in M$. Then, $x \in Q \vee y \in Q$.

CASE 2: Suppose $xy \notin M$. If x is M -unsaturated, then y must be M -saturated (or else we would add the edge to the matching). We can create an M -alternating path that starts at x and ends at y , meaning that $y \in Q$.

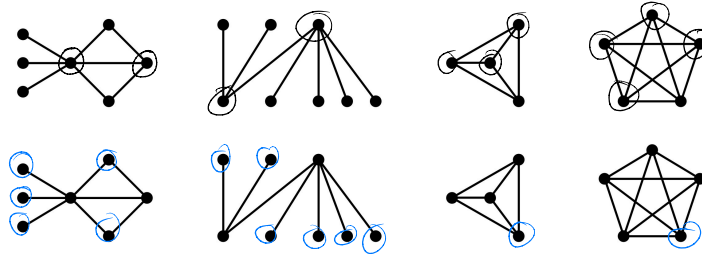
Otherwise, if x is M -saturated, let $xv \in M$ — then, we can find that $v \in Q$ because we create an M -alternating path that starts at $u \in X$ and ends at v . If P does not contain y , then we extend by going from v to x to y .

Edge Covers

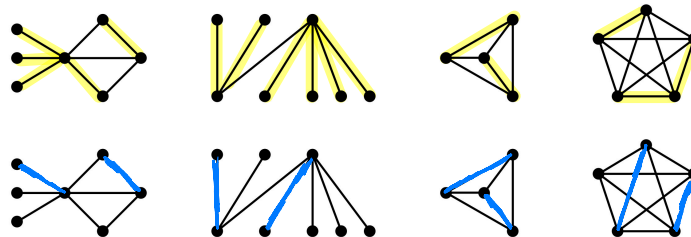
An **edge cover** is a set of edges A such that every vertex is incident to some edge in A . If A is an edge cover, then $\bigcup_{e \in A} e = V(G)$.

The **edge cover number** is the size of a minimum edge cover of G . It is denoted $\beta'(G)$.

We can see a relationship between vertex covers and independent sets as follows, where the top graph shows the minimum vertex cover and the bottom graph shows the maximum independent set.



Similarly, we can see a relationship between minimum edge covers and maximum matchings.



Theorems and Lemmas

Following from the previous result regarding independent sets and vertex covers, we might be able to state the following propositions:

Lemma 3.1.21

In a graph G , $S \subseteq V(G)$ is an independent set if and only if \bar{S} is a vertex cover. Therefore, we get the following result:

$$\alpha(G) + \beta(G) = n(G)$$

S is an independent set.

$$\Leftrightarrow \forall u, v \in S, u \not\leftrightarrow v.$$

$$\Leftrightarrow \forall e \in e(G), \text{ both endpoints of } e \text{ are not in } S.$$

$$\Leftrightarrow \text{Both endpoints of } e \text{ are in } \bar{S}.$$

$$\Leftrightarrow \bar{S} \text{ is a vertex cover.}$$

Thus, if S is a maximum independent set, then \bar{S} is a minimum vertex cover, and since $|S| + |\bar{S}| = n(G)$, this means $\alpha(G) + \beta(G) = n(G)$.

Theorem 3.1.22 (Gallai's Theorem)

If G is a graph without isolated vertices, then $\alpha'(G) + \beta'(G) = n(G)$. (i.e., if G is a graph without isolated vertices, then the size of a minimum edge cover summed with the size of the maximum matching equals the number of vertices in G).

We will prove the forward direction first, as follows:

Let M be a maximum matching, and let $A = V(M)$ and $B = \overline{A}$. Since M is a maximum matching, $\nexists e \in E(G)$ such that e connects two edges in B (or else you could simply add it to the matching). Additionally, since there are no isolated vertices in G , $\forall v \in B, \exists w \in A$ such that $v \leftrightarrow w$.

Select one edge such that $e = vw$ for each $v \in B$, and let the set of these edges be N . Then, $C = M \cup N$ is an edge cover, as every vertex in A is covered by M , and every vertex in B is covered by N . So, we have the following:

$$\begin{aligned} |M| + |C| &= |M| + (|M| + |N|) && \text{Since } M \text{ and } N \text{ are disjoint} \\ &= \left(\frac{|A|}{2}\right) + \left(\frac{|A|}{2} + |B|\right) && \text{By the definition of } M \text{ and } N \\ &= n(G) \end{aligned}$$

So, we have that $|M| + |C| = n(G) \geq \alpha'(G) + \beta'(G)$.

Now, we will prove the reverse direction:

Let C be a minimum edge cover. The components of C are stars (there are no cycles or copies of P_4). Let M be a matching formed by taking one edge from each edge in each component of C . Assign each "center" vertex in with an edge from M and assign each "leaf" edge with an edge from C .

Each vertex in $V(G)$ appears exactly once in each star, we have that $n(G) = |M| + |C|$. So, since $|C| = \beta'(G)$ and $|M| \leq \alpha'(G)$, we have that $|M| + |C| = n(G) \leq \alpha'(G) + \beta'(G)$.

Since $n(G) \geq \alpha'(G) + \beta'(G)$ and $n(G) \leq \alpha'(G) + \beta'(G)$, we have that $n(G) = \alpha'(G) + \beta'(G)$.

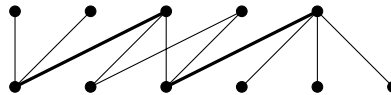
Corollary 3.1.24

If G is a bipartite graph with no isolated vertices, then $\alpha(G) = \beta'(G)$.

We previously showed that $\alpha(G) + \beta(G) = n(G)$ and $\alpha'(G) + \beta'(G) = n(G)$, and by the König-Egerváry theorem, we know that $\alpha'(G) = \beta(G)$ in bipartite graphs with non-isolated vertices, so $\alpha(G) + \beta(G) - \beta(G) = \alpha'(G) - \alpha'(G) + \beta'(G)$, so $\alpha(G) = \beta'(G)$.

Our Hungarian Method for Maximum Matchings

This is a method for finding a maximum matching in a bipartite graph that builds off the augmenting path algorithm. Consider the graph below:



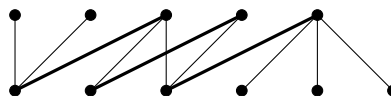
We will find the maximum matching by creating a “Hungarian Forest” as follows:

INITIALIZATION In our graph, we set $X_0 = \{x_1, x_2, x_4\}$ and $Y_0 = \{y_2, y_4, y_5, y_6\}$, where $X_0 \cup Y_0$ are the set of M -unsaturated vertices in G .

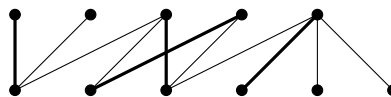
ITERATION Set $X_1 = X_0$ and $Y_1 = N(X_1)$. Then, we do $X_2 = N_M(X_1)$, and $Y_2 = N(X_2) - Y_1$, and so on. The general case is $X_n = N_M(Y_{n-1})$, and $Y_n = N(X_n) - Y_{n-1}$.

END CONDITION If Y_n is empty, then we stop the algorithm.

Let $Y^* = \bigcup Y_i$ (the vertices of Y that are in our Hungarian Forest). If $Y_0 \cap Y^* = \emptyset$, then M is a maximum matching. Otherwise, select $y \in Y_0 \cap Y^*$, and follow the path from y to the root of the tree to create an M -augmenting path, and “flip” the path (exchange matching edges with non-matching edges). In this case, after flipping the path starting between x_4 and y_2 , yielding the following graph.



After building our Hungarian forest in the same process, we find that $Y_0 \cap Y^* \neq \emptyset$, so we pick an M -augmenting path starting at y_4 , and flip the path starting at y_4 following up to x_1 . Afterwards, we find the graph as follows:



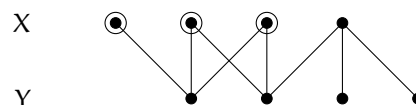
3.3

Preliminary

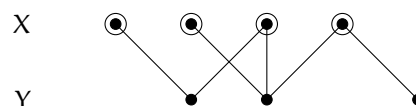
Q: Suppose $G = (X, E, Y)$ is a bipartite graph that *does not* have a X -perfect matching. What does Hall’s Theorem imply about G ?

A: There must exist a subset $A \subseteq X$ such that $|A| > |N(A)|$.

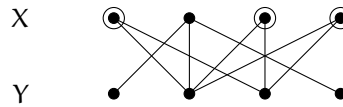
If you have a bipartite graph, you can *either* find an X -perfect matching, or you can find $A \subseteq X$ such that $|A| > |N(A)|$. We can test this proposition on the following set of bipartite graphs:



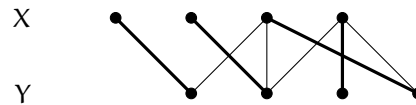
In the above graph, we can see that there is a set of vertices X such that the set of neighbors is of smaller cardinality.



In the above graph, we can see that the entirety of X is of a larger cardinality than the entirety of $N(X)$, so there cannot be an X -perfect matching.



In the above graph, we can see that there is a subset of X such that $N(S)$ is of smaller cardinality than S .

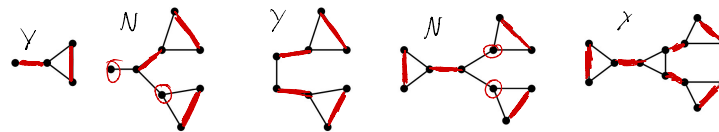


In the above graph, we have found an X -perfect matching.

We can infer from the previous exercise that if $G = (X, E, Y)$ and $|X| = |Y|$, and G does *not* have a perfect matching, then we can find both $A \subseteq X$ and $B \subseteq Y$ such that $|A| > |N(A)|$ and $|B| > |N(B)|$.

Conditions for a Perfect Matching

We can find a set of perfect matchings in the following graphs as follows:



From this, we can find some conditions for a perfect matching as follows:

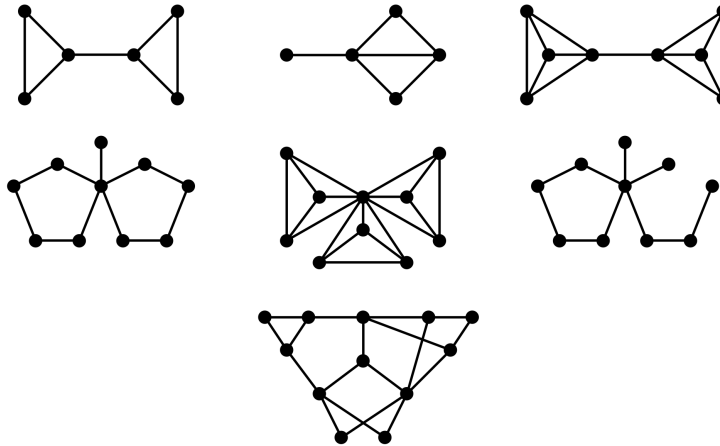
- $n(G)$ is even
- For each component H in G , $n(H)$ is even
- $\nexists v \in G$ such that $G - v$ has 2 or more odd components

We will let $q(G)$ denote the number of *odd* components in G . Our necessary condition is as follows:

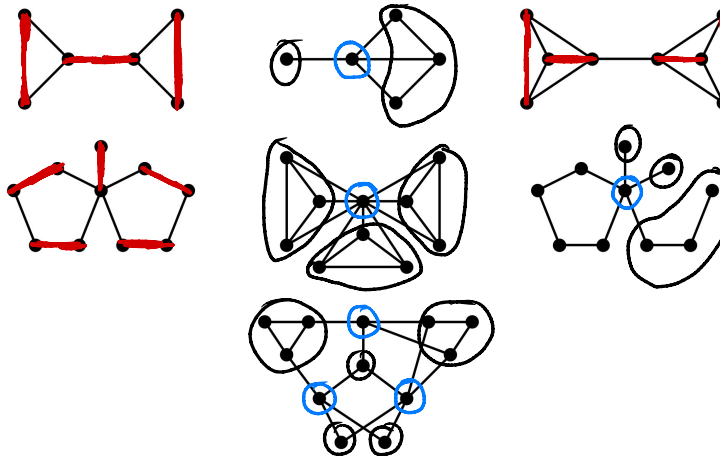
$$\forall S \subseteq V, q(G - S) \leq |S|$$

Exercise

For each of the following graphs, find a perfect matching or find S such that $q(G - S) > |S|$.



The solutions are as follows — matchings are denoted in red, deleted vertices in blue, and components after deletion in black.



Tutte's Theorem

A **factor** of a graph is a spanning subgraph of G , and a k -factor is a spanning, k -regular subgraph of G . Note that a 1-factor is essentially a perfect matching, but a matching is a set of edges while a factor is a graph. If M is a matching in G , then $H = (V, M)$ is a 1-factor.

A graph G has a 1-factor if and only if $q(G - S) \leq |S|$ for all $S \subseteq V$. A graph satisfies Tutte's Condition if $\forall S \subseteq V, q(G - S) \leq |S|$.

(\Rightarrow) To prove the forward direction, we know that the odd components of $G - S$ must have vertices matched to distinct vertices in S — this is because in any maximum matching within an odd component of $G - S$, there must be one vertex remaining, and since G has a 1-factor, each of

these spare vertices must be matched to a vertex in S .

(\Leftarrow) To prove the reverse direction, we will do so by contradiction. When we add an edge between two components of $G - S$, the number of *odd* components does not increase. So, Tutte's condition is preserved by the addition of edges, meaning that if $G' = G + e$, then $q(G' - S) \leq q(G - S) \leq |S|$, because we are assuming $q(G - S) \leq |S|$, and $q(G' - S) \leq q(G - S)$.

Additionally, if $G' = G + e$ has no 1-factor, then G has no 1-factor (since $q(G - S) \geq q(G' - S) > |S|$ in that scenario).

So, if there exists a simple graph G does not have a 1-factor and satisfies Tutte's condition, the theorem does not hold, nor does it hold if G would have a 1 factor upon adding an edge (or else G would have a 1-factor). Assume towards contradiction that G satisfies the Tutte condition or G' has a 1-factor while G does not. We will show that G must contain a 1-factor.

Let U be the set of vertices of degree $|V| - 1$.

CASE 1 Suppose $G - U$ consists of disjoint complete graphs. In this case, the vertices in each component of $G - U$ can be paired away, with one extra in the odd components. Since $q(G - U) \leq |U|$ by assumption, and each vertex of U is adjacent to every element of $G - U$, we can match the leftover vertices to U .

The remaining vertices are in U , which is a clique (since, by our assumption, every vertex in U is adjacent to every other vertex in G), meaning we only need to show that there are an even number of vertices in U . Seeing as we have matched an even number of vertices in G so far, it suffices to show that $|V|$ is even, which we can see if we apply the Tutte condition for $S = \emptyset$, which would mean that $q(G) \leq 0$, meaning $q(G) = 0$. Therefore, for the case where $G - U$ is a set of disjoint cliques, we have proved sufficiency.

CASE 2 Suppose $G - U$ is not a disjoint union of cliques. Then, there must be at least one component in $G - U$ that is not a clique, meaning that there is one pair of vertices $x, z \in V(G - U)$ such that $x \not\leftrightarrow z$, while $\exists y \in V(G - U)$ such that $x \leftrightarrow y$ and $y \leftrightarrow z$. Furthermore, $\exists w \in V(G - U)$ such that w and y are not adjacent to each other — if $y \in U$, then y would be adjacent to every vertex in $G - U$, but $y \notin U$.

By our choice of G , adding an edge to G creates a 1-factor. Let M_1 and M_2 be 1-factors defined in $G + xz$ and $G + yw$ respectively. All we need to show is that $M_1 \Delta M_2 \cup \{xy, yz\}$ contains a 1-factor avoiding xz and yw , implying that G contains a 1-factor. Let $F = M_1 \Delta M_2$. Then, since $xz \in M_1 - M_2$ and $yw \in M_2 - M_1$, we know that xz and yw are in F . Additionally, every edge in M_1 and M_2 is of degree 1, so if $ab \in F$, then $ac \in F$, implying that every vertex is of degree 0 or 2.

The components of F are either even cycles or isolated vertices. Let C be the cycle of F containing xz . If C does not contain yw , then we can create a 1-factor by taking all the edges of M_2 in C (as M_2 does not contain xz) and all the edges of M_1 not in C (as M_1 does not contain yw). Because this set contains every edge either in M_1 or M_2 , it is a 1-factor.

If C contains both yw and xz , then we use yx and xz to avoid them while constructing our 1-factor. We use M_1 to avoid yw , and then when we reach x or z , we use xy or zy respectively. In the remainder, we use the edges of C in M_2 . Seeing as we have avoided both xz and yw in this arrangement, and as seen previously we created a 1-factor by using the edges of M_2 in C , we know that this contains neither xz or yw . Finally, when we use either M_1 or M_2 outside of C , we have finished our 1-factor of G .

Corollary 3.3.7

The Berge-Tutte formula for the largest number of vertices saturated by a matching in G is as follows.

$$\min_{S \subseteq V(G)} \{n(G) - d(S)\}$$

where $d(S) = q(G - S) - |S|$.

If G has a perfect matching, then every unmatched vertex in a component of $G - S$ is matched to a vertex in S , meaning that $d(S) = 0$. If G does not have a perfect matching, then the meaning of minimization of $q(G - S) - |S|$ finds the excess number of vertices in S that cannot be matched to the odd components in $G - S$. Therefore, we have found a lowest bound for the number of vertices in the maximum matching.

Let $d = \max\{q(G - S) - |S|\}$. If $S = \emptyset$, then $d \geq 0$. Let $G' = G \vee K_d$ (the join of G to a complete graph on d vertices). Since $d(S)$ has the same parity as $n(G)$, we know that $n(G')$ is even. If G' satisfies Tutte's condition, then we know that G' has a perfect matching — by deleting d vertices, we obtain a matching of our desired size in G because deleting the d added vertices eliminates edges that saturate at most d vertices in G .