

## 2.1.22

Let  $T$  be an  $n$ -vertex tree with one vertex of each degree  $2 \leq i \leq k$ ; the remaining  $n - k + 1$  vertices are leaves. Determine  $n$  in terms of  $k$ .

We will find the number of vertices in  $T$  by finding the number of edges in  $T$  and adding 1. For  $2 \leq i \leq k$  corresponding to each of the non-leaf vertices, summation yields  $\frac{k(k+1)}{2} - 1$  edges. However, this scheme double-counts each edge, so we have to subtract the  $k - 2$  edges connecting the  $k - 1$  non-leaf vertices, yielding  $\frac{k(k+1)}{2} - k + 1$  edges. Finally, because  $T$  is a tree, we get that  $T$  has  $\frac{k(k+1)}{2} - k + 2$  vertices.

## 2.1.27

Let  $d_1, \dots, d_n$  be positive integers with  $n \geq 2$ . Prove that there exists a tree with vertex degrees  $d_1, \dots, d_n$  if and only if  $\sum d_i = 2n - 2$ .

- ( $\Rightarrow$ ) Suppose that for some tree  $T$ ,  $d_1, \dots, d_n$  are the degrees of the vertices of the tree. Since  $T$  is a tree, this means  $e(G) = n - 1$ , and  $\sum d_i = 2e(G)$ , meaning  $\sum d_i = 2(n - 1) = 2n - 2$ .
- ( $\Leftarrow$ ) Suppose that  $\sum d_i = 2n - 2$  for  $d_1, \dots, d_n$  corresponding to the degrees of the vertices in  $G$ . By a previous result, we know that  $\sum d_i = 2e(G)$ , meaning that  $\sum d_i = 2(n - 1)$ , implying that  $e(G) = n - 1$ . We can find a tree  $G$  with  $n - 1$  edges by letting  $G$  be connected with  $n - 1$  edges.

## 2.1.33

Let  $G$  be a connected  $n$ -vertex graph. Prove that  $G$  has exactly one cycle if and only if  $G$  has exactly  $n$  edges.

- ( $\Rightarrow$ ) Let  $G$  be a connected  $n$ -vertex graph with exactly one cycle. If we delete an edge from this cycle, then  $G - e$  is acyclic, as well as connected (since  $e$  is not a cut-edge), so  $G - e$  has  $n - 1$  edges. Adding back  $e$ , we get that  $G$  has  $n$  edges.
- ( $\Leftarrow$ ) Let  $G$  be a connected  $n$ -vertex graph with  $n$  edges. Then,  $G$  contains a spanning tree that contains all  $n$  vertices. Therefore,  $T \subseteq G$  contains  $n - 1$  edges. By adding another edge, we get that  $e(G) = e(T) + 1 = n - 1 + 1$ . Thus,  $G$  has exactly one cycle.

## 2.1.34

Let  $T$  be a tree with  $k$  edges, and let  $G$  be a  $n$ -vertex simple graph with more than  $n(k - 1) - \binom{k}{2}$  edges. Use Proposition 2.1.8 to prove that  $T \subseteq G$  if  $n > k$ .

We will use induction to prove that  $T \subseteq G$  as follows:

BASE CASE Suppose  $n = k + 1$ . Then, we can find the following:

$$\begin{aligned}
 e(G) &> (k+1)(k-1) - \binom{k}{2} \\
 e(G) &> (k^2 - 1) - \frac{k(k-1)}{2} \\
 e(G) &> \frac{k^2 - 1}{2} + \frac{k^2 - 1 - (k^2 - k)}{2} \\
 e(G) &> \frac{k^2 + k}{2} - 1 \\
 e(G) &> \frac{k(k+1)}{2} - 1
 \end{aligned}$$

This means  $e(G) = \frac{k(k+1)}{2}$  in the base case, meaning  $G$  is the complete graph on  $k+1$  vertices, where  $\delta(G) = k$ . By Theorem 2.1.8, we know that  $T \subseteq G$ .

INDUCTIVE HYPOTHESIS If  $n > k + 1$ ,  $e(G) > n(k-1) - \binom{k}{2}$ , then either  $\delta(G) \geq k$  or, if  $\delta(G) < k$ , then  $e(G-x) > (n-1)(k-1) - \binom{k}{2}$  for  $\delta(G) = d(x)$ .

PROOF If  $\delta(G) \geq k$ , then we know by Theorem 2.1.8 that  $T \subseteq G$ . Otherwise, suppose  $\delta(G) < k$ , and let  $d(x) = \delta(G)$ . Let  $G' = G - x$ .

$$\begin{aligned}
 e(G') &= e(G) - \delta(G) \\
 e(G') &\geq e(G) - (k-1) \\
 e(G') &> n(k-1) - \binom{k}{2} - (k-1) \\
 e(G') &> (n-1)(k-1) - \binom{k}{2}
 \end{aligned}$$

Therefore, the inductive hypothesis is proven.