

## 4.1

**Problem:** Write a Turing machine that computes the constant function  $C_0(n) = 0$ .

**Solution:** The Turing machine with the instructions

$$\begin{array}{l} q_1 1 B q_2 \\ q_2 B R q_3 \\ q_3 1 B q_2 \end{array}$$

computes  $C_0(n)$ .

- The Turing machine deletes the first 1, then moves into state  $q_2$ .
- The Turing machine reads the blank, moves right, and enters state  $q_3$ .
- The Turing machine reads the next input in state  $q_3$  — if 1, then blank, and go back into  $q_2$ , and if blank, it stops.

## 4.2

**Problem:** Write a Turing machine that computes the identity function  $f(m) = m$  on  $\mathbb{N}$ . Then write a machine for the identity function on  $\mathbb{N}^n$ .

**Solution:** The Turing machine with instructions

$$q_1 1 B q_2$$

computes  $f(m) = m$ .

The Turing machine with instructions

$$\begin{array}{l} q_1 1 B q_2 \\ q_2 B R q_3 \\ q_3 1 R q_3 \\ q_3 B R q_4 \\ q_4 B B q_5 \\ q_4 1 B q_2 \end{array}$$

computes  $\text{id}_{\mathbb{N}^n}$ .

## 4.3

**Problem:** Write a Turing machine that computes the function

$$Z(n) = \begin{cases} 1 & n = 0 \\ 0 & n \geq 1 \end{cases}.$$

**Solution:** The Turing machine with the instructions

$q_1 1 R q_2$   
 $q_2 1 B q_3$   
 $q_3 B L q_4$   
 $q_4 1 B q_5$   
 $q_5 B R q_5$   
 $q_5 1 1 q_6$   
 $q_6 1 B q_7$   
 $q_7 B R q_6$

computes  $Z(n)$ .

## 4.4

**Problem:** Write a Turing machine that computes  $f(m, n) = |m - n|$ .

**Solution:** The Turing machine with the instructions

$q_1 1 B q_2$   
 $q_2 B R q_3$   
 $q_3 1 R q_3$   
 $q_3 B R q_4$   
 $q_4 1 R q_5$   
 $q_5 1 R q_5$   
 $q_5 B L q_6$   
 $q_6 1 B q_7$   
 $q_7 B L q_8$   
 $q_8 1 L q_8$   
 $q_8 B L q_9$   
 $q_9 1 1 q_{10}$   
 $q_{10} 1 L q_{10}$   
 $q_{10} B R q_1$

computes  $f(m, n) = |m - n|$ . In short, for an input tape

$$B, \underbrace{1, 1, \dots, 1}_{m+1}, B, \underbrace{1, 1, \dots, 1}_{n+1}, B$$

the machine successively deletes the left-most 1 on the tape denoting  $m$  and deletes the right-most 1 on the tape denoting  $n$ , until it reaches the end of one of the domain elements.

## 4.6

**Problem:** Show that there is no Turing machine that can determine if a given Turing machine acting on a given input  $m$  will yield an output that contains the symbol  $s_k$  for a fixed  $k \geq 1$ .

**Solution:** Suppose  $V$  is computable. Define a Turing machine

$$G(n) = \begin{cases} 0 & V(n, n) = 0 \\ \text{undefined} & V(n, n) = 1 \end{cases}$$

Then,  $G$  is computable. I don't know how to do this problem.

## Extra Problem 1

**Problem:** Write the partial function computed by each of the following Turing machines.

- (a)  $q_1 1 R q_2, q_2 B R q_3, q_2 1 1 q_2.$
- (b)  $q_1 1 R q_2, q_2 B 1 q_3, q_2 1 R q_4, q_4 1 1 q_1.$

**Solution:**

- (a) This machine computes  $f(0) = 2$ , and does not halt for any other inputs.
- (b) This machine computes  $f(0) = 2$ , and does not halt for any other inputs.

## Extra Problem 2

**Problem:** Write a Turing machine that computes the function  $f(n) = 0$  if  $n$  is even and  $f(n) = 1$  if  $n$  is odd.

**Solution:**

$q_1 1 R q_2$   
 $q_2 1 B q_3$   
 $q_3 B R q_4$   
 $q_4 B B q_5$   
 $q_5 B L q_5$   
 $q_5 1 1 q_6$   
 $q_4 1 B q_7$   
 $q_7 B R q_8$   
 $q_8 B B q_9$   
 $q_9 B L q_9$   
 $q_9 1 B q_{10}$   
 $q_8 1 1 q_2.$