

## Problem 1.1.1

Determine which bipartite graphs are complete graphs

The graph  $K_{1,1}$  is the only bipartite graph that is complete.

## Problem 1.1.3

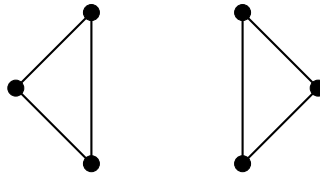
Using rectangular blocks whose entries are equal, write down an adjacency matrix for  $K_{m,n}$

$$K_{m,n} = \begin{array}{c} \begin{array}{cccccccc} & a_1 & a_2 & \cdots & a_m & b_1 & b_2 & \cdots & b_n \\ \begin{array}{c} a_1 \\ a_2 \\ \vdots \\ a_m \\ b_1 \\ b_2 \\ \vdots \\ b_n \end{array} & \begin{bmatrix} 0 & 0 & \cdots & 0 & 1 & 1 & \cdots & 1 \\ 0 & 0 & \cdots & 0 & 1 & 1 & \cdots & 1 \\ \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & 0 & 1 & 1 & \cdots & 1 \\ 1 & 1 & \cdots & 1 & 0 & 0 & \cdots & 0 \\ 1 & 1 & \cdots & 1 & 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 1 & 1 & \cdots & 1 & 0 & 0 & \cdots & 0 \end{bmatrix} \end{array} \end{array}$$

## Problem 1.1.5

Prove or disprove: If every vertex of a simple graph  $G$  has degree 2, then  $G$  is a cycle.

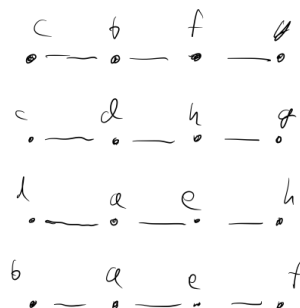
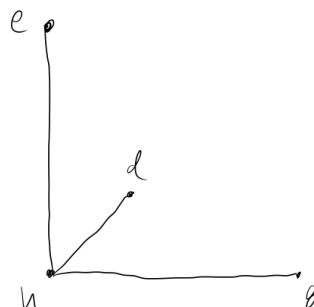
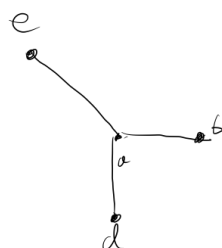
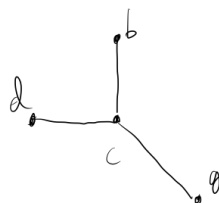
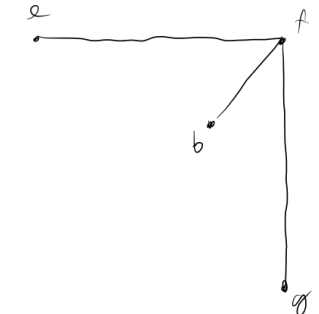
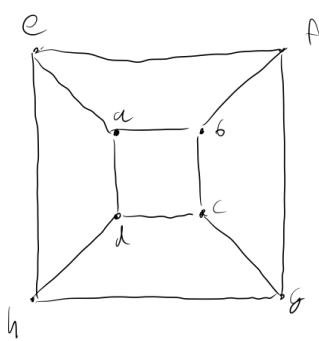
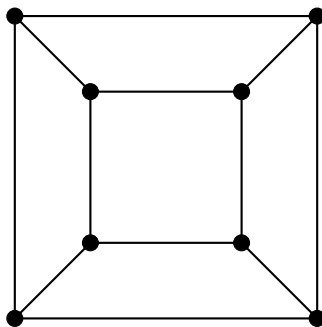
Let  $G$  be the following graph:



Every vertex in  $G$  has a degree 2, yet  $G$  is not a cycle.

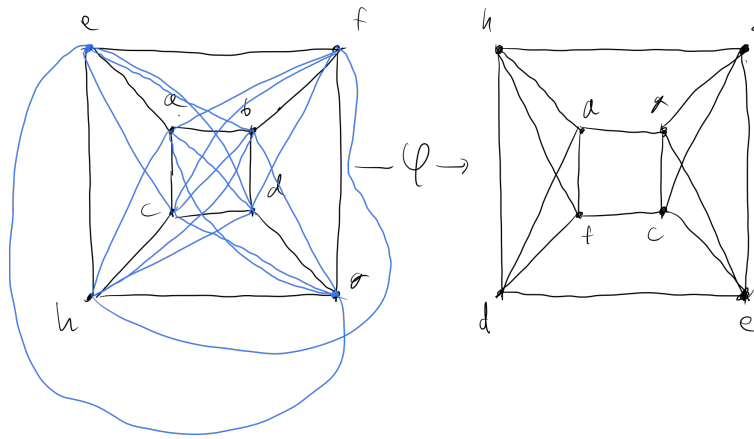
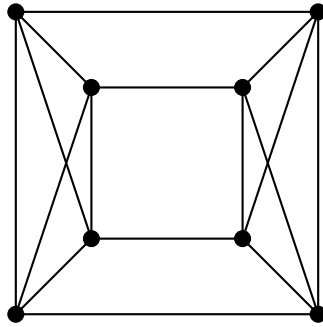
## Problem 1.1.8

Prove that the 8 vertex graph below decomposes into copies of  $K_{1,3}$  and also into copies of  $P_4$



## Problem 1.1.9

Prove that the graph below is isomorphic to the complement of the previous graph



## Problem 1.1.10

Prove or disprove: the complement of a simple disconnected graph must be connected.

Let  $G$  be a graph that is disconnected. We want to show that  $\forall x, y \in V(G), \exists x, z \text{ path}$ . We can split into two cases.

- Suppose  $x \leftrightarrow y$  in  $G$ . Then, in  $\overline{G}$ ,  $x \leftrightarrow y$  by the definition of a graph complement.
- Suppose  $x \nleftrightarrow y$  in  $G$ . Then, since  $G$  is disconnected, we know that there must be some  $z \in V(G)$  such that there is no  $x, z$  path. Since there is no  $x, z$  path, then there is no  $y, z$  path. In particular, this means  $x \nleftrightarrow z$  and  $y \nleftrightarrow z$  in  $G$ . Therefore, in  $\overline{G}$ , we have that  $x \leftrightarrow z$  and  $y \leftrightarrow z$ , meaning there is a path between  $x$  and  $y$ .