



LDL-precoded FTN Signaling with Power Allocation in The Block Fading Channel

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Outline

- **Motivation**
- **System Model**
- **Maximum average information rate**
- **Outage capacity**
- **Simulation Results**
- **Summary**



Outline

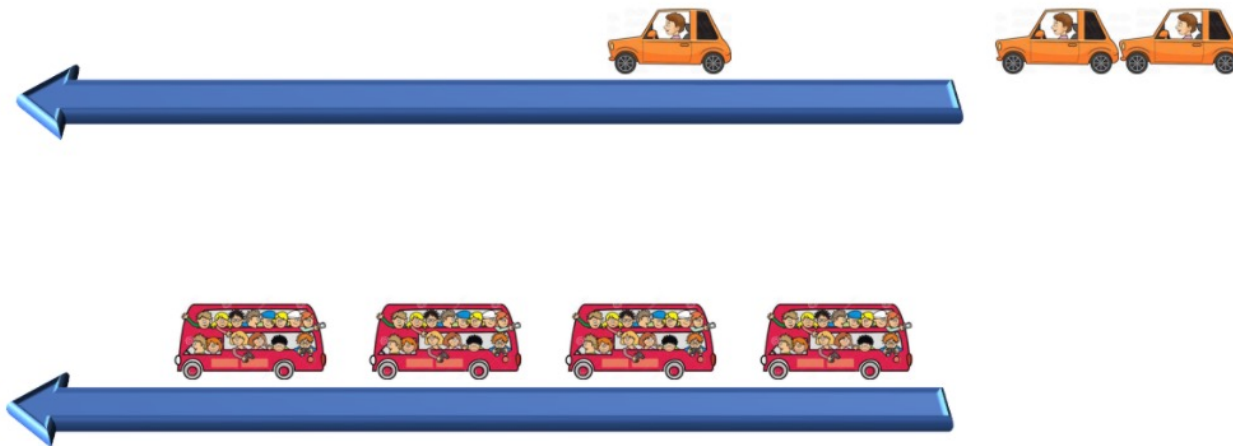
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Background

Spectrum efficiency is a crucial quota in communication systems.

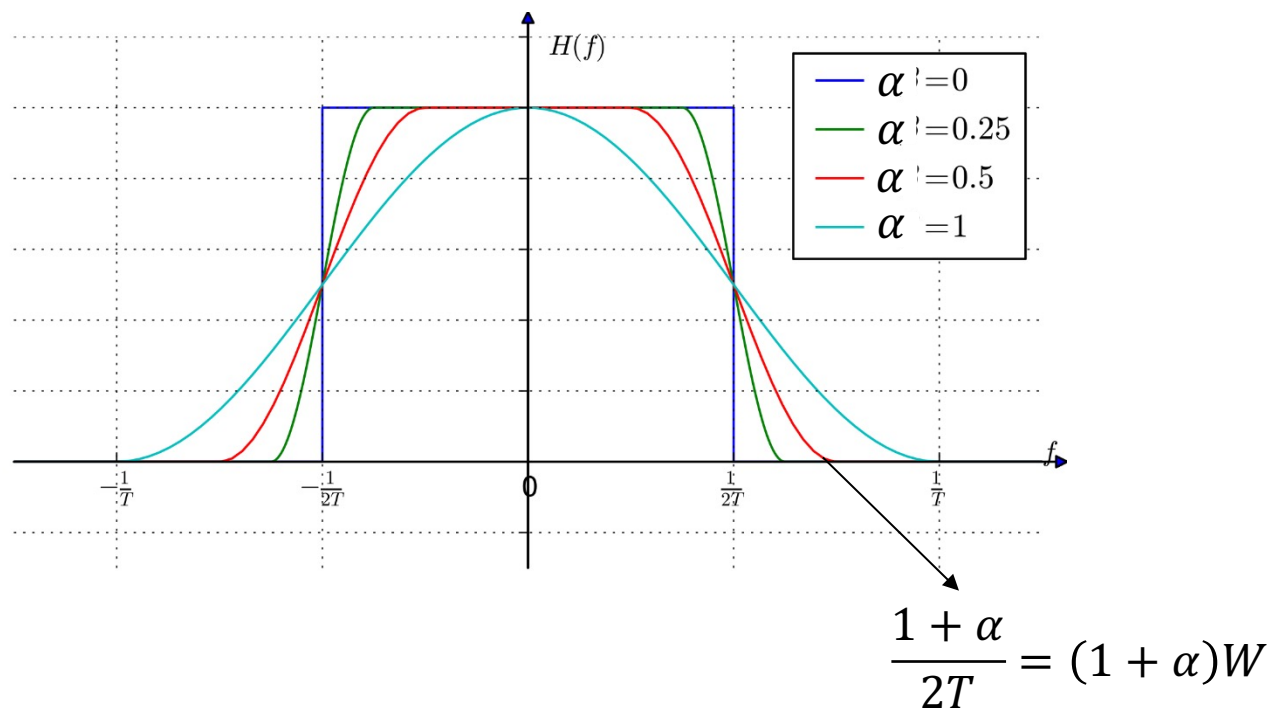
$$SE = \frac{R}{B}$$

where R is the information rate in bps and B is the bandwidth in Hz.



Just like using high-order modulations in communication systems!

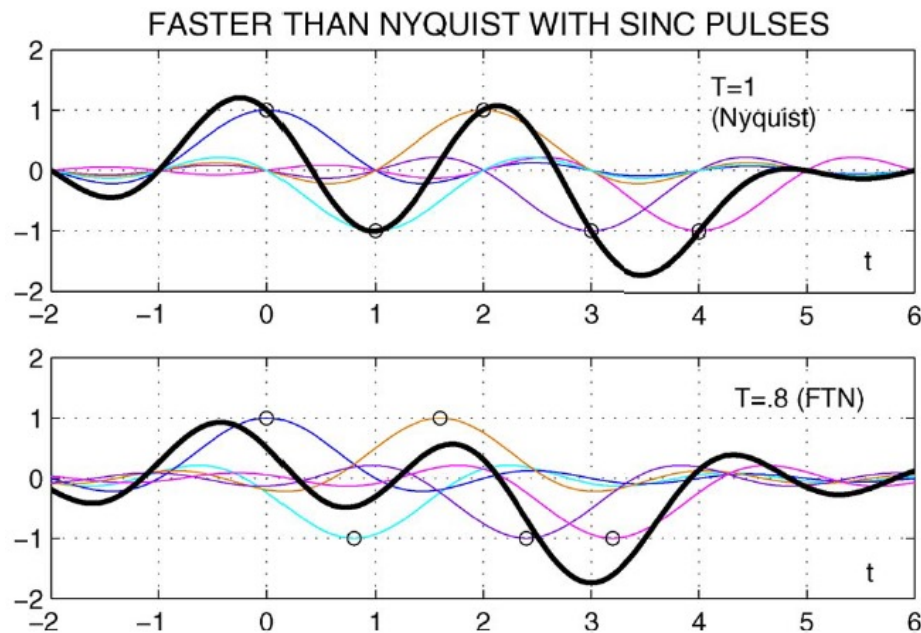
Background



Background

- What is faster-than-Nyquist (FTN) signaling?

$$x(t) = \sum_n x_n \phi(t - n\tau T) \quad \tau \in (0,1)$$



Background

- **Drawbacks of decoding algorithms for FTN signaling at receiver side**
 - Complexity increases exponentially as the considered ISI taps
 - Detection complexity may be excessive for low τ
- **Typical schemes for precoded FTN signaling**
 - Singular-value decomposition (SVD) based method
 - Square root decomposition based method
 - Cholesky decomposition based method

Background

- Complexity of typical schemes for precoded FTN signaling

Scheme	Formulation	flops
SVD	$\mathbf{G} = \mathbf{V}\mathbf{\Lambda}\mathbf{V}^T$	$O(2N^3)$
Square-root decomposition	$\mathbf{G} = \mathbf{G}^{\frac{1}{2}}\mathbf{G}^{\frac{1}{2}}$	$O(N^3)$
Cholesky decomposition	$\mathbf{G} = \mathbf{L}\mathbf{L}^T$	$O(N^3/3)$

Wide application range, but high complexity

limited application range, but low complexity

- ✓ LDL decomposition

Background

- **LDL decomposition**

For a given positive definite matrix \mathbf{G} , it can be decomposed as $\mathbf{G} = \mathbf{LDL}^T$, where \mathbf{L} is a lower triangular matrix and \mathbf{D} is a diagonal matrix.

In contrast, if \mathbf{G} is indefinite, $\mathbf{MGM}^T = \mathbf{LBL}^T$ holds, where \mathbf{M} is the permutation matrix and \mathbf{B} is the block diagonal matrix with 1×1 and 2×2 blocks on its diagonal.

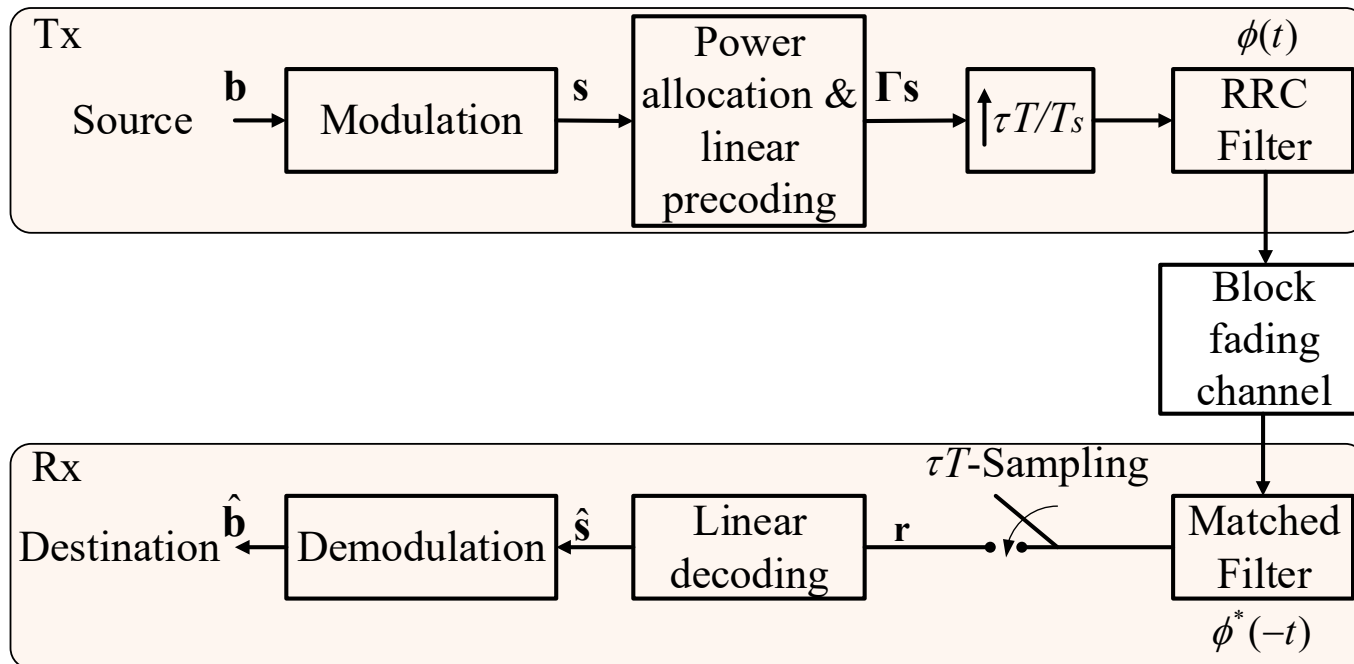
For the sake of unity, we use $\mathbf{G} = \mathbf{LDL}^T$ to represent these two forms.

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System Model

- ✓ An LDL-precoded FTN communication system with power allocation
- ✓ Block fading channel



System Model

- Power allocation and linear precoding

$$\mathbf{\Gamma} = (\mathbf{L}^T)^{-1} \mathbf{P},$$

where \mathbf{L} is obtained by $\mathbf{G} = \mathbf{L}\mathbf{D}\mathbf{L}^T$, $\mathbf{P} = \text{diag}\{\sqrt{p_0}, \dots, \sqrt{p_{N-1}}\}$.

- Received signal

$$r(t) = h \sum_n x_n g(t - n\tau T) + \eta(t)$$

$$\mathbf{r} = h\mathbf{G}\mathbf{x} + \boldsymbol{\eta} = h\mathbf{G}\mathbf{\Gamma}\mathbf{s} + \boldsymbol{\eta} \longrightarrow \text{colored noise, } E[\boldsymbol{\eta}\boldsymbol{\eta}^H] = N_0\mathbf{G}$$

- Linear decoding

$$\mathbf{\Phi} = \mathbf{\Gamma}^T$$

- Estimate of symbol vector

$$\begin{aligned} \hat{\mathbf{s}} &= \mathbf{\Phi}\mathbf{r} \\ &= h\mathbf{P}\mathbf{D}\mathbf{P}\mathbf{s} + \boldsymbol{\eta}_u \longrightarrow E[\boldsymbol{\eta}_u\boldsymbol{\eta}_u^H] = N_0\mathbf{P}\mathbf{D}\mathbf{P} \end{aligned}$$

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System Model

- **The positive definiteness of \mathbf{G}**

As the block length $N \rightarrow \infty$, \mathbf{G} is positive definite for $\tau \in \left[\frac{1}{1+\alpha}, 1\right)$, \mathbf{G} is indefinite for $\tau \in (0, \frac{1}{1+\alpha})$.

As mentioned before, $\mathbf{G} = \mathbf{L}\mathbf{D}\mathbf{L}^T$, where \mathbf{L} is a (permuted) lower triangular matrix and \mathbf{D} is the block diagonal matrix with 1×1 and 2×2 blocks on its diagonal. The 1×1 and 2×2 diagonal blocks can be represented as d_k and $\begin{bmatrix} d_l & d_{l,l+1} \\ d_{l+1,l} & d_{l+1} \end{bmatrix}$, respectively. $k \in A, l \in B$, and $|A| + 2|B| = N$

Lagrange Method

- Problem formulation**

- Objective: maximize the average information rate
- Variables: allocated power for each symbol
- Constraint: block energy N

- Instantaneous mutual information**

$$\begin{aligned}
 I_{ins}(\mathbf{s}; \hat{\mathbf{s}}) &= h(\hat{\mathbf{s}}) - h(\boldsymbol{\eta}_u) \\
 &\leq \frac{1}{2} \log_2 \left((2\pi e)^N \det(\text{Cov}(\hat{\mathbf{s}})) \right) - \frac{1}{2} \log_2 \left((2\pi e)^N \det(N_0 \mathbf{PDP}) \right) \\
 &= \log_2 \left(\det \left(\frac{\|h\|_2^2 \mathbf{PD}^T \mathbf{P}}{N_0} + \mathbf{I}_N \right) \right) \\
 &= \sum_{k=0}^{|A|-1} \log_2 \left(1 + \frac{\|h\|_2^2 d_k p_k}{N_0} \right) + \sum_{k=0}^{|B|-1} \log_2 (\rho_1 - \rho_2),
 \end{aligned}$$

$$\text{where } \rho_1 = \left(1 + \frac{\|h\|_2^2 d_l p_l}{N_0} \right) \left(1 + \frac{\|h\|_2^2 d_{l+1} p_{l+1}}{N_0} \right) \text{ and } \rho_2 = \frac{\|h\|_2^4 d_{l+1,l} d_{l,l+1} p_l p_{l+1}}{N_0^2}$$

Lagrange Method

- Transmit block energy

$$E_{blk} = E[(\mathbf{\Gamma}\mathbf{s})^H \mathbf{G}\mathbf{\Gamma}\mathbf{s}] = E[\mathbf{s}^H \mathbf{P}\mathbf{D}\mathbf{P}\mathbf{s}] = \sum_{i=0}^{N-1} d_i p_i$$

- Lagrange Function

$$\sum_{k=0}^{|A|-1} \log_2 \left(1 + \frac{\|h\|_2^2 d_k p_k}{N_0} \right) + \sum_{k=0}^{|B|-1} \log_2(\rho_1 - \rho_2) + \beta \left(\sum_{i=0}^{N-1} d_i p_i - N \right)$$

Maximum average information rate

✓ $\tau \in \left[\frac{1}{1+\alpha}, 1\right)$ **LDL-PA-FTN**

$|B| = 0$. Solving the Lagrange function, we have

$$p_n = \frac{1}{d_n}$$

Let $\gamma = \frac{\|h\|_2^2 P \tau T}{N_0}$ be the instantaneous signal to noise ratio (SNR) and substitute $p_n = \frac{1}{d_n}$ into $I_{ins}(\mathbf{s}; \hat{\mathbf{s}})$, we have

$$\begin{aligned} \bar{R} &= \int_0^\infty \lim_{N \rightarrow \infty} \frac{\max\{I_{ins}(\mathbf{s}; \hat{\mathbf{s}})\}}{N \tau T} \Pr(\gamma) d\gamma \\ &= \int_0^\infty \frac{1}{\tau T} \log_2(1 + \gamma) \Pr(\gamma) d\gamma \\ &\leq \frac{2W}{\tau} \log_2(1 + \bar{\gamma}) \text{ bits/s} \end{aligned}$$

where $W = \frac{1}{2T}$ and $\bar{\gamma} = E[\gamma]$.

Maximum average information rate

$$\sum_{k=0}^{|A|-1} \log_2 \left(1 + \frac{\|h\|_2^2 d_k p_k}{N_0} \right) + \sum_{k=0}^{|B|-1} \log_2(\rho_1 - \rho_2) + \beta \left(\sum_{i=0}^{N-1} d_i p_i - N \right)$$

✓ $\tau \in (0, \frac{1}{1+\alpha})$ **LDL-TPA-FTN**

B is not empty. Ignore B and set a threshold in A !

$$\sum_{k=0}^{|A|-1} \log_2 \left(1 + \frac{\|h\|_2^2 d_k p_k}{N_0} \right) + \beta \left(\sum_{k=0}^{|A|-1} d_k p_k - N \right)$$

In the same manner, we have

$$p_k = \begin{cases} \frac{N}{N_a d_k}, & d_k \geq \theta \\ 0, & d_k < \theta \end{cases}$$

$$\bar{R}_1 \leq \frac{2W}{\tau N} \log_2(1 + \bar{\gamma}) \text{ bits/s}$$

Number of active symbols

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Outage capacity

$$\gamma = \frac{||h||_2^2 P \tau T}{N_0}$$

$$p(\gamma) = \frac{1}{\bar{\gamma}} e^{-\frac{\gamma}{\bar{\gamma}}}$$

- LDL-PA-FTN**

Given required minimum information rate R_{min} , we have

$$\begin{aligned} P_o(R_{min}) &= \Pr \left[\frac{2W}{\tau} \log_2(1 + \gamma) \leq R_{min} \right] \\ &= \int_0^{2^{\frac{R_{min}\tau}{2W}} - 1} p(\gamma) d\gamma \\ &= 1 - \exp\left(-\frac{2^{\frac{R_{min}\tau}{2W}} - 1}{\bar{\gamma}}\right) \end{aligned}$$

Given the outage probability ε , we have

$$C_\varepsilon = \frac{2W}{\tau} \log_2 \left(1 + \frac{F^{-1}(1 - \varepsilon) P \tau T}{N_0} \right) \text{ bits/s}$$

where $F(v) = \Pr[||h||_2^2 > v]$.

Outage capacity

- **LDL-TPA-FTN**

Given required minimum information rate R_{min} , we have

$$Pt_o(R_{min}) = 1 - \exp\left(-\frac{2^{\frac{R_{min}\tau N}{2WN_a}-1}}{\bar{\gamma}}\right)$$

Given the outage probability ε , we have

$$Ct_\varepsilon = \frac{2WN_a}{\tau N} \log_2 \left(1 + \frac{F^{-1}(1 - \varepsilon)P\tau T}{N_0} \right) \text{ bits/s}$$

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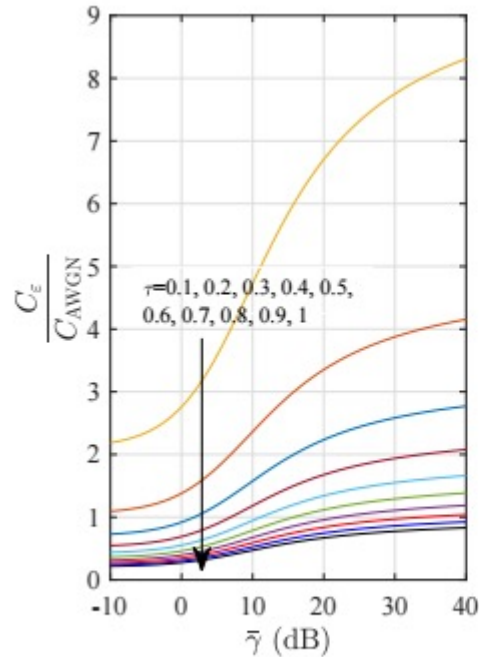
Simulation: Parameters

Name	Value
Orthogonal symbol duration (T)	1s
Sample interval (T_s)	0.01s
Modulation type	BPSK
Acceleration factor (τ)	{0.1, 0.2, ... , 1}
Shape filter type	Normal RRC filter/ Blackman windowed RRC filter
Shape filter order	1000/10000
Roll-off factor (α)	0.3
Channel type	Rayleigh block fading channel
Block length (N)	1000
Outage probability (ε)	0.1

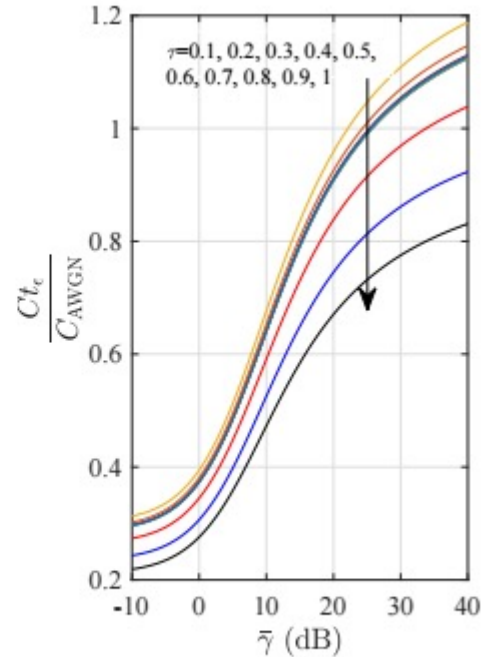
Simulation: ε – outage capacity to AWGN capacity ratio



$$\varepsilon = 0.1$$



(a) LDL-PA-FTN

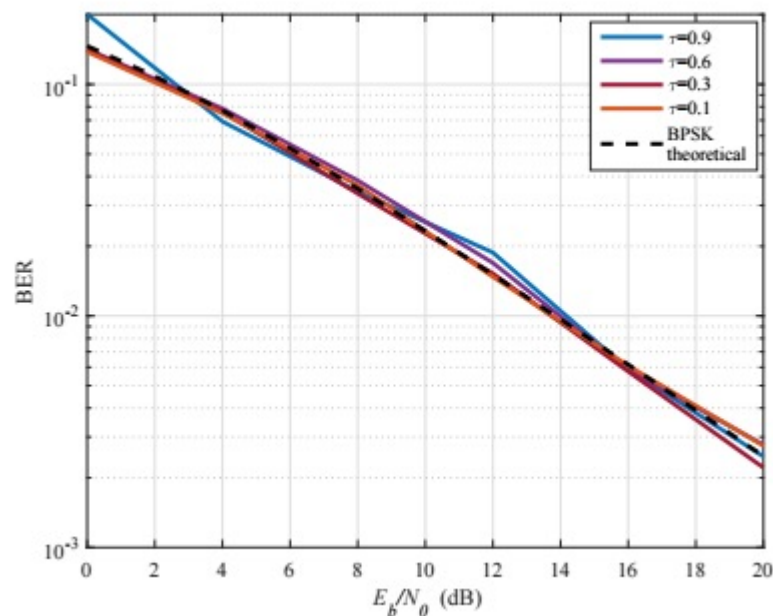


(b) LDL-TPA-FTN

$$\theta = 10^{-8}$$

- The gain in (a) is not able to obtain.
- Ct_ε/C_{AWGN} increases as τ decreases, but its improvement is small for $\tau \leq 0.7$

Simulation: BER of LDL-PA/TPA-FTN in the block fading channel



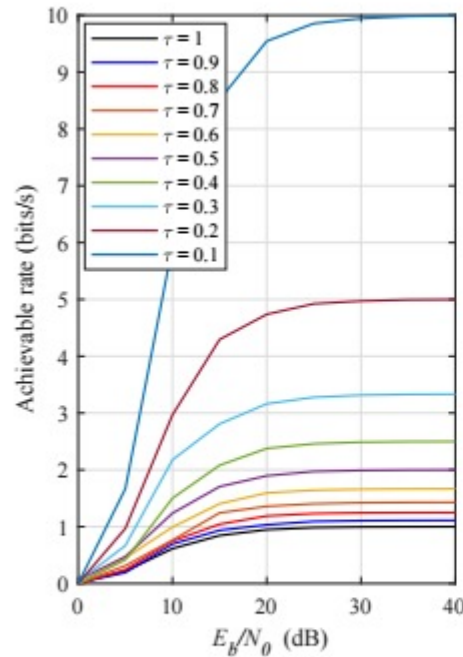
- The incurred ISI by FTN is well handled.

Simulation: Achievable rate

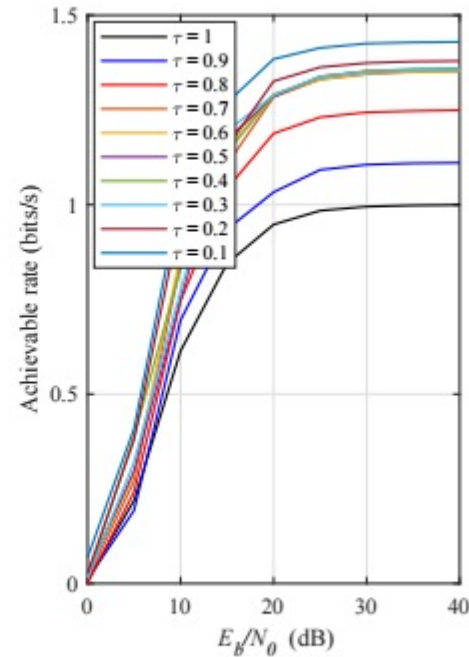
Packet error rate

$$\frac{N_a(1 - \text{PER})}{\tau NT}$$

$\theta = 10^{-8}$



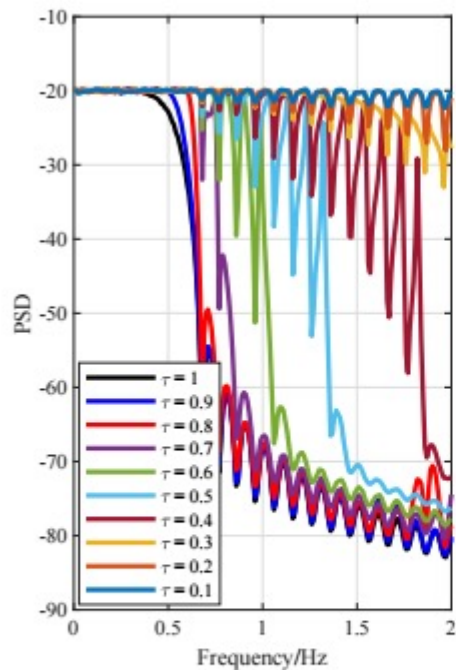
(a) Normal RRC filter with the order of 1000.



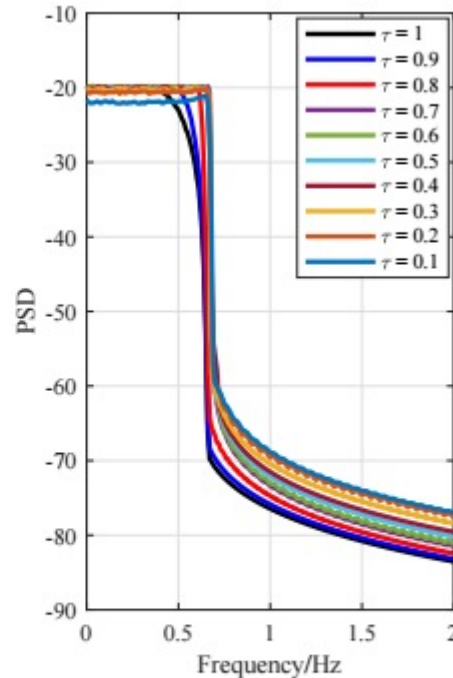
(b) Blackman windowed RRC filter with the order of 10000.

Simulation: PSD

$$\theta = 10^{-8}$$



(a) Normal RRC filter with the order of 1000.

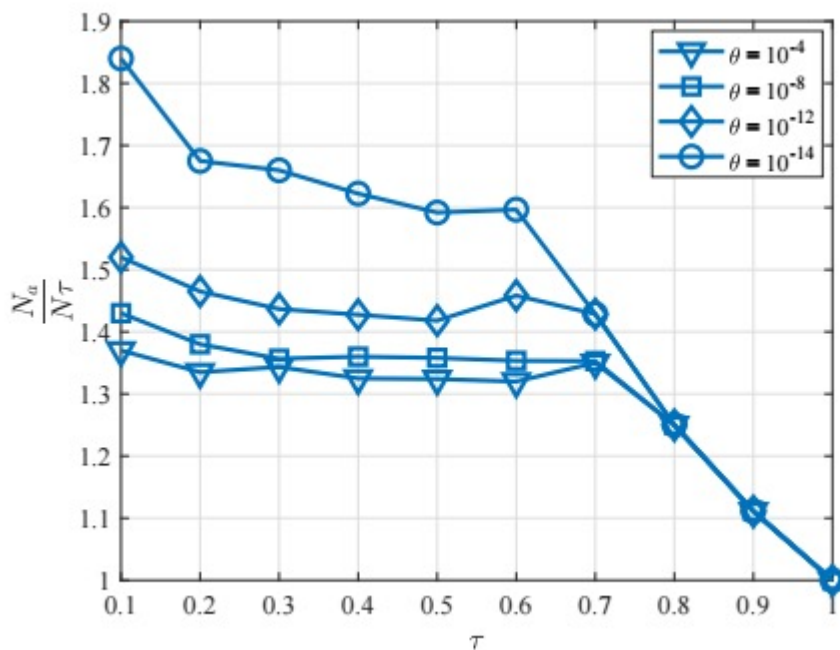


(b) Blackman windowed RRC filter with the order of 10000.

- A trade-off should be made between the side-lobe suppression of the PSD and achievable rate.

Simulation: Achievable rate

$$\frac{N_a}{\tau N}$$



- The performance gain of LDL-TPA-FTN depends on how small θ is.
- The precision of the decomposition of G should be considered carefully for LDL-TPA-FTN for small τ .

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Summary

- ✓ A low complexity precoded FTN scheme with power allocation
- ✓ The threshold selection in LDL-TPA-FTN may be troublesome in practice.
- ✓ The proposed scheme needs further evaluation, e.g., PAPR, IBI, high-order modulation.

Thank you!