## Unsupervised Data Mining: From Batch to Stream Mining Algorithms

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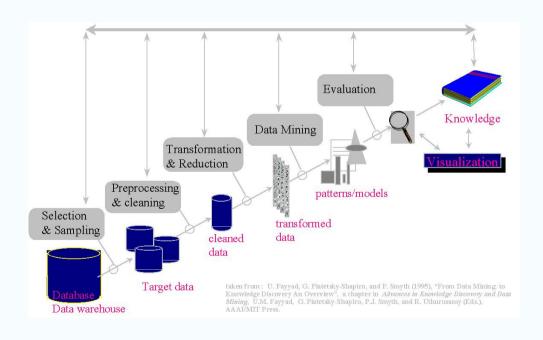
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## A Brief Introduction to Data Mining and KDD

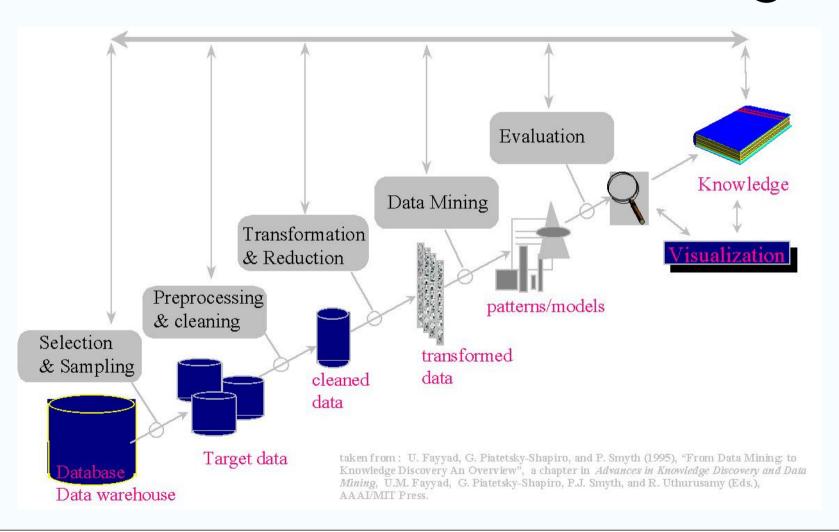
## Knowledge Discovery in Databases

"... is the process of identifying valid, novel, potentially useful and ultimately understandable structure in data."



(Fayyad & Uthurusamy, 1996)
Structure = pattern or model

# Knowledge Discovery in Databases and Data Mining



## **Data Mining**

- Knowledge Discovery in Databases (KDD) (Fayyad 96): "KDD is the non-trivial process of identifying valid, novel, potentially useful, and ultimately understandable patterns in data."
- Data Mining: data analysis step within the KDD process

### **Machine Learning**

- Learning = improving with experience at some task
  - Improve on task T
  - With respect to performance measure P
  - Based on experience E.
- Learn to play checkers:
  - T: Play checkers
  - P: % of games won
  - E: opportunity to play against oneself

### **Machine Learning**

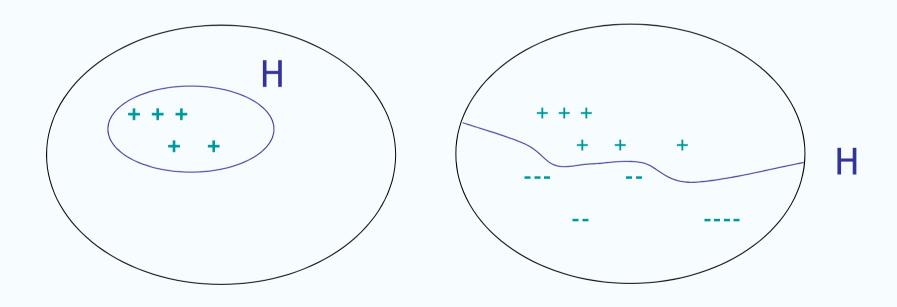
- Learning to classify examples (e.g., gene expression profiles into two subtypes):
  - T: Classifying examples
  - P: % of examples classified correctly
  - E: Training set of examples to learn from
- Machine learning algorithms (such as for classification) often used in Data Mining

#### Alternative Definitions...

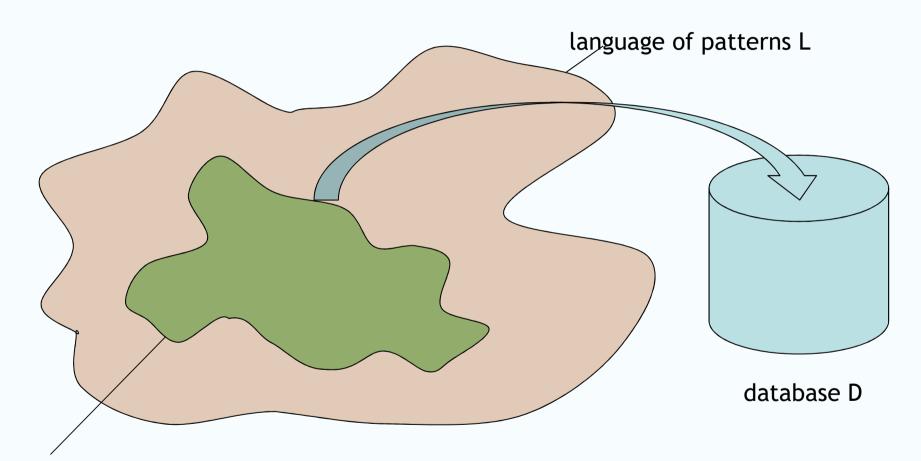
#### Heikki Mannila:

- "Knowledge Discovery in Databases is finding the joint probability distribution"
- "Data Mining is the technology of fast counting"

## Descriptive Data Mining, Predictive Data Mining

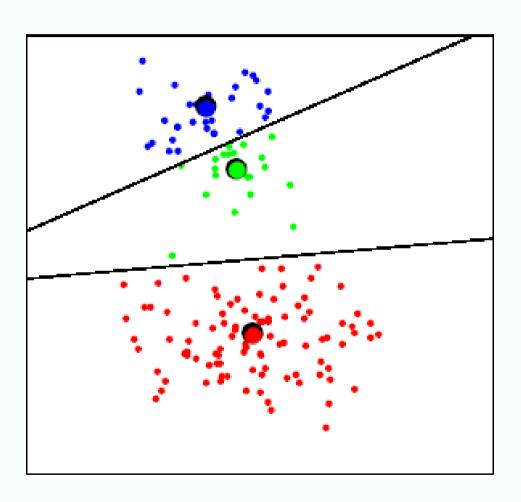


## Pattern Mining

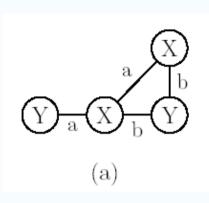


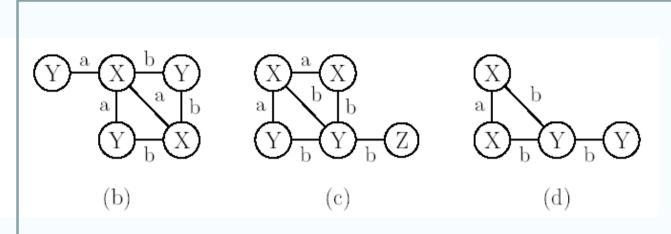
q(p, D) ... interestingness predicate: a pattern p from L is interesting wrt. database D what is interesting? frequent, non-redundant, class correlated, structurally diverse, ...

## Clustering



### **Graph Mining**



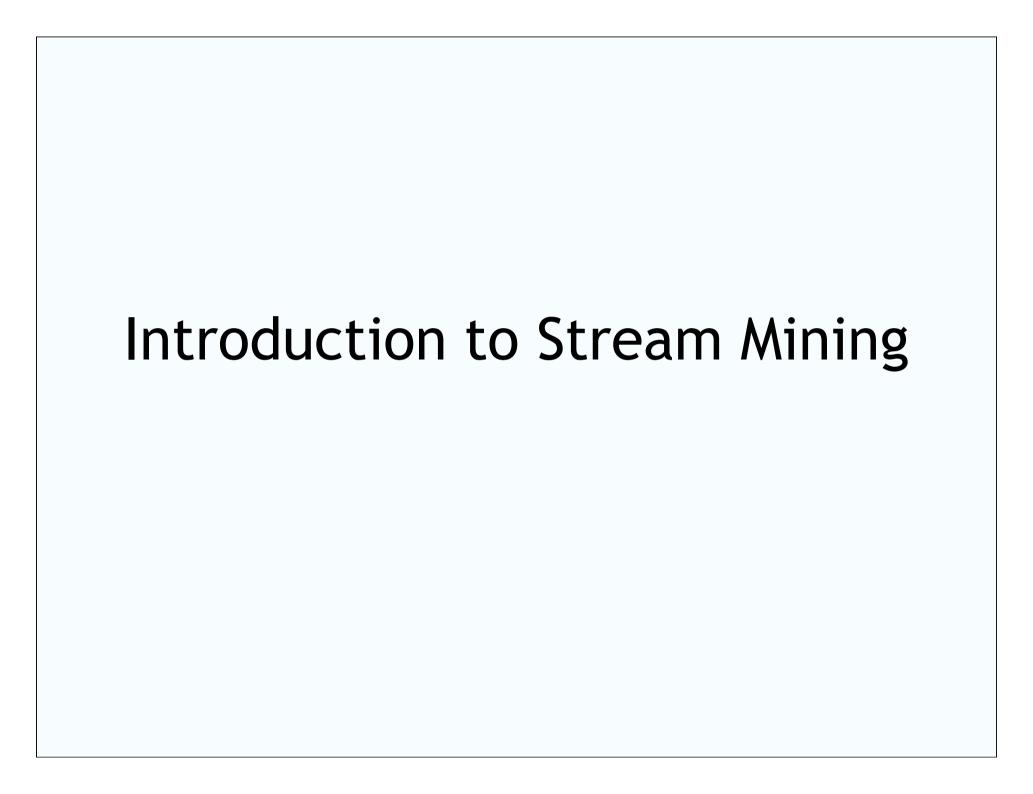


- Graph database D (graphs (b) to (d))
- Find all subgraphs (patterns) that occur in at least two of the three graphs (examples)
- Example subgraph pattern p shown in (a)<sub>12</sub>

## Stream Mining

Algorithm / Bounded Resource Analyzer

- Data stream model and philosophy
- Not one model
- Approximations



## Big (Data) Vs

#### Volume

- depends also on

   preprocessing or on
   operations on them (e.g.,
   pairwise comparisons)
   11 to 100 PB
   1.1 to 10 PB
   101 TB to 1 Petabyte
   1.1 to 100 TB
   1.1 to 10 TB
   101 GB to 1 Terabyte (TB)
   11 to 100 GB
   1.1 to 10 GB
   1.1 to 10 GB
- how much volume really?

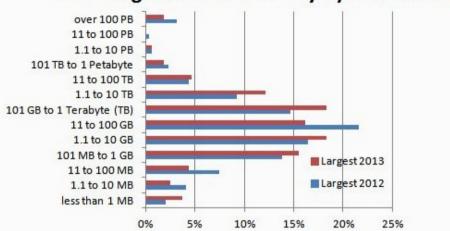
#### Variety

- often underestimated
- Velocity
  - analyzing data as they are generated ("one-touch"), real-time, anytime, ...

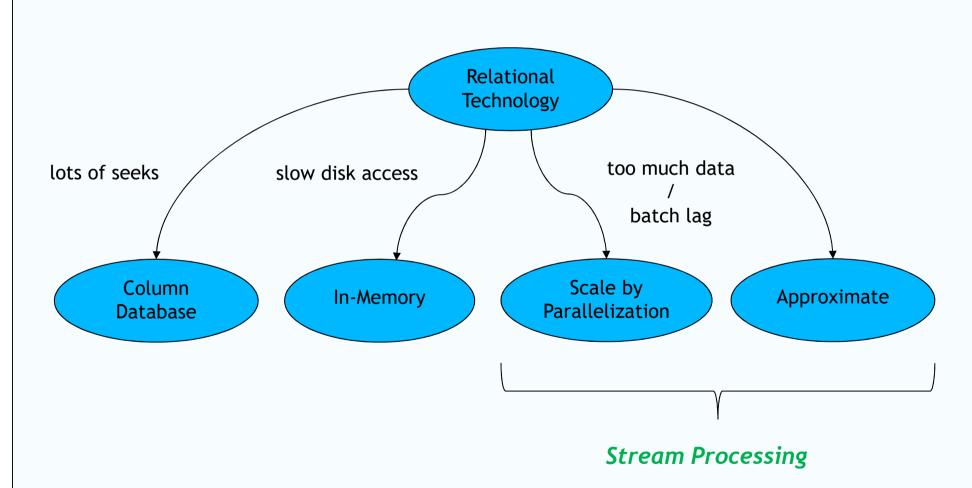
#### Veracity

uncertainty in data, data quality, trust, but also prediction quality

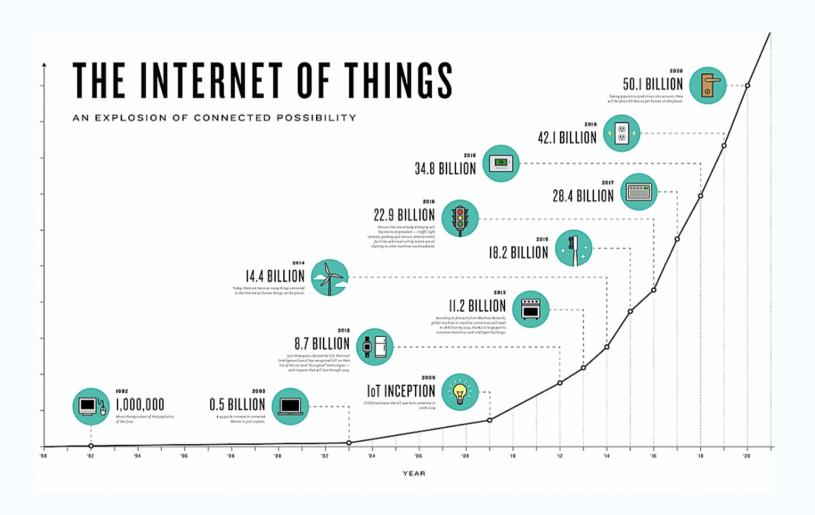
#### 2013 Largest Database Analyze/Data Mined



# Options for Scaling Up / Out (Partly Inspired by Mikio Braun)



## Expected Growth of Connected Devices



## Stream Mining

Algorithm / Bounded Resource Analyzer

Model

- Data stream model and philosophy
- Not one model
- Approximations

# Basic Stream Mining Algorithmics

#### Mean and Variance

Given a stream  $x_1, x_2, \ldots, x_n$ 

$$\bar{x}_n = \frac{1}{n} \cdot \sum_{i=1}^n x_i$$

$$\sigma_n^2 = \frac{1}{n-1} \cdot \sum_{i=1}^n (x_i - \bar{x}_i)^2.$$

#### Mean and Variance

Given a stream  $x_1, x_2, \ldots, x_n$ 

$$s_n = \sum_{i=1}^n x_i, \ q_n = \sum_{i=1}^n x_i^2$$

$$s_n = s_{n-1} + x_n, \ q_n = q_{n-1} + x_n^2$$

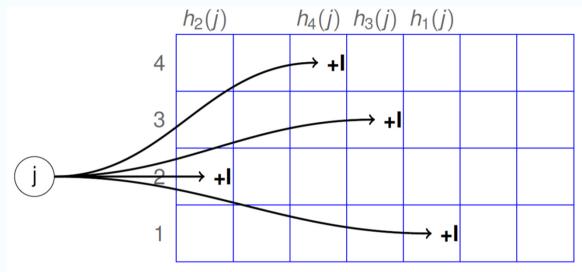
$$\bar{x}_n = s_n/n$$

$$\sigma_n^2 = \frac{1}{n-1} \cdot (\sum_{i=1}^n x_i^2 - n\bar{x}_i^2) = \frac{1}{n-1} \cdot (q_n - s_n^2/n)$$

#### Count Min Sketch

- Data structure for fast and memory-efficient counting on data streams
- Multiple hash tables and *pairwise independent* hash functions are used to update counts, effect of collisions is alleviated by taking the minimum of the results (i.e., the minimum of the counts from the table)
- Collisions still lead to overcounting, but this is upper-bounded, where the bound depends on the number d and the dimension w of the hash tables used.

## CM Sketch Example Structure



- Width of table = dimension of hash tables = w = 7
- Depth of table = number of hash tables = d = 4
- d and w are *derived* from a bound (see below)
- Assume parameters are set as follows  $\varepsilon = 0.4$ ,  $\delta = 0.02$
- Bound tells us that the count resulting from CM sketch  $\hat{a}_i \leq a_i + N^* \epsilon$  in all but  $\delta$  cases, with N being the length of the stream up to that point

# Determining Size of Table and Determining Count from Table

• If you want overcounting only by maximally N\* $\epsilon$  with a probability of 1-  $\delta$ , then you dimension the table by:

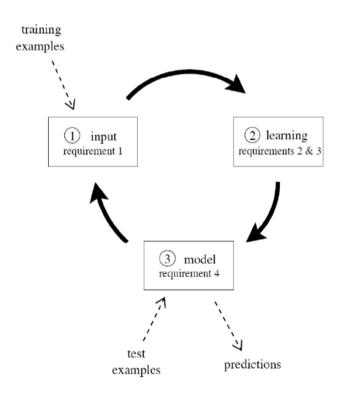
$$w = \left\lceil \frac{e}{\epsilon} \right\rceil, \qquad d = \left\lceil \ln \frac{1}{\delta} \right\rceil$$

- CM Sketch uses space w\*d and update time d
- Counts are determined by:
   â<sub>i</sub> = min<sub>j</sub> count[j, h<sub>j</sub>(i)]

# Classification on Data Streams and Hoeffding Trees

# Data Stream Classification Cycle

- Process an example at a time, and inspect it only once (at most)
- Use a limited amount of memory
- Work in a limited amount of time
- Be ready to predict at any point



## Prequential Testing

 First use new instance from stream to predict/test, then update the model based on it

 Approximates hold-out evaluation (testing on an "external" test set)

 Estimate accuracy using sliding windows or fading factors

## Hoeffding Tree Algorithm (Domingos & Hulten, KDD 2000)

Procedure HoeffdingTree(Stream,  $\delta$ ) Let HT = Tree with single leaf (root) Initialize sufficient statistics at root For each example (X, Y) in Stream Sort (X, Y) to leaf using HT Update sufficient statistics at leaf

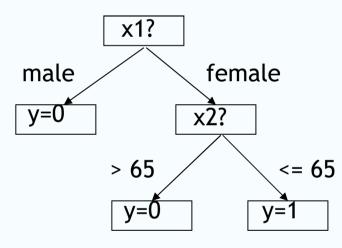
Compute G for each attribute  $R^2 \ln \left(\frac{1}{\delta}\right)$ If G(best) - G(2<sup>nd</sup> best) >  $\varepsilon = \sqrt{\frac{R^2 \ln \left(\frac{1}{\delta}\right)}{2\pi}}$ 

then Split leaf on best attribute

For each branch

Start new leaf, init sufficient statistics

Return HT



### Hoeffding Trees Preliminaries

- Suppose we have n observations of a real-valued random variable whose range is R (we assume R is 1 in the following) and mean is X.
- The Hoeffding bound states that with probability 1- $\delta$ , the true mean of the variable is at least  $\overline{X}$ - $\epsilon$ , where

$$\varepsilon = \sqrt{\frac{R^2 \ln(1/\delta)}{2n}}$$

## **Hoeffding Trees**

- Assume we have some evaluation function G like information gain or Gini index to assess the goodness of a split up to some number of examples n, and the difference between the best evaluated and the second best evaluated attribute is  $\Delta \overline{G} = \overline{G}(best) \cdot \overline{G}(2^{nd} best) \geq 0$ .
- Then, given a desired  $\delta$ , the Hoeffding bound guarantees that best is the correct choice with probability 1- $\delta$  and  $\Delta \overline{G}$ > $\epsilon$ .

## (Concept) Drift

- One of the main problems with stream mining methods
- Gradual drift, abrupt drift
- Drift in class or in attributes/features
- Recurrence of distributions
- Methods range from simple sliding window based ones (moving average type) to classifier based