Priority Queues and Heaps

A **priority queue** is an abstract data structure where each element has a priority. Elements are served based on priority rather than just order of insertion.

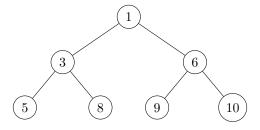
Heaps are a concrete implementation of priority queues. A heap is a complete binary tree satisfying the *heap property*:

• Min-Heap: Parent \leq children

• Max-Heap: Parent \geq children

1 Use Case: Task Scheduling

A heap can represent task priorities, where lower numbers indicate higher priority (Min-Heap). The following figure shows a binary Min-Heap storing integers:



Min-Heap Property: Each parent node is less than or equal to its children.

- Root node = minimum element
- Efficient for implementing priority queues
- Balanced due to complete binary tree structure

2 Terminology

• Heap: Complete binary tree with ordering property

• Priority: Determines service order

• Insert (Push): Add element and restore heap

• Extract (Pop): Remove min/max and restore heap

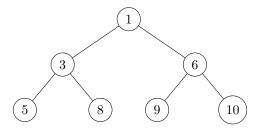
2.1 Array-Based Heap Representation

A binary heap can be stored efficiently using a **zero-based array** without using explicit pointers. Given a node at index i, its position relative to other nodes is determined by:

• Left Child: at index 2i + 1

- Right Child: at index 2i + 2
- Parent: at index $\left| \frac{i-1}{2} \right|$

Example: For the array [1, 3, 6, 5, 8, 9, 10] This corresponds to the following binary Min-Heap:



Benefits of Array Representation:

- Compact memory layout—no pointer overhead
- Easy index arithmetic for navigation
- Suited for cache-efficient access

Note: Array-based heaps are limited to complete binary trees. They must be restructured or resized carefully during insertions and deletions to maintain the heap shape and property.

3 C++ Implementation of a Min-Heap

```
#include <iostream>
#include <vector>
#include <stdexcept>
using namespace std;
class MinHeap {
  vector < int > heap;
  // Moves a node up the tree until the heap property is restored
  void heapifyUp(int i) {
    // While node is not the root and its value is less than its parent
    while (i > 0 && heap[i] < heap[(i - 1) / 2]) {</pre>
      \ensuremath{//} Swap the node with its parent
      swap(heap[i], heap[(i - 1) / 2]);
      // Move to the parent index
      i = (i - 1) / 2;
    }
  // Moves a node down the tree until the heap property is restored
  void heapifyDown(int i) {
    int n = heap.size();
    while (2*i + 1 < n) { // While at least one child exists
      int left = 2*i + 1;
      int right = 2*i + 2;
      int smallest = left;
      // Choose the smaller of the two children
      if (right < n && heap[right] < heap[left])</pre>
        smallest = right;
      // If the current node is smaller than both children, stop
      if (heap[i] <= heap[smallest])</pre>
        break;
```

```
// Otherwise, swap with the smaller child
swap(heap[i], heap[smallest]);
     i = smallest; // Continue at the child index
 }
public:
 void insert(int val) {
                              // Insert at the end
   heap.push_back(val);
   heapifyUp(heap.size() - 1); // Restore heap property
 int extractMin() {
   if (heap.empty()) throw runtime_error("Empty heap");
   heapifyDown(0);
                              // Restore heap property
   return minVal;
 int peek() const {
   if (heap.empty()) throw runtime_error("Empty heap");
   return heap[0];
                              // Return root without removing
 bool empty() const { return heap.empty(); }
```

Usage:

```
MinHeap h;
h.insert(5);
h.insert(2);
h.insert(8);
cout << h.extractMin(); // prints 2</pre>
```

3.1 STL priority_queue

C++ STL provides a container adapter for priority queues backed by a heap.

```
#include <queue>
#include <vector>
using namespace std;

priority_queue<int> maxpq; // default: max-heap
priority_queue<int, vector<int>, greater<int>> minpq; // min-heap

minpq.push(4);
minpq.push(1);
cout << minpq.top(); // prints 1
minpq.pop(); // removes 1</pre>
```

4 Time Complexities

- Insert: $\mathcal{O}(\log n)$
- Extract-Min/Max: $\mathcal{O}(\log n)$
- Peek: $\mathcal{O}(1)$
- Build-Heap (from array): O(n)

Why Are Heap Operations $O(\log n)$?

A binary heap is a complete binary tree stored as an array. Its height is $\log n$ (base 2), since each level doubles the number of nodes.

Operations like insert, extract-min, or decrease-key involve traversing from a leaf to the root or vice versa, adjusting positions to maintain the heap property. These traversals take at most $\log n$ steps, resulting in $O(\log n)$ time complexity.

5 Practice problem

Leetcode Problem 295: Find median from data stream