

$$a) \quad 5n^3 + 2n^2 + 3n = O(n^3)$$

$$5n^3 + 2n^2 + 3n < 5n^3 + 2n^3 + 3n^3 = C_1 \cdot n^3$$

$$5n^3 + 2n^3 + 3n^3 < C_1 \cdot n^3 \quad n > 1$$

$$10n^3 < C_1 \cdot n^3$$

$$C_1 = 10 \quad [n, n > 1]$$

$$b) \quad \sqrt{7n^2 + 2n + 8} = O(n)$$

$$C_1 \cdot n < \sqrt{7n^2 + 2n + 8} < C_2 \cdot n$$

$$C_1 \cdot n < \sqrt{7n^2 + 2n^2} < C_2 \cdot n$$

$$C_1 \cdot n < \sqrt{9n^2} < C_2 \cdot n$$

$$C_1 \cdot n < 3n < C_2 \cdot n$$

$$C_1 = 2 \quad n > 1$$

$$C_2 = 4$$

$$c) O(f(n)) \cdot O(g(n)) = O(f(n)g(n))$$

$$0 < f(n) < C_1 f(n) \cdot 0 < g(n) < C_2 g(n) = 0 < f(n)g(n) < C_3 \cdot f(n)g(n)$$

$$C_1 f(n) \cdot C_2 g(n) = C_3 f(n)g(n)$$

$$C_1 C_2 = C_3 \quad n \geq 1$$

2) If example 1:

$n = \text{len list} \quad \Theta(1)$

$\text{total} = 0 \quad \Theta(1)$

For j in range (n) :

For k in range $(i+j)$

$\text{total} += \text{lst}[k]$

return $\text{total} \quad \Theta(1)$

$\left. \begin{array}{l} \text{for } n \text{ times} \\ n-1 \text{ times} \end{array} \right\} n \text{ work}$

$\Theta(n^2)$

Example 1 $\Theta(n^2)$

def example 2:

$n = \text{len lst} \quad \Theta(1)$

$\text{prefix} = 0$

$\text{total} = 0$

For j in range (n)

$\text{prefix} += \text{lst}[j]$

$\text{total} += \text{prefix}$

return total

$\left. \begin{array}{l} \text{prefix} += \text{lst}[j] \\ \text{total} += \text{prefix} \end{array} \right\} \text{constant operation}$

$\left. \begin{array}{l} \text{for } n \\ \text{times} \end{array} \right\} \Theta(n)$

mp2, $\Theta(n)$

c) $1, 2, 4, 8, 16, 32, 64$ cost steps at 64
 to find ~~with~~ n work bit by
 $\text{avg } \log_2(n) \cdot n = 2 \log(n) = \Theta(\log^2 n)$

d) def `change n(n):`
 $i = n$ $sm = 0$
 while $(i \geq 1)$:
 for j in `range(i)`:
 $sm += i \cdot j$
 $i = i // 2$
 return sm

$\left[n + \frac{n}{2} + \frac{n}{4} + \frac{n}{8} + \dots \right]$ calculate how much work
 is done

$$\sum_{i=0}^{\log_2 n} n = n \cdot \left(\frac{1 - \frac{1}{2^{\log_2 n + 1}}}{1 - \frac{1}{2}} \right)$$

$$= \frac{n \cdot (1 - 2^{-\log_2 n - 1})}{1 - \frac{1}{2}}$$

$$2n \cdot \left(1 - \frac{1}{2^{\log_2 n + 1}} \right) \approx 2n$$

$$2n \left(1 - \frac{1}{n} \right) = 2n - 2 = \Theta(n)$$