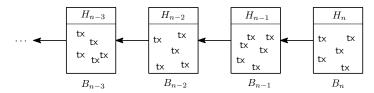
Introduction to IOTA – a feeless cryptocurrency

Serguei Popov

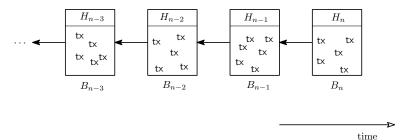
IOTA Foundation

Blockchain: decentralized consensus (Satoshi Nakamoto, 2008)





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Bitcoin mining: find a *nonce* N_{n+1} such that

$$\operatorname{hash}(\textit{H}_\textit{n},\textit{N}_\textit{n+1}) \leq 0.00\dots01$$

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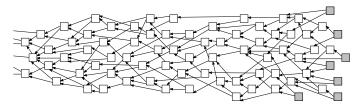
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- so: "help the others, and the others will help you; however, if you don't help the others, the others won't help you"

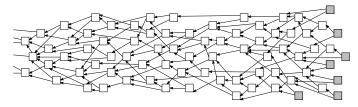
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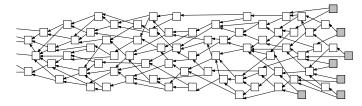


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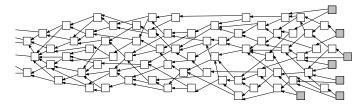
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- actors will behave in a "reasonable" way because it's a good idea to do so, and the designer's role is to propose a "good" set of rules (e.g. for tip selection)
- possible applications of AI in the future versions of the protocol for example, to approach better reputation systems (i.e., one can think about using AI for detecting node's malicious behavior)

Deeper look: consensus in IOTA Tangle

The Coo:

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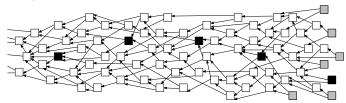
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- consensus as an attracting state

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- test implementation: github.com/iotaledger/goshimmer

Majority dynamics (threshold Voter Models):

- there is a graph, each site of which has an "opinion", 0 or 1
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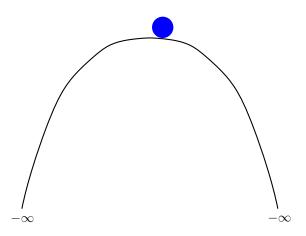
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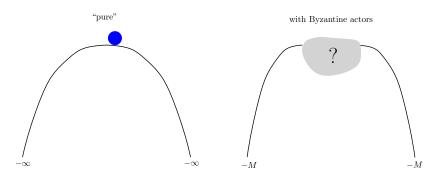
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- in particular: only two extremal invariant measures ("all-0" and "all-1"), which are consensus states.

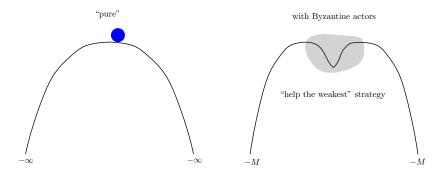
Majority dynamics as a RW on a potential:



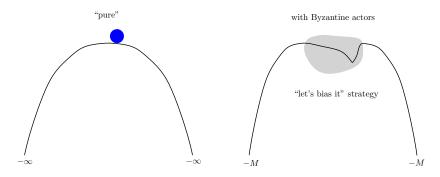
Majority dynamics with Byzantine actors:



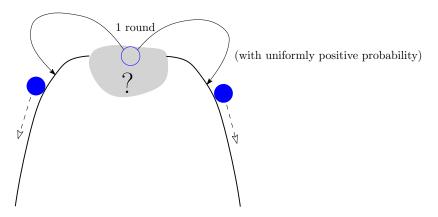
Majority dynamics with Byzantine actors: the curse of metastability



Majority dynamics with Byzantine actors: the curse of metastability



Fast Probabilistic Consensus (FPC): defeating the metastability with turn-based common random thresholds:



Questions?