# The Art of Being Lazy: 2

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#### New abstraction

freeze( $\langle expr \rangle$ )  $\rightarrow$   $\langle expr \rangle$  is not evaluated now, but postponed indefinitely.  $\langle expr \rangle$  is wrapped in a *thunk*.

freeze is not a normal function. Its argument is not evaluated immediately.

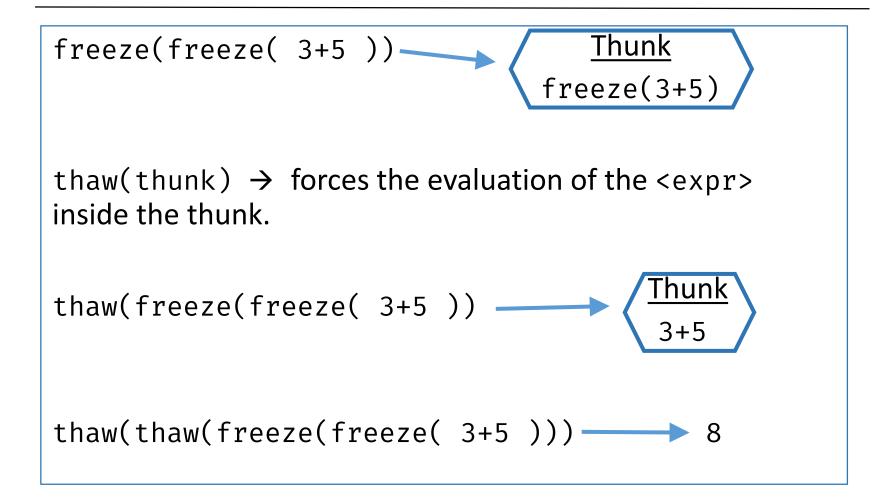
Its secret will be revealed shortly.

```
compare: int foo(int n) \{ \dots \} foo( 3+5 ) \rightarrow foo( 8 )
```

The argument (which is an expression) is eagerly evaluated BEFORE foo is called. But,

freeze( 
$$3+5$$
 ) The expression  $3+5$  is not evaluated but frozen in a thunk

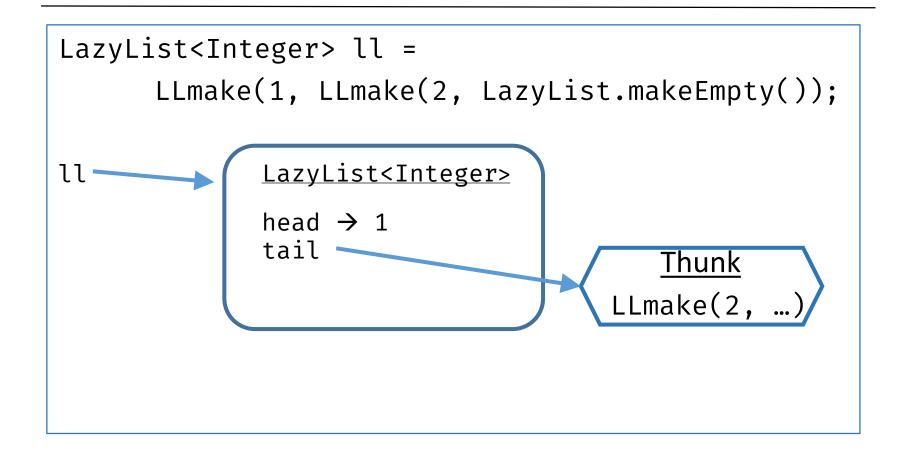
#### New abstraction



#### Lazy Lists

```
public class LazyList<T> {
   private final T head;
   private final Thunk(LazyList<T>) tail;
   . . . }
Constructors:
LLmake(a, b) = new LazyList<>(a, freeze(b))
Note that a is eagerly evaluated, but b is frozen.
LLmake is syntactic sugar, and not a normal function.
```

#### Lazy Lists: example



# Integers from a (inclusive) to b (exclusive)

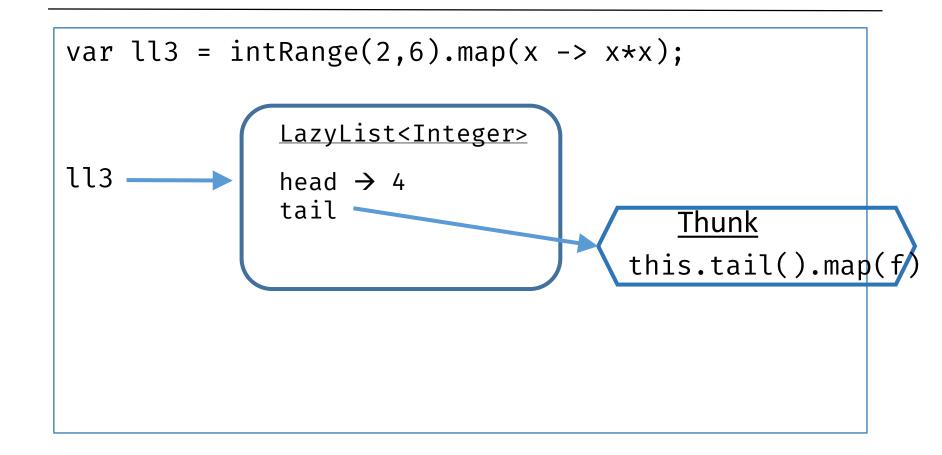
```
LazyList<Integer> intRange(int a, int b) {
    if (a >= b)
        return LazyList.makeEmpty();
    return LLmake(a, intRange(a+1, b)); }
var ll2 = intRange(4, 12);
              LazyList<Integer>
              head \rightarrow 4
              tail
                                           Thunk
                                       intRange(a+1, b)
```

## map: apply function to every element

Note the use of recursion: we recursively apply map to the tail of the list.

But this is frozen because it is the 2<sup>nd</sup> argument of LLmake.

# Compute the squares of 2,3,4,5



#### filter: keep elements that satisfy predicate

```
public LazyList<T> filter(Predicate<T> pred) {
    if (this.isEmpty())
        return this;
    else if (pred.test(this.head()))
        return LLmake(this.head(),
                      this.tail().filter(pred));
    else
        return this.tail().filter(pred); }
```

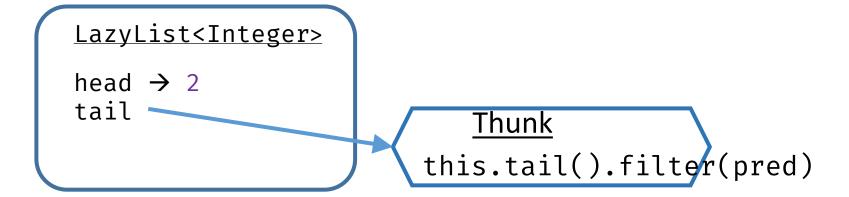
#### filter: keep elements that satisfy predicate

```
intRange(1,21).filter(x-> x%2==0)
    LazyList<Integer>
    head \rightarrow 1
    tail
                                    Thunk
                                intRange(a+1, b)
pred.test(1) \rightarrow false
so, intRange(2, 21) will be called.
       LazyList<Integer>
       head \rightarrow 2
        tail
                                       Thunk
```

intRange(a+1, b)

#### filter: keep elements that satisfy predicate

```
pred.test(2) → true
so, return LLmake(2, this.tail().filter(pred))
```



#### elementWiseCombine

```
if (this.isEmpty() || other.isEmpty())
   return LazyList.makeEmpty();
 else
   return LLmake(binOp.apply(this.head(),
                            other.'head()).
              this.tail()
          var ll = intRange(1, 11);
var twos = ll.map(x->2);
ll.elementWiseCombine(twos, (x,y)->x+y).print();
=> (* 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, *)
```

#### reduce

This combines all the elements, two at a time, via the generic accumulator. The resulting value is the "aggregate" of all elements.

### eg: add the squares of 1, 2, ..., 20

#### This computes:

$$1^2 + 2^2 + \ldots + 19^2 + 20^2$$

## eg2: factorial may be expressed as

```
int factorial(int n) {
        return LazyList.intRange(1, n+1)
            .reduce(1, (x,y) -> x*y);
```

#### Demo

# Infinite LazyLists!

```
LazyList<Integer> integersFrom(int n) {
    return LLmake(n, integersFrom(n+1)); }
integersFrom(1).limit(50).print();

=> (* 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 20, 21, 22, 23, 24, 25, 26, 27, 28, 29, 30, 31, 32, 33, 34, 35, 36, 37, 38, 39, 40, 41, 42, 43, 44, 45, 46, 47, 48, 49, 50, *)
```

#### Prime numbers

Of course, one way to check if a number n is prime is to see if n can be found in a list, L, of primes. To generate L, an ancient method called the Sieve of Eratosthenes may be used:

- ① Start by listing all integers greater than 1. Call this list C. Also, let L be an empty list.
- 2 Take the first number p=2 in C, and add it to L. This is the first prime.
- In C, cross out all multiples of p.
- 4 Let p be the next uncrossed number in C. This is the next prime. Add it to L, and repeat from Step 3.

	2	3	A	5	8	7	<b>X</b>	<b>X</b>	M
11	)22	13	14	<b>X</b> 5	<b>X</b> 6	17	<b>1</b> %	19	20
24	22	23	24	25	26	27	28	29	30

```
LazyList<Integer> sieve(LazyList<Integer> s) {
   return LLmake(s.head(),
                 sieve(s.tail()
                         .filter(x->
                                 x%s.head()!=0)));
var primes = sieve(integersFrom(2));
primes.limit(50).print();
=> (* 2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41,
43, 47, 53, 59, 61, 67, 71, 73, 79, 83, 89, 97, 101,
103, 107, 109, 113, 127, 131, 137, 139, 149, 151,
157, 163, 167, 173, 179, 181, 191, 193, 197, 199,
211, 223, 227, 229, *)
```

## Another way to get integers

Demo

```
LazyList<Integer> integers;
integers = LLmake(1, integers.map(x->x+1));
```

#### Fibonacci numbers

#### The infinite sequence:

The next number is the sum of the previous two.

0	1	1	2	3	5			<b>←</b> fib
		0	1	1	2	3	5	<b>←</b> fib
		1	1	2	3	5		fib.tail()

# The BIG Reveal!

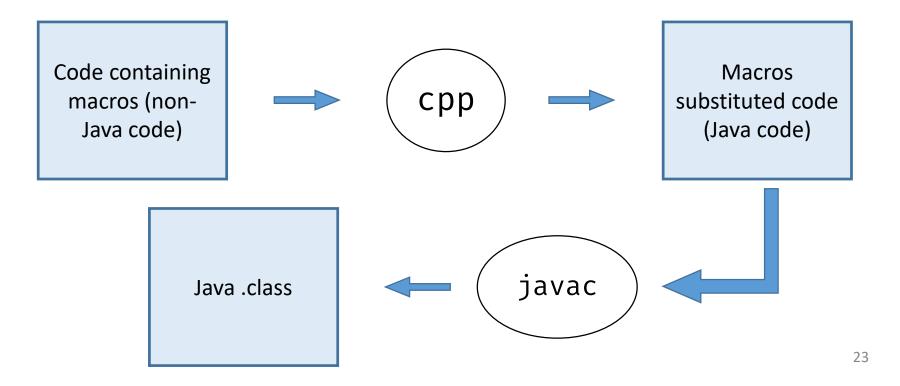
# C macros to effect syntactic sugar

Use a nullary lambda to freeze an expression:

```
freeze(3+5) \equiv (() -> (3+5))
```

```
T thaw(Thunk(T) ice) {
    return ice.get();
}
```

### Using C pre-processor to convert macros



### jpp script

```
PROC DIR=processed
TMPFILE=$(mktemp /tmp/jpp.XXXX)
FNAME=$1
cat <<EOF > $TMPFILE
#define freeze(x) (()->(x))
#define LLmake(a, b) new LazyList<>((a), freeze(b))
#define Thunk(T) Supplier<T>
EOF
cat $TMPFILE $FNAME | cpp -P - -o $PROC_DIR/$FNAME
javac -classpath $PROC_DIR -d $PROC_DIR $PROC_DIR/$FNAME
rm -f $TMPFILE
```

# Let's see an example

#### Before jpp

```
LazyList<Integer> integersFrom(int n) {
   return LLmake(n, integersFrom(n + 1));
}
```

#### After jpp

```
LazyList<Integer> integersFrom(int n) {
   return new LazyList<>((n), (()->(integersFrom(n +
1)) ));
}
```

### Example 2

#### Before jpp

#### After jpp

```
LazyList<Integer> sieve(LazyList<Integer> s) {
   return new LazyList<>((s.head()), (()->
   (sieve(s.tail().filter(x-> (x%s.head()) != 0 )))
  ));
}
```

#### Memoization

In <u>computing</u>, **memoization** is an <u>optimization</u> technique used primarily to speed up <u>computer programs</u> by storing the results of expensive <u>function calls</u> and returning the cached result when the same inputs occur again.

https://en.wikipedia.org/wiki/Memoization

#### Memo.java

```
public class Memo<T> {
    private boolean hasBeenRun;
    private T value;
    private Supplier<T> supplier;
    private Memo(Supplier<T> s) {
        this.hasBeenRun = false;
        this.value = null;
        this.supplier = s; }
    public static <T> Memo<T> make(Supplier<T> s) {
        return new Memo<>(Objects.requireNonNull(s));
```

#### Memo.java

```
public T get() {
    if (!hasBeenRun) {
        this.hasBeenRun = true;
        this.value = this.supplier.get();
    }
    return this.value;
}
```

First time: invoke supplier to evaluate the expression, and store the result in value

Subsequent times: just return the value

### Change the macros to use Memo

```
#define freeze(x) Memo.make(()->(x))
#define LLmake(a, b) new LazyList<>((a), freeze(b))
#define Thunk(T) Memo<T>
```

#### Demo

#### Let's time the difference

```
timeIt(()-> fib.get(10), "fib 10");
timeIt(()-> fib.get(10), "fib 10");
timeIt(()-> fib.get(20), "fib 20");
timeIt(()-> fib.get(15), "fib 15");
```

```
fib 10 took 24 ms
fib 10 took 0 ms
fib 20 took 1 ms
fib 15 took 0 ms
```

# Tracking #get

```
Memo Get Table:
head: 1, tail: thunk!Memo@546a03af: count-> 6
head: 1, tail: thunk!Memo@7b3300e5: count-> 6
head: 1597, tail: thunk!Memo@256216b3: count-> 3
head: 2, tail: thunk!Memo@721e0f4f: count-> 6
head: 21, tail: thunk!Memo@725bef66: count-> 6
head: 233, tail: thunk!Memo@649d209a: count-> 4
head: 55, tail: thunk!Memo@6e3c1e69: count-> 6 (fib 10)
head: 610, tail: thunk!Memo@357246de: count-> 4 (fib 15)
head: 6765, tail: thunk!Memo@16c0663d: count-> 1 (fib 20)
```

# java.util.stream.Stream<T>

https://docs.oracle.com/en/java/javase/11/docs/api/java.base/java/util/stream/Stream.html

### LazyList<T> vs Stream<T>: similarities

Both are lazily evaluated
Both can be finite or infinite
Both have these functions:
map, flatmap, filter, for Each, empty, limit, reduce

### LazyList<T> vs Stream<T>: differences

#### LazyList<T>

Multiple use

Can be memorized

Can retrieve element by index

Has Constructor, no converters

#### Stream<T>

Single use only
stream is used up when
terminal operation is
applied

No memorization
Cannot retrieve by index
No constructor, but
converters

```
List.stream(),
String.lines(),
String.chars(),
Scanner.tokens()
```

This computes: 
$$\frac{1}{20}(1^2 + 2^2 + \ldots + 20^2)$$

2 categories of stream operations:

<u>Intermediate</u>: This produces another lazily evaluated stream from the given stream, eg: map, flatmap, filter, limit

<u>Terminal</u>: This forces the eager evaluation of all the elements in the stream, and uses it up, eg: reduce, sum, collect, allMatch

## Summary

Lazy Evaluation means a computation is postponed until the value is needed.

Advantages: avoids unnecessary computation, amortizes time complexity, allows infinite data structures

Lazy Evaluation may be achieved using nullary lambda expressions (ie. Supplier<T>)

LazyList<T> and Stream<T> are similar in many ways, and yet are different in important ways.

#### A Tutorial on Java Stream:

http://tutorials.jenkov.com/java-functional-programming/streams.html