The good of this notorook is to produce an article out of the 11/30/12 REU worth.

The theoretical approach has gotten me nowhere, so I want to take a statistical approach instead. First we need to establish some notation.

An adjavency matrix A stands in as a digraph on [n], n ?. l. Entry (iii) of A is I if note in projects an edge to node in Oif not. There may be 100ps.
[k] = {1,..., k}, k EIN

[k] = {0, 1, ..., k}, keln. Xn = [1]o; this is the state space of the dynamical process

to be invertigated, # Define f

XL+1 = f(AxL+bL), Xo given # Include some graphics demonstrating a so XL EXNN for all LE INO. couple iterations.

be is another stochastic process, the input vectors the message vector, the efferent messages, something like that.

It is uniformly distributed over

BL = { b ∈ [1] o: || b || 1 = L}

where LE[n]o. Entries of bt are called newages.

Our desire is to estimate the distribution of survival times for messages injected at node i. for each node it [n]. Let's denote this variable Ti. Before doing this, I want to know if T. Marker chain. Actually the most important thing is to just get an empirical distribution. If we denote the theoretical part of T; by p; then we are estimating point his we need: a message object that:

· A way to identify message starting nodes.
· A way to compute the age of a message.
· Knowledge of when the message goes extinct.

Now. instead of having a fixed time horizon, we'll let the simulation keep runing vatil we've collected "enough data" to be confident that, for all it[n] and for all TEIN, |p:(T)-p:(T)|< E. This is a fair theoretical grestion to ask, and to answer it, we need to do some research to see . What mode of convergence this is · How many messages injected at a given node have to
go extint to get good consequence; i.e., what published
younds exist? Once this is implemented, we should compute this empirical distribution for an ensemble of graphs A sampled from the graphon eA = ER(1/2). We have to play with some math now. We to know that for a fixed TEIN and annumber of messages let M be the total number of messages injected and Mr be the proportion number of those newsges... #{messages injected at i which died it age of N-200 p; (I) I'm not artain the is correct because what about nessages that never die? Here's am idea: if there are mellages that never die then we stort by putting all of that mass on the maximum age +1. Then, we portion this mass out over morage 1. Then, we portion this mass out over morage 1. It less than the mass over marage 1 is less than the mass over max age. I think a geometric distribution over \{max age \forall 1..., \ood \} we viold work.

Take adjacency matrix A,

$$X \xrightarrow{A \mapsto b_t} Y = [n+1]^n \xrightarrow{f(\cdot)} Z = X$$

The goal is to find a formula for the statutary probabilities. It best to matter if the formula is terrible.

Fix z \ Z, and let [z] = f - 1/2 \]

let pre be transition propositions.

PXZ

Compute Ax. For all valves of b&BL, compute b

[2-(ax-by) (ax+by) VO = (2(ax+by)-(ax+by)2) VD

$$(2-x)x \vee 0 = (2x-x^2) \vee 0$$

 $2(2x-x^2)v0 - [(2x-x^2)v0]^2$

- elementer that the <u>neal good</u> is to compute calculate the distribution of message survival times for a message injected at node i, at time 0, under load in network A. L, in network A.
 - Note that this does not care about the initial state, x_{-1} , which is for the best because a message sender is unlikely to know that.
- The first thing I should do is take a statistical approach. So, when you have more free time:

 - Design a simulation to collect data about survival times

 Find a simple probability distribution to model the distributions

 you see
 - Jeatures from the nodewise and graphwise graph wide features and attempt to predict the parameters of that distribution.

The goal here is to produce an unbiased estimator $\hat{p}_{\pm}(\tau)$ for the distribution of extinction dimes, $p_{\pm}(\tau)$.

Loose with notation for now, we say we rejected N messages at node I.

The time at which this happens is random, but at least N, since only 1 message can be injected at node I at a time. Let's call this time $T^*(N)$. So: $T^*(N)$ is the time at which N nodes have been injected at node I.

- Then there's a time horizon T. We need $T \ge T^*(N)$, but we still have to make a choice. Essentially, T is the point at which we stop collecting Jata about the extinction times.
- Now. If we inject by time T, it's possible that more than N menages have been injected at node 1; call this number N. So, the set of all menages which were injected at node 1 time T be indexed as $\{m_0\}_{n=1}^{N}=:M$.

be indexed as $\{m_{\alpha}\}_{\alpha=1}^{N}=:M$.

"M can be partitioned into M° and M°, where

M° is the set of messages that were extinct by time Trand

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For the elements of M°, the extinction time is known, but for the elements of M°, we only have a lower bound on the extinction

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time.

graph preparties.