

Readings...

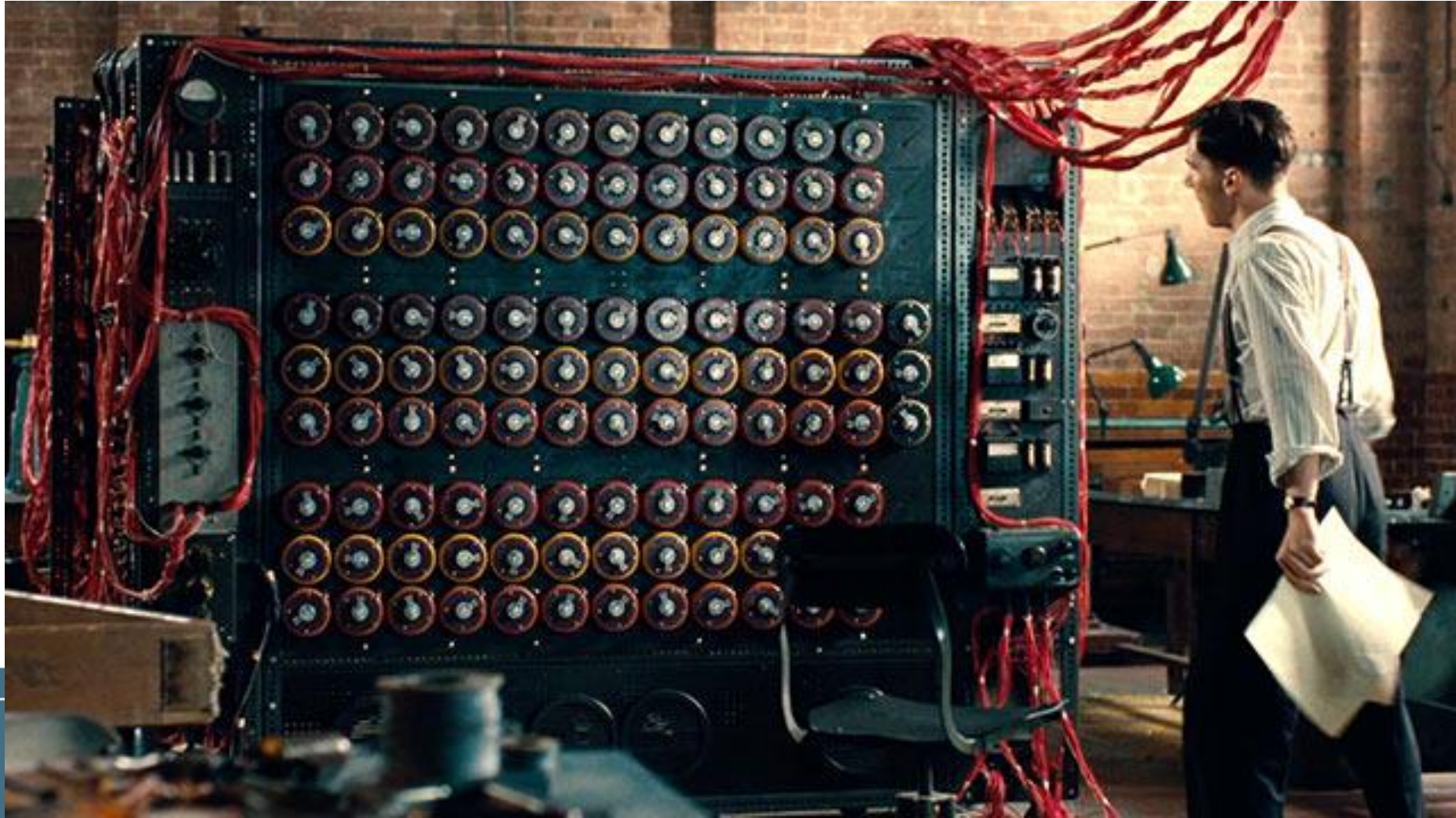
- Sections 3.1 – 3.3 from Sipser
- Wikipedia article:

https://en.wikipedia.org/wiki/Turing_machine

- Program your own Turing Machine!

<https://turingmachinesimulator.com/>

What is a Turing Machine?



A Bit of History

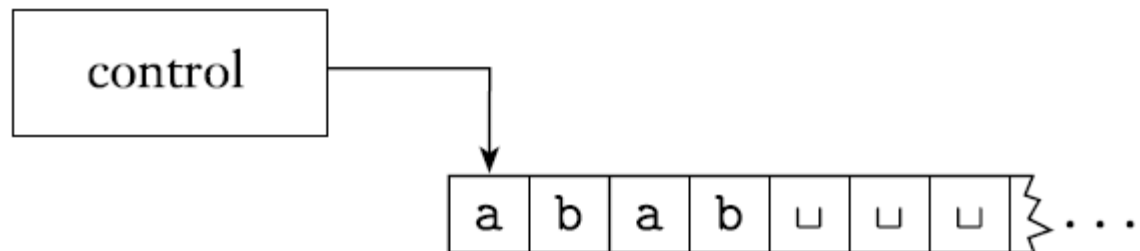
- 1928 – David Hilbert’s Entscheidungsproblem (*Decision Problem*):
“Is there an algorithm that can prove whether a statement in first-order logic is true or false?”
- Technical Problem: What exactly do we mean by “algorithm” here?

Church's Characterization

- In the 1930s, Alonzo Church characterized “effective calculability” as being anything that could be computed using his λ -calculus
- The functional language LISP was influenced by the λ -calculus
- Very “mathy”, and doesn't immediately resemble our sequential notions of algorithm

Turing's Characterization

- Alan Turing saw algorithms as a “mechanical process”, and so used a machine-based characterization
 - These became known as Turing Machines
- Turing Machines can be in one of a finite number of states, and are deterministic, very much like a DFA
- However, TMs also have an infinite tape for memory on which it can read and write characters



Church-Turing Thesis

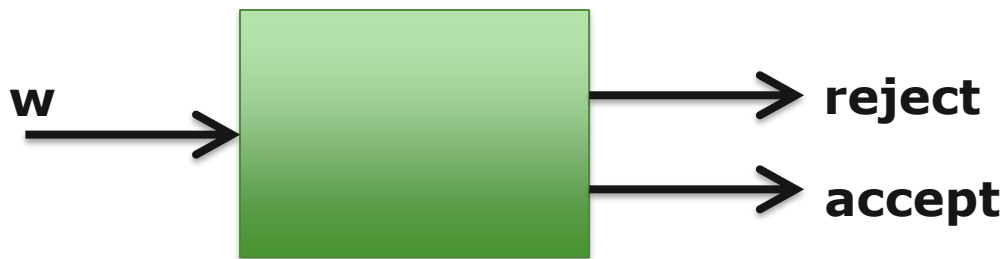
- **Church-Turing Thesis:** Both of these notions are equivalent in terms of the ability to compute things, and can be used to describe any algorithm!

Ramifications of Church-Turing

- Any computer program you can write in your favorite language (e.g., C/C++, Java, Lisp) can be converted into an equivalent Turing Machine!
- Thus, if someone asks you “Is this function computable?” or “Can a Turing Machine do this?” you can simply ask yourself if you can write a program in any other language to do it
- For this reason, we will not spend time on how to write a formal Turing Machine definition this semester

Turing Machines as Language Recognizers

- Turing Machines are meant to act as language recognizers like DFAs – they attempt to either accept or reject a given input string, though they might loop forever!



Recognizers vs. Deciders

- We say that a Turing Machine M **recognizes** a language L , if M accepts exactly those strings in L
- We say that M **decides** a language L , if M accepts all strings in L , and rejects all other strings
- In other words, deciders always give an accept or reject answer, and will never loop forever

Recognizable vs. Decidable

- We say that a language L is **recognizable** if there exists some Turing Machine that recognizes it
- We say that L is **decidable** if there exists some Turing Machine that decides it
- Most languages that you can generally think of are decidable
 - All regular languages are decidable, as are many non-regular languages

Transducer Turing Machines

- A **Transducer Turing Machine** is a Turing Machine that is used to compute functions, instead of recognizing or deciding languages
- TTM's operate in the same way as standard TMs, except that they produce an output string for a given input string, instead of accepting or rejecting it



Transducer Turing Machines

- TTM's must always terminate with an answer, they are not allowed to loop forever
- We say that a function f is **computable** if there is a TTM that outputs the correct value for $f(w)$ for every possible input w
- All arithmetic operations on integers are computable, as are most functions that we commonly use (given that the input/output can be represented using finite strings)