

# Languages (Not Programming Languages)

- An **alphabet** is a finite set of characters, usually denoted  $\Sigma$
- A **string** or **word** is a finite sequence of characters
- The **empty string** ( $\epsilon$ ) has no characters
- A **language** is a set, which could be infinite, of strings over a given alphabet

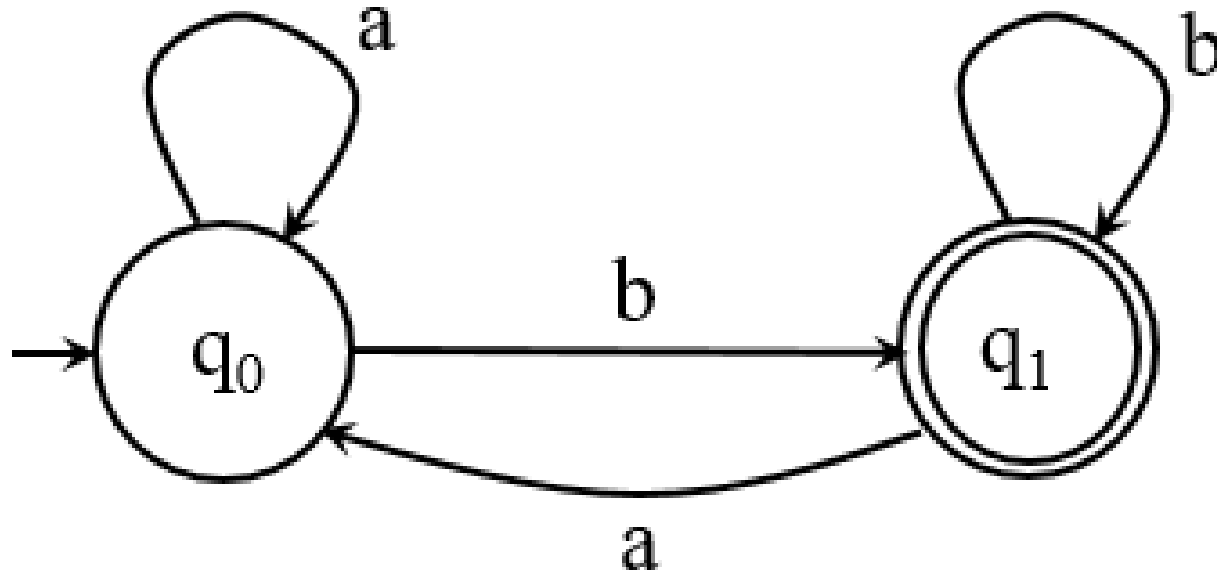
# Deterministic Finite Automata (DFA)

- A DFA is a model of computation that can be used to define a language
- Formal parts of a DFA:
  - States:  $Q = \{q_0, q_1, q_2, \dots, q_k\}$
  - Alphabet:  $\Sigma$
  - Transition function:  $\delta : Q \times \Sigma \rightarrow Q$
  - Start state:  $q_0 \in Q$
  - Accept/final states:  $F \subseteq Q$

# DFA Characterizations

- A DFA could be thought of as:
  - An **oracle** that answers the question “Does a given string  $w$  belong to the language  $L$ ?”
  - A **language generator**, which can be used to create all of the words in a language  $L$  by doing a breadth-first traversal over transition sequences

# A Simple Example

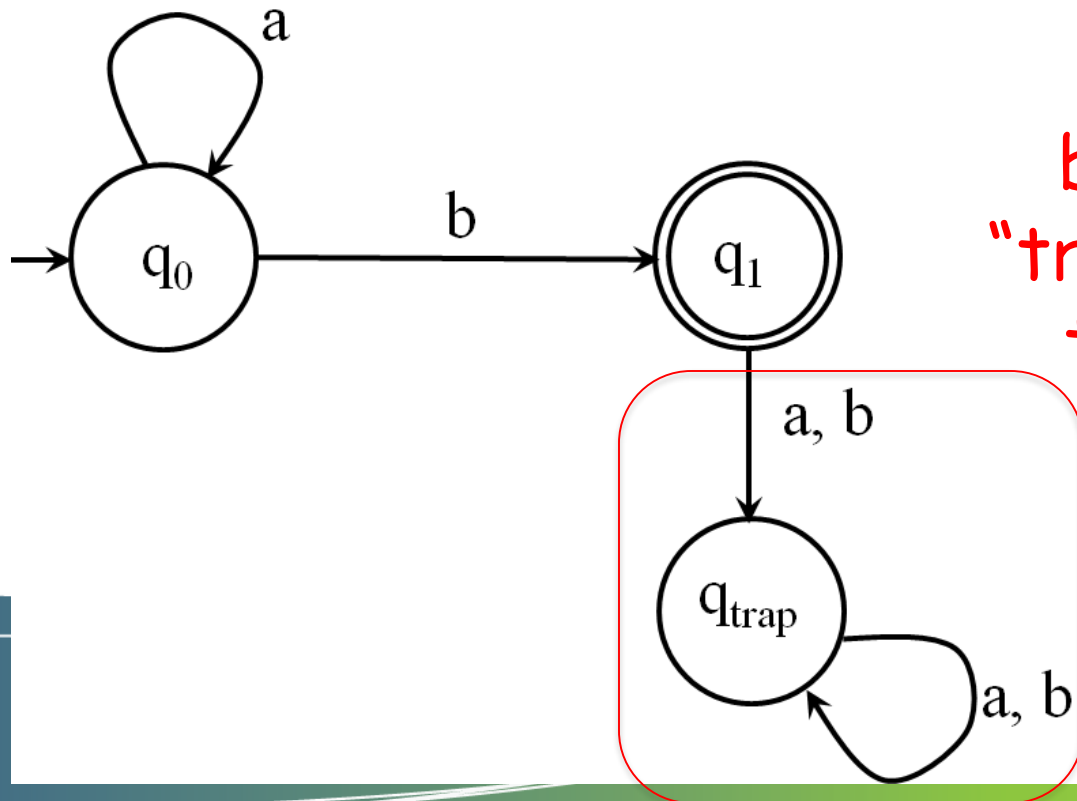


- This DFA accepts words such as **ab**, **aab**, and **aaabbb**, but also **ababbaab** and **b**



# Another Example

- Create a DFA that only accepts those words that begin with zero or more **a**'s, and ends with a single **b**



People usually don't bother to include this "trap state" part, but it technically makes the DFA complete by showing all possible transitions

# More Terminology

- If a DFA  $M$  ends in an accepting state after processing a given word  $w$ , we say that  $M$  accepts  $w$
- Otherwise, we say that  $M$  rejects  $w$
- If  $M$  accepts all words in a language  $L$ , and rejects all others, then we say that  $M$  accepts  $L$  or  $M$  recognizes  $L$