

Chapter 7: Quicksort

Quicksort is an algorithm that solves the sorting problem. It uses the divide-and-conquer approach. Although the quicksort has a worst-case running time of $\Theta(n^2)$, it is commonly used in practice as it has $\Theta(n \lg n)$ running time on average. Like Insertion sort and Heapsort it is an in-place algorithm. Not a *stable* algorithm.

QUICKSORT(A, p, r)

```

1  if  $p < r$ 
2       $q = \text{PARTITION}(A, p, r)$ 
3      QUICKSORT( $A, p, q - 1$ )
4      QUICKSORT( $A, q + 1, r$ )

```

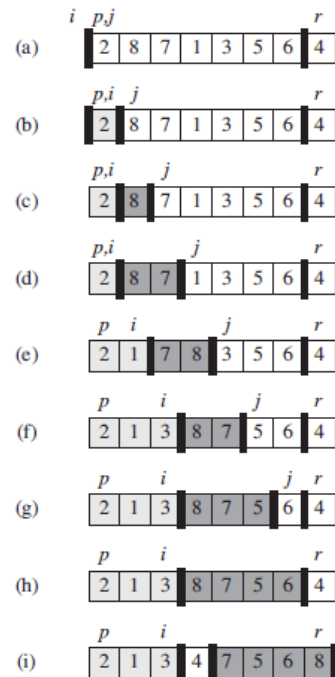
PARTITION(A, p, r)

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1   $x = A[r]$ 
2   $i = p - 1$ 
3  for  $j = p$  to  $r - 1$ 
4      if  $A[j] \leq x$ 
5           $i = i + 1$ 
6      exchange  $A[i]$  with  $A[j]$ 
7  exchange  $A[i + 1]$  with  $A[r]$ 
8  return  $i + 1$ 

```

Example:



Performance of quicksort

The running time of quicksort depends on how the input is portioned: if the portioning is balance or not. For example, if it is well balanced (equally) then $T(n) = 2T(n/2) + \Theta(n)$.

Worst-case partitioning

The worst case occurs when the partition algorithms divides the input into two of which one has $n-1$ elements, and the other has 0 elements. $T(n) = T(n-1) + T(0) + \Theta(n)$.

$$T(n) = T(n-1) + cn =$$

$$= cn + c(n-1) + T(n-2) =$$

$$= cn + c(n-1) + c(n-2) + T(n-3) =$$

...

$$= cn + c(n-1) + c(n-2) + \dots + c \cdot 1 =$$

$$= c \cdot \sum_{i=1}^n i = c \cdot n \cdot (n+1)/2 = \Theta(n^2)$$

Let's verify this by the substitution method:

$$T(n) = T(n-1) + cn$$

$$T(n) \leq dn^2$$

$$T(n) \leq d(n-1)^2 + cn = d(n^2 - 2n + 1) + cn = dn^2 - 2dn + d + cn =$$

$$= dn^2 - (2dn - cn - d). \text{ We need } 2dn - cn - d \geq 0 \text{ so that we can say } T(n) \leq dn^2.$$

$$2dn - cn - d \geq 0 \rightarrow n(2d - c) - d \geq 0$$

$$\text{For } n=1, 2d - c \geq d \rightarrow d \geq c. \text{ So, we found constants } d \geq c > 0.$$

Best-case partitioning

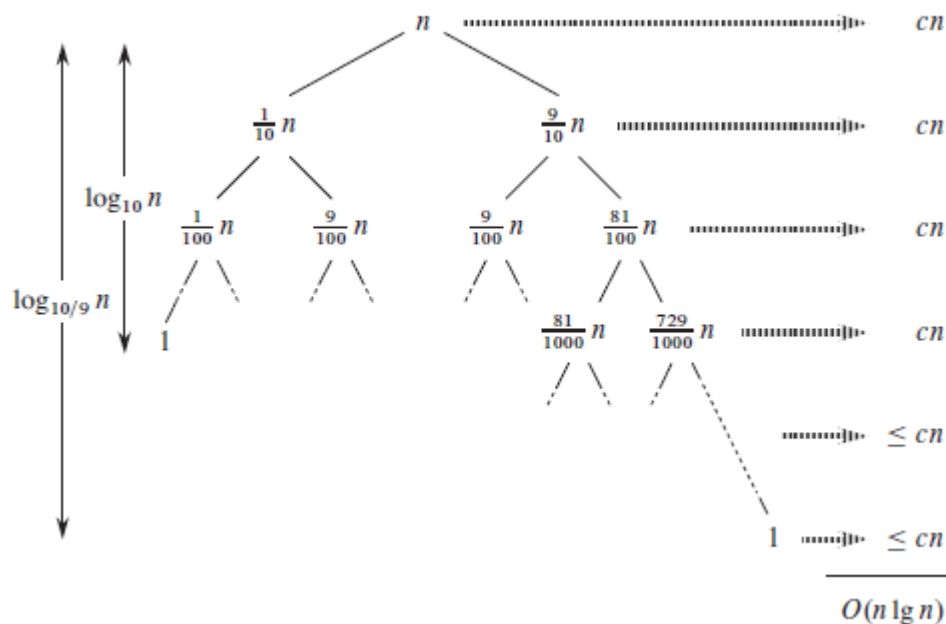
The best case occurs when the partitioning is well balanced, i.e., $T(n) = 2T(n/2) + \Theta(n)$.

By the master method (case 2), $T(n) = \Theta(n \lg n)$.

Average case

Closer to the best case = $O(n \lg n)$

$$T(n) = T(9n/10) + T(n/10) + \Theta(n).$$



A randomized version of quicksort

RANDOMIZED-PARTITION(A, p, r)

- 1 $i = \text{RANDOM}(p, r)$
- 2 exchange $A[r]$ with $A[i]$
- 3 return **PARTITION**(A, p, r)

RANDOMIZED-QUICKSORT(A, p, r)

- 1 if $p < r$
- 2 $q = \text{RANDOMIZED-PARTITION}(A, p, r)$
- 3 **RANDOMIZED-QUICKSORT**($A, p, q - 1$)
- 4 **RANDOMIZED-QUICKSORT**($A, q + 1, r$)