INTRODUCTION

Preference-inclusion comparison under partial information

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OUTLINE

Introduction

Model

PRELIMINARIES

MAIN RESULTS

Conclusions

Reforms¹

► Seattle Public Schools in 1999

Model.

- ▶ Boston Public Schools in 2005
- ► Ghanaian Secondary Public Schools in 2007
- ► Chicago Selective High Schools in 2009 and 2010
- ► Primary Public Schools in more than 50 cities and provinces in England and Wales in 2005-2011
- ► Denver Public Schools in 2012

change of mechanism:

- ► change of algorithm Deferred Acceptance $(DA) \Rightarrow Boston (BM)$
- ► change of constraint $k \uparrow$ in DA^k or BM^k

Problem - moving from one manipulable mechanism to another manipulable mechanism or stable mechanism?

¹Abdulkadiroğlu and Sönmez (2003); Abdulkadiroğlu et al. (2005); Pathak and Sönmez (2008); Pathak and Sönmez (2013)

LITERATURE

INTRODUCTION

(parallel results)

- superiority of *DA* to *BM*
- superiority of M^{k+1} to M^k

comparison under complete information

- ▶ less manipulable (Pathak and Sönmez, 2013)
- ▶ more truthful (Decerf and Van der Linden, 2016)
- ▶ more immune (Bonkoungou and Nesterov, 2020)

comparison under partial information

- ▶ obviously manipulable (Bonkoungou and Nesterov, 2019)
- BM is less obviously manipulable with more information
- DA is more obviously manipulable with more information
 - ▶ obviously manipulable (Troyan and Morrill, 2020)
- BM is the only obviously manipulable mechanism

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- ► In reality, agents have **partial information**: something is known for sure (or fixed), something is unknown
- ▶ When agents begin to report their *truthful* preferences more often:
- under which mechanism?
- under which information?
 - Criterion of truthfulness obvious dominance (Li, 2017)

INTRODUCTION

- ▶ set of students $S := \{s_1, s_2, ..., s_n\}$
- ▶ preference profile $R_S := (R_{s_1}, R_{s_2}, ..., R_{s_n})$
- ▶ set of schools $C := \{c_1, c_2, ..., c_m\}$
- $\blacktriangleright \text{ priority profile } R_C := (R_{c_1}, R_{c_2}, ..., R_{c_m})$
- capacity vector $q = (q_1, q_2, ..., q_m)$
- ▶ information structure I = 0, 1, ..., m

If each student knows his own preferences R_s and priority profile R_C but only the top $I \ge 0$ rows of the preference profile R_{-s} , then the information is denoted as I (Bonkoungou and Nesterov, 2019)

 R_{c_2}

 S_2

 S_1

 S_3

 R_{c_3}

 S_3

s₁ s₂

COMPLETE VS PARTIAL INFORMATION

3 students, 3 schools

INTRODUCTION

ightharpoonup complete information I = 3

$\overline{\mathbf{R}_{s_1}}$	R_{s_2}	R_{s_3}	$\overline{R_{c_1}}$
c_1	c ₂	C 3	s_1
c ₂	c_1	c ₂	s_2
C 3	C 3	c_1	s_3

ightharpoonup partial information I = 1

R_{s_1}	R_{s_2}	R_{s_3}	$\overline{R_{c_1}}$	R_{c_2}	R_{c_3}
c ₁	c ₂	C 3	$\frac{s_{t_1}}{s_1}$	S ₂	$\frac{14_3}{s_3}$
÷	:	:	s_2	s_1	s_1
:	:	:	s_3	s_3	s_2

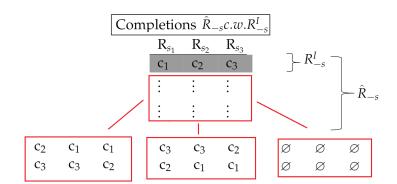
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(1)

PARTIAL INFORMATION STRUCTURE

ightharpoonup examples fixing I=1

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Abdulkadiroğlu et al. (2005)

Round 1:

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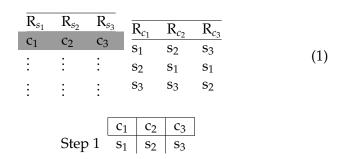
- 1) Each *student* applies to the school he reported as his most preferred acceptable school (if any)
- 2) Every *school* rejects the students in excess of its capacity according to its priority
- 3) Each *student* who is not rejected is **assigned** to the *school* he applied to and *capacities are adjusted* accordingly

:

Round 1:

- 1) Each rejected *student* applies to the best acceptable school that did not reject him yet
- 2) and 3) as in Round 1

BOSTON (IMMEDIATE ACCEPTANCE) MECHANISM



assignment is uniquely determined given I = 1

▶ no competition

INTRODUCTION

- ▶ assigned to the most preferred school
- ⇒ no improvement by misreport is possible

DEFERRED ACCEPTANCE MECHANISM

GALE AND SHAPLEY (1962)

Round 1:

INTRODUCTION

- 1) Each *student* applies to the school he reported as his most preferred acceptable school (if any)
- 2) Every *school* rejects the students in excess of its capacity according to its priority
- 3) Each *student* who is not rejected is **TEMPORARILY assigned** to the *school* he applied (could be rejected on any further step)

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Round 1:

- 1) Each rejected *student* applies to the best acceptable school that did not reject him yet
- 2) and 3) as in Round 1

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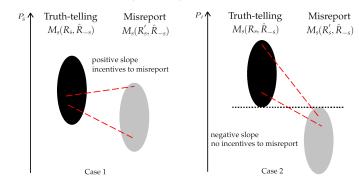
DEFERRED ACCEPTANCE MECHANISM

assignment is uniquely determined given I = 1

- ► assigned to both top priority and most preferred school
- ⇒ no improvement by misreport is possible

OBVIOUS TRUTHFULNESS

▶ obvious dominance (Li, 2017)

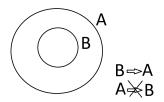


Given mechanism M, preference R_i is **obviously truthful** under information structure *I* for student t_i if for any R'_i :

$$\min_{\hat{R}_{-i}c.w.R_{-i}^{I}} M_{i}(R_{i},\hat{R}_{-i}) \; R_{i} \max_{\hat{R}_{-i}c.w.R_{-i}^{I}} M_{i}(R_{i}^{'},\hat{R}_{-i})$$

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▶ preference-inclusion (Arribillaga and Massó, 2015)



Mechanism M_A is **more obviously truthful** than mechanism M_B under information structure I for student s if:

- (i) R_s is o. t. under I for s in $M_B \Rightarrow R_s$ is o. t. under I for s in M_A
- (ii) R_s is o. t. under I for s in $\mathbf{M_A} \neq R_s$ is o. t. under I for s in $\mathbf{M_B}$

Pre-assignment under partial information

$\overline{\mathbf{R}_{s_1}}$	R_{s_2}	R_{s_3}	R_{s_4}	$\overline{\mathbf{R}_{c_1}}$	R_{c_2}	R_{c_3}	R_{c_4}	
c ₁	c_2	c ₃	c_1	$\overline{s_1}$	s_1	s_1	s_1	
:	:	:	c_2	:	s_2	s_2	s_2	(2
:	:	:	c ₃	:	•	s_3	s_3	
:	:	:	C4	:	:	:	S ₄	

 \triangleright $s_1 - c_1$

INTRODUCTION

- \triangleright $s_2 c_2$
- $ightharpoonup s_3 c_3$
- \triangleright $s_4 c_4$

A student s is **pre-assigned** to a school c in mechanism M if c is the first school to which s is (temporarily) assigned without the threat of being rejected on any step given information.

DEFERRED ACCEPTANCE MECHANISM

Proposition 1

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For any $k \in \{1,...,m\}$, truth-telling is an obviously dominant strategy in DA^k for student s if and only if under information structure $I \in \{1,...,k\}$ in DA^k :

- (i) *s* is pre-assigned to an *acceptable* school *c* under truth-telling or
- (ii) there exists no $\hat{C} \subseteq C$ of acceptable and guaranteed for s schools

Suppose a student

- ► not pre-assigned ⇒ min(true report)=unassigned
- ► set of guaranteed schools ⇒ max(misreport)=assigned
- ⇒ improvement by misreport

BOSTON MECHANISM

Proposition 2

For any $k \in \{1, ..., m\}$, truth-telling is an obviously dominant **strategy in BM**^k for student s if and only if under information structure $I \in \{1, ..., k\}$ in BM^k :

- (i) s is pre-assigned to *most preferred feasible* school c under truth-telling or
- (ii) there exists no $\hat{C} \subseteq C$ of acceptable and guaranteed for sschools
 - ► top-ranking misreport
 - ► *feasibility* in terms of top-ranking strategies

DEFERRED ACCEPTANCE VS BOSTON

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Proposition 1 and 2: student has an obviously truthful strategy in *BM*^k \implies student has an obviously truthful strategy in DA^k

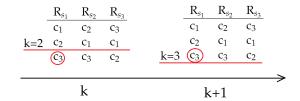
Theorem 1 For any $k \in \{1, ..., m\}$ and fixed information structure $I \in \{1, ..., k\}$, DA^k is more obviously truthful than BM^k under I

COMPARING MECHANISMS UNDER FIXED I

DIFFERENT CONSTRAINTS

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constraint $\uparrow \implies$ risk of running out of schools \downarrow



Theorems 2 and 3

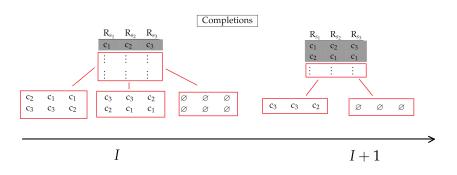
For any $k \in \{1, ..., m\}$ and fixed information structure $I \in \{1, ..., k\},$

- ▶ DA^{k+1} is more obviously truthful than DA^k under I
- ► BM^{k+1} is more obviously truthful than BM^k under *I*

COMPARING MECHANISMS UNDER DIFFERENT I

DEFERRED ACCEPTANCE AND BOSTON

INTRODUCTION



Theorems 4 and 5

For any $k \in \{1, ..., m\}$ and information structure $I \in \{1, ..., k\}$

- ▶ DA^k is more obviously truthful under I + 1 than under I
- ▶ BM^k is more obviously truthful under I + 1 than under I

CONSISTENCY WITH EXISTING LITERATURE

COMPARISON UNDER FIXED INFORMATION

INTRODUCTION

- 1. DA^k is more obviously truthful than BM^k
- 2. DA^{k+1} is more obviously truthful than DA^k
- 3. BM^{k+1} is more obviously truthful than BM^k
 - reinforced the main conclusion regarding the Boston and constrained DA comparison
 - ► reinforced conclusions under *fixed* partial information

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Consistency with existing literature

COMPARISON UNDER DIFFERENT INFORMATION

- 4. BM^k is more obviously truthful under I + 1 than under I
- 5. DA^k is more obviously truthful under I + 1 than under I
 - $ightharpoonup DA^k$ is more obviously manipulable (Bonkoungou and Nesterov, 2019) and more obviously truthful under I + 1than under I
 - new results under *not fixed* partial information

- ► Deferred Acceptance mechanism incentivize students to be more truthful compared to Boston mechanism
- longer constraint incentivize students to be more truthful under both mechanisms
- ► more **information** announced/available incentivize student to be more truthful under Boston mechanism
- more information announced/available provides no clear incentive to students under Deferred Acceptance mechanism (accordingly to current research)

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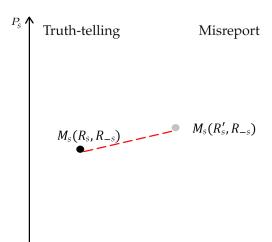
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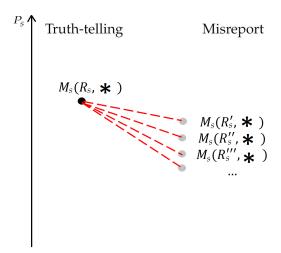
PATHAK AND SÖNMEZ (2013)

MANIPULABLE



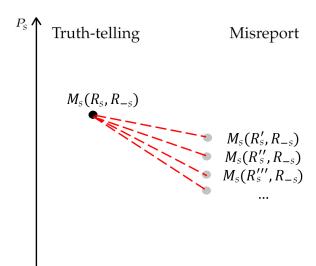
DECERF AND VAN DER LINDEN (2016)

TRUTHFUL



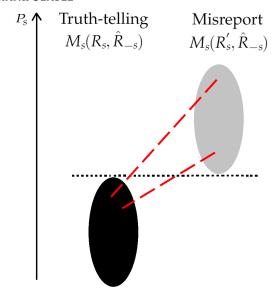
BONKOUNGOU AND NESTEROV (2020)

IMMUNE



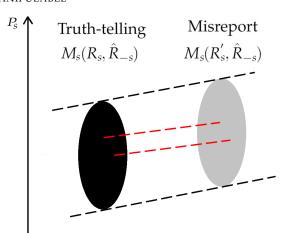
BONKOUNGOU AND NESTEROV (2019)

OBVIOUSLY MANIPULABLE



Troyan and Morrill (2020)

OBVIOUSLY MANIPULABLE



DA^k vs BM^{k+1}

<u>Case 1</u>. If BM^{k+1} is obviously truthful under given I it does not imply that DA^k is obviously truthful under I.

- ▶ Suppose that student s is pre-assigned to his k + 1'th preferred school in BM^{k+1} under $I \Longrightarrow BM^{k+1}$ is obviously truthful for s under I.
- ▶ It can be the case that student s is rejected from every school under truth-telling in DA^k under I but he cannot be rejected from his k + 1'th preferred school which provides an incentive to misreport $\implies DA^k$ is not obviously truthful for s under I.

DA^k vs BM^{k+1}

Case 2. If DA^k is obviously truthful under given I it does not imply that BM^{k+1} is obviously truthful under I.

- ▶ Suppose that student *s* is pre-assigned to some school *c* in DA^k under $I \Longrightarrow DA^k$ is obviously truthful for s under I.
- ▶ It can be the case that student *s* is rejected from all schools in BM^{k+1} under I because he applies too late but he has a strong incentive to misreport c as his most preferred school $\implies BM^{k+1}$ is not obviously truthful for s under I.

SAFE SET

$\overline{\mathbf{R}_{s_1}}$	R_{s_2}	R_{s_3}	R_{s_4}	$\overline{\mathbf{R}_{c_1}}$	R_{c_2}	R_{c_3}	R_{c_4}		
c_1	c_2	c ₃	c_1	s_1	s_1	s_1	s_1		
÷	:	:	c_2	:	s_2	s_2	s_2	((2)
:	:	:	c_3	÷	:	s_3	s_3		
:	÷	:	C 4	:	:	:	S ₄		

$ightharpoonup DA^k$ or BM^k

	c ₁	c ₂	c ₃	c ₄
Step 1	s_1, s_4	s_2	S 3	
Step 2	$\overline{s_1}$	s_2, s_4	S 3	
Step 3	$\overline{s_1}$	s ₂	S ₃ , S ₄	
Step 4	s_1	s ₂	s_3	S ₄

SAFE SET

$\overline{\mathrm{R}_{s_1}}$	R_{s_2}	R_{s_3}	R_{s_4}	$\overline{\mathbf{R}_{c_1}}$	R_{c_2}	R_{c_3}	R_{c_4}	
c_1	c_2	c ₃	c_1	$\overline{s_1}$	s_1	s_1	s_1	
:	:	:	c_2	:	s_2	s_2	s_2	(2)
÷	:	:	c ₃	:	:	s_3	s_3	
:	:	:	C4	:	:	:	Sa	

- $ightharpoonup s_1 c_1$
- $ightharpoonup s_2 c_2$
- \triangleright $s_3 c_3$
- \triangleright s_4 c_4

A set \hat{C} forms **safe set** for student s in mechanism M if $\hat{C} \subseteq C$ protects s from being unassigned under information structure I when report of s includes \hat{C} (Decerf and Van der Linden, 2016)

Deferred Acceptance Mechanism DA^3

R_{s_1}	R_{s_2}	R_{s_3}	R_{s_4}	R_{c_1}	R_{c_2}	R_{c_3}	R_{c_4}
c ₁	c_1	c_1	c ₄	s_1	s_1	s_1	s_4
:	:	c ₂	÷	:	s_2	s_3	:
:	:	c_3	:	:	s_3	÷	:
:	:	C ₄	:	:	:	:	:

Completion 1								
R_{s_1}	R_{s_2}	R_{s_3}	R_{s_4}					
c_1	c_1	c_1	C ₄					
:	c_2	c_2	:					
:	:	C3	:					

	c_1	c ₂	c ₃	C4
Step 1	s_1, s_2, s_3			S ₄
Step 2	s ₁	s_2, s_3		S ₄
Step 3	s ₁	S ₂	S 3	S4

DEFERRED ACCEPTANCE MECHANISM

R_{s_1}				R_{c_1}	R_{c_2}	R_{c_3}	R_{c_4}
c_1	c_1	c_1	C ₄	s_1	s_1	s_1	s_4
:	:	c_2	÷ :	÷	s_2	s_3	:
:	:	c ₃	:	:	s_3	:	:
÷	÷	c_4	÷	÷	÷	÷	:

Completion 2								
R_{s_1}	R_{s_2}	R'_{s_3}	R_{s_4}					
c_1	c_1	c ₂	c_4					
:	c_3	:	:					
:	:	:	:					

	c ₁	c ₂	C 3	C4
Step 1	s_1, s_2	S 3		S ₄
Step 2	s_1	s_3	s_2	S ₄

DEFERRED ACCEPTANCE MECHANISM

R_{s_1}	R_{s_2}	R_{s_3}	R_{s_4}	R_{c_1}	R_{c_2}	R_{c_3}	R_{c_4}
c_1	c_1	c ₁	C4	s_1	s_1	s_1	S ₄
:	:	c_2	÷	:	s_2	s_3	:
:	:	c ₃	÷	:	s_3	:	:
:	÷	C ₄	:	:	:	:	÷

- ► Completion $1 s_3$ to c_3
- ► Completion 2 s_3 to c_2
- \implies trut-telling is not obviously truthful for s_3

BOSTON MECHANISM as in Completion 1

R_{s_1}	R_{s_2}	R_{s_3}	R_{s_4}	R_{c_1}	R_{c_2}	R_{c_3}	R_{c_4}
c_1	c_1	c_1	c ₄	s_1	s_1	s_1	s_4
÷	c ₂	c ₂	÷	:	s_2	s_3	:
÷	:	c ₃	:	÷	s_3	÷	÷
:	:	C ₄	:	:	:	:	:

R_{s_1}	R_{s_2}	R_{s_3}	R_{s_4}
c_1	c_1	c ₂	C ₄
:	c ₂	:	:
:	:	:	:

	c ₁	C2	C 3	c ₄
Step 1	$\overline{s_1, s_2}$	S 3		S ₄
Step 2	$\overline{s_1}$	s ₃ , s ₂		