

Lab3  
The Power operator

Ai Nguyen

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# 1 power math

The power formula is given by this specification:

$$power : (\mathbb{R} \times \mathbb{N}) \rightarrow \mathbb{R}$$

$$power(r, p) = r^p = \prod_{i=1}^p r$$

This can be implemented by multiplying  $r$  times itself  $p$  times. One algorithm for calculating the power more efficiently than that can be visualized like this:

$$r^p = \overbrace{\underbrace{r \times r \cdots r}_{p \text{ div } 2} \times \underbrace{r \times r \cdots r}_{p \text{ div } 2} (\times r)}^p$$

Here the  $p$  multiplications are split into two groups of  $p \text{ div } 2$  multiplications, and, if  $p$  is odd, one more multiplication. This is a recursive definition where  $r : \mathbb{R}$  and  $p : \mathbb{N}$ :

$$\begin{cases} 1, & \text{if } p = 0 \\ r \times r^{p-1}, & \text{if } p \text{ is odd} \\ \text{sqr}(r^{p \text{ div } 2}), & \text{otherwise} \end{cases}$$

$$\text{where } \text{sqr}(x) = x \times x$$

## 2 power sml

SML

```
fun square x: real = x * x;  
square 2.0;  
fun odd x = x mod 2 = 1;  
odd 3;  
odd 4;  
fun power r p = if p = 0 then 1.0  
                 else if odd p then r * power r (p-1)  
                 else square (power r (p div 2));  
power 5.0 3; (* 125 *)  
power 2.0 8; (* 256 *)
```

```
debian@debian:~/labs/lab3$ poly < lab3.sml  
Poly/ML 5.5.2 Release  
val square = fn: real -> real  
val it = 4.0: real  
val odd = fn: int -> bool  
val it = true: bool  
val it = false: bool  
val power = fn: real -> int -> real  
val it = 125.0: real  
val it = 256.0: real
```

### 3 power c without using function composition

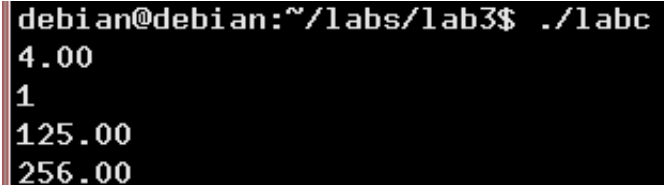
It will help us transition to ASM if we rewrite the C code using only one calculation per line instead of combining them into complex expressions. In fact even the returns are done differently—they are all expected to be in EAX(for integers) and ST(0) for floats.

C source code written to file lab3.c

```
#include <stdio.h>
#include <stdbool.h>
float ST0;
int EAX;
float square(float x) {ST0 = x * x;}
int odd(int x) { EAX = x % 2;}
// return value 0 or 1 in EAX
void power(float r, int p)
{
    if(p!=0) goto first_else;
    ST0 = 1.0;
    goto end_if;
first_else:
    odd(p);
    if(EAX != 1) goto second_else;
    int t = p-1;
    power(r,t);
    ST0 = r * ST0; // ST0 *=r
    goto end_if;
```

C++ source code appended to file lab3.c

```
second_else:
    t = p/2;
    power(r,t);
    square(ST0);
end_if:
    return;
} //result ST0
int main()
{
    power(5.0,3);
    printf("%.2f\n", ST0);
    power(2.0,8);
    printf("%.2f\n", ST0);
}
```



```
debian@debian:~/labs/lab3$ ./lab3
4.00
1
125.00
256.00
```

## 4 power asm

### 4.1 square function

asm source code written to file lab3.s

```
.text
square:
    push %ebp
    mov %esp,%ebp
    fmul $8(%ebp)           # multiply st(0) with the parameter
                             # st(0) should be equal to 8(%ebp)
    mov %ebp,%esp
    pop %ebp
    ret
```

Here are using 4 offset from esp because we didn't put ebp on the stack, so the parameter is only 4 bytes above esp.

### 4.2 odd function

asm source code appended to file lab3.s

```
odd:
    mov 4(%esp), %eax
    and $1, %eax           # compare 1 and eax by bit-wise arithmetic
                           # return 1 to %eax if the number is odd
                           # or return 0 to %eax if the number is even
    ret
```

### 4.3 power function

The power function will return the result in  $st(0)$ , the first parameter is  $r$ , a real number, and second is  $p$ , an integer to calculate  $r^p$

First, compare  $p$  stored at  $8(\%ebp)$  with 0, if  $p$  equals to 0, push 1 onto floating point stack,  $st(0) = 1$  and then return. If  $p$  is not equal to 0, jump to first\_if.

I found out that the code we did together in class did not work for the power that greater than 4. We will have some weird output -nan. So I fix square and power function a little bit by deleting flds line. The idea is that fld1 already push 1 to floating-point stack, the function basically just modify  $st(0)$  by multiplying by  $r$  each time.

I think we do not need to use flds to push  $r$  onto floating-point stack in the beginning of power function, because it will cause stack overflow if we call power function too many time and we did not pop anything out.

asm source code appended to file lab3.s

```
power:
    push %ebp
    mov %esp, %ebp
    sub $4, %esp                # room for local variable t

    cmp $0, 8(%ebp)             # p == 0
    jne first_else              # if p != 0 go to first else

    # if p = 0
    fldl                        # push 1 to FP-stack, st(0) = 1
    jmp end_if                  # exit function
```

The first\_if is to do calculation if  $p$  is odd. If  $p$  is odd, decrement  $p$  and call power function for  $r^{p-1}$

asm source code appended to file lab3.s

```
first_else:
    push 8(%ebp)          # push p to stack for odd function's parameter
    call odd              # store 1 to %eax if p is odd
                           # or store 0 to %eax if p is even

    add $4, %esp
    cmp $1, %eax          # eax == 1
    jne second_else      # if p is even, go to second_else

    # if p is odd
    mov 8(%ebp), %eax      # eax = p
    dec %eax              # eax = p - 1
    mov %eax, -4(%ebp)     # t = eax = p - 1

    # call power function with two parameters r and t = p-1
    push 12(%ebp)          # r
    push -4(%ebp)          # t
    call power
    add $8, %esp
    fmul 12(%ebp)          # st(0) *= r
    jmp end_if
```

The second\_if is to do calculation if  $p$  is even. If  $p$  is even, divide  $p$  by 2 and call power function for  $r^{p/2}$ . We will use bit-shift to divide  $p$ .

asm source code appended to file lab3.s

```
|second_else:
|    mov 8(%ebp), %eax          # eax = p
|    shr %eax                  # eax = eax/2, bit-shift right will divide by 2
|    mov %eax, -4(%ebp)         # t = p/2
|
|    #call power function with two parameters r and t = p/2
|    push 12(%ebp)              # r
|    push -4(%ebp)              # t
|    call power
|    add $8, %esp
|
|    #push st(0) to stack to prepare parameter for calling square function
|    fsts (%esp)
|    call square
|    add $4, %esp
|
|end_if:
|    mov %ebp, %esp
|    pop %ebp
|    ret
```



This is another version of power that I wrote without using %ebp, but making use of variable r, p and ST0. It works fine, but the problem is that this power function can just be used for only one set of r and p.

asm source code appended to file lab3.s

```
power_more:
    cmp $0, 4(%esp)
    jne first_else_more
    fldl
    jmp end_if_more
first_else_more:
    push 4(%esp)
    call odd
    add $4, %esp
    cmp $1, %eax
    #jump to second_else if p is even
    jne second_else_more

    sub $1, p
    mov p, %eax
    push r
    push p
    call power_more
    add $8, %esp

    fmul r
    jmp end_if_more
```

asm source code appended to file lab3.s

```
second_else_more:
    mov p, %eax
    shr %eax
    mov %eax, p

    push r
    push p
    call power_more
    add $8, %esp

    fst ST0
    push ST0
    call square
    add $4, %esp
    jmp end_if_more
end_if_more:
    ret
```

## 4.4 Start program

asm source code appended to file lab3.s

```
.data
r: .float 5.0
p: .int 3
r2: .float 2.0
p2: .int 8
fmt: .string "%f\n"
fmt2: .string "%i\n"
ST0 : .float 0.0
.text
.globl _start
_start:
```

Testing square function with  $r = 5$ , expect result will be 25

asm source code appended to file lab3.s

```
flds r
push r
call square
add $4,%esp
#expect result in %st(0)

add $-8,%esp
fstpl (%esp) #push 64-bits st(0) onto the stack and pop st(0)
push $fmt
call printf
add $12, %esp
```

Testing odd function with  $p = 3$  and  $p = 8$ , expect result will be 1 and 0

asm source code appended to file lab3.s

```
push p
call odd
add $4,%esp
#expect result in %eax

push %eax
push $fmt2
call printf
add $8, %esp
```

asm source code appended to file lab3.s

```
push p2
call odd
add $4,%esp
#expect result in %eax

push %eax
push $fmt2
call printf
add $8, %esp
```

Testing power function with  $r = 5$ ,  $p = 3$  and  $r = 2$ ,  $p = 8$ , expect result will be 125 and 256

asm source code appended to file lab3.s

```
push r
push p
call power
add $8,%esp
#expect result in %eax
```

```
add $-8,%esp
fstpl (%esp) #push 64-bits st(0) onto the stack and pop st(0)
push $fmt
call printf
add $12, %esp
```

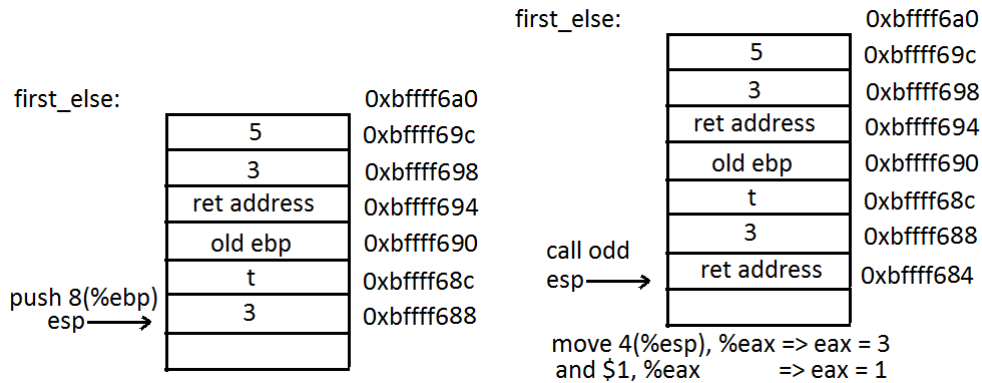
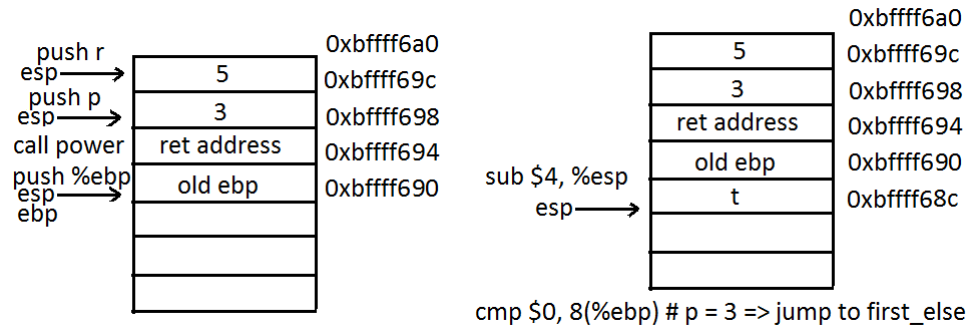
```
push r2
push p2
call power
add $8,%esp
#expect result in %eax
```

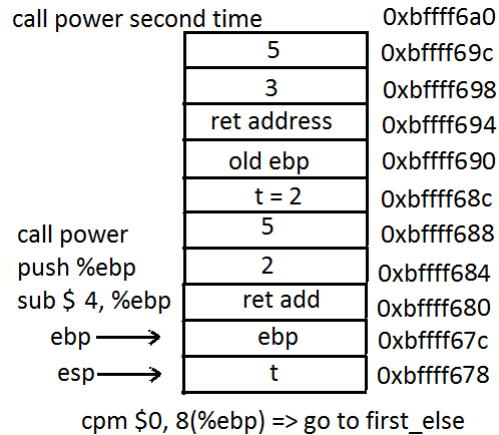
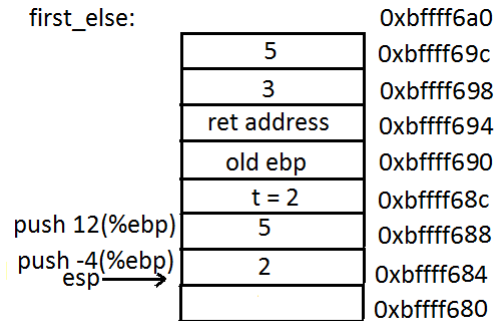
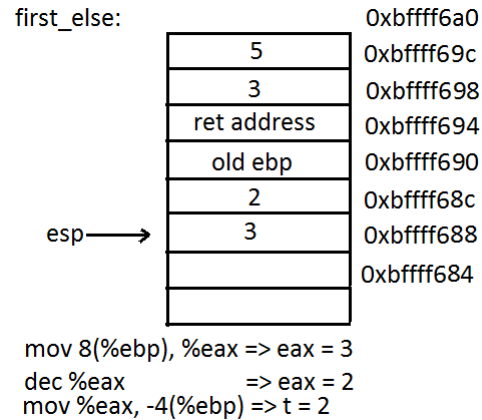
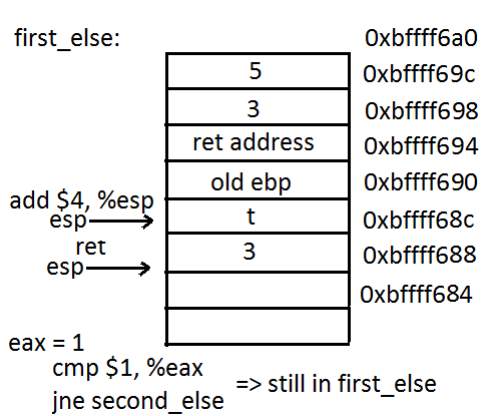
```
add $-8,%esp
fstpl (%esp) #push 64-bits st(0) onto the stack and pop st(0)
push $fmt
call printf
add $12, %esp
```

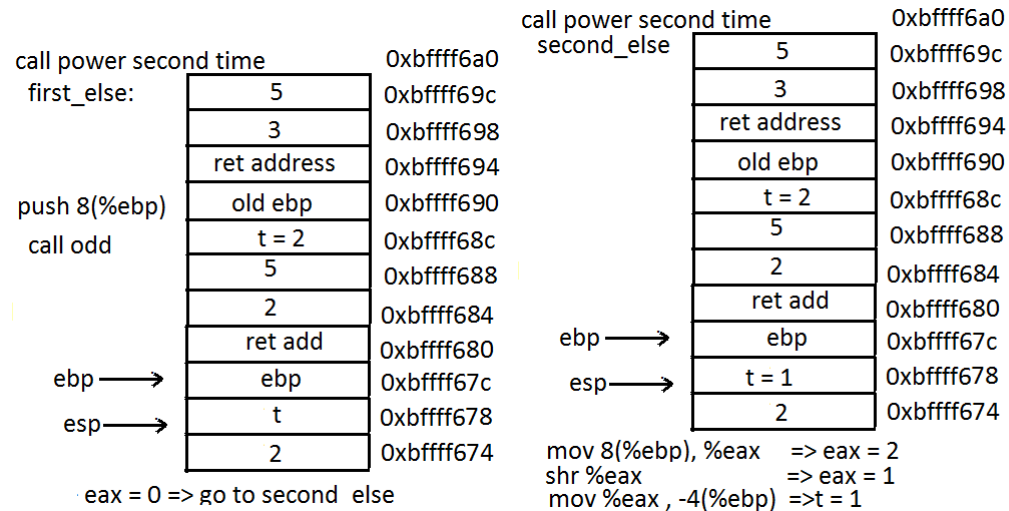
```
mov $1, %eax
mov $0, %ebx
int $0x80
```

```
debian@debian:~/labs/lab3$ ./lab3asm
25.000000
1
0
125.000000
256.000000
```

## 5 Stack Frame







call power third time		0xbffff6a0	call power fourth time		
	5	0xbffff69c	call odd	5	0xbffff688
	3	0xbffff698	p is odd	2	0xbffff684
	ret address	0xbffff694	=> call power	ret add	0xbffff680
	old ebp	0xbffff690	the fourth time	old ebp	0xbffff67c
	t = 2	0xbffff68c	eax = 1- 1 =0	t = 1	0xbffff678
	5	0xbffff688	=> t = 0	5	0xbffff674
	2	0xbffff684	push 12(%ebp)	1	0xbffff670
	ret add	0xbffff680	push -4(%ebp)	ret add	0xbffff66c
push 12(%ebp)	old ebp	0xbffff67c	call power	ebp	0xbffff668
push -4(%ebp)	t = 1	0xbffff678	push %ebp	t = 0	0xbffff664
call power	5	0xbffff674	sub \$ 4, %ebp	5	0xbffff670
push %ebp	1	0xbffff670		0	0xbffff65c
sub \$ 4, %ebp	ret add	0xbffff66c		ret add	0xbffff658
ebp →	ebp	0xbffff668	ebp →	ebp	0xbffff654
esp →	t	0xbffff664	esp →	t	0xbffff650



call power fourth time

p = 0	5	0xbffff688
=> fldz	2	0xbffff684
=> st(0) = 1	ret add	0xbffff680
Then ret from	old ebp	0xbffff67c
fourth power	t = 1	0xbffff678
FP - stack	5	0xbffff674
<div>1</div>	1	0xbffff670
st(0)	ret add	0xbffff66c
ebp →	ebp	0xbffff668
esp →	t = 0	0xbffff664
	5	0xbffff670
	0	0xbffff65c

continue third power function call

fmul 12(%ebp)	5	0xbffff6a0
=> st(0) = 5*1	3	0xbffff69c
= 5	ret address	0xbffff694
then return	old ebp	0xbffff690
FP - stack	t	0xbffff68c
<div>5</div>	5	0xbffff688
st(0)	2	0xbffff684
ebp →	ret add	0xbffff680
esp →	old ebp	0xbffff67c
	t = 1	0xbffff678
	5	0xbffff674
	1	0xbffff670

continue second power call:

	5	0xbffff6a0
	3	0xbffff69c
fsts (%ebp)	ret address	0xbffff694
=> push st(0)	old ebp	0xbffff690
onto stack	t	0xbffff68c
	5	0xbffff688
FP - stack	2	0xbffff684
5	ret add	0xbffff680
st(0)	old ebp	0xbffff67c
ebp →	t = 1	0xbffff678
	st(0) = 5	0xbffff674
esp →	1	0xbffff670

continue second power call:

	5	0xbffff6a0
call square	3	0xbffff69c
fmul 8(%ebp)	ret address	0xbffff694
=> st(0) = 5*5	old ebp	0xbffff690
= 25	t	0xbffff68c
Then return	5	0xbffff688
second power	2	0xbffff684
FP - stack	ret add	0xbffff680
25	old ebp	0xbffff67c
st(0)	t = 1	0xbffff678
ebp →	st(0) = 5	0xbffff674
esp →		0xbffff670
(after return from square)		

continue first power call:

	5	0xbffff6a0
FP - stack	3	0xbffff69c
125	ret address	0xbffff694
st(0)	old ebp	0xbffff690
ebp →	t	0xbffff68c
esp →	5	0xbffff688
fmul 12(%ebp)	2	0xbffff684
=> st(0) = 25 * 5		
= 125		

Then we finish calling power function for  $5^3$