

Chapter 10: Free type constructions

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The interpreter pattern in functional programming

- Consider

Free constructions in mathematics: Example I

- Consider the Russian letter \mathfrak{u} (tsè) and the Chinese word 水 (shuǐ)
- We want to *multiply* \mathfrak{u} by 水. Multiply how?
- Say, we want an associative (but noncommutative) product of them
 - ▶ So we want to define a *semigroup* that *contains* \mathfrak{u} and 水 as elements
 - ★ while we still know nothing about \mathfrak{u} and 水
- Consider the set of all *unevaluated expressions* such as $\mathfrak{u} \cdot \text{水} \cdot \text{水} \cdot \mathfrak{u} \cdot \text{水}$
 - ▶ Here $\mathfrak{u} \cdot \text{水}$ is different from $\text{水} \cdot \mathfrak{u}$ but $(a \cdot b) \cdot c = a \cdot (b \cdot c)$
- All these expressions form a **free semigroup** generated by \mathfrak{u} and 水
- Example calculation: $(\text{水} \cdot \text{水}) \cdot (\mathfrak{u} \cdot \text{水}) \cdot \mathfrak{u} = \text{水} \cdot \text{水} \cdot \mathfrak{u} \cdot \text{水} \cdot \mathfrak{u}$

How to represent this as a data type:

- Redundant encoding, as the full expression tree: $((\text{水}, \text{水}), (\mathfrak{u}, \text{水})), \mathfrak{u})$
 - ▶ Implement the operation $a \cdot b$ as pair constructor (easy and cheap)
- Reduced encoding, as a “smart” structure: $\text{List}(\text{水}, \text{水}, \mathfrak{u}, \text{水}, \mathfrak{u})$
 - ▶ Implement the operation $a \cdot b$ by concatenating the lists (more expensive)

Free constructions in mathematics: Example II

- Want to define a product operation for n -dimensional vectors: $\mathbf{v}_1 \otimes \mathbf{v}_2$
- The \otimes must be linear and distributive (but not commutative):

$$\mathbf{u}_1 \otimes \mathbf{v}_1 + (\mathbf{u}_2 \otimes \mathbf{v}_2 + \mathbf{u}_3 \otimes \mathbf{v}_3) = (\mathbf{u}_1 \otimes \mathbf{v}_1 + \mathbf{u}_2 \otimes \mathbf{v}_2) + \mathbf{u}_3 \otimes \mathbf{v}_3$$

$$\mathbf{u} \otimes (a_1 \mathbf{v}_1 + a_2 \mathbf{v}_2) = a_1 (\mathbf{u} \otimes \mathbf{v}_1) + a_2 (\mathbf{u} \otimes \mathbf{v}_2)$$

$$(a_1 \mathbf{v}_1 + a_2 \mathbf{v}_2) \otimes \mathbf{u} = a_1 (\mathbf{v}_1 \otimes \mathbf{u}) + a_2 (\mathbf{v}_2 \otimes \mathbf{u})$$

- ▶ We have such a product for 3-dimensional vectors only; ignore that
- Consider *unevaluated expressions* of the form $\mathbf{u}_1 \otimes \mathbf{v}_1 + \mathbf{u}_2 \otimes \mathbf{v}_2 + \dots$
 - ▶ A free vector space generated by pairs of vectors
- Impose the equivalence relationships shown above
 - ▶ The result is known as the **tensor product**
- Redundant encoding: unevaluated expression tree
 - ▶ A list of any number of vector pairs $\sum_i \mathbf{u}_i \otimes \mathbf{v}_i$
- Reduced encoding: a matrix
 - ▶ Reduced encoding requires proof and more complex operations

- 1 Show that