

Heaven's Light is Our Guide



# **Rajshahi University of Engineering and Technology**

## **Department of Computer Science and Engineering**

**Course No:** CSE.2202

**Course Title:** Sessional based on CSE.2201 (Computer Algorithms)

**Lab Report No:** 04

**Lab Report On:** Greedy Knapsack Algorithm.

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Title: A solution to Greedy knapsack problem for finding the optimal subset of the input data.

Introduction: The greedy method is the most straight-forward design technique and can be applied to a wide variety of problems. Most of the problems will have  $n$  inputs and required to find a subset that satisfies some constraints. Any subset that satisfies the constraints is called feasible solution. A feasible solution that either maximizes or minimizes a given objective function is called an optimal solution.

Description: In a knapsack problem we will be given  $n$  objects and a knapsack or bag. Object  $i$  will have a weight  $w_i$  and a profit  $P_i$  for  $w_i$ . The knapsack will have a capacity  $m$ . If a

Fraction  $x_i$ ,  $0 \leq x_i \leq 1$ , of object  $i$  is placed into the knapsack, then a profit  $p_i x_i$  is earned. The objective of knapsack problem is to maximize the profit by filling the knapsack and the total weight of all chosen objects to be at most  $m$ . Formally, the problem can be stated as,

$$\text{maximum } \sum p_i x_i$$

$$\text{subject to, } \sum w_i x_i \leq m$$

$$\text{where } 0 \leq x_i \leq 1, \quad 1 \leq i \leq n$$

$w_i$  = weight of object  $i$

$p_i$  = Profit of object  $i$

$x_i$  = fraction of  $w_i$ .

Algorithm:

Greedy knapsack ( $m, n$ ) {

// Objects are ordered as  $P[i]/w[i] \geq P[i+1]/w[i+1]$

//  $m$  is the size of knapsack

//  $n$  is the number of objects.

for  $i := 1$  to  $n$  do  $x[i] = 0.0$

$U := m$

for  $i := 1$  to  $n$  do {

if ( $w[i] > U$ ) then break;

$x[i] := 1.0$  ;  $U = U - w[i]$ ;

}

if ( $i \leq n$ ) then  $x[i] := U/w[i]$

}

Input: From the "greedy.txt" file.

Output: The output will be like this -



Knapsack: 20

Item	Weight	Profit
------	--------	--------

1	2	3
2	4	5
3	5	8
4	3	4
5	9	10

Solution: 1

Fractional Amount: 1 1 1 1 0.666667

Total Weight: 20

Total Profit: 26.6667

Knapsack: 21

Item	Weight	Profit
------	--------	--------

1	4	8
2	6	12
3	9	4
4	8	11
5	5	12
6	7	4
7	4	12
8	8	12

Solution: 2

Fractional Amount: 1 1 0 0 1 0 1 0.25

Total Weight: 21

Total Profit: 47

Knapsack: 26

Item	Weight	Profit
------	--------	--------

1	8	3
2	5	9
3	3	8
4	8	6
5	8	3
6	7	10
7	8	10
8	9	4

Solution: 3

Fractional Amount: 0 1 1 0.375 0 1 1 0

Total Weight: 26

Total Profit: 39.25

Knapsack: 22

Item	Weight	Profit
------	--------	--------

1	9	11
2	8	12
3	3	8
4	3	12
5	3	12
6	6	6
7	4	8
8	4	9
9	9	9

Solution: 4

Fractional Amount: 0 0.625 1 1 1 0 1 1 0

Total Weight: 22

Total Profit: 56.5

Knapsack: 21

Item	Weight	Profit
------	--------	--------

1	8	9
2	4	9
3	7	8
4	7	10
5	4	10

Solution: 5

Fractional Amount: 0 1 0.857143 1 1

Total Weight: 21

Total Profit: 35.8571

Knapsack: 29

Item	Weight	Profit
------	--------	--------

1	4	6
2	7	9
3	6	3
4	5	10
5	4	10
6	3	9
7	8	9
8	5	8

Solution: 6

Fractional Amount: 1 1 0 1 1 1 0.125 1

Total Weight: 29

Total Profit: 53.125

Knapsack: 25

Item	Weight	Profit
------	--------	--------

1	8	6
2	4	9
3	4	12
4	9	7
5	7	5

Solution: 7

Fractional Amount: 1 1 1 1 0

Total Weight: 25

Total Profit: 34

Knapsack: 26

Item	Weight	Profit
------	--------	--------

1	7	12
2	7	6
3	7	12
4	5	9
5	7	10
6	4	5
7	7	12

Solution: 8

Fractional Amount: 1 0 1 1 0 0 1

Total Weight: 26

Total Profit: 45

Knapsack: 27

Item	Weight	Profit
------	--------	--------

1	8	12
2	8	10
3	4	11
4	8	12
5	8	6
6	8	10

Solution: 9

Fractional Amount: 1 0.875 1 1 0 0

Total Weight: 27

Total Profit: 43.75

Knapsack: 25

Item	Weight	Profit
------	--------	--------

1	9	9
2	5	6
3	6	6
4	3	8
5	3	7
6	4	12
7	8	6

Solution: 10

Fractional Amount: 0.444444 1 1 1 1 1 0

Total Weight: 25

Total Profit: 43

Process returned 0 (0x0) execution time : 0.347 s

Press any key to continue.



Discussion and Conclusion : In the problem we had to determine the profit per weight first. Then from the maximum values we started filling the knapsack.

We could have successfully implemented the GreedyKnapsack algorithm in the code and get the maximized profit for  $m$  sized knapsack.