

Rajshahi University of Engineering & Technology
Department of Computer Science of Engineering

EXPERIMENT NO: 03

NAME OF EXPERIMENT: Arrays, Records and Pointers

SUBMITTED TO:

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GROUP: 2ND THIRTY

DATE OF EXP.: 16-09-2019

DATE OF SUB. : 23-09-2019

SERIES: 18 SERIES

MACHINE CONFIGURATION:

ASUS X510UF

CORE I5 8TH GEN PROCESSOR

UP To 3.4 GHZ

8 GB RAM

OS WIN 10

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Preblem No.	Name of Problem
01.	Finding the Complexity of Bubble Sort
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07.	<empty>
08	<empty>
09	<empty>
10.	<empty>

THEORY: A LINEAR ARRAY IS A LIST OF FINITE NUMBER N OF HOMOGENEOUS DATA ELEMENTS (I.E., DATA ELEMENTS OF THE SAME TYPE) SUCH THAT:

1. The elements of the array are referenced respectively by an Index set consisting of n consecutive numbers.
2. The elements of the **Array** are stored respectively in successive memory locations.

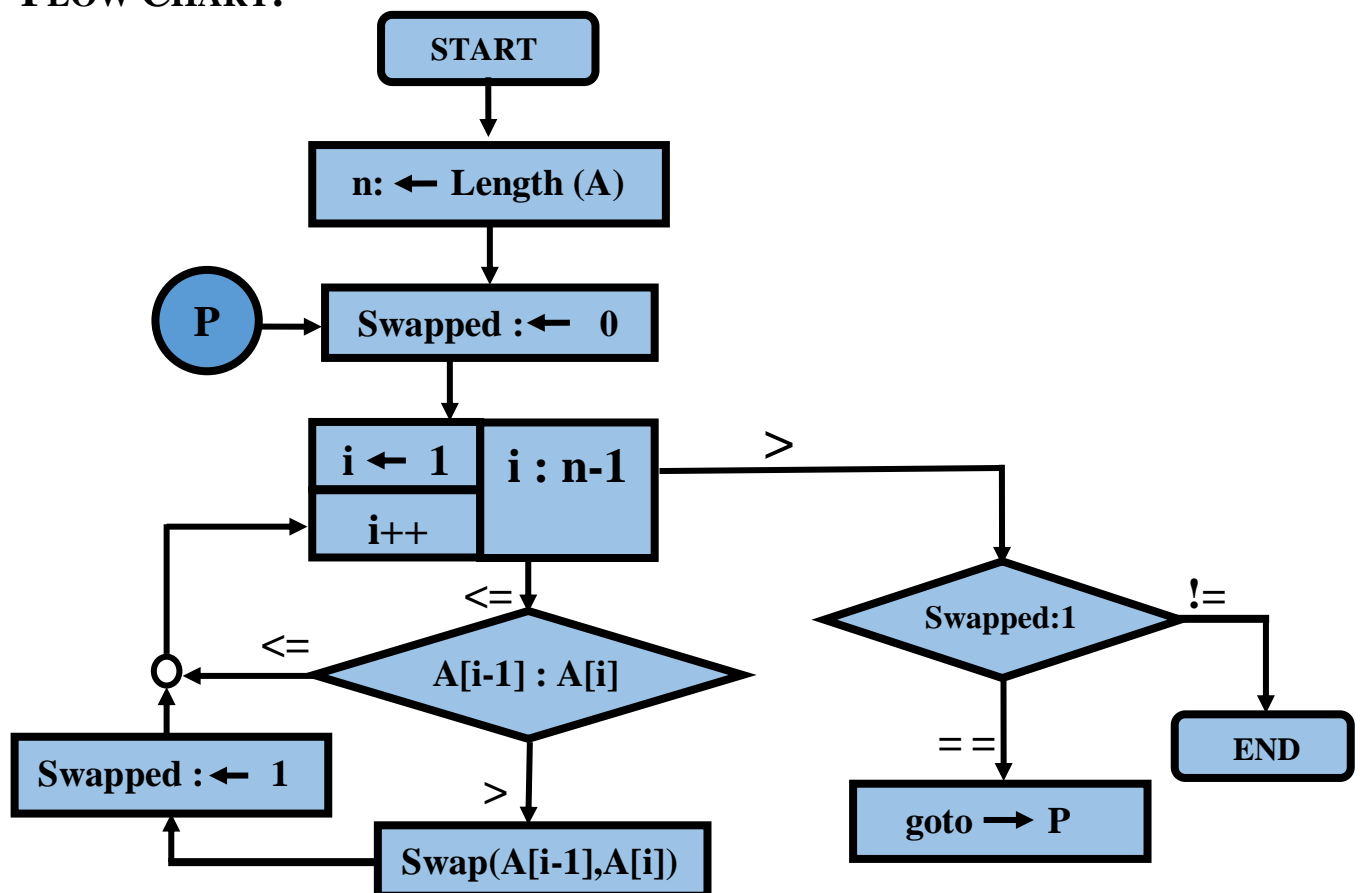
The number n of elements is called the **Length** or **Size** of the Array.

In the lab we -

1. Find the **Complexity of Bubble Sort**. In Bubble Sorting we can sort an array in Ascending or Descending order.
2. Searched an element from an array by **Linear Search**. In Linear Search we check all the elements of an array one by one until we find out our searched data.
3. Searched an element from an array by **Binary Search**. For binary search the **Array must be Sorted**. In Binary Search we check the elements of an array part by part by divideing the array in some condition.
4. Multiplied two matrixes (**Matrix Multiplication**).

PROBLEM 1: Find the Complexity of the bubble Sort algorithm.

FLOW CHART:



ALGORITHM:

(Given a nonempty array A with n numerical values. This algorithm sorts the values)

1. Set $n := \text{length}(A)$
2. swapped = false
3. Repeat i = 1 to n-1 by 1
4. If $A[i-1] > A[i]$ then:
 Swap ($A[i-1]$, $A[i]$) and swapped = true.
 [End of If Structure]
 [End of Repeat 4 loop]
5. If swapped = true then:
 Go to step 2
 [End of If structure]
6. Exit

CODE:

```
#include<stdio.h>

int main()
{
    int A[]={43,21,67,98,76,47,23,12,69,55},n,i,j=0,temp,swapped,p;
    n=sizeof(A)/4;
    p: swapped=0;
    for(i=1;i<n;i++)
    {
        if(A[i-1]>A[i])
        {
            temp=A[i-1];
            A[i-1]=A[i];
            A[i]=temp;
            swapped=1;
        }
    }

    if(swapped==1)
        goto p;

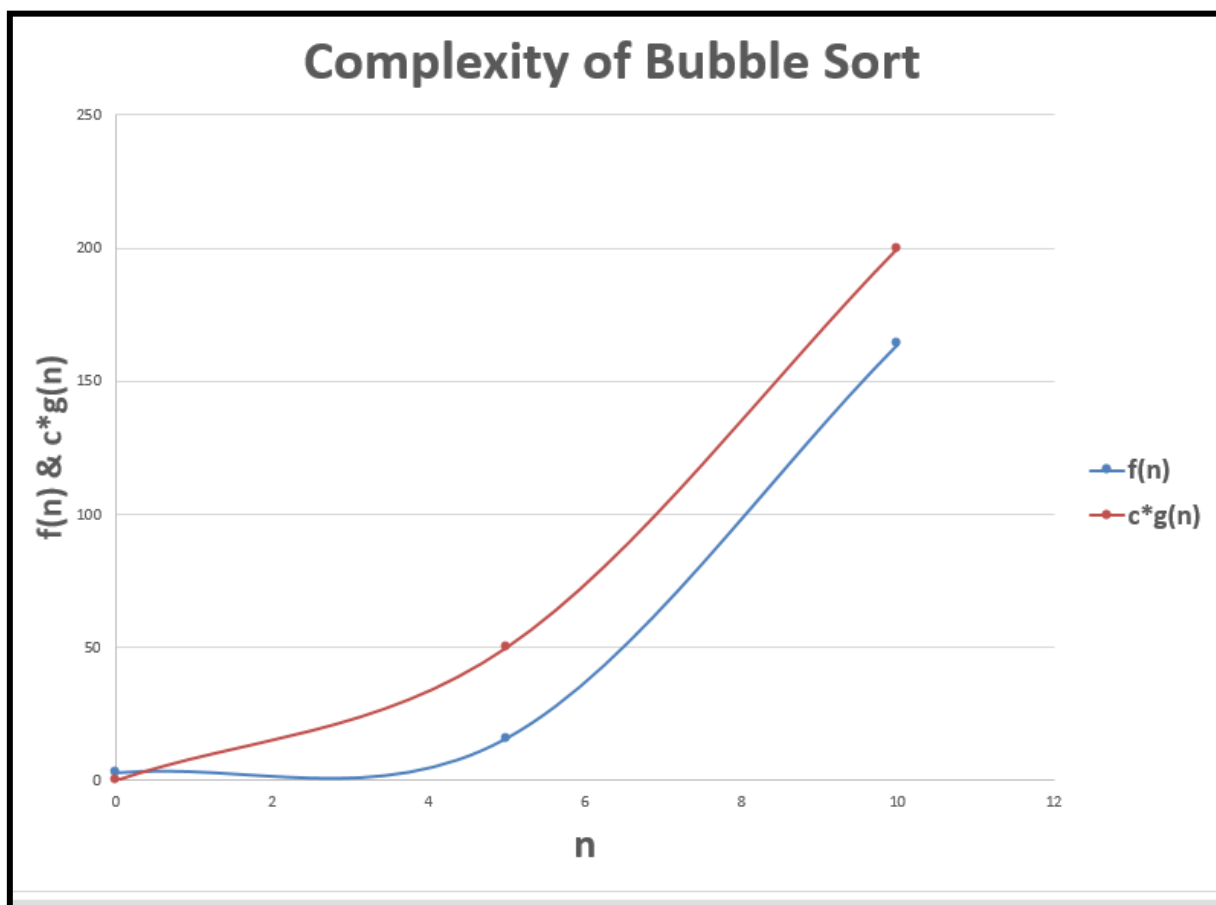
    return 0;
}
```

Complexity: The complexity of the above algorithm is $O(n^2)$. (Let $c = 4$)

Complexity Table:

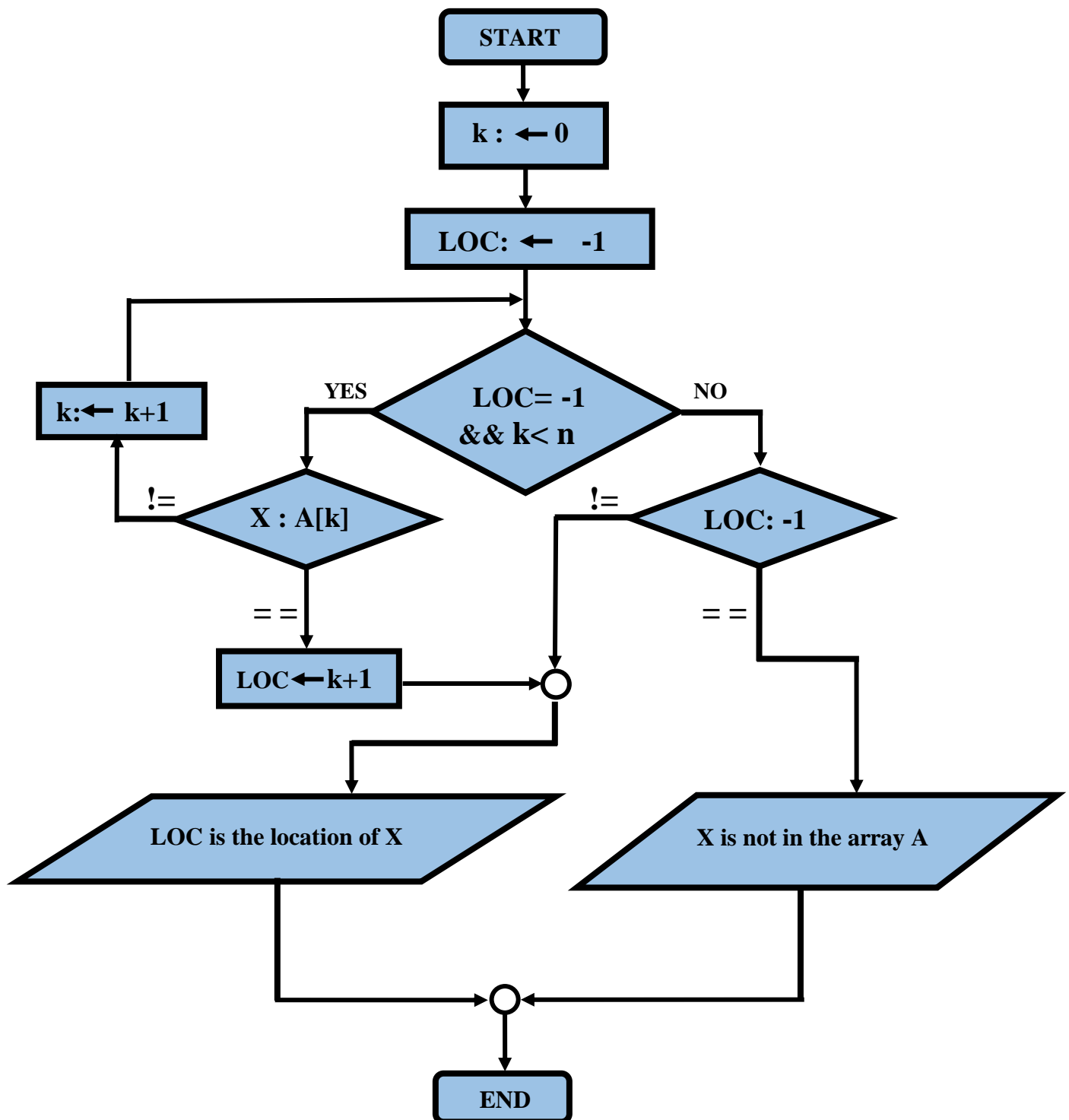
n	f(n)	c*g(n)
0	3	0
5	16	50
10	164	200

GRAPH:



PROBLEM 2: Linear Search from an Linear Array.

FLOW CHART:



ALGORITHM:

Given a nonempty array A with n numerical values and a specific x of information is given. This algorithm finds the location LOC of x in the array A or Sets LOC=-1)

1. Set K:=1, LOC:=-1
2. Repeat steps 3 and 4 while LOC = -1 and $K \leq n$
3. IF $x = A[K]$ then:
 Set LOC:=K.
 [End of If structure]
4. K:=K+1.
 [End of step 2 loop]
5. If LOC = -1 then: Write: x is not in the array A.
 Else: Write: LOC is the location of x
 [End of If structure]
6. Exit

CODE:

```
#include<stdio.h>

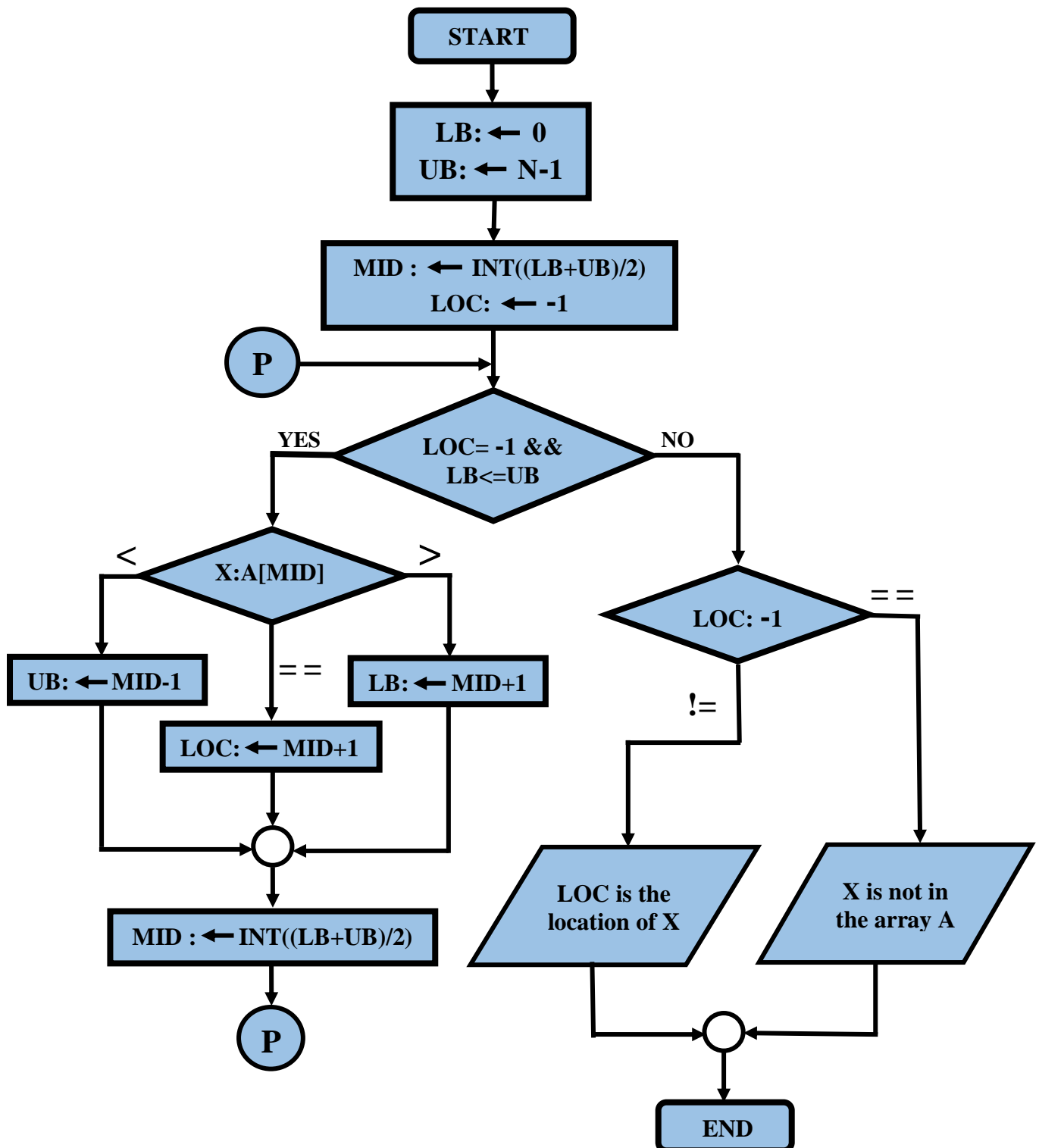
int main()
{
    int A[]={43,21,67,98,76,47,23,12,69,55},x,k=0,n,loc=-1;
    n=sizeof(A)/4;
    scanf("%d",&x);
    while(loc==-1&&k<n)
    {
        if(x==A[k])
            loc=k+1;
        k++;
    }
    if(loc==-1)
        printf("%d is not in the Array\n",x);
    else
        printf("%d is the location of %d\n",loc,x);

    return 0;
}
```

Complexity: The worst case complexity is $C(n) = n$ when X is the last element of the array and the average case complexity is $C(n) = (n+1)/2$.

PROBLEM 3: Binary Search from an Linear Array.

FLOW CHART:



ALGORITHM:

(Given a sorted array A with n numerical values and a specific x of information is given. This algorithm finds the location LOC of x in the array A or Sets LOC=-1)

1. Set LB:=1, UB:=n, MID = INT((LB+UP)/2) and LOC:=-1
2. Repeat steps 3 and 4 while LOC = -1 and LB≤UB
3. IF x < A[MID] then:
 Set UB:=MID-1.
 Else If x > A[MID] then:
 Set LB:=MID+1.
 Else:
 Set LOC:=MID
 [End of If structure]
4. MID = INT((LB+UP)/2).
 [End of step 2 loop]
5. If LOC = -1 then: Write: x is not in the array A.
 Else: Write: LOC is the location of x
 [End of If structure]
6. Exit

CODE:

```
#include<stdio.h>

int main()
{
    int A[]={23,26,31,37,45,53,65,69,74,78,81,85,89,93,99};
    int n,x,lb,ub,mid,loc=-1;

    n=sizeof(A)/4;
    lb=0; ub=n-1;
    mid=(lb+ub)/2;
    scanf("%d",&x);
```

```

while(loc==-1&&lb<=ub)
{
    if(x<A[mid])
        ub=mid-1;
    else if(x>A[mid])
        lb=mid+1;
    else
        loc=mid+1;
    mid=(lb+ub)/2;
}
if(loc==-1)
    printf("%d is not in the Array\n",x);
else
    printf("%d is the location of %d\n",loc,x);

return 0;
}

```

Complexity:

After,

Loop 1	-----	$n / 2^1$
Loop 2	-----	$n / 2^2$
Loop 3	-----	$n / 2^3$
.....		
.....		

After the last loop (say, k) there will be only one element in the array.

So after,

Loop k	-----	$1 = n / 2^k$
--------	-------	---------------

As, $1 = n / 2^k$

or, $2^k = n$

or, $\log_2 2^k = \log_2 n$

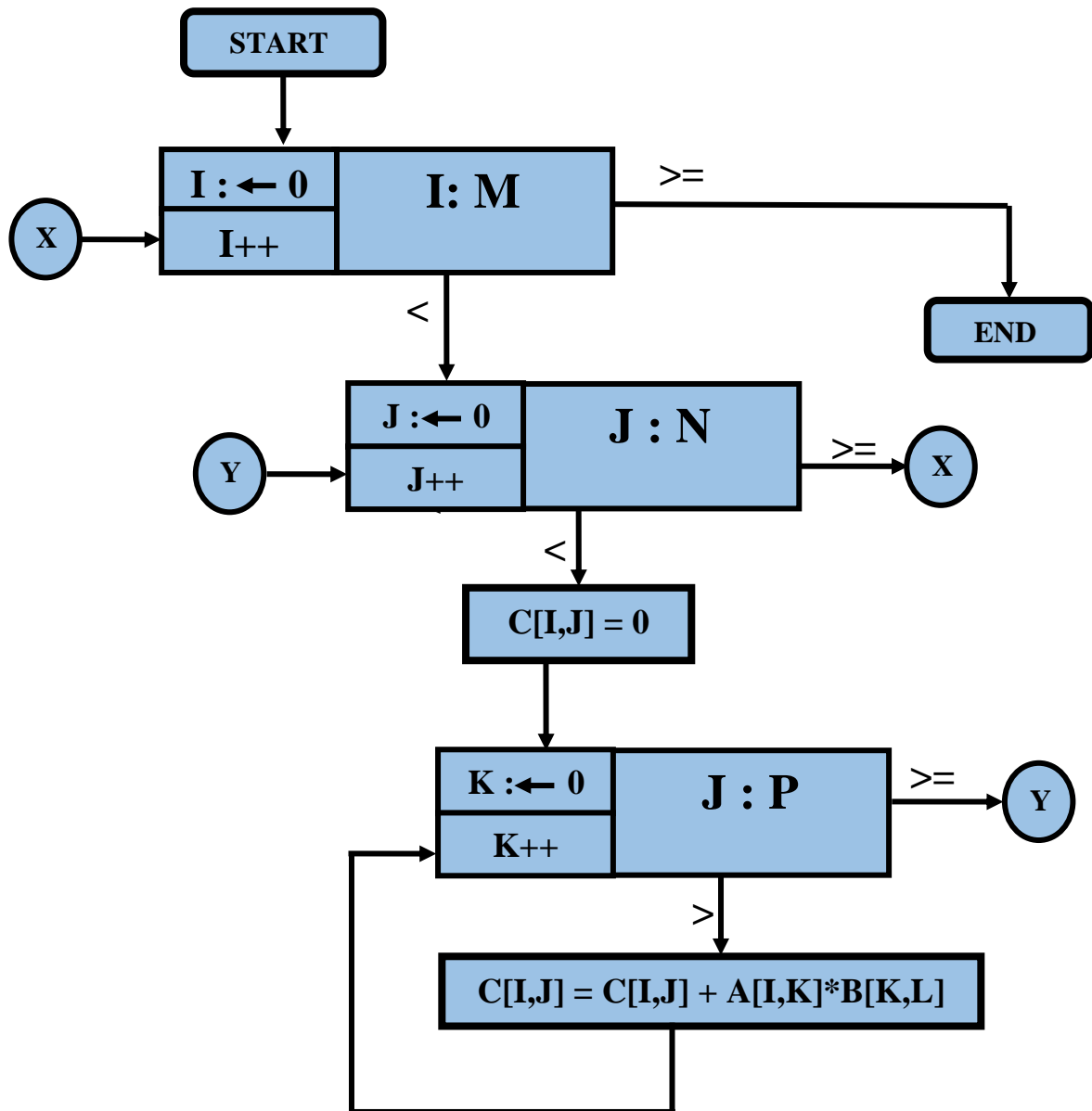
or, $k \log_2 2 = \log_2 n$

or, $k = \log_2 n$

So the complexity is **$O(\log_2 n)$** .

PROBLEM 4: Multiplication of Matrix.

FLOW CHART:



ALGORITHM:

(Let A be an $M \times P$ matrix array and B be an $P \times N$ matrix array. This algorithm stores the product of A and B in an $M \times N$ matrix array C)

1. Repeat steps 2 to 4 for $I = 1$ to M :
2. Repeat steps 3 and 4 for $J = 1$ to N :
3. Set $C[I,J] := 0$.
4. Repeat for $K = 1$ to P : $C[I,J] := C[I,J] + A[I,K] * B[K,J]$.

[End of step 4 loop]
[End of step 2 loop]
[End of step 1 loop]
5. Exit

CODE:

```
#include<stdio.h>

int main()
{
    int A[3][2]={ 2,3,1,2,5,3 },B[2][3]={ 1,4,2,3,1,3 };
    int C[3][3],i,j,k;

    for(i=0;i<3;i++)
        for(j=0;j<3;j++)
        {
            C[i][j]=0;
            for(k=0;k<2;k++)
                C[i][j]=C[i][j]+A[i][k]*B[k][j];
        }

    return 0;
}
```

Complexity: In the programme there is running 3 nested 'for' loop. We know each loop's complexity is n. So for 3 nested for loop the complexity will be $n*n*n = n^3$.

So the complexity is **$O(n^3)$** .

DISCUSSION: The problems were based on Array. After trying sometime I had solved the problems.

The End