

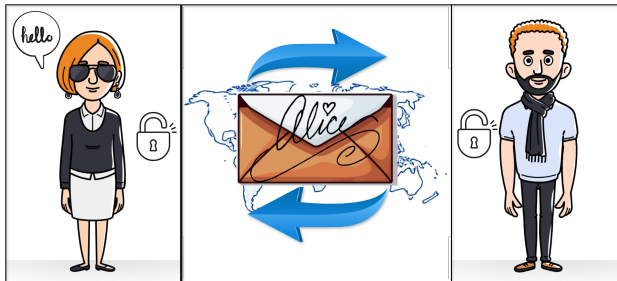
Polynomial Arithmetic Tools

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Introduction: Quantum Computers and Cryptography



Multiplication Comparison [Kan]

| Technique | Constraints on q | Constraints on n |
|------------|--|--------------------|
| Schoolbook | None | None |
| Karatsuba | None | Divisible by 2 |
| NTT | Primitive $2n$ -th root of unity in \mathbb{Z}_q | Power of 2 |

Ring Example

Multiplication in $\mathbb{Z}_7[x]/(x^3 + 1)$

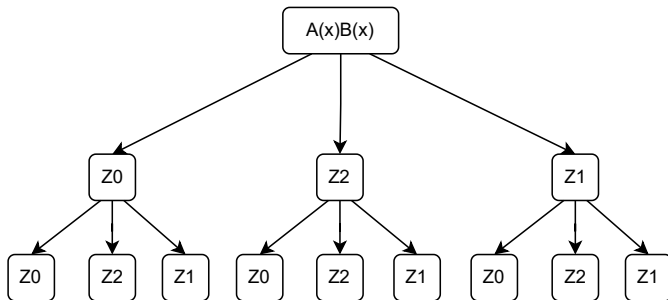
$$(1x^2 + 2x + 3) \cdot (4x^2 + 5x + 6)$$

| | x^4 | x^3 | x^2 | $x^1, -x^4$ | $x^0, -x^3$ |
|--|----------------------|----------------------|----------------------|----------------------|----------------------|
| | | | $3 \cdot 4 \equiv 5$ | $3 \cdot 5 \equiv 1$ | $3 \cdot 6 \equiv 4$ |
| | | $2 \cdot 4 \equiv 1$ | $2 \cdot 5 \equiv 3$ | $2 \cdot 6 \equiv 5$ | |
| | $1 \cdot 4 \equiv 4$ | $1 \cdot 5 \equiv 5$ | $1 \cdot 6 \equiv 6$ | | |
| | | | | $1 \cdot 4 \equiv 4$ | $2 \cdot 4 \equiv 1$ |
| | | | | | $1 \cdot 5 \equiv 5$ |
| | | | $14 \equiv 0$ | $2 \equiv 2$ | $-2 \equiv 5$ |

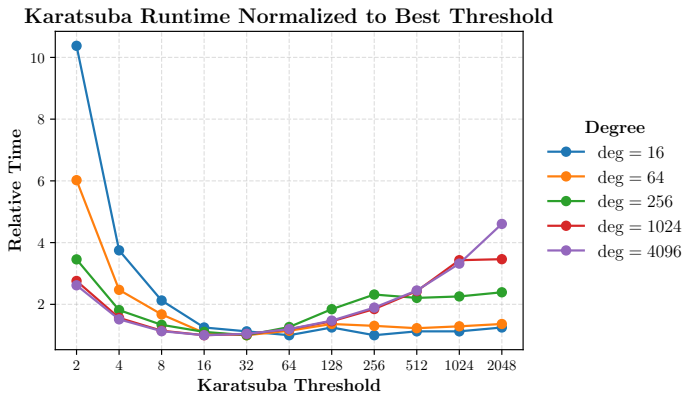
$$0x^2 + 2x + 5$$

Karatsuba

Split into 3 multiplications: $O(n^{1.585})$

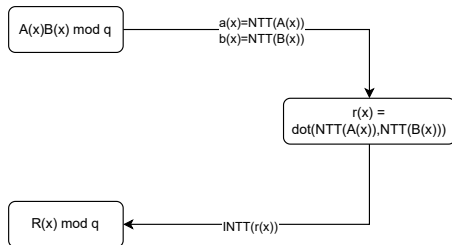


Karatsuba Recursion

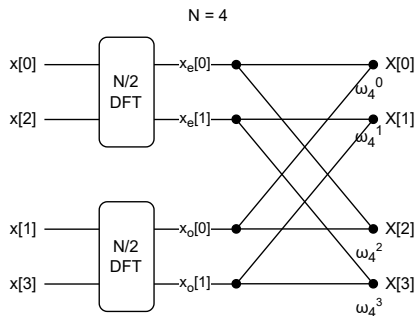


Number-Theoretic Transform (NTT)

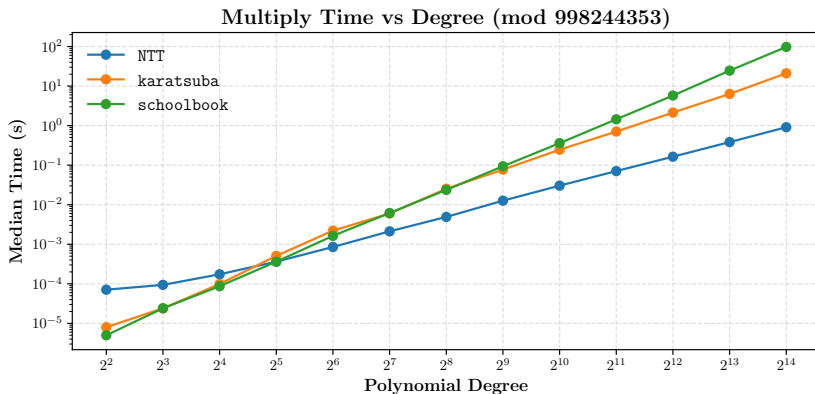
- Analagous to the DFT.
- FFT like acceleration possible. $\approx O(n \log(n))$



NTT with Cooley-Tukey Butterfly



Multiplication Comparison



Modular Reduction: overview

- Division-based
- Barrett
- Montgomery
- Mersenne and Pseudo-Mersenne primes

Division-based Reduction

- Division instructions are slow [Fog23].
- Avoid at all costs.

Barrett Reduction

- Precomputes a constant $m = \lfloor 2^k / q \rfloor$ to approximate division.
- Avoids division during runtime by replacing it with shifts and multiplications.

Montgomery Reduction

- Avoids division by transforming into a special "Montgomery domain."
- Uses $R = 2^k$, where $k > \log_2(q)$, to simplify calculations.
- Efficient for large q , but requires q coprime with R .
- Example:

$T \bmod q$ is computed as $(T + m \cdot q)/R$

Mersenne Prime Reduction

If $q = 2^k - 1$: $x \bmod q = (x \& (2^k - 1)) + (x \gg k)$.

- Since $2^k \equiv 1 \pmod{q}$, each “high-bits” chunk adds back to the low part.
- If the sum $((x \& (2^k - 1)) + (x \gg k))$ still $\geq q$, subtract q once more.
- Example: $q = 31 = 2^5 - 1$, $x = 45$:

$$45 \bmod 31 = (45 \& 31) + (45 \gg 5) = (13) + (1) = 14.$$

Dilithium Pseudo-Mersenne Reduction [AMI+23]

If $q = 2^{23} - 2^{13} + 1 \Rightarrow 2^{23} \equiv 2^{13} - 1 \pmod{q}$,

Require: x

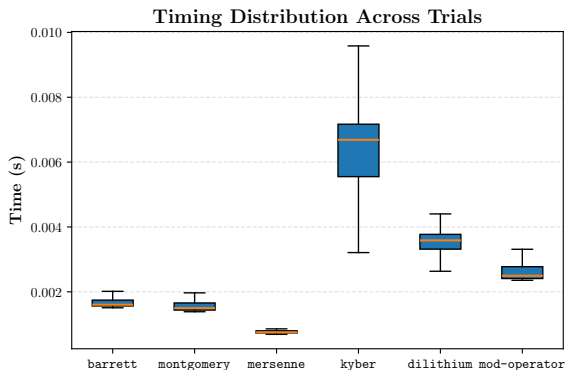
▷ integer to reduce

```

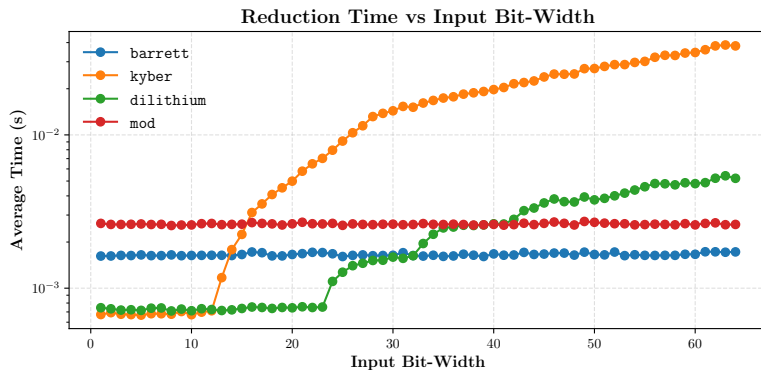
1: while  $x \gg 23 \neq 0$  do
2:    $hi \leftarrow x \gg 23$ 
3:    $lo \leftarrow x \& ((1 \ll 23) - 1)$ 
4:    $x \leftarrow lo + (hi \ll 13) - hi$ 
5: end while
6: if  $x \geq 8380417$  then
7:    $x \leftarrow x - 8380417$ 
8: end if
9: return  $x$ 

```

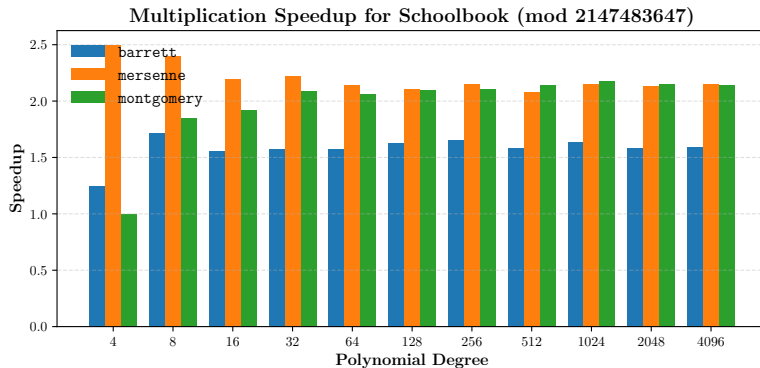
Modular Reduction comparison



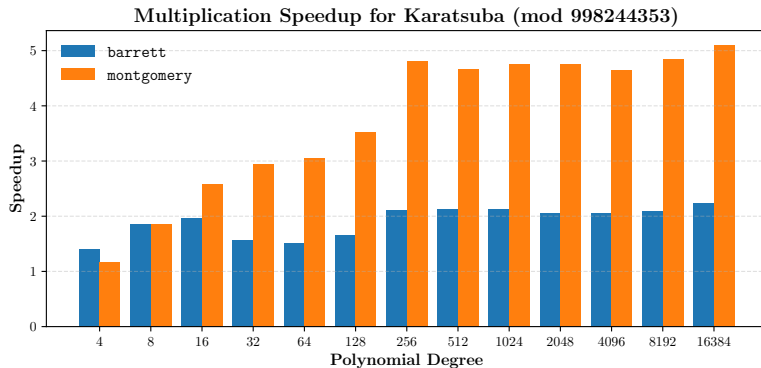
Modular Reduction comparison



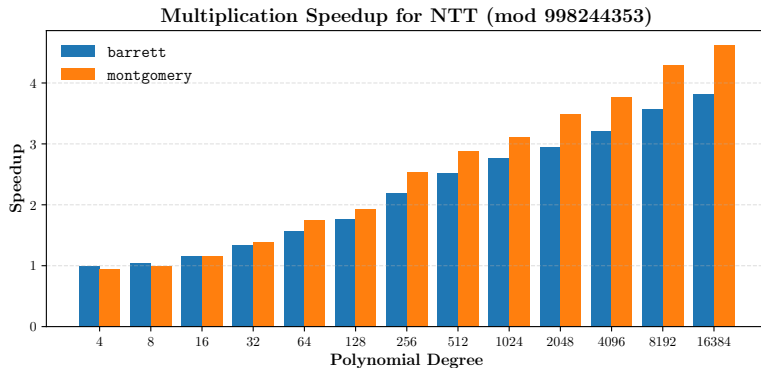
Modular Reduction Results: Schoolbook



Modular Reduction Results: Karatsuba



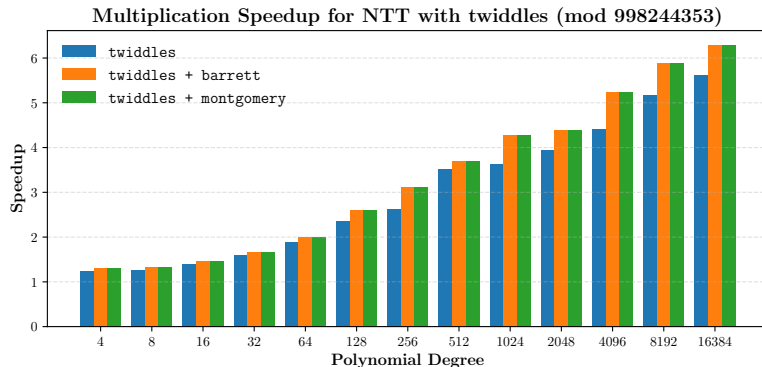
Modular Reduction Results: NTT



Twiddle Precomputation

1. Twiddles can be precomputed for (n, q, ω)
2. For montgomery they can be computed in montgomery-form.

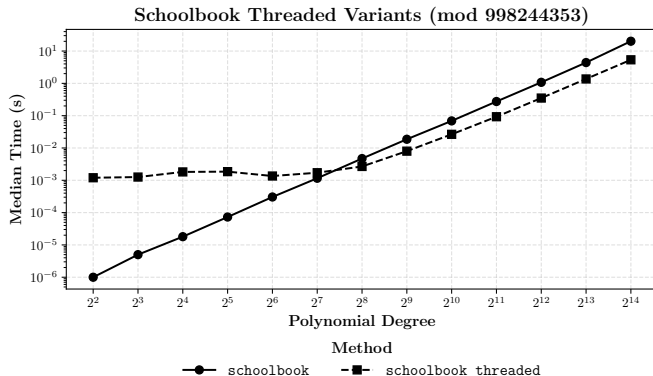
Modular Reduction Results: NTT + Twiddles



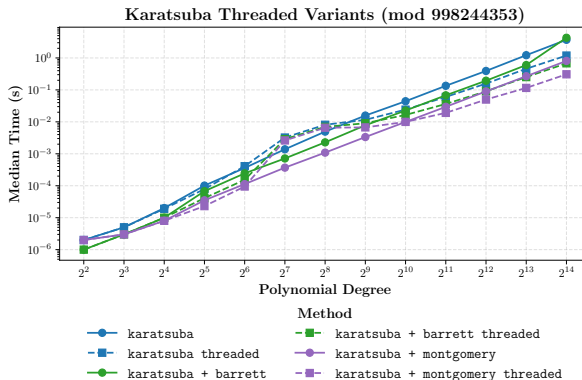
Multithreading

1. For each method split the workload.
2. For recursive calls turn function call into a thread.

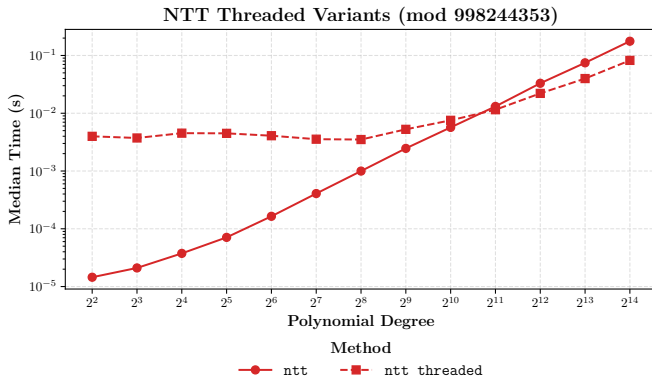
Multithreading: Schoolbook



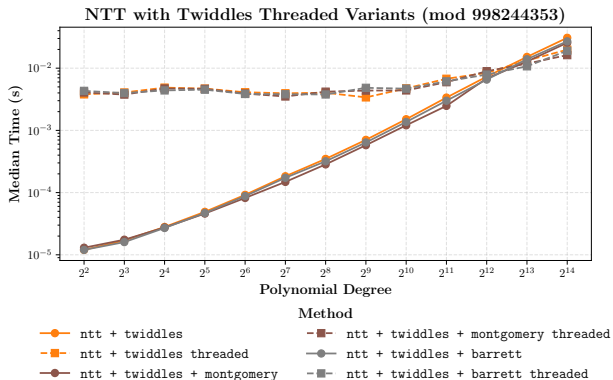
Multithreading: Karatsuba



Multithreading: NTT



Multithreading: NTT + Twiddles



Tool Capabilities

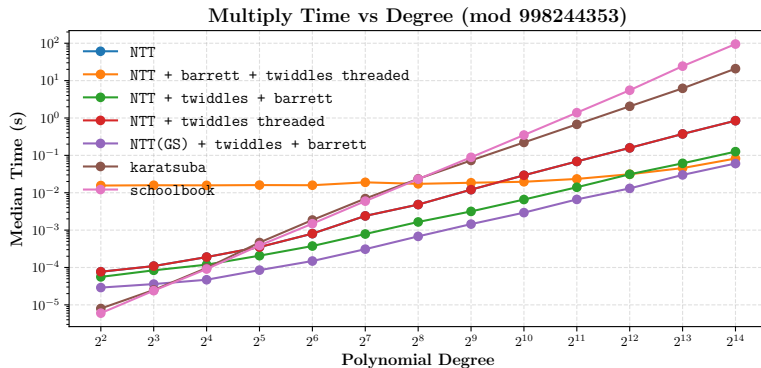
- Basic polynomial Arithmetic
- Variety of multiplication and reduction techniques.
- Support for testing.

Further Improvements and Future Work

- Montgomery only lift one operand into representation. [Sei18].
- CPU Specifics: AVX2 vectorization.
- Avoiding memory allocations.
- Gentleman-Sande butterfly for INTT.
- Branchless code.

Summary

Summary: Results



References I

- [AMI+23] Aikata Aikata, Ahmet Can Mert, Malik Imran, Samuel Pagliarini, and Sujoy Sinha Roy. **KaLi: A Crystal for Post-Quantum Security Using Kyber and Dilithium.** *IEEE Transactions on Circuits and Systems I: Regular Papers* 70.2 (Feb. 2023). Conference Name: IEEE Transactions on Circuits and Systems I: Regular Papers, pp. 747–758. ISSN: 1558-0806. DOI: [10.1109/TCSI.2022.3219555](https://doi.org/10.1109/TCSI.2022.3219555). URL: <https://ieeexplore.ieee.org/abstract/document/9946370> (visited on 11/17/2024).

References II

- [Fog23] Agner Fog. **Instruction Tables: Lists of instruction latencies, throughputs and micro-operation decomposition.**
https://www.agner.org/optimize/instruction_tables.pdf. Accessed: 2025-06-01. 2023.
- [Kan] Matthias J Kannwischer. **Polynomial Multiplication for Post-Quantum Cryptography.** en ().
- [Sei18] Gregor Seiler. **Faster AVX2 optimized NTT multiplication for Ring-LWE lattice cryptography.** Cryptology ePrint Archive, Paper 2018/039. 2018.
URL: <https://eprint.iacr.org/2018/039>.