

First Order Masking for Hybrid Homomorphic Encryption Schemes

Nedžma Mušović

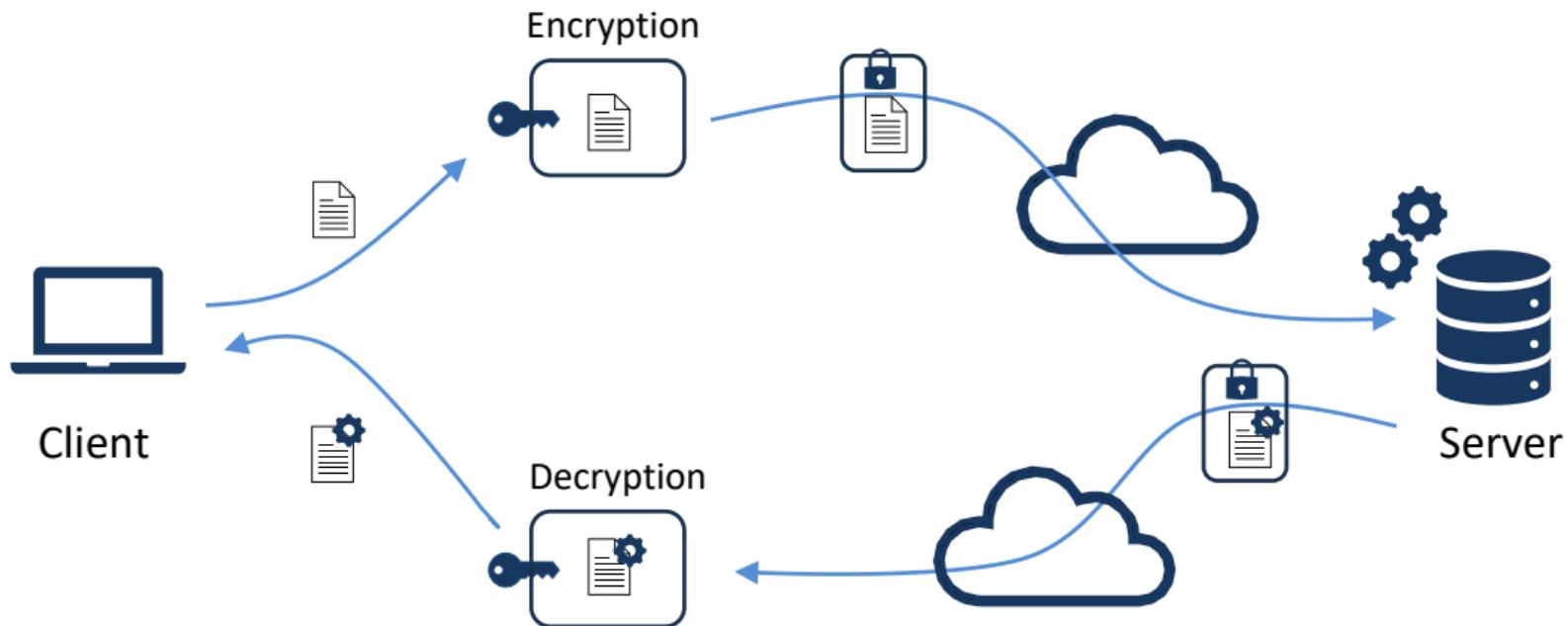
Master Project Presentation

Outline

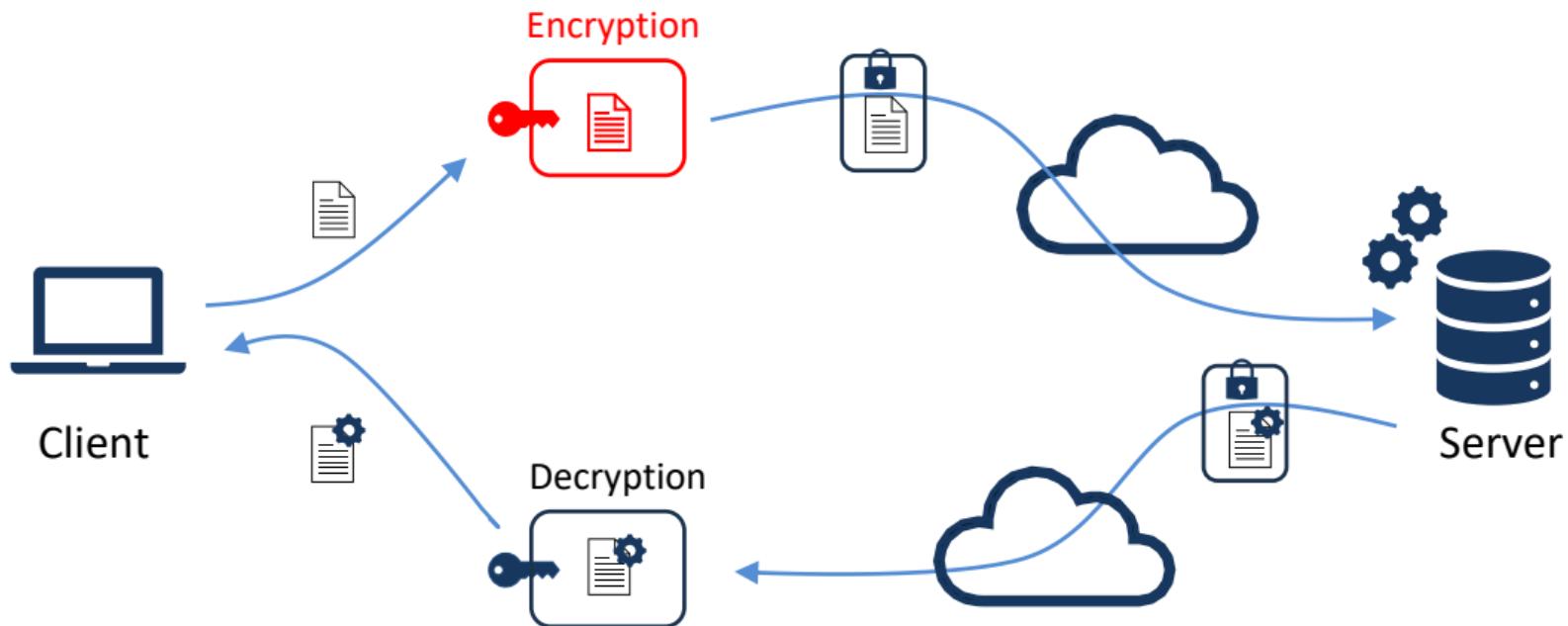
Outline

- i Introduction
- i **HERA - First Order Masking**
- i Application to **PASTA**
- i Higher Order Masking and Conclusion

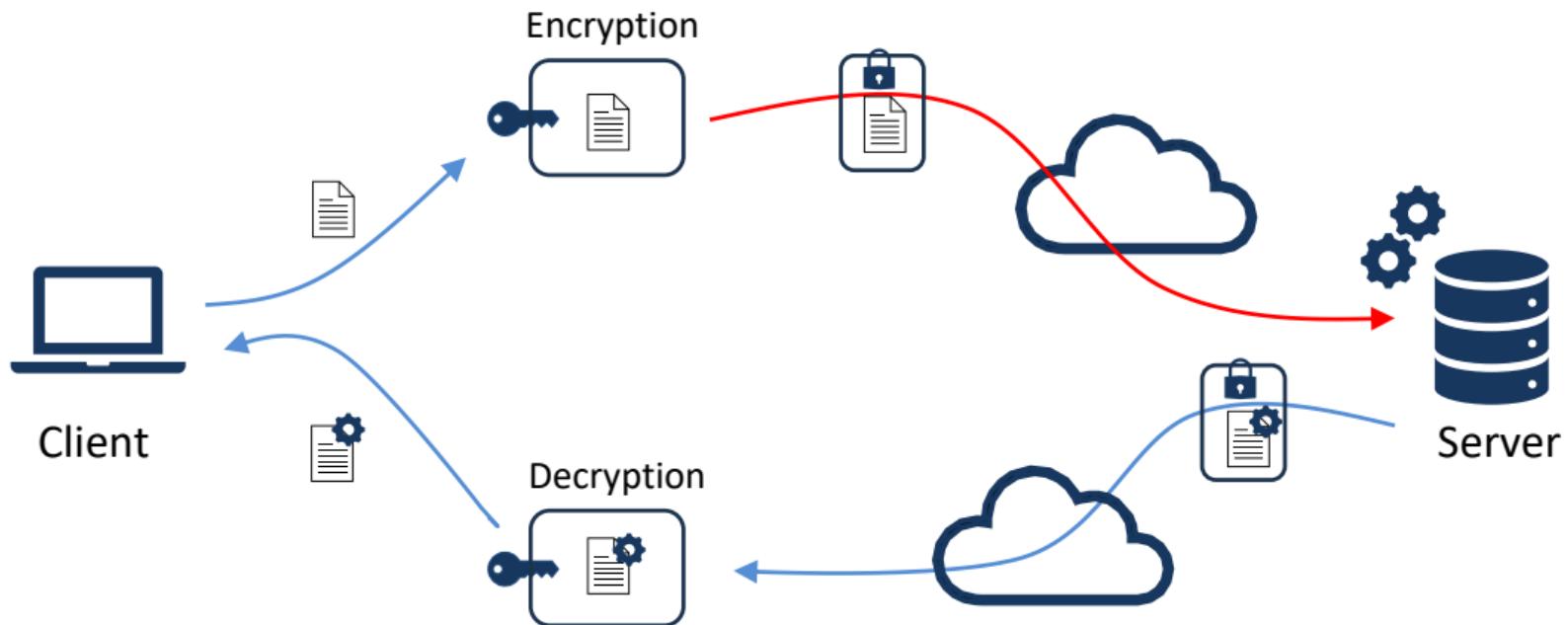
Homomorphic Encryption



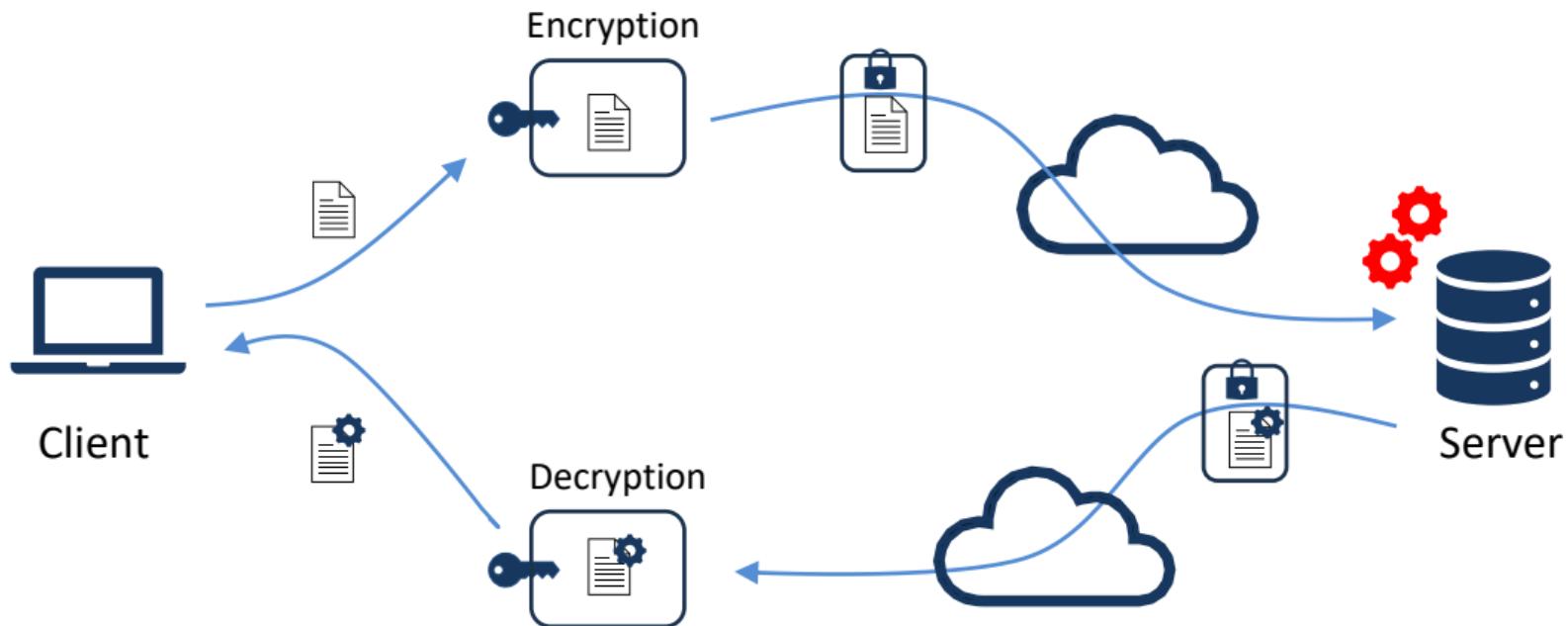
Homomorphic Encryption



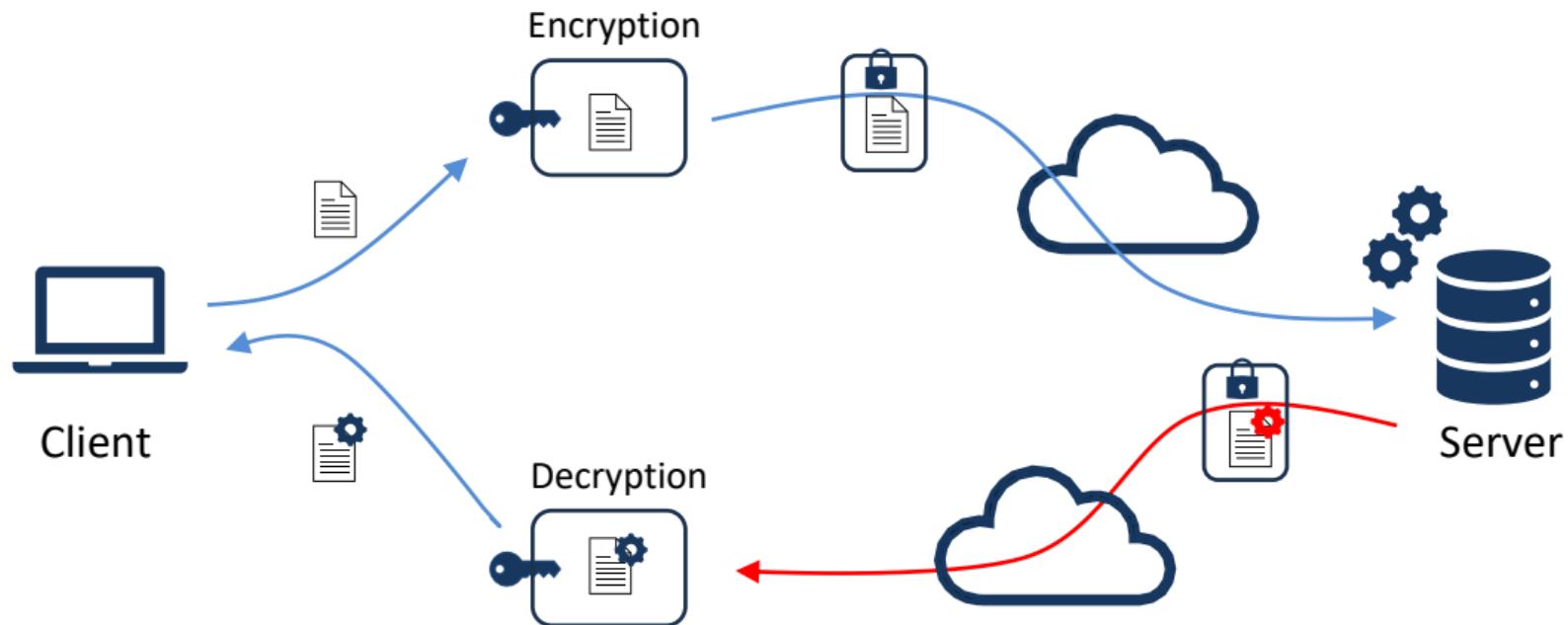
Homomorphic Encryption



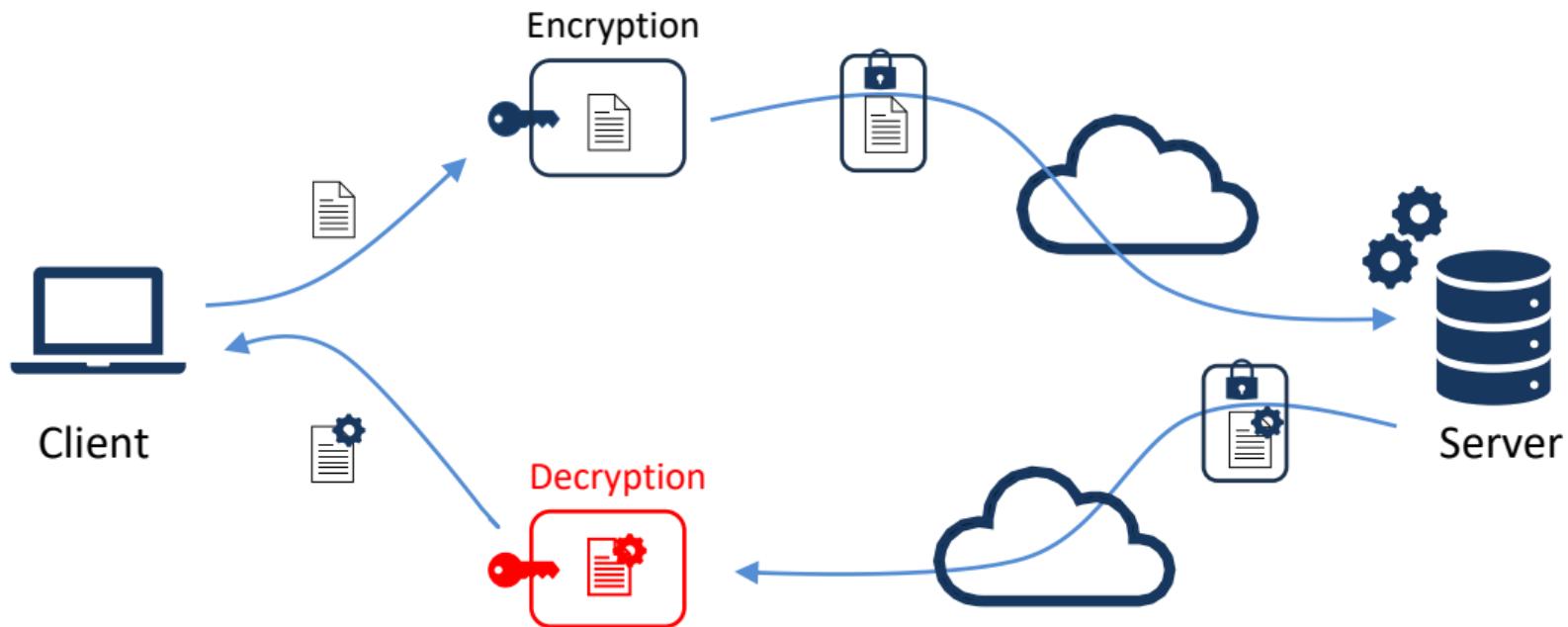
Homomorphic Encryption



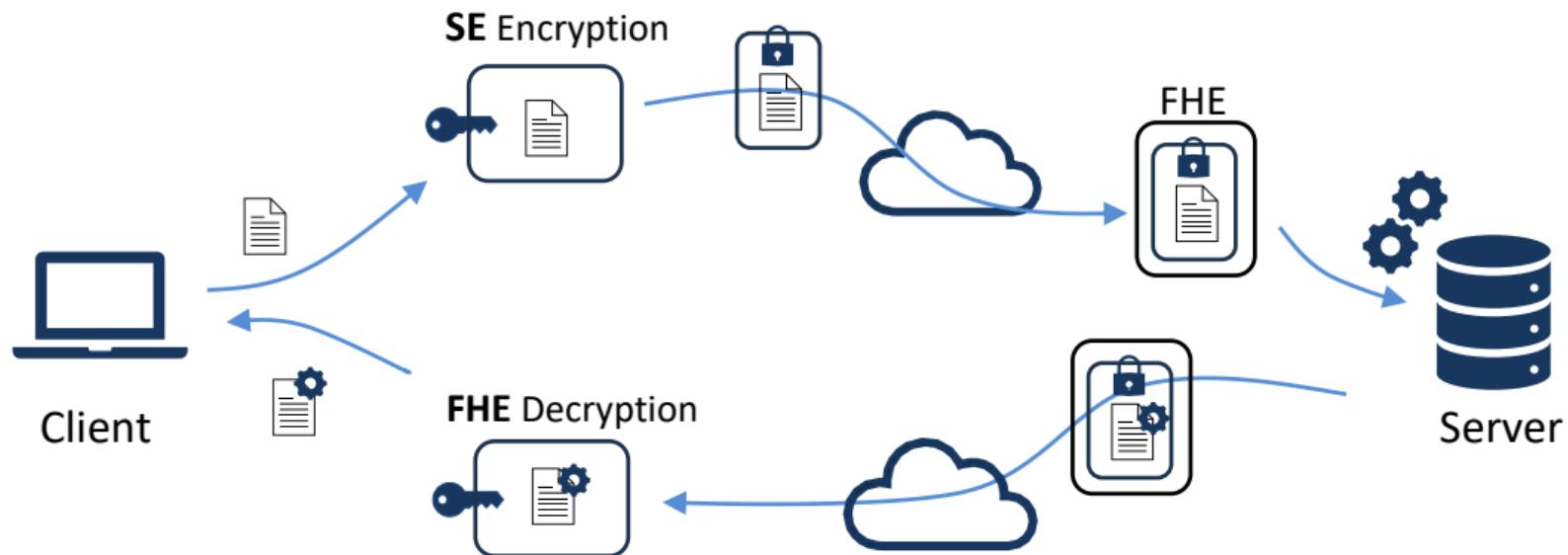
Homomorphic Encryption



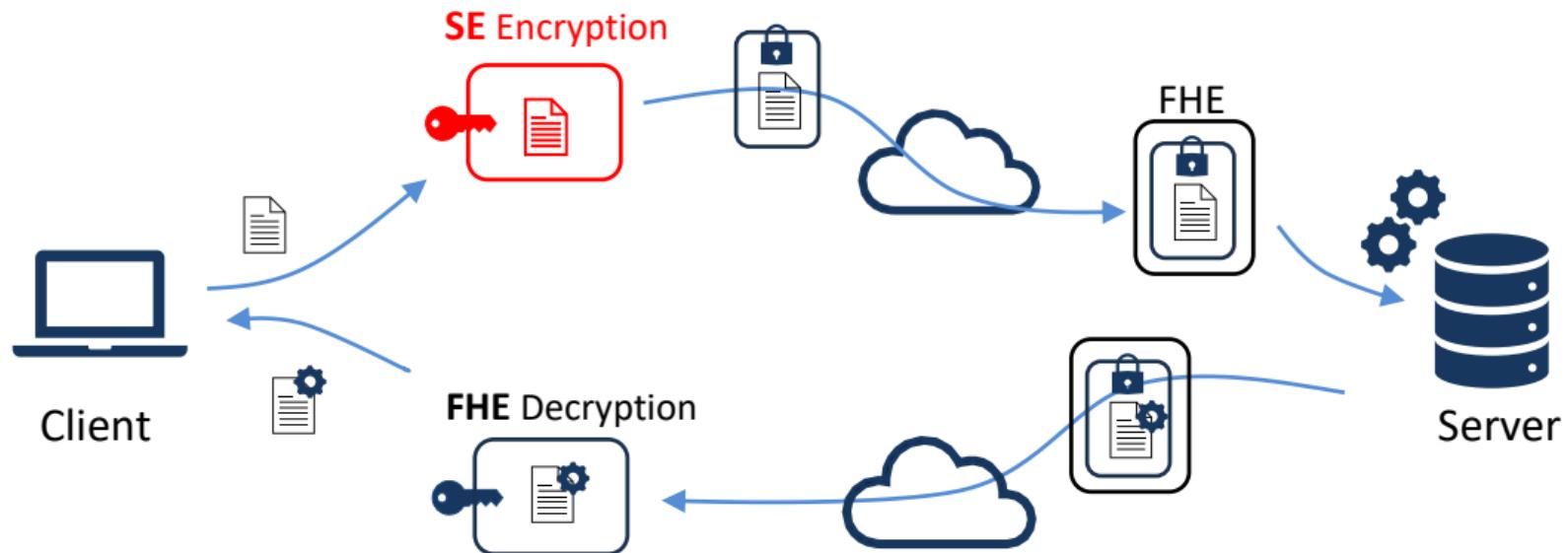
Homomorphic Encryption



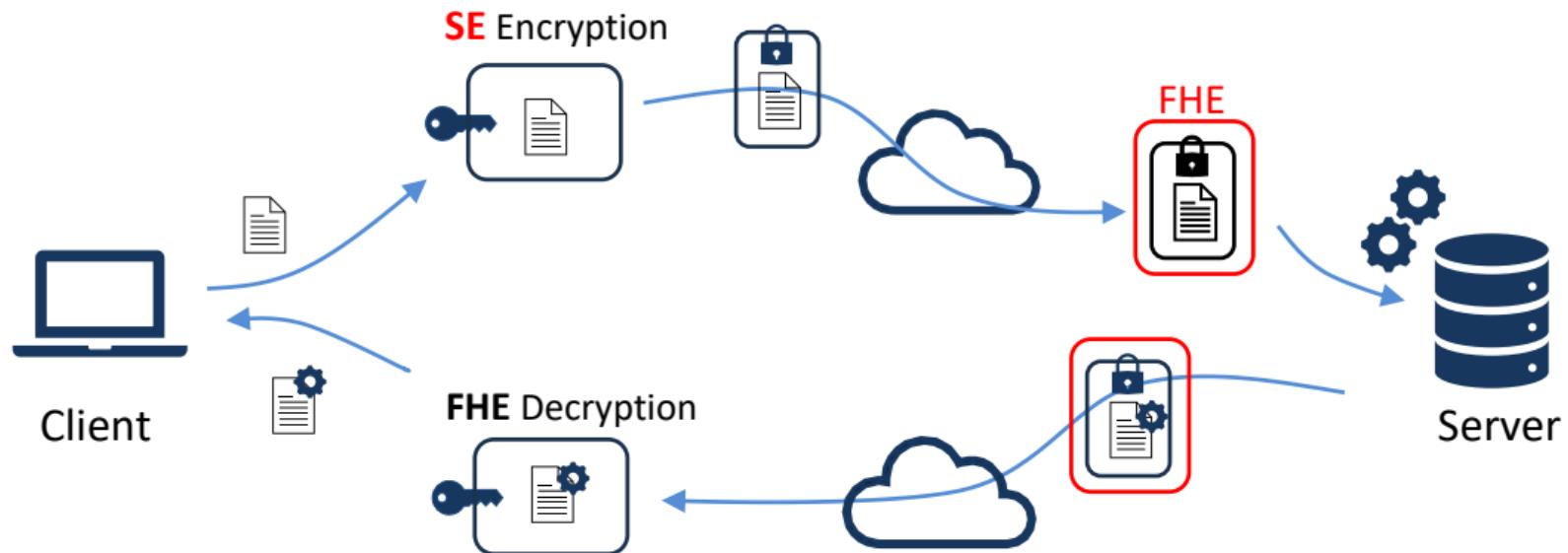
Hybrid Homomorphic Encryption



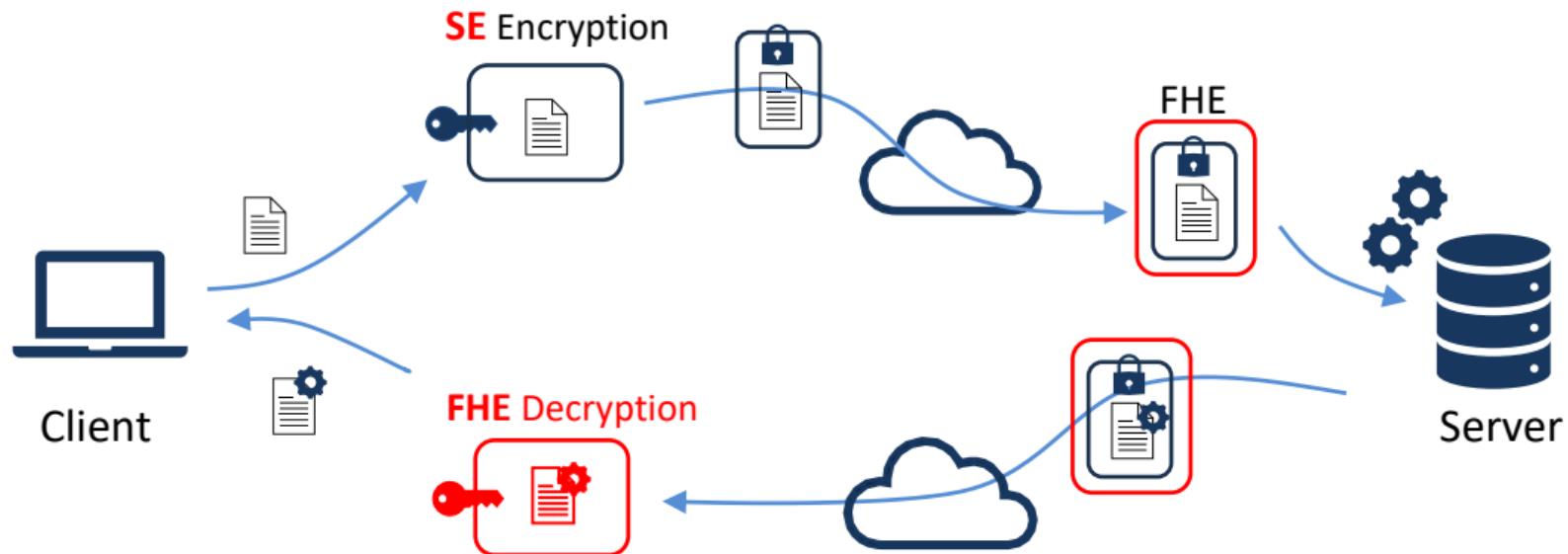
Hybrid Homomorphic Encryption



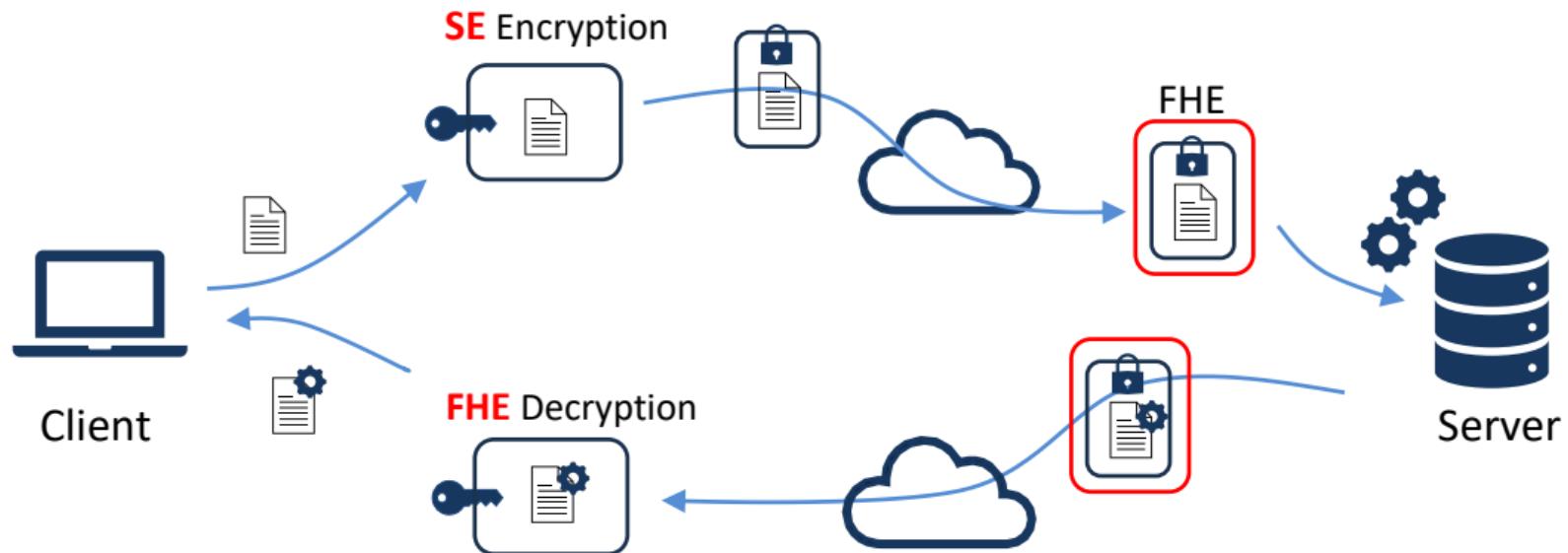
Hybrid Homomorphic Encryption



Hybrid Homomorphic Encryption



Hybrid Homomorphic Encryption



HERA Encryption

- i Prime-field arithmetic; 16 to 60 bits prime number p .
- i Generates a keystream from the input key K and adds it to the plaintext.
- i Typical **HERA** permutation consists out of 5 rounds.
- i State size: $n = 16$; hence state $X \in \mathbb{Z}_p^n$.

HERA Encryption

- i Prime-field arithmetic; 16 to 60 bits prime number p .
- i Generates a keystream from the input key K and adds it to the plaintext.
- i Typical HERA permutation consists out of 5 rounds.
- i State size: $n = 16$; hence state $X \in \mathbb{Z}_p^n$.

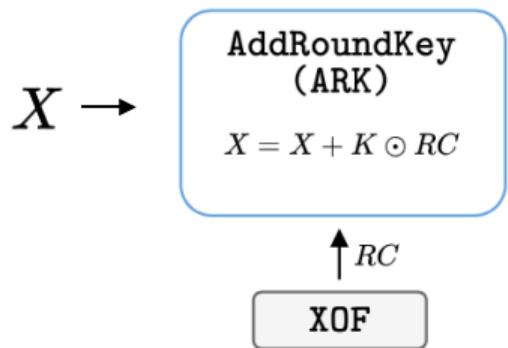
HERA Encryption

- i Prime-field arithmetic; 16 to 60 bits prime number p .
- i Generates a keystream from the input key K and adds it to the plaintext.
- i Typical **HERA** permutation consists out of 5 rounds.
- i State size: $n = 16$; hence state $X \in \mathbb{Z}_p^n$.

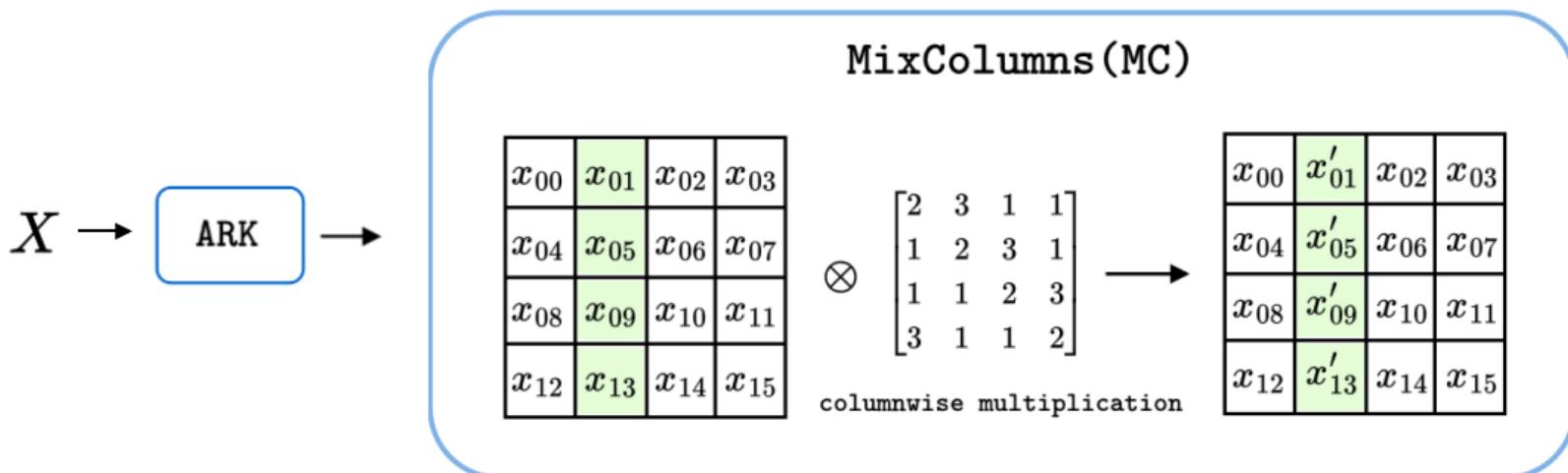
HERA Encryption

- ① Prime-field arithmetic; 16 to 60 bits prime number p .
- ① Generates a keystream from the input key K and adds it to the plaintext.
- ① Typical **HERA** permutation consists out of 5 rounds.
- ① State size: $n = 16$; hence state $X \in \mathbb{Z}_p^n$.

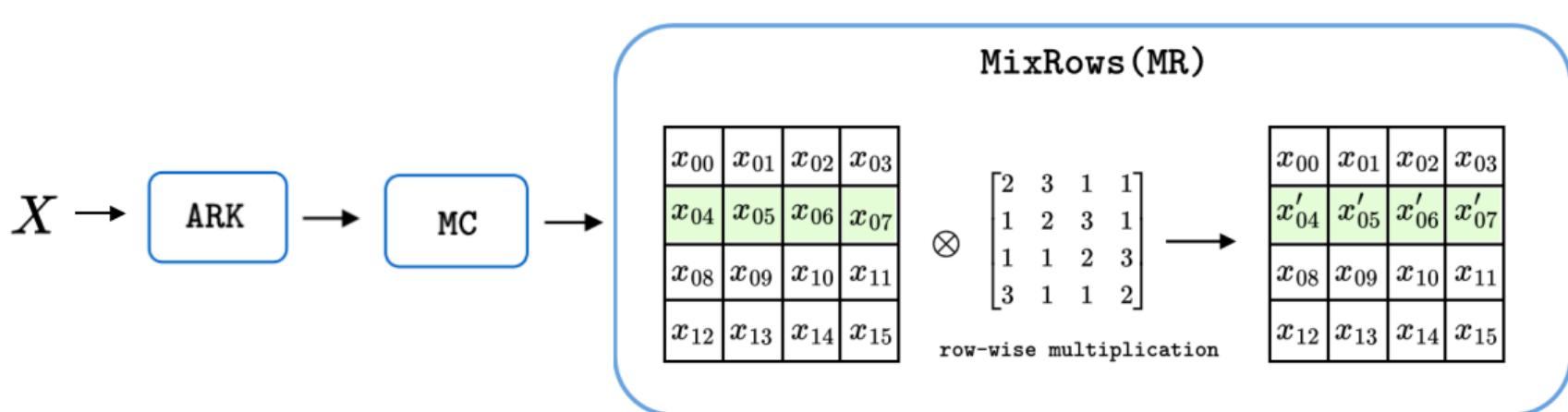
HERA Encryption



HERA Encryption



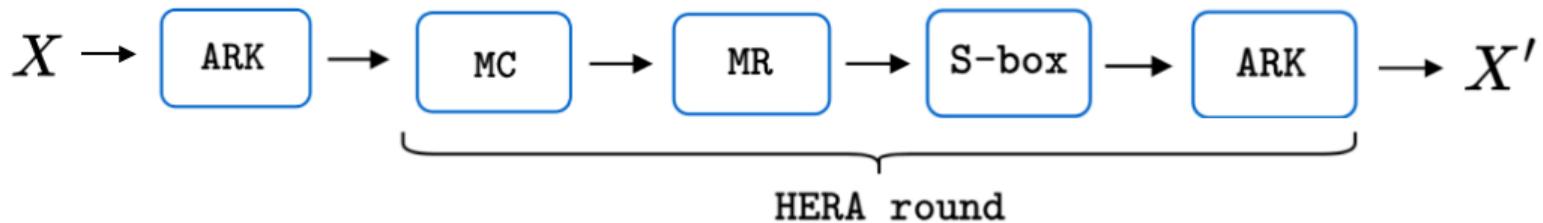
HERA Encryption



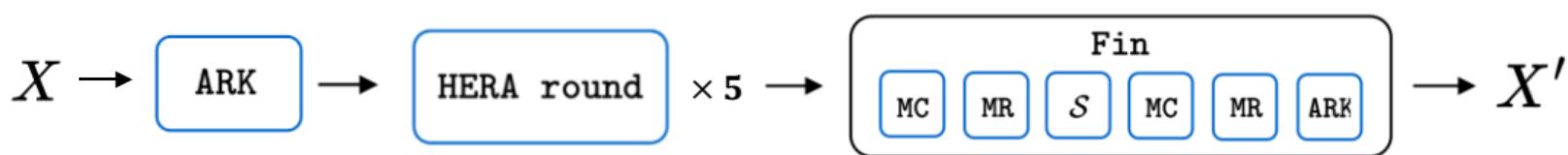
HERA Encryption



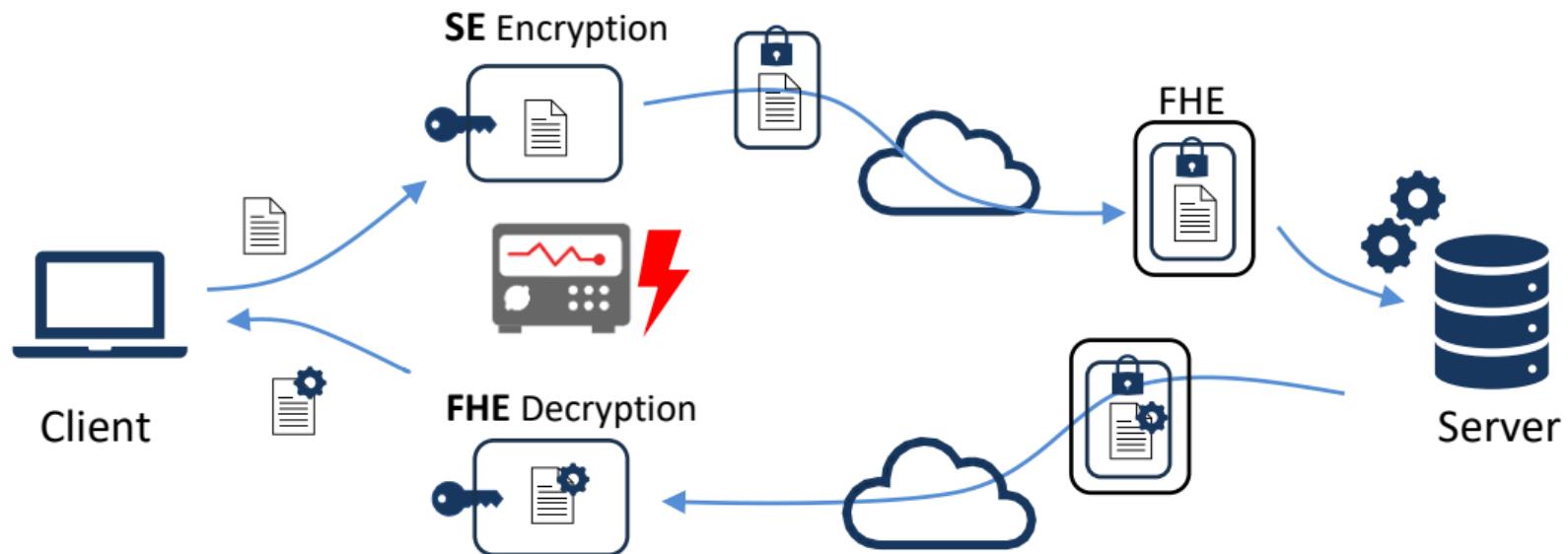
HERA Encryption



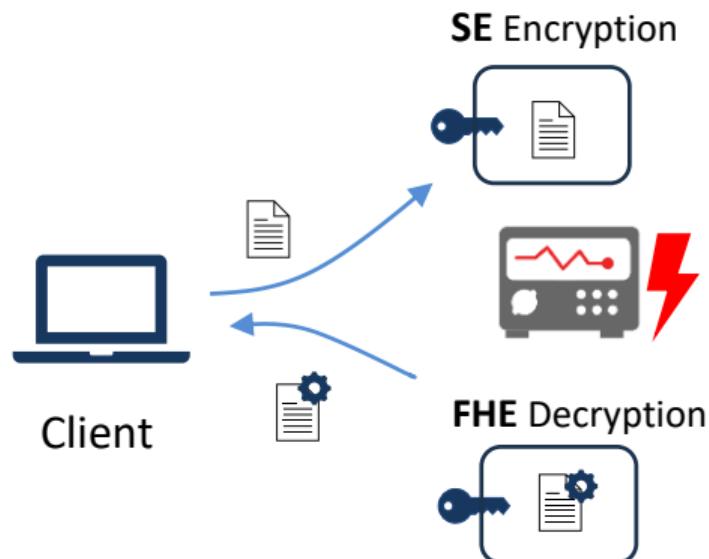
HERA Encryption



HERA Side-Channels Potential



HERA Side-Channels Potential

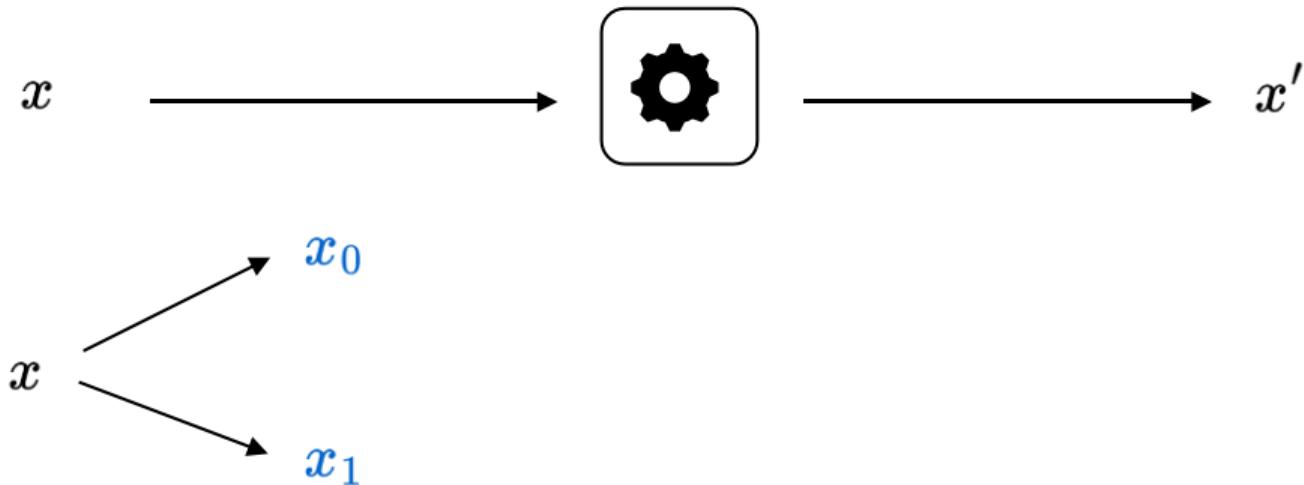


- Exploit relation between power consumption and processed data.
- Statistical methods to guess secret key bits - **DPA/CPA** possible.
- Modular arithmetics spreads side-channel leakage.

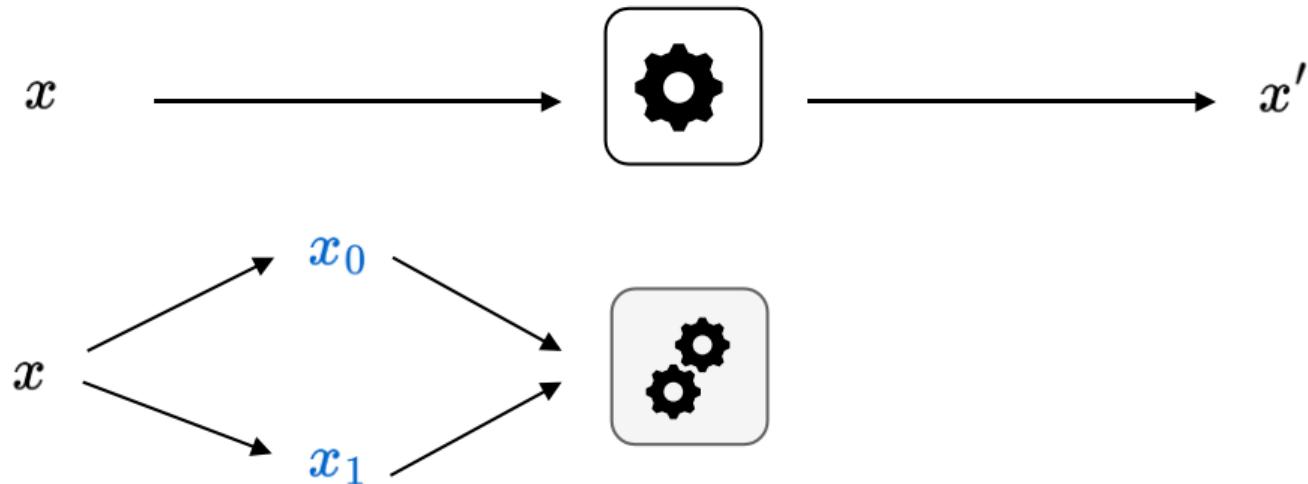
First Order Masking



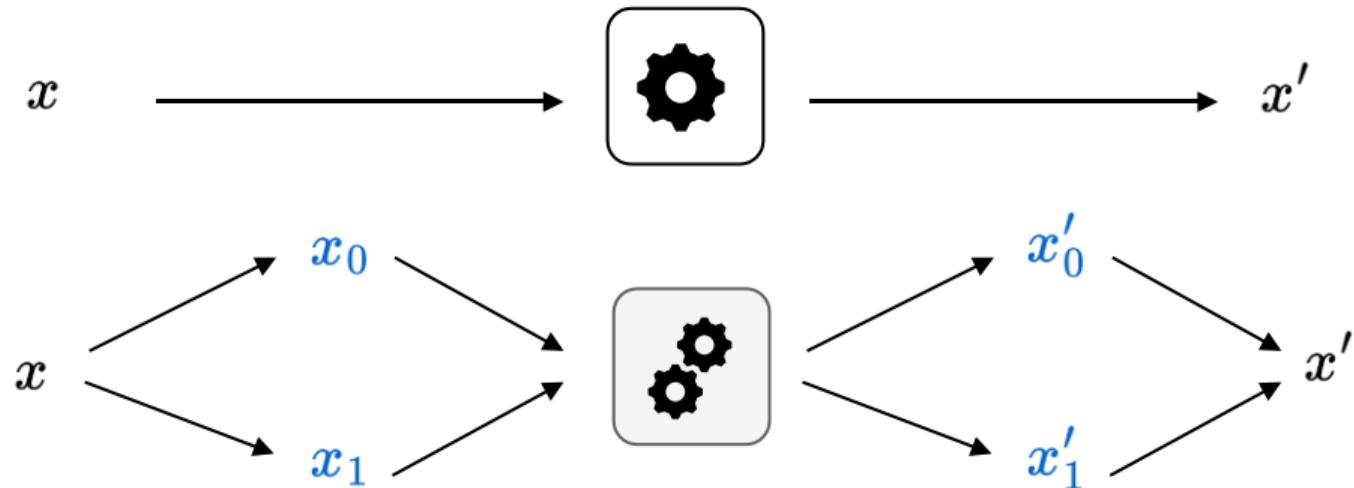
First Order Masking



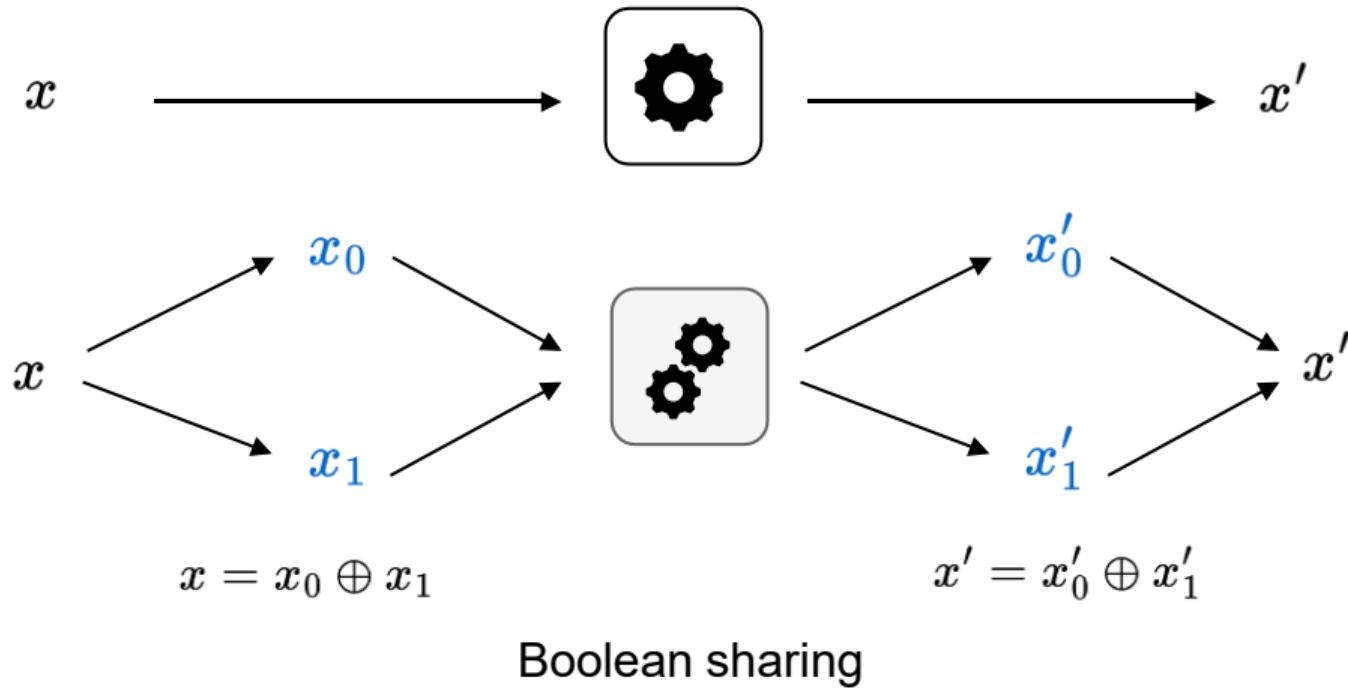
First Order Masking



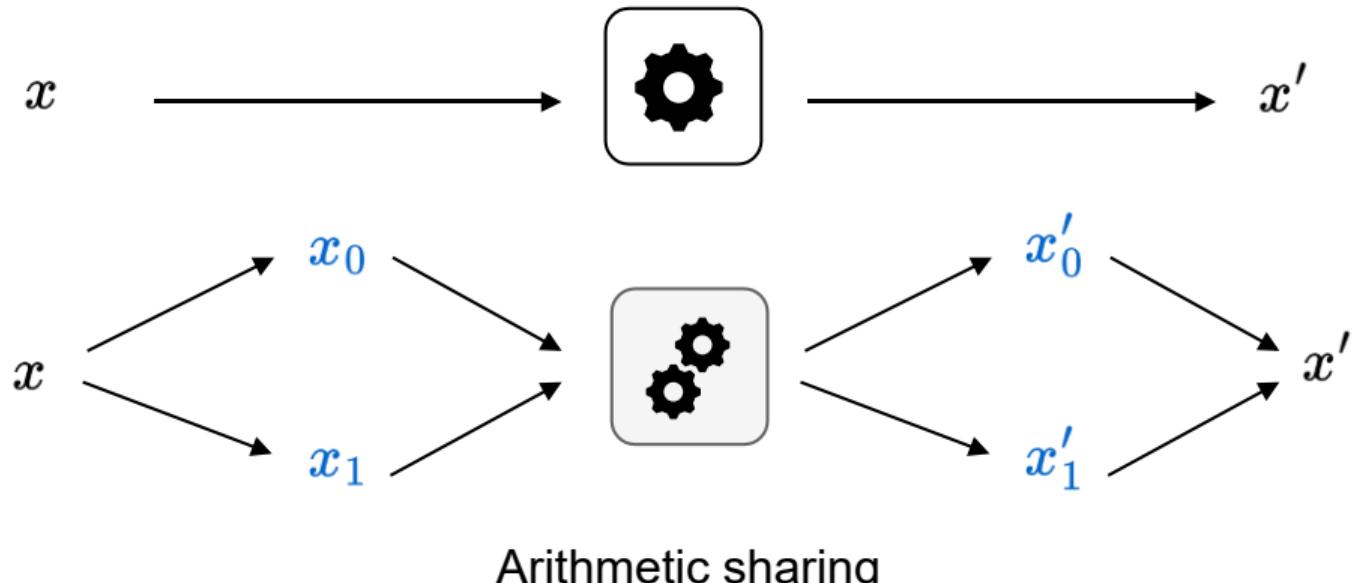
First Order Masking



First Order Masking



First Order Masking



$$x = x_0 + x_1$$

$$x' = x'_0 + x'_1$$

Test Vector Leakage Assessment (TVLA)

- ① Standard statistical method to detect **side-channel leakage**.
- ① Compares device traces for **fixed vs random inputs**.
- ① Applies Welch's t-test at each time point.
- ① Evidence of leakage:
 $|t| > 4.5$.

Test Vector Leakage Assessment (TVLA)

- ① Standard statistical method to detect **side-channel leakage**.
- ① Compares device traces for **fixed vs random inputs**.
- ① Applies Welch's t-test at each time point.
- ① Evidence of leakage:
 $|t| > 4.5$.

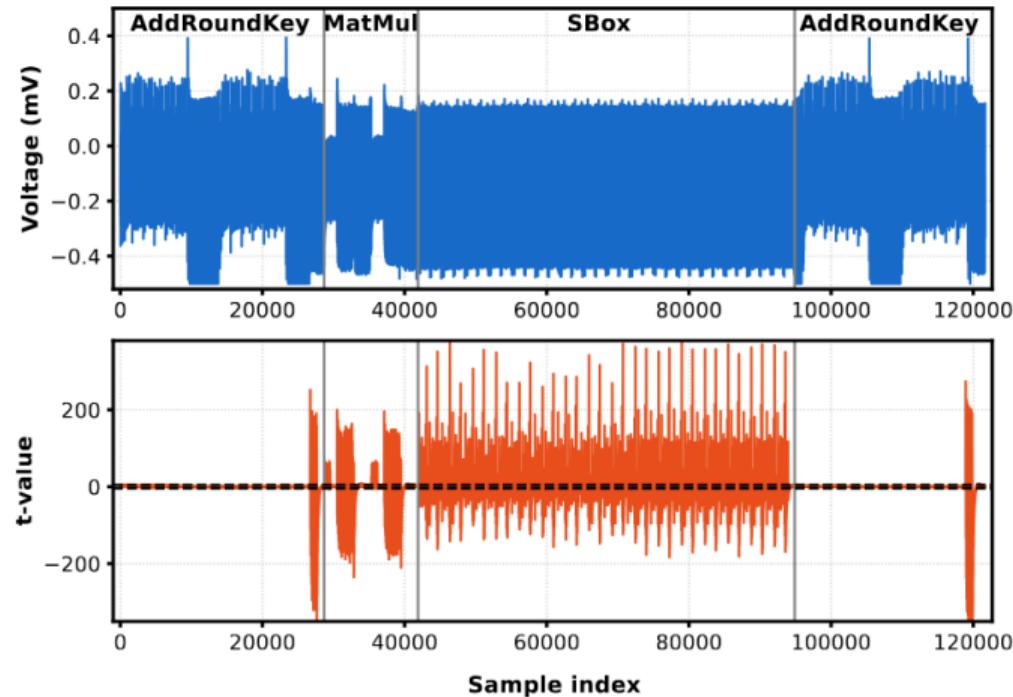
Test Vector Leakage Assessment (TVLA)

- ① Standard statistical method to detect **side-channel leakage**.
- ① Compares device traces for **fixed vs random inputs**.
- ① Applies **Welch's t-test** at each time point.
- ① Evidence of leakage:
 $|t| > 4.5$.

Test Vector Leakage Assessment (TVLA)

- ① Standard statistical method to detect **side-channel leakage**.
- ① Compares device traces for **fixed vs random inputs**.
- ① Applies **Welch's t-test** at each time point.
- ① Evidence of leakage:
 $|t| > 4.5$.

TVLA Results for HERA



Masking

Masking a HERA round

$$X = X_0 + X_1 \pmod{p}$$

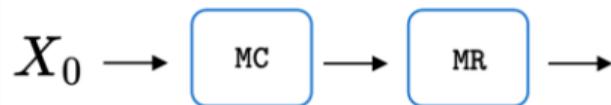
$X_0 \rightarrow$

$X_1 \rightarrow$

Masking

Masking a HERA round

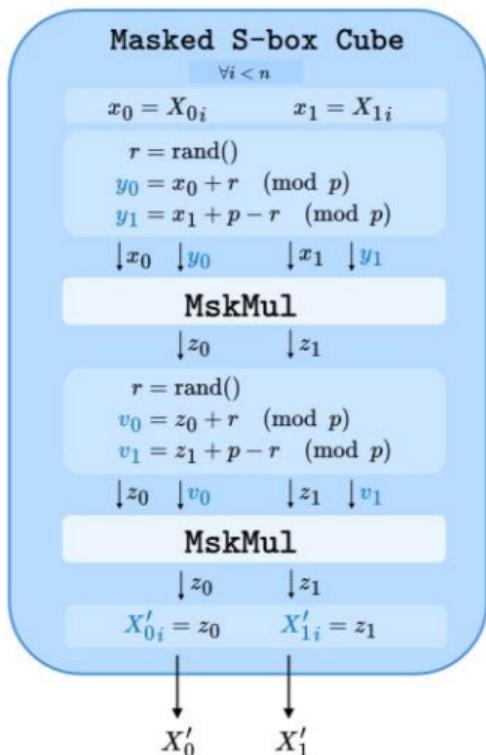
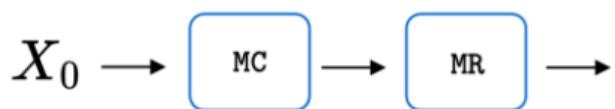
$$X = X_0 + X_1 \pmod{p}$$



Masking

Masking a HERA round

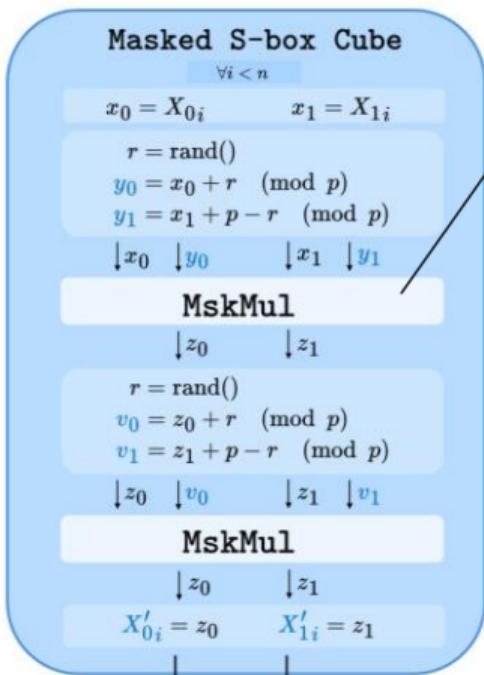
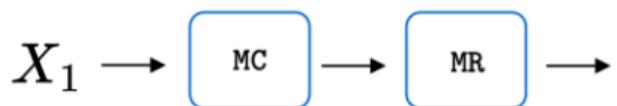
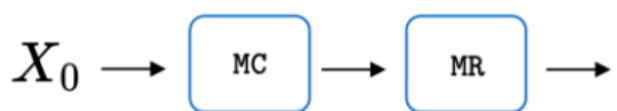
$$X = X_0 + X_1 \pmod{p}$$



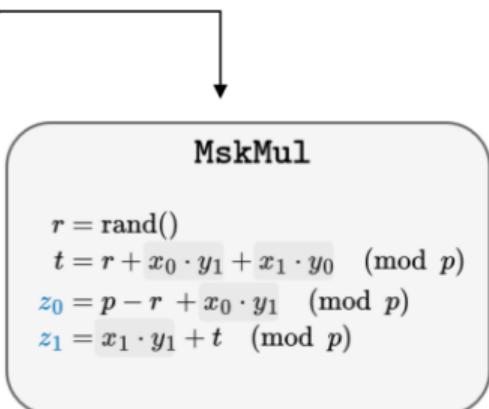
Masking

Masking a HERA round

$$X = X_0 + X_1 \pmod{p}$$



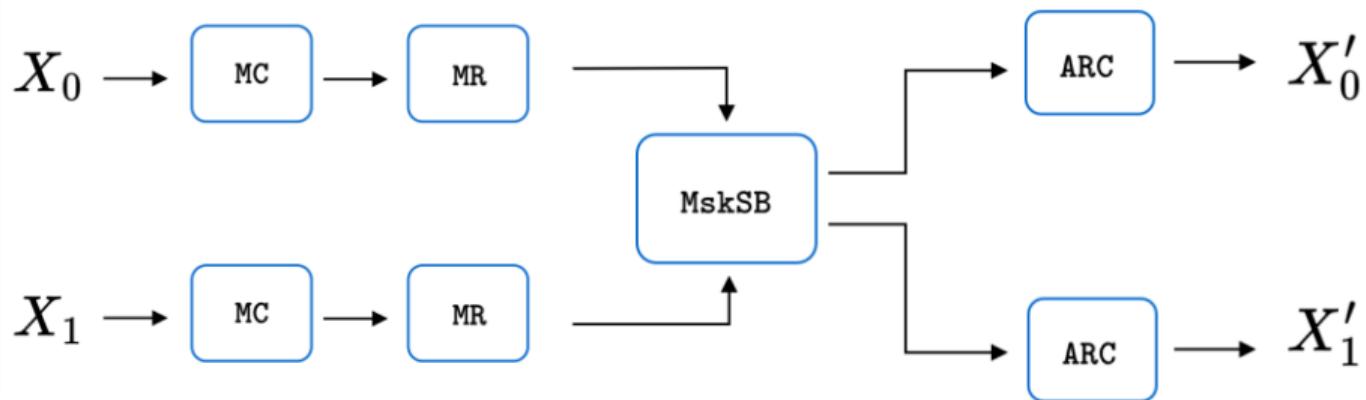
X'_0 X'_1



Masking

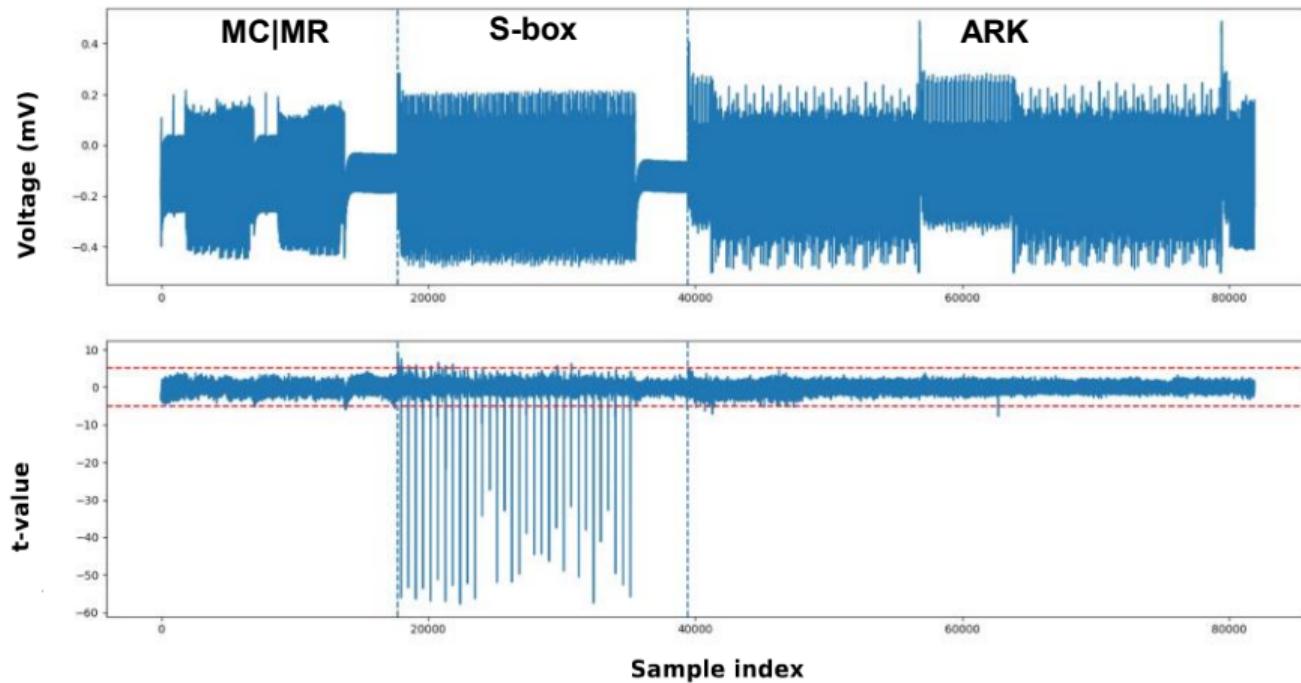
Masking a HERA round

$$X = X_0 + X_1 \pmod{p}$$



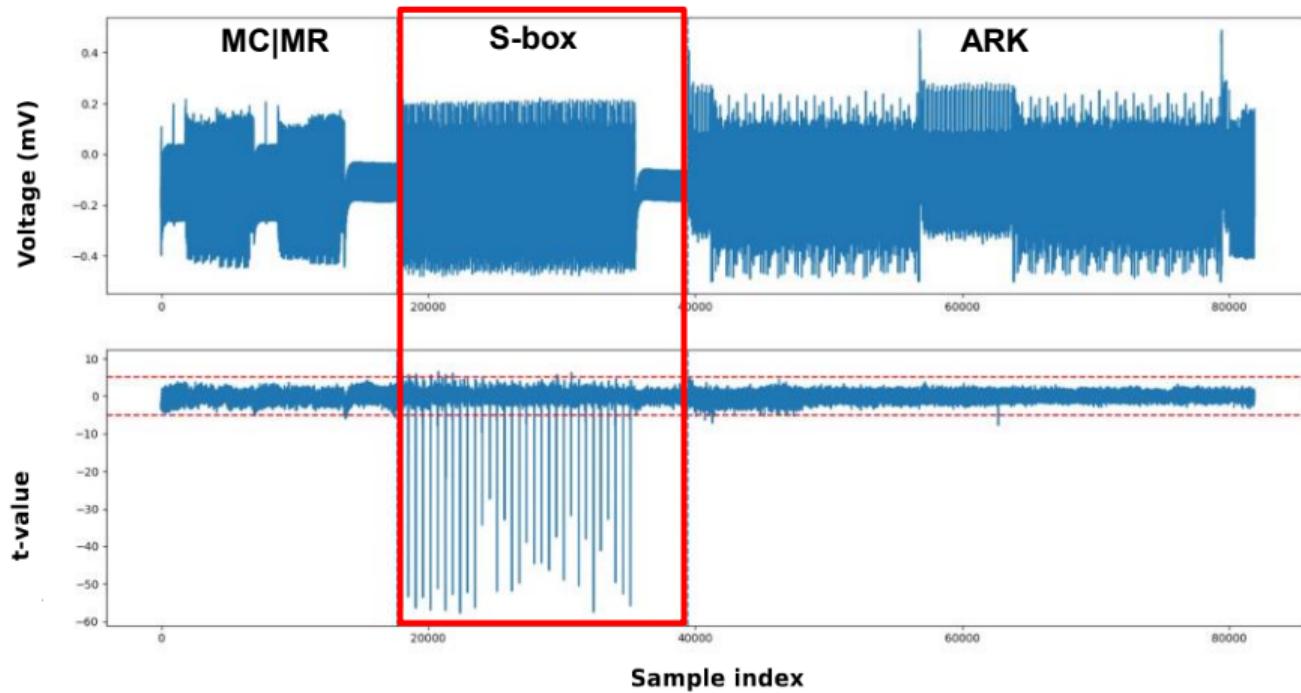
Preliminary Results

TVLA Results



Preliminary Results

TVLA Results



Implementation Heuristic

- ① Ensure the random source is random indeed.
- ② Separate shares usage within the C code.
- ③ Employ assembly code for full register control.
- ④ Address critical points which could lead to the statistical leakage despite the masked model.

Implementation Heuristic

- ① Ensure the random source is random indeed.
- ① Separate shares usage within the C code.
- ① Employ assembly code for full register control.
- ① Address critical points which could lead to the statistical leakage despite the masked model.

Implementation Heuristic

- ① Ensure the random source is random indeed.
- ① Separate shares usage within the C code.
- ① Employ assembly code for full register control.
- ① Address critical points which could lead to the statistical leakage despite the masked model.

Implementation Heuristic

- ① Ensure the random source is random indeed.
- ① Separate shares usage within the C code.
- ① Employ assembly code for full register control.
- ① Address critical points which could lead to the statistical leakage despite the masked model.

Implementation Heuristic

A large red 'X' symbol.A large green checkmark symbol. $r_4 \leftarrow x_0$ $r_4 \leftarrow x_1$

Implementation Heuristic

X

✓

$$\left. \begin{array}{l} r_4 \leftarrow x_0 \\ r_4 \leftarrow x_1 \end{array} \right\} \begin{array}{l} \text{HD}(x_0, x_1) = \text{HW}(x_0 \oplus x_1) = \text{HW}(x) \\ \text{leakage through Hamming distance} \end{array}$$

Implementation Heuristic

X

$$\left. \begin{array}{l} r_4 \leftarrow x_0 \\ r_4 \leftarrow x_1 \end{array} \right\} \text{HD}(x_0, x_1) = \text{HW}(x_0 \oplus x_1) = \text{HW}(x)$$

leakage through Hamming distance

$$\left. \right\} \begin{array}{l} r_3 \leftarrow x_0 \\ r_4 \leftarrow x_1 \end{array}$$

✓

Implementation Heuristic

X

✓

→
unknown $r_3 = ?$

$r_3 \leftarrow x_0$

Implementation Heuristic

X

:

✓

$$\left. \begin{array}{l} \xrightarrow{\text{unknown}} r_3 = ? \\ r_3 \leftarrow x_0 \end{array} \right\} \begin{array}{l} \text{HD}(x_0, x_1) = \text{HW}(x_0 \oplus x_1) = \text{HW}(x) \\ \text{potential leakage through Hamming distance} \end{array}$$

Implementation Heuristic

X

:

✓

$$\left. \begin{array}{l} \xrightarrow{\text{unknown}} r_3 = ? \\ r_3 \leftarrow x_0 \end{array} \right\} \begin{array}{l} \text{HD}(x_0, x_1) = \text{HW}(x_0 \oplus x_1) = \text{HW}(x) \\ \text{potential leakage through Hamming distance} \end{array} \right\} \begin{array}{l} r_3 \leftarrow \text{rand} \\ r_3 \leftarrow x_0 \end{array}$$

Implementation Heuristic

X

:

✓

$$\left. \begin{array}{l} r_3 \leftarrow x \\ \xrightarrow{\text{unknown}} r_3 \leftarrow ? \end{array} \right\} \begin{array}{l} \text{HD}(x_0, x_1) = \text{HW}(x_0 \oplus x_1) = \text{HW}(x) \\ \text{potential leakage through Hamming distance} \end{array} \right\} \begin{array}{l} r_3 \leftarrow x_0 \\ r_3 \leftarrow \text{rand} \end{array}$$

Implementation Heuristic

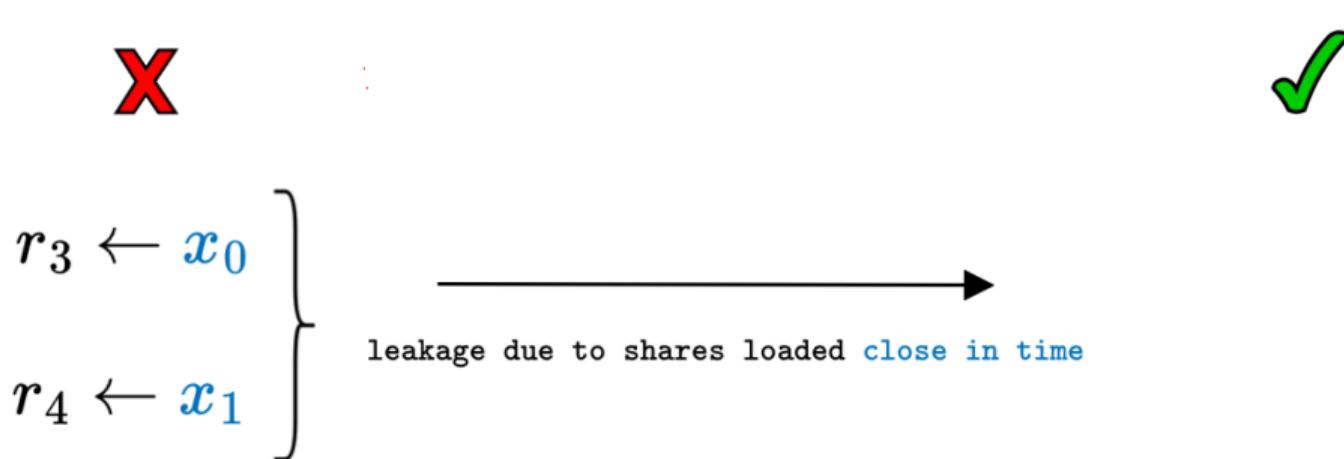
X

:

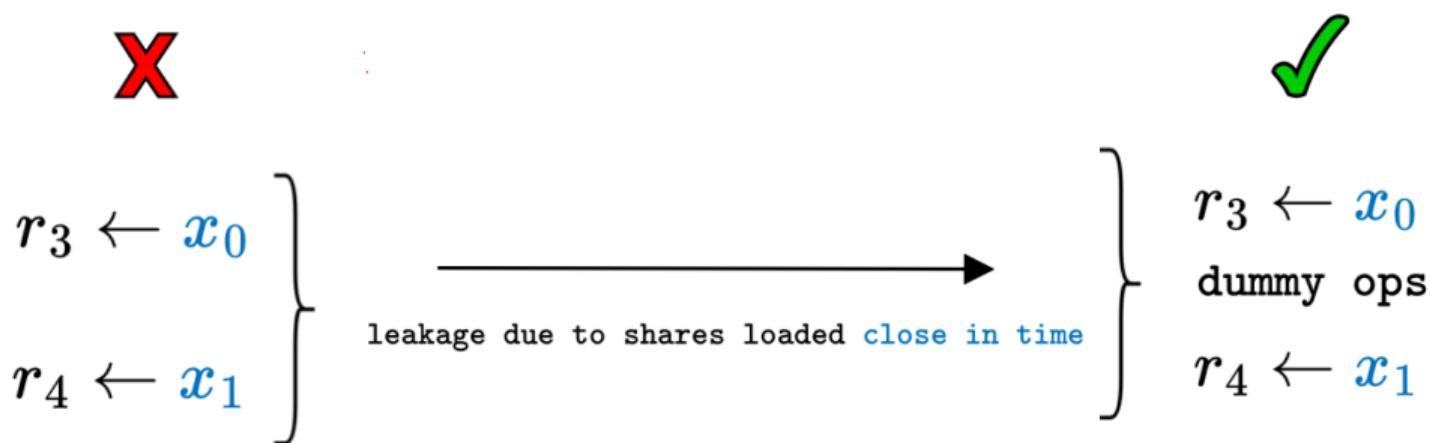
✓

$$\left. \begin{array}{l} r_3 \leftarrow x_0 \\ r_4 \leftarrow x_1 \end{array} \right\}$$

Implementation Heuristic



Implementation Heuristic



Implementation Heuristic



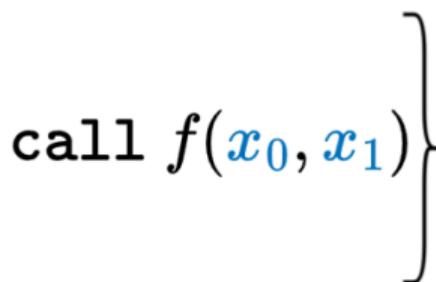
call $f(x_0, x_1)$

Implementation Heuristic

X

:

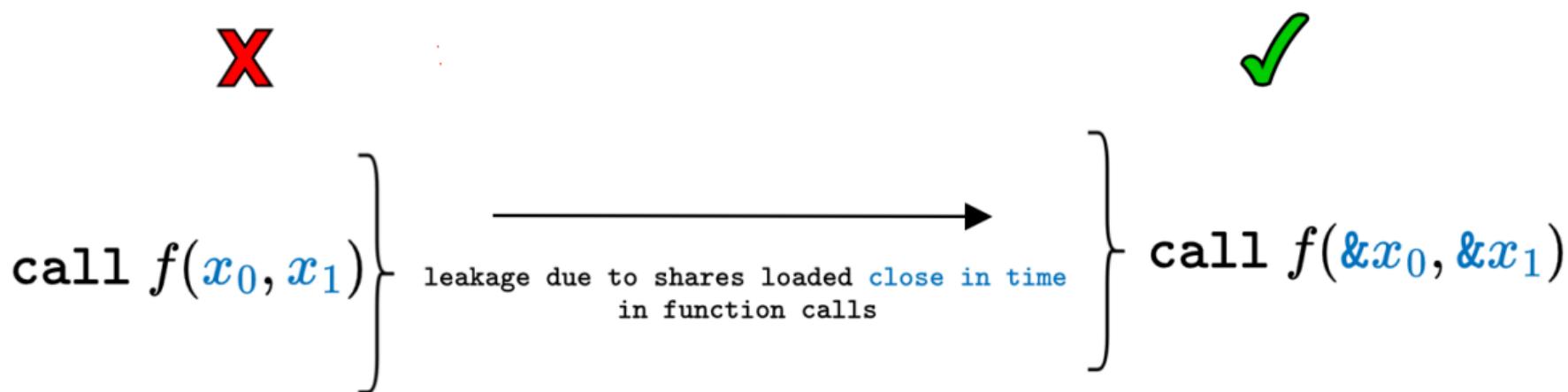
call $f(x_0, x_1)$



leakage due to shares loaded close in time
in function calls

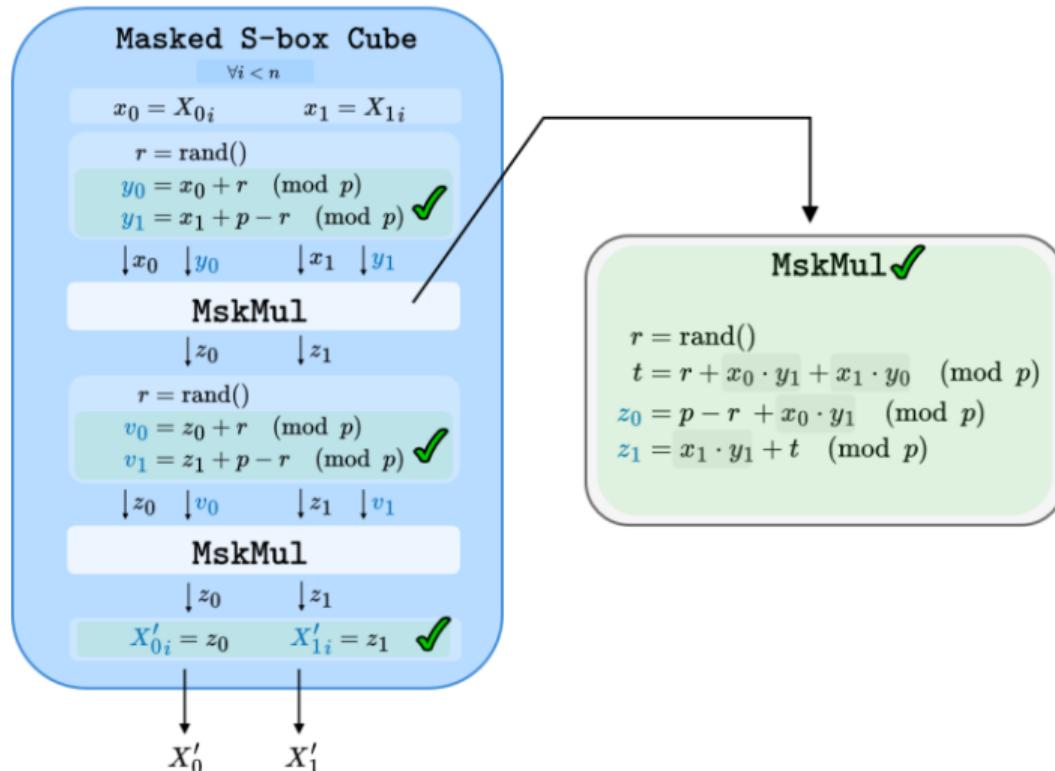


Implementation Heuristic



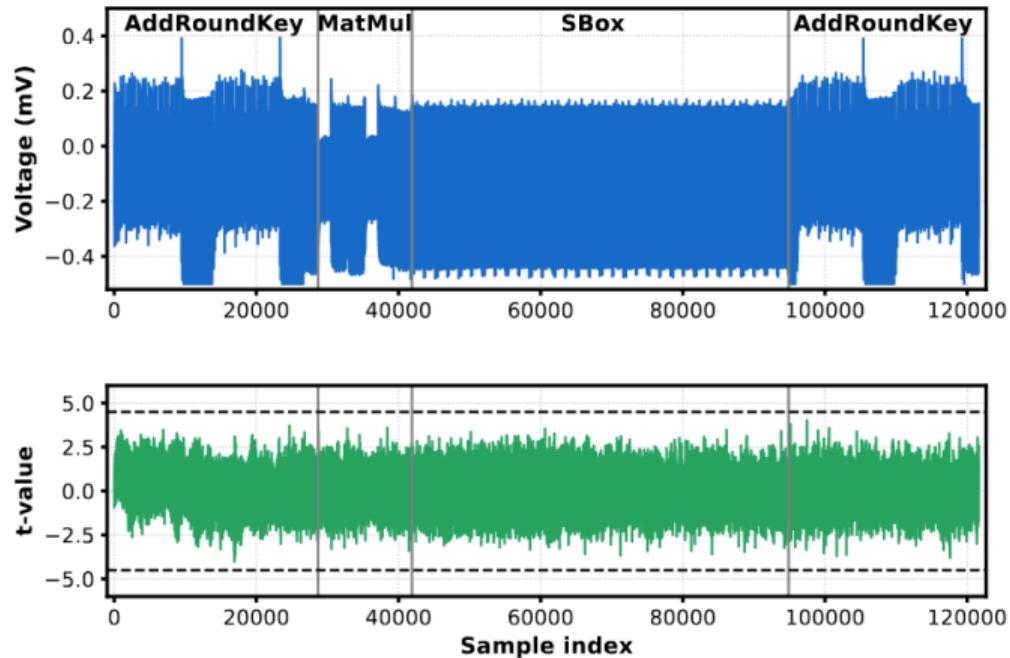
Implementation

Application to HERA S-box



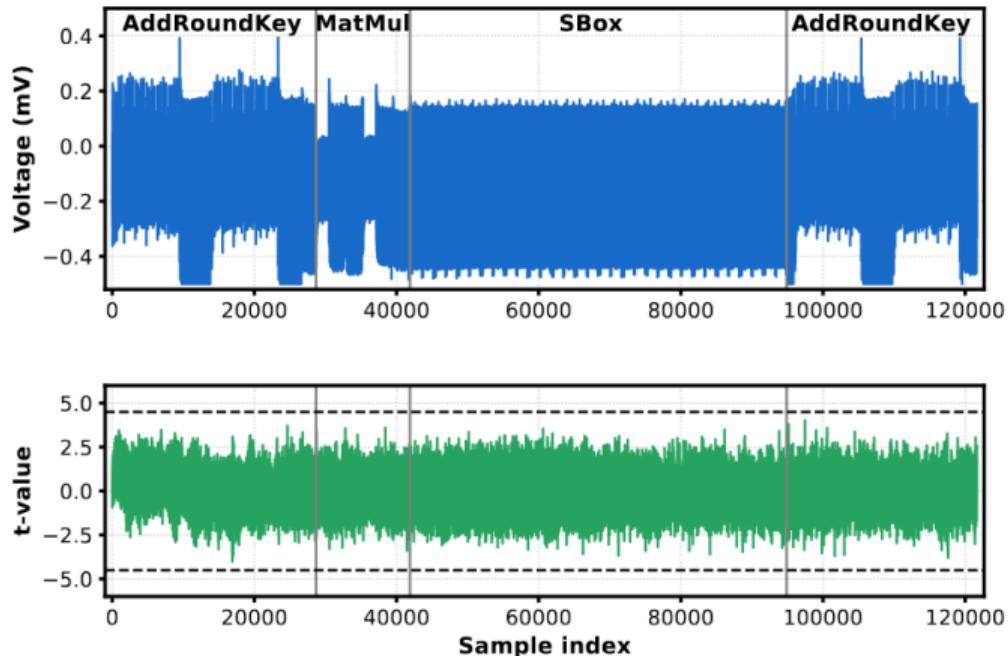
HERA Results

TVLA Results

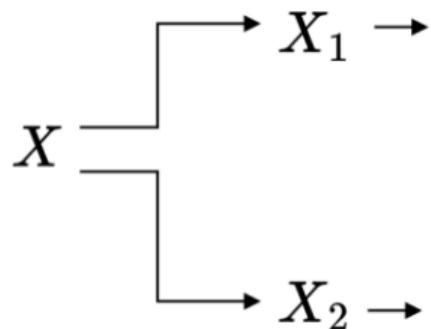


HERA Results

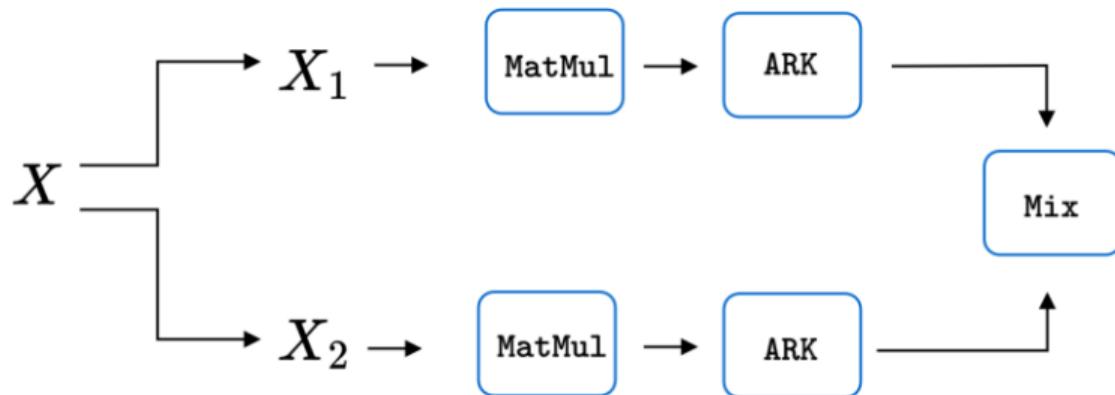
TVLA Results



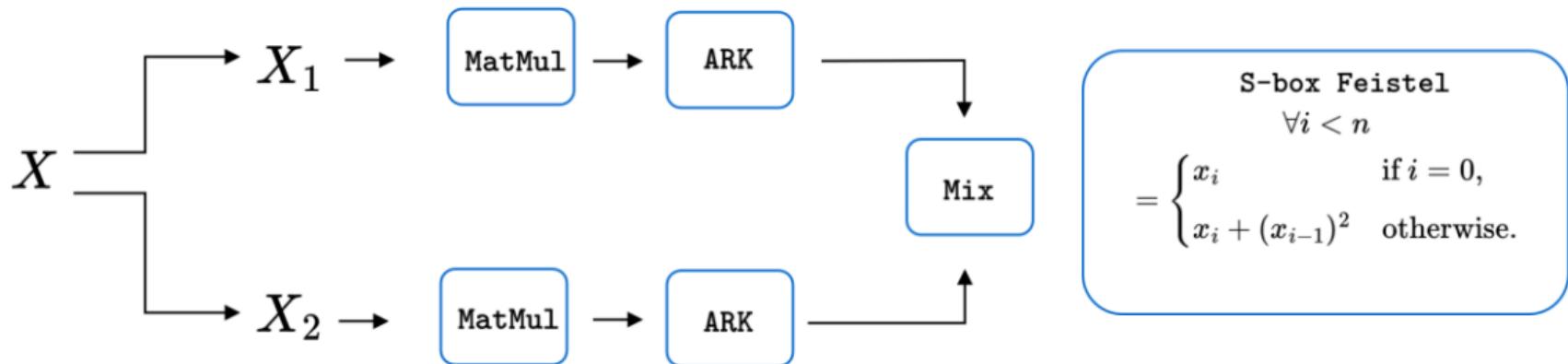
PASTA Encryption



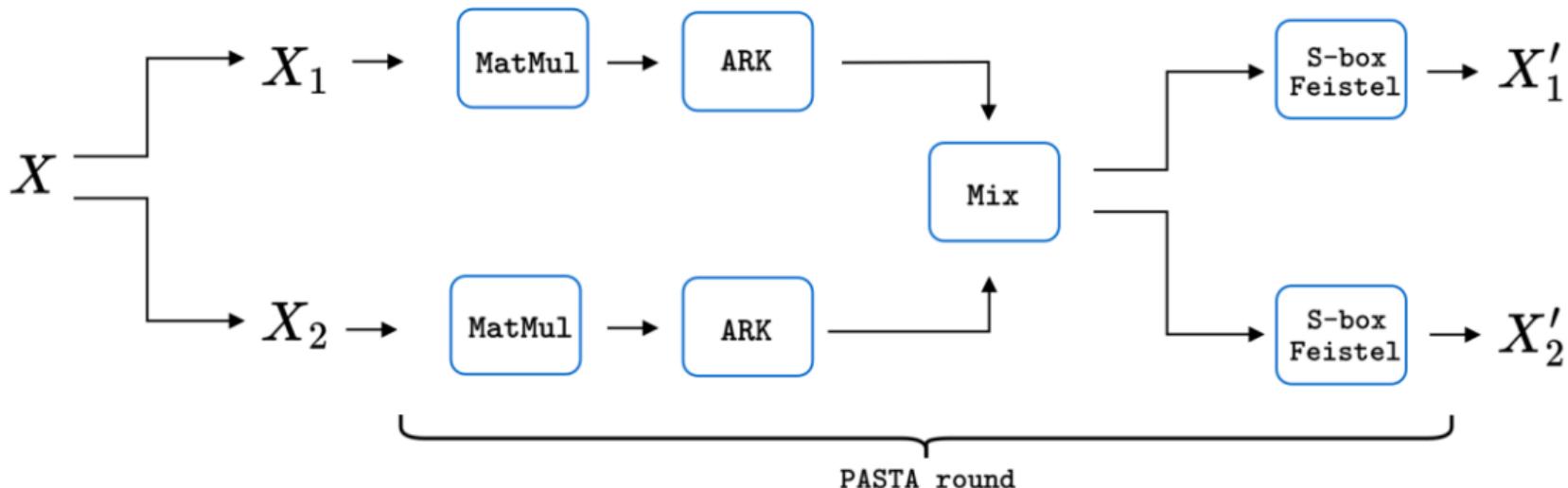
PASTA Encryption



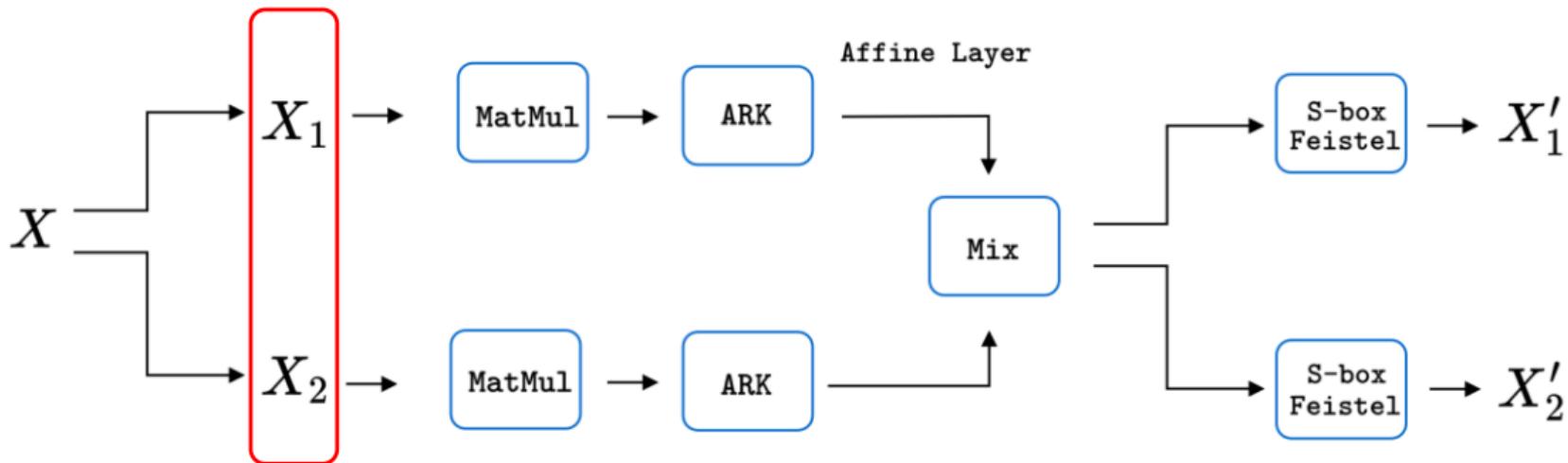
PASTA Encryption



PASTA Encryption

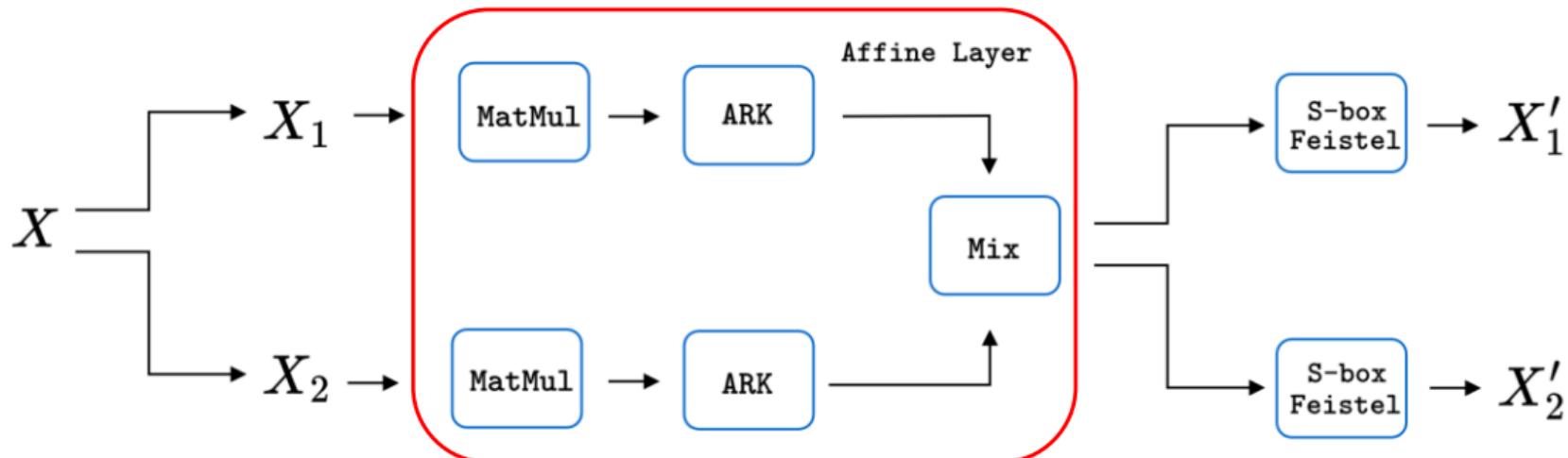


Masking PASTA Round



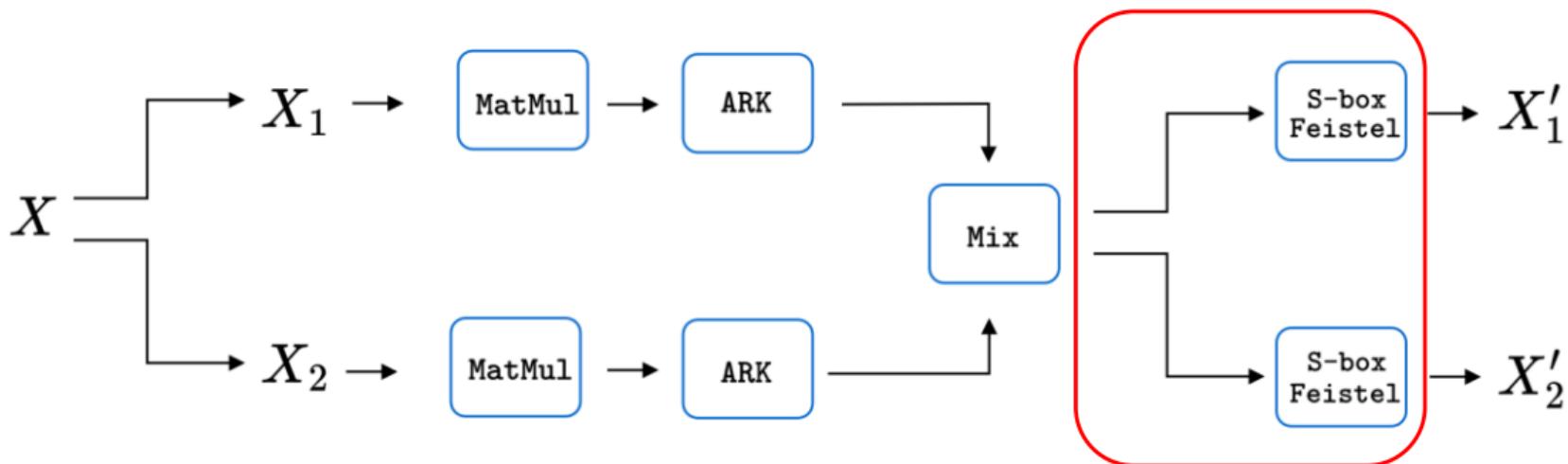
- ✓ mask both parts

Masking PASTA Round



- ✓ apply linear layers on each share

Masking PASTA Round



✓ Use secure multiplication

Masking PASTA Round

Masked S-box Feistel

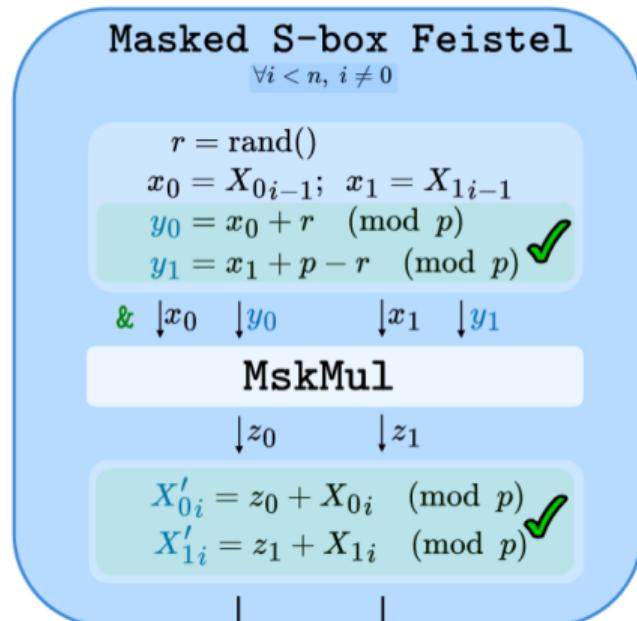
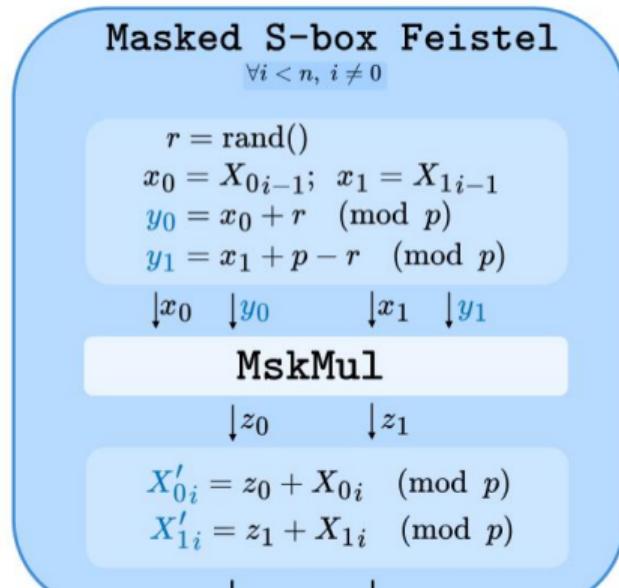
 $\forall i < n, i \neq 0$ $r = \text{rand}()$ $x_0 = X_{0i-1}; x_1 = X_{1i-1}$ $y_0 = x_0 + r \pmod{p}$ $y_1 = x_1 + p - r \pmod{p}$ $\downarrow x_0 \quad \downarrow y_0 \quad \downarrow x_1 \quad \downarrow y_1$

MskMul

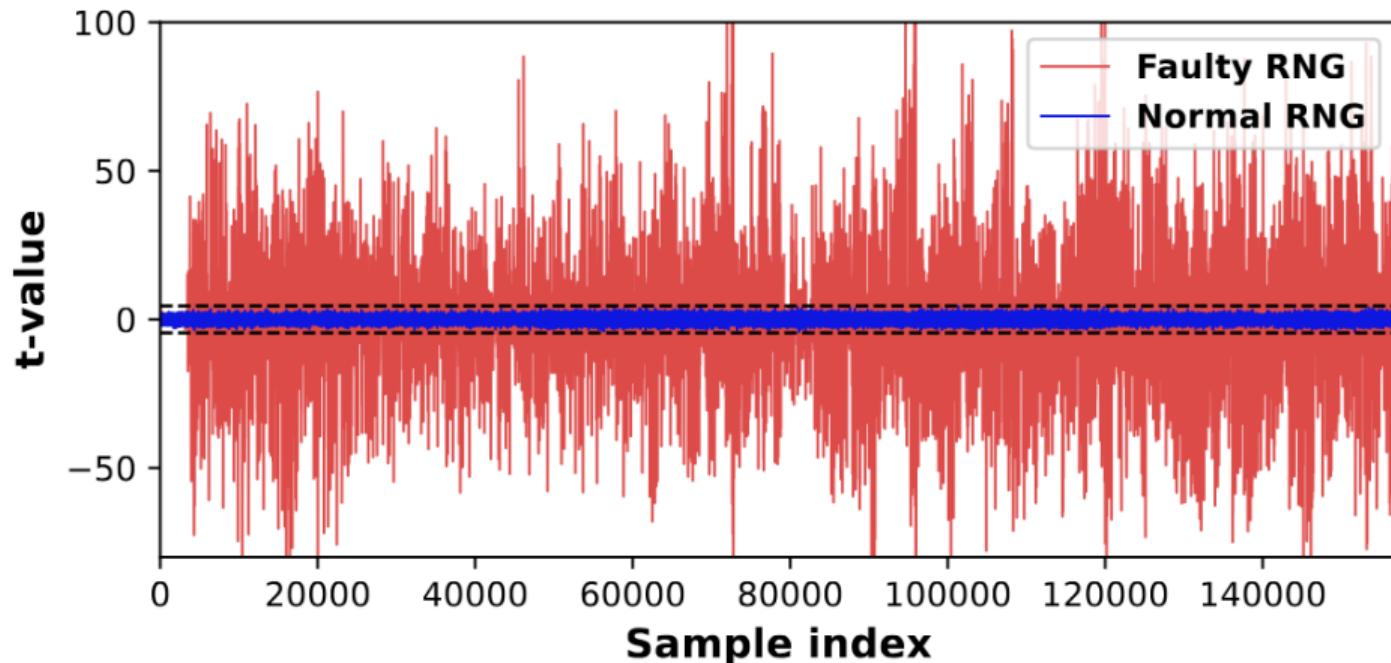
 $\downarrow z_0 \quad \downarrow z_1$ $X'_{0i} = z_0 + X_{0i} \pmod{p}$ $X'_{1i} = z_1 + X_{1i} \pmod{p}$ $\downarrow \quad \downarrow$
 $X'_0 \quad X'_1$

Application to Similar Encryption Schemes

Masking PASTA Round



TVLA Results



Higher Order Masking

- ① First order masking protects only against single-probe attacks.
- ① Extend with **higher order masking** for more security, by expanding the MskMul accordingly.

Conclusion

- ① HERA vulnerable to side-channel attacks.
- ① Correct masking model can leak due to implementation specificities.
- ① The proposed approach application concretely verified on PASTA.
- ① First order masking protects only against single-probe attacks.