

First Order Masking for Hybrid Homomorphic Encryption Schemes

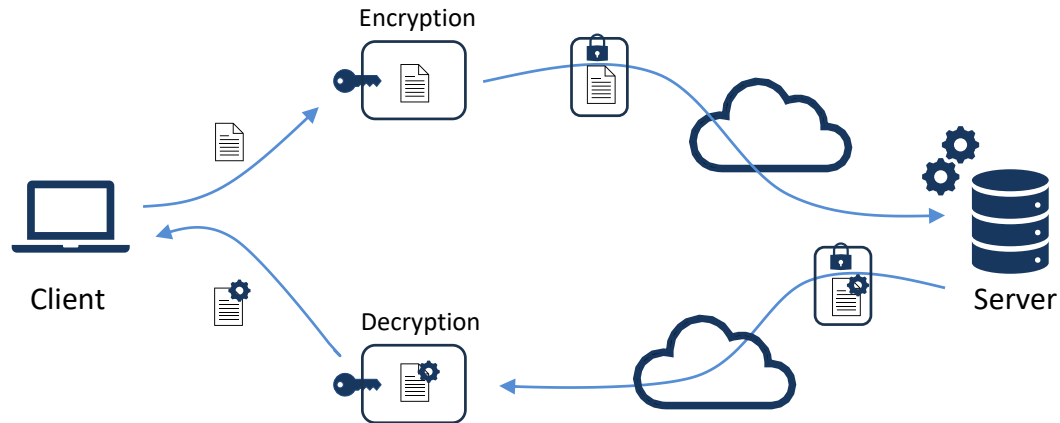
Nedžma Mušović

Master Project Presentation

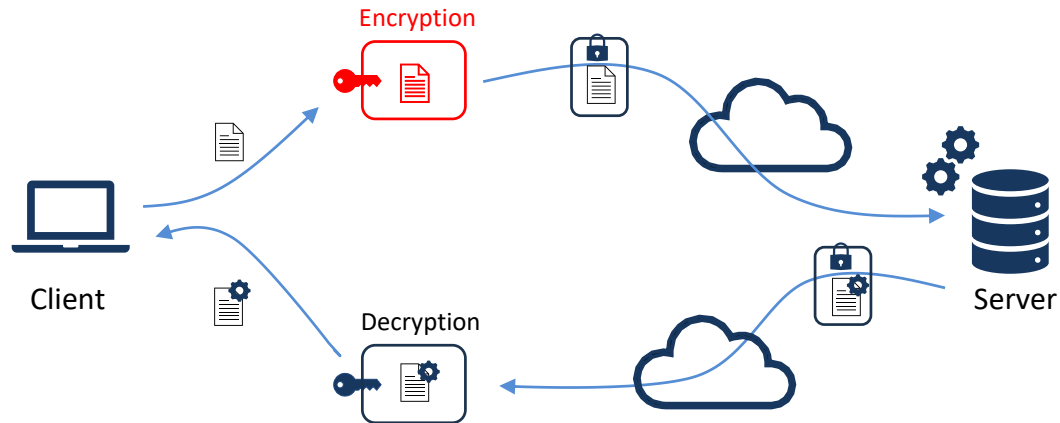
Outline

- i Introduction
- i **HERA** - First Order Masking
- i Application to **PASTA**
- i Higher Order Masking and Conclusion

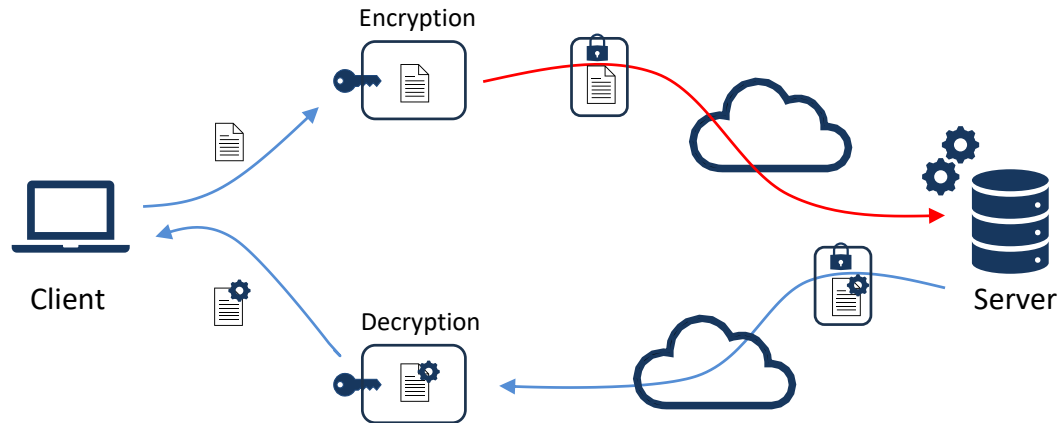
Homomorphic Encryption



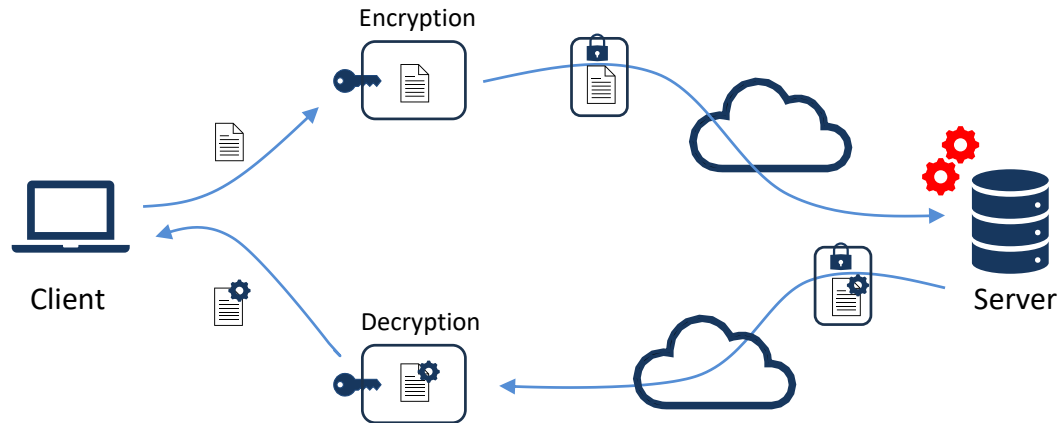
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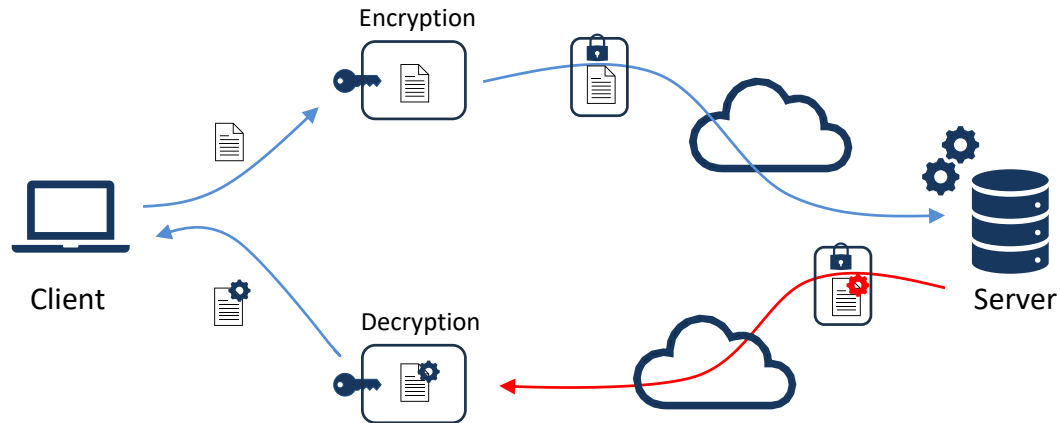
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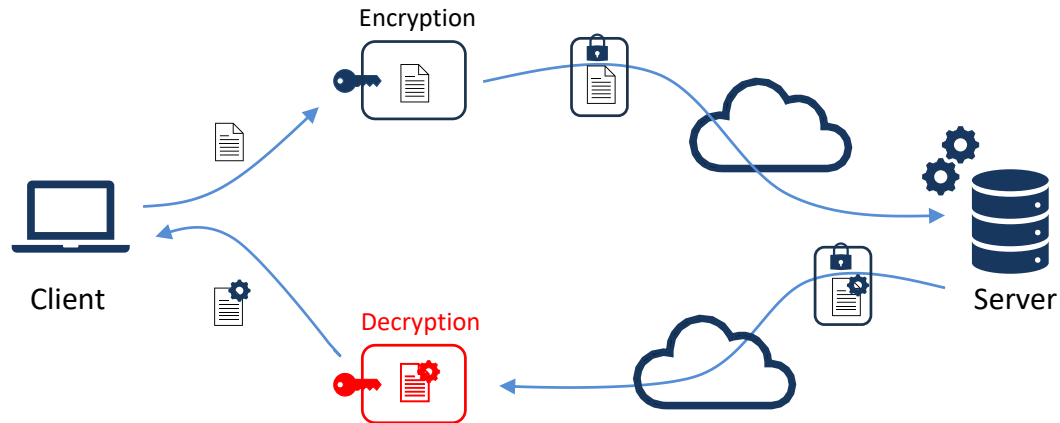
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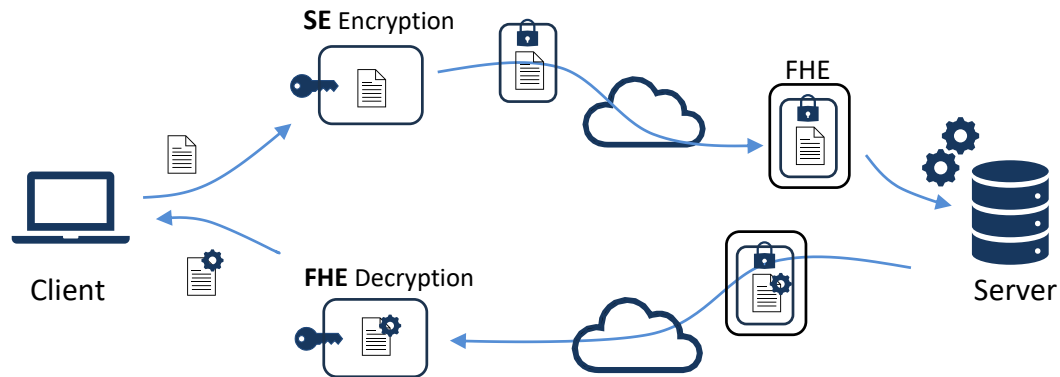
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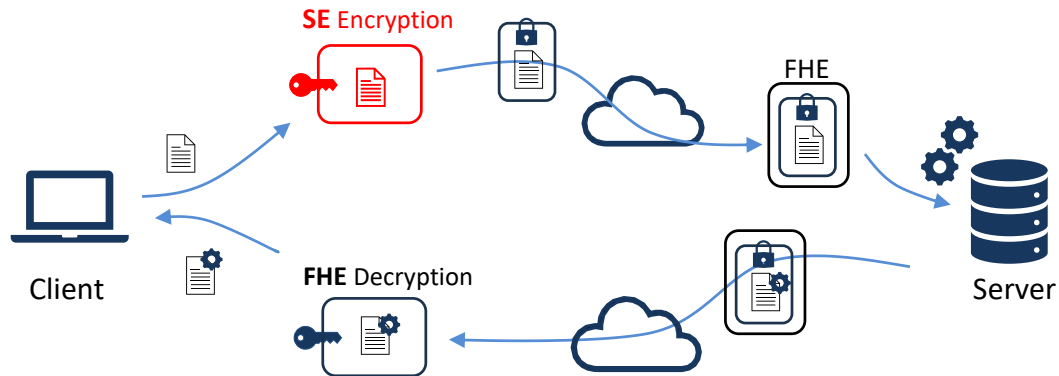
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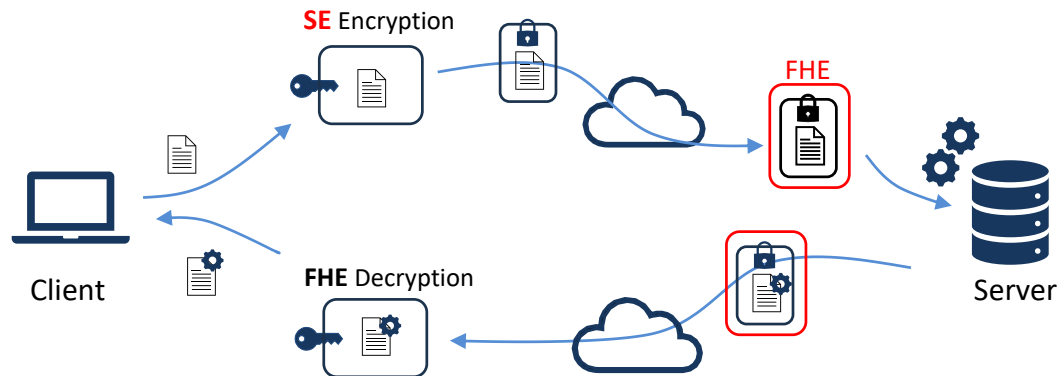
Hybrid Homomorphic Encryption



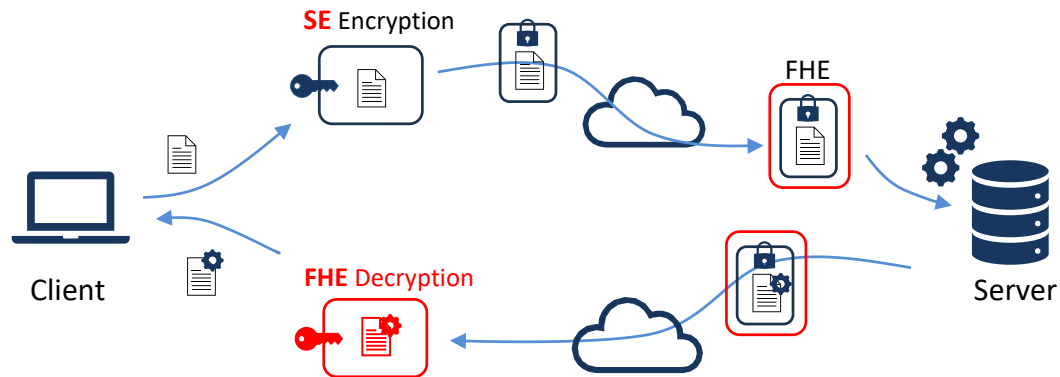
Hybrid Homomorphic Encryption



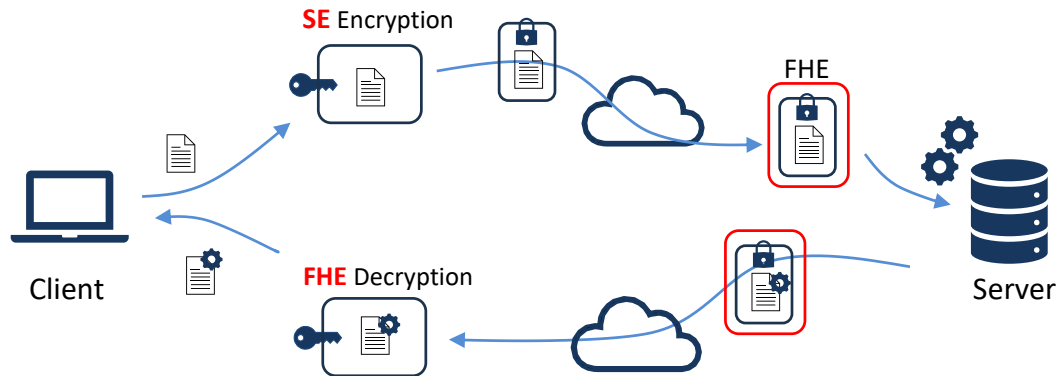
Hybrid Homomorphic Encryption



Hybrid Homomorphic Encryption



Hybrid Homomorphic Encryption



HERA Encryption

- i Prime-field arithmetic; 16 to 60 bits prime number p .
- i Generates a keystream from the input key K and adds it to the plaintext.
- i Typical **HERA** permutation consists out of 5 rounds.
- i State size: $n = 16$; hence state $X \in \mathbb{Z}_p^n$.

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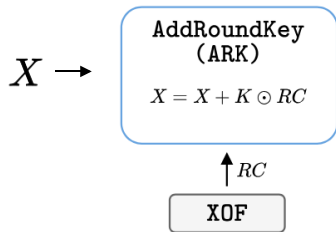
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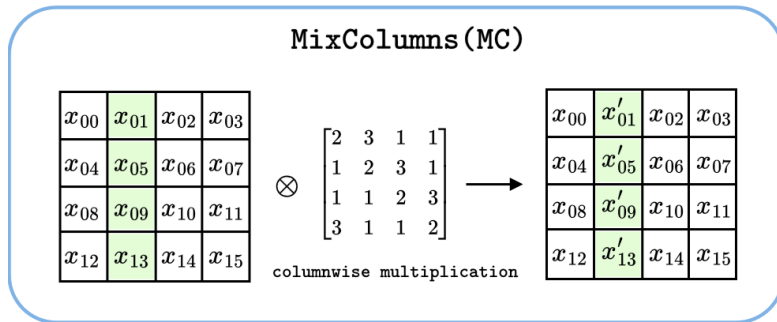
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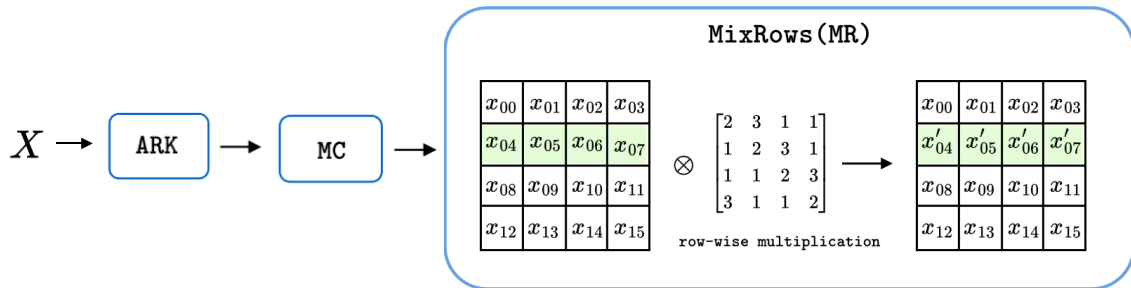
HERA Encryption



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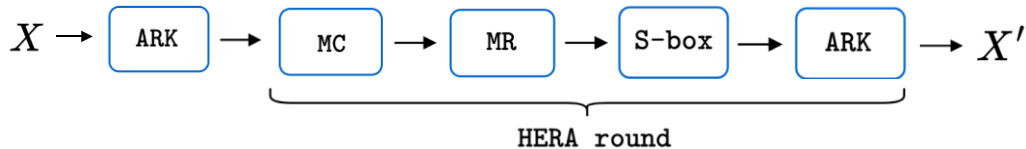
HERA Encryption



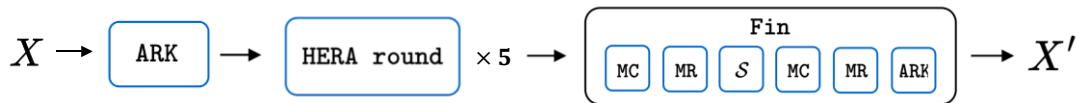
HERA Encryption



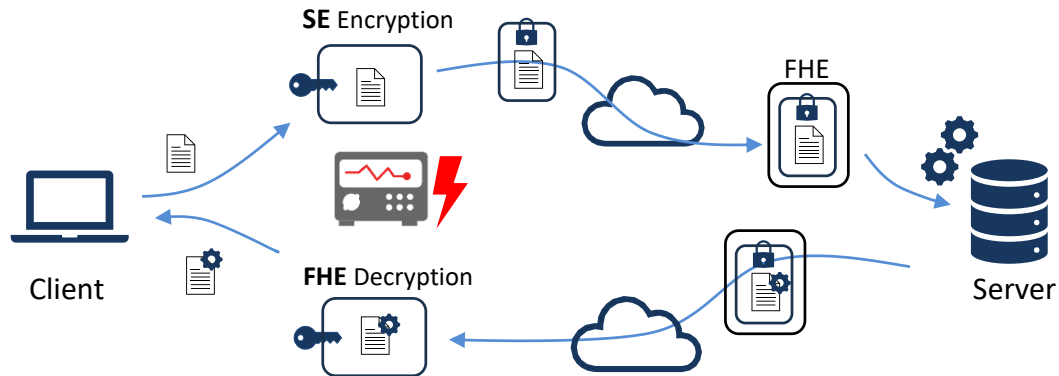
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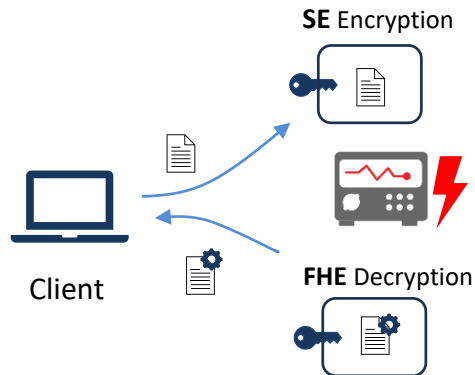
HERA Encryption



HERA Side-Channels Potential



HERA Side-Channels Potential

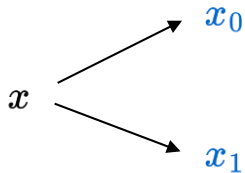


- i Exploit relation between power consumption and processed data.
- i Statistical methods to guess secret key bits - **DPA/CPA** possible.
- i Modular arithmetics spreads side-channel leakage.

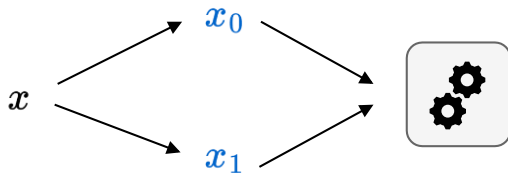
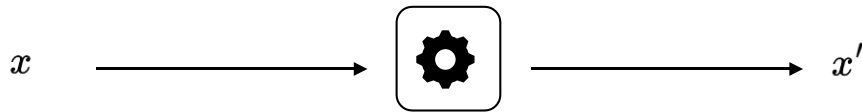
First Order Masking



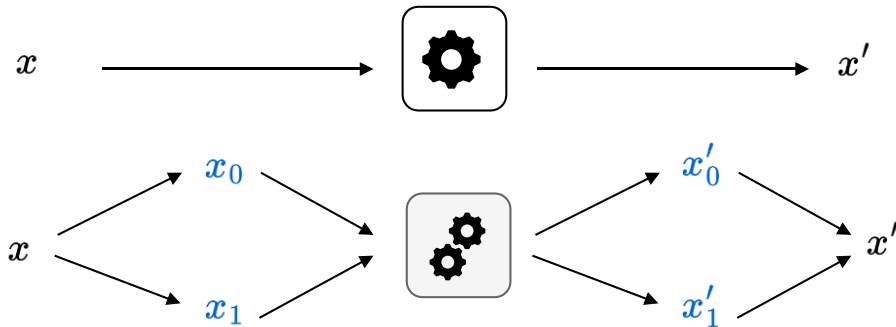
First Order Masking



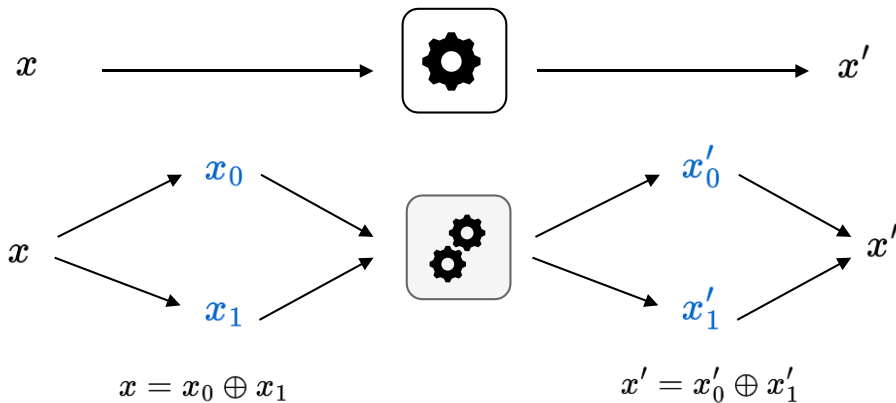
First Order Masking



First Order Masking

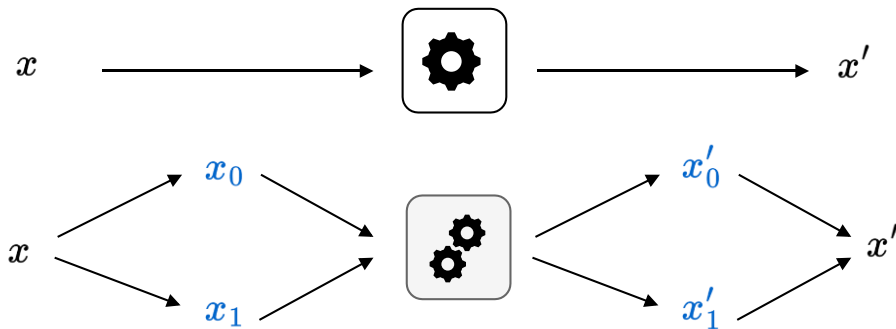


First Order Masking



Boolean sharing

First Order Masking



Arithmetic sharing

$$x = x_0 + x_1$$

$$x' = x'_0 + x'_1$$

Test Vector Leakage Assessment (TVLA)

- i Standard statistical method to detect **side-channel leakage**.
- i Applies **Welch's t-test** at each time point.
- i Compares device traces for **fixed vs random inputs**.
- i Evidence of leakage: $|t| > 4.5$.

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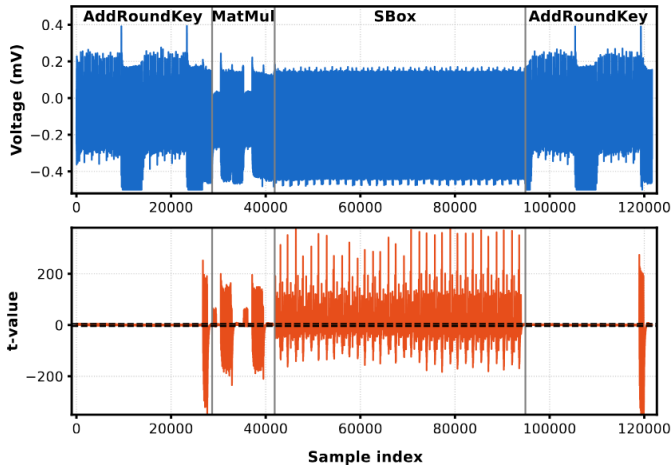
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TVLA Results for HERA



Masking

Masking a HERA round

$$X = X_0 + X_1 \pmod{p}$$

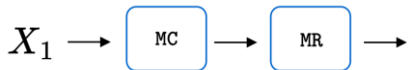
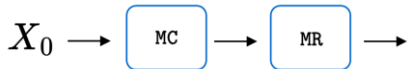
$$X_0 \rightarrow$$

$$X_1 \rightarrow$$

Masking

Masking a HERA round

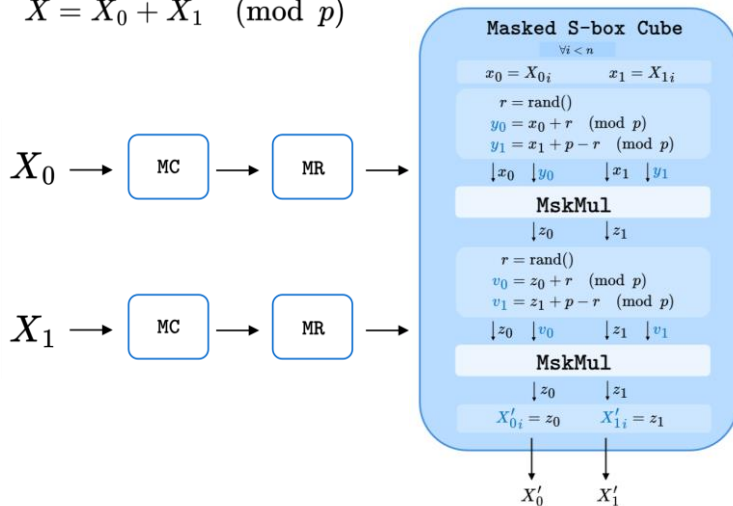
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Masking

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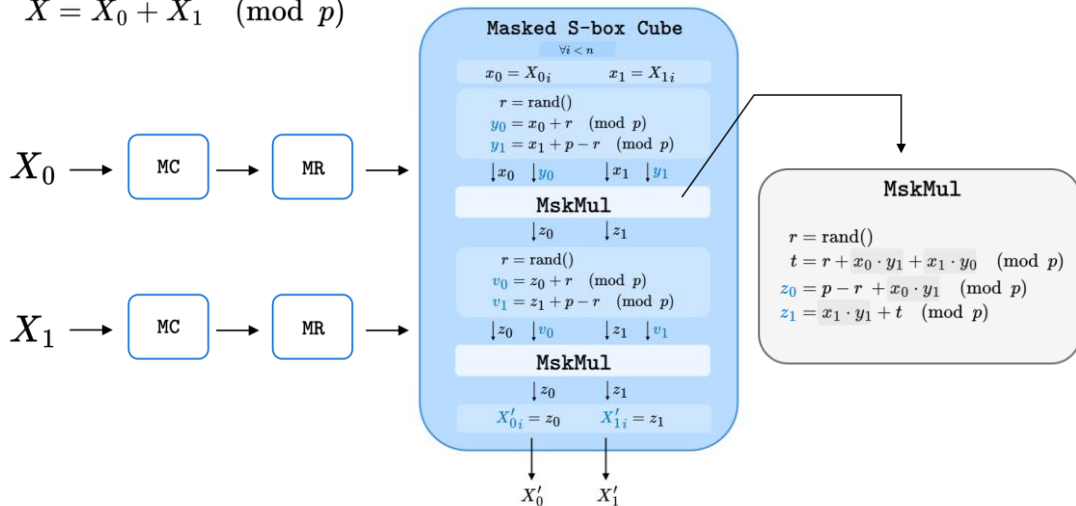
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Masking

Masking a HERA round

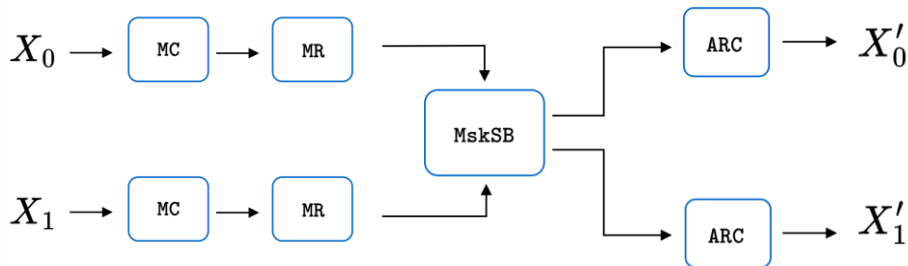
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Masking

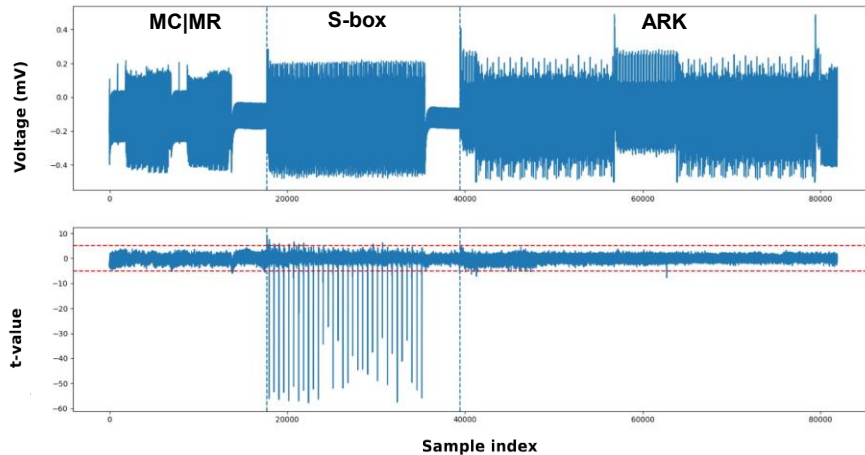
Masking a HERA round

$$X = X_0 + X_1 \pmod{p}$$



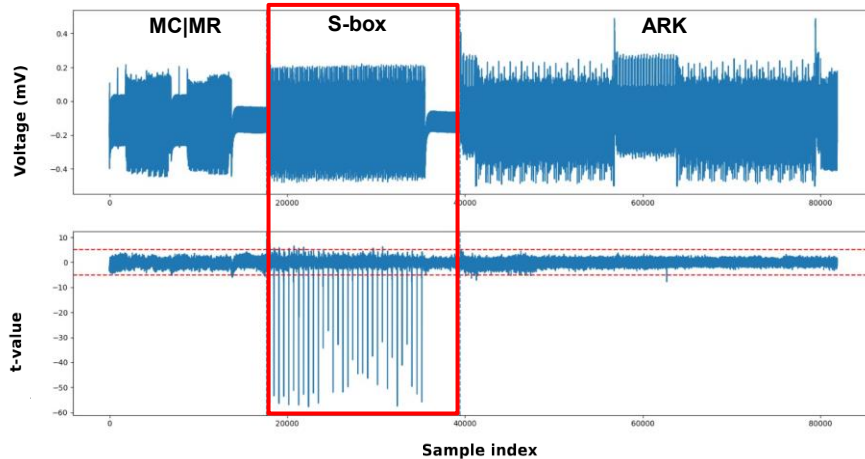
Preliminary Results

TVLA Results



Preliminary Results

TVLA Results



Implementation Heuristic

- i Ensure the random source is random indeed.
- i Separate shares usage within the C code.
- i Employ assembly code for full register control.
- i Address critical points which could lead to the statistical leakage despite the masked model.

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Implementation Heuristic

X

$$r_4 \leftarrow x_0$$

$$r_4 \leftarrow x_1$$



Implementation Heuristic

X

$$\left. \begin{array}{l} r_4 \leftarrow x_0 \\ r_4 \leftarrow x_1 \end{array} \right\} \begin{array}{l} \text{HD}(x_0, x_1) = \text{HW}(x_0 \oplus x_1) = \text{HW}(x) \\ \text{leakage through Hamming distance} \end{array}$$

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leakage through Hamming distance

✓

$$r_3 \leftarrow x_0$$

$$r_4 \leftarrow x_1$$

Implementation Heuristic

X

unknown $\rightarrow r_3 = ?$

$r_3 \leftarrow x_0$



Implementation Heuristic

X

:



$\xrightarrow{\text{unknown}} r_3 = ?$
 $r_3 \leftarrow x_0$

$$\text{HD}(x_0, x_1) = \text{HW}(x_0 \oplus x_1) = \text{HW}(x)$$

potential leakage through Hamming distance

Implementation Heuristic

X

$$\left. \begin{array}{l} \text{unknown} \rightarrow r_3 = ? \\ r_3 \leftarrow x_0 \end{array} \right\}$$

$$\left. \begin{array}{l} \text{HD}(x_0, x_1) = \text{HW}(x_0 \oplus x_1) = \text{HW}(x) \\ \text{potential leakage through Hamming distance} \end{array} \right\}$$

✓

$$\left. \begin{array}{l} r_3 \leftarrow \text{rand} \\ r_3 \leftarrow x_0 \end{array} \right\}$$

Implementation Heuristic

X

$r_3 \leftarrow x$

$r_3 \leftarrow ?$

$$\text{HD}(x_0, x_1) = \text{HW}(x_0 \oplus x_1) = \text{HW}(x)$$

potential leakage through Hamming distance

✓

$r_3 \leftarrow x_0$

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unknown

Implementation Heuristic

X

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Implementation Heuristic

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leakage due to shares loaded close in time



Implementation Heuristic

X
$$\left. \begin{array}{l} r_3 \leftarrow x_0 \\ r_4 \leftarrow x_1 \end{array} \right\}$$

leakage due to shares loaded close in time

✓
$$\left. \begin{array}{l} r_3 \leftarrow x_0 \\ \text{dummy ops} \\ r_4 \leftarrow x_1 \end{array} \right\}$$

Implementation Heuristic


X



call $f(x_0, x_1)$

Implementation Heuristic

X

`call $f(x_0, x_1)$` }  leakage due to shares loaded close in time
in function calls

Implementation Heuristic

X

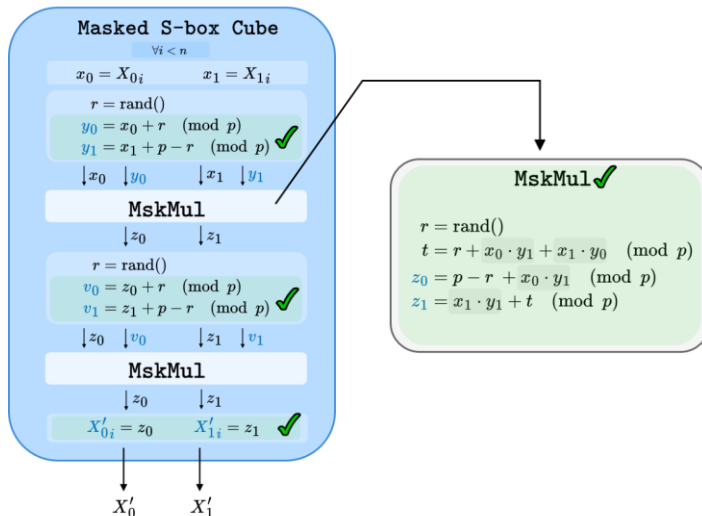
`call f(x_0 , x_1)`

leakage due to shares loaded close in time
in function calls

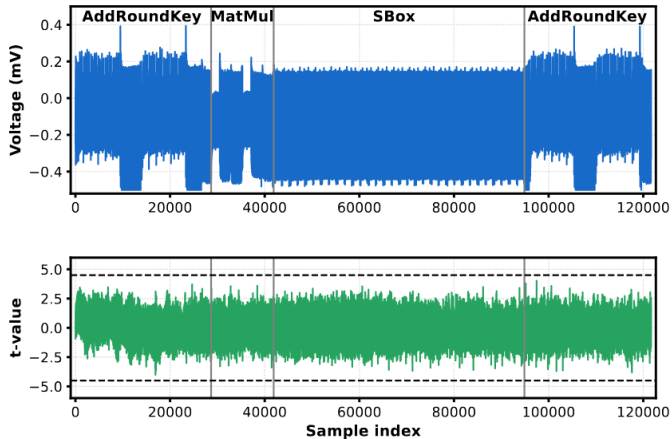
✓

`call f(& x_0 , & x_1)`

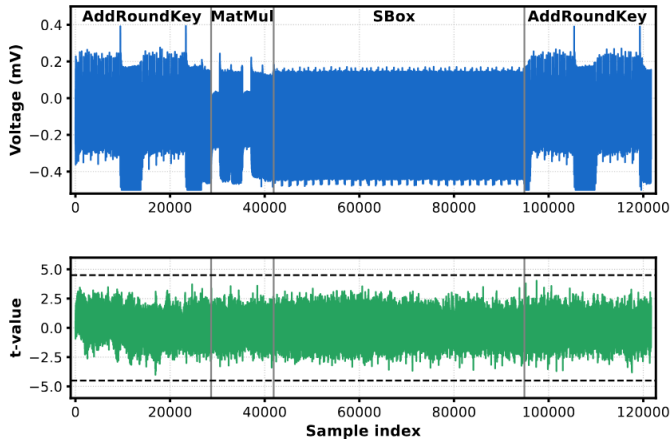
Application to HERA S-box



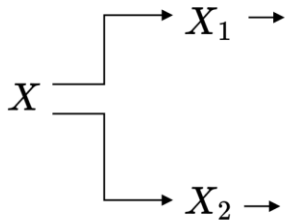
TVLA Results



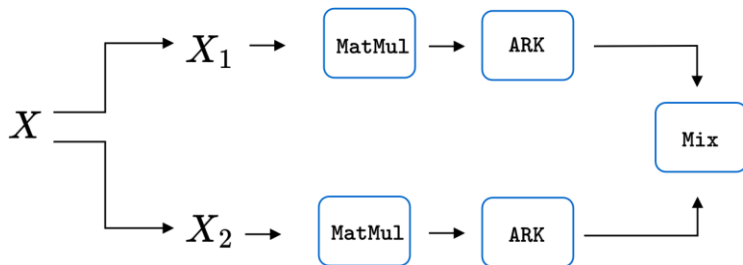
TVLA Results



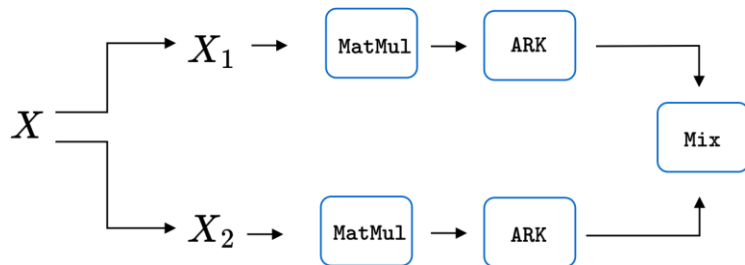
PASTA Encryption



PASTA Encryption

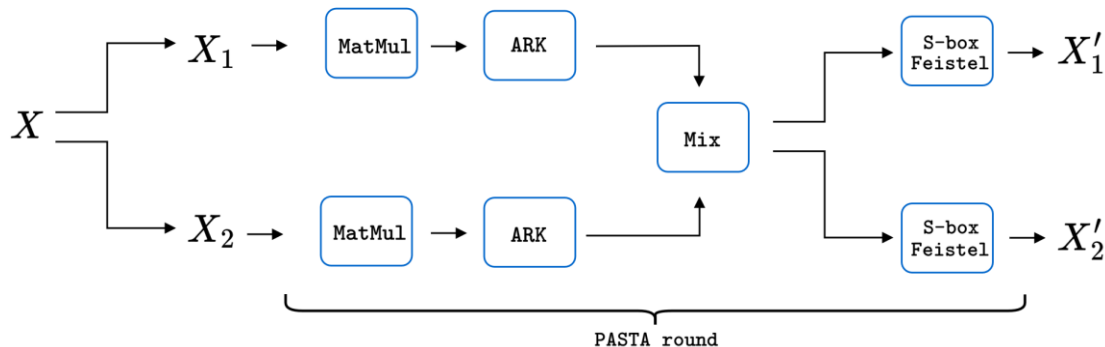


PASTA Encryption

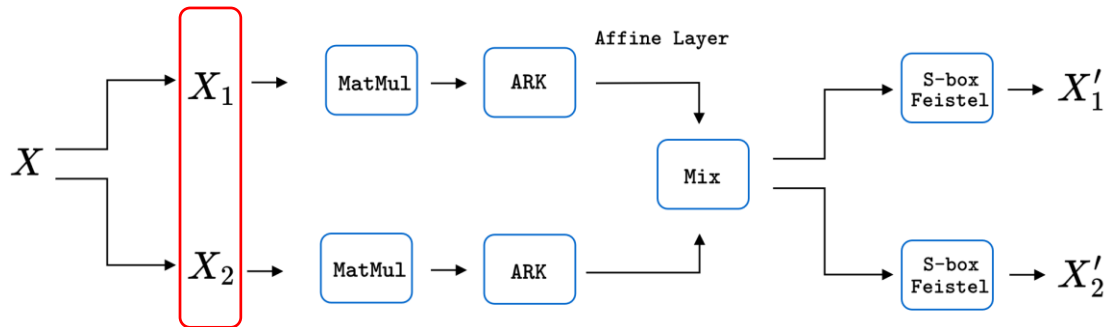


$$\begin{aligned} &\text{S-box Feistel} \\ &\quad \forall i < n \\ &= \begin{cases} x_i & \text{if } i = 0, \\ x_i + (x_{i-1})^2 & \text{otherwise.} \end{cases} \end{aligned}$$

PASTA Encryption

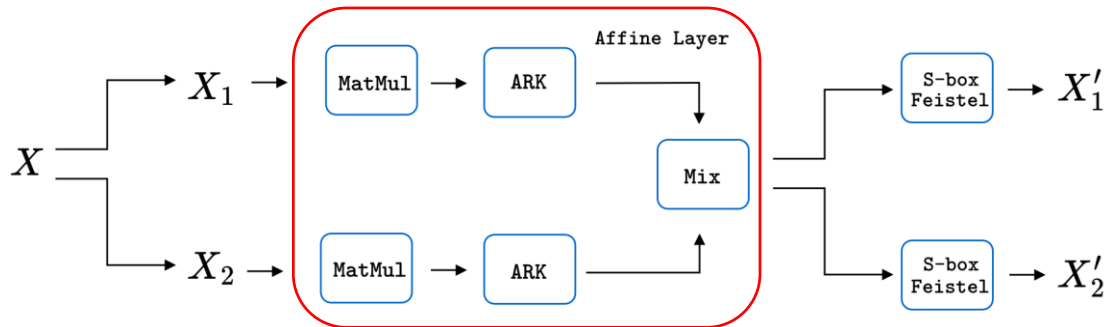


Masking PASTA Round



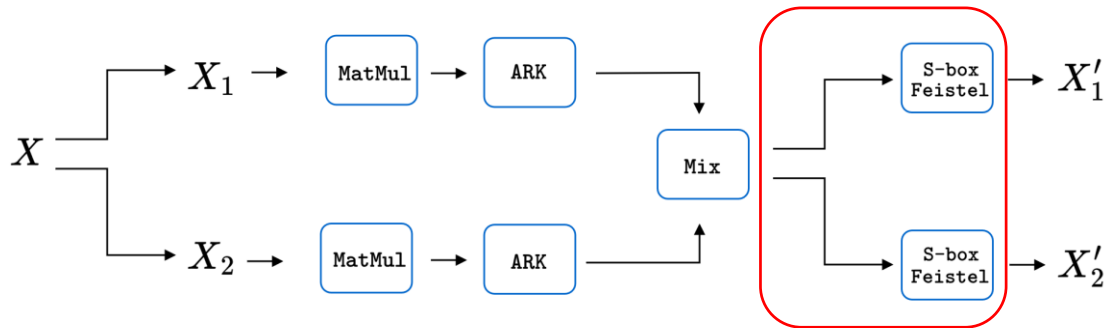
✓ mask both parts

Masking PASTA Round



✓ apply linear layers on each share

Masking PASTA Round



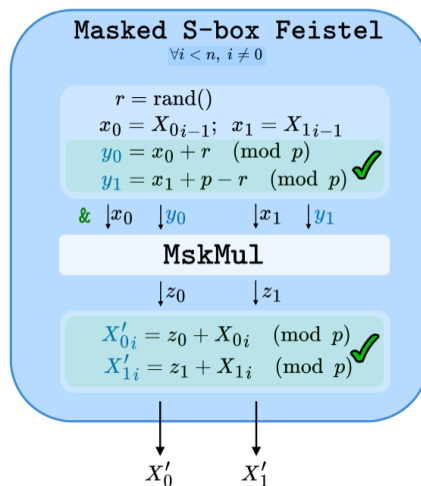
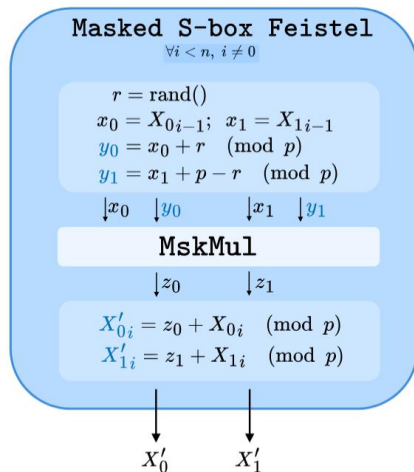
✓ Use secure multiplication

Masking PASTA Round

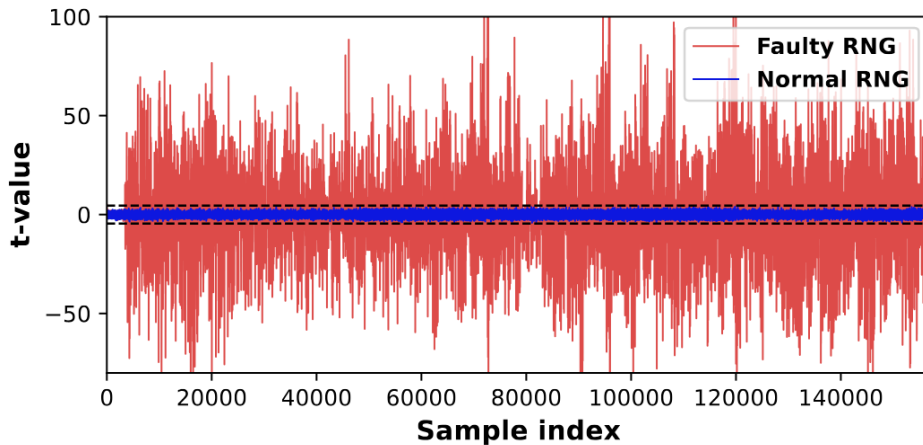
Masked S-box Feistel

 $\forall i < n, i \neq 0$ $r = \text{rand}()$ $x_0 = X_{0i-1}; x_1 = X_{1i-1}$ $y_0 = x_0 + r \pmod{p}$ $y_1 = x_1 + p - r \pmod{p}$ $\downarrow x_0 \quad \downarrow y_0 \quad \downarrow x_1 \quad \downarrow y_1$ **MskMul** $\downarrow z_0 \quad \downarrow z_1$ $X'_{0i} = z_0 + X_{0i} \pmod{p}$ $X'_{1i} = z_1 + X_{1i} \pmod{p}$ $\downarrow \quad \downarrow$
 $X'_0 \quad X'_1$

Masking PASTA Round



TVLA Results



Higher Order Masking

- i First order masking protects only against single-probe attacks.
- i Extend with **higher order masking** for more security, by expanding the MskMul accordingly.

Conclusion

- i HERA vulnerable to side-channel attacks.
- i Correct masking model can leak due to implementation specificities.
- i The proposed approach application concretely verified on PASTA.
- i First order masking protects only against single-probe attacks.