

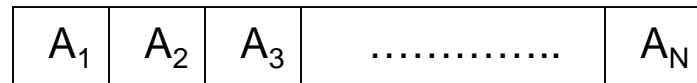
Chap. 3 Lists, Stacks, and Queues

3.1 Abstract Data Types (ADTs)

- An ADT is a set of objects (of the same type) with a set of operations
- ADT can be naturally implemented by C++ classes
- Conceptually, a data structure (list, set, graph) is an ADT

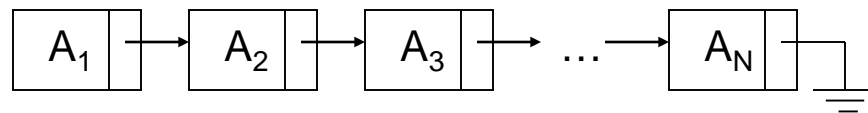
3.2 Lists

- List ADT
 - A sequence of objects A_1, A_2, \dots, A_N
 - Common list operations include: *printList*, *makeEmpty*, *find*, *insert*, *remove*, *findKth*, *next*, *findPrevious*, etc.
- Array implementation of lists
 - Size estimate
 - *printList* & *find*: linear time $O(N)$
 - *findKth*: constant time $O(1)$



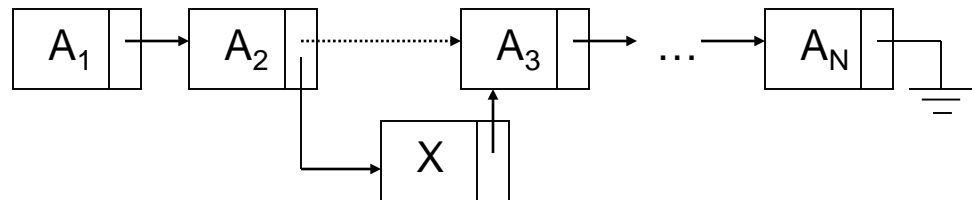
- *insert & remove*: involving memory management, linear cost.
- A generally inefficient implementation for dynamic lists, but good for relatively static lists.

– Linked lists

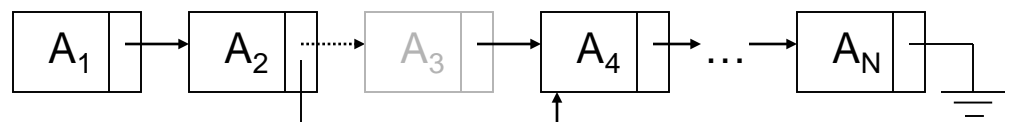


The nodes are not necessarily contiguous memory cells.

- *printList, findKth*: linear time
- *insert*: constant time

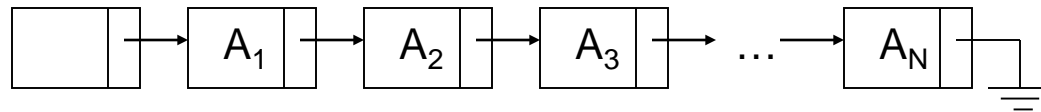


- *remove*: constant time

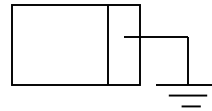


– Header

- A dummy node pointing to the first node of the list

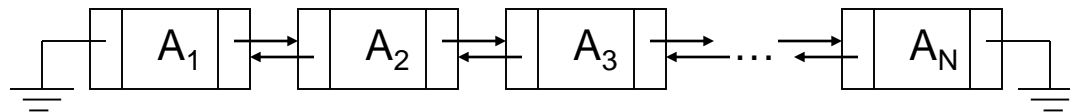


- Can reprint empty list



- Can avoid failures in special cases for list operations
- C++ implementations (Textbook, page 73–79)

– Doubly linked lists



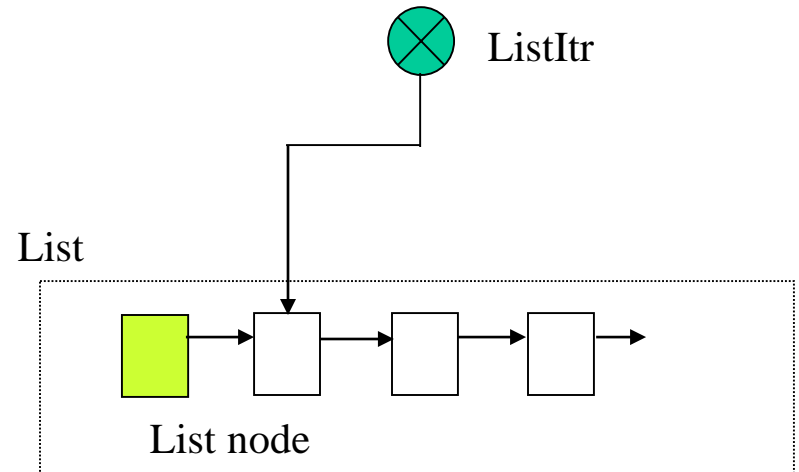
- More memory
- Simplifies deletion (*findPrevious*)

```

#ifndef _LinkedList_H
#define _LinkedList_H
#include "dexceptions.h"
#include <iostream.h> // For NULL
// List class CONSTRUCTION: with no initializer
// Access is via ListItr class
// boolean isEmpty()
//
//      --> Return true if empty; else false
// void makeEmpty()    --> Remove all items
// ListItr zeroth()
//
//      --> Return position to prior to first
// ListItr first()     --> Return first position
// void insert( x, p )
//
//      --> Insert x after current iterator position p
// void remove( x )    --> Remove x
// ListItr find( x )   --> Return position that views x
// ListItr findPrevious( x )
//
//      --> Return position prior to x
template <class Object>
class List; // Incomplete declaration.

template <class Object>
class ListItr; // Incomplete declaration.

```



```

template <class Object>
class ListNode
{
    ListNode( const Object & theElement
= Object( ), ListNode * n = NULL ) : element(
theElement ), next( n ) { }
    Object element;
    ListNode *next;
    friend class List<Object>;
    friend class ListItr<Object>;
};

```

```

template <class Object>
class List
{
public:
    List();
    List( const List & rhs );
    ~List();

    bool isEmpty() const;
    void makeEmpty();
    ListItr<Object> zeroth() const;
    ListItr<Object> first() const;
    void insert( const Object & x, const ListItr<Object> & p );
    ListItr<Object> find( const Object & x ) const;
    ListItr<Object> findPrevious( const Object & x ) const;
    void remove( const Object & x );
    const List & operator=( const List & rhs );
private:
    ListNode<Object> *header;
};

// ListItr class; maintains "current position" CONSTRUCTION:
// Package friendly only, with a ListNode

```

```

// *****PUBLIC OPERATIONS*****

// bool isPastEnd()    --> True if past end position in list
// void advance()      --> Advance (if not already null)
// Object retrieve      --> Return item in current position

template <class Object>
class ListItr
{
public:
    ListItr() : current( NULL ) { }
    bool isPastEnd() const
        { return current == NULL; }
    void advance()
        { if( !isPastEnd() ) current = current->next; }
    const Object & retrieve() const
        { if( isPastEnd() ) throw BadIterator();
          return current->element; }
private:
    ListNode<Object> *current; // Current position

    ListItr( ListNode<Object> *theNode )
        : current( theNode ) { }
    friend class List<Object>; // Grant access to constructor
};

#include "LinkedList.cpp"
#endif

```

```

#include "LinkedList.h"

/* Construct the list. */
template <class Object>
List<Object>::List( )
{
    header = new ListNode<Object>;
}

/* Copy constructor. */
template <class Object>
List<Object>::List( const List<Object> & rhs )
{
    header = new ListNode<Object>;
    *this = rhs;
}

/* Destructor. */
template <class Object>
List<Object>::~~List( )
{
    makeEmpty( );
    delete header;
}

/* Test if the list is logically empty. Return true if empty, false,
otherwise.*/
template <class Object>
bool List<Object>::isEmpty( ) const
{
    return header->next == NULL;
}

/* Make the list logically empty. */
template <class Object>
void List<Object>::makeEmpty( )
{
    while( !isEmpty( ) )
        remove( first( ).retrieve( ) );
}

/* Return an iterator representing the header node. */
template <class Object>
ListItr<Object> List<Object>::zeroth( ) const
{
    return ListItr<Object>( header );
}

/* Return an iterator representing the first node in the list.
This operation is valid for empty lists. */
template <class Object>
ListItr<Object> List<Object>::first( ) const
{
    return ListItr<Object>( header->next );
}

```

```

/* Insert item x after p. */
template <class Object>
void List<Object>::insert( const Object & x, const ListItr<Object>
& p )
    {   if( p.current != NULL )
        p.current->next = new ListNode<Object>( x, p.current-
>next );
    }

/* Return iterator corresponding to the first node containing an
item x. Iterator isPastEnd if item is not found. */
template <class Object>
ListItr<Object> List<Object>::find( const Object & x ) const {
/* 1*/   ListNode<Object> *itr = header->next;
/* 2*/   while( itr != NULL && itr->element != x )
/* 3*/       itr = itr->next;
/* 4*/   return ListItr<Object>( itr );
    }

/* Return iterator prior to the first node containing an item x. */
template <class Object>
ListItr<Object> List<Object>::findPrevious( const Object & x )
const
    {
/* 1*/   ListNode<Object> *itr = header;
/* 2*/   while( itr->next != NULL && itr->next->element != x )
/* 3*/       itr = itr->next;
/* 4*/   return ListItr<Object>( itr );
    }

```

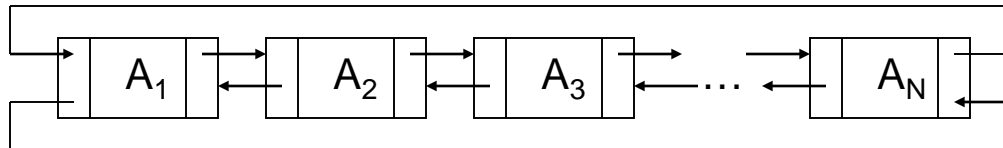
```

/* Remove the first occurrence of an item x. */
template <class Object>
void List<Object>::remove( const Object & x )
{
    ListItr<Object> p = findPrevious( x );
    if( p.current->next != NULL )
    {
        ListNode<Object> *oldNode = p.current->next;
        p.current->next = p.current->next->next;
        // Bypass deleted node
        delete oldNode;
    }
}

/* Deep copy of linked lists.*/
template <class Object>
const List<Object>&List<Object>::operator=( const List<Object>
& rhs )
{   if( this != &rhs )
    {
        makeEmpty( );
        ListItr<Object> ritr = rhs.first();
        ListItr<Object> itr = zeroth( );
        for( ; !ritr.isPastEnd(); ritr.advance( ), itr.advance( ) )
            insert( ritr.retrieve( ), itr );
    }
    return *this;
}

```

– Circular linked lists



- No need for header
- No special case for “*next*” & “*findPrevious*”

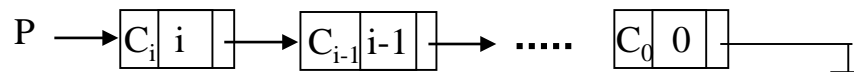
– Examples

- Polynomials

(Textbook, page 81-83)

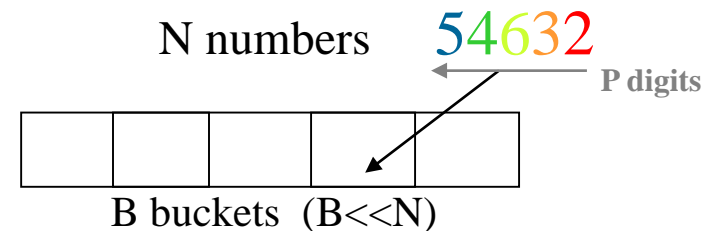
$$F(x) = \sum_{i=0}^N c_i x^i$$

$$c_n x^n + c_{n-1} x^{n-1} + \dots + c_1 x + c_0$$



- Radix Sort - $O(P(N+B))$

(Textbook, page 83-85)



- Multi-lists

(Textbook, page 85-86)

Students

Classes

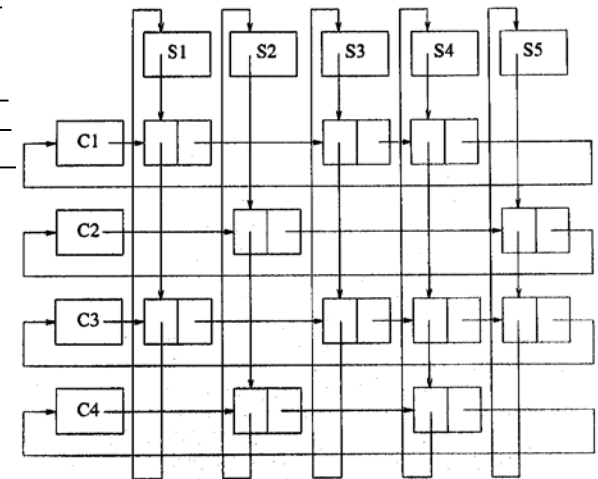
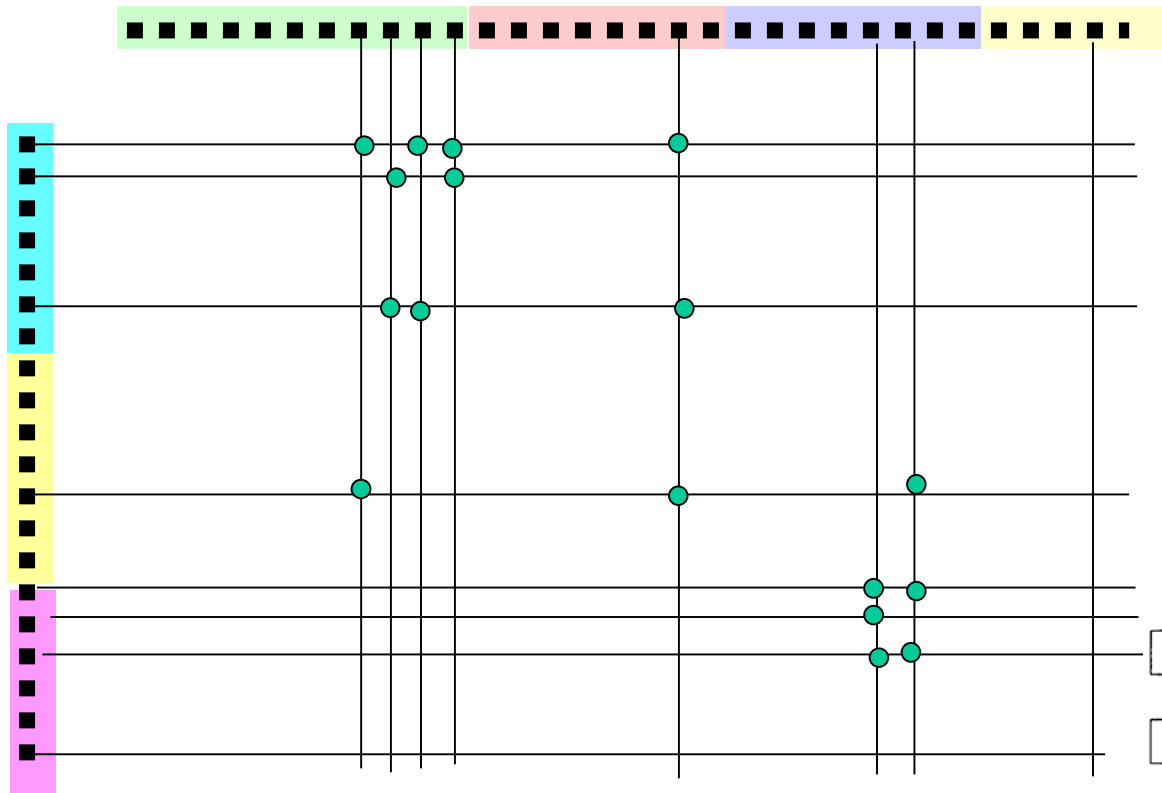


Figure 3.28 Multilist implementation for registration problem

$$F(x) = \sum_{i=0}^n a_i x^i$$

$$G(x) = \sum_{i=0}^m b_i x^i$$

$$a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$$

$$b_m x^m + b_{m-1} x^{m-1} + \dots + b_1 x + b_0$$

$$F(x) + G(x), \quad F(x) * G(x)$$

/* This code doesn't really do much, and abstraction is not built in.

*** Thus, I haven't bothered testing it exhaustively.***

#include <iostream.h>

#include "vector.h"

class Polynomial

{ enum { MAX_DEGREE = 100 };

friend int main(); // So I can do a quick test.

public:

Polynomial();

void zeroPolynomial();

Polynomial operator+(const Polynomial & rhs) const;

Polynomial operator*(const Polynomial & rhs) const;

void print(ostream & out) const;

private:

vector<int> coeffArray;

int highPower;

};

int max(int a, int b)

{ return a > b ? a : b;

}

Polynomial::Polynomial() : coeffArray(MAX_DEGREE + 1)

{ zeroPolynomial();

}

void Polynomial::zeroPolynomial()

{ for(int i = 0; i <= MAX_DEGREE; i++)

coeffArray[i] = 0;

highPower = 0;

}

```
Polynomial Polynomial::operator+( const Polynomial & rhs )  
const
```

```
{ Polynomial sum;  
  sum.highPower = max( highPower, rhs.highPower );  
  for( int i = sum.highPower; i >= 0; i-- )  
    sum.coeffArray[ i ] = coeffArray[ i ] + rhs.coeffArray[ i ];  
  return sum;  
}
```

```
Polynomial Polynomial::operator*( const Polynomial & rhs ) const
```

```
{ Polynomial product;  
  product.highPower = highPower + rhs.highPower;  
  if( product.highPower > MAX_DEGREE )  
    cerr << "operator* exceeded MAX_DEGREE" << endl;  
  for( int i = 0; i <= highPower; i++ )  
    for( int j = 0; j <= rhs.highPower; j++ )  
      product.coeffArray[ i + j ] +=  
        coeffArray[ i ] * rhs.coeffArray[ j ];  
  return product;  
}
```

```
void Polynomial::print( ostream & out ) const
```

```
{ for( int i = highPower; i > 0; i-- )  
  out << coeffArray[ i ] << "x^" << i << " + ";  
  out << coeffArray[ 0 ] << endl;  
}
```

```
ostream & operator<<( ostream & out, const  
Polynomial & rhs )
```

```
{ rhs.print( out );  
  return out;  
}
```

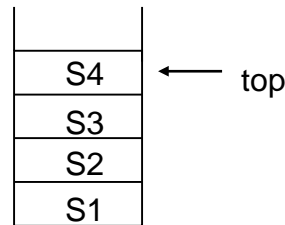
```
int main( )
```

```
{ Polynomial p;  
  Polynomial q;  
  
  p.highPower = 1;  
  p.coeffArray[ 0 ] = 1;  
  p.coeffArray[ 1 ] = 1;  
  
  q = p + p;  
  p = q * q;  
  q = p + p;  
  cout << q << endl;  
  return 0;  
}
```

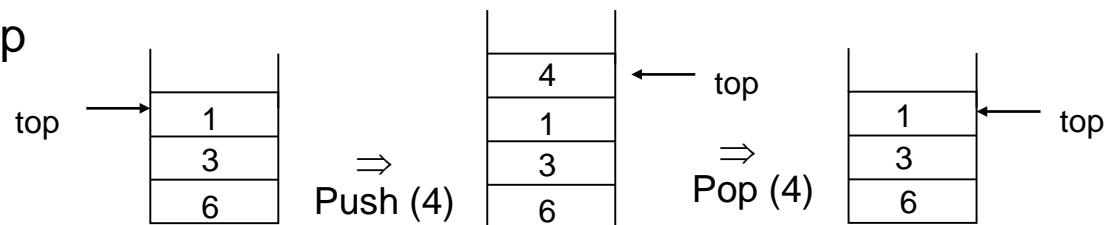
3.3 Stacks

– Stack model

- A stack is a list with the restriction that insertion & deletion can be performed only at the end (or top) of the list.

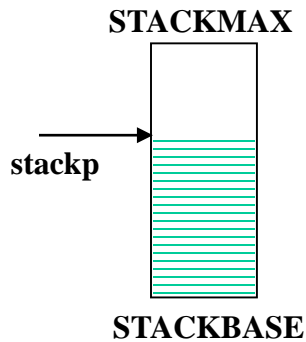


- Only the top node is accessible
- Last in, first out (LIFO)
- Push
- Pop



- A stack can be empty, “pop” from an empty stack is an error
- A stack can never be full (assuming infinite memory)

- Implementation by linked lists (Fig. 3.41, page 94)
 - Methods implementation (Fig. 3.42 – 3.47)
- Implementation by array (Fig. 3.48, page 99)
 - Need to set the maximum stack size
 - Pop & push: constant time (fast)
 - More commonly used than list implementation



```
PUSH: if(stackp)>STACKMAX then
{
    (stackp)+1 → stackp
    x → (stackp)
}
```

```
POP: if (stackp)<STACKBASE then
{
    ((stackp)) → x;
    (stackp)-1 → stackp
}
```

Push and Pop in C or Assembly

```
// Stack class -- array implementation
// CONSTRUCTION: with or without a capacity; default is 10
// *****PUBLIC OPERATIONS*****

// void push( x )      --> Insert x
// void pop( )          --> Remove most recently inserted item
// Object top( )        --> Return most recently inserted item
// Object topAndPop( )  --> Return and remove most
//      recently inserted item
// bool isEmpty( )      --> Return true if empty; else false
// bool isFull( )       --> Return true if full; else false
// void makeEmpty( )    --> Remove all items
// *****ERRORS*****
// Overflow and Underflow thrown as needed
```

```
#ifndef _STACKAR_H
#define _STACKAR_H
#include "vector.h"
#include "dsexceptions.h"

template <class Object>
class Stack
{
public:
    explicit Stack( int capacity = 10 );
    bool isEmpty( ) const;
    bool isFull( ) const;
    const Object & top( ) const;

    void makeEmpty( );
    void pop( );
    void push( const Object & x );
    Object topAndPop( );
private:
    vector<Object> theArray;
    int topOfStack;
};
#include "StackAr.cpp"
#endif
```

```

#include "StackAr.h"

/**Construct the stack. */

    template <class Object>
    Stack<Object>::Stack( int capacity ) : theArray( capacity )
    {
        topOfStack = -1;
    }

/** Test if the stack is logically empty. Return true if empty, false, otherwise. */

    template <class Object>
    bool Stack<Object>::isEmpty( ) const
    {
        return topOfStack == -1;
    }

/**Test if the stack is logically full. Return true if full, false otherwise. */

    template <class Object>
    bool Stack<Object>::isFull( ) const
    {
        return topOfStack == theArray.size( ) - 1;    }

/**      * Make the stack logically empty.      */

    template <class Object>
    void Stack<Object>::makeEmpty( )
    {
        topOfStack = -1;    }

/**Get the most recently inserted item in the stack. Does not alter the stack. Return the most recently inserted item in the stack. Exception Underflow if stack is already empty. */

    template <class Object>
    const Object & Stack<Object>::top( ) const
    {
        if( isEmpty( ) )
            throw Underflow( );

        return theArray[ topOfStack ];
    }

/** Remove the most recently inserted item from the stack.Exception Underflow if stack is already empty. */

    template <class Object>
    void Stack<Object>::pop( )
    {
        if( isEmpty( ) )
            throw new Underflow( );

        topOfStack--;
    }

/**Insert x into the stack, if not already full.Exception Overflow if stack is already full.*/

    template <class Object>
    void Stack<Object>::push( const Object & x )
    {
        if( isFull( ) )
            throw Overflow( );

        theArray[ ++topOfStack ] = x;
    }

/**Return and remove most recently inserted item from the stack.Return most recently inserted item.Exception Underflow if stack is already empty. */

    template <class Object>
    Object Stack<Object>::topAndPop( )
    {
        if( isEmpty( ) )
            throw Underflow( );

        return theArray[ topOfStack-- ];
    }

```

– Stack applications

- Balancing symbols: linear time (page 101 - 102)

`[xxx(xxx)xx(x)x]xxxx{xx[x(xxxx)x(x)]xxx}`

Read characters until end-of-file

Push opening symbol to stack,

pop counterpart if closing symbol is read

- Postfix expressions (inverse Polish notation) (page 104)

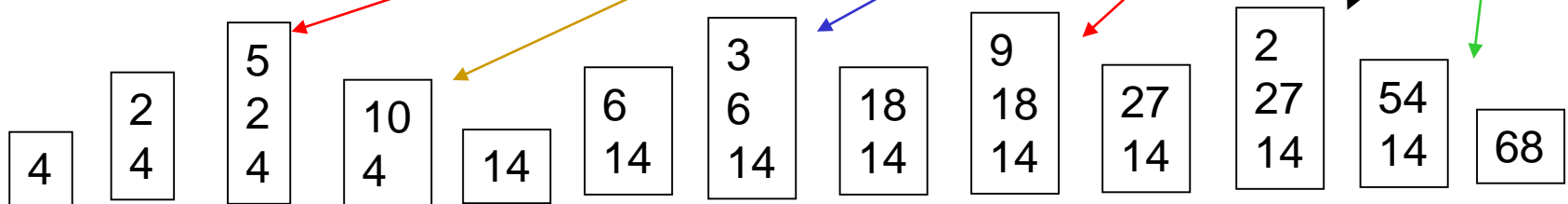
$a*b+c+d*e \Rightarrow ab*c+de*+$

Polish notation

Normal arithmetic form

4 + 2 * 5 + (6 * 3 + 9) * 2 → 4 2 5 * + 6 3 * 9 + 2 * +

Evaluation



- Infix to Postfix Conversion (page 106)

e.g. $a + b * c + (d * e + f) * g \rightarrow a b c * + d e * f + g * +$

Conversion

Operator precedence: (, * / , + - ,)

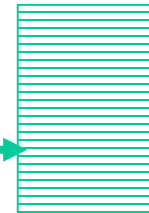
If '(' then push
 else if ')' then pop entries until '('
 else if 'operator' then {
 pop operators with **higher** or **same** precedence
 push
 }
 else if end of input then pop until stack is empty

Input sequence

a + b * c + (d * e + f) * g

↑
pointer

Stack of pending operator



↓
pointer

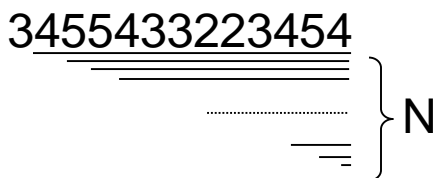
a b c * + d e * f + g * +

Output sequence

–Function calls using stacks

- Saving local variables using stack
- Recursion: stack implementation
- Stack overflow with runaway recursion

tail recursion 3455433223454



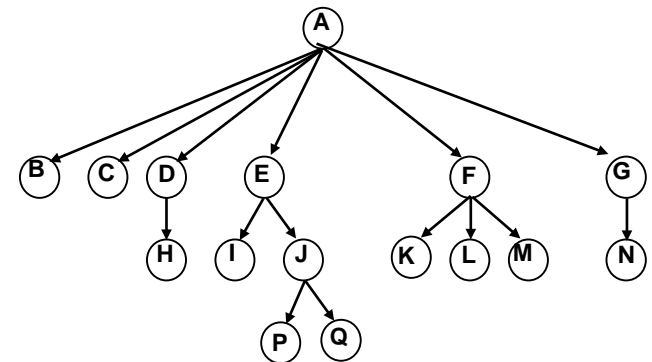
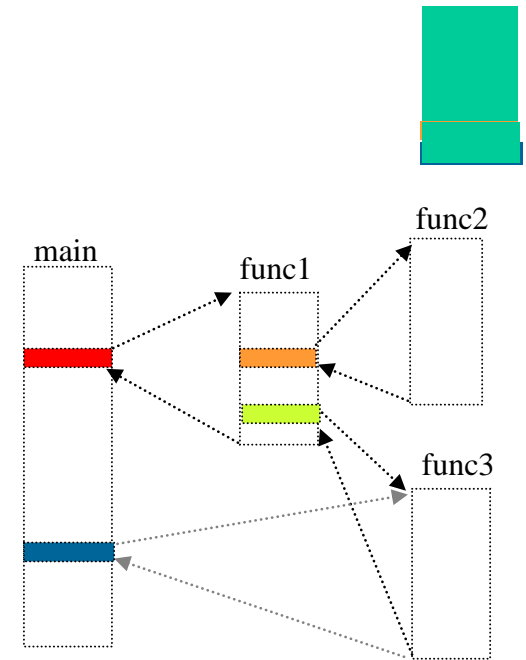
The diagram shows a sequence of numbers: 3, 4, 5, 5, 4, 3, 3, 2, 2, 3, 4, 5, 4. The last five numbers (5, 4, 3, 3, 2) are underlined. A bracket to the right of the underlined numbers is labeled 'N', indicating the number of elements in the stack.

Tail recursion: N elements \rightarrow N layers of recursion

Tree: N elements \rightarrow $\log N$ layers of recursion

example: fig.3.55

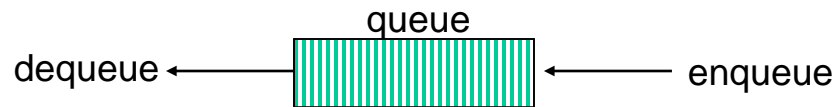
Stack for activation record



3.4 Queues

– Queue Model

- queue is a list, with insertion done only at one end and deletion done at the other end.
- enqueue: insert an element at the end of the queue
- dequeue: delete (and return) the element at the start of the queue
- first in first out model



– Linked list implementation of queues

- operating as a list
- constant time for enqueue & dequeue (keeping pointer to both the head and tail of the list)

- Array implementation of queues
 - front pointer, back pointer, current size
 - circular array (Fig. in page 111)
 - C++ implementation (Fig. 3.58 - 3.61)
- Applications of queues
 - printer job queues
 - telephone queues
 - class waiting list

