FEAT: Functional Enumeration of Algebraic Types

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Get the Haskell lib.: cabal install testing-feat

Example

```
*Main> index (10<sup>100</sup>)::([[Bool]],[Bool])
([[],[],[False,False, True, True],...[False],[]],
[True, True, False, False])
```

Use case: test pretty-print + parse for Template Haskell.

data Exp = VarE Name | CaseE Exp [Match] | ... -- 18 Con

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data $Exp = VarE \ Name \ | \ CaseE \ Exp \ [Match] \ | \ ... - 18 \ Con$

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data $Exp = VarE \ Name \mid CaseE \ Exp \ [Match] \mid ... -- 18 \ Con$

data Match = Match Pat Body [Dec]

```
\begin{array}{lll} \textbf{data} \ \textit{Pat} & = \textit{LitP Lit} \ | \ \textit{ViewP Exp Pat} \ | \ \dots & -- \ 14 \ \mathsf{Con} \\ \textbf{data} \ \textit{Body} & = \textit{GuardedB} \left[ \left( \textit{Guard}, \textit{Exp} \right) \right] \ | \ \textit{NormalB Exp} \\ \textbf{data} \ \textit{Dec} & = \textit{FunD Name} \left[ \textit{Clause} \right] \ | \ \dots & -- \ 14 \ \mathsf{Con} \\ \end{array}
```

Use case: test pretty-print + parse for Template Haskell.

data $Exp = VarE Name \mid CaseE Exp [Match] \mid ...$ -- 18 Con

data Match = Match Pat Body [Dec]

```
data Pat= LitP Lit | ViewP Exp Pat | ...-- 14 Condata Body= GuardedB [(Guard, Exp)] | NormalB Expdata Dec= FunD Name [Clause] | ...-- 14 Con
```

data Clause = Clause [Pat] Body [Dec]

Around 10 datatypes, around 100 constructors.

Enumerating by size

```
Enumeration of [Bool]
  { [], [False], [True], [False, False], [False, True] ...
 index 0 \mapsto []
 index 1 \mapsto [False]
 index 2 \mapsto [True]
```

Partitioning by size

```
Partitioning of [Bool]:
                                     0 Constructors
    {[]},
                                     1 Constructor
                                     2 Constructors
    { [False], [True] },
    { [False, False], [False, True]...
```

An infinite sequence of finite sequences

```
type E a -- fun. enum. of values of type a type F a -- similar but for finite enum. sel :: E a 	o Part 	o F a -- pick one finite enum. card :: F a 	o \mathbb{N} -- get its cardinality index_F:: F a 	o Index 	o a -- ... and its elements -- where Part = Index = \mathbb{N}
```

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Possible implementations of F and E:

type
$$F$$
 $a = (\mathbb{N}, Index \rightarrow a)$ -- we will use this **type** F' $a = [a]$ -- not this

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Possible implementations of F and E:

type
$$F$$
 $a = (\mathbb{N}, Index \rightarrow a)$ -- we will use this **type** F' $a = [a]$ -- not this **type** E $a = [F a]$ -- but we use (roughly) this **type** E' $a = Part \rightarrow F a$ -- not this

Set operations for E and for F:

$$empty :: E \ a$$
 $sing :: a \rightarrow E \ a$
 $(\oplus) :: E \ a \rightarrow E \ a \rightarrow E \ b \rightarrow E \ (a \times b)$ -- diagonalisation
 $pay :: E \ a \rightarrow E \ a$ -- increase cost (size)
 $fmap :: (a \rightarrow b) \rightarrow E \ a \rightarrow E \ b$

$$sing_F$$
 :: $a o F$ a (\oplus_F) :: F $a o F$ $a o F$ $a o F$ (\otimes_F) :: F $a o F$ $b o F$ $(a imes b)$

Example enumeration of a simple datatype

data
$$T = L \mid S T \mid B T T$$

instance Enumerable T where

enumerate = pay (pure L
 \oplus pure $S \circledast$ enumerate

 \oplus pure $B \circledast$ enumerate \circledast enumerate)

 $[(0,[]),$
 $(1,[L]),$
 $(1,[S L]),$

(4,[S(S(SL)),S(BLL),BL(SL),B(SL)L])

• • •

(2,[S(SL),BLL]),

```
sel:: E a \rightarrow Part \rightarrow F a
sel(pay e) 0 = empty_{r}
sel(pay e)(p+1) = sel e p
sel\ empty\ p\ =\ empty_r
sel(sing x) 0 = sing_{E} x
sel(sing x)(p+1) = empty_{E}
sel(a \oplus b) p
                         = sel a p \oplus_{\mathsf{F}} sel b p
                         = \bigoplus (sel\ a\ k \otimes_{\scriptscriptstyle \mathsf{F}} sel\ b\ (p-k))
sel(a \otimes b) p
                             k=0
```

Functional finite sequences

```
data Finite a = F

\{ card_F :: Integer

, index_F :: Integer \rightarrow a \} -- between 0 and card_F - 1
```

Functional finite sequences

```
Finite a \approx [a]

card_F \approx length

index_F \approx (!!)
```

 $empty_{_F}$:: Finite a $empty_{_F} = F \ 0 \ (\lambda i \rightarrow \bot)$

$$sing_F :: a \rightarrow Finite\ a$$
 $sing_F\ a = F\ 1\ one\ {\bf where}$
 $one\ 0 = a$
 $one\ _ = \bot$

$$empty_{_F} \approx []$$
 $sing_{_F} \approx (:[])$

 $map_F :: (a \rightarrow b) \rightarrow Finite \ a \rightarrow Finite \ b$

 $map_{F} f (F c ix) = F c (f \circ ix)$

$$(\bigoplus_F)$$
:: Finite $a \to F$ inite $a \to F$ inite a
 $f1 \bigoplus_F f2 = F$ car ix **where**
 $car = card_F f1 + card_F f2$
 $ix \ i = \mathbf{if} \ i < card_F f1$
 $\mathbf{then} \ index_F f1 \ i$
 $\mathbf{else} \ index_F f2 \ (i - card_F f1)$

$$(\oplus_F) \approx (++)$$

 $(xs ++ ys) !! i \equiv xs !! i$ for $i < length xs$
 $\equiv ys !! (i - length xs)$ for $i \geqslant length xs$

$$f1 \otimes_F f2 = F \ car \ ix \ where$$

 $car = card_F \ f1 * card_F \ f2$
 $ix \ i = let \ (d, m) = (i 'divMod' \ card_F \ f2)$
 $in \ (index_F \ f1 \ d, index_F \ f2 \ m)$

$$[(x,y) | x \leftarrow xs, y \leftarrow ys] \approx xs \otimes_F ys$$
$$[(x,y) | x \leftarrow xs, y \leftarrow ys] !! i \equiv$$
$$(xs !! div i lys, ys !! mod i lys)$$
$$\mathbf{where} \ lys = length \ ys$$

Infinite enumerations

```
newtype Enumerate a = E \{ parts :: [Finite a] \}
```

- -- Actual implementation also store
- -- the reversals of all initial segments of the list

empty :: Enumerate a

 $empty = E (repeat empty_{\epsilon})$

 $sing :: a \rightarrow Enumerate a$ sing $a = E([sing_{E} \ a] ++ repeat \ empty_{E})$

 $fmap :: (a \rightarrow b) \rightarrow Enumerate \ a \rightarrow Enumerate \ b$ $fmap \ f = E \circ fmap \ (map_r \ f) \circ parts$

(⊕):: Enumerate $a \to E$ numerate $a \to E$ numerate $a \to E$ numerate $a \to E$ 0 (⊕) $e1 \ e2 = E \ zipWith (⊕_F) \ (parts \ e1) \ (parts \ e2)$

Each part of a product is a finite convolution sum:

```
(e_1 \otimes e_2) ! n
    \equiv e_1 ! 0 \otimes_{\scriptscriptstyle F} e_2 ! n
          \bigoplus_{\Gamma} e_1 ! 1 \otimes_{\Gamma} e_2 ! (n-1)
          ⊕_ ...
                e_1 ! n \otimes_{\epsilon} e_2 ! 0
     \equiv concat_F
        (zipWith (\otimes_r)
             (parts e_1)
             (reverse (take n (parts e_2))))
```

$$e! n = parts e!! n$$

 $concat_F = foldr(\oplus_F) empty_F$

(⊗):: Enumerate
$$a \rightarrow$$
 Enumerate $b \rightarrow$ Enumerate (a, b)
 $xs \otimes ys = E \$ map (conv (parts xs))$
(reversals (parts ys))

(reversals (parts ys))

 $conv :: [Finite a] \rightarrow [Finite b] \rightarrow Finite (a, b)$ $conv \ xs \ ys = concat_F \ (zipWith \ (\otimes_F) \ xs \ ys)$

reversals :: $[a] \rightarrow [[a]]$ reversals = go [] where go rev (x:xs) =**let** rev' = x: revin rev': go rev' xs

Assigning costs

```
pay :: Enumerate a \rightarrow Enumerate a
pay e = E \$ empty_F : parts e
```

Assigning costs

pay :: Enumerate $a \rightarrow Enumerate a$ pay $e = E \$ empty_e$: parts e

Guarded recursion operator:

fix pay ≡ empty

Examples

 $enum_{Bool} = pay \$ sing False \oplus sing True$

*Main> $map\ card_F\ (parts\ enum_{Bool})$ $[0,2,0,0,0,0\dots$

Examples

```
data N = Ze \mid Su \ N \ deriving Show enum_N = pay $ sing Ze <math> \oplus  fmap Su \ enum_N
```

```
*Main> map\ card_F\ (parts\ enum_N)
[0,1,1,1,1,1,1,1,1,1,1...
```

Examples

```
enum_{[Bool]} = pay $
sing [] \oplus fmap (uncurry (:)) (enum_{Bool} \otimes enum_{[Bool]})

*Main> map \ card_F \ (parts \ enum_{[Bool]})
[0,1,0,2,0,4,0,8,0,16,0,32,0,64,0,128,0,256,0...]
```

Absolute indexing

index' :: Enumerate $a \rightarrow Integer \rightarrow a$ $index' = index_F \circ concat_F \circ parts$

Absolute indexing

```
index' :: Enumerate \ a \rightarrow Integer \rightarrow a index' = index_F \circ concat_F \circ parts

Consider: length \ (repeat \ ()) \equiv \bot repeat \ ()!! \ i \equiv ()
```

Random selection

uniform_F:: Finite $a \rightarrow Gen\ a$ uniform_F $f = liftM\ (index_F\ f)\ (choose\ (0, card_F\ f - 1))$

Random selection

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uniform :: Enumerate $a \rightarrow Int \rightarrow Gen \ a$ uniform $e \ n = uniform_F \ concat_F \ take \ n \ parts \ e$

- An algebra of Functional Enumerations
 - Concise definition (complete implementation in slides)
 - Nice algebraic properties
 - Automatically derivable for algebraic data types

- An algebra of Functional Enumerations
 - Concise definition (complete implementation in slides)
 - Nice algebraic properties
- Automatically derivable for algebraic data types
- Actually finds bugs
- Particularly useful for large abstract syntax tree types
 - Found errors in haskell-src-exts and template-haskell
 - Neither SmallCheck nor QuickCheck finds these bugs

Try it

Have some untested code lying around?

- cabal install testing-feat
- import Test.Feat

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Thank you!