

题目讲解

- Leetcode 207. Course Schedule

图搜索初探

先导课

- 递归 = 递 + 归
- 递归奥义——复制自己
 - 如何造纳米机器人

先导课

- 递归框架

```
• int robot(int x, int y)           // 机器人的输入
    if ( 边界条件 )                 // 什么时候不用造了 ( 自己就能干完 )
        return 0;

    int a = robot(x1, y1);           // 造一个小的自己帮忙干活
    int b = robot(x2, y2);           // 再造一个小的自己帮忙干活
    return a + b;                    // 自己要做的就是别人的成果组装起来
```

Leetcode 78


Outline

- 何谓图？
 - 图是描述世间万物关系的一种方式
 - 节点+边
- 隐式图
 - 状态（结点）不确定（明显）
 - 关系（边）不确定（明显）
 - 如何确定状态和关系（重点）

Outline

- 图搜索
 - 深度优先遍历 (DFS)
 - 广度优先遍历 (BFS)
- 隐式图搜索
 - N皇后问题、骑士游历问题、八数码

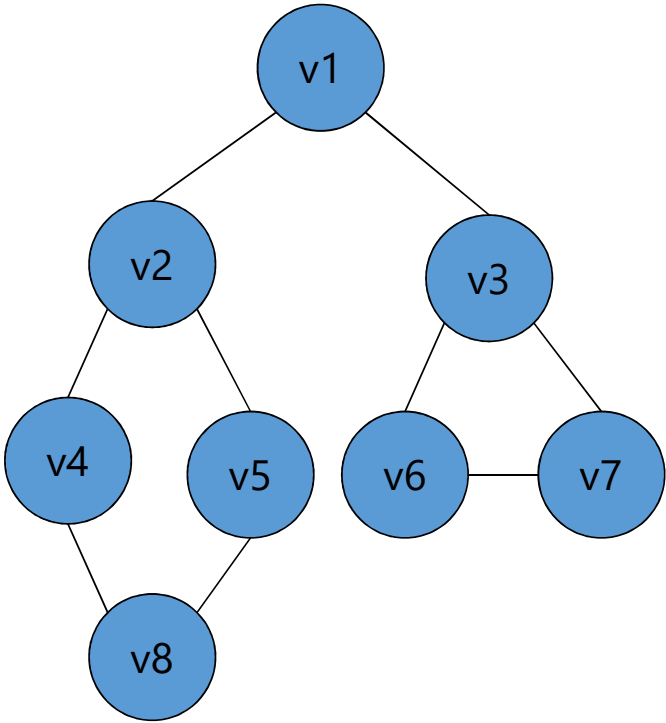
遍历：定义

- 按某种顺序访问 “” 中所有的节点
- 顺序
 - 深度优先（优先往深处走）
 - 广度优先（优先走最近的）
- 时间复杂度 $O(n + m)$
- 空间复杂度？

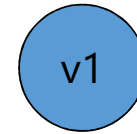
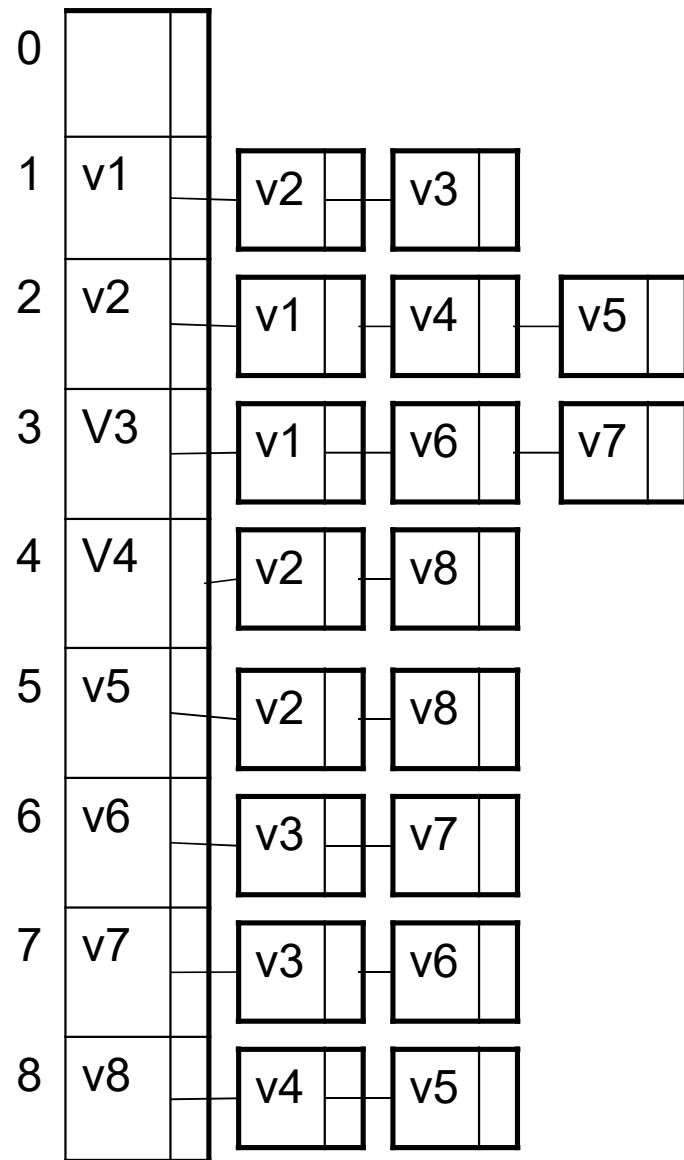
遍历

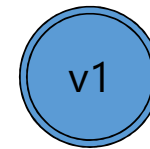
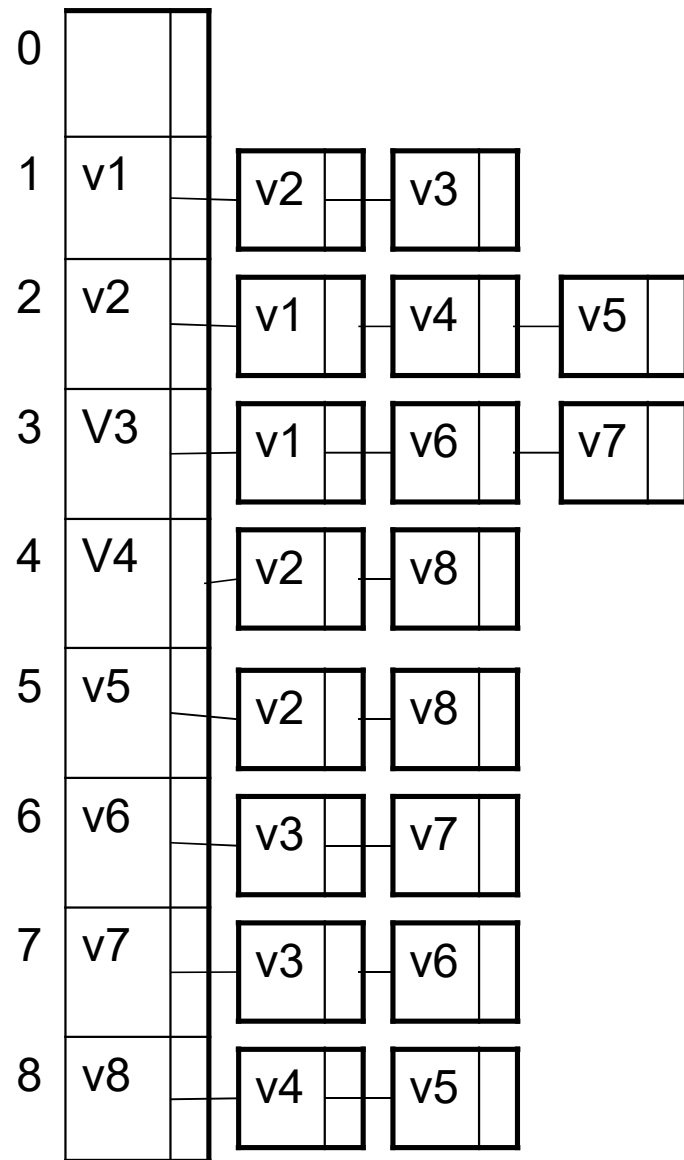
- 给出图G，要求求从入口v1访问到每一个点
- 两种遍历方式的数据结构
 - 栈（递归，深度优先）
 - 队列（广度优先）
- 广度优先找出的路径，经过节点数最少

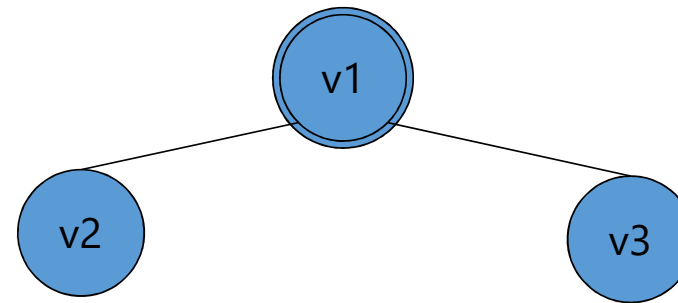
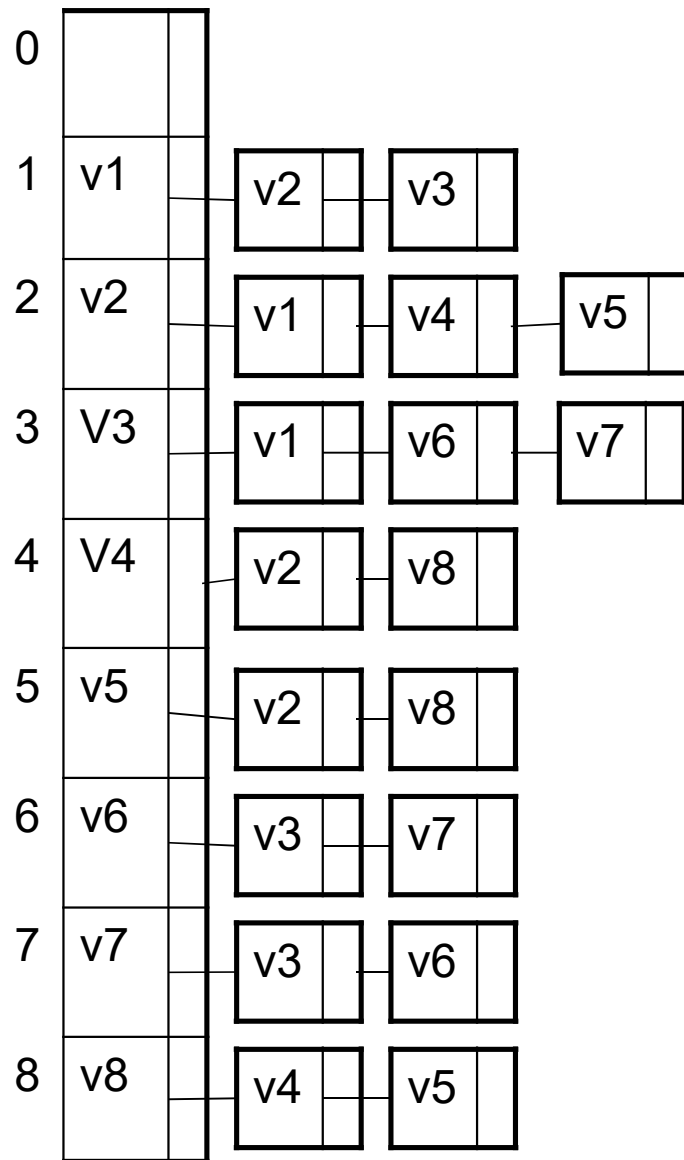
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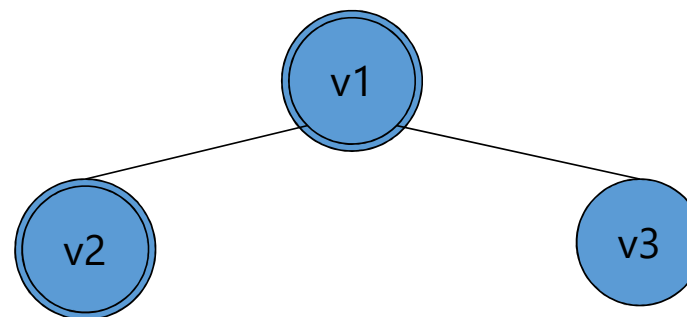
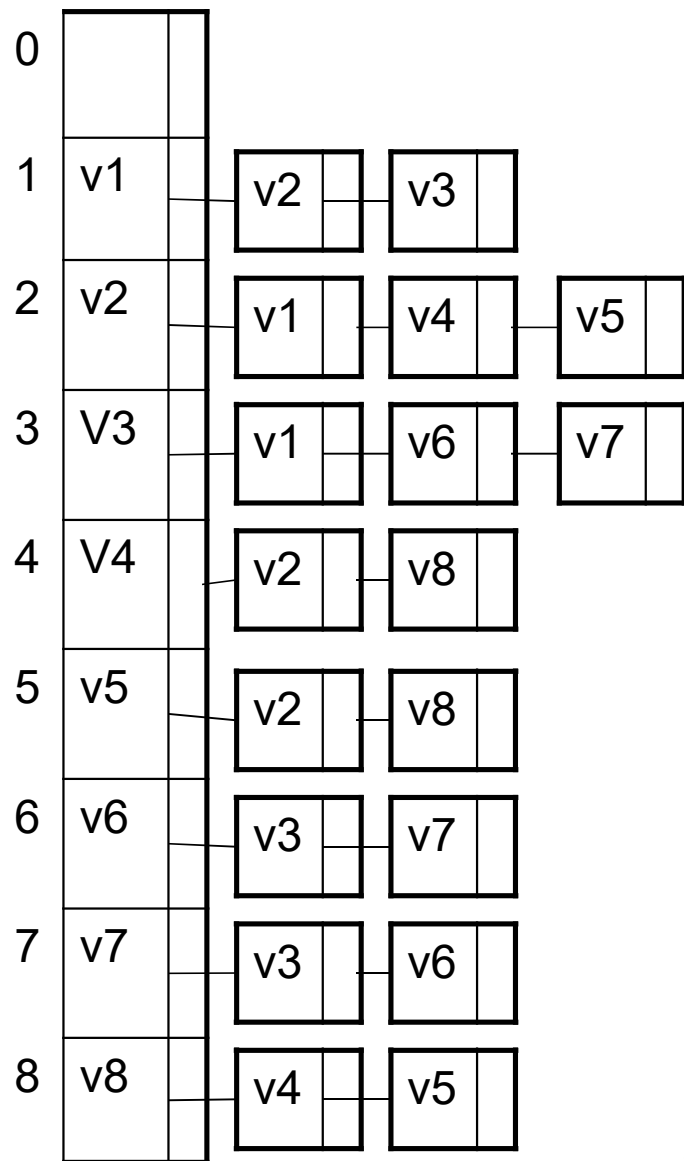


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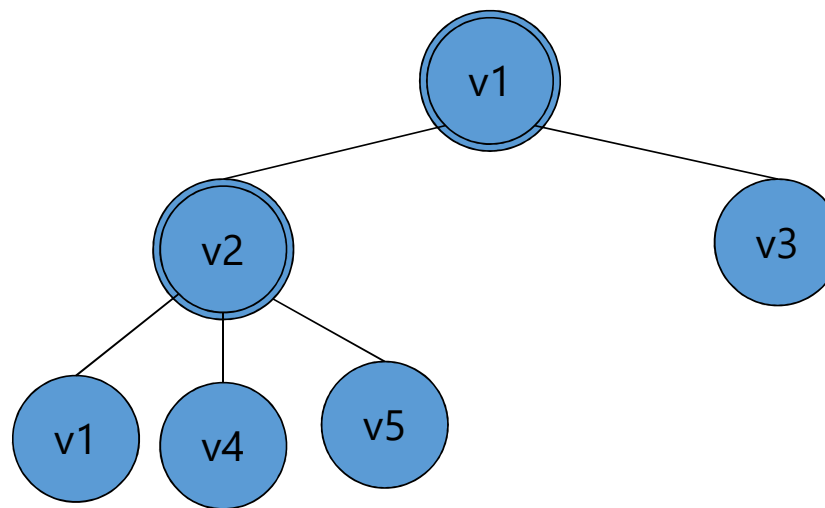


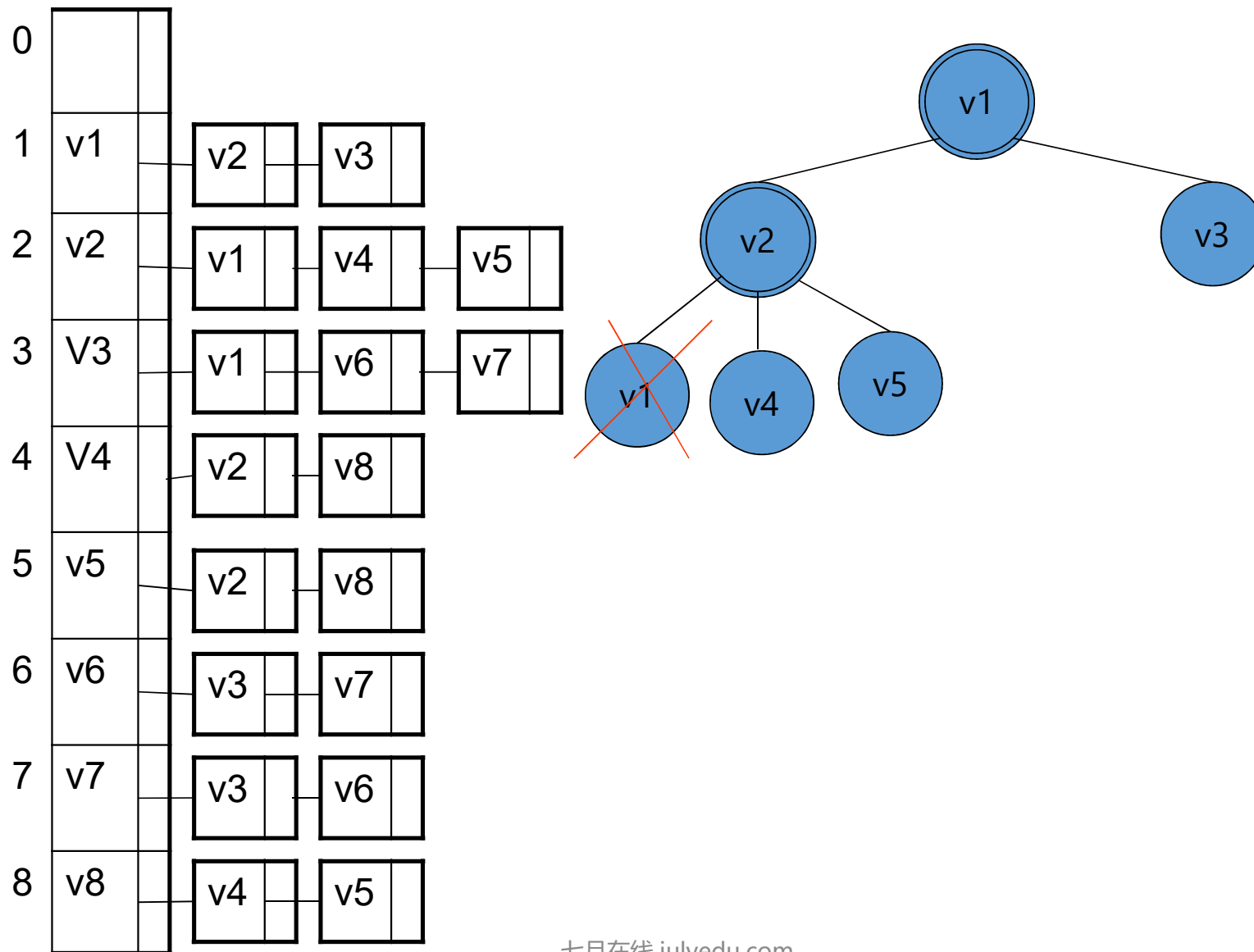


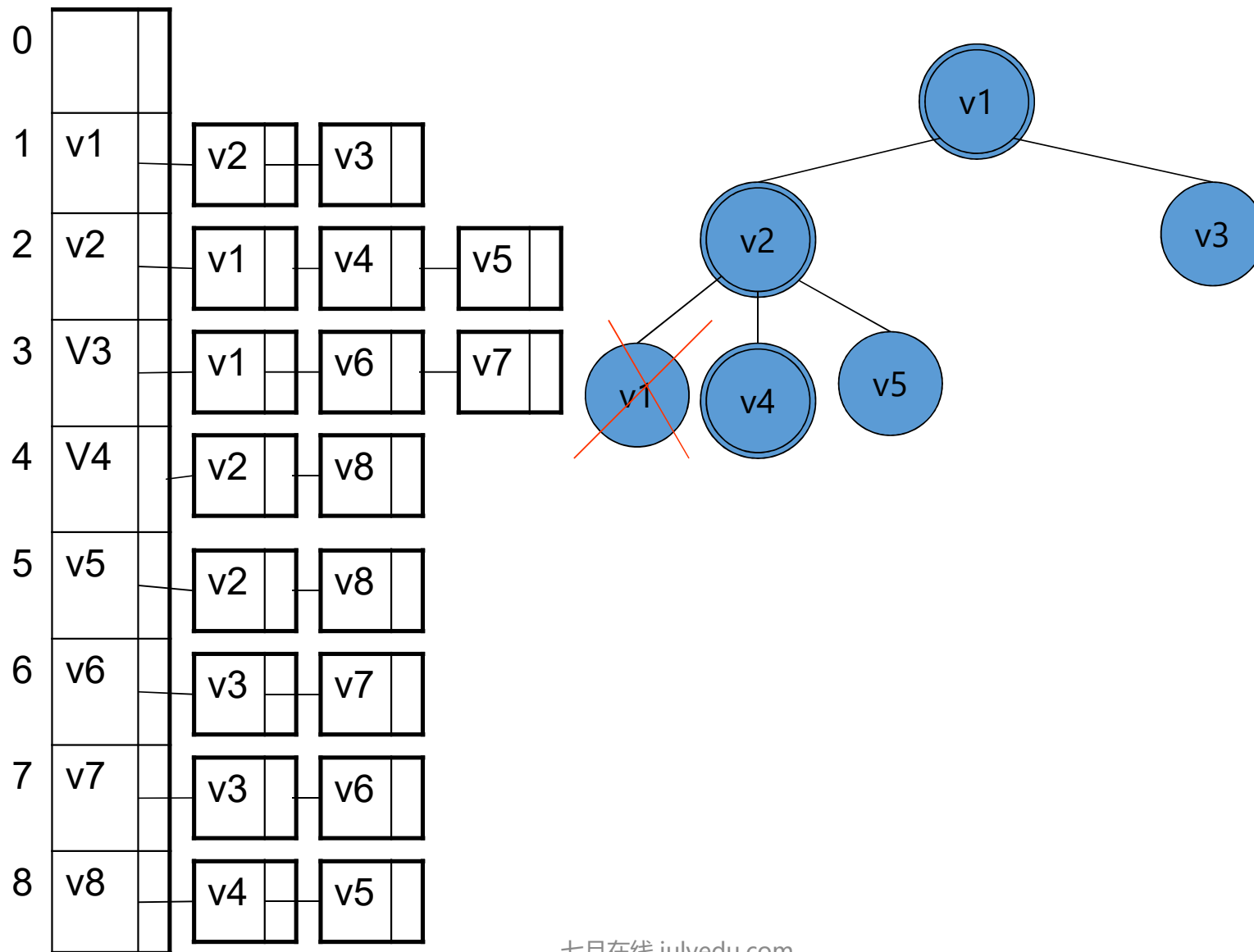


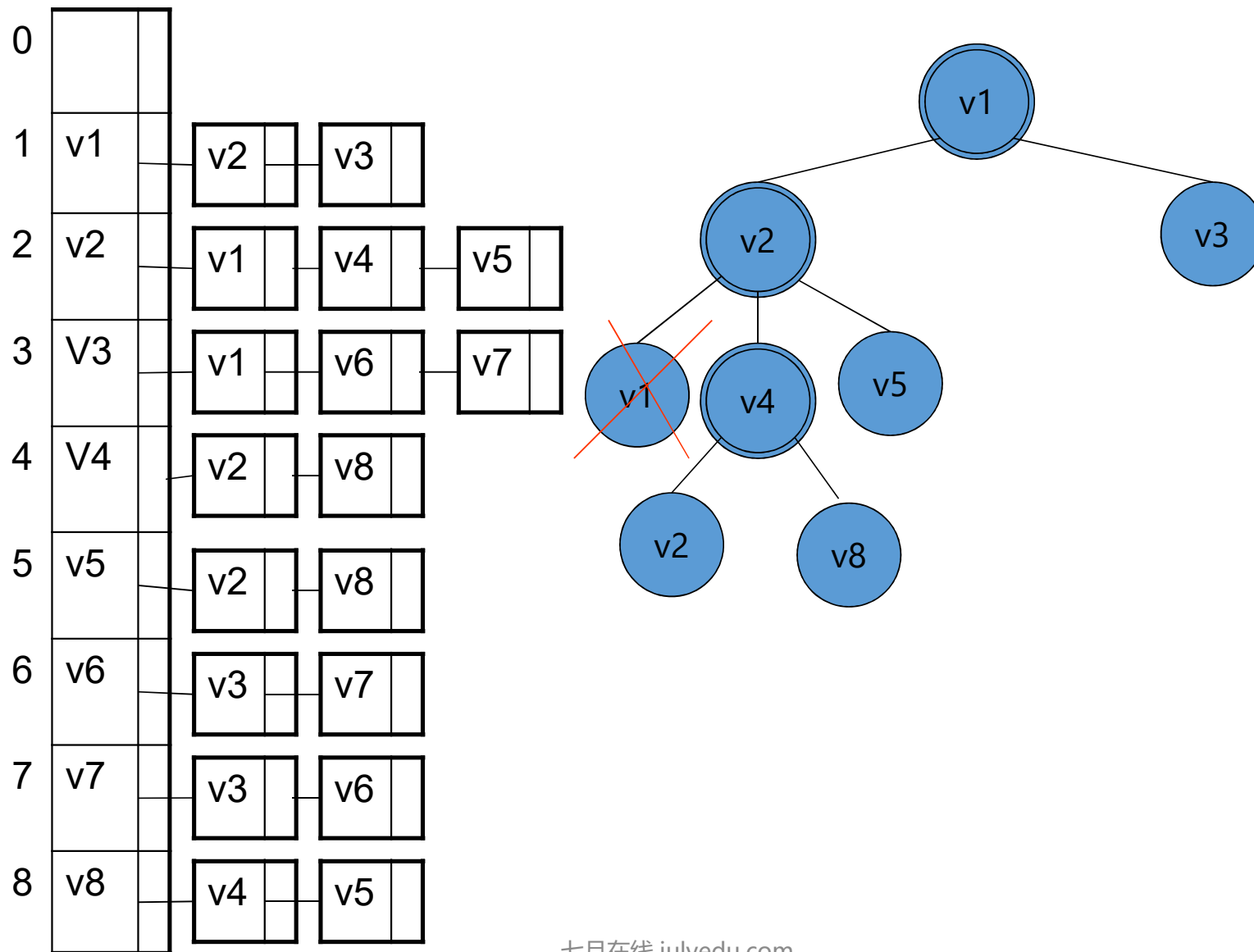


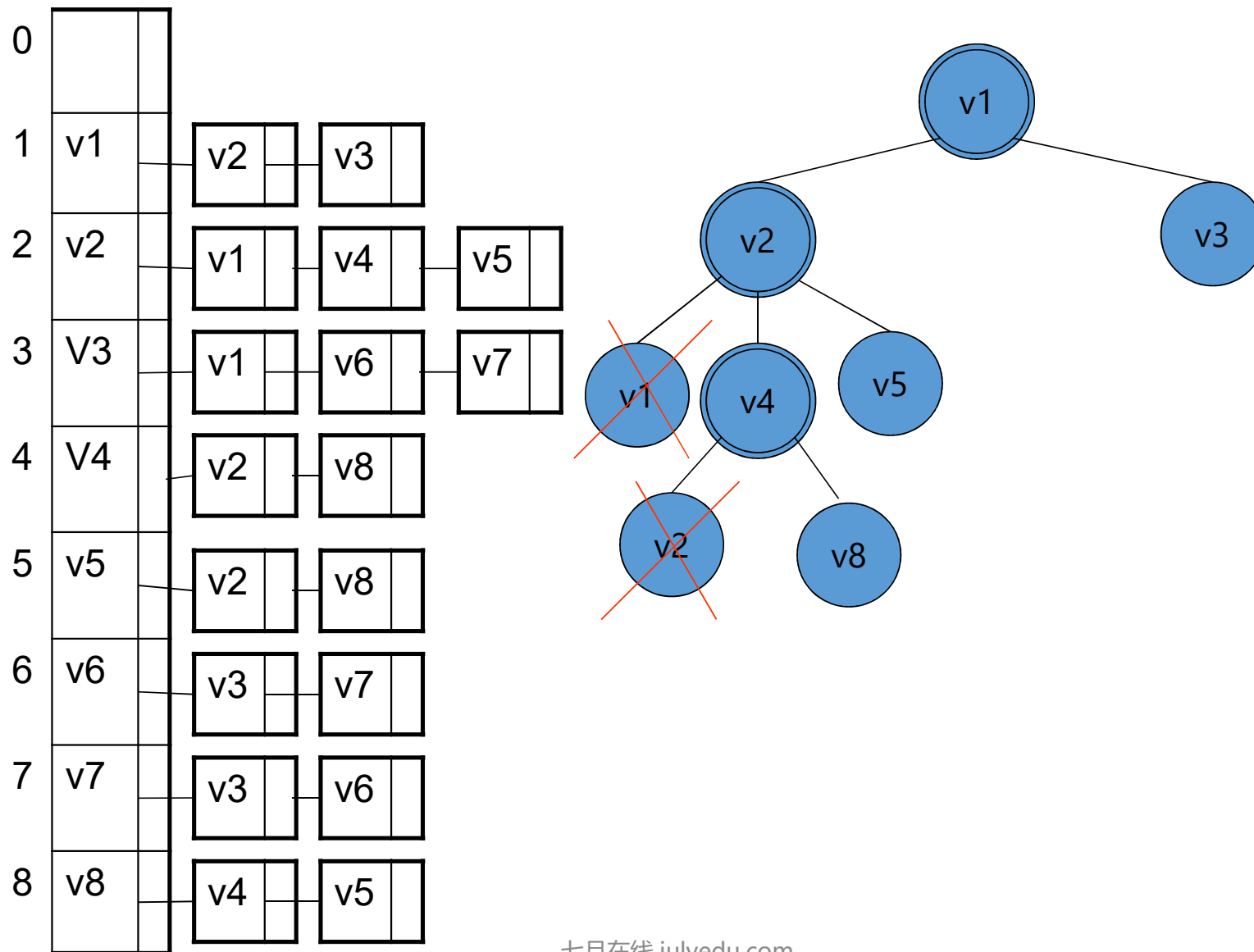
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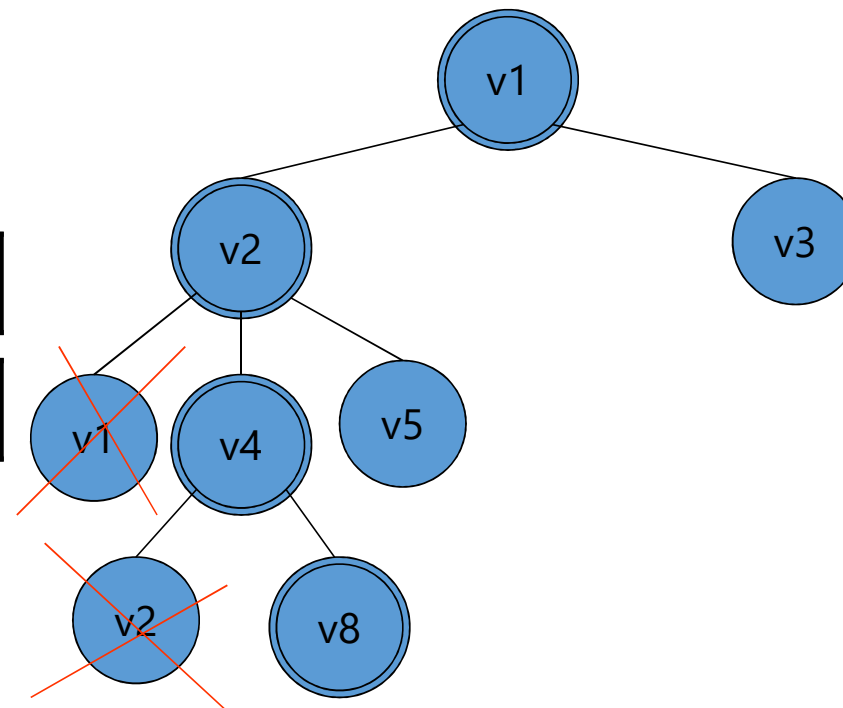


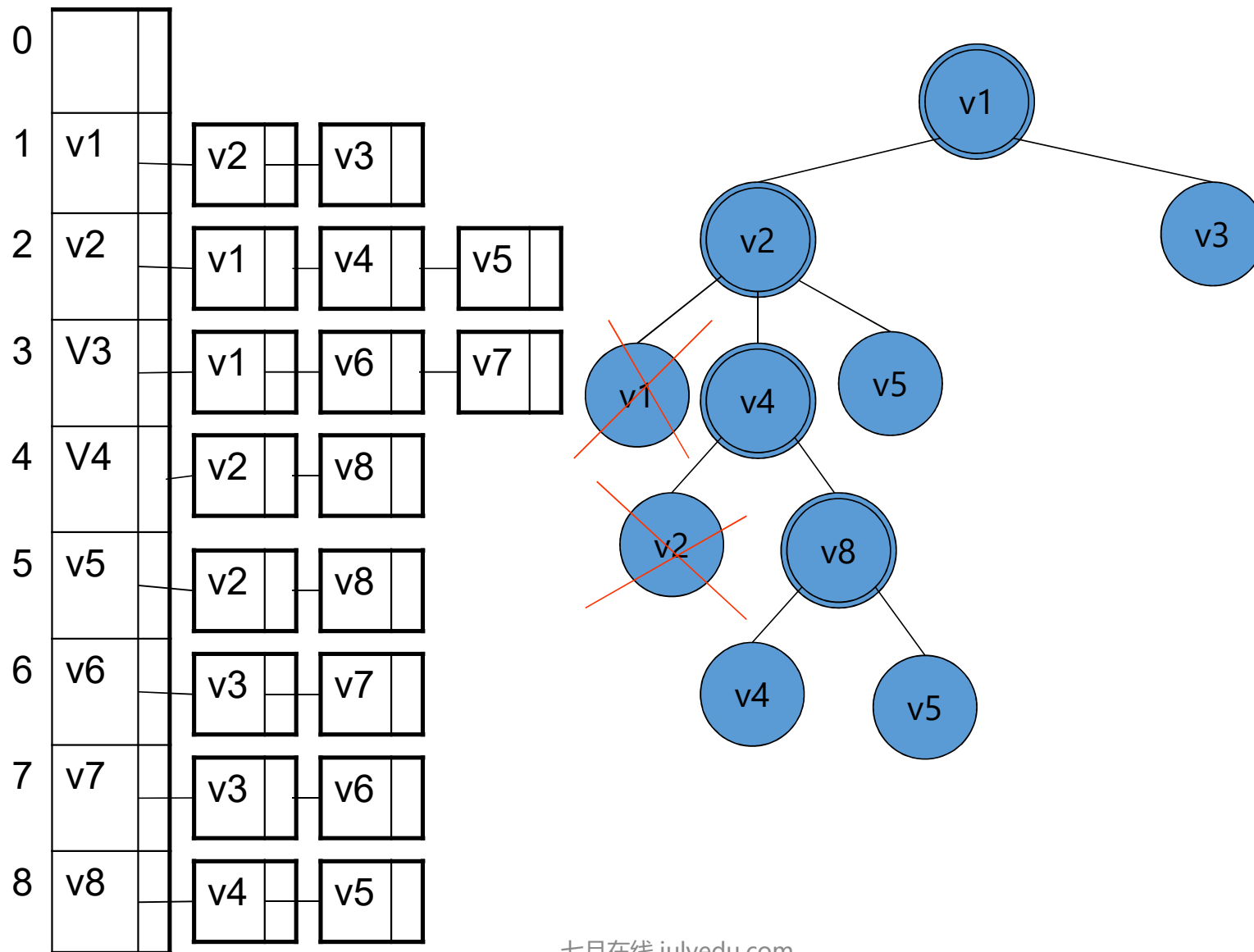


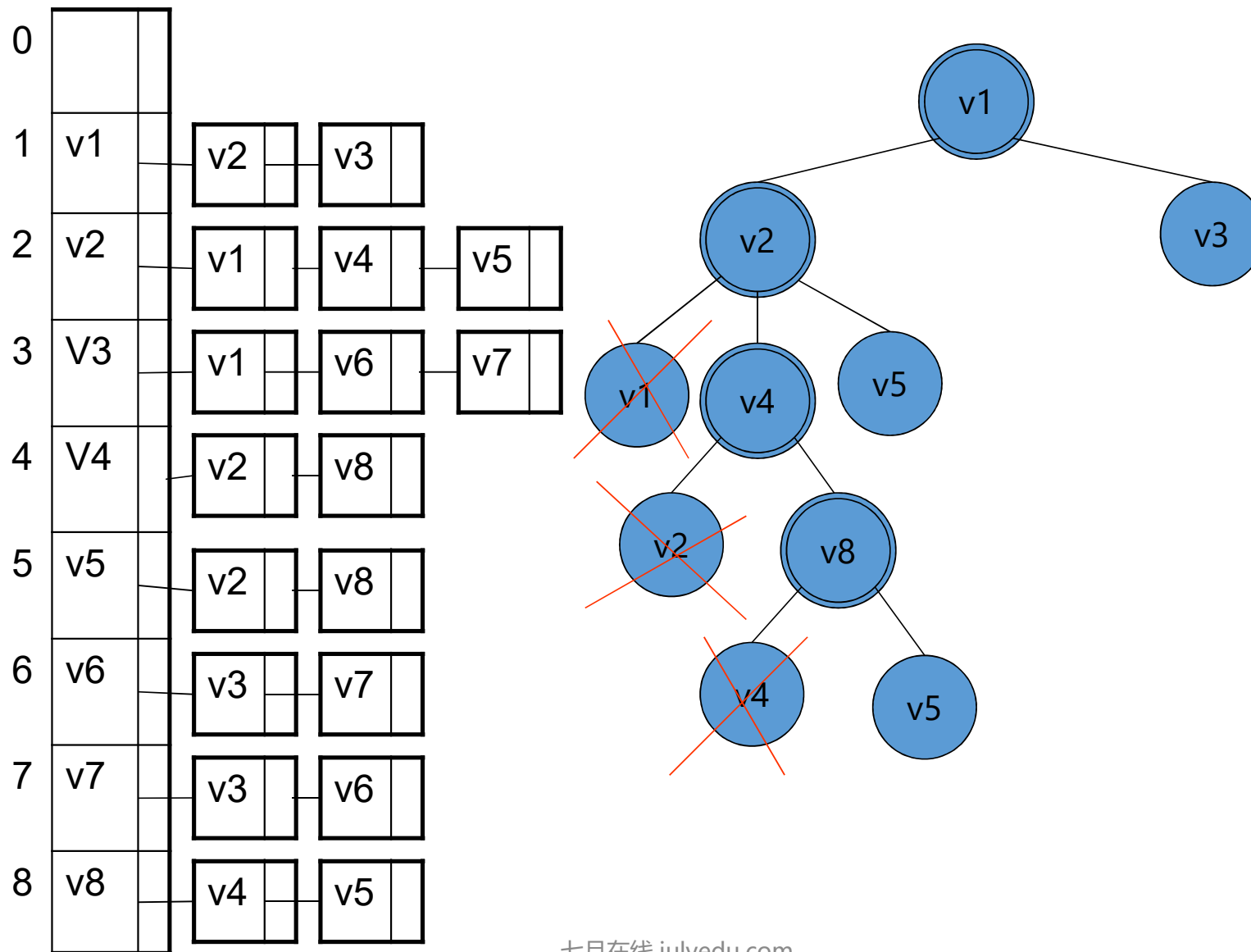


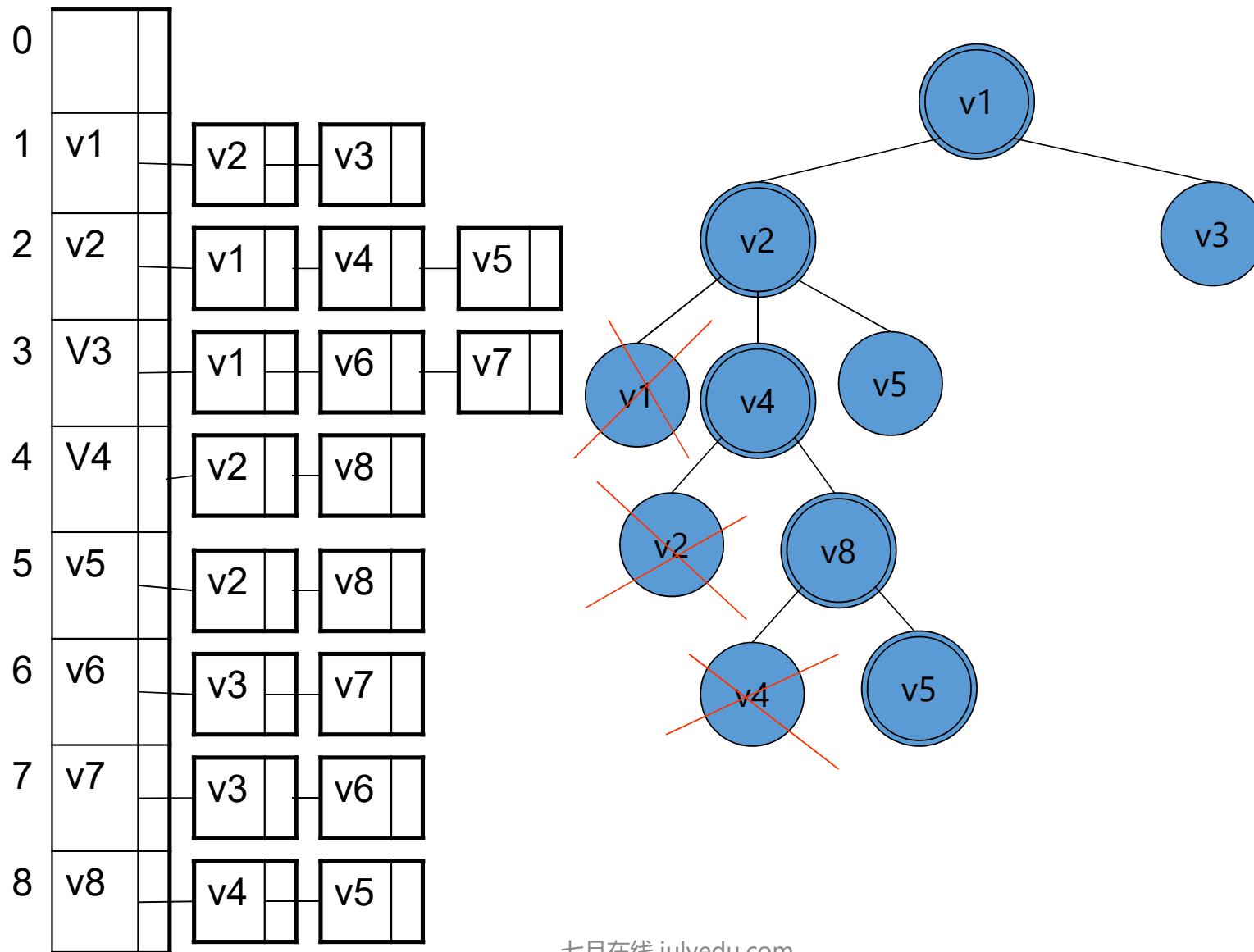


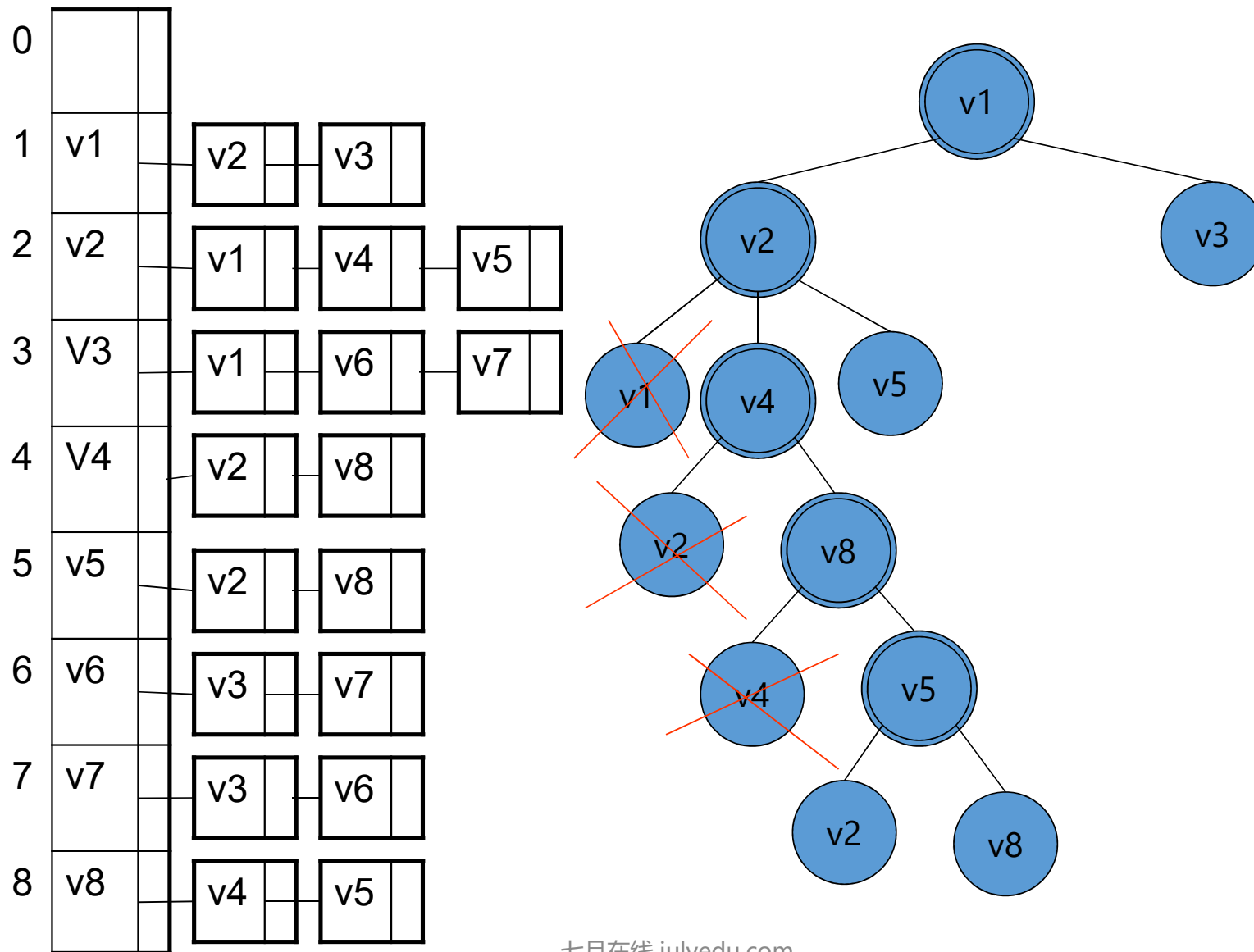
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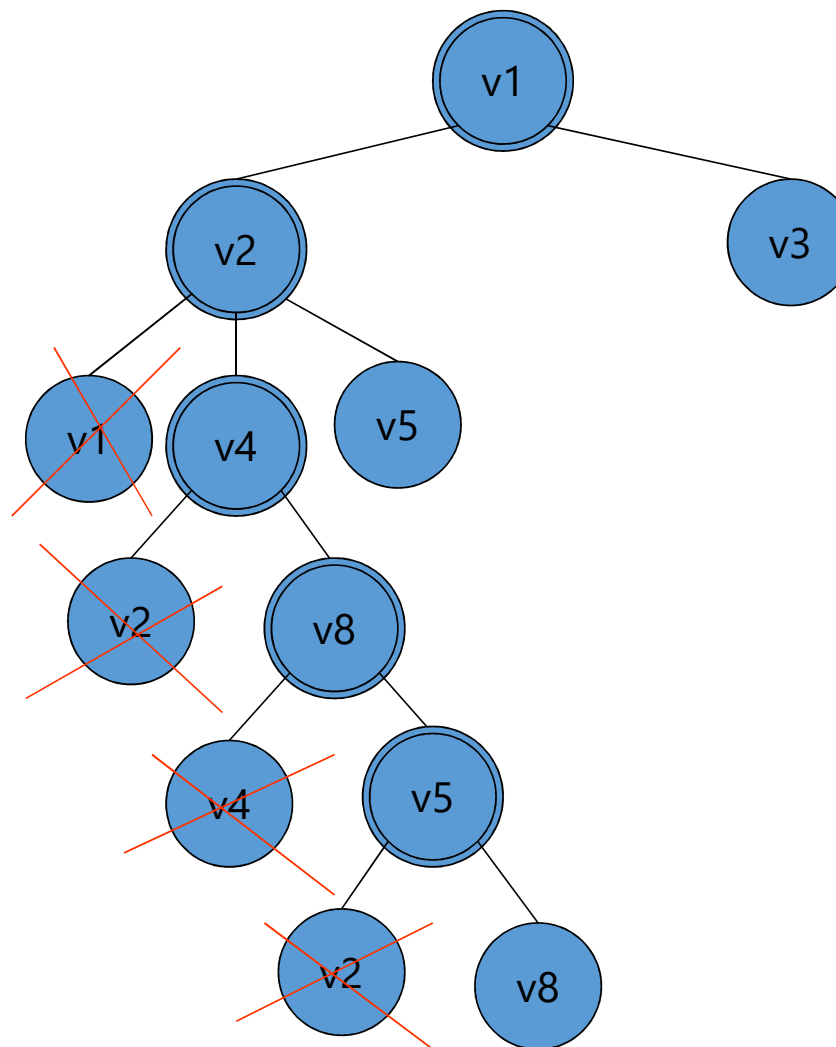




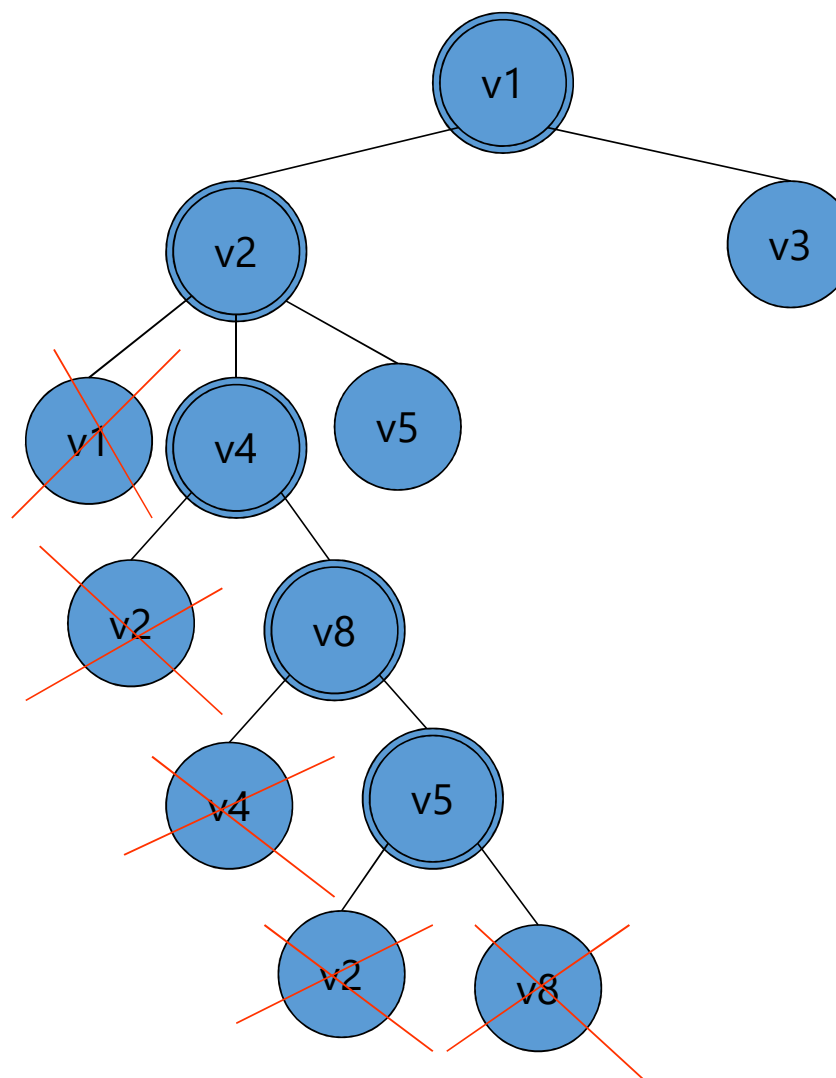




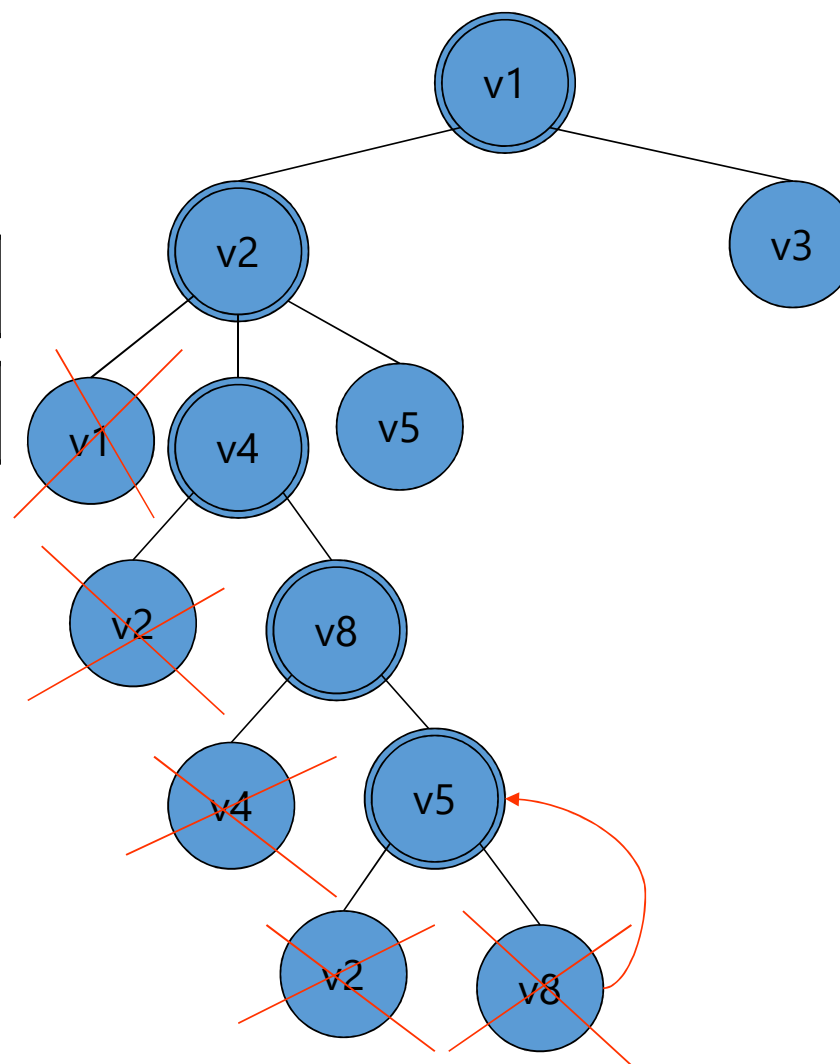
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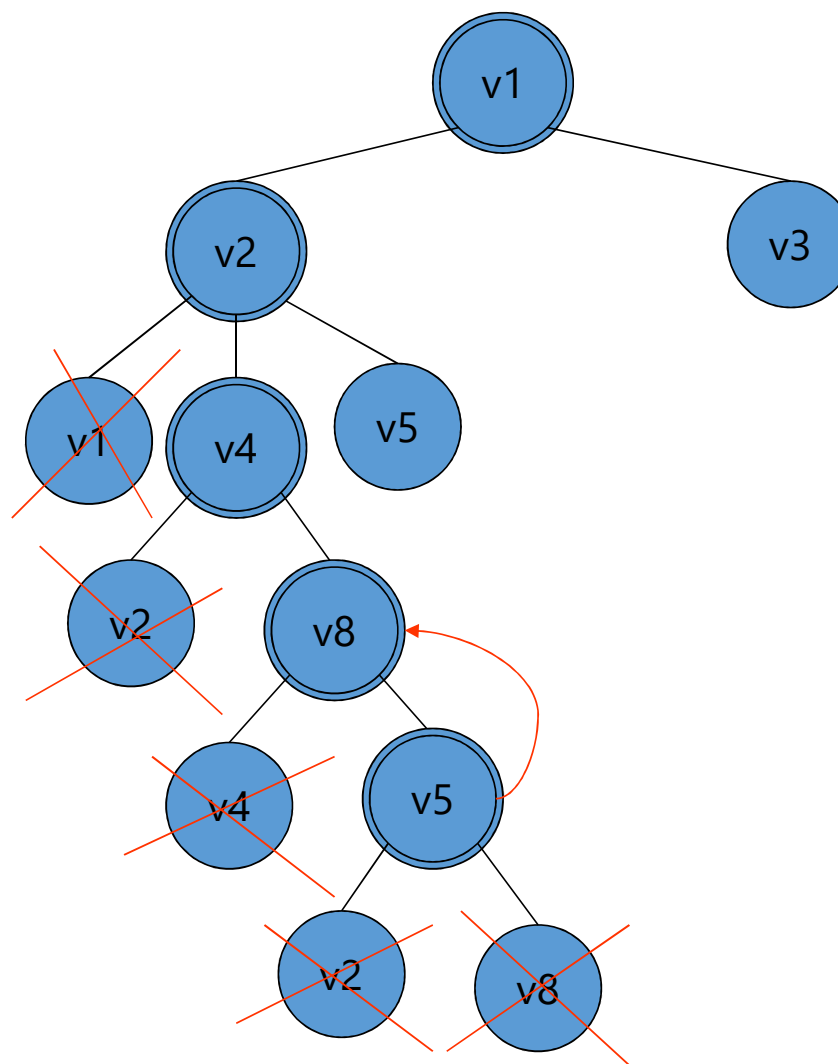
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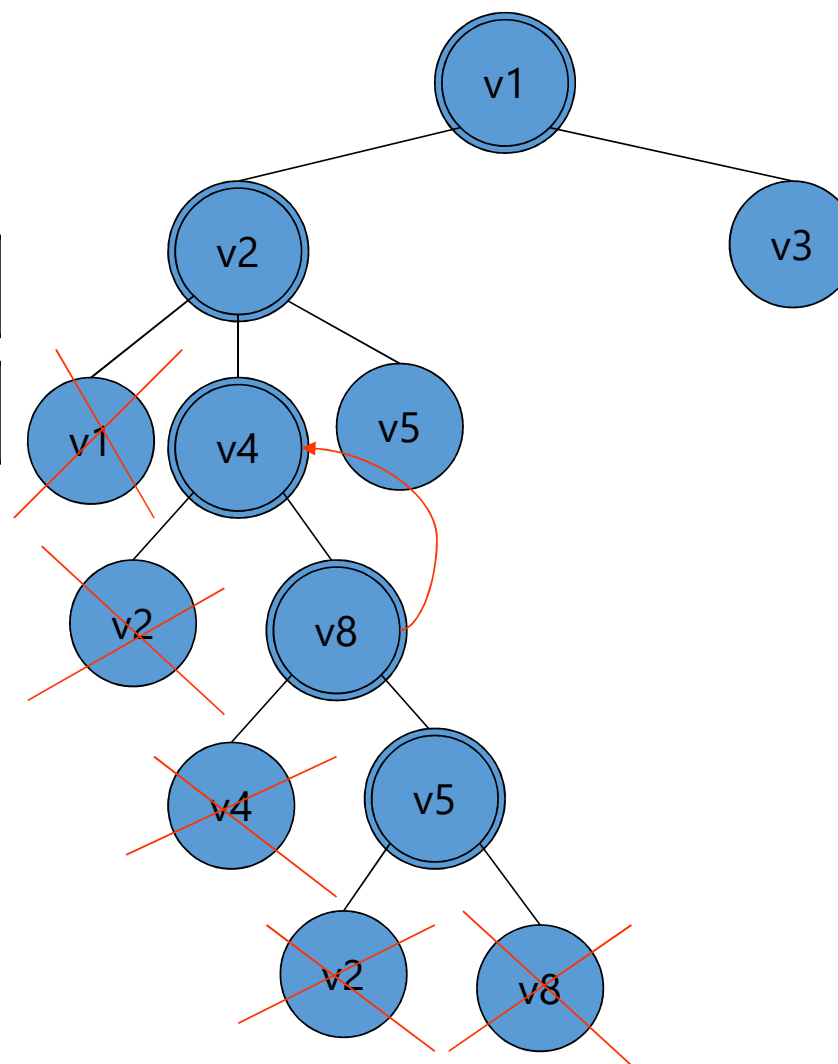
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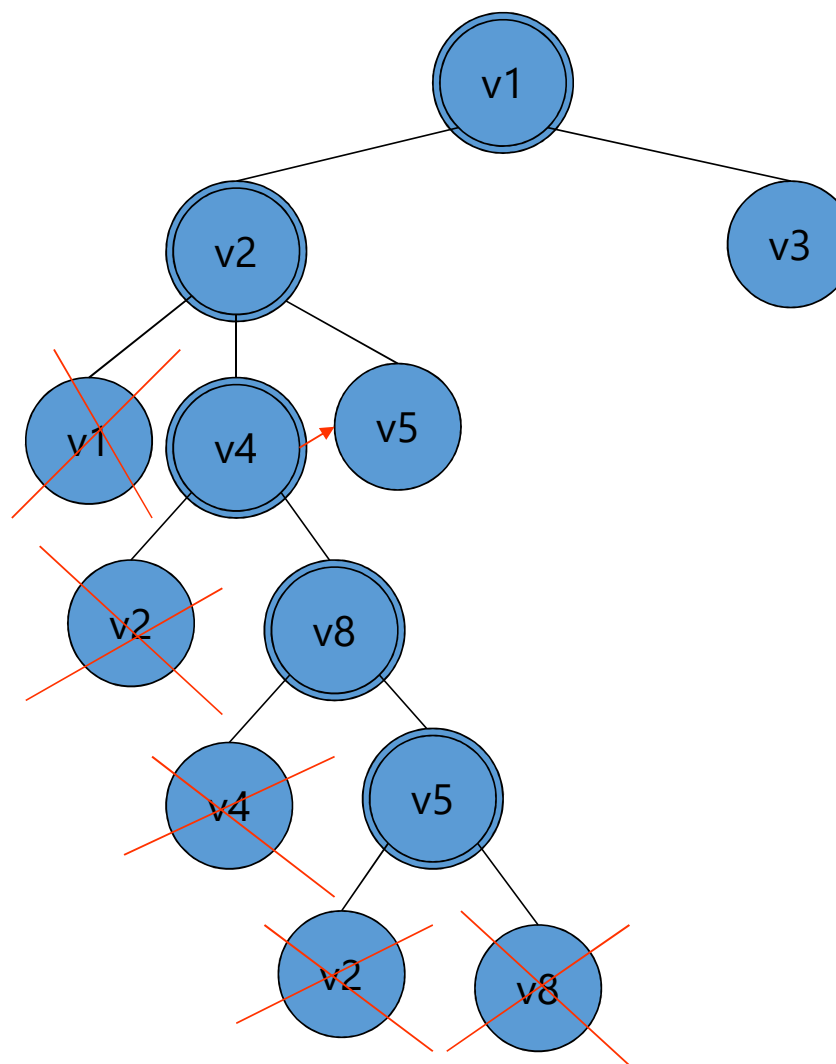
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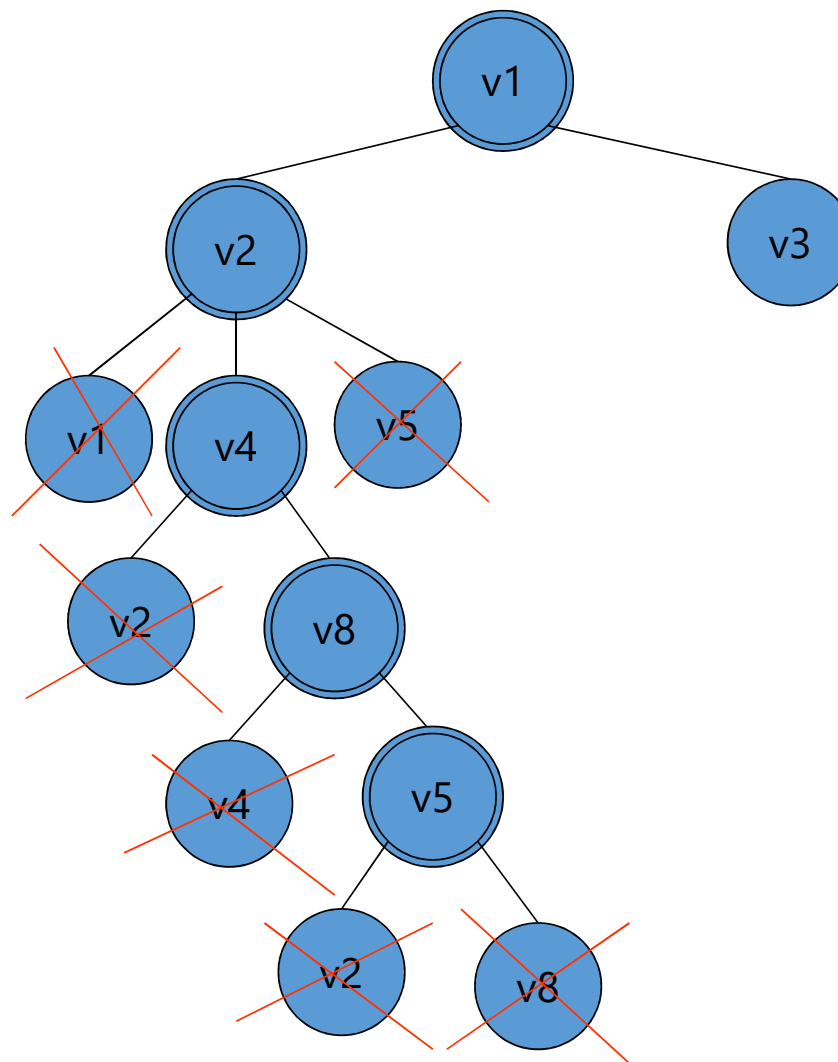
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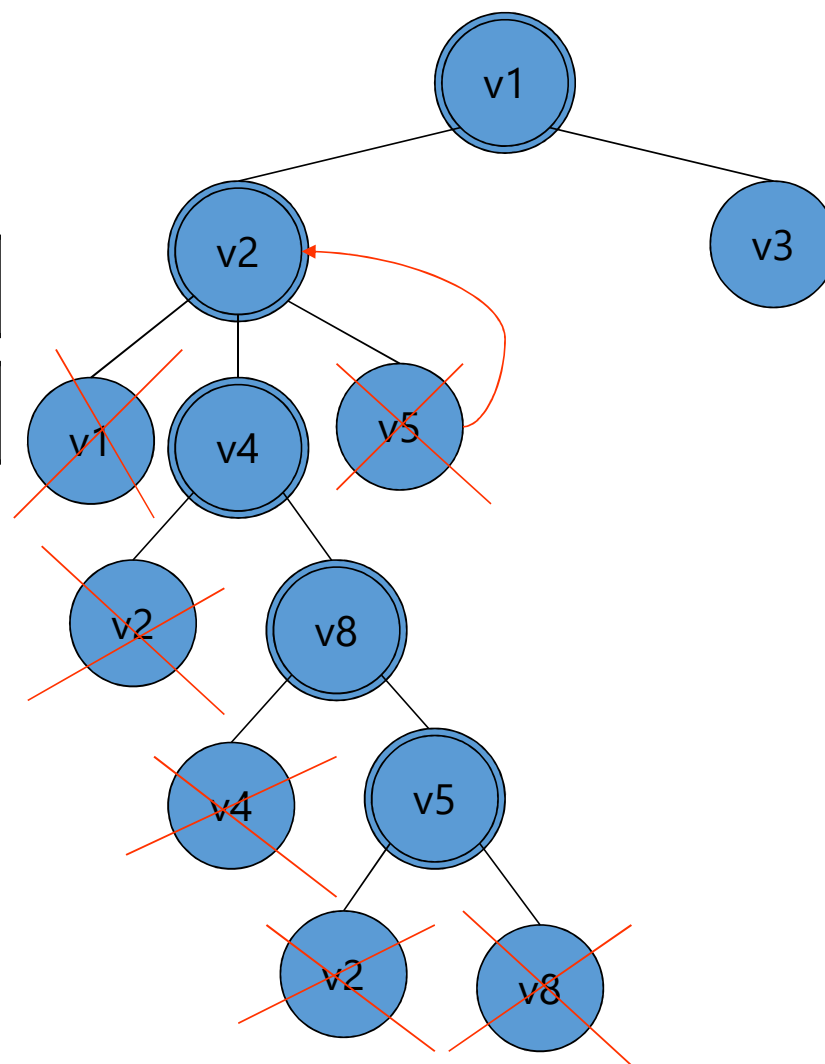
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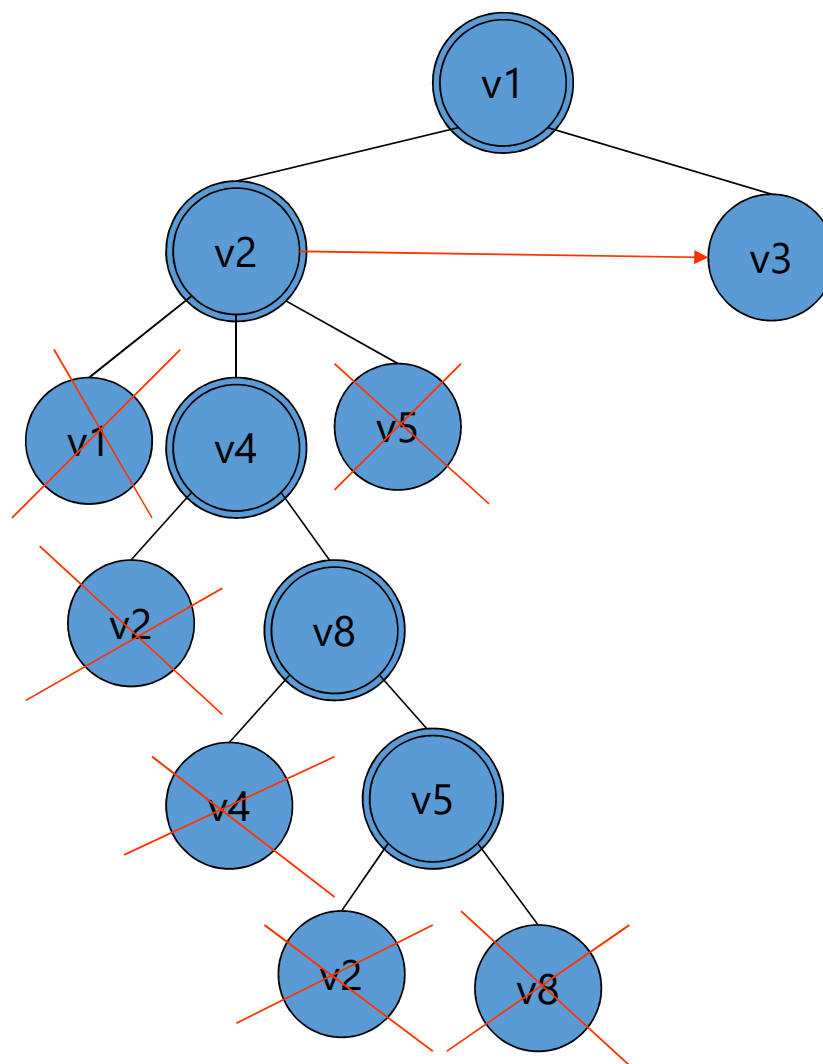
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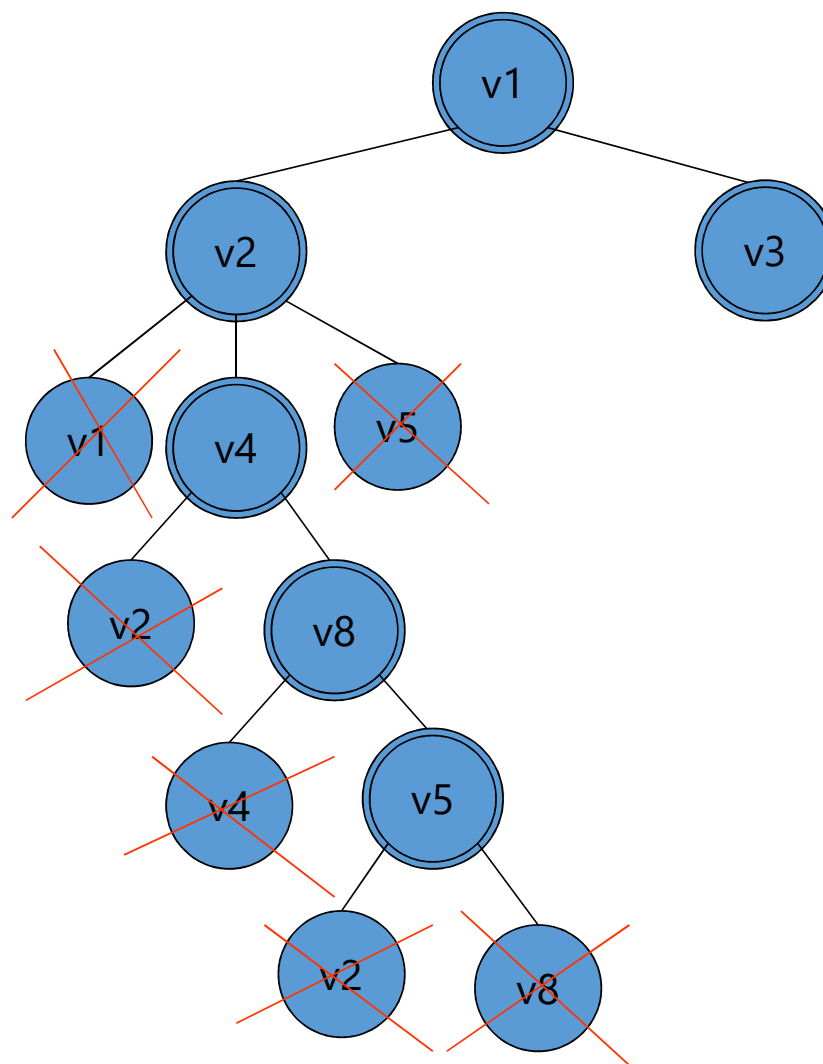
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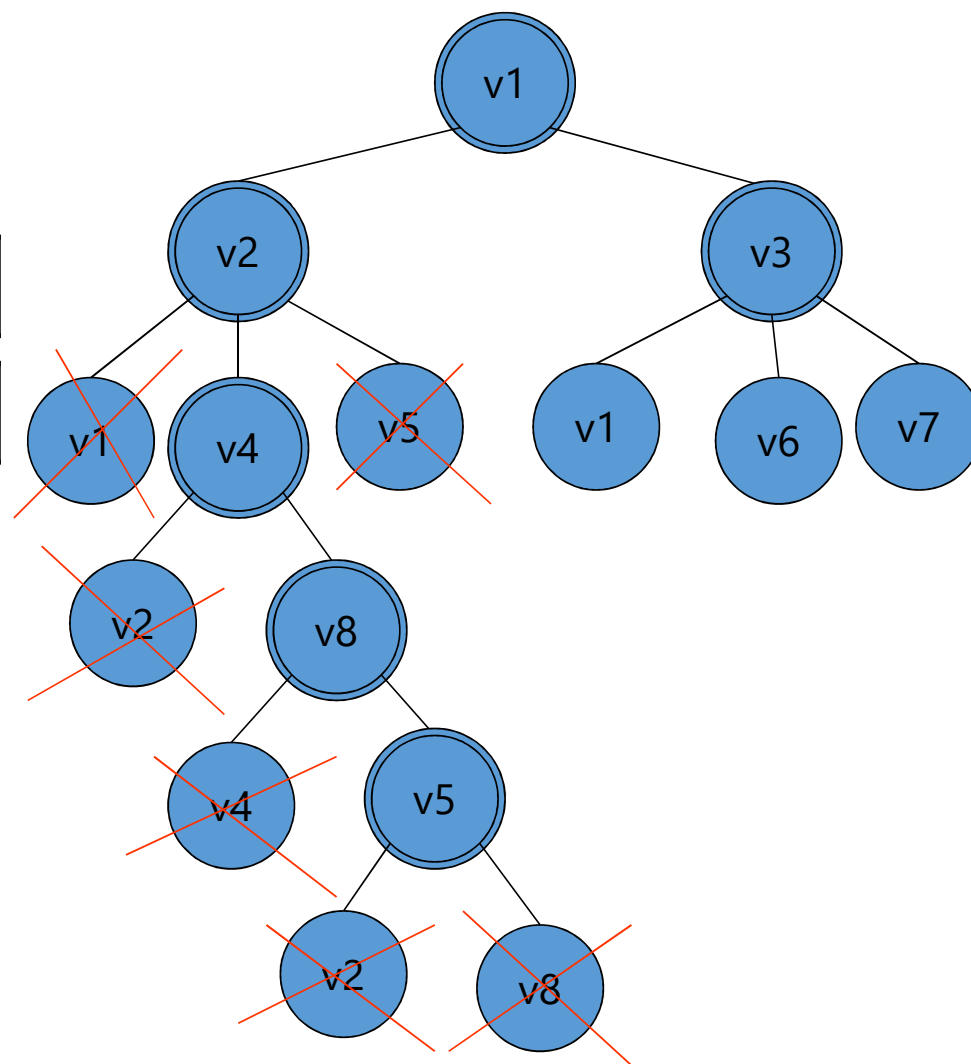
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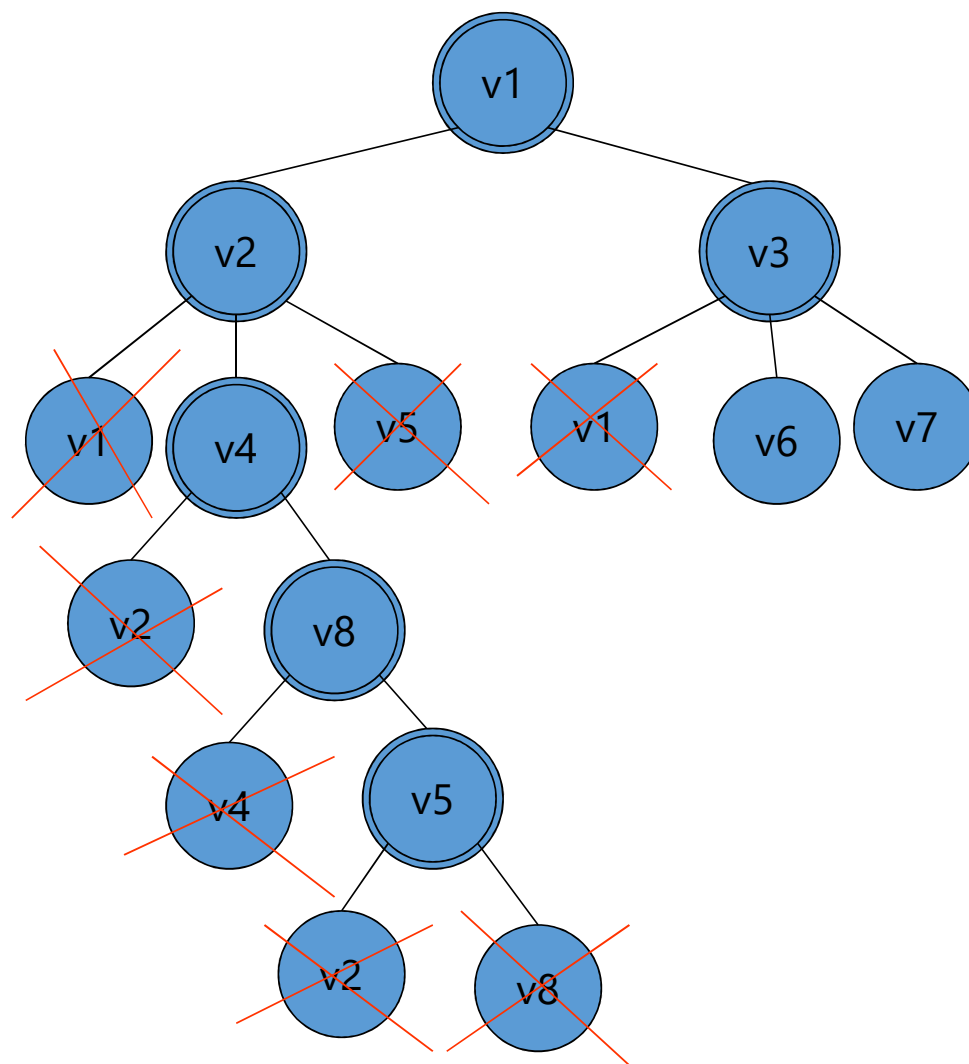
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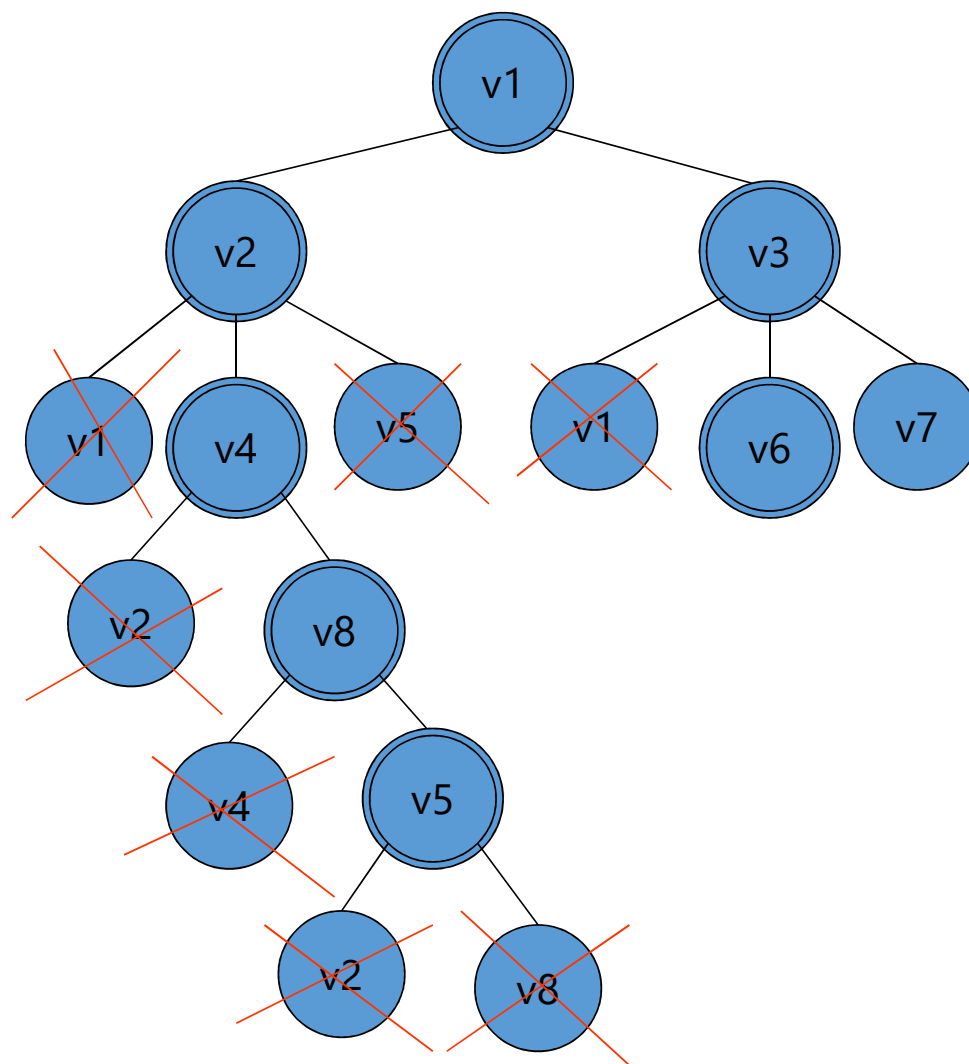
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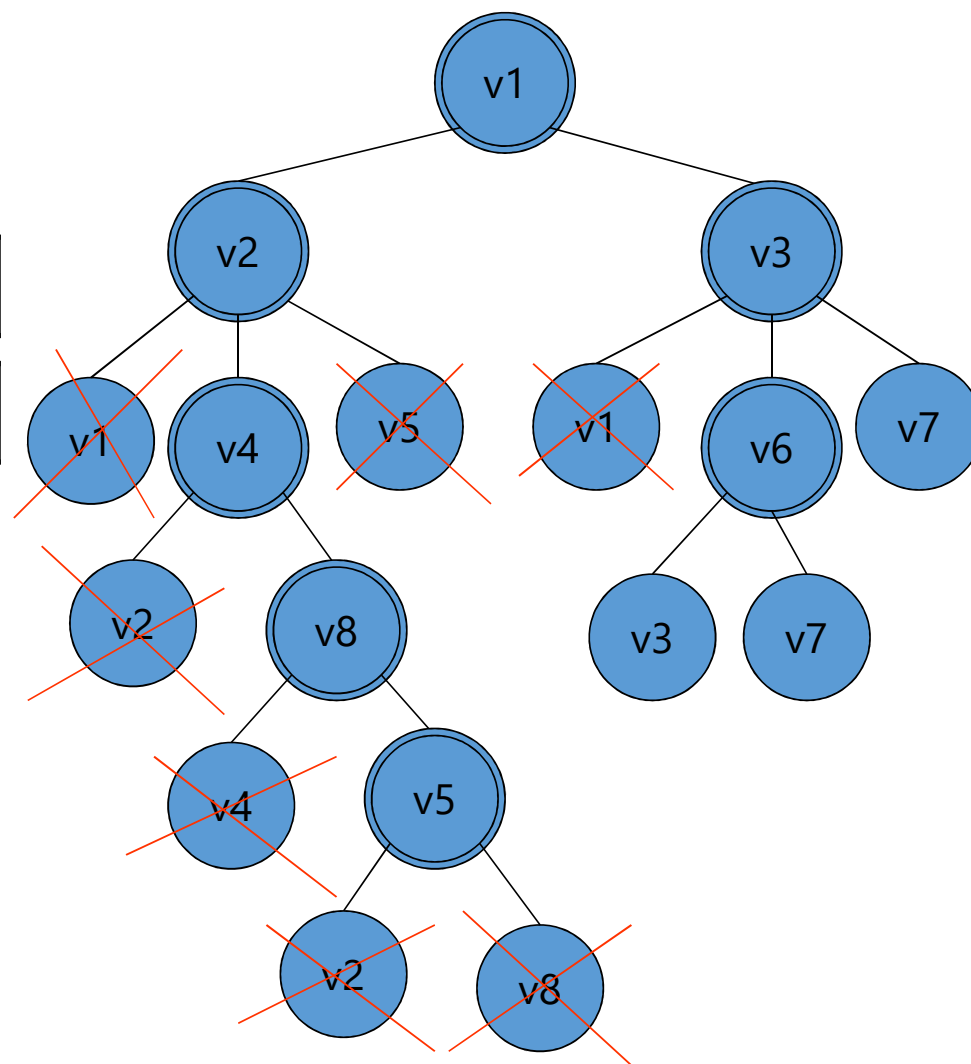
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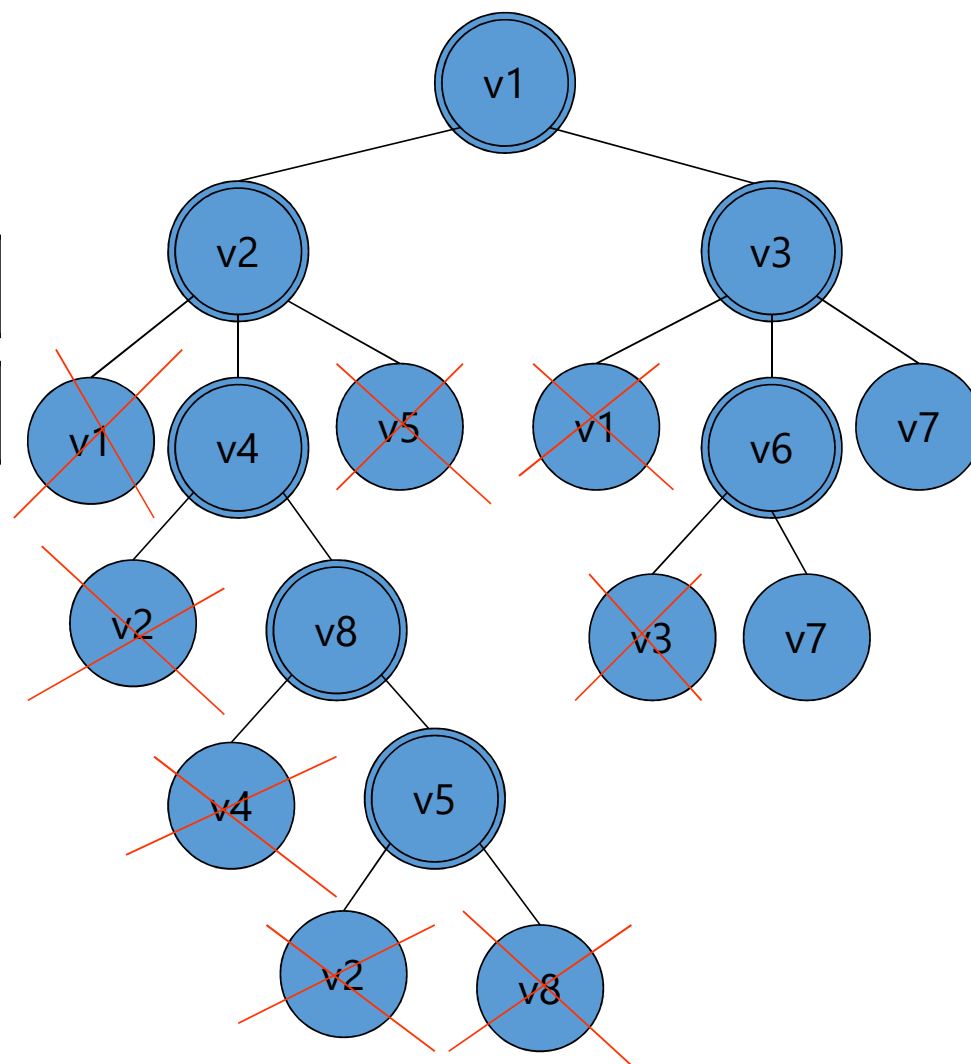
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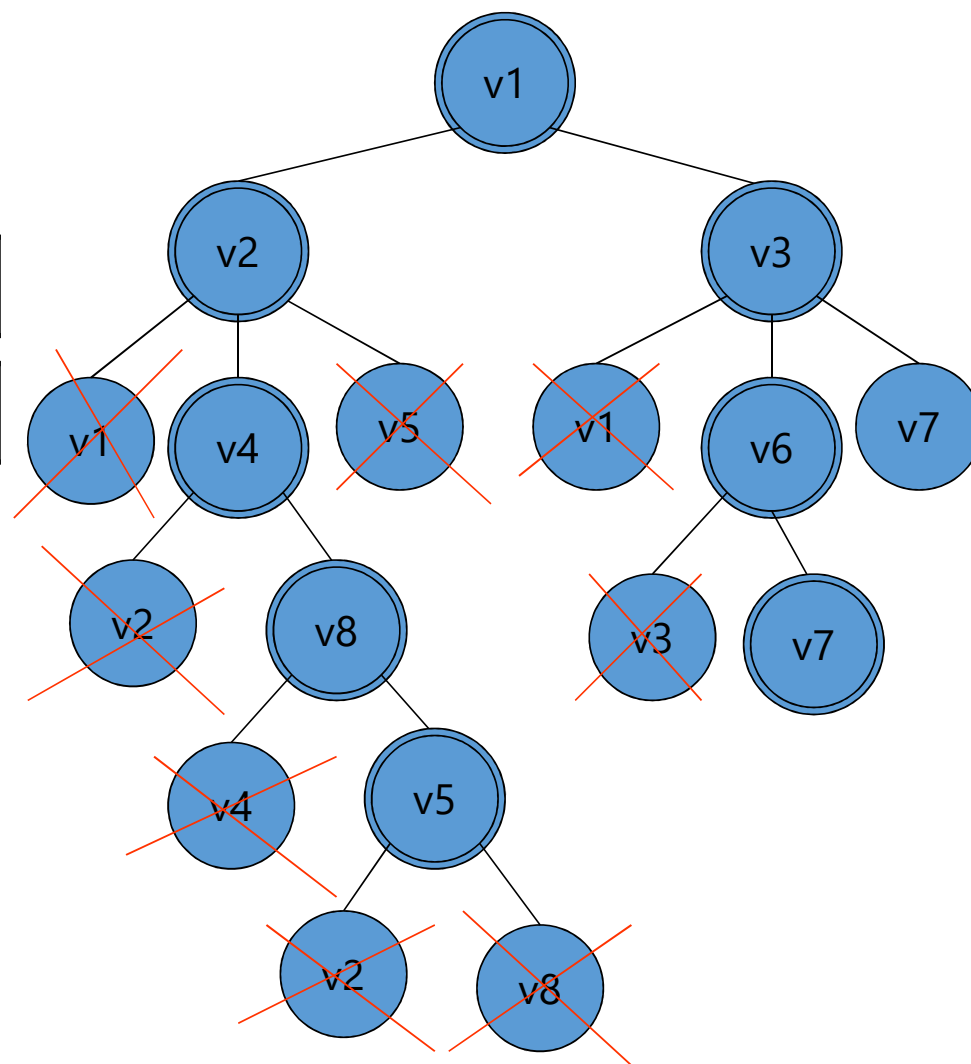
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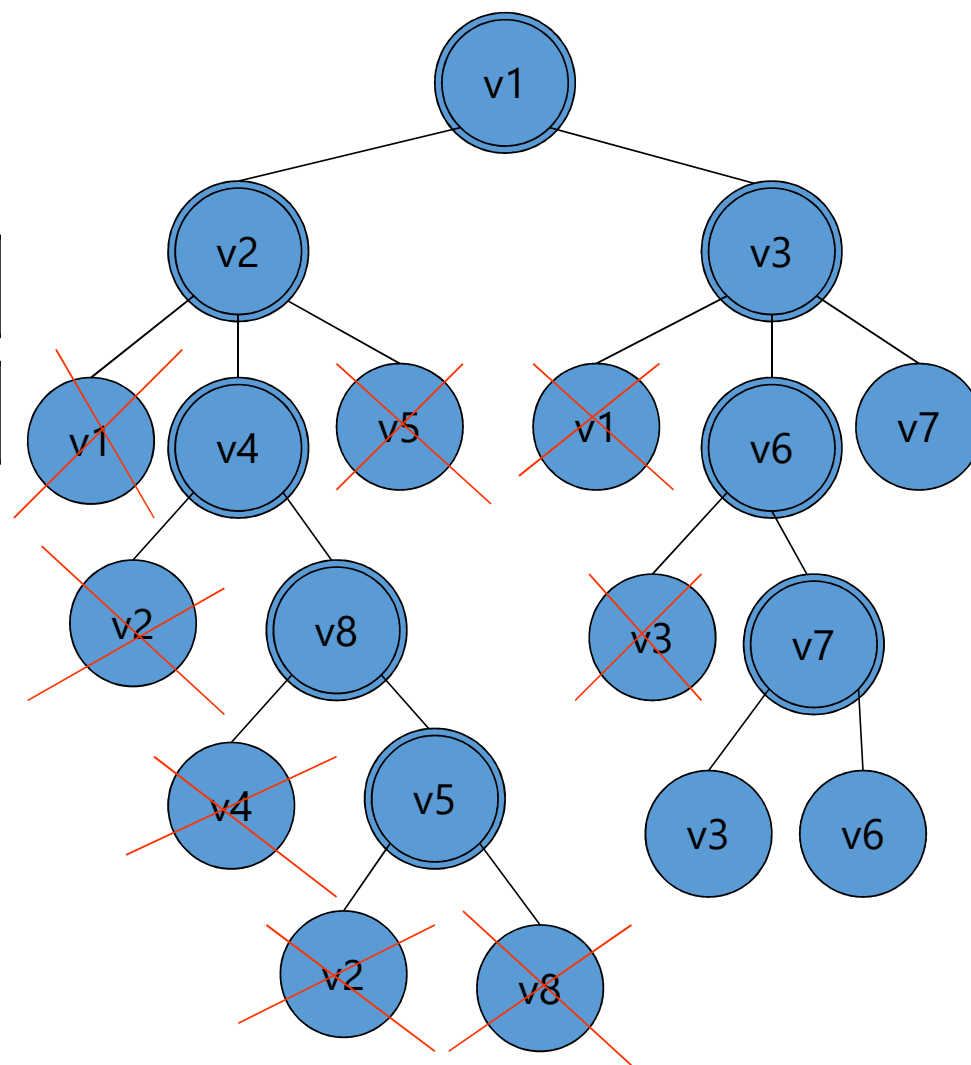
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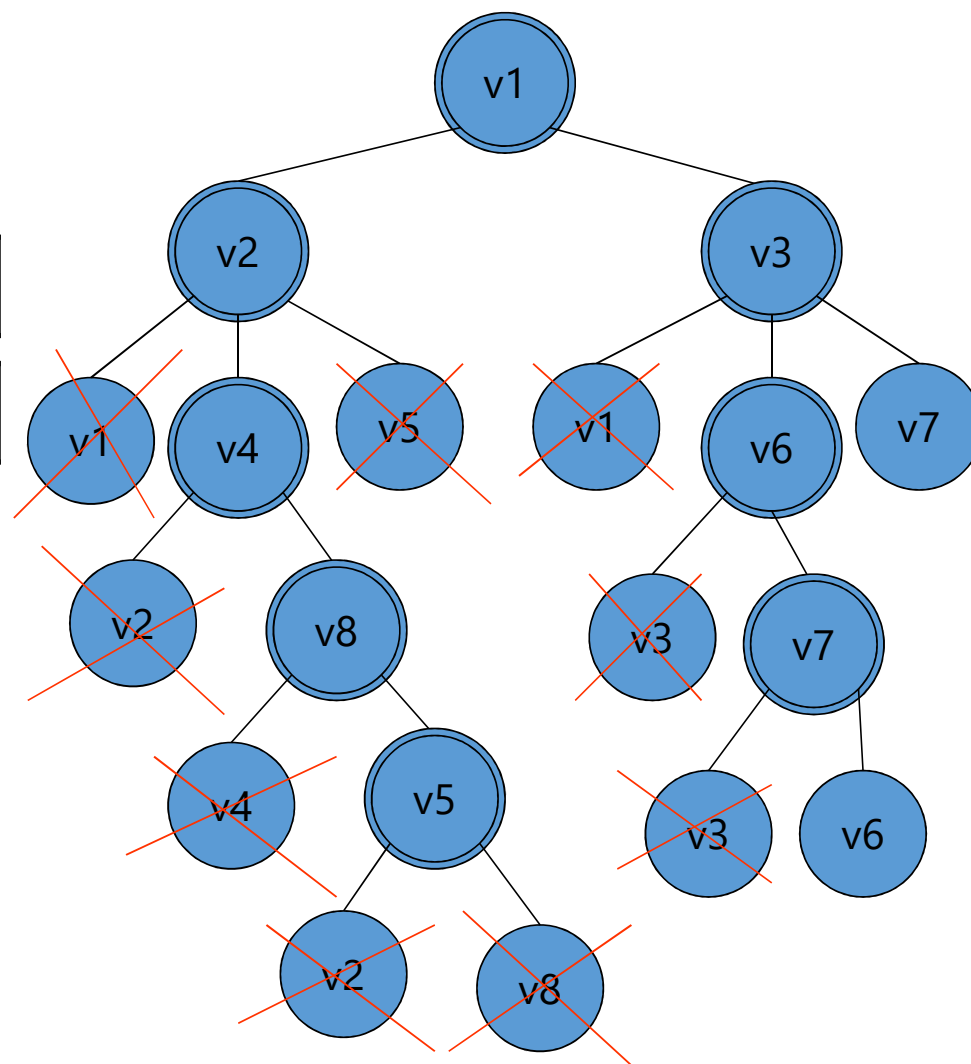
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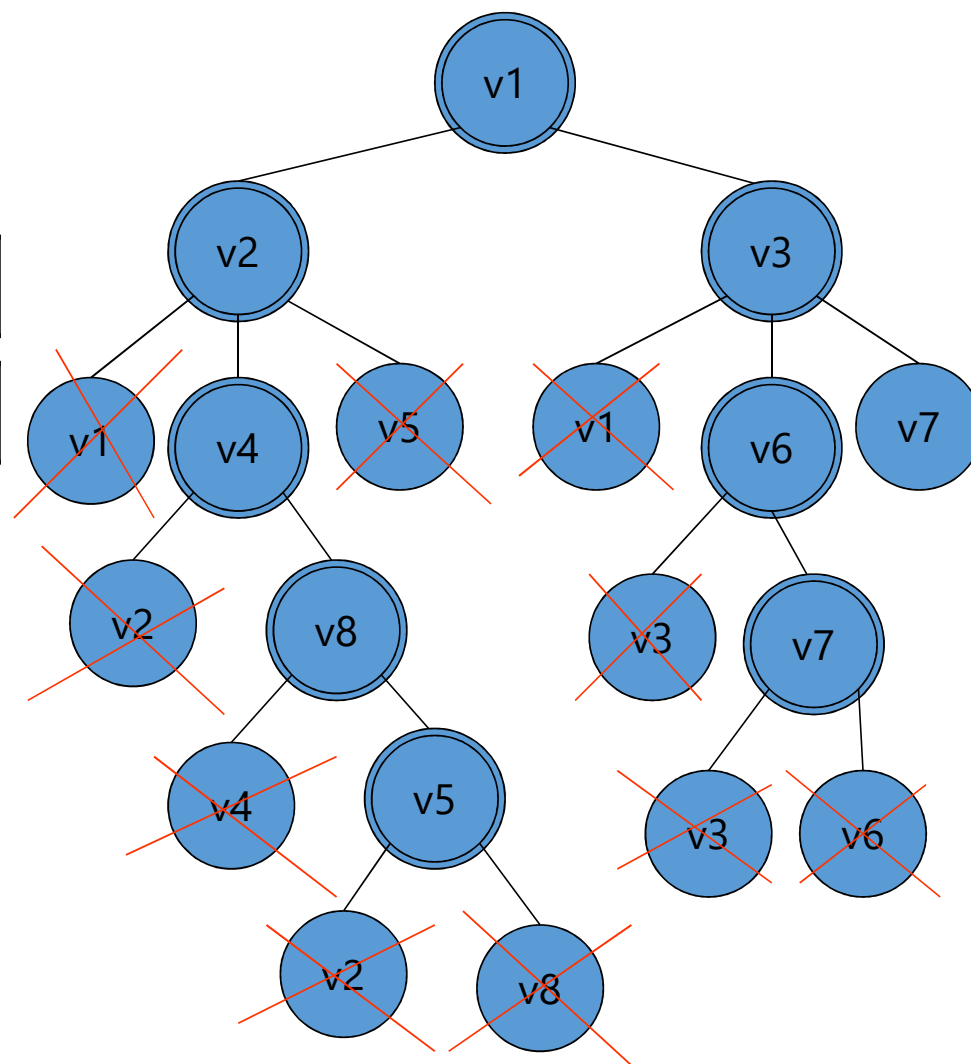
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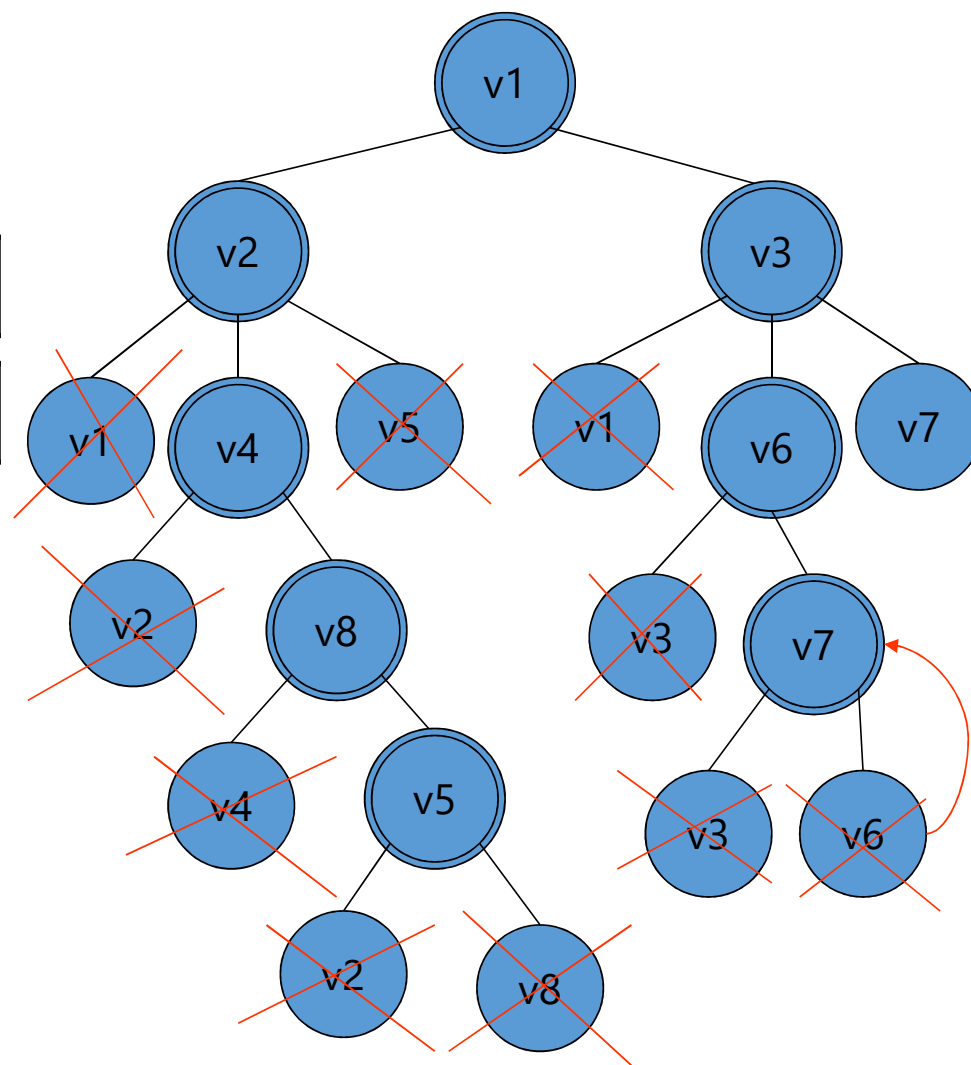
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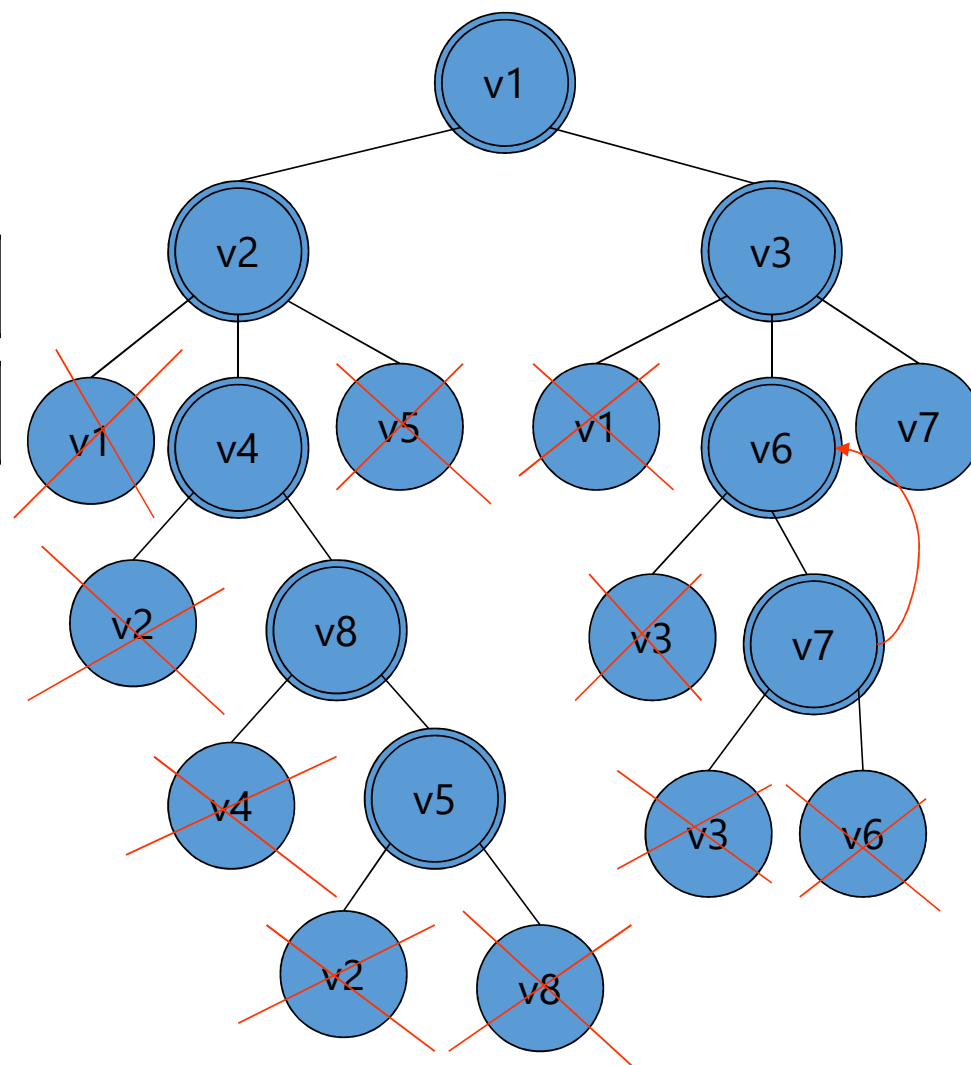
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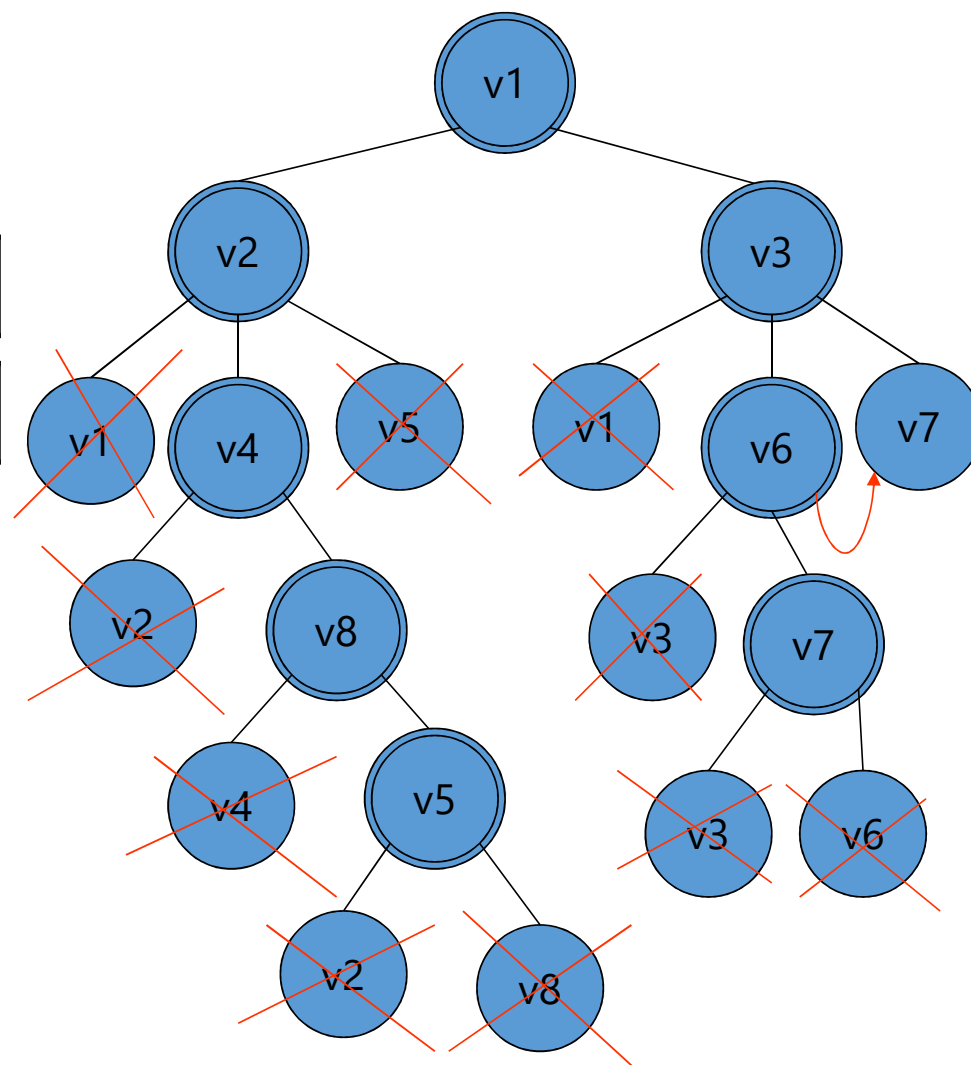
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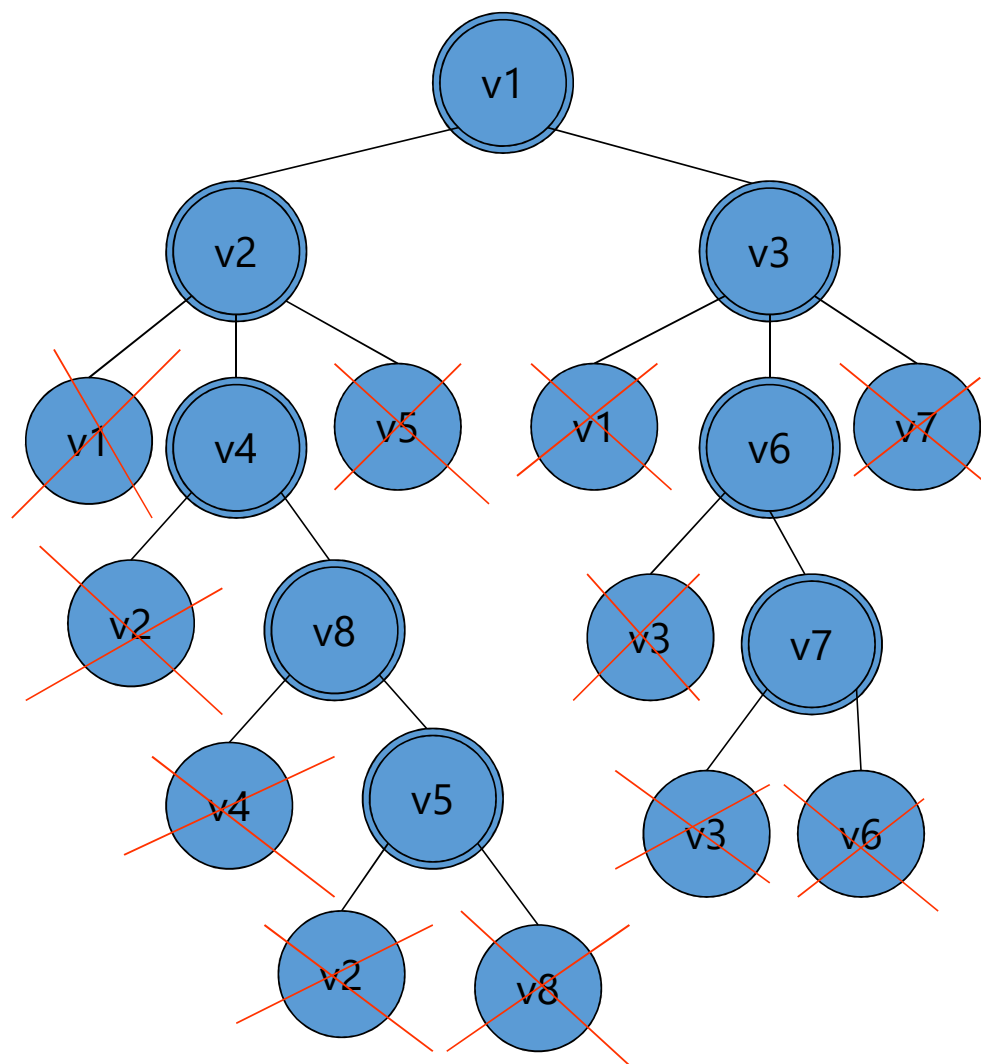
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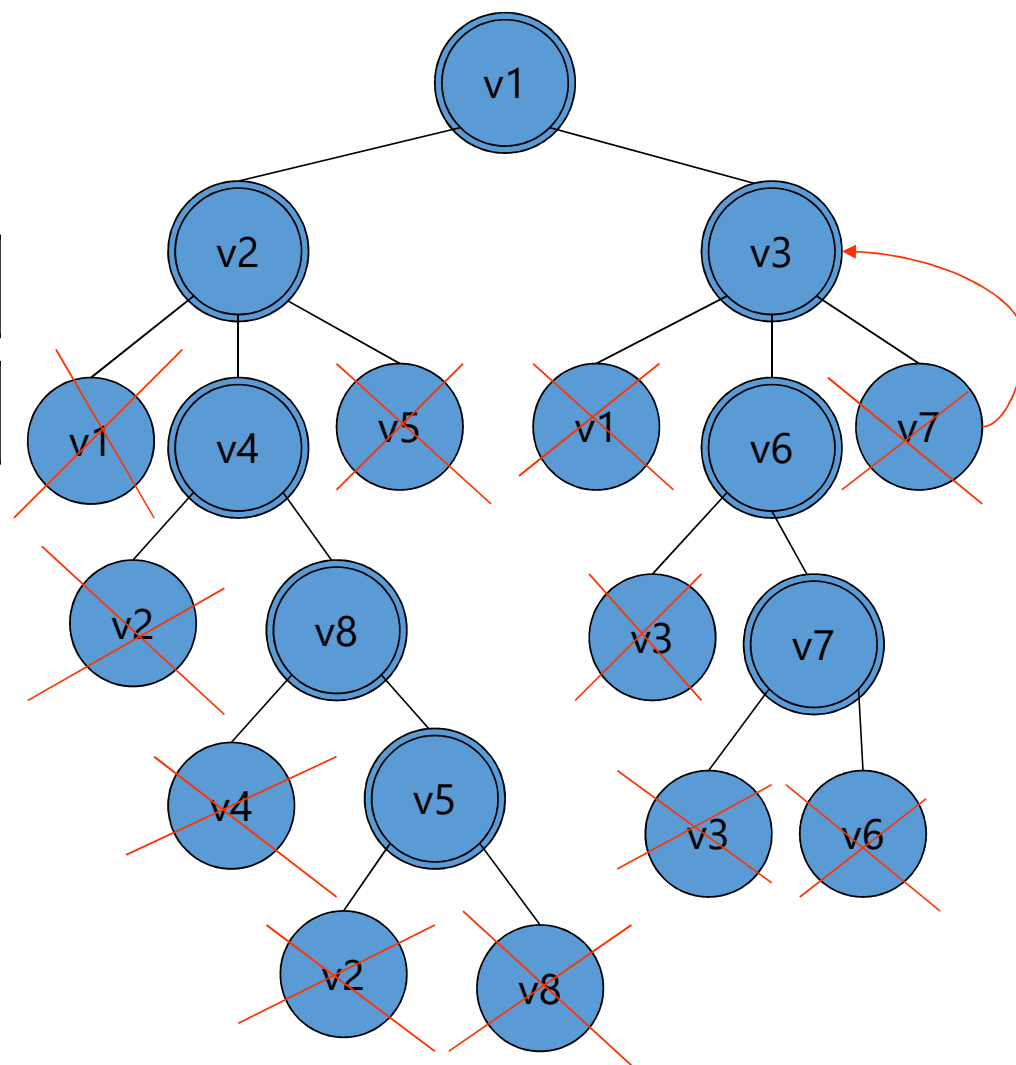
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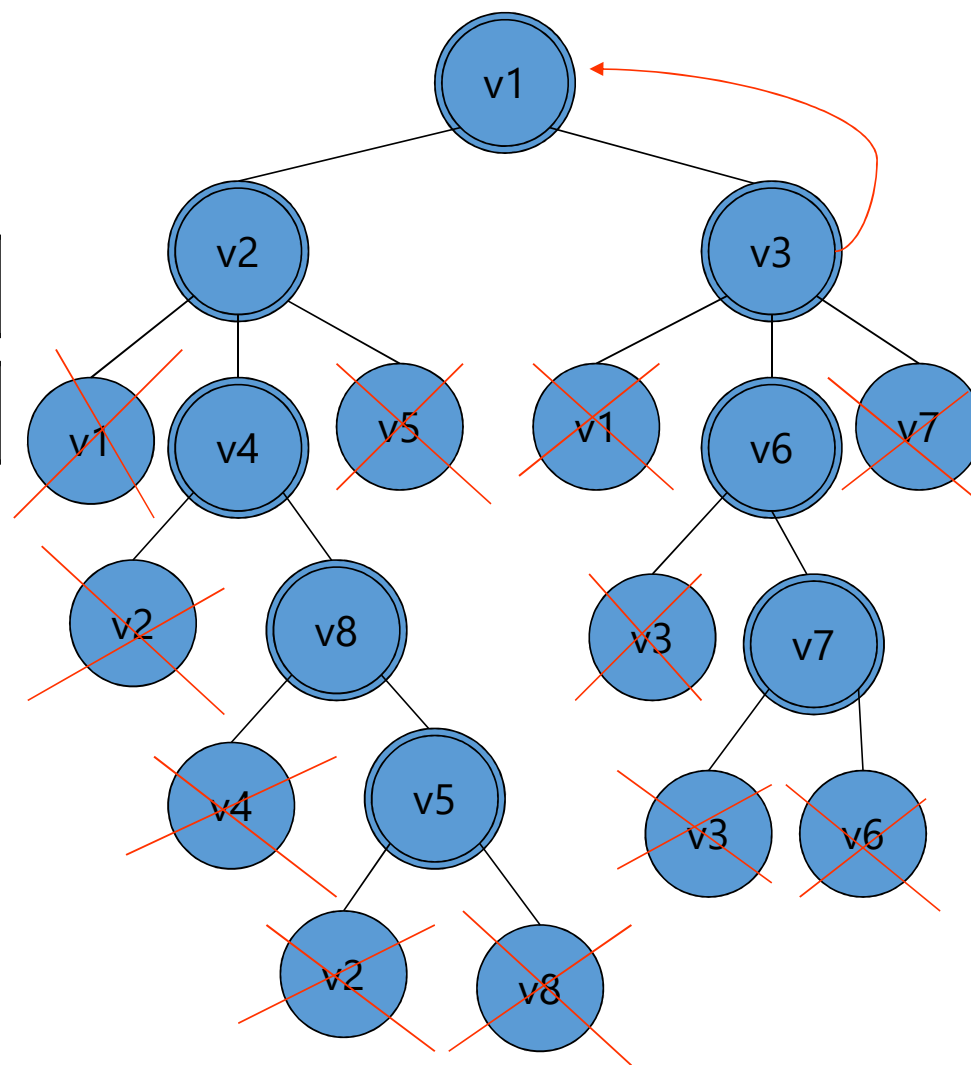
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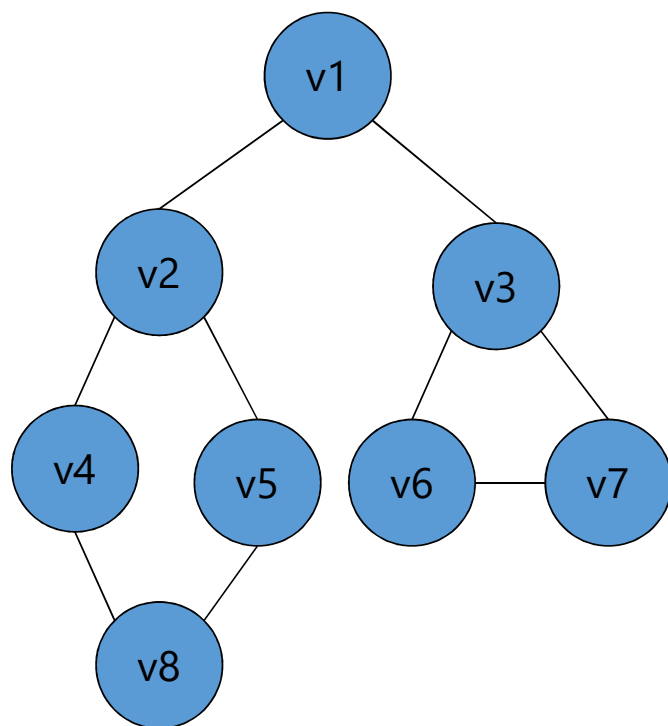
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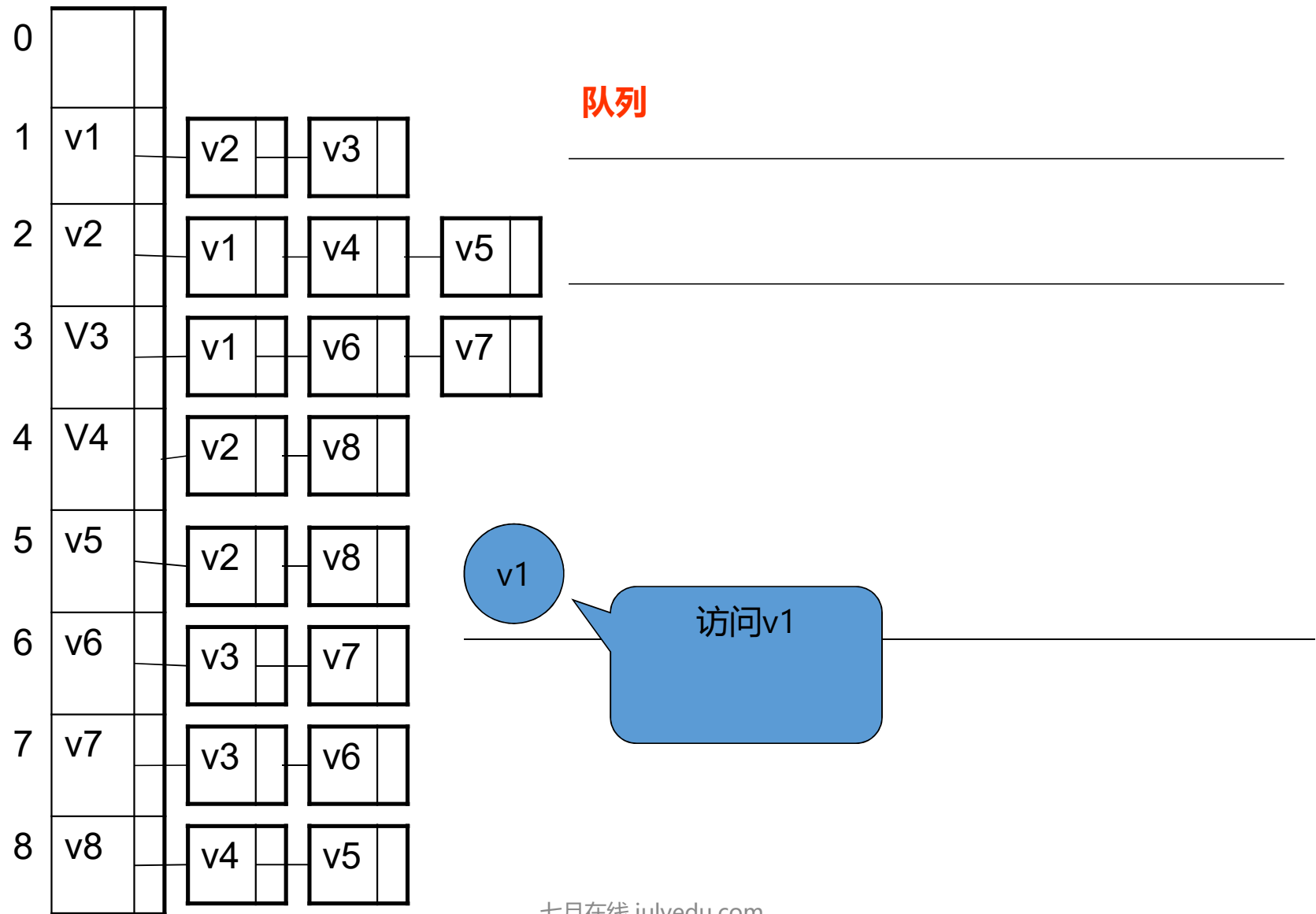
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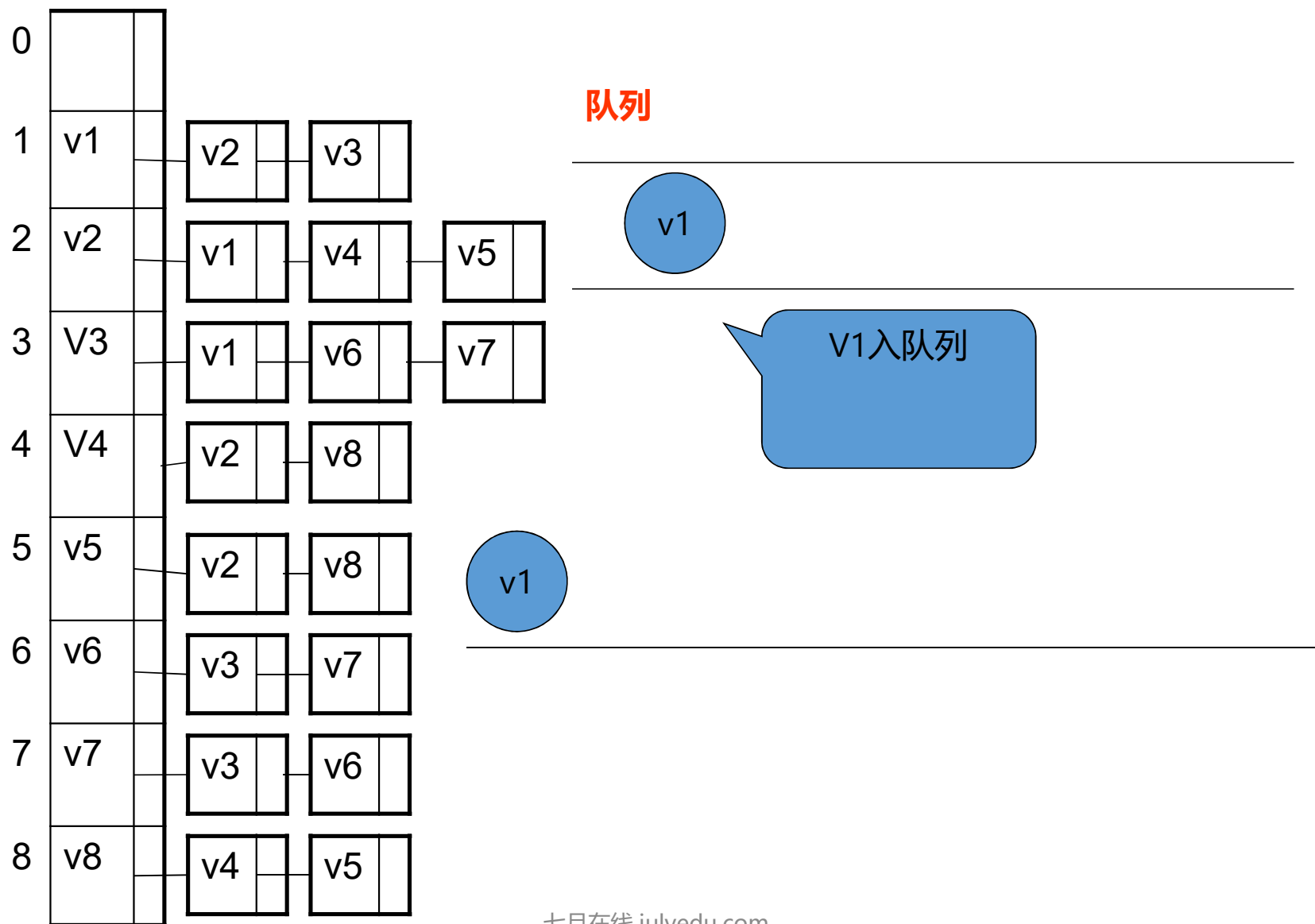


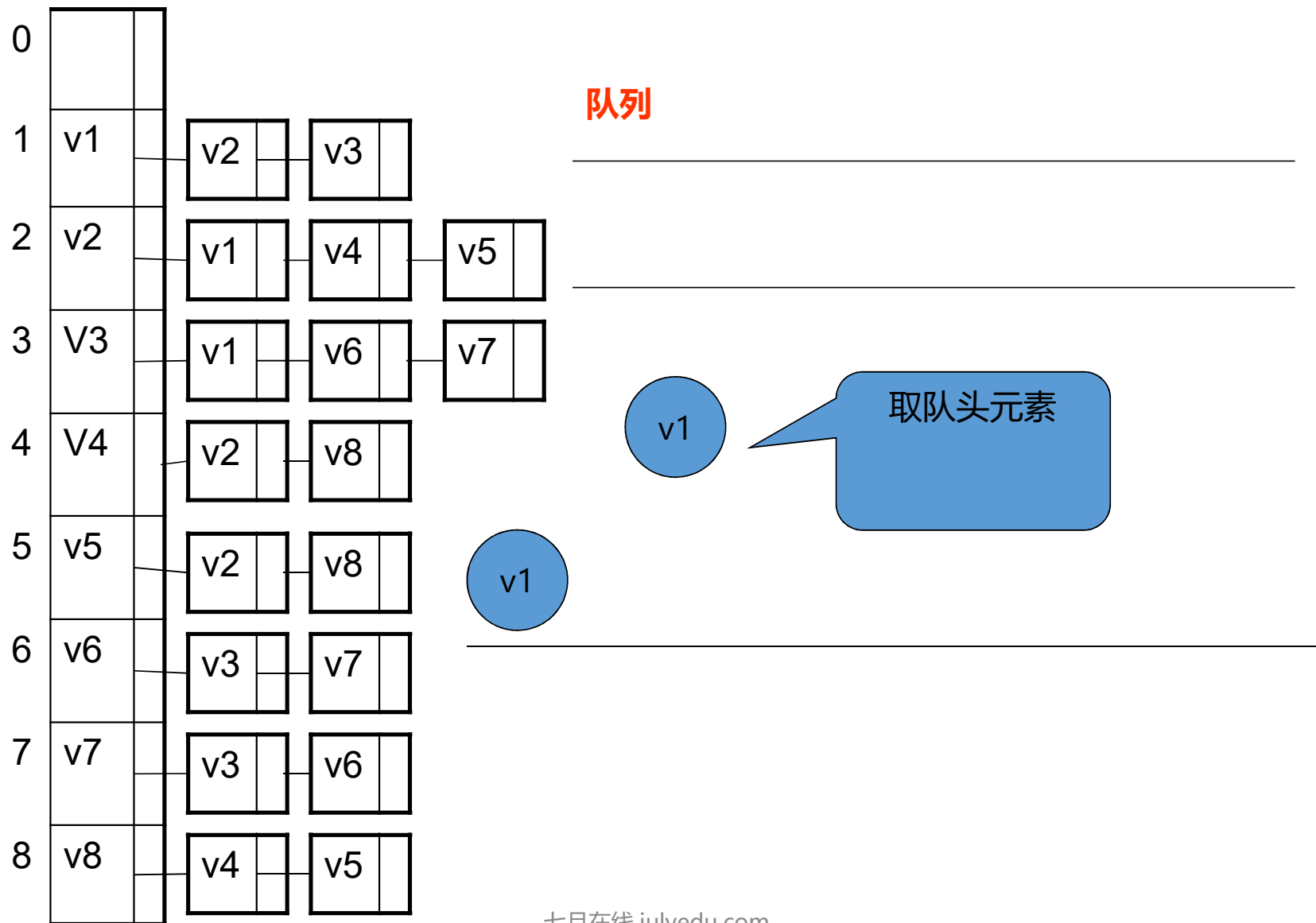
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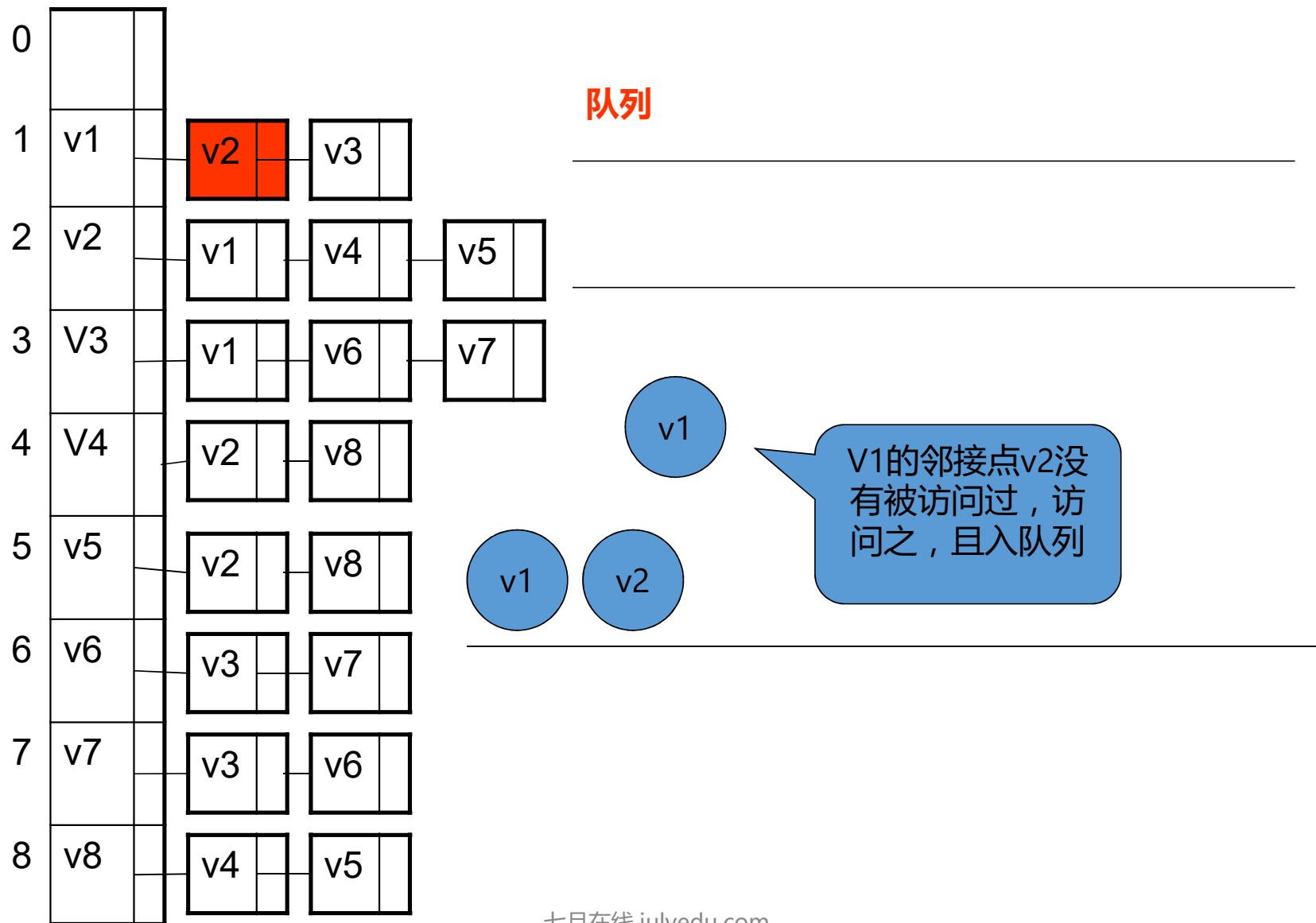


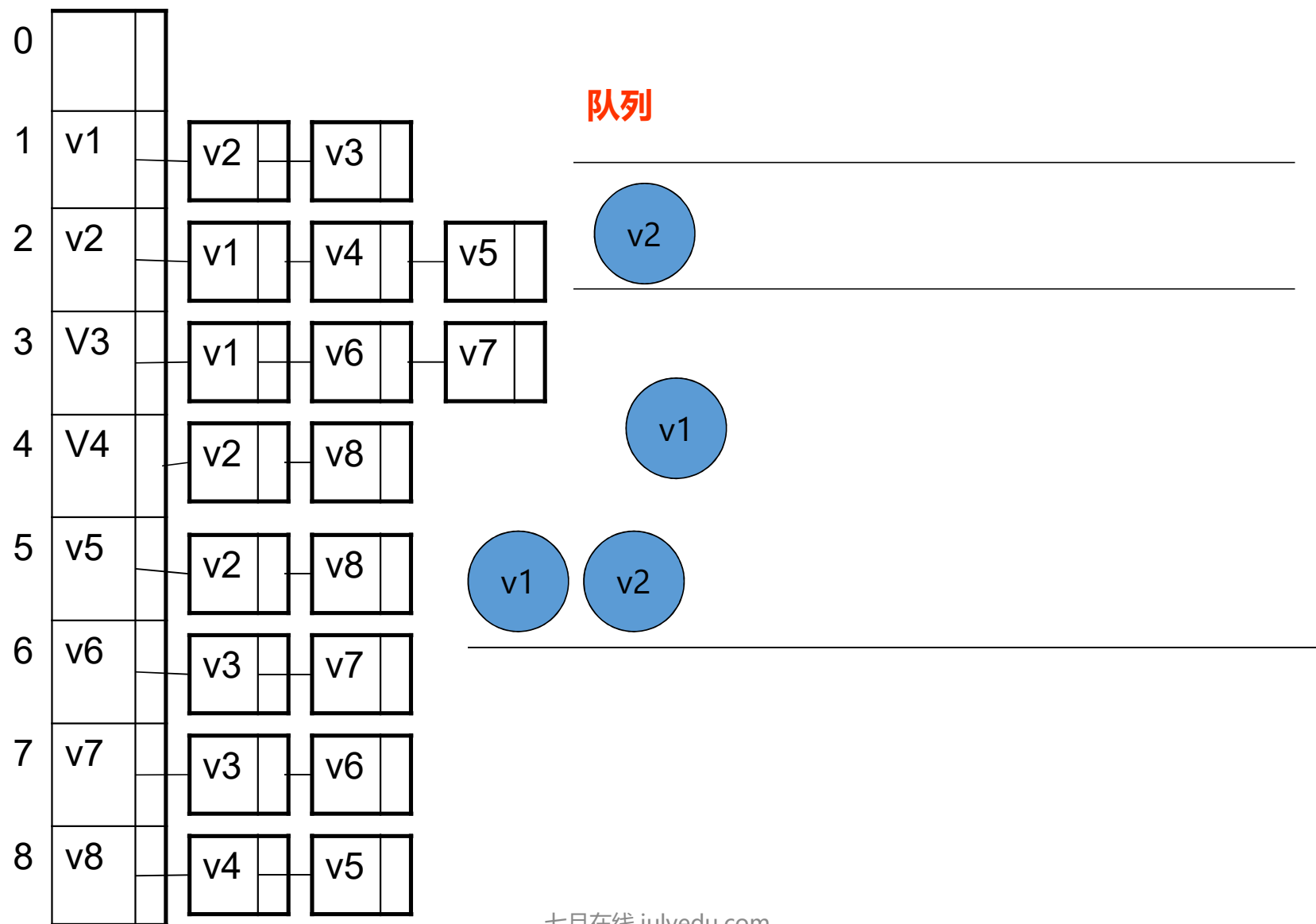
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1	v1		v2	v3	
2	v2		v1	v4	v5
3	v3		v1	v6	v7
4	v4		v2	v8	
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7	v7		v3	v6	
8	v8		v4	v5	

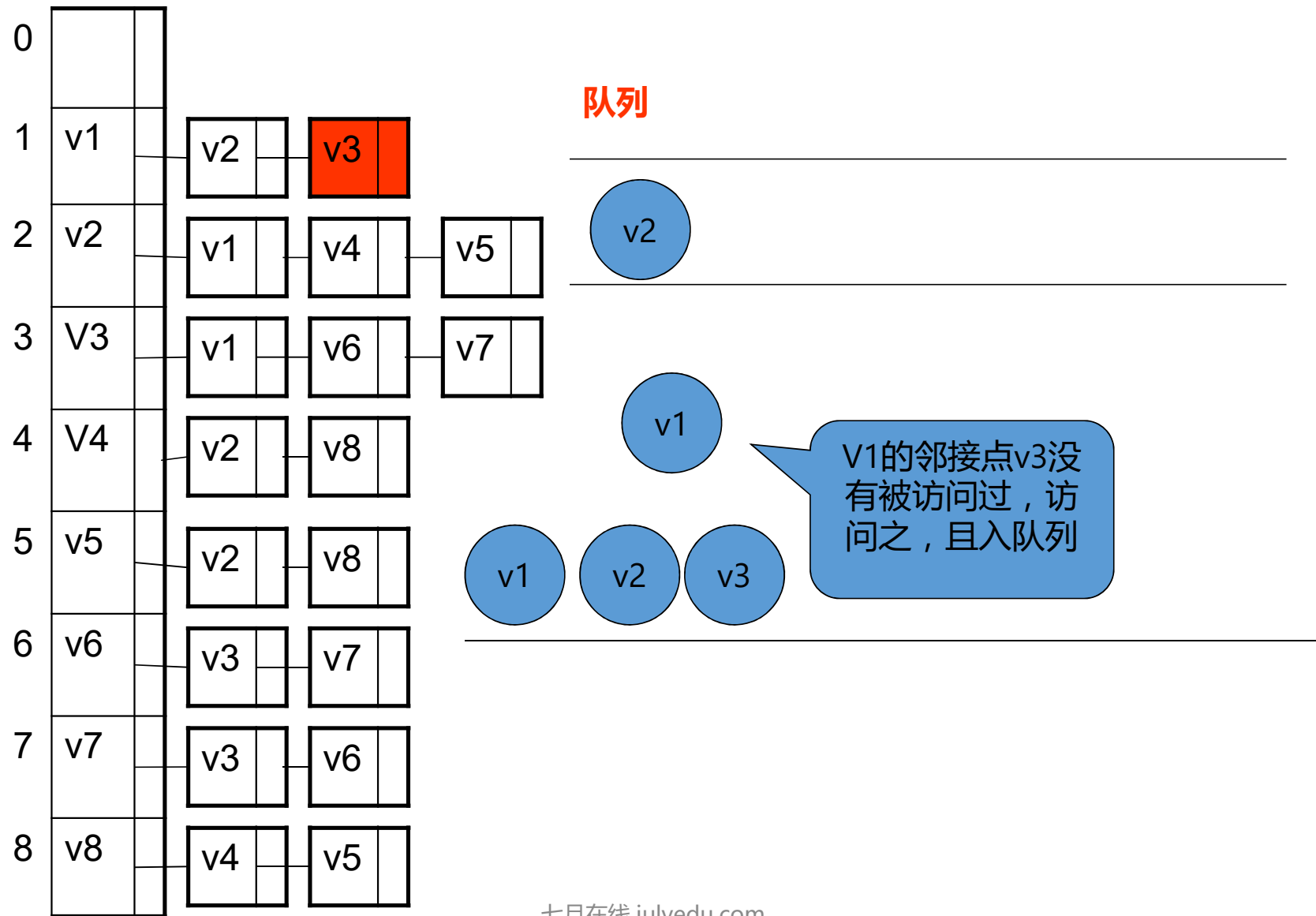


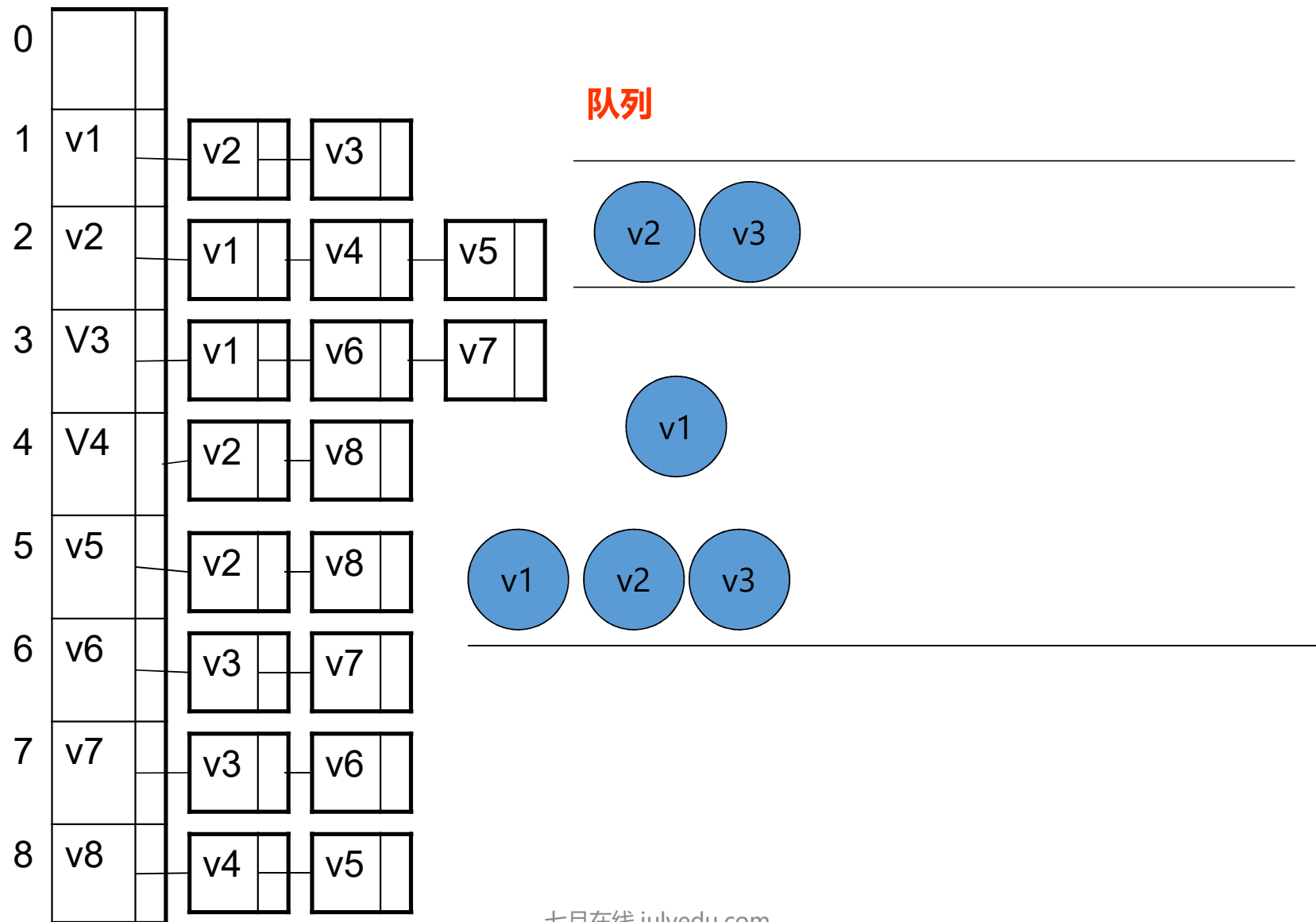


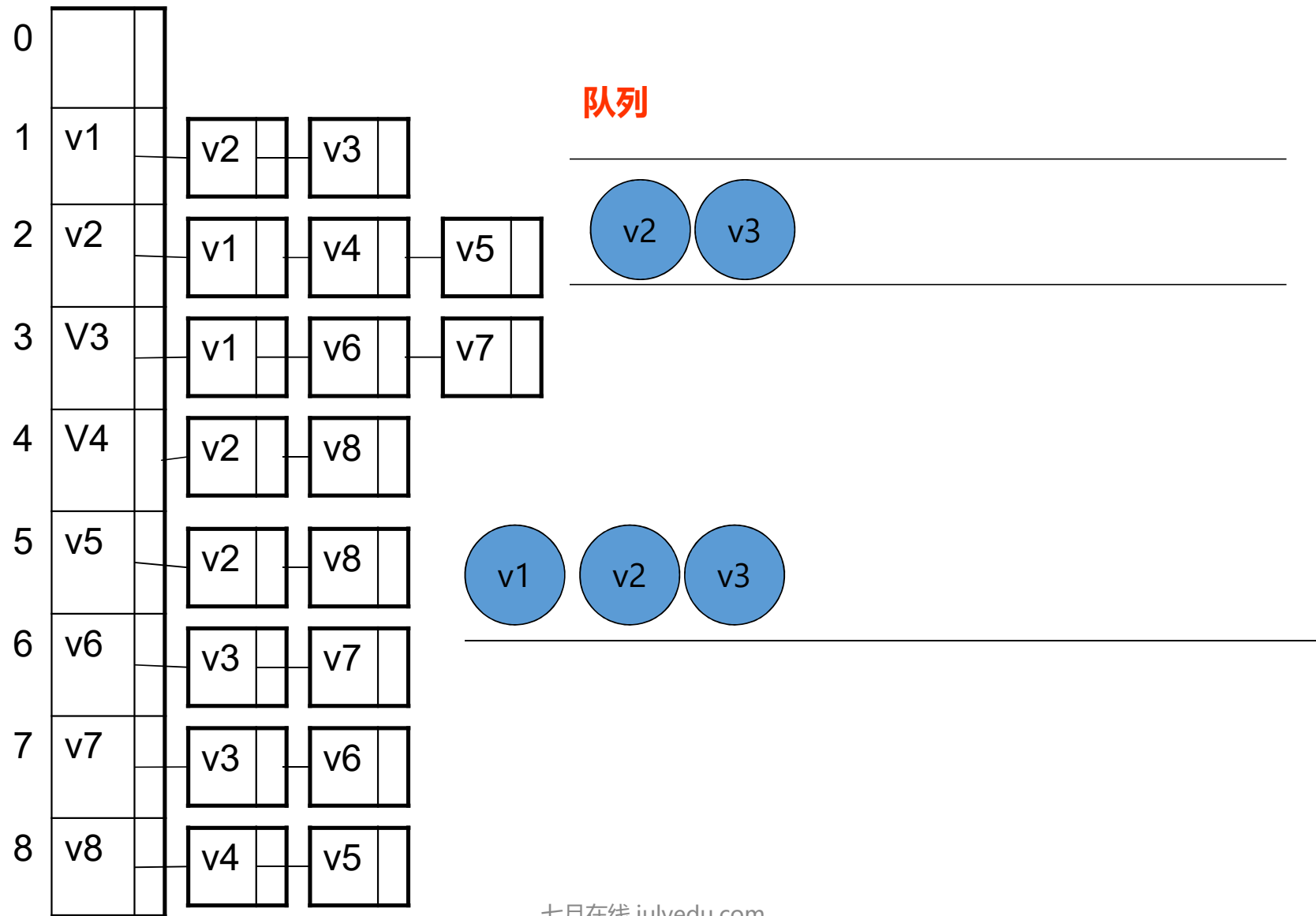


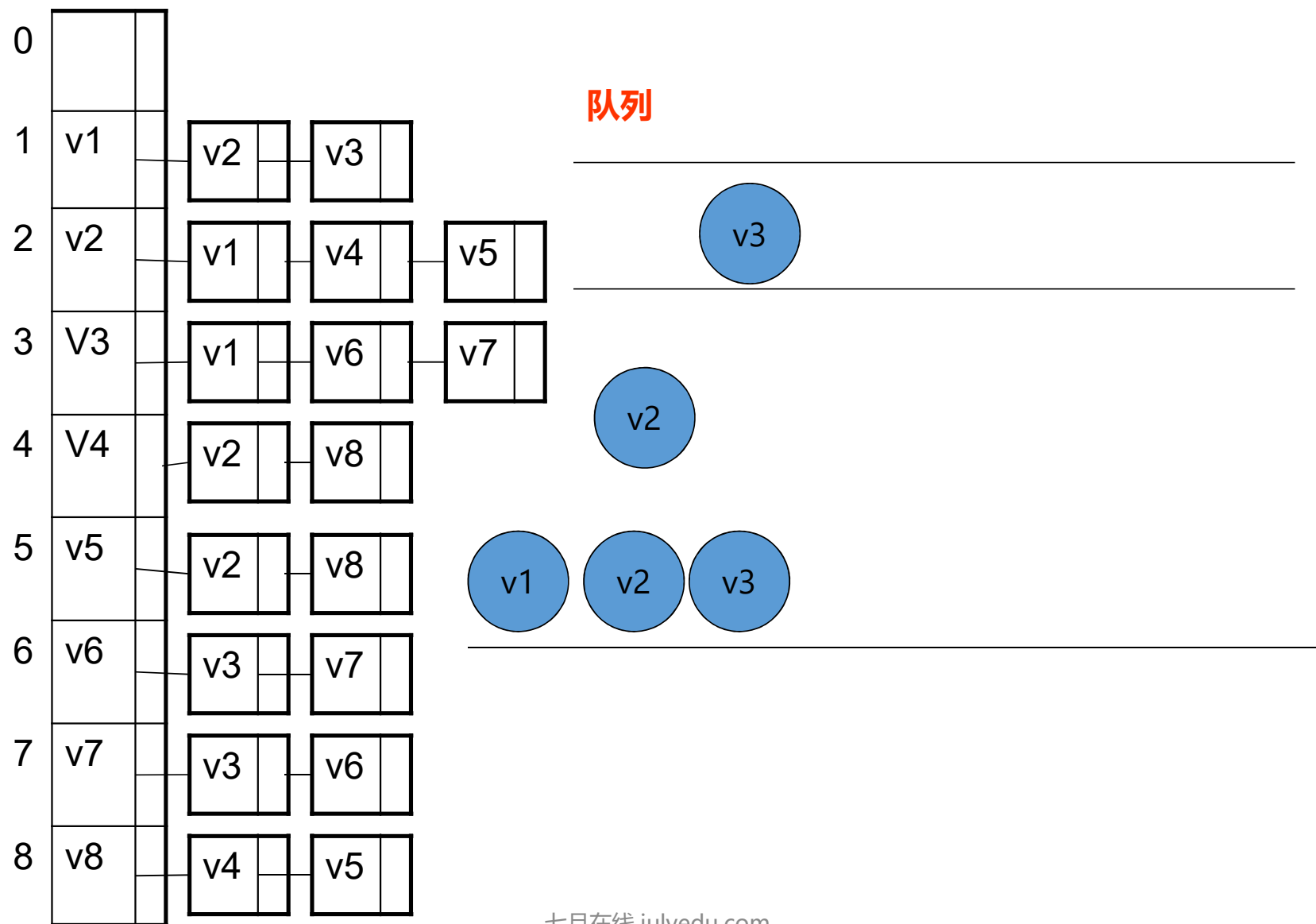


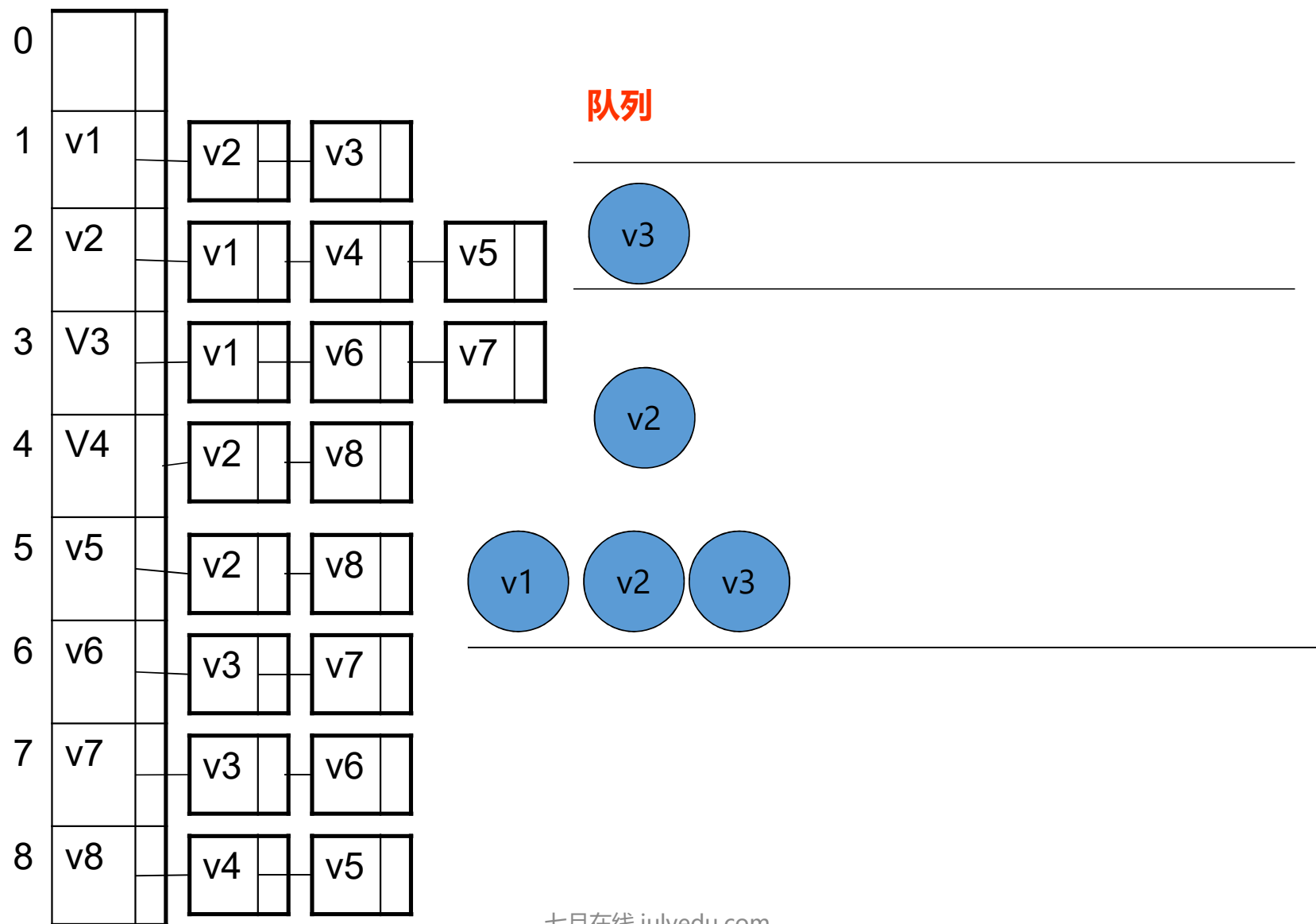


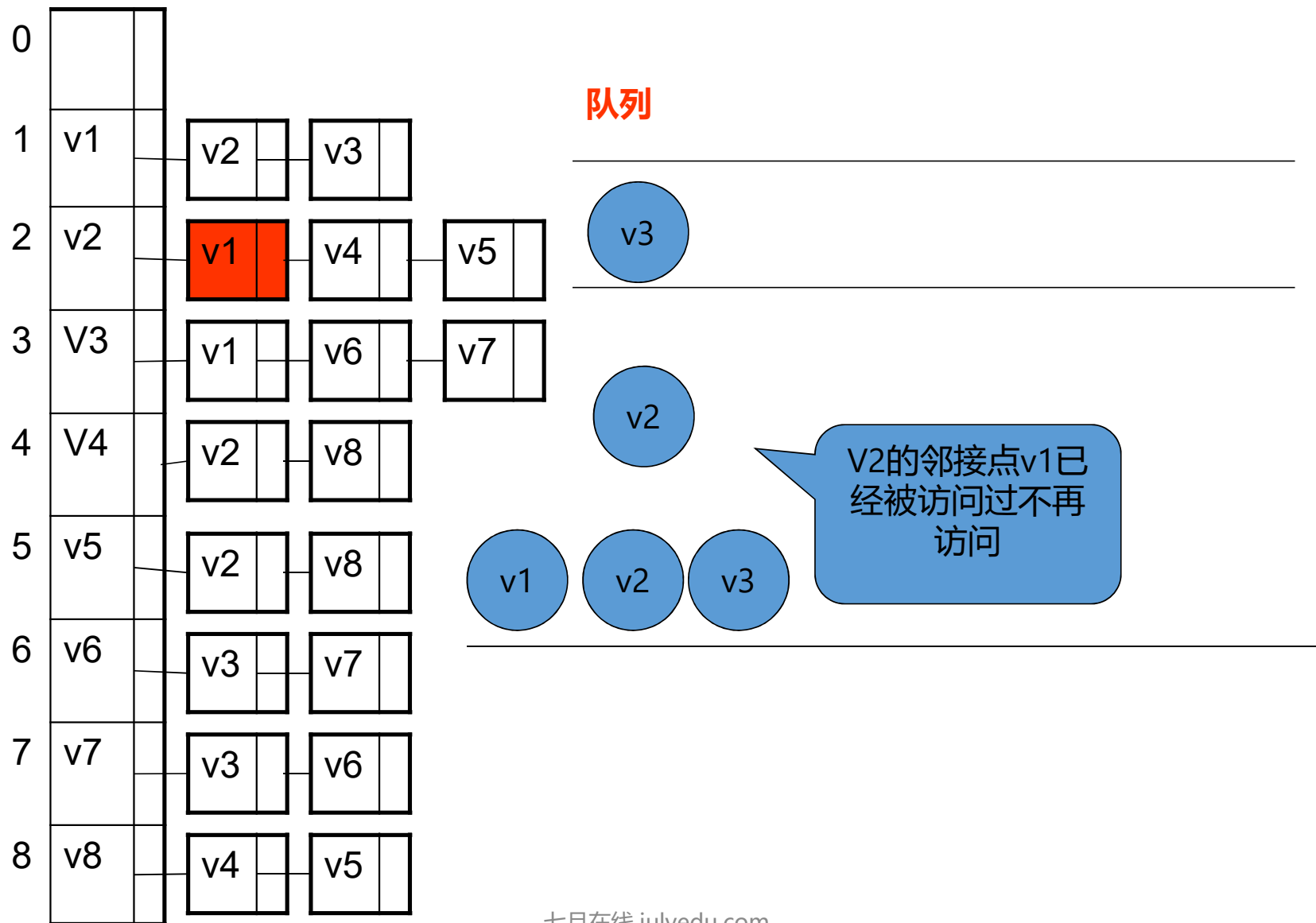


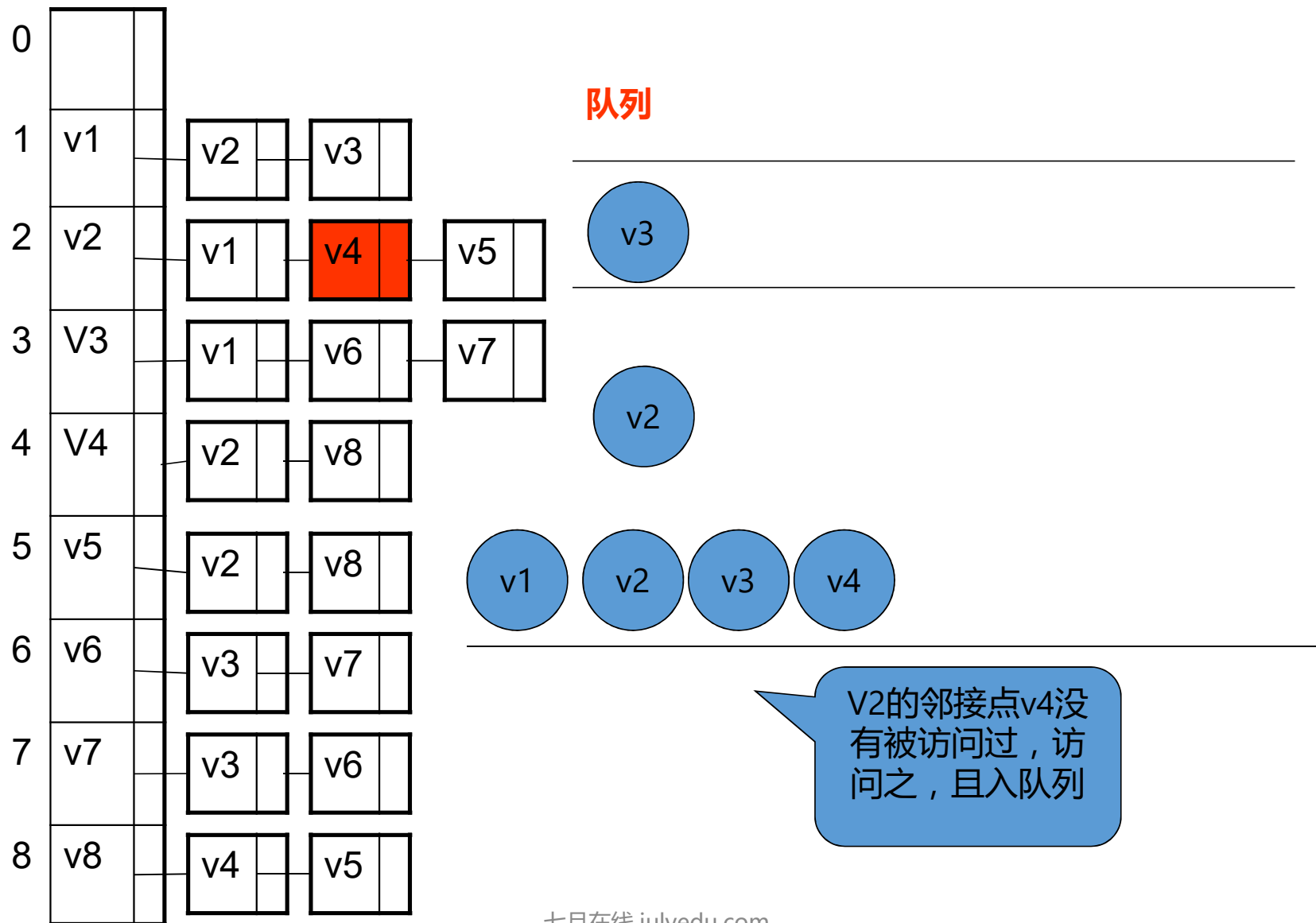


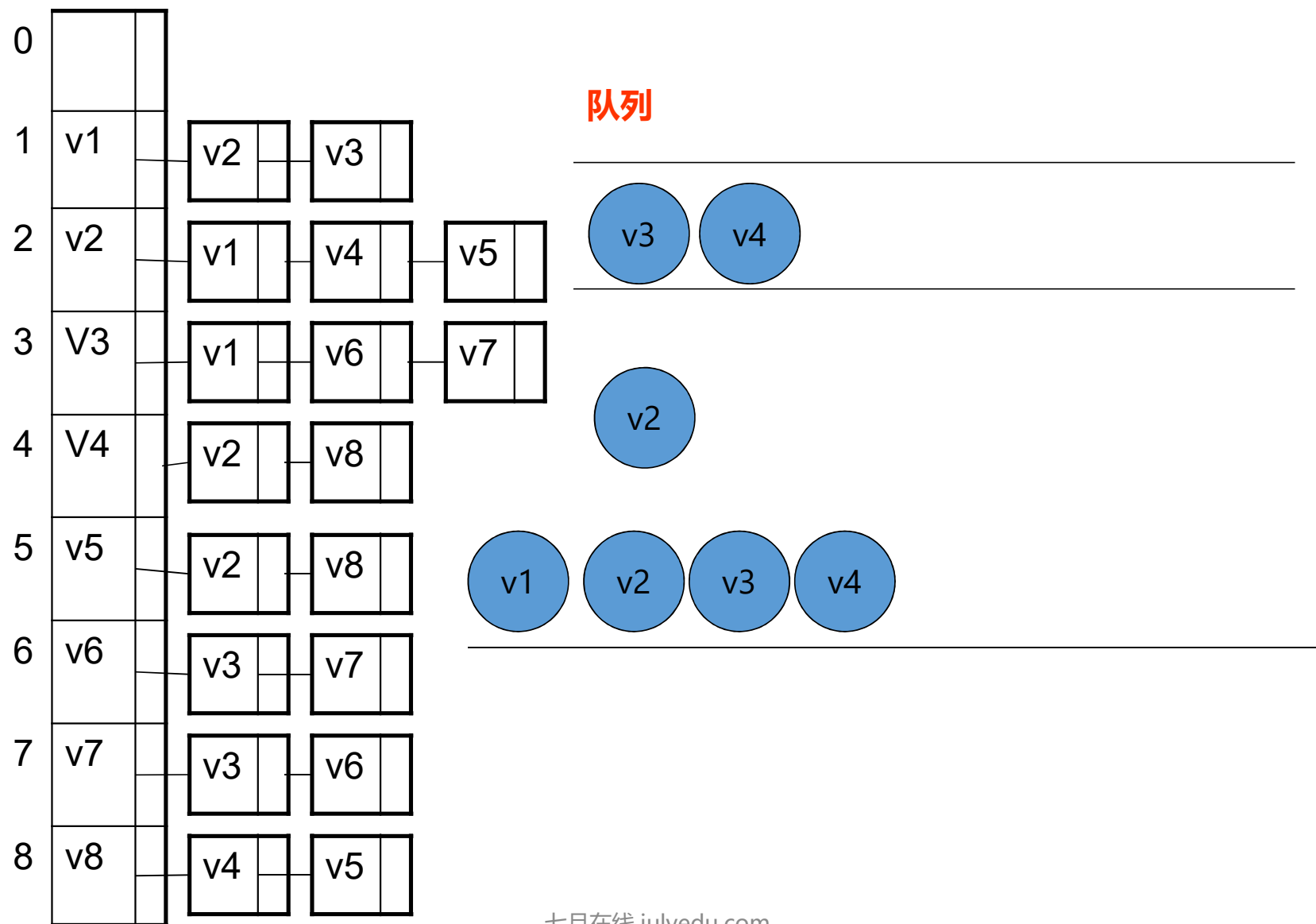


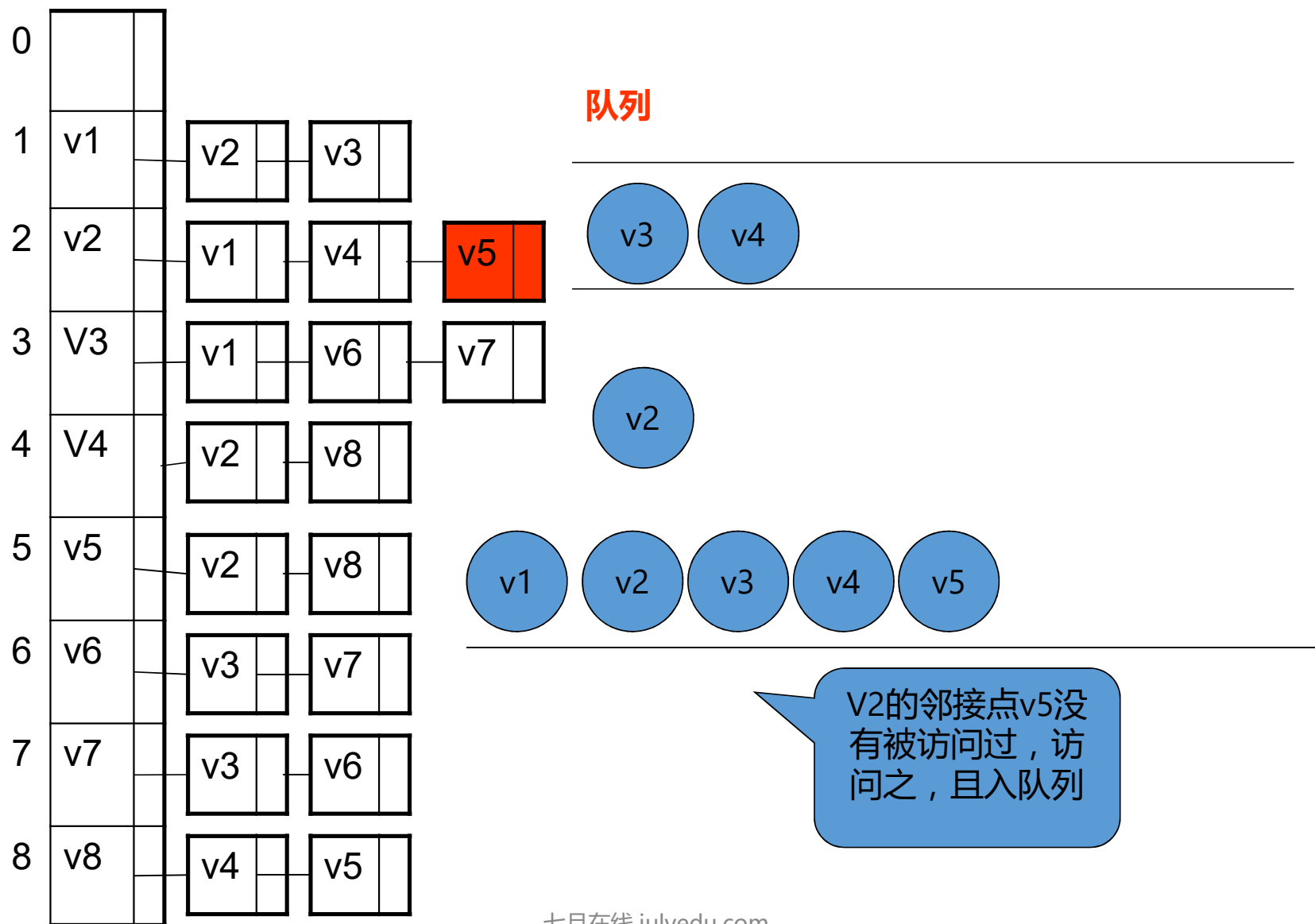


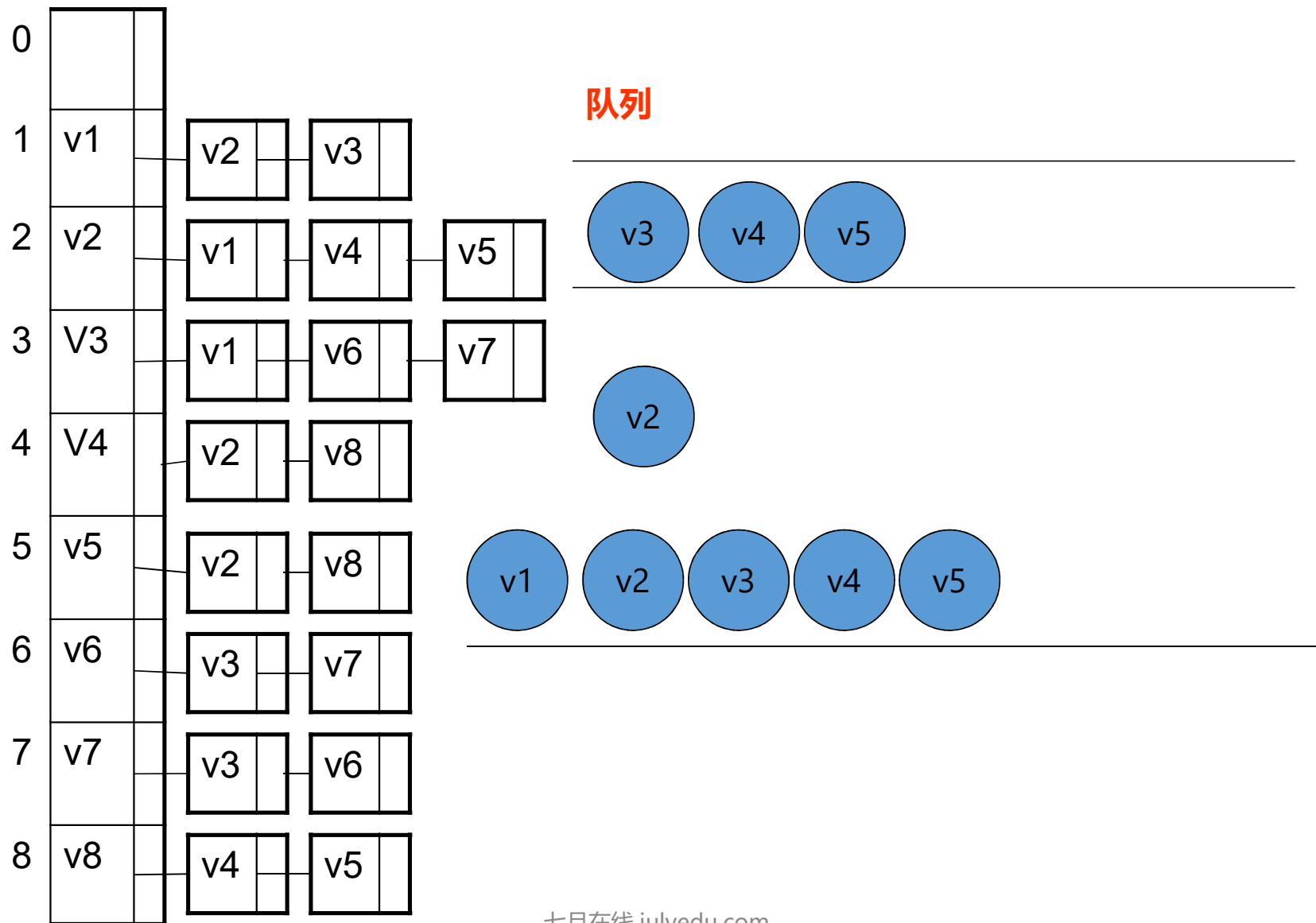


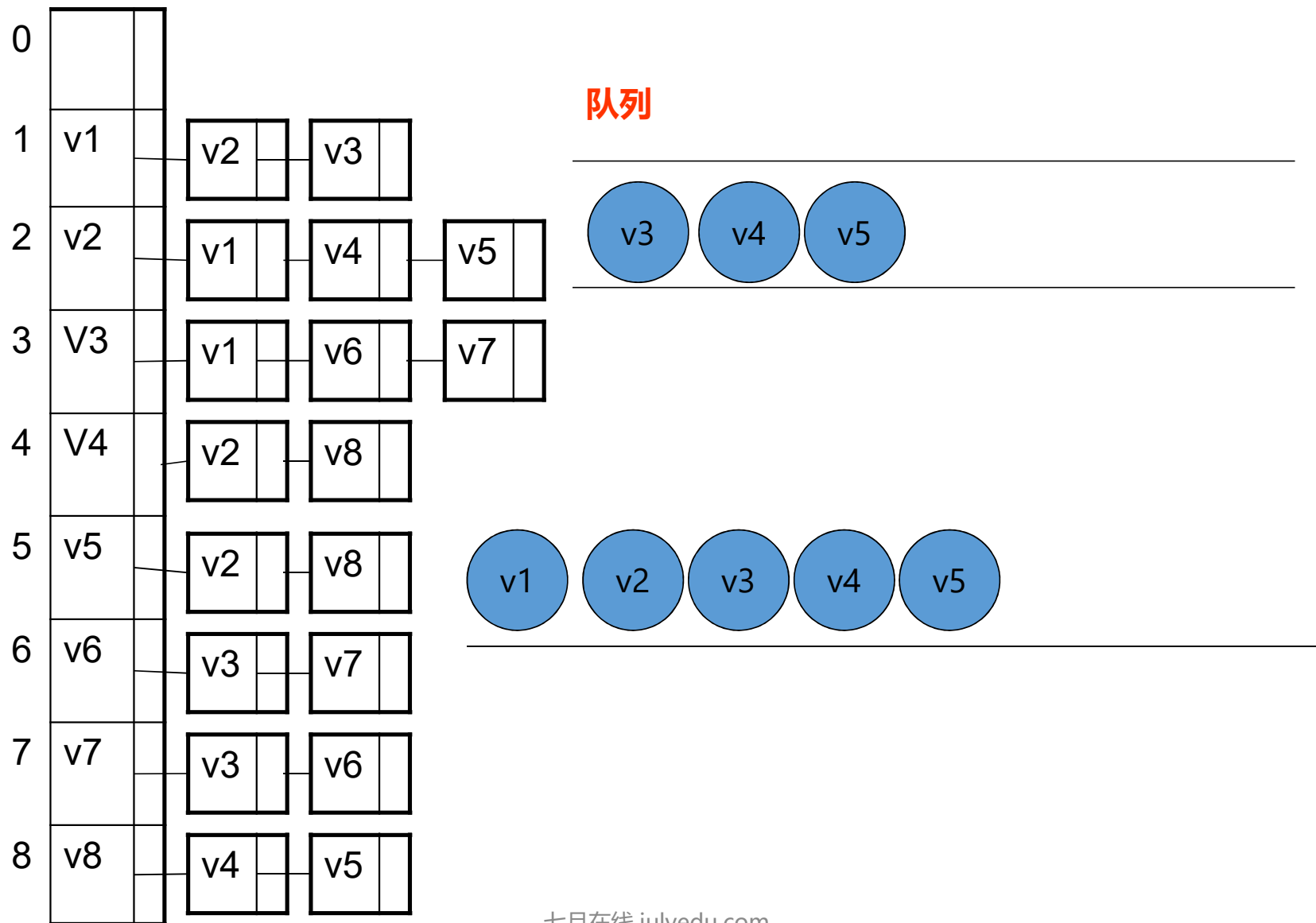


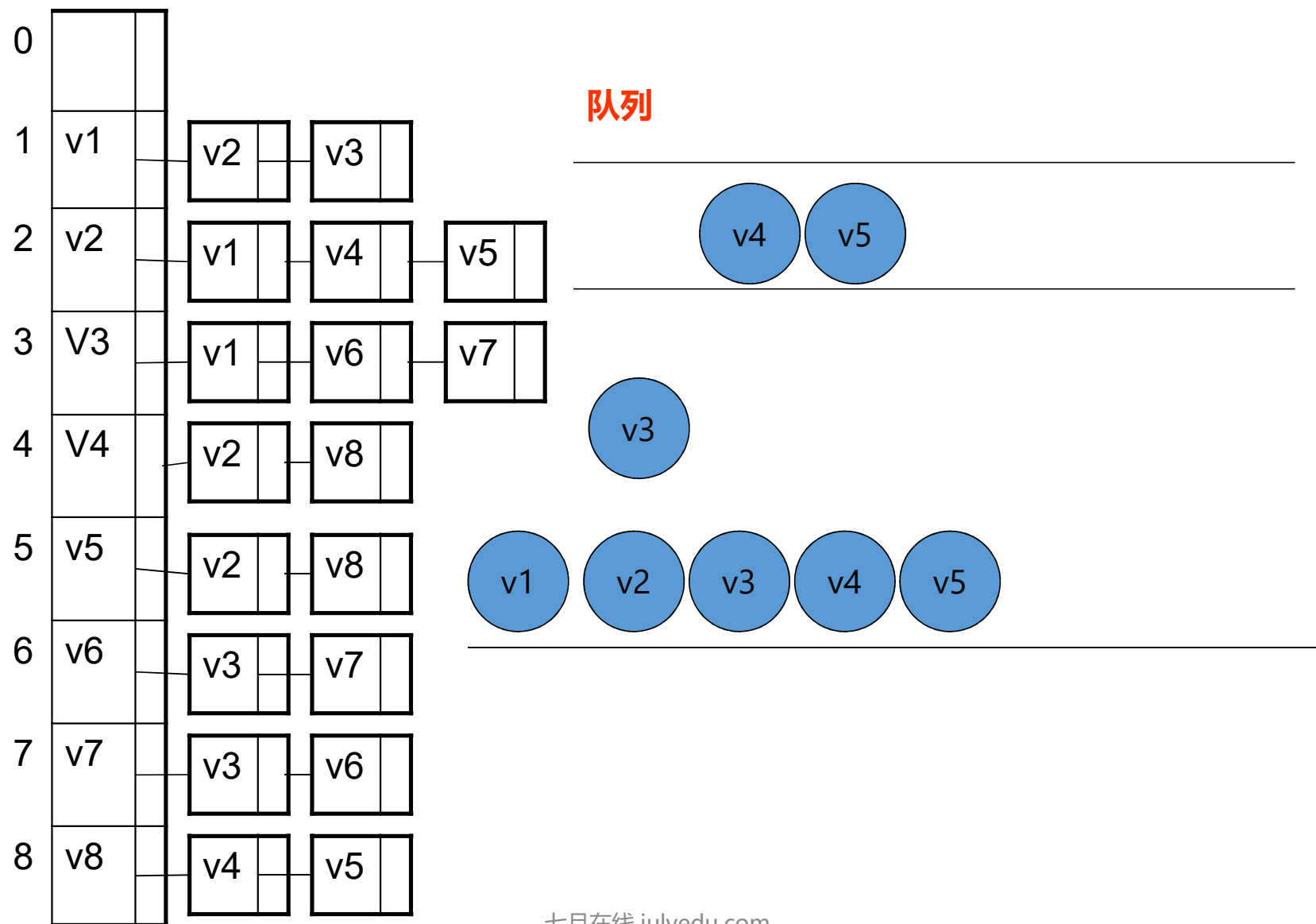


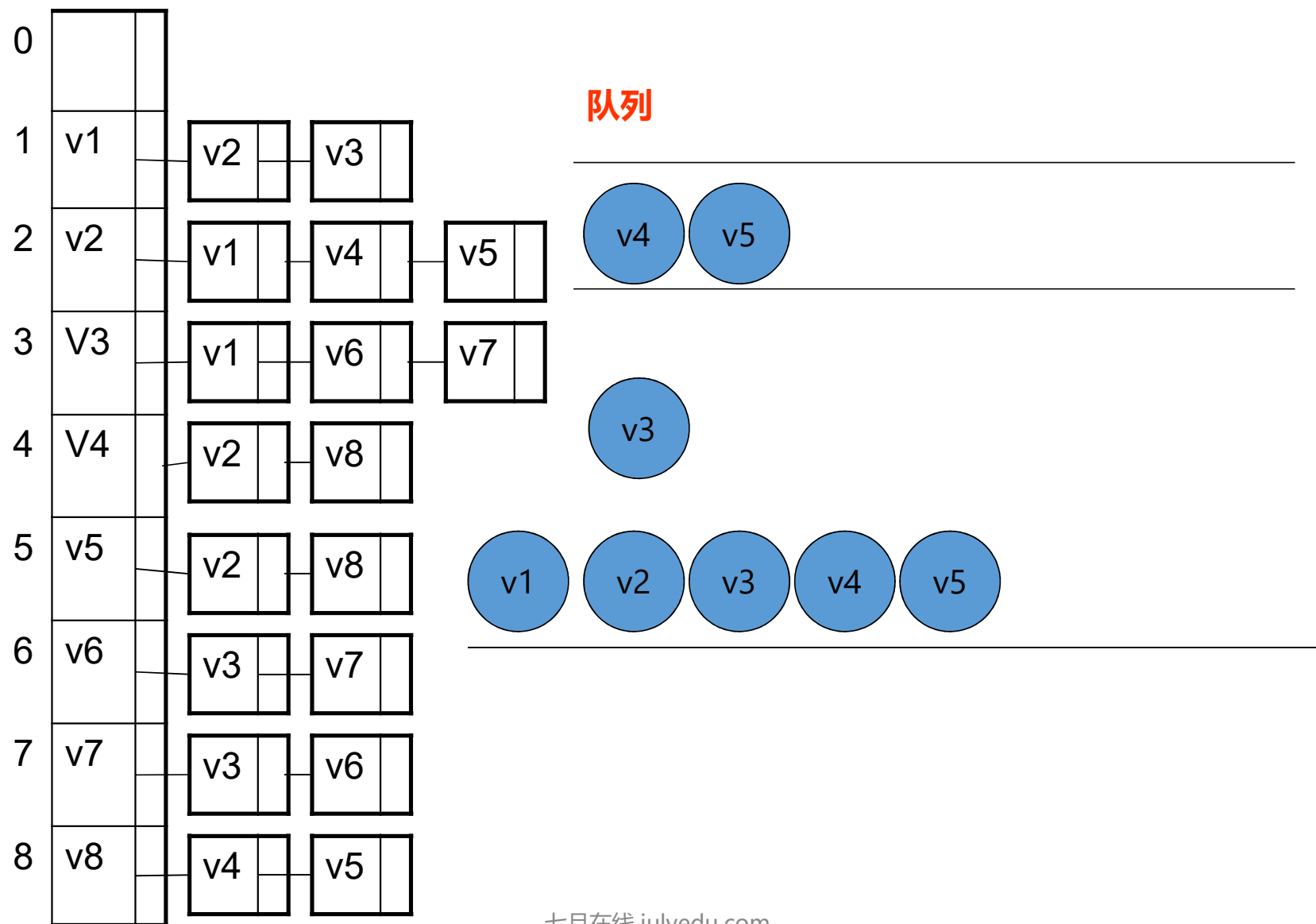


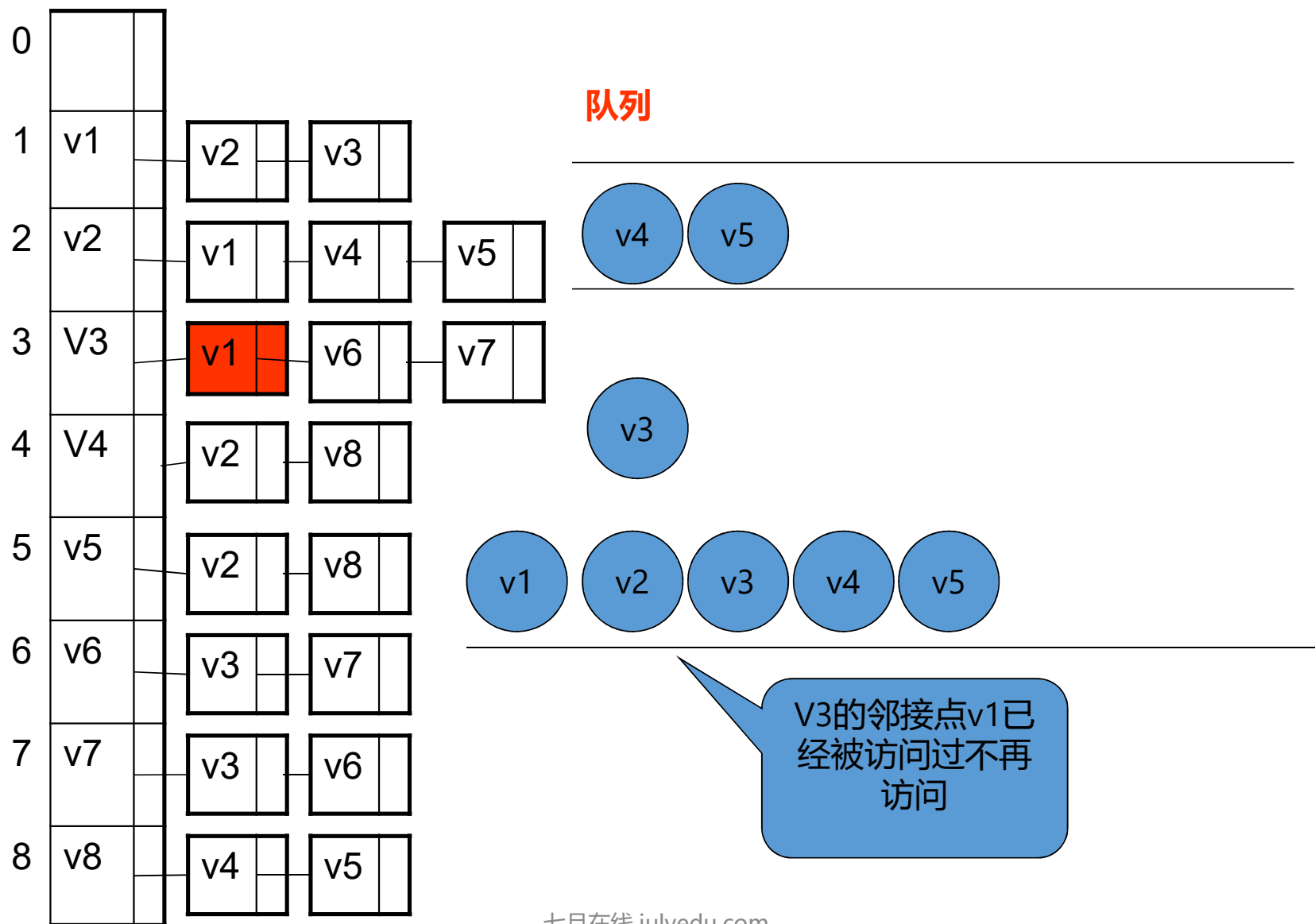


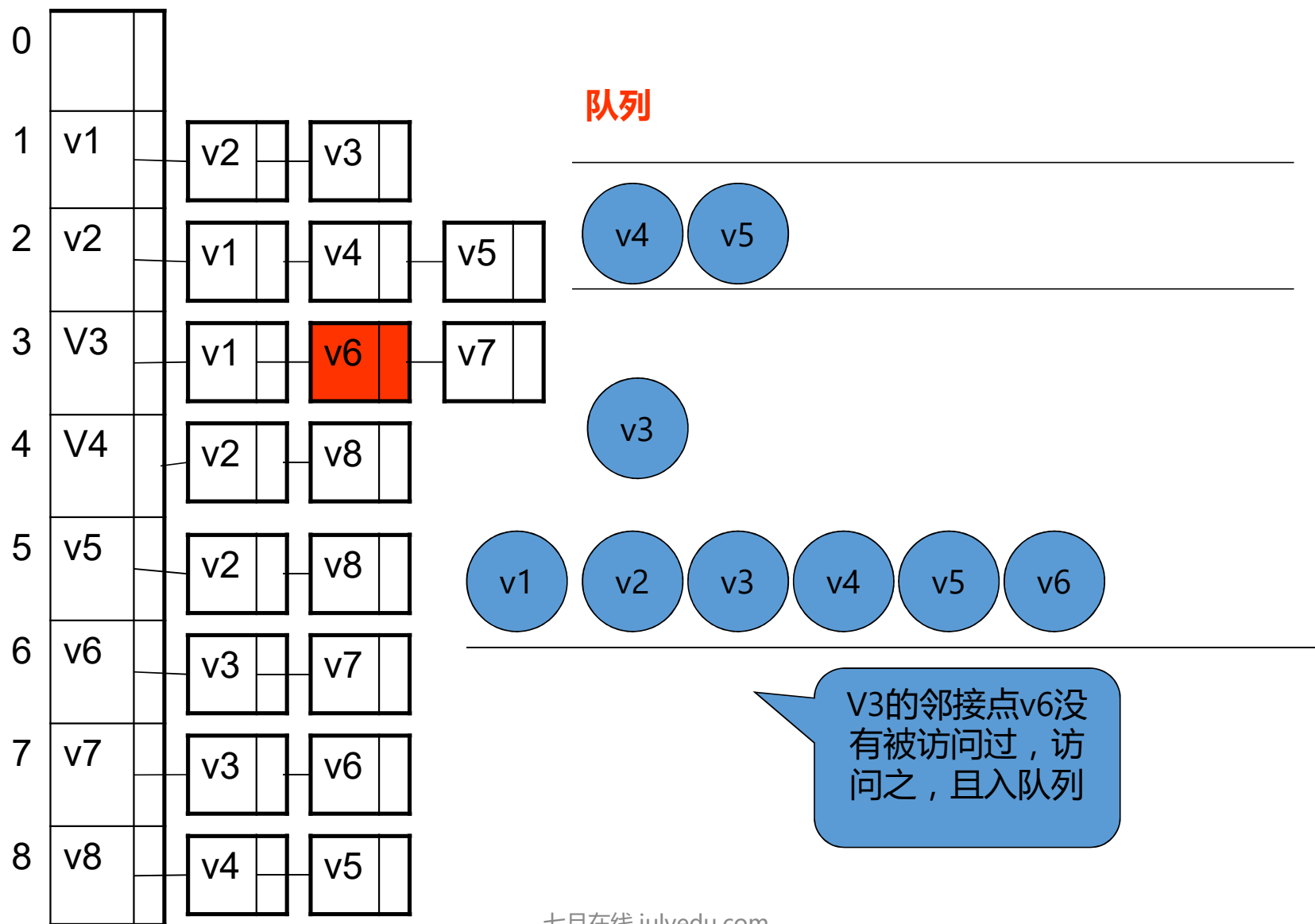


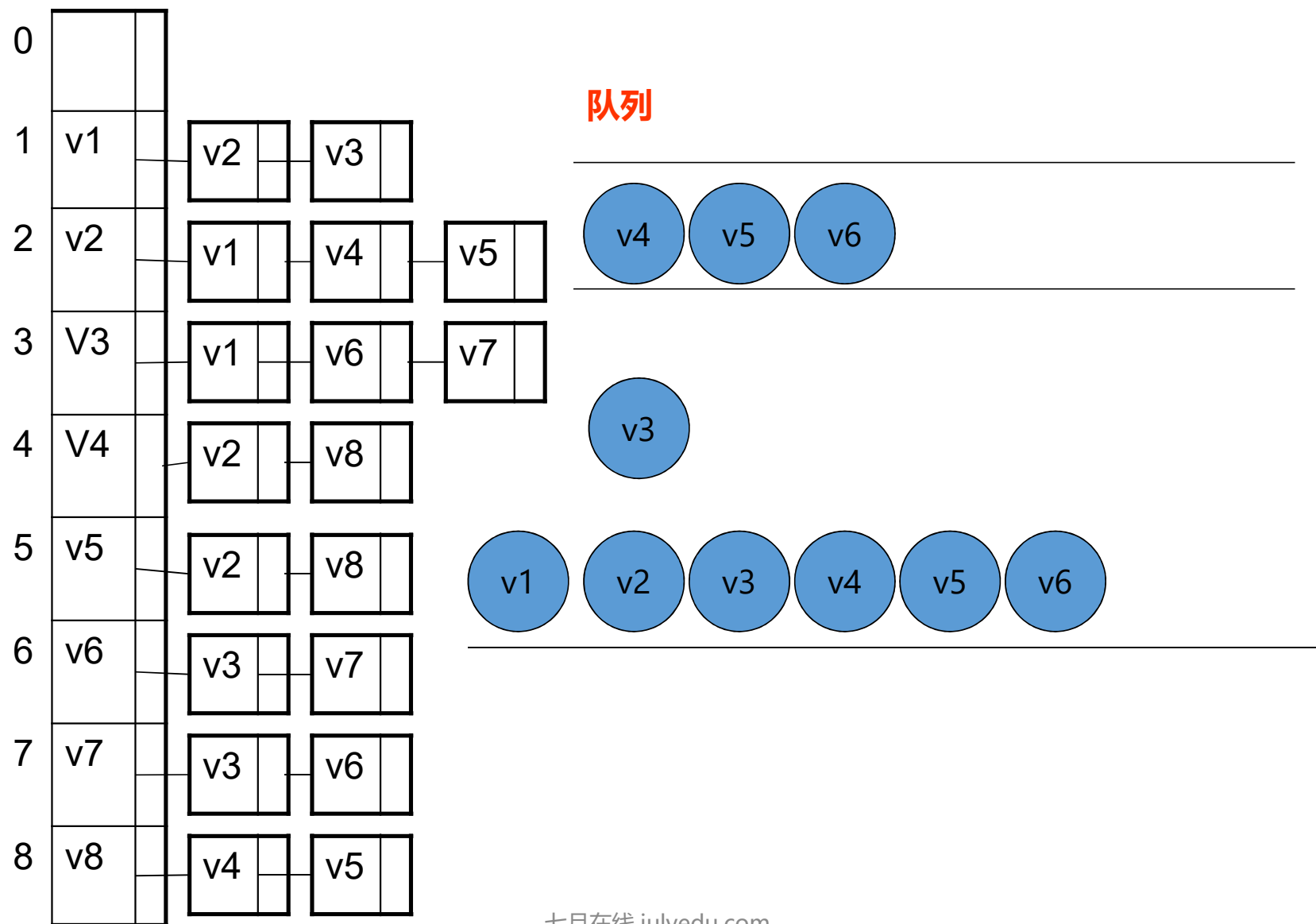


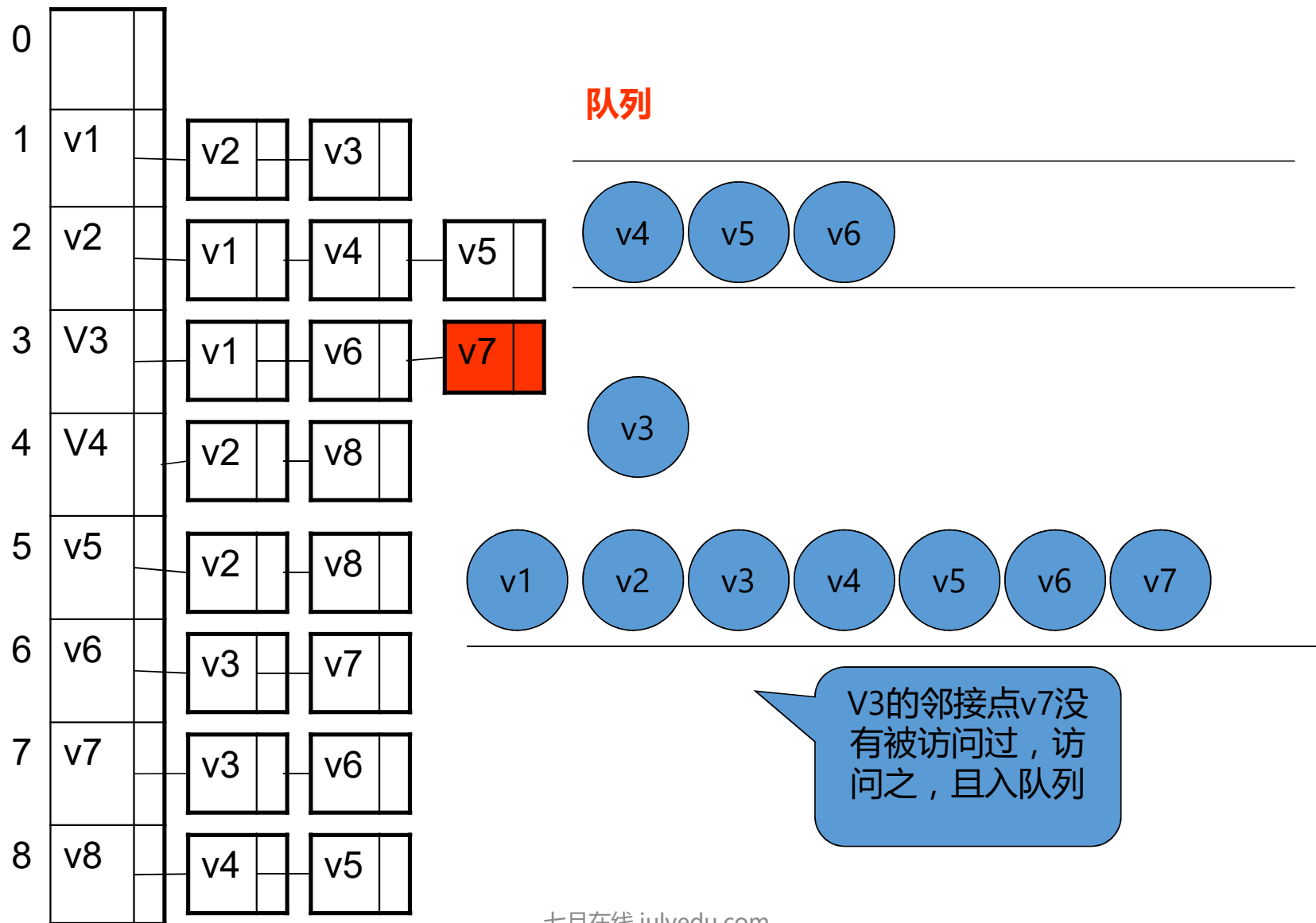


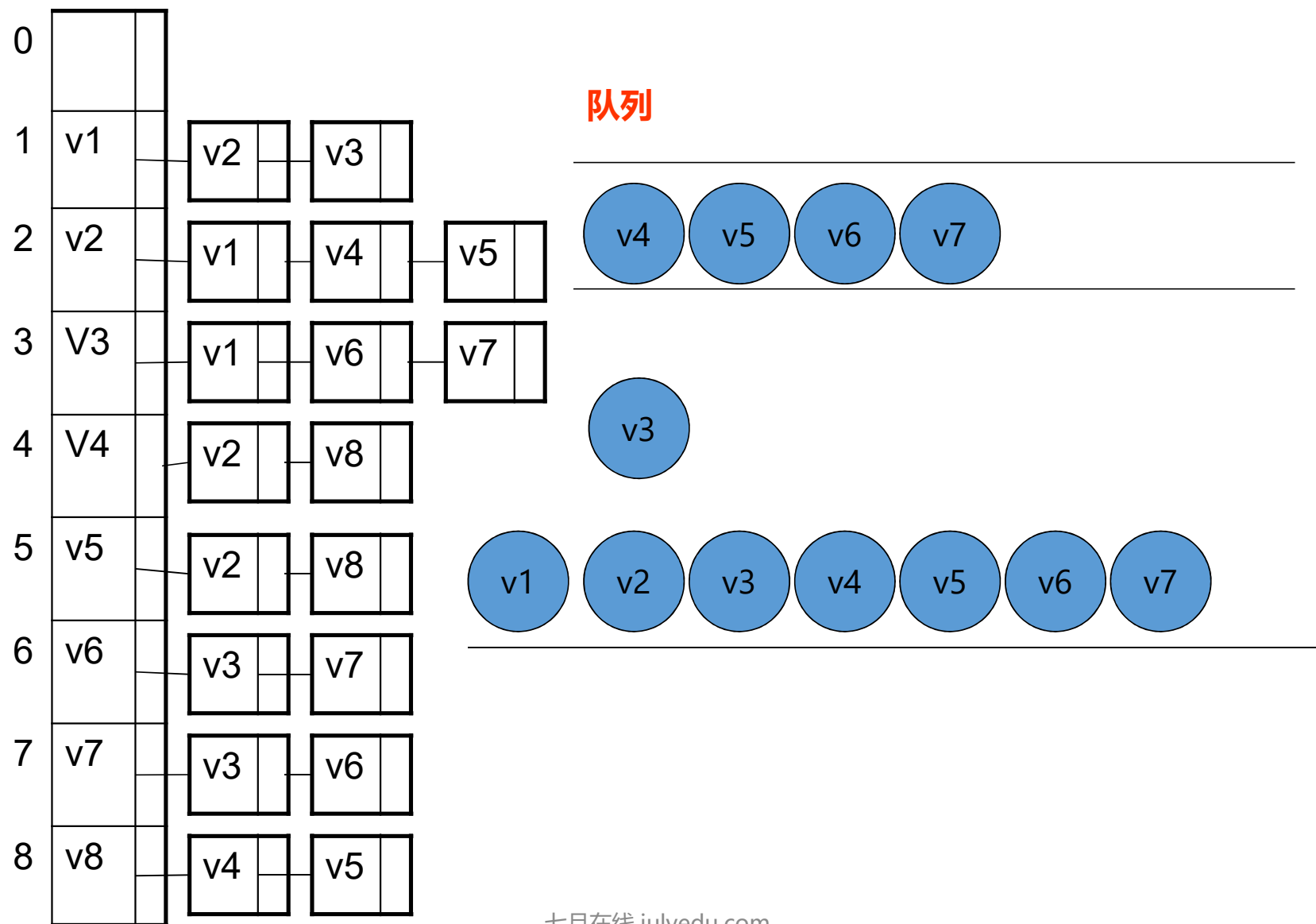


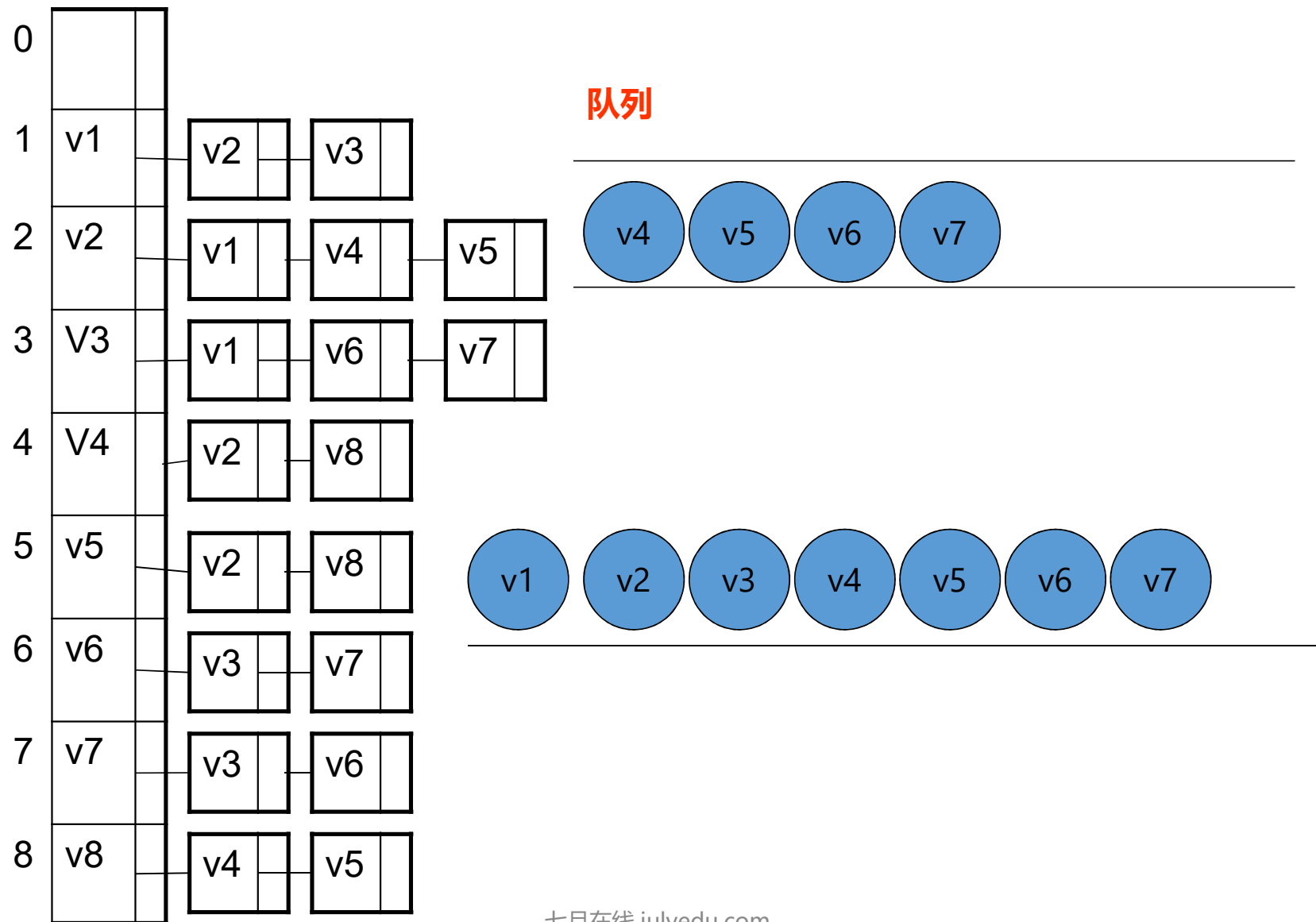


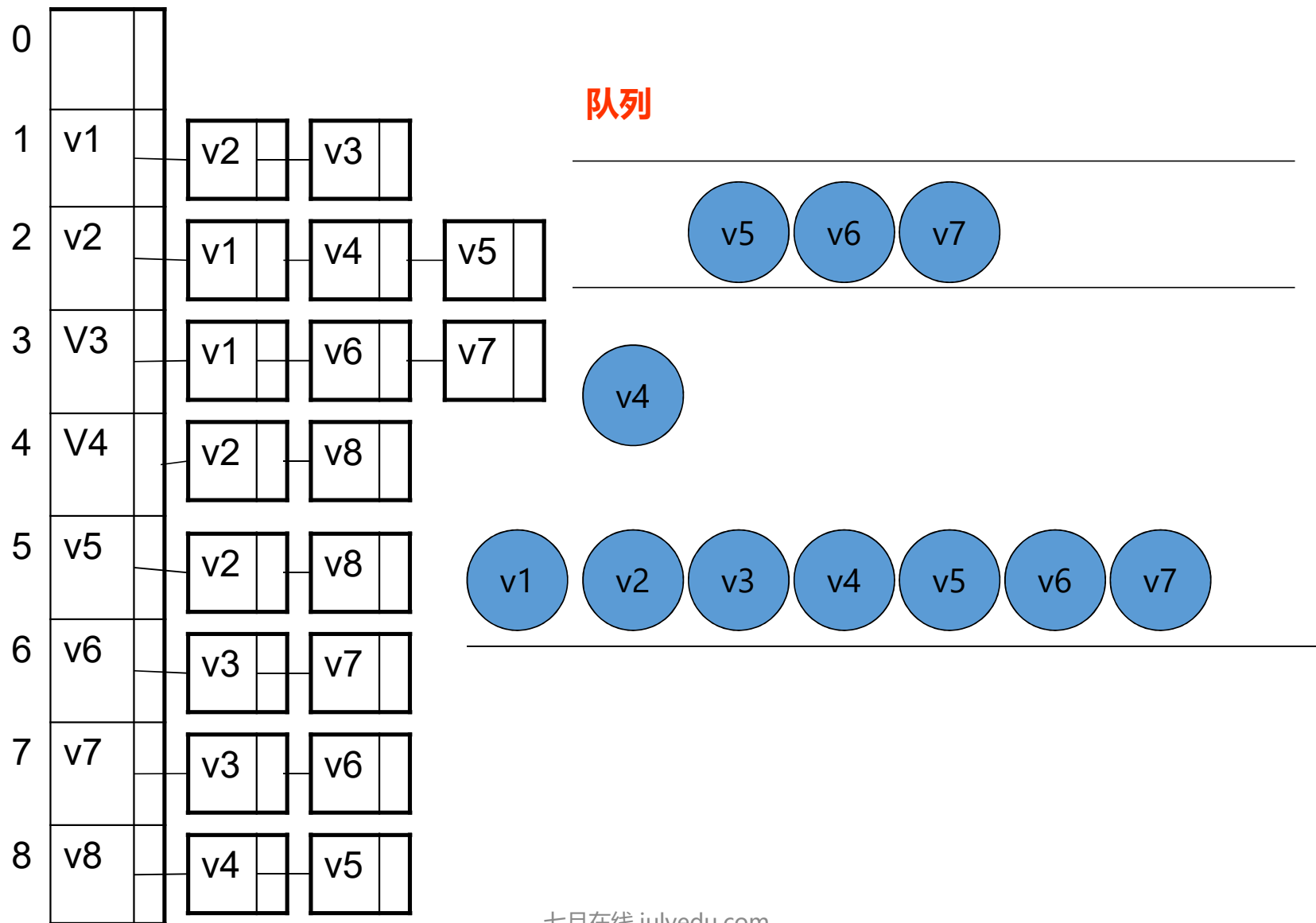


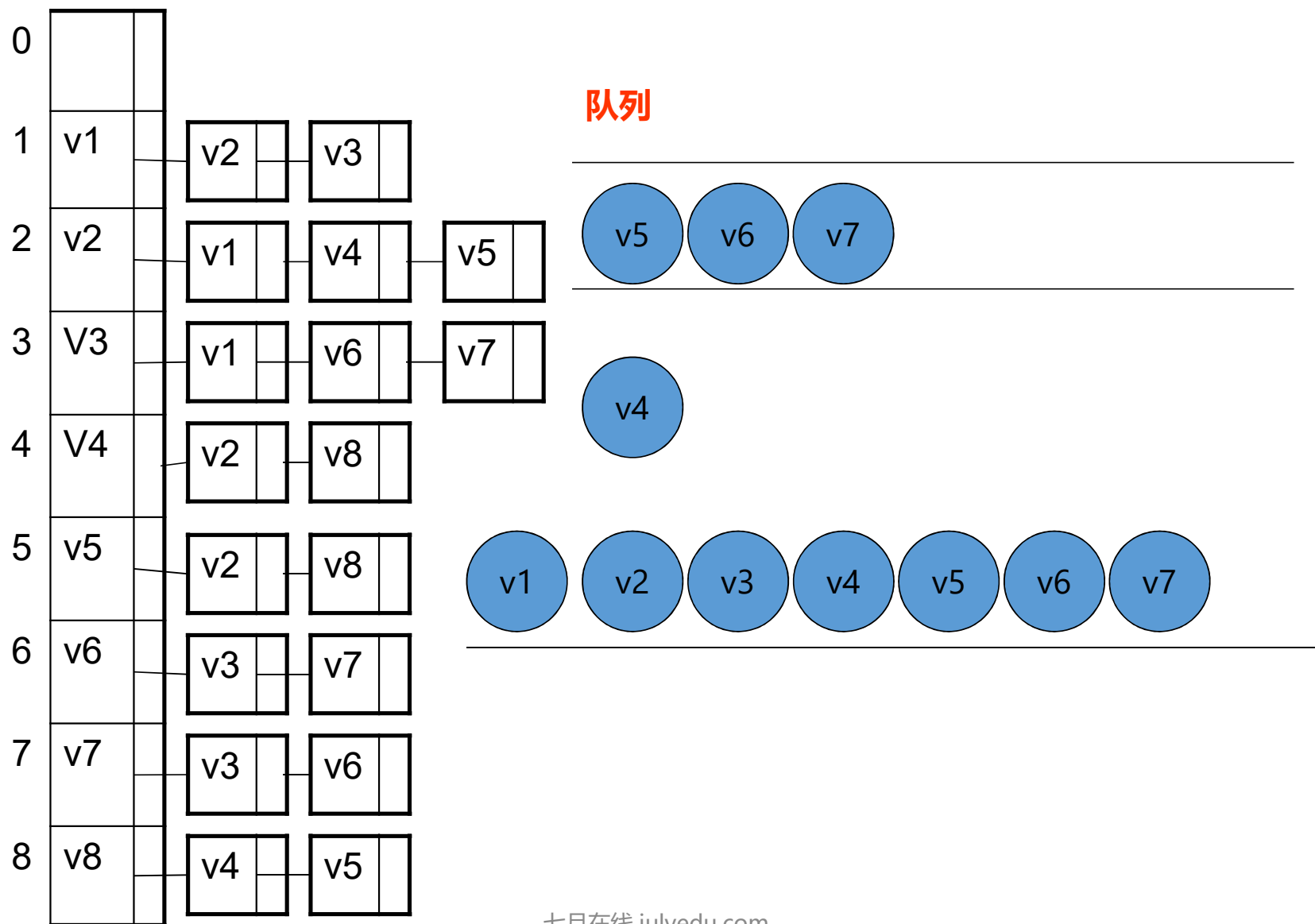


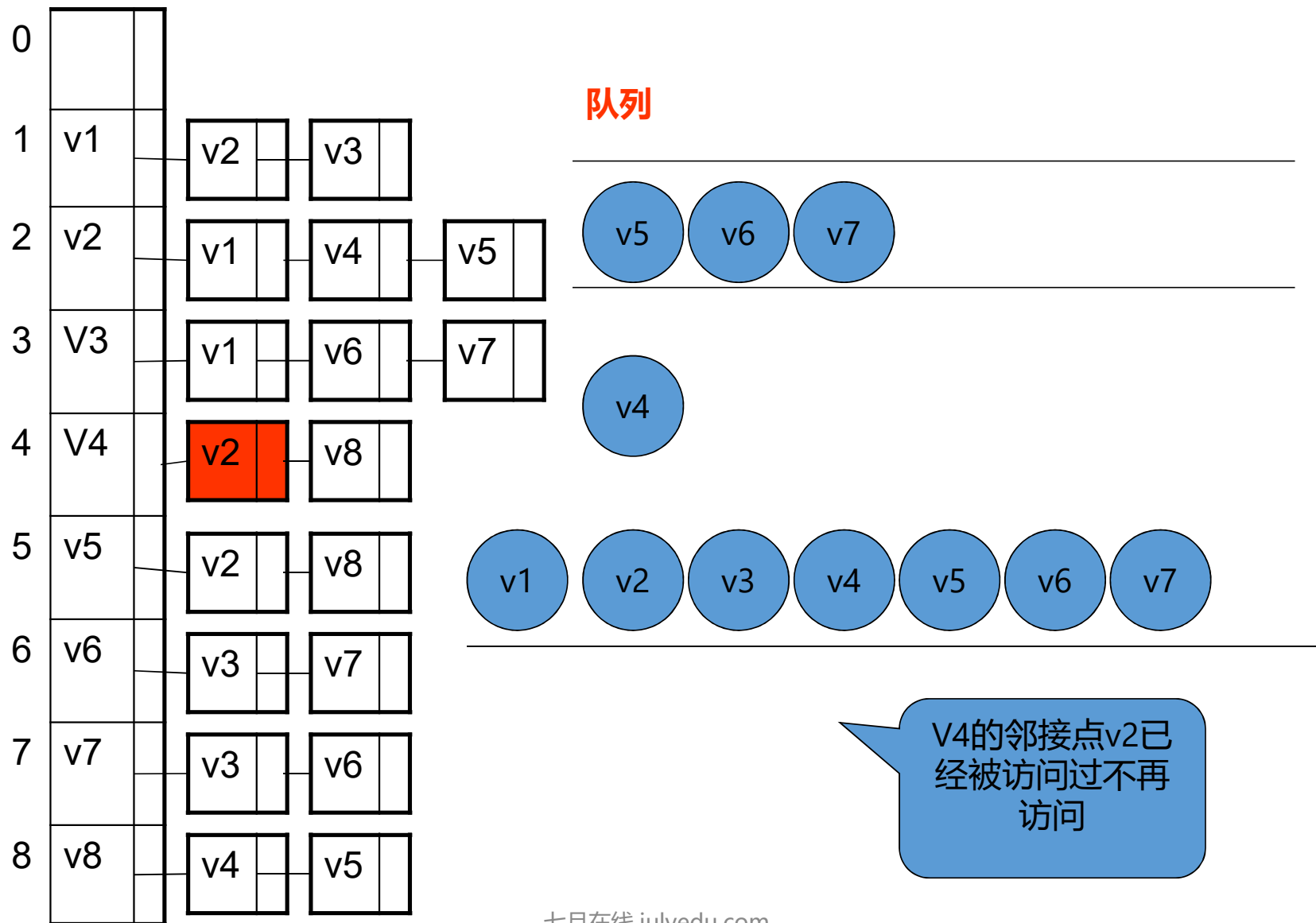


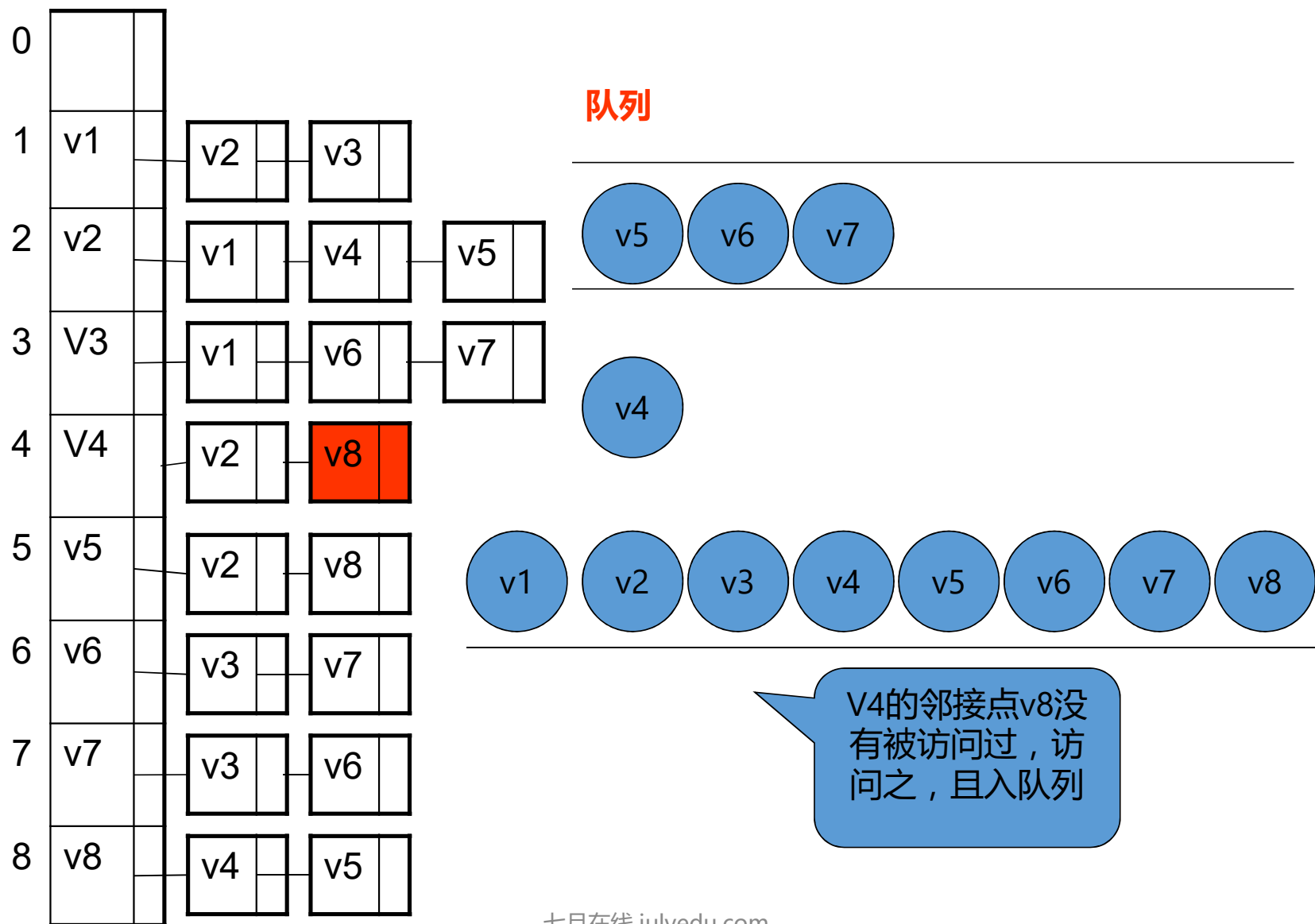


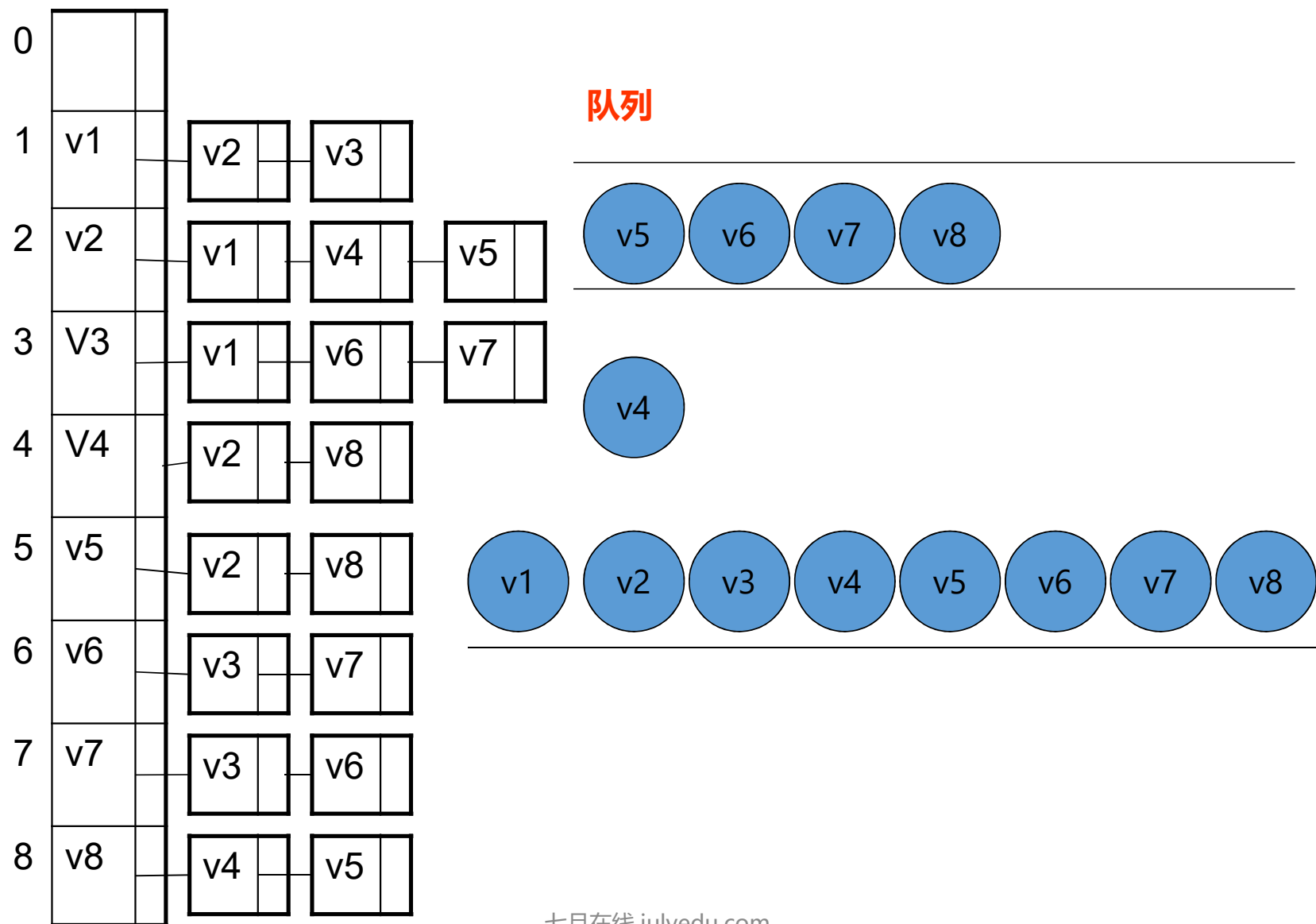


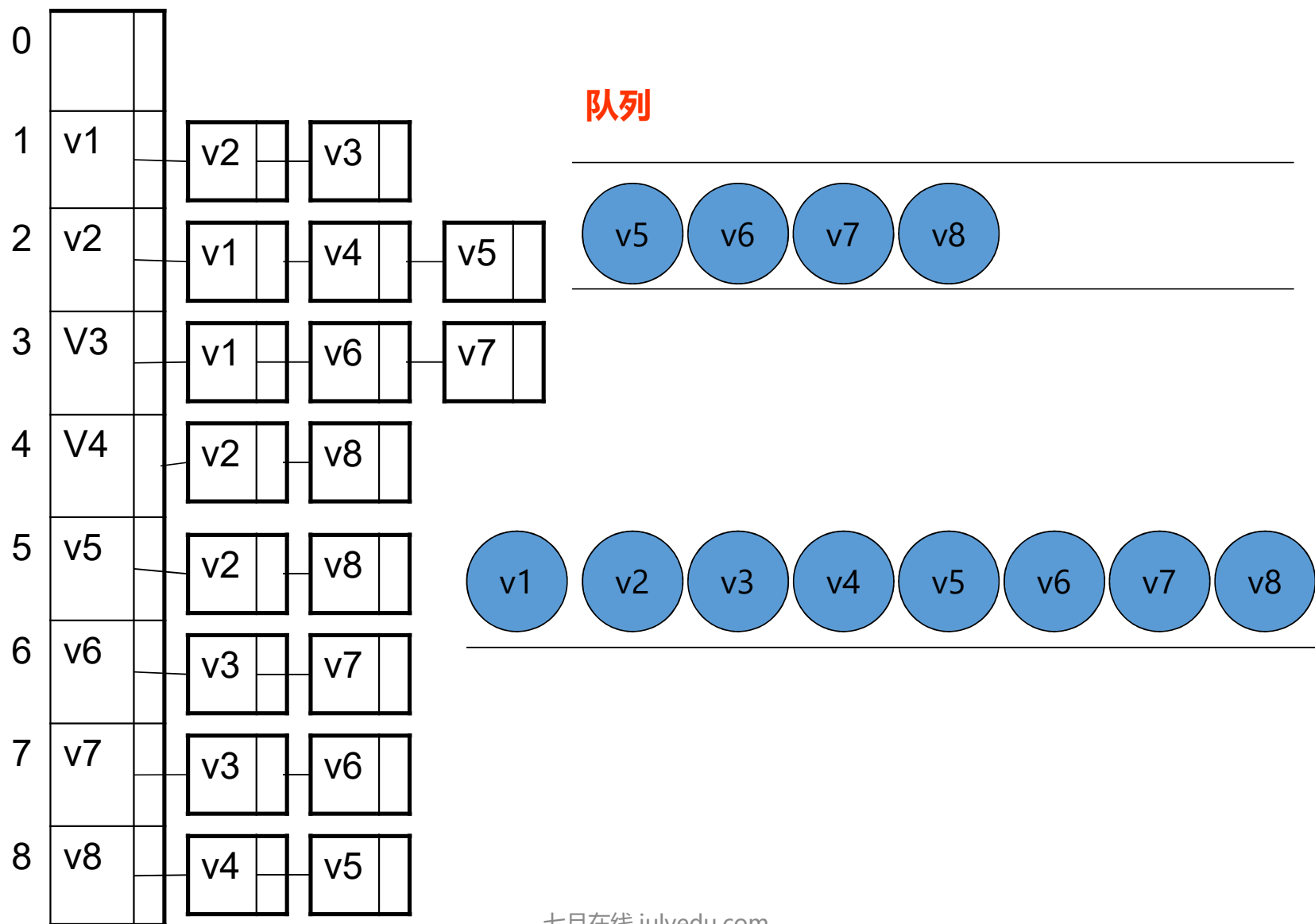


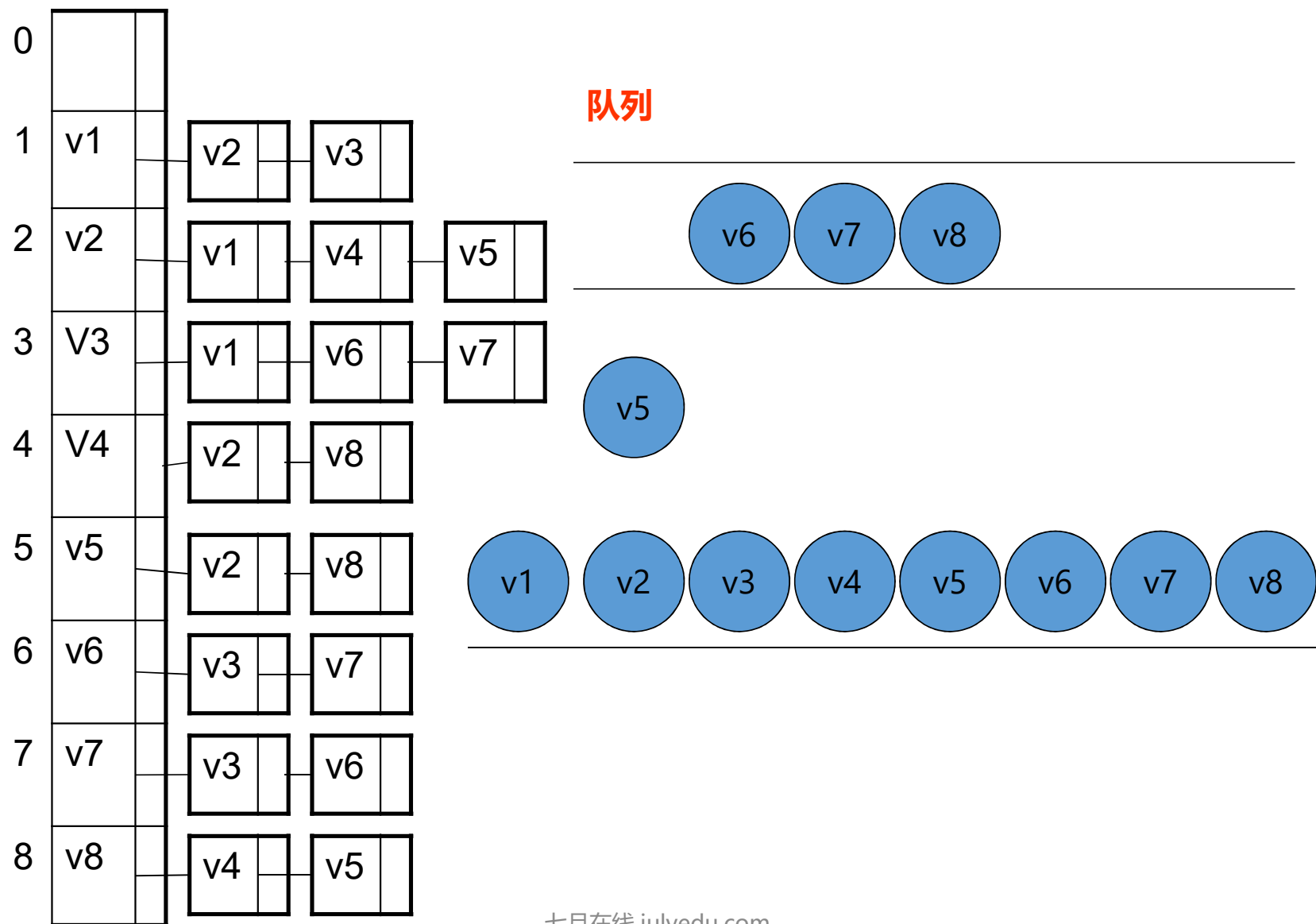


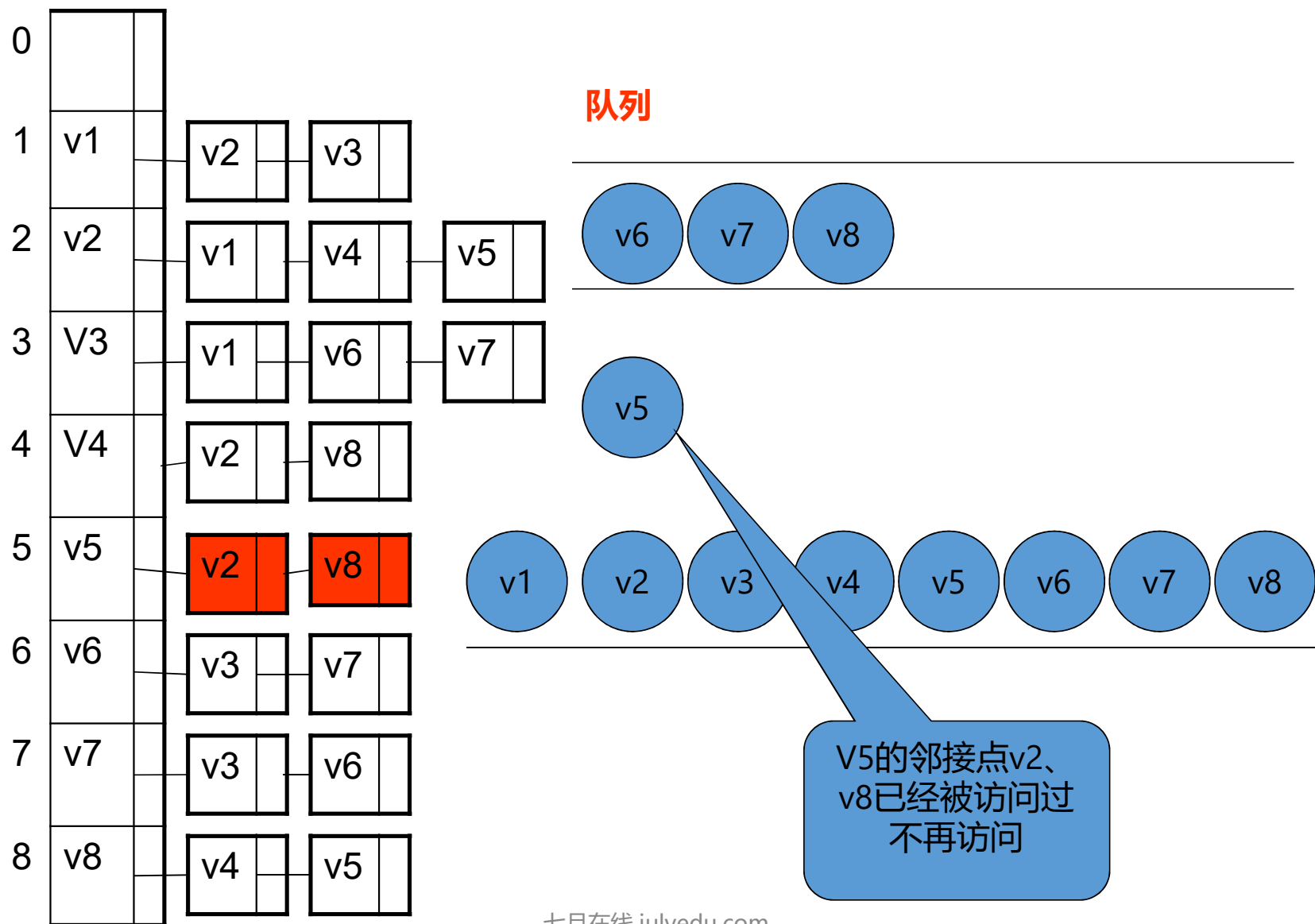


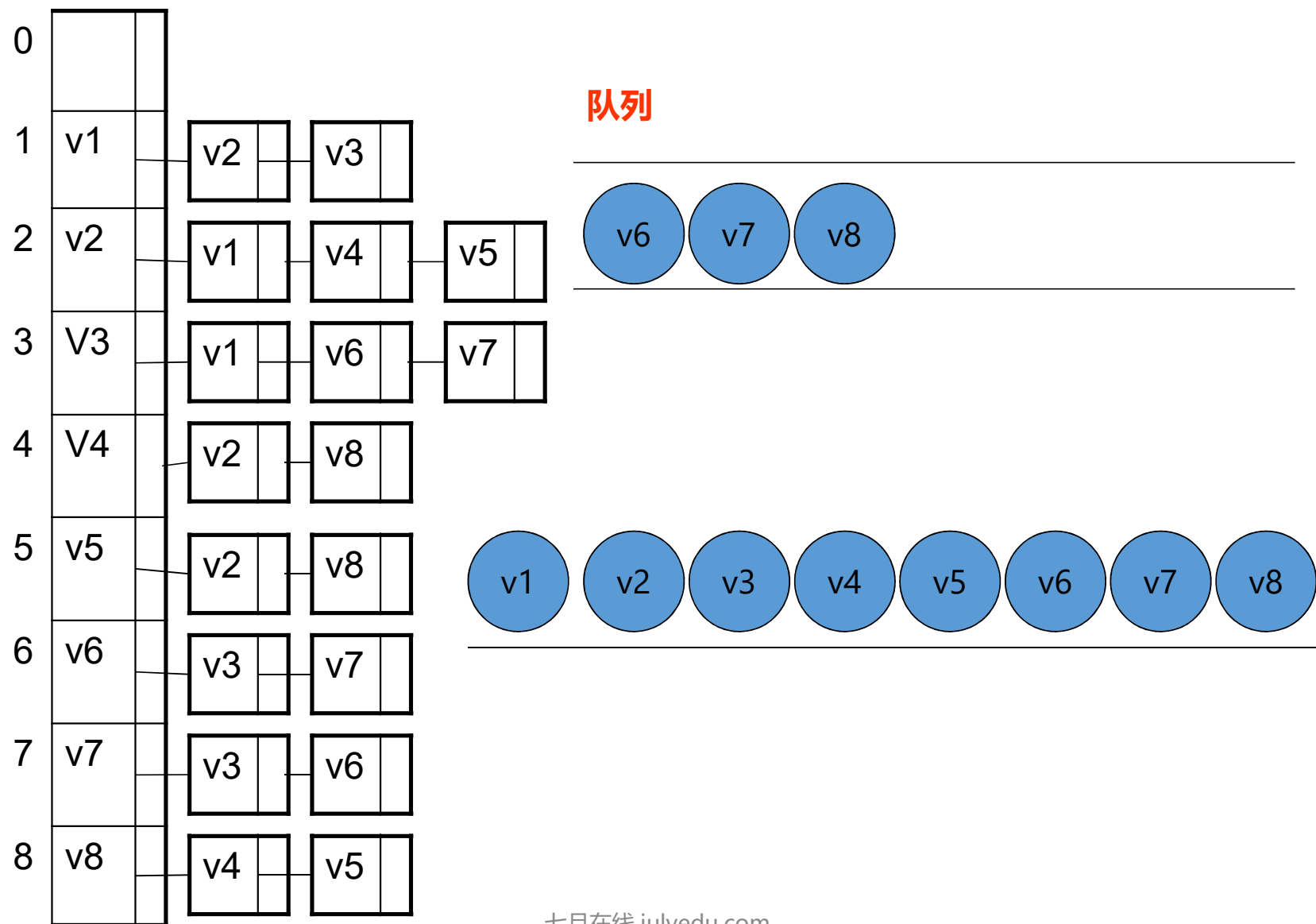


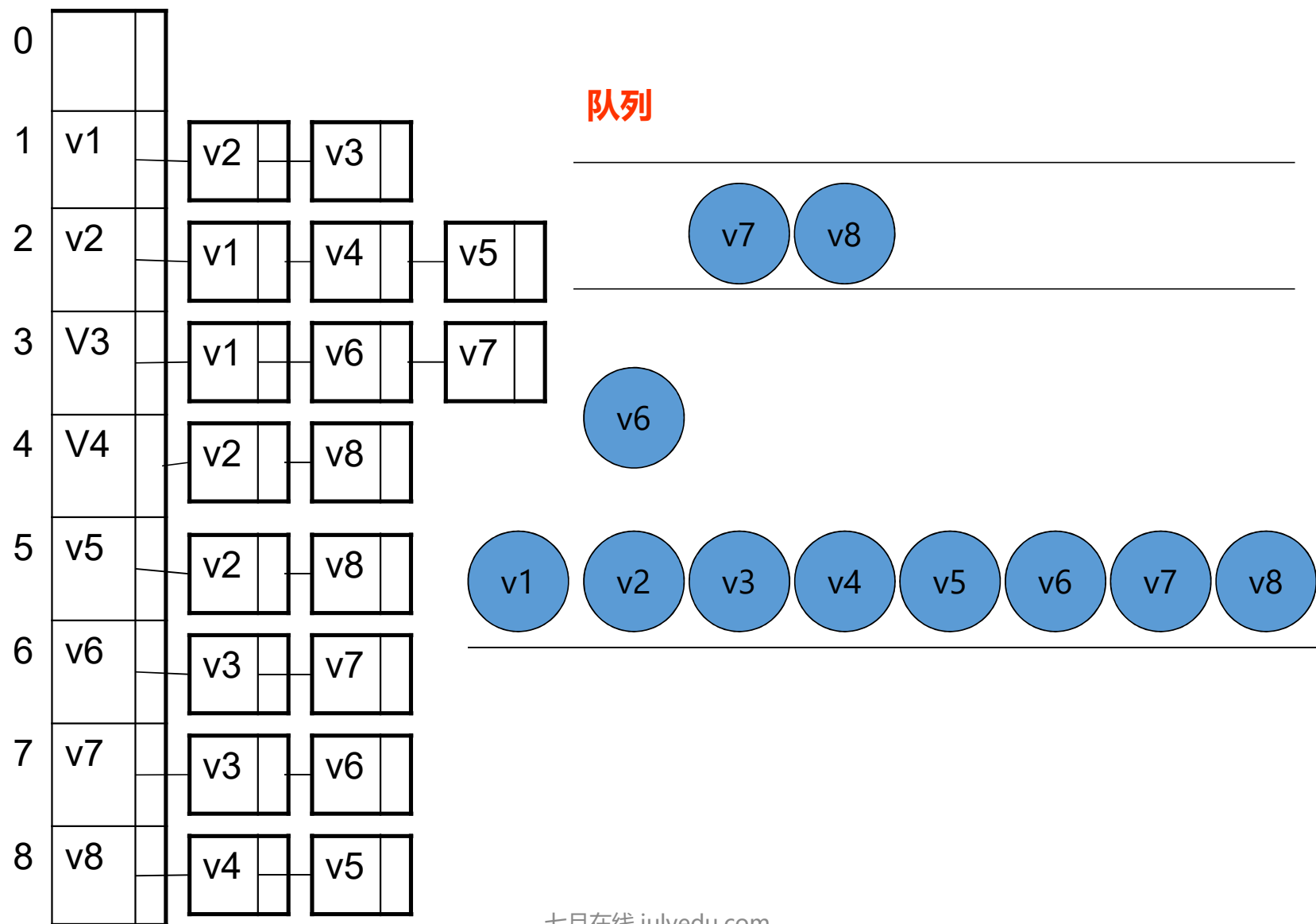


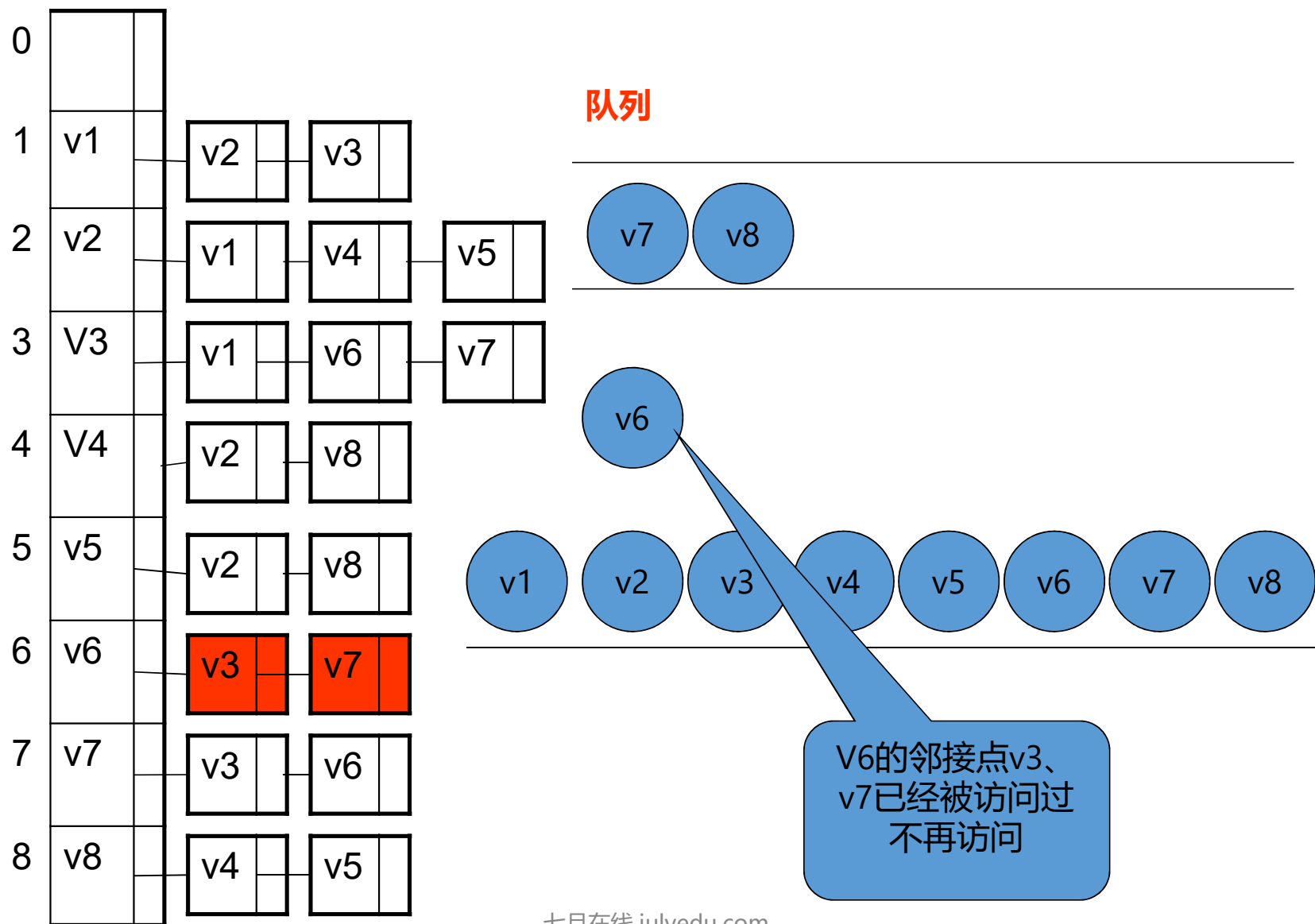


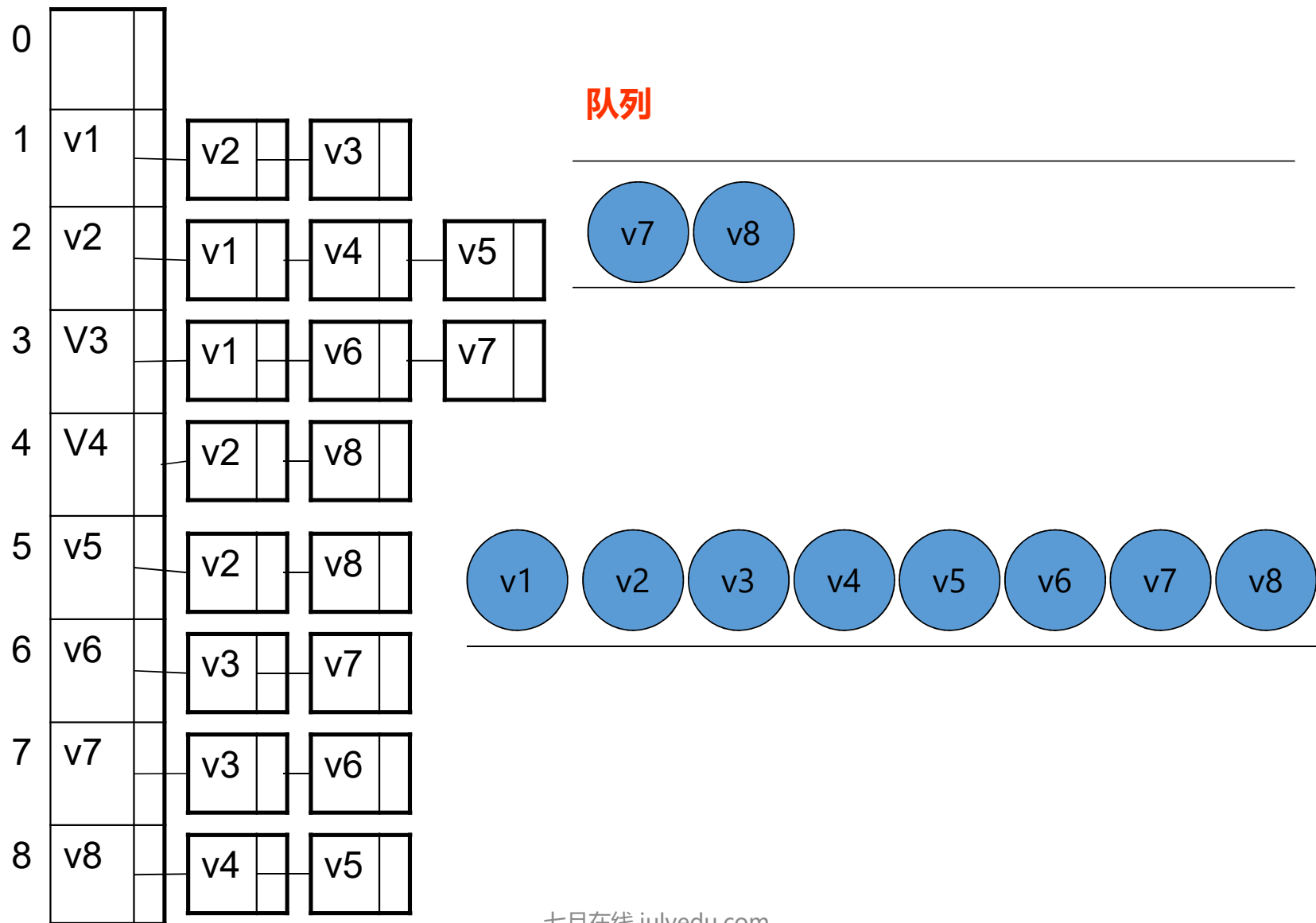


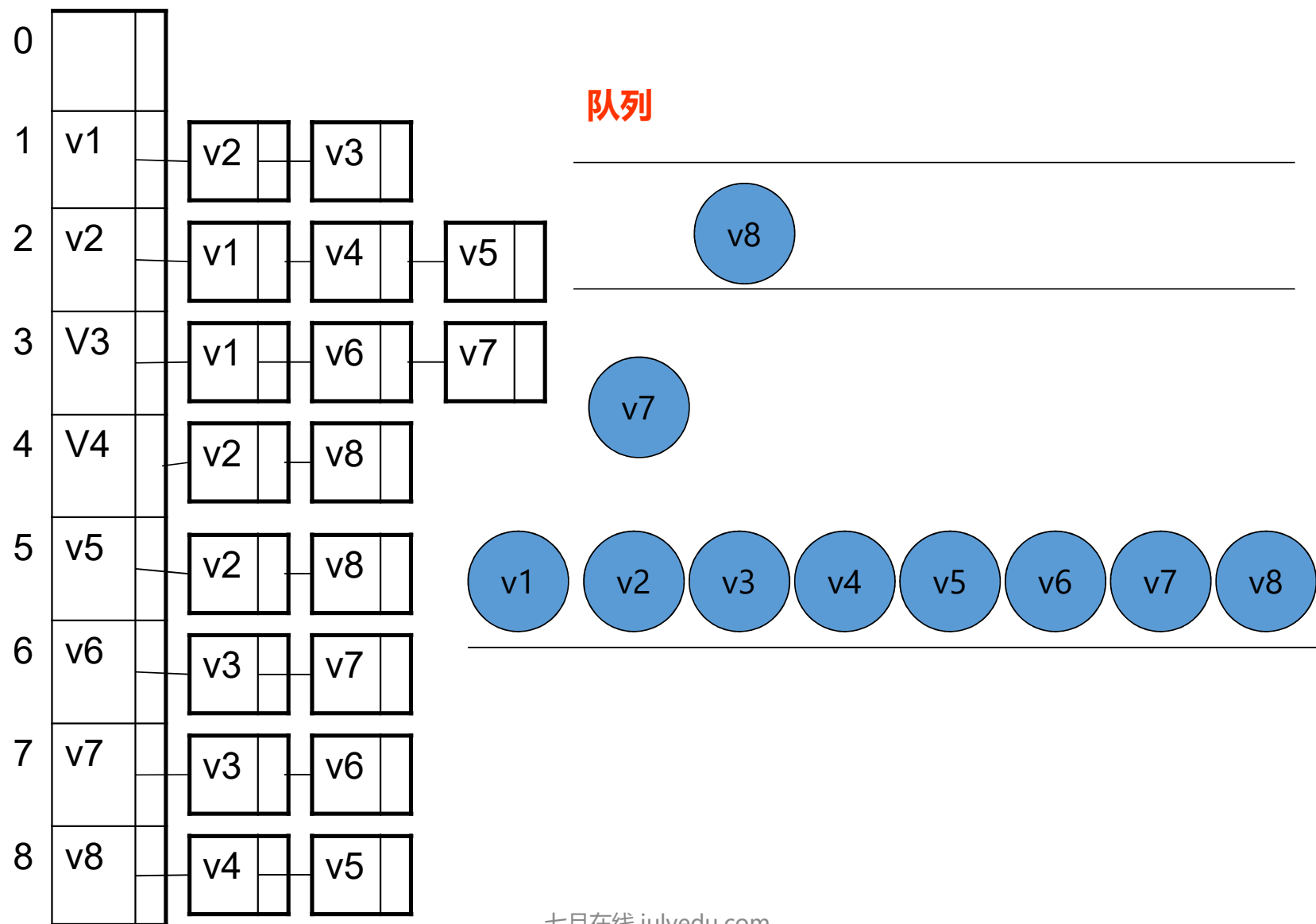


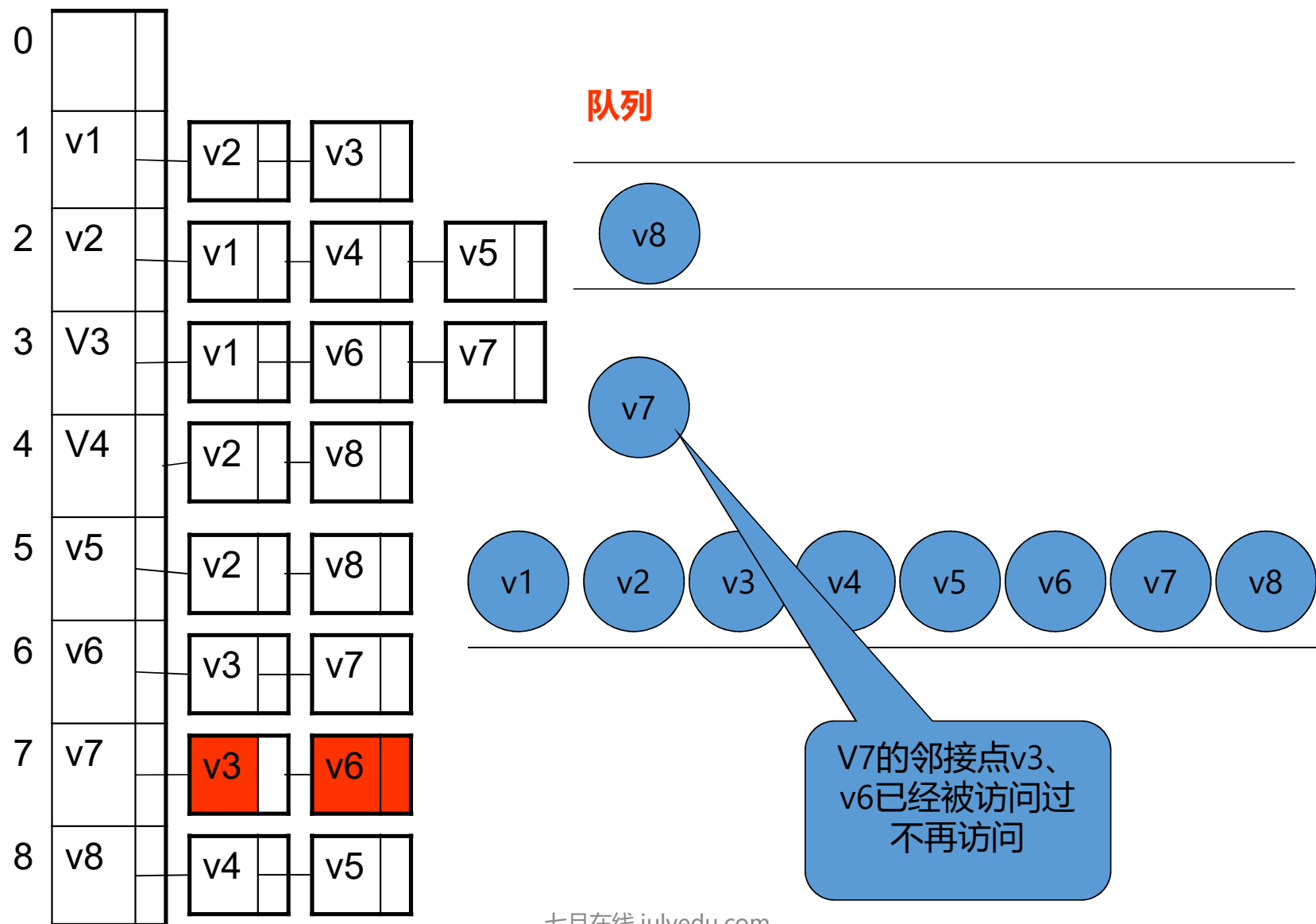


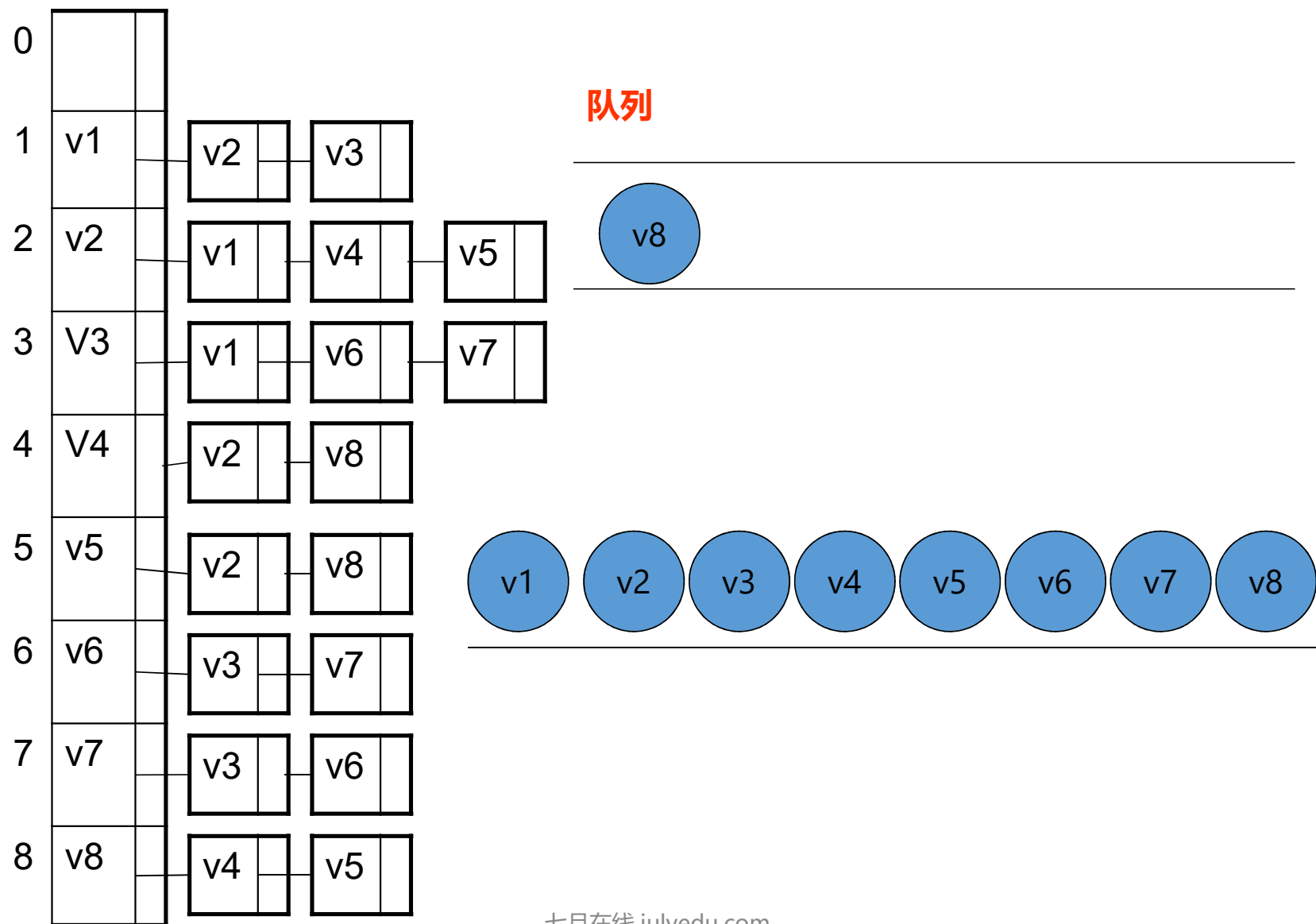


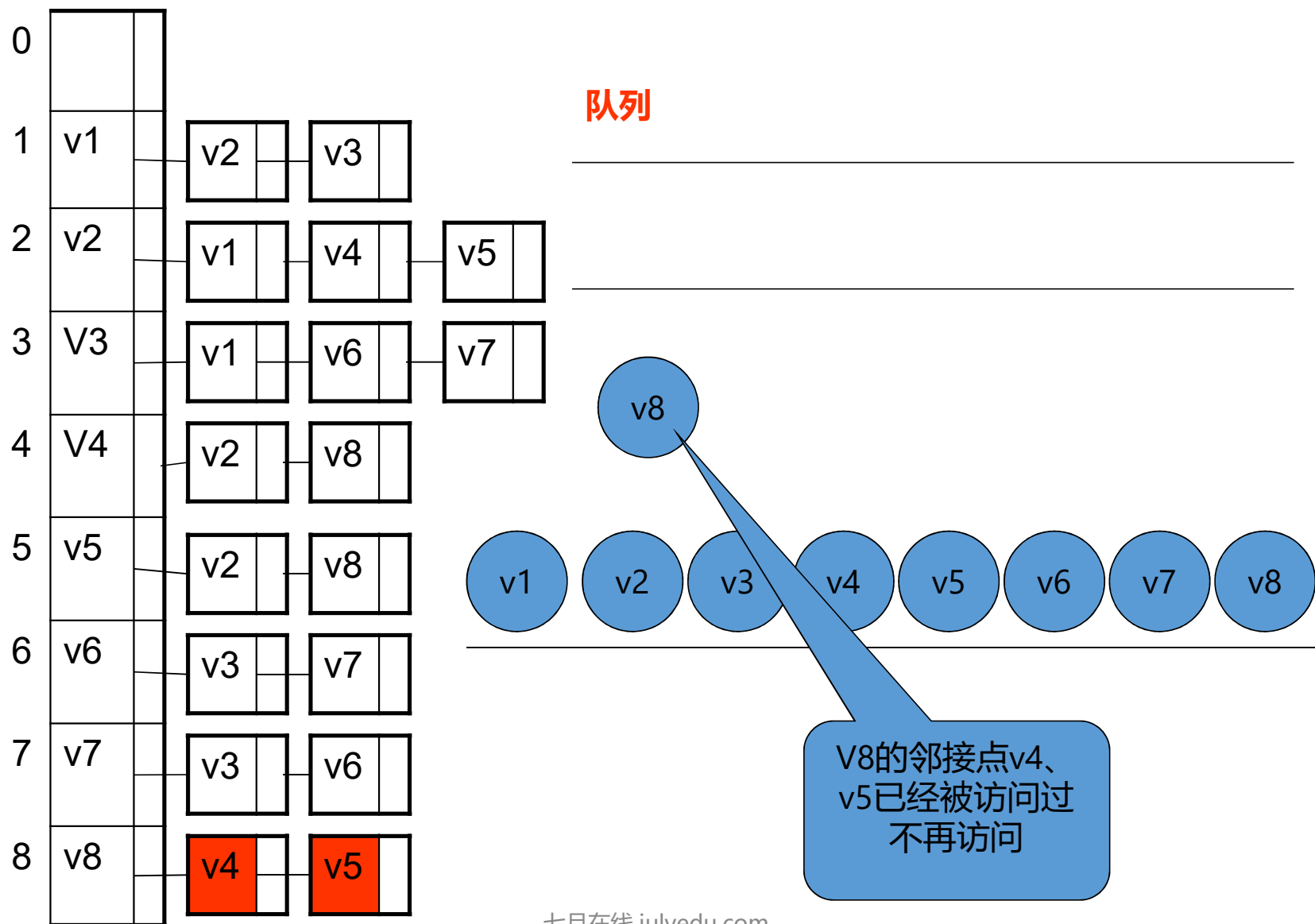


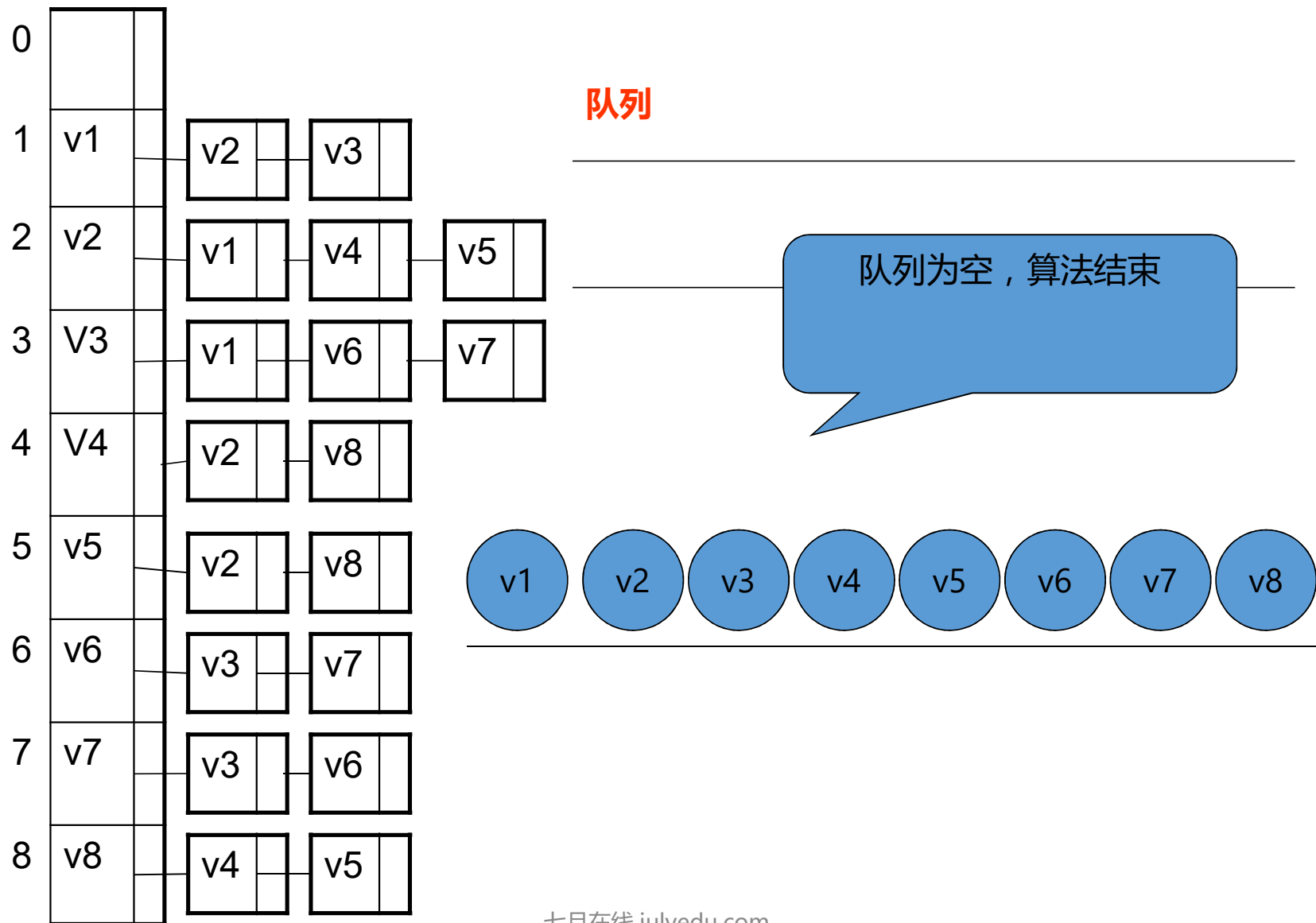












深度优先：伪代码

```
void DFS(int v)
```

```
    visited[v] = true
```

```
    for (v的每一个邻接点w)
```

```
        if (!visited[w])           //如果没有被访问过
```

```
            DFS(w)
```

N皇后问题

- 在 $N \times N$ 的棋盘上摆放 N 个皇后，使得任意两个皇后都不能处于同一行、同一列或同一斜线上
- 回溯法（暴力搜索**万金油**）
- 本质：深度优先（隐式图搜索）

N皇后问题

- 如何定义状态，关系？
- 时间复杂度 $O(N^N)$ （状态空间）
- 剪枝（确定性）
- 作业题：你写的代码能跑多少个皇后？
- Leetcode 51

骑士游历问题

- 国际象棋棋盘上，有一个骑士（马）从左下角出发，是否能不重复的遍历每一个格子
- 没啥好的办法，考虑暴力
- 如何定义状态，关系？
- 除了剪枝，还有什么办法？（求任意解和所有解）
- 启发式(A^*)
 - 改变搜索顺序
 - 不确定性

迭代加深

- 求最优 -> 求判定
- 不需要判重
- 搜到即是最优
- 前一个阶段相对下一个阶段仅仅是常数
- 迭代加深 + 启发式 (IDA*)

广度优先：伪代码

```
void BFS(int x)
    visited[x] = true
    Queue.push(x)
    While (!Q.empty())
        v = Queue.pop()
        for (v的每个邻接点w)
            if (!visited[w])
                visited[w] = true
                Queue.push(w)
```

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种子填充法

- Flood Fill 洪水填充法
- 目标：标记某块封闭的区域，并找出其边界
- 如何定义状态，关系？
- BFS犹如墨汁滴入清水

种子填充法

- Leetcode 200. Number of Islands
- Leetcode 130. Surrounded Regions

八数码

- 3*3的方格内有编号1-8的方块，求最少步数，恢复这些方块的顺序
- 深度优先 or 广度优先？
- 判重（Hash）

八数码

- 双向搜索
 - 起始点和目标点，轮流扩展
 - Hash表判断相遇
 - 复杂度
- 启发式
 - 价值函数（启发函数）
 - 优先队列（堆）

总结

- DFS vs BFS
 - 都为暴力搜索，但搜索顺序不同
 - 栈 vs 队列
 - 可行解 vs 最优解
 - 递归 vs 非递归
 - 空间占用，BFS需要存储状态，DFS无需
- 事无绝对，仅供参考

Q&A