EDI: Third Lab Report

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Abstract

1 Parallel Connections

Page Load Time (PLT) is one of the most critical performance evaluation metrics for a website, as it correlates directly with user experience, taking into account both the time to download the html of a page, as well as the time to download the whole set of its embedded objects, as a browser would typically do when directed to navigate to a page.

A good Page Load Time in 2022 is between 1 and 2 seconds, but it depends, and having an even lower Page Load Time is desirable; users obviously tend to abandon websites that take more time to load than their competitors. [1]

1.1 HTTP/1 vs HTTP/2

The first experiment performed here pertains to the effect of increased parallel connections on Page Load Time. As power users will know, modern web browsers allow their users to peek into various configuration settings and tweak them, this is especially true of Mozilla Firefox, where the user can access various technical settings by typing in 'about:config' into the navigation bar and hitting enter. More specifically, one of these configuration parameters is called 'network.http.max-persistent-connections-per-server', and as the name suggests, it governs the maximum amount of parallel persistent TCP connections that the browser is allowed to open for any particular web server on the Internet. [2]

By default, the value for this option is set to '6', but can increasing it arbitrarily lead to an overall improvement in the performance of the web browser when loading pages?

Here comes to play a very important distinction between websites, and that is difference between HTTP/1 and HTTP/2: the first two major versions of the ubiquitous Hypertext Transfer Protocol. There are various ways in which HTTP/2 tried to improve on HTTP/1, but what is of specific interest to us in this particular inquiry is that: websites running on HTTP/1 tend to favor connections to multiple separate servers, to dowload the objects of a page in parallel, and overcome the limitation imposed by browsers on the number of parallel connections to a single server, a technique known as Domain Sharding [3].

HTTP/2 helped overcome this very necessity, by introducing Streams that could convey multiple concurrent (and bi-directional) HTTP requests within a single persistent TCP connection; thus eliminating the need to connect to separate resource servers, and even eliminating the overhead associated to having to open multiple TCP connections with them. This is something that, at least on the long run, will render the technique of Domain Sharding obsolete.

So, to go back to the original question: resource intensive HTTP/1 websites may benefit from an increased number of parallel TCP connections, as they only support a single HTTP request per TCP connection. On the other hand, HTTP/2 websites support multiple concurrent HTTP requests over the same TCP connection, hence, from a theoritical viewpoint, it shouldn't make much of a difference to them if the client supports a greater number of TCP connections.

1.2 Methodology and experimental setup

The experiment was performed in practice, by measuring the PLTs of 6 different websites (3 HTTP/1 and 3 HTTP/2 websites) under two different conditions: 1 and 6 maximum connections per server.

The browser in question was Firefox, and the number of maximum connections was set tweaking the aforementioned: "network.http.max-persistent-connections-per-server" parameter.

Although the results aren't exhaustive, they clearly show that HTTP/1 websites fare much better when granted a higher number of connections per sever.

		Average PLT (seconds)
HTTP Version	Connections	
1	1	3.017778
1	6	2.212222
2	1	3.178889
	6	3.486667

Table 1: ...

2 Caching Policies

Cache is crucial to the correct operation web infrastructure, as its valid usage provides advantages for the origin servers, as well as for the end users of a particular web-service.

From the service provider's point of view, cache helps to optimize the number of accesses made by clients to the origin web server, easing up the load on them; and, from the client's side, cache helps in reducing PLT, especially if the assets (images, stylesheets, scripts) that have to be served with the page are many.

Caches can be remote or local. Remote caches can either fall under the category of Proxy or Managed caches. Proxy caches follow a "pull model", where any client that accesses content X within a particular local network, will trigger the caching of content X on the local network's proxy, so as to avoid having to download the content again from the wider Internet. Managed caches follow a "push model", in that the service providers manually load content to the cache, with the goal of easing up the load on their origin servers by distributing client requests over multiple servers, an example of a Managed cache is a CDN (Content Delivery Network).

Caches can also be private or public, the idea behind private caches is to avoid disclosing private information that has to be cached on the user's side.

All of these, and many more, caching policies and behaviors are governed by specific HTTP request and response headers, such as the Cache-Control, If-Modified-Since and ETag headers. [5]

2.1 Methodology and experimental setup

The goal of this experiment was to assess whether the caching policies of various websites proved to be effective, at least on the client's side, by comparing the performances of said websites with and without cache enabled.

This was done on a real browser (Firefox), with the help of a custom bash script with the xdotool [6] utility, to automate the process of reloading the web-pages, and collecting the data consistently, taking several readings. The script in question, dubbed "scripton.sh", is available on the Github repository associated to this report [7], but it has to be fine tuned to the screen-dimensions of the particular host computer, as it leverages the position of the cursor on the screen to interact with Firefox and copy the textual output from the Network Monitor's interface; this process of fine-tuning can be performed manually, with the help of the included "calibrate.sh" script.

9 websites (3 secure HTTP/1, 3 unsecure HTTP/1 and 3 HTTP/2), were chosen for this test, and the data collected for each was:

- 1. the PLT, and
- 2. the number of requests made to the server

As it is evident from the results, in all cases, caching does prove to be advantageous for the client, as it saves quite a bit of time loading the page, decreasing the number of requests that have to be performed towards the server.

		number_requests	plt_seconds
http_version	cache		
http1	False	39.950000	2.343500
	True	29.923077	1.543077
http2	False	80.083333	3.708333
	True	56.333333	2.652667
unsecure-http1	False	78.666667	4.010000
	True	45.500000	3.580000

Table 2: ...

3 Apache Benchmark

There are several automated CLI benchmarking tools to measure the performance of webservers, and Apache Benchmark (AB) is one of them. Unfortunately, AB doesn't support the second version of the HTTP protocol, but it provides several options to test the performance of HTTP/1 servers, and especially Apache webservers.

3.1 Methodology and experimental setup

The idea behind this experiment was to test the effects of the concurrency (-c) and keep-alive (-k) options of Apache Benchmark, on the performance of 3 websites, when issuing a fixed amount of requests (20).

The concurrency (-c) option regulates the number of multiple requests to perform at a time; since the servers in questions are HTTP/1 only servers, concurrent requests will have to be performed over multiple TCP connections.

The keep-alive option (-k) specifies whether to request the usage of the HTTP KeepAlive feature (aka: persistent connections), which means that a single TCP connection will be re-used for multiple sequential HTTP requests, thus reducing the overhead of TCP connection creation [4]. Obviously, since this is HTTP/1, concurrent requests cannot share the same TCP connection, as already stated.

Each website was tested with the following 4 commands:

- 1. ab -c 1 -n 20 http://example.com/
 - 20 requests, one at a time, without reusing the TCP connection.
- 2. ab -c 10 -n 20 http://example.com/
 - 20 requests, 10 at a time, without reusing TCP connections.
- 3. ab -k -n 20 http://example.com/
 - 20 requests, one at a time, reusing the TCP connection.
- 4. ab -c 10 -k -n 20 http://example.com/
 - 20 requests, 10 at a time, reusing TCP connections.

The first command (-c 1 and no -k) is expected to be the slowest, as the requests are all performed sequentially, and, for each request, a new TCP connection must be opened, with all of the overhead associated to that costly operation.

The last command instead (-c 10 and k), is expected to be the fastest, as it combines the benefits from issuing multiple parallel requests, with the benefits of reusing already opened TCP connections.

Sure enough, these are the average times that it took to run these tests:

Options	Average Time (seconds)
-c1	5.191
-c1 && -k	3.194
-c10	1.036
-c10 && -k	0.939

Table 3: ...

4 nghttp and h2load

Among the tools that can be used to measure the performance of HTTP/2 servers are nghttp, which is an HTTP/2 client, and h2load, which is a benchmarking tool.

4.1 The Role of Warmup Time in h2load

Warmup time is used by the benchmarking tool's client to start a connection with the webserver, before actually beginning to transfer benchmark data, you can verify that some traffic is exchanged during the warmup phase using a sniffer such as Wireshark, as figure 1 shows:

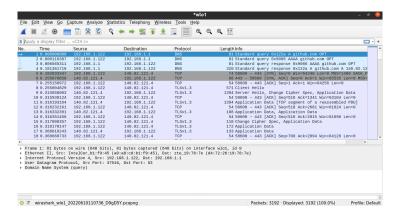


Figure 1: captured during warmup time

The packets were captured during the warmup phase of the command:

h2load -warm-up-time 1 https://github.com -duration 5

with 1 second of warmup time and 5 seconds for the duration of the actual benchmark.

When warmup time is greater than 0, "time for connect" and "time to 1st byte" tend to become 0, as the following tables show:

Without warmup time (warmup time = 0)

	min	max	mean	sd	+/- sd
time for request:	Ous	Ous	Ous	Ous	0.00%
time for connect:	157.24ms	157.24ms	157.24ms	Ous	100.00%
time to 1st byte:	227.56ms	227.56ms	227.56ms	Ous	100.00%
req/s	0.00	0.00	0.00	0.00	100.00%

Table 4: ...

With warmup time = 1 second

	min	max	mean	sd	+/- sd
time for request:	106.22ms	164.90ms	134.89ms	14.35ms	67.57%
time for connect:	Ous	Ous	Ous	0us	0.00%
time to 1st byte:	Ous	Ous	Ous	0us	0.00%
req/s	7.40	7.40	7.40	0.00	100.00%

Table 5: ...

5 References

- 1. https://www.hobo-web.co.uk/your-website-design-should-load-in-4-seconds/
- 2. https://www.computerworld.com/article/2541429/hacking-firefox--the-secrets-of-about html?page=5
- 3. https://blog.stackpath.com/glossary-domain-sharding/
- 4. https://httpd.apache.org/docs/2.4/programs/ab.html#synopsis
- 5. https://developer.mozilla.org/en-US/docs/Web/HTTP/Caching
- 6. http://manpages.ubuntu.com/manpages/trusty/man1/xdotool.1.html
- 7. https://github.com/aiman-al-masoud/edi_reports