

Verbose proof for **vc**.

vc:

$$\frac{}{\{1\} \quad \forall (\text{value}: \mathbb{Z}, \text{length}: \mathbb{N}, x: [\mathbb{N} \rightarrow \mathbb{Z}], y: [\mathbb{N} \rightarrow \mathbb{Z}]): \\ (0 < \text{length}) \supset \text{rep}(\text{length}, 0, x, y) = \text{DP_RESULT}(0, \text{length}, x, y)}$$

Inducting on length on formula 1,

we get 2 subgoals:

vc.1:

$$\frac{}{\{1\} \quad \forall (\text{value}: \mathbb{Z}, x: [\mathbb{N} \rightarrow \mathbb{Z}], y: [\mathbb{N} \rightarrow \mathbb{Z}]): \\ (0 < 0) \supset \text{rep}(0, 0, x, y) = \text{DP_RESULT}(0, 0, x, y)}$$

Simplifying, rewriting, and recording with decision procedures,

This completes the proof of **vc.1**.

vc.2:

$$\frac{}{\{1\} \quad \forall j: \\ (\forall (\text{value}: \mathbb{Z}, x: [\mathbb{N} \rightarrow \mathbb{Z}], y: [\mathbb{N} \rightarrow \mathbb{Z}]): \\ (0 < j) \supset \text{rep}(j, 0, x, y) = \text{DP_RESULT}(0, j, x, y)) \\ \supset \\ (\forall (\text{value}: \mathbb{Z}, x: [\mathbb{N} \rightarrow \mathbb{Z}], y: [\mathbb{N} \rightarrow \mathbb{Z}]): \\ (0 < j+1) \supset \\ \text{rep}(j+1, 0, x, y) = \text{DP_RESULT}(0, j+1, x, y))}$$

Skolemizing,

vc.2:

$$\frac{}{\{1\} \quad (\forall (\text{value}: \mathbb{Z}, x: [\mathbb{N} \rightarrow \mathbb{Z}], y: [\mathbb{N} \rightarrow \mathbb{Z}]): \\ (0 < j') \supset \text{rep}(j', 0, x, y) = \text{DP_RESULT}(0, j', x, y)) \\ \supset \\ (\forall (\text{value}: \mathbb{Z}, x: [\mathbb{N} \rightarrow \mathbb{Z}], y: [\mathbb{N} \rightarrow \mathbb{Z}]): \\ (0 < j'+1) \supset \\ \text{rep}(j'+1, 0, x, y) = \text{DP_RESULT}(0, j'+1, x, y))}$$

Applying disjunctive simplification to flatten sequent,

vc.2:

$$\frac{\frac{}{\{-1\} \quad \forall (\text{value}: \mathbb{Z}, x: [\mathbb{N} \rightarrow \mathbb{Z}], y: [\mathbb{N} \rightarrow \mathbb{Z}]): \\ (0 < j') \supset \text{rep}(j', 0, x, y) = \text{DP_RESULT}(0, j', x, y)}}{\{1\} \quad \forall (\text{value}: \mathbb{Z}, x: [\mathbb{N} \rightarrow \mathbb{Z}], y: [\mathbb{N} \rightarrow \mathbb{Z}]): \\ (0 < j'+1) \supset \\ \text{rep}(j'+1, 0, x, y) = \text{DP_RESULT}(0, j'+1, x, y)}}$$

Skolemizing,

vc.2:

$$\frac{\begin{array}{l} \{-1\} \quad \forall (\text{value: } \mathbb{Z}, x: [\mathbb{N} \rightarrow \mathbb{Z}], y: [\mathbb{N} \rightarrow \mathbb{Z}]): \\ \quad (0 < j') \supset \text{rep}(j', 0, x, y) = \text{DP_RESULT}(0, j', x, y) \end{array}}{\begin{array}{l} \{1\} \quad (0 < j' + 1) \supset \\ \quad \text{rep}(j' + 1, 0, x', y') = \text{DP_RESULT}(0, j' + 1, x', y') \end{array}}$$

Applying disjunctive simplification to flatten sequent,

vc.2:

$$\frac{\begin{array}{l} \{-1\} \quad \forall (\text{value: } \mathbb{Z}, x: [\mathbb{N} \rightarrow \mathbb{Z}], y: [\mathbb{N} \rightarrow \mathbb{Z}]): \\ \quad (0 < j') \supset \text{rep}(j', 0, x, y) = \text{DP_RESULT}(0, j', x, y) \\ \{-2\} \quad (0 < j' + 1) \end{array}}{\begin{array}{l} \{1\} \quad \text{rep}(j' + 1, 0, x', y') = \text{DP_RESULT}(0, j' + 1, x', y') \end{array}}$$

Instantiating the top quantifier in -1 with the terms: $0, x', y'$,

vc.2:

$$\frac{\begin{array}{l} \{-1\} \quad (0 < j') \supset \text{rep}(j', 0, x', y') = \text{DP_RESULT}(0, j', x', y') \\ \{-2\} \quad (0 < j' + 1) \end{array}}{\begin{array}{l} \{1\} \quad \text{rep}(j' + 1, 0, x', y') = \text{DP_RESULT}(0, j' + 1, x', y') \end{array}}$$

Splitting conjunctions,

we get 2 subgoals:

vc.2.1:

$$\frac{\begin{array}{l} \{-1\} \quad \text{rep}(j', 0, x', y') = \text{DP_RESULT}(0, j', x', y') \\ \{-2\} \quad (0 < j' + 1) \end{array}}{\begin{array}{l} \{1\} \quad \text{rep}(j' + 1, 0, x', y') = \text{DP_RESULT}(0, j' + 1, x', y') \end{array}}$$

Expanding the definition of rep,

vc.2.1:

$$\frac{\begin{array}{l} \{-1\} \quad \text{rep}(j', 0, x', y') = \text{DP_RESULT}(0, j', x', y') \\ \{-2\} \quad (0 < j' + 1) \end{array}}{\begin{array}{l} \{1\} \quad \text{rep}(j', 0, x', y') + x'(j') \times y'(j') = \\ \quad \text{DP_RESULT}(0, 1 + j', x', y') \end{array}}$$

Expanding the definition of DP_RESULT,

vc.2.1:

$$\frac{\begin{array}{l} \{-1\} \quad \text{rep}(j', 0, x', y') = \text{DP_RESULT}(0, j', x', y') \\ \{-2\} \quad (0 < j' + 1) \end{array}}{\begin{array}{l} \{1\} \quad \text{rep}(j', 0, x', y') = \text{DP_RESULT}(0, j', x', y') \end{array}}$$

which is trivially true.

This completes the proof of **vc.2.1**.

vc.2.2:

{-1}	$(0 < j' + 1)$
{1}	$(0 < j')$
{2}	$\text{rep}(j' + 1, 0, x', y') = \text{DP_RESULT}(0, j' + 1, x', y')$

Simplifying, rewriting, and recording with decision procedures,

vc.2.2:

{-1}	$(0 < 1 + j')$
{1}	$(0 < j')$
{2}	$\text{rep}(1 + j', 0, x', y') = \text{DP_RESULT}(0, 1 + j', x', y')$

Simplifying with decision procedures,

vc.2.2:

{-1}	$(0 < 1 + j')$
{1}	$(0 < j')$
{2}	$\text{rep}(1 + j', 0, x', y') = \text{DP_RESULT}(0, 1 + j', x', y')$

Case splitting on $0 \leq j' \leq 1$,

we get 2 subgoals:

vc.2.2.1:

{-1}	$0 < j'$
{-2}	$(0 < 1 + j')$
{1}	$(0 < j')$
{2}	$\text{rep}(1 + j', 0, x', y') = \text{DP_RESULT}(0, 1 + j', x', y')$

which is trivially true.

This completes the proof of vc.2.2.1.

vc.2.2.2:

{-1}	$(0 < 1 + j')$
{1}	$0 < j'$
{2}	$(0 < j')$
{3}	$\text{rep}(1 + j', 0, x', y') = \text{DP_RESULT}(0, 1 + j', x', y')$

Case splitting on $j' \leq 0$,

we get 2 subgoals:

vc.2.2.2.1:

{-1}	$j' < 0$
{-2}	$(0 < 1 + j')$
{1}	$0 < j'$
{2}	$(0 < j')$
{3}	$\text{rep}(1 + j', 0, x', y') = \text{DP_RESULT}(0, 1 + j', x', y')$

Simplifying, rewriting, and recording with decision procedures,

This completes the proof of vc.2.2.2.1.

vc.2.2.2.2:

{-1}	$(0 < 1 + j')$
{1}	$j' < 0$
{2}	$0 < j'$
{3}	$(0 < j')$
{4}	$\text{rep}(1 + j', 0, x', y') = \text{DP_RESULT}(0, 1 + j', x', y')$

Case splitting on NOT $j!1 = 0$,

we get 2 subgoals:

vc.2.2.2.2.1:

{-1}	$(0 < 1 + j')$
{1}	$j' = 0$
{2}	$j' < 0$
{3}	$0 < j'$
{4}	$(0 < j')$
{5}	$\text{rep}(1 + j', 0, x', y') = \text{DP_RESULT}(0, 1 + j', x', y')$

Simplifying, rewriting, and recording with decision procedures,

This completes the proof of vc.2.2.2.2.1.

vc.2.2.2.2.2:

{-1}	$j' = 0$
{-2}	$(0 < 1 + j')$
{1}	$j' < 0$
{2}	$0 < j'$
{3}	$(0 < j')$
{4}	$\text{rep}(1 + j', 0, x', y') = \text{DP_RESULT}(0, 1 + j', x', y')$

Expanding the definition of rep,

vc.2.2.2.2.2.2:

{-1}	$j' = 0$
{-2}	$(0 < 1 + j')$
{1}	$j' < 0$
{2}	$0 < j'$
{3}	$(0 < j')$
{4}	$\text{rep}(j', 0, x', y') + x'(j') \times y'(j') =$ $\text{DP_RESULT}(0, 1 + j', x', y')$

Expanding the definition of DP_RESULT,

vc.2.2.2.2.2:

{-1}	$j' = 0$
{-2}	$(0 < 1 + j')$
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{1}	$j' < 0$
{2}	$0 < j'$
{3}	$(0 < j')$
{4}	$\text{rep}(j', 0, x', y') = \text{DP_RESULT}(0, j', x', y')$

Expanding the definition of rep,

vc.2.2.2.2.2:

{-1}	$j' = 0$
{-2}	$(0 < 1 + j')$
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{1}	$j' < 0$
{2}	$0 < j'$
{3}	$(0 < j')$
{4}	$0 = \text{DP_RESULT}(0, j', x', y')$

Expanding the definition of DP_RESULT,

vc.2.2.2.2.2:

{-1}	$j' = 0$
{-2}	$(0 < 1 + j')$
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{1}	$j' < 0$
{2}	$0 < j'$
{3}	$(0 < j')$
{4}	TRUE

which is trivially true.

This completes the proof of vc.2.2.2.2.2.

Q.E.D.