Verifying filesystems in ACL2 Project report, CS380L Fall 2017

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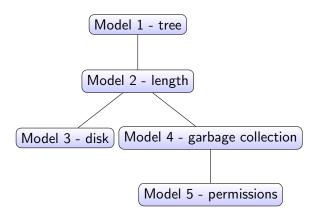
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Why we need a verified filesystem

- ▶ Modern filesystems have become increasingly complex, and so have the tools to analyse and recover data from them.
- ► It would be worthwhile to specify and formally verify, in the ACL2 theorem prover, the guarantees claimed by filesystems and tools.
- ▶ Work on this project started last year since then I've built 5 increasingly complex models.
- This talk will be focussed on the 4th and 5th models.

File system models

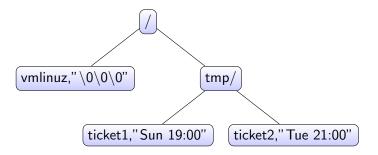


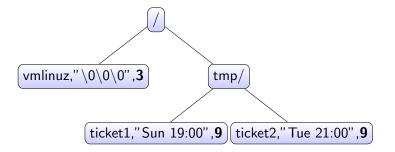
Minimal set of operations?

- ► The Google filesystem suggests a minimal set of operations:
 - create
 - ▶ delete
 - open
 - ▶ close
 - read
 - write
- Of these, open and close require the maintenance of file descriptor state - so they can wait.
- ► However, they are essential when describing concurrency and multiprogramming behaviour.
- Thus, we can start modelling a filesystem, and several refinements thereof.

Quick overview of models

- What follows is a sequence of increasingly refined models.
- ▶ Model 1: Tree representation of directory structure with unbounded file size and unbounded filesystem size.
- Model 2: Model 1 with file length as metadata.
- Model 3: Tree representation of directory structure with file contents stored in a "disk" (unbounded in length).
- ► Model 4: Model 3 with bounded filesystem size and garbage collection.
- Model 5: Model 4 with permissions for read/write for the user and others (no groups as yet)





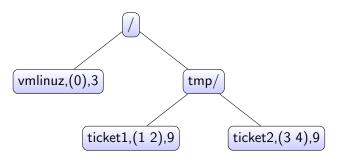


Table: Disk

0	\0\0\0
1	Sun 19:0
2	0
3	Tue 21:0
4	0

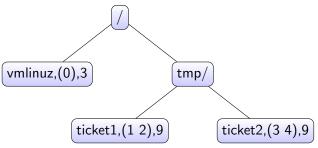


Table: Disk and allocation vector

0	\0\0\0	true
1	Sun 19:0	true
2	0	true
3	Tue 21:0	true
4	0	true
5		false

Proof approaches and techniques

- ► There are many properties that could be considered for correctness, but we choose to focus on the read-over-write theorems from the first-order theory of arrays.
- Read n characters starting at position start in the file at path hns in filesystem fs:
 - 11-rdchs(hns, fs, start, n)
- Write string text characters starting at position start in the file at path hns in filesystem fs:
 - 11-wrchs(hns, fs, start, text)

Proof approaches and techniques

► The first read-over-write theorem defines the semantics of reading from a location after writing to the same location. Formally, assuming n = length(text) and suitable "type" hypotheses (omitted here):

► The second read-over-write theorem defines the semantics of reading from a location after writing to a different location. Formally, assuming hns1 != hns2 and suitable "type" hypotheses (omitted here):

Proof approaches and techniques

- ► For each of the models 1 through 5, we have proofs of correctness of the two read-after-write properties, making use of the proofs of equivalence between models and their successors.
- ► For model 5, the proof assumes the permissions are granted.
- ► The proof of the converse property that reads and writes fail when permissions are not granted remains to be done.

Proof example: first read-over-write in model 2

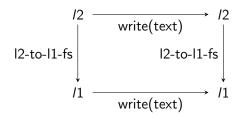


Figure: I2-wrchs-correctness-1

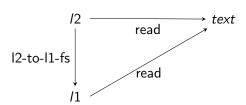


Figure: I2-rdchs-correctness-1 $\rightarrow \sim \sim 13/15$

Proof example: first read-over-write in model 2

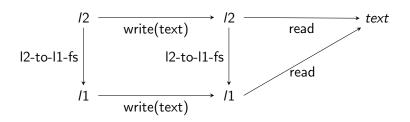


Figure: I2-read-over-write-1

Source analysis

Table: Code written for this project

Source lines of code (ACL2)	
defun events (function definitions)	
defthm events (lemmas and proofs)	