

Tools and Algorithms for Deciding Timed Relations

B Tech Project, 2012-2013

Mihir Mehta

Department of Computer Science and Engineering
Indian Institute of Technology, Delhi.

`cs1090197@cse.iitd.ac.in`

December 2012

Overview

- ▶ Bisimilarity and related notions
- ▶ Kanellakis and Smolka's algorithm
- ▶ Fernandez' algorithm
- ▶ Paige and Tarjan's algorithm
- ▶ Timed Automata
- ▶ Equivalences on Timed Automata
- ▶ Code written so far

Bisimilarity and related notions

- ▶ Labeled Transition System: This is a triple $(Proc, Act, \{\xrightarrow{a} \mid a \in Act\})$ where
 - ▶ $Proc$ is a set of states (also called processes or configurations.)
 - ▶ Act is a set of actions (also called labels.)
 - ▶ $\xrightarrow{a} \subseteq Proc \times Proc$ is a transition relation.
- ▶ CCS expression: Defined by the following grammar:
 - ▶ $P ::= K$
 - ▶ $P ::= \alpha.P$
 - ▶ $P ::=_{i \in I} P_i$
 - ▶ $P ::= P|Q$
 - ▶ $P ::= P[f]$
 - ▶ $P ::= P\mathbf{L}$

Bisimilarity and related notions

- ▶ Equivalence for CCS processes.
- ▶ Trace equivalence: $Traces(P) = Traces(Q)$
- ▶ Unsatisfactory (differences in deadlock behaviour.)
- ▶ Strong bisimulation: A binary relation R is a *strong bisimulation* if and only if, for all $(s_1, s_2) \in R$ and $a \in Act$.
$$\forall s'_1 (s_1 \xrightarrow{a} s'_1 \Rightarrow \exists s'_2. (s_2 \xrightarrow{a} s'_2 \wedge (s'_1, s'_2) \in R)) \wedge$$
$$\forall s'_2 (s_2 \xrightarrow{a} s'_2 \Rightarrow \exists s'_1. (s_1 \xrightarrow{a} s'_1 \wedge (s'_1, s'_2) \in R))$$
- ▶ It can be shown that the union of all strong bisimulations over the set of states is a strong bisimulation. This binary relation is called *strong bisimilarity*, denoted by \sim .

Bisimilarity and related notions

- ▶ Better notion of equivalence than trace equivalence: picks up differences in the deadlock behaviour of processes under study.
- ▶ Failing: does not account for the invisible nature of τ transitions in CCS processes.
- ▶ Weak bisimulation: A binary relation R is a *weak bisimulation* if and only if, for all $(s_1, s_2) \in R$ and $a \in \text{Act}$.
$$\forall s'_1 (s_1 \xrightarrow{a} s'_1 \Rightarrow \exists s'_2. (s_2 \xRightarrow{a} s'_2 \wedge (s'_1, s'_2) \in R)) \wedge$$
$$\forall s'_2 (s_2 \xrightarrow{a} s'_2 \Rightarrow \exists s'_1. (s_1 \xRightarrow{a} s'_1 \wedge (s'_1, s'_2) \in R))$$
- ▶ It can be shown that the union of all weak bisimulations over the set of states is a weak bisimulation. This binary relation is called *weak bisimilarity*, denoted by \approx .
- ▶ Better suited to CCS processes, as it ignores τ transitions, thus disregarding hidden behaviour within a process.

Kanellakis and Smolka's algorithm

- ▶ This is a naive algorithm for determining the bisimilarity relation for the set of processes in a labelled transition system.
- ▶ This relies on the notion of a splitter.
- ▶ Let $\pi = \{B_0, \dots, B_k\}$, $k \geq 0$ be a partition of the set of states Pr in a labeled transition system.
- ▶ A splitter for a block $B_i \in \pi$ is a block $B_j \in \pi$ such that for some action $a \in Act$, some states in B_i have a -labelled transitions whose targets lie in B_j while other states in B_i do not.
- ▶ This suggests a refinement of π : replace block B_i with
$$B_i^1 = B_i \cap T_a^{-1}[B_j]$$
$$B_i^2 = B_i - B_i^1$$
- ▶ Refinements of this kind constitute the steps of this algorithm.

Kanellakis and Smolka's algorithm

Initialise π to Pr ;

while *there exist splitters among the elements of π* **do**

 Pick a splitter B ;

foreach $B_j \in \pi$ **do**

foreach $a \in Act$ **do**

 Split B_j with respect to B for action a ;

end

end

end

Return π ;

- ▶ The time complexity of this algorithm is $O(mn)$, since there can be at most n iterations, and all m edges are scanned in each iteration.

Fernandez' algorithm

- ▶ More efficient algorithm ($O(m \log n)$).
- ▶ Relies on Paige and Tarjan's technique of three-way splitting.
- ▶ Splitters can now be 'simple' or 'compound'.
- ▶ Stability: A partition π is said to be stable with respect to a compound block S if S is not a splitter for any block in π for any action.
- ▶ For a compound block S , having a constituent simple block B satisfying $n(B) \leq 0.5 * n(S)$, and with respect to which π is stable, we can split a block B_i on an action a as follows:

$$B_i^1 = (B_i \cap T_a^{-1}[B]) - T_a^{-1}[S - B]$$

$$B_i^2 = (B_i \cap T_a^{-1}[S - B]) - T_a^{-1}[B]$$

$$B_i^3 = B_i \cap T_a^{-1}[B] \cap T_a^{-1}[S - B]$$

Fernandez' algorithm

Initialise $\pi = \{Pr\}$;

Initialise $W = \{Pr\}$ **while** W is not empty **do**

 Choose a splitter B from W , removing it;

if B is a simple splitter **then**

 Perform a two-way split on each action with respect to B
 and update W ;

else

 Perform a three-way split on each action with respect to B
 and update W ;

end

end

Paige and Tarjan's Algorithm

- ▶ Technique of three way splitting came from here.
- ▶ Special case of Fernandez' algorithm when there's only one kind of action.

Timed Automata

- ▶ Formally, a timed automaton over a finite set of clocks C and a finite set of actions Act is a 4-tuple (L, l_0, E, I) .
- ▶ L is a finite set of locations.
- ▶ l_0 is the initial location.
- ▶ $E \subseteq L \times B(C) \times Act \times 2^C \times L$ is a finite set of edges.
- ▶ $I : L \rightarrow B(C)$ assigns invariants to each edge location.
- ▶ $B(C)$ is the set of clock constraints over C . An element of $B(C)$ can be an equality, a slack inequality, a strict inequality, or an AND combination of such constraints.

Equivalences on Timed Automata

- ▶ *Timed bisimulation*: A binary relation R is a timed bisimulation if and only if, for all $(s_1, s_2) \in R$, $a \in \text{Act}$, $d \in \mathbb{R}_{\geq 0}$
 $\forall s'_1 (s_1 \xrightarrow{a} s'_1 \Rightarrow \exists s'_2. (s_2 \xrightarrow{a} s'_2 \wedge (s'_1, s'_2) \in R)) \wedge$
 $\forall s'_2 (s_2 \xrightarrow{a} s'_2 \Rightarrow \exists s'_1. (s_1 \xrightarrow{a} s'_1 \wedge (s'_1, s'_2) \in R)) \wedge$
 $\forall s'_1 (s_1 \xrightarrow{d} s'_1 \Rightarrow \exists s'_2. (s_2 \xrightarrow{d} s'_2 \wedge (s'_1, s'_2) \in R)) \wedge$
 $\forall s'_2 (s_2 \xrightarrow{d} s'_2 \Rightarrow \exists s'_1. (s_1 \xrightarrow{d} s'_1 \wedge (s'_1, s'_2) \in R))$
- ▶ It can be shown that the union of all timed bisimulations over the set of (location, valuation) pairs is a timed bisimulation. This binary relation is called *timed bisimilarity*, denoted by \sim .

Equivalences on Timed Automata

- *Time abstracted bisimulation*: A binary relation R is a time abstracted bisimulation if and only if, for all $(s_1, s_2) \in R$,

$$a \in \text{Act}, d \in R_{\geq 0}$$

$$\forall s'_1 (s_1 \xrightarrow{a} s'_1 \Rightarrow \exists s'_2. (s_2 \xrightarrow{a} s'_2 \wedge (s'_1, s'_2) \in R)) \wedge$$

$$\forall s'_2 (s_2 \xrightarrow{a} s'_2 \Rightarrow \exists s'_1. (s_1 \xrightarrow{a} s'_1 \wedge (s'_1, s'_2) \in R)) \wedge$$

$$\forall s'_1 (s_1 \xrightarrow{d} s'_1 \Rightarrow \exists (s'_2, d'). (s_2 \xrightarrow{d} s'_2 \wedge (s'_1, s'_2) \in R)) \wedge$$

$$\forall s'_2 (s_2 \xrightarrow{d} s'_2 \Rightarrow \exists (s'_1, d'). (s_1 \xrightarrow{d} s'_1 \wedge (s'_1, s'_2) \in R))$$

- It can be shown that the union of all time abstracted bisimulations over the set of (location, valuation) pairs is a time abstracted bisimulation. This binary relation is called *time abstracted bisimilarity*, denoted by \sim_u .

Equivalences on Timed Automata

- ▶ *Equivalence of valuations:* Two valuations v and v' of a timed automaton are said to be equivalent ($v \equiv v'$) if and only if:
 - ▶ For each $x \in C$, either both $v(x)$ and $v'(x)$ are greater than c_x or $\lfloor v(x) \rfloor = \lfloor v'(x) \rfloor$
 - ▶ For each $x \in C$ such that $v(x) \leq c_x$, $\text{frac}(v(x)) = 0$ if and only if $\text{frac}(v'(x)) = 0$.
 - ▶ For all $x, y \in C$ such that $v(x) \leq c_x$ and $v(y) \leq c_y$, we have $\text{frac}(v(x)) \leq \text{frac}(v(y))$ if and only if $\text{frac}(v'(x)) \leq \text{frac}(v'(y))$.
- ▶ Under this notion of equivalence, an equivalence class is known as a *region*. The equivalence class containing a valuation v is denoted by $[v]_{\equiv}$.

Equivalences on Timed Automata

- *Region graph*: the region graph of a timed automaton A with clock set C and action set Act is an LTS

$$T_r(A) = (Proc, Act \cup \{\varepsilon\}, \{\xrightarrow{a} \mid a \in Act \cup \{\varepsilon\}\})$$

where $Proc = \{(l, [v]_{\equiv}) \mid l \in L, v : C \rightarrow R_{\geq 0}\}$ (these states are called symbolic states)

The transitions are defined as follows:

- For each $a \in A$, $(l, [v]_{\equiv}) \xrightarrow{a} (l', [v']_{\equiv})$ iff $(l, v) \xrightarrow{a} (l', v')$
- $(l, [v]_{\equiv}) \xrightarrow{d} (l, [v']_{\equiv})$ if for some $d \in R_{\geq 0}$, $(l, v) \xrightarrow{d} (l, v')$

Equivalences on Timed Automata

- ▶ *Extended Clock Constraints*: These are described by the grammar

$$g ::= x \bowtie n \mid x - y \bowtie n \mid g_1 \wedge g_2$$

- ▶ The set of extended clock constraints is denoted by $B^+(C)$
- ▶ *Zones*: A zone is a set of valuations represented by an extended clock constraint.

$$Z = \{v \mid v \in g_Z\}$$

- ▶ *Symbolic state*: For a timed automaton A , a symbolic state is an ordered pair of a location and a zone, that is, (l, Z) .

Equivalences on Timed Automata

- ▶ *Reachability graph*: The reachability graph of a timed automaton is a directed graph in which the nodes are symbolic states and the edges are given by the *reachability relation*:
 - ▶ (l, Z') is reachable from (l, Z) if there exists a valuation belonging Z' which can be obtained from Z by applying a time delay which does not violate the invariant at l .
 - ▶ (l', Z') is reachable from (l, Z) if l has an action transition to l' and there exists a valuation in Z' which violates neither the guard of the transition nor the invariant at l' .

Code written so far

- ▶ *Fernandez' algorithm*: We expect to need this algorithm while building the software, so we built a module implementing it in Ocaml (primary language for this project.)
- ▶ *Grammar*: We implemented a lexer and parser in Ocaml to build a representation of a timed automaton from a specification in a text file.

References

- ▶ Reactive Systems: Modeling, Specification and Verification -
Luca Aceto, Anna Ingolfssdottir, Kim Larsen, Jiri Srba.