

Exploring Optimized Locking

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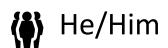




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Objectives



- Current problems
- Optimized Locking Components
- How does it work
- Limitations
- Q&A



Current Problems



ACID Compliance

Lock escalation

Long-term locking

Excessive blocking

Excessive memory utilization



Traditional Locking



Update lock taken on rows to allow for predicate evaluation

 If predicate is satisfied, then an exclusive lock is taken on the row

Locks are held until the end of the transaction

Traditional Locking



```
ALTER DATABASE [Locking] SET READ_COMMITTED_SNAPSHOT OFF;
SET TRANSACTION ISOLATION LEVEL READ COMMITTED;

CREATE TABLE t1 (a int, b int);
INSERT INTO t1 VALUES (1,10), (2,20), (3,30);

-- TID2: Increase b by 10
BEGIN TRAN UPDATE t1 SET b=b+10;

Table Update

Compute Scalar

Table Scan
```

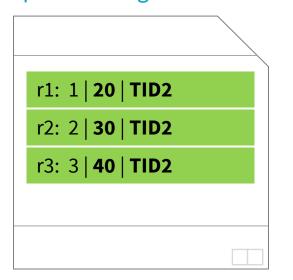
Cost: 0 %

p1: Data Page for t1

Cost: 0 %

[t1]

Cost: 75 %



Lock Manager

[t1]

Cost: 25 %

Lock Mode	Lock Type	Lock Resource
IX	OBJECT	t1
IX	PAGE	p1
Х	RID	r1
Х	RID	r2
Х	RID	r3



Old School

Updating 1 million rows might require 1 million exclusive (X) row locks held until the end of the transaction.

Optimized Locking Components



Accelerated Database Recovery

Transaction ID (TID)

Lock After Qualification (LAQ)

Transaction ID (TID)



A unique identifier

Each row will contain the TID that last modified it

Lock will be held on the transaction ID versus row key

Lock After Qualification (LAQ)



 Predicate is applied to the row using the latest version of the row

If the predicate is not satisfied, move to the next row

• If the predicate is satisfied, an exclusive (X) lock is placed on the row

Can retry predicate evaluation as needed due to previous exclusive lock

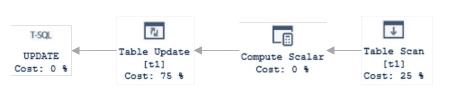
Traditional Blocking



```
ALTER DATABASE [db1] SET READ_COMMITTED_SNAPSHOT ON; SET TRANSACTION ISOLATION LEVEL READ COMMITTED;
```

```
CREATE TABLE t1 (a int, b int);
INSERT INTO t1 VALUES (1,10), (2,20), (3,30);
```

```
-- TID2 [SESSION 1]: Increase b by 10 where a=1 BEGIN TRAN UPDATE t1 SET b=b+10 where a=1;
```

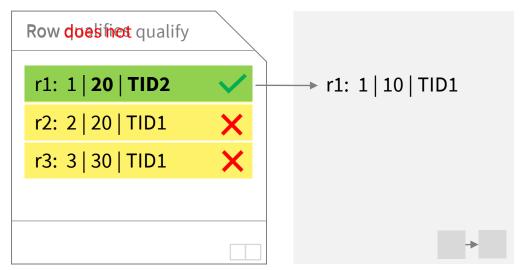


-- TID3 [SESSION 2]: Increase b by 10 where a=2 BEGIN TRAN UPDATE t1 SET b=b+10 where a=2;



p1: Data Page for t1





Slide Credit: Perry Skountrianos/Prashanth Purnananda & the Microsoft team who brought you optimized locking

Lock Manager

Lock Mode	Lock Type	Resource	Owner	Status
IX	OBJECT	t1	TID2, <mark>TID3</mark>	GRANT
IX	PAGE	p1	TID2	GRANT
Χ	RID	r1	TID2	GRANT
IU	PAGE	p1	TID3	GRANT
U	RID	r1	TID3	WAIT

Session 2 is blocked waiting for Session 1 to commit

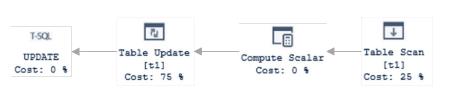
Optimized Locking



```
ALTER DATABASE [db1] SET READ_COMMITTED_SNAPSHOT ON; SET TRANSACTION ISOLATION LEVEL READ COMMITTED;
```

```
CREATE TABLE t1 (a int, b int);
INSERT INTO t1 VALUES (1,10), (2,20), (3,30);
```

```
-- TID2 [SESSION 1]: Increase b by 10 where a=1
BEGIN TRAN UPDATE t1 SET b=b+10 where a=1;
```

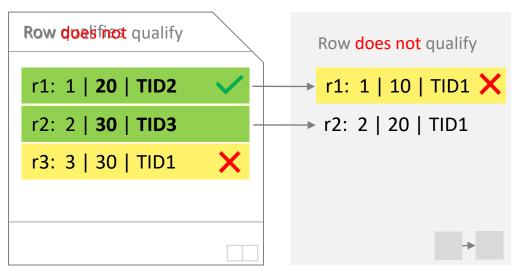


-- TID3 [SESSION 2]: Increase b by 10 where a=2 BEGIN TRAN UPDATE t1 SET b=b+10 where a=2;



p1: Data Page for t1

Row version store



Slide Credit: Perry Skountrianos/Prashanth Purnananda & the Microsoft team who brought you optimized locking

Lock Manager

	Lock Mode	Lock Type	Resource	Owner	Status
\Rightarrow	X	XACT	TID2	TID2	GRANT
\Rightarrow	IX	OBJECT	t1	TID2, <mark>TID3</mark>	GRANT
	IX	PAGE	p1	TID2	GRANT
	IX	PAGE	p1	TID3	GRANT
	Х	RID	r2	TID3	GRANT

Session 2 is not blocked by Session 1



New School Updating 1 million rows might require 1 million X row locks but each lock is released as soon as each row is updated,

Only one TID lock will be held until the end of the transaction.

Warning





Even without LAQ, applications should not assume that SQL Server (under versioning isolation levels) will guarantee strict ordering, without using locking hints.

https://learn.microsoft.com/en-us/sql/relational-databases/performance/optimized-locking?view=azuresqldb-current

Best Practices & Troubleshooting



- Avoid locking hints
- Make sure RCSI is enabled
- New entries for Deadlock Graphs
- New Waits introduced:
 - LCK M S XACT READ
 - LCK M S XACT MODIFY
 - LCK_M_S_XACT

LCK_M_S_XACT	Occurs when a task is waiting for a shared lock on an XACT wait_resource type, where the intent cannot be inferred. Rare. Related to optimized locking	
LCK_M_S_XACT_READ	Occurs when a task is waiting for a shared lock on an XACT wait_resource type, with an intent to read. Related to optimized locking.	
LCK_M_S_XACT_MODIFY	Occurs when a task is waiting for a shared lock on an XACT wait_resource type, with an intent to modify. Related to optimized locking.	

Limitations



- Azure SQL Database Only (for now)
 - DTU service tier
 - vCore Provisioned tier (including serverless)

Accelerated Database Recovery Required

 Repeatable Read & Serializable Isolation forces the lock to be held on the row or page until the end of the transaction

Summary



Better concurrency

Significantly reduced locking and lock memory

- On by default in Azure SQL Database
 - ADR/RCSI is enabled by default as well

 Every SQL Server DBA should be watching this feature arrive to the box product

Resources



Optimized Locking - https://learn.microsoft.com/en-us/sql/relational-databases/sql-server-transaction-locking-and-row-versioning-guide?view=sql-server-ver16&source=recommendations

Article - https://www.red-gate.com/simple-talk/databases/sql-server/optimized-locking-in-azure-sql-database/ (Simple Talk - Aaron Bertrand)



Questions? Answers!

Got Questions?

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Slides & Demos



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