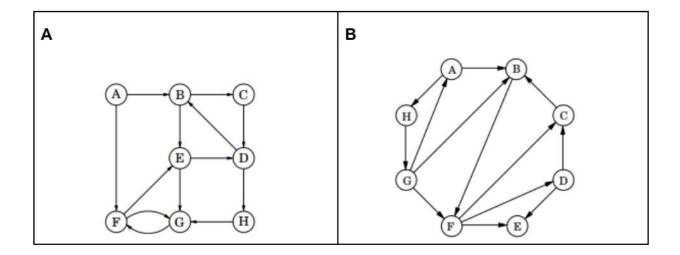
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## **Design and Analysis of Algorithms**

## **CS302**

**Total Marks = 60, each question carries 10 marks.** 

**Q1**) Run the DFS algorithm on the following graph. Draw the DFS tree (or forest). Label each node with pre and post numbers. Indicate the backward, forward and cross edges (See slides on DFS to learn about edge classifications). Wherever multiple options are available, visit the nodes in the alphabetic order.



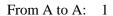
**Q2**) You are given a DAG G=(V, E) of the courses in your curriculum (nodes), with the edges between them showing the pre-requisite relationship, i.e. an edge going from A to B means that A must be studied before B: so you take the course A in one semester and then B must come in the following semester. Imagine that other than the prerequisite relationship, there is no bound on the number of courses you can take in a semester. Design an O(|V|+|E|) algorithm to find the minimum number of semesters needed to graduate.

**Hint:** notice that this problem asks you to find the longest path in the DAG, since there will be at least that many semesters needed to graduate.

**Q3**) The reverse of a directed graph G = (V, E) is another directed graph  $G^R = (V, E^R)$  on the same vertex set, but with all edges reversed; that is,  $E^R = \{(v, u): (u, v) \in E\}$ . Give a linear-time algorithm for computing the reverse of a graph in the adjacency list format. Give a clearly specified pseudo-code.

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**Q4)** Recall that BFS can be used to solve the problem of single source shortest path for unweighted graph G(V, E). Also BFS only finds one shortest paths between s and x, where s is the source vertex and x is any other vertex of the graph. But shortest path may not be unique. Suppose we want to count the number of distinct shortest paths between s and x for each x, where  $x \in V$ . For example in the graph given below if A is the source vertex then the number of distinct shortest paths for each pair are:



From A to B: 1

From A to C: 1

From A to D: 1

From A to E: 1

From A to F: 2

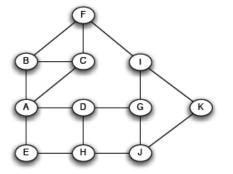
From A to G: 1

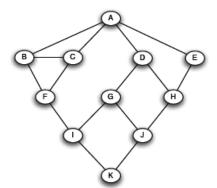
From A to H: 2

From A to I: 3

From A to J: 3

From A to K: 6





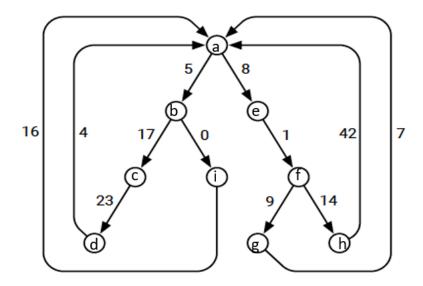
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You are required to design a linear time algorithm that counts the number of distinct shortest paths between a given source vertex s and all the other vertices in the graph. You can assume that graph is represented using adjacency list representation.

Hint: you can use the tree structure of BFS to count the distinct shortest paths.

Give detailed explanation of your algorithm

**Q5**) The following picture shows an example of a Korchoff Graph. It is a weighted, directed graph. As you can see the graph has a binary tree component at is center. Moreover, there are edges going from each leaf of the binary tree backto the root. All edges are weighted and the weights are non-negative. We wish to compute single-source shortest paths in a Korchoff graph.



For example, the shortest path costs from vertex 'b' to all other vertices are as follows:

 $b \rightarrow a$ : 16, (b, i, a)

 $b \rightarrow c: 17 (b,c)$ 

 $b \rightarrow i: 0 (b,i)$ 

 $b \to d: 40 (b,c,d)$ 

 $b \rightarrow e: 24, (b, i, a, e)$ 

 $b \rightarrow f: 25, (b, i, a, e, f)$ 

 $b \rightarrow g: 34, (b, i, a, e, f, g)$ 

 $b \rightarrow h: 39, (b, i, a, e, f, h)$ 

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- a) Is the following statement true or false? Justify your answer: Dijkstra's shortest path algorithm takes O(|V|lg|V|) time on a Korchoff graph.
- b) Design a Single Source shortest path algorithm for the Korchoff graph which works in linear time. The call to the algorithm looks like: spKorchoff(G=(V, E), s, r), where G is a Korchoff graph, s is the source node used to measure the shortest paths, and r is the root node of the binary tree component of the graph. First describe the algorithm in few lines (or steps) in English, then write the pseudocode.
- **Q6**) Agent Bob is living in his houseat A in the enemy territory. There are n other agents in the area, stationed at hotels  $h_1, h_2, ..., h_n$ . They all must visit Bob's house at A for a top secret meeting. However, the area is constantly monitored and there is a risk of being monitored associated with every road. You have a map, G=(V, E), of the n hotels  $h_1, h_2, ..., h_n$  and A. This map shows every road between any two hotels or between a hotel and A and also shows its associated risk value, which is a positive weight. The total risk of a path taken by an agent is the sum of the risk values of all the edges in that path. You need to write an algorithm that produces paths for all agents, such that the total risk for each agent is minimized. Your algorithm should work in O((|V|+|E|)lg|V|).