

Name: Aisha Muhammael Nawaz

Roll# 202-0921

Section :- BSCS 4A1

DAA Dr. Maryam Bashir

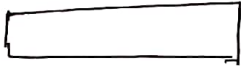
Homework#5 Rod cutting Problem

Total pages = 3

TASK#1

Price

17

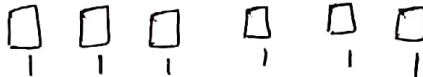


Length

6

ca).

6



cb)

14



cc)

16



3

3

cd)

14

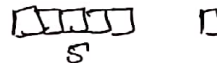


4

2

ce)

11



cf)

TASK#2

Dry Run Dynamic Programming Solution

starts here:• $i=1$
 $\rightarrow K=1 \quad v_K + c[i-K] = 1$
 $\text{max} = 1$
• $c[1] = 1$ • $i=2$
 $\rightarrow K=1 \quad v_K + c[i-K] = 2$
 $\text{max} = 2$
 $\rightarrow K=2 \quad v_K + c[i-K] = 5$
 $\text{max} = 5$
• $c[2] = 5$ • $i=3$
 $\rightarrow K=1 \quad v_K + c[i-K] = 6$
 $\text{max} = 6$

continued

 $\rightarrow K=2 \quad v_K + c[i-K] = 6$
 $\text{max} = 6$
 $\rightarrow K=3 \quad v_K + c[i-K] = 8$
 $\text{max} = 8$
• $c[3] = 8$ • $i=4$
 $\rightarrow K=1 \quad v_K + c[i-K] = 9$
 $\text{max} = 9$
 $\rightarrow K=2 \quad v_K + c[i-K] = 10$
 $\text{max} = 10$
 $\rightarrow K=3 \quad v_K + c[i-K] = 9$
 $\text{max} = 10$
 $\rightarrow K=4 \quad v_K + c[i-K] = 9$
 $\text{max} = 10$
• $c[4] = 10$

- $i = 5$
 $\rightarrow K=1$
 $v_K + c \sum_{i=K}^i = 11$
 $\max = 11$

- $\rightarrow K=2$
 $v_K + c \sum_{i=K}^i = 13$
 $\max = 13$

- $\rightarrow K=3$
 $v_K + c \sum_{i=K}^i = 13$
 $\max = 13$

- $\rightarrow K=4$
 $v_K + c \sum_{i=K}^i = 10$
 $\max = 13$

- $\rightarrow K=5$
 $v_K + c \sum_{i=K}^i = 10$
 $\max = 13$

- $c \sum 5 = 13$

- $i = 6$
 $\rightarrow K=1$
 $v_K + c \sum_{i=K}^i = 14$
 $\max = 14$

- $\rightarrow K=2$
 $v_K + c \sum_{i=K}^i = 15$
 $\max = 15$

- $\rightarrow K=3$
 $v_K + c \sum_{i=K}^i = 16$
 $\max = 16$

- $\rightarrow K=4$
 $v_K + c \sum_{i=K}^i = 14$
 $\max = 16$

- $\rightarrow K=5$
 $v_K + c \sum_{i=K}^i = 11$
 $\max = 16$

- $\rightarrow K=6$
 $v_K + c \sum_{i=K}^i = 17$
 $\max = 17$

- $c \sum 6 = 17$

\hookrightarrow final Answer

SK # 3

Q. When should we cut the rod to get optimal revenue?

Rod-Cutting - DP (V, n)

Σ int $j = 0, j_2 = 0$; // edited
 int arr[n];
 arr[0] = 0;

Σ for $i = 1$ to n
 max = $-\infty$

for $k = 1$ to i

if (max < $V_k + C[i-k]$)

Σ max = $V_k + C[i-k]$

$j = k$; // edited

$j_2 = \text{arr}[i-k]$; // edited

$C[i] = \text{max};$
 arr[i] = j ;

3

// edited
 cout << " For optimal revenue cut from: " << j << " and " << j_2 << endl;

3