

# VSComp Verification

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# Introduction

- Yearly competition with different variants
  - VSComp, VSTTE
- Aim is to write code for a set of challenge problems then verify properties in limited time
- Anyone can participate
- Any verification tools can be used
  - Dafny, VCC, etc.

# Chosen Tools and Problems

- Used Dafny for coding and verifying
- Chose four problems:
  1. *Sum-Max*
  2. *Inversion*
  3. *Two-Duplets*
  4. *TwoWaySort*
  - Will give quick overview of first 3 problems
  - Will discuss 4<sup>th</sup> one in some detail

# 1. Sum-Max

- Problem:
  - Given N-element array A, computes sum of all elements, and the maximum element
- Verify:
  - Given  **$N \geq 0$**  and  **$A[i] \geq 0$**  *for*  $0 \leq i < N$ 
    - prove the post-condition that  **$\text{sum} \leq N * \text{Max}$**

## 2. Inversion

- Problem: Invert an injective array  $A$  on  $N$  elements in the subrange from 0 to  $N - 1$ .
- Verify:
  - Verify output array  $B$  is also injective:
    - $B[A[i]] = i$  ( $0 \leq i < N$ )
  - Example:

➤  $A =$ 

2	0	3	1
---	---	---	---

➤  $B =$ 

0	1	2	3
1	3	0	2

# 3. Two-Duplets

- Problem:
  - Find two duplets in given array
  - Assume at least two unique duplets exist in array
  - Assume input array  $a$  of length  $N+2$  with  $N \geq 2$ .
  - May assume array values between 0 and  $N-1$ .
- Verify:
  - Verify post-condition:
    - Program finds 2 duplets from given array,
      - $\text{Duplet1} \neq \text{Duplet2}$

# 4. TwoWay Sort

- Problem:
  - Sort boolean array
    - using two pointers moving from beginning and end of array
      - swapping of values when **TRUE** seen before **FALSE**
- Verify:
  1. *Safety*: every array access within bounds
  2. *Termination*: Prove function always terminates
  3. *Behavior*: After execution, following hold:
    - (a) array is a permutation of initial contents
    - (b) array is sorted in increasing order.

```

method TwoWaySort(a:array<bool>)
  requires a != null; modifies a;
  ensures compare(a, old(a)); //ensuring permutation
  ensures forall k:: forall l:: 0 <= k < l < a.Length
    ==> !a[k] || a[l]; //sortedness

{
  i := 0
  j := a.Length - 1
  while (i <= j)
    invariant 0 <= i <= (j+1);
    invariant (i-1) <= j < a.Length;
    decreases (j+1) - i;
    decreases j - (i-1);
    invariant forall k:: 0 <= k < i ==> !a[k];
    invariant forall l:: j < l < a.Length ==> a[l];
    {
      ... //move pointers i,j and swap when TRUE before FALSE
    }
}

```

1. Safety!



```

method TwoWaySort(a:array<bool>)
  requires a != null; modifies a;
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    invariant forall k:: 0 <= k < i ==> !a[k];
    invariant forall l:: j < l < a.Length ==> a[l];
    {
      ... //move pointers i,j and swap values as needed
    }
}

```

## 2. Termination!

method TwoWaySort(a:array<bool>)

requires a != null; modifies a;

ensures compare(a, old(a)); //ensuring permutation

ensures forall k:: forall l::  $0 \leq k < l < a.Length$

$\implies !a[k] \vee a[l]$ ; //sortedness

{

i := 0

j := a.Length - 1

while (i <= j)

invariant  $0 \leq i \leq (j+1)$ ;

invariant  $(i-1) \leq j < a.Length$ ;

decreases  $(j+1) - i$ ;

decreases  $j - (i-1)$ ;

invariant forall k::  $0 \leq k < i \implies !a[k]$ ;

invariant forall l::  $j < l < a.Length \implies a[l]$ ;

{

... //move pointers i,j and swap values as needed

}

}

### 3. Behavior!

method TwoWaySort(a:array<bool>)

requires a != null; modifies a;

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decreases  $(j+1) - i$ ;

decreases  $j - (i-1)$ ;

invariant forall k::  $0 \leq k < i \implies !a[k]$ ;

invariant forall l::  $j < l < a.Length \implies a[l]$ ;

{

... //move pointers i,j and swap values as needed

}

}

### 3. Behavior!

TRUE – FALSE X

method TwoWaySort(a:array<bool>)

requires a != null; modifies a;

ensures compare(a, old(a)); //ensuring permutation

ensures forall k:: forall l::  $0 \leq k < l < a.Length$

$\implies !a[k] \parallel a[l]$ ; //sortedness

{

i := 0

j := a.Length - 1

while (i <= j)

invariant  $0 \leq i \leq (j+1)$ ;

invariant  $(i-1) \leq j < a.Length$ ;

decreases  $(j+1) - i$ ;

decreases  $j - (i-1)$ ;

invariant forall k::  $0 \leq k < i \implies !a[k]$ ;

invariant forall l::  $j < l < a.Length \implies a[l]$ ;

{

... //move pointers i,j and swap values as needed

}

}

### 3. Behavior!

**predicate compare (oldArray: array<bool>,  
newArray: array<bool>)**

*reads a; reads a1;*

*requires a != null; requires a1 != null;*

{

forall key:bool :: **total**(oldArray, oldArray.Length, key)  
==

**total**(newArray, newArray.Length, key)

}

```
function total(a: array<bool>, i:int, key: bool): int  
requires a != null; reads a;
```

```
{
```

```
    if i < 0
```

```
        then 0
```

```
    else if (a.Length == 0)
```

```
        then 0
```

```
    else if 0<=i<a.Length && a[i] == key
```

```
        then total(a, i-1, key) + 1
```

```
    else total(a, i-1, key)
```

```
}
```

**method TwoWaySort(a:array<bool>)**

requires a != null; modifies a;

ensures compare(a, old(a)); //ensuring permutation

ensures forall k:: forall l::  $0 \leq k < l < a.Length$

$\implies !a[k] \parallel a[l]$ ; //sortedness

{

i := 0

j := a.Length - 1

while (i <= j)

*invariant*  $0 \leq i \leq (j+1)$ ;

*invariant*  $(i-1) \leq j < a.Length$ ;

*decreases*  $(j+1) - i$ ;

*decreases*  $j - (i-1)$ ;

*invariant forall k::  $0 \leq k < i \implies !a[k]$ ;*

*invariant forall l::  $j < l < a.Length \implies a[l]$ ;*

{

... //move pointers i,j and swap values as needed

}

}

### 3. Behavior!

Thank You