

# Recursion

- Is method to solve complex problem by breaking them down into simpler ones. Recursion is process by which a function calls itself directly or indirectly by applying some subroutines on parameters by keeping an extra space overhead (stack)

Basic eg

factorial

$$f(n) = n * f(n-1)$$

eg  $5! = 5 \times 4!$

Subroutine on parameter

fn call itself

break problem in small parts.

$$f(3) = 3 \times f(2)$$
$$f(2) = 2 \times f(1)$$
$$f(1) = \underline{\underline{1! = 1}}$$

- This recursion is related to principal of mathematical Induction.

## Steps for recursion

- 1) Find out smallest subproblem for which we know the ans (base case)
- 2) Assume for given problem recursion will work correctly & calculate a subproblem.
- 3) Find how smaller subproblem contribute to final ans. (Self work)  
like  $f(n) = n \times f(n-1)$

## eg 1 Find factorial

```
def fact(n):  
    // base case  
    if (n == 1): return 1  
    // recursive assumption  
    sub = fact(n-1)  
    // self work  
    return n * sub
```

```
def fact(n):  
    if (n == 1): return 1  
  
    else :  
        return n * fact(n-1)
```

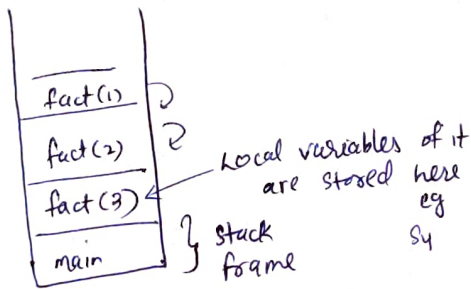
← can compress it

Q How memory is distributed for process?

Stack (Linear)

# call stack is extra space used in recursion.

call stack



Heap (big pool of data)

# is not heap data structure

# slower than stack

# size is quite large

eg `int * ptr = new int [10];`

# Whenever a function is called, its variables get memory allocated on stack. And whenever function call is over, its variables & memory is deallocated from stack.

A programmer does not have to worry about memory allocation & de-allocation of stack variables.

eg `int a`, `int b[10]`, `int c=10`, `int d[c]` are stored in stack.

# Whenever an object is created it's always stored in heap space & stack memory contains reference to it.

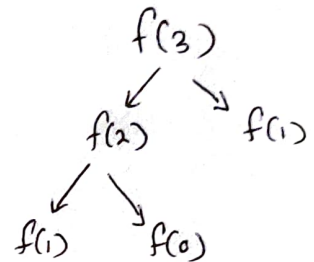
## Basic Questions

Q1. fibonacci series, return nth value

0, 1, 1, 2, 3, 5 - - - - -  
0th 1st 2nd term

```
Code def fib(n):  
    if (n == 0 or n == 1): // base case  
        return n  
    else:  
        return fib(n-1) + fib(n-2)
```

print(fib(6))



# note

relations like  $\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$

this are called recurrence relation

Q2 Print 1 to n natural number in normal & reverse order.

```
def sub(n):  
    if (n == 0): # base case  
        return  
    else:  
        sub(n-1) # recursive assumption  
        print(n) # self work
```

sub(5) → 1  
2  
3  
4  
5

reverse order

```
def sub(n):  
    if (n == 0):  
        return  
    else:  
        print(n)  
        sub(n-1)
```

sub(5) → 5  
4  
3  
2  
1

# make call stack to see  
how its working.

Q3) Count number of ways to make binary string of length  $n$  where there is no adjacent 1.

for such questions there is usually a pattern.

eg

$n=1$	$\rightarrow 0, 1$	(2)
$n=2$	$\rightarrow 00, 01, 10$	(3)
$n=3$	$\rightarrow 000, 001, 010, 100$ 101	(5)
$n=4$	$\rightarrow 0000, 0001, 0010, 0100$ 1000, 1001, 1010, 0101	(8)
$n=5$	$\rightarrow 00000, 00001, 00010, 00100, 01000$ 10000, 10100, 10101, 10010, 10001 01010, 01001, 00101	(13)

So pattern here is 2, 3, 5, 8, 13 - -  
is fibonacci series

soln

```
def fib(n):  
    if (n == 1):  
        return 2  
    if (n == 2):  
        return 3  
    else:  
        return fib(n-1) + fib(n-2)
```

Q4 There are  $N$  persons, who want to go to party. There is a constraint that any person can either go alone or can go in pair. Calculate no. of ways in which  $N$  persons will go to party.

Sol Like  $ABC \rightarrow A, B, C \mid AB, C \mid A, BC \mid AC, B$   
for  $N=3$  4 ways

here concept of making recursive relation.




$f(n)$  depends on 2 things

$\rightarrow$  Let if  $A$  goes alone then  $f(n) = 1 * f(n-1)$  ways

$\rightarrow$  if  $A$  decides to go in pair so  $f(n) = (\text{no. of ways } A \text{ can make pair}) * f(n-2)$   
 $\uparrow$   
since  $n-2$  persons are left

$$\text{So } f(n) = f(n-1) + (n-1) * f(n-2)$$

Base case

$n=1$   ways = 1  
 $n=2$   pair,  alone = 2

Code

```
def sub(n):
    if (n == 1 or n == 2):
        return n
    else:
        return sub(n-1) + (n-1) * sub(n-2)
```

print(sub(4))  $\Rightarrow$  10

Q5) Print pattern

N=4

```
* * * *
* * *
* *
*
```

# Concept (handling 2 variables)

Code

```
def sub(n, n):
    if n <= 0:          // base case
        return

    if n == n:
        print(" ")      // self work
        sub(n-1, 0)      // recursive call

    else:
        print("*", end=" ") // self work
        sub(n, n+1)        // recursive call

sub(4, 0)
```

Q6) Reverse a string

```
class Solution:
    def reverseString(self, s: List[str]) -> None:
        n = len(s)
        def sub(s, l, r):
            if (l >= r):
                return
            else:
                s[l], s[r] = s[r], s[l]
                sub(s, l+1, r-1)
            return s

        s = sub(s, 0, n-1)
```



# Intermediate

1) Concept of output so far.  
? take or not take cases.

Q1) Print all subsets of array.

eg ['a', 'b', 'c'] → [a, b, c], [a, b], [a, c], [a], [b, c], [b], [c], []

Code

```
def sub(a, i, n, s):
```

# s = output so far.

```
    if i == n:                // base case
```

```
        print(list(s))
```

```
    else:
```

```
        sub(a, i+1, n, s+a[i])  // recursive assumption  
        sub(a, i+1, n, s)      } self work
```

# main

```
a = ['a', 'b', 'c']
```

```
sub(a, 0, len(a), "")
```

Q2) Print n length binary string not having any consecutive 1's.  
eg n=2      10, 01, 00

code

```
def sub(i, n, s, t):
```

```
    if (i == n):
```

```
        print(s)
```

```
    else:
```

```
        if (t == 1):
```

```
            sub(i+1, n, '0'+s, 0)
```

```
        else:
```

```
            sub(i+1, n, '0'+s, 0)
```

```
            sub(i+1, n, '1'+s, 1)
```

```
sub(0, 3, "", 0)
```

s → output so far

t → keeps track of last digit

eg if last digit added is 0,  
we can add both 0 & 1  
(t=0)

## Concept

### Recursion on Grid

Q3) You are given a 2D grid, you are at top left point and you need to reach bottom right point. (can go only right or down)

1<sup>st</sup>) find total number of paths possible.

```
def sub(n, m):
```

```
    if (n == 1 or m == 1):
```

```
        return 1
```

```
    else:
```

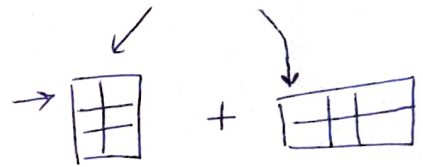
```
        return sub(n, m-1) + sub(n-1, m)
```

```
print(sub(3, 3))
```

n = rows  
m = columns



A → B



2<sup>nd</sup>) Print all paths also

```
def sub(n, m, s, l):
```

```
    if (n == 1):
```

```
        print(s + "R" * (l - len(s)))
```

```
        return(1)
```

```
    elif (m == 1):
```

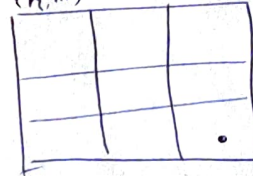
```
        print(s + "D" * (l - len(s)))
```

```
        return 1
```

```
    else:
```

```
        return sub(n, m-1, s+"R", l) +  
               sub(n-1, m, s+"D", l)
```

(n, m)



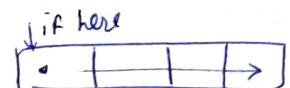
(1, 1)

can go Right & Down

also see

Size of all paths  
= (n-1) + (m-1)

also



So only  
"RRR"

if



only "DDD"

n=4

m=5

print(sub(n, m, "", n-1+m-1))



3rd) Lets even allow Diagonal move

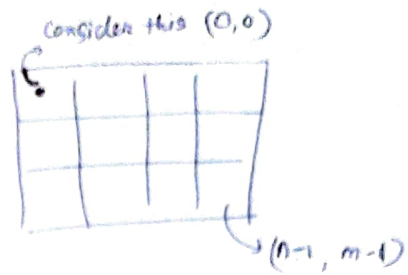
Using another method

if



so if  $i \geq n$   
return

Same for  $j \geq m$   
return.



```
def sub(i, j, n, m, s):
```

```
    if (i == n-1 and j == m-1):
```

```
        print(s)
```

```
        return 1
```

```
    elif (i >= n or j >= m):
```

```
        return 0
```

```
    else:
```

```
        return sub(i, j+1, n, m, s+"R") + sub(i+1, j, n, m, s+"B")  
               + sub(i+1, j+1, n, m, s+"D")
```

// R → Right

// B → Down (bottom)

// D → Diagonal

n=3

m=3

print(sub(0,0,n,m,""))

// total 13 paths