Lecture 4 Radu Marculescu CMV

Outline 1. Phase transitions in network traffic. 2. Phase transitions in graphs.

Note: The paper today addresses the first topic.

1. Phase transitions

· First, we consider phase transitions in network traffic (from a low to a highly congested traffic)

· Mordel

Processor v/ unlimited
que ues; gerrerates (Poisson
traffic to candom dest

router (only)
node w/
least #

use a fitness

Routing: Intermediate made selected based on shortest path tollowing either a determ or pub. routing

 $P(A) = \frac{e^{-\beta \times A}}{e^{-\beta \times A} + e^{-\beta \times B}}$

P(A) + P(B) = 1

 $P(R) = \frac{e^{-\beta X_R}}{e^{-\beta X_A} - \rho X_B}$

· Results

N=100 N=25 Phase transitions 2c, < 2c,

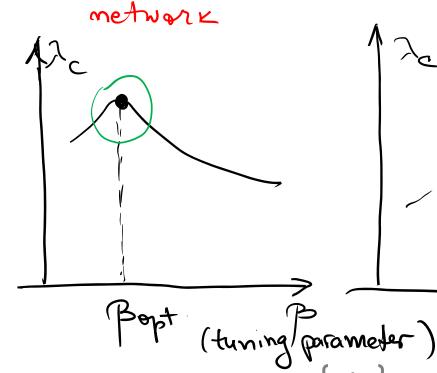
determ routing

average lifetime of a packet (avg. time between (L)) sending and receiving a packet)

de term.

N=25; prob. us. determ.

· An emergent (collective) be havior of routers decides the congestion in the



potential for fault-tolerant communication

% pws. routers

3

2. Kandom graphs

Basic idea: Consider n disconnected rertices (nordes). Each pair of modes
gets connected with a certain probability
P>0. This produces a statistical
ensemble of all possible graphs
of n modes of n modes.

· Typically, large n are considered while keeping the mean obeque z constant:

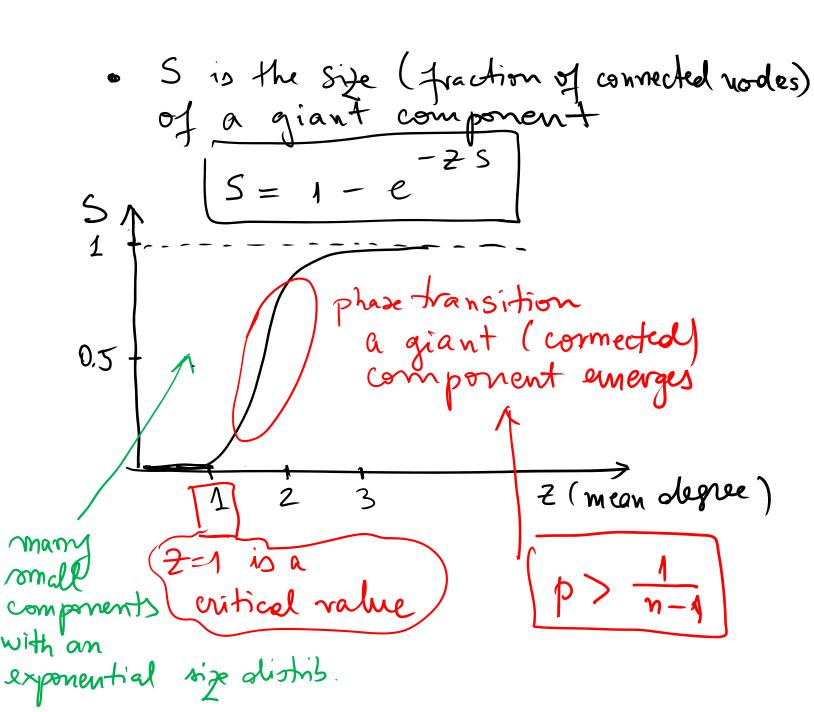
lach node gets potentially (n-1) links

2 = p(n-1) (mean degue)

 $P_{k} = \binom{n}{k} p^{k} (1-p)^{n-k} \approx \left(\frac{2^{k}-2}{k!}\right)^{n-k}$

prob. of having (random a node of k degree graph)

Note: This decreases very fast with k, prob. of highly connected resolutions is very small!



· S plays the role of an order parameter in this phase bransition;

S ~ 12-1 B Power law size distribution