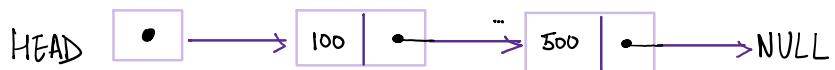


# Linked List

19 March 2024 00:17

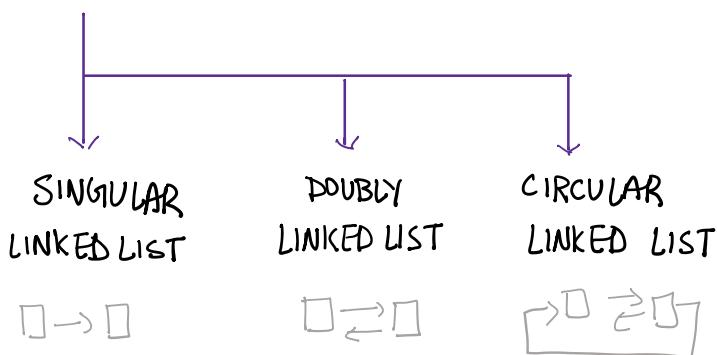
## LINKED LISTS

↳ collection of objects each containing DATA and REFERENCE to  
at least one other node.  
NODE



- A linked list is a series of connected nodes
  - Each node contains — data & pointer to the next node in list
- HEAD → [ ] → [ ] → TAIL | NULL
- It can grow & shrink in size during execution of a program
  - Doesn't waste any memory space
  - Can be made just as long as required

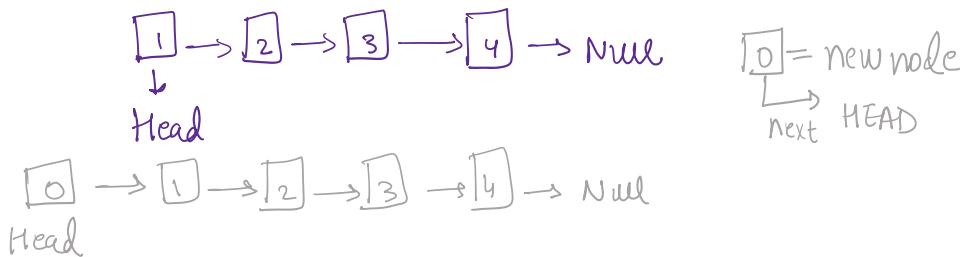
## TYPES OF LIST



## OPERATIONS OF LIST

IsEmpty	Determine whether the list is empty or not
InsertNode	Insert new node at particular position
FindNode	Find a new node with a given value
DeleteNode	Delete a node with a given value
DisplayList	Print all the nodes in the list

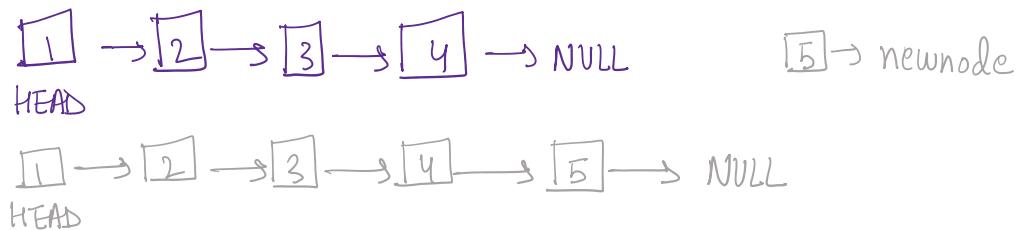
## ① Inserting a Node at the Beginning of a List



i) Create Newnode

- ii) If list is empty the head will point to the newnode , Else:
- iii) newnode → next = head
- iv) head = newnode

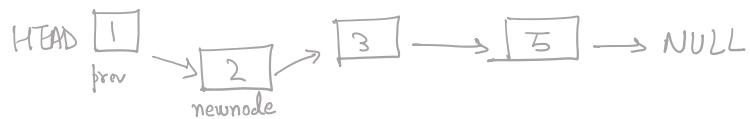
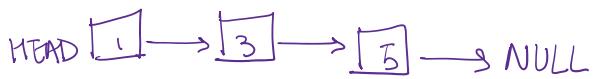
## ② Insert a Node at the End of a List



i) Create Newnode

- ii) Traverse the list starting from HEAD until the node where next pointer is NULL
- iii) Add newnode at the next pointer of last node
- iv) Set next pointer of newnode as NULL.

## ③ Insert a Node after a given Node



③ Create newnode

- (ii) Declare node after which the newnode is to be inserted as 'prev'
- (iii) newnode → next = prev → next
- (iv) prev → next = newnode

④ Print the contents of Each Node of the list

- (i) Start
- (ii) Set Ptr = Start
- (iii) Repeat (iv) of (v) until Ptr ≠ NULL
- (iv)      print Ptr → Data
- (v)      set ptr = ptr → Link
- (vi) Stop

⑤ Search for an Item

- (i) Set Ptr = Start
- (ii) While Ptr ≠ NULL, repeat step 3 -
- (iii)      if item == Ptr → Data, then
- (iv)            set loc = Ptr & Exit
- (v)      else
- (vi)            set ptr = ptr → Link
- (vii) Set loc = NULL
- (viii) Stop

Algorithms -

- (i) Put/Display all the elements of list
- (ii) Insert Node at - Beginning, End, Middle
- (iii) Search for a Node
- (iv) Delete node from Beginning, End or Middle

## SINGLE LINKED LIST

① Algorithm to Insert a node at the beginning of a Single Linked List

## SINGLE LINKED LIST

i) Algorithm to Insert a node at the beginning of a Single Linked List

Step 1 : Start

Step 2 : Create newnode and store the desired data in newnode

Step 3 : If list is empty -

    3.1 : Set the head of the list as new node  
                head → newnode

Step 4 : If list is not empty -

    4.1 : Set the newnode to point to the current head of the list  
                newnode → head

    4.2 : Set newnode as the new head of the list  
                newnode = head

Step 5 : Stop / End

ii) Algorithm to Insert a node at the End of a Single Linked List

Step 1 : Start

Step 2 : Create a newnode and store the desired data in newnode

Step 3 : Traverse the list until the lastnode is reached

Step 4 : Set lastnode to point to the newnode  
                lastnode → newnode

Step 5 : Set newnode to point to NULL  
                newnode → NULL

Step 6 : Stop / End

iii) Algorithm to Insert a node in the middle of a Single Linked List

Step 1 : Start

Step 2 : Create newnode and store the desired data in newnode

Step 3 : Traverse the list to find the two nodes - prenode and nextnode;  
                between where the newnode is to be inserted

Step 4 : Set prenode to point to newnode  
                prenode → newnode

Final E : Set nextnode to point to nextnode.

$\text{prevnode} \rightarrow \text{newnode}$

Step 5 : Set newnode to point to nextnode  
 $\text{newnode} \rightarrow \text{nextnode}$

Step 6 : Stop/End

(iv) Algorithm to print all items of a single linked list

Step 1 : Start

Step 2 : Set head as currentnode.

Step 3 : While currentnode is NOT NULL :

    3.1 : Print the data of currentnode

    3.2 : Move to the nextnode

    3.3 : Set nextnode as currentnode

Step 4 : Stop/End

(v) Algorithm to search for a particular node of a single linked list

Step 1 : Start

Step 2 : Set currentnode to the head of the list

Step 3 : While the currentnode is NOT NULL

    3.1 : If the data of the currentnode matches the key

        3.1.1 : Return true

    3.2 : Move to the nextnode

    3.3 : Set currentnode to nextnode

Step 4 : If the key is not found after traversing the entire list

    4.1 : Return False

Step 5 : Stop/End

(vi) Algorithm to delete a node from the beginning of a single linked list

Step 1 : Start

Step 2 : If the list is empty

    2.1 : Return without performing any operation

Step 3 : Traverse to the nextnode ; node that head points next to.

else : Set nextnode as head

Step 3 : Traverse to the nextnode ; node that head points next to.

Step 4 : Set nextnode as head

Step 5 : Delete the previous head

5.1 : Release the memory allocated to the previous head

Step 6 : Stop/End

(vii) Algorithm to delete a node from the end of a Single Linked List

Step 1 : Start

Step 2 : If the list is empty

2.1 : Return without performing any operation

Step 3 : Traverse to the End of the list

3.1 : Set the second-lastnode to point to NULL

Step 4 : Delete the lastnode

4.1 : Release the memory allocated to the previous lastnode

Step 5 : Stop/End

(viii) Algorithm to delete a node from the middle of a Single Linked List

Step 1 : Start

Step 2 : If the list is empty

2.1 : Return without performing any operation

Step 3 : Traverse to the node that is to be deleted ; say currentnode.

Step 4 : Set the node before the currentnode to point to the node after the currentnode

Step 5 : Delete currentnode

5.1 : Release the memory allocated to the deleted node

Step 6 : Stop/End

## DOUBLE LINKED LIST

(i) Algorithm to insert a node at the beginning of a double linked list

Step 1 : Start

- v v  
 Step1 : Start  
 Step2 : Create newnode and add the desired data to newnode  
 Step3 : If the list is empty  
     3.1 : Set the head and tail to newnode  
 Step4 : If the list is not empty  
     4.1 : Set the next pointer of newnode to point to the current head  
             newnode → next = head  
     4.2 : Set the previous pointer of current head to point to newnode  
             head → prev = newnode  
     4.3 : Set newnode as head  
 Step5 : Stop/End

(ii) Algorithm to Insert a node at the end of a double linked list

- Step1 : Start  
 Step2 : Create newnode and store desired data in newnode  
 Step3 : If the list is empty  
     3.1 : Set the head and tail to newnode  
 Step4 : If the list is not empty  
     4.1 : Traverse to the end of the list to lastnode  
     4.2 : Set lastnode next pointer to point to newnode  
             lastnode → next = newnode  
     4.3 : Set newnode previous pointer to point to the lastnode  
             newnode → prev = lastnode  
     4.4 : Set newnode next pointer to point to NULL  
             newnode → next = NULL  
 Step5 : Stop/End

(iii) Algorithm to Insert a node in the middle of a double linked list

- Step1 : Start  
 Step2 : Create newnode & store the desired data in newnode  
 Step3 : Traverse the list to the nodes where the newnode is to be inserted  
         . . . . . " "

Step 3 : Traverse the list to the nodes where the newnode is to be inserted between; say prenode and nextnode

- 3.1 : Set prenode next pointer to point to newnode  
 $\text{prenode} \rightarrow \text{next} = \text{newnode}$
- 3.2 : Set newnode previous pointer to point to prenode  
 $\text{newnode} \rightarrow \text{prev} = \text{prenode}$
- 3.3 : Set newnode next pointer to point to nextnode  
 $\text{newnode} \rightarrow \text{next} = \text{nextnode}$
- 3.4 : Set nextnode previous pointer to point to newnode  
 $\text{nextnode} \rightarrow \text{prev} = \text{newnode}$

Step 4 : Stop / End

(iv) Algorithm to print all items of a doubly linked list

Step 1 : Start

Step 2 : Set the head to currentnode

Step 3 : While currentnode is NOT NULL

- 3.1 : Print all the data of the currentnode

- 3.2 : Move to the next node

- 3.3 : Set nextnode to currentnode

Step 4 : Stop / End

(v) Algorithm to search for a node in double linked list

Step 1 : Start

Step 2 : Set head to currentnode

Step 3 : while currentnode is NOT NULL

- 3.1 : If the data of currentnode matches the key

- 3.1.1 : Return true

- 3.2 : Move to the next node

- 3.3 : Set nextnode to currentnode

Step 4 : If the key is not found while traversing the entire list

- 4.1 : Return False

Step 5 : Stop / End

(vii) Algorithm to delete a node from the beginning of a Double Linked List

Step 1 : Start

Step 2 : If the list is empty

2.1 : Return without performing any operation

Step 3 : If the list is not empty

3.1 : Traverse to the second node of the list

3.2 : Set previous pointer of second node as NULL  
 $\text{secondnode} \rightarrow \text{prer} = \text{NULL}$

3.3 : Set secondnode as head

3.4 : Delete the previous head

3.4.1 : Release the memory allocated for the deleted node

Step 4 : Stop / End

(viii) Algorithm to delete a node from the end of a Double Linked List

Step 1 : Start

Step 2 : If the list is empty

2.1 : Return without performing any operation

Step 3 : If the list is not empty

3.1 : Traverse to the end of the list to the secondlast node

3.2 : Set the secondlastnode next pointer to NULL

$\text{secondlastnode} \rightarrow \text{next} = \text{NULL}$

3.3 : Delete the lastnode

3.3.1 : Release the memory allocated for the deleted node

Step 4 : End / Stop

(ix) Algorithm to delete a node from the middle of a Double Linked List

Step 1 : Start

Step 2 : If the list is empty

2.1 : Return without performing any operation

" " " " "

- 1
- 2.0 : Return without performing any operation
- Step 3 : If the list is not empty
- 3.1 : Traverse the list to the node that is to be deleted; say curnode
  - 3.2 : Set the node previous to curnode, say prevnode to point next to the node after the curnode, say nextnode  
 $\text{prevnode} \rightarrow \text{next} = \text{nextnode}$
  - 3.3 : Set nextnode to point previous to prevnode  
 $\text{nextnode} \rightarrow \text{prev} = \text{prevnode}$
  - 3.4 : Delete curnode
  - 3.4.1 : Release the memory allocated to curnode
- Step 4 : Stop / End

## CIRCULAR SINGLY LINKED LIST

- i) Algorithm to insert a node at the beginning of Circular Singly Linked List
- Step 1 : Start
- Step 2 : Create newnode & store the desired data in newnode
- Step 3 : If the list is empty
- 3.1 : Set next pointer of newnode to point to itself  
 $\text{newnode} \rightarrow \text{next} = \text{newnode}$
  - 3.2 : Set newnode as head
- Step 4 : If the list is not empty
- 4.1 : Set the next pointer of newnode to point to the current head  
 $\text{newnode} \rightarrow \text{next} = \text{head}$
  - 4.2 : Traverse the list and set the next pointer of lastnode to point to newnode  
 $\text{lastnode} \rightarrow \text{next} = \text{newnode}$
  - 4.3 : Set newnode as Head
- Step 5 : End / Stop

ii) Algorithm to insert a node at the end of a Singly Circular List

Step 5 is the last step

(ii) Algorithm to Insert a node at the End of a Single Circular Linked List

Step 1 : Start

Step 2 : Create newnode & store the desired data in newnode

Step 3 : If list is empty

3.1 : Set the next pointer of new node to point to itself  
 $\text{newnode} \rightarrow \text{next} = \text{newnode}$

3.2 : Set newnode as head

Step 4 : If list is not empty

4.1 : Traverse to the end of the list and set the next pointer of lastnode to newnode

$\text{lastnode} \rightarrow \text{next} = \text{newnode}$

4.2 : Set the next pointer of new node to point to head  
 $\text{newnode} \rightarrow \text{next} = \text{head}$

Step 5 : Stop / End

(iii) Algorithm to Insert a node in the middle of a Single Circular Linked list

→ Same as the Algorithm to Insert a node in the middle of a Singly Linked list

(iv) Algorithm to print all the data from a single circular linked list

→ same everytime

(v) Algorithm to search for a particular node in Single circular Linked list

→ same as well

(vi) Algorithm to delete a node from the Beginning of a Single Circular Linked list

Step 1 : Start

Step 2 : If list is empty

2.1 : Return without performing any operation

Step 3 : If list is not empty

3.1 : Set the next pointer of the lastnode to the secondnode of the list  
 $\text{lastnode} \rightarrow \text{next} = \text{secondnode}$

3.2 : Set second node as head

3.3 :  $\text{K} \rightarrow \text{L} \rightarrow \text{M} \dots \text{O} \rightarrow \text{N}$

3.2 : Set second node as head

3.3 : Delete the previous head

3.3.1 : Release the memory allocated to previous head

Step 4 : Stop / End

(vii) Algorithm to delete a node from the End of the Single Circular Linked List

Step 1 : Start

Step 2 : If list is empty

2.1 : Return without performing any operation

Step 3 : If list is not empty

3.1 : Set next pointer of the secondlastnode to head

Secondlastnode  $\rightarrow$  next = head

3.2 : Delete lastnode

3.2.1 : Release the memory allocated to lastnode

Step 4 : End / Stop

## DOUBLE CIRCULAR LINKED LIST

(i) Algorithm to Insert a Node at the Beginning of Double Circular Linked List

Step 1 : Start

Step 2 : Create newnode & insert desired data in newnode

Step 3 : If the list is empty

3.1 : Set the next and previous pointer of newnode to point to itself

newnode  $\rightarrow$  next = newnode

newnode  $\rightarrow$  prev = newnode

3.2 : Set newnode as head

Step 4 : If the list is not empty

4.1 : Set the next pointer of lastnode to newnode

lastnode  $\rightarrow$  next = newnode

4.2 : Set the previous pointer of newnode to lastnode

newnode  $\rightarrow$  prev = lastnode

4.3 : Set the next pointer of newnode to head

newnode  $\dots$  L<sub>1</sub>

4.2 : Set the next pointer of newnode to head

$$\text{newnode} \rightarrow \text{next} = \text{head}$$

4.4 : Set the previous pointer of head to newnode

$$\text{head} \rightarrow \text{prev} = \text{newnode}$$

4.5 : set newnode as head

Step 5 : Stop/End

(ii) Algorithm to Insert a node at the end of Double Circular Linked List

Step 1 : Start

Step 2 : Create newnode & store the desired data in newnode

Step 3 : If the list is empty

3.1 : Set the next and previous pointer of newnode to point to itself

$$\text{newnode} \rightarrow \text{next} = \text{newnode}$$

$$\text{newnode} \rightarrow \text{prev} = \text{newnode}$$

3.2 : Set newnode as head

Step 4 : If the list is not empty

4.1 : Set the next pointer of lastnode to newnode

$$\text{lastnode} \rightarrow \text{next} = \text{newnode}$$

4.2 : Set the previous pointer of newnode to lastnode

$$\text{newnode} \rightarrow \text{prev} = \text{lastnode}$$

4.3 : Set the next pointer of newnode to head

$$\text{newnode} \rightarrow \text{next} = \text{head}$$

4.4 : Set the previous pointer of head to newnode

$$\text{head} \rightarrow \text{prev} = \text{newnode}$$

Step 5 : Stop/End

(iii) Algorithm to delete a node from the Beginning of a Double Circular Linked list.

Step 1 : Start

Step 2 : If the list is empty

2.1 : Return without performing any operations

Step 3 : If the list is not empty

3.1 : Set nextpointer of lastnode to secondnode

$$\text{lastnode} \rightarrow \text{next} = \text{secondnode}$$

- 3.1 : Set nextpointer of lastnode to secondnode  
 $\text{lastnode} \rightarrow \text{next} = \text{secondnode}$
- 3.2 : Set previous pointer of secondnode to lastnode  
 $\text{secondnode} \rightarrow \text{prev} = \text{lastnode}$
- 3.3 : Set secondnode as head

Step 4 : Delete previous head

- 4.1 : Release the memory allocated to the deleted node

Step 5 : Stop/End

(iv) Algorithm to delete a node from the End of a Double Circular Linked List

Step 1 : Start

Step 2 : If the list is empty

- 2.1 : Return without performing any operations

Step 3 : If the list is not empty

- 3.1 : Set the next pointer of secondlast node to head  
 $\text{secondlastnode} \rightarrow \text{next} = \text{head}$

- 3.2 : Set the previous pointer of head to secondlastnode  
 $\text{head} \rightarrow \text{prev} = \text{secondlastnode}$

Step 4 : Delete lastnode

- 4.1 : Release the memory allocated to the deleted node

Step 5 : Stop/End