Bounded-Buffer Problem

- ☐ How many semaphores do we need?
- Again -- Three requirements:
 - The Producer (P) must not insert item when buffer is full
 - The Consumer (C) must not remove item when buffer is empty
 - P and C should not insert and remove must at the same time

```
Producer:
                                     Consumer:
while (true) {
                                     while (true) {
    /* produce an item in
                                         while (counter = 0);
                                         /* do nothing */
    next produced */
    while (counter == BUFFER SIZE);
                                         next consumed = buffer[out];
    /* do nothing */
                                         out = (out + 1) % BUFFER SIZE;
    buffer[in] = next produced;
                                         counter--;
    in = (in + 1) % BUFFER SIZE;
                                         /* consume the item in next
    counter++;
                                         consumed */
```

Bounded-Buffer Problem

☐ Semaphore mutex

☐ Semaphore **full**

Initial values for three semaphores?

- ☐ Semaphore **empty**
 - Again -- Three requirements:
 - The Producer (P) must not insert item when buffer is full
 - The Consumer (C) must not remove item when buffer is empty
 - P and C should not insert and remove must at the same time

Bounded-Buffer Problem

- ☐ Semaphore mutex initialized to the value 1
 - > To protect count operation
- Semaphore full initialized to the value 0
 - Number of full buffers
- ☐ Semaphore **empty** initialized to the value *n*
 - Number of empty buffers

```
int n;
semaphore mutex = 1;
semaphore empty = n;
semaphore full = 0
```

Let's Make Pseudocode First

Producer with Semaphore

Producer

```
while (true) {
    /* produce an item in
    next produced */

    while (counter == BUFFER_SIZE);
    /* do nothing */

    buffer[in] = next_produced;
    in = (in + 1) % BUFFER_SIZE;
    counter++;
}
```

Consumer with Semaphore

Consumer

```
while (true) {
    while (counter == 0);
    /* do nothing */

    next_consumed = buffer[out];
    out = (out + 1) % BUFFER_SIZE;

    counter--;
    /* consume the item in next
    consumed */
}
```

Consumer (w/ Semaphore)

Producer and Consumer

```
while (true) { Producer (w/ Semaphore)
  /* produce an item in next_produced */
  wait(empty);
  wait(mutex);
  /* add next_produced to the buffer */
  signal(mutex);
                         while (true) {    Consumer (w/ Semaphore)
  signal(full);
                           wait(full);
                           wait(mutex);
                           /* remove an item from buffer to next_consumed */
                           signal(mutex);
                           signal(empty);
                           /* consume the item in next_consumed */
```

Let's Implement "Bounded-Buffer Problem"

☐ Let's use two semaphores and one mutex lock

```
int n;
semaphore mutex = 1;
semaphore empty = n;
semaphore full = 0
```

```
8 int n_items_in_buffer = 0;
9 pthread_mutex_t lock;
10 sem_t sem_empty;
11 sem_t sem_full;
```

APIs

- > pthread mutex init
- > pthread_mutex_lock & pthread_mutex_unlock
- > sem_init & sem_destroy
- > sem_wait & sem_post

Base Code

👫 nike.cs.uga.edu - PuTTY

```
39 int main(int argc, char *argv[])
                                                                                    int i;
🗬 nike.cs.uga.edu - PuTTY
 1 #include <stdio.h>
                                                                                    pthread_t p;
 2 #include <string.h>
                                                                               44
                                                                                    pthread t c[10];
 3 #include <stdlib.h>
 4 #include <pthread.h>
 5 #include <unistd.h>
                                                                                    for(i = 0; i < NUM CONSUMER; i++)</pre>
 6 #include <semaphore.h>
                                                                                       pthread_create(&(c[i]), NULL, consumer, NULL);
 8 int n items in buffer = 0;
                                                                               49
 9 pthread_mutex_t lock;
                                                                                    pthread_create(&p, NULL, producer, NULL);
10 sem t sem empty;
                                                                              51
11 sem t sem full;
                                                                                    for(i = 0; i < NUM_CONSUMER; i++)</pre>
13 int MAX SLOT = 5;
                                                                                      pthread join(c[i], NULL);
14 int NUM CONSUMER = 3;
                                                                               54
                                                                                    pthread join(p, NULL);
16 void *producer() {
                                                                                    return 0;
    while(1) {
                                                                               57 }
       printf("[Producer] produce an item\n");
                                                                                                                          57,1
       n items in buffer++;
       printf("[Producer] added the item to the buffer. # items in buffer = %d\n"
             , n items in buffer);
       usleep(rand()%500000); // processing overhead
26 }
28 void *consumer(void *arg) {
    printf("[Consumer %d] joined.\n", syscall(186)-getpid());
    while(1) {
       n items in buffer--;
       printf("[Consumer %d] removed an item from the buffer. # items in buffer = %d\n"
             , syscall(186)-getpid(), n_items_in_buffer);
       usleep(rand()%800000); // processing overhead
37 }
                                                                    14,1
                                                                                  Top
```

Bot

Real-World Example of Bounded Buffer Problem

- □ IoT Data Processing
- Audio/Video player:
 - Thread #1: Network/data transmission
 - Thread #2: Decode packets and display data from shared buffer
- Web Server: Master threads and slave threads
 - Listener Threads: Producers
 - Worker Threads: Consumer

Classical Problems of Synchronization

- Classical Problems in Synchronization
 - Bounded-Buffer Problem
 - Readers and Writers Problem
 - Dining-Philosophers Problem

Readers-Writers Problem

- ☐ Classic synchronization problem in data sharing/management system. e.g., DBMS
- □ A data set is shared among a number of concurrent processes
 - Readers only read the data set; they do not perform any updates
 - Writers can both read and write

Readers-Writers Problem - Requirements

- Two or more readers access the shared data at the same time, no adverse effects will result
- No simultaneous access from a writer and some other readers/writers -- no reading while data is being updated
- Writers should have exclusive access to the shared data.
 - In other words, if a writer is updating (writing) data, no other readers or writers are allowed to access the data.

Readers-Writers Problem

- You have a global variable count
- Reader
 - Do "printf count"
- Writer
 - > Do "count = count*10;"
- Let's use mutex or binary semaphore